Microcomputers I – CE 320

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Announcement

- You are going to have quiz no.4 next Friday, Feb 12.
 - Covers homework exercise 5
 - Stack and subroutines

Lecture 16: Parameter Passing

Today's Topics

Use different methods to pass parameters to subroutines

Use stack frames to track data in use by a subroutine

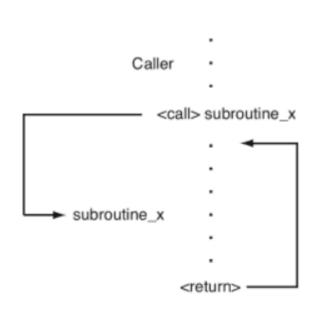
Subroutine

 A subroutine is a sequence of instructions that can be called from many different places in a program.

• The program unit that makes the subroutine call is referred to as a **caller**.

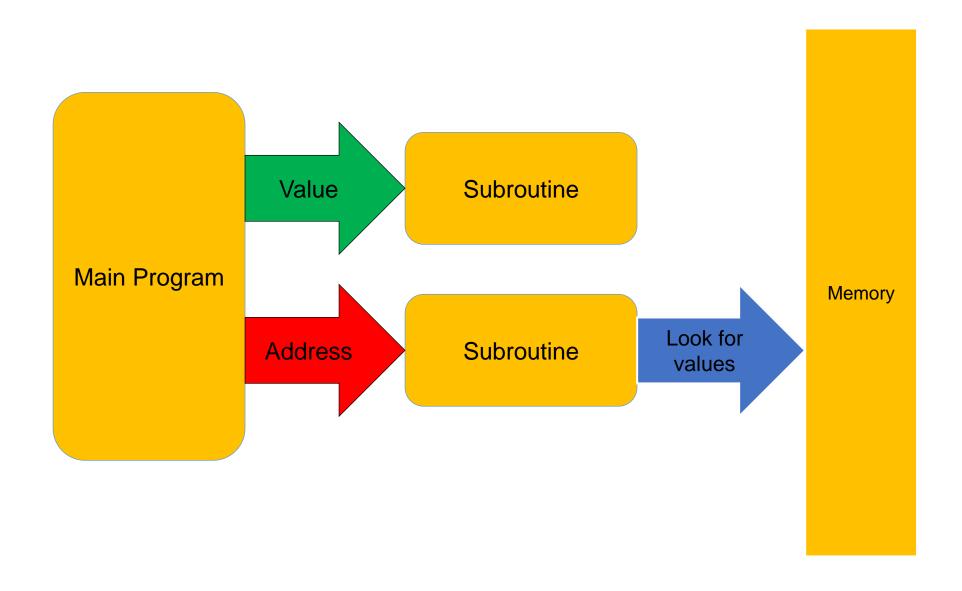
• The subroutine called by other program units is referred to as a **callee**.

 The caller usually wants the subroutine to perform a computation using the <u>parameters</u> <u>passed to it</u>.



Program flow during a subroutine call

Call by Value vs. Call by Reference



 The caller may use the following methods to pass parameters to the subroutine:

Use Registers

- In this method, parameters are placed in CPU registers before the subroutine is called.
- This method is very convenient when there are only a few parameters to be passed.

Use the stack

- In this method, parameters are pushed into the stack before the subroutine is called.
- The <u>stack must be cleaned up after the computation is completed</u>. This can be done by either the caller or the callee.

Use the global memory

- Global memory is accessible to both the caller and the callee.
- As long as the caller places parameters in global memory before it calls the subroutine, the callee will be able to access them.

- Passing parameters using memory:
 - The caller sets memory locations and the callee reads them.

Example:

```
movb #2,$1000
movb #4,$1001
bsr addition
```

The caller stores the data in memory locations, e.g., \$1000 and \$1001

```
addition:

ldaa $1000

adaa $1001
```

The subroutine reads the data from memory locations, e.g., \$1000 and \$1001

- Passing parameters using registers:
 - The caller sets registers and the callee reads them.

