07-XB

BST Application

Interval Tree

Your instinct, rather than precision stabbing, is more about just random bludgeoning.

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Stabbing Query

- Given a set of intervals in general position

  on the x-axis:

  $$S = \{s_i = [x_i, x_i'] \mid 1 \le i \le n\}$$

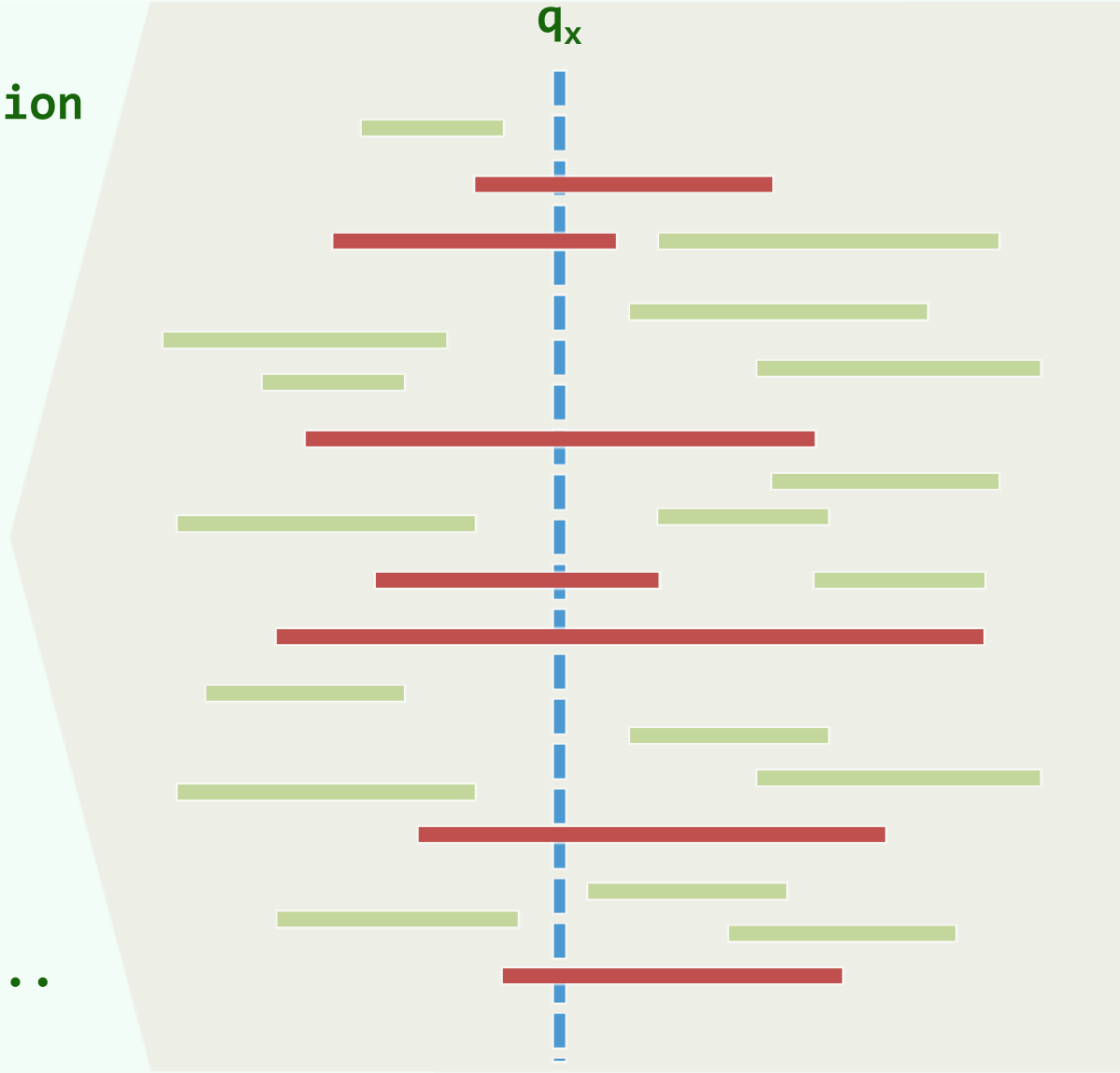
  and a query point q_x

- Find all intervals that contain q_x

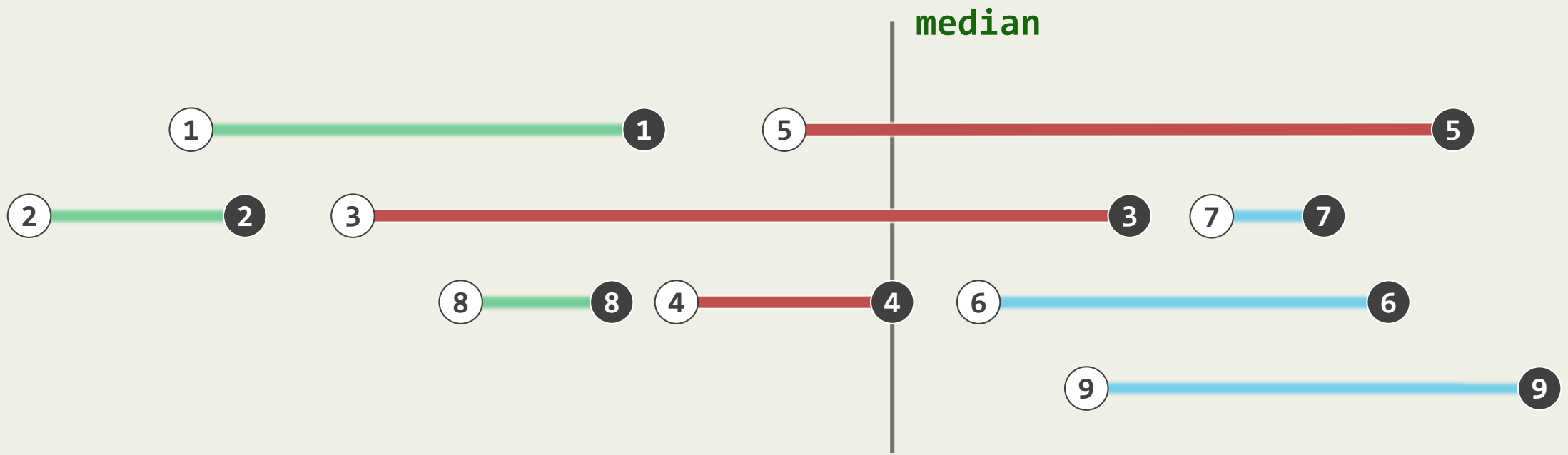
  $$\{s_i = [x_i, x_i'] \mid x_i \le q_x \le x_i'\}$$

- To solve this query,

  we will use the so-called interval tree ...



