



ORACLE®

ASH Outliers: Detecting Unusual Events in Active Session History

John Beresniewicz, Oracle USA

ORACLE®



Agenda

- ASH Fundamentals
- Wait Events and ASH
- Event Histograms
- Event Significance Levels
- SQL to Find Outliers in ASH
- Handling the GV\$ problem



Motivating Use Case

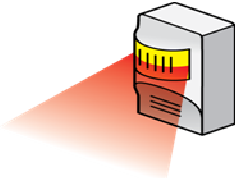
- Unusually long wait events (outliers) suspected to trigger cascade-effect performance incidents
- RAC performance experts claim EM Top Activity not helpful as aggregation masks outlier events
- Can we see if ASH has sampled any such events?
 - If none observed does not mean they have not happened
 - However ASH sampling is biased to longer events

ASH Fundamentals



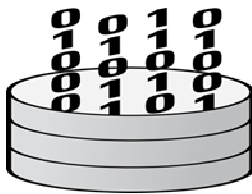
All 'Active' sessions captured every second

- Foregrounds and backgrounds are sampled
- Active foregrounds contribute to DB Time



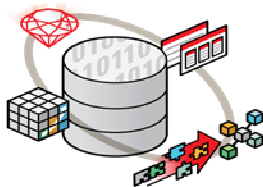
In-memory: *V\$ACTIVE_SESSION_HISTORY*

- Sampling interval = 1 second
- Circular SGA buffer with latchless query access



On-disk: *DBA_HIST_ACTIVE_SESS_HISTORY*

- Sub-sampling interval = 10 seconds



ASH is a system-wide record of database activity

- A **FACT** table with multiple dimensions that help diagnose performance issues

