

# Lecture 7: Searching Genomes and Genome Indexing

## Student Handout & In-Class Exercises

**Course:** BINF301 — Computational Biology

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**Date:** 4/2/2026

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## 1 Text Search Basics (Slides 3–6)

Genome search is defined as: Given a pattern  $P$  of length  $n$  and a text  $T$  of length  $m$ , find all occurrences of  $P$  in  $T$ .

Examples:

- grep-like text search
- text editors (Ctrl+F)
- mapping reads to a reference genome

**Naïve Approach (Slides 4–6):** Try every alignment of  $P$  in  $T$ .

- Worst case:  $O(mn)$
- Best case:  $O(m)$

## 2 Boyer–Moore Algorithm (Slides 8–15)

A fast exact matching algorithm using:

- **Bad character rule:** After mismatch, shift pattern so mismatch aligns with rightmost matching character.
- **Good suffix rule:** Reuse matched suffix to shift pattern optimally.

## 3 Genome Indexing (Slides 17–18)

Indexing preprocesses the text  $T$  to support fast lookups for many patterns. Essential for:

- read mapping
- repeated queries on a reference genome

## 4 k-mer Tables (Slides 19–23)

A basic genome index storing each  $k$ -mer and the positions where it appears.

Advantages:

- fast lookup (especially via hashing)

Limitations:

- best when pattern length equals  $k$
- not suitable for all substring lengths

## 5 Hash Tables (Slides 28–29)

Hash tables store key-value pairs with near-constant lookup time.

Used for:

- $k$ -mer count tables
- presence/absence queries

## 6 Suffix Trees and Suffix Arrays (Slides 32–38)

### 6.1 Suffix Trees

Compressed tries of all suffixes. Fast but extremely memory-heavy (on the order of 15 bytes per base).

### 6.2 Suffix Arrays

A space-efficient alternative:

- store sorted suffix positions
- support fast binary search
- require much less space

## 7 Burrows–Wheeler Transform (Slides 39–41)

A reversible transform that groups identical characters together, making the text more compressible.

## 8 FM-index (Slides 44–52)

FM-index = BWT + rank/select + partial suffix array sampling.

Properties:

- extremely low memory usage (about 0.5 bytes/character)
- supports fast backward search

## 9 Hands-On Exercises with Solutions

### 9.1 Exercise 1 — Naïve Search Simulation (Slides 3–6)

Text: ATGATCATGAC Pattern: ATG

**Task:** List match positions.

### 9.2 Exercise 2 — Boyer–Moore Skip (Slides 9–12)

Pattern: ATCGA Mismatch char: T

**Task:** Compute skip distance using the bad character rule.

### 9.3 Exercise 3 — Build a 3-mer Table (Slides 19–21)

Text: GATATAGA

**Task:** List all 3-mers and positions.

### 9.4 Exercise 4 — Suffix Array Binary Search (Slides 36–38)

0	\$
1	A\$
2	CTA\$
3	GCTA\$
4	TA\$

**Task:** Does “TA” occur?

### 9.5 Exercise 5 — BWT Construction (Slides 39–41)

Text: GAT\$

**Task:** Build BWT.

### 9.6 Exercise 6 — FM-index Backward Step (Slides 45–48)

BWT L = T G \$ A A Pattern = GA

**Task:** Compute initial range for final character A.