

This diagram (on the following page) shows the interaction of the Marlin prover and verifier. It is similar to the diagrams in the paper (Figure 5 in Section 5 and Figure 7 in Appendix E, in the latest ePrint version), but with two changes: it shows not just the AHP but also the use of the polynomial commitments (the cryptography layer); and it aims to be fully up-to-date with the recent optimizations to the codebase. This diagram, together with the diagrams in the paper, can act as a “bridge” between the codebase and the theory that the paper describes.

## 1 Glossary of notation

|  |   |
|--|---|
| $\mathbb{F}$   | the finite field over which the R1CS instance is defined  |
| $x$  | public input  |
| $w$  | secret witness  |
| $H$  | variable domain   |
| $K$  | matrix domain   |
| $X$  | domain sized for input (not including witness)  |
| $v_D(X)$   | vanishing polynomial over domain $D$  |
| $A, B, C$  | R1CS instance matrices  |
| $A^*, B^*, C^*$  | shifted transpose of $A, B, C$ matrices given by $M_{a,b}^* := M_{b,a} \cdot u_H(b, b) \forall a, b \in H$<br>(optimization from Fractal, explained in Claim 6.7 of that paper)   |
| $\{\hat{\text{val}}, \hat{\text{row}}, \hat{\text{col}}\}_{\{A^*, B^*, C^*\}}$ | preprocessed polynomials from $A^*, B^*, C^*$ matrices containing LDEs of (respectively)<br>row positions, column positions, and values of non-zero matrix elements               |
| $\hat{\text{rowcol}}_{\{A^*, B^*, C^*\}}$                                      | the product polynomial of $\hat{\text{row}}$ and $\hat{\text{col}}$ , given separately for efficiency (namely<br>to allow this product to be part of a <i>linear</i> combination) |
| $\mathcal{P}$  | prover  |
| $\mathcal{V}$  | verifier  |
| $\mathcal{V}^p$  | $\mathcal{V}$ with “oracle” access to polynomial $p$ (via commitments provided<br>by the indexer, later opened as necessary by $\mathcal{P}$ )                                    |

## 2 Diagram