ASH and DB Time

Active sessions contribute to DB Time

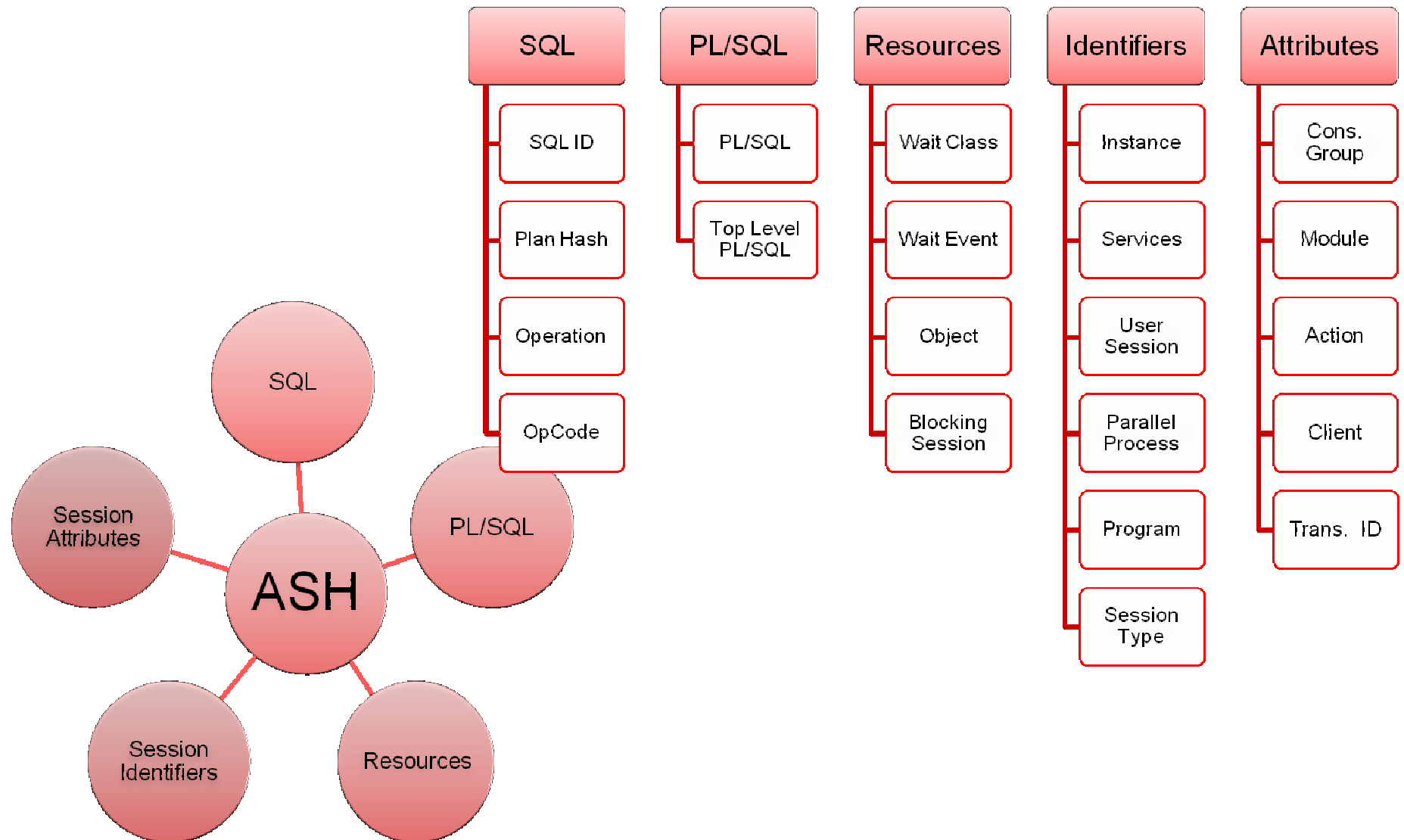
ASH samples active sessions

ASH Math = estimate DB Time by counting
ASH samples

COUNT of ASH Samples = Total DB Time in
seconds for that time interval

Group by over 70+ performance dimensions

ASH Fact Table Dimensions



