

Problems:

1. path coverage guarantee
2. sensitive analysis guarantee

Constraints:

1. language extensibility
2. performance (minimize the instrumentation)

P1

1. get the block-level CFG of a function, encode each CFG node as a block ID
2. the CFG is language-independent, satisfying the constraint 1.
3. apply the instrumentation algorithm on the CFG to compute PCG guidance, satisfying the constraint 2.

P2

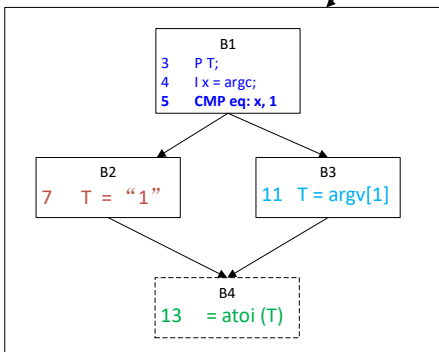
1. define a language-independent IR to retain the necessary information for SAG, satisfying the constraint 1.
2. extend the block-level CFG to instruction (or statement)-level CFG: each node represents a statement in IR
3. Def-Use computation on the use of each CMP and SWITCH node of the CFG: get the definition statement for each branch variables as SAG
4. merge SAG and PCGG, satisfying the constraint 2.

```
1 int main (int argc, char* argv[])
2 {
3     char *T;
4     int x = argc;
5     if (x == 1)
6     {
7         T = "1" ;
8     }
9     else
10    {
11        T = argv[1];
12    }
13    return atoi (T);
14}
```

IR

```
1 | main (I argc, P argv[])
2 {
3     P T;
4     I x = argc;
5     CMP eq: x, 1
6     {
7         T = "1"
8     }
9     else
10    {
11        T = argv[1]
12    }
13    = atoi (T)
14}
```

CFG



Source: 5 CMP eq: x, 1

Definition: 4 I x = argc;

PCGG: B1, B2, B3
SAG: B1

=> PCGG & SAG:
B1, B2, B3

