# Graph

- Introduction
- Graph API
- Elementary Graph Operations
  - DFS: Depth first search
  - BFS: Breadth first search
  - CC: Connected Components

#### Major references:

- 1. Fundamentals of Data Structures by Horowitz, Sahni, Anderson-Freed,
- 2. Algorithms 4<sup>th</sup> edition Part 1 & Part 2 by Robert Sedgewick and Kevin Wayne
- 3. Wikipedia and many resources available from internet

Prof. Youngsup Kim, idebtor@gmail.com, Data Structures, CSEE Dept., Handong Global University

### Design pattern for graph processing

Design pattern: Decouple graph data type

Idea: Mimic maze exploration

## DFS (to visit a vertex v)

- Mark v as visited.
- Recursively visit all unmarked vertices w adjacent to v.

### Typical applications:

- Find all vertices connected to a given source vertex.
- Find a path between two vertices.

### Challenge:

How to implement?

#### Breadth-first search

**Depth-first search:** Put unvisited vertices on a **stack**. Breadth-first search: Put unvisited vertices on a queue.

**Shortest path:** Find path from s to t that uses fewest number of edges.

### **BFS:** (from source vertex s)

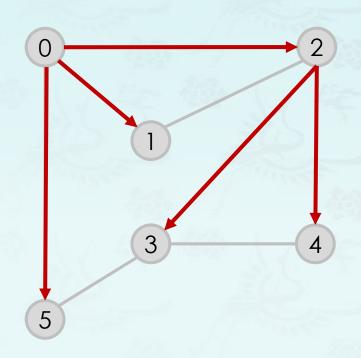
- Put s onto a FIFO queue, and mark s as visited.
- Repeat until the queue is empty:
  - remove the least recently added vertex v
  - add each of v's unvisited neighbors to the queue,
    - and mark them as visited.



**Intuition:** BFS examines vertices in increasing distance from s.

### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.



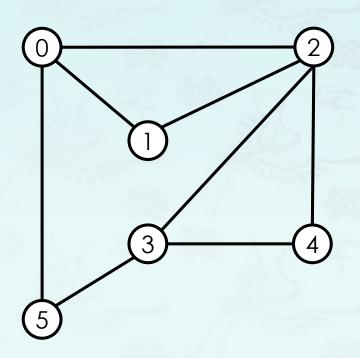
## graph4.txt

#### v parent[v] distTo[]

0	1-4	0
1	0	1
2	0	1
3	2	2
4 5	2	2
5	0	1

#### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.



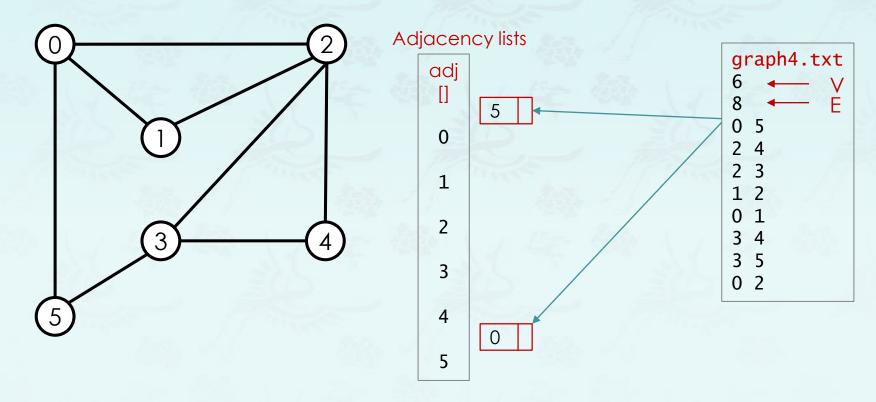
gı	raph4.txt
6	<b>←</b> ∨
8	<b>←</b> E
0	5
2	4
2	3
1	2
0	1
3	4
3	5
0	2

Graph g:

**Challenge:** build adjacency lists?

### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.

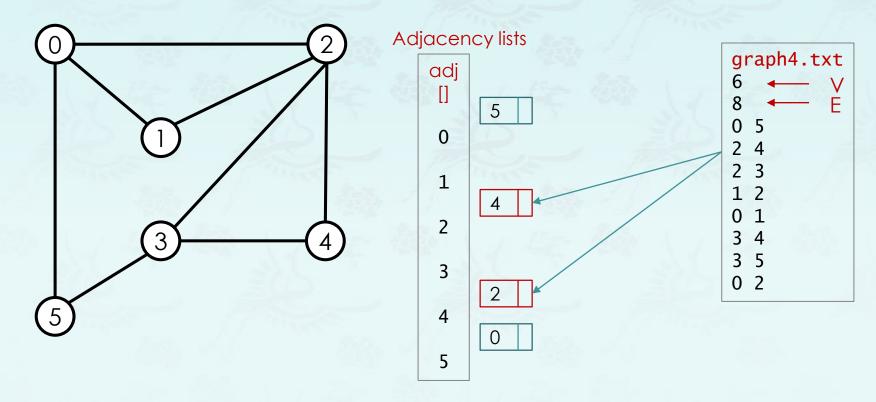


Graph g:

**Challenge:** build adjacency lists?

### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.

