

Mark3 Realtime Kernel

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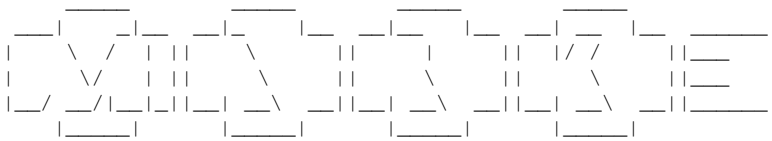
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Chapter 1

The Mark3 Realtime Kernel



--[Mark3 Realtime Platform]-----

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The Mark3 Realtime [Kernel](#) is a completely free, open-source, real-time operating system aimed at bringing multi-tasking to microcontroller systems without MMUs.

It uses modern programming languages and concepts (it's written entirely in C++) to minimize code duplication, and its object-oriented design enhances readability. The API is simple - there are only six functions required to set up the kernel, initialize threads, and start the scheduler.

The source is fully-documented with example code provided to illustrate concepts. The result is a performant RTOS, which is easy to read, easy to understand, and easy to extend to fit your needs.

But Mark3 is bigger than just a real-time kernel, it also contains a number of class-leading features:

- Device driver HAL which provides a meaningful abstraction around device-specific peripherals.
- Capable recursive-make driven build system which can be used to build all libraries, examples, tests, documentation, and user-projects for any number of targets from the command-line.
- Graphics and UI code designed to simplify the implementation of systems using displays, keypads, joysticks, and touchscreens
- Standards-based custom communications protocol used to simplify the creation of host tools
- A bulletproof, well-documented bootloader for AVR microcontrollers
- Support for kernel-aware simulators, specifically, Funkenstein Software's own fIAVR AVR simulator

Chapter 2

Preface

2.1 Who should read this

As the cover clearly states, this is a book about the Mark3 real-time kernel. I assume that if you're reading this book you have an interest in some, if not all, of the following subjects:

- Embedded systems
- Real-time systems
- Operating system kernel design

And if you're interested in those topics, you're likely familiar with C and C++ and the more you know, the easier you'll find this book to read. And if C++ scares you, and you don't like embedded, real-time systems, you're probably looking for another book. If you're unfamiliar with RTOS fundamentals, I highly suggest searching through the vast amount of RTOS-related articles on the internet to familiarize yourself with the concepts.

2.2 Why Mark3?

My first job after graduating from university in 2005 was with a small company that had a very old-school, low-budget philosophy when it came to software development. Every make-or-buy decision ended with "make" when it came to tools. It was the kind of environment where vendors cost us money, but manpower was free. In retrospect, we didn't have a ton of business during the time that I worked there, and that may have had something to do with the fact that we were constantly short on ready cash for things we could code ourselves.

Early on, I asked why we didn't use industry-standard tools - like JTAG debuggers or IDEs. One senior engineer scoffed that debuggers were tools for wimps - and something that a good programmer should be able to do without. After all - we had serial ports, GPIOs, and a bi-color LED on our boards. Since these were built into the hardware, they didn't cost us a thing. We also had a single software "build" server that took 5 minutes to build a 32k binary on its best days, so when we had to debug code, it was a painful process of trial and error, with lots of Youtube between iterations. We complained that tens of thousands of dollars of productivity was being flushed away that could have been solved by implementing a proper build server - and while we eventually got our wish, it took far more time than it should have.

Needless to say, software development was painful at that company. We made life hard on ourselves purely out of pride, and for the right to say that we walked "up-hills both ways through 3 feet of snow, everyday". Our code was tied ever-so-tightly to our hardware platform, and the system code was indistinguishable from the application. While we didn't use an RTOS, we had effectively implemented a 3-priority threading scheme using a carefully designed interrupt nesting scheme with event flags and a while(1) superloop running as a background thread. Nothing was abstracted, and the code was always optimized for the platform, presumably in an effort to save on code size and wasted cycles. I asked why we didn't use an RTOS in any of our systems and received dismissive scoffs - the overhead from thread switching and maintaining multiple threads could not be tolerated in our systems according

to our chief engineers. In retrospect, our ad-hoc system was likely as large as my smallest kernel, and had just as much context switching (although it was hidden by the compiler).

And every time a new iteration of our product was developed, the firmware took far too long to bring up, because the algorithms and data structures had to be re-tooled to work with the peripherals and sensors attached to the new boards. We worked very hard in an attempt to reinvent the wheel, all in the name of producing "efficient" code.

Regardless, I learned a lot about software development.

Most important, I learned that good design is the key to good software; and good design doesn't have to come at a price. In all but the smallest of projects, the well-designed, well-abstracted code is not only more portable, but it's usually smaller, easier to read, and easier to reuse.

Also, since we had all the time in the world to invest in developing our own tools, I gained a lot of experience building them, and making use of good, free PC tools that could be used to develop and debug a large portion of our code. I ended up writing PC-based device and peripheral simulators, state-machine frameworks, and abstractions for our horrible ad-hoc system code. At the end of the day, I had developed enough tools that I could solve a lot of our development problems without having to re-inventing the wheel at each turn. Gaining a background in how these tools worked gave me a better understanding of how to use them - making me more productive at the jobs that I've had since.

I am convinced that designing good software takes honest effort up-front, and that good application code cannot be written unless it is based on a solid framework. Just as the wise man builds his house on rocks, and not on sand, wise developers write applications based on a well-defined platforms. And while you can probably build a house using nothing but a hammer and sheer will, you can certainly build one a lot faster with all the right tools.

This conviction lead me to development my first RTOS kernel in 2009 - FunkOS. It is a small, yet surprisingly full-featured kernel. It has all the basics (semaphores, mutexes, round-robin and preemptive scheduling), and some pretty advanced features as well (device drivers and other middleware). However, it had two major problems - it doesn't scale well, and it doesn't support many devices.

While I had modest success with this kernel (it has been featured on some blogs, and still gets around 125 downloads a month), it was nothing like the success of other RTOS kernels like uC/OS-II and FreeRTOS. To be honest, as a one-man show, I just don't have the resources to support all of the devices, toolchains, and evaluation boards that a real vendor can. I had never expected my kernel to compete with the likes of them, and I don't expect Mark3 to change the embedded landscape either.

My main goal with Mark3 was to solve the technical shortfalls in the FunkOS kernel by applying my experience in kernel development. As a result, Mark3 is better than FunkOS in almost every way; it scales better, has lower interrupt latency, and is generally more thoughtfully designed (all at a small cost to code size).

Another goal I had was to create something easy to understand, that could be documented and serve as a good introduction to RTOS kernel design. The end result of these goals is the kernel as presented in this book - a full source listing of a working OS kernel, with each module completely documented and explained in detail.

Finally, I wanted to prove that a kernel written entirely in C++ could perform just as well as one written in C, without incurring any extra overhead. Comparing the same configuration of Mark2 to Mark3, the code size is remarkably similar, and the execution performance is just as good. Not only that, but there are fewer lines of code. The code is more readable and easier to understand as a result of making use of object-oriented concepts provided by C++. Applications are easier to write because common concepts are encapsulated into objects (Threads, Semaphores, Mutexes, etc.) with their own methods and data, as opposed to APIs which rely on lots of explicit pointer-passing, type casting, and other operations that are typically considered "unsafe" or "advanced topics" in C.

Chapter 3

Can you Afford an RTOS?

Of course, since you're reading the manual for an RTOS that I've been developing for the last few years, you can guess that the conclusion that I draw is a resounding "yes".

If your code is of any sort of non-trivial complexity (say, at least a few-thousand lines), then a more appropriate question would be "can you afford **not** to use an RTOS in your system?".

In short, there are simply too many benefits of an RTOS to ignore.

- Sophisticated synchronization objects
- The ability to efficiently block and wait
- Enhanced responsiveness for high-priority tasks
- Built in timers
- Built in efficient memory management

Sure, these features have a cost in code space and RAM, but from my experience the cost of trying to code around a lack of these features will cost you as much - if not more. The results are often far less maintainable, error prone, and complex. And that simply adds time and cost. Real developers ship, and the RTOS is quickly becoming one of the standard tools that help keep developers shipping.

3.1 Intro

(Note - this article was written for the C-based Mark2 kernel, which is slightly different. While the general principles are the same, the numbers are not an 100% accurate reflection of the current costs of the Mark3 kernel.)

One of the main arguments against using an RTOS in an embedded project is that the overhead incurred is too great to be justified. Concerns over "wasted" RAM caused by using multiple stacks, added CPU utilization, and the "large" code footprint from the kernel cause a large number of developers to shun using a preemptive RTOS, instead favoring a non-preemptive, application-specific solution.

I believe that not only is the impact negligible in most cases, but that the benefits of writing an application with an RTOS can lead to savings around the board (code size, quality, reliability, and development time). While these other benefits provide the most compelling case for using an RTOS, they are far more challenging to demonstrate in a quantitative way, and are clearly documented in numerous industry-based case studies.

While there is some overhead associated with an RTOS, the typical arguments are largely unfounded when an RTOS is correctly implemented in a system. By measuring the true overhead of a preemptive RTOS in a typical application, we will demonstrate that the impact to code space, RAM, and CPU usage is minimal, and indeed acceptable for a wide range of CPU targets.

To illustrate just how little an RTOS impacts the size of an embedded software design we will look at a typical microcontroller project and analyze the various types of overhead associated with using a pre-emptive realtime kernel versus a similar non-preemptive event-based framework.

RTOS overhead can be broken into three distinct areas:

- Code space: The amount of code space eaten up by the kernel (static)
- Memory overhead: The RAM associated with running the kernel and application threads.
- Runtime overhead: The CPU cycles required for the kernel's functionality (primarily scheduling and thread switching)

While there are other notable reasons to include or avoid the use of an RTOS in certain applications (determinism, responsiveness, and interrupt latency among others), these are not considered in this discussion - as they are difficult to consider for the scope of our "canned" application. Application description:

For the purpose of this comparison, we first create an application using the standard preemptive Mark3 kernel with 2 system threads running: A foreground thread and a background thread. This gives three total priority levels in the system - the interrupt level (high), and two application priority threads (medium and low), which is quite a common paradigm for microcontroller firmware designs. The foreground thread processes a variety of time-critical events at a fixed frequency, while the background thread processes lower priority, aperiodic events. When there are no background thread events to process, the processor enters its low-power mode until the next interrupt is acknowledged.

The contents of the threads themselves are unimportant for this comparison, but we can assume they perform a variety of I/O using various user-input devices and a serial graphics display. As a result, a number of Mark3 device drivers are also implemented.

The application is compiled for an ATmega328p processor which contains 32kB of code space in flash, and 2kB of RAM, which is a lower-mid-range microcontroller in Atmel's 8-bit AVR line of microcontrollers. Using the WinAVR GCC compiler with -O2 level optimizations, an executable is produced with the following code/RAM utilization:

31600 Bytes Code Space 2014 Bytes RAM

An alternate version of this project is created using a custom "super-loop" kernel, which uses a single application thread and provides 2 levels of priority (interrupt and application). In this case, the event handler processes the different priority application events to completion from highest to lowest priority.

This approach leaves the application itself largely unchanged. Using the same optimization levels as the preemptive kernel, the code compiles as follows:

29904 Bytes Code Space 1648 Bytes RAM

3.2 Memory overhead:

At first glance, the difference in RAM utilization seems quite a lot higher for the preemptive mode version of the application, but the raw numbers don't tell the whole story.

The first issue is that the cooperative-mode total does not take into account the system stack - whereas these values are included in the totals for RTOS version of the project. As a result, some further analysis is required to determine how the stack sizes truly compare.

In cooperative mode, there is only one thread of execution - so considering that multiple event handlers are executed in turn, the stack requirements for cooperative mode is simply determined by those of the most stack-intensive event handler.

In contrast, the preemptive kernel requires a separate stack for each active thread, and as a result the stack usage of the system is the sum of the stacks for all threads.

Since the application and idle events are the same for both preemptive and cooperative mode, we know that their (independent) stack requirements will be the same in both cases.

For cooperative mode, we see that the idle thread stack utilization is lower than that of the application thread, and so the application thread's determines the stack size requirement. Again, with the preemptive kernel the stack utilization is the sum of the stacks defined for both threads.

As a result, the difference in overhead between the two cases becomes the extra stack required for the idle thread - which in our case is (a somewhat generous) 64 bytes.

The numbers still don't add up completely, but looking into the linker output we see that the rest of the difference comes from the extra data structures used to declare the threads in preemptive mode.

With this taken into account, the true memory cost of a 2-thread system ends up being around 150 bytes of RAM - which is less than 8% of the total memory available on this particular microcontroller. Whether or not this is reasonable certainly depends on the application, but more importantly, it is not so unreasonable as to eliminate an RTOS-based solution from being considered.

3.3 Code Space Overhead:

The difference in code space overhead between the preemptive and cooperative mode solutions is less of an issue. Part of this reason is that both the preemptive and cooperative kernels are relatively small, and even an average target device (like the Atmega328 we've chosen) has plenty of room.

Mark3 can be configured so that only features necessary for the application are included in the RTOS - you only pay for the parts of the system that you use. In this way, we can measure the overhead on a feature-by-feature basis, which is shown below for the kernel as configured for this application:

3466 Bytes

The configuration tested in this comparison uses the thread/port module with timers, drivers, and semaphores, for a total kernel size of ~3.5KB, with the rest of the code space occupied by the application.

The custom cooperative-mode framework has a similar structure which is broken down by module as follows:

1850 Bytes

As can be seen from the compiler's output, the difference in code space between the two versions of the application is about 1.7kB - or about 5% of the available code space on the selected processor. While nearly all of this comes from the added overhead of the kernel, the rest of the difference comes the changes to the application necessary to facilitate the different frameworks.

3.4 Runtime Overhead

On the cooperative kernel, the overhead associated with running the thread is the time it takes the kernel to notice a pending event flag and launch the appropriate event handler, plus the timer interrupt execution time.

Similarly, on the preemptive kernel, the overhead is the time it takes to switch contexts to the application thread, plus the timer interrupt execution time.

The timer interrupt overhead is similar for both cases, so the overhead then becomes the difference between the following:

Preemptive mode:

- Posting the semaphore that wakes the high-priority thread
- Performing a context switch to the high-priority thread

Cooperative mode:

- Setting the high-priority thread's event flag
- Acknowledging the event from the event loop

Using the cycle-accurate AVR simulator, we find the end-to-end event sequence time to be 20.4us for the cooperative mode scheduler and 44.2us for the preemptive, giving a difference of 23.8us.

With a fixed high-priority event frequency of 33Hz, we achieve a runtime overhead of 983.4us per second, or 0.0983% of the total available CPU time. Now, obviously this value would expand at higher event frequencies and/or slower CPU frequencies, but for this typical application we find the difference in runtime overhead to be negligible for a preemptive system. Analysis:

For the selected test application and platform, including a preemptive RTOS is entirely reasonable, as the costs are low relative to a non-preemptive kernel solution. But these costs scale relative to the speed, memory and code space of the target processor. Because of these variables, there is no "magic bullet" environment suitable for every application, but Mark3 attempts to provide a framework suitable for a wide range of targets.

On the one hand, if these tests had been performed on a higher-end microcontroller such as the ATmega1284p (containing 128kB of code space and 16kB of RAM), the overhead would be in the noise. For this type of resource-rich microcontroller, there would be no reason to avoid using the Mark3 preemptive kernel.

Conversely, using a lower-end microcontroller like an ATmega88pa (which has only 8kB of code space and 1kB of RAM), the added overhead would likely be prohibitive for including a preemptive kernel. In this case, the cooperative-mode kernel would be a better choice.

As a rule of thumb, if one budgets 10% of a microcontroller's code space/RAM for a preemptive kernel's overhead, you should only require at minimum a microcontroller with 16k of code space and 2kB of RAM as a base platform for an RTOS. Unless there are serious constraints on the system that require much better latency or responsiveness than can be achieved with RTOS overhead, almost any modern platform is sufficient for hosting a kernel. In the event you find yourself with a microprocessor with external memory, there should be no reason to avoid using an RTOS at all.

Chapter 4

Superloops

4.1 Intro to Superloops

Before we start taking a look at designing a real-time operating system, it's worthwhile taking a look through one of the most-common design patterns that developers use to manage task execution in embedded systems - Superloops.

Systems based on superloops favor the system control logic baked directly into the application code, usually under the guise of simplicity, or memory (code and RAM) efficiency. For simple systems, superloops can definitely get the job done. However, they have some serious limitations, and are not suitable for every kind of project. In a lot of cases you can squeak by using superloops - especially in extremely constrained systems, but in general they are not a solid basis for reusable, portable code.

Nonetheless, a variety of examples are presented here- from the extremely simple, to cooperative and limited-preemptive multitasking systems, all of which are examples are representative of real-world systems that I've either written the firmware for, or have seen in my experience.

4.2 The simplest loop

Let's start with the simplest embedded system design possible - an infinite loop that performs a single task repeatedly:

```
int main()
{
    while(1)
    {
        Do_Something();
    }
}
```

Here, the code inside the loop will run a single function forever and ever. Not much to it, is there? But you might be surprised at just how much embedded system firmware is implemented using essentially the same mechanism - there isn't anything wrong with that, but it's just not that interesting.

While the execution timeline for this program is equally boring, for the sake of completeness it would look like this:

Despite its simplicity we can see the beginnings of some core OS concepts. Here, the `while(1)` statement can be logically seen as the operating system kernel - this one control statement determines what tasks can run in the system, and defines the constraints that could modify their execution. But at the end of the day, that's a big part of what a kernel is - a mechanism that controls the execution of application code.

The second concept here is the task. This is application code provided by the user to perform some useful purpose in a system. In this case `Do_something()` represents that task - it could be monitoring blood pressure, reading a sensor and writing its data to a terminal, or playing an MP3; anything you can think of for an embedded system to do. A simple round-robin multi-tasking system can be built off of this example by simply adding additional tasks in

sequence in the main while-loop. Note that in this example the CPU is always busy running tasks - at no time is the CPU idle, meaning that it is likely burning a lot of power.

While we conceptually have two separate pieces of code involved here (an operating system kernel and a set of running tasks), they are not logically separate. The OS code is indistinguishable from the application. It's like a single-celled organism - everything is crammed together within the walls of an indivisible unit; and specialized to perform its given function relying solely on instinct.

4.3 Interrupt-Driven Super-loop

In the previous example, we had a system without any way to control the execution of the task- it just runs forever. There's no way to control when the task can (or more importantly can't) run, which greatly limits the usefulness of the system. Say you only want your task to run every 100 milliseconds - in the previous code, you have to add a hard-coded delay at the end of your task's execution to ensure your code runs only when it should.

Fortunately, there is a much more elegant way to do this. In this example, we introduce the concept of the synchronization object. A Synchronization object is some data structure which works within the bounds of the operating system to tell tasks when they can run, and in many cases includes special data unique to the synchronization event. There are a whole family of synchronization objects, which we'll get into later. In this example, we make use of the simplest synchronization primitive - the global flag.

With the addition of synchronization brings the addition of event-driven systems. If you're programming a microcontroller system, you generally have scores of peripherals available to you - timers, GPIOs, ADCs, UARTs, ethernet, USB, etc. All of which can be configured to provide a stimulus to your system by means of interrupts. This stimulus gives us the ability not only to program our micros to do_something(), but to do_something() if-and-only-if a corresponding trigger has occurred.

The following concepts are shown in the example below:

```
volatile K_BOOL something_to_do = false;

__interrupt__ My_Interrupt_Source(void)
{
    something_to_do = true;
}

int main()
{
    while(1)
    {
        if( something_to_do )
        {
            Do_something();
            something_to_do = false;
        }
        else
        {
            Idle();
        }
    }
}
```

So there you have it - an event driven system which uses a global variable to synchronize the execution of our task based on the occurrence of an interrupt. It's still just a bare-metal, OS-baked-into-the-application system, but it's introduced a whole bunch of added complexity (and control!) into the system.

The first thing to notice in the source is that the global variable, something_to_do, is used as a synchronization object. When an interrupt occurs from some external event, triggering the My_Interrupt_Source() ISR, program flow in main() is interrupted, the interrupt handler is run, and something_to_do is set to true, letting us know that when we get back to main(), that we should run our Do_something() task.

Another new concept at play here is that of the idle function. In general, when running an event driven system, there are times when the CPU has no application tasks to run. In order to minimize power consumption, CPUs usually contain instructions or registers that can be set up to disable non-essential subsets of the system when there's nothing to do. In general, the sleeping system can be re-activated quickly as a result of an interrupt or other external stimulus, allowing normal processing to resume.

Now, we could just call `Do_something()` from the interrupt itself - but that's generally not a great solution. In general, the more time we spend inside an interrupt, the more time we spend with at least some interrupts disabled. As a result, we end up with interrupt latency. Now, in this system, with only one interrupt source and only one task this might not be a big deal, but say that `Do_something()` takes several seconds to complete, and in that time several other interrupts occur from other sources. While executing in our long-running interrupt, no other interrupts can be processed - in many cases, if two interrupts of the same type occur before the first is processed, one of these interrupt events will be lost. This can be utterly disastrous in a real-time system and should be avoided at all costs. As a result, it's generally preferable to use synchronization objects whenever possible to defer processing outside of the ISR.

Another OS concept that is implicitly introduced in this example is that of task priority. When an interrupt occurs, the normal execution of code in `main()` is preempted: control is swapped over to the ISR (which runs to completion), and then control is given back to `main()` where it left off. The very fact that interrupts take precedence over what's running shows that `main` is conceptually a "low-priority" task, and that all ISRs are "high-priority" tasks. In this example, our "high-priority" task is setting a variable to tell our "low-priority" task that it can do something useful. We will investigate the concept of task priority further in the next example.

Preemption is another key principle in embedded systems. This is the notion that whatever the CPU is doing when an interrupt occurs, it should stop, cache its current state (referred to as its context), and allow the high-priority event to be processed. The context of the previous task is then restored its state before the interrupt, and resumes processing. We'll come back to preemption frequently, since the concept comes up frequently in RTOS-based systems.

4.4 Cooperative multi-tasking

Our next example takes the previous example one step further by introducing cooperative multi-tasking:

```
// Bitfield values used to represent three distinct tasks
#define TASK_1_EVENT (0x01)
#define TASK_2_EVENT (0x02)
#define TASK_3_EVENT (0x04)

volatile K_UCHAR event_flags = 0;

// Interrupt sources used to trigger event execution

__interrupt__ My_Interrupt_1(void)
{
    event_flags |= TASK_1_EVENT;
}

__interrupt__ My_Interrupt_2(void)
{
    event_flags |= TASK_2_EVENT;
}

__interrupt__ My_Interrupt_3(void)
{
    event_flags |= TASK_3_EVENT;
}

// Main tasks
int main(void)
{
    while(1)
    {
        while(event_flags)
        {
            if( event_flags & TASK_1_EVENT)
            {
                Do_Task_1();
                event_flags &= ~TASK_1_EVENT;
            } else if( event_flags & TASK_2_EVENT) {
                Do_Task_2();
                event_flags &= ~TASK_2_EVENT;
            } else if( event_flags & TASK_3_EVENT) {
                Do_Task_3();
                event_flags &= ~TASK_3_EVENT;
            }
        }
        Idle();
    }
}
```

This system is very similar to what we had before - however the differences are worth discussing. First, we have stimulus from multiple interrupt sources: each ISR is responsible for setting a single bit in our global event flag, which is then used to control execution of individual tasks from within main().

Next, we can see that tasks are explicitly given priorities inside the main loop based on the logic of the if/else if structure. As long as there is something set in the event flag, we will always try to execute Task1 first, and only when Task1 isn't set will we attempt to execute Task2, and then Task 3. This added logic provides the notion of priority. However, because each of these tasks exist within the same context (they're just different functions called from our main control loop), we don't have the same notion of preemption that we have when dealing with interrupts.

That means that even through we may be running Task2 and an event flag for Task1 is set by an interrupt, the CPU still has to finish processing Task2 to completion before Task1 can be run. And that's why this kind of scheduling is referred to ascooperative multitasking: we can have as many tasks as we want, but unless they cooperate by means of returning back to main, the system can end up with high-priority tasks getting starved for CPU time by lower-priority, long-running tasks.

This is one of the more popular Os-baked-into-the-application approaches, and is widely used in a variety of real-time embedded systems.

