1 Interactive v2 Workload

This chapter is based on the TPCTC 2023 paper "The LDBC Social Network Benchmark Interactive Workload v2: A Transactional Graph Query Benchmark with Deep Delete Operations" [7], co-authored by several members of the SNB task force.

Work-in-Progress

The Interactive v2 workload is currently work-in-progress. As of January 2024, commissioning audits for this workload is not yet possible.

Related Software Components

- Datagen (Spark-based): https://github.com/ldbc/ldbc_snb_datagen_spark
- Driver: https://github.com/ldbc/ldbc_snb_interactive_v2_driver
- Reference implementations: https://github.com/ldbc/ldbc_snb_interactive_v2_impls

1.1 Overview

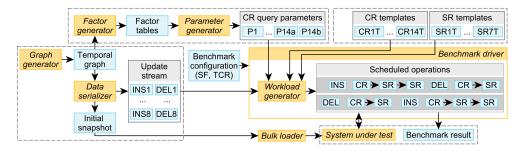


Figure 1.1: Components and workflow of the Interactive v2 workload. The corresponding sections are shown in green circles (§). Legend: Software component Data artifact

1.2 Operations

The LDBC SNB Interactive v2 workload uses four types of operations. There are 14 complex and 7 short read queries. Update operations include 8 inserts and, newly introduced in the Interactive v2 workload, 8 deletes. The workload mix consists of approximately 8% complex read, 72% short read, 20% insert, and 0.2% delete operations. The complex reads and the short reads are identical to the ones in Interactive v1, except for query 14, which was replaced to cover the *Cheapest path-finding* choke point. 1

Cheapest path-finding While we strived to keep the changes to the queries minimal, we replaced Q14 due to two reasons. First, we found the original query in Interactive v1 to be ill-suited to the workload as it required the enumeration of *all shortest paths* between two Persons, which can be prohibitively expensive on large scale factors. Second, we introduced a new choke point, CP-7.6 *Cheapest path-finding*, a key computational kernel and a language opportunity for GQL [2]. Therefore, we changed Q14 to use *cheapest paths* instead of *all shortest paths*.

¹The term *shortest paths* refers to the problem of finding *unweighted shortest paths*, which can be computed with BFS. The term *cheapest paths* refers to the *weighted shortest paths* problem, which can be solved using e.g. Dijkstra's algorithm.

1.2.1 Complex Reads

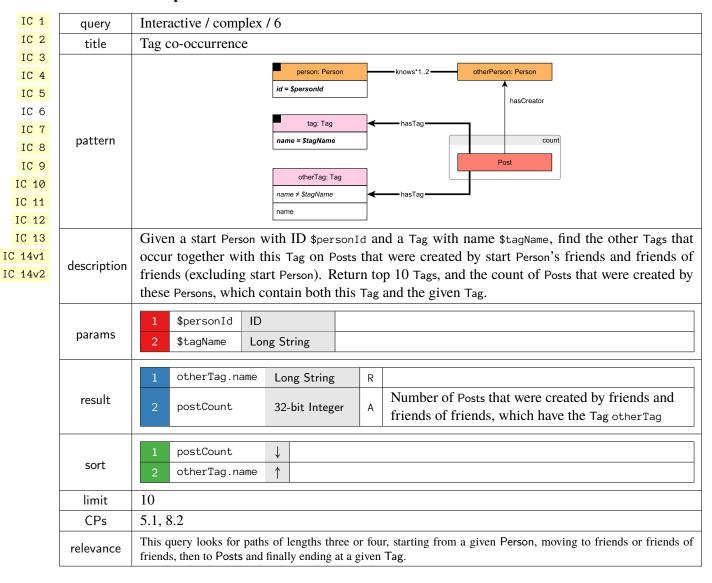
	Interactive / complex / 1				
IC 1	query	Interactive / complex / 1			
IC 2	title	Transitive friends with a certain n	ame		
IC 3 IC 4 IC 5 IC 6 IC 7 IC 8 IC 9	pattern	person: Person id = \$personId id = \$personId id lastName = id lastName birthday creationDate gender browserUse locationIP email speaks	«opt» workAt «opt»	name company: Company isLocatedIn companyCountry: Country name	
IC 10 IC 11 IC 12 IC 13 IC 14v1	description	start Person is connected to (exclud	ling start Person) b	with a given first name (\$firstName) that the by at most 3 steps via the knows relationships. aries of the Persons workplaces and places of	
IC 14v2	params	1 \$personId ID 2 \$firstName String			
	result	2 otherPerson.lastName 3 distanceFromPerson 4 otherPerson.birthday 5 otherPerson.creationDate 6 otherPerson.gender 7 otherPerson.browserUsed 8 otherPerson.locationIP 9 otherPerson.email 10 otherPerson.speaks 11 locationCity.name 12 universities	ID R String R 32-bit Integer C Date R DateTime R String R {String A String R String R String R String R String A String A String A String>}	{ <university.name, studyat.classyear,="" universitycity.name="">}</university.name,>	
	sort	1 distanceFromPerson ↑ 2 otherPerson.lastName ↑ 3 otherPerson.id ↑			
	limit	20			
	CPs	2.1, 5.3, 8.2			
	relevance	a complex aggregation for returning the of the Person. (2) It tests the ability of the second control of the s	concatenation of univ he optimizer to move e top-k. (3) Its perfo	It is interesting for several aspects. (1) It requires for ersities, companies, languages and email information the evaluation of sub-queries functionally dependant rmance is highly sensitive to properly estimating the explore already visited Persons.	

IC 1	query	Interactive / complex / 2							
IC 2	title	Recent messages by your fr	iends						
IC 3	0.0.0	person: Person		: Person		V Message			
IC 4	nattorn	id = \$personId	knows id		← hasCreator —	creationDate < \$maxDate			
IC 6	pattern		firstName lastName			id content / imageFile			
IC 7						creationDate			
IC 8		Given a start Person with II	given a start Person with ID \$personId, find the most recent Messages from all of that Person						
IC 9	description		iends (friend nodes). Only consider Messages created before the given \$maxDate (excluding the						
IC 10		day).	,		2				
IC 11		1 h T1 ID							
IC 12	params	1 \$personId ID							
IC 13		2 \$maxDate Date							
IC 14v1		1 friend.id	ID	R					
IC 14v2									
		2 friend.firstName	String	R					
		3 friend.lastName	String	R					
	result	4 message.id	ID	R					
		message.content or	_						
		5 message.imageFile (f	or Text	R					
		photos)							
		6 message.creationDate	DateTime	R					
		1 message.creationDate	·						
	sort	2 message.id	↑						
	limit	20							
	CPs								
	CPS	1.1, 2.2, 2.3, 3.2, 8.5							
	relevance	This is a navigational query looking for paths of length two, starting from a given Person, going to their friends and from them, moving to their published Posts and Comments. This query exercices both the optimizer and how data is stored. It tests the ability to create execution plans taking advantage of the orderings induced by some operators to avoid performing expensive sorts. This query requires selecting Posts and Comments based on their creation date, which might be correlated with their identifier and therefore, having intermediate results with interesting orders. Also, messages could be stored in an order correlated with their creation date to improve data access locality. Finally, as many of the attributes required in the projection are not needed for the execution of the query, it is expected that the query optimizer will move the projection to the end.							

