## CS156 – Introduction to Artificial Intelligence Final Exam Review

By: Zayd Hammoudeh

## **Introduction to Agents**

An agent perceives its environment through sensors and acts upon the environment through actuators.

Turing Test – A test where a human poses a series of questions to the computer and after seeing the responses cannot distinguish the responses from those of a human.

## **Components Needs to Pass the Turing Test:**

- 1. Natural Language Processing
- 2. Knowledge Representation (i.e. storage paradigm)
- 3. Automated Reasoning
- 4. Machine Learning

**Total Turing Test** – A variant of the Turing Test where the robot passes entirely as a human.

## Additional Requirements Over Standard Turing Test:

- 1. Computer Vision
- 2. Robotics

Rational Agent – For every possible percept sequence, the rational agent selects the action it expects to maximize its performance measure given the information in the percept sequence and whatever built-in knowledge it has.

## The maximizing action depends on:

- 1. Performance Measure
- 2. Any prior/built-in knowledge of the agent
- 3. Percept sequence to date.
- 4. Set of possible actions.

Percept – An agent's perceptual inputs through sensors at any given instant.

**Percept Sequence** – Set of all percepts to date.

Agent Function: Map from percept sequences to an agent action. Example: An agent action table.

Agents run an agent program. The agent program runs on the agent architecture. The combination of the agent program and agent architecture is called a complete agent.

Cognitive Science: Brings together computer models from Al and experimental techniques from psychology to construct precise and testable theories of the human mind.

## Task Environment (PEAS)

Performance Measure (P) – Targets/goals	Environment (E) – Objects that interact	Actuators (A) – Tool(s) used by the	Sensors (S) – Tool(s) used by the agent
the agent will try to achieve.	with the agent or the agent interacts with	agent to interact with the environment.	to perceive the environment.

#### **Properties of a Task Environment**

Fully Observable vs. Partially Observable	Deterministic vs. Stochastic	Single-Agent vs. Multi-agent	Episodic vs. Sequential
Can the agent see the entire	Is the next state completely determined	Do objects in the environment need to be	In an episodic environment, the agent's
environment at once (e.g. chess)? If not,	by the current state and the action	treated as other agents? Multi-agent	experience is divided into episodes. In an
it may keep a history of what it has	(chess)? Otherwise it is stochastic (taxi-	environments can be competitive (chess)	episode the agent receives one percept
observed (taxi-driver).	driver).	or cooperative (taxi-driving).	and performs one action (e.g. quality
		Communication between agents is	control robot). In sequential
		possible as is randomized behavior to	environments, current actions affect
		avoid predictability.	future actions.
Static vs. Dynamic	Discrete vs. Continuous	Known versus Unknown	
Does the environment change while the	Time, percepts, and actions divided into	In a known environment, all outcomes of	
agent is making a decision? Chess is	a fixed, finite set (e.g. chess)? A	actions are known. In an unknown	
static while taxi driving is dynamic.	continuous environment is taxi-driving.	environment, the agent needs to figure	
		out how it works to make good decisions.	

## **Example Episodic Agent**

Quality Assurance robot.

- Performance Measure: Fixed minimum and maximum tolerances for a widget. (Example ball board min/max weight, diameter, roundness)
- Environment: Widget (example ball bearing) received for inspection on an input system. Good bin and discard bins.
- Actuator: Arm to place widget in either discard bin or good bin.
- Sensor: Check ball bearing weight, diameter, roundness etc.

## **Types of Agent Programs**

Simple Reflex Agent – Select actions based off the current percept only. Often defined by condition-action rules (i.e. productions)	Model-Based Reflex Agent – Similar to a Finite State Automata. Uses internal states to keep track of the environment. Updates the internal state based off how the environment evolves independently and how the agent's action affect the environment. This is called the agent model.	
Goal Based Agents – A goal is a binary condition (i.e. either met or not met). A goal based	Utility Based Agent – Agent applies a utility function to its performance. Agent	
agent tries to reach a target goal. Search and planning agents may be goal based agents.	tries to maximize its overall utility function.	

## **Additional Definitions**

Problem solving agents deal with atomic	Planning agents deal with factored or	Search – Process of looking over a	Solution – A sequence of actions that
environments (i.e. the environment is	structured environments (i.e. the	sequence of actions.	takes the agent from the initial state to
treated as a single whole and is	environment has attributes/variables		the goal state.
indivisible).	each of which has a value).		

## Search Problems

#### Classical search problems are deterministic, fully-observable, known, and the solution is a sequence of actions.

<b>Solution:</b> A sequence of actions that takes the agent from the initial state to the goal state.	Root: Initial State Edge/Branches: Actions Node/Vertices: States in the state space Leaf: A node with no children	Node Expansion – Applying all legal actions to the node and generating all successor states.	Frontier or Open List – Set of successor nodes that have not yet been expanded.
Search Strategy: Method for choosing the node on the frontier to next expand.	Repeated State: Any state visited more than once during a search. Redundant Path: Any two or more paths that go to the same state.	Closed or Explored Set: States that have already been expanded.	Loopy Path – Where a repeated state is expanded causing you not to continue to explore the same section of a graph.

## **Definitions:**

<b>Uniformed Search</b> – Also known as <b>(Blind Search)</b> is any search that has no information on the search space.	<b>Informed Search</b> – Uses <b>heuristics</b> that inspect the state space to prioritize moves.	Explored Set – Set of all nodes already visited.
Branching Factor (b) – Number of branches/children/successors from a given node. Generally lists as the maximum branching factor.	Depth (d) – Number of branches/children/successors from a given node.	Frontier Set – Set of all nodes available for expansion.

## A Problem consists of five attributes:

- 1. Initial State
- 2. Set of possible actions (ACTIONS)
- 3. Successor Function/Transitional Model (RESULTS)
- 4. Goal test (TERMINAL-TEST)
- 5. Cost Function

## Four Ways to Rate/Measure a Search Strategy:

- 1. Completeness If a solution exists, does the algorithm always find it?
- 2. Optimal Is the solution found by the algorithm always optimal (i.e. have the lowest cost).
- 3. Time Complexity Amount of time required by the algorithm to perform the search.4. Space Complexity Amount of memory required by the algorithm to perform the search.
- Memory Queue **Time Complexity** Complete **Optimal** Comments Name Complexity Type Used l is the maximum allowed depth. **Depth Limited Search**  $O(b^l)$ 1. Incomplete if d > l0(l) Nο Nο Stack 2. Can be non-optimal if l > d1. Not complete because of the infinite branching problem (e.g. loop). Yes if the graph is Depth-First Search 0(d)  $O(b^d)$ No Stack 2. Can be considered special case of depth-limited search finite, No with  $l = \infty$ otherwise Always expand left most node that can be expanded. Iterative Deepening  $O(b) + O(b^2) + \cdots$ 0(d) Yes Yes Stack Calls Depth Limited Search algorithm d times  $+ O(b^d) = O(b^{d+1})$ Depth First Search Can be considered a variant of uniform cost search where Yes if each step cost is the same. uniform **Breadth First Search**  $O(b^d)$  $O(b^d)$ Queue Yes Expand the root node and then expand all children of the step root node in the order they are encountered until all nodes cost are expanded or a goal is reached. Variant of Breadth-First Search where two breadth first Yes if searches (one from start and one from the goal) are initiated and carried out simultaneously.  $O\left(b^{\frac{d}{2}}\right)$ uniform  $O\left(b^{\frac{d}{2}}\right)$ **Bidirectional Search** Yes Queue step Generalization of Breadth-First where the root (i.e. initiate state) node is expanded first and nodes are expanded based cost of their non-decreasing distance/cost from the root. Variant of Breadth-First Search where the step cost is not  $O\left(b^{1+\frac{C^*}{\epsilon}}\right)$ Priority uniform.  $O\left(b^{1+\frac{C^*}{\epsilon}}\right)$ **Uniform Cost Search** Yes Yes  $C^*$  - Minimum (optimal) cost to the goal. Queue  $\epsilon$  - Minimum step cost Selects node for expansion based off the one with the lowest **Greedy Best First** heuristic cost. N/A N/A No Nο None Search f(n) = h(n)Can oscillate in a dead end condition. Based off Yes with Based off quality of Priority Α\* quality of Yes heuristic heuristic Queue heuristic conditions Yes if Based off quality of **Recursive Best First** O(d)Yes heuristic Stack Search heuristic

admissible

 $Completeness\ above\ assumes\ the\ branching\ factor\ is\ \textbf{finite}.$ 

Iterative Deepening Depth First Search (also known as Iterative Lengthening Search)

```
def Depth_Limited_Search(node, problem, depth):
                                                                                             if(problem.GOAL_TEST(node)):
def ID_DFS(problem, limit):
                                                                                                  return SOLUTION(node)
     # Incrementally increase the maximum depth
                                                                                             if(depth == 0):
     for maximum depth in range(0, limit):
                                                                                                  return None
          result = Depth_Limited_Search(problem.INITIAL_STATE(),
                                                                                             for action in problem.ACTIONS(node):
                                         problem, maximum_depth)
                                                                                                  child = problem.RESULT(node, action)
          # If solution found return it.
                                                                                                  result = Depth_Limited_Search(child, problem, depth - 1)
          if(result is not None):
                                                                                                  if(result is not None):
                return result
                                                                                                        return result
                                                                                             return None
```

Space Complexity: O(d) since at one time only keeping in memory at most d nodes.

Time Complexity: Depth-Limited-Search is called up to d times. Each call to Depth-Limited-Search takes  $O(b^m)$  time. Given:  $\sum_{i=m}^{n-1}a^i=\frac{a^m-a^n}{1-a}$ , Then  $b+b^2+b^3+\cdots+b^d=O(b^{d+1})$ 

Complete: Yes since all nodes are explored if  $d \leq limit$ 

Optimal: Yes if all steps have uniform cost.

## **Uniform Cost Search (Uniformed Search)**

Uniform cost search explores nodes on the frontier based of a monotonically increase cost function. Hence its evaluation function is:

```
f(n) = c(n) also referred to as f(n) = g(n)
```

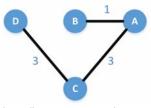
```
def UCS(problem):
     initial_state = problem.INITIAL_STATE()
     priority queue = {}
     explored set = {}
     priority_queue.enqueue(initial_state)
     # Continue until either a solution is found or all nodes explored.
     while( len(priority_queue) > 0):
          node = priority_queue.pop()
          # Must only check AFTER dequeueing the item to ensure it is optimal.
          if(problem.GOAL_TEST(node)): return SOLUTION(node)
          # Add the node to the explored set.
          explored_set.append(result)
          for action in problem.ACTIONS(node):
                result = problem.RESULT(node, action)
                # If not in the priority queue then enqueue it.
                if( result not in priority_queue and result not in explored_set):
                      priority_queue.enqueue(result)
                # Current version of node has lower cost than version in priority queue
                elif( result in priority_queue and result.COST() < priority_queue[result]. COST()):</pre>
                      priority_queue.remove(result)
                     priority_queue.enqueue(result)
     # No path found
     return None
```

Pseudo code for A\* and UCS is the same with the implementation of the COST() method.

## A\* Algorithm

A* algorithm is a combination of the benefits of Greedy-Best First Search and Uniform Cost Search. Evaluation Function $f(n)$ : $f(n) = g(n) + h(n)$ Also written as: $f(n) = c(n) + h(n)$	Only performs the GOAL-TEST after the node has been dequeued from the priority queue. Similar to Uniform Cost Search.	Derives from Dijkstra's Algorithm.
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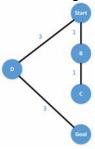
# Example of A\* Performing Better than Greedy Best First Search



Greedy Best First Search Oscillates Between Nodes A and B so it is Incomplete. This graph is solvable by  $A^*$ .

Greedy Best First Search is **memory efficient** since it does not need to remember where it has been.

## **Example of DFS Performing Better than A\***



Heuristic for A\* is Euclidean distance. In this case, A\* adds B then D to the frontier. It next expands B and adds C to the frontier. It next explores C and finds no solutions so it explores D then finds the goal.

#### **Recursive Best-First Search**

This algorithm is **optimal** when the heuristic is **admissible for trees**. The heuristic needs to be **consistent for tree search** to be optimal.

f\_limit/min\_eval\_func\_val – Best alternative path available from the any ancestor of the current node.

#### Simplified Description of Recursive Best First Search

- 1. Start from initial state and set the initial minimum cost of ∞
- 2. Generate all successors of current node. Set successor cost to either current node evaluation function value (f(n)) or the successors evaluation function cost.
- 3. Select successor node with minimum evaluation function (f(n)) cost.
- 4. If current node is a goal state, then return the solution.
- 5. If this cost is more than the current minimum, backtrack to find node with current minimum.
- 6. Extract the evaluation function cost (f(n)) of the second best successor of the current node.
- 7. Recurse using best successor found in step #3 and the minimum of the current minimum cost that was passed to the function and the second best successor of this node. This function results either a solution or None and updates the current best node's evaluation function cost (f(n)).
- 8. If step #7 returned a solution, then return that, otherwise, jump to step #3.

```
# Continues to recurse until current best cost is more than

def RBFS(problem, state, min_eval_func_val):

# Check if a goal was reached. If so, return it.

if( problem.GOAL_TEST(state) ):

return SOLUTION(state)

# Get set of successors
```

for a in problem.ACTIONS(state):
 successors.append(problem.RESULT(state, a))
# Check a successor exists

return RBFS(problem, problem.INITIAL\_STATE(), inf)

If(len(successors) == 0): return None, ∞

# Update all successor eval function values

def RECURSIVE BEST FIRST SEARCH(problem):

for s in successors:

s.eval\_func\_val = max(state.eval\_func\_val, s.g + s.h)

while(True):

# Best successor is a node with min eval cost from successors

best\_successor = node with least eval function value from the **successors** 

# If the best successor is not better than current best, backtrack to current best

if(best\_succesor.eval\_func\_value > min\_eval\_func\_value):

return None, best\_successor.eval\_func\_value

# May need to recurse back to current level so store second best value for this level. second\_best\_successor\_eval\_func\_val = Eval func value for second best successor of state

# Run RBFS again from current node with the new min value the minimum of the current

# minimum and the second best successor (i.e. alternative) for this current state/node.

result, best\_successor.eval\_func\_val = \

RBFS(problem, best\_successor, min(min\_eval\_func\_val,

 $second\_best\_succesor\_eval\_func\_val)$ 

# If solution found, return it. if(result is not None):

return result

## **Memory Bounded Heuristic Search**

## Iterative Deepening A\* (IDA\*) Algorithm

Variant of the A\* algorithm that *generally* slower but uses less memory. Sets a maximum total cost (i.e. f(n)) to a starting value of  $\mu$ . In each round, any node whose total cost (i.e. f(n)) is greater than the maximum is ignored. Perform A\* for thresholds:

$$\mu < 2\mu < 3\mu < \cdots$$

```
def IDA_Star(problem, initial_max_cost, maximum_cost):
    current_max_cost = initial_max_cost
    while(current_max_cost < maximum_cost):
        result = A_Star_Search(problem, current_max_cost)
        if(problem.GOAL_TEST(result)):
            return result
            current_maximum_cost += initial_max_cost
        return None</pre>
```

## Simplified Memory Bounded A\*

Approach to save memory in A\* algorithm. **Procedure:** 

- Perform A\* until you run out of memory.
- 2. Delete fringe or explored set node with the worst cost.

## Evaluation Functions f(n) for Three Related Search Algorithms:

Uniform Cost Search: f(n) = c(n)

Greedy Best First Search: f(n) = h(n)

A\* Search Algorithm: f(n) = c(n) + h(n)

A\* algorithm is the only one of the three whose evaluation function estimates the cost of the total solution.

Admissible (Optimistic) Heuristic: Any heuristic that never over estimates the cost of a solution.

Consistent (Monotonic) Heuristic: For every node, n, every successor, n', that is reached by action, a, then the cost to reach the goal from n is less than or equal to the actual cost to go from n to n' by action a(c(n, a, n')) plus the heuristic cost of n'.

 $h(n) \le c(n, a, n') + h(n')$ 

Note: Any heuristic that is consistent is also admissible.

Example: Triangle Inequality when the heuristic is straight-line distance.

The tree-search version of A\* (i.e. DAG) is optimal if h(n) is admissible, while the graph search version of A\* is optimal if h(n) is consistent.

## Lemma #1

If h(n) was a consistent heuristic, then the values of f(n) are nondecreasing.

Given a node n' is a successor of n through action a, then:

$$g(n') = g(n) + c(n, a, n')$$

If h(n) is consistent, then:

$$h(n') + c(n, a, n') \ge h(n)$$

Then:

$$g(n') + h(n') + c(n, a, n') \ge g(n) + h(n) + c(n, a, n')$$

$$f(n') + c(n, a, n') \ge f(n) + c(n, a, n')$$

$$f(n') \ge f(n)$$

Lemma #2: Whenever A\* selects a node for expansion, the optimal path to that node has been found.

Had lemma #2 not been the case, then there would have been another node n' on the path from the start to nthat would have been on the optimal path. Because f(n) is non-decreasing, this node would have had a lower value of f(n) and would be expanded before nin A\*. Hence, this is a contradiction.

## Combining Lemma #1 and Lemma #2

By Lemma #2: If a goal node is explored, it is the optimal path to that goal node.

By Invariant of A\*: A\* algorithm explores nodes in nondecreasing order of f(n).

By Lemma #1: f(n) is nondecreasing.

Combining Lemma #1, Lemma #2, and Invariant of A\*: Paths to any other unexplored states, including goal states, will have evaluation function values (f(n))greater than the first one explored. Hence, the optimal path to the first explored goal state is the optimal solution to the entire problem.

Since by lemma #2 A\* returns the optimal path to the first goal state, it returns the optimal path to the entire problem.

## **Choosing a Heuristic**

Effective Branching Factor (b\*): For a set of N moves, it is the equivalent number of uniform branches for a depth d. It is a way to quantify the quality of a heuristic.

$$N+1=1+b+b^2+\cdots+b^d$$
 
$$N+1=\frac{b^{*d+1}-1}{b^*-1}$$
 Derives from:

$$\sum_{i=m}^{m-1} a^i = \frac{a^m - a^n}{1 - a}$$

Best branch possible factor is 1.

Relaxed Problem: A version of the actual problem with fewer restrictions.

An exact solution to a relaxed problem is an admissible heuristic for the original problem.

**Dominating Heuristic:** A heuristic that always has a lower branching factor than another heuristic.

Composite Heuristic: Given a set of admissible heuristics  $\{h_1, h_2, ..., h_n\}$  none of which is dominating, then the best heuristic is the composite heuristic:

 $h_{composite} = \max\{h_1, h_2, ..., h_n\}$ 

Subproblem: A reduced version of the actual problem. Admissible heuristics can be derived from the solution to subproblems.

Pattern Database: Stores the exact solution for all versions of a particular subproblem.

To determine the heuristic cost for a version of the subproblem, look up the solution in the database and calculate the heuristic cost.

Disjoint Patterns: A problem can be divided into disjoint (i.e. nonoverlapping) subproblems. The disjoint solution to the problem is referred to as a disjoint pattern.

**Disjoint Pattern Database: Stores** solution to disjoint (non-overlapping, non-dependent) subproblems.

Using multiple disjoint subproblems in a disjoint pattern database, you can come up with a composite heuristic by summing the cost to solve each individual subproblem.

## **Local Search**

**Local search** generally operates using a single **current node** and generally moves to neighbors of that node.

If the local search problem is an **optimization problem**, then it is accompanied by an **objective function** that is to be maximized or minimized.

**Complete Algorithm:** Always finds a solution if it exists.

**Optimal Algorithm:** Always finds a global maximum or minimum.

State Space Landscape: Landscape has a location (i.e. state) and an elevation (utility from the objective function)

## **Hill Climbing Algorithm**

Local search algorithm that always proceeds to the next successor state with maximum utility. If two successors have the same utility, algorithm randomly chooses between them. Susceptible to local maxima.

Also referred to as Greedy Local Search.

## **Variants of Hill Climbing**

**Sideways Move:** Allow hill climbing algorithm to move to a state of equal value. Helps to move past flat area in a graph. However, in a plateau, it can lead to an infinite loop so a limit on the number of consecutive sideways moves is common.