4.5 Hybrid cooperative/preemptive multi-tasking

The final variation on the superloop design utilizes software-triggered interrupts to simulate a hybrid cooperative/preemptive multitasking system. Consider the example code below.

```
// Bitfields used to represent high-priority tasks. Tasks in this group
// can preempt tasks in the group below - but not eachother.
#define HP_TASK_1      (0x01)
#define HP_TASK_2      (0x02)

volatile K_UCHAR hp_tasks = 0;

// Bitfields used to represent low-priority tasks.
#define LP_TASK_1      (0x01)
#define LP_TASK_2      (0x02)

volatile K_UCHAR lp_tasks = 0;

// Interrupt sources, used to trigger both high and low priority tasks.
__interrupt__ System_Interrupt_1(void)
{
    // Set any of the other tasks from here...
    hp_tasks |= HP_TASK_1;
    // Trigger the SWI that calls the High_Priority_Tasks interrupt handler
    SWI();
}

__interrupt__ System_Interrupt_n...(void)
{
    // Set any of the other tasks from here...
}

// Interrupt handler that is used to implement the high-priority event context
__interrupt__ High_Priority_Tasks(void)
{
    // Enabled every interrupt except this one
    Disable_My_Interrupt();
    Enable_Interrupts();
    while( hp_tasks)
    {
        if( hp_tasks & HP_TASK_1)
        {
            HP_Task1();
            hp_tasks &= ~HP_TASK_1;
        }
        else if (hp_tasks & HP_TASK_2)
        {
            HP_Task2();
            hp_tasks &= ~HP_TASK_2;
        }
    }
    Restore_Interrupts();
    Enable_My_Interrupt();
}
```

```
// Main loop, used to implement the low-priority events
int main(void)
{
    // Set the function to run when a SWI is triggered
    Set_SWI(High_Priority_Tasks);

    // Run our super-loop
    while(1)
    {
        while (lp_tasks)
        {
            if (lp_tasks & LP_TASK_1)
            {
                LP_Task1();
                lp_tasks &= ~LP_TASK_1;
            }
            else if (lp_tasks & LP_TASK_2)
            {
                LP_Task2();
                lp_tasks &= ~LP_TASK_2;
            }
        }
        Idle();
    }
}
```

In this example, `High_Priority_Tasks()` can be triggered at any time as a result of a software interrupt (SWI). When a high-priority event is set, the code that sets the event calls the SWI as well, which instantly preempts whatever is happening in main, switching to the high-priority interrupt handler. If the CPU is executing in an interrupt handler already, the current ISR completes, at which point control is given to the high priority interrupt handler.

Once inside the HP ISR, all interrupts (except the software interrupt) are re-enabled, which allows this interrupt to be preempted by other interrupt sources, which is called interrupt nesting. As a result, we end up with two distinct execution contexts (main and `HighPriorityTasks()`), in which all tasks in the high-priority group are guaranteed to preempt main() tasks, and will run to completion before returning control back to tasks in main(). This is a very basic preemptive multitasking scenario, approximating a "real" RTOS system with two threads of different priorities.

4.6 Problems with superloops

As mentioned earlier, a lot of real-world systems are implemented using a superloop design; and while they are simple to understand due to the limited and obvious control logic involved, they are not without their problems.

Hidden Costs

It's difficult to calculate the overhead of the superloop and the code required to implement workarounds for blocking calls, scheduling, and preemption. There's a cost in both the logic used to implement workarounds (usually involving state machines), as well as a cost to maintainability that comes with breaking up into chunks based on execution time instead of logical operations. In moderate firmware systems, this size cost can exceed the overhead of a reasonably well-featured RTOS, and the deficit in maintainability is something that is measurable in terms of lost productivity through debugging and profiling.

Tightly-coupled code

Because the control logic is integrated so closely with the application logic, a lot of care must be taken not to compromise the separation between application and system code. The timing loops, state machines, and architecture-specific control mechanisms used to avoid (or simulate) preemption can all contribute to the problem. As a result, a lot of superloop code ends up being difficult to port without effectively simulating or replicating the underlying system for which the application was written. Abstraction layers can mitigate the risks, but a lot of care should be taken to fully decouple the application code from the system code.

No blocking calls

In a super-loop environment, there's no such thing as a blocking call or blocking objects. Tasks cannot stop mid-execution for event-driven I/O from other contexts - they must always run to completion. If busy-waiting and polling are used as a substitute, it increases latency and wastes cycles. As a result, extra code complexity is often times necessary to work-around this lack of blocking objects, often times through implementing additional state machines. In a large enough system, the added overhead in code size and cycles can add up.

Difficult to guarantee responsiveness

Without multiple levels of priority, it may be difficult to guarantee a certain degree of real-time responsiveness without added profiling and tweaking. The latency of a given task in a priority-based cooperative multitasking system is the length of the longest task. Care must be taken to break tasks up into appropriate sized chunks in order to ensure that higher-priority tasks can run in a timely fashion - a manual process that must be repeated as new tasks are added in the system. Once again, this adds extra complexity that makes code larger, more difficult to understand and maintain due to the artificial subdivision of tasks into time-based components.

Limited preemption capability

As shown in the example code, the way to gain preemption in a superloop is through the use of nested interrupts. While this isn't unwieldy for two levels of priority, adding more levels beyond this becomes complicated. In this case, it becomes necessary to track interrupt nesting manually, and separate sets of tasks that can run within given priority loops - and deadlock becomes more difficult to avoid.

Chapter 5

Mark3 Overview

5.1 Intro

The following section details the overall design of Mark3, the goals I've set out to achieve, the features that I've intended to provide, as well as an introduction to the programming concepts used to make it happen.

5.2 Features

Mark3 is a fully-featured real-time kernel, and is feature-competitive with other open-source and commercial RTOS's in the embedded arena.

The key features of this RTOS are:

- Flexible [Scheduler](#)
 - Unlimited number of threads with 8 priority levels
 - Unlimited threads per priority level
 - Round-robin scheduling for threads at each priority level
 - Time quantum scheduling for each thread in a given priority level
- Configurable stacks for each [Thread](#)
- Resource protection:
 - Integrated mutual-exclusion semaphores ([Mutex](#))
 - Priority-inheritance on [Mutex](#) objects to prevent priority inversion
- Synchronization Objects
 - Binary and counting [Semaphore](#) to coordinate thread execution
 - Event flags with 16-bit bitfields for complex thread synchronization
- Efficient Timers
 - The RTOS is tickless, the OS only wakes up when a timer expires, not at a regular interval
 - One-shot and periodic timers with event callbacks
 - Timers are high-precision and long-counting (about 68000 seconds when used with a 16us resolution timer)
- [Driver](#) API
 - A hardware abstraction layer is provided to simplify driver development

- Robust Interprocess Communications
 - Threadsafe global [Message](#) pool and configurable message queues
- Support for kernel-aware simulation
 - Provides advanced test and verification functionality, allowing for easy integration into continuous-integration systems
 - Provide accurate engineering data on key metrics like stack usage and realtime performance, with easy-to-use APIs and little overhead

5.3 Design Goals

Lightweight

Mark3 can be configured to have an extremely low static memory footprint. Each thread is defined with its own stack, and each thread structure can be configured to take as little as 26 bytes of RAM. The complete Mark3 kernel with all features, setup code, a serial driver, and the Mark3 protocol libraries comes in at under 9K of code space and 1K of RAM on atmel AVR.

Modular

Each system feature can be enabled or disabled by modifying the kernel configuration header file. Include what you want, and ignore the rest to save code space and RAM.

Easily Portable

Mark3 should be portable to a variety of 8, 16 and 32 bit architectures without MMUs. Porting the OS to a new architecture is relatively straightforward, requiring only device-specific implementations for the lowest-level operations such as context switching and timer setup.

Easy To Use

Mark3 is small by design - which gives it the advantage that it's also easy to develop for. This manual, the code itself, and the Doxygen documentation in the code provide ample documentation to get you up to speed quickly. Because you get to see the source, there's nothing left to assumption.

Simple to Understand

Not only is the Mark3 API rigorously documented (hey - that's what this book is for!), but the architecture and naming conventions are intuitive - it's easy to figure out where code lives, and how it works. Individual modules are small due to the "one feature per file" rule used in development. This makes Mark3 an ideal platform for learning about aspects of RTOS design.

Chapter 6

Getting Started

6.1 Kernel Setup

This section details the process of defining threads, initializing the kernel, and adding threads to the scheduler.

If you're at all familiar with real-time operating systems, then these setup and initialization steps should be familiar. I've tried very hard to ensure that as much of the heavy lifting is hidden from the user, so that only the bare minimum of calls are required to get things started.

The examples presented in this chapter are real, working examples taken from the ATmega328p port.

First, you'll need to create the necessary data structures and functions for the threads:

1. Create a [Thread](#) object for all of the "root" or "initial" tasks.
2. Allocate stacks for each of the Threads
3. Define an entry-point function for each [Thread](#)

This is shown in the example code below:

```
//-----  
#include "thread.h"  
#include "kernel.h"  
  
//1) Create a thread object for all of the "root" or "initial" tasks  
static Thread AppThread;  
static Thread IdleThread;  
  
//2) Allocate stacks for each thread  
#define STACK_SIZE_APP      (192)  
#define STACK_SIZE_IDLE     (128)  
  
static K_UCHAR aucAppStack[STACK_SIZE_APP];  
static K_UCHAR aucIdleStack[STACK_SIZE_IDLE];  
  
//3) Define entry point functions for each thread  
void AppThread(void);  
void IdleThread(void);
```

Next, we'll need to add the required kernel initialization code to main. This consists of running the [Kernel's](#) init routine, initializing all of the threads we defined, adding the threads to the scheduler, and finally calling [Kernel::Start\(\)](#), which transfers control of the system to the RTOS.

These steps are illustrated in the following example.

```
int main(void)  
{  
    //1) Initialize the kernel prior to use  
    Kernel::Init();           // MUST be before other kernel ops  
  
    //2) Initialize all of the threads we've defined
```

```

AppThread.Init( aucAppStack,      // Pointer to the stack
                STACK_SIZE_APP,   // Size of the stack
                1,                // Thread priority
                (void*)AppEntry,   // Entry function
                NULL );           // Entry function argument

IdleThread.Init( aucIdleStack,    // Pointer to the stack
                 STACK_SIZE_IDLE, // Size of the stack
                 0,               // Thread priority
                 (void*)IdleEntry, // Entry function
                 NULL );          // Entry function argument

//3) Add the threads to the scheduler
AppThread.Start();               // Actively schedule the threads
IdleThread.Start();

//4) Give control of the system to the kernel
Kernel::Start();                // Start the kernel!
}

```

Not much to it, is there? There are a few noteworthy points in this code, though.

In order for the kernel to work properly, a system must always contain an idle thread; that is, a thread at priority level 0 that never blocks. This thread is responsible for performing any of the low-level power management on the CPU in order to maximize battery life in an embedded device. The idle thread must also never block, and it must never exit. Either of these operations will cause undefined behavior in the system.

The App thread is at a priority level greater-than 0. This ensures that as long as the App thread has something useful to do, it will be given control of the CPU. In this case, if the app thread blocks, control will be given back to the Idle thread, which will put the CPU into a power-saving mode until an interrupt occurs.

Stack sizes must be large enough to accommodate not only the requirements of the threads, but also the requirements of interrupts - up to the maximum interrupt-nesting level used. Stack overflows are super-easy to run into in an embedded system; if you encounter strange and unexplained behavior in your code, chances are good that one of your threads is blowing its stack.

6.2 Threads

Mark3 Threads act as independent tasks in the system. While they share the same address-space, global data, device-drivers, and system peripherals, each thread has its own set of CPU registers and stack, collectively known as the thread's **context**. The context is what allows the RTOS kernel to rapidly switch between threads at a high rate, giving the illusion that multiple things are happening in a system, when really, only one thread is executing at a time.

6.2.1 Thread Setup

Each instance of the [Thread](#) class represents a thread, its stack, its CPU context, and all of the state and metadata maintained by the kernel. Before a [Thread](#) will be scheduled to run, it must first be initialized with the necessary configuration data.

The Init function gives the user the opportunity to set the stack, stack size, thread priority, entry-point function, entry-function argument, and round-robin time quantum:

[Thread](#) stacks are pointers to blobs of memory (usually K_CHAR arrays) carved out of the system's address space. Each thread must have a stack defined that's large enough to handle not only the requirements of local variables in the thread's code path, but also the maximum depth of the ISR stack.

Priorities should be chosen carefully such that the shortest tasks with the most strict determinism requirements are executed first - and are thus located in the highest priorities. Tasks that take the longest to execute (and require the least degree of responsiveness) must occupy the lower thread priorities. The idle thread must be the only thread occupying the lowest priority level.

The thread quantum only applies when there are multiple threads in the ready queue at the same priority level. This interval is used to kick-off a timer that will cycle execution between the threads in the priority list so that they each get a fair chance to execute.

The entry function is the function that the kernel calls first when the thread instance is first started. Entry functions have at most one argument - a pointer to a data-object specified by the user during initialization.

An example thread initialization is shown below:

```
Thread clMyThread;
K_UCHAR aucStack[192];

void AppEntry(void)
{
    while(1)
    {
        // Do something!
    }
}

...
{
    clMyThread.Init(aucStack,    // Pointer to the stack to use by this thread
                    192,        // Size of the stack in bytes
                    1,          // Thread priority (0 = idle, 7 = max)
                    (void*)AppEntry, // Function where the thread starts executing
                    NULL );      // Argument passed into the entry function
}
```

Once a thread has been initialized, it can be added to the scheduler by calling:

```
clMyThread.Start();
```

The thread will be placed into the [Scheduler's](#) queue at the designated priority, where it will wait its turn for execution.

6.2.2 Entry Functions

Mark3 Threads should not run-to-completion - they should execute as infinite loops that perform a series of tasks, appropriately partitioned to provide the responsiveness characteristics desired in the system.

The most basic [Thread](#) loop is shown below:

```
void Thread( void *param )
{
    while(1)
    {
        // Do Something
    }
}
```

Threads can interact with eachother in the system by means of synchronization objects ([Semaphore](#)), mutual-exclusion objects ([Mutex](#)), Inter-process messaging ([MessageQueue](#)), and timers ([Timer](#)).

Threads can suspend their own execution for a predetermined period of time by using the static [Thread::Sleep\(\)](#) method. Calling this will block the [Thread's](#) executin until the amount of time specified has ellapsed. Upon expiry, the thread will be placed back into the ready queue for its priority level, where it awaits its next turn to run.

6.3 Timers

[Timer](#) objects are used to trigger callback events periodic or on a one-shot (alarm) basis.

While extremely simple to use, they provide one of the most powerful execution contexts in the system. The timer callbacks execute from within the timer callback ISR in an interrupt-enabled context. As such, timer callbacks are considered higher-priority than any thread in the system, but lower priority than other interrupts. Care must be taken to ensure that timer callbacks execute as quickly as possible to minimize the impact of processing on the throughput of tasks in the system. Wherever possible, heavy-lifting should be deferred to the threads by way of semaphores or messages.

Below is an example showing how to start a periodic system timer which will trigger every second:

```

{
    Timer clTimer;
    clTimer.Init();

    clTimer.Start( 1000,
                  1,
                  MyCallback,
                  (void*)&my_data );

    ... // Keep doing work in the thread
}

// Callback function, executed from the timer-expiry context.
void MyCallback( Thread *pclOwner_, void *pvData_ )
{
    LED.Flash(); // Flash an LED.
}

```

6.4 Semaphores

Semaphores are used to synchronized execution of threads based on the availability (and quantity) of application-specific resources in the system. They are extremely useful for solving producer-consumer problems, and are the method-of-choice for creating efficient, low latency systems, where ISRs post semaphores that are handled from within the context of individual threads. (Yes, Semaphores can be posted - but not pended - from the interrupt context).

The following is an example of the producer-consumer usage of a binary semaphore:

```

Semaphore clSemaphore; // Declare a semaphore shared between a producer and a consumer thread.

void Producer()
{
    clSemaphore.Init(0, 1);
    while(1)
    {
        // Do some work, create something to be consumed

        // Post a semaphore, allowing another thread to consume the data
        clSemaphore.Post();
    }
}

void Consumer()
{
    // Assumes semaphore initialized before use...
    While(1)
    {
        // Wait for new data from the producer thread
        clSemaphore.Pend();

        // Consume the data!
    }
}

```

And an example of using semaphores from the ISR context to perform event- driven processing.

```

Semaphore clSemaphore;

__interrupt__ MyISR()
{
    clSemaphore.Post(); // Post the interrupt. Lightweight when uncontested.
}

void MyThread()
{
    clSemaphore.Init(0, 1); // Ensure this is initialized before the MyISR interrupt is enabled.
    while(1)
    {
        // Wait until we get notification from the interrupt
        clSemaphore.Pend();

        // Interrupt has fired, do the necessary work in this thread's context
        HeavyLifting();
    }
}

```

6.5 Mutexes

Mutexes (Mutual exclusion objects) are provided as a means of creating "protected sections" around a particular resource, allowing for access of these objects to be serialized. Only one thread can hold the mutex at a time - other threads have to wait until the region is released by the owner thread before they can take their turn operating on the protected resource. Note that mutexes can only be owned by threads - they are not available to other contexts (i.e. interrupts). Calling the mutex APIs from an interrupt will cause catastrophic system failures.

Note that these objects are also not recursive- that is, the owner thread can not attempt to claim a mutex more than once.

Priority inheritance is provided with these objects as a means to avoid priority inversions. Whenever a thread at a priority than the mutex owner blocks on a mutex, the priority of the current thread is boosted to the highest-priority waiter to ensure that other tasks at intermediate priorities cannot artificially prevent progress from being made.

Mutex objects are very easy to use, as there are only three operations supported: Initialize, Claim and Release. An example is shown below.

```

Mutex clMutex; // Create a mutex globally.

void Init()
{
    // Initialize the mutex before use.
    clMutex.Init();
}

// Some function called from a thread
void Thread1Function()
{
    clMutex.Claim();

    // Once the mutex is owned, no other thread can
    // enter a block protect by the same mutex

    my_protected_resource.do_something();
    my_protected_resource.do_something_else();

    clMutex.Release();
}

// Some function called from another thread
void Thread2Function()
{
    clMutex.Claim();

    // Once the mutex is owned, no other thread can
    // enter a block protect by the same mutex

    my_protected_resource.do_something();
    my_protected_resource.do_different_things();

    clMutex.Release();
}

```

6.6 Event Flags

Event Flags are another synchronization object, conceptually similar to a semaphore.

Unlike a semaphore, however, the condition on which threads are unblocked is determined by a more complex set of rules. Each Event Flag object contains a 16-bit field, and threads block, waiting for combinations of bits within this field to become set.

A thread can wait on any pattern of bits from this field to be set, and any number of threads can wait on any number of different patterns. Threads can wait on a single bit, multiple bits, or bits from within a subset of bits within the field.

As a result, setting a single value in the flag can result in any number of threads becoming unblocked simultaneously. This mechanism is extremely powerful, allowing for all sorts of complex, yet efficient, thread synchronization schemes that can be created using a single shared object.

Note that Event Flags can be set from interrupts, but you cannot wait on an event flag from within an interrupt.

Examples demonstrating the use of event flags are shown below.

```
// Simple example showing a thread blocking on a multiple bits in the
// fields within an event flag.

EventFlag clEventFlag;

int main()
{
    ...
    clEventFlag.Init(); // Initialize event flag prior to use
    ...
}

void MyInterrupt()
{
    // Some interrupt corresponds to event 0x0020
    clEventFlag.Set(0x0020);
}

void MyThreadFunc()
{
    ...
    while(1)
    {
        ...
        K_USHORT usWakeCondition;

        // Allow this thread to block on multiple flags
        usWakeCondition = clEventFlag.Wait(0x00FF, EVENT_FLAG_ANY);

        // Clear the event condition that caused the thread to wake (in this case,
        // usWakeCondition will equal 0x20 when triggered from the interrupt above)
        clEventFlag.Clear(usWakeCondition);

        // <do something>
    }
}
```

6.7 Messages

Sending messages between threads is the key means of synchronizing access to data, and the primary mechanism to perform asynchronous data processing operations.

Sending a message consists of the following operations:

- Obtain a [Message](#) object from the global message pool
- Set the message data and event fields
- Send the message to the destination message queue

While receiving a message consists of the following steps:

- Wait for a messages in the destination message queue
- Process the message data
- Return the message back to the global message pool

These operations, and the various data objects involved are discussed in more detail in the following section.

6.7.1 Message Objects

[Message](#) objects are used to communicate arbitrary data between threads in a safe and synchronous way.

The message object consists of an event code field and a data field. The event code is used to provide context to the message object, while the data field (essentially a void * data pointer) is used to provide a payload of data corresponding to the particular event.

Access to these fields is marshalled by accessors - the transmitting thread uses the `SetData()` and `SetCode()` methods to seed the data, while the receiving thread uses the `GetData()` and `GetCode()` methods to retrieve it.

By providing the data as a void data pointer instead of a fixed-size message, we achieve an unprecedented measure of simplicity and flexibility. Data can be either statically or dynamically allocated, and sized appropriately for the event without having to format and reformat data by both sending and receiving threads. The choices here are left to the user - and the kernel doesn't get in the way of efficiency.

It is worth noting that you can send messages to message queues from within ISR context. This helps maintain consistency, since the same APIs can be used to provide event-driven programming facilities throughout the whole of the OS.

6.7.2 Global Message Pool

To maintain efficiency in the messaging system (and to prevent over-allocation of data), a global pool of message objects is provided. The size of this message pool is specified in the implementation, and can be adjusted depending on the requirements of the target application as a compile-time option.

Allocating a message from the message pool is as simple as calling the `GlobalMessagePool::Pop()` Method.

Messages are returned back to the `GlobalMessagePool::Push()` method once the message contents are no longer required.

One must be careful to ensure that discarded messages always are returned to the pool, otherwise a resource leak can occur, which may cripple the operating system's ability to pass data between threads.

6.7.3 Message Queues

`Message` objects specify data with context, but do not specify where the messages will be sent. For this purpose we have a `MessageQueue` object. Sending an object to a message queue involves calling the `MessageQueue::Send()` method, passing in a pointer to the `Message` object as an argument.

When a message is sent to the queue, the first thread blocked on the queue (as a result of calling the `MessageQueue::Receive()` method) will wake up, with a pointer to the `Message` object returned.

It's worth noting that multiple threads can block on the same message queue, providing a means for multiple threads to share work in parallel.

6.7.4 Messaging Example

```
// Message queue object shared between threads
MessageQueue cMsgQ;

// Function that initializes the shared message queue
void MsgQInit()
{
    cMsgQ.Init();
}

// Function called by one thread to send message data to
// another
void TxMessage()
{
    // Get a message, initialize its data
    Message *pclMsg = GlobalMessagePool::Pop();

    pclMsg->SetCode(0xAB);
    pclMsg->SetData((void*)some_data);

    // Send the data to the message queue
    cMsgQ.Send(pclMsg);
}

// Function called in the other thread to block until
// a message is received in the message queue.
void RxMessage()
{
    Message *pclMsg;
```

```

// Block until we have a message in the queue
pclMsg = clMsgQ.Receive();

// Do something with the data once the message is received
pclMsg->GetCode();

// Free the message once we're done with it.
GlobalMessagePool::Push(pclMsg);
}

```

6.8 Sleep

There are instances where it may be necessary for a thread to poll a resource, or wait a specific amount of time before proceeding to operate on a peripheral or volatile piece of data.

While the [Timer](#) object is generally a better choice for performing time-sensitive operations (and certainly a better choice for periodic operations), the [Thread::Sleep\(\)](#) method provides a convenient (and efficient) mechanism that allows for a thread to suspend its execution for a specified interval.

Note that when a thread is sleeping it is blocked, during which other threads can operate, or the system can enter its idle state.

```

int GetPeripheralData()
{
    int value;
    // The hardware manual for a peripheral specifies that
    // the "foo()" method will result in data being generated
    // that can be captured using the "bar()" method.
    // However, the value only becomes valid after 10ms

    peripheral.foo();
    Thread::Sleep(10); // Wait 10ms for data to become valid
    value = peripheral.bar();
    return value;
}

```

6.9 Round-Robin Quantum

Threads at the same thread priority are scheduled using a round-robin scheme. Each thread is given a timeslice (which can be configured) of which it shares time amongst ready threads in the group. Once a thread's timeslice has expired, the next thread in the priority group is chosen to run until its quantum has expired - the cycle continues over and over so long as each thread has work to be done.

By default, the round-robin interval is set at 4ms.

This value can be overridden by calling the thread's [SetQuantum\(\)](#) with a new interval specified in milliseconds.

Chapter 7

Build System

Mark3 is distributed with a recursive makefile build system, allowing the entire source tree to be built into a series of libraries with simple make commands.

The way the scripts work, every directory with a valid makefile is scanned, as well as all of its subdirectories. The build then generates binary components for all of the components it finds -libraries and executables. All libraries that are generated can then be imported into an application using the linker without having to copy-and-paste files on a module-by-module basis. Applications built during this process can then be loaded onto a device directly, without requiring a GUI-based IDE. As a result, Mark3 integrates well with 3rd party tools for continuous-integration and automated testing.

This modular framework allows for large volumes of libraries and binaries to be built at once - the default build script leverages this to build all of the examples and unit tests at once, linking against the pre-built kernel, services, and drivers. Whatever can be built as a library is built as a library, promoting reuse throughout the platform, and enabling Mark3 to be used as a platform, with an ecosystem of libraries, services, drivers and applications.

7.1 Source Layout

One key aspect of Mark3 is that system features are organized into their own separate modules. These modules are further grouped together into folders based on the type of features represented:

Root	Base folder, contains recursive makefiles for build system
arduino	Arduino-specific headers and API documentation files
bootloader	Mark3 Bootloader code for AVR microcontrollers
build	Makefiles and device-configuraton data for various platforms
docs	Documentation (including this)
drivers	Device driver code for various supported devices
example	Example applications
fonts	Bitmap fonts converted from TTF, used by Mark3 graphics library
kernel	Basic Mark3 Components (the focus of this manual)
cpu	CPU-specific porting code
scripts	Scripts used to simplify build, documentation, and profiling
services	Utility code and services, extended system features
stage	Staging directory, where the build system places artifacts
tests	Unit tests, written as C/C++ applications
util	.net-based utils: font conversion, terminal, programmer, and configuration

7.2 Building the kernel

The base.mak file determines how the kernel, drivers, and libraries are built, for what targets, and with what options. Most of these options can be copied directly from the options found in your IDE managed projects. Below is an overview of the main variables used to configure the build.

