

# Chapter 15 Solusion

frc6.com

i@frc6.com

<https://github.com/frc123/CLRS-code-solution>

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## 15.1

### 15.1-1

**Proof.** We prove by substitution method. For  $n = 0$ ,  $T(0) = 2^0 = 1$ . For  $n > 0$ ,

$$T(n) = 1 + \sum_{j=0}^{n-1} T(j) = 1 + \sum_{j=0}^{n-1} 2^j = 1 + (2^n - 1) = 2^n$$

□

### 15.1-2

Consider the following case:

length $i$	1	2	3
price $p_i$	1	6	8
density $p_i/i$	1	3	2.67

If we use “greedy” strategy, our solution will be “2 1”, and the total price will be 7. However, the optimal way is “3”, and the total price is 8.

### 15.1-3

```
1      /**
2      * running time:  $O(n^2)$ 
3      *  $p$ : table of prices (index start from 0)
4      *  $n$ : length of rod
5      *  $c$ : cost of each cut
6      * return maximum revenue
7      */
8      int BottomUpCutRodWithCost(const std::vector<int>& p, int n, int c)
9      {
10         int *r, q, i, j;
11         r = new int[n + 1];
12         r[0] = 0;
```

```

13     for (j = 1; j <= n; ++j)
14     {
15         q = p[j - 1];
16         for (i = 0; i < j - 1; ++i)
17             q = std::max(q, p[i] + r[j - i - 1] - c);
18         r[j] = q;
19     }
20     delete[] r;
21     return q;
22 }

```

#### 15.1-4

```

1     /**
2     * p: table of prices (index start from 0)
3     * n: length of rod
4     * r: table of maximum revenue (index start from 1)
5     * s: table of optimal size i of the first piece to cut off (index start from 1)
6     * return maximum revenue
7     */
8     int ExtendedMemoizedCutRodAux(const std::vector<int>& p, int n, int *r, int *s)
9     {
10         int q, i, reminder_r;
11         if (r[n] >= 0) return r[n];
12         q = INT_MIN;
13         for (i = 0; i < n; ++i)
14         {
15             reminder_r = ExtendedMemoizedCutRodAux(p, n - i - 1, r, s);
16             if (q < p[i] + reminder_r)
17             {
18                 q = p[i] + reminder_r;
19                 s[n] = i + 1;
20             }
21         }
22         r[n] = q;
23         return q;
24     }
25
26     /**
27     * running time:  $O(n^2)$ 

```

```

28  * p: table of prices (index start from 0)
29  * n: length of rod
30  * return (r, s)
31  * r: table of maximum revenue (index start from 1)
32  * s: table of optimal size i of the first piece to cut off (index start from 1)
33  * caller is responsible to deallocate return value r and s
34  */
35  std::pair<int*, int*> ExtendedMemoizedCutRod(const std::vector<int>& p, int n)
36  {
37      int *r, *s, i;
38      r = new int[n + 1];
39      s = new int[n + 1];
40      r[0] = 0;
41      s[0] = 0;
42      for (i = 1; i <= n; ++i) r[i] = INT_MIN;
43      ExtendedMemoizedCutRodAux(p, n, r, s);
44      return std::make_pair(r, s);
45  }

```

## 15.1-5

```

1  /**
2  * running time:  $O(n)$ 
3  * n: n-th fibonacci number (must greater than 0)
4  */
5  int FibonacciNumber(int n)
6  {
7      int *f, i, result;
8      f = new int[n + 1];
9      f[0] = 0;
10     f[1] = 1;
11     for (i = 2; i <= n; ++i)
12         f[i] = f[i - 1] + f[i - 2];
13     result = f[n];
14     delete[] f;
15     return result;
16 }

```

## 15.2

### 15.2-1

Optimal parenthesization:  $((1, 2), ((3, 4), (5, 6)))$

Minimum cost: 2010

### 15.2-2

```

1      /**
2      * s: table (2d) storing index of k achieved the optimal cost
3      *      (index start by 1)
4      * caller is responsible to deallocate the return value
5      */
6      Matrix* MatrixChainMultiply
7          (const std::vector<Matrix*>& matrices, const Table* s, int i, int j)
8      {
9          Matrix *matrix_a, *matrix_b, *matrix_c;
10         if (i == j)
11         {
12             return matrices[i - 1];
13         }
14         matrix_a = MatrixChainMultiply(matrices, s, i, (*s)[i][j]);
15         matrix_b = MatrixChainMultiply(matrices, s, (*s)[i][j] + 1, j);
16         matrix_c = MatrixMultiply(matrix_a, matrix_b);
17         if (i != (*s)[i][j]) delete matrix_a;
18         if ((*s)[i][j] + 1 != j) delete matrix_b;
19         return matrix_c;
20     }
```

### 15.2-3

**Proof.** We prove by substitution method. For  $n = 1$ ,  $P(1) = 1 \geq 2^k$  for  $k \leq 0$ . For  $n \geq 2$ ,

$$P(n) = \sum_{k=1}^{n-1} P(k)P(n-k) \geq \sum_{k=1}^{n-1} (c \cdot 2^k)(c \cdot 2^{n-k}) = \sum_{k=1}^{n-1} (c^2 \cdot 2^n) = (n-1)(c^2 \cdot 2^n) \geq c^2 \cdot 2^n$$

for some constant  $c$ . □

### 15.2-4

For all vertices  $v_{i,j}$  in the graph, it contains edge  $(v_{i,j}, v_{i,k})$  and  $(v_{i,j}, v_{k+1,j})$  for all  $i \leq k < j$ .