Example:

```
ldaa #2
ldab #4
bsr
addition
```

The caller stores the data in registers, e.g., a and b

```
addition:
aba
```

The subroutine reads the data from the registers

- Passing parameters using stack:
 - The caller pushes the parameters in the stack and the callee reads them.

Example:

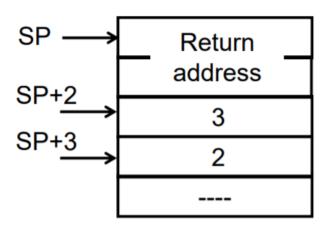
```
ldaa #2
psha
ldaa #3
psha
bsr addition
leas 2,+sp
```

The caller pushes the data in stack before calling the subroutine

The stack is cleaned after the subroutine call

addition: tfr sp,X ldaa 2,+x ldaa 1,+x

The subroutine reads the locations sp+2 and sp+3



Passing Parameters in Registers

 This is often the preferred method of passing data to and from the subroutine.

- This method has the following advantages and disadvantages:
 - Data is immediately available to the subroutine.
 - Registers must usually be used to move data using the other methods, and this saves the extra preparation.
 - Often fastest execution, smallest code size
 - There are **only a limited number of registers available**, however, pass-by-reference may be used to point to a list of input/output values.

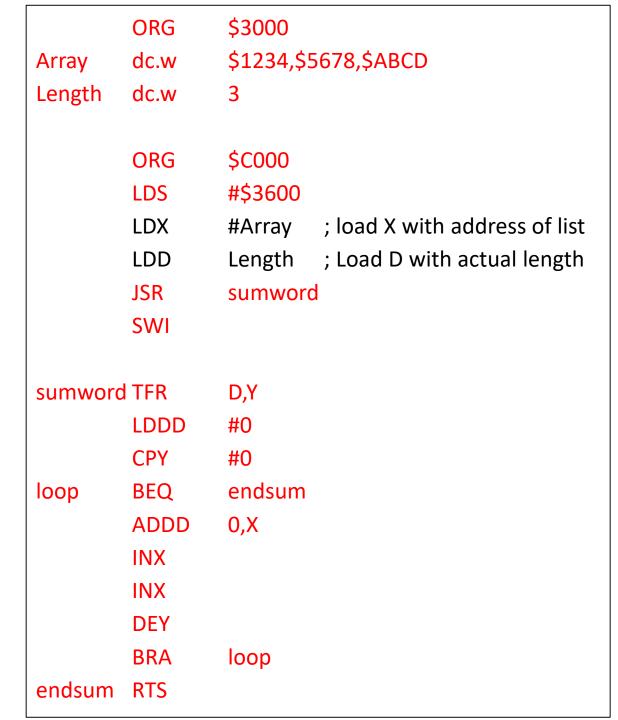
Example

- Write a subroutine that adds an array of two-byte numbers and a sample main program that calls it.
- The array is passed by reference in X, the length is passed by value in D.
- The sum should be returned by value in D.
- Do not worry about indicating signed or unsigned overflow.

```
ORG
                  $3000
                  $1234,$5678,$ABCD
         dc.w
Array
Length
                  3
         dc.w
                  $C000
         ORG
         LDS
                  #$3600
         LDX
                  #Array
                           ; load X with address of list
         LDD
                  Length
                           ; Load D with actual length
         ISR
                  sumword
         SWI
sumword TFR
                  D,Y
         LDDD
                  #0
         CPY
                  #0
                  endsum
qool
         BEQ
         ADDD
                  0.X
         INX
         INX
         DFY
         BRA
                  loop
endsum
```

Preserving Registers

- The previous example used Y within the subroutine, which destroyed the original value in that register, and it modified X.
- If Y was in use by the main program, the main program would continue after the subroutine with an incorrect value in Y.
- To avoid this, registers used by a subroutine may be saved to the stack before a subroutine uses them, and they are restored before the main program resumes using them.
- For this, there are two options:
 - The caller (main program) does this.
 - The callee (subroutine) does this.