STAGE - Location in the filesystem where the build output is stored

```

ROOT_DIR    - The location of the root source tree
ARCH        - The CPU architecture to build against
VARIANT     - The variant of the above CPU to target
TOOLCHAIN   - Which toolchain to build with (dependent on ARCH and VARIANT)

```

Build.mak contains the logic which is used to perform the recursive make in all directories. Unless you really know what you're doing, it's best to leave this as-is.

You must make sure that all required paths are set in your system environment variables so that they are accessible through from the command-line.

Once configured, you can build the source tree using the various make targets:

- make headers
 - copy all headers in each module's /public subdirectory to the location specified by STAGE environment variable's ./inc subdirectory.
- make library
 - regenerate all objects copy marked as libraries (i.e. the kernel + drivers). Resulting binaries are copied into STAGE's ./lib subdirectory.
- make binary
 - build all executable projects in the root directory structure. In the default distribution, this includes the basic set of demos.

These steps are chained together automatically as part of the build.sh script found under the /scripts subdirectory. Running ./scripts/build.sh from the root of the embedded source directory will result in all headers being exported, libraries built, and applications built. This script will also default to building for atmega328p using GCC if none of the required environment variables have previously been configured.

To add new components to the recursive build system, simply add your code into a new folder beneath the root install location.

Source files, the module makefile and private header files go directly in the new folder, while public headers are placed in a ./public subdirectory. Create a ./obj directory to hold the output from the builds.

The contents of the module makefile looks something like this:

```

# Include common prelude make file
include $(ROOT_DIR)base.mak

# If we're building a library, set IS_LIB and LIBNAME
# If we're building an app, set IS_APP and APPNAME
IS_LIB=1
LIBNAME=mylib

#this is the list of the source modules required to build the kernel
CPP_SOURCE = mylib.cpp \
             someotherfile.cpp

# Similarly, C-language source would be under the C_SOURCE variable.

# Include the rest of the script that is actually used for building the
# outputs
include $(ROOT_DIR)build.mak

```

Once you've placed your code files in the right place, and configured the makefile appropriately, a fresh call to make headers, make library, then make binary will guarantee that your code is built.

Now, you can still copy-and-paste the required kernel, port, and drivers, directly into your application avoiding the whole process of using make from the command line. To do this, run "make source" from the root directory in svn, and copy the contents of /stage/src into your project. This should contain the source to the kernel, all drivers, and all services that are in the tree - along with the necessary header files.

7.3 Building on Windows

Building Mark3 on Windows is the same as on Linux, but there are a few prerequisites that need to be taken into consideration before the build scripts and makefiles will work as expected.

Step 1 - Install Latest Atmel Studio IDE

Atmel Studio contains the AVR8 GCC toolchain, which contains the necessary compilers, assemblers, and platform support required to turn the source modules into libraries and executables.

To get Atmel Studio, go to the Atmel website (<http://www.atmel.com>) and register to download the latest version. This is a free download (and rather large). The included IDE (if you choose to use it) is very slick, as it's based on Visual Studio, and contains a wonderful cycle-accurate simulator for AVR devices. In fact, the simulator is so good that most of the kernel and its drivers were developed using this tool.

Once you have downloaded and installed Atmel Studio, you will need to add the location of the AVR toolchain to the PATH environment variable.

To do this, go to Control Panel -> System and Security -> System -> Advanced System Settings, and edit the PATH variable. Append the location of the toolchain bin folder to the end of the variable.

On Windows 7 x64, it should look something like this:

C: Files (x86) Toolchain GCC\Native\3.4.2.1002-gnu-toolchain

Step 2 - Install MinGW and MinSys

MinGW (and MinSys in particular) provide a unix-like environment that runs under windows. Some of the utilities provided include a version of the bash shell, and GNU standard make - both which are required by the Mark3 recursive build system.

The MinGW installer can be downloaded from its project page on SourceForge. When installing, be sure to select the "MinSys" component.

Once installed, add the MinSys binary path to the PATH environment variable, in a similar fashion as with Atmel Studio in Step 1.

Step 3 - Setup Include Paths in Platform Makefile

The AVR header file path must be added to the "platform.mak" makefile for each AVR Target you are attempting to build for. These files can be located under /embedded/build/avr/atmegaXXX/. The path to the includes directory should be added to the end of the CFLAGS and CPPFLAGS variables, as shown in the following:

```
TEST_INC="/c/Program Files (x86)/Atmel/Atmel Toolchain/AVR8
GCC/Native/3.4.2.1002/avr8-gnu-toolchain/include"
CFLAGS += -I$(TEST_INC)
CPPFLAGS += -I$(TEST_INC)
```

Step 4 - Build Mark3 using Bash

Launch a terminal to your Mark3 base directory, and cd into the "embedded" folder. You should now be able to build Mark3 by running "bash ./build.sh" from the command-line.

Alternately, you can run bash itself, building Mark3 by running ./build.sh or the various make targets using the same syntax as documented previously.

Note - building on Windows is *slow*. This has a lot to do with how "make" performs under windows. There are faster substitutes for make (such as cs-make) that are exponentially quicker, and approach the performance of make on Linux. Other mechanisms, such as running make with multiple concurrent jobs (i.e. "make -j4") also helps significantly, especially on systems with multicore CPUs.

Chapter 8

License

8.1 License

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Chapter 9

Profiling Results

The following profiling results were obtained using an ATmega328p @ 16MHz.

The test cases are designed to make use of the kernel profiler, which accurately measures the performance of the fundamental system APIs, in order to provide information for user comparison, as well as to ensure that regressions are not being introduced into the system.

9.1 Date Performed

Wed May 27 21:18:21 EDT 2015

9.2 Compiler Information

The kernel and test code used in these results were built using the following compiler: `./scripts/profile.sh: 55: ./scripts/profile.sh: /home/moslevin/atmel/bin/avr-gcc: not found`

9.3 Profiling Results

- Semaphore Initialization: cycles (averaged over 0 iterations)
- Semaphore Post (uncontested): cycles (averaged over 0 iterations)
- Semaphore Pend (uncontested): cycles (averaged over 0 iterations)
- Semaphore Flyback Time (Contested Pend): cycles (averaged over 0 iterations)
- Mutex Init: cycles (averaged over 0 iterations)
- Mutex Claim: cycles (averaged over 0 iterations)
- Mutex Release: cycles (averaged over 0 iterations)
- Thread Initialize: cycles (averaged over 0 iterations)
- Thread Start: cycles (averaged over 0 iterations)
- Context Switch: cycles (averaged over 0 iterations)
- Thread Schedule: cycles (averaged over 0 iterations)

Chapter 10

Code Size Profiling

The following report details the size of each module compiled into the kernel.

The size of each component is dependent on the flags specified in [mark3cfg.h](#) at compile time. Note that these sizes represent the maximum size of each module before dead code elimination and any additional link-time optimization, and represent the maximum possible size that any module can take.

The results below are for profiling on Atmel AVR atmega328p-based targets using gcc. Results are not necessarily indicative of relative or absolute performance on other platforms or toolchains.

10.1 Information

Subversion Repository Information:

- Repository Root: `svn+ssh://m0slevin.code.sf.net/p/mark3/source`
- Revision: 211
- URL: `svn+ssh://m0slevin.code.sf.net/p/mark3/source/branch/release/R1/embedded` Relative URL: `^/branch/release/R1/embedded`

Date Profiled: Mon May 18 22:24:26 EDT 2015

10.2 Compiler Version

avr-gcc (GCC) 4.8.2 Copyright (C) 2013 Free Software Foundation, Inc. This is free software; see the source for copying conditions. There is NO warranty; not even for MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE.

10.3 Profiling Results

Mark3 Module Size Report:

- Synchronization Objects - Base Class..... : 92 Bytes
- Device [Driver](#) Framework (including /dev/null)... : 226 Bytes
- Synchronization Object - Event Flag..... : 770 Bytes
- Fundamental [Kernel](#) Linked-List Classes..... : 496 Bytes

- Message-based IPC..... : 426 Bytes
- [Mutex](#) (Synchronization Object)..... : 658 Bytes
- Performance-profiling timers..... : 546 Bytes
- Round-Robin Scheduling Support..... : 264 Bytes
- [Thread](#) Scheduling..... : 452 Bytes
- [Semaphore](#) (Synchronization Object)..... : 536 Bytes
- [Thread](#) Implementation..... : 1613 Bytes
- Fundamental [Kernel](#) Thread-list Data Structures.. : 210 Bytes
- Mark3 [Kernel](#) Base Class..... : 110 Bytes
- Software [Timer Kernel](#) Object..... : 378 Bytes
- Software [Timer](#) Management..... : 645 Bytes
- Runtime [Kernel](#) Trace Implementation..... : 0 Bytes
- Circular Logging Buffer Base Class..... : 0 Bytes
- Atmel AVR - [Kernel](#) Aware Simulation Support..... : 296 Bytes
- Atmel AVR - Basic Threading Support..... : 598 Bytes
- Atmel AVR - [Kernel](#) Interrupt Implemenation..... : 56 Bytes
- Atmel AVR - [Kernel Timer](#) Implementation..... : 322 Bytes
- kernelprofile.cpp.o : 256 Bytes

Mark3 [Kernel](#) Size Summary:

- [Kernel](#) : 2973 Bytes
- Synchronization Objects : 2390 Bytes
- Port : 1272 Bytes
- Features : 2059 Bytes
- Total Size : 8694 Bytes

Chapter 11

Hierarchical Index

11.1 Class Hierarchy

This inheritance list is sorted roughly, but not completely, alphabetically:

BlockingObject	43
EventFlag	53
Mutex	74
Semaphore	84
DriverList	52
FakeThread_t	56
GlobalMessagePool	57
Kernel	58
KernelAware	60
KernelAwareData_t	64
KernelSWI	64
KernelTimer	65
LinkList	67
CircularLinkList	44
ThreadList	96
DoubleLinkList	48
TimerList	103
LinkListNode	69
Driver	49
DevNull	45
Message	70
Thread	88
Timer	99
MessageQueue	72
Profiler	76
ProfileTimer	77
Quantum	79
Scheduler	81
ThreadPort	98
TimerScheduler	105

Chapter 12

Class Index

12.1 Class List

Here are the classes, structs, unions and interfaces with brief descriptions:

BlockingObject	Class implementing thread-blocking primitives	43
CircularLinkedList	Circular-linked-list data type, inherited from the base LinkedList type	44
DevNull	This class implements the "default" driver (/dev/null)	45
DoubleLinkedList	Doubly-linked-list data type, inherited from the base LinkedList type	48
Driver	Base device-driver class used in hardware abstraction	49
DriverList	List of Driver objects used to keep track of all device drivers in the system	52
EventFlag	Blocking object, similar to a semaphore or mutex, commonly used for synchronizing thread execution based on events occurring within the system	53
FakeThread_t	If the kernel is set up to use an idle function instead of an idle thread, we use a placeholder data structure to "simulate" the effect of having an idle thread in the system	56
GlobalMessagePool	Implements a list of message objects shared between all threads	57
Kernel	Class that encapsulates all of the kernel startup functions	58
KernelAware	The KernelAware class	60
KernelAwareData_t	This structure is used to communicate between the kernel and a kernel-aware host	64
KernelSWI	Class providing the software-interrupt required for context-switching in the kernel	64
KernelTimer	Hardware timer interface, used by all scheduling/timer subsystems	65
LinkedList	Abstract-data-type from which all other linked-lists are derived	67
LinkedListNode	Basic linked-list node data structure	69
Message	Class to provide message-based IPC services in the kernel	70
MessageQueue	List of messages, used as the channel for sending and receiving messages between threads	72

Mutex	Mutual-exclusion locks, based on BlockingObject	74
Profiler	System profiling timer interface	76
ProfileTimer	Profiling timer	77
Quantum	Static-class used to implement Thread quantum functionality, which is a key part of round-robin scheduling	79
Scheduler	Priority-based round-robin Thread scheduling, using ThreadLists for housekeeping	81
Semaphore	Counting semaphore, based on BlockingObject base class	84
Thread	Object providing fundamental multitasking support in the kernel	88
ThreadList	This class is used for building thread-management facilities, such as schedulers, and blocking objects	96
ThreadPort	Class defining the architecture specific functions required by the kernel	98
Timer	Timer - an event-driven execution context based on a specified time interval	99
TimerList	TimerList class - a doubly-linked-list of timer objects	103
TimerScheduler	"Static" Class used to interface a global TimerList with the rest of the kernel	105

Chapter 13

File Index

13.1 File List

Here is a list of all documented files with brief descriptions:

/home/moslevin/Project/R1/kernel/ atomic.cpp	
Basic Atomic Operations	107
/home/moslevin/Project/R1/kernel/ blocking.cpp	
Implementation of base class for blocking objects	109
/home/moslevin/Project/R1/kernel/ driver.cpp	
Device driver/hardware abstraction layer	125
/home/moslevin/Project/R1/kernel/ eventflag.cpp	
Event Flag Blocking Object/IPC-Object implementation	127
/home/moslevin/Project/R1/kernel/ kernel.cpp	
Kernel initialization and startup code	131
/home/moslevin/Project/R1/kernel/ kernelaware.cpp	
Kernel aware simulation support	133
/home/moslevin/Project/R1/kernel/ ksemaphore.cpp	
Semaphore Blocking-Object Implemenation	136
/home/moslevin/Project/R1/kernel/ ll.cpp	
Core Linked-List implementation, from which all kernel objects are derived	139
/home/moslevin/Project/R1/kernel/ message.cpp	
Inter-thread communications via message passing	141
/home/moslevin/Project/R1/kernel/ mutex.cpp	
Mutual-exclusion object	144
/home/moslevin/Project/R1/kernel/ profile.cpp	
Code profiling utilities	148
/home/moslevin/Project/R1/kernel/ quantum.cpp	
Thread Quantum Implementation for Round-Robin Scheduling	194
/home/moslevin/Project/R1/kernel/ scheduler.cpp	
Strict-Priority + Round-Robin thread scheduler implementation	197
/home/moslevin/Project/R1/kernel/ thread.cpp	
Platform-Independent thread class Definition	199
/home/moslevin/Project/R1/kernel/ threadlist.cpp	
Thread linked-list definitions	204
/home/moslevin/Project/R1/kernel/ timer.cpp	
Timer implementations	206
/home/moslevin/Project/R1/kernel/ timerlist.cpp	
Implements timer list processing algorithms, responsible for all timer tick and expiry logic	208
/home/moslevin/Project/R1/kernel/ tracebuffer.cpp	
Kernel trace buffer class definition	211
/home/moslevin/Project/R1/kernel/ writebuf16.cpp	
16 bit circular buffer implementation with callbacks	212

/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/ kernelprofile.cpp	
ATMega328p Profiling timer implementation	110
/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/ kernelswi.cpp	
Kernel Software interrupt implementation for ATMega328p	112
/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/ kerneltimer.cpp	
Kernel Timer Implementation for ATMega328p	113
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ATMega328p Multithreading	122
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Profiling timer hardware interface	115
/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/ kernelswi.h	
Kernel Software interrupt declarations	116
/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/ kerneltimer.h	
Kernel Timer Class declaration	117
/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/ threadport.h	
ATMega328p Multithreading support	119
/home/moslevin/Project/R1/kernel/public/ atomic.h	
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Hex codes/translation tables used for efficient string tokenization	152
/home/moslevin/Project/R1/kernel/public/ driver.h	
Driver abstraction framework	154
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Event Flag Blocking Object/IPC-Object definition	156
/home/moslevin/Project/R1/kernel/public/ kernel.h	
Kernel initialization and startup class	157
/home/moslevin/Project/R1/kernel/public/ kernelaware.h	
Kernel aware simulation support	159
/home/moslevin/Project/R1/kernel/public/ kerneldebug.h	
Macros and functions used for assertions, kernel traces, etc	160
/home/moslevin/Project/R1/kernel/public/ kerneltypes.h	
Basic data type primitives used throughout the OS	163
/home/moslevin/Project/R1/kernel/public/ ksemaphore.h	
Semaphore Blocking Object class declarations	164
/home/moslevin/Project/R1/kernel/public/ ll.h	
Core linked-list declarations, used by all kernel list types	166
/home/moslevin/Project/R1/kernel/public/ manual.h	
Ascii-format documentation, used by doxygen to create various printable and viewable forms	167
/home/moslevin/Project/R1/kernel/public/ mark3.h	
Single include file given to users of the Mark3 Kernel API	168
/home/moslevin/Project/R1/kernel/public/ mark3cfg.h	
Mark3 Kernel Configuration	172
/home/moslevin/Project/R1/kernel/public/ message.h	
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/home/moslevin/Project/R1/kernel/public/ paniccodes.h	
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Timer scheduler declarations	191
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Kernel trace buffer class declaration	192
/home/moslevin/Project/R1/kernel/public/writebuf16.h	
Thread-safe circular buffer implementation with 16-bit elements	193

Chapter 14

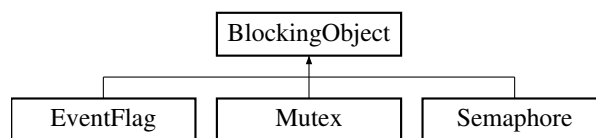
Class Documentation

14.1 BlockingObject Class Reference

Class implementing thread-blocking primitives.

```
#include <blocking.h>
```

Inheritance diagram for BlockingObject:



Protected Member Functions

- void [Block](#) ([Thread](#) *[pclThread_](#))
- void [UnBlock](#) ([Thread](#) *[pclThread_](#))

Protected Attributes

- [ThreadList](#) [m_clBlockList](#)
[ThreadList](#) which is used to hold the list of threads blocked on a given object.

14.1.1 Detailed Description

Class implementing thread-blocking primitives.

Used for implementing things like semaphores, mutexes, message queues, or anything else that could cause a thread to suspend execution on some external stimulus.

Definition at line 65 of file [blocking.h](#).

14.1.2 Member Function Documentation

14.1.2.1 void [BlockingObject::Block](#) ([Thread](#) * [pclThread_](#)) [protected]

Parameters

<code>pciThread_</code>	Pointer to the thread object that will be blocked.
-------------------------	--

Blocks a thread on this object. This is the fundamental operation performed by any sort of blocking operation in the operating system. All semaphores/mutexes/sleeping/messaging/etc ends up going through the blocking code at some point as part of the code that manages a transition from an "active" or "waiting" thread to a "blocked" thread.

The steps involved in blocking a thread (which are performed in the function itself) are as follows;

1) Remove the specified thread from the current owner's list (which is likely one of the scheduler's thread lists) 2) Add the thread to this object's thread list 3) Setting the thread's "current thread-list" point to reference this object's threadlist.

Definition at line 36 of file [blocking.cpp](#).

14.1.2.2 void BlockingObject::UnBlock (Thread * `pciThread_`) [protected]

Parameters

<code>pciThread_</code>	Pointer to the thread to unblock.
-------------------------	-----------------------------------

Unblock a thread that is already blocked on this object, returning it to the "ready" state by performing the following steps:

1) Removing the thread from this object's threadlist 2) Restoring the thread to its "original" owner's list

Definition at line 52 of file [blocking.cpp](#).

The documentation for this class was generated from the following files:

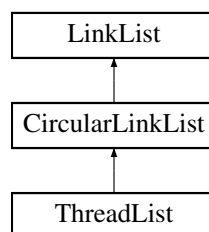
- [/home/moslevin/Project/R1/kernel/public/blocking.h](#)
- [/home/moslevin/Project/R1/kernel/blocking.cpp](#)

14.2 CircularLinkedList Class Reference

Circular-linked-list data type, inherited from the base [LinkedList](#) type.

```
#include <ll.h>
```

Inheritance diagram for CircularLinkedList:



Public Member Functions

- virtual void [Add](#) ([LinkedListNode](#) *node_)
Add the linked list node to this linked list.
- virtual void [Remove](#) ([LinkedListNode](#) *node_)
Add the linked list node to this linked list.
- void [PivotForward](#) ()
Pivot the head of the circularly linked list forward (Head = Head->next, Tail = Tail->next)
- void [PivotBackward](#) ()
Pivot the head of the circularly linked list backward (Head = Head->prev, Tail = Tail->prev)

Additional Inherited Members

14.2.1 Detailed Description

Circular-linked-list data type, inherited from the base [LinkedList](#) type.

Definition at line 196 of file [ll.h](#).

14.2.2 Member Function Documentation

14.2.2.1 `void CircularLinkedList::Add (LinkedListNode * node_) [virtual]`

Add the linked list node to this linked list.

Parameters

<i>node_</i>	Pointer to the node to add
--------------	----------------------------

Implements [LinkedList](#).

Reimplemented in [ThreadList](#).

Definition at line 102 of file [ll.cpp](#).

14.2.2.2 `void CircularLinkedList::Remove (LinkedListNode * node_) [virtual]`

Add the linked list node to this linked list.

Parameters

<i>node_</i>	Pointer to the node to remove
--------------	-------------------------------

Implements [LinkedList](#).

Reimplemented in [ThreadList](#).

Definition at line 127 of file [ll.cpp](#).

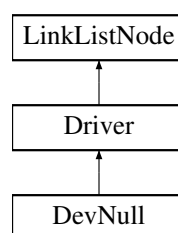
The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/ll.h](#)
- [/home/moslevin/Project/R1/kernel/ll.cpp](#)

14.3 DevNull Class Reference

This class implements the "default" driver (/dev/null)

Inheritance diagram for DevNull:



Public Member Functions

- virtual void [Init](#) ()
Initialize a driver, must be called prior to use.
- virtual [K_UCHAR Open](#) ()
Open a device driver prior to use.
- virtual [K_UCHAR Close](#) ()
Close a previously-opened device driver.
- virtual [K_USHORT Read](#) ([K_USHORT](#) usBytes_, [K_UCHAR](#) *pucData_)
Read a specified number of bytes from the device into a specific buffer.
- virtual [K_USHORT Write](#) ([K_USHORT](#) usBytes_, [K_UCHAR](#) *pucData_)
Write a payload of data of a given length to the device.
- virtual [K_USHORT Control](#) ([K_USHORT](#) usEvent_, void *pvDataIn_, [K_USHORT](#) usSizeIn_, void *pvDataOut_, [K_USHORT](#) usSizeOut_)
This is the main entry-point for device-specific io and control operations.

Additional Inherited Members

14.3.1 Detailed Description

This class implements the "default" driver (/dev/null)

Definition at line 40 of file [driver.cpp](#).

14.3.2 Member Function Documentation

14.3.2.1 virtual [K_UCHAR](#) DevNull::Close () [inline],[virtual]

Close a previously-opened device driver.

Returns

Driver-specific return code, 0 = OK, non-0 = error

Implements [Driver](#).

Definition at line 45 of file [driver.cpp](#).

14.3.2.2 virtual [K_USHORT](#) DevNull::Control ([K_USHORT](#) usEvent_, void * pvDataIn_, [K_USHORT](#) usSizeIn_, void * pvDataOut_, [K_USHORT](#) usSizeOut_) [inline],[virtual]

This is the main entry-point for device-specific io and control operations.

This is used for implementing all "side-channel" communications with a device, and any device-specific IO operations that do not conform to the typical POSIX read/write paradigm. Use of this function is analagous to the non-POSIX (yet still common) devctl() or ioctl().

Parameters

<i>usEvent_</i>	Code defining the io event (driver-specific)
<i>pvDataIn_</i>	Pointer to the input data
<i>usSizeIn_</i>	Size of the input data (in bytes)

<i>pvDataOut_</i>	Pointer to the output data
<i>usSizeOut_</i>	Size of the output data (in bytes)

Returns

Driver-specific return code, 0 = OK, non-0 = error

Implements [Driver](#).

Definition at line 53 of file [driver.cpp](#).

14.3.2.3 `virtual K_UCHAR DevNull::Open () [inline],[virtual]`

Open a device driver prior to use.

Returns

Driver-specific return code, 0 = OK, non-0 = error

Implements [Driver](#).

Definition at line 44 of file [driver.cpp](#).

14.3.2.4 `virtual K_USHORT DevNull::Read (K_USHORT usBytes_, K_UCHAR * pucData_) [inline],[virtual]`

Read a specified number of bytes from the device into a specific buffer.

Depending on the driver-specific implementation, this may be a number less than the requested number of bytes read, indicating that there there was less input than desired, or that as a result of buffering, the data may not be available.

Parameters

<i>usBytes_</i>	Number of bytes to read (<= size of the buffer)
<i>pucData_</i>	Pointer to a data buffer receiving the read data

Returns

Number of bytes actually read

Implements [Driver](#).

Definition at line 47 of file [driver.cpp](#).

14.3.2.5 `virtual K_USHORT DevNull::Write (K_USHORT usBytes_, K_UCHAR * pucData_) [inline],[virtual]`

Write a payload of data of a given length to the device.

Depending on the implementation of the driver, the amount of data written to the device may be less than the requested number of bytes. A result less than the requested size may indicate that the device buffer is full, indicating that the user must retry the write at a later point with the remaining data.

Parameters

<i>usBytes_</i>	Number of bytes to write (<= size of the buffer)
<i>pucData_</i>	Pointer to a data buffer containing the data to write

Returns

Number of bytes actually written

Implements [Driver](#).

Definition at line 50 of file [driver.cpp](#).

The documentation for this class was generated from the following file:

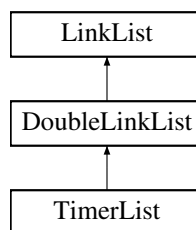
- [/home/moslevin/Project/R1/kernel/driver.cpp](#)

14.4 DoubleLinkedList Class Reference

Doubly-linked-list data type, inherited from the base [LinkedList](#) type.

```
#include <ll.h>
```

Inheritance diagram for DoubleLinkedList:



Public Member Functions

- [DoubleLinkedList](#) ()
Default constructor - initializes the head/tail nodes to NULL.
- virtual void [Add](#) ([LinkedListNode](#) *node_)
Add the linked list node to this linked list.
- virtual void [Remove](#) ([LinkedListNode](#) *node_)
Add the linked list node to this linked list.