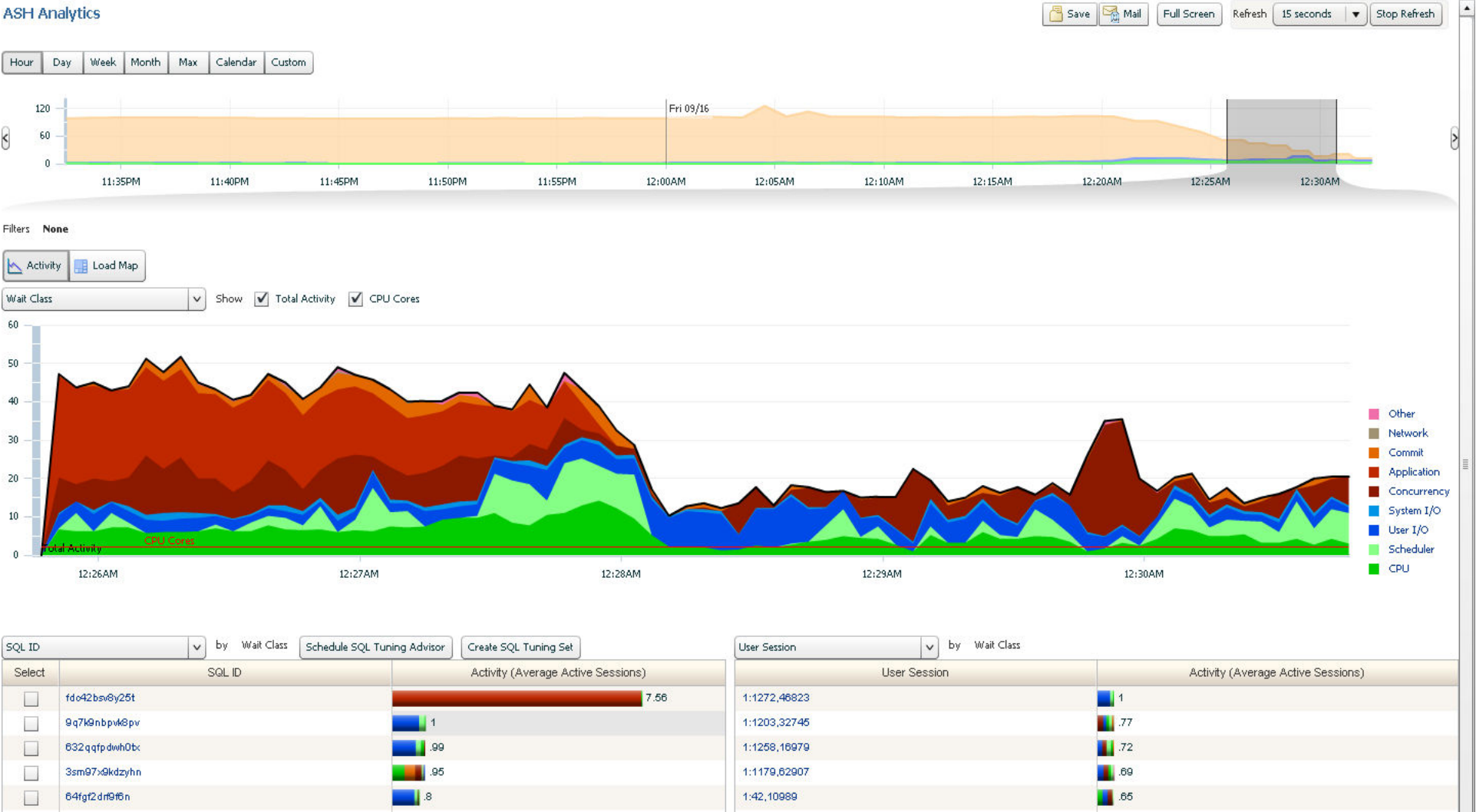


ASH Math

- $\text{COUNT}(\ast) = \text{DB TIME (secs)}$
 - Basic in-memory (V\$ASH) formula
- The Kurtz Construct is nice
 - $\text{SUM}(1) = \text{DB Time (secs)}$ for in-memory
 - $\text{SUM}(10) = \text{DB TIME (secs)}$ for on-disk
- DB Time Method analysis:
 - Dimensional GROUP BY over $\text{COUNT}(\ast)$