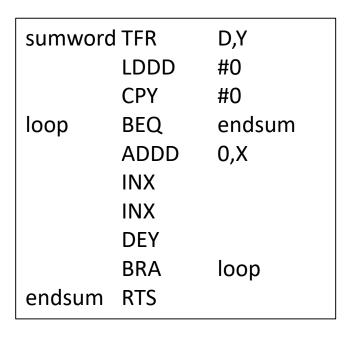


Preserving Registers

- Responsibility of the Caller:
 - The calling main program assumes all registers are destroyed by the subroutine.
 - Calling program saves only those registers that are in use.
 - If the registers used by the subroutine are unknown (i.e. using a sub. provided by someone else), may save registers that the subroutine wouldn't affect.
 - Code to save/restore registers duplicated with every subroutine call.
- Responsibility of the Callee:
 - Saves only the registers that will be used by the subroutine.
 - Save/restore code only occurs once.
 - May waste time saving registers not in use by the main program.
- Notes:
 - Caller responsible is the safest bet if we do not know how the subroutine is implemented.

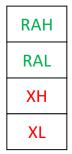
Example – Caller Responsible

	SWI		
	PULX		
	JSR	sumword	
	PSHX		
	LDD	Length ; Load actual length	
	LDX	#Array ; load list address	
	LDS	#\$3600	
	ORG	\$C000	
Length	dc.w	3	
Array	dc.w	\$1234,\$5678,\$ABCD	
	ORG	\$3000	



We push index register X onto stack before we call a subroutine.

Stack Frame



Example – Callee Responsible



Array	ORG dc.w	\$3000 \$1234,\$5678,\$ABCD
Length	dc.w	3
	ORG LDS LDX LDD JSR SWI	\$C000 #\$3600 #Array ; load list address Length ; Load actual length sumword

sumword PSHX				
	PSHY			
	TFR	D,Y		
	LDDD	#0		
	CPY	#0		
loop	BEQ	endsum		
	ADDD	0,X		
	INX			
	INX			
	DEY			
	BRA	loop		
endsum	PULY			
	PULX			
	RTS			

We push index registers X and Y onto stack at the beginning of a subroutine.

Q: Why isn't register D saved and restored?

A: It is used to return the answer!

Stack Frame

YH
YL
XH
XL
RAH

Passing Parameters in the Stack

- The stack is the next preferred location to pass data. This method has the following advantages and disadvantages.
 - The stack is already in use for <u>saving the return address</u>
 - <u>Indexed addressing</u> can easily access data stored on the stack
 - The amount of data passed is not limited by the register set
 - The data passed on the stack must be removed, and this is usually the responsibility of the caller (although this can be done by the callee)
 - However, it takes up more CPU clocks to pass parameters on stack and for the subroutine to access them compared to data passed in registers
- Note:
 - If the caller will save registers on the stack, it should be done before passing parameters.

Example – Parameter Passing in the Stack

	ORG	\$3000	
Array	dc.w	\$1234,\$5678,\$ABCD	
Length	dc.w	3	
	ORG	\$2000	
	LDS	#\$3600	
	LDX	#\$1234 ; something to save	
	PSHX		
	LDD	Length	
	PSHD		
	LDD	#Array	
	PSHD		
	JSR	sumword	
	LEAS	4,SP	
	PULX		
	SWI		

sumword	LDDD	#0
	LDX	2,SP
	LDY	4,SP
Loop	BEQ	endsum
	ADDD	0,X
	INX	
	INX	
	DEY	
	BRA	loop
endsum	RTS	

RetH
RetL
ArrH
ArrL
LenH
LenL
XH
XL

The **address of the array** and the **two-byte length** are passed on the stack. The sum should be returned by value in D.

Questions?

Wrap-up What we've learned

Parameter Passing

Pros and Cons of each method

What to Come

More about subroutines