Additional Inherited Members

14.4.1 Detailed Description

Doubly-linked-list data type, inherited from the base [LinkedList](#) type.

Definition at line 165 of file [ll.h](#).

14.4.2 Member Function Documentation

14.4.2.1 void [DoubleLinkedList::Add](#) ([LinkedListNode](#) * node_) [virtual]

Add the linked list node to this linked list.

Parameters

<code>node_</code>	Pointer to the node to add
--------------------	----------------------------

Implements [LinkedList](#).

Definition at line 41 of file [ll.cpp](#).

14.4.2.2 `void DoubleLinkedList::Remove (LinkedListNode * node_) [virtual]`

Add the linked list node to this linked list.

Parameters

<code>node_</code>	Pointer to the node to remove
--------------------	-------------------------------

Implements [LinkedList](#).

Definition at line 65 of file [ll.cpp](#).

The documentation for this class was generated from the following files:

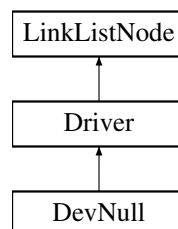
- [/home/moslevin/Project/R1/kernel/public/ll.h](#)
- [/home/moslevin/Project/R1/kernel/ll.cpp](#)

14.5 Driver Class Reference

Base device-driver class used in hardware abstraction.

```
#include <driver.h>
```

Inheritance diagram for Driver:



Public Member Functions

- virtual void [Init](#) ()=0
Initialize a driver, must be called prior to use.
- virtual [K_UCHAR](#) [Open](#) ()=0
Open a device driver prior to use.
- virtual [K_UCHAR](#) [Close](#) ()=0
Close a previously-opened device driver.
- virtual [K_USHORT](#) [Read](#) ([K_USHORT](#) usBytes_, [K_UCHAR](#) *pucData_)=0
Read a specified number of bytes from the device into a specific buffer.
- virtual [K_USHORT](#) [Write](#) ([K_USHORT](#) usBytes_, [K_UCHAR](#) *pucData_)=0
Write a payload of data of a given length to the device.
- virtual [K_USHORT](#) [Control](#) ([K_USHORT](#) usEvent_, void *pvDataIn_, [K_USHORT](#) usSizeIn_, void *pvDataOut_, [K_USHORT](#) usSizeOut_)=0
This is the main entry-point for device-specific io and control operations.

- void [SetName](#) (const [K_CHAR](#) *pcName_)
Set the path for the driver.
- const [K_CHAR](#) * [GetPath](#) ()
Returns a string containing the device path.

Private Attributes

- const [K_CHAR](#) * [m_pcPath](#)
string pointer that holds the driver path (name)

Additional Inherited Members

14.5.1 Detailed Description

Base device-driver class used in hardware abstraction.

All other device drivers inherit from this class

Definition at line 121 of file [driver.h](#).

14.5.2 Member Function Documentation

14.5.2.1 [K_UCHAR Driver::Close](#) () [pure virtual]

Close a previously-opened device driver.

Returns

Driver-specific return code, 0 = OK, non-0 = error

Implemented in [DevNull](#).

14.5.2.2 [K_USHORT Driver::Control](#) ([K_USHORT usEvent_](#), void * [pvDataIn_](#), [K_USHORT usSizeIn_](#), void * [pvDataOut_](#), [K_USHORT usSizeOut_](#)) [pure virtual]

This is the main entry-point for device-specific io and control operations.

This is used for implementing all "side-channel" communications with a device, and any device-specific IO operations that do not conform to the typical POSIX read/write paradigm. Use of this function is analagous to the non-POSIX (yet still common) `devctl()` or `ioctl()`.

Parameters

usEvent_	Code defining the io event (driver-specific)
pvDataIn_	Pointer to the input data
usSizeIn_	Size of the input data (in bytes)
pvDataOut_	Pointer to the output data
usSizeOut_	Size of the output data (in bytes)

Returns

Driver-specific return code, 0 = OK, non-0 = error

Implemented in [DevNull](#).

14.5.2.3 `const K_CHAR * Driver::GetPath () [inline]`

Returns a string containing the device path.

Returns

`pcName_` Return the string constant representing the device path

Definition at line 231 of file [driver.h](#).

14.5.2.4 `K_UCHAR Driver::Open () [pure virtual]`

Open a device driver prior to use.

Returns

Driver-specific return code, 0 = OK, non-0 = error

Implemented in [DevNull](#).

14.5.2.5 `K_USHORT Driver::Read (K_USHORT usBytes_, K_UCHAR * pucData_) [pure virtual]`

Read a specified number of bytes from the device into a specific buffer.

Depending on the driver-specific implementation, this may be a number less than the requested number of bytes read, indicating that there was less input than desired, or that as a result of buffering, the data may not be available.

Parameters

<i>usBytes_</i>	Number of bytes to read (<= size of the buffer)
<i>pucData_</i>	Pointer to a data buffer receiving the read data

Returns

Number of bytes actually read

Implemented in [DevNull](#).

14.5.2.6 `void Driver::SetName (const K_CHAR * pcName_) [inline]`

Set the path for the driver.

Name must be set prior to access (since driver access is name-based).

Parameters

<i>pcName_</i>	String constant containing the device path
----------------	--

Definition at line 222 of file [driver.h](#).

14.5.2.7 `K_USHORT Driver::Write (K_USHORT usBytes_, K_UCHAR * pucData_) [pure virtual]`

Write a payload of data of a given length to the device.

Depending on the implementation of the driver, the amount of data written to the device may be less than the requested number of bytes. A result less than the requested size may indicate that the device buffer is full, indicating that the user must retry the write at a later point with the remaining data.

Parameters

<i>usBytes_</i>	Number of bytes to write (<= size of the buffer)
<i>pucData_</i>	Pointer to a data buffer containing the data to write

Returns

Number of bytes actually written

Implemented in [DevNull](#).

The documentation for this class was generated from the following file:

- [/home/moslevin/Project/R1/kernel/public/driver.h](#)

14.6 DriverList Class Reference

List of [Driver](#) objects used to keep track of all device drivers in the system.

```
#include <driver.h>
```

Static Public Member Functions

- static void [Init](#) ()
Initialize the list of drivers.
- static void [Add](#) ([Driver](#) *pclDriver_)
Add a [Driver](#) object to the managed global driver-list.
- static void [Remove](#) ([Driver](#) *pclDriver_)
Remove a driver from the global driver list.
- static [Driver](#) * [FindByPath](#) (const [K_CHAR](#) *m_pcPath)
Look-up a driver in the global driver-list based on its path.

Static Private Attributes

- static [DoubleLinkedList](#) [m_clDriverList](#)
LinkedList object used to implementing the driver object management.

14.6.1 Detailed Description

List of [Driver](#) objects used to keep track of all device drivers in the system.

By default, the list contains a single entity, "/dev/null".

Definition at line [244](#) of file [driver.h](#).

14.6.2 Member Function Documentation

14.6.2.1 [DriverList::Add](#) ([Driver](#) * *pclDriver_*) `[inline], [static]`

Add a [Driver](#) object to the managed global driver-list.

Parameters

<code>pciDriver_</code>	pointer to the driver object to add to the global driver list.
-------------------------	--

Definition at line 264 of file [driver.h](#).

14.6.2.2 `Driver * DriverList::FindByPath (const K_CHAR * m_pcPath) [static]`

Look-up a driver in the global driver-list based on its path.

In the event that the driver is not found in the list, a pointer to the default "/dev/null" object is returned. In this way, unimplemented drivers are automatically stubbed out.

Definition at line 107 of file [driver.cpp](#).

14.6.2.3 `void DriverList::Init () [static]`

Initialize the list of drivers.

Must be called prior to using the device driver library.

Definition at line 98 of file [driver.cpp](#).

14.6.2.4 `void DriverList::Remove (Driver * pciDriver_) [inline],[static]`

Remove a driver from the global driver list.

Parameters

<code>pciDriver_</code>	Pointer to the driver object to remove from the global table
-------------------------	--

Definition at line 274 of file [driver.h](#).

The documentation for this class was generated from the following files:

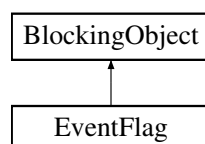
- [/home/moslevin/Project/R1/kernel/public/driver.h](#)
- [/home/moslevin/Project/R1/kernel/driver.cpp](#)

14.7 EventFlag Class Reference

The [EventFlag](#) class is a blocking object, similar to a semaphore or mutex, commonly used for synchronizing thread execution based on events occurring within the system.

```
#include <eventflag.h>
```

Inheritance diagram for EventFlag:



Public Member Functions

- `void Init ()`
Init initializes the [EventFlag](#) object prior to use.
- `K_USHORT Wait (K_USHORT usMask_, EventFlagOperation_t eMode_)`

Wait - Block a thread on the specific flags in this event flag group.

- [K_USHORT Wait](#) ([K_USHORT](#) usMask_, [EventFlagOperation_t](#) eMode_, [K_ULONG](#) ulTimeMS_)

Wait - Block a thread on the specific flags in this event flag group.

- void [WakeMe](#) ([Thread](#) *pclOwner_)

WakeMe.

- void [Set](#) ([K_USHORT](#) usMask_)

Set - Set additional flags in this object (logical OR).

- void [Clear](#) ([K_USHORT](#) usMask_)

ClearFlags - Clear a specific set of flags within this object, specific by bitmask.

- [K_USHORT GetMask](#) ()

GetMask Returns the state of the 16-bit bitmask within this object.

Private Member Functions

- [K_USHORT Wait_i](#) ([K_USHORT](#) usMask_, [EventFlagOperation_t](#) eMode_, [K_ULONG](#) ulTimeMS_)

Wait_i.

Private Attributes

- [K_USHORT m_usSetMask](#)

Event flags currently set in this object.

Additional Inherited Members

14.7.1 Detailed Description

The [EventFlag](#) class is a blocking object, similar to a semaphore or mutex, commonly used for synchronizing thread execution based on events occurring within the system.

Each [EventFlag](#) object contains a 16-bit bitmask, which is used to trigger events on associated threads. Threads wishing to block, waiting for a specific event to occur can wait on any pattern within this 16-bit bitmask to be set. Here, we provide the ability for a thread to block, waiting for ANY bits in a specified mask to be set, or for ALL bits within a specific mask to be set. Depending on how the object is configured, the bits that triggered the wakeup can be automatically cleared once a match has occurred.

Definition at line 46 of file [eventflag.h](#).

14.7.2 Member Function Documentation

14.7.2.1 void EventFlag::Clear ([K_USHORT](#) usMask_)

ClearFlags - Clear a specific set of flags within this object, specific by bitmask.

Parameters

usMask_	- Bitmask of flags to clear
-------------------------	-----------------------------

Definition at line 295 of file [eventflag.cpp](#).

14.7.2.2 [K_USHORT](#) EventFlag::GetMask ()

GetMask Returns the state of the 16-bit bitmask within this object.

Returns

The state of the 16-bit bitmask

Definition at line 304 of file [eventflag.cpp](#).

14.7.2.3 void EventFlag::Set (K_USHORT *usMask_*)

Set - Set additional flags in this object (logical OR).

This API can potentially result in threads blocked on [Wait\(\)](#) to be unblocked.

Parameters

<i>usMask_</i>	- Bitmask of flags to set.
----------------	----------------------------

Definition at line 176 of file [eventflag.cpp](#).

14.7.2.4 K_USHORT EventFlag::Wait (K_USHORT *usMask_*, EventFlagOperation_t *eMode_*)

Wait - Block a thread on the specific flags in this event flag group.

Parameters

<i>usMask_</i>	- 16-bit bitmask to block on
<i>eMode_</i>	- EVENT_FLAG_ANY: Thread will block on any of the bits in the mask <ul style="list-style-type: none"> • EVENT_FLAG_ALL: Thread will block on all of the bits in the mask

Returns

Bitmask condition that caused the thread to unblock, or 0 on error or timeout

Definition at line 158 of file [eventflag.cpp](#).

14.7.2.5 K_USHORT EventFlag::Wait (K_USHORT *usMask_*, EventFlagOperation_t *eMode_*, K_ULONG *ulTimeMS_*)

Wait - Block a thread on the specific flags in this event flag group.

Parameters

<i>usMask_</i>	- 16-bit bitmask to block on
<i>eMode_</i>	- EVENT_FLAG_ANY: Thread will block on any of the bits in the mask <ul style="list-style-type: none"> • EVENT_FLAG_ALL: Thread will block on all of the bits in the mask
<i>ulTimeMS_</i>	- Time to block (in ms)

Returns

Bitmask condition that caused the thread to unblock, or 0 on error or timeout

Definition at line 169 of file [eventflag.cpp](#).

14.7.2.6 K_USHORT EventFlag::Wait_i (K_USHORT *usMask_*, EventFlagOperation_t *eMode_*, K_ULONG *ulTimeMS_*) [private]

Wait_i.

Internal abstraction used to manage both timed and untimed wait operations

Parameters

<i>usMask_</i>	- 16-bit bitmask to block on
<i>eMode_</i>	- EVENT_FLAG_ANY: Thread will block on any of the bits in the mask • EVENT_FLAG_ALL: Thread will block on all of the bits in the mask
<i>ulTimeMS_</i>	- Time to block (in ms)

Returns

Bitmask condition that caused the thread to unblock, or 0 on error or timeout

! If the Yield operation causes a new thread to be chosen, there will ! Be a context switch at the above [CS_EXIT\(\)](#). The original calling ! thread will not return back until a matching SetFlags call is made ! or a timeout occurs.

Definition at line 65 of file [eventflag.cpp](#).

14.7.2.7 void EventFlag::WakeMe (Thread * *pclOwner_*)

WakeMe.

Wake the given thread, currently blocking on this object

Parameters

<i>pclOwner_</i>	Pointer to the owner thread to unblock.
------------------	---

Definition at line 57 of file [eventflag.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/eventflag.h](#)
- [/home/moslevin/Project/R1/kernel/eventflag.cpp](#)

14.8 FakeThread_t Struct Reference

If the kernel is set up to use an idle function instead of an idle thread, we use a placeholder data structure to "simulate" the effect of having an idle thread in the system.

```
#include <thread.h>
```

Public Attributes

- [K_WORD * m_pwStackTop](#)
Pointer to the top of the thread's stack.
- [K_WORD * m_pwStack](#)
Pointer to the thread's stack.
- [K_UCHAR m_ucThreadID](#)
Thread ID.
- [K_UCHAR m_ucPriority](#)
Default priority of the thread.
- [K_UCHAR m_ucCurPriority](#)
Current priority of the thread (priority inheritance)
- [ThreadState_t m_eState](#)
Enum indicating the thread's current state.

14.8.1 Detailed Description

If the kernel is set up to use an idle function instead of an idle thread, we use a placeholder data structure to "simulate" the effect of having an idle thread in the system.

When cast to a [Thread](#), this data structure will still result in `GetPriority()` calls being valid, which is all that is needed to support the tick-based/tickless times – while saving a fairly decent chunk of RAM on a small micro.

Note that this struct must have the same memory layout as the [Thread](#) class up to the last item.

Definition at line 494 of file [thread.h](#).

The documentation for this struct was generated from the following file:

- [/home/moslevin/Project/R1/kernel/public/thread.h](#)

14.9 GlobalMessagePool Class Reference

Implements a list of message objects shared between all threads.

```
#include <message.h>
```

Static Public Member Functions

- static void [Init](#) ()
Initialize the message queue prior to use.
- static void [Push](#) ([Message](#) *pclMessage_)
Return a previously-claimed message object back to the global queue.
- static [Message](#) * [Pop](#) ()
Pop a message from the global queue, returning it to the user to be populated before sending by a transmitter.

Static Private Attributes

- static [Message](#) [m_aclMessagePool](#) [[GLOBAL_MESSAGE_POOL_SIZE](#)]
Array of message objects that make up the message pool.
- static [DoubleLinkedList](#) [m_clList](#)
Linked list used to manage the [Message](#) objects.

14.9.1 Detailed Description

Implements a list of message objects shared between all threads.

Definition at line 157 of file [message.h](#).

14.9.2 Member Function Documentation

14.9.2.1 [Message](#) * [GlobalMessagePool::Pop](#) () [static]

Pop a message from the global queue, returning it to the user to be populated before sending by a transmitter.

Returns

Pointer to a [Message](#) object

Definition at line 70 of file [message.cpp](#).

14.9.2.2 void GlobalMessagePool::Push (Message * pclMessage_) [static]

Return a previously-claimed message object back to the global queue.

Used once the message has been processed by a receiver.

Parameters

<code>pclMessage_</code>	Pointer to the Message object to return back to the global queue
--------------------------	--

Definition at line 58 of file [message.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/message.h](#)
- [/home/moslevin/Project/R1/kernel/message.cpp](#)

14.10 Kernel Class Reference

Class that encapsulates all of the kernel startup functions.

```
#include <kernel.h>
```

Static Public Member Functions

- static void [Init](#) (void)
Kernel Initialization Function, call before any other OS function.
- static void [Start](#) (void)
Start the kernel; function never returns.
- static [K_BOOL](#) [IsStarted](#) ()
IsStarted.
- static void [SetPanic](#) ([panic_func_t](#) pfPanic_)
SetPanic Set a function to be called when a kernel panic occurs, giving the user to determine the behavior when a catastrophic failure is observed.
- static [K_BOOL](#) [IsPanic](#) ()
IsPanic Returns whether or not the kernel is in a panic state.
- static void [Panic](#) ([K_USHORT](#) usCause_)
Panic Cause the kernel to enter its panic state.
- static void [SetIdleFunc](#) ([idle_func_t](#) pfIdle_)
SetIdleFunc Set the function to be called when no active threads are available to be scheduled by the scheduler.
- static void [IdleFunc](#) (void)
IdleFunc Call the low-priority idle function when no active threads are available to be scheduled.
- static [Thread](#) * [GetIdleThread](#) (void)
GetIdleThread Return a pointer to the Kernel's idle thread object to the user.

Static Private Attributes

- static [K_BOOL](#) [m_bIsStarted](#)
true if kernel is running, false otherwise
- static [K_BOOL](#) [m_bIsPanic](#)
true if kernel is in panic state, false otherwise
- static [panic_func_t](#) [m_pfPanic](#)
user-set panic function
- static [idle_func_t](#) [m_pfIdle](#)

user-set idle function

- static [FakeThread_t m_clIdle](#)

Idle thread object (note: not a real thread)

14.10.1 Detailed Description

Class that encapsulates all of the kernel startup functions.

Definition at line 48 of file [kernel.h](#).

14.10.2 Member Function Documentation

14.10.2.1 static Thread* Kernel::GetIdleThread(void) [inline],[static]

GetIdleThread Return a pointer to the [Kernel](#)'s idle thread object to the user.

Note that the [Thread](#) object involved is to be used for comparisons only – the thread itself is "virtual", and doesn't represent a unique execution context with its own stack.

Returns

Pointer to the [Kernel](#)'s idle thread object

Definition at line 124 of file [kernel.h](#).

14.10.2.2 Kernel::Init(void) [static]

[Kernel](#) Initialization Function, call before any other OS function.

Initializes all global resources used by the operating system. This must be called before any other kernel function is invoked.

Definition at line 52 of file [kernel.cpp](#).

14.10.2.3 static K_BOOL Kernel::IsPanic() [inline],[static]

IsPanic Returns whether or not the kernel is in a panic state.

Returns

Whether or not the kernel is in a panic state

Definition at line 95 of file [kernel.h](#).

14.10.2.4 static K_BOOL Kernel::IsStarted() [inline],[static]

IsStarted.

Returns

Whether or not the kernel has started - true = running, false = not started

Definition at line 80 of file [kernel.h](#).

14.10.2.5 void Kernel::Panic(K_USHORT usCause_) [static]

Panic Cause the kernel to enter its panic state.

Parameters

<i>usCause_</i>	Reason for the kernel panic
-----------------	-----------------------------

Definition at line 95 of file [kernel.cpp](#).

14.10.2.6 `static void Kernel::SetIdleFunc (idle_func_t pfIdle_) [inline],[static]`

SetIdleFunc Set the function to be called when no active threads are available to be scheduled by the scheduler.

Parameters

<i>pfIdle_</i>	Pointer to the idle function
----------------	------------------------------

Definition at line 109 of file [kernel.h](#).

14.10.2.7 `static void Kernel::SetPanic (panic_func_t pfPanic_) [inline],[static]`

SetPanic Set a function to be called when a kernel panic occurs, giving the user to determine the behavior when a catastrophic failure is observed.

Parameters

<i>pfPanic_</i>	Panic function pointer
-----------------	------------------------

Definition at line 89 of file [kernel.h](#).

14.10.2.8 `Kernel::Start (void) [static]`

Start the kernel; function never returns.

Start the operating system kernel - the current execution context is cancelled, all kernel services are started, and the processor resumes execution at the entrypoint for the highest-priority thread.

You must have at least one thread added to the kernel before calling this function, otherwise the behavior is undefined.

Definition at line 86 of file [kernel.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/kernel.h](#)
- [/home/moslevin/Project/R1/kernel/kernel.cpp](#)

14.11 KernelAware Class Reference

The [KernelAware](#) class.

```
#include <kernelaware.h>
```

Static Public Member Functions

- static void [ProfileInit](#) (const [K_CHAR](#) *szStr_)
ProfileInit.
- static void [ProfileStart](#) (void)
ProfileStart.
- static void [ProfileStop](#) (void)
ProfileStop.

- static void [ProfileReport](#) (void)
ProfileReport.
- static void [ExitSimulator](#) (void)
ExitSimulator.
- static void [Print](#) (const [K_CHAR](#) *szStr_)
Print.
- static void [Trace](#) ([K_USHORT](#) usFile_, [K_USHORT](#) usLine_, [K_USHORT](#) usCode_)
Trace.
- static void [Trace](#) ([K_USHORT](#) usFile_, [K_USHORT](#) usLine_, [K_USHORT](#) usCode_, [K_USHORT](#) usArg1_)
Trace.
- static void [Trace](#) ([K_USHORT](#) usFile_, [K_USHORT](#) usLine_, [K_USHORT](#) usCode_, [K_USHORT](#) usArg1_, [K_USHORT](#) usArg2_)
Trace.
- static [K_BOOL](#) [IsSimulatorAware](#) (void)
IsSimulatorAware.

Static Private Member Functions

- static void [Trace_i](#) ([K_USHORT](#) usFile_, [K_USHORT](#) usLine_, [K_USHORT](#) usCode_, [K_USHORT](#) usArg1_, [K_USHORT](#) usArg2_, [KernelAwareCommand_t](#) eCmd_)
Trace_i.

14.11.1 Detailed Description

The [KernelAware](#) class.

This class contains functions that are used to trigger kernel-aware functionality within a supported simulation environment (i.e. fIAVR).

These static methods operate on a singleton set of global variables, which are monitored for changes from within the simulator. The simulator hooks into these variables by looking for the correctly-named symbols in an elf-formatted binary being run and registering callbacks that are called whenever the variables are changed. On each change of the command variable, the kernel-aware data is analyzed and interpreted appropriately.

If these methods are run in an unsupported simulator or on actual hardware the commands generally have no effect (except for the exit-on-reset command, which will result in a jump-to-0 reset).

Definition at line 65 of file [kernelaware.h](#).

14.11.2 Member Function Documentation

14.11.2.1 void [KernelAware::ExitSimulator](#) (void) [static]

[ExitSimulator.](#)

Instruct the kernel-aware simulator to terminate (destroying the virtual CPU).

Definition at line 104 of file [kernelaware.cpp](#).

14.11.2.2 [K_BOOL](#) [KernelAware::IsSimulatorAware](#) (void) [static]

[IsSimulatorAware.](#)

Use this function to determine whether or not the code is running on a simulator that is aware of the kernel.

Returns

true - the application is being run in a kernel-aware simulator. false - otherwise.

Definition at line 164 of file [kernelaware.cpp](#).

14.11.2.3 void KernelAware::Print (const K_CHAR * szStr_) [static]

Print.

Instruct the kernel-aware simulator to print a char string

Parameters

<i>szStr_</i>	
---------------	--

Definition at line 155 of file [kernelaware.cpp](#).

14.11.2.4 void KernelAware::ProfileInit (const K_CHAR * szStr_) [static]

ProfileInit.

Initializes the kernel-aware profiler. This function instructs the kernel-aware simulator to reset its accounting variables, and prepare to start counting profiling data tagged to the given string. How this is handled is the responsibility of the simulator.

Parameters

<i>szStr_</i>	String to use as a tag for the profiling session.
---------------	---

Definition at line 77 of file [kernelaware.cpp](#).

14.11.2.5 void KernelAware::ProfileReport (void) [static]

ProfileReport.

Instruct the kernel-aware simulator to print a report for its current profiling data.

Definition at line 98 of file [kernelaware.cpp](#).

14.11.2.6 void KernelAware::ProfileStart (void) [static]

ProfileStart.

Instruct the kernel-aware simulator to begin counting cycles towards the current profiling counter.

Definition at line 86 of file [kernelaware.cpp](#).

14.11.2.7 void KernelAware::ProfileStop (void) [static]

ProfileStop.

Instruct the kernel-aware simulator to end counting cycles relative to the current profiling counter's iteration.