IC 1	query	Interactive / complex / 3						
IC 2	title	Friends and friends or		that hav	e been to	o giv	ven countries	
IC 3	0.0.0							
IC 4					- hasCreator		xCount = count countryX: Country Message isLocatedIn	
IC 5		_		<u> </u>			name = \$countryXName	
IC 6	nottorn.	id = \$personId k	nows*12 — id	otherPerson	Person	< \$3	isPartOf	
IC 7	pattern	ia – spersonia	fir	stName stName			yCount = count «neg»	
IC 8				↑		V	Message siperior country: country	
IC 40				L	-hasCreator-		artDate ≤ creationDate startDate + \$durationDays isLocatedIn name = \$countryYName	
IC 10 IC 11								
IC 12							s that are their friends and friends of friends	
IC 13			-				omments in both of the given Countries (named	
IC 14v1	description		countryXName and \$countryYName), within [\$startDate, \$startDate + \$durationDays) (closed					
IC 14v2		•	open interval). Only Persons that are foreign to these Countries are considered, that is Person whose location Country is neither named \$countryXName nor \$countryYName.					
		whose location Countr	y is neiti	ier nam	eu \$count	tryx	Name HOI \$countryIName.	
		1 \$personId	ID					
					In SNE	3 Int	eractive v2, this query has two variants:	
		2 \$countryXName	String		` '		ated Countries	
					(b) An	ti-co	orrelated Countries	
	params	3 \$countryYName	\$countryYName String					
	4 \$startDate Date Beginning of req		of requested period					
			32-bit Integer				f requested period, in days. The interval	
		5 \$durationDays					e, \$startDate + \$durationDays) is	
					closed-	ope	n	
		1 otherPerson.id		ID		R		
		2 otherPerson.fi	rstName	String		R		
		3 otherPerson.la	stName	String		R		
							Number of Messages from Country named	
	1.	4 xCount	32-bit		Integer A		\$countryXName created by the Person within	
	result						the given time	
							Number of Messages from Country named	
		5 yCount		32-bit	Integer	Α	\$countryYName created by the Person within	
							the given time	
		6 count		32-bit	Integer	Α	count = xCount + yCount	
		1 count	 					
	sort	2 otherPerson.id	1					
	limit	20						
	CPs	2.1, 3.1, 5.1, 8.2, 8.5						
							om a Person, going to friends or friends of friends, and uery optimizer to select the most efficient join ordering,	
							sults. Many friends of friends can be duplicate, then it	
	relevance	is expected to eliminate d	uplicates a	nd those	people pri	or to	access the Post and Comments, as well as eliminate	
							htryYName, as the size of the intermediate results can be be to materialize the number of Posts and Comments	
							at could not even fall in the top 20 even having all their	
		posts in the Countries nan	ned \$count:	ryXName a	nd \$countr	yYNan	ne.	

T.C. 4									
IC 1	query	Interactive / complex / 4							
IC 2	title	New topics							
IC 3		Person	knows	person: Person knows friend: Person id = \$personId «opt»					
IC 5		hasCrea		hasCreator					
IC 6	pattern			postCount = count					
IC 7		Post	— «neg» hasTag ➤	tag: Tag ← Post					
IC 8		creationDate < \$startDate		\$startDate ≤ creationDate < \$startDate + \$durationDays					
IC 9									
IC 10		Given a start Person with II	D\$personId,	find Tags that are attached to Posts that were created by that					
IC 11		Person's friends. Only incl	ude Tags that	were attached to friends' Posts created within a given time					
IC 12	description	<pre>interval [\$startDate, \$sta</pre>	artDate + \$d	urationDays) (closed-open) and that were never attached					
IC 13 IC 14v1		to friends' Posts created be	eated before this interval.						
IC 14v2		1 \$personId ID							
		2 \$startDate Da	te						
	params	3 \$durationDays 32-	-bit Integer	Duration of requested period, in days. The interval [\$startDate, \$startDate + \$durationDays) is closed-open					
		1 tag.name Long St	tring R						
	result	2 postCount 32-bit li	nteger A	Number of Posts made within the given time interval that have tag					
		10 1							
	sort	1 postCount ↓							
	5011	2 tag.name ↑							
	limit	10							
	CPs	2.3, 8.2, 8.5							
	relevance	This query looks for paths of lethe ability of the query optimize	zer to properly s results. These	ing from a given Person, moving to Posts and then to Tags. It tests select the usage of hash joins or index based joins, depending on the cardinalities are clearly affected by the input Person, the number of nd the number of Posts.					

IC 1	query	Interactive / complex / 5						
IC 2	title	New groups						
IC 3 IC 4 IC 5 IC 6 IC 7 IC 8	pattern	person: Person id = \$personId white the second id = \$person \$\text{knows*12} \tag{otherPerson: Person} \tag{count}						
IC 9 IC 10 IC 11 IC 12	description	Given a start Person with ID \$personId, denote their friends and friends of friends (excluding the start Person) as otherPerson. Find Forums that any Person otherPerson became a member of after a given date (\$minDate). For each of those Forums, count the number of Posts that were created by the Person otherPerson.						
IC 13 IC 14v1 IC 14v2	params	1 \$personId ID 2 \$minDate Date						
	result	1 forum.title Long String R 2 postCount 32-bit Integer A Number of Posts made in forum that were created by the Person otherPerson						
	sort	1 postCount ↓ 2 forum.id ↑						
	limit	20						
	CPs	2.3, 3.3, 8.2, 8.5						
	relevance	This query looks for paths of length two and three, starting from a given Person, moving to friends and friends of friends, and then getting the Forums they are members of. Besides testing the ability of the query optimizer to select the proper join operator, it rewards the usage of indices, but their accesses will be presumably scattered due to the two/three-hop search space of the query, leading to unpredictable and scattered index accesses. Having efficient implementations of such indices will be highly beneficial.						



