

# Solution 8.4

## T1 字符串

签到题。反转串后按字符开头 `sort`，然后相同字符开头的按评分 `sort`。  
或者读入时按字符结尾归类，然后一类的按评分 `sort`。0(1) 查询输出。

写法一：

```
#include <bits/stdc++.h>
using namespace std;

const int A = 1e5 + 5;
int n, m;
struct node {
    int val, id;
    string x;
    inline void reverse() {
        int len = x.size() - 1;
        for (int i = 0; i <= len / 2; i++) swap(x[i], x[len - i]); return;
    }
} a[A];
int w[A];

inline bool cmp1(node u, node v) { return u.x < v.x; }

inline bool cmp2(node u, node v) {
    if (u.val != v.val) return u.val > v.val;
    return u.id < v.id;
}

signed main() {

    cin >> n >> m;
    for (int i = 1; i <= n; i++) {
        cin >> a[i].x >> a[i].val;
        a[i].id = i;
        a[i].reverse();
    }
    sort(a + 1, a + 1 + n, cmp1);
    for (int i = 1; i <= n; i++) {
```

```

        w[a[i].x[0] - 'a'] = i;
        if (a[i].x[0] != a[i - 1].x[0]) {
            int pos = i;
            while (a[pos + 1].x[0] == a[pos].x[0]) pos++;
            sort(a + i, a + pos + 1, cmp2);
            i = pos;
        }
    }
}

w[26] = n + 1;
for (int i = 26; ~i; i--) if (!w[i]) w[i] = w[i + 1];

for (int i = 1; i <= n; i++) a[i].reverse();
while (m--) {
    char x; cin >> x;
    int k; cin >> k;
    int num = x - 'a';
    if (w[num + 1] - w[num] < k)
        puts("Orz YYR tq!");
    else
        cout << a[w[num] + k - 1].x << '\n';
}
return 0;
}

```

写法二:

```

#include <bits/stdc++.h>
#define ll long long
using namespace std;
int n, m;
struct pr {
    int id, sc;
    string c;
};
vector<pr> a[28];

char s[55];
bool cmp(pr x, pr y) { return x.sc != y.sc ? x.sc > y.sc : x.id < y.id; }
int main() {
    scanf("%d%d", &n, &m);
    for (int x, i = 1; i <= n; ++i) {
        scanf("%s%d", s, &x);
        a[s[strlen(s) - 1] - 'a'].push_back(pr{i, x, s});
    }
    for (int i = 0; i < 26; ++i) sort(a[i].begin(), a[i].end(), cmp);
    int p;

```

```

while (m--) {
    scanf("%s%d", s, &p);
    s[0] -= 'a';
    if (p > a[s[0]].size())
        printf("Orz YZR tq!\n");
    else
        cout << (a[s[0]][p - 1].c) << endl;
}
return 0;
}

```

## T2 公约数

首先枚举  $a, b$ ，预处理出每个  $\text{gcd}$  的出现次数，然后枚举每个  $\text{gcd}$  与  $[1, n]$  中的数，将贡献相加即可。

【核心代码】

```

for(int i = 1; i <= n; i++)
    for(int j = 1; j <= n; j++)
        ++tong[gcd(i, j)];
for(int i = 1; i <= n; i++)
    for(int j = 1; j <= n; j++)
        ans += tong[i] * gcd(i, j);

```

## T3 数独游戏

爆搜 + 合理的剪枝，类似 NOIP 靶形数独

## T4 入侵攻击

并查集，我们可以对  $d$  数组排序，离线处理这个问题。

$n$  最大 1000，因此我们可以预处理  $n^2$  条边然后排序，把边从小到大依次加入并查集中。

对于当前  $d$  询问和  $(0, 0, 0)$  点相连的连通块大小

【核心代码】

```

cin >> n >> m;
for (int i = 1; i <= n; i++) {
    cin >> p[i].x >> p[i].y >> p[i].z;
}
p[++n] = (point){0, 0, 0};

for (int i = 1; i <= n; i++) {
    for (int j = 1; j < i; j++) {

```

```

        insert(i, j, dist(p[i], p[j]));
    }
}

sort(e + 1, e + cnt + 1, cmp);
for (int i = 1; i <= n; i++) {
    father[i] = i;
    sz[i] = 1;
}

for (int i = 1; i <= m; i++) {
    cin >> q[i].d;
    q[i].id = i;
}

sort(q + 1, q + m + 1, cmp);

int now = 1;
for (int i = 1; i <= m; i++) {
    while (now <= cnt && e[now].w <= q[i].d) {
        int x = find(e[now].u);
        int y = find(e[now].v);
        if (x != y) {
            father[x] = y;
            sz[y] += sz[x];
        }
        now++;
    }

    ans[q[i].id] = sz[find(n)] - 1;
}

for (int i = 1; i <= m; i++) {
    printf("%d\n", ans[i]);
}

```