Definition at line 92 of file [kernelaware.cpp](#).

14.11.2.8 void KernelAware::Trace (K_USHORT usFile_, K_USHORT usLine_, K_USHORT usCode_) [static]

Trace.

Insert a kernel trace statement into the kernel-aware simulator's debug data stream.

Parameters

<i>usFile_</i>	16-bit code representing the file
<i>usLine_</i>	16-bit code representing the line in the file
<i>usCode_</i>	16-bit data code, which indicates the line's format.

Definition at line 110 of file [kernelaware.cpp](#).

14.11.2.9 void KernelAware::Trace (K_USHORT *usFile_*, K_USHORT *usLine_*, K_USHORT *usCode_*, K_USHORT *usArg1_*) [static]

Trace.

Insert a kernel trace statement into the kernel-aware simulator's debug data stream.

Parameters

<i>usFile_</i>	16-bit code representing the file
<i>usLine_</i>	16-bit code representing the line in the file
<i>usCode_</i>	16-bit data code, which indicates the line's format
<i>usArg1_</i>	16-bit argument to the format string.

Definition at line 118 of file [kernelaware.cpp](#).

14.11.2.10 void KernelAware::Trace (K_USHORT *usFile_*, K_USHORT *usLine_*, K_USHORT *usCode_*, K_USHORT *usArg1_*, K_USHORT *usArg2_*) [static]

Trace.

Insert a kernel trace statement into the kernel-aware simulator's debug data stream.

Parameters

<i>usFile_</i>	16-bit code representing the file
<i>usLine_</i>	16-bit code representing the line in the file
<i>usCode_</i>	16-bit data code, which indicates the line's format
<i>usArg1_</i>	16-bit argument to the format string.
<i>usArg2_</i>	16-bit argument to the format string.

Definition at line 127 of file [kernelaware.cpp](#).

14.11.2.11 void KernelAware::Trace_i (K_USHORT *usFile_*, K_USHORT *usLine_*, K_USHORT *usCode_*, K_USHORT *usArg1_*, K_USHORT *usArg2_*, KernelAwareCommand_t *eCmd_*) [static],[private]

Trace_i.

Private function by which the class's [Trace\(\)](#) methods are reflected, which allows us to realize a modest code saving.

Parameters

<i>usFile_</i>	16-bit code representing the file
<i>usLine_</i>	16-bit code representing the line in the file
<i>usCode_</i>	16-bit data code, which indicates the line's format
<i>usArg1_</i>	16-bit argument to the format string.
<i>usArg2_</i>	16-bit argument to the format string.
<i>eCmd_</i>	Code indicating the number of arguments to emit.

Definition at line 137 of file [kernelaware.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/kernelaware.h](#)

- [/home/moslevin/Project/R1/kernel/kernelaware.cpp](#)

14.12 KernelAwareData_t Union Reference

This structure is used to communicate between the kernel and a kernel- aware host.

Public Attributes

- volatile [K_USHORT](#) [ausBuffer](#) [5]
Raw binary contents of the struct.

- The [Profiler](#) struct contains data related to the code-execution profiling functionality provided by a kernel-aware host simulator.

The Trace struct contains data related to the display and output of kernel-trace strings on a kernel-aware host.

The Print struct contains data related to the display of arbitrary null-terminated ASCII strings on the kernel-aware host.

14.12.1 Detailed Description

This structure is used to communicate between the kernel and a kernel- aware host.

Its data contents is interpreted differently depending on the command executed (by means of setting the `g_ucKA-Command` variable, as is done in the command handlers in this module). As a result, any changes to this struct by way of modifying or adding data must be mirrored in the kernel-aware simulator.

Definition at line 37 of file [kernelaware.cpp](#).

The documentation for this union was generated from the following file:

- [/home/moslevin/Project/R1/kernel/kernelaware.cpp](#)

14.13 KernelSWI Class Reference

Class providing the software-interrupt required for context-switching in the kernel.

```
#include <kernelswi.h>
```

Static Public Member Functions

- static void [Config](#) (void)
Configure the software interrupt - must be called before any other software interrupt functions are called.
- static void [Start](#) (void)
Enable ("Start") the software interrupt functionality.
- static void [Stop](#) (void)
Disable the software interrupt functionality.
- static void [Clear](#) (void)

- *Clear the software interrupt.*
- static void [Trigger](#) (void)
Call the software interrupt.
- static [K_UCHAR](#) [DI](#) ()
Disable the SWI flag itself.
- static void [RI](#) ([K_BOOL](#) [bEnable_](#))
Restore the state of the SWI to the value specified.

14.13.1 Detailed Description

Class providing the software-interrupt required for context-switching in the kernel.

Definition at line 32 of file [kernelswi.h](#).

14.13.2 Member Function Documentation

14.13.2.1 [K_UCHAR](#) [KernelSWI::DI](#) () [static]

Disable the SWI flag itself.

Returns

previous status of the SWI, prior to the DI call

Definition at line 50 of file [kernelswi.cpp](#).

14.13.2.2 void [KernelSWI::RI](#) ([K_BOOL](#) [bEnable_](#)) [static]

Restore the state of the SWI to the value specified.

Parameters

bEnable_	true - enable the SWI, false - disable SWI
--------------------------	--

Definition at line 58 of file [kernelswi.cpp](#).

The documentation for this class was generated from the following files:

- /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/[kernelswi.h](#)
- /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/[kernelswi.cpp](#)

14.14 KernelTimer Class Reference

Hardware timer interface, used by all scheduling/timer subsystems.

```
#include <kerneltimer.h>
```

Static Public Member Functions

- static void [Config](#) (void)
Initializes the kernel timer before use.
- static void [Start](#) (void)
Starts the kernel time (must be configured first)
- static void [Stop](#) (void)

- Shut down the kernel timer, used when no timers are scheduled.*

 - static [K_UCHAR](#) [DI](#) (void)

Disable the kernel timer's expiry interrupt.

 - static void [RI](#) ([K_BOOL](#) bEnable_)

Retstore the state of the kernel timer's expiry interrupt.

 - static void [EI](#) (void)

Enable the kernel timer's expiry interrupt.

 - static [K_ULONG](#) [SubtractExpiry](#) ([K_ULONG](#) ullInterval_)

Subtract the specified number of ticks from the timer's expiry count register.

 - static [K_ULONG](#) [TimeToExpiry](#) (void)

Returns the number of ticks remaining before the next timer expiry.

 - static [K_ULONG](#) [SetExpiry](#) ([K_ULONG](#) ullInterval_)

Resets the kernel timer's expiry interval to the specified value.

 - static [K_ULONG](#) [GetOvertime](#) (void)

Return the number of ticks that have elapsed since the last expiry.

 - static void [ClearExpiry](#) (void)

Clear the hardware timer expiry register.

Static Private Member Functions

- static [K_USHORT](#) [Read](#) (void)

Safely read the current value in the timer register.

14.14.1 Detailed Description

Hardware timer interface, used by all scheduling/timer subsystems.

Definition at line 33 of file [kerneltimer.h](#).

14.14.2 Member Function Documentation

14.14.2.1 [K_ULONG](#) [KernelTimer::GetOvertime](#) (void) [static]

Return the number of ticks that have elapsed since the last expiry.

Returns

Number of ticks that have elapsed after last timer expiration

Definition at line 115 of file [kerneltimer.cpp](#).

14.14.2.2 [K_USHORT](#) [KernelTimer::Read](#) (void) [static], [private]

Safely read the current value in the timer register.

Returns

Value held in the timer register

Definition at line 66 of file [kerneltimer.cpp](#).

14.14.2.3 void [KernelTimer::RI](#) ([K_BOOL](#) bEnable_) [static]

Retstore the state of the kernel timer's expiry interrupt.

Parameters

<i>bEnable_</i>	1 enable, 0 disable
-----------------	---------------------

Definition at line 169 of file [kerneltimer.cpp](#).

14.14.2.4 K_ULONG KernelTimer::SetExpiry (K_ULONG ulInterval_) [static]

Resets the kernel timer's expiry interval to the specified value.

Parameters

<i>ulInterval_</i>	Desired interval in ticks to set the timer for
--------------------	--

Returns

Actual number of ticks set (may be less than desired)

Definition at line 121 of file [kerneltimer.cpp](#).

14.14.2.5 K_ULONG KernelTimer::SubtractExpiry (K_ULONG ulInterval_) [static]

Subtract the specified number of ticks from the timer's expiry count register.

Returns the new expiry value stored in the register.

Parameters

<i>ulInterval_</i>	Time (in HW-specific) ticks to subtract
--------------------	---

Returns

Value in ticks stored in the timer's expiry register

Definition at line 84 of file [kerneltimer.cpp](#).

14.14.2.6 K_ULONG KernelTimer::TimeToExpiry (void) [static]

Returns the number of ticks remaining before the next timer expiry.

Returns

Time before next expiry in platform-specific ticks

Definition at line 95 of file [kerneltimer.cpp](#).

The documentation for this class was generated from the following files:

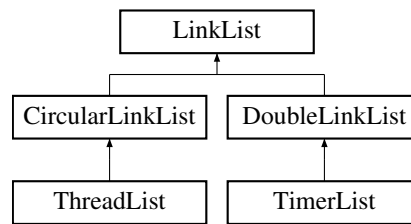
- [/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/kerneltimer.h](#)
- [/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/kerneltimer.cpp](#)

14.15 LinkList Class Reference

Abstract-data-type from which all other linked-lists are derived.

```
#include <ll.h>
```

Inheritance diagram for LinkList:



Public Member Functions

- `void Init ()`
Clear the linked list.
- `virtual void Add (LinkListNode *node_)=0`
Add the linked list node to this linked list.
- `virtual void Remove (LinkListNode *node_)=0`
Add the linked list node to this linked list.
- `LinkListNode * GetHead ()`
Get the head node in the linked list.
- `LinkListNode * GetTail ()`
Get the tail node of the linked list.

Protected Attributes

- `LinkListNode * m_pstHead`
Pointer to the head node in the list.
- `LinkListNode * m_pstTail`
Pointer to the tail node in the list.

14.15.1 Detailed Description

Abstract-data-type from which all other linked-lists are derived.

Definition at line 112 of file [ll.h](#).

14.15.2 Member Function Documentation

14.15.2.1 `void LinkList::Add (LinkListNode * node_) [pure virtual]`

Add the linked list node to this linked list.

Parameters

<code>node_</code>	Pointer to the node to add
--------------------	----------------------------

Implemented in [CircularLinkList](#), [DoubleLinkList](#), and [ThreadList](#).

14.15.2.2 `LinkListNode * LinkList::GetHead () [inline]`

Get the head node in the linked list.

Returns

Pointer to the head node in the list

Definition at line 149 of file [ll.h](#).

14.15.2.3 `LinkedListNode * LinkedList::GetTail ()` `[inline]`

Get the tail node of the linked list.

Returns

Pointer to the tail node in the list

Definition at line 158 of file [ll.h](#).

14.15.2.4 `void LinkedList::Remove (LinkedListNode * node_)` `[pure virtual]`

Add the linked list node to this linked list.

Parameters

<code>node_</code>	Pointer to the node to remove
--------------------	-------------------------------

Implemented in [CircularLinkedList](#), [DoubleLinkedList](#), and [ThreadList](#).

The documentation for this class was generated from the following file:

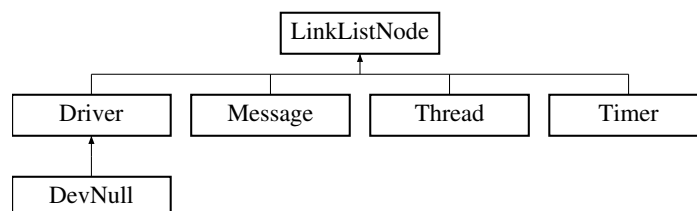
- [/home/moslevin/Project/R1/kernel/public/ll.h](#)

14.16 LinkedListNode Class Reference

Basic linked-list node data structure.

```
#include <ll.h>
```

Inheritance diagram for LinkedListNode:



Public Member Functions

- `LinkedListNode * GetNext (void)`
Returns a pointer to the next node in the list.
- `LinkedListNode * GetPrev (void)`
Returns a pointer to the previous node in the list.

Protected Member Functions

- `void ClearNode ()`
Initialize the linked list node, clearing its next and previous node.

Protected Attributes

- [LinkedListNode](#) * `next`
Pointer to the next node in the list.
- [LinkedListNode](#) * `prev`
Pointer to the previous node in the list.

Friends

- class **LinkedList**
- class **DoubleLinkedList**
- class **CircularLinkedList**

14.16.1 Detailed Description

Basic linked-list node data structure.

This data is managed by the linked-list class types, and can be used transparently between them.

Definition at line 68 of file [ll.h](#).

14.16.2 Member Function Documentation

14.16.2.1 [LinkedListNode](#) * [LinkedListNode::GetNext](#) (void) [\[inline\]](#)

Returns a pointer to the next node in the list.

Returns

a pointer to the next node in the list.

Definition at line 92 of file [ll.h](#).

14.16.2.2 [LinkedListNode](#) * [LinkedListNode::GetPrev](#) (void) [\[inline\]](#)

Returns a pointer to the previous node in the list.

Returns

a pointer to the previous node in the list.

Definition at line 101 of file [ll.h](#).

The documentation for this class was generated from the following files:

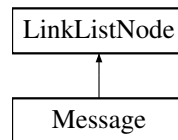
- [/home/moslevin/Project/R1/kernel/public/ll.h](#)
- [/home/moslevin/Project/R1/kernel/ll.cpp](#)

14.17 Message Class Reference

Class to provide message-based IPC services in the kernel.

```
#include <message.h>
```

Inheritance diagram for Message:



Public Member Functions

- void `Init` ()
Initialize the data and code in the message.
- void `SetData` (void *pvData_)
Set the data pointer for the message before transmission.
- void * `GetData` ()
Get the data pointer stored in the message upon receipt.
- void `SetCode` (K_USHORT usCode_)
Set the code in the message before transmission.
- K_USHORT `GetCode` ()
Return the code set in the message upon receipt.

Private Attributes

- void * `m_pvData`
Pointer to the message data.
- K_USHORT `m_usCode`
Message code, providing context for the message.

Additional Inherited Members

14.17.1 Detailed Description

Class to provide message-based IPC services in the kernel.

Definition at line 99 of file [message.h](#).

14.17.2 Member Function Documentation

14.17.2.1 K_USHORT Message::GetCode () [inline]

Return the code set in the message upon receipt.

Returns

User code set in the object

Definition at line 143 of file [message.h](#).

14.17.2.2 void * Message::GetData () [inline]

Get the data pointer stored in the message upon receipt.

Returns

Pointer to the data set in the message object

Definition at line 125 of file [message.h](#).

14.17.2.3 Message::SetCode (K_USHORT usCode_) [inline]

Set the code in the message before transmission.

Parameters

<code>usCode_</code>	Data code to set in the object
----------------------	--------------------------------

Definition at line 134 of file [message.h](#).

14.17.2.4 void Message::SetData (void * pvData_) [inline]

Set the data pointer for the message before transmission.

Parameters

<code>pvData_</code>	Pointer to the data object to send in the message
----------------------	---

Definition at line 116 of file [message.h](#).

The documentation for this class was generated from the following file:

- [/home/moslevin/Project/R1/kernel/public/message.h](#)

14.18 MessageQueue Class Reference

List of messages, used as the channel for sending and receiving messages between threads.

```
#include <message.h>
```

Public Member Functions

- `void Init ()`
Initialize the message queue prior to use.
- `Message * Receive ()`
Receive a message from the message queue.
- `Message * Receive (K_ULONG ulTimeWaitMS_)`
Receive a message from the message queue.
- `void Send (Message *pclSrc_)`
Send a message object into this message queue.
- `K_USHORT GetCount ()`
Return the number of messages pending in the "receive" queue.

Private Member Functions

- `Message * Receive_i (K_ULONG ulTimeWaitMS_)`
Receive_i.

Private Attributes

- `Semaphore m_clSemaphore`
Counting semaphore used to manage thread blocking.
- `DoubleLinkedList m_clLinkList`
List object used to store messages.

14.18.1 Detailed Description

List of messages, used as the channel for sending and receiving messages between threads.

Definition at line 201 of file [message.h](#).

14.18.2 Member Function Documentation

14.18.2.1 K_USHORT MessageQueue::GetCount ()

Return the number of messages pending in the "receive" queue.

Returns

Count of pending messages in the queue.

Definition at line 156 of file [message.cpp](#).

14.18.2.2 Message * MessageQueue::Receive ()

Receive a message from the message queue.

If the message queue is empty, the thread will block until a message is available.

Returns

Pointer to a message object at the head of the queue

Definition at line 92 of file [message.cpp](#).

14.18.2.3 Message * MessageQueue::Receive (K_ULONG ulWaitTimeMS_)

Receive a message from the message queue.

If the message queue is empty, the thread will block until a message is available for the duration specified. If no message arrives within that duration, the call will return with NULL.

Parameters

<i>ulWaitTimeMS_</i>	The amount of time in ms to wait for a message before timing out and unblocking the waiting thread.
----------------------	---

Returns

Pointer to a message object at the head of the queue or NULL on timeout.

Definition at line 103 of file [message.cpp](#).

14.18.2.4 Message * MessageQueue::Receive_i (K_ULONG ulTimeWaitMS_) [private]

Receive_i.

Internal function used to abstract timed and un-timed Receive calls.

Parameters

<code>ulWaitMS_</code>	Time (in ms) to block, 0 for un-timed call.
------------------------	---

Returns

Pointer to a message, or 0 on timeout.

Definition at line 111 of file [message.cpp](#).

14.18.2.5 void MessageQueue::Send (Message * pclSrc_)

Send a message object into this message queue.

Will un-block the first waiting thread blocked on this queue if that occurs.

Parameters

<code>pclSrc_</code>	Pointer to the message object to add to the queue
----------------------	---

Definition at line 140 of file [message.cpp](#).

The documentation for this class was generated from the following files:

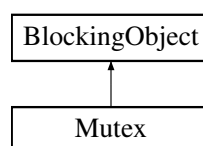
- [/home/moslevin/Project/R1/kernel/public/message.h](#)
- [/home/moslevin/Project/R1/kernel/message.cpp](#)

14.19 Mutex Class Reference

Mutual-exclusion locks, based on [BlockingObject](#).

```
#include <mutex.h>
```

Inheritance diagram for Mutex:



Public Member Functions

- void [Init](#) ()
Initialize a mutex object for use - must call this function before using the object.
- void [Claim](#) ()
Claim the mutex.
- [K_BOOL Claim](#) ([K_ULONG ulWaitTimeMS_](#))
- void [WakeMe](#) ([Thread *pclOwner_](#))
Wake a thread blocked on the mutex.
- void [Release](#) ()
Release the mutex.

Private Member Functions

- [K_UCHAR WakeNext \(\)](#)
Wake the next thread waiting on the [Mutex](#).
- [K_BOOL Claim_i \(K_ULONG ulWaitTimeMS_\)](#)
Claim_i.

Private Attributes

- [K_UCHAR m_ucRecurse](#)
The recursive lock-count when a mutex is claimed multiple times by the same owner.
- [K_UCHAR m_bReady](#)
State of the mutex - true = ready, false = claimed.
- [K_UCHAR m_ucMaxPri](#)
Maximum priority of thread in queue, used for priority inheritance.
- [Thread * m_pclOwner](#)
Pointer to the thread that owns the mutex (when claimed)

Additional Inherited Members

14.19.1 Detailed Description

Mutual-exclusion locks, based on [BlockingObject](#).

Definition at line 68 of file [mutex.h](#).

14.19.2 Member Function Documentation

14.19.2.1 void Mutex::Claim (void)

Claim the mutex.

When the mutex is claimed, no other thread can claim a region protected by the object.

Definition at line 209 of file [mutex.cpp](#).

14.19.2.2 K_BOOL Mutex::Claim (K_ULONG ulWaitTimeMS_)

Parameters

ulWaitTimeMS_	
-------------------------------	--

Returns

true - mutex was claimed within the time period specified
false - mutex operation timed-out before the claim operation.

Definition at line 220 of file [mutex.cpp](#).

14.19.2.3 K_BOOL Mutex::Claim_i (K_ULONG ulWaitTimeMS_) [private]

Claim_i.

Abstracts out timed/non-timed mutex claim operations.

Parameters

<code>uWaitTimeMS_</code>	Time in MS to wait, 0 for infinite
---------------------------	------------------------------------

Returns

true on successful claim, false otherwise

Definition at line 107 of file [mutex.cpp](#).

14.19.2.4 void Mutex::Release ()

Release the mutex.

When the mutex is released, another object can enter the mutex-protected region.

Definition at line 227 of file [mutex.cpp](#).

14.19.2.5 void Mutex::WakeMe (Thread * pOwner_)

Wake a thread blocked on the mutex.

This is an internal function used for implementing timed mutexes relying on timer callbacks. Since these do not have access to the private data of the mutex and its base classes, we have to wrap this as a public method - do not use this for any other purposes.

Parameters

<code>pOwner_</code>	Thread to unblock from this object.
----------------------	---

Definition at line 65 of file [mutex.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/mutex.h](#)
- [/home/moslevin/Project/R1/kernel/mutex.cpp](#)

14.20 Profiler Class Reference

System profiling timer interface.

```
#include <kernelprofile.h>
```

Static Public Member Functions

- static void [Init](#) ()
Initialize the global system profiler.
- static void [Start](#) ()
Start the global profiling timer service.
- static void [Stop](#) ()
Stop the global profiling timer service.
- static [K_USHORT Read](#) ()
Read the current tick count in the timer.
- static void [Process](#) ()
Process the profiling counters from ISR.
- static [K_ULONG GetEpoch](#) ()
Return the current timer epoch.

14.20.1 Detailed Description

System profiling timer interface.

Definition at line 37 of file [kernelprofile.h](#).

14.20.2 Member Function Documentation

14.20.2.1 void Profiler::Init (void) [static]

Initialize the global system profiler.

Must be called prior to use.

Definition at line 32 of file [kernelprofile.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/kernelprofile.h](#)
- [/home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/kernelprofile.cpp](#)

14.21 ProfileTimer Class Reference

Profiling timer.

```
#include <profile.h>
```

Public Member Functions

- void [Init](#) ()
Initialize the profiling timer prior to use.
- void [Start](#) ()
Start a profiling session, if the timer is not already active.
- void [Stop](#) ()
Stop the current profiling session, adding to the cumulative time for this timer, and the total iteration count.
- [K_ULONG](#) [GetAverage](#) ()
Get the average time associated with this operation.
- [K_ULONG](#) [GetCurrent](#) ()
Return the current tick count held by the profiler.

Private Member Functions

- [K_ULONG](#) [ComputeCurrentTicks](#) ([K_USHORT](#) usCount_, [K_ULONG](#) ulEpoch_)
Figure out how many ticks have elapsed in this iteration.

Private Attributes

- [K_ULONG](#) [m_ulCumulative](#)
Cumulative tick-count for this timer.
- [K_ULONG](#) [m_ulCurrentIteration](#)
Tick-count for the current iteration.
- [K_USHORT](#) [m_usInitial](#)
Initial count.

- [K_ULONG m_ulInitialEpoch](#)
Initial Epoch.
- [K_USHORT m_usIterations](#)
Number of iterations executed for this profiling timer.
- [K_UCHAR m_bActive](#)
Whether or not the timer is active or stopped.

14.21.1 Detailed Description

Profiling timer.

This class is used to perform high-performance profiling of code to see how K_LONG certain operations take. Useful in instrumenting the performance of key algorithms and time-critical operations to ensure real-time behavior.

Definition at line 69 of file [profile.h](#).

14.21.2 Member Function Documentation

14.21.2.1 **K_ULONG ProfileTimer::ComputeCurrentTicks** ([K_USHORT usCount_](#), [K_ULONG ulEpoch_](#))
[private]

Figure out how many ticks have elapsed in this iteration.

Parameters

usCount_	Current timer count
ulEpoch_	Current timer epoch

Returns

Current tick count

Definition at line 106 of file [profile.cpp](#).

14.21.2.2 **K_ULONG ProfileTimer::GetAverage** ()

Get the average time associated with this operation.

Returns

Average tick count normalized over all iterations

Definition at line 79 of file [profile.cpp](#).

14.21.2.3 **K_ULONG ProfileTimer::GetCurrent** ()

Return the current tick count held by the profiler.

Valid for both active and stopped timers.

Returns

The currently held tick count.

Definition at line 89 of file [profile.cpp](#).

14.21.2.4 void ProfileTimer::Init (void)

Initialize the profiling timer prior to use.

Can also be used to reset a timer that's been used previously.

Definition at line 37 of file [profile.cpp](#).

14.21.2.5 void ProfileTimer::Start (void)

Start a profiling session, if the timer is not already active.

Has no effect if the timer is already active.

Definition at line 46 of file [profile.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/profile.h](#)
- [/home/moslevin/Project/R1/kernel/profile.cpp](#)

14.22 Quantum Class Reference

Static-class used to implement [Thread](#) quantum functionality, which is a key part of round-robin scheduling.

```
#include <quantum.h>
```

Static Public Member Functions

- static void [UpdateTimer](#) ()
This function is called to update the thread quantum timer whenever something in the scheduler has changed.
- static void [AddThread](#) ([Thread](#) *pclThread_)
Add the thread to the quantum timer.
- static void [RemoveThread](#) ()
Remove the thread from the quantum timer.
- static void [SetInTimer](#) (void)
SetInTimer.
- static void [ClearInTimer](#) (void)
ClearInTimer.