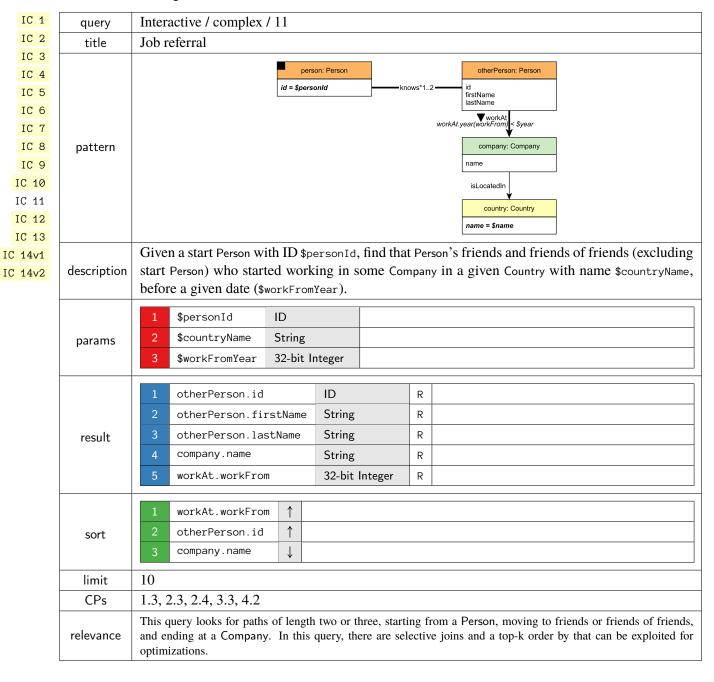
IC

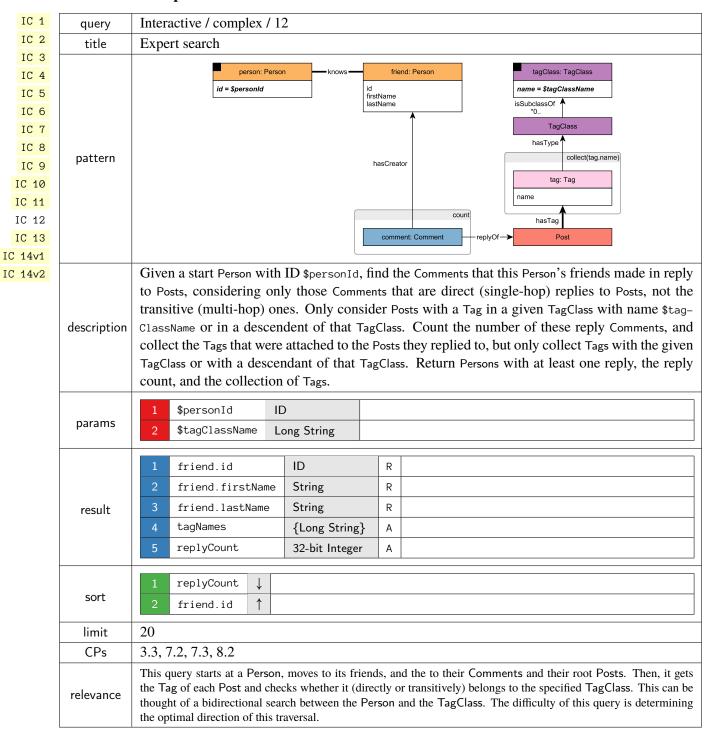
query	Interactive / complex / 7						
title	Recent likers						
pattern	id = \$perso	ge: Message	id firstName lastName				
description	sages. Find Persons that liked (like most recently, the creation date of creation of Messages and like. Add whether the liker is a friend of st same time, return the Message with Validation rule: Depending on VUTC-SLS (UTC with Smoothed IminutesLatency results of two corr	Given a start Person with ID \$personId, find the most recent likes on any of start Person's Messages. Find Persons that liked (likes edge) any of start Person's Messages, the Messages they liked most recently, the creation date of that like, and the latency in minutes (minutesLatency) between creation of Messages and like. Additionally, for each Person found return a flag indicating (isNew) whether the liker is a friend of start Person. In case that a Person liked multiple Messages at the same time, return the Message with lowest identifier. Validation rule: Depending on whether the system-under-test supports leap seconds or uses UTC-SLS (UTC with Smoothed Leap Seconds), a difference of 1 minute can occur between the minutesLatency results of two correct implementations when the time interval includes June 30, 2012, when there was a leap second. Therefore, the minutesLatency value is validated using a					
params	1 \$personId ID						
result	2 friend.firstName S 3 friend.lastName S 4 likes.creationDate D 5 message.id III message.content or 6 message.imageFile (for photos) 7 minutesLatency 3	D R String R String R DateTime R D R Fext R S2-bit Integer C Soolean C	Duration between the creation of the Message and the creation of the like, in minutes. False if person and friend know each				
sort	1 likes.creationDate ↓ 2 friend.id ↑						
limit	20						
CPs	2.2, 2.3, 3.3, 5.1, 8.1, 8.3						
relevance	This query looks for paths of length two, starting from a given Person, moving to its published messages and then to Persons who liked them. It tests several aspects related to join optimization, both at query optimization plan level and execution engine level. On the one hand, many of the columns needed for the projection are only needed in the last stages of the query, so the optimizer is expected to delay the projection until the end. This query implies accessing two-hop data, and as a consequence, index accesses are expected to be scattered. We expect to observe variate cardinalities, depending on the characteristics of the input parameter, so properly selecting the join operators will be crucial. This query has a lot of correlated sub-queries, so it is testing the ability to flatten the query execution plans.						

IC 1	query	Interactive / complex / 8
IC 2	title	Recent replies
IC 3 IC 4 IC 5 IC 6 IC 7 IC 8 IC 9 IC 10	pattern	person: Person id = \$personId firstName lastName hasCreator Message replyOf comment: Comment id content creationDate
IC 12 IC 13 IC 14v1	description	Given a start Person with ID \$personId, find the most recent Comments that are replies to Messages of the start Person. Only consider direct (single-hop) replies, not the transitive (multi-hop) ones. Return the reply Comments, and the Person that created each reply Comment.
IC 14v2	params	1 \$personId ID
	result	1 commentAuthor.id ID R 2 commentAuthor.firstName String R 3 commentAuthor.lastName String R 4 comment.creationDate DateTime R 5 comment.id ID R 6 comment.content Text R
	sort	1 comment.creationDate ↓ 2 comment.id ↑
	limit	20
	CPs	2.4, 3.3, 5.3
	relevance	This query looks for paths of length two, starting from a given Person, going through its created Messages and finishing at their replies. In this query there is temporal locality between the replies being accessed. Thus the top-k order by this can interact with the selection, i.e. do not consider older Posts than the 20th oldest seen so far.

IC 1	query	Interactive / complex / 9					
IC 2	title	Recent messages by friends or friends of friends					
IC 3 IC 4 IC 5 IC 6 IC 7 IC 8 IC 9 IC 10	pattern	person: Person id = \$personId knows*12 id firstName lastName hasCreator message: Message creationDate < \$maxDate} id content / imageFile creationDate					
IC 12 IC 13 IC 14v1	description	Given a start Person with ID \$personId, find the most recent Messages created by that Person's friends or friends (excluding the start Person). Only consider Messages created before the given \$maxDate (excluding that day).					
IC 14v2	params	1 \$personId ID 2 \$maxDate Date					
	result	1 otherPerson.id ID R 2 otherPerson.firstName String R 3 otherPerson.lastName String R 4 message.id ID R message.content or message.imageFile (for photos) 6 message.creationDate DateTime R					
	sort	<pre>1 message.creationDate ↓ 2 message.id ↑</pre>					
	limit	20					
	CPs	1.1, 1.2, 2.2, 2.3, 3.2, 3.3, 8.5					
	relevance	This query looks for paths of length two or three, starting from a given Person, moving to its friends and friends of friends, and ending at their created Messages. This is one of the most complex queries, as the list of choke points indicates. This query is expected to touch variable amounts of data with entities of different characteristics, and therefore, properly estimating cardinalities and selecting the proper operators will be crucial.					