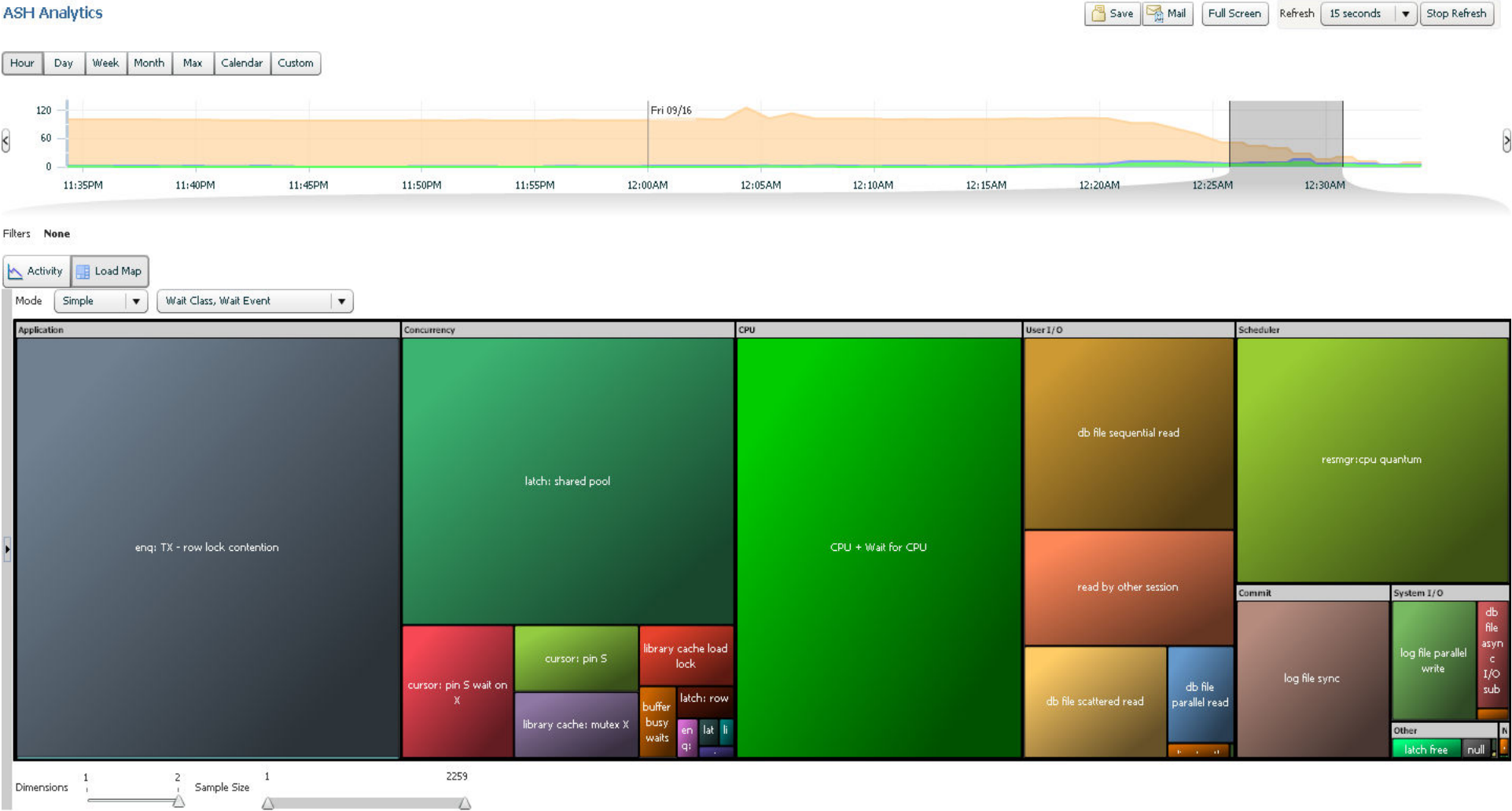


Good ASH Math: ASH Analytics





Group by Wait Event within Wait Class

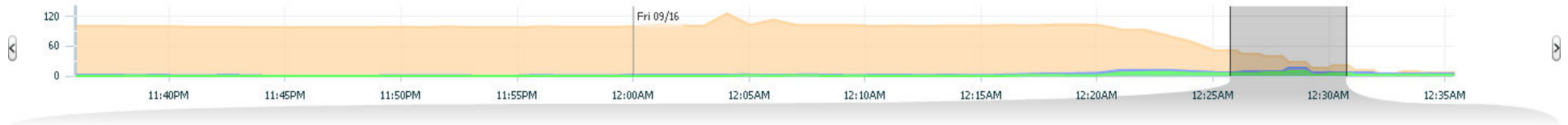


Add SQL_ID Grouping Dimension

ASH Analytics

Save Mail Full Screen Refresh 15 seconds Stop Refresh

Hour Day Week Month Max Calendar Custom



Filters None

Activity Load Map

Mode Advanced Wait Class Wait Event SQL ID



Dimensions 1 2 3 Sample Size 1 2258

ORACLE

Bad ASH Math

- SQL observed using 9 secs of CPU every 10 secs

ORACLE Enterprise Manager Cloud Control 12c

Enterprise Targets Favorites History

lin22.us.oracle.com

Oracle Database Performance Availability Schema Administration

Top Activity > SQL Details: 9q7k9nbpvk8pv

SQL Details: 9q7k9nbpvk8pv

Switch to SQL ID Go View Data Real Time: Manual Refresh Refresh SQL Worl

Text

```
SELECT NVL(SUM(TIME_WAITED/1000), 0)
FROM SYS.DBA_HIST_ACTIVE_SESS_HISTORY S, SYS.WRH$_SEG_STAT_OBJ SO, SYS.WRM$_SNAPSHOT SN, SYS.OBJ$ OF
WHERE SN.SNAP_ID = S.SNAP_ID AND SN.DBID = S.DBID AND SN.BEGIN_INTERVAL_TIME >= (CURRENT_TIMESTAMP -
AND S.WAIT_CLASS = 'User I/O' AND S.WAIT_TIME = 0 AND SO.DATAOBJ# = :B1 AND OBJ.OBJ# = SO.OBJ# AND OF
```