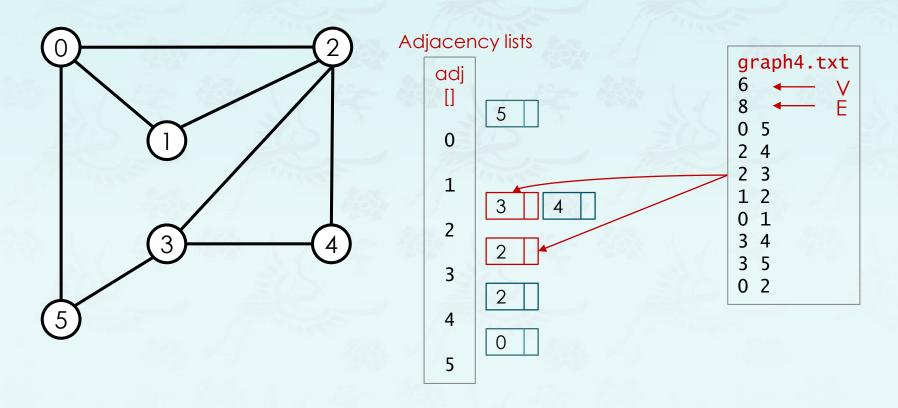


Graph g:

**Challenge:** build adjacency lists?

### Repeat until queue is empty:

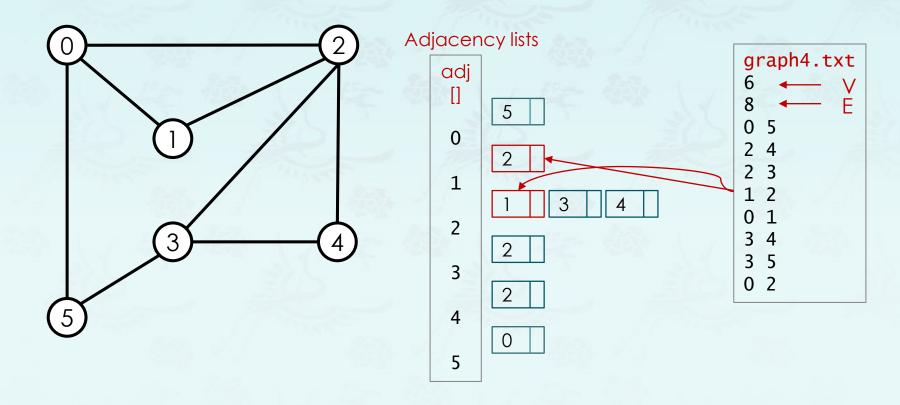
- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.



#### Graph g:

### Repeat until queue is empty:

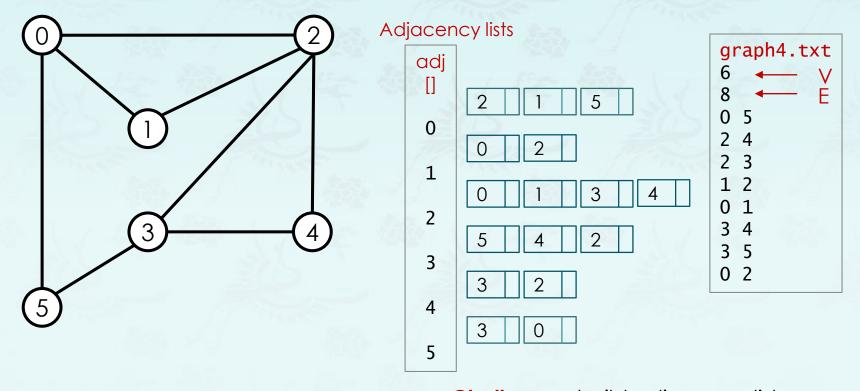
- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.



#### Graph g:

### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.

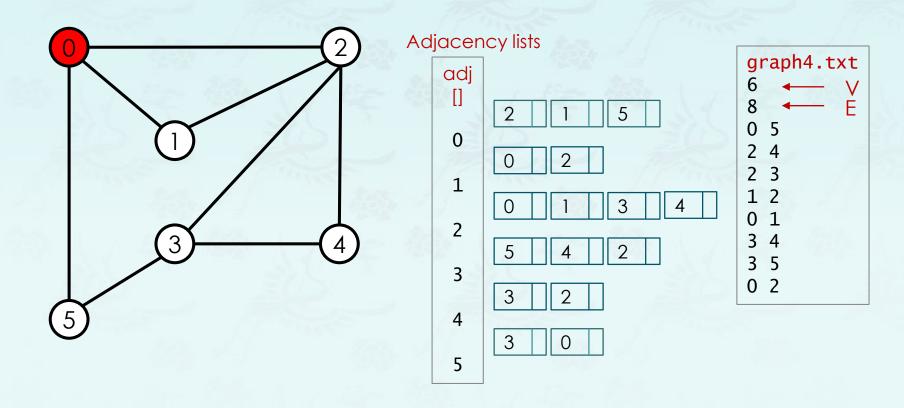


Graph g:

**Challenge:** build adjacency lists – Job done

#### Repeat until queue is empty:

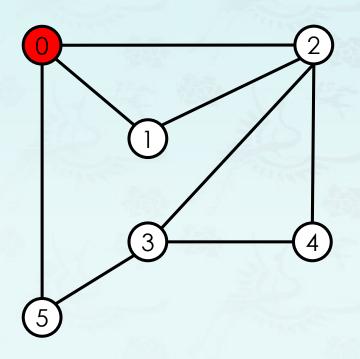
- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.



#### add 0 to queue:

# Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.

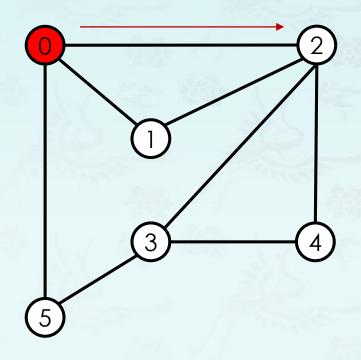


queue	v parent[v] distTo[]			
	0	1-4	0	
	1	<u> </u>	_	
	2	_	_	
	3	_	-	
	4	- 2		
0	5	_	_ =	

### add 0 to queue:

### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.