Vertices:

$$\binom{n}{2} + n = \frac{n(n-1)}{2} + n = \frac{n(n+1)}{2}$$

Edges:

$$\begin{aligned}
\sum_{i=1}^n \sum_{j=i}^n (j-i) &= \sum_{i=1}^n \left( \sum_{j=i}^n (j) - \sum_{j=i}^n (i) \right) = \sum_{i=1}^n \left( \sum_{j=1}^n (j) - \sum_{j=1}^{i-1} (j) - (n-i+1)i \right) \\
&= \sum_{i=1}^n \left( \frac{n(n+1)}{2} - \frac{(i-1)i}{2} - (n-i+1)i \right) = \frac{n^2(n+1)}{2} - \sum_{i=1}^n \left( \frac{i-i^2+2ni}{2} \right) \\
&= \frac{n^2(n+1)}{2} + \frac{1}{2} \sum_{i=1}^n (i^2) - \frac{1}{2} (1+2n) \sum_{i=1}^n (i) = \frac{n^2(n+1)}{2} + \frac{n(n+1)(2n+1)}{12} - \frac{n(n+1)(2n+1)}{4} \\
&= \frac{n^2(n+1)}{2} - \frac{n(n+1)(2n+1)}{6} = \frac{(n-1)n(n+1)}{6}
\end{aligned}$$

### 15.2-5

**Proof.** Notice that  $\sum_{i=1}^n \sum_{j=i}^n R(i, j)$  is equal to the total times of any entries are referenced during the entire call of MATRIX-CHAIN-ORDER. In other words, it is equal to twice the times of line 10 was executed during the entire call.

Hence, we have

$$\begin{aligned}
\sum_{i=1}^n \sum_{j=i}^n R(i, j) &= \sum_{l=2}^n \sum_{i=1}^{n-l+1} \sum_{k=i}^{i+l-2} 2 = 2 \sum_{l=2}^n (n-l+1)(l-1) = 2((n+2) \sum_{l=2}^n l - \sum_{l=2}^n l^2 - (n-1)(n+1)) \\
&= 2((n+2)(\frac{n(n+1)}{2} - 1) - (\frac{n(n+1)(2n+1)}{6} - 1)) = \frac{n^3-n}{3} \quad \square
\end{aligned}$$

### 15.2-6

**Proof.** We prove by induction. Let  $P(n)$  be the claim: A full parenthesization of an  $n$ -element expression has exactly  $n-1$  pairs of parentheses.

(Base Case) A 2-element full parenthesization  $(A_1, A_2)$  has only one pair of parentheses clearly. Hence, we have proved  $P(2)$  is true.

(Induction Step) Suppose that  $P(n)$  is true. Let  $C$  be a sequence with  $n+1$  elements:  $A_1 A_2 \dots A_n A_{n+1}$ . Delete one arbitrary element from  $C$ , we have a sequence with  $n$  elements. By induction hypothesis,  $C$  ( $n$ -element) has exactly  $n-1$  pairs of parentheses now. Add the deleted element back, we can add one pair of parentheses to surround the deleted element and one of the element's neighbor element or one of the element's neighbor parenthesization. This says,  $C$  ( $n+1$ -element) has exactly  $n$  pairs of parentheses now. We have proved  $P(n+1)$  is true.  $\square$

## 15.3

### 15.3-1

RECURSIVE-MATRIX-CHAIN is a more efficient way.

**Proof.** By recurrence (15.6) in section 15.2, there are  $P(n)$  alternative parenthesizations of a sequence of matrices where

$$P(n) = \begin{cases} 1 & \text{if } n = 1. \\ \sum_{k=1}^{n-1} P(k)(n-k) & \text{if } n \geq 2. \end{cases} \quad (15.6)$$

By problem 12-4, we proved that  $P(n) = \Omega(4^n/n^{3/2})$ . This says enumerating takes  $\Omega(4^n/n^{3/2})$  time.

In order to prove that RECURSIVE-MATRIX-CHAIN is a more efficient than enumerating, we just need to prove that RECURSIVE-MATRIX-CHAIN takes  $o(4^n/n^{3/2})$  time.