Details

Select the plan hash value to see the details below. Plan Hash Value 2106978214

Statistics Activity Plan Plan Control Tuning History SQL Monitoring

Summary

ORACLE



Wait Events and ASH

- ASH samples actively waiting sessions
 - Wait time is unknown at sample time
 - The “fix up” writes actual wait time into TIME_WAITED
- ASH is a biased sampler of wait events
 - Longer events have higher probability of being sampled
- Avoid the temptation of TIME_WAITED
 - AVG(time_waited) DOES NOT estimate avg wait times
 - MIN and MAX do not work either
 - Except when MAX exceeds 1-second



The ASH “fix up”

- ASH columns may be unknown at sampling time
 - TIME_WAITED: session is still waiting
 - PLAN_HASH: session is still optimizing SQL
 - GC events: event details unknown at event initiation
 - Certain time model bit vector columns
- ASH “fixes up” missing data during subsequent sample processing
 - TIME_WAITED fixed up in last event sample
- Querying current ASH may return un-fixed rows
 - Should not be a problem generally



ON CPU and ASH

- ASH row status “ON CPU” derived, not observed
 - Session is in a database call
 - Session is NOT in a wait event (idle or non-idle)
- Un-instrumented waits => “ON CPU”
 - These are bugs and should be rare, but have happened
- Session on run queue may be WAITING or ON CPU
 - Depends on state prior to going onto run queue



V\$EVENT_HISTOGRAM

- Histogram buckets of event wait times
- Captures statistical distribution of wait times
- All events since instance startup counted in some bucket
- Exponential time bucketing scheme captures long-tail distributions efficiently



V\$EVENT_HISTOGRAM

```
SQL> desc v$event_histogram
```

Name	Type
------	------

EVENT#	NUMBER
--------	--------

EVENT	VARCHAR2 (64)
-------	---------------

WAIT_TIME_MILLI	NUMBER
-----------------	--------

WAIT_COUNT	NUMBER
------------	--------

LAST_UPDATE_TIME	VARCHAR2 (64)
------------------	---------------

Event Histogram Time Buckets

```
SQL> select distinct
        wait_time_milli
        ,log(2,wait_time_milli)
from
        v$event_histogram
order by 1;
```

WAIT_TIME_MILLI	LOG(2,WAIT_TIME_MILLI)
1	0
2	1
4	2
8	3
16	4
32	5
64	6
128	7
256	8
512	9
1024	10



I/O Event Histogram

[Database: database](#) > [Active Sessions Waiting: User I/O](#) > Histogram for Wait Event: db file scattered read

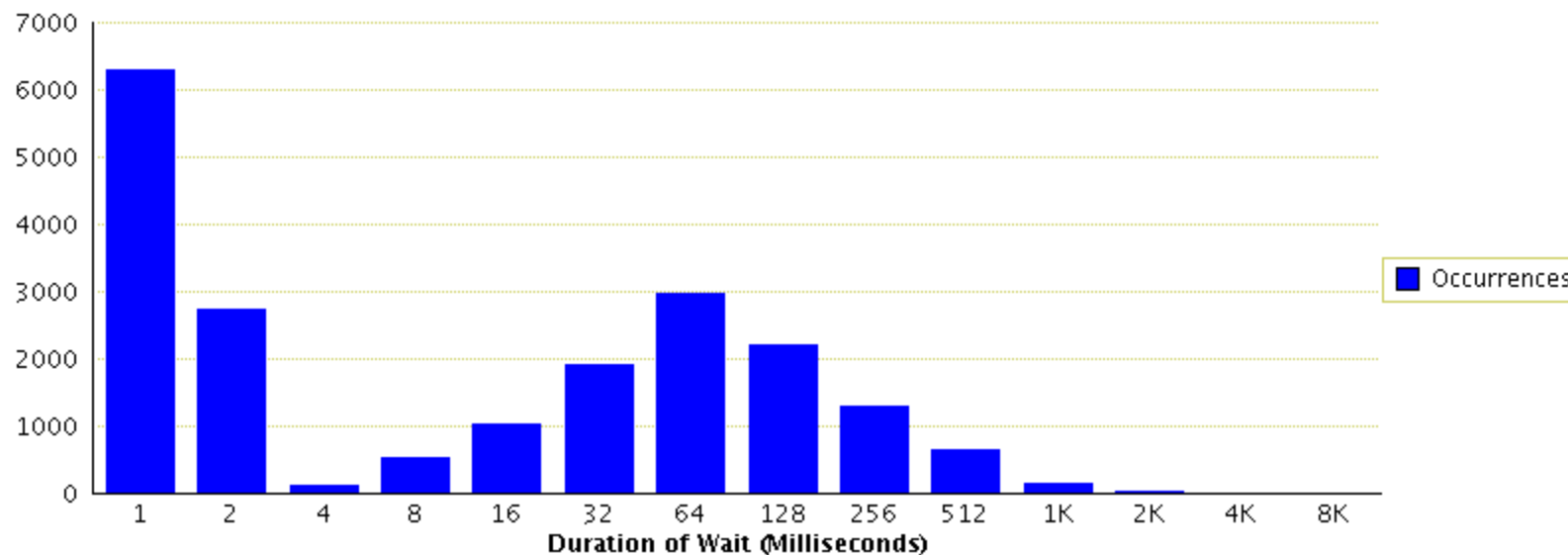
[Logged in As](#)

