

# 8. Data Structure

九章算法IT求职面试培训 第8章

[www.jiuzhang.com](http://www.jiuzhang.com)

# Outline

## Linear Data Structure

- Queue
- Stack
- Hash

## Tree Data Structure

- Heap
- Trie

# Definition

Data Structure is a way to organize data. It provides some methods to handle data stream, e.g. insert, delete, etc.

# Queue

# Queue

Operations:

$O(1)$  Push

$O(1)$  Pop

$O(1)$  Top

Always used for BFS

# Stack

# Stack

## Operations:

Push  $O(1)$

Pop  $O(1)$

Top  $O(1)$

# Min-Stack

<http://www.lintcode.com/en/problem/min-stack/>

<http://www.jiuzhang.com/solutions/min-stack/>

Implement a stack, enable  $O(1)$  Push, Pop, Top, Min.  
Where Min() will return the value of minimum number in the stack.



# Min-Stack

Using two stacks.

The first one is the regular stack.

The second one only store minimum numbers if a smaller number comes.

# Min-Stack

## Push(x)

1. `stack.push(x)`
2. 如果 $\leq \text{minStack}$ 的最小值, 那么就`minStack.push(x)`

## Pop()

1. `stack.pop()`
2. 如果 $= \text{minStack.top}()$ , 那么就`minStack.pop()`

# Implement a queue by two stacks

<http://www.lintcode.com/en/problem/implement-queue-by-two-stacks/>

<http://www.jiuzhang.com/solutions/implement-queue-by-two-stacks/>

Implement a queue by two stacks. Support  $O(1)$  push, pop, top.

# Implement a queue by two stacks

Q.Push(x): S1-Push(x)

Q.Pop(): if S2.empty() --> S1->S2; S2.pop()

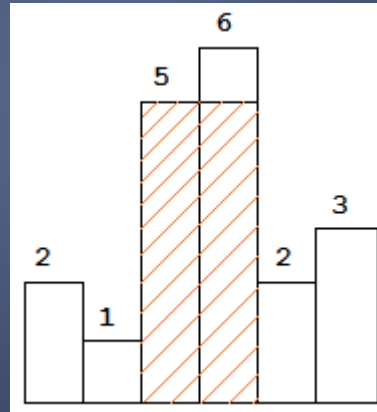
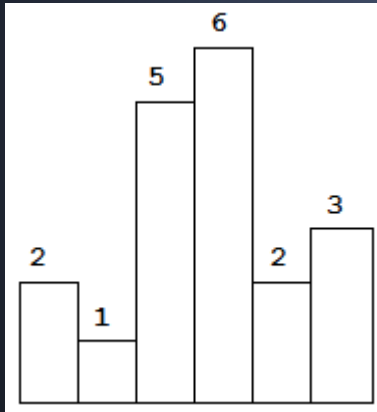
Q.Top(): Similar with Q.Pop()

# Largest Rectangle in histogram

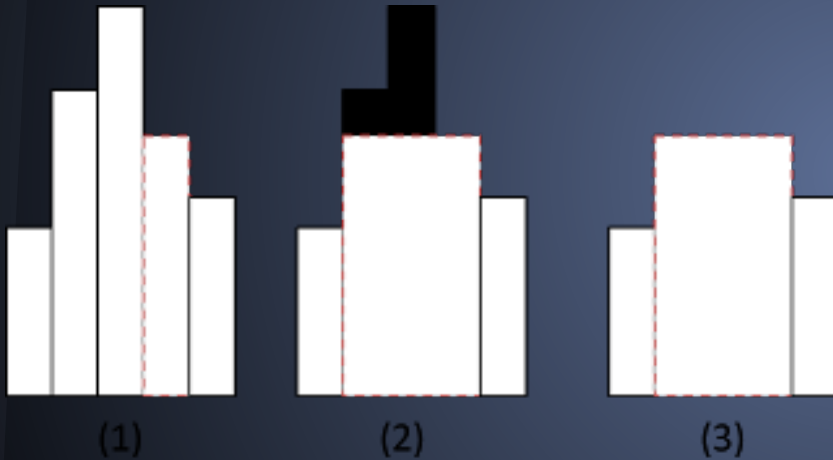
<http://www.lintcode.com/en/problem/largest-rectangle-in-histogram/>

<http://www.jiuzhang.com/solutions/largest-rectangle-in-histogram/>

# Largest rectangle in histogram



# Largest rectangle in histogram



Maintain an incremental stack:

(1) if  $a[i] > \text{stack top}$ :

push  $a[i]$  to stack

(2) if  $a[i] \leq \text{stack top}$ :

keep popping element out from stack until the top of stack is smaller than current.