**Stochastic Hill Climbing:** Choose a successor state at random with the probability each successor is selected proportional to its utility.

**Hill Climbing with Restarts:** Hill climbing runs from a randomly chosen initial state. If it gets a solution, it returns. Otherwise, it generates another random initial state and repeats the process. Repeated *n* times or until a solution is found.

**Example:** If the probability of finding a solution from an initial state is p, then it is expected  $\frac{1}{n}$  restarts will be required.

See page 122.

```
def HILL CLIMBING WITH RESTART(problem, max restarts):
     while( max restarts > 0):
          max_restarts -= 1
          problem.INITIAL_STATE = problem.RANDOMIZE_STATE()
          result = Hill Climbing(problem)
          if(problem.GOAL TEST(result)):
                return result
     return None
def HILL_CLIMBING(problem):
     current_state = problem.INITIAL_STATE()
     while( True ):
          # Update the previous utility
          best successor = None
          # Iterate through set of possible actions
          for action in state. ACTIONS():
                new_state = problem.RESULTS(state, action)
                if(best_successor is None
                  or problem.UTILITY(new_state) > problem.UTILITY(current_state)):
                     best_successor = new_state
          # Determine if the best successor is better than the current state
          if(problem.UTILITY(best_successor) > problem. UTILITY(current_state)):
                current_state = best_successor
          else:
                return current_state
```

Note: This is a goal based version of Hill Climbing. If you are simply searching for a maximum or minimum, you would need to modify the algorithm to return "current state" at the end.

## **Simulated Annealing**

Can be used for either maximization or minimization problems.

Algorithm is designed to allow the current\_node to move to a worse state with decreasing probability as time progresses.

Probability of Moving to a Lower Value Solution is:

$$P = e^{-\frac{\Delta k}{schedule(t)}}$$

Simulated annealing chooses a random successor.

```
import math
import random
def SIMULATED ANNEALING(problem, schedule, limit, t_min):
     current state = problem.INITIAL STATE()
     t = 0
     while( True ):
           t += 1
           T = schedule(T)
           if(T < t_min or problem.GOAL_TEST(current_state) ):
                return current_state
           # Get the set of actions.
           actions = current_state.ACTIONS()
           # If no successors possible, terminate
           if(len(actions) == 0):
                return current_state
           # Randomly select a successor
           a = actions[random.randint(0, len(actions) - 1]
           # Get the successor state
           next_state = problem.RESULT(current_state, a)
           # Calculate the error
           error = problem.UTILITY(next_state) - problem.UTILITY(current_state)
           # If error is positive or probability less than specified number, then update the current state.
           if(error > 0 or random.random() < math.exp( error/ T ):</pre>
                current_state = next_state
```

Note: This version of the code is a maximization problem. Would need to modify slightly for a

## **Local Beam Search**

minimization problem.

Type of local search.

#### Procedure:

- 1. Begin with k randomly generated states.
- 2. Check if any descendent states at the goal. If so, return state.
- 3. Order all successors from the k states and sort them by decreasing performance.
- 4. Choose the best *k* successors. If any successor has performance measure better than the current best, return to step #2.

The k successors are considered a **pool of candidates**. The successors are considered **offspring**.

#### Variant of Local Beam Search

Stochastic Local Beam Search: Choose *k* successors stochastically based off some metric.

## **Genetic Algorithm**

def GENETIC\_ALGORITHM(problem, FITNESS\_FUNCTION, t\_max)

```
# Generate the population.
                                                                     population = problem.GENERATE_POPULATION()
                                                                     # Start at time 0.
                                                                     t = 0
                                                                     while(t < t_max or Not problem.GOAL_TEST(best_solution)):</pre>
                                                                           # Increment current time.
A genetic algorithm is a stochastic beam search algorithm
with one key modification:
                                                                          new_population = {}
      In local beam search, successors come from modifying
                                                                          best_solution = None
       a single state (asexual reproduction).
                                                                          for i in range(0, problem.POPULATION_SIZE()):
      In genetic algorithm, successors come from combing
                                                                                # Select two parent solutions.
                                                                                x = RANDOM_SELECTION(population, FITNESS_FUNCTION)
       two parent states (sexual reproduction).
                                                                                y = RANDOM_SELECTION(population, FITNESS_FUNCTION)
Population: Set of k solutions. The initial population is k
                                                                                # Merge the two solutions
                                                                                child = REPRODUCE(x, y)
randomly generated solutions.
Individual: One solution/state in the population.
                                                                                # Mutate on a low probability
                                                                                if(random.random() < problem.MUTATION PROBABILITY):</pre>
                                                                                     problem.MUTATE(child)
Fitness Function: Evaluation function that rates the quality
                                                                                if(best_solution is None or problem. UTILTY(best_solution) < problem.UTILTY(child)):</pre>
(i.e. fitness of a solution) generally with general condition
                                                                                     best solution = child
that better states have higher fitness function value.
                                                                                # Add the child solution to the new population.
                                                                                new_population.append(child)
Crossover: Process of merging two solution states to form a
                                                                          # Set the population to the newly created set.
new successor.
Mutation: Random change to a successor solution.
                                                                          population = new_population
                                                                     return best_solution
                                                                def REPRODUCE(x, y):
                                                                     # Pick a random cross over point
```

## 8-Puzzle Goal State:

# Crossover the two halves

 $crossover_point = random.randint(0, len(x) - 1)$ 

return x[0:crossover\_point] + y[crossover\_point:len(y)]

X	1	2
3	4	5
6	7	8

## Minimax (Adversarial Search)

Adversarial search problems are those search problems that arise in multiagent, competitive environments. Adversarial search problems are also known as games.

In a zero-sum game, the results for the two players are always equal and opposite.

**Optimal Strategy** – A sequence of contingent decisions that will lead to outcomes as least as good as any other sequence of decisions against an infallible player.

**Perfect Information** – Any situation where an agent has all relevant information with which to make a decision and the results of actions are **deterministic**.

Minimax Value — Utility of being in a current state assuming both players play optimally until the end of the game.

```
H - MINIMAX(s, d) = \begin{cases} \max_{a \in ACTIONS(s)} H - MINIMAX(RESULT(s, a), d + 1), \text{ if } PLAYER(s) \text{ is MAX} \\ \max_{a \in ACTIONS(s)} H - MINIMAX(RESULT(s, a), d + 1), \text{ if } PLAYER(s) \text{ is MIN} \end{cases}
```

#### Initial State in Minimax - 50

Given a state, s, the six key methods used on that state are:

- 1. PLAYER(s) Returns active player for the current state
- 2. ACTIONS() Set of all possible actions/moves that can be made.
- 3. RESULTS(s,a) Given a state, s, and an action a, it returns the successor state. It is also called a Transitional Model.
- 4. CUTOFF\_TEST(s,d) Used in Heuristic minimax. Given a state, s, and a recursive depth, d, it determines if the cutoff condition of either a maximum depth or goal state has been reached.
- 5. TERMINAL\_TEST(s) Used in standard minimax. Given a state, s, this function returns whether a goal state has been met. Terminal states are leaf nodes in the search tree.
- 6. UTILITY(s) Given a state, s, this function returns the state's utility score. It is also called a Utility Function.
- 7. EVAL(s) Given a state, s, this function estimates how good (i.e. the utility) of a given state.

```
Time Complexity with Alpha-Beta Pruning: O\left(b^{\frac{d}{2}}\right)
```

Time Complexity without Alpha-Beta Pruning:  $O(b^d)$ 

```
def Minimax_Algorithm(state, is_max):
     alpha max = -inf
     beta min = inf
    best_successor = None
     # Iterate through all possible actions from this state
     for a in state. ACTIONS():
          # Get the successor state
          next state = state.RESULT(state,a)
          # Call heuristic minimax with starting depth 0
          score = H-Minimax(next_state, 0, !is_max,
                            alpha max, beta min)
          if(is max and score > alpha max):
                best_successor = a
                alpha_max = score
          elif(not is max and score < beta min):
                best successor = a
                beta_min = score
     # Return the move with the best score
     return best_move
```

```
def H-Minimax(state, depth, is_max, alpha_max, beta_min)
     # p is the reference player for the utility function. Typically max.
    if ( state.CUTOFF-TEST(depth) ):
          return state. UTILITY(PLAYER(p))
    for a in state. ACTIONS():
          next_state = state.RESULT(state, a)
          if(is_max):
               # Perform beta pruning
                alpha_max = max(alpha_max, H-Minimax(next_state, depth+1,
                                  not is_max, alpha_max, beta_min))
               if(alpha max ≥ beta min):
                     return alpha max
          else:
               beta_min = min (beta_min, H-Minimax(next_state, depth+1,
                                not is max, alpha max, beta min))
                # Perform alpha pruning
               if(alpha_max ≥ beta_min):
                     return beta min
     # After all actions tested, return score.
    if(is_max):
          return alpha max
          return beta_min
```

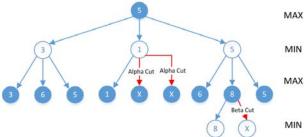
## **Alpha Beta Pruning**

Alpha (α) – Maximum value found along the path by the MAX player.

**Alpha Cut/Alpha Pruning** – Performed by the **MIN player**. When the MIN player's minimum score is already less than a previous MAX player's maximum score, stop investigating subsequent paths and return the **current minimum score**.

Beta  $(\beta)$  – Minimum value found along the path by the MIN player.

**Beta Cut/Beta Pruning** – Performed by the **MAX player**. When the MAX player's maximum score is greater than a previous MIN player's minimum score, stop investigating subsequent paths and return the **current maximum score**.



Minimax Search Tree Example with Alpha and Beta Cuts.

This is a three move/ply search tree.

## **Constraint Satisfaction Problem**

Search problems deal with states that are atomic (i.e. indivisible).

Often a state has field variables. Such field values are called a **factored representation** of the problem. A state **solves** a factored representation if each field variable satisfies all constraints on that variable.

A factored representation can allow you to eliminate large areas of the search space by identifying then ignoring variable/value combinations that violate constraints.

A constraint satisfaction problem solution is an assignment of values to variables that satisfies all constraints.

Assignment of values to variables in CSPs is **commutative**. Hence, the order that the values are assigned do not matter. If you consider the problem a search tree, there are at most *d* children from each node leaving a total of *d*<sup>n</sup> solutions for a finite domain

## **Components of a Constraint Satisfaction Problem:**

- 1. X Set of variables  $\{X_1, X_2, X_3, \dots, X_n\}$
- 2. D Set of Domains  $\{D_1, D_2, D_3, \dots, D_n\}$
- 3. C Set of Constraints  $\{C_1, C_2, C_3, \dots, C_n\}$

#### **Optional Definition:**

assignment.

R - Relation of multiple variables  $R(X_i, ..., X_m)$ 

**Definition of a Constraint** 

A constraint is a pair: < scope, relation >

Scope: Tuple of variables that participate in the constraint

**Relation:** A relation that the variables can take on.

**Assignment** – Allocation of values to variables.

Solution: A complete and consistent

**Consistent Assignment** – An assignment of values that does not violate any constraints.

This leads to the term **consistency** which is the **satisfaction of constraints**.

**Complete Assignment** – Every variable is assigned a value.

**Partial Assignment** – Only a subset of variables are assigned a value.

Domain

A variable's domain can be either discrete or continuous. If it is discrete, it can be either finite or infinite (e.g. set of integers).

Simplest CSP Type: Finite, discrete domain

**Constraint Language** 

Defines the allowed relations between variables. It eliminates the need to enumerate allowed value lists.

**Linear Programming Problem:** Continuous CSP with **linear constraint function**(s).

Constraint functions can also be nonlinear.

Constraint Types

 $X = \{X_1, X_2\} \text{ and } D = \{A, B\}$ 

2, 2,

Example Constraint:  $C = \{(X, X_n) | r\rho\}$ 

 $C = \langle (X_1, X_2), rel \rangle$ 

 $rel = \{(A, B), \qquad (B, A)\}$ 

**Precedence Constraint:** A

constraint that forces one variable to occur before (i.e. be less than) another variable.

xample:

 $T_1 + d \le T_2$ 

Disjunctive Constraint: A

constraint that two variables do not overlap (i.e. are not equal):

Example:

 $T_1 + d \le T_2 \text{ or } T_2 + d \le T_1$ 

Absolute Constraint: Any constraint that must be met.

Preference Constraint: A constraint which guides the solution to preferred values.

Problems that optimized preference constraints are called **constraint optimization problems**.

**Unary Constraint** – A constrain involving only a single variable.

**Binary Constraint** – A constraint involving exactly two variables.

Higher Order Constraint: A constraint that involves a fixed number of variables that is more than two.

All higher order constrains can be reformed as a set of binary constraints.

Global Constraint: A constraint that takes an arbitrary number of variables. It does not need to be all variables. It just needs to be **not fixed** (i.e. arbitrary).

Example: Alldiff

Constraint Graph/CSP Network: Representation of a CSP as a graph. Each node is a variable and the arcs are binary constraints.

**Inference:** Using known/assigned values for a set of variables to select the values for other variables.

Constraint Propagation: Using the constraints to reduce the number of legal values for a variable. This in turn reduces the number of legal values for other variables in a cycle.

**Local Consistency:** Given a constraint graph, enforcing consistency (i.e. ensuring variables satisfy constraints) locally **in each part of the graph** leads to invalid values being eliminated throughout the graph.

**Node Consistency** 

Node Consistent Variable – Any variable where every value in the variable's domain satisfies all of its unary constraints in a CSP network.

Node Consistent Network – Any CSP network where all variables are node consistent.

Node consistency can be done as a **preprocessing step** to eliminate invalid values.

## **Arc Consistency**

Arc Consistent Variable – Any variable where every value in the variable's domain satisfies all of its binary constraints in a CSP network.

Variables are arc-consistent with respect to one another. Example: X being arc consistent with respect to Y does **NOT** imply Y is arc consistent with respect to X.

Arc Consistent Network - Any CSP network where all variables are arc consistent.

## AC-3 (Arc Consistency Algorithm #3)

Algorithm used to solve for Arc consistency Only possible with finite domains.

## **Constraints in Arc Consistency Algorithm**

In each iteration of AC-3 algorithm, it only checks the variable being arc-constrained (example in constraint (X,Y), X is being constrained by Y). To have a two directional constraint for X and Y, arc queue would need to contain (X, Y) and (Y, X)

After reducing the domain of X from constraint (X, Y), algorithm needs to recheck any domains that were constrained by X to ensure its domain values are still valid.

## **Running Time of AC-3 Algorithm**

## 1. REVISE Function: $O(d^2)$

For each value in the domain of  $X_i$  (up to delements), you iterate overall elements in the domain of  $X_i$ . Hence the running time is:

$$O(d*d) = O(d^2)$$

## 2. Number of Times REVISE function is Run Per Constraint: O(d)

The REVISE function is run whenever a constraint is popped off the queue. If the domain size is d, it can be popped off the queue up to d times (once for each element in the domain).

## 3. Number of Constraints: c

## **Total Running Time:**

$$\mathbf{O}(c) \cdot \mathbf{O}(d) \cdot \mathbf{O}(d^2) = \mathbf{O}(cd^3)$$

```
def AC_3(csp):
     arc_queue = []
     # Add all binary constraints to the queue.
     for b constraint in csp.BINARY CONSTRAINTS:
           arc_queue.append( (b_constraint.X_i, b_constraint.X_j )
     # Iterate until all arcs have been made consistent or an inconsistency is found.
     while(len(arc queue) > 0):
           (X_i, X_j) = arc_queue.pop()
           # Check if the domain of X i is revised.
           if( REVISE(csp, X_i, X_j) ):
                 if(len(X_i) == 0):
                      return False
                 # Only X_i's domain is reduced in function "REVISE" so only check relative to that.
                 # Since X_i's domain is reduced, any variable that is constrained by X_i may need to be reduced
                 for X k in X i.NEIGHBORS() - {X j}:
                      # Only add back to domain if not X j
                      if( (X_k, X_i) not in arc_queue):
                            arc_queue.append((X_k, X_i))
     return True
def REVISE(csp, X_i, X_j):
     revised = False # Confirmed in loop
     # Verify all elements in the domain of X_i have a corresponding value in X_j.
     for x in csp.D_i:
           constraining_value_exists = False
           # Iterate through all elements in X_j's domain to see if it constrains x in X_i.
           for y in csp.D_j:
                 if( (x,y) in csp.C(X_i, X_j)) :
                      constraining_value_exists = True
           # If no constraining value exists in X_j, then remove the value from X_i.
           if(not constraining_value_exists):
```

## **Path Consistency**

# Return whether the domain of X\_i was revised (i.e. reduced)

csp.D\_i.remove(d)

revised = True

return revised

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Path Consistency – A two variable set  $(X_i, X_i)$  are path consistent with respect to a third variable  $X_m$  if for every assignment of values to  $X_i$  and  $X_i$ consistent with the constraint  $\{X_i, X_i\}$ , there is a valid assignment to  $X_m$ that satisfies the constraints  $\{X_i, X_m\}$  and  $\{X_m, X_i\}$ .

## Origin of the Term "Path Consistency"

Given a two variable set  $\{X_i, X_i\}$  that is path consistent with respect to a variable  $X_m$ , then it is like  $X_m$  is on the path between  $X_i$  and  $X_i$ .

Algorithm to Solve to Check for Path Consistency: PC-2

## k-Consistency

A CSP is k-consistent if for any set of k-1 variables and for any consistent assignment to those variables, a consistent value can always be assigned to any k-th variable.

Proving k-consistency takes exponential and space in the worst case.

1-consistency is node consistency.

2-consistency is arc consistency. Strongly k-consistent: Any CSP that is 1-consistent and 2-consistent and 3-consistent through k-consistent. Hence it is consistent for variable sets of size 1 through k.

Given *n* variables and a CSP that is strongly *n*-consistent, then an assignment of values is possible for this CSP.

Running Time to Solve n-Consistent CSP

## Time Complexity: $O(n^2d)$

Running time derives since for every i-th variable to assign, you must check all i-1 variables for every d elements in the

$$d \cdot \sum_{i=1}^{n} i - 1 = d \cdot \left( \frac{n \cdot (n+1)}{2} - n \right) = O(dn^2)$$

## **Consistency Checks for Global Constraints**

**Global Constraint** – A constraint with an arbitrary number of variables.

**Example Global Constraint:** Alldiff

## Alldiff Consistency Algorithm

- 1. Delete a variable that has a singleton domain.
- 2. Remove the value from the domains of all other variables.
- 3. If any singleton domain variables still exists, jump to step #1.
- 4. If a domain has no values or there are more values than there are variables, the Alldiff constraint fails.

## **Simplified Explanation of Alldiff Consistency Check**

If there are *m* variables and *n* possible values and m > n, then an inconsistency exists.

## Sudoku

Square grid of n by n cells. All numbers in a row must be unique and all numbers in a column must be unique. For every  $\sqrt{n}$  by  $\sqrt{n}$  subgrid, all numbers must be unique. Each section of the board where all numbers must be unique (e.g. row, column, subgrid) is called a **unit**.

Formal Definition of Sudoku as a CSP:

Variables:  $n^2$  total variables (one for each cell).

**Domain:**  $\{1, 2, 3, ..., n\}$ 

Constraints: 3n Alldiff constraints for each unit.

AC-3 Algorithm can be used to infer the value of cells and to reduce the domains of cells.

## **CSPs and Backtracking**

Backtracking Search – Variant of Depth First Search where values are assigned to variables until no consistent, legal assignments are possible for a given variable at which point the algorithm backtracks to try to reassign a previous variable to a new value.