[View Data](#) [Real Time: Manual Refresh](#)

Histogram for Wait Event: db file scattered read

Page Refreshed **Nov 17, 2004 12:58:43**

Wait Event Occurrences Per Duration



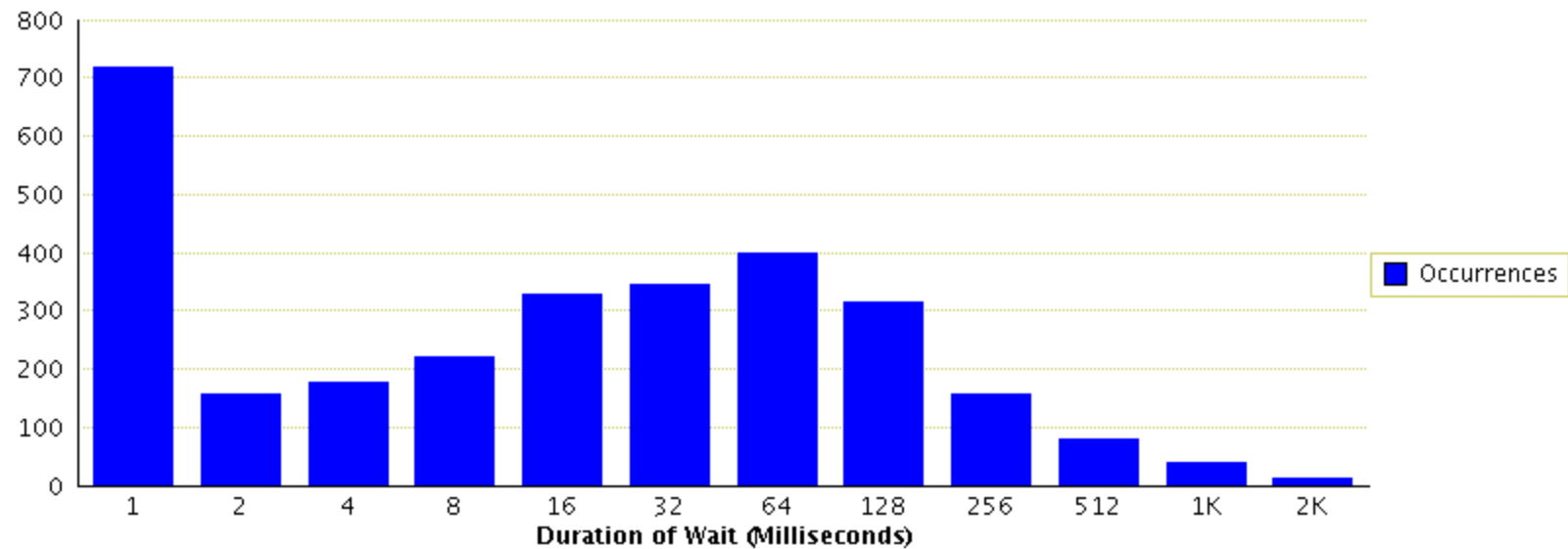


Latch Wait Event Histogram

Histogram for Wait Event: latch: library cache

Page Refreshed Nov 17, 2004 1:01:36

Wait Event Occurrences Per Duration





Histogram Math

- Histograms capture probability distribution of TIME_WAITED by event over this startup cycle