Median

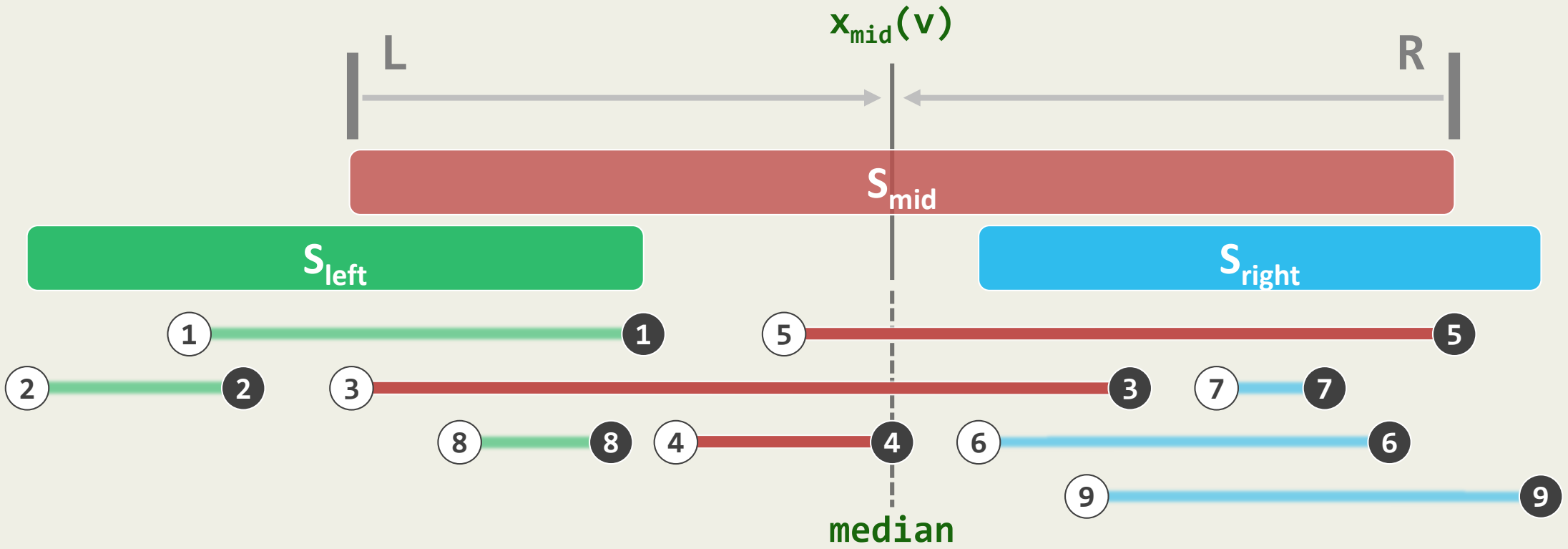


- Let $P = \partial S$ be the set of all endpoints

  (By general position assumption, $|P| = 2n$)

- Let $x_{mid} = median(P)$ be the median of P

Partitioning



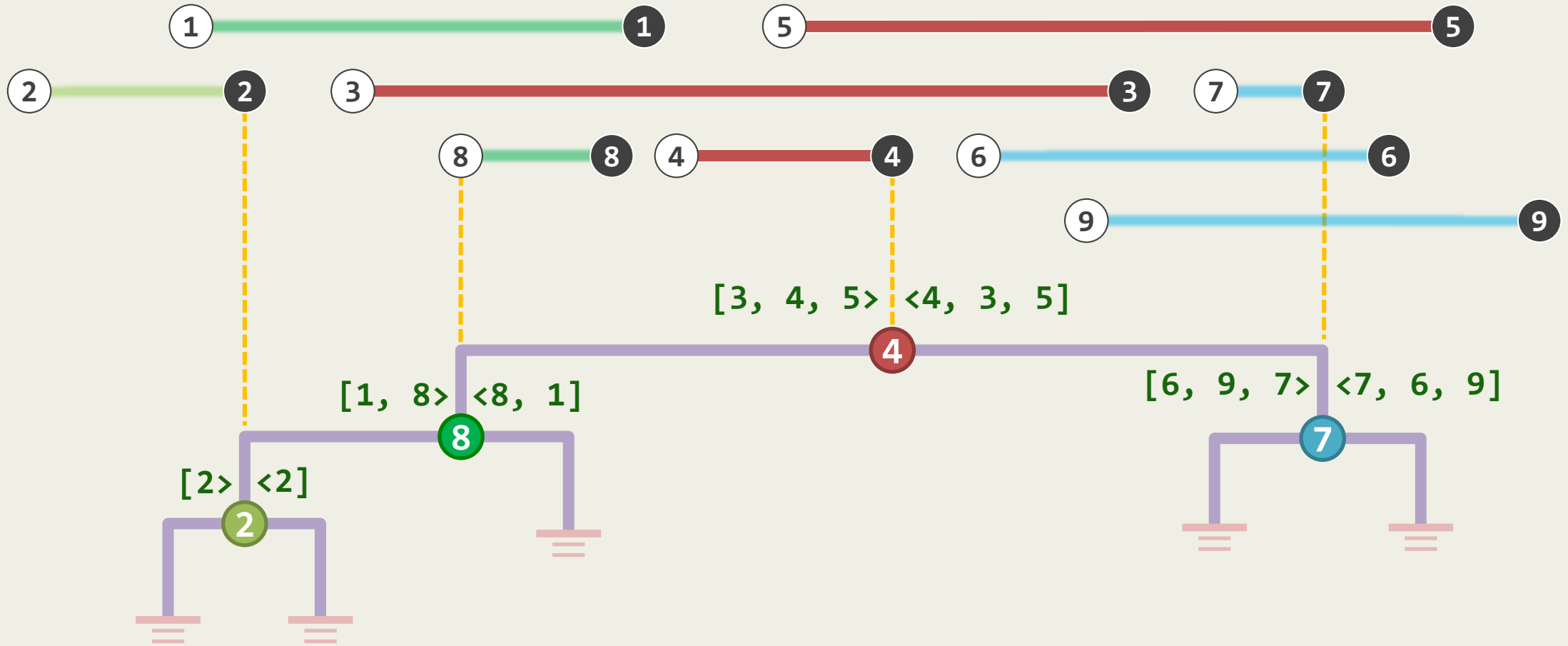
- All intervals can be then categorized into 3 subsets:

$$S_{left} = \{ S_i \mid x_i' < x_{mid} \} \quad S_{mid} = \{ S_i \mid x_i \le x_{mid} \le x_i' \} \quad S_{right} = \{ S_i \mid x_{mid} < x_i \}$$

- $S_{left/right}$ will be recursively partitioned until they are empty (leaves)

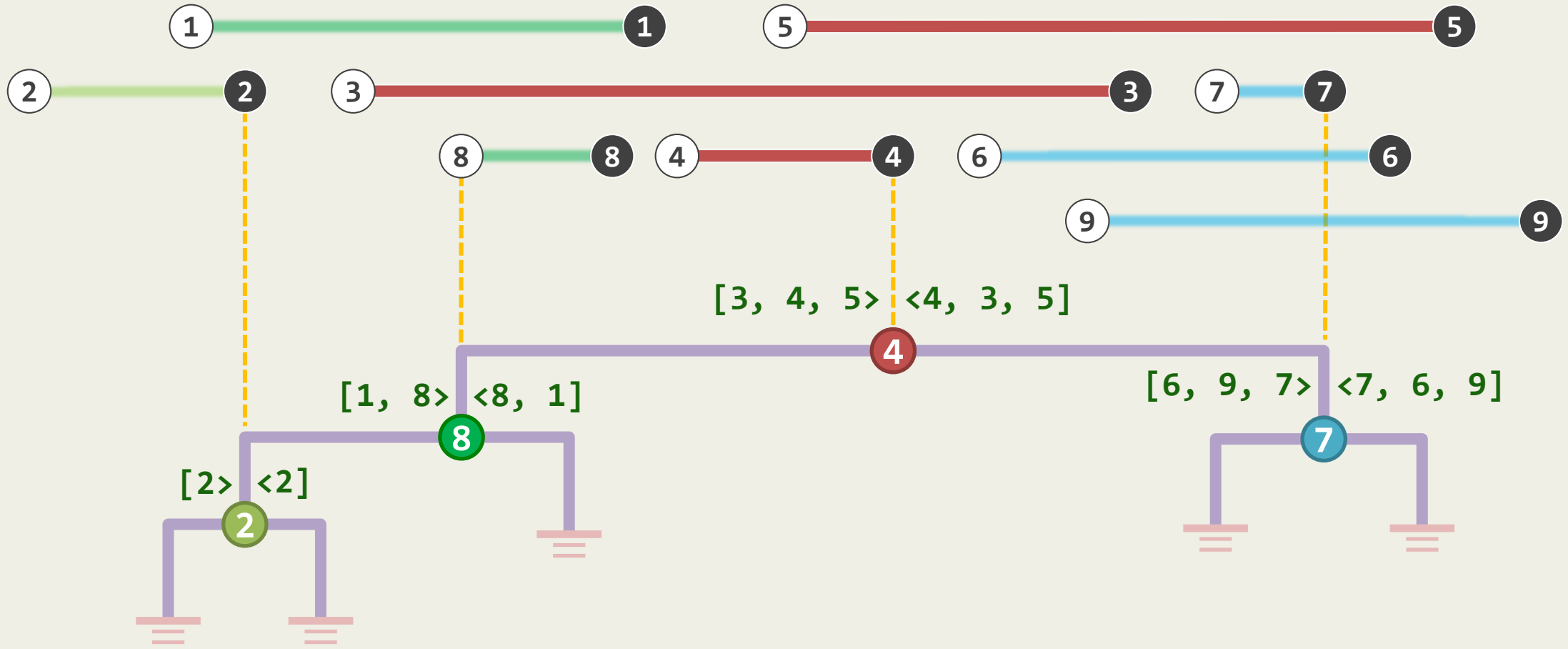
Balance & O(logn) Depth

$\max\{\ |S_{left}|, |S_{right}|\ \} \leq n/2$ **Best case:** $|S_{mid}| = n$ **Worst case:** $|S_{mid}| = 1$