**Key Functions in Backtracking Search** 

- 1. SELECT UNASSIGNED VARIABLE
- 2. ORDER\_DOMAIN\_VALUES
- 3. INFERENCE
- 4. BACKTRACK (recursion)

See page 215.

```
def BACKTRACKING SEARCH(csp):
     return BACKTRACK({}, csp)
def BACKTRACK(assignment, CSP):
     # Consistency of all variable assignment checked so if assignment is complete, it is a solution.
     if(csp.COMPLETE_ASSIGNMENT(assignment)) return assignment
     # Select the next variable to assign
     next var = csp.SELECT UNASSIGNED VARIABLE()
     # Order the domain values based off which want to check first
     var_doman = csp.ORDER_DOMAIN_VARIABLES(assignment, next_var)
     # Iterate through all domain values.
     for d in var_domain:
          # Ensure the assignment is consistent.
          if(csp.CONSISTENT_ASSIGNMENT(assignment, d)):
               # Add the variable value to the assignment
               assignment[var_domain] = d
               # Get and apply any inferences
               inferences = csp.INFERENCE(assignment)
               # Only recurse if valid inferences found.
               if(inferences is not None):
                    assignment.APPLY_INFERENCES(inference)
                     result = BACKTRACK(assignment, csp)
                    if( result is not None):
                          return result
                     assignment.REMOVE INFERENCES(inference)
               # Since no solution found using this assignment and variable value
               # remove this variable value from the assignment.
               remove( assignment[var_domain] )
     # No solution found so return None for failure.
     return None
```

## **Making Backtracking Search More Efficient and Sophisticated**

## **Variable Ordering**

By selecting a variable most likely to fail earliest, you are prune the search tree and reduce the effective branching factor.

Minimum Remaining Value (MRV), Fail First, Most Constrained Variable Heuristic: Select the variable to assign next that has the smallest inferred domain (i.e. least remaining legal values).

Degree Heuristic: Select the variable for expansion that has the largest number of constraints on other variables. Most commonly used heuristic to select the first variable for assignment.

Degree heuristic can be used as a tie breaker for the more powerful MRV heuristic.

## Value Ordering

Constraining

Least-

Value
Heuristic:
Select the
value that
rules out the
least number
of values for
neighboring
variables in
the graph.

# Interleaving Search and Inference

AC-3 can be used to infer reductions in the search domain both **before and during** search.

Forward Checking – One way to implement "Inference" in Backtracking algorithm. Whenever a variable is assigned, establish arc consistency for it on all unassigned variables. If arc consistency checking was done in preprocessing, forward checking adds no value.

MRV can be combined with forward checking to further prune the search tree.

Chronological Backtracking: Simplest form of backtracking. Revisit the last assigned variable (i.e. most recent decision) before the current variable. If the previous variable does not constrain the current variable, backtracking to only that level is wasteful.

## **Intelligent Backtracking**

Better to backtrack to a variable that may fix the consistency issue.

**Conflict Set:** Set of value assignments that conflict with a some value for a variable. **Note:** This is value assignments not variables since a variable that can conflict for one value does not conflict for the currently assigned value.

**Backjumping:** Backtracking to the most recent variable in the conflict set.

Variable ordering is fail-first ordering while value ordering is fail-last. This is because when you are trying to fail-first by selecting a variable, the order you inspect the values does not matter as you need to inspect them all anyway. As such, it makes the most sense to inspect the best solutions first in case one of them does actually succeed.

## **Logical and Knowledge Based Agents**

Knowledge Base (KB) - Central component of a knowledge based agent. Composed of a set of sentences. Similar to a database.

Knowledge Representation Language - Formal notation used to express sentences in the knowledge base (KB). Sentence – Statements that define the knowledge based. They have a specific notation called a syntax and their value (i.e. true or false) is defined by the semantics.

Axiom – A sentence that is taken as given without being derived from other sentences.

Inference – Deriving new sentences from existing sentences.

#### Valid Knowledge Base Operations:

- TELL
- ASK 2.

#### **Supporting Knowledge Based Agent Commands:**

- 1. MAKE PERCEPT SENTENCE
- 2. MAKE\_ACTION\_QUERY
- MAKE\_ACTION\_SENTENCE

Background Knowledge – Initial knowledge in the knowledge base.

#### Four Step Procedure for a Knowledge Based Agent:

- Tell the knowledge base what it perceives.
- Ask the knowledge base it should perform.
- Tell the knowledge base the action it will perform. 3.
- Executive the action.

Knowledge Level – What the agent knows at a give point in time. Given an agent's knowledge level and goals, you can predict its actions.

Declarative Approach – Tell the knowledge base all it needs to know. Describe what you want the agent to do and get it via a black box. Procedural Approach – Procedures for desired behaviors and actions are hard coded into the agent. Teach the agent how to do things and it finds the solution.

def KNOWLEDGE BASED AGENT()

t += 1 # Increment time

return action

# Return the selected action.

# KB is the persistent knowledge base.

# t a time counter initially starting at 0.

TELL( KB, MAKE\_PERCEPT\_SENTENCE(t) )

TELL (KB, MAKE ACTION SENTENCE(t))

action = ASK(KB, MAKE ACTION QUERY(t))

**Wumpus World** 

The knowledge based agent is in an environment consisting of rooms connected by passageways. Some rooms contain bottomless pits while others contain goal. One wumpus lives in the cave in one room. Wumpus eats anyone who enters its room but does not move. Player has one arrow that can kill the wumpus.

## **Performance Measure**

- +1000 points for getting gold. -1000 points for falling into a pit or eating a wumpus. -1 for each action taken.
- -10 for using an arrow.

## **Actuators**

Move forward one room. Turn left 90 degrees. Turn right 90 degrees. Shoot the arrow

Climb out (if in starting space)

#### **Sensors**

Stench: A wumpus is in an adjacent room. Breeze: A pit is in an adjacent room. Glitter: Gold is in the player's room Scream: Wumpus is killed. Bump: Player walks into a wall.

## Logic

Syntax – Sentence formatting to make all knowledge sentences well formed.

Semantics – Provide meaning to sentences. It defines truth for every possible world.

**Example:** For the sentence, x + y = 4 is true in the world where x = 2 and y = 2. Model - Substitute for the phrase "possible world." A model fixes the truth or falsehood for every relevant sentence.

Satisfaction: Making a sentence true using an allowed model/possible world.

**Example:** If sentence  $\alpha$  is true in model m, then model m satisfies sentence  $\alpha$ .

## **Entailment**

**Entailment Between Sentences: When one sentence** logically follows from another sentence or set of sentences. It is similar to implies in philosophy.

## Symbol: ⊨

Given two sentences  $\alpha$  and  $\beta$ , then sentence  $\alpha$  entails the sentence  $\beta$  if and only if:

$$\alpha \vDash \beta \Leftrightarrow \forall M(M(\alpha) \subseteq M(\beta))$$

The knowledge base is a set of sentences. The knowledge base is false in models that conflict with the knowledge base.

Model: Fixes the truth value (i.e. true or false) for each

Atomic Sentence: Simplest type of sentence and contains a

Naming Convention: First letter is capitalized followed by

Positional symbols with fixed meaning: True (always true

Syntax: Defines allowable sentences.

proposition symbol.

Semantics: Defines what a sentence means.

single propositional symbol (i.e. variable)

statement that can be either true of false.

position) and False (always false proposition)

lower case letters and subscripts.

Propositional Symbol: Represents a proposition or

Model Checking: Given a knowledge base, KB, and verify it is a model of  $\alpha$ . Hence:

## $M(KB) \subseteq M(\alpha)$

Model checking entails enumerating all possible models to determine whenever KB is true that  $\alpha$  is also true. It only works on finite domains.

Logical Inference: Process of drawing conclusions (i.e. new sentences) through entailment.

## Symbol of Inference: H

Given a knowledge base, KB, and a sentence  $\alpha$ , if an inference algorithm, *i*, inferred  $\alpha$  from *KB* then:

 $KB \vdash_i \alpha$ 

Sound or Truth Preserving Inference Algorithm: Can only derive entailed sentences. Hence it cannot prove any sentence that is wrong.

Example: Model checking is a sound algorithm since it does not work on infinite spaces.

Complete Inference Algorithm: Can derive any entailed sentence. A complete inference algorithm can prove anything that is right.

## **Syntax**

## **Logical Connectives**

Symbols that operate on propositional logic symbols.

- →: Not (Negation)
- V: Or (Disjunction). Individual terms are called disjuncts.
- A: And (Conjunction). Individual terms are called conjuncts.
- ⇒: Imply (Implication)
- ⇔ or ≡: Biconditional. "If and only if"

 $A \Rightarrow B$  is True unless A is true and B is false.  $A \Leftrightarrow B$  is true only if A and B are both true or are both false.

If  $A \Rightarrow B$ , then:

- A is the premise or antecedent
  - B is the conclusion or consequent.

#### **Valid Sentence**

AtomicSentence := True|False|P|O|RSentence := AtomicSentence | SentenceComplexSentence := (Sentence) | [Sentence]

| → Sentence | Sentence V Sentence

| Sentence ∧ Sentence | Sentence  $\Rightarrow$  Sentence | Sentence  $\Leftrightarrow$  Sentence

## **Operator Precedence**

 $\neg$ ,  $\lor$ ,  $\land$ ,  $\Rightarrow$ ,  $\Leftrightarrow$ 

## **Inference Proving**

Checking if  $KB \models \alpha$ 

**Model Checking:** Enumerate all the models and check if all for all possible models where KB is that  $\alpha$  is also true. **Model checking is very similar to a truth table.** 

**Theorem Proving:** Using sentences already in the model, apply rules of inference to construct a proof of the desired sentence without consulting models.

**Literal:** In a complex sentence, a literal is either an atomic sentence (i.e. **positive literal**) or its negation (i.e. **negative literal**).

**Logical Connectives:** Used to construct complex sentences out of atomic sentences.

**Logical Equivalence:** Two sentences  $\alpha$  and  $\beta$  that are true in the same set of models. **Notation:**  $\alpha \equiv \beta$ 

Validity: A sentence that is valid (true) in all models.

Tautology: A valid sentence.

## **Common Logical Equivalences**

Commutative of ∧	$(\alpha \land \beta) \equiv (\beta \land \alpha)$	Commutative of ∨	$(\alpha \vee \beta) \equiv (\beta \vee \alpha)$
Associativity of ∧	$((\alpha \land \beta) \land \gamma) \equiv (\alpha \land (\beta \land \gamma))$	Associativity of ∨	$((\alpha \vee \beta) \vee \gamma) \equiv (\alpha \vee (\beta \vee \gamma))$
<b>Double Negation</b>	$\neg (\neg \alpha) \equiv \alpha$	Contraposition	$(\alpha \Rightarrow \beta) \equiv (\neg \beta \Rightarrow \neg \alpha)$
Implication Elimination	$(\alpha \Rightarrow \beta) \equiv \neg \ \alpha \lor \beta$	Biconditional Elimination	$(\alpha \Leftrightarrow \beta) \equiv ((\alpha \land \beta) \lor (\neg \alpha \land \neg \beta))$
DeMorgan's Law	$\neg (\alpha \land \beta) \equiv (\neg \alpha \lor \neg \beta)$	DeMorgan's Law	$\neg (\alpha \lor \beta) \equiv (\neg \alpha \land \neg \beta)$
Distributivity of ∧ and ∨	$(\alpha \land (\beta \lor \gamma)) \equiv ((\alpha \land \beta) \lor (\alpha \land \gamma))$	Distributivity of ∧ and ∨	$(\alpha \vee (\beta \wedge \gamma)) \equiv ((\alpha \vee \beta) \wedge (\alpha \vee \gamma))$
Modus Ponens	$(\alpha, \alpha \Rightarrow \beta) \equiv \beta$	Modus Tollens	$(\neg \beta, \alpha \Rightarrow \beta) \equiv \neg \alpha$
And Elimination	$(\alpha \land \beta) \Rightarrow \alpha$		

**Satisfiability:** A sentence that can be made true with some model. For a finite environment, satisfiability can be by enumerating all possible models and seeing if any leads to the statement being true. CSP consistency checking is a type of satisfiability problem.

**Validity and Satisfiability:** A sentence is valid if and only if its negation is not satisfiable.

**Reduction ad absurdum/Proof By Reduction/Proof by Contradiction**: Given a logical expression, assume the opposite of the expression and determine if it is satisfiable.

**Example:**  $(\alpha \land \beta)$  is true in the model:  $M = \{\alpha = True, \beta = True\}$ 

**Proof:** A chain of conclusions that leads to the establishing some statement following from the knowledge base.

## **Example**

Consider a situation where four light switches on a control panel. Define a knowledge base for this system with conditions defined in **Part A** and **Part B**.

## **Definition:**

 $S_1$ : Propositional symbol for the first switch and is true if the switch is on and false otherwise.

S<sub>2</sub>: Propositional symbol for the second switch and is true if the switch is on and false otherwise.

 $S_3$ : Propositional symbol for the third switch and is true if the switch is on and false otherwise.

S4: Propositional symbol for the fourth (i.e. last) switch and is true if the switch is on and false otherwise.

Part A: The first and last switches are never both on.

$$\neg (S_1 \land S_4)$$
  
$$\neg S_1 \lor \neg S_4$$

Part B: At least one switch must be on.

$$S_1 \vee S_2 \vee S_3 \vee S_4$$

## **Python Review**

## **Python Basics**

**Command Line Call to Run Python:** python filename.py **Python File Extension:** \*.py

**Command to Print to Console:** print "Hello World!" Printing without Inserting a Newline: Use "," (Comma)

**Command to Get Last Result:** (Underscore) Example: >>> 2/3 + 7.9 >>> print \_ + 1 # prints 8.9

**Valid Python Operators:** +, \*, -, /, \*=, /=, -=, +=, %, ==, != // (Integer Division), \*\* (Power)

Math Functions:

math.exp( value ): e^value random.randint(n,m): Integer  $n \le x \le m$ random.random(): Float  $0 \le x < 1$ 

**Invalid Operators:** 

++, --

**Minimum and Maximum Value:** 

inf -inf

Conditionals: if( expr ): # Do something

elif( expr ): # Do something

# Do something

**Boolean Arithmetic:** 

print "Hello World",

is, and, or, not

**Boolean Literals:** True, False

**Check Membership in List:** 

File IO:

f = open("filename.txt", "w") line = f.readline() f.close()

# Iterate over a file line by line for line in open("my\_file.txt"):

#Do something

**Formatted Printing:** 

Use the % symbol similar to C/C++

print "%3d %0.2f" % (10, .9799)

# Prints "10 0.98"

**Python String Manipulation** 

**Python String Implementation** Immutable list of characters.

**String Concatenation:** 

+ (plus sign)

**Converting from a String:** 

int("38")

float("46.456")

Converting to a String:

str(7)

repr(32.9)

**Substring Manipulation** 

print a[3:8] # Prints "lo Wo"

Use [] like a list with the first character index 0 a = "Hello World" print a[4] # Prints "o" print a[:5] # Prints "Hello" print a[6:] # Prints "World"

Use the *in* operator:

**Checking for Substring:** 

if( "hello" in "hello world"): print "It's in there."

**Get Index of Substring:** 

x = "hello world".index("llo") print x # Prints "2"

## **Element Containers**

List (Array) Basics:

Able to hold data of different types in the same list including other lists. Uses [] x =[ 5, 4, "hello", "world" ] print x[1] # Prints "4" print x[1:] # Prints "[4, "hello", "world"]" print x[0:2] # Prints "[4, 5]" y = [3, 2], [1, 0]

**Nested (Two-Dimensional) Lists:** 

y = [[3, 2], [1, 0]]print y[1][0] # Prints 1 **Concatenating Lists:** 

y = [4, 5]

print z # Prints "[1, 2, 3, 4, 5]"

x = [1, 2, 3]z = x + v

List Length: Use len()

x = [1, 2, 5, 10]print len(x) # Prints "4"

**Extracting List Properties:** 

max( list ) - Gets Maximum Value in List min( list ) - Gets Maximum Value in List

Tuple:

Immutable list. Created used () parenthesis.

**Accessing Tuple Elements:** 

c = (4, 5)print c[1] # Prints "5" a, b = c # a = 4 and b = 5

**Creating a Tuple:** 

print y[1][0] # Prints 1

a = (1, 2, 3) # Tuple of size 3 b = (x, y) # Tuple made of two variables c = "Hello", "World" # Tuple of size 2 d = () # Empty Tuple e = "yo", # Tuple of size 1 f = ("yo", ) # Equal to e

g = (d, ) # Tuple of empty tuple ((), )

Sets:

Unordered collection of unique elements. x = set([3, 6, 9, 2])my\_set = set("goodness")

print my set # Prints ["g", "o", "d", "n",

"e", "s"] with no duplicates

Frozenset:

An immutable set. x = frozenset([4, 5, 6])

**Set Operations:** 

| Union, & Intersection, - Difference, ^ Symmetric Difference (XOR)

Dictionary:

Associative Array (i.e. hash table). Uses {} curly brackets. person ={

"name": "bob". "age": "27", "sex": "Male"

print person["name"] # prints "Bob" **Deleting from a Dictionary:** del person["name"]

**Dictionary Membership Test:** 

Use the keyword "in" if( "name" in person ): print person["name"] # Prints "bob"

**Accessing Dictionary Elements:** 

person.kevs() # Gets all dict keys person.values() # Gets all dict values person.len() # Gets all dict length

**Looping and Iteration** 

While Loop:

while( expr ): # Do something For Loop:

for x in [2, 4, 5, 6, 9]: print x for y in range(1, 10): print y # Only prints 9 lines range:

Iterable object in Python. range(0, 10) - Creates list of 0 to 9 in steps of 1 range(10) – Starting 0 not needed. Same as range(0,10) range(0, 5, 2) - Starts 0 and steps by 2 until 5 range(7, 2, -1) - Starts at 7 and decrements by 1 until 3

range vs. xrange:

range creates an array that Python iterates over. This is memory inefficient. xrange acts like a real for loop without the memory overhead of range.

**Iterable Objects in Python:** 

set, frozenset List, Tuple Dictionary key File (open("filename") String (letter by letter) Generator

## **Functions**

## Creating a Function: Keyword: def def my\_func(params): # Do something Keyword to Return: return **Supports Recursion: Yes Taking an Arbitrary Number of Input Variables Keyword: \*args** def my\_function(\*args):

```
Scope:
Default scope in python is local.
def print_i():
 i = 4
 print i
```

print\_i() # Prints "4"

print I # Prints "5"

```
Keyword to Add to Global Scope: global
def assign_i():
 global i
 i = 3
```

```
Storing a Function in a Variable:
def print i()
 i = 4
 print i
a = print_i
a() # Prints "4"
```

```
Anonymous Function:
Keyword: lambda
g = lambda x: x**3
print g(10) # Prints "1000"
h = lambda y, z: z + 2*y
print h(2, 3) # Prints "8"
def make_adder(n):
  return lambda z: z+n
f = make adder(2)
print f(3) # Prints "5"
print f(6) # Prints "8"
g = make adder(4)
print f(3) # Prints "7"
print f(6) # Prints "10"
LAMBDA NEVER HAS A RETURN
```

## Generator

Uses the yield construct and the object method next.

Allows you to get a sequence of objects in a dedicated routine.

```
def countdown(n):
     while(n > 0):
           yield n
           n -= 1
```

# Creates the function call as object but does NOT run it yet x = countdown(3)

```
print x.next() # First runs "countdown(3)" then prints "3"
print x.next () # Prints "2"
print x.next () # Prints "1"
```

## Coroutine

Uses the yield construct and the object method send and next.

Allows you to pass a sequence of values one at a time to a function (e.g. log file printer)

```
def print_matches(text):
     print "Trying to find text: " + text
      while(True):
            line = (yield)
            if(text in line):
                 print line
```

# Creates the function call as object but does NOT run it yet x = print\_matches("hello") x.next() # Runs to first yield. print x.send("lalalala") # Prints nothing print x. send ("hello world") # Prints "hello world"

## Classes

```
class ClassName(inherited class1, inherited closs2):
     # Class variables
     class name = "Class Name"
     # Constructor
     def init (self):
          self.attribute1 = 1
          self.attribute2 = [3, 4]
          self.length value = 1
     # Called without parenthesis for methods
     @property
     def length(self)
          return self.length_value
     # Called by ClassName.static_method(arg)
     @staticmethod
     def print_class_name()
          print class_name
             Calling Supercass Methods
Option #1
super(SuperClassName, self).methodName(variables)
Option #2
 _ClassName__method_name(variables)
```

**Invoking a Class Constructor:** Use the class name followed by two parenthesis. Example for class "Stack":

Example: my\_stack = Stack()

## **Class Special Methods:**

name Always preceded and proceeded by two underscores.

@property: Class methods that do not require parenthesis when called. Typically return an object or primitive.

Static Method: @staticmethod Called using the class name not an object name.

## Example:

ClassName.static method()

Inheritance and Classes: Python class can inherit multiple classes.

