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"Brian Beckman: Don't Fear the Monad"

Available on Youtube [3].

Dr. Beckman, astrophysicist and senior software engineer, begins with a basic introduction to functional programming as a concept. Most notably, he focuses on the concept of functions as being **replaceable** by table-lookups.

1.1 Outline (7:50)

- 1. Functions
- 2. Monoids
- 3. Functions
- 4. Monads

1.2 Notation (8:25)

ONAL NOTATION:

- $int \ x = x \in int$
- x :: int
- int f(int x)
- $f :: int \rightarrow int$

PE VARIABLE $a: \forall a$

- A x
- x :: a
- static A f < A > (A x)
- $f :: a \rightarrow a$

1.2.1 Composition

Given:

in imperative style function composition might appear as: f(g(a)) or in reverse: g(f(a)).

In functional style, function application appears as: g a and composition can be shown as: f(g a). Parenthesis are necessary due to partial application being left associative. For example, f h g is applied as though (f h) g.

It is also possible to use a composition operator, \circ , to imply composition: $(f \circ g)$ a. So, given the above 1.2.1, we can deduce:

```
h = (f \circ g) \ a = f \circ g
h :: a \rightarrow a
```

This does confuse the concepts of a as argument and a as type, but the point remains clear, I think.

1.3 Monoids (20:40)

In abstract algebra, a branch of mathematics, a monoid is an algebraic structure with a single associative binary operation and an identity element.

Monoids are studied in semigroup theory, because they are semigroups with identity. $\ddot{}$

 a Wikipedia

A Monoid is a Set with:

- 1. an associative binary operator (generally composition)
- 2. an identity value

The operator need not be commutative.

In a programming context, a Monoid guarantees type-consistency over function composition.

1.4 Monads (30:39)

Given:

```
x :: a
f :: a \rightarrow M a
g :: a \rightarrow M a
g :: a \rightarrow M a
```

 ${\it M}$ is described as a "Type Constructor."

Again, Dr. Beckman is using a to represent both a value of type a as well as the type itself a. Here he introduced the Monad "bind" operator: >>=, which he likes to call "shove":

```
f:: a \to M a g:: a \to M a -- the right hand side is g, but written with -- a lambda to preserve symmetry \lambda a \to (f \ a) >>= \lambda a \to (g \ a)
```

The reason to preserve symmetry in the above expression is that the desired expression is "bracketed" as: $\lambda a \to [(fa)>>=\lambda a \to (ga)]$ because the bind operator has type:

```
\langle \rangle \Rightarrow :: Monad m \Rightarrow m a \rightarrow (a \rightarrow m b) \rightarrow m b
```

That is, >>= accepts a Monad (M a) and returns a function from a \rightarrow M a.

The functions f, and g live in a Monoid. M a (the data) lives in a Monad.

(>>=) is the analog of function composition and, therefore, obeys the rules of a *Monoid*. Including associativity and identity.

In a Monad, identity is—in Haskell—written as:

```
return :: Monad m \Rightarrow a \rightarrow m a
```

Extended to non-uniform types:

```
g :: a \to Mb

f :: b \to Mc

\lambda a \to (g \ a) >>= \lambda b \to (f \ b) :: a \to Mc

g >>= \lambda b \to (f \ b)
```

The Maybe Monad

2.1 As Described by Computerphile [2]

```
-- Type and 2 type constuctors: Val and Div

data Expr = Val Int | Div Expr Expr

-- Val 1

-- Div (Val 6) (Val 2)

unsafe_eval :: Expr → Int
unsafe_eval (Val n) = n
unsafe_eval (Div x y) = div (eval x) (eval y)

-- eval (Div (Val 6) (Val 2))
```

What if Div is passed zero? The program will crash. So, error-checking is necessary.

Abstracting the pattern of case checking Maybe values can be represented as:

```
Without 'do' notation

eval' :: Expr \rightarrow Maybe Int
eval' (Val n) = return n
eval' (Div x y) =
eval' x >>= (\lambdan \rightarrow
eval' y >>= (\lambdam \rightarrow
eval' y >>= (\lambdam \rightarrow
safediv n m))

With 'do' notation

eval' :: Expr \rightarrow Maybe Int
eval', (Val n) = return n
eval', (Div x y) = do
n \leftarrow eval', x
m \leftarrow eval', x
safediv n m
```

Other Monads

The Monad type [4, p. 402]:

Monad comes with some default definitions:

```
m >> k = m >>= \lambda_{-} \rightarrow k
fail s = \text{error } s
```

From this definition it can be seen that >> acts like >>=, except that the value returned by the first argument is discarded rather than being passed to the second argument. [4, p. 403]