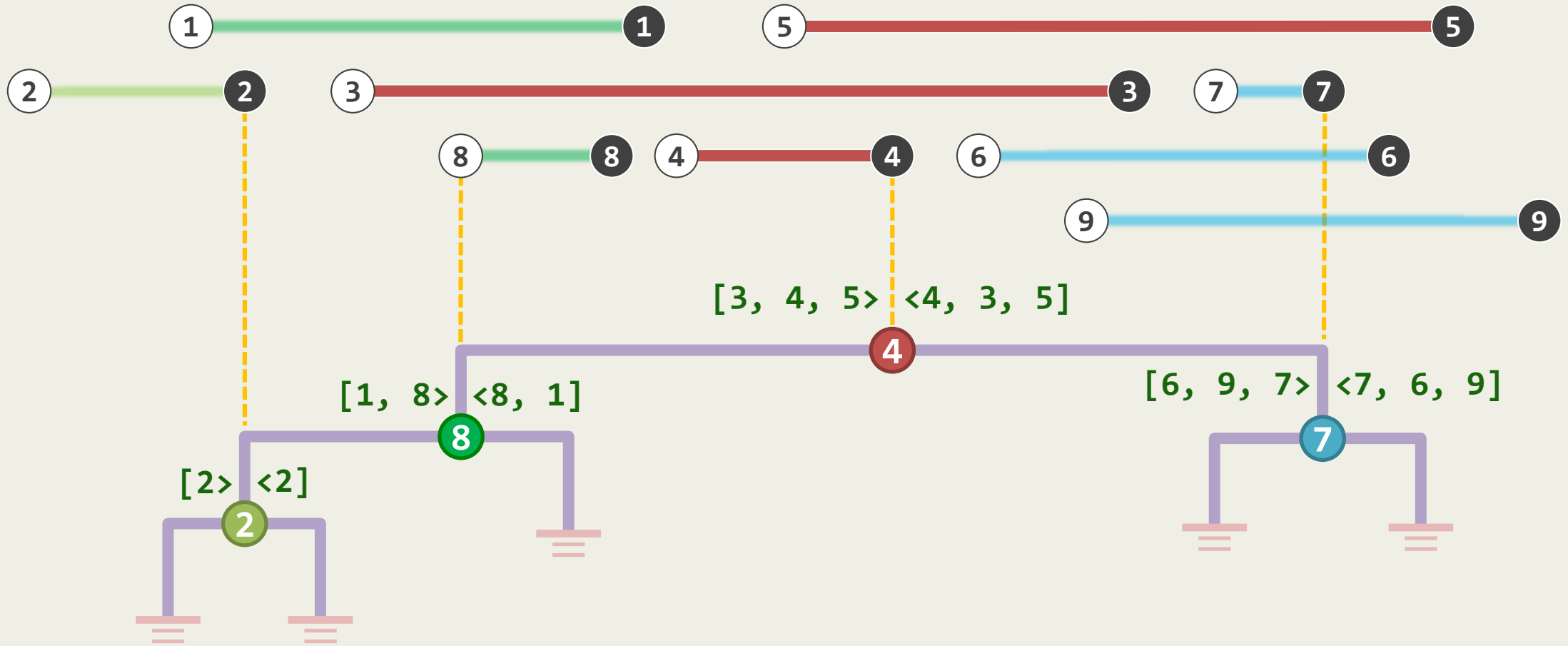
Associative Lists

❖ $L_{left/right}$ = all intervals of S_{mid} sorted by the