# A.另一个爱与希望的故事

简单递推,对斐波那契数列略微修改,开一个数组记录当前台阶有没有损坏,循环的时候如果发现损坏就continue

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
typedef pair<int,int> P;
const int maxn=100000+10;
11 x[maxn],ans[maxn];
11 mod=1000000007;
int main()
    //freopen("test0.in","r",stdin);
    //freopen("test0.out","w",stdout);
    int t; cin>>t;
    while(t--)
        int n,k; cin>>n>>k;
        memset(x,0,sizeof(x));
        memset(ans,0,sizeof(ans));
        for(int i=0;i< k;i++)
            int num; cin>>num;
            x[num]=1;
        }
        ans[0]=1;
        for(int i=1;i<=n;i++)</pre>
            if(x[i]) continue;
            if(i==1) ans[i]=ans[i-1];
            else ans[i]=ans[i-1]+ans[i-2];
            ans[i]%=mod;
        }
        cout<<ans[n]<<"\n";</pre>
    return 0;
```

#### **B.BanGosu!**

按照题目意思一步一步模拟即可,我们判断距离时可以使用不开根号的形式,这样就能减少误差

```
#include<bits/stdc++.h>
using namespace std;
typedef long long ll;
vector<pair<double, double> > circle;
double distence(double x1 , double y1 , double x2 , double y2) {
```

```
return (x1 - x2) * (x1 - x2) + (y1 - y2) * (y1 - y2);
}
double getcombo(int x) {
   if(x >= 0 \& x < 100) return 1.0;
    else if(x >= 100 \& x < 200) return 1.01;
    else if(x \ge 200 \& x < 300) return 1.02;
    else if(x >= 300 \& x < 400)return 1.03;
    else return 1.04;
}
int main() {
   int n;cin >> n;
    double r;cin >> r;
    for(int i = 0; i < n; i ++) {
        double x,y;cin >> x >> y;
        circle.push_back(make_pair(x , y));
    }
    int combo = 0;
    double res = 0;
    for(int i = 0; i < n; i ++) {
        double x,y; cin >> x >> y;
        double dis = distence(x , y , circle[i].first , circle[i].second);
        if(dis < 0.04 * r * r) {
            combo ++;
            double add = getcombo(combo);
            res += 300.0*add;
        }
        else if(dis < 0.25 * r * r) {
            combo ++;
            double add = getcombo(combo);
            res += 200.0*add;
        }
        else if(dis < r * r) {
            combo ++;
            double add = getcombo(combo);
            res += 100.0*add;
        }
        else {
            combo = 0;
        }
    }
    int ans = (int)res;
    cout << ans;</pre>
}
```

# C.奇怪的引擎

对W(t)求导,得到 $P(t)=\frac{1}{2}x^{\frac{3}{2}}+sin(x)$ ,不难发现这个函数在 $(0,+\infty)$ 上单调递增,所以我们对着这个函数二分即可

```
#include <bits/stdc++.h>
using namespace std;
typedef long long l1;

double cal(double t) { return sqrt(t * t * t) / 2.0 + sin(t); }
double solve(double p)
{
```

```
int cnt = 1000;
    double l = 0, r = 1e18;
    while (cnt--)
        double mid = (1 + r) / 2.0;
        if (cal(mid) < p)</pre>
            1 = mid;
        else
            r = mid;
    }
    return 1;
}
int main()
    int t;
    cin >> t;
    while (t--)
        double p;
        cin >> p;
        double ans = solve(p);
        cout << fixed << setprecision(10) << ans<<"\n";</pre>
    return 0;
}
```

### D.Diana压缩算法

由于题目要求三个相邻的字符表示一个英文字母,并且要求字典序最小。那么我们定义一个字符变量 Ch = 'a' , 之后我们从前到后遍历一遍数组,如果发现某三个相邻字符没有对应的英文字母,那么就让 Ch 去对应它,然后输出Ch , 并且令Ch = Ch + 1; 反之,如果发现某三个相邻字符已经有了对应的英文字母,那么直接输出其对应的字母即可。这样就能保证字典序最小,而且我们最多只会使用 8 个英文字母。

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
typedef pair<int,int> P;
const int maxn=100000+10;
char ans[10];
int getnum(string s)
{
    for(int i=2;i>=0;i--) res=res*2+(s[i]-'0');
   return res;
}
int main()
    for(int i=0; i<10; i++) ans[i]='#';
    int n; cin>>n;
    n*=3;
    string s; cin>>s;
```

```
int alpha=0;
for(int i=0;i<n;i+=3)
{
    string ss;
    ss+=s[i],ss+=s[i+1],ss+=s[i+2];
    int num=getnum(ss);
    if(ans[num]=='#')
    {
        ans[num]='a'+alpha;
        alpha++;
    }
    cout<<ans[num];
}
cout<<"\n";
return 0;
}</pre>
```