IC 1	query	Interactive / complex / 10									
IC 2	title	Friend recommen									
IC 3					V			1			
IC 4		id = \$personId		knows*22	(month)		Person = \$month	isl ocatedin	name	city: City	
IC 5		1 1,21.2			and day (month(and day(birthday) ≥ 21) or (month(birthday) = \$month+1 (and day(birthday) < 22)					
IC 6					id	(birtnaay) < 22)				
IC 8					firstNam lastNam						
IC 9					gender						
IC 10	pattern				com	nmon				unce	ommon
IC 11		person: Person		foa	ıf: Person		person:	Porcon		foaf: Person	
IC 12		person. Person		hasCreato	_			reison		hasCreator 1	-
IC 13		hasInterest			cou	nt	«neg» hasInterest	,		C	ount
IC 14v1		Tag	← hasTag		Post		Ta	ıg 🗲	−hasTag −	Post	
IC 14v2						$\;\; \; \; \; \; \; \; \; \; \; \; \; \; \; \; \; \; \; \;$					
	description	start Person and his/her immediate friends –, who were born on or after the 21st of a given (in any year) and before the 22nd of the following month. Calculate the similarity betwee friend and the start person, where commonInterestScore is defined as follows: • common = number of Posts created by friend, such that the Post has a Tag that the start is interested in • uncommon = number of Posts created by friend, such that the Post has no Tag that the person is interested in • commonInterestScore = common - uncommon					nilarity between was: The start per	each			
		1 \$personId	ID								
	params	2 \$month	32-bit Inte	hit Integer		etween 1 and 12. Implementations may also pass the next north as an additional \$nextMonth parameter					ext
		1 foaf.id		ID		R					\neg
		2 foaf.first	Name	Strin	String						
		3 foaf.lastNa	ame	Strin	ıg	R					
	result	4 commonInte	restScore	32-b	it Integer	А					
		5 foaf.gende	r	Strin	ıg	R					
		6 city.name		Strin	ıg	R					
											=
	a a wt	1 commonInte	restScore	↓							
	sort 2 foaf.id		↑								
	limit	10									
	CPs	2.3, 3.3, 4.1, 4.2,	5.1, 5.2, 6.	1, 7.1	, 8.6						
	relevance	This query looks for widely scattered grap a long time and have in my friends" is bett the candidate's Posts	paths of leng h traversal, a widely scatte er with hash	gth two nd one red ide . Also	expects no le entifiers. The	ocality join o ern in	y of in friend order is simp the scalar su	ls of friends ble but one i	, as thes nust see	e have been acquired that the anti-join for	d over r "not





IC 1	query	Interactive / complex / 13					
IC 2	title	Single shortest path					
IC 3							
IC 4	pattern	Person Person					
IC 5		id = \$person1ld knows*0 id = \$person2ld					
IC 6		Given two Persons with IDs \$person1Id and \$person2Id, find the shortest path between these two					
IC 7		Persons in the subgraph induced by the knows edges. Return the length of this path:					
IC 8		• 1. no noth found					
IC 9	description	 -1: no path found 0: start person = end person 					
IC 10		 Start person = end person > 0: path found (start person ≠ end person) 					
IC 11		7 0. paul found (start person)					
IC 12 IC 13							
IC 13		In SNB Interactive v2, this query has two variants:					
IC 14v1		(b) Guaranteed that there is no path between the two					
10 1402		1 \$person1Id ID Persons					
	params	(b) Guaranteed that there is a 4-hop path between the two					
		Persons					
		2 \$person2Id ID					
	result	1 shortestPathLength 32-bit Integer C					
	CPs	3.3, 7.2, 7.3, 7.5, 7.8, 8.1, 8.6					
	relevance	This query looks for a variable length path, starting at a given Person and finishing at an another given Person. Proper cardinality estimation and search space pruning, will be crucial. This query also allows for possible parallel implementations.					

IC 1	query	Interactive / complex / 14v2
IC 2	title	Trusted connection paths (v2)
IC 3 IC 4 IC 5 IC 6		Find a cheapest path on edges where numInteractions ≥ 1, using edge weight = max(round(40 - sqrt(numInteractions)), 1) person1: Person id = \$person2!d numInteractions = count(c) personA: Person hasCreator hasCreator c: Comment replyOf → m: Message
IC 8 IC 9 IC 10 IC 11 IC 12 IC 13 IC 14v1 IC 14v2	pattern	Example for finding a path between person1 and person2 pt knows px knows py knows py knows py replyOf
	description	This query is used in SNB Interactive v2. Find a cheapest path between two given Persons with IDs \$person1Id and \$person2Id in the interaction subgraph. If there are multiple cheapest paths, any of them can be returned. Do not return any rows if there is no path between the Persons. The interaction subgraph is based on a projection of the Person-knows-Person graph. In this projection, only those knows edges are kept whose endpoint Persons have at least one interaction between them. An interaction is defined as a direct reply Comment (by one of the Persons) to a Message (by the other Person). The weights are defined as: $\max(\text{round}(40 - \sqrt{numInteractions}), 1)$ Note: Interactions are counted both ways, e.g. if Alice knows Bob, Alice writes 2 reply Comments to Bob's Messages and Bob writes 3 reply Comments to Alice's Messages, their total number of interactions is 5 and the weight of the knows edge is 38. Remark: Determinism is ensured by using square root followed by rounding. For all integers between 1 and 100000 , the square root's fractional part is more than $10e-5$ from 0.5 , where the rounding could be non-deterministic based on floating point inaccuracies. As $10e-5$ is significantly larger than the machine epsilon of IEEE 754 floats (both 32- and 64-bit), the floating point inaccuracies have no chance to affect the derived integer edge weights.
	params	\$\text{person1Id}\$ ID (b) There are no paths between the two Persons (b) There is a 4-hop path between the two Persons 2 \$\text{\$person2Id}\$ ID
	result	personIdsInPath [ID] C Identifiers representing an ordered sequence of the Persons in the path pathWeight 64-bit Integer C
	CPs	3.3, 5.3, 7.6, 7.7, 7.8, 8.1, 8.2, 8.3, 8.6
	relevance	This query tests the performance of cheapest path (weighted shortest path) computation.