By recurrence (15.7) in section 15.2, RECURSIVE-MATRIX-CHAIN takes  $T(n)$  times where

$$T(n) = \begin{cases} 1 & \text{if } n = 1. \\ 1 + \sum_{k=1}^{n-1} (T(k) + T(n-k) + 1) & \text{if } n \geq 2. \end{cases}$$

$$T(n) = 2 \sum_{i=1}^{n-1} T(i) + n$$

This says we want to prove that  $T(n) = o(4^n/n^{3/2})$ . Since  $\lim_{n \rightarrow \infty} \frac{4^n/n^{3/2}}{3.5^n} = \infty$ , we just need to prove that  $T(n) = O(3.5^n)$ . ( $T(n) = O(3^n)$  is false, so we try  $T(n) = O(3.5^n)$ .)

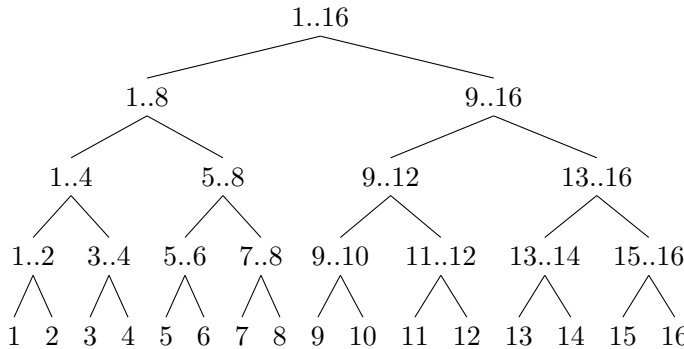
We claim that  $T(n) = O(3.5^n)$  and prove this by substitution method. Let  $c$  be some constant.

Assume  $T(n) \leq c \cdot 3.5^n$ .

$$\begin{aligned} T(n) &= 2 \sum_{i=1}^{n-1} T(i) + n \\ &\leq 2 \sum_{i=1}^{n-1} (c \cdot 3.5^i) + n \\ &= 2c \left( \frac{3.5^n - 1}{3.5 - 1} - 1 \right) + n \\ &= 2c \left( \frac{3.5^n - 3.5}{2.5} \right) + n \\ &= 0.8c \cdot 3.5^n - 2.8c + n \end{aligned}$$

Let  $c = 1$ . We have  $0.8c \cdot 3.5^n - 2.8c + n \leq c \cdot 3.5^n$ . We have proved  $T(n) = O(3.5^n)$ . Hence,  $T(n) = o(4^n/n^{3/2})$ .  $\square$

### 15.3-2



We notice that there is no overlapping subproblem, so memoization does not help to speed up the algorithm.

### 15.3-3

Yes.

**Proof.** Let  $A_{i\dots j}$  denotes sequece of matrices  $A_i A_{i+1} \dots A_j$ .

The subproblems in maximize multiplication are independent. (for more information, refer to page 383) An optimal parenthesization of  $A_{i\dots j}$  that splits the product between  $A_k$  and  $A_{k+1}$  contains within it optimal solutions to the problems of parenthesizing  $A_{i\dots k}$  and  $A_{k+1\dots j}$ .

Given parenthesization  $P_{ij}$  maximize the number of scalar multiplications to  $A_{i\dots j}$ , and  $P_{ij}$  splits the product between  $A_k$  and  $A_{k+1}$ . Making the choice (splits the product between  $A_k$  and  $A_{k+1}$ ) leaves subproblems  $A_{i\dots k}$  and  $A_{k+1\dots j}$  to solve. Let  $P_{ik}$  and  $P_{k+1,j}$  be the parenthesization on  $A_{i\dots k}$  and  $A_{k+1\dots j}$  respectively. We want to prove that  $P_{ik}$  and  $P_{k+1,j}$  maximize the number of scalar multiplications to  $A_{i\dots k}$  and  $A_{k+1\dots j}$  by contradiction. Suppose that  $P_{ik}$  and  $P_{k+1,j}$  does not maximize the number of scalar multiplications to  $A_{i\dots k}$  and  $A_{k+1\dots j}$ . Let  $Q_{ik}$  and  $Q_{k+1,j}$  be the optimal parenthesization (maximize number of multiplications) to  $A_{i\dots k}$  and  $A_{k+1\dots j}$ . Then, by “cutting out”  $P_{ik}$  and  $P_{k+1,j}$  and “pasting in”  $Q_{ik}$  and  $Q_{k+1,j}$ , we get a better solution (more number of multiplications) to the original problem than  $P_{ij}$ . This contradicts to parenthesization  $P_{ij}$  maximize the number of scalar multiplications.  $\square$

### 15.3-4

Consider the following  $p$ 's:

$p_0$	$p_1$	$p_2$	$p_3$
1	10	20	100

By the approach of greedy algorithm, we choose  $k = 1$  for  $[i, j] = [1, 3]$  since  $p_0 p_1 p_3 = 1000$  and  $p_0 p_2 p_3 = 2000$ . Hence, the solution of greedy algorithm is  $(A_1(A_2 A_3))$ .