Static Private Member Functions

- static void [SetTimer](#) ([Thread](#) *pclThread_)
Set up the quantum timer in the timer scheduler.

14.22.1 Detailed Description

Static-class used to implement [Thread](#) quantum functionality, which is a key part of round-robin scheduling.

Definition at line 41 of file [quantum.h](#).

14.22.2 Member Function Documentation

14.22.2.1 void Quantum::AddThread (Thread * *pciThread_*) [static]

Add the thread to the quantum timer.

Only one thread can own the quantum, since only one thread can be running on a core at a time.

Definition at line 82 of file [quantum.cpp](#).

14.22.2.2 static void Quantum::ClearInTimer (void) [inline],[static]

ClearInTimer.

Clear the flag once the timer callback function has been completed.

Definition at line 84 of file [quantum.h](#).

14.22.2.3 void Quantum::RemoveThread (void) [static]

Remove the thread from the quantum timer.

This will cancel the timer.

Definition at line 111 of file [quantum.cpp](#).

14.22.2.4 static void Quantum::SetInTimer (void) [inline],[static]

SetInTimer.

Set a flag to indicate that the CPU is currently running within the timer-callback routine. This prevents the [Quantum](#) timer from being updated in the middle of a callback cycle, potentially resulting in the kernel timer becoming disabled.

Definition at line 77 of file [quantum.h](#).

14.22.2.5 void Quantum::SetTimer (Thread * *pciThread_*) [static],[private]

Set up the quantum timer in the timer scheduler.

This creates a one-shot timer, which calls a static callback in [quantum.cpp](#) that on expiry will pivot the head of the threadlist for the thread's priority. This is the mechanism that provides round-robin scheduling in the system.

Parameters

<i>pciThread_</i>	Pointer to the thread to set the Quantum timer on
-------------------	---

Definition at line 72 of file [quantum.cpp](#).

14.22.2.6 void Quantum::UpdateTimer (void) [static]

This function is called to update the thread quantum timer whenever something in the scheduler has changed.

This can result in the timer being re-loaded or started. The timer is never stopped, but it may be ignored on expiry.

Definition at line 124 of file [quantum.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/quantum.h](#)
- [/home/moslevin/Project/R1/kernel/quantum.cpp](#)

14.23 Scheduler Class Reference

Priority-based round-robin [Thread](#) scheduling, using ThreadLists for housekeeping.

```
#include <scheduler.h>
```

Static Public Member Functions

- static void [Init](#) ()
Intialize the scheduler, must be called before use.
- static void [Schedule](#) ()
Run the scheduler, determines the next thread to run based on the current state of the threads.
- static void [Add](#) ([Thread](#) *pclThread_)
Add a thread to the scheduler at its current priority level.
- static void [Remove](#) ([Thread](#) *pclThread_)
Remove a thread from the scheduler at its current priority level.
- static [K_BOOL](#) [SetScheduler](#) ([K_BOOL](#) bEnable_)
Set the active state of the scheduler.
- static [Thread](#) * [GetCurrentThread](#) ()
Return the pointer to the currently-running thread.
- static volatile [Thread](#) * [GetNextThread](#) ()
Return the pointer to the thread that should run next, according to the last run of the scheduler.
- static [ThreadList](#) * [GetThreadList](#) ([K_UCHAR](#) ucPriority_)
Return the pointer to the active list of threads that are at the given priority level in the scheduler.
- static [ThreadList](#) * [GetStopList](#) ()
Return the pointer to the list of threads that are in the scheduler's stopped state.
- static [K_UCHAR](#) [IsEnabled](#) ()
Return the current state of the scheduler - whether or not scheudling is enabled or disabled.
- static void [QueueScheduler](#) ()
QueueScheduler.

Static Private Attributes

- static [K_BOOL](#) [m_bEnabled](#)
Scheduler's state - enabled or disabled.
- static [K_BOOL](#) [m_bQueuedSchedule](#)
Variable representing whether or not there's a queued scheduler operation.
- static [ThreadList](#) [m_clStopList](#)
ThreadList for all stopped threads.
- static [ThreadList](#) [m_aclPriorities](#) [[NUM_PRIORITIES](#)]
ThreadLists for all threads at all priorities.
- static [K_UCHAR](#) [m_ucPriFlag](#)
Bitmap flag for each.

14.23.1 Detailed Description

Priority-based round-robin [Thread](#) scheduling, using ThreadLists for housekeeping.

Definition at line 62 of file [scheduler.h](#).

14.23.2 Member Function Documentation

14.23.2.1 `void Scheduler::Add (Thread * pclThread_) [static]`

Add a thread to the scheduler at its current priority level.

Parameters

<i>pcThread_</i>	Pointer to the thread to add to the scheduler
------------------	---

Definition at line 108 of file [scheduler.cpp](#).

14.23.2.2 static Thread* Scheduler::GetCurrentThread () [inline],[static]

Return the pointer to the currently-running thread.

Returns

Pointer to the currently-running thread

Definition at line 119 of file [scheduler.h](#).

14.23.2.3 static volatile Thread* Scheduler::GetNextThread () [inline],[static]

Return the pointer to the thread that should run next, according to the last run of the scheduler.

Returns

Pointer to the next-running thread

Definition at line 127 of file [scheduler.h](#).

14.23.2.4 static ThreadList* Scheduler::GetStopList () [inline],[static]

Return the pointer to the list of threads that are in the scheduler's stopped state.

Returns

Pointer to the [ThreadList](#) containing the stopped threads

Definition at line 145 of file [scheduler.h](#).

14.23.2.5 static ThreadList* Scheduler::GetThreadList (K_UCHAR ucPriority_) [inline],[static]

Return the pointer to the active list of threads that are at the given priority level in the scheduler.

Parameters

<i>ucPriority_</i>	Priority level of
--------------------	-------------------

Returns

Pointer to the [ThreadList](#) for the given priority level

Definition at line 137 of file [scheduler.h](#).

14.23.2.6 K_UCHAR Scheduler::IsEnabled () [inline],[static]

Return the current state of the scheduler - whether or not scheduling is enabled or disabled.

Returns

true - scheduler enabled, false - disabled

Definition at line 155 of file [scheduler.h](#).

14.23.2.7 `static void Scheduler::QueueScheduler () [inline],[static]`

QueueScheduler.

Tell the kernel to perform a scheduling operation as soon as the scheduler is re-enabled.

Definition at line 163 of file [scheduler.h](#).

14.23.2.8 `void Scheduler::Remove (Thread * pcIThread_) [static]`

Remove a thread from the scheduler at its current priority level.

Parameters

<i>pcIThread_</i>	Pointer to the thread to be removed from the scheduler
-------------------	--

Definition at line 114 of file [scheduler.cpp](#).

14.23.2.9 `Scheduler::Schedule () [static]`

Run the scheduler, determines the next thread to run based on the current state of the threads.

Note that the next-thread chosen from this function is only valid while in a critical section.

Definition at line 71 of file [scheduler.cpp](#).

14.23.2.10 `void Scheduler::SetScheduler (K_BOOL bEnable_) [static]`

Set the active state of the scheduler.

When the scheduler is disabled, the *next thread* is never set; the currently running thread will run forever until the scheduler is enabled again. Care must be taken to ensure that we don't end up trying to block while the scheduler is disabled, otherwise the system ends up in an unusable state.

Parameters

<i>bEnable_</i>	true to enable, false to disable the scheduler
-----------------	--

Definition at line 120 of file [scheduler.cpp](#).

The documentation for this class was generated from the following files:

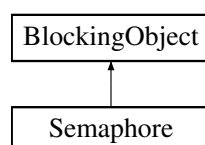
- [/home/moslevin/Project/R1/kernel/public/scheduler.h](#)
- [/home/moslevin/Project/R1/kernel/scheduler.cpp](#)

14.24 Semaphore Class Reference

Counting semaphore, based on [BlockingObject](#) base class.

```
#include <ksemaphore.h>
```

Inheritance diagram for Semaphore:



Public Member Functions

- void [Init](#) ([K_USHORT](#) usInitVal_, [K_USHORT](#) usMaxVal_)
Initialize a semaphore before use.
- [K_BOOL](#) [Post](#) ()
Increment the semaphore count.
- void [Pend](#) ()
Decrement the semaphore count.
- [K_USHORT](#) [GetCount](#) ()
Return the current semaphore counter.
- [K_BOOL](#) [Pend](#) ([K_ULONG](#) ulWaitTimeMS_)
Decrement the semaphore count.
- void [WakeMe](#) ([Thread](#) *pclChosenOne_)
Wake a thread blocked on the semaphore.

Private Member Functions

- [K_UCHAR](#) [WakeNext](#) ()
Wake the next thread waiting on the semaphore.
- [K_BOOL](#) [Pend_i](#) ([K_ULONG](#) ulWaitTimeMS_)
Pend_i.

Private Attributes

- [K_USHORT](#) [m_usValue](#)
Current count held by the semaphore.
- [K_USHORT](#) [m_usMaxValue](#)
Maximum count that can be held by this semaphore.

Additional Inherited Members

14.24.1 Detailed Description

Counting semaphore, based on [BlockingObject](#) base class.

Definition at line 37 of file [ksemaphore.h](#).

14.24.2 Member Function Documentation

14.24.2.1 [K_USHORT](#) Semaphore::GetCount ()

Return the current semaphore counter.

This can be used by a thread to bypass blocking on a semaphore - allowing it to do other things until a non-zero count is returned, instead of blocking until the semaphore is posted.

Returns

The current semaphore counter value.

Definition at line 234 of file [ksemaphore.cpp](#).

14.24.2.2 void Semaphore::Init (K_USHORT *usInitVal_*, K_USHORT *usMaxVal_*)

Initialize a semaphore before use.

Must be called before post/pend operations.

Parameters

<i>usInitVal_</i>	Initial value held by the semaphore
<i>usMaxVal_</i>	Maximum value for the semaphore

Definition at line 95 of file [ksemaphore.cpp](#).

14.24.2.3 void Semaphore::Pend ()

Decrement the semaphore count.

If the count is zero, the thread will block until the semaphore is pended.

Definition at line 216 of file [ksemaphore.cpp](#).

14.24.2.4 K_BOOL Semaphore::Pend (K_ULONG ulWaitTimeMS_)

Decrement the semaphore count.

If the count is zero, the thread will block until the semaphore is pended. If the specified interval expires before the thread is unblocked, then the status is returned back to the user.

Returns

true - semaphore was acquired before the timeout false - timeout occurred before the semaphore was claimed.

Definition at line 227 of file [ksemaphore.cpp](#).

14.24.2.5 K_BOOL Semaphore::Pend_i (K_ULONG ulWaitTimeMS_) [private]

Pend_i.

Internal function used to abstract timed and untimed semaphore pend operations.

Parameters

<i>ulWaitTimeMS_</i>	Time in MS to wait
----------------------	--------------------

Returns

true on success, false on failure.

Definition at line 160 of file [ksemaphore.cpp](#).

14.24.2.6 void Semaphore::Post ()

Increment the semaphore count.

Returns

true if the semaphore was posted, false if the count is already maxed out.

Definition at line 107 of file [ksemaphore.cpp](#).

14.24.2.7 void Semaphore::WakeMe (Thread * pChosenOne_)

Wake a thread blocked on the semaphore.

This is an internal function used for implementing timed semaphores relying on timer callbacks. Since these do not have access to the private data of the semaphore and its base classes, we have to wrap this as a public method - do not use this for any other purposes.

Definition at line 68 of file [ksemaphore.cpp](#).

The documentation for this class was generated from the following files:

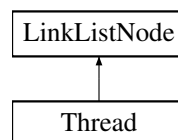
- [/home/moslevin/Project/R1/kernel/public/ksemaphore.h](#)
- [/home/moslevin/Project/R1/kernel/ksemaphore.cpp](#)

14.25 Thread Class Reference

Object providing fundamental multitasking support in the kernel.

```
#include <thread.h>
```

Inheritance diagram for Thread:



Public Member Functions

- void [Init](#) ([K_WORD](#) *paucStack_, [K_USHORT](#) usStackSize_, [K_UCHAR](#) ucPriority_, [ThreadEntry_t](#) pfEntryPoint_, void *pvArg_)
Initialize a thread prior to its use.
- void [Start](#) ()
Start the thread - remove it from the stopped list, add it to the scheduler's list of threads (at the thread's set priority), and continue along.
- void [Stop](#) ()
Stop a thread that's actively scheduled without destroying its stacks.
- [ThreadList](#) * [GetOwner](#) (void)
Return the [ThreadList](#) where the thread belongs when it's in the active/ready state in the scheduler.
- [ThreadList](#) * [GetCurrent](#) (void)
Return the [ThreadList](#) where the thread is currently located.
- [K_UCHAR](#) [GetPriority](#) (void)
Return the priority of the current thread.
- [K_UCHAR](#) [GetCurPriority](#) (void)
Return the priority of the current thread.
- void [SetQuantum](#) ([K_USHORT](#) usQuantum_)
Set the thread's round-robin execution quantum.
- [K_USHORT](#) [GetQuantum](#) (void)
Get the thread's round-robin execution quantum.
- void [SetCurrent](#) ([ThreadList](#) *pclNewList_)
Set the thread's current to the specified thread list.
- void [SetOwner](#) ([ThreadList](#) *pclNewList_)
Set the thread's owner to the specified thread list.
- void [SetPriority](#) ([K_UCHAR](#) ucPriority_)
Set the priority of the [Thread](#) (running or otherwise) to a different level.

- void [InheritPriority](#) ([K_UCHAR](#) ucPriority_)
Allow the thread to run at a different priority level (temporarily) for the purpose of avoiding priority inversions.
- void [Exit](#) ()
Remove the thread from being scheduled again.
- void [SetID](#) ([K_UCHAR](#) ucID_)
Set an 8-bit ID to uniquely identify this thread.
- [K_UCHAR](#) [GetID](#) ()
Return the 8-bit ID corresponding to this thread.
- [K_USHORT](#) [GetStackSlack](#) ()
Performs a (somewhat lengthy) check on the thread stack to check the amount of stack margin (or "slack") remaining on the stack.
- [K_USHORT](#) [GetEventFlagMask](#) ()
GetEventFlagMask returns the thread's current event-flag mask, which is used in conjunction with the [EventFlag](#) blocking object type.
- void [SetEventFlagMask](#) ([K_USHORT](#) usMask_)
SetEventFlagMask Sets the active event flag bitfield mask.
- void [SetEventFlagMode](#) ([EventFlagOperation_t](#) eMode_)
SetEventFlagMode Sets the active event flag operation mode.
- [EventFlagOperation_t](#) [GetEventFlagMode](#) ()
GetEventFlagMode Returns the thread's event flag's operating mode.
- [Timer *](#) [GetTimer](#) ()
Return a pointer to the thread's timer object.
- void [SetExpired](#) ([K_BOOL](#) bExpired_)
SetExpired.
- [K_BOOL](#) [GetExpired](#) ()
GetExpired.
- void [InitIdle](#) ()
InitIdle Initialize this [Thread](#) object as the [Kernel](#)'s idle thread.
- [ThreadState_t](#) [GetState](#) ()
GetState Returns the current state of the thread to the caller.
- void [SetState](#) ([ThreadState_t](#) eState_)
SetState Set the thread's state to a new value.

Static Public Member Functions

- static void [Sleep](#) ([K_ULONG](#) ulTimeMs_)
Put the thread to sleep for the specified time (in milliseconds).
- static void [USleep](#) ([K_ULONG](#) ulTimeUs_)
Put the thread to sleep for the specified time (in microseconds).
- static void [Yield](#) (void)
Yield the thread - this forces the system to call the scheduler and determine what thread should run next.

Private Member Functions

- void [SetPriorityBase](#) ([K_UCHAR](#) ucPriority_)

Static Private Member Functions

- static void [ContextSwitchSWI](#) (void)
This code is used to trigger the context switch interrupt.

Private Attributes

- [K_WORD * m_pwStackTop](#)
Pointer to the top of the thread's stack.
- [K_WORD * m_pwStack](#)
Pointer to the thread's stack.
- [K_UCHAR m_ucThreadID](#)
Thread ID.
- [K_UCHAR m_ucPriority](#)
Default priority of the thread.
- [K_UCHAR m_ucCurPriority](#)
Current priority of the thread (priority inheritance)
- [ThreadState_t m_eState](#)
Enum indicating the thread's current state.
- [K_USHORT m_usStackSize](#)
Size of the stack (in bytes)
- [ThreadList * m_pclCurrent](#)
Pointer to the thread-list where the thread currently resides.
- [ThreadList * m_pclOwner](#)
Pointer to the thread-list where the thread resides when active.
- [ThreadEntry_t m_pfEntryPoint](#)
The entry-point function called when the thread starts.
- [void * m_pvArg](#)
Pointer to the argument passed into the thread's entrypoint.
- [K_USHORT m_usQuantum](#)
Thread quantum (in milliseconds)
- [K_USHORT m_usFlagMask](#)
Event-flag mask.
- [EventFlagOperation_t m_eFlagMode](#)
Event-flag mode.
- [Timer m_clTimer](#)
Timer used for blocking-object timeouts.
- [K_BOOL m_bExpired](#)
Indicate whether or not a blocking-object timeout has occurred.

Friends

- class **ThreadPort**

Additional Inherited Members

14.25.1 Detailed Description

Object providing fundamental multitasking support in the kernel.

Definition at line 71 of file [thread.h](#).

14.25.2 Member Function Documentation

14.25.2.1 void Thread::ContextSwitchSWI(void) [static], [private]

This code is used to trigger the context switch interrupt.

Called whenever the kernel decides that it is necessary to swap out the current thread for the "next" thread.

Definition at line 395 of file [thread.cpp](#).

14.25.2.2 void Thread::Exit()

Remove the thread from being scheduled again.

The thread is effectively destroyed when this occurs. This is extremely useful for cases where a thread encounters an unrecoverable error and needs to be restarted, or in the context of systems where threads need to be created and destroyed dynamically.

This must not be called on the idle thread.

Definition at line 174 of file [thread.cpp](#).

14.25.2.3 K_UCHAR Thread::GetCurPriority(void) [inline]

Return the priority of the current thread.

Returns

Priority of the current thread

Definition at line 174 of file [thread.h](#).

14.25.2.4 ThreadList * Thread::GetCurrent(void) [inline]

Return the [ThreadList](#) where the thread is currently located.

Returns

Pointer to the thread's current list

Definition at line 155 of file [thread.h](#).

14.25.2.5 K_USHORT Thread::GetEventFlagMask() [inline]

GetEventFlagMask returns the thread's current event-flag mask, which is used in conjunction with the [EventFlag](#) blocking object type.

Returns

A copy of the thread's event flag mask

Definition at line 327 of file [thread.h](#).

14.25.2.6 EventFlagOperation_t Thread::GetEventFlagMode() [inline]

GetEventFlagMode Returns the thread's event flag's operating mode.

Returns

The thread's event flag mode.

Definition at line 346 of file [thread.h](#).

14.25.2.7 K_BOOL Thread::GetExpired ()

GetExpired.

Return the status of the most-recent blocking call on the thread.

Returns

true - call expired, false - call did not expire

Definition at line 413 of file [thread.cpp](#).

14.25.2.8 K_UCHAR Thread::GetID () [inline]

Return the 8-bit ID corresponding to this thread.

Returns

[Thread](#)'s 8-bit ID, set by the user

Definition at line 302 of file [thread.h](#).

14.25.2.9 ThreadList * Thread::GetOwner (void) [inline]

Return the [ThreadList](#) where the thread belongs when it's in the active/ready state in the scheduler.

Returns

Pointer to the [Thread](#)'s owner list

Definition at line 146 of file [thread.h](#).

14.25.2.10 K_UCHAR Thread::GetPriority (void) [inline]

Return the priority of the current thread.

Returns

Priority of the current thread

Definition at line 165 of file [thread.h](#).

14.25.2.11 K_USHORT Thread::GetQuantum (void) [inline]

Get the thread's round-robin execution quantum.

Returns

The thread's quantum

Definition at line 193 of file [thread.h](#).

14.25.2.12 `K_USHORT Thread::GetStackSlack ()`

Performs a (somewhat lengthy) check on the thread stack to check the amount of stack margin (or "slack") remaining on the stack.

If you're having problems with blowing your stack, you can run this function at points in your code during development to see what operations cause problems. Also useful during development as a tool to optimally size thread stacks.

Returns

The amount of slack (unused bytes) on the stack

! ToDo: Take into account stacks that grow up

Definition at line 284 of file [thread.cpp](#).

14.25.2.13 `ThreadState_t Thread::GetState ()` `[inline]`

GetState Returns the current state of the thread to the caller.

Can be used to determine whether or not a thread is ready (or running), stopped, or terminated/exit'd.

Returns

ThreadState_t representing the thread's current state

Definition at line 390 of file [thread.h](#).

14.25.2.14 `void Thread::InheritPriority (K_UCHAR ucPriority_)`

Allow the thread to run at a different priority level (temporarily) for the purpose of avoiding priority inversions.

This should only be called from within the implementation of blocking-objects.

Parameters

<i>ucPriority_</i>	New Priority to boost to.
--------------------	---------------------------

Definition at line 388 of file [thread.cpp](#).

14.25.2.15 `void Thread::Init (K_WORD * paucStack_, K_USHORT usStackSize_, K_UCHAR ucPriority_, ThreadEntry_t pfEntryPoint_, void * pvArg_)`

Initialize a thread prior to its use.

Initialized threads are placed in the stopped state, and are not scheduled until the thread's start method has been invoked first.

Parameters

<i>paucStack_</i>	Pointer to the stack to use for the thread
<i>usStackSize_</i>	Size of the stack (in bytes)
<i>ucPriority_</i>	Priority of the thread (0 = idle, 7 = max)
<i>pfEntryPoint_</i>	This is the function that gets called when the thread is started
<i>pvArg_</i>	Pointer to the argument passed into the thread's entrypt function.

Definition at line 41 of file [thread.cpp](#).

14.25.2.16 `void Thread::InitIdle (void)`

InitIdle Initialize this [Thread](#) object as the [Kernel](#)'s idle thread.

There should only be one of these, maximum, in a given system.

Definition at line 418 of file [thread.cpp](#).

14.25.2.17 void Thread::SetCurrent (ThreadList * *pcNewList_*) [inline]

Set the thread's current to the specified thread list.

Parameters

<i>pcNewList_</i>	Pointer to the threadlist to apply thread ownership
-------------------	---

Definition at line 203 of file [thread.h](#).

14.25.2.18 void Thread::SetEventFlagMask (K_USHORT *usMask_*) [inline]

SetEventFlagMask Sets the active event flag bitfield mask.

Parameters

<i>usMask_</i>	
----------------	--

Definition at line 333 of file [thread.h](#).

14.25.2.19 void Thread::SetEventFlagMode (EventFlagOperation_t *eMode_*) [inline]

SetEventFlagMode Sets the active event flag operation mode.

Parameters

<i>eMode_</i>	Event flag operation mode, defines the logical operator to apply to the event flag.
---------------	---

Definition at line 340 of file [thread.h](#).

14.25.2.20 void Thread::SetExpired (K_BOOL *bExpired_*)

SetExpired.

Set the status of the current blocking call on the thread.

Parameters

<i>bExpired_</i>	true - call expired, false - call did not expire
------------------	--

Definition at line 410 of file [thread.cpp](#).

14.25.2.21 void Thread::SetID (K_UCHAR *ucID_*) [inline]

Set an 8-bit ID to uniquely identify this thread.

Parameters

<i>ucID_</i>	8-bit Thread ID, set by the user
--------------	--

Definition at line 293 of file [thread.h](#).

14.25.2.22 void Thread::SetOwner (ThreadList * *pcNewList_*) [inline]

Set the thread's owner to the specified thread list.

Parameters

<i>pc/NewList_</i>	Pointer to the threadlist to apply thread ownership
--------------------	---

Definition at line 212 of file [thread.h](#).

14.25.2.23 void Thread::SetPriority (K_UCHAR ucPriority_)

Set the priority of the [Thread](#) (running or otherwise) to a different level.

This activity involves re-scheduling, and must be done so with due caution, as it may effect the determinism of the system.

This should *always* be called from within a critical section to prevent system issues.

Parameters

<i>ucPriority_</i>	New priority of the thread
--------------------	----------------------------

Definition at line 344 of file [thread.cpp](#).

14.25.2.24 void Thread::SetPriorityBase (K_UCHAR ucPriority_) [private]

Parameters

<i>ucPriority_</i>	
--------------------	--

Definition at line 334 of file [thread.cpp](#).

14.25.2.25 void Thread::SetQuantum (K_USHORT usQuantum_) [inline]

Set the thread's round-robin execution quantum.

Parameters

<i>usQuantum_</i>	Thread 's execution quantum (in milliseconds)
-------------------	---

Definition at line 184 of file [thread.h](#).

14.25.2.26 void Thread::SetState (ThreadState_t eState_) [inline]

SetState Set the thread's state to a new value.

This is only to be used by code within the kernel, and is not intended for use by an end-user.

Parameters

<i>eState_</i>	New thread state to set.
----------------	--------------------------

Definition at line 399 of file [thread.h](#).

14.25.2.27 void Thread::Sleep (K_ULONG ulTimeMs_) [static]

Put the thread to sleep for the specified time (in milliseconds).

Actual time slept may be longer (but not less than) the interval specified.

Parameters

<code>ulTimeMs_</code>	Time to sleep (in ms)
------------------------	-----------------------

Definition at line 239 of file [thread.cpp](#).

14.25.2.28 void Thread::Stop ()

Stop a thread that's actively scheduled without destroying its stacks.