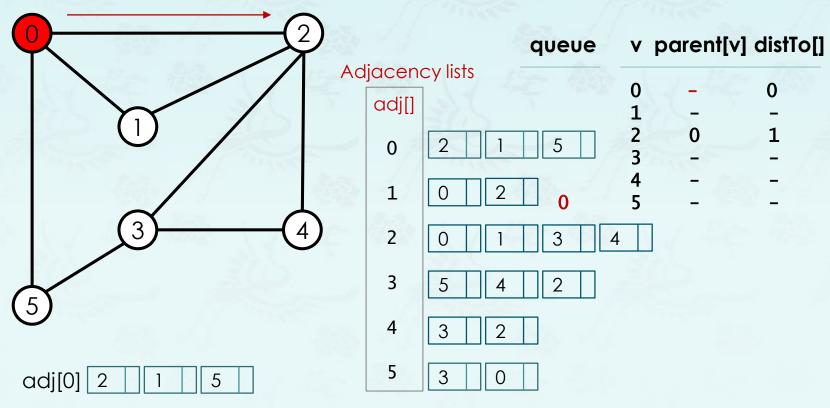


queue	v parent[v] distTo[]		
	0	1-4	0
	1	<u> </u>	_
	2	0	1
	3	111	- //
	4	- 2	- / /
0	5	-	- E

adj[0] 2 1 5

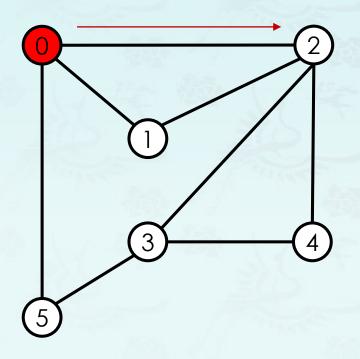
#### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.



### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.

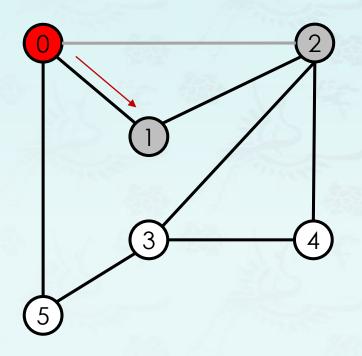


queue	v parent[v] distTo[]			
	0	1-4	0	
	1	\ <u>-</u>	_	
	2	0	1	
	3	_	<u>-</u>	
	4	_	2000 <b>-</b> 7 7	
2	5	_	_ =	

adj[0] 2 1 5

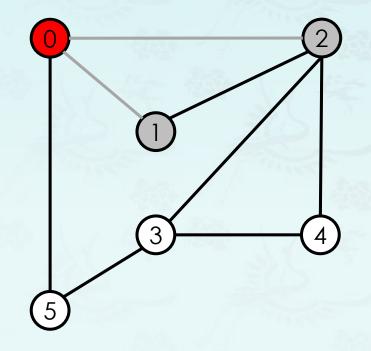
### Repeat until queue is empty:

- Remove vertex v from queue.
- Add to queue all unmarked vertices adjacent to v and mark them.



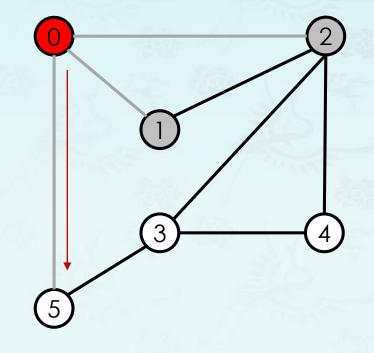
queue	v parent[v] distTo[]		
	0	1-4	0
	1	\ <u>-</u>	_
	2	0	1
	3	_	- //
	4	_	-//
2	5	-	- F

adj[0] 2 1 5



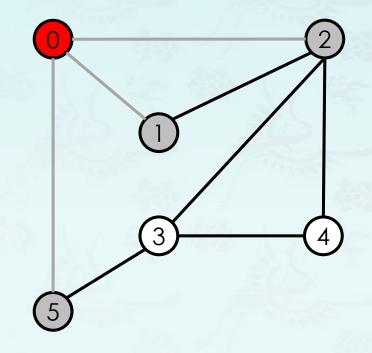
v parent[v] distTo[]		
0	1-4	0
1	0	1
2	0	1
3	_	_
4	- 2	- / /
5	_	- 1
	0 1 2	0 - 1 0 2 0

adj[0] 2 1 5



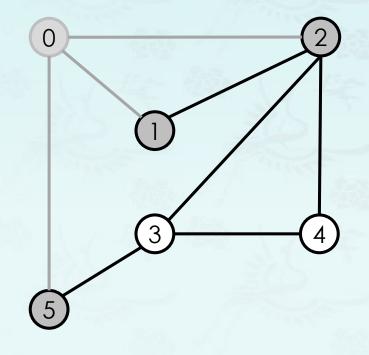
queue	v parent[v] distTo[]		
	0	1-4	0
	1	0	1
	2	0	1
	3	-	-
1	4	- 0	- / I
2	5	_	_

adj[0] 2 1 5



queue	v parent[v] distTo[]			
	0	1-4	0	
	1	0	1	
	2	0	1	
5	3	-	- //	
201	4	- 9	- ·	
2	5	0	1	