However,  $((A_1 A_2) A_3)$  ( $k = 2$ ) is the optimal solution, which takes  $p_0 p_1 p_2 + p_0 p_2 p_3 = 200 + 2000 = 2200$  multiplications. Greedy solution  $(A_1(A_2 A_3))$  ( $k = 1$ ) takes  $p_1 p_2 p_3 + p_0 p_1 p_3 = 20000 + 1000 = 21000$  multiplications.

How can we find the counterexample? We start to try to find a counterexample in a sequence with 3 matrices. This says  $[i, j] = [1, 3]$ , and there are two choices for  $k$ : 1 or 2.

The algorithm perform  $m[1, 3]$  times multiplication.

$$m[1, 3] = \begin{cases} m[1, 1] + m[2, 3] + p_0 p_1 p_3 & k = 1 \\ m[1, 2] + m[3, 3] + p_0 p_2 p_3 & k = 2 \end{cases} = \begin{cases} p_1 p_2 p_3 + p_0 p_1 p_3 & k = 1 \\ p_0 p_1 p_2 + p_0 p_2 p_3 & k = 2 \end{cases}$$

We try to make the greedy algorithm choose  $k = 1$ . This says we want  $p_0 p_1 p_3 < p_0 p_2 p_3 \iff p_1 < p_2$ .

In order to make the greedy approach ( $k = 1$ ) yields a suboptimal solution, we want  $k = 2$  to be the optimal approach. This says we want  $p_1 p_2 p_3 + p_0 p_1 p_3 > p_0 p_1 p_2 + p_0 p_2 p_3$ .

Hence, our goal is to find  $p_0, p_1, p_2, p_3$  such that  $p_1 < p_2$  and  $p_1 p_2 p_3 + p_0 p_1 p_3 > p_0 p_1 p_2 + p_0 p_2 p_3$ . We try to let  $p_1 = 10$  and  $p_2 = 20$ . By a sloppy way, we can try to let  $p_3$  much larger than  $p_0$  since  $p_3$  appears twice and  $p_0$  appears once on the LHS, and  $p_0$  appears twice and  $p_3$  appears once on the RHS. We try to let  $p_0 = 1$  and  $p_3 = 100$ . After testing, we find that this is a good counterexample.

### 15.3-5

If we have limit  $l_i$  on the number of pieces of length  $i$  that we are allowed to produce, We can not find the optimal subproblems indenpendently. We show the optimal-substructure property does not hold by providing a counterexample. (Recall optimal substructure on page 374)

Consider the following case:

$i$	1	2	3
$p_i$	5	8	9
$l_i i$	2	2	1

The optimal solution of cutting the rod where  $i = 3$  is lengths 1 and 2 with price  $5+8 = 13$ . However, the optimal solution of cutting the rod where  $i = 2$  is lengths 1 and 1 with price  $5 + 5 = 10$ . We have showed that there is a way of cutting that does not contain in the optimal solution to the original problem but does contain in the optimal solution to the subproblem, which violates the optimal-substructure property.

### 15.3-6

Note: For this question, we assume that  $r_{ij}r_{ji} = 1$  for any  $1 \leq i, j \leq n$ .

**Claim 1.** If  $c_k = 0$  for all  $k = 1, 2, \dots, n$ , then the problem of finding the best sequence of exchanges from currency 1 to currency  $n$  exhibits optimal substructure.

**Proof.** Let the sequence of currencies  $k_1, k_2, k_3, \dots, k_{n-1}, k_n$  be the best sequence to exchange from currency from  $k_1$  to currency  $k_n$ , which means  $r_{k_1 k_2} r_{k_2 k_3} \dots r_{k_{n-1} k_n}$  is maximized. We show that the sequence of currencies  $k_i, k_{i+1}, \dots, k_{j-1}, k_j$  is the best sequence to exchange from currency from  $k_i$  to currency  $k_j$  by contradiction. Assume  $k_i, q_{i+1}, \dots, q_{j-1}, k_j$  is the best sequence to exchange from currency from  $k_i$  to currency  $k_j$ . By using the “cut-and-paste” technique to replace  $k_i, k_{i+1}, \dots, k_{j-1}, k_j$  with  $k_i, q_{i+1}, \dots, q_{j-1}, k_j$ , we get a better sequence of currencies to exchange from currency from  $k_1$  to currency  $k_n$  (e.g.  $k_1, k_2, k_3, \dots, k_i, q_{i+1}, \dots, q_{j-1}, k_j, \dots, k_{n-1}, k_n$ ), which contradicts to  $k_1, k_2, k_3, \dots, k_{n-1}, k_n$  is the best sequence.  $\square$

Now, we show that if  $c_k$  are arbitrary values, then the problem of finding the best sequence of exchanges from currency 1 to currency  $n$  does not necessarily exhibit optimal substructure.

Consider the following case:

We want to exchange from currency from  $k_1$  to currency  $k_4$ .