# Max Tree

<http://www.lintcode.com/en/problem/max-tree/>  
<http://www.jiuzhang.com/solutions/max-tree/>



**5 min break**

# Hash

# Hash

## Operations

Insert -  $O(1)$

Delete -  $O(1)$

Find -  $O(1)$

## Hash Function

## Collision

Open Hashing (LinkedList)

Closed Hashing (Array)

# Hash Function

Typical: From string to int.

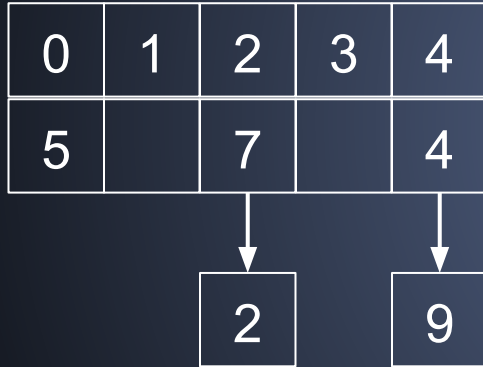
```
int hashfunc(String key) {  
    // do something to key  
    // return a deterministic integer number  
    return md5(key) % hash_table_size;  
}
```

# APR hashfunc - Magic Number 33

```
int hashfunc(String key) {  
    int sum = 0;  
    for (int i = 0; i < key.length(); i++) {  
        sum = sum * 33 + (int)(key.charAt(i));  
        sum = sum % HASH_TABLE_SIZE;  
    }  
    return sum  
}
```

# Collision - Open Hashing vs Closed Hashing

Insert: 5 7 4 2 9



0	1	2	3	4
5	9	7	2	4

Diagram illustrating Closed Hashing (Probing). The hash table has slots 0 through 4. Slot 0 contains 5, slot 1 contains 9, slot 2 contains 7, slot 3 contains 2, and slot 4 contains 4. All elements are stored directly in the hash table slots.

# Rehashing

<http://www.lintcode.com/en/problem/rehashing/>

# Java

What's Differences of:

HashTable

HashSet

HashMap

Which one is thread Safe?



# LRU Cache

<http://www.lintcode.com/zh-cn/problem/lru-cache/>  
<http://www.jiuzhang.com/solutions/lru-cache/>

Example: [2 1 3 2 5 3 6 7]

# LRU Cache

LinkedHashMap = DoublyLinkedList + HashMap

```
HashMap<key, DoublyListNode>
```

```
DoublyListNode {
```

```
    prev, next, key, value;
```

```
}
```

Newest node append to tail.

Eldest node remove from head.

# Longest Consecutive Sequence

<http://www.lintcode.com/zh-cn/problem/longest-consecutive-sequence/>

<http://www.jiuzhang.com/solutions/longest-consecutive-sequence/>

# Hash Related Questions

<http://www.lintcode.com/en/problem/subarray-sum/>

<http://www.lintcode.com/en/problem/copy-list-with-random-pointer/>

<http://www.lintcode.com/en/problem/anagrams/>

# Heap

# Heap

## Operations

Add  $O(\log N)$

Remove  $O(\log N)$

Min/Max  $O(1)$

# Heap - Implementation

Low level data structure: Dynamic Array

```
Heap {
```

```
    elems[], size;
```

```
}
```

elems[1] - root, also the minimum elem in elems.

i's left child:  $i*2$ , right child:  $i*2+1$

Internal Method:

siftup, siftdown

# Heap - Implementation

Add:

Push back to elems; size ++; Siftup;

Remove:

Replace the elem to be removed with the last elem (elems[size]); size --; Siftup and Siftdown.



# Median Number

<http://www.lintcode.com/en/problem/data-stream-median/>

# Heap Related Questions

<http://www.lintcode.com/en/problem/heapify/>

<http://www.lintcode.com/en/problem/merge-k-sorted-lists/>

# Trie

# Word Search II

<http://www.lintcode.com/en/problem/word-search-ii/>

<http://www.jiuzhang.com/solutions/word-search-ii/>

# Word Search II

Given a matrix of upper alphabets. e.g.

ACAF

ACAD

ACAE

and a dictionary. Find all words in the dictionary that can be found in the matrix.