#### **Class and Inheritance Functions:**

- type(variable name): Returns a formatted string of object's class name.
- isinstance(variable\_name, ClassName): Returns True if variable is of type ClassName, False otherwise.

Example: isinstance(my\_stack, Stack) returns True.

issubclass(SubclassName, ClassName): Returns true if SubclassName is a subclass of ClassName.

Example: issubclass(Stack, object) returns True.

#### **Abstract Classes**

Requires the import: from abc import ABCMeta, abstractmethod, abstractproperty

# Required first line for abstract class \_\_meta\_class\_\_ = ABCMeta

@abstractmethod def my\_method(args):

pass

@abstractproperty def my\_method(args): pass

**Abstract classes do NOT inherit** ABCMeta.

## **Exceptions**

## Format for an Exception try: pass except ErrorTypeName as error object: # Catches only error of type ErrorName pass except: # Catches all exceptions finally: # Always run pass

```
Creating Your Own Exception
```

```
class MyException(exception):
     def init (self, errno, msg):
           self.args = (errno, msg)
           self.errno = errno
           self.msg = msg
```

class MyException2(exception):

pass

**Throwing an Exception** Use the raise keyword

raise MyException(404, "Access Forbidden")

## Modules, Importing, and the sys Toolset

## Importing From a Module with Normal Namespace Syntax: import filename

Filename is the python filename without the file extension (.py). When importing in this fashion, it uses the file name as the namespace for the functions/classes in that file.

Example: Python file div.py has a function called divide that divides to integers.

import div

print div.divide(4,2)

Importing From a Module with a New Namespace

Syntax: import filename as namespace Use a custom namespace name for

Example: Python file div.py has a function called divide that divides to integers. New namespace is named "foo"

import div as foo print foo.divide(4,2)

## sys - Common System Functions

import sys

**Command Line Arguments:** 

svs.argv

**Quitting Python:** 

sys.exit(0)

Printing to the Console (Substitute for print): sys.stdout("Hello World")

**Getting User Input from the Console:** input = sys.stdin.readline()

## **Function to Add Set of Integers** Passed by Command Line

import sys

```
def sum command line args()
     input_args = sys.argv
     sum = 0
     try:
          # Skip element one since module name
          for i in range(1, len(input_args)):
               sum += int(input_args[i])
          print "Input argument not an integer"
```

sys.exit(0) # Print the sum to the console.

**Unit Testing** 

print "The sum of the input arguments is: ", print sum\_command\_line\_args()

## **Documentation String**

Documentation String: First statement of a module, class, or function.

**Extracting Documentation String for a Function, Class, or Module:** 

Use the method \_\_doc\_\_

**Example:** A function exists called fact. To print its documentation string, call:

print fact. doc

**Accessing Documentation String Outside a Python Program** 

Example: Function fact exists in module MyModule.py

Interpretative Mode:

import(MyModule) help(MyModule.fact)

**Command Line:** 

pydoc MyModule.fact

## Included in **Documentation String**.

**Module Name: doctest** 

Unit Test Function Name: testmod()

Format:

>>> function\_name(args)

result

Example:

def multiply(a, b):

>>> multiply (0, 1) >>> multiply (2, 1) >>> multiply (3, -1)

-3

return a \* b

doctest.testmod()

## **Setting Up doctest in Supporting Modules**

```
# Check to see if this module is main
if( __name__ == 'main'):
     # Import doctest module then run testmod()
     import doctest
```

## **Benefits of Python**

Good string and list processing functionality which minimizes awkward additional	Scripted/interpreted coding available for testing
coding.	
Higher order function support (e.g. functions can take other functions as arguments)	Syntax is comparable to other languages.
Good set of built-in libraries.	Wide range of free libraries and projects to build off.
People outside AI use it so others can appreciate your code.	

## **Midterm Special Notes**

## Python:

- 1. Do not forget colons in Python code including after function definitions, for, while, and if statements.
- 2. Do not forget to call imports in Python code for modules such as math, random, and sys.
- 3. Printing a formatted string of numbers can be written:

print "%3d %0.2f" % (10, .9799) # Prints 10 with a preceding space and 0.98

4. It is possible to have Tuples of size 0 by doing:

x = ()

5. It is possible to have Tuples of size 1 by doing:

x = "Hello World", x = ("Hello World".)

6. For an abstract class, you need the line:

\_\_metaclass\_\_ = ABCMeta

## **General Agents:**

- 7. Components Needs to Pass the Turing Test:
  - a. Natural Language Processing
  - b. Knowledge Representation (i.e. storage paradigm)
  - c. Automated Reasoning
  - d. Machine Learning
- 8. Cognitive Science: Brings together computer models from AI and experimental techniques from psychology to construct precise and testable theories of

## the human mind.

- 9. Agent Function Maps percept sequence to agent action.
- 10. Simple Reflex Agent Select actions based off the current percept only. Often defined by condition-action rules (i.e. productions)
- 11. Goal Based Agents A goal is a binary condition (i.e. either met or not met). A goal based agent tries to reach a target goal. Search and planning agents may be goal based agents.
- 12. Problem solving agents deal with atomic environments (i.e. the environment is treated as a single whole and is indivisible).

#### Search:

- 13. In Recursive Best First Search code, remember to do the Goal\_Test at the beginning of the function and to check if the successors list is empty after creating it.
- 14. Effective Branch Factor:  $b^*$  Equivalent branch factor if the search tree was modelled as a balanced tree (i.e. where the number of children for each node is equivalent for all nodes).

#### **Constraint Satisfaction:**

- 15. Node Consistent Variable Any variable where every value in the variable's domain satisfies all of its unary constraints in a CSP network.
- 16. In AC-3, only excluding the current paired variable are expanded.
- 17. Local Consistency: Given a constraint graph, enforcing consistency (i.e. ensuring variables satisfy constraints) locally in each part of the graph leads to invalid values being eliminated throughout the graph.
- 18. Path Consistency A two variable set  $(X_i, X_j)$  are path consistent with respect to a third variable  $X_m$  if for every assignment of values to  $X_i$  and  $X_j$  consistent with the constraint  $\{X_i, X_j\}$ , there is a valid assignment to  $X_m$  that satisfies the constraints  $\{X_i, X_m\}$  and  $\{X_m, X_j\}$ .
- 19. Interleaving Search and Inference AC-3 can be used to infer reductions in the search domain both before and during search.
- 20. **Forward Checking** One way to implement "Inference" in Backtracking algorithm. Whenever a variable is assigned, establish arc consistency for it on all unassigned variables. If arc consistency checking was done in preprocessing, forward checking adds no value.
- 21. **Minimum Remaining Value (MRV),** Fail First, Most Constrained Variable Heuristic: Select the variable to assign next that has the smallest inferred domain (i.e. least remaining legal values).

## **Logic and Logic Agents**

- 22. Declarative Programming: Provide information to the agent on information it needs to know and it figures out how to achieve the solution. De Procedural approach: Teach the agent how to do certain actions and it uses that information to figure out a solution to what you intend for it to do.
- 23. Background Knowledge Initial knowledge in the knowledge base.
- 24. Inference Deriving new sentences from existing sentences.
- 25. Logical Connectives: Used to construct complex sentences out of atomic sentences.
- 26. Theorem Proving: Using sentences already in the model, apply rules of inference to construct a proof of the desired sentence without consulting models.
- 27. Entailment Between Sentences: When one sentence logically follows from another sentence or set of sentences. It is similar to implies in philosophy.
- 28. Logical Inference: Process of drawing conclusions (i.e. new sentences) through entailment. Symbol of Inference:  $\vdash$  Given a knowledge base, KB, and a sentence  $\alpha$ , if an inference algorithm, i, inferred  $\alpha$  from KB then:  $KB \vdash_i \alpha$
- 29. Sound or Truth Preserving Inference Algorithm: Can only **Cerive** entailed sentences. Hence it cannot prove any sentence that is wrong. Example: Model checking is a sound algorithm since it does not work on infinite spaces.
- 30. Complete Inference Algorithm: Can derive any entailed sentence. A complete inference algorithm can prove anything that is right.
- 31. Literal: In a complex sentence, a literal is either an atomic sentence (i.e. positive literal) or its negation (i.e. negative literal).
- 32. Proof: A chain of conclusions that leads to the establishing some statement following from the knowledge base.

## **General Rule:**

$$\sum_{i=m}^{n-1} a^i = \frac{a^m - a^n}{1 - a}$$

## Inferences, Proofs, and Resolution

**Three Key Notions in Propositional Logic** 

**Logical Equivalence:**  $a \equiv b \Leftrightarrow (b + a \land a + b)$  **Validity** – A statement that is true in all models. **Satisfiability** – A statement where at least one model can make the statement true.

Propositional Proof – A series of steps where each statement is either from the knowledge base, a valid propositional statement, or a statement follows previous statements via some rule of propositional inference.

## Framing a Proof as a Search Problem

A propositional logic proof can be		
treated as search problem and		
existing search algorithms can be		
used to find a valid proof		

Initial State: The initial knowledge base

Actions: Set of all inference rules applied to all the sentences that match the first half of an inference rule

Results: Add the bottom half of all applicable inference rules (see actions) to the knowledge base. Goal: A knowledge base that contains the statement that is trying to be proven.

Monotonicity – Property of some knowledge bases where the set of entailed sentences only increases as sentences are added to the knowledge base.

Nonmonotonic logics – Common in the study of human AI. Set of entailed sentences may decrease.

**Literal** – A propositional variable or its negation. Example: X or  $\overline{X}$ 

## Resolution

Resolution is a sound and valid inference rule.

Requires two disjunctive clauses. If the clauses contain complimentary variables, the two clauses are combined with complementary literals excluded.

Example of Resolution:

 $\frac{A \lor B \lor C, \ \bar{C} \lor D \lor E}{A \lor B \lor D \lor E}$ 

**Resolvent:** Clause produced by resolution. (i.e. bottom line of inference specifically:  $A \lor B \lor D \lor E$ )

**Complementary Literals** – One literal is the negation of the other literal.

**Unit Resolution:** Right hand clause contains a single literal whose complement is in the left clause.

Clause Set Notation:  $\{L_1, L_2, \dots, L_m\}$  is the same as a disjunction of those literals.

**Conjunctive Normal Form (CNF):** Conjunction (ANDs) of disjunctions (ORs).

Resolution works best on propositional knowledge bases in CNF.

**Using CNF with Resolution** 

Goal: Prove  $KB \Rightarrow \alpha$ 

Step #1: Use implication elimination

 $\overline{KB} \vee \alpha$ 

Step#2: Negate the goal  $KB \wedge \overline{\alpha}$ 

KD A U

Step #3: Convert to CNF

**Step #4:** Prove the statement is not satisfiable (i.e. the empty clause is found through resolution).

**Truth Table Approach to Convert to CNF** 

- Enumerate all models.
- For any model that is false, take a disjunction of the literals negation.

## Example:

Α	В	Result
True	True	False
True	False	True
False	True	False
False	False	True

Result  $\Leftrightarrow$   $(\bar{A} \lor \bar{B}) \land (A \lor \bar{B})$ 

#### **Key Inference Steps:**

- Double negation
- DeMorgan's Theorem
- **Biconditional Elimination**

$$(A \Leftrightarrow B) \Leftrightarrow ((A \Rightarrow B) \land (B \Rightarrow A))$$

Inference Algorithm Approach to Convert to CNF

- Distributivity
- Implication Elimination

 $(A \Rightarrow B) \Leftrightarrow (\bar{A} \lor B)$ 

Example:

$$(A \wedge B) \vee (\bar{A} \wedge \bar{B}) \vee (A \wedge \bar{B})$$

$$\neg \neg ((A \wedge B) \vee (\bar{A} \wedge \bar{B}) \vee (A \wedge \bar{B}))$$

$$\neg ((\bar{A} \vee \bar{B}) \wedge (A \vee B) \wedge (\bar{A} \vee B))$$

$$\neg ((\bar{A} \vee \bar{B}) \wedge (A \vee B) \wedge (\bar{A} \vee B))$$

$$\neg (((\bar{A} \wedge B) \vee (A \wedge \bar{B})) \wedge (\bar{A} \vee B))$$

$$\neg (((\bar{A} \wedge B) \vee (A \wedge \bar{B})) \wedge (\bar{A} \vee B))$$

$$\neg (((\bar{A} \wedge B) \vee (A \wedge \bar{B})) \wedge (\bar{A} \vee B))$$

$$\neg (\bar{A} \wedge B)$$

$$(A \vee \bar{B})$$

**Resolution Closure:** Set of all statements that derive from the knowledge base through resolution.

**Resolution Refutation Stops in Two Cases:** 

- 1. Empty clause found
- 2. No new clauses are possible in the resolution closure.

**Refutation** – Empty clause found when performing resolution.

**Definite Clause** – Disjunctive (OR) clause with **exactly one positive literal**.

Example:  $(L \lor \bar{B} \lor \bar{C})$ 

**Notation for Definite Clause:** 

Positive Literal: -Negative Literals

Example:

L: -B, C

**ASCII Notation:** 

 $(B \wedge C) \Rightarrow L$ 

**Head:** Positive literal in the clause (e.g. L)

Tail: Negative literals if any (e.g. B, C)

Rule: Entire clause.

Horn clause: Disjunctive clause with at most one positive literal.

Example Horn Clause:  $\bar{B}$ 

Alternative Notation: -B

**Propositional Logic Notation:**  $B \Rightarrow False$ 

**Horn clause:** Collection of Horn clauses. A type of **logic program**.

Importance of Horn Clauses and Program: Knowledge bases that are Horn programs can decide if a clause is entailed in linear time and space.

**Goal Clause** – A horn clause with no positive literals. (Example:-b or  $b \Rightarrow False$ )

**Goal:** See if  $KB \Rightarrow B$ 

**Backward Chaining:** If KB is a Horn program, look for a clause where B is the head. Check for a rule where the head is true. If one is found, then continue search.

Forward Chaining: If KB is a Horn program, start from the facts and search forward until no possible change to KB or the goal is found.

$R_1$	A (Fact)
$R_2$	C (Fact)
$R_3$	$\bar{A} \vee B$ (i.e. $A \Rightarrow B$ )
$R_4$	$\bar{B} \vee \bar{C} \vee D$ (i.e. $(B \wedge C) \Rightarrow D$

**Backward Chaining** 

Finds  $R_4$  then  $R_3$  then  $R_1$  then  $R_2$ 

Finds  $R_1$  then  $R_2$  then  $R_3$  then  $R_4$ 

Closed World Assumption (CWA) – Facts that are not known are assumed to be false.

This favors minimal models.

Open World Assumption (OWA) – Facts that are not known are assumed to be true. This favors maximal models.

#### **DPLL** – Resolution Finding Algorithm

## Three Optimizations Over the Basic Resolution Algorithm:

- Early Termination: If all clauses are satisfied (have at least one positive literal) or any clause is false, terminate the algorithm.
- Pure Symbol Heuristic: A pure symbol is any symbol that has the same sign in all clauses. Pure symbols are set to true if they exist.
- Unit Clause: A unit clause contains only a single literal. The variable in the unit clause is set to true to satisfy the clause.

## def DPLL\_Satisfiable(s): # Returns True or False

clauses = set of clauses from CNF representation of s symbols = list of symbols in s

return DPLL(clauses, symbols, {})

def DPLL(clauses, symbols, model):

# Check Early Termination

if every clause is true in model:

return True

elif some clause is false in model:

return False

## # Check Pure Symbol Heuristic

P, value = FIND\_PURE\_SYMBOL(clauses, symbol, model)

if P is not None:

return DPLL(clauses, symbols - P, model U {P=value})

#### # Check Unit Clause Heuristic

P, value = FIND\_UNIT\_CLAUSE(clauses, model)

if P is not None:

return DPLL(clauses, symbols - P, model U {P=value})

## # Select first symbol and check both true and false

P = FIRST(symbols)

rest = **REST**(symbols)

return DPLL(clauses, rest, model U {P = True})

or DPLL(clauses, rest, model U {P = False})

## **Prolog**

#### a. - Fact A in Prolog.

b:- a – Horn Clause  $(\neg a \lor b)$  or  $a \Rightarrow b$ . Since a is true, then b is also true. c:- b – Horn Clause  $(\neg b \lor c)$  or  $b \Rightarrow c$ . Since b is true, then so is c

d :- a, b – Horn Clause  $(\neg a \lor \neg b \lor d)$ . Since a and b are both true, so is d

This is the same as:

a  $a \Rightarrow b$ 

 $b \Rightarrow c$  $(a \land b) \Rightarrow d$ 

Prolog supports non-Horn clauses like:

e:-not(a) and f:-false

## Question #1 from Practice Final

$$\bigwedge_{i=1}^{6} \left( \overline{x_i} \vee \bigvee_{1 \leq j \leq 6, i \neq j} x_j \right) \wedge \bigwedge_{i=1}^{6} \left( \bigwedge_{j=i+1}^{6} \left( \overline{x_i} \vee \overline{x_j} \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} x_k \right) \right) \wedge \bigwedge_{i=1}^{6} \left( \bigwedge_{j=i+1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigvee_{1 \leq k \leq 6, k \neq i, k \neq i} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i} \overline{x_k} \right) \wedge \bigvee_{1 \leq k \leq 6, k \neq i} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i} \overline{x_k} \right) \wedge \bigvee_{1 \leq k \leq 6, k \neq i} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i}$$

## **First Order Logic**

Logic based agents tell the knowledge base about their percepts.

First Order Logic - Logic system Predicate: Takes inputs and Variables: Range over sets. Constants: Fixed values from a set function symbols and return a where variable domains is outputs True/False constant greater than solely "True" and Usual notation: x, y, z Usual notation: a, b, c "False" Usual notation: P, Q, R Usual notation: f, g, h, Term: A variable, a constant, or Atomic Formula: Predicate Universal Quantifier: Symbol ∀ **Existential Quantifier:** Symbol ∃ where each of the predicate built up from these using Formula: An atomic formula or a function symbols and slots is filled by a term. composite of simpler formula.  $\forall x F_1$  - For all x,  $F_1$  is true.  $\exists y F_2$  - For some y,  $F_2$  is true. composition. Example: IsPrime(X \* X + 3)

First Order Logic Semantics

Universe (M) - A set M over which all variables range over.

Constant  $(c^M) - A$  value in the universe M

Constant  $(c^M) - A$  value defined as:  $f^M: M \times M \times ... \times M \to M$ Predicate  $(P^M)$ : Returns True or false and is defined as:  $P^M: M \times M \times ... \times M \to \{T/F\}$ Language: Set of all constants in the universe and all function symbols.

Structure/Model (M): Bound Variable: A variable in a Example: Unbound Variable: A variable in Variable/Object Assignment (v): A Combination of the universe. first order function that is within  $(\exists x)F(x,y)$ a first order function that has no map from unbound variables to constants, functions, and the scope of an existential or **Unbound Variable:** y quantifier. elements in the universe (M) predicates. universal quantifier. **Bound Variable:** x

Logic Equations with Quantifiers:  $(\forall x)(\neg P) \Rightarrow \neg(\exists x)P$   $A(x) \Rightarrow (\forall y)(A(y))$   $A(x) \Rightarrow (\forall y)($ 

Interacting with a First Order Knowledge Base

TELL(KB, King(John)) – Tells the knowledge base if return true.

**TELL**(KB, Person(Richard)) – Tells the knowledge base that Richard is a person.

Ask(KB, King(John)) – Predicate that asks the knowledge base if John is a King. Would return true.

Ask(KB, King(Zayd)) – Returns false since Zayd is not a king.

This command is referred to as query or goal.

AskVars(KB, Person(x)) – Asks questions that returns a constant.

Function: Take variables with

Query response is known as a **binding list** or **substitution**. Example return is {x/Richard}

**Example First Order Knowledge Bases** 

1. Any relational database

2. Basic set theory

- No function symbols
- No function symbols
- = operator checks for equality
- Constant is the empty set Ø

**Theorem Proving in First Order Logic** 

## **Procedure**

1. Convert each formula in  $KB \cup \{\neg \alpha\}$  into **prenex normal** form. Prenex normal form is:

 $\forall x \forall y \exists z \ F(x,y,z) \land G(x,y) \Rightarrow H(x,y,z)$ 

- 2. Skolemize the equation to remove any existential quantifiers. Each existentially quantified variable gets its own function. (i.e.  $f_1(x, y), f_2(x, y), f_3(x, y)$ , etc.)
- 3. If all variables are bound and only universal quantifiers, the quantifers can be dropped and all variables are free.
- 4. Convert the open formula to CNF and use resolution to prove refutation

**Skolemization** – Process of removing existential quantifiers by making the existentially quantifier variables functions of universally quantified variables.