$r_{k_1 k_2}$	$r_{k_2 k_3}$	$r_{k_3 k_4}$	$r_{k_1 k_3}$	$r_{k_2 k_4}$	$r_{k_1 k_4}$	$c_1$	$c_2$	$c_3$
6	2	5	2	8	10	4	5	20

We try to find the optimal sequence from  $k_1$  to  $k_4$  by list all possible sequences:

$$k_1 k_4: r_{k_1 k_4} - c_1 = 10 - 4 = 6$$

$$k_1 k_2 k_4: r_{k_1 k_2} r_{k_2 k_4} - c_2 = 6 \cdot 8 - 5 = 43$$

$$k_1 k_3 k_4: r_{k_1 k_3} r_{k_3 k_4} - c_2 = 2 \cdot 5 - 5 = 5$$

$$k_1 k_2 k_3 k_4: r_{k_1 k_2} r_{k_2 k_3} r_{k_3 k_4} - c_3 = 6 \cdot 2 \cdot 5 - 20 = 40$$

$$k_1 k_3 k_2 k_4: r_{k_1 k_3} r_{k_3 k_2} r_{k_2 k_4} - c_3 = 2 \cdot \frac{1}{2} \cdot 8 - 20 = -12$$



The optimal sequence from  $k_1$  to  $k_4$  is  $k_1k_2k_4$ .

Now, we try to show that  $k_2k_4$  is not the optimal sequence from  $k_2$  to  $k_4$  by list all possible sequences:

$$k_2k_4: r_{k_2k_4} - c_1 = 8 - 4 = 4$$

$$k_2k_3k_4: r_{k_2k_3}r_{k_3k_4} - c_2 = 2 \cdot 5 - 5 = 5$$

$k_2k_1k_4$ : (unnecessary)

$k_2k_1k_3k_4$ : (unnecessary)

$k_2k_3k_1k_4$ : (unnecessary)

It is unnecessary to solve results for  $k_2k_1k_4$ ,  $k_2k_1k_3k_4$ , and  $k_2k_3k_1k_4$  since we already find that  $k_2k_3k_4$  is a better sequence than  $k_4$  is  $k_2k_4$ .

## 15.4

### 15.4-1

$\langle 1, 0, 0, 1, 1, 0 \rangle$

(the solution is not unique)

### 15.4-2

```

1     template <typename T>
2     std::list<T> OutputLCS(LengthTable length_table,
3         const std::vector<T>& sequence_x, const std::vector<T>& sequence_y)
4     {
5         size_t it_x, it_y;
6         std::list<T> common_sequence;
7         it_x = sequence_x.size();
8         it_y = sequence_y.size();
9         while (it_x > 0 && it_y > 0)
10        {
11            const T& element_x = sequence_x[it_x - 1];
12            const T& element_y = sequence_y[it_y - 1];
13            if (element_x == element_y)
14            {
15                common_sequence.push_front(element_y);
16                --it_x;
17                --it_y;
18            }
19            else if (length_table[it_x - 1][it_y] >= length_table[it_x][it_y - 1])
20            {
21                --it_x;
22            }

```

```

23         else
24         {
25             --it_y;
26         }
27     }
28     return common_sequence;
29 }

```

### 15.4-3

```

1     template <typename T>
2     int RecursiveLCSLengthMemoized
3         (const std::vector<T>& sequence_x, const std::vector<T>& sequence_y,
4          size_t it_x, size_t it_y,
5          LengthTable& length_table, PointerTable& pointer_table)
6     {
7         int sequence_length_x_prev, sequence_length_y_prev;
8         if (length_table[it_x][it_y] >= 0)
9             return length_table[it_x][it_y];
10        if (sequence_x[it_x - 1] == sequence_y[it_y - 1])
11        {
12            length_table[it_x][it_y] = RecursiveLCSLengthMemoized(sequence_x, sequence_y,
13                it_x - 1, it_y - 1, length_table, pointer_table) + 1;
14            pointer_table[it_x][it_y] = PointerType::X_Y_PREV;
15        }
16        else
17        {
18            sequence_length_x_prev = RecursiveLCSLengthMemoized(sequence_x, sequence_y,
19                it_x - 1, it_y, length_table, pointer_table);
20            sequence_length_y_prev = RecursiveLCSLengthMemoized(sequence_x, sequence_y,
21                it_x, it_y - 1, length_table, pointer_table);
22            if (sequence_length_x_prev >= sequence_length_y_prev)
23            {
24                length_table[it_x][it_y] = sequence_length_x_prev;
25                pointer_table[it_x][it_y] = PointerType::X_PREV;
26            }
27            else
28            {
29                length_table[it_x][it_y] = sequence_length_y_prev;
30                pointer_table[it_x][it_y] = PointerType::Y_PREV;