left/right endpoints



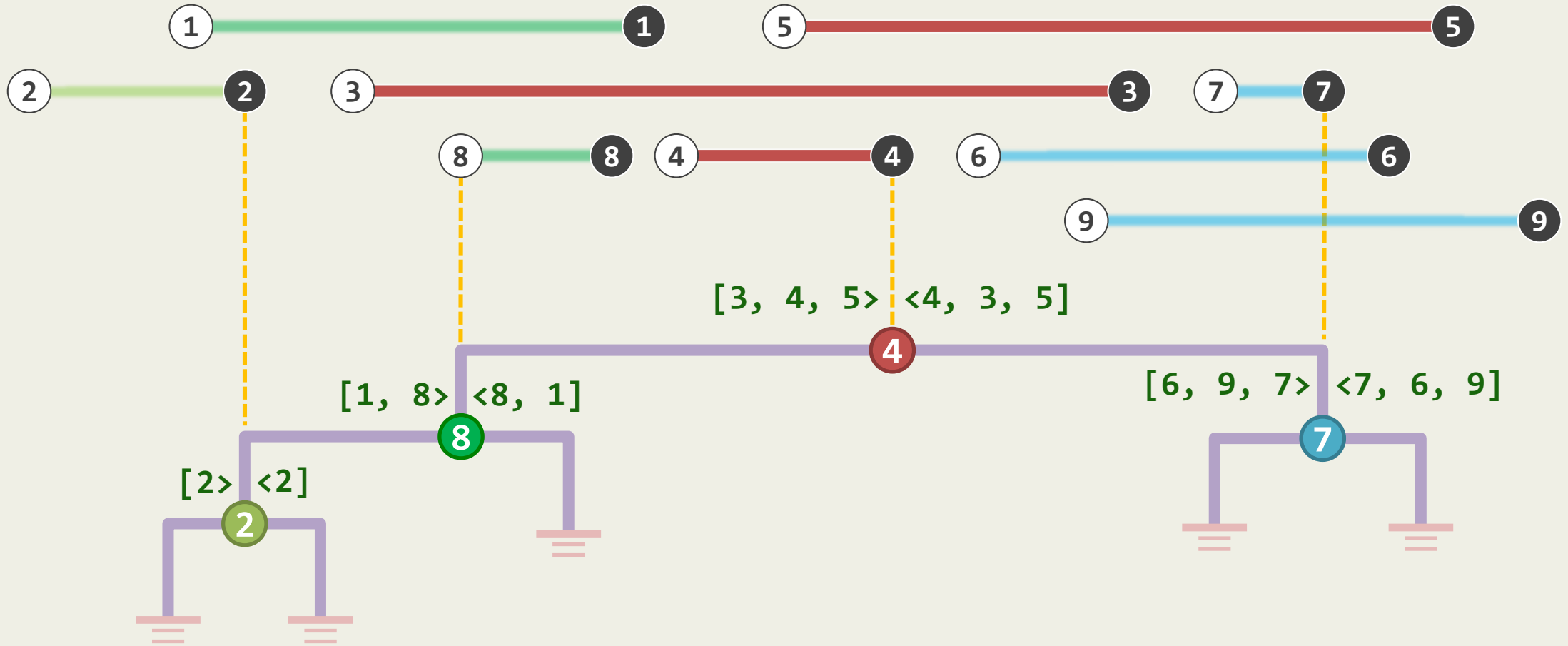
$O(n)$ Size

❖ Each segment appears twice (one in each list)