$$\Pr(\textit{time_waited} < \textit{bucket}_N) = \frac{\sum_{\textit{bucket} < N} \textit{WaitCount}}{\sum_{\textit{allbuckets}} \textit{WaitCount}}$$

Significance of Histogram Buckets

$$\textit{Significance}_{\textit{bucket}N} = 1 - \left(\frac{\sum_{\textit{bucket} \geq N} \textit{WaitCount}}{\sum_{\textit{allbuckets}} \textit{WaitCount}} \right)$$

- Measures the cumulative distribution function of TIME_WAITED probabilities represented by the histograms (per bucket)
- Every event in the bucket has at least this significance



Defining “Outlier Events”

- Events with low probability of occurrence
- Events with high significance value
- Q: Has ASH sampled any such events?



“Outlier” = “Unusual”

Database: database > [Active Sessions Waiting: Concurrency](#) > Histogram for Wait Event: latch: library cache

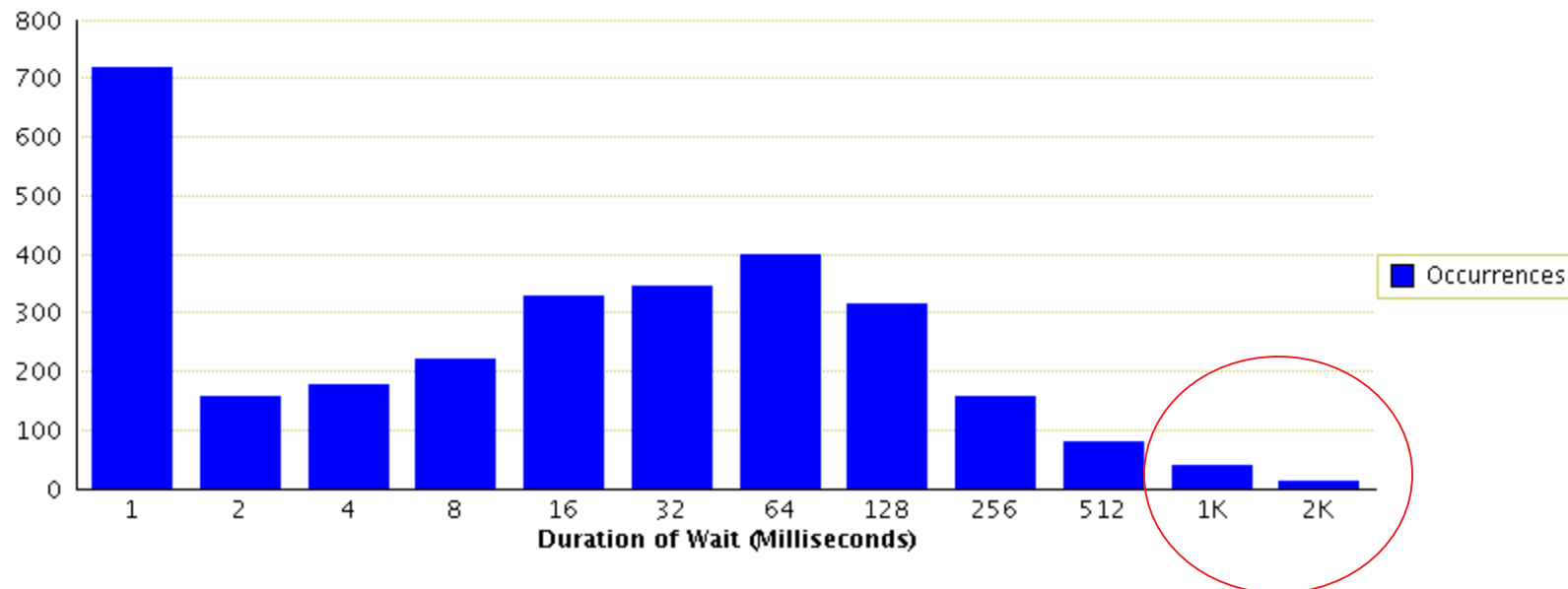
[Logged in As](#)

