

## 1. Graph Basics

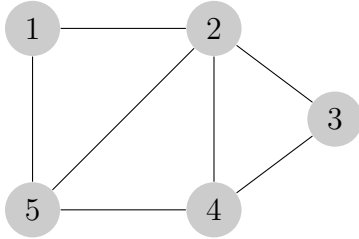


Figure 1: An undirected graph  $G_1$  with 5 vertices and 7 edges

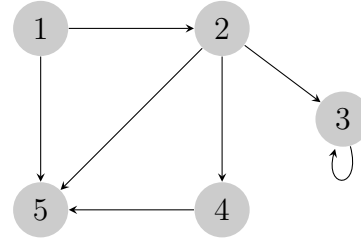
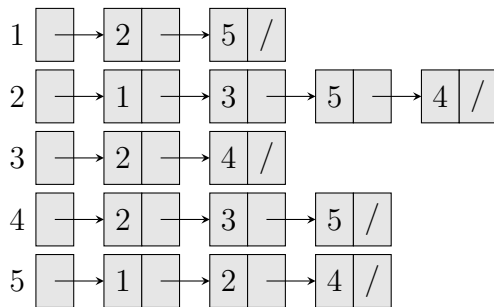


Figure 2: A directed graph  $G_2$  with 5 vertices and 7 edges.

### a. Representation: Adjacency List, Adjacency Matrix

What is the adjacency matrix and adjacency list of  $G_1$  and  $G_2$ , respectively?



	1	2	3	4	5
1	0	1	0	0	0
2	1	0	1	1	1
3	0	1	0	1	0
4	0	1	1	0	1
5	1	1	0	1	0

Figure 3: Adjacency list and adjacency matrix representation of  $G_1$

How about  $G_2$ ? It's your turn now!

1. Adjacency matrix needs  $\Theta(|V| \times |E|)$  storage space while adjacency list requires  $\Theta(|V| + |E|)$ .
2. If  $G$  is undirected, its adjacency matrix  $A$  is symmetric. Namely,  $A^T = A$ . Further, the main diagonal entries of  $A$  are all zeros.
3. **Self-loops**—edges from a vertex to itself—are possible in a directed graph, but are forbidden in an undirected graph.

b. Degree

1.  $\sum_{u \in V} \text{degree}(u) = 2|E|$ , where  $G$  is an undirected graph.
2.  $\sum_{u \in V} \text{out-degree}(u) = \sum_{u \in V} \text{in-degree}(u) = |E|$ , where  $G$  is a directed graph.
3.  $\text{degree}(u) = \text{out-degree}(u) + \text{in-degree}(u)$ , where  $u \in V$  and  $G$  is a directed graph.
4. A vertex whose degree is 0 is **isolated**.

c. Path, Connected Component

1. A **path** of length  $k$  from a vertex  $u$  to a vertex  $u'$  in a graph  $G = \langle V, E \rangle$  is a sequence  $\langle v_0, v_1, v_2, \dots, v_k \rangle$  of vertices such that  $u = v_0$ ,  $u' = v_k$ .
2. An undirected graph is **connected** if every vertex is reachable from all other vertices.
3. A directed graph is **strongly connected** if every two vertices are reachable from each other.

## 2. BFS and DFS

BFS( $G, s$ )

```

1  // G: input graph; s: source vertex
2  for each vertex  $u \in V - \{s\}$ 
3       $d[u] = +\text{INFTY}$ 
4   $d[s] = 0$ 
5
6  // create FIFO queue
7   $Q = \text{EMPTY}$ 
8  ENQUEUE( $G, s$ )
9  while  $Q$  not EMPTY
10      $u = \text{DEQUEUE}(G)$ 
11     for each  $v \in \text{Adj}[u]$ 
12         if  $d[v] = +\text{INFTY}$ 
13              $d[v] = d[u] + 1$ 
14             ENQUEUE( $G, v$ )
15 return  $d$ 
```

DFS( $G$ )

```
1  for each vertex  $u \in V$ 
2       $u.color = \text{WHITE}$ 
3   $time = 0$ 
4  for each vertex  $u \in V$ 
5      if  $u.color = \text{WHITE}$ 
6          DFS-VISIT( $u$ )
```

DFS-VISIT( $u$ )

```
1  // white vertex  $u$  has just been discovered
2   $time = time + 1$ 
3   $u.discover = time$ 
4   $u.color = \text{GRAY}$ 
5
6  // explore edge  $(u, v)$ 
7  for each vertex  $u \in Adj[u]$ 
8      if  $u.color = \text{WHITE}$ 
9          DFS-VISIT( $u$ )
10
11 // blacken  $u$ , it is finished
12  $u.color = \text{BLACK}$ 
13  $time = time + 1$ 
14  $u.finish = time$ 
```