3.1 The Identity Monad [4, p. 404]

The identity monad takes a type to itself with definitions:

```
m >>= f = f m
return = id
```

3.2 Definition of Maybe

```
instance Monad Maybe where
  (Just x) >>= k = k x
Nothing >>= k = Nothing
  return = Just
  fail s = Nothing
```

3.3 The List Monad

```
instance Monad [] where
    xs >>= f = concat (map f xs)
    return x = [x]
    fail s = []
```

Lists are, in fact, themselves instances of Monad.

```
fmap :: Functor f \Rightarrow (a \rightarrow b) \rightarrow f \ a \rightarrow f \ b

\langle > \Rightarrow \rangle :: Monad m \Rightarrow m \ a \rightarrow (a \rightarrow m \ b) \rightarrow m \ b

map :: (a \rightarrow b) \rightarrow [a] \rightarrow [b]
```

Using Parsec

4.1 Parsing CSV

many (noneOf ", λ n")

```
import Text.ParserCombinators.Parsec
```

[5, Ch. 16]

Input type is a sequence of characters, i.e., a Haskell *String*. *String* is the same as [*Char*]. The return value is [[*String*]]; a list of a list of Strings. We'll ignore *st* for now.

The do block implies a Monad. GenParser is a parsing monad.

many is a higher-order function that passes input repeatedly to the function passed as its argument. It collects the return values and treturns them in a list.

```
csvFile :: GenParser Char st [[String]]
csvFile =
  do result ← many line
   eof
   return result
```

A line is a list of cells followed by eol.

```
line :: GenParser Char st [String]
line =
  do result ← cells
   eol
   return result
```

```
cells :: GenParser Char st [String]
cells =
  do first ← cellContent
    next ← remainingCells
    return (first : next)
```

The choice operator, (<|>), tries the parser on the left and tries The parser on the right if the left consumes no input.

```
remainingCells :: GenParser Char st [String]
remainingCells =
   (char ',' >> cells) <|> (return [])

cellContent :: GenParser Char st String
cellContent =
```

```
eol :: GenParser Char st Char
eol = char 'λn'

parseCSV :: String → Either ParseError [[String]]
parseCSV input = parse csvFile "(unknown)" input
```

4.2 sepBy and endBy Combinators

```
import Text.ParserCombinators.Parsec

csvFile = endBy line eol
line = sepBy cell (char ',')
cell = quotedCell <|> many (noneOf ", \lambda n \lambda r")
```

A CSV cell may be either a bare cell or a quoted cell. Since a quoted cell may, itself, contain quotes (doubled for escape) a quotedCell is many quotedChar.

```
quotedCell =
   do char '"'

LULULUCONTENTU ← LMANYLYQUOTEdChar
LULULUCCHATU''' > <?> "quoteLatLendLofLcell" -- see eol below
   return content
```

The function *quotedChar* begins by consuming any character that is *not* itself a quote. If it is a quoted character, the stream must be checked for a second consecutive quote. If so, a single quote mark is returned to the result string.

Notice that try in quotedChar on the right side of <|>. Recall that I said that try only has an effect if it is on the left side of <|>. This try does occur on the left side of a <|>, but on the left of one that must be within the implementation of many.

This try is important. Let's say we are parsing a quoted cell, and are getting towards the end of it. There will be another cell following. So we will expect to see a quote to end the current cell, followed by a comma. When we hit quotedChar, we will fail the noneOf test and proceed to the test that looks for two quotes in a row. We'll also fail that one because we'll have a quote, then a comma. If we hadn't used try, we'd crash with an error at this point, saying that it was expecting the second quote, because the first quote was already consumed. Since we use try, this is properly recognized as not a character that's part of the cell, so it terminates the many quotedChar expression as expected. Lookahead has once again proven very useful, and the fact that it is so easy to add makes it a remarkable tool in Parsec.

```
\begin{array}{l} quotedChar = \\ noneOf \ "\lambda"" \\ \\ \sqcup \sqcup <|>_{\bot} try_{\sqcup}(string_{\sqcup}"\lambda"\lambda""_{\sqcup}>>_{\sqcup} return_{\sqcup}`") \end{array}
```

Parsec also includes combinators for error handling and reporting. A first attempt at an eol implementation that handles multiple line-ending styles might appear as:

```
> parseCSV "line1" Left "(unknown)" (line 1, column 6): unexpected end of input expecting ",", "\lambda n\lambda r", "\lambda r\lambda n", "\lambda r\lambda
```

The failure above is unclear and requires knowlege of the parser implementation to debug fully. The monad fail function can be used to add messaging:

```
eol' = try (string "λnλr")
  <|> try (string "λrλn")
  <|> string "λn"
  <|> string "λr"
  <|> fail "Couldn't_find_EOL"
```

This adds messaging to the result, but is still noisy and unclear:

```
> parseCSV "line1" Left "(unknown)" (line 1, column 6): unexpected end of input expecting ",", "\lambda n n", "\lambda n n", "\lambda n n" or "\lambda n" Couldn't find EOL
```

The Parsec <?> operator is designed to help here.