adj[0] 2 1 5

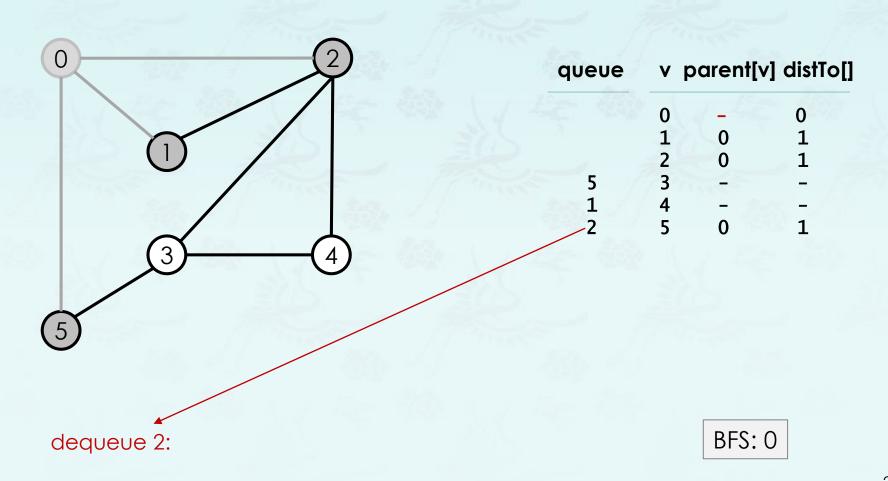


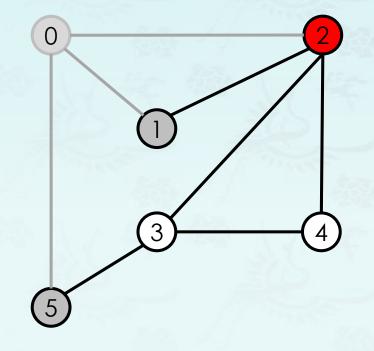
queue	v parent[v] distTo[]		
	0	1-4	0
	1	0	1
	2	0	1
5	3	141	- //
201	4	- 9	- / I
2	5	0	1

adj[0] 2 1 5

0 done

BFS: 0





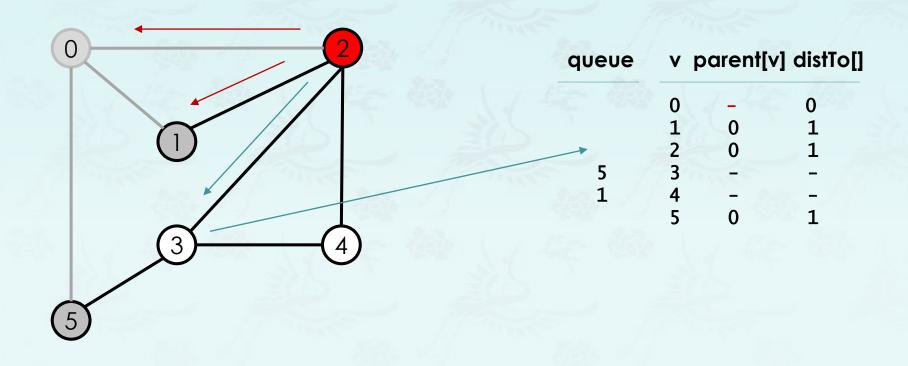
v parent[v] distTo[]		

adj[2] 0 1 3 4

dequeue 2: check 0, check 1, check 3 and check 4

BFS: 0

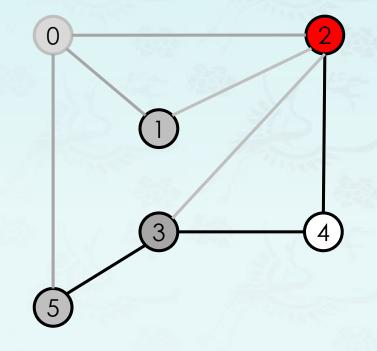
adj[2] 0



23

BFS: 0

dequeue 2: check 0, check 1, check 3 and check 4



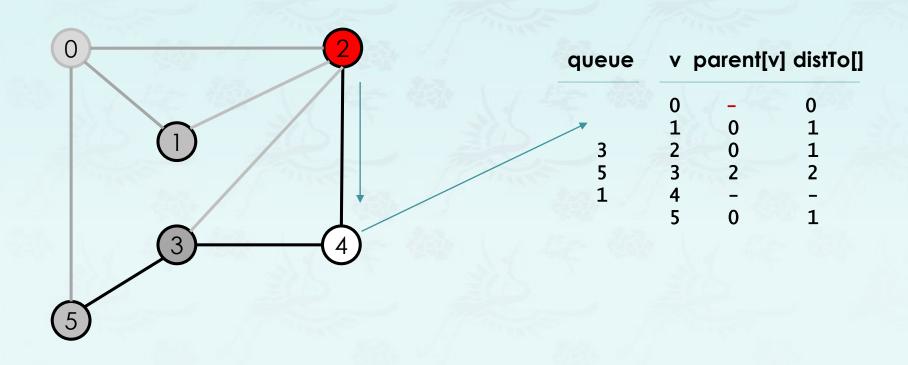
v parent[v] distTo[]		
0	1-4	0
1	0	1
2	0	1
3	2	2
4	- 2	- / I
5	0	1
	0 1 2	0 - 1 0 2 0

