绪论

动态规划:记忆法

圣人不记事, 所以常记得; 今人忘事, 以其记事

有人建议不妨置备一本签名簿,供来访者留下自己的名字,就像怀特山那样;可是,天哪!我的记性非常好,用不着那个玩意儿。

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fib(): 递归

$$fib(n) = fib(n-1) + fib(n-2)$$
 Ø 1 1 2 3 5 8 13 21 34 55 89 ...

int fib(n) { return (2 > n) ? n : fib(n - 1) + fib(n - 2); } //为何这么慢?

令复杂度:
$$T(0) = T(1) = 1;$$
 $T(n) = T(n-1) + T(n-2) + 1, \forall n > 1$

- $\Leftrightarrow S(n) = [T(n) + 1]/2$
- \mathbb{N} S(0) = 1 = fib(1), S(1) = 1 = fib(2)
- 故 S(n) = S(n-1) + S(n-2) = fib(n+1)

$$T(n) = 2 \cdot S(n) - 1 = 2 \cdot fib(n+1) - 1 = \mathcal{O}(fib(n+1)) = \mathcal{O}(\phi^n)$$

- 其中 $\phi = (1+\sqrt{5})/2 \approx 1.618$

封底估算

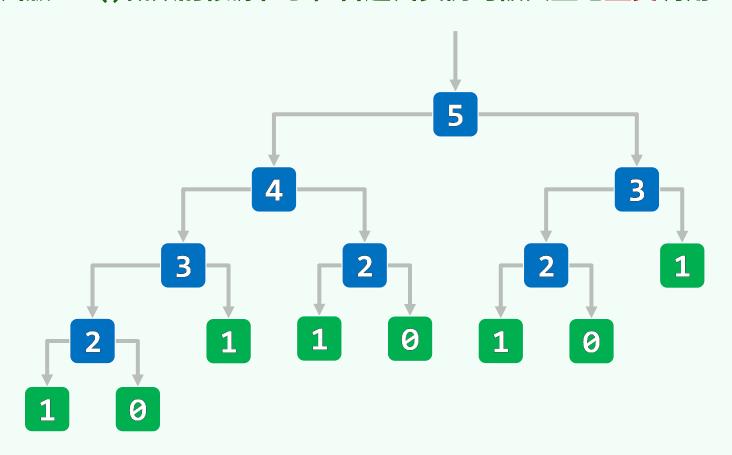
$$\phi^{36} \approx 2^{25} \quad \phi^{43} \approx 2^{30} \approx 10^9 \, flo = 1 \, sec$$

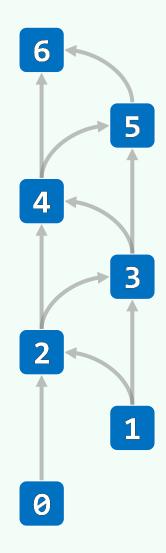
$$\phi^5 \approx 10 \quad \phi^{67} \approx 10^{14} flo = 10^5 sec \approx 1 day$$

$$\phi^{92} \approx 10^{19} flo = 10^{10} sec \approx 10^5 day \approx 3 century$$

递归

❖ 递归版fib()低效的根源在于,各递归实例均被大量地重复调用





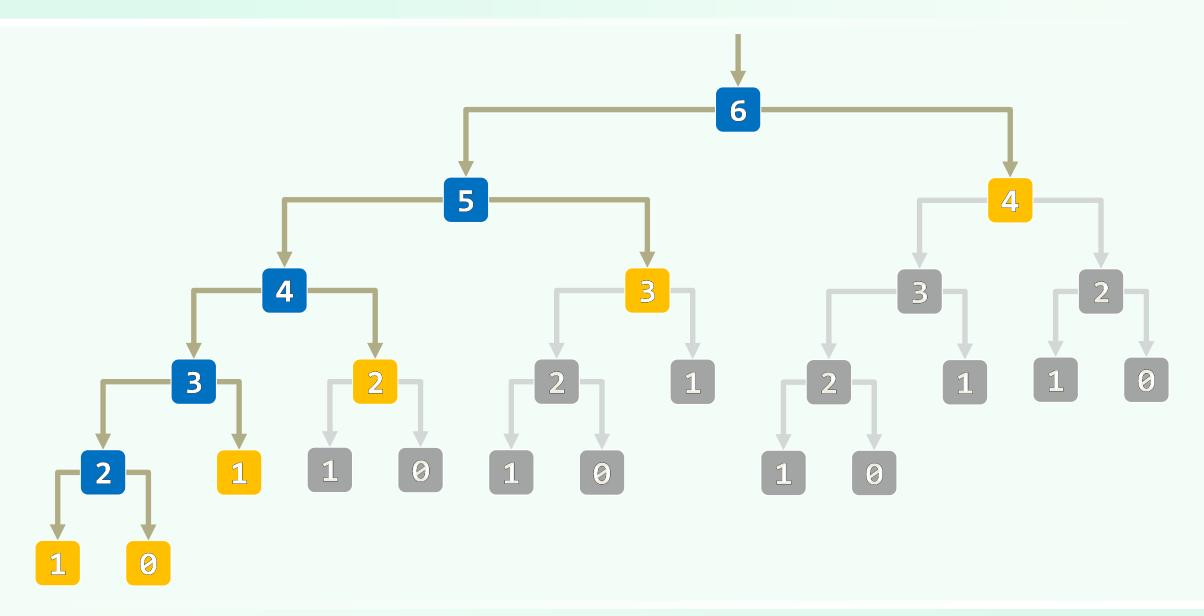
* 先后出现的递归实例,共计 $\mathcal{O}(\phi^n)$ 个;而去除重复之后,总共不过 $\mathcal{O}(n)$ 种

Memoization

```
def f(n)
   if ( n < 1 ) return trivial( n );</pre>
   return f(n-X) + f(n-Y)*f(n-Z);
```

```
T M[ N ]; #init. with UNDEFINED
def f(n)
   if ( n < 1 ) return trivial( n );</pre>
# recur only when necessary &
   always write down the result
   if ( M[n] == UNDEFINED )
      M[n] = f(n-X) + f(n-Y)*f(n-Z);
   return M[n];
```

Memoization: fib()



动态规划

❖ Dynamic programming, 颠倒计算方向: 由自顶而下递归, 改为自底而上迭代

```
while ( 0 < n-- ) {
  g = g + f;
   f = g - f;
 return g;
```

❖ T(n) = O(n), 而且仅需O(1)空间!