It is similar to <|> in that it first tries the parser on its left. Instead of trying another parser in the event of a failure, it presents an error message. Here's how we'd use it:

```
eol = try (string "<math>\lambda n \lambda r") 
 <|> try (string "\lambda r \lambda n") 
 <|> string "\lambda r" 
 <|> string "\lambda r" 
 <?> "end_of_line"
```

This has a more pleasing result:

```
> parseCSV "line1"

Left "(unknown)" (line 1, column 6):

unexpected end of input

expecting "," or end of line
```

The general rule of thumb is that you put a human description of what you're looking for to the right of <?>.

```
parseCSV :: String → Either ParseError [[String]]
parseCSV input = parse csvFile "(unknown)" input
```

Using HUnit to test the Parser

5.1 Helper Functions

```
module SpecHelpers where
import Test.HUnit

import Text.Printf

import Data.Time.Clock
import Data.Time.Calendar
import Data.Time.Format

import Text.Parsec
import Text.Parsec.String
```

5.1.1 Date Helpers

Create a date-parser via partial application of parseTimeOrError:

```
acceptOuterWhitespace = True

parseDateF :: String → String → Either String Day
parseDateF fmt s =
    case parseTimeM acceptOuterWhitespace defaultTimeLocale fmt s :: Maybe UTCTime of
    Nothing → Left $ printf "Failed_to_parse_"(%s'_with_"(%s'" s fmt
    (Just x) → Right $ utctDay x

parseTimeF :: String → String → Either String UTCTime
parseTimeF fmt s =
    case parseTimeM acceptOuterWhitespace defaultTimeLocale fmt s :: Maybe UTCTime of
    Nothing → Left $ printf "Failed_to_parse_"(%s'_with_"(%s'" s fmt
    (Just x) → Right x

showTime :: UTCTime → String
showTime = formatTime defaultTimeLocale "%e_"%b_"%Y_"%R"
```

5.1.2 Parser Testing Helpers

```
nullState = ()
testParse :: Parser a → String → String → Either ParseError a
testParse p src = runParser p nullState src

-- assertParsesTo expected p input = do
-- let result = testParse p input
-- case result of
-- (Right actual) → expected @=? show actual
-- (Left error) → assertFailure
-- (printf "Failed to parse input '%s': %s" input (show error))
```

A Type for expressing parser assertions

```
data ParserContext a =
                            :: Parser a
  ParserContext { parser
                   , input :: String
                   , expected :: String
parses :: (Show a) \Rightarrow ParserContext a \rightarrow Assertion
parses pa =
  case testParse (parser pa) ("expected:" ++ (expected pa)) (input pa) of
    (Right actual) \rightarrow assertEqual "" (expected pa) (show actual)
    (Left error) → assertFailure (show error)
instance (Show a) ⇒ Testable (ParserContext a)
  where
    test = TestCase o parses
assertParser message expected parser input =
  parses pa
   pa = ParserContext parser input expected
```

5.1.3 Writing the Tests

Table 5.1: Test Types

```
recordTests :: Test

test :: Testable t ⇒t →Test

(~:) :: Testable t ⇒String →t →Test

Creates a test from the specified Testable, with the specified label attached to it.

Since Test is Testable, this can be used as a shorthand way of attaching a STestLabel to one or more tests.

assertEqual :: (Show a, Eq a) ⇒String →a →a →Assertion

(@=?) :: (Show a, Eq a) ⇒a →a →Assertion

Asserts that the specified actual value is equal to the expected value (with the actual value on the left-hand side).
```

5.1.4 The Preamble

Declare the test module and export its tests.

```
module EmacsDiary.RecordSpec (recordTests) where

import Test.HUnit
import Text.Parsec
import Text.Parsec.String
import Data.Time.Clock
import Data.Time.Format

import SpecHelpers

import EmacsDiary.Record

recordTests = test [
   "create_ua_DiaryEntry" ~: do
   let (Right d) = parseTimeF "%e_u%b_u%Y" "7_July_2008"
   "_7_July_2008_Chicago,_IL" @=? show (DiaryEntry d "Chicago,_IL")
```

5.2 Testing the Parser

```
module EmacsDiary.ParserSpec (tests) where
import Test.HUnit
import SpecHelpers
import EmacsDiary.Parser (diary)

tests = test [
```