```

```

31         }
32     }
33     return length_table[it_x][it_y];
34 }
35
36 template <typename T>
37 std::pair<LengthTable, PointerTable> LCSLengthMemoized
38     (const std::vector<T>& sequence_x, const std::vector<T>& sequence_y)
39 {
40     size_t size_x, size_y, it_x, it_y,
41         sequence_length_x_prev, sequence_length_y_prev;
42     size_x = sequence_x.size();
43     size_y = sequence_y.size();
44     // note this line init all elements in length_table with -1
45     LengthTable length_table(size_x + 1, LengthTableRow(size_y + 1, -1));
46     PointerTable pointer_table(size_x + 1, PointerTableRow(size_y + 1, NIL));
47     for (it_x = 0; it_x <= size_x; ++it_x)
48         length_table[it_x][0] = 0;
49     for (it_y = 1; it_y <= size_y; ++it_y)
50         length_table[0][it_y] = 0;
51     RecursiveLCSLengthMemoized(sequence_x, sequence_y, size_x, size_y,
52         length_table, pointer_table);
53     return std::make_pair(std::move(length_table), std::move(pointer_table));
54 }

```

#### 15.4-4

Solution for space of  $2 \cdot \min(m, n) + O(1)$ :

```

1     template <typename T>
2     std::pair<int, PointerTable> LCSLengthByTwoRowLengthTable
3         (const std::vector<T>& sequence_x, const std::vector<T>& sequence_y)
4     {
5         size_t size_x, size_y, it_x, it_y;
6         int sequence_length_x_prev, sequence_length_y_prev;
7         size_x = sequence_x.size();
8         size_y = sequence_y.size();
9         if (size_x < size_y)
10             return LCSLengthByTwoRowLengthTable(sequence_y, sequence_x);
11         LengthTableRow length_table_row_1(size_y + 1), length_table_row_2(size_y + 1);
12         PointerTable pointer_table(size_x + 1, PointerTableRow(size_y + 1));

```

```

13     for (it_y = 0; it_y <= size_y; ++it_y)
14         length_table_row_1[it_y] = 0;
15     for (it_x = 1; it_x <= size_x; ++it_x)
16     {
17         for (it_y = 1; it_y <= size_y; ++it_y)
18         {
19             if (sequence_x[it_x - 1] == sequence_y[it_y - 1])
20             {
21                 length_table_row_2[it_y] =
22                     length_table_row_1[it_y - 1] + 1;
23                 pointer_table[it_x][it_y] = PointerType::X_Y_PREV;
24             }
25             else
26             {
27                 sequence_length_x_prev = length_table_row_1[it_y];
28                 sequence_length_y_prev = length_table_row_2[it_y - 1];
29                 if (sequence_length_x_prev >= sequence_length_y_prev)
30                 {
31                     length_table_row_2[it_y] = sequence_length_x_prev;
32                     pointer_table[it_x][it_y] = PointerType::X_PREV;
33                 }
34                 else
35                 {
36                     length_table_row_2[it_y] = sequence_length_y_prev;
37                     pointer_table[it_x][it_y] = PointerType::Y_PREV;
38                 }
39             }
40         }
41         length_table_row_1.swap(length_table_row_2);
42     }
43     return std::make_pair(length_table_row_2[size_y], std::move(pointer_table));
44 }

```

Solution for space of  $\min(m, n) + O(1)$ :

```

1     template <typename T>
2     std::pair<int, PointerTable> LCSLengthByOneRowLengthTable
3         (const std::vector<T>& sequence_x, const std::vector<T>& sequence_y)
4     {
5         size_t size_x, size_y, it_x, it_y;
6         int sequence_length_x_y_prev, sequence_length_x_y_now;

```

```

7      size_x = sequence_x.size();
8      size_y = sequence_y.size();
9      if (size_x < size_y)
10         return LCSLengthByOneRowLengthTable(sequence_y, sequence_x);
11      LengthTableRow length_table_row(size_y + 1);
12      PointerTable pointer_table(size_x + 1, PointerTableRow(size_y + 1));
13      for (it_y = 0; it_y <= size_y; ++it_y)
14         length_table_row[it_y] = 0;
15      for (it_x = 1; it_x <= size_x; ++it_x)
16      {
17         sequence_length_x_y_prev = 0;
18         for (it_y = 1; it_y <= size_y; ++it_y)
19         {
20             sequence_length_x_y_now = length_table_row[it_y];
21             if (sequence_x[it_x - 1] == sequence_y[it_y - 1])
22             {
23                 length_table_row[it_y] =
24                     sequence_length_x_y_prev + 1;
25                 pointer_table[it_x][it_y] = PointerType::X_Y_PREV;
26             }
27             else
28             {
29                 if (length_table_row[it_y] >= length_table_row[it_y - 1])
30                 {
31                     // length_table_row[it_y] = length_table_row[it_y];
32                     pointer_table[it_x][it_y] = PointerType::X_PREV;
33                 }
34                 else
35                 {
36                     length_table_row[it_y] = length_table_row[it_y - 1];
37                     pointer_table[it_x][it_y] = PointerType::Y_PREV;
38                 }
39             }
40             sequence_length_x_y_prev = sequence_length_x_y_now;
41         }
42     }
43     return std::make_pair(length_table_row[size_y], std::move(pointer_table));
44 }