adj[2] 0 1 3 4

dequeue 2: check 0, check 1, check 3 and check 4

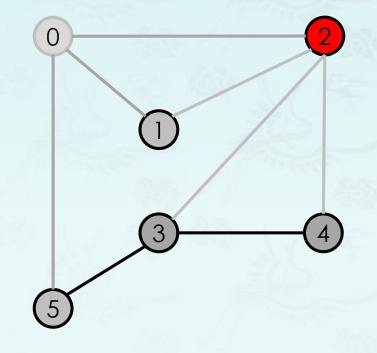
BFS: 0

adj[2] 0



BFS: 0

dequeue 2: check 0, check 1, check 3 and check 4

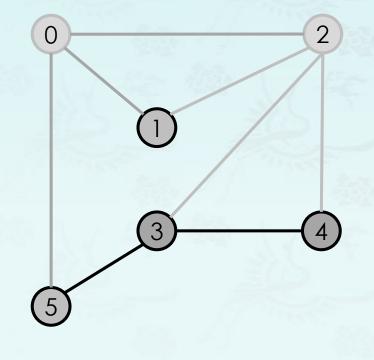


queue	v parent[v] distTo[]			
	0	1-4	0	
4	1	0	1	
3	2	0	1	
5	3	2	2	
1	4	- 2	-	
	5	0	1	

adj[2] 0 1 3 4

dequeue 2: check 0, check 1, check 3 and check 4

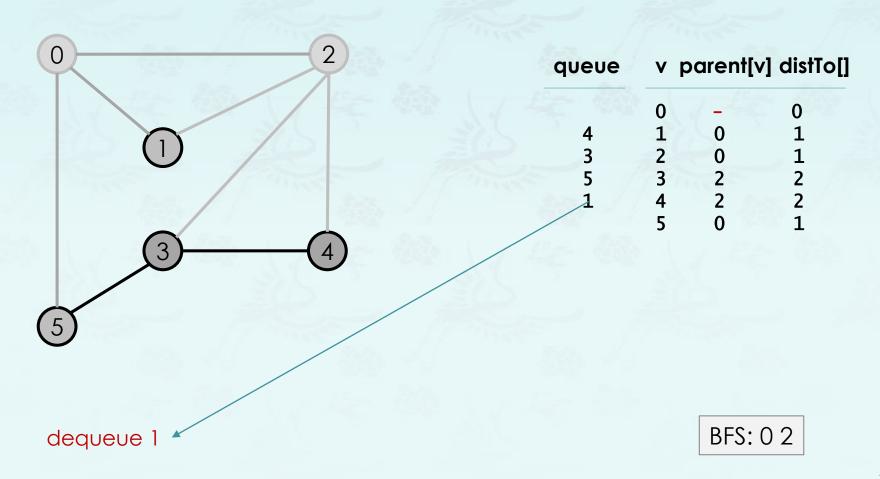
BFS: 0

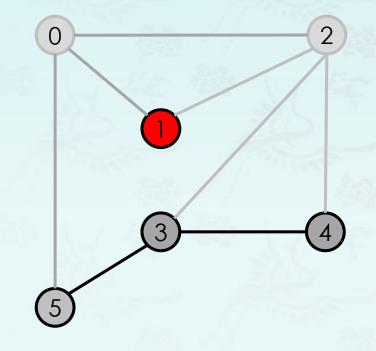


queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
3	2	0	1
5	3	2	2
201	4	2	2
	5	0	1

adj[2] 0 1 3 4

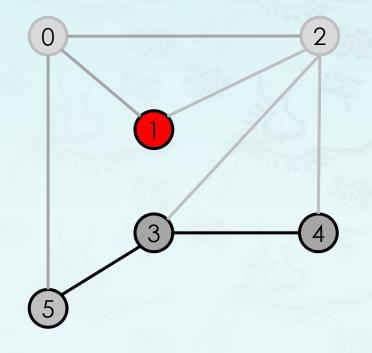
2 done





queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
3	2	0	1
5	3	2	2
	4	2	2
	5	0	1

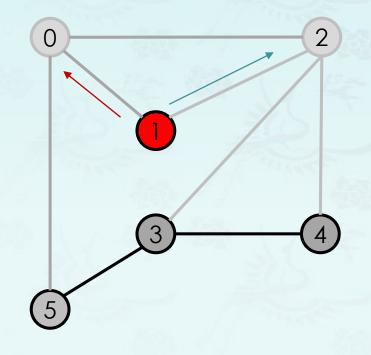
dequeue 1



queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
3	2	0	1
5	3	2	2
	4	2	2
	5	0	1

adj[1]0 2

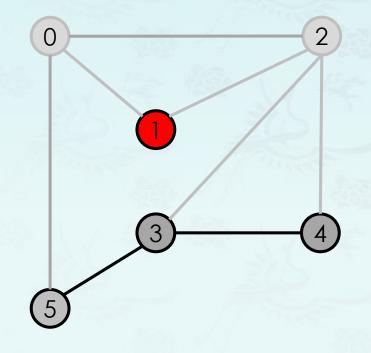
dequeue 1: check 0, and check 2



queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
3	2	0	1
5	3	2	2
	4	2	2
	5	0	1

adj[1] 0 2

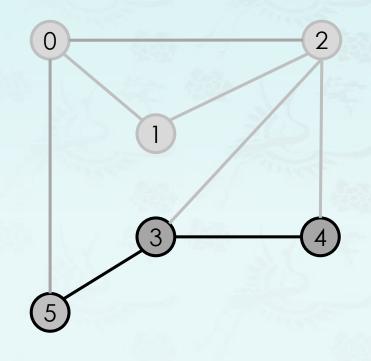
dequeue 1: check 0, and check 2



queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
3	2	0	1
5	3	2	2
	4	2	2
	5	0	1