1.2.2 Short Reads

IS 1
IS 2
IS 3
IS 4
IS 5
IS 6
IS 7

Interactive / short / 1

query	Interactive / short / 1					
title	Profile of a person					
pattern	id = \$person! d if = \$personId firstName lastName birthday locationIP browserUsed gender creationDate id id id					
description	Given a start Person with ID \$personId, retrieve their first name, last name, birthday, IP address, browser, and city of residence.					
params	1 \$personId ID					
result	1 person.firstName String R 2 person.lastName String R 3 person.birthday Date R 4 person.locationIP String R 5 person.browserUsed String R 6 city.id ID R 7 person.gender String R 8 person.creationDate DateTime R					

IS 1	query	Interactive / short / 2		
IS 2	title	Recent messages of a person		
IS 3 IS 4 IS 5 IS 6 IS 7	pattern	person: Person id = \$personId hasCreator message: Message id content / imageFile creationDate replyOf*0 originalPoster: Person id firstName lastName		
	description	Given a start Person with ID \$personId, retrieve the last 10 Messages created by that user. For each Message, return that Message, the original Post in its conversation (post), and the author of that Post (originalPoster). If any of the Messages is a Post, then the original Post (post) will be the same Message, i.e. that Message will appear twice in that result.		
	params	1 \$personId D		
	result	1 message.id ID R message.content or 2 message.imageFile (for photos) 3 message.creationDate DateTime R 4 post.id ID R 5 originalPoster.id ID R 6 originalPoster.firstName String R 7 originalPoster.lastName String R		
	sort	<pre>1 message.creationDate ↓ 2 message.id ↓</pre>		
	limit	10		

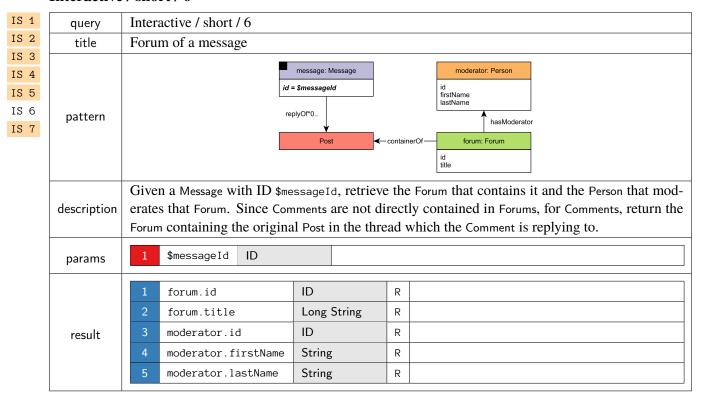
IS 1	query	Interactive / short / 3
IS 2	title	Friends of a person
IS 3		
IS 4		person: Person knows friend: Person creationDate
IS 5	pattern	id = \$personId id firstName
IS 6		lastName
IS 7	description	Given a start Person with ID \$personId, retrieve all of their friends, and the date at which they became friends.
	params	1 \$personId ID
		1 friend.id ID R
		2 friend.firstName String R
	result	3 friend.lastName String R
		4 knows.creationDate DateTime R
		1 knows.creationDate ↓
	sort	2 friend.id ↑

Interactive / short / 4

IS 1	query	Interactive / short / 4
IS 2	title	Content of a message
IS 3 IS 4 IS 5 IS 6 IS 7	pattern	message: Message id = \$messageId creationDate content / imageFile
15 /	description	Given a Message with ID \$messageId, retrieve its content and creation date.
	params	1 \$messageId ID
	result	1 message.creationDate

IS 1	query	Interactive / short / 5
IS 2	title	Creator of a message
IS 3		
IS 4		message: Message
IS 5	pattern	id = \$messageId id firstName
IS 6		lastName
IS 7	description	Given a Message with ID \$messageId, retrieve its author.
	params	1 \$messageId ID
		1 person.id ID R
	result	2 person.firstName String R
		3 person.lastName String R

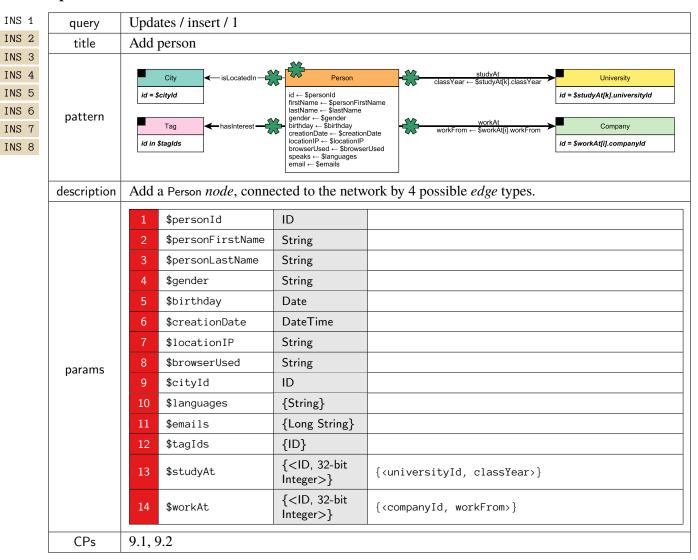
Interactive / short / 6



IS 1	query	Interactive / short / 7
IS 2	title	Replies of a message
IS 3 IS 4 IS 5 IS 6 IS 7	pattern	message: Message id = \$messageId replyOf comment: Comment id content creationDate hasCreator messageAuthor: Person id firstName lastName
	description	Given a Message with ID \$messageId, retrieve the (1-hop) Comments that reply to it. In addition, return a boolean flag knows indicating if the author of the reply (replyAuthor) knows the author of the original message (messageAuthor). If author is same as original author, return False for knows flag.
	params	1 \$messageId ID
	result	1 comment.id ID R 2 comment.content Text R 3 comment.creationDate DateTime R 4 replyAuthor.id ID R 5 replyAuthor.firstName String R 6 replyAuthor.lastName String R 7 knows Boolean C C True if the knows edge exists between the replyAuthor and the messageAuthor nodes, False otherwise (including the case when the two nodes are the same)
	sort	1 comment.creationDate ↓ 2 replyAuthor.id ↑

1.2.3 Insert Operations

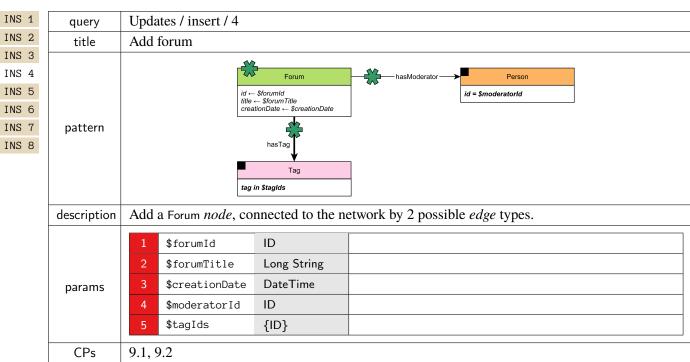
Updates / insert / 1



TNG 4				
INS 1	query	Updates / insert / 2		
INS 2	title	Add like to post		
INS 3		History (1997)		
INS 4	pattern	Person creationDate → \$creationDate → Post		
INS 5		id = \$personId id = \$postId		
INS 6	description	Add a likes <i>edge</i> to a Post.		
INS 7				
INS 8		1 \$personId ID		
	params	2 \$postId ID		
		3 \$creationDate DateTime		
	CPs	9.2		