Examples

 $\exists x \exists y \ F(x) \Rightarrow G(y)$ skolemizes to  $F(a) \Rightarrow G(b)$ 

 $\forall x \forall y \exists z \ F(x, y, z) \Rightarrow G(x, y, z)$ skolemizes to  $\forall x \forall y \ F(x, y, f(x, y)) \Rightarrow G(x, y, f(x, y))$ 

## **Additional Notes**

If there are only existential quantifiers, the variables are turned into constants and existential quantifiers dropped.

To perform refutation, a substitution list may be required to ensure the terms in the predicate match. This can be checked using the **unification algorithm**.

Model checking is possible to prove entailment in first order knowledge bases. However, the time complexity is just as bad or worse than it is for propositional logic.

\* = operator for checking two values are the same

First Order Logic Database Commands

PDDL – Planning Domain Definition Language

Successor of Strips language.

Planning – Application of first order logic. Develop a sequence of actions to achieve a goal while at each step in time satisfying all constraints.

#### **Necessary Functions for Unify Function**

- is var(z) Checks if z is a variable.
- is\_term(z) Checks if parameter z is a term.
- args(z) Extracts a list of arguments in z
   args((z\*z)+35) Returns (z\*z, 35)
- op(z) Gets the outermost function symbol in z
   op((z\*z)+35) Returns "+"
- is\_list(z) Checks if parameter z is a list.
- head(z) Returns first element in list z
- tail(z) Returns all elements after the first element in z.

#### **Necessary Functions for Unify Var Function**

- occur\_ck(var, z) Checks if z is function containing var
  - o occur\_ck(z, (z\*z)+35) Returns **True**
- o occur\_ck(y, (z\*z)+35) Returns False
- append(new\_sub, sub\_list) Appends the new substitution new\_sub to the sub\_list.

```
Unify(x, y, S):
  #x-a variable, constant, term, or list
  #y-a variable, constant, term, or list
  #S-substitution so far
  # returns a Substitution list or "None"
  # Check for previous failure
  if(S == None):
     return False
  # If with substitution the two parameters are the same
  # then return the substitution.
  if( x(S) == y(S)):
     return S
  # If x or v are variables, try to create a new substitution
  if(is_var(x)):
     return Unify_Var(x, y, S)
  elif(is_var(y)):
     return Unify_Var(y, x, S)
  elif( is_term(x) and is_term(y) ):
     return Unify(args(x), args(y), Unify(op(x), op(y), S) )
  elif( is list(x) and is list(y) ):
     return Unify( tail(x), tail(y), Unify( head(x), head(y), S) )
  else:
     return None
```

```
Unify_Var(var, y, S):
  # var - A variable
  #y-a variable, constant, term, or list
  #S-substitution so far
  # returns a Substitution list or "None"
  # Check if substitution exists for var (i.e. sub_val1)
  if( var |-> sub_val1 ) in S:
     return Unify( sub_val1, y, S)
  # Check if substitution exists for y (i.e. sub_val2)
  elif( y |-> sub_val2 ) in S:
     return Unify( var, sub_val2, S)
  # Check if y is a function f(var)
  elif( occur_ck(var, y) ):
     return None
  else:
     return append( var |-> y, S)
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```

#### **Unification Examples**