adj[1] 0 2

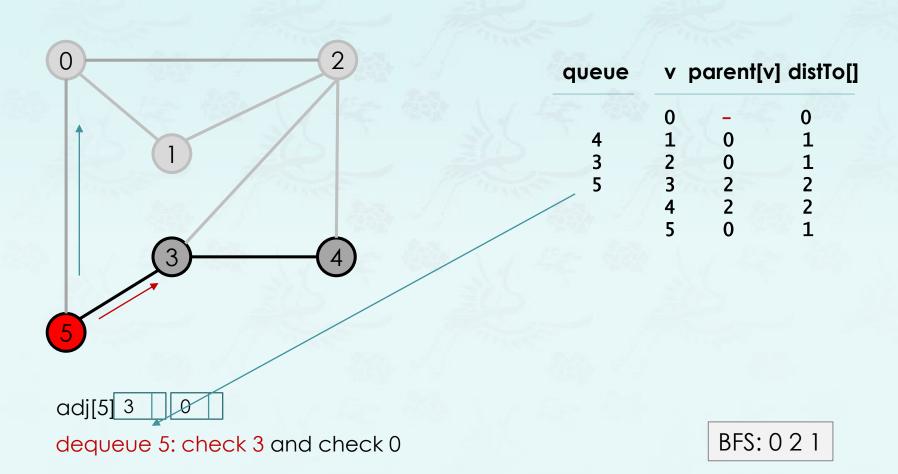
dequeue 1: check 0, and check 2



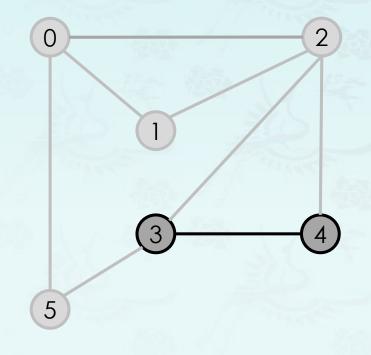
queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
3	2	0	1
5	3	2	2
	4	2	2
	5	0	1

adj[1] 0 2 1 done

BFS: 0 2 1



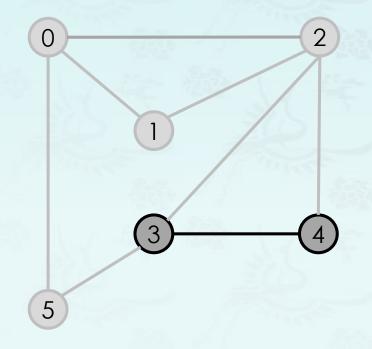
34



queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
3	2	0	1
	3	2	2
	4	2	2
	5	0	1

adj[5] 3 0 5 done

BFS: 0 2 1 5

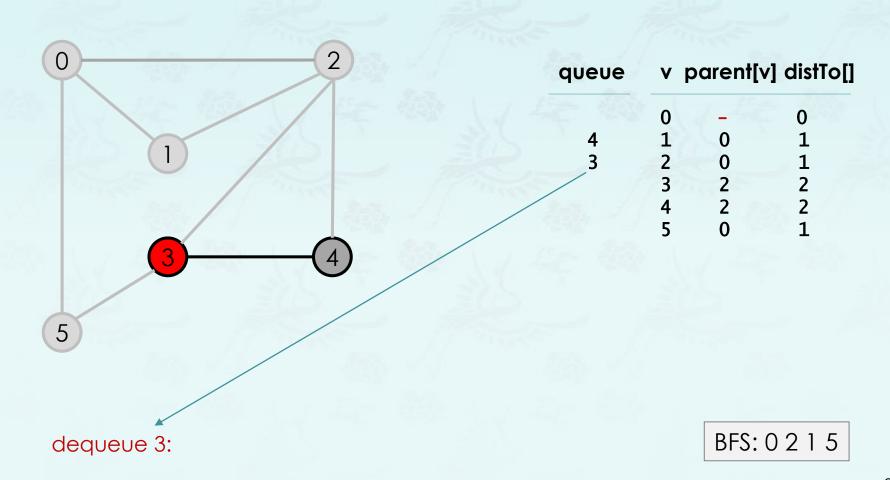


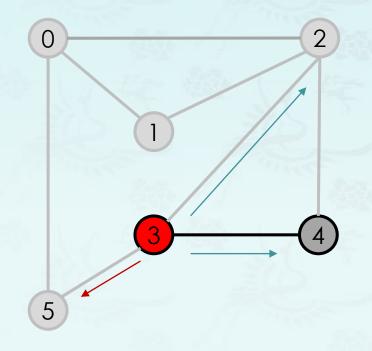
queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
3	2	0	1
	3	2	2
	4	2	2
	5	0	1

adj[3] 5 4 2

dequeue 3: Check 5, Check 4, and Check 2

BFS: 0 2 1 5



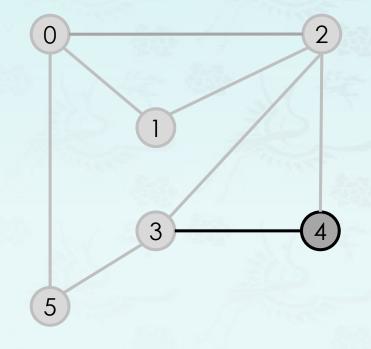


queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
	2	0	1
	3	2	2
	4	2	2
	5	0	1

