

Microcomputers I – CE 320

Mohammad Ghamari, Ph.D.
Electrical and Computer Engineering
Kettering University

Announcement

- You are going to have quiz no.4 next Friday, Feb 12.
 - Covers homework exercise 5
 - Stack and subroutines

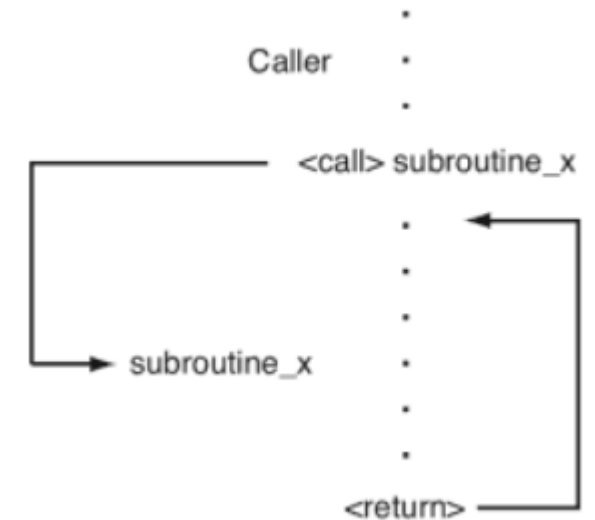
Lecture 16: Parameter Passing

Today's Topics

- Use different methods to pass parameters to subroutines
- Use stack frames to track data in use by a subroutine