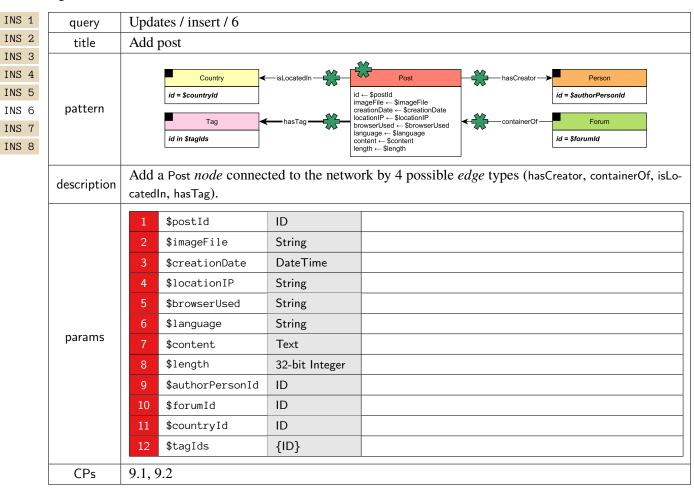
Updates / insert / 3

INS 1	query	Updates / insert / 3		
INS 2	title	Add like to comment		
INS 3				
INS 4	pattern	Person likes Comment CreationDate ← \$creationDate		
INS 5		id = \$personId id = \$commentId		
INS 6	description	Add a likes <i>edge</i> to a Comment.		
INS 7				
INS 8		1 \$personId ID		
	params	2 \$commentId ID		
		3 \$creationDate DateTime		
	CPs	9.2		



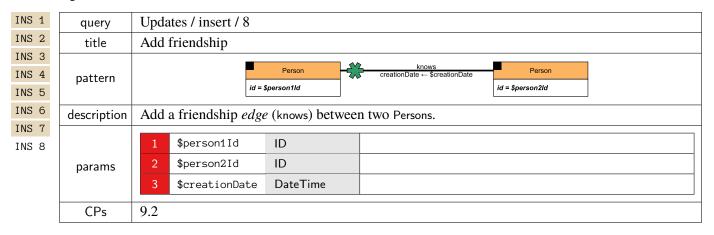
Updates / insert / 5

INS 1	query	Updates / insert / 5		
INS 2	title	Add forum membership		
INS 3		The state of the s		
INS 4	pattern	Person hasMember creationDate ← \$creationDate Forum		
INS 5		id = \$personId id = \$forumId		
INS 6	description	Add a Forum membership <i>edge</i> (hasMember) to a Person.		
INS 7	•	1 U V		
INS 8		1 \$personId ID		
	params	2 \$forumId ID		
		3 \$creationDate DateTime		
	CPs	9.1, 9.2		



Updates / insert / 7

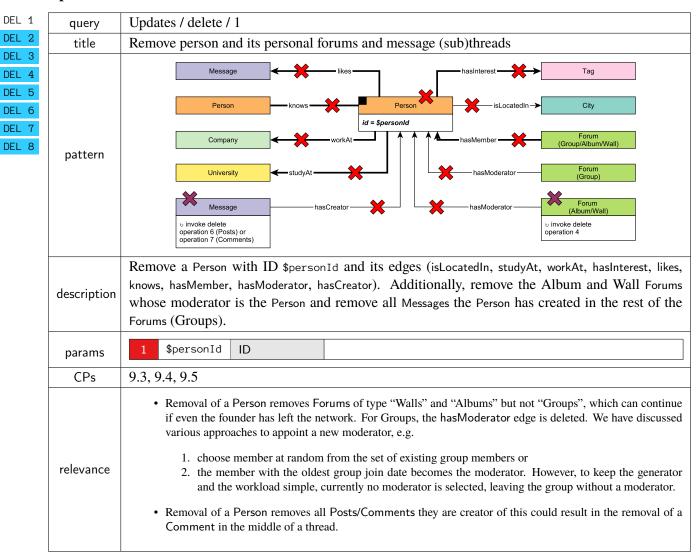
INS 1	query	Updates / insert / 7		
INS 2	title	Add comment		
INS 3 INS 4 INS 5 INS 6 INS 7 INS 8	pattern	Country id = \$countryid Tag id in \$taglds	Post id = \$replyToPostId replyOf- isLocatedIn hasTag Post id ← \$ creatic locatio brows:	Comment id = \$replyToCommentId id = \$replyToCommentId comment comment id = \$authorPersonId
	description	Add a Comment <i>node</i> repl types (replyOf, hasCreator, i		omment, connected to the network by 4 possible <i>edge</i> g).
	params	1 \$commentId 2 \$creationDate 3 \$locationIP 4 \$browserUsed 5 \$content 6 \$length 7 \$authorPersonId 8 \$countryId 9 \$replyToPostId 10 \$replyToCommentId	ID DateTime String String Text 32-bit Integer ID ID ID ID	old version: -1 if the Comment is a reply of a Comment; new version: null if the Comment is a reply of a Post old version: -1 if the Comment is a reply of a Post; new version: null if the Comment is a reply of a Post; Post
	CPs	9.1, 9.2		



1.2.4 Delete Operations

Updates / delete / 1

DEL



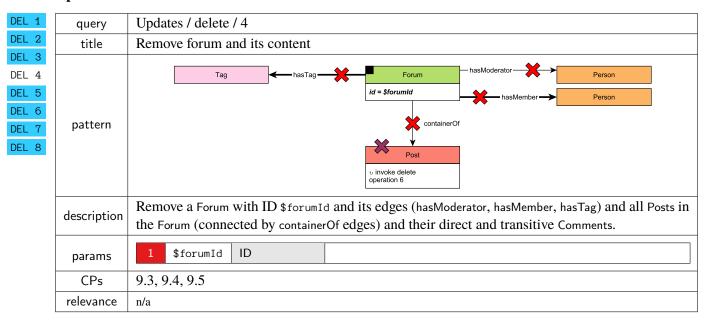
Updates / delete / 2

DEL 1	query	Updates / delete / 2
DEL 2	title	Remove post like
DEL 3		Person Post
DEL 4	pattern	id = \$personId id = \$postId
DEL 5		id – sposid
DEL 6	description	Given a Person with ID \$personId and a Post with ID \$postId, remove the likes edge between them.
DEL 7		A 71 ID
DEL 8	narame	1 \$personId ID
	params	2 \$postId ID
	CPs	9.4
	relevance	Removal of a likes edge is a rare event, e.g. people accidently liking a Post, this can be reflected by the relative frequency of the operation.

Updates / delete / 3

DEL 1	query	Updates / delete / 3	
DEL 2	title	Remove comment like	
DEL 3 DEL 4 DEL 5	pattern	Person id = \$personId id = \$commentId	
DEL 6	description	Given a Person with ID \$personId and a Comment with ID \$commentId, remove the likes edge between them.	
DEL 8	params	1 \$personId ID 2 \$commentId ID	
	CPs	9.4	
	relevance	Removal of a likes edge is a rare event, e.g. people accidently liking a Comment, this can be reflected by the relative frequency of the operation.	