adj[3] 5 4 2

dequeue 3: Check 5, Check 4, and Check 2

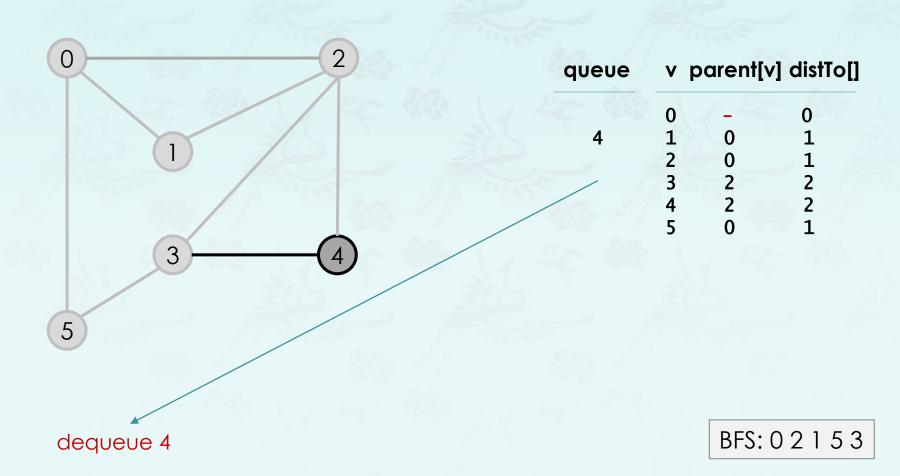
BFS: 0 2 1 5

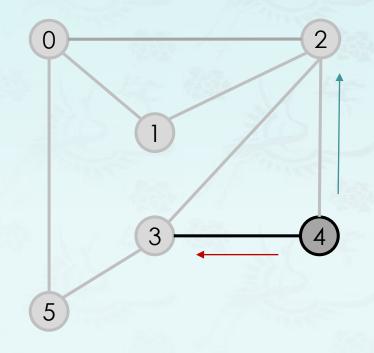


queue	v parent[v] distTo[]		
	0	1-4	0
4	1	0	1
	2	0	1
	3	2	2
	4	2	2
	5	0	1

adj[3] 5 4 2 3 done

BFS: 0 2 1 5 3



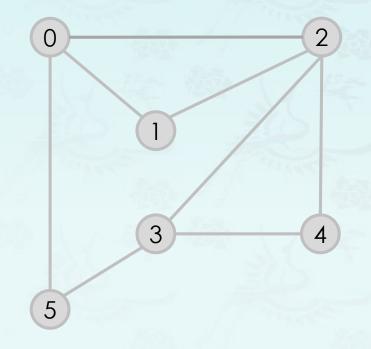


queue	v parent[v] distTo[]		
	0	1-4	0
	1	0	1
	2	0	1
	3	2	2
	4	2	2
	5	0	1

adj[4] 3 2

dequeue 4: Check 3 and Check 2

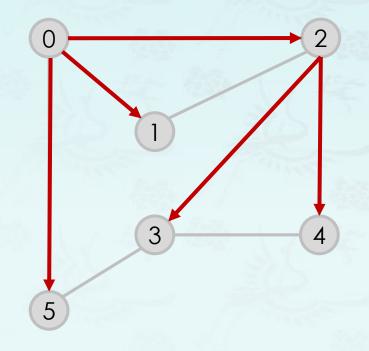
BFS: 0 2 1 5 3



queue	v parent[v] distTo[]		
<b>是一颗</b>	0	1-4	0
	1	0	1
	2	0	1
	3	2	2
	4	2	2
	5	0	1

4 done

BFS: 0 2 1 5 3 4



٧	parent[v]	distTo	[]
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0	1-4	0
1	0	1
2	0	1
3	2	2
4	2	2
5	0	1

done

BFS: 0 2 1 5 3 4

#### **Breadth-first search**

**Depth-first search:** Put unvisited vertices on a **stack**. **Breadth-first search:** Put unvisited vertices on a **queue**.

Shortest path: Find path from s to t that uses fewest number of edges.

## **BFS:** (from source vertex s)

- Put s onto a FIFO queue, and mark s as visited.
- Repeat until the queue is empty:
  - remove the least recently added vertex v
  - add each of v's unvisited neighbors to the queue, and mark them as visited.

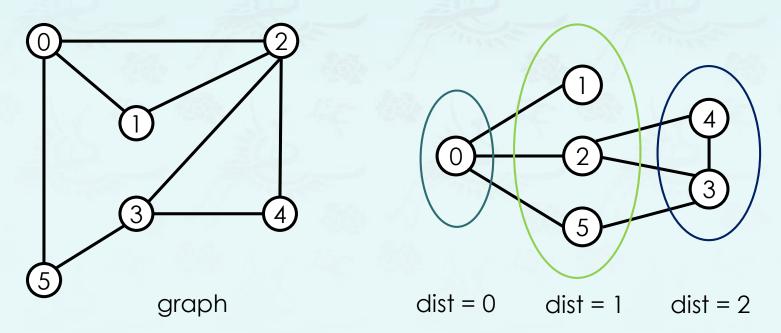
Intuition: BFS examines vertices in increasing distance from s.

# **Breadth-first search properties**

**Proposition:** BFS computes shortest paths (fewest number of edges) from s to all other vertices in a graph in time proportional to E + V.

**Proof:** [correctness] Queue always consists of zero or more vertices of distance k from s, followed by zero or more vertices of distance k + 1.

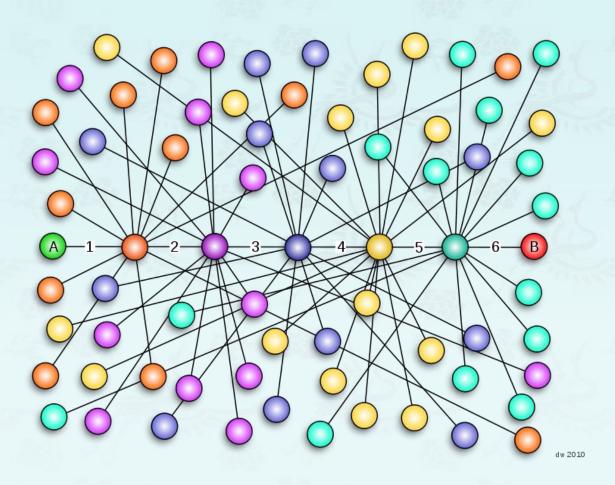
**Proof:** [running time] Each vertex connected to s is visited once.