### E.子矩形

我们要注意到最多只有10中字符,所以子矩形的大小不可能超过10,否则一定有重复元素;所以我们可以直接暴力求解,控制子矩形的大小,然后定位子矩形,判断其中是否有重复元素。我们可以使用标程的6重循环,也可以使用分类讨论求解

```
#include<bits/stdc++.h>
using namespace std;
typedef long long 11;
const int N = 510;
int a[N][N], numUse[10], n, m;
//判断以(x,y)为左上角,长宽分别为lenx和leny的子矩形是否有重复元素
int getNum(int x,int y,int lenx,int leny) {
   memset(numUse , 0 , sizeof numUse);
   int flag = 0;
    for(int i = x ; i \le x + lenx - 1 ; i ++) {
        for(int j = y ; j \le y + leny - 1 ; j ++) {
            if(i \le n \& j \le m \& numUse[a[i][j]] == 0) {
               numUse[a[i][j]] = 1;
            }
            else {
               flag = 1;
               break;
            }
        if(flag)break;
   if(!flag)return 1;
    return 0;
/*分类讨论,只有
1*1、1*2.....1*10
2*1、3*1.....10*1
2*2、2*3、3*2、2*4、4*2、2*5、5*2
3*3
种子矩形
*/
int getDifferent(int x,int y){
   int res = 0;
```

```
for(int i = 1; i \leftarrow 10; i \leftrightarrow 1) {
        res += getNum(x , y , 1 , i);
        res += getNum(x , y , i , 1);
    }
    res += getNum(x , y , 2 , 2);
    res += getNum(x , y , 2 , 3);
    res += getNum(x , y , 3 , 2);
    res += getNum(x , y , 2 , 4);
    res += getNum(x , y , 4 , 2);
    res += getNum(x , y , 2 , 5);
    res += getNum(x , y , 5 , 2);
    res += getNum(x , y , 3 , 3);
    return res;
int main() {
    cin >> n >> m;
    for(int i = 1; i \le n; i \leftrightarrow ++)
        for(int j = 1 ; j <= m ; j ++)
            cin >> a[i][j];
    int res = 0;
    for(int i = 1; i <= n; i ++) {
        for(int j = 1; j <= m; j ++) {
            res += getDifferent(i , j);
    }
    //前面重复算了n*m个大小为1*1的子矩形,所以减去
    cout << res - n * m << endl;</pre>
    return 0:
}
```

### F.天天爱跑步

结构体存一下信息,然后按用时升序排序,若用时相同,则按照输入顺序排序即可。很多同学排序的时候没有考虑用时相同的时候按输入顺序排序,这个可以多加一个关键字进行排序,也可以直接使用稳定排序,题目里的n最大5000,其实就是方便大家冒泡排序的。

```
#include <bits/stdc++.h>
using namespace std;
const int maxn = 5e3+10;
struct INFO {
    string name;
    int time;
    bool operator > (const INFO &a) const {
        return time>a.time;
    }
} info[maxn];
int main() {
   int n; cin >> n;
    for (int i = 1; i <= n; ++i) {
        string name; cin >> name;
        int h, m, s; scanf("%d:%d:%d", &h, &m, &s);
        info[i] = {name, h*3600+m*60+s};
    for (int i = 1; i < n; ++i)
        for (int j = 1; j <= n-i; ++j) {
```

```
if (info[j]>info[j+1]) swap(info[j], info[j+1]);
}
for (int i = 1; i<=n; ++i) cout << info[i].name << endl;
return 0;
}</pre>
```