Stopped threads can be restarted using the [Start\(\)](#) API.

Definition at line 129 of file [thread.cpp](#).

14.25.2.29 void Thread::USleep (K_ULONG ulTimeUs_) [static]

Put the thread to sleep for the specified time (in microseconds).

Actual time slept may be longer (but not less than) the interval specified.

Parameters

<code>ulTimeUs_</code>	Time to sleep (in microseconds)
------------------------	---------------------------------

Definition at line 261 of file [thread.cpp](#).

14.25.2.30 void Thread::Yield (void) [static]

Yield the thread - this forces the system to call the scheduler and determine what thread should run next.

This is typically used when threads are moved in and out of the scheduler.

Definition at line 305 of file [thread.cpp](#).

The documentation for this class was generated from the following files:

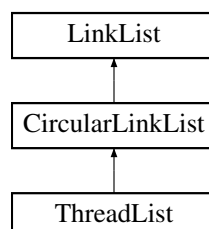
- [/home/moslevin/Project/R1/kernel/public/thread.h](#)
- [/home/moslevin/Project/R1/kernel/thread.cpp](#)

14.26 ThreadList Class Reference

This class is used for building thread-management facilities, such as schedulers, and blocking objects.

```
#include <threadlist.h>
```

Inheritance diagram for ThreadList:



Public Member Functions

- [ThreadList \(\)](#)

Default constructor - zero-initializes the data.

- void [SetPriority](#) ([K_UCHAR](#) ucPriority_)
Set the priority of this threadlist (if used for a scheduler).
- void [SetFlagPointer](#) ([K_UCHAR](#) *pucFlag_)
Set the pointer to a bitmap to use for this threadlist.
- void [Add](#) ([LinkListNode](#) *node_)
Add a thread to the threadlist.
- void [Add](#) ([LinkListNode](#) *node_, [K_UCHAR](#) *pucFlag_, [K_UCHAR](#) ucPriority_)
Add a thread to the threadlist, specifying the flag and priority at the same time.
- void [Remove](#) ([LinkListNode](#) *node_)
Remove the specified thread from the threadlist.
- [Thread](#) * [HighestWaiter](#) ()
Return a pointer to the highest-priority thread in the thread-list.

Private Attributes

- [K_UCHAR](#) m_ucPriority
Priority of the threadlist.
- [K_UCHAR](#) * m_pucFlag
Pointer to the bitmap/flag to set when used for scheduling.

Additional Inherited Members

14.26.1 Detailed Description

This class is used for building thread-management facilities, such as schedulers, and blocking objects.

Definition at line 34 of file [threadlist.h](#).

14.26.2 Member Function Documentation

14.26.2.1 void ThreadList::Add ([LinkListNode](#) * node_) [virtual]

Add a thread to the threadlist.

Parameters

<i>node_</i>	Pointer to the thread (link list node) to add to the list
--------------	---

Reimplemented from [CircularLinkList](#).

Definition at line 46 of file [threadlist.cpp](#).

14.26.2.2 void ThreadList::Add ([LinkListNode](#) * node_, [K_UCHAR](#) * pucFlag_, [K_UCHAR](#) ucPriority_)

Add a thread to the threadlist, specifying the flag and priority at the same time.

Parameters

<i>node_</i>	Pointer to the thread to add (link list node)
<i>pucFlag_</i>	Pointer to the bitmap flag to set (if used in a scheduler context), or NULL for non-scheduler.
<i>ucPriority_</i>	Priority of the threadlist

Definition at line 59 of file [threadlist.cpp](#).

14.26.2.3 Thread * ThreadList::HighestWaiter ()

Return a pointer to the highest-priority thread in the thread-list.

Returns

Pointer to the highest-priority thread

Definition at line 84 of file [threadlist.cpp](#).

14.26.2.4 void ThreadList::Remove (LinkListNode * node_) [virtual]

Remove the specified thread from the threadlist.

Parameters

<i>node_</i>	Pointer to the thread to remove
--------------	---------------------------------

Reimplemented from [CircularLinkList](#).

Definition at line 68 of file [threadlist.cpp](#).

14.26.2.5 void ThreadList::SetFlagPointer (K_UCHAR * pucFlag_)

Set the pointer to a bitmap to use for this threadlist.

Once again, only needed when the threadlist is being used for scheduling purposes.

Parameters

<i>pucFlag_</i>	Pointer to the bitmap flag
-----------------	----------------------------

Definition at line 40 of file [threadlist.cpp](#).

14.26.2.6 void ThreadList::SetPriority (K_UCHAR ucPriority_)

Set the priority of this threadlist (if used for a scheduler).

Parameters

<i>ucPriority_</i>	Priority level of the thread list
--------------------	-----------------------------------

Definition at line 34 of file [threadlist.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/threadlist.h](#)
- [/home/moslevin/Project/R1/kernel/threadlist.cpp](#)

14.27 ThreadPort Class Reference

Class defining the architecture specific functions required by the kernel.

```
#include <threadport.h>
```

Static Public Member Functions

- static void [StartThreads](#) ()
Function to start the scheduler, initial threads, etc.

Static Private Member Functions

- static void [InitStack](#) ([Thread](#) *pstThread_)
Initialize the thread's stack.

Friends

- class [Thread](#)

14.27.1 Detailed Description

Class defining the architecture specific functions required by the kernel.

This is limited (at this point) to a function to start the scheduler, and a function to initialize the default stack-frame for a thread.

Definition at line 167 of file [threadport.h](#).

14.27.2 Member Function Documentation

14.27.2.1 void [ThreadPort::InitStack](#) ([Thread](#) * *pstThread_*) [static], [private]

Initialize the thread's stack.

Parameters

<i>pstThread_</i>	Pointer to the thread to initialize
-------------------	-------------------------------------

Definition at line 39 of file [threadport.cpp](#).

The documentation for this class was generated from the following files:

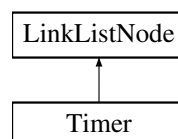
- /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/[threadport.h](#)
- /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/[threadport.cpp](#)

14.28 Timer Class Reference

[Timer](#) - an event-driven execution context based on a specified time interval.

```
#include <timer.h>
```

Inheritance diagram for [Timer](#):



Public Member Functions

- [Timer](#) ()
Default Constructor - zero-initializes all internal data.
- void [Init](#) ()
Re-initialize the [Timer](#) to default values.

- void [Start](#) ([K_BOOL](#) bRepeat_, [K_ULONG](#) ulIntervalMs_, [TimerCallback_t](#) pfCallback_, void *pvData_)
Start a timer using default ownership, using repeats as an option, and millisecond resolution.
- void [Start](#) ([K_BOOL](#) bRepeat_, [K_ULONG](#) ulIntervalMs_, [K_ULONG](#) ulToleranceMs_, [TimerCallback_t](#) pfCallback_, void *pvData_)
Start a timer using default ownership, using repeats as an option, and millisecond resolution.
- void [Stop](#) ()
Stop a timer already in progress.
- void [SetFlags](#) ([K_UCHAR](#) ucFlags_)
Set the timer's flags based on the bits in the ucFlags_ argument.
- void [SetCallback](#) ([TimerCallback_t](#) pfCallback_)
Define the callback function to be executed on expiry of the timer.
- void [SetData](#) (void *pvData_)
Define a pointer to be sent to the timer callback on timer expiry.
- void [SetOwner](#) ([Thread](#) *pOwner_)
Set the owner-thread of this timer object (all timers must be owned by a thread).
- void [SetIntervalTicks](#) ([K_ULONG](#) ulTicks_)
Set the timer expiry in system-ticks (platform specific!)
- void [SetIntervalSeconds](#) ([K_ULONG](#) ulSeconds_)
! The next three cost us 330 bytes of flash on AVR...
- void [SetIntervalMSeconds](#) ([K_ULONG](#) ulMSeconds_)
Set the timer expiry interval in milliseconds (platform agnostic)
- void [SetIntervalUSeconds](#) ([K_ULONG](#) ulUSeconds_)
Set the timer expiry interval in microseconds (platform agnostic)
- void [SetTolerance](#) ([K_ULONG](#) ulTicks_)
Set the timer's maximum tolerance in order to synchronize timer processing with other timers in the system.

Private Attributes

- [K_UCHAR](#) m_ucFlags
Flags for the timer, defining if the timer is one-shot or repeated.
- [TimerCallback_t](#) m_pfCallback
Pointer to the callback function.
- [K_ULONG](#) m_ulInterval
Interval of the timer in timer ticks.
- [K_ULONG](#) m_ulTimeLeft
Time remaining on the timer.
- [K_ULONG](#) m_ulTimerTolerance
Maximum tolerance (used for timer harmonization)
- [Thread](#) * m_pOwner
Pointer to the owner thread.
- void * m_pvData
Pointer to the callback data.

Friends

- class [TimerList](#)

Additional Inherited Members

14.28.1 Detailed Description

Timer - an event-driven execution context based on a specified time interval.

This inherits from a [LinkedListNode](#) for ease of management by a global [TimerList](#) object.

Definition at line 105 of file [timer.h](#).

14.28.2 Member Function Documentation

14.28.2.1 void Timer::SetCallback (TimerCallback_t pfCallback_) [inline]

Define the callback function to be executed on expiry of the timer.

Parameters

<i>pfCallback_</i>	Pointer to the callback function to call
--------------------	--

Definition at line 166 of file [timer.h](#).

14.28.2.2 void Timer::SetData (void * pvData_) [inline]

Define a pointer to be sent to the timer callbacak on timer expiry.

Parameters

<i>pvData_</i>	Pointer to data to pass as argument into the callback
----------------	---

Definition at line 175 of file [timer.h](#).

14.28.2.3 void Timer::SetFlags (K_UCHAR ucFlags_) [inline]

Set the timer's flags based on the bits in the ucFlags_ argument.

Parameters

<i>ucFlags_</i>	Flags to assign to the timer object. TIMERLIST_FLAG_ONE_SHOT for a one-shot timer, 0 for a continuous timer.
-----------------	--

Definition at line 157 of file [timer.h](#).

14.28.2.4 void Timer::SetIntervalMSeconds (K_ULONG ulMSeconds_)

Set the timer expiry interval in milliseconds (platform agnostic)

Parameters

<i>ulMSeconds_</i>	Time in milliseconds
--------------------	----------------------

Definition at line 88 of file [timer.cpp](#).

14.28.2.5 void Timer::SetIntervalSeconds (K_ULONG ulSeconds_)

! The next three cost us 330 bytes of flash on AVR...

Set the timer expiry interval in seconds (platform agnostic)

Parameters

<i>ulSeconds_</i>	Time in seconds
-------------------	-----------------

Definition at line 82 of file [timer.cpp](#).

14.28.2.6 void Timer::SetIntervalTicks (K_ULONG ulTicks_)

Set the timer expiry in system-ticks (platform specific!)

Parameters

<i>ulTicks_</i>	Time in ticks
-----------------	---------------

Definition at line 74 of file [timer.cpp](#).

14.28.2.7 void Timer::SetIntervalUSeconds (K_ULONG ulUSeconds_)

Set the timer expiry interval in microseconds (platform agnostic)

Parameters

<i>ulUSeconds_</i>	Time in microseconds
--------------------	----------------------

Definition at line 94 of file [timer.cpp](#).

14.28.2.8 void Timer::SetOwner (Thread * pOwner_) [inline]

Set the owner-thread of this timer object (all timers must be owned by a thread).

Parameters

<i>pOwner_</i>	Owner thread of this timer object
----------------	-----------------------------------

Definition at line 185 of file [timer.h](#).

14.28.2.9 void Timer::SetTolerance (K_ULONG ulTicks_)

Set the timer's maximum tolerance in order to synchronize timer processing with other timers in the system.

Parameters

<i>ulTicks_</i>	Maximum tolerance in ticks
-----------------	----------------------------

Definition at line 100 of file [timer.cpp](#).

14.28.2.10 void Timer::Start (K_BOOL bRepeat_, K_ULONG ulIntervalMs_, TimerCallback_t pfCallback_, void * pvData_)

Start a timer using default ownership, using repeats as an option, and millisecond resolution.

Parameters

<i>bRepeat_</i>	0 - timer is one-shot. 1 - timer is repeating.
<i>ulIntervalMs_</i>	- Interval of the timer in milliseconds
<i>pfCallback_</i>	- Function to call on timer expiry

<i>pvData_</i>	- Data to pass into the callback function
----------------	---

Definition at line 42 of file [timer.cpp](#).

14.28.2.11 void Timer::Start (K_BOOL bRepeat_, K_ULONG ulIntervalMs_, K_ULONG ulToleranceMs_, TimerCallback_t pfCallback_, void * pvData_)

Start a timer using default ownership, using repeats as an option, and millisecond resolution.

Parameters

<i>bRepeat_</i>	0 - timer is one-shot. 1 - timer is repeating.
<i>ulIntervalMs_</i>	- Interval of the timer in milliseconds
<i>ulToleranceMs_</i>	- Allow the timer expiry to be delayed by an additional maximum time, in order to have as many timers expire at the same time as possible.
<i>pfCallback_</i>	- Function to call on timer expiry
<i>pvData_</i>	- Data to pass into the callback function

Definition at line 61 of file [timer.cpp](#).

14.28.2.12 void Timer::Stop ()

Stop a timer already in progress.

Has no effect on timers that have already been stopped.

Definition at line 68 of file [timer.cpp](#).

The documentation for this class was generated from the following files:

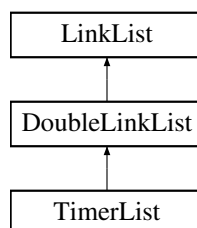
- [/home/moslevin/Project/R1/kernel/public/timer.h](#)
- [/home/moslevin/Project/R1/kernel/timer.cpp](#)

14.29 TimerList Class Reference

[TimerList](#) class - a doubly-linked-list of timer objects.

```
#include <timerlist.h>
```

Inheritance diagram for TimerList:



Public Member Functions

- void [Init](#) ()
Initialize the [TimerList](#) object.
- void [Add](#) (Timer *pclListNode_)
Add a timer to the [TimerList](#).
- void [Remove](#) (Timer *pclListNode_)

Remove a timer from the [TimerList](#), cancelling its expiry.

- void [Process](#) ()

Process all timers in the timerlist as a result of the timer expiring.

Private Attributes

- [K_ULONG m_ulNextWakeup](#)

The time (in system clock ticks) of the next wakeup event.

- [K_UCHAR m_bTimerActive](#)

Whether or not the timer is active.

Additional Inherited Members

14.29.1 Detailed Description

[TimerList](#) class - a doubly-linked-list of timer objects.

Definition at line 37 of file [timerlist.h](#).

14.29.2 Member Function Documentation

14.29.2.1 void [TimerList::Add](#) ([Timer](#) * [pclListNode_](#))

Add a timer to the [TimerList](#).

Parameters

pclListNode_	Pointer to the Timer to Add
------------------------------	---

Definition at line 51 of file [timerlist.cpp](#).

14.29.2.2 void [TimerList::Init](#) (void)

Initialize the [TimerList](#) object.

Must be called before using the object.

Definition at line 44 of file [timerlist.cpp](#).

14.29.2.3 void [TimerList::Process](#) (void)

Process all timers in the timerlist as a result of the timer expiring.

This will select a new timer epoch based on the next timer to expire. ToDo - figure out if we need to deal with any overtime here.

Definition at line 116 of file [timerlist.cpp](#).

14.29.2.4 void [TimerList::Remove](#) ([Timer](#) * [pclListNode_](#))

Remove a timer from the [TimerList](#), cancelling its expiry.

Parameters

<code>pcListNode_</code>	Pointer to the Timer to remove
--------------------------	--

Definition at line 99 of file [timerlist.cpp](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/timerlist.h](#)
- [/home/moslevin/Project/R1/kernel/timerlist.cpp](#)

14.30 TimerScheduler Class Reference

"Static" Class used to interface a global [TimerList](#) with the rest of the kernel.

```
#include <timerscheduler.h>
```

Static Public Member Functions

- static void [Init](#) ()
Initialize the timer scheduler.
- static void [Add](#) ([Timer](#) *`pcListNode_`)
Add a timer to the timer scheduler.
- static void [Remove](#) ([Timer](#) *`pcListNode_`)
Remove a timer from the timer scheduler.
- static void [Process](#) ()
This function must be called on timer expiry (from the timer's ISR context).

Static Private Attributes

- static [TimerList](#) `m_cTimerList`
[TimerList](#) object manipulated by the [Timer Scheduler](#).

14.30.1 Detailed Description

"Static" Class used to interface a global [TimerList](#) with the rest of the kernel.

Definition at line 38 of file [timerscheduler.h](#).

14.30.2 Member Function Documentation

14.30.2.1 void [TimerScheduler::Add](#) ([Timer](#) * `pcListNode_`) [inline], [static]

Add a timer to the timer scheduler.

Adding a timer implicitly starts the timer as well.

Parameters

<code>pcListNode_</code>	Pointer to the timer list node to add
--------------------------	---------------------------------------

Definition at line 57 of file [timerscheduler.h](#).

14.30.2.2 `void TimerScheduler::Init(void) [inline],[static]`

Initialize the timer scheduler.

Must be called before any timer, or timer-derived functions are used.

Definition at line 47 of file [timerscheduler.h](#).

14.30.2.3 `void TimerScheduler::Process(void) [inline],[static]`

This function must be called on timer expiry (from the timer's ISR context).

This will result in all timers being updated based on the epoch that just elapsed. The next timer epoch is set based on the next [Timer](#) object to expire.

Definition at line 79 of file [timerscheduler.h](#).

14.30.2.4 `void TimerScheduler::Remove(Timer * pclListNode_) [inline],[static]`

Remove a timer from the timer scheduler.

May implicitly stop the timer if this is the only active timer scheduled.

Parameters

<i>pclListNode_</i>	Pointer to the timer list node to remove
---------------------	--

Definition at line 68 of file [timerscheduler.h](#).

The documentation for this class was generated from the following files:

- [/home/moslevin/Project/R1/kernel/public/timerscheduler.h](#)
- [/home/moslevin/Project/R1/kernel/timerlist.cpp](#)

File Documentation

Basic Atomic Operations.

Definition in file [atomic.cpp](#).

15.2 atomic.cpp

```

00001 /*-----*/
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00021 #include "kerneltypes.h"
00022 #include "mark3cfg.h"
00023 #include "atomic.h"
00024 #include "threadport.h"
00025
00026 #if KERNEL_USE_ATOMIC
00027
00028 //-----
00029 K_UCHAR Atomic::Set( K_UCHAR *pucSource_, K_UCHAR ucVal_ )
00030 {
00031     K_UCHAR ucRet;
00032     CS_ENTER();
00033     ucRet = *pucSource_;
00034     *pucSource_ = ucVal_;
00035     CS_EXIT();
00036     return ucRet;
00037 }
00038 //-----
00039 K_USHORT Atomic::Set( K_USHORT *pusSource_, K_USHORT usVal_ )
00040 {

```

```

00041     K_USHORT usRet;
00042     CS_ENTER();
00043     usRet = *pusSource_;
00044     *pusSource_ = usVal_;
00045     CS_EXIT();
00046     return usRet;
00047 }
00048 //-----
00049 K_ULONG Atomic::Set( K_ULONG *pulSource_, K_ULONG ulVal_ )
00050 {
00051     K_ULONG ulRet;
00052     CS_ENTER();
00053     ulRet = *pulSource_;
00054     *pulSource_ = ulVal_;
00055     CS_EXIT();
00056     return ulRet;
00057 }
00058
00059 //-----
00060 K_UCHAR Atomic::Add( K_UCHAR *pucSource_, K_UCHAR ucVal_ )
00061 {
00062     K_UCHAR ucRet;
00063     CS_ENTER();
00064     ucRet = *pucSource_;
00065     *pucSource_ += ucVal_;
00066     CS_EXIT();
00067     return ucRet;
00068 }
00069
00070 //-----
00071 K_USHORT Atomic::Add( K_USHORT *pusSource_, K_USHORT usVal_ )
00072 {
00073     K_USHORT usRet;
00074     CS_ENTER();
00075     usRet = *pusSource_;
00076     *pusSource_ += usVal_;
00077     CS_EXIT();
00078     return usRet;
00079 }
00080
00081 //-----
00082 K_ULONG Atomic::Add( K_ULONG *pulSource_, K_ULONG ulVal_ )
00083 {
00084     K_ULONG ulRet;
00085     CS_ENTER();
00086     ulRet = *pulSource_;
00087     *pulSource_ += ulVal_;
00088     CS_EXIT();
00089     return ulRet;
00090 }
00091
00092 //-----
00093 K_UCHAR Atomic::Sub( K_UCHAR *pucSource_, K_UCHAR ucVal_ )
00094 {
00095     K_UCHAR ucRet;
00096     CS_ENTER();
00097     ucRet = *pucSource_;
00098     *pucSource_ -= ucVal_;
00099     CS_EXIT();
00100     return ucRet;
00101 }
00102
00103 //-----
00104 K_USHORT Atomic::Sub( K_USHORT *pusSource_, K_USHORT usVal_ )
00105 {
00106     K_USHORT usRet;
00107     CS_ENTER();
00108     usRet = *pusSource_;
00109     *pusSource_ -= usVal_;
00110     CS_EXIT();
00111     return usRet;
00112 }
00113
00114 //-----
00115 K_ULONG Atomic::Sub( K_ULONG *pulSource_, K_ULONG ulVal_ )
00116 {
00117     K_ULONG ulRet;
00118     CS_ENTER();
00119     ulRet = *pulSource_;
00120     *pulSource_ -= ulVal_;
00121     CS_EXIT();
00122     return ulRet;
00123 }
00124
00125 //-----
00126 K_BOOL Atomic::TestAndSet( K_BOOL *pbLock_ )
00127 {

```

```

00128     K_UCHAR ucRet;
00129     CS_ENTER();
00130     ucRet = *pbLock_;
00131     if (!ucRet)
00132     {
00133         *pbLock_ = 1;
00134     }
00135     CS_EXIT();
00136     return ucRet;
00137 }
00138
00139 #endif // KERNEL_USE_ATOMIC

```

15.3 /home/moslevin/Project/R1/kernel/blocking.cpp File Reference

Implementation of base class for blocking objects.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "kerneldebug.h"
#include "blocking.h"
#include "thread.h"

```

Macros

- `#define __FILE_ID__ BLOCKING_CPP`
File ID used in kernel trace calls.

15.3.1 Detailed Description

Implementation of base class for blocking objects.

Definition in file [blocking.cpp](#).

15.4 blocking.cpp

```

00001  /*=====
00002
00003  _____
00004  |   \   /   |   \   /   |   \   /   |   \   /   |   \   /   |
00005  |  / \ / \  |  / \ / \  |  / \ / \  |  / \ / \  |  / \ / \  |
00006  |_/___\___\_|_/___\___\_|_/___\___\_|_/___\___\_|_/___\___\_|
00007  |_____|_____|_____|_____|_____|_____|_____|_____|_____|_____|
00008
00009  --[Mark3 Realtime Platform]-----
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information
00013  =====*/
00021  #include "kerneltypes.h"
00022  #include "mark3cfg.h"
00023  #include "kerneldebug.h"
00024
00025  #include "blocking.h"
00026  #include "thread.h"
00027
00028  //-----
00029  #if defined __FILE_ID__
00030      #undef __FILE_ID__
00031  #endif
00032  #define __FILE_ID__      BLOCKING_CPP
00033
00034  #if KERNEL_USE_SEMAPHORE || KERNEL_USE_MUTEX
00035  //-----
00036  void BlockingObject::Block(Thread *pclThread_)
00037  {
00038      KERNEL_ASSERT( pclThread_ );

```

```

00039     KERNEL_TRACE_1( STR_THREAD_BLOCK_1, (K_USHORT)pclThread->
GetID() );
00040
00041     // Remove the thread from its current thread list (the "owner" list)
00042     // ... And add the thread to this object's block list
00043     Scheduler::Remove(pclThread_);
00044     m_clBlockList.Add(pclThread_);
00045
00046     // Set the "current" list location to the blocklist for this thread
00047     pclThread->SetCurrent(&m_clBlockList);
00048     pclThread->SetState(THREAD_STATE_BLOCKED);
00049 }
00050
00051 //-----
00052 void BlockingObject::UnBlock(Thread *pclThread_)
00053 {
00054     KERNEL_ASSERT( pclThread_ );
00055     KERNEL_TRACE_1( STR_THREAD_UNBLOCK_1, (K_USHORT)pclThread_->
GetID() );
00056
00057     // Remove the thread from its current thread list (the "owner" list)
00058     pclThread->GetCurrent()->Remove(pclThread_);
00059
00060     // Put the thread back in its active owner's list. This is usually
00061     // the ready-queue at the thread's original priority.
00062     Scheduler::Add(pclThread_);
00063
00064     // Tag the thread's current list location to its owner
00065     pclThread->SetCurrent(pclThread->GetOwner());
00066     pclThread->SetState(THREAD_STATE_READY);
00067 }
00068
00069 #endif

```

```

00021 #include "mark3cfg.h"
00022 #include "profile.h"
00023 #include "kernelprofile.h"
00024 #include "threadport.h"
00025 #include <avr/io.h>
00026 #include <avr/interrupt.h>
00027
00028 #if KERNEL_USE_PROFILER
00029 K_ULONG Profiler::m_ulEpoch;
00030
00031 //-----
00032 void Profiler::Init()
00033 {
00034     TCCR0A = 0;
00035     TCCR0B = 0;
00036     TIFR0 = 0;
00037     TIMSK0 = 0;
00038     m_ulEpoch = 0;
00039 }
00040
00041 //-----
00042 void Profiler::Start()
00043 {
00044     TIFR0 = 0;
00045     TCNT0 = 0;
00046     TCCR0B |= (1 << CS01);
00047     TIMSK0 |= (1 << TOIE0);
00048 }
00049
00050 //-----
00051 void Profiler::Stop()
00052 {
00053     TIFR0 = 0;
00054     TCCR0B &= ~(1 << CS01);
00055     TIMSK0 &= ~(1 << TOIE0);
00056 }
00057 //-----
00058 K_USHORT Profiler::Read()
00059 {
00060     K_USHORT usRet;
00061     CS_ENTER();
00062     TCCR0B &= ~(1 << CS01);
00063     usRet = TCNT0;
00064     TCCR0B |= (1 << CS01);
00065     CS_EXIT();
00066     return usRet;
00067 }
00068
00069 //-----
00070 void Profiler::Process()
00071 {
00072     CS_ENTER();
00073     m_ulEpoch++;
00074     CS_EXIT();
00075 }
00076
00077 //-----
00078 ISR(TIMER0_OVF_vect)
00079 {
00080     Profiler::Process();
00081 }
00082
00083 #endif

```

15.7 /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/kernelswi.cpp File Reference

Kernel Software interrupt implementation for ATmega328p.