```
Step #1: Unify("f(z)", "g(w)", {})
                                                                                       Step #1: Unify("[g(v), f(g(z))]", "[g(f(w)), f(w)]", {}) # Remove the head of the lists.
Step #2: Unify("z", "w", Unify("f", "g", {})) # Remove operator
                                                                                       Step #2: Unify("[f(g(z))]", "[f(w)]", Unify( "g(v)", "g(f(w))", {} )) # Remove outermost
                                                                                       function symbols g.
Step #3: Returns False # Unification terminated since it was not possible to
unify f and g since they are different operators.
                                                                                       Step #3: Unify("[f(g(z))]", "[f(w)]", Unify( "v", "f(w)", Unify("g", "g", {} )) # No unification
                                                                                       required since function operators are identical
                                                                                       Step #4: Unify("[f(g(z))]", "[f(w)]", Unify("v", "f(w)", {}) # Unify on variable v
                                                                                       Step #5: Unify("[f(g(z))]", "[f(w)]", Unify_Var( "v", "f(w)", {} ) # Append to substitution list
                                                                                       Step #6: Unify("[f(g(z))]", "[f(w)]", { v |-> f(w) }) # Extract the first item in each list.
                                                                                      Step #7: Unify("[]", "[]", Unify( "f(g(z))", "f(w)", { v |-> f(w) } ) ) # Extract function symbol
                                                                                      f on the two functions
                                                                                       Step #8: Unify("[]", "[]", Unify( "g(z)", "w", Unify( "f", "f", { v | -> f(w) } ) ) ) # Unify on
                                                                                      identical function symbols f
                                                                                       Step #9: Unify("[]", "[]", Unify( "g(z)", "w", {v |-> f(w) })) # Perform Unify var on
                                                                                       Step #10: Unify("[]", "[]", Unify_Var( "w", "g(z)", { v |-> f(w) } ) ) # Append substitution
                                                                                      list for variable w
                                                                                      Step #11: Unify("[]", "[]", "w", "g(z)", { v |-> f(w), w|-> g(z) } ) # Identical unification lists
                                                                                       so no step here
                                                                                       Step #12: { v |-> f(w), w |-> g(z) } # Final Substitution
```

## **Planning**

PDDL – Planning Domain Definition Language			
Heavily influenced by earlier planning languages including	Fluent – Facts that may change from situation to situation.	Ground Fluent – Fluent contain no variable (i.e. only constants). They are	State – Conjunction of fluents that are ground.
STRIPS and ADL.	Situation.	functionless atoms.	States cannot contain negative atoms.
	Unique Names Assumption – Any objects that have	Illegal Fluents in a State Description	
Closed World Assumption –	different names are assumed to be different.	1. Fluents containing variables.	Fluents are a conjunction so fluent order
Fluents not in the knowledge		Example: $At(x, y)$	does not matter.
base are false. (Used in PDDL)	Different names refer to different entities in the	2. Fluents containing negations.	uoes not matter.
	world.	Example: Poor	

## **Actions in Planning**

**Actions** need to clearly define what aspect of the state changes and what stays the same.

atomic domains.

Problem Solving Agent – Goal based agent that is focused on solving problems with

Frame Problem: In classical planning, most aspects of the state remain the same after an action. It can be prohibitive to detail the countless aspect of a state that stayed the same after an action Solution to the Frame Problem in PDDL: PDDL only enumerates the aspects of the state that change as a result of an action. Any unmentioned aspects are assumed not to change.

## **PDDL Action Schema**

Three Components in PDDL Action Schema	Action Name and Input Variables Name of the action performed	Preconditions Aspects of the state that must be true before an action can be	Effects Action results. Changes in	Complete Example  Action(Fly(p, from, to),
1. Action Name and Input Variables	and any input variables.	performed. Cannot contain negated atoms.	state.	<b>Precond</b> : $At(p, from) \land Plane(p)$
2. Precondition(s) if any	Example: Fly(p, from, to) Action Name: Fly Variables: p, from, to	Example: $At(p, from) \land Plane(p)$	Example: $\neg At(p, From) \land At(p, To)$	$\land$ Airport(from) $\land$ Airport(to)
3. Effect(s)	p, y, r e.i.s, ce	∧ Airport(from) ∧ Airport(to)		$\pmb{\textit{Effect}} : \neg \textit{At}(p, From) \land \textit{At}(p, To))$
			If an action h	has $v$ variables and the variable have a

**Applicable Action** – An action a is applicable in state s if all of action a's preconditions are satisfied in state s.

In any given state, multiple instances of a given action could be applicable. Example: plane  $P_1$  could fly from SFO to LAX or from SFO to JFK.

If an action has v variables and the variable have a maximum domain of size d, then it takes  $\boldsymbol{O}(\boldsymbol{d}^v)$  to find all applicable ground actions worst case.

**Result of an Action** – Conjunction of fluents.

Delete List – Negative literals in the result of an action. These negative literals correspond to fluents deleted from the state.

Add List – Positive literals in the result of an action. These positive literals correspond to fluents added to the state.

Note: Actions do not refer to time. Precondition refers to time t and results refer to time t + 1.

**Planning Domain** – Set of action schemas.

Initial State – Conjunction of ground atoms. Hence, every slot in the fluent must be filled.

**Goal** – Conjunction of Literals. **Goals can** have variables, which are treated as existentially quantified.

Goal Example:  $At(p,SFO) \land Plane(p)$ In this case, p could be any plane. **Solution** – Sequence of actions from the initial state to a **state that ENTAILS the goal**.

Planning Agents – Goal based agents that work on factored domains.

**Inequality Condition** – Used to prevent illegal conditions in actions where two input variables have the same value.

**Example Planning Algorithms** 

**PlanSat** – Given a planning problem, it determines whether a plan exists that solves the problem.

**BoundedPlanSat** – Given a planning problem, it determines whether a plan exists that solves the problem in **k steps or less**.

Both algorithms are **PSPACE** but **NP-Hard** (hard as any other problem in NP).

For problems without negative preconditions, PlanSet is polynomial time (P).

## **Planning as a Search Problem**

**Forward State Space Search** – Start from the initial condition and search towards the solution.

**Backward (Regression) Relevant States Search** – Start from the goal and try to search backwards until a state **IMPLIED** by the start state is found.

## **Negatives of Forward State Space Search**

Referred to as relevant-states search since only states relevant to the goal are explored. At each step, there may be a set of relevant states (not just a single state).

Negatives of Backward Relevant States Search

Prone to search irrelevant states. Example: Planning problem trying to go from Buy(ISBN) and OWN(ISBN). Would involve searching many irrelevant states.

Requires domain-independent heuristics since planning problems can have large state spaces.

Partially uninstantiated actions and states – Since a goal will not always detail a complete state, negative relevant states search often involves handling only partially instantiated actions and states. It must also handle ground states.

**Heuristics in Planning** – When trying to come up with a plan through search, heuristics may be helpful.

**Example Heuristic Type** – Come up with plans to relaxed problems.

#### **Planning Problem and Search**

Nodes – States in the state space.

Edges – Actions in the planning domain (i.e. set of schemas)

Solution – Path (i.e. sequence of actions) to go from the initial state to a state entailed by the goal state.

## **Heuristics for Search**

#### **Ignore Preconditions Heuristic**

Drop all preconditions from actions. By itself, this is NOT an admissible heuristic as it may over estimate the solution.

Modified Approach – Delete all effects except those literals that are in the goal. Then count the number of actions needed to reach the goal.

When combined with the cost to get to the current node, this heuristic allows you to use A\* search to find a plan.

Exact Count: NP-Hard since it does not reduce the number of states to search.

In P time, can approximate the cost within log(n) factor where n is the number of literals in the goal

#### **Ignore Delete Lists**

Remove all delete lists (i.e. set of negated literals) from all actions.

Literals in the state are monotonically increase and if the goal is possible, it is eventually found.

Still leaves a problem that is **NP Hard since it** does not reduce the number of states to search.

Approximate solution can be found using this heuristic in polynomial time using hillclimbing.

#### **State Abstraction**

A many-to-one mapping from states in the ground representation of the problem to the abstract representation.

Example: In the plane cargo problem, require that all packages have the same destination (e.g. a hub) and that packages can only start in one of five airports.

This usually entails ignoring some fluents.

**Decomposition** – Key ideal in defining heuristics. It entails dividing a problem into parts, solving each part individually and then combining the parts. Similar to divide and conquer algorithm.

**Subgoal Independence Assumption** – Cost of solving a conjunction of subgoals is **approximated** by the sum of the costs to solve each subgoal independently.

**This assumption can be optimistic or pessimistic.** Optimistic if when solving each subgoal, actions that would otherwise cancel each other do not. Pessimistic as there **may be redundant actions**.

## **Planning Graph**

**Planning Graph** – Special data structure used to give better heuristic estimates for the cost of a plan.

**Polynomial size approximation** of the tree one would get by exploring all actions.

Useable for propositional planning problems only.

Level – Organizational structure for a planning graph.

Each level is denoted as  $S_i$ 

•  $S_0$  – Initial State

Each level is linked by a set of possible actions.

•  $A_0$  – Set of all possible actions possible in level  $S_0$ 

In each level, the set of achieveable literals are shown. For a given  $S_i$ , both the positive (P) and negative  $(\neg P)$  could hold given different sets of actions.

 $S_i$  and  $A_i$  alternate in the tree.

Action  $A_i$ :

- Preconditions:  $S_i$
- Effects:  $S_{i+1}$

Persistence Action – Type of noop. Used to preserve/persist any literal which is not negated by an action.

Every literal has a persistence action (small square in action) from  $S_i$  to  $S_{i+1}$  in the planning graph.

Once a literal appears in a level  $\boldsymbol{S}_i$  it remains present for all future levels of the planning graph.

Mutex – Mutual exclusion.
Curved links to indicate things (e.g. actions, literals, etc.) that cannot occur at the same time.

Leveled Off – When two consecutive levels of a graph are identical.

This is the termination condition of the planning graph.

Given graph with  $\boldsymbol{l}$  literals and  $\boldsymbol{a}$  actions:

- O(l) Nodes maximum in each  $S_i$
- $O(l^2)$  Mutex links in each  $S_i$
- O(a + l) Maximum number of nodes in each  $A_i$
- $O((a+l)^2)$  Mutex links in each  $A_i$
- O(2 \* (al + l)) Effect and precondition links because each persistence action goes
  to one effect and one precondition link and every standard action a could go to l
  precondition and l effect links.

Hence, for an n level planning graph, the maximum size is  $O(n(a+l)^2)$  which is polynomial space. The construction time is equivalent.

## **Using Planning Graphs for Heuristic Estimation**

Unsolvable Planning Problem: Goal does not appear in the final step in the planning graph.

**Level Cost of g\_i** - First level in the planning graph where goal literal  $g_i$  first appears.

Three Methods to Estimate Conjunction of Goals

- 1. Max-Level
- 2. Level-Sum
- 3. Set Level

Max Level – Largest level cost amongst all the goal literals.

Admissible.

Level Sum – Sum of the level costs of all goal literals.

Inadmissible

Set-Level – Level in the graph where all literals in the SET of goal literals first appear without any pair being mutually exclusive.

issible Admissible.

## Graphplan

**Graphplan** – An algorithm that uses a planning graph to find a solution to a planning problem.

nogoods – Hash table containing level and goal combinations in the planning graph that failed to yield solutions. This prevents unnecessary repeat searching of the graph.

<code>INITIAL\_PLANNING\_GRAPH</code> — Builds the planning graph for the initial state of the CSP (i.e.  $S_0$ ).

Conjuncts – Returns the goal statement as a conjunction of literals.

Extract\_Solution – Search through the planning graph to try to find the solution using either a constraint satisfaction approach or a search. There are two common implementations of this function.

NumLevels - Number of levels in the planning graph

**Expand\_Graph** – Expand the graph to include action  $A_t$ , state  $S_{t+1}$ , and all mutex relations.

def GraphPlan(problem): # Returns a solution or None

graph = INITIAL\_PLANNING\_GRAPH(problem) # Build planning graph for initial state goals = CONJUNCTS(problem.GOAL)

nogoods = {} # Empty hash table in Python

for t = 0 to infty:

# Only try to find a solution if the goal is achievable

if all goals present and non-mutex in S t of graph:

solution = Extract\_Solution(graph, goals, NumLevels(graph), nogoods) if(solution is not None):

return solution

# Termination found as graph has leveled off and no goods unchanged

if graph and nogoods unchanged since t-1:

return None

graph = Expand\_Graph(graph, problem)

## Possible Implementations of the EXTRACT SOLUTION Function

## **Extract\_Solution** as a Constraint Satisfaction Problem

**Variables:** Actions at each level of the graph. Hence a given action may appear as multiple different actions if it is present in multiple levels of the graph.

**Domain:** An action is either **IN** or **OUT** of the plan.

**Constraints:** Mutex relations (between literals and actions), goal literals, and preconditions.

Given this definition, any CSP solver can be used to find a plan if it exists.

## **Extract\_Solution** as a Backward Relevant State Search Problem

Each state in the planning graph contains a pointer to the previous level in the planning graph as well as a set of unsatisfied goals.

Search Criteria:

- 1. Initial State  $S_n$  in the planning graph (since working backwards)
- 2. Actions Set of actions  $A_{i-1}$  available in state  $S_i$ . User selects a **subset of** conflict-free actions whose effects cover the goals in that state. Conflict free means actions which are not mutex and whose preconditions are not mutex.
- 3. Result A set of preconditions at state  $S_{i-1}$  based off the actions in  $A_{i-1}$  that must be fulfilled.
- 4. Goal Reach  $S_0$  with all goal literals satisfied
- 5. Cost 1 for each action.

Level Cost – The level in the graph where a literal first appears. Example: Any literal in the initial state has a level cost of 0.

No-good – A set of goals at a particular level in the graph for which no planning solution exists. Graphplan is PSpace-Complete and making the planning graph can be done in polynomial time.

Heuristics are needed to choose among actions in the planning graph.

#### **Graphplan Heuristic**

- 1. Select the literal with the highest level cost
- 2. Prefer the actions where the sum of the level costs of their preconditions is smallest.

This is a greedy based approach.

## Types of Mutexes – Between Both Actions and Literals

**Inconsistent Effects** – One action negates the effect of another action.

**Interference** – One action's effect is the negation of the precondition of another action.

Competing Needs – One action's precondition is mutually exclusive (not only negated) with the precondition of another action.

Inconsistent Support – Two literals in a state can only be achieve through mutually excluded actions.

## Practice Final Question #4

## **Predicates**

- 1. Shoe(shoe) Returns whether "shoe" is a shoe.
- 2. Sock(sock) Returns whether "sock" is a sock.
- 3. Foot(foot) Returns whether "foot" is a foot.
- 4. Bare(foot) Returns whether "foot" is bare (i.e. has no socks or shoes)
- 5. HasSock(foot) Returns whether "foot" has a sock on already.
- 6. HasShoe(foot) Returns whether "foot" has a shoe on already.7. OnGround(sock) Returns whether "sock" is on the ground.
- 8. *OnGround(shoe)* Returns whether "shoe" is on the ground.
- 9. SameFoot(foot, shoe) Returns whether "foot" and "shoe" go on the same side (e.g. left or right)

## Constants

Foot:  $foot_{Left}$ ,  $foot_{Right}$ Sock:  $sock_1$ ,  $sock_2$ Shoe:  $shoe_{Left}$ ,  $shoe_{Right}$   $Init\Big(Bare\big(foot_{Left}\big) \land Bare\big(foot_{Right}\big) \land Foot\big(foot_{Left}\big) \land Foot\big(foot_{Right}\big) \\ \land Sock\big(sock_1\big) \land Sock\big(sock_2\big) \land Shoe\big(shoe_{Left}\big)$ 

 $\land Shoe(shoe_{Right}) \land OnGround(sock_1) \\ \land OnGround(sock_2) \land OnGround(shoe_{Left}) \\ \land OnGround(shoe_{Right}) )$ 

 $\begin{aligned} \textit{Goal}: \textit{HasSock}(foot_{\textit{Left}}) \land \textit{HasSock}(foot_{\textit{Right}}) \land \textit{HasShoe}(foot_{\textit{Left}}) \\ \land \textit{HasShoe}(foot_{\textit{Right}}) \end{aligned}$ 

 $\begin{array}{c} \textbf{Action}(\ PutOnSock(foot,sock),\\ Precond: Foot(foot) \land Sock(sock) \land Bare(foot) \land OnGround(sock)\\ Effect: \neg Bare(foot) \land HasSock(foot) \land \neg OnGround(sock)) \end{array}$ 

 $\begin{array}{ll} \textbf{Action}(PutOnShoe(foot, shoe),\\ Precond: Foot(foot) \land Shoe(shoe) \land HasSock(foot) \land \neg HasShoe(foot) \\ \land OnGround(shoe),\\ Effect: HasShoe(foot) \land \neg OnGround(shoe)) \end{array}$ 

## Example Plan

 $PutOnSock(foot_{Left}, sock_1) \\ PutOnSock(foot_{Right}, sock_2) \\ PutOnShoe(foot_{Left}, shoe_{Left}) \\ PutOnShoe(foot_{Right}, shoe_{Right}) \\$ 

## **Knowledge Representation**

Complex domains require more general and flexible knowledge representation paradigms than "toy" domains like the Wumpus World.

Common items that need to represented

- Events
- Time
- Physical Objects
- Beliefs

Ontological engineering - A

field that studies the methods and methodologies for representing knowledge specifically in ontologies.

Ontology – Determine what kinds of things exist without determining the things properties or their interrelationships. It decides on a vocabulary of predicates, functions, and constants.

Frameworks which can be used to represent facts about the world so they can be used by knowledge based agents.

Upper Ontology - Hierarchical ontology in an Object-Oriented like style. The most general representation/concept is at the top of the tree/hierarchy. This is further subdivided into more specific classifications/concepts.

Not able to handle exceptions to rules well.

Most successful ontologies are specific to a certain domain.

Example: Create an ontology for circuits so that theorem provers could be developed to check the circuits.

Ontology - Organizes everything in the world into a hierarchy of categories.

**Knowledge Reasoning Systems** 

**First Order Predicates** 

Objects in the world are often group into Predicate checks for membership of an object within a category.

> Example: A category of objects could be Basketballs. The first-order predicate to check if something is a basketball:

> > Basketball(b)

Reification

Creating a category as an object itself. Hence, Basketballs is an object that all basketballs are a component of. It turns a proposition into an object. Hence:

 $\forall b[Basketball(b) \Rightarrow b \in Basketballs]$ 

Subset/Subcategory/Subclass - A

category that is a subset of a parent class. The subclass inherits a set of features from the parent class.

**Example:** Basketballs is a subclass of Balls

Inheritance – Entails that all properties that are true about the parent class are true about the subclass as well.

**Example:** If all *Food* is edible, then if *Fruit* is a subclass of *Food*, then all Fruit is edible too.

Exhaustive Decomposition – Every object in an original set is assigned to a subcategory.

Partition - An exhaustive decomposition where all subcategories are disjoint (i.e. non-overlapping)

Taxonomy/Taxonomy Hierarchy -Organizational structure for representing subclass relationships.

**Facts Taxonomies Can State** 

An object is a member of a category. Example:

Two ways to create categories

2. Objects

1. First Order Predicates

categories.

 $BB_9 \in Basketballs$ 

A category is a subclass of another category. Example:

 $Basketballs \in Balls$ 

All members of a category have some property. Example:

> $(x \in Basketballs)$  $\Rightarrow$  Spherical(x)

Members of a category can be recognized by some set of properties. Example:

 $Round(x) \wedge Orange(x)$  $\wedge$  (Diameter(x) = 9.5")  $\land (x \in Ball) \Rightarrow Basketball(x)$  Category as whole has some properties. Example:

 $Dogs \in DomesticatedSpecies$ 

**Physical Composition** 

Physical composition is useful to represent knowledge of physical objects.

PartOf – Relation that categorizes objects by saying they are part of another object.

**Example:** PartOf (Bucharest, Romania)

Composite Object - A representation of an object by asserting the existence of its parts and their relationships.

Example: Define a Biped as having two legs that are attached to a body.

BunchOf - A relation that forms a composite object of definite parts but no structure.

Objects in the bunch are parts of the object not elements in it.

**Relationship Between PartOf and BunchOf** 

If an element is part of a category, then it is a PartOf a bunch containing that category.

 $\forall x [x \in s \Rightarrow PartOf(x, BunchOf(s))]$ 

 $\forall y[[\forall x, x \in s \Rightarrow PartOf(x, y) \Rightarrow PartOf(BunchOf(s), y)]$ 

If all members of a set/object are part of a bunch y, then a bunch of that set is also a part of the larger bunch y.

Logical Minimization - Define an object as the smallest one possible while still satisfying certain conditions.

Measurements

Lengths and measures are turned into abstract measure objects.

Unit Functions – Used to represent lengths/measurements in terms of a unit (e.g. inches, hours, dollars, etc.). Note these do not return a normalized value rather are essentially a representation of units.

Example #1:

 $Length(L_1) = Inches(d) = Centimeters(2.54 \times d)$ 

Example #2:

 $[b \in Basketballs] \Rightarrow [Diameter(b) = Inches(9.5)]$ 

Measures in a knowledge system are not numbers, but they can be used for ordering using symbols such as: >.

## **Events**

Event Calculus - A logical language that deals with time rather than situations. Event calculus reifies (i.e. groups) fluents and events.

PDDL uses a situation calculus that cannot say anything except before and after events.

Predicate – A new predicate that tests whether some point in time or during some interval in time.

**Example:** T(At(Shankar, Berkeley), t)A test of whether Shankar is at Berkeley at some point t.

Specific events are part of an events category.

Example: Describing the event  $E_1$  of Shankar flying from SF to LAX could be:

 $(E_1 \in Flyings) \land Flyer(E_1, Shankar)$  $\land Origin(E_1, SF)$  $\land$  Destination( $E_1$ , LAX) Time Interval - Has a start and end time  $(t_1, t_2)$ . Also can be used in Event Calculus in the same way as individual points in time t.

## **Event Calculus Fluents**

T(f,t) - Predicate for whether fluent f is true at time or interval t

Happens(e,i) -Predicate for whether event e happened over interval i

Initiates(e, f, t) -Predicate for whether event e caused fluent f to start to hold at time t.

Terminates(e, f, t) -Predicate for whether event e caused fluent f to stop holding at time t.

Clipped(f, i) – Predicate for whether fluent f ceased to be true sometime during interval i.

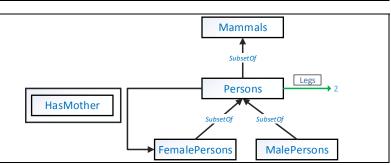
Restored(f, i) – Predicate for whether fluent f became true sometime during interval i.

#### **Semantic Networks**

- Used to integrate reasoning with knowledge systems based on categories.
- Do not support n-ary relations.
- Used to perform inheritance reasoning on graphs to determine what properties an object has.

#### Notation

- Objects/Categories Ovals or boxes
- Connections Between Categories Themselves Labeled Links
- Connections Between Objects in the Same Category Outlined Box
- **Default Values** Arrows to nothing



## **Description Logics**

**Description Logics** – First order logics geared toward making it easier to describe definitions and properties of categories.

**Example Description Logic: CLASSIC** 

Three Primary Inference Tasks in a Description Logic

Subsumption - Checking if one category Classification - Checking if one object is a subset of another.

belongs to a category.

Consistency - Checking whether the membership criteria of a category is satisfiable.

## Syntax in CLASSIC

 $And(Concept_1, Concept_2, ...)$ 

Example:

Bachelor

= And(Unmarried, Adult, Male)

**All**(RoleName, Concept)

Example: All of a person's daughters are married and Unemployed:

All(Daughters, And(Married, Unemployed)

AtLeast(Integer, RoleName)

Example: AtLeast(2, Daughters)

AtMost(Integer, RoleName)

Example: AtMost(3, Sons)

OR is not possible in classic nor negation so it is weaker than first order logic.

Fills(RoleName, IndividualName1, ...)

Role name fills one of the individual names.

Example: Fills(Department, Physics, Math)

 $OneOf(IndividualName_1, ...)$ 

Often is P-time and involves unification.

Selects one of the individual name objects. It is

a limited type of disjunction.

## **Reasoning with Default Information**

## **Jumping To Conclusions -**

Assuming default values for an object to be true without verification and if it is later shown to be untrue taking it back.

Nonmonotonic Logic - Set of beliefs in the knowledge base does not grow overtime. Rather new evidence can change existing beliefs.

Circumscription - More powerful and precise version of the closed world assumption. Every particular predicate is assumed to be "as false as possible" for every object except those objects for which it is known to be true.

Circumscribed Predicate - A specific predicate that a circumscribed reasoner is allowed to assume is false unless it is known to be true.

Example:  $Abnormal_1(x)$ which entails every object is normal unless known otherwise. Circumscription deals with preferred models of the knowledge base rather than the requirement of truth of all models.

A model is preferred to another if it has fewer abnormal objects.

## Default Logic – A type of non-monotonic logic. It is a formalism of default rules which have the form:

## Prequisite: Justification

Conclusion

## **Example of a Default Rule:**

Bird(x): Flies(x)Flies(x)

The conclusion Flies(x) is true unless it is known the justification is false. This rule can be rewritten:

 $Bird(x) \land \neg Abnormal(x) \Rightarrow Flies(x)$ 

## Extension (S) -

Maximum set of consequences from the default theory/rules and the facts.

For a given set of default rules and facts, there may be multiple possible extensions.

## **Nixon Diamond Problem**

Fact:  $Republican(Nixon) \land Quaker(Nixon)$ 

**Consider Two Default Rules** Quaker(x):Pacifist(x) $Republican(x) : \neg Pacifist(x)$ Pacifist(x) $\neg Pacifist(x)$ 

## **Two Possible Extensions**

**Extension #1:** { $Republican(Nixon) \land Quaker(Nixon) \land Pacifist(Nixon)$ } **Extension #2:** { $Republican(Nixon) \land Quaker(Nixon) \land \neg Pacifist(Nixon)$ }

Both have the same number of abnormal objects so neither is preferred. Using these extensions, you can derive the abnormal object in the default rules.

## **Decision Theory and Decision Theory Agents**

Previous agents dealt with the world assuming everything was either: true, false, or unknown.

Rational Decision - Dependent on the relative importance of various goals and the likelihood (i.e. degree/extent to which) these goals can be achieved.

Example:

 $\forall p \ Symptom(p, Toothache)$  $\Rightarrow$  Disease(p, Cavity)

Not always true since you can have a toothache for reasons other than a cavity.

Possible Solution:

 $\forall p \ Symptom(p, Toothache)$ 

 $\Rightarrow$  Disease(p, Cavity) V Disease(p, GumDisease) V ...

This solution can be prohibitive since not all causes may be known or there are too many to enumerate individually.

Decision Theory – Takes the utility of all possible agents and adds it to some calculation based on the probability of achieving each of the possible goals.

DecisionTheory = ProbabilityTheory + UtilityTheory

Preference – The extent to which the agents prefers certain goal states/outcomes to others.

Example: An agent may prefer coffee twice as much as tea.

Utility Theory – Used to represent and reason about preferences. Utility theory says that every state has a degree of usefulness (i.e. utility) to an agent and that the agent prefers states with high utility. Maximum Expected Utility (MEU) – Agent should choose the action which yields the highest expected payoff among the available choices def DT\_Agent(percept):

# Returns an action

persistent belief\_state # Probability belies about the current state of the world

persistent actions # Set of agent actions.

# Update set of probabilities based on percept and set of available actions update belief\_state based on action and percept

for action\_i in actions:

calculate outcome\_probability\_i based off action description and belief state

select action with highest expected utility given outcome\_probability and utility information

return action

## **Probability**

Sample Space  $(\Omega)$  – Set of all possible worlds. In other words, they are the set of all things that could be the world.

Elements in the sample space are mutually exclusive.

Example: Sample space for rolling two six sided dice is: (1,1), (1,2), ... (2,1), ...(6,6)

Probability Model/Probability Distribution -Associate a number,  $P(\omega)$ , between 0 and 1 with each element ( $\omega$ ) in the sample space ( $\Omega$ ) with the condition that:

$$\sum_{\omega\in\Omega}P(\omega)=1$$

Probabilistic assertions may not be about individual worlds. Rather, it may deal with sets of them.

Example: Probability of the sum of a two dice roll equaling 11 entails the case of (6,5) and (5,6).

Events – Set of all possible worlds where a corresponding proposition holds (e.g. rolling 11).

Unconditional or Prior Probability - Degree of belief that a proposition holds in the absence of any other probability.

**Example:** P(11) or P(doubles) for two dice roll.

**Evidence** – Additional information that may reveal information about the probability of other events.

Conditional Probability – Probability factoring in evidence from events. It is defined as:

$$P(A|B) = \frac{P(A \land B)}{P(B)}$$
 for  $P(B) > 0$ 

Example:  $P(11|Dice_1 = 5)$  is:

$$\frac{P(11 \land Dice_1 = 5)}{P(Dice_1 = 5)} = \frac{\frac{1}{36}}{\frac{1}{6}} = \frac{1}{6}$$

**Notation for Probability Theory** 

- Variables Initial capital letter
- Value from Domain/Sample Space - Initial lower case letter

Random Variable (X) – A function that maps elements  $(\omega)$ from the sample space  $(\Omega)$  to the set of real numbers.

$$X:\Omega\to\mathbb{R}$$

Example Random Variable: Bet \$3 on whether a coin flip is heads or tails. The random variable could be:

$$X(HEADS) = 3$$

X(Tails) = -3

**Domain Variable** – Enumerate possible elements in the state space. There are **DIFFERENT from** random variables as they may not map to the set of real numbers.

**Example Domain Variable:** 

 $Weather = \{cloudy, rainy, sunny, snowy\}$ 

Joint Probability Distribution - Probability distribution of the Cartesian product of two or more random variables.

**Example:** P(Cavity, Weather)

Joint probability distributions allow us to discuss probability for sentences involving AND (A):

**Example:**  $P(Cavity = true \land Weather = sunny)$ 

## **Random Variable Classifications**

**Boolean Random Variable (also** known as an Indicator Random Variable) - Random variable where each point in the sample space is mapped to one of two values.

Discrete Random Variable -

Random variable where the sample space is finite or if the image of the random variable is a subset of the integers.

Continuous Random Variable - Usually has a domain that consists an infinite number of states and where the function is continuous on the domain.

Example: Sample space could be points in a room and random variable could be the temperature at a point in degrees Celsius.

Calculating probabilities of continuous random variables usually involves computing an integral.

## **Important Probability Functions**

$$P(A|B) = \frac{P(A \land B)}{P(B)}$$

$$P(A) = 1 - P(\neg A)$$

$$P(A \lor B) = P(A) + P(B) - P(A \land B)$$

$$P(A \lor B) = P(A) + P(B) - P(A \land B)$$

Random variables often have interrelated values. Example: Probability of a toothache, cavity, and dentist catching your gums are related.

	Toothache		¬Too	thache
	Catch	¬Catch	Catch	¬Catch
Cavity	0.108	0.012	0.072	0.008
¬Cavity	0.016	0.064	0.144	0.576

Joint Probability Distribution for Random Variables *Cavity*, *Catch*, and *Toothache*.

This approach is not scalable with large numbers of random variables as it grows a rate of  $2^n$  for n random variables.

Unconditional/Marginal Probability – Probability of a single random variable or single random variable's state without dependence on other random variables. Marginalization/Summing Out – Given a joint probability function, P(Y, Z) of two random variables, Y and Z, the marginal probability of random variable Y is found by:

$$\vec{P}(Y) = \sum_{z \in Z} \vec{P}(Y, z)$$

Example:  $\vec{P}(Cavity) = \{0.108 + 0.012 + 0.072 + .008, 0.016 + 0.064 + 0.144 + 0.576\}$ 

Note the resulting probability is a VECTOR.

**Conditioning** – Dependent on the conditional probabilities to find the marginal probability of *Y* through:

$$\vec{P}(Y) = \sum_{z \in Z} \vec{P}(Y|z) * P(z)$$

Note the resulting probability is a VECTOR.

## **Conditioning Example**

Example: Find the conditional probabilities P(Cavity|Toothache) and  $P(\neg Cavity|Toothache)$ .

 $P(Cavity|Toothache) = \frac{P(Cavity \land Toothache)}{P(Toothache)}$ 

 $P(\neg Cavity | Toothache) = \frac{P(\neg Cavity \land Toothache)}{P(Toothache)}$ 

Both P(Cavity|Toothache) and  $P(\neg Cavity|Toothache)$  contain  $\frac{1}{P(Toothache)}$ . Hence this can be simplified to:

 $\alpha(P(Cavity \land Toothache) + P(\neg Cavity \land Toothache)) = 1$  $P(Cavity \land Toothache) = 0.108 + 0.012 = 0.12$ 

 $P(\neg Cavity \land Toothache) = 0.016 + 0.064 = 0.08$ 

Hence,  $\alpha$  is:

 $\alpha(0.12 + 0.08) = 1$   $\alpha = 5$ 

 $P(Cavity \land Toothache) = 0.6$  $P(\neg Cavity \land Toothache) = 0.4$  Normalization
Constant (α) –
Used to simplify
calculations and to
ensure the results
each the expected
value (e.g. 1 for
probabilities)

Independence – Random variable states do not affect one another. Hence, joint probability distributions can be factored into separate disjoint distributions.

When A and B are independent: P(A|B) = P(A)

**Bayes' Rule** 

Given two random variables A and B, then:

 $P(A \land B) = P(A|B)P(B)$  $P(A \land B) = P(B|A)P(A)$ 

Hence:

P(A|B)P(B) = P(B|A)P(A)

**Importance of Bayes' Rule:** If you need to know P(A|B), it is hard to find but you know, P(B|A), you can use Baye's rule in combination with marginal probabilities to solve for P(A|B)

**Example:** 70% of people with meningitis have a stiff neck. Odds of meningitis are 1/50000 (0.00002) and the odds of a stiff neck are 1/100 (0.01). The probability of P(M|SN) is:

$$P(M|SN) = \frac{P(SN|M)P(M)}{P(SN)} = \frac{0.7 * 0.00002}{0.01} = 0.0014$$

## Learning

Learning – Process by which an agent improves its performance on future tasks after making observations about the world.

## **Applications of Learning**

- 1. Programmer could not predict all possible situations an agent could encounter.
- 2. Programmer cannot predict changes over time.
- 3. Programmer might not have any idea to program a solution to the same problem themselves.

Component Improvements and Available Learning Techniques Depend On

- 1. Component to be improved.
- 2. Prior knowledge the agent has.
- 3. The representation used for the data and the component.
- 4. Available feedback to learn from.

Inductive Learning – Learning from a set of input/output pairs and generating a general function that governs those pairs.

Input is usually a vector of attribute values.

Deductive/Analytic Learning – Start from a set of general rules and derive things logically entailed from these general reules. **Unsupervised Learning** – Agent learns patterns from the input although no explicit feedback is supplied.

Example: Clustering – Input examples are grouped into potentially useful clusters.

Reinforcement Learning – Agent learns through a series of reinforcements (rewards or punishments).

Example #1: Lack of a tip at the end of a journey gives the taxi agent it did something wrong.

**Example #2:** Winning for a chess playing agent is a reinforcement it did something right.

## **Supervised Learning**

Supervised Learning – Agent observes input-output pairs and learns a function that maps from the input to the output. Semi-supervised Learning – Given a few labeled examples, the agent must make what it can from a large set of unlabelled examples.

$$(x_1, y_1), (x_2, y_2), ..., (x_n, y_n)$$

These are generated by some unknown function, f, defined as: y = f(x)

Hypothesis – A learned function h that approximates the unknown function f

Test Set – Disjoint from the training set. Used to test the quality of the hypothesis function h Classification –
Learning problem
where the output *y*is a finite collection
of values.

Regression – When the output y is always a number (often an infinite range)

**Consistent Hypothesis** – Any hypothesis function h that agrees (i.e. is consistent) with all input-output pairs.

A given set of data may have multiple consistent hypotheses.

Ockham's Razor – Always prefer the simplest hypothesis.

Definition of "Simplest" may vary.

**Example of a Simpler Hypothesis** – First order polynomial versus degree 7 polynomial.

## **Decision Trees**

## A supervised learning algorithm

Takes a vector of input attributes and returns a **single output value**.

Input attributes can be either continuous or discrete.

Focus of this class is on Boolean decision trees. Hence the outputs are either:

- Positive Examples return true
- Negative Examples return false

Leaf nodes correspond to the decision tree's result. Internal nodes corresponding to one of the input attributes.

Not all paths (branches) in a decision tree need to be the same length.

#### Important Pseudocode Functions and Methods

Plurality\_Value(examples) – Returns the most common boolean result from the set of examples

Importance(attribute, examples) – Returns a value quantifying the importance of attribute for the set of examples.

**DecisionTree()** – Python style constructor for an object of class DecisionTree

add\_branch(attribute\_value, subtree) - Method to append a subtree to the tree with the edge having the value "attribute\_value". Initial Call: decision\_tree = Decision\_Tree\_Algorithm(all\_examples, all\_attributes, {})

# Builds a decision tree

 ${\color{red} \textbf{def Decision\_Tree\_Algorithm}(examples, attributes, parent\_examples):}$ 

# examples – Remaining unclassified examples

# attributes – Remaining attributes not yet in the tree

# parent\_examples - Set of all examples in this node's parent.

 $\hbox{\tt\# No examples match this classification so return most common value for set of parents } \\$ 

if( len(examples) == 0 ):

return PLURALITY\_VALUE(parent\_examples)

# All examples agree so return the agreed upon classification

elif( all examples have same classification ): return classification

# Since no attributes remaining, take most common value from remaining examples

elif( len(attributes) == 0): return PLURALITY\_VALUE(examples)
else:

# Find the most important attribute

A = argmax\_(a in attributes)Importance(a, examples)

# Create a new tree

tree = DecisionTree()

# Iterate through all attribute values.

for v\_k in A:

subset\_examples = { exs in examples and E.A == v\_k }

subtree = Decision\_Tree\_Algorithm( subset\_examples, attributes - A, examples)

# Add the subtree to the tree

tree.add\_branch( v\_k, subtree )

return tree

## **Decision Tree Importance Function**

**Importance** function in the decision tree algorithm selects the next attribute in the tree.

Good attribute selections result in example sets that contain either only positive or only negative examples.

Bad attribute selections result in example sets that have the same proportion of positive and negative examples.

**Information Gain** – Quantifies the quality of an attribute selection.

Entropy H(v) – Fundamental quantity in information theory. It is a measure of the uncertainty of a domain variable.

The higher the entropy, the higher the uncertainty.

$$H(v) = -\sum_{v_k \in V} (P(v)log_2(P(v)))$$

Entropy of a Boolean Random Variable B(q):

$$B(q) = -1 * (q * \lg(q) + (1 - q) * \lg(1 - q))$$

**Entropy of an Attribute in a Decision Tree:** 

$$H(Goal) = B\left(\frac{p_k}{p_k + n_k}\right)$$

where  $p_k$  is the number of positive examples and  $n_k$  is the number of negative examples.

Information gain is defined as:

$$Gain = H(S_i) - H(S_{i+1})$$

For an attribute, A, in a decision tree, this simplifies to:

$$Gain(A) = B\left(\frac{p}{p+n}\right) - Remainder(A)$$

Remainder(A) is a weight sum of the entropy of each random variable and its likelihood of occurring:

$$Remainder(A) = \sum_{v_k \in A} \left( rac{p_k + n_k}{p + n} * B\left(rac{p_k}{p_k + n_k}
ight) 
ight)$$

## **Neural Networks**

Neurons - Type of brain cell. Electrochemical activity in the network of neurons is responsible for most mental activity.

## **Benefits of Neural Networks**

- 1. Perform distributed computation.
- 2. Tolerate noisy inputs

Neural networks are composed of nodes or units called neurons.

Input Activation Function Output Function Σ(wi\*in\_i)

## **Basic Neuron Structure**

Each input link has a different weight as shown as the thickness of the input arrow.

## **Neuron Structure**

A neuron is a link from unit i to unit j that propagates the activation signal  $a_i$  from i to j.

Note the activation signal is different than the activation function.

Weight  $(w_{i,i})$  – Numeric value which determines the strength and sign of the connection.

Output of the Unit is derived from the weighted sum function  $(in_i)$  which is defined for unit j as:

$$in_j = \sum_{i=0}^n (w_{i,j} * a_i)$$

Activation Function  $(g_i)$  – From the weight function, it derives the neuron j's output  $(a_i)$ . It is defined as:

$$a_j = g(in_j) = g\left(\sum_{i=0}^n (w_{i,j} * a_i)\right)$$

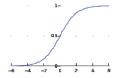
Each neuron has a single output that can be fed into several other neurons.

Types of Activation Functions (g)

Logistic Function – A sigmoid curve in an "S" to mark the

Threshold Activation Function – Output is binary (i.e. 0 or 1) depending on the weighted sum function's  $(in_i)$ value and the threshold value.

transition as more gradual.



**Common Function for Logistic Function:** 

$$L = \frac{1}{1 + e^{-(x - x_0)}}$$

Perceptron – Uses a threshold activation function in the neural network's neurons.

Sigmoid Perceptron - Uses the logistic function in the neural network's neurons.

Feed-forward network - Neuron connections are only in a single direction. Hence, back connections are not permitted.

**Recurrent Networks** – Outputs of neurons are allowed to go back and serve eventually as its own input.

Can lead to oscillation in result but are more realistic.

Feed-forward networks are usually arranged into layers where each layer only receives inputs from the previous laver.

Hidden Layer - Any layer that is not connected to either an input or an output.

**Perceptron/Single-Layer Neural Network** 

 All inputs are connected to nodes whose outputs are the final outputs.

**Neuron Weighted Sum Function: Given** inputs  $x_1$  and  $x_2$  with activation signals  $a_1$  and  $a_2$  respectively, then the weight function for neuron 0 is:

$$w_0 + w_{1,0} * a_1 + w_{2,0} * a_2$$

**Capabilities of a Feed-Forward Perceptron Network** 

- 1. Can calculate AND as well as OR
- 2. Cannot calculate XOR (parity) or binary summation.
- 3. Generally can learn only linearly separable functions.

There are some problems where a perceptron will perform better than a decision tree (e.g. majority function) while there are others where the perceptron will perform worse (e.g. restaurant seating problem).

Logistic Function - A type of function used as an activation function for a neuron. It has a general equal of the form:

$$LogisticFunction(t) = \frac{1}{1 + e^{-t}}$$

Logistic Function for Sigmoid Perceptron - The logistic function rewritten as a hypothesis function for a sigmoid perceptron is:

$$h_{\overrightarrow{w}}(\overrightarrow{a}) = LogisticFunction(\overrightarrow{w} \cdot \overrightarrow{a}) = \frac{1}{1 + e^{-\overrightarrow{w} \cdot \overrightarrow{a}}}$$

Value of Logistic Function **Over Threshold Function:** 

The logistic function is differentiable in the real number space while a threshold function is not.

## Loss Function in a Perceptron

Given a training set, E, with examples in the form  $(\vec{x}, y)$  where y is a binary output (i.e. 0 or 1) and a hypothesis function  $h_{\overrightarrow{w}}(\overrightarrow{x})$ , then the error or loss of  $h_{\overrightarrow{w}}(\overrightarrow{x})$  is:

$$Loss(\overrightarrow{w}) = \sum_{(x,y)\in F} (y - h_{\overrightarrow{w}}(\overrightarrow{x}))^2$$

Logistic Regression – Processing of fitting weights to a sigmoid perception.

Squaring in the loss function is used to prevent negative errors skewing the results.

Goal - Make the loss/error function as close to 0 as possible.

## Finding the Loss Function for a Perceptron

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$\frac{\partial}{\partial w_i} Loss(\vec{w}) = \frac{\partial}{\partial w_i} (y - h_{\vec{w}}(\vec{x}))^2$	Take the partial derivative with respect to one dimension (i.e. input to the perceptron) in the weight vector $\vec{w}$ . Hence, if there is 10 inputs into the perceptron, then there is 11 elements in $w$ (one for each input and one for the offset). This would require 11 partial derivatives.		
$\frac{\partial}{\partial w_i} Loss(\vec{w}) = 2 * (y - h_{\vec{w}}(\vec{x})) \times \frac{\partial}{\partial w_i} (y - h_{\vec{w}}(\vec{x}))$	Derivate of $x^2 = 2x * \frac{dx}{dx}$ . A similar implementation of the chain rule is followed here.		
$\frac{\partial}{\partial w_i} Loss(\overrightarrow{w}) = 2 * (y - h_{\overrightarrow{w}}(\overrightarrow{x})) \times g'(\overrightarrow{w} \cdot \overrightarrow{x}) \times \frac{\partial}{\partial w_i} (\overrightarrow{w} \cdot \overrightarrow{x})$	$y$ is a constant so its derivative is 0. Define $g'(\vec{x})$ as the derivative of $h_{\overrightarrow{w}}(\overrightarrow{x_t})$ which is the activation function $g$ .		
$\frac{\partial}{\partial w_i} Loss(\overrightarrow{w}) = 2 * (y - h_{\overrightarrow{w}}(\overrightarrow{x})) \times g'(\overrightarrow{w} \cdot \overrightarrow{x}) \times \overrightarrow{x_i}$	The partial derivative of $\overrightarrow{w}$ is 1 resulting in the final equation.		
$\frac{\partial}{\partial w_i} Loss(\overrightarrow{w}) = 2 * (y - h_{\overrightarrow{w}}(\overrightarrow{x})) \times g(\overrightarrow{w} \cdot \overrightarrow{x}) (1 - g(\overrightarrow{w} \cdot \overrightarrow{x})) \times \overrightarrow{x_i}$	Given the logistic function $g(t)=rac{1}{1+e^{-t}}$ , then $g'(t)=g(t)ig(1-g(t)ig)$		
$\frac{\partial}{\partial w_i} Loss(\overrightarrow{w}) = 2 * (y - h_{\overrightarrow{w}}(\overrightarrow{x})) \times h_{\overrightarrow{w}}(\overrightarrow{w} \cdot \overrightarrow{x}) (1 - h_{\overrightarrow{w}}(\overrightarrow{w} \cdot \overrightarrow{x})) \times \overrightarrow{x_i}$	The activation function $g$ is also known as $h_{\overrightarrow{w}}.$		

## **Perceptron Update Rule**

## Perceptron Update Rule

For each element in the training set  $(\vec{x}, y)$ ,  $newW_i$  is calculated for each dimension in  $\vec{w}$ . For the next training example to be considered,  $newW_i$ 's become the  $w_i$ 's.

$$newW_i = w_i + \alpha (y - h_{\overrightarrow{w}}(\overrightarrow{x})) h_{\overrightarrow{w}}(\overrightarrow{x}) (1 - h_{\overrightarrow{w}}(\overrightarrow{x})) x_i$$

**Fixed Rate Ratio** –  $\alpha$  is constant value.

 $\alpha$  can be a temperature like schedule. The textbook gives as an example:

$$lpha(t) = rac{1000}{1000 + t}$$
 where  $t$  is the number of iterations.

## **Feed-Forward Learning**

Feed-forward networks can compute more complicated networks than perceptron networks.

Example: At most a four level feedforward network can create a spike which is otherwise impossible with a perceptron network. In a feed-forward network, the activation function of the output unit is a composite of the activation function of many other units. Example: unit 5 has as inputs units 3 and 4. Unit #3 and #4 both have as inputs the initial inputs 1 and 2 (i.e. inputs to the network).

$$a_5 = g(w_{05} + w_{35}a_3 + w_{45}a_4)$$

This can be rewritten as:

$$a_5 = g(w_{05} + w_{35}g(w_{03} + w_{13}x_1 + w_{23}x_2) + w_{45}g(w_{04} + w_{14}x_1 + w_{24}x_2))$$

Hence, to solve for the individual weights requires **nonlinear regression**.

## **Back Propagation**

Back Propagation – The error at the output of the neural network is easily determined. However, this is not the case for the output at hidden layers. Back propagation involves propagating the error at the output nodes to the inner nodes in order to update the weights in the inner nodes.

 $\Delta_k = Err_k \cdot g'(in_k)$ 

$$in_k = w_k \cdot \overrightarrow{x_k}$$

$$Err_k = (y - a_k)$$

Where  $\overrightarrow{x_k}$  is the inputs into output node k

The new weight for the  $j^{th}$  input into output node k is:

$$newW_{jk} = w_{jk} + \alpha \cdot a_j \cdot \Delta_k$$

## **Error Back Propagation**

For a hidden node j,  $\Delta_i$  is defined as:

$$\Delta_j = g'(in_j) \cdot \sum_k w_{jk} \cdot \Delta_k$$

This would then be used in an

In back propagation, the values of  $a_k$  can be found be expanding back to the inputs as shown under "Feed-Forward Learning".

After minimizing the error at the output, the error is driven back in the network. This is through a process called back propagation.

k is the error of the k

 $in_k$  is the dot product of the weights and inputs into neuron k (i.e.  $\vec{w} \cdot \vec{x}$ ).

 $||y-h_{\overrightarrow{w}}(\overrightarrow{x})||$  - Double bars represent the magnitude of the vector.

## **Non-Parametric Learning**

Parametric Model – The learning model summarizes the data with a set of parameters of a fixed size

**Nonparametric Model** – A model that is not characterized by a bounded set of parameters.

Memory Based Learning/Instance Based Learning –
Learning based off a look-up table of learned examples.

## **K-Nearest Neighbors**

Nearest Neighbor's learning is nonparametric since all training data is used to determine the nearest neighbor.

Decides classification by a "majority vote" approach.

Given a query vector,  $\overrightarrow{x_q}$ , look up the k nearest neighbors. Denote these k nearest neighbors of  $\overrightarrow{x_a}$  as:

$$NN(k, \overrightarrow{x_q})$$

If k is too small, the algorithm is susceptible to overfitting where the classification can be skewed by outliers.

If k is too large, the algorithm is susceptible to underfitting where the classification just becomes a majority function of the entire dataset.

**Usual Range of Ideal Value of** *k***:** Between 1 and the square root of the dataset size.

## Quantifying "Nearest"

 ${\it L^{P} Norm}$  - Commonly used distance equation.

$$L^{P}(\overrightarrow{x_{q}},\overrightarrow{x_{e}}) = \left(\sum_{i} \left| \overrightarrow{x_{e}} - \overrightarrow{x_{q}} \right|^{P} \right)^{\frac{1}{P}}$$

If P = 1, then it is the Manhattan distance. If P = 2, then it is the Euclidean distance. Hamming Distance – Given two strings of equal length, it is the number of positions where the symbols are different.

Example: 010 and 110 have a Hamming distance of 1.

Not all dimensions in the vector vary over the same range of numbers. As such, some degree of normalization of variation is required to prevent skewing the nearest neighbor approach.

A common variation normalization scheme:

$$\frac{x_i - \mu_i}{\sigma_i}$$

#### Where:

- $i i^{th}$  dimension
- $\mu_i$  Mean of the  $i^{\text{th}}$  dimension
- $\sigma_i$  Standard deviation of measurements in the  $i^{\text{th}}$  dimension

## **Algorithms to Find the Nearest Neighbor**

For each new point classification, cycle over all *N* elements in the data set and find the *K* nearest neighbors.

Time Complexity: O(N)

**k-d (i.e.** *k*-**Dimensional) Tree** – A balanced binary tree with an arbitrary number of dimensions.

Cycle over each dimension and split the elements in the data set where units less than median are on the right side of the split and those greater than the median are on the left side of the split.

Time Complexity:  $O(\lg(n))$ 

**Locality Sensitive Hashing** – A leads to a variation of k-Nearest Neighbors called **approximate nearest neighbors**. The hash function maps points in the space onto a single line. The hash function is designed such that if two points,  $x_j$  and  $x_j'$ , will be in the same hash bin if they are close to each other in the space. There may be some points that are not close that map into the same bin.

**Support Vector Machine** 

**Support Vectors** – The closest vectors to the separator (usually in each dimension)

Maximum Margin Separator – Output of a support vector machine. Minimizes generalization loss by having the maximum margin (distance from the training data).

Margin – Twice the distance between the separating hyperplane and the nearest example point. **Kernel Trick** – A list of functions that can be used to map a vector of real numbers and maps it to another vector of real numbers (often of with an increased number of dimensions) to make the data linearly separable.

## **Final Exam Additional Notes Page**

- Breadth first search, bidirectional search, and iterative deepening depth first search are only optimal if the graph has uniform step size.
- Optimal Strategy A sequence of contingent decisions that will lead to outcomes as least as good as any other sequence of decisions against an infallible player.
- Perfect Information Any situation where an agent has all relevant information with which to make a decision and the results of actions are deterministic 3.
- Path Consistency A two variable set  $(X_i, X_j)$  are path consistent with respect to a third variable  $X_m$  if for every assignment of values to  $X_i$  and  $X_j$  consistent with the constraint  $\{X_i, X_i\}$ , there is a valid assignment to  $X_m$  that satisfies the constraints  $\{X_i, X_m\}$  and  $\{X_m, X_i\}$ .
- Deleting from a Dictionary: del person["name"] 5.
- \*args: Arbitrary number of arguments in Python function
- 7. raise - Keyword to raise an exception.
- pass Needed in abstract class methods and properties 8.
- Horn program can derive entailed sentences in linear time and space.
- 10. Universe (M) Set of all values variables can range over.
- 11. Structure/Model (M) Combination of predicates, functions, universe, and constants.
- 12. State Conjunction of ground fluents.
- 13. In a planning graph, there are a maximum 2(al+l) precondition and effect lines per level.
- 14. State Abstraction Many to one mapping of states in the ground representation to the abstract representation.
- 15. Inheritance Entails that all properties that are true about the parent class are true for subclasses.
- 16. Description Logics First order logic geared toward representing definitions or properties of categories.
- 17. Circumscribed Predicate A predicate a circumscribed reasoner can assume to be false unless it is known to be true.
- 18. Sample Space  $(\Omega)$  Defines possible worlds in a probabilistic scenario.
- 19. Joint Probability Distribution Probability distribution of the Cartesian product of two or more random variables.
- 20. Test Set Set of input output pairs used to verify the quality of a hypothesis function.
- 21. Unsupervised Leaning Learning patterns in data without explicit feedback on the quality of the results.
- 22. Given a training set, E, with examples in the form  $(\vec{x}, y)$  where y is a binary output (i.e. 0 or 1) and a hypothesis function  $h_{\vec{w}}(\vec{x})$ , then the **error or loss of**  $h_{\vec{w}}(\vec{x})$  is:

$$Loss(\overrightarrow{w}) = \sum_{(x,y) \in E} (y - h_{\overrightarrow{w}}(\overrightarrow{x}))^2$$

23. Perceptron Update Rule – For each element in the training set  $(\vec{x}, y)$ ,  $newW_i$  is calculated for each dimension in  $\vec{w}$ . For the next training example to be considered,  $newW_i$ 's become the  $w_i$ 's.

$$newW_i = w_i + \alpha (y - h_{\overrightarrow{w}}(\overrightarrow{x})) h_{\overrightarrow{w}}(\overrightarrow{x}) (1 - h_{\overrightarrow{w}}(\overrightarrow{x})) x_i$$

Back Propagation - The error at the output of the neural network is easily determined. However, this is not the case for the output at hidden layers. Back propagation involves propagating the error at the output nodes to the inner nodes in order to undate the weights in the inner nodes.

$$\begin{array}{c|c} & \mathsf{Back} \, \mathsf{Propagation} \\ \hline \Delta_k = \mathit{Err}_k \cdot g'(in_k) \end{array}$$

$$in_k = w_k \cdot \overrightarrow{x_k}$$

$$Err_k = (y - a_k)$$

Where  $\overrightarrow{x_k}$  is the inputs into output node k

The new weight for the  $j^{\rm th}$  input into output

$$newW_{jk} = w_{jk} + \alpha \cdot a_j \cdot \Delta_k$$

## **Error Back Propagation**

For a hidden node j,  $\Delta_i$  is defined as:

$$\Delta_j = g'(in_j) \cdot \sum_i w_{jk} \cdot \Delta_k$$

This would then be used in an

## **Fall 2014 Practice Midterm Questions**

1. Briefly explain what the Turing Test is. Also, define the term agent in the context of Al.

The Turing Test involves a human posing a series of questions to a terminal and based off the responses it receives, the human is unable to determine whether the respondent is a human or computer. An agent perceives its environment through sensors and acts upon its environment through actuators.

2. What is the PEAS description of a task environment? Give a fully spelt-out example of an episodic task environment?

PEAS is an acronym for performance measure, environment, actuators, and sensors. The performance measure is the goals/tasks the agent is trying to achieve. The environment entails everything the agent interacts with or that interact with the agent. The actuators are the tools the agent uses to interact with the environment. Sensors are the tools the agent uses to perceive its environment.

An episodic agent is one where its current decisions have no effects on future decisions. An example would be a quality assurance robot that investigates the quality of objects coming off an assembly line. It determines whether an object has sufficient quality or is defective. Each object the episodic agent interacts with has no bearing on the QA's decision of quality for the next or future objects.

3. Brief describe iterative deepening depth first search and analyze its time complexity.

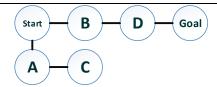
In standard depth first search, the agent explores the left most path in the graph until it reaches a leaf node (assuming the graph is finite). After reaching the leaf node, the agent then recurses one level in the graph and then tries to explore the next left most path for that node. This process continues until the graph is fully explored. Iterative deepening depth first search is slightly different than normal depth first search in that traverses the left most path in the graph until it reaches a left node or until the maximum specified depth is reached.

Iterative deepening begins with a maximum depth d. If no solution is found at depth 1, it repeats iterative depth first search with a maximum depth 2. If no solution is found at depth 2d, the process is repeated with maximum depth 3. This process continues until the graph is fully explored or the solution is found, whichever comes first.

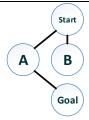
Assuming a finite graph, the runtime of depth first search is:  $O(b^d)$ , where b is the maximum branching factor out of any node and d is the recursion depth. Hence, IDFS has a running time of:

$$b^1 + b^2 + b^3 + \dots + b^d = O(b^{d+1})$$

4. Give a concrete example where a problem solving agent using A\* search would not traverse a graph in the same way as one using breadth-first search. Given an example where depth first search outputs A\*-search.



Breadth first search explores the frontier uniformly. In this case, BFS would add all successors of the start state to the frontier (i.e. A and B). It would then explore D which adds "D" to the frontier and then explores "A" which adds "C" to the frontier. Next it explores "D" which adds "Goal" to the frontier. It then explores "C" and adds nothing to the frontier before exploring "Goal" where it then terminates. In contrast, A\* explores the start state where it adds "B" and "A" to the frontier. Assuming the heuristic was euclidean distance, it would next explore "B" since its distance (actual + heursitic) is shorter than the distance from "A". It then adds "D" to the frontier. Next it explores D since its combined actual and heuristic distance to the goal is less than "A" where "Goal" is added to the priority queue. It finally explores "Goal" since its distance (actual distance only since heuristic distance is 0) is less than the combined distance from A and it returns a solution.



Depth-first search always takes the left most path in the graph so it would traverse this graph in two steps: Start  $\rightarrow$  A  $\rightarrow$  Goal. In contrast, A\* would add A and B to the frontier. It would then explore B first (assuming the heuristic is euclidean distance) since it is closer to the goal (including its actual cost than A). However, B is a dead end so it would then explore A then goal making it three steps.

Write a short python program which takes its command line arguments and sums them together. This program should make use of at least one function definition.

6. Given two admissible A\* heuristics, explain how to make a new heuristic which performs at least as well as either of them.

An admissible heuristic is one that is optimistic (i.e. it underestimates the remaining cost to the solution). The best admissible heuristic estimates as closely as possible the actual cost without exceeding it. If there are two or more admissible heuristics, you can form a composite heuristic which returns the maximum of the set of admissible heuristics. This will ensure that you select the best of the two heuristics in each situation which is in turn at least as good or better than each of the individual heuristics alone.

- 7. Briefly explain how genetic programming and local beam search are related hill climbing algorithms.
- 8. What is the minimax function? Given of an example where a beta cut might arise while running the minimax algorithm with alpha-beta pruning.

The minimax function is:

$$Minimax(s) = \begin{cases} Utility(s), & \text{If } TERMINAL\_TEST(s) \text{ is True} \\ \max_{a \in ACTIONS(s)} Minimax(RESULT(s, a)), \text{If } PLAYER(S) \text{ is MAX} \\ \min_{a \in ACTIONS(s)} Minimax(RESULT(s, a)), \text{If } PLAYER(S) \text{ is MIN} \end{cases}$$

- 9. Give pseudo-code for the AC-3 algorithm.
- 10. Consider the following situation for four light switches on the control panel of a nuclear power plant. (a) The first and last switch can never both be on. (b) At least one light must be on.

Define four variables  $L_1$ ,  $L_2$ ,  $L_3$ , and  $L_4$  which represent whether the first through fourth light switches are on respectively. The two relations for this are:

$$R_a: \neg(L_1 \land L_2)$$
  
$$R_b: (L_1 \lor L_2 \lor L_3 \lor L_4)$$

The combined relation is:

$$R_a \wedge R_b$$

## **Fall 2012 Practice Midterm Questions**

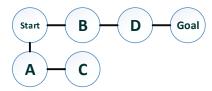
1. Briefly say what the Total Turing Test is and what a Rational Agent is.

The Total Turing Test is for a robot with artificial intelligence to pass entirely as a human being. To achieve this would require robotics and computer vision in addition to features in the standard Turing Test of knowledge representation, natural language processing, etc. A rational agent is one that for every percept sequence, it chooses the action that is expected to maximum the agent's utility given the percept sequence and whatever built-in knowledge it has.

2. Give the formal way to specify a problem solving agent.

There are five components needed to specify a problem solving agent. They are:

- a. Initial State: The initial state of the agent.
- b. Actions: The set of possible actions the agent can perform in a specific state.
- c. Results: Describes the change in the agent's state after an action has been performed.
- d. Goal Test: A test for whether the agent has reached one of its goal states.
- e. Cost Function: Describes the cost to perform any action.
- 3. Given an example problem and then explain how iterative deepening search might search the environment of this problem to find a goal. Explain the runtime and space complexity of IDS.



Iterative deepening explores the graph by incrementally increasing the maximum depth the algorithm can search. Hence, in this graph, it would start at depth = 1. It would explore A first, then reach its maximum depth of 1, then explore B, reach its maximum depth then terminate that round of searching. Since it did not find the goal, it would increase the depth to two. It would explore A->C, hit its maximum depth, then explore B->D, hit its maximum depth then terminate. It would then increase its depth to three, explore A->C find no more children, then recurse and explore B->D->Goal.

Space complexity for IDDFS is O(d) where d is the maximum depth of that round. This is the same as depth first search up to depth d since it only needs to keep track of the current node and its parents. The time complexity of IDDFS is:

$$\sum_{i=1}^{d} b^{i} = \frac{b^{1} - b^{d+1}}{1 - b} = \frac{(b^{d+1} - b)}{b - 1} = O(b^{d+1})$$

4. Given an example of the following programming language features in Python: generators, coroutines, and lambda.