# **G.link-cut-graph**

题目的灵感来自于abc中的一道题。首先求出1到n的最短路径经过的边并标记起来,由于总共最多删除m条边,并且每条边被删除的概率是随机的,所以如果我们每次删掉被标记的边,就重新跑一遍最短路的话,并且假设一个比较坏的情况,从1到n的最短路径要经过n条边的话,跑最短路的期望次数也不过n+1次,总的时间复杂度的期望是 $n^3$ 。

```
#include <bits/stdc++.h>
using namespace std;
#define INF 0x3f3f3f3f
#define endl '\n'
typedef long long 11;
typedef pair<int,int> P;
const int maxn = 5e2+10;
const int maxm = maxn*maxn;
int g[maxn][maxn], mp[maxn][maxn];
struct INFO {
    int u, v, w;
} info[maxm];
bool vis[maxn], flag[maxm];
int n, m, s, d[maxn], p[maxn];
void dij() {
    memset(d, 0x3f, sizeof(d));
    memset(vis, 0, sizeof(vis));
    d[1] = 0;
    for (int i = 1; i < n; ++i) {
        int id = -1;
        for (int j = 1; j \le n; ++j)
            if (!vis[j] && (id==-1 || d[id]>d[j])) id = j;
        vis[id] = 1;
        for (int j = 1; j <= n; ++j)
            if (d[j]>d[id]+g[id][j]) {
                d[j] = d[id]+g[id][j];
                p[j] = id;
            }
    }
}
int main() {
    ios::sync_with_stdio(false);
    cin.tie();
    cout.tie();
    cin >> n >> m >> s;
    memset(g, 0x3f, sizeof(g));
    for (int i = 1, u, v, w; i <= m; ++i) {
        cin >> u >> v >> w;
        info[i] = \{u, v, w\};;
        mp[u][v] = mp[v][u] = i;
        g[u][v] = g[v][u] = w;
    }
    dij();
```

```
int x = d[n];
    if (x==INF) x = -1;
    int ed = n;
    while(ed) {
        if (p[ed]) flag[mp[ed][p[ed]]] = 1;
        ed = p[ed];
    }
    for (int i = 1; i <= s; ++i) {
        vector<int> tmp;
        int num;
        cin >> num;
        int edge, f = 0;
        while(num--) {
            cin >> edge;
            if (flag[edge]) f = 1;
            INFO &t = info[edge];
            g[t.u][t.v] = g[t.v][t.u] = INF;
            tmp.push_back(edge);
        }
        if (!f) cout << x << end1;
        else {
            dij();
            if (d[n] == INF) d[n] = -1;
            cout << d[n] << endl;</pre>
        for (auto v: tmp) {
            INFO \&t = info[v];
            g[t.u][t.v] = g[t.v][t.u] = t.w;
        }
   return 0;
}
```

# H.奇怪的加法问题

xor可以理解为不进位的加法,而且本题求得是MOD2,所以考虑最后一位即可。

思路一:观察该公式,发现每个元素都出现了n - 1次,那么直接把所有元素加起来乘以n - 1然后再判奇偶即可

```
return 0;
}
```

思路二: 计算有多少个 $a_i + a_j$ 是奇数,因为只有奇数与奇数进行xor运算才会影响结果;假如有ans个(ai + aj)是奇数,如果ans为奇数,那么结果就是1,反之为0

```
#include<bits/stdc++.h>
using namespace std;
typedef long long ll;
int main() {
    int n;cin >> n;
    int sum = 0;
    for(int i = 1 ; i <= n ; i ++) {
        int x;cin >> x;
        if(x % 2)
            sum ++;//计算输入奇数的个数
    }
    ll ans = (n - sum) * sum;//计算ai + aj中奇数的个数
    if(ans % 2)cout << 1 << endl;
    else cout << 0 << endl;
}
```

# I.奇怪的加法问题

可以枚举b的质因数x,然后根据x算出比a大的x的最小倍数。

```
# include "bits/stdc++.h"
using namespace std;
using 11 = long long;
const int maxn = 1e5+10;
const int MOD = 1e9+7;
int dp[maxn], flag[maxn];
int main () {
    ios::sync_with_stdio(0);
    int __; cin >> __;
    while(__--) {
        11 a, b; cin >> a >> b;
        11 ans = (a+b-1)/b*b;
        for (11 i = 2; i*i <= b; ++i) {
            if (b%i==0) {
                ans = min(ans, (a+i-1)/i*i);
                11 j = b/i;
                ans = min(ans, (a+j-1)/j*j);
                //cout << i << ' ' << ans << endl;
            }
        }
        cout << ans-a << endl;</pre>
    return 0;
}
```