[View Data](#) [Real Time: Manual Refresh](#)

Histogram for Wait Event: latch: library cache

Page Refreshed Nov 17, 2004 1:01:36

Wait Event Occurrences Per Duration





Finding Outlier Events in ASH

- Which ASH rows (if any) represent wait events with significantly long `TIME_WAITED` against the event histogram record?
- Two step process:
 1. Compute event histogram bucket significance
 2. Join ASH to histograms and filter by significance

Step 1: Compute Bucket Significance

```
WITH EH$stats
as
(select
    EH.*
    ,ROUND(1 - (tot_count - bucket_tot_count + wait_count) / tot_count,6)
                                              as event_bucket_siglevel
from
    (select event#
        ,event
        ,wait_time_milli
        ,wait_count
        ,ROUND(LOG(2,wait_time_milli))           as event_bucket
        ,SUM(wait_count) OVER (PARTITION BY event#) as tot_count
        ,SUM(wait_count) OVER (PARTITION BY event# ORDER BY wait_time_milli
                                RANGE UNBOUNDED PRECEDING)
                                              as bucket_tot_count
        from v$event_histogram
    ) EH
)
```

Step 2: Join ASH to Buckets and Filter

```
select
    EH.event_bucket
  ,ASH.sample_id
  ,ASH.session_id
  ,EH.event_bucket_siglevel as bucket_siglevel
  ,ASH.event
  ,ASH.time_waited/1000 ASH_time_waited_milli
  ,ASH.sql_id
from
    EH$stats EH
  ,v$active_session_history ASH
where
    EH.event# = ASH.event#
  and EH.event_bucket_siglevel > &siglevel
  and EH.event_bucket = CASE ASH.time_waited
                        WHEN 0 THEN null
                        ELSE TRUNC(LOG(2,ASH.time_waited/1000))+1
                        END
order by
sample_id, event, session_id
;
```

A man in a dark suit, light blue shirt, and striped tie is sitting in a black leather office chair. He is gesturing with his right hand, palm facing up. Behind him is a large, silver, perforated metal rack of Oracle hardware. The rack has various labels and controls, including 'TAPE', 'AC', 'DC', 'STANDBY', 'XSCF', and a power button. The background is a blurred office setting with large windows.

DEMO?

SOFTWARE.
HARDWARE.
COMPLETE.

ORACLE



The GV\$ Problem

- Motivating use case is for RAC
- Execute V\$ query on all instances?
 - Too much effort
- Convert V\$ query to GV\$ query?
 - GV\$ remote joins not optimized
 - QC will pull V\$ASH and V\$EVENT_HISTOGRAM and join locally
- Neither of these “solutions” is acceptable



GV\$ Table Function

- Table function distributes V\$ cursor across RAC nodes and marshals result sets
- Perfect solution for this use case

```
SELECT ... FROM TABLE GV$( (CURSOR( SELECT FROM V$)) )
```

A man in a dark suit, light blue shirt, and striped tie is sitting in a black leather office chair. He is gesturing with his right hand, palm facing up. Behind him is a large, silver, perforated metal rack of Oracle hardware. The rack has various labels and controls, including 'TAPE', 'AC', 'DC', 'STANDBY', 'XSCF', and a power button. The background is a blurred office setting with large windows.

DEMO?

**SOFTWARE.
HARDWARE.
COMPLETE.**

ORACLE



Comments and Caveats

- Highly significant events may not be sampled
- Works best for long-tailed distributions
 - Bi-modal, single-bucket, timeout events not well-behaved
- Exponential bucket sizing gets coarse quickly
 - Significance levels increase in big jumps
 - Important distinctions may hide inside large buckets
- An interesting and unusual application of ASH where `TIME_WAITED` is the key
 - Does it help with the motivating use case?