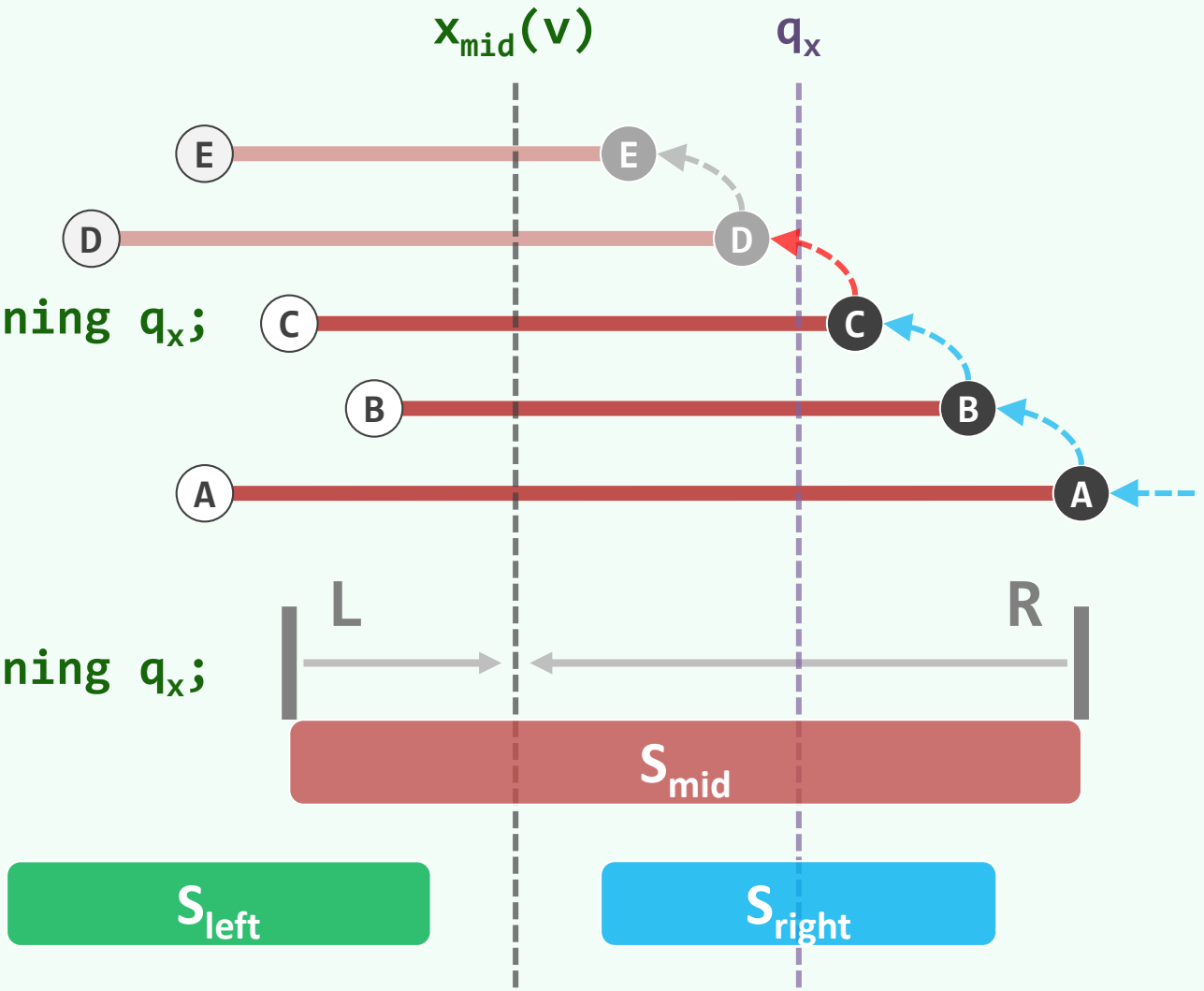
O(nlogn) Construction Time

❖ Hint: avoid repeatedly sorting



queryIntervalTree(v, q_x)

```
if ( ! v ) return;   //base
if ( q_x < x_mid(v) )
   report all segments of S_mid(v) containing q_x;
   queryIntervalTree( lc(v), q_x );
else if ( x_mid(v) < q_x )
   report all segments of S_mid(v) containing q_x;
   queryIntervalTree( rc(v), q_x );
else //with a probability ≈ 0
   report all segments of S_mid( v ); //both rc(v) & lc(v) can be ignored
```



O(r + logn) Query Time

❖ Each query visits O(logn) nodes //LINEAR recursion