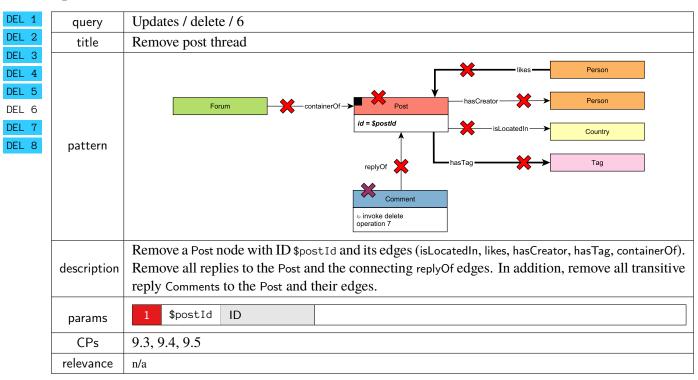
Updates / delete / 4



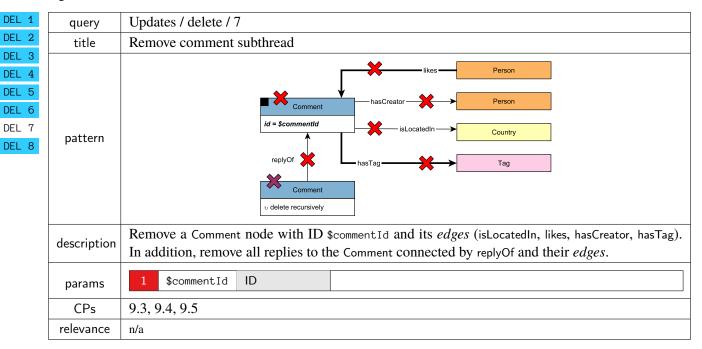
Updates / delete / 5

DEL 1	query	Updates / delete / 5		
DEL 2	title	Remove forum membership		
DEL 3				
DEL 4	pattern	Forum hasMember Person		
DEL 5	•	id = \$forumId id = \$personId		
DEL 6		Given a Forum with ID \$forumId and a Person with ID \$personId, remove the hasMember edge		
DEL 7	description	between them.		
DEL 8				
		1 \$forumId ID		
	params	2 \$personId D		
	CPs	9.4		
	relevance	n/a		

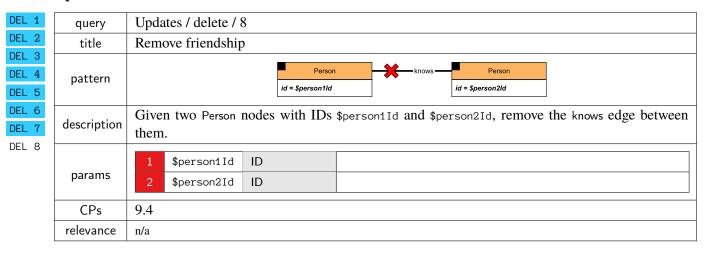
Updates / delete / 6



Updates / delete / 7



Updates / delete / 8



1.3 Parameter Curation

To prevent caching query results, the SNB Interactive v2 driver instantiates the parameterized complex read (IC) query templates with different *substitution parameters* (a.k.a. parameter bindings). However, the naïve approach (using a uniform random sampling of parameters and ignoring updates) leads to unstable runtimes, which compromise both the benchmark's understandability and reproducibility. To ensure stable runtimes, LDBC invented *parameter curation* techniques, which select parameters that produce query runtimes with a unimodal (preferably Gaussian) distribution [5, 10].

1.3.1 Building Blocks for Parameter Curation

Temporal bucketing To ensure that operations are always executable, i.e. they avoid targeting nodes that are yet to be inserted or ones that are already deleted, the parameter curation process in Interactive v2

employs *temporal bucketing*. Namely, we create a parameter bucket for *each day in the simulation time* of the update streams, i.e. each day in the simulation time has its own distinct set of parameters. This is a novel feature in Interactive v2 – previous SNB benchmarks lacked this feature and only selected parameters from the *initial snapshot*.

Factor tables As shown in Figure 1.1, the parameter generation is a two-step process. The *factor generator* produces *factor tables*, which contain data cube-like summary statistics [4] of the temporal graph such as the number of Messages for friends. The factor generator is executed in a distributed setup using Spark as this computation includes expensive joins over large tables, e.g. knows(person, friend) \bowtie hasCreator(person, comment).

1.3.2 Parameter Curation for Relational Queries

For relational queries (without path-finding), we based our parameter generation on two techniques.

- (1) **Selecting windows** To select the parameters that are expected to yield similar runtimes, we look for windows with the smallest variance for a given value using SQL window functions. The parameters are first sorted and grouped together based on their difference in frequency. Groups that are smaller than a given minimum threshold are discarded to select a group of parameters large enough to generate a sufficient amount of parameters. From the latter, we select the group with the smallest standard deviation.
- (2) **Selecting distributions** For queries where we want to select parameters that are correlated or anti-correlated, we use factor tables encoding possible combinations (e.g. countryPairsNumFriends for IC 3): we select values near a high percentile for the correlated and a low percentile for the anti-correlated case.

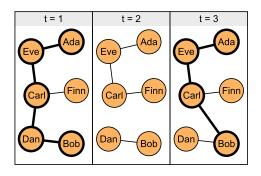
Generating the parameters The parameter candidates discovered by the previous approaches are stored in temporary tables. The parameter generation step uses these tables to select parameters for each day in the update stream.

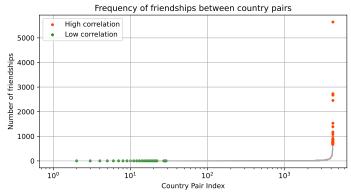
1.3.3 Parameter Curation for Path-Finding Queries

The effect of deletes A key distinguishing feature of graph data management systems is their first-class support for path queries [1]. We demonstrate why ensuring stable query runtimes for path queries is particularly challenging through the example of Figure 1.2a, where we query for the (unweighted) shortest path between Ada and Bob over a dynamic graph. Initially, at t=1, the length of the shortest path is 4 hops. Then, the edge between Carl and Dan is deleted, making Ada and Bob unreachable from each other at t=2. Finally, a new edge is inserted between Carl and Bob, yielding a shortest path of length 3 at t=3. This illustrates how a given input parameter (a pair of Persons) can oscillate between being reachable and being in disjoint connected components over a short period. To ensure stable query runtimes for path queries in the presence of inserts and deletes, Interactive v2 introduces a novel path curation algorithm, which produces pairs of Person nodes whose shortest path length from each other is guaranteed to be exactly k hops at any point during a given day.