```

## 15.4-5

```

1      // O(n^2)
2      std::vector<int> LongestIncreasingSubsequence
3          (const std::vector<int>& sequence)
4      {
5          size_t i, j, size, sequence_length_it, max_length_index;
6          size = sequence.size();
7          std::vector<size_t> length(size), prev(size);
8          // compute
9          max_length_index = 0;
10         for (i = 0; i < size; ++i)
11         {
12             length[i] = 1;
13             for (j = 0; j < i; ++j)
14             {
15                 if (sequence[j] < sequence[i])
16                 {
17                     sequence_length_it = length[j] + 1;
18                     if (sequence_length_it > length[i])
19                     {
20                         length[i] = sequence_length_it;
21                         prev[i] = j;
22                     }
23                 }
24             }
25             if (length[i] > length[max_length_index])
26                 max_length_index = i;
27         }
28         // output
29         std::vector<int> result(length[max_length_index]);
30         for (i = length[max_length_index] - 1; i >= 0; --i)
31         {
32             result[i] = sequence[max_length_index];
33             max_length_index = prev[max_length_index];
34         }
35         return result;
36     }

```

We also can sort the sequence and find the longest common sequence between the sequence and the sorted sequence.

## 15.4-6

```

1      // O(nlgn)
2      std::vector<int> LongestIncreasingSubsequenceBinarySerach
3          (const std::vector<int>& sequence)
4      {
5          size_t i, size, seq_index;
6          int lower, upper, middle;
7          size = sequence.size();
8          /**
9           * prev
10          * index is the element index of sequence
11          * value is the element index of sequence
12          *      such that the element is the prev element
13          *      in the increasing subsequence
14          * -----
15          * sub_seq
16          * index is length of the increasing subsequence
17          * value is the element index of sequence
18          */
19          std::vector<size_t> prev(size), sub_seq;
20          sub_seq.reserve(size);
21          // compute
22          for (i = 0; i < size; ++i)
23          {
24              lower = 0;
25              upper = sub_seq.size() - 1;
26              while (lower <= upper)
27              {
28                  middle = lower + ((upper - lower) >> 1);
29                  if (sequence[sub_seq[middle]] > sequence[i])
30                  {
31                      upper = middle - 1;
32                  }
33                  else if (sequence[sub_seq[middle]] < sequence[i])
34                  {
35                      lower = middle + 1;
36                  }
37                  else
38                  {

```

```

39         lower = middle;
40         break;
41     }
42 }
43 // sequence[i] <= sequence[sub_seq[lower]] must be true
44 if (lower >= sub_seq.size())
45     sub_seq.push_back(i);
46 else
47     sub_seq[lower] = i;
48 if (lower > 0)
49     prev[i] = sub_seq[lower - 1];
50 }
51 // output
52 size = sub_seq.size();
53 std::vector<int> result(size);
54 seq_index = sub_seq[size - 1];
55 for (i = size - 1; i < size; --i)
56 {
57     result[i] = sequence[seq_index];
58     seq_index = prev[seq_index];
59 }
60 return result;
61 }

```

## 15.5

### 15.5-1

```

1  void ConstructOptimalBST(std::ostream& os, const Table& root,
2      const std::string& suffix, size_t i, size_t j)
3  {
4      if (j == i - 1)
5      {
6          os << "d_" << j << " is the " << suffix << "\n";
7      }
8      else
9      {
10         size_t r = root[i][j];
11         os << "k_" << r << " is the " << suffix << "\n";
12         ConstructOptimalBST(os, root,

```



```

13         "left child of k_" + std::to_string(r), i, r - 1);
14         ConstructOptimalBST(os, root,
15         "right child of k_" + std::to_string(r), r + 1, j);
16     }
17 }
18
19 void ConstructOptimalBST(std::ostream& os, const Table& root)
20 {
21     ConstructOptimalBST(os, root, "root", 1, root.size() - 1);
22 }

```

## 15.5-2

cost: 3.12

structure:

$k_5$  is the root

$k_2$  is the left child of  $k_5$

$k_1$  is the left child of  $k_2$

$d_0$  is the left child of  $k_1$

$d_1$  is the right child of  $k_1$

$k_3$  is the right child of  $k_2$

$d_2$  is the left child of  $k_3$

$k_4$  is the right child of  $k_3$

$d_3$  is the left child of  $k_4$

$d_4$  is the right child of  $k_4$

$k_7$  is the right child of  $k_5$

$k_6$  is the left child of  $k_7$

$d_5$  is the left child of  $k_6$

$d_6$  is the right child of  $k_6$

$d_7$  is the right child of  $k_7$

## 15.5-3

The OPTIMAL-BST would still take  $\Theta(n^3)$ .