```
Generator:
```

5. Give the resolution refutation of the following clauses.

$$(a \lor \neg b) \land (\neg a) \Rightarrow (\neg b)$$
$$(b \lor \neg c) \land (c) \Rightarrow (b)$$
$$(b) \land (\neg b) \Rightarrow ()$$

Hence the resolution refutation is proven since the empty clause was yielded.

## **Practice Final Questions**

1. Mod3(x\_1, ..., x\_n) is the propositional formula which returns true if the number of variables `x\_i` which are true in a truth assignment is exactly 0 mod 3. Write down a CNF formula for Mod3 in the case where `n=6`.

This uses the truth table\Karnaugh map approach to solve the problem. In the truth table, any assignment that makes the result false is added to the CNF as single clause that is the disjunction of the literals in the assignment but negated.

$$\bigwedge_{i=1}^{6} \left( \overline{x_i} \vee \bigvee_{1 \leq j \leq 6, i \neq j} x_j \right) \wedge \bigwedge_{i=1}^{6} \left( \bigwedge_{j=i+1}^{6} \left( \overline{x_i} \vee \overline{x_j} \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} x_k \right) \right) \wedge \bigwedge_{i=1}^{6} \left( \bigwedge_{j=i+1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigwedge_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq j} \overline{x_k} \right) \wedge \bigvee_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i, k \neq i} \overline{x_k} \right) \wedge \bigvee_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i} \overline{x_k} \right) \wedge \bigvee_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i} \overline{x_k} \right) \wedge \bigvee_{i=1}^{6} \left( x_i \vee x_j \vee \bigvee_{1 \leq k \leq 6, k \neq i} \overline{x_k} \right) \wedge \bigvee_{i=1}^{6} \left( x_i \vee x$$

2. Give the DPLL algorithm and explain each of the three main "shortcuts" it checks for.

**DPLL** – Resolution Finding Algorithm

Three Optimizations Over the Basic Resolution Algorithm:

- Early Termination: If all clauses are satisfied (have at least one positive literal) or any clause is false, terminate the algorithm.
- Pure Symbol Heuristic: A pure symbol is any symbol that has the same sign in all clauses. Pure symbols are set to true if they exist.
- 3. Unit Clause: A unit clause contains on a single literal. The variable in the unit clause is set to true to satisfy the clause.