The Emacs Diary Parser

6.1 Parsing Dates and Times

```
import qualified EmacsDiary.Parser.Tokens as T

import Text.Parsec (
    many, string, count, spaces, digit, many1, choice, (<|>), try, sepBy1, unexpected)
import Text.Parsec.String (Parser)
import Data.Time.Clock (
    addUTCTime,
    NominalDiffTime,
    utctDayTime,
    utctDay,
    UTCTime)
import Data.Time.Calendar (fromGregorian, showGregorian, Day)
import Data.Time.Format (iso8601DateFormat, formatTime, parseTimeM, defaultTimeLocale)
import Text.Printf (printf)
import Control.Monad (void)
```

```
showTimeFormat = "%H:%M:%S%z"
locale = defaultTimeLocale
```

6.1.1 Type Definitions

A newtype declaration creates a new type in much the same way as data. The syntax and usage of newtypes is virtually identical to that of data declarations - in fact, you can replace the newtype keyword with data and it'll still compile, indeed there's even a good chance your program will still work. The converse is not true, however - data can only be replaced with newtype if the type has exactly one constructor with exactly one field inside it. ([1, NewType])

```
newtype Date = Date Day deriving (Eq)

Ill show Date using the ISO-8601 format:

instance Show Date where
show (Date d) = showGregorian d
```

Time will represent instances of UTC times.

```
data Time = Time { timeHour :: Integer
                   , timeMinute :: Integer
                   } deriving (Eq)
instance Show Time where
        show t = iso8601DateFormat (Just showTimeFormat)
instance Num Time where
  a + b = fromInteger $ (seconds a) + (seconds b)
  a * b = fromInteger $ (seconds a) * (seconds b)
  abs = fromInteger o abs o seconds
  signum = fromInteger o signum o seconds
  negate = fromInteger o negate o seconds
  fromInteger a = Time h m
   where
      i = fromInteger a
      h = \text{quot } i 3600
      m = div s 60
      s = rem i 3600
timeFromList :: [Integer] \rightarrow Time
timeFromList (h:m:_) = Time h m
seconds :: Time \rightarrow Integer
seconds (Time h m)= 3600*h + 60*m
\texttt{timeToNDiff} \; :: \; \texttt{Time} \; \rightarrow \; \texttt{NominalDiffTime}
timeToNDiff (Time h m) = fromInteger (3600*h + 60*m)
\texttt{addTime} \; :: \; \texttt{UTCTime} \; \rightarrow \; \texttt{Time} \; \rightarrow \; \texttt{UTCTime}
addTime a t = addUTCTime (timeToNDiff t) a
```

6.1.2 Parsers

```
decided :: Parser (Either String a) \rightarrow Parser a
decided p = do
        e \leftarrow p
         case e of
          (Left s) \rightarrow unexpected s
           (Right r) \rightarrow return r
timeS :: Parser Time
timeS = timeFromList <$> sepBy1 T.numeric (T.symbol ":")
dateFormats = [
  "%d_%b_%Y",
  "d_{\square}B_{\square}Y",
  "%e_%b_%Y",
  "%e_%B_%Y",
  "%F"]
--date :: Parser UTCTime
--date = decided p
-- where
-- p = asUTC <  dateS
```

```
-- dateString :: Parser String
-- dateString = do
-- void spaces
-- d \leftarrow digits
\text{--} \quad \text{m} \, \leftarrow \, \text{word}
-- y \leftarrow digits
-- return (d:(m:y))
     -- $ parseDateS (printf "%s %s %s" d m y)
dayP :: Parser Int
dayP = fromIntegral <$> T.numeric
yearP :: Parser Integer
yearP = T.numeric
months = [
  ("January", 1), ("Jan",
                                     1),
  ("February", 2), ("Feb",
                                    2),
  ("March", 3), ("Mar", ("April", 4), ("Apr", 5),
                                     3),
                                     4),
  ("June",
               6), ("Jun",
                                     6),
              7), ("Jul",
  ("July",
                                    7),
  ("August", 8), ("Aug",
                                     8),
  ("September", 9), ("Sep",
                                    9),
  ("October", 10), ("Oct",
                                    10),
  ("November", 11), ("Nov",
                                     11),
  ("December", 12), ("Dec",
                                     12)
keyValueParser :: (String, Int) \rightarrow Parser Int
keyValueParser (m,n) = try (string m) >> return n
monthP :: Parser Int
monthP = choice $ map keyValueParser months
-- date :: Parser Date
date :: Parser Day
date = do
  d \leftarrow dayP
 m \leftarrow T.lexeme monthP
 y \leftarrow yearP
 return $ fromGregorian y m d
time :: UTCTime \rightarrow Parser \ UTCTime
time datetime = addTime datetime <$> timeS
```

Bibliography

- [1] The Haskell Wiki April 2018
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