Chapter 16 Solusion

github.com/frc123/CLRS

1/10/2022

16.1

16.1-1

```
void OutputAux(const std::vector< std::vector<int> >& dp_selection,
       int i, int j, std::list<int>& output)
   {
3
       if (dp_selection[i][j] > 0)
       {
            OutputAux(dp_selection, i, dp_selection[i][j], output);
            output.push_back(dp_selection[i][j] - 1);
            OutputAux(dp_selection, dp_selection[i][j], j, output);
       }
   }
10
11
   // assume intervals are sorted by finish time
   std::list<int> DpActivitySelector(const std::vector<Activity>& intervals)
13
   {
14
       int n, i, j, l, k, l_size;
15
       n = (int)(intervals.size());
16
       // dp_size index start by 1
17
       std::vector< std::vector<int> > dp_size(n + 2, std::vector<int>(n + 2, 0)),
18
            dp_selection(n + 2, std::vector<int>(n + 2, -1));
       // compute
       for (1 = 2; 1 \le n + 1; ++1)
21
22
            for (i = 0; i \le n + 1 - 1; ++i)
23
24
                j = i + 1;
25
                for (k = i + 1; k \le j - 1; ++k)
26
                {
27
```

```
if ((i == 0 \mid | intervals[k - 1].s >= intervals[i - 1].f) &&
                          (j == n + 1 \mid \mid intervals[k - 1].f \le intervals[j - 1].s))
29
30
                          l_size = dp_size[i][k] + dp_size[k][j] + 1;
31
                          if (dp_size[i][j] < l_size)</pre>
32
                          {
                               dp_size[i][j] = l_size;
34
                               dp_selection[i][j] = k;
35
                          }
36
                      }
                 }
             }
39
40
        // output
41
        std::list<int> output;
^{42}
        OutputAux(dp_selection, 0, n + 1, output);
43
        return output;
44
   }
45
```

The dynamic-programming algorithm runs in $O(n^3)$.

16.1 - 2

```
// assume intervals are sorted by start time
   std::list<int> GreedyActivitySelector(const std::vector<Activity>& intervals)
   {
3
       int k, m, n;
       std::list<int> activities;
       n = (int)(intervals.size());
       activities.push_front(n - 1);
       k = n - 1;
       for (m = n - 2; m >= 0; --m)
10
            if (intervals[k].s >= intervals[m].f)
11
            {
                activities.push_front(m);
13
                k = m;
14
            }
15
       }
16
       return activities;
   }
18
```

Claim 1. Consider any nonempty subproblem S_k , and let a_m be an activity in S_k with the latest start time. Then a_m is included in some maximum-size subset of mutually compatible activities of S_k .

Proof. Let A_k be a maximum-size subset of mutually compatible activities in S_k , and let a_j be the activity in A_k with the latest start time. If $a_j = a_m$, we are done, since we have shown that a_m is in some maximum-size subset of mutually compatible activities of S_k . If $a_j \neq a_m$, let the set $A'_k = A_k - \{a_j\} \cup \{a_m\}$ be A_k but substituting a_m for a_j . The activities in A'_k are disjoint, which follows because the activities in A_k are disjoint, a_j is the last activity in A_k to start, and $s_m \geq s_j$. Since $|A'_k| = |A_k|$, we conclude that A'_k is a maximum-size subset of mutually compatible activities of S_k , and it includes a_m .

16.1 - 3

1. selecting the compatible activity of least duration

By this approach, the solution will be $\{a_2\}$. However, the optimal solution is $\{a_1, a_3\}$.

2. selecting the compatible activity that overlaps the fewest other remaining activ-

ities

By this approach, the solution will include a_7 . However, the optimal solution is $\{a_1, a_6, a_8, a_{11}\}$, and a_7 is not compatible with the optimal solution.

3. selecting the compatible activity with the earliest start time

By this approach, the solution will be $\{a_3\}$. However, the optimal solution is $\{a_1, a_2\}$.

16.1 - 4

```
struct Element
{
    int interval;
    std::list< std::list<int> >::iterator list;

Element(int interval, std::list< std::list<iint> >::iterator list)
    : interval(interval), list(list) {}
};
```

```
9
   // assume intervals are sorted by finish time
   std::list< std::list<int> > IntervalGraphColoring(const std::vector<Activity>& intervals)
   {
12
       int i, n;
13
       n = (int)(intervals.size());
       std::list< std::list<int> > collection;
       auto heap_cmp = [&intervals](const Element& a, const Element& b) {
16
            return intervals[a.interval].s < intervals[b.interval].s;</pre>
       };
       std::priority_queue<Element, std::vector<Element>, decltype(heap_cmp)> heap(heap_cmp);
       std::list< std::list<int> >::iterator curr_list;
20
       collection.emplace_front();
21
       curr_list = collection.begin();
22
       curr_list->push_front(n - 1);
23
       heap.emplace(n - 1, curr_list);
       for (i = n - 2; i >= 0; --i)
25
       {
26
            if (intervals[i].f <= intervals[heap.top().interval].s)</pre>
27
            {
                curr_list = heap.top().list;
                curr_list->push_front(i);
30
                heap.pop();
31
            }
32
            else
33
            {
                collection.emplace_front();
35
                curr_list = collection.begin();
36
                curr_list->push_front(i);
            }
            heap.emplace(i, curr_list);
40
       return collection;
41
   }
42
16.1-5
    This algorithm is actually a revision from 16.1-1.
   // assume intervals are sorted by finish time
   std::list<int> DpActivitySelector(const std::vector<Activity>& activities)
```

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Thank you very much for starring and contributing

```
{
3
        int n, i, j, l, k, l_size;
        n = (int)(activities.size());
        // dp_size index start by 1
        std::vector< std::vector<int> > dp_size(n + 2, std::vector<int>(n + 2, 0)),
            dp_selection(n + 2, std::vector<int>(n + 2, -1));
        // compute
        for (1 = 2; 1 \le n + 1; ++1)
10
11
            for (i = 0; i \le n + 1 - 1; ++i)
^{12}
            {
                j = i + 1;
14
                for (k = i + 1; k \le j - 1; ++k)
15
16
                     if ((i == 0 || activities[k - 1].s >= activities[i - 1].f) &&
                         (j == n + 1 \mid \mid activities[k - 1].f \le activities[j - 1].s))
18
                     {
19
                         l_size = dp_size[i][k] + dp_size[k][j] + activities[k - 1].v;
20
                         if (dp_size[i][j] < l_size)</pre>
21
                         {
                              dp_size[i][j] = l_size;
                              dp_selection[i][j] = k;
^{24}
                         }
25
                     }
26
                }
27
            }
        }
29
        // output
30
        std::list<int> output;
31
        OutputAux(dp_selection, 0, n + 1, output);
32
        return output;
   }
34
```

Updating...