## Breadth-first search implementation in Java

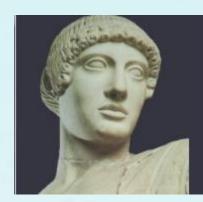
```
// runs BFS at v and produces distTo[] & parentBFS[]
void BFS(graph g, int v) {
 queue<int> que; // to process each vertex
 g->marked[v] = true;
 q \rightarrow distTo[v] = 0;
 que.push(v);
 q.push(v);
 while (!que.empty()) {
   int cur = que.front(); que.pop(); // remove it since processed
   for (gnode w = g->adj[cur].next; w; w = w->next) {
       if (!g->marked[w->item]) {
        que.push(w->item); // queued to process next
        q.push(w->item);
        g->distTo[w->item] = g->distTo[cur] + 1;
        g->parentBFS[w->item] = cur;
        g->marked[w->item] = true;
 g \rightarrow BFSV = q;
                        // save the result at v
 setBFSx(g, v, q);
                          // save the result
```





six degrees of separation?





# THE ORACLE OF BACON



#### About the Oracle of Bacon

This is the most comprehensive version of the Kevin Bacon game on the web. The object of the game is to start with any actor or actress who has been in a movie and connect them to Kevin Bacon in the smallest number of **links** possible. Two people are linked if they've been in a movie together. We do not consider links through television shows, made-for-tv movies, writers, producers, directors, etc. For example, you might wonder how Mary Pickford can be connected to Kevin Bacon. One answer is that:

```
Mary Pickford

was in

Coquette (1929)

with

Louise Beavers

was in

All the Fine Young Cannibals (1960)

with

Robert Wagner (I)

was in

Wild Things (1998)

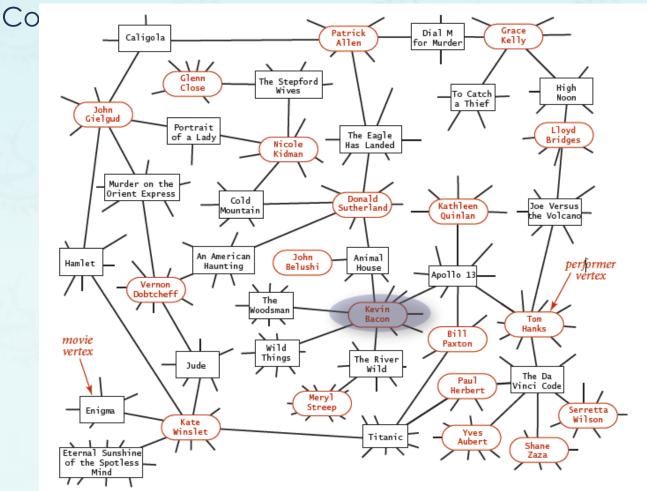
with

Kevin Bacon
```

Prof. Youngsup Kim, idebtor@gmail.com, Data Structures, CSEE Dept., Handong Global University

# **Breadth-first search properties**

- Include one vertex for each performer and one for each movie.
- Connect a movie to all performers that appear in that movie.





http://www.bbc.co.uk/newsbeat/article/35500398/how-facebook-update d-six-degrees-of-separation-its-now-357

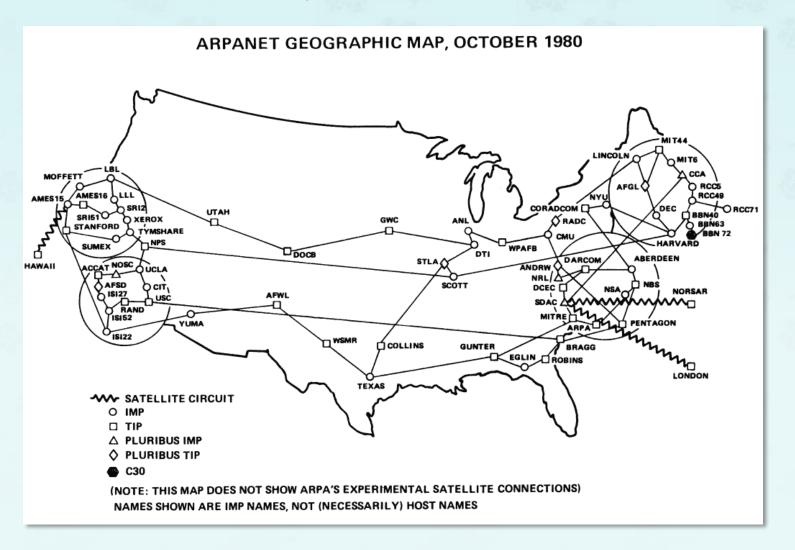


 $2008: 5.28 \rightarrow 2011: 4.74 \rightarrow 2016.2: 3.57$ 

http://www.bbc.co.uk/newsbeat/article/35500398/how-facebook-update d-six-degrees-of-separation-its-now-357

# Breadth-first search application: routing

Fewest number of hops in a communication network.



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#### Major references:

- Fundamentals of Data Structures by Horowitz, Sahni, Anderson-Freed, Algorithms 4<sup>th</sup> edition Part 1 & Part 2 by Robert Sedgewick and Kevin Wayne
- Wikipedia and many resources available from internet