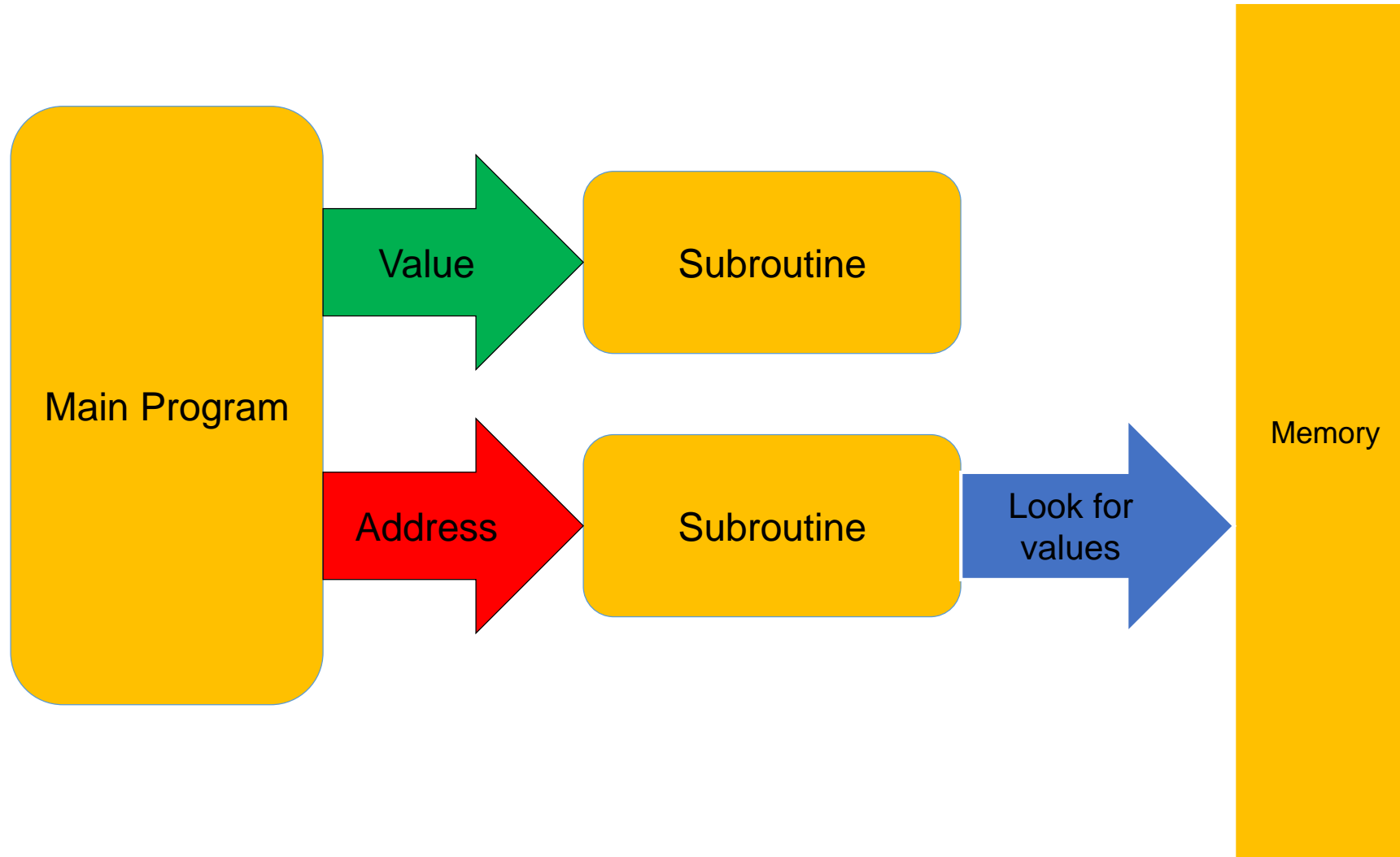
Subroutine

- A subroutine is a sequence of instructions that can be called from many different places in a program.
- The program unit that makes the subroutine call is referred to as a **caller**.
- The subroutine called by other program units is referred to as a **callee**.
- The caller usually wants the subroutine to perform a computation using the parameters passed to it.



Program flow during a subroutine call

Call by Value vs. Call by Reference



Parameter Passing Methods

- The caller may use the following methods to pass parameters to the subroutine:
 - **Use Registers**
 - In this method, parameters are **placed in CPU registers** before the subroutine is called.
 - This method is very convenient when there are only a few parameters to be passed.
 - **Use the stack**
 - In this method, parameters are **pushed into the stack** before the subroutine is called.
 - The stack must be cleaned up after the computation is completed. This can be done by either the caller or the callee.
 - **Use the global memory**
 - **Global memory is accessible to both the caller and the callee.**
 - As long as the caller places parameters in global memory before it calls the subroutine, the callee will be able to access them.

Parameter Passing Methods

- Passing parameters using memory:
 - The caller sets memory locations and the callee reads them.

Example:

```
movb #2,$1000  
movb #4,$1001  
bsr addition
```

The caller stores the data in memory locations, e.g., \$1000 and \$1001

```
addition:  
    ldaa $1000  
    adaa $1001
```

The subroutine reads the data from memory locations, e.g., \$1000 and \$1001

Parameter Passing Methods

- Passing parameters using registers:
 - The caller sets registers and the callee reads them.

Example:

```
ldaa #2  
ldab #4  
bsr  
addition
```

```
addition:  
    aba
```

The caller stores the data in registers, e.g., a and b

The subroutine reads the data from the registers

Parameter Passing Methods

- Passing parameters using stack:
 - The caller pushes the parameters in the stack and the callee reads them.

Example:

```
ldaa #2  
psha  
ldaa #3  
psha  
bsr addition  
leas 2,+sp
```

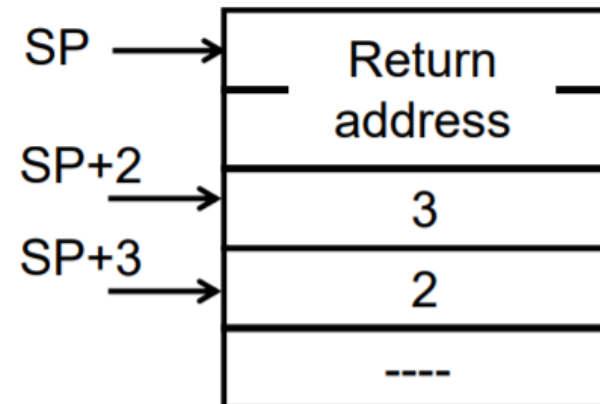
The caller pushes the data in stack before calling the subroutine

The stack is cleaned after the subroutine call

addition:

```
tfr sp,X  
ldaa 2,+x  
ldaa 1,+x
```

The subroutine reads the locations sp+2 and sp+3



Passing Parameters in Registers

- This is often the **preferred method** of passing data to and from the subroutine.
- This method has the following advantages and disadvantages:
 - Data is **immediately available** to the subroutine.
 - Registers must usually be used to move data using the other methods, and this saves the extra preparation.
 - **Often fastest execution**, smallest code size
 - There are **only a limited number of registers available**, however, pass-by-reference may be used to point to a list of input/output values.

Example

- Write a subroutine that adds an array of two-byte numbers and a sample main program that calls it.
- The array is ***passed by reference*** in X, the length is ***passed by value*** in D.
- The sum should be ***returned by value*** in D.
- Do not worry about indicating signed or unsigned overflow.

	ORG	\$3000
Array	dc.w	\$1234,\$5678,\$ABCD
Length	dc.w	3
	ORG	\$C000
	LDS	#\$3600
	LDX	#Array ; load X with address of list
	LDD	Length ; Load D with actual length
	JSR	sumword
	SWI	
sumword	TFR	D,Y
	LDDD	#0
	CPY	#0
loop	BEQ	endsum
	ADDD	0,X
	INX	
	INX	
	DEY	
	BRA	loop
endsum	RTS	

Preserving Registers

- The previous example used Y within the subroutine, which destroyed the original value in that register.
 - Additionally, X is used as an index to access array elements.
- If Y was in use by the main program, the main program would continue after the subroutine with an incorrect value in Y.
- To avoid this, registers used by a subroutine may be **saved to the stack** before a subroutine uses them, and they are restored before the main program resumes using them.
- For this, there are two options:
 - The caller (main program) does this.
 - The callee (subroutine) does this.