Graph construction The parameter curation algorithm builds two variants of the Person–knows–Person subgraph for each day based on the *temporal graph*: graph G_1 has the inserts applied until the beginning of the day and the deletes applied until the end of the day, while G_2 has the deletes applied until the beginning of the day and the inserts applied until the end of the day. For a given pair of Person nodes, their shortest path length in G_1 is an upper bound $k_{\rm upper}$ on their shortest path length at any point in the day – when the inserts during the day are gradually applied, the shortest path length can only become shorter. Conversely, G_2 gives a lower bound $k_{\rm lower}$ for the shortest path – the deletes can only make the shortest path length become longer.

Parameter selection The bounds provided by G_1 and G_2 guarantee for the shortest path length k that $k_{\mathrm{lower}} \leq k \leq k_{\mathrm{upper}}$ will hold at any point during the day. We can ensure that k will stay constant during the day by selecting Person pairs where $k_{\mathrm{lower}} = k_{\mathrm{upper}}$ holds. To this end, we select pairs who are exactly 4 hops apart in both G_1 and G_2 , hence they will be always 4 hops apart during the given day. Unreachable pairs of nodes can be generated by calculating the connected components of G_2 and selecting nodes from disjoint components. The path curation for both the reachable and the unreachable cases is implemented using the NetworKit graph algorithm library [9].





- (a) Shortest path (denoted with thick lines) between *Ada* and *Bob* in the presence of updates.
- (b) Pairs of Countries in the countryPairsNumFriends factor table representing the number of friendships between both Countries.

Figure 1.2: Example graph and distribution for path curation.

1.3.4 Query Variants

The new workload introduces variants for three queries: IC 3, IC 13, IC 14v2.

Complex read 3: Correlated vs. anti-correlated Countries For IC 3 , variant IC 3(a) starts from Countries that have a high correlation in the friendship network, while variant IC 3(b) starts from Countries that have a low correlation of friendships between. To generate these inputs, we use the country-PairsNumFriends factor table visualized in Figure 1.2b and select values at percentile 1.00 for variant (a) and percentile 0.01 for variant (b).

Complex reads 13 and 14: Reachable vs. unreachable Persons Path queries are expected to have different runtimes if there is a path vs. when there is no path. While the performance characteristics vary highly between systems, in principle, the "no path" case should be simpler in the SNB graph, where one of the nodes is always in a small connected component. To distinguish between these cases, we have two variants for the two path queries IC 13 and IC 14v2. For variants (a) we select Person pairs which *do not have a path*, and for variants (b) we select pairs which *have* a path of length 4.

1.3.5 Parameter Generator Implementation

The parameter generator is implemented in Python using NetworKit [9] and SQL queries executed by DuckDB [8]. Based on our experiments in [6, Figure 4.3], the new parameter generator is scalable. Even with the significant extra work performed for temporal bucketing, it outperforms the old parameter generator by more than $100 \times$ on SF1 000, and finishes in less than 1.5 hours on SF10 000.

1.4 Workload Scheduling and Benchmark Driver

In this section, we explain how operations are scheduled in the SNB Interactive workload, how the driver operates, and how the final *throughput* metric is determined. In all cases, we assume that the system-

under-test has been populated with the *initial snapshot* using a *bulk loader* before the driver runs the operations.

1.4.1 Scheduling Operations

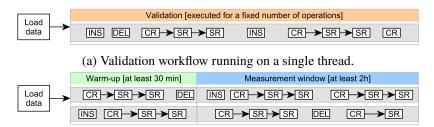
TCR (total compression ratio) The scheduling follows the *simulation time* of the temporal social network graph. The user-provided *total compression ratio* (TCR) value controls the speed at which the simulation is replayed. For example, a TCR value of 0.02 means that the simulation is replayed $50 \times$ faster, i.e. for every 20 milliseconds in wall clock time, 1 second passes in the simulation time.

Update operations The driver replays the update operations starting from the cutoff date, Nov 29, 2012. The operations are scheduled according to the distance of their start time from this date, adjusted by the TCR. They are then used to set the cadence of the schedule for the complex reads and, in turn, the short read queries, as we will explain momentarily.

Complex read queries The *complex read queries* differ significantly in their expected runtimes as they touch on different amounts of data. As each query instance contributes equally to the output metric,² we balance them such that each query type is expected to take the same amount of time to execute. For example, IC 14 (new) is expected to be more difficult than IC 13, therefore it is scheduled less frequently. Frequencies vary based on the SF as the relative difficulties of queries change with the data size (e.g. three-hop neighbourhood queries grow faster on larger SFs than one-hop ones).

Short read queries Short read queries are triggered by complex read queries and other short read queries, and use their output as their input. For example, both IC 3 and IC 14 trigger IS 2, which also triggers itself. This mimics the real-life scenario of a user retrieving more information about Person profiles based on the result of the earlier queries. To see which short read queries are potentially triggered after given short read and complex read queries, see ??.

1.4.2 Driver



(b) Benchmark workflow using multiple threads.

Figure 1.3: Workflow of driver modes in SNB Interactive v2.

Driver modes The SNB driver has two key modes of operation. In *cross-validation mode* (Figure 1.3a)m the driver tests an implementation against the output of another implementation. To ensure deterministic results, operations in this mode are executed sequentially with no overlap between queries and updates. In *benchmark mode* (Figure 1.3b), the driver performs a benchmark run where queries and updates are issued concurrently from multiple threads. The run starts with a 30-minute warm-up period, followed by a 2-hour *measurement window*. This mode does not perform validation as query results may differ (slightly) due to concurrent updates.

²Unlike in TPC-H [11] and SNB BI [10], which use *geometric mean* in their metrics.

Dependency tracking To ensure that updates are executable, concurrent threads must be synchronized so that an operation is only executed when its dependencies exist in the network (e.g. two Persons can only become friends if both of them already exist). This is achieved via maintaining a global clock in the driver and performing *dependency tracking* for the updates [3]: each update operation has a timestamp denoting the creation time of the last operation it depends on. The data generator calculates these timestamp during generation and ensures that there is a minimum time separation, T_{safe} , between dependent entities to reduce synchronization overhead in the driver when executing operations. The driver then only needs to check every T_{safe} time whether a given update operation can be executed. By default, T_{safe} is set to 10 seconds in the simulation time.

Latency requirements The workload simulates a highly transactional scenario where operations are subject to (soft) latency requirements. To incorporate this in the workload, it prescribes the 95% on-time requirement: for a benchmark run to be successful, 95% of the operations must start on-time, i.e. within 1 second of their scheduled start time. Benchmark runs where the system-under-test falls behind too much from the schedule are considered invalid.

Throughput The throughput of a run is the total number of operations (IC , IS , INS, DEL) executed per second. A lower TCR value implies a higher throughput.

Individual execution times To facilitate deeper analyis, the benchmark driver also collects all individual query execution times. Based on these, the benchmark reports must include statics for each operation type (min, max, mean, P_{50} , P_{90} , P_{95} , and P_{99} of the execution times).

Driver implementation in v2 The Interactive v2 is implemented in Java 17. It consists of 26 500 lines of code for the core project and an additional 18 000 lines of test code. The new version contains several patches including bug fixes, usability improvements, and performance optimizations.

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