

## 1. Mathematical Induction

**Example:** Prove  $T(n) = \lg n + 1$  is the solution of the recurrence relation:

$$T(n) = \begin{cases} 1 & \text{if } n = 1 \\ T(\frac{n}{2}) + 1 & \text{otherwise} \end{cases}$$

Assume that  $n = 2^k$ ,  $k \geq 0$ .

1. Base Case:

2. Inductive Hypothesis:

3. Inductive Step:

4. Conclusion: We can prove by induction that ...

**2a. Write pseudocode to compute the  $n$ -th Fibonacci number using recursion.**

FIB( $n$ )

```
1  // Base Case
2
3
4  // Recursive Call
5
6
```

2b. Draw a recursion tree for  $n = 4$ .

### 3. Running Time Analysis: Merge Sort

MERGE-SORT( $A, i, j$ )

```
1  if  $i < j$ 
2    then  $k \leftarrow (i + j)/2$ 
3         MERGE-SORT( $A, i, k$ )
4         MERGE-SORT( $A, k + 1, j$ )
5         MERGE( $A, i, k, j$ )
```

Let us denote  $T(n)$  as the maximum number of comparisons of merge sort on a list of size  $n = j - i + 1$ . What is the recurrence relation of  $T(n)$ ?

Likewise, We can construct a recursion tree for merge sort.