	ORG	\$3000	
Array	dc.w	\$1234,\$5678,\$ABCD	
Length	dc.w	3	
	ORG	\$C000	
	LDS	#\$3600	
	LDX	#Array ; load X with address of list	
	LDD	Length ; Load D with actual length	
	JSR	sumword	
	SWI		
sumword	TFR	D,Y	Transfers the contents of D to Y
	LDDD	#0	
	CPY	#0	
loop	BEQ	endsum	
	ADDD	0,X	
	INX		
	INX		X is used as an index to access array elements
	DEY		Decrements the value in Y register
	BRA	loop	
endsum	RTS		

Preserving Registers

- Responsibility of the Caller:
 - The calling main program assumes all registers are destroyed by the subroutine.
 - Calling program saves only those registers that are in use.
 - If the registers used by the subroutine are unknown (i.e. using a sub. provided by someone else), may save registers that the subroutine wouldn't affect.
 - Code to save/restore registers duplicated with every subroutine call.
- Responsibility of the Callee:
 - Saves only the registers that will be used by the subroutine.
 - Save/restore code only occurs once.
 - May waste time saving registers not in use by the main program.
- Notes:
 - Caller responsible is the **safest bet** if we do not know how the subroutine is implemented.

Example – Caller Responsible

```

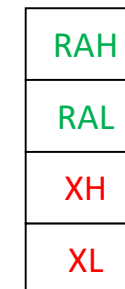
                ORG    $3000
Array          dc.w    $1234,$5678,$ABCD
Length         dc.w    3

                ORG    $C000
                LDS     #$3600
                LDX     #Array    ; load list address
                LDD     Length    ; Load actual length
                PSHX
                JSR      sumword
                PULX
                SWI
```

```

sumword TFR      D,Y
                LDDD     #0
                CPY      #0
loop        BEQ      endsum
                ADDD     0,X
                INX
                INX
                DEY
                BRA      loop
endsum      RTS
```

Stack Frame



We push index register X onto stack before we call a subroutine.

Example – Callee Responsible

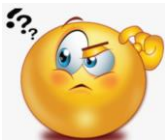
```

Array  ORG    $3000
       dc.w   $1234,$5678,$ABCD
Length dc.w   3

       ORG    $C000
       LDS    # $3600
       LDX    #Array    ; load list address
       LDD    Length    ; Load actual length
       JSR    sumword
       SWI
```

```

sumword PSHX
       PSHY
       TFR    D,Y
       LDDD   #0
       CPY    #0
loop    BEQ    endsum
       ADDD   0,X
       INX
       INX
       DEY
       BRA    loop
endsum  PULY
       PULX
       RTS
```



Q: Why isn't register D saved and restored?

A: It is used to return the answer!

Stack Frame

YH
YL
XH
XL
RAH
RAL

We push index registers X and Y onto stack at the beginning of a subroutine.

Passing Parameters in the Stack

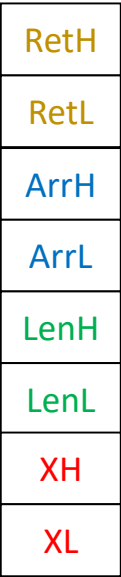
- The stack is the **next preferred** location to pass data. This method has the following advantages and disadvantages.
 - The stack is already in use for saving the return address
 - Indexed addressing can easily access data stored on the stack
 - The amount of data passed is **not limited by the register set**
 - The data passed on the stack must be removed, and this is usually the responsibility of the caller (although this can be done by the callee)
 - However, it takes up more CPU clocks to pass parameters on stack and for the subroutine to access them compared to data passed in registers
- Note:
 - If the caller will save registers on the stack, it should be done before passing parameters.

Example – Parameter Passing in the Stack

	ORG	\$3000
Array	dc.w	\$1234,\$5678,\$ABCD
Length	dc.w	3
	ORG	\$2000
	LDS	#\$3600
	LDX	#\$1234 ; something to save
	PSHX	
	LDD	Length
	PSHD	
	LDD	#Array
	PSHD	
	JSR	sumword
	LEAS	4,SP
	PULX	
	SWI	

sumword	LDDD	#0
	LDX	2,SP
	LDY	4,SP
Loop	BEQ	endsum
	ADDD	0,X
	INX	
	INX	
	DEY	
	BRA	loop
endsum	RTS	

Stack Frame



The **address of the array** and the **two-byte length** are passed on the stack. The sum should be returned by value in D.

Questions?

Wrap-up

What we've learned

- Parameter Passing
- Pros and Cons of each method

What to Come

- More about subroutines