# J.扫描线

```
#include <bits/stdc++.h>
using namespace std;
const int maxn = 1e3+10;
int n, g[maxn][maxn];
int calc0(int i, int j) {
    //cout << "----" << endl;
    int sum = 0;
    while(i <= n \&\& j >= 1) {
         sum += g[i][j];
         // \texttt{cout} \, << \, \texttt{i} \, << \, \texttt{i} \, << \, \texttt{j} \, << \, \texttt{endl};
         ++i, --j;
    }
    return sum;
int calc1(int i, int j) {
    //cout << "----2-----" << endl;
    int sum = 0;
    while(i<=n && j<=n) {
         sum += g[i][j];
         // \texttt{cout} \, << \, \texttt{i} \, << \, \texttt{i} \, << \, \texttt{j} \, << \, \texttt{endl};
         ++i, ++j;
    }
    return sum;
}
int main() {
    ios::sync_with_stdio(0);
    cin.tie(0);
    cout.tie(0);
    cin >> n;
    int sum0 = 0;
    for (int i = 1; i <= n; ++i)
         for (int j = 1; j <= n; ++j)
             cin >> g[i][j], sum0 += g[i][j];
    int ans = 1e9;
    int sum1 = 0;
    for (int i = 1; i <= n; ++i) {
         int x = calco(1, i);
         sum1 += x;
         //cout << sum0-sum1 << ' ' << sum1-x << end1;
         ans = min(ans, abs(sum0-sum1*2+x));
    }
    for (int i = 2; i <= n; ++i) {
         int x = calco(i, n);
         sum1 += x;
         //cout << sum0-sum1 << ' ' << sum1-x << end1;
         ans = min(ans, abs(sum0-sum1*2+x));
    for (int i = 1; i <= n; ++i) reverse(g[i]+1, g[i]+n+1);
    sum1 = 0;
    for (int i = 1; i <= n; ++i) {
         int x = calco(1, i);
         sum1 += x;
         //cout << sum0-sum1 << ' ' << sum1-x << end1;
         ans = min(ans, abs(sum0-sum1*2+x));
    }
```

```
for (int i = 2; i<=n; ++i) {
    int x = calcO(i, n);
    sum1 += x;
    //cout << sum0-sum1 << ' ' << sum1-x << end1;
    ans = min(ans, abs(sum0-sum1*2+x));
}
cout << ans << end1;
return 0;
}</pre>
```

解法二: 优秀选手的做法, 把矩阵分别左旋45°和右旋45°, 然后用前缀和计算

```
#include <bits/stdc++.h>
using namespace std;
int n;
int a[1010][1010];
int 1[2050],r[2050];
int main(){
    scanf("%d",&n);
    for(int i=1;i<=n;i++){</pre>
        for(int j=1;j<=n;j++){
            scanf("%d",&a[i][j]);
            1[i+j]+=a[i][j];
            r[j-i+n+1]+=a[i][j];
        }
    }
    for(int i=2;i<=2*n;i++)][i]+=][i-1];
    for(int i=2;i<=2*n;i++)r[i]+=r[i-1];
    int ans=0x7fffffff;
    for(int i=2; i<=2*n; i++) ans=min(ans, abs(l[i-1]-(l[2*n]-l[i]));
    for(int i=2; i<=2*n; i++) ans=min(ans, abs(r[i-1]-(r[2*n]-r[i])));
    printf("%d",ans);
    return 0;
}
```

# K.黄金戟

语法题

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
typedef pair<int,int> P;
const int maxn=100000+10;

int main()
{
    //freopen("test0.in","r",stdin);
    //freopen("test0.out","w",stdout);
    int t; cin>>t;
    while(t--)
    {
        int a,b,c; cin>>a>>b>>c;
        if(a>=30&&b>=14&&c>=12) cout<<"Yes\n";
        else if(a>=20&&b>=14&&c>=12) cout<<"ullin";
        else cout<<"No\n";</pre>
```

```
}
return 0;
}
```