According to equation (15.12), it takes  $\Theta(j - i)$  to compute  $w(i, j)$  each time. Hence the total time contributed by line 9 in OPTIMAL-BST is

$$\begin{aligned}
\sum_{l=1}^n \sum_{i=1}^{n-l+1} (c \cdot (j-i)) &= c \cdot \sum_{l=1}^n \sum_{i=1}^{n-l+1} ((i+l-1) - i) \\
&= c \cdot \sum_{l=1}^n (n - (l-1))(l-1) \\
&= c \cdot \sum_{z=0}^{n-1} (n-z)z \\
&= c \cdot \sum_{z=0}^{n-1} (n-z)z \\
&= \Theta(n^3)
\end{aligned}$$

#### 15.5-4

```

1      // O(n^2)
2      std::pair<Table, Table> OptimalBST(const Row& p, const Row& q)
3      {
4          size_t q_size, i, j, l, r;
5          int t;
6          q_size = q.size(); // n + 1
7          Table e(q_size + 1, Row(q_size)), w(q_size + 1, Row(q_size)),
8              root(q_size, Row(q_size));
9          for (i = 1; i <= q_size; ++i) // 1 to (n + 1)
10         {
11             e[i][i - 1] = q[i - 1];
12             // w[i][i - 1] = q[i - 1];
13         }
14         // l = 1
15         for (i = 1; i < q_size; ++i) // 1 to n
16         {
17             // w[i][i] = w[i][i - 1] + p[i] + q[i];
18             w[i][i] = q[i - 1] + p[i] + q[i];
19             // e[i][i] = e[i][i - 1] + e[i + 1][i] + w[i][i];
20             e[i][i] = q[i - 1] + q[i] + w[i][i];
21             root[i][i] = i;
22         }
23         // l = 2 ... n
24         for (l = 2; l < q_size; ++l) // 2 to n
25         {
26             for (i = 1; i <= q_size - l; ++i) // 1 to (n - l + 1)
27             {

```

```

28         j = i + 1 - 1;
29         e[i][j] = INT_MAX;
30         // root[i][j - 1] to root[i + 1][j]
31         for (r = root[i][j - 1]; r <= root[i + 1][j]; ++r)
32         {
33             t = e[i][r - 1] + e[r + 1][j];
34             if (t < e[i][j])
35             {
36                 e[i][j] = t;
37                 root[i][j] = r;
38             }
39         }
40         w[i][j] = w[i][j - 1] + p[j] + q[j];
41         e[i][j] += w[i][j];
42     }
43 }
44 return std::make_pair(std::move(e), std::move(root));
45 }

```

The main principle is to reduce the loop range of line 10 of OPTIMAL-BST from  $i \dots j$  to  $root[i, j - 1] \dots root[i + 1, j]$ . This says modify line 10 of OPTIMAL-BST to

**for**  $r = root[i, j - 1]$  **to**  $root[i + 1, j]$

Note the modification of case of  $l = 1$  is skipped here. We primary consider  $l \geq 2$ .

**Claim 2.** OPTIMAL-BST run in  $\Theta(n^2)$  time

**Proof.** Denote the **for** loop of line 5-14 as  $L_1$ , and denote the **for** loop of line 6-14 as  $L_2$ . In order to show OPTIMAL-BST run in  $\Theta(n^2)$  time, we just need to show each iteration of  $L_1$  takes  $\Theta(n)$  time.

Let  $t(l, i)$  be the running time of each iteration of  $L_2$  with variable  $l$  and  $i$ .

$$\begin{aligned}
 t(l, i) &= r[i + 1, j] - r[i, j - 1] + 1 \\
 &= r[i + 1, i + l - 1] - r[i, i + l - 2] + 1
 \end{aligned}$$

Each iteration of  $L_1$  takes  $\sum_{i=1}^{n-l+1} t(l, i)$  time. Notice  $1 \leq root[a, b] \leq n$  for all  $a, b$ . Also, by Knuth [212],  $root[a + k, b - 1 + k] \leq root[a + h, b - 1 + h]$  for all  $k \leq h$ . Hence we have  $root[a + h, b - 1 + h] - root[a + k, b - 1 + k] \leq n$  for all  $k \leq h$ .

$$\begin{aligned}
\sum_{l=1}^n \sum_{i=1}^{n-l+1} t(l, i) &= \sum_{l=1}^n \sum_{i=1}^{n-l+1} (r[i+1, i+l-1] - r[i, i+l-2] + 1) \\
&= \sum_{l=1}^n (r[n-l+2, n] - r[1, l-1] + (n-l+1))
\end{aligned}$$

We claim  $r[n-l+2, n] - r[1, l-1] \leq n$  by setting  $a = 1$ ,  $b = l$ ,  $k = 0$ , and  $h = n-l+1$ . Hence

$$\sum_{l=1}^n \sum_{i=1}^{n-l+1} t(l, i) = \Theta(n^2).$$

□