```
def DPLL Satisfiable(s): # Returns True or False
   clauses = set of clauses from CNF representation of s
   symbols = list of symbols in s
   return DPLL(clauses, symbols, {})
def DPLL(clauses, symbols, model):
  # Check Early Termination
   if every clause is true in model:
     return True
   elif some clause is false in model:
     return False
   # Check Pure Symbol Heuristic
   P, value = FIND_PURE_SYMBOL(clauses, symbol, model)
   if P is not None:
     return DPLL(clauses, symbols - P, model U {P=value})
   # Check Unit Clause Heuristic
   P, value = FIND_UNIT_CLAUSE(clauses, model)
   if P is not None:
     return DPLL(clauses, symbols - P, model U {P=value})
   # Select first symbol and check both true and false
   P = FIRST(symbols)
   rest = REST(symbols)
   return DPLL(clauses, rest, model U {P = True})
     or DPLL(clauses, rest, model U {P = False})
```

(a) Let `x:= f(z)` and `y:= g(w)` explain how the unification algorithm from class would work on these inputs. (b) Now suppose `x:= [g(v), f(g(z))]` and `y:= [g(f(w)), f(w)]`. Explain how the unification algorithm from class would work on these inputs

```
Step #1: Unify("[g(v), f(g(z))]", "[g(f(w)), f(w)]", {}) # Remove the head of the lists.
                                                                                              Step #2: Unify("[f(g(z))]", "[f(w)]", Unify( "g(v)", "g(f(w))", {} )) # Remove outermost
                                                                                              function symbols g.
                                                                                              Step #3: Unify("[f(g(z))]", "[f(w)]", Unify("v", "f(w)", Unify("g", "g", {} )) # No unification
                                                                                              required since function operators are identical
                                                                                              Step #4: Unify("[f(g(z))]", "[f(w)]", Unify( "v", "f(w)", {} )) # Unify on variable v
                                                                                              Step #5: Unify("[f(g(z))]", "[f(w)]", Unify_Var("v", "f(w)", {} )) # Append to substitution
                                                                                              list for variable v
Step #1: Unify("f(z)", "g(w)", {})
                                                                                              Step #6: Unify("[f(g(z))]", "[f(w)]", { v \rightarrow f(w) } ) # Extract the first item in each list.
Step #2: Unify("z", "w", Unify("f", "g", {})) # Remove operator
                                                                                              Step #7: Unify("[]", "[]", Unify("f(g(z))", "f(w)", { v | -> f(w) } ) ) # Extract function symbol
Step #3: Returns False # Unification terminated since it was not possible to
                                                                                              f on the two functions
                                                                                              Step #8: Unify("[]", "[]", Unify( "g(z)", "w", Unify( "f", "f", { v |-> f(w) } ) ) ) # Unify on
unify f and g since they are different operators.
                                                                                              identical function symbols f
                                                                                              Step #9: Unify("[]", "[]", Unify( "g(z)", "w", { v |-> f(w) } ) ) # Perform Unify var on
                                                                                              variable w
                                                                                               \begin{tabular}{ll} \textbf{Step #10:} & Unify("[]", "[]", Unify\_Var("w", "g(z)", \{ v \mid -> f(w) \} ) ) \end{tabular} \begin{tabular}{ll} \textbf{Append substitution} \end{tabular} 
                                                                                              list for variable w
                                                                                              Step #11: Unify("[]", "[]", { v \rightarrow f(w), w \rightarrow g(z) } ) # Identical unification lists so no step
                                                                                              Step #12: { v |-> f(w), w |-> g(z) } # Final Substitution
```

4. Consider the problem where you have two socks and two shoes all of which are on the ground. You also have two feet. Your goal is to put on your shoes. Your feet can wear socks, but not shoes directly. Your available actions are to put on socks and put on shoes. Formulate this problem reasonably in PDDL. Then give an example plan solving it.

#### Predicates

10.Shoe(shoe) - Returns whether "shoe" is a shoe.
11.Sock(sock) - Returns whether "sock" is a sock.

12.Foot(foot) - Returns whether "foot" is a foot.

- 13.Bare(foot) Returns whether "foot" is bare (i.e. has no socks or shoes)
- 14. HasSock(foot) Returns whether "foot" has a sock on already. 15. HasShoe(foot) Returns whether "foot" has a shoe on already.

16. On Ground (sock) - Returns whether "sock" is on the ground.

17. On Ground (shoe) - Returns whether "shoe" is on the ground.

18. Same Foot (foot, shoe) - Returns whether "foot" and "shoe" go on the same side (e.g. left or right)

#### **Constants**

Foot:  $foot_{Left}$ ,  $foot_{Right}$ Sock:  $sock_1$ ,  $sock_2$ Shoe:  $shoe_{Left}$ ,  $shoe_{Right}$   $Init\Big(Bare\big(foot_{Left}\big) \land Bare\big(foot_{Right}\big) \land Foot\big(foot_{Left}\big) \land Foot\big(foot_{Right}\big) \\ \land Sock\big(sock_1\big) \land Sock\big(sock_2\big) \land Shoe\big(shoe_{Left}\big)$ 

 $\land Sock(sock_1) \land Sock(sock_2) \land Shoe(shoe_L \land Shoe(shoe_{Right}) \land OnGround(sock_1)$ 

 $\land OnGround(sock_2) \land OnGround(shoe_{Left})$  $\land OnGround(shoe_{Right})$ 

Goal:  $HasSock(foot_{Left}) \land HasSock(foot_{Right}) \land HasShoe(foot_{Left}) \land HasShoe(foot_{Right})$ 

 $\begin{array}{c} \textbf{Action}(\ PutOnSock(foot,sock),\\ Precond: Foot(foot) \land Sock(sock) \land Bare(foot) \land OnGround(sock)\\ Effect: \neg Bare(foot) \land HasSock(foot) \land \neg OnGround(sock)) \end{array}$ 

Action(PutOnShoe(foot, shoe),  $Precond: Foot(foot) \land Shoe(shoe) \land HasSock(foot) \land \neg HasShoe(foot)$  $\land OnGround(shoe).$ 

 $\land$  OnGround(shoe), Effect: HasShoe(foot)  $\land \neg$ OnGround(shoe))

#### **Example Plan**

 $PutOnSock(foot_{Left}, sock_1) \\ PutOnSock(foot_{Right}, sock_2) \\ PutOnShoe(foot_{Left}, shoe_{Left}) \\ PutOnShoe(foot_{Right}, shoe_{Right}) \\$ 

- 5. Show how the Graphplan algorithm would work on the example of the previous problem.
- 6. Define the following terms related to knowledge engineering: (a) ontology, (b) reification, (c) taxonomy.

Ontology – A formal naming and definition of the types, properties, and interrelationships of the entities that exist for a particular domain or discourse. It is a framework which can be used to represent facts about the world so they can be used by knowledge based agents.

Reification is the process of turning a predicate into an object. For example, all basketballs are reified into an object Basketball such that  $\forall b[Basketball(b) \Rightarrow b \in Basketballs$ .

Taxonomy/Taxonomy Hierarchy - Organizational structure for representing subclass relationships.

7. Explain and give an example of the following concepts from probability theory: (a) random variable, (b) marginalization, (c) Bayes' rule.

A **random variable** (X) maps elements in the state space ( $\Omega$ , i.e. the set of possible, disjoint worlds), to the set of real numbers. Hence:  $X:\Omega\to\mathbb{R}$ 

Example: We bet \$3 on the result of a coin flip. A random variable X could be:

$$X(HEADS) = 3$$
  
 $X(TAILS) = -3$ 

Marginalization is the extraction of probability of a single random variable from a joint probability distribution function.

$$\vec{P}(Y) = \sum_{z \in \mathcal{I}} \vec{P}(Y, z)$$

Note these are vectors.

Example:

	Toothache	¬Toothache
Cavity	0.1	0.2
¬Cavity	0.3	0.4

$$P(Cavity) = \{0.1 + 0.2, 0.3 + 0.4\} = \{.3, .7\}$$

**Bayes' Rule** comes from conditional probability which is defined as P(A) given B or:

$$P(A|B) = \frac{P(A \land B)}{P(B)}$$

Using  $P(A \land B)$ , it can be shown:

$$P(A|B)P(B) = P(B|A)P(A)$$

Example: Probability of a stiff neck if you have meningitis is 0.7. Probability of meningitis is  $\frac{1}{50000}$  and the probability of a stiff neck is 0.01. Hence the probability you have meningitis given a stiff neck is:

$$P(M|SN) = \frac{P(SN|M)P(M)}{P(SN)} = \frac{0.7 * 0.00002}{0.01} = 0.0014$$

8. Consider the following training set of 4-tuples.

(T,T,T,F)

(T,T,F, F)

(T,F,T, F)

(T,F,F, F)

(F,T,T,T)

Here `T` is short for true, `F` is short for false. The first three columns correspond to the variables  $x_1$ ,  $x_2$ ,  $x_3$ , the last column is the output of some function `f`. Calculate `Gain( $x_i$ )` for `i=1,2,3`. Which variable should we use as the top of a decision tree for `f`?

To calculate the information gain for each parameter, you only need to calculate the  $Remainder(x_i)$  of each attribute  $x_i$ . The attribute with the lowest Remainder is the one to selected. Hence:

$$Remainder(x_1) = \frac{1}{5} * B\left(\frac{1}{1}\right) + \frac{4}{5} * B\left(\frac{0}{4}\right)$$

$$Remainder(x_1) = 0$$

$$Remainder(x_2) = \frac{2}{5} * B(0) + \frac{3}{5}B\left(\frac{1}{3}\right)$$

$$Remainder(x_2) = 0 + \frac{3}{5}B\left(\frac{1}{3}\right)$$

$$Remainder(x_2) = -\frac{3}{5} * \left(\frac{1}{3}\lg\left(\frac{1}{3}\right) + \frac{2}{3}\lg\left(\frac{2}{3}\right)\right)$$

$$Remainder(x_2) > 0$$

$$Remainder(x_3) = \frac{2}{5} * B(0) + \frac{3}{5} * B\left(\frac{1}{3}\right)$$

$$Remainder(x_3) = 0 + \frac{3}{5} B\left(\frac{1}{3}\right)$$

$$Remainder(x_3) = -\frac{3}{5} * \left(\frac{1}{3} \lg\left(\frac{1}{3}\right) + \frac{2}{3} \lg\left(\frac{2}{3}\right)\right)$$

$$Remainder(x_3) > 0$$

Since  $x_1$  has the lowest remainder, it is the attribute that should be expanded at the top of the tree. This also makes intuitive sense as  $x_1$  results in sets containing only positive or only negative examples.

9. Give the formal definition of a perceptron. Explain and give an example of a feed forward network is and what a recurrent network is.

A **perceptron** is a neural network where the activation function (g) is exclusively a threshold function (e.g. 0 if below the threshold and 1 if above or equal to the threshold). It cannot have a logistic function as its activation function which has a more gradual turn-on profile. Such networks are called **sigmoid perceptrons**. A perceptron network is a single layer networks meaning the inputs are connected to **units** that are exclusively connected to final outputs.

A **feed-forward network** is a neural networks where the outputs neuron's only move in a single direction (i.e. forward). No back lines are allowed in the network so a neuron's output can never form part of its own input signals.

**Example of a Feed-Forward Network With a Single Neuron** 



A recurrent network is a neural networks where the outputs of neurons are looped back to eventually form part of the neuron's inputs (either directly or through a predecessor node).

**Example of a Recurrent Network With a Single Neuron** 



10. Give and explain the update rule for learning neuron weight from class.