Discussion 2

1. Review: Polynomials, Exponentials, Logarithms, and their Derivatives

a. Polynomials

$p(n) = \sum_{i=0}^{d} a_i n^i$ is called a polynomial in n of degree d where $a_d \neq 0$.

A polynomial is asymptotically positive if and only if $a_d > 0$, and we have $p(n) = \Theta(n^d)$.

b. Exponentials

a^n is the basic form of an exponential, where a is called base.

$(a^n)^m = a^{mn} = (a^m)^n$.

For all real constants a and b such that $a > 1$,

$$\lim_{x \to \infty} \frac{n^b}{a^n} = 0$$

from which we can conclude that any exponential function with a base strictly greater than 1 grows faster than any polynomial function.

c. Logarithms

$\lg n = \log_2 n$; $\ln n = \log_e n$ (natural logarithm).

For all real $a > 0$, $b > 0$, $c > 0$, and n, we have the following properties:

i. $a = b^{\log_b a}$

ii. $\log_c(ab) = \log_c a + \log_c b$

iii. $\log_b a^n = n \log_b a$

iv. $\log_b a = \frac{\log_c a}{\log_c b}$

d. Derivatives

$f(x) = ax^n$, $\dfrac{\mathrm{d}f}{\mathrm{d}x} = anx^{n-1}$.

$f(x) = a^x$, $\dfrac{\mathrm{d}f}{\mathrm{d}x} = a^x \ln a$.

$f(x) = \ln x$, $\dfrac{\mathrm{d}f}{\mathrm{d}x} = \dfrac{1}{x}$, $x > 0$; $f(x) = \log_a x$, $\dfrac{\mathrm{d}f}{\mathrm{d}x} = \dfrac{1}{x \ln a}$, $x > 0$.

2. Asymptotic Notation: O, Ω, Θ

Example: We have $f(n) = \lg n$, $g(n) = \ln n$, what is their relation represented in asymptotic notation? There are two ways.

Discussion 2

a. By definition

b. By using limits (L'Hopital's rule)

3. Solving Recurrence Relation

Recall the recurrence relation of merge sort: $T(n) = 2T(\frac{n}{2}) + n - 1, n >= 2$.

a. Substitution Method

b. Master Theorem

Followup: Which one should we use?

1. Substitution always works.
2. Master Theorem works only when $T(n) = aT(\frac{n}{b}) + f(n)$. Think about $T(n) = 2T(n - 1) + n$.

Discussion 2

4. Divide and Conquer: n-Queens Problem

You are asked to place queens so that no two queens are in the same row, column or diagonal on a $n - by - n$ chessboard. Figure 1 is an example.

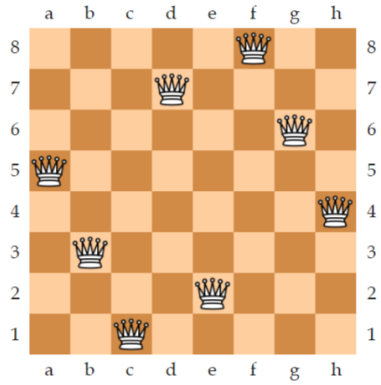

Figure 1: A solution for 8-Queens

Devise an algorithm which takes n as input and outputs all distinct placements of n queens on an $n - by - n$ chessboard. Analyze its running time.

Hint: If the first queen is placed at (i, j), where can't any of the other queens be placed?