```

#include "kerneltypes.h"
#include "kernelswi.h"
#include <avr/io.h>
#include <avr/interrupt.h>

```

15.7.1 Detailed Description

Kernel Software interrupt implementation for ATmega328p.

Definition in file [kernelswi.cpp](#).

15.8 kernelswi.cpp

```

00001 /*=====
00002
00003
00004 | | | | | | | | | | | | | | | | | |
00005 | | | | | | | | | | | | | | | | | |
00006 | | | | | | | | | | | | | | | | | |
00007 | | | | | | | | | | | | | | | | | |
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00022 #include "kerneltypes.h"
00023 #include "kernelswi.h"
00024
00025 #include <avr/io.h>
00026 #include <avr/interrupt.h>
00027
00028 //-----
00029 void KernelSWI::Config(void)
00030 {
00031     PORTD &= ~0x04; // Clear INT0
00032     DDRD |= 0x04;    // Set PortD, bit 2 (INT0) As Output
00033     EICRA |= (1 << ISC00) | (1 << ISC01); // Rising edge on INT0
00034 }
00035
00036 //-----
00037 void KernelSWI::Start(void)
00038 {
00039     EIFR &= ~(1 << INTF0); // Clear any pending interrupts on INT0
00040     EIMSK |= (1 << INT0);  // Enable INT0 interrupt (as K_LONG as I-bit is set)
00041 }
00042
00043 //-----
00044 void KernelSWI::Stop(void)
00045 {
00046     EIMSK &= ~(1 << INT0); // Disable INT0 interrupts
00047 }
00048
00049 //-----
00050 K_UCHAR KernelSWI::DI()
00051 {
00052     K_BOOL bEnabled = ((EIMSK & (1 << INT0)) != 0);
00053     EIMSK &= ~(1 << INT0);
00054     return bEnabled;
00055 }
00056
00057 //-----
00058 void KernelSWI::RI(K_BOOL bEnable_)
00059 {
00060     if (bEnable_)
00061     {
00062         EIMSK |= (1 << INT0);
00063     }
00064     else
00065     {
00066         EIMSK &= ~(1 << INT0);
00067     }
00068 }
00069
00070 //-----
00071 void KernelSWI::Clear(void)
00072 {
00073     EIFR &= ~(1 << INTF0); // Clear the interrupt flag for INT0
00074 }
00075
00076 //-----
00077 void KernelSWI::Trigger(void)
00078 {
00079     //if(Thread_IsSchedulerEnabled())
00080     {
00081         PORTD &= ~0x04;
00082         PORTD |= 0x04;

```



```
00083     }
00084 }
```

15.9 /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/kerneltimer.cpp File Reference

[Kernel Timer](#) Implementation for ATmega328p.

```
#include "kerneltypes.h"
#include "kerneltimer.h"
#include "mark3cfg.h"
#include <avr/io.h>
#include <avr/interrupt.h>
```

15.9.1 Detailed Description

[Kernel Timer](#) Implementation for ATmega328p.

Definition in file [kerneltimer.cpp](#).

15.10 kerneltimer.cpp

```
00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00021 #include "kerneltypes.h"
00022 #include "kerneltimer.h"
00023 #include "mark3cfg.h"
00024
00025 #include <avr/io.h>
00026 #include <avr/interrupt.h>
00027
00028 #define TCCR1B_INIT ((1 << WGM12) | (1 << CS12))
00029 #define TIMER_IMSK (1 << OCIE1A)
00030 #define TIMER_IFR (1 << OCF1A)
00031
00032 //-----
00033 void KernelTimer::Config(void)
00034 {
00035     TCCR1B = TCCR1B_INIT;
00036 }
00037
00038 //-----
00039 void KernelTimer::Start(void)
00040 {
00041     #if !KERNEL_TIMERS_TICKLESS
00042         TCCR1B = ((1 << WGM12) | (1 << CS11) | (1 << CS10));
00043         OCR1A = (SYSTEM_FREQ / 1000) / 64;
00044     #else
00045         TCCR1B |= (1 << CS12);
00046     #endif
00047
00048     TCNT1 = 0;
00049     TIFR1 &= ~TIMER_IFR;
00050     TIMSK1 |= TIMER_IMSK;
00051 }
00052
00053 //-----
00054 void KernelTimer::Stop(void)
00055 {
```

```

00056 #if KERNEL_TIMERS_TICKLESS
00057     TIFR1 &= ~TIMER_IFR;
00058     TIMSK1 &= ~TIMER_IMSK;
00059     TCCR1B &= ~(1 << CS12);    // Disable count...
00060     TCNT1 = 0;
00061     OCR1A = 0;
00062 #endif
00063 }
00064
00065 //-----
00066 K_USHORT KernelTimer::Read(void)
00067 {
00068     #if KERNEL_TIMERS_TICKLESS
00069         volatile K_USHORT usRead1;
00070         volatile K_USHORT usRead2;
00071
00072         do {
00073             usRead1 = TCNT1;
00074             usRead2 = TCNT1;
00075         } while (usRead1 != usRead2);
00076
00077         return usRead1;
00078     #else
00079         return 0;
00080     #endif
00081 }
00082
00083 //-----
00084 K_ULONG KernelTimer::SubtractExpiry(K_ULONG ulInterval_)
00085 {
00086     #if KERNEL_TIMERS_TICKLESS
00087         OCR1A -= (K_USHORT)ulInterval_;
00088         return (K_ULONG)OCR1A;
00089     #else
00090         return 0;
00091     #endif
00092 }
00093
00094 //-----
00095 K_ULONG KernelTimer::TimeToExpiry(void)
00096 {
00097     #if KERNEL_TIMERS_TICKLESS
00098         K_USHORT usRead = KernelTimer::Read();
00099         K_USHORT usOCR1A = OCR1A;
00100
00101         if (usRead >= usOCR1A)
00102         {
00103             return 0;
00104         }
00105         else
00106         {
00107             return (K_ULONG)(usOCR1A - usRead);
00108         }
00109     #else
00110         return 0;
00111     #endif
00112 }
00113
00114 //-----
00115 K_ULONG KernelTimer::GetOvertime(void)
00116 {
00117     return KernelTimer::Read();
00118 }
00119
00120 //-----
00121 K_ULONG KernelTimer::SetExpiry(K_ULONG ulInterval_)
00122 {
00123     #if KERNEL_TIMERS_TICKLESS
00124         K_USHORT usSetInterval;
00125         if (ulInterval_ > 65535)
00126         {
00127             usSetInterval = 65535;
00128         }
00129         else
00130         {
00131             usSetInterval = (K_USHORT)ulInterval_ ;
00132         }
00133
00134         OCR1A = usSetInterval;
00135         return (K_ULONG)usSetInterval;
00136     #else
00137         return 0;
00138     #endif
00139 }
00140
00141 //-----
00142 void KernelTimer::ClearExpiry(void)

```

```

00143 {
00144 #if KERNEL_TIMERS_TICKLESS
00145     OCR1A = 65535;                // Clear the compare value
00146 #endif
00147 }
00148
00149 //-----
00150 K_UCHAR KernelTimer::DI(void)
00151 {
00152 #if KERNEL_TIMERS_TICKLESS
00153     K_BOOL bEnabled = ((TIMSK1 & (TIMER_IMSK)) != 0);
00154     TIFR1 &= ~TIMER_IFR;         // Clear interrupt flags
00155     TIMSK1 &= ~TIMER_IMSK;       // Disable interrupt
00156     return bEnabled;
00157 #else
00158     return 0;
00159 #endif
00160 }
00161
00162 //-----
00163 void KernelTimer::EI(void)
00164 {
00165     KernelTimer::RI(0);
00166 }
00167
00168 //-----
00169 void KernelTimer::RI(K_BOOL bEnable_)
00170 {
00171 #if KERNEL_TIMERS_TICKLESS
00172     if (bEnable_)
00173     {
00174         TIMSK1 |= (1 << OCIE1A);    // Enable interrupt
00175     }
00176     else
00177     {
00178         TIMSK1 &= ~(1 << OCIE1A);
00179     }
00180 #endif
00181 }

```

15.11 /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/kernelprofile.h File Reference

Profiling timer hardware interface.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "ll.h"

```

Classes

- class [Profiler](#)
System profiling timer interface.

15.11.1 Detailed Description

Profiling timer hardware interface.

Definition in file [kernelprofile.h](#).

15.12 kernelprofile.h

```

00001 /*=====
00002
00003 |-----|-----|-----|-----|-----|
00004 |  \  /  |  \  /  |  \  /  |  \  /  |  \  /  |
00005 |  \  /  |  \  /  |  \  /  |  \  /  |  \  /  |

```



```

00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00023 #include "kerneltypes.h"
00024 #ifndef __KERNELSWI_H_
00025 #define __KERNELSWI_H_
00026
00027 //-----
00032 class KernelSWI
00033 {
00034 public:
00041     static void Config(void);
00042
00048     static void Start(void);
00049
00055     static void Stop(void);
00056
00062     static void Clear(void);
00063
00069     static void Trigger(void);
00070
00078     static K_UCHAR DI();
00079
00087     static void RI(K_BOOL bEnable_);
00088 };
00089
00090
00091 #endif // __KERNELSWI_H_

```

15.15 /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/kerneltimer.h File Reference

[Kernel Timer](#) Class declaration.

```
#include "kerneltypes.h"
```

Classes

- class [KernelTimer](#)
Hardware timer interface, used by all scheduling/timer subsystems.

15.15.1 Detailed Description

[Kernel Timer](#) Class declaration.

Definition in file [kerneltimer.h](#).

15.16 kerneltimer.h

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.

```

```

00012 See license.txt for more information
00013 =====*/
00021 #include "kerneltypes.h"
00022 #ifndef __KERNELTIMER_H_
00023 #define __KERNELTIMER_H_
00024
00025 //-----
00026 #define SYSTEM_FREQ      ((K_ULONG)16000000)
00027 #define TIMER_FREQ      ((K_ULONG)(SYSTEM_FREQ / 256)) // Timer ticks per second...
00028
00029 //-----
00033 class KernelTimer
00034 {
00035 public:
00041     static void Config(void);
00042
00048     static void Start(void);
00049
00055     static void Stop(void);
00056
00062     static K_UCHAR DI(void);
00063
00071     static void RI(K_BOOL bEnable_);
00072
00078     static void EI(void);
00079
00090     static K_ULONG SubtractExpiry(K_ULONG ulInterval_);
00091
00100     static K_ULONG TimeToExpiry(void);
00101
00110     static K_ULONG SetExpiry(K_ULONG ulInterval_);
00111
00120     static K_ULONG GetOvertime(void);
00121
00127     static void ClearExpiry(void);
00128
00129 private:
00137     static K_USHORT Read(void);
00138
00139 };
00140
00141 #endif //__KERNELTIMER_H_

```

15.17 /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/public/threadport.h File Reference

ATMega328p Multithreading support.

```

#include "kerneltypes.h"
#include "thread.h"
#include <avr/io.h>
#include <avr/interrupt.h>

```

Classes

- class [ThreadPort](#)
Class defining the architecture specific functions required by the kernel.

Macros

- #define [ASM](#)(x) asm volatile(x);
ASM Macro - simplify the use of ASM directive in C.
- #define [SR_](#) 0x3F
Status register define - map to 0x003F.
- #define [SPH_](#) 0x3E
Stack pointer define.
- #define [TOP_OF_STACK](#)(x, y) (K_UCHAR*) (((K_USHORT)x) + (y-1))

- Macro to find the top of a stack given its size and top address.
- `#define PUSH_TO_STACK(x, y) *x = y; x--;`
Push a value *y* to the stack pointer *x* and decrement the stack pointer.
- `#define Thread_SaveContext()`
Save the context of the *Thread*.
- `#define Thread_RestoreContext()`
Restore the context of the *Thread*.
- `#define CS_ENTER()`
These macros must be used in pairs !
- `#define CS_EXIT()`
Exit critical section (restore status register)
- `#define ENABLE_INTS() ASM("sei");`
Initiate a contex switch without using the SWI.

15.17.1 Detailed Description

ATMega328p Multithreading support.

Definition in file [threadport.h](#).

15.17.2 Macro Definition Documentation

15.17.2.1 `#define CS_ENTER()`

Value:

```
{ \
volatile K_UCHAR x; \
x = _SFR_IO8(SR_); \
ASM("cli");
```

These macros *must* be used in pairs !

Enter critical section (copy status register, disable interrupts)

Definition at line 142 of file [threadport.h](#).

15.18 threadport.h

```
00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00021 #ifndef __THREADPORT_H_
00022 #define __THREADPORT_H_
00023
00024 #include "kerneltypes.h"
00025 #include "thread.h"
00026
00027 #include <avr/io.h>
00028 #include <avr/interrupt.h>
00029
00030 //-----
00032 #define ASM(x) asm volatile(x);
```

```

00033 #define SR_          0x3F
00035 #define SPH_          0x3E
00037 #define SPL_          0x3D
00038
00039
00040 //-----
00042 #define TOP_OF_STACK(x, y)      (K_UCHAR*) ( ((K_USHORT)x) + (y-1) )
00043 #define PUSH_TO_STACK(x, y)    *x = y; x--;
00045
00046 //-----
00048 #define Thread_SaveContext() \
00049 ASM("push r0"); \
00050 ASM("in r0, __SREG__"); \
00051 ASM("cli"); \
00052 ASM("push r0"); \
00053 ASM("push r1"); \
00054 ASM("clr r1"); \
00055 ASM("push r2"); \
00056 ASM("push r3"); \
00057 ASM("push r4"); \
00058 ASM("push r5"); \
00059 ASM("push r6"); \
00060 ASM("push r7"); \
00061 ASM("push r8"); \
00062 ASM("push r9"); \
00063 ASM("push r10"); \
00064 ASM("push r11"); \
00065 ASM("push r12"); \
00066 ASM("push r13"); \
00067 ASM("push r14"); \
00068 ASM("push r15"); \
00069 ASM("push r16"); \
00070 ASM("push r17"); \
00071 ASM("push r18"); \
00072 ASM("push r19"); \
00073 ASM("push r20"); \
00074 ASM("push r21"); \
00075 ASM("push r22"); \
00076 ASM("push r23"); \
00077 ASM("push r24"); \
00078 ASM("push r25"); \
00079 ASM("push r26"); \
00080 ASM("push r27"); \
00081 ASM("push r28"); \
00082 ASM("push r29"); \
00083 ASM("push r30"); \
00084 ASM("push r31"); \
00085 ASM("lds r26, g_pstCurrent"); \
00086 ASM("lds r27, g_pstCurrent + 1"); \
00087 ASM("adiw r26, 4"); \
00088 ASM("in r0, 0x3D"); \
00089 ASM("st x+, r0"); \
00090 ASM("in r0, 0x3E"); \
00091 ASM("st x+, r0");
00092
00093 //-----
00095 #define Thread_RestoreContext() \
00096 ASM("lds r26, g_pstCurrent"); \
00097 ASM("lds r27, g_pstCurrent + 1"); \
00098 ASM("adiw r26, 4"); \
00099 ASM("ld r28, x+"); \
00100 ASM("out 0x3D, r28"); \
00101 ASM("ld r29, x+"); \
00102 ASM("out 0x3E, r29"); \
00103 ASM("pop r31"); \
00104 ASM("pop r30"); \
00105 ASM("pop r29"); \
00106 ASM("pop r28"); \
00107 ASM("pop r27"); \
00108 ASM("pop r26"); \
00109 ASM("pop r25"); \
00110 ASM("pop r24"); \
00111 ASM("pop r23"); \
00112 ASM("pop r22"); \
00113 ASM("pop r21"); \
00114 ASM("pop r20"); \
00115 ASM("pop r19"); \
00116 ASM("pop r18"); \
00117 ASM("pop r17"); \
00118 ASM("pop r16"); \
00119 ASM("pop r15"); \
00120 ASM("pop r14"); \
00121 ASM("pop r13"); \
00122 ASM("pop r12"); \
00123 ASM("pop r11"); \
00124 ASM("pop r10"); \
00125 ASM("pop r9"); \

```



```

00126 ASM("pop r8"); \
00127 ASM("pop r7"); \
00128 ASM("pop r6"); \
00129 ASM("pop r5"); \
00130 ASM("pop r4"); \
00131 ASM("pop r3"); \
00132 ASM("pop r2"); \
00133 ASM("pop r1"); \
00134 ASM("pop r0"); \
00135 ASM("out __SREG__, r0"); \
00136 ASM("pop r0");
00137
00138 //-----
00140 //-----
00142 #define CS_ENTER() \
00143 { \
00144 volatile K_UCHAR x; \
00145 x = _SFR_IO8(SR_); \
00146 ASM("cli");
00147 //-----
00149 #define CS_EXIT() \
00150 _SFR_IO8(SR_) = x;\
00151 }
00152
00153 //-----
00155 #define ENABLE_INTS() ASM("sei");
00156 #define DISABLE_INTS() ASM("cli");
00157
00158 //-----
00159 class Thread;
00167 class ThreadPort
00168 {
00169 public:
00175 static void StartThreads();
00176 friend class Thread;
00177 private:
00178
00186 static void InitStack(Thread *pstThread_);
00187 };
00188
00189 #endif //__ThreadPORT_H_

```

15.19 /home/moslevin/Project/R1/kernel/cpu/avr/atmega328p/gcc/threadport.cpp File Reference

ATMega328p Multithreading.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "thread.h"
#include "threadport.h"
#include "kernelswi.h"
#include "kerneltimer.h"
#include "timerlist.h"
#include "quantum.h"
#include "kernel.h"
#include "kernelaware.h"
#include <avr/io.h>
#include <avr/interrupt.h>

```

Functions

- [ISR \(INT0_vect\) __attribute__\(\(signal](#)
SWI using INT0 - used to trigger a context switch.
- [ISR \(TIMER1_COMPA_vect\)](#)
Timer interrupt ISR - causes a tick, which may cause a context switch.


```

00081     for (i = 26; i <=31; i++)
00082     {
00083         PUSH_TO_STACK(pucStack, i);
00084     }
00085
00086     // Set the top o' the stack.
00087     pclThread->m_pwStackTop = (K_UCHAR*)pucStack;
00088
00089     // That's it!  the thread is ready to run now.
00090 }
00091
00092 //-----
00093 static void Thread_Switch(void)
00094 {
00095     #if KERNEL_USE_IDLE_FUNC
00096         // If there's no next-thread-to-run...
00097         if (g_pstNext == Kernel::GetIdleThread())
00098         {
00099             g_pstCurrent = Kernel::GetIdleThread();
00100
00101             // Disable the SWI, and re-enable interrupts -- enter nested interrupt
00102             // mode.
00103             KernelSWI::DI();
00104
00105             K_UCHAR ucSR = _SFR_IO8(SR_);
00106
00107             // So long as there's no "next-to-run" thread, keep executing the Idle
00108             // function to conclusion...
00109
00110             while (g_pstNext == Kernel::GetIdleThread())
00111             {
00112                 // Ensure that we run this block in an interrupt enabled context (but
00113                 // with the rest of the checks being performed in an interrupt disabled
00114                 // context).
00115                 ASM( "sei" );
00116                 Kernel::IdleFunc();
00117                 ASM( "cli" );
00118             }
00119
00120             // Progress has been achieved -- an interrupt-triggered event has caused
00121             // the scheduler to run, and choose a new thread.  Since we've already
00122             // saved the context of the thread we've hijacked to run idle, we can
00123             // proceed to disable the nested interrupt context and switch to the
00124             // new thread.
00125
00126             _SFR_IO8(SR_) = ucSR;
00127             KernelSWI::RI( true );
00128         }
00129     #endif
00130     g_pstCurrent = (Thread*)g_pstNext;
00131 }
00132
00133 //-----
00134 //-----
00135 void ThreadPort::StartThreads()
00136 {
00137     KernelSWI::Config();           // configure the task switch SWI
00138     KernelTimer::Config();         // configure the kernel timer
00139
00140     Scheduler::SetScheduler(1);    // enable the scheduler
00141     Scheduler::Schedule();         // run the scheduler - determine the first
00142                                     thread to run
00143
00144     Thread_Switch();               // Set the next scheduled thread to the current thread
00145
00146     KernelTimer::Start();          // enable the kernel timer
00147     KernelSWI::Start();            // enable the task switch SWI
00148
00149     // Restore the context...
00150     Thread_RestoreContext();       // restore the context of the first running thread
00151     ASM("reti");                  // return from interrupt - will return to the first scheduled thread
00152 }
00153 //-----
00154 //-----
00155 ISR(INT0_vect) __attribute__((signal, naked));
00160 ISR(INT0_vect)
00161 {
00162     Thread_SaveContext();          // Push the context (registers) of the current task
00163     Thread_Switch();              // Switch to the next task
00164     Thread_RestoreContext();       // Pop the context (registers) of the next task
00165     ASM("reti");                  // Return to the next task
00166 }
00167
00168 //-----
00173 //-----
00174 ISR(TIMER1_COMPA_vect)

```

```

00175 {
00176     #if KERNEL_USE_TIMERS
00177         TimerScheduler::Process();
00178     #endif
00179     #if KERNEL_USE_QUANTUM
00180         Quantum::UpdateTimer();
00181     #endif
00182 }

```

15.21 /home/moslevin/Project/R1/kernel/driver.cpp File Reference

Device driver/hardware abstraction layer.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "kerneldebug.h"
#include "driver.h"

```

Classes

- class [DevNull](#)
This class implements the "default" driver (/dev/null)

Macros

- #define [__FILE_ID__](#) DRIVER_CPP
File ID used in kernel trace calls.

Functions

- static [K_UCHAR](#) [DrvCmp](#) (const [K_CHAR](#) *szStr1_, const [K_CHAR](#) *szStr2_)
DrvCmp.

Variables

- static [DevNull](#) [clDevNull](#)
Default driver included to allow for run-time "stubbing".

15.21.1 Detailed Description

Device driver/hardware abstraction layer.

Definition in file [driver.cpp](#).

15.21.2 Function Documentation

15.21.2.1 static [K_UCHAR](#) [DrvCmp](#) (const [K_CHAR](#) * [szStr1_](#), const [K_CHAR](#) * [szStr2_](#)) [static]

[DrvCmp](#).

String comparison function used to compare input driver name against a known driver name in the existing driver list.

Parameters

<i>szStr1_</i>	user-specified driver name
<i>szStr2_</i>	name of a driver, provided from the driver table

Returns

1 on match, 0 on no-match

Definition at line 75 of file [driver.cpp](#).

15.22 driver.cpp

```

00001  /*=====
00002
00003  _____
00004  |   /   \   |   /   \   |   /   \   |   /   \   |
00005  |  /     \  |  /     \  |  /     \  |  /     \  |
00006  | /       \ | /       \ | /       \ | /       \ |
00007  |_____|   |_____|   |_____|   |_____|   |
00008
00009  --[Mark3 Realtime Platform]-----
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information
00013  =====*/
00021  #include "kerneltypes.h"
00022  #include "mark3cfg.h"
00023  #include "kerneldebug.h"
00024  #include "driver.h"
00025
00026  //-----
00027  #if defined __FILE_ID__
00028      #undef __FILE_ID__
00029  #endif
00030  #define __FILE_ID__      DRIVER_CPP
00031
00032  //-----
00033  #if KERNEL_USE_DRIVER
00034
00035  DoubleLinkedList DriverList::m_clDriverList;
00036
00040  class DevNull : public Driver
00041  {
00042  public:
00043      virtual void Init() { SetName("/dev/null"); };
00044      virtual K_UCHAR Open() { return 0; }
00045      virtual K_UCHAR Close() { return 0; }
00046
00047      virtual K_USHORT Read( K_USHORT usBytes_,
00048          K_UCHAR *pucData_){ return 0; }
00049
00050      virtual K_USHORT Write( K_USHORT usBytes_,
00051          K_UCHAR *pucData_ ) { return 0; }
00052
00053      virtual K_USHORT Control( K_USHORT usEvent_,
00054          void *pvDataIn_,
00055          K_USHORT usSizeIn_,
00056          void *pvDataOut_,
00057          K_USHORT usSizeOut_ ) { return 0; }
00058
00059  };
00060
00061  //-----
00062  static DevNull clDevNull;
00063
00064  //-----
00075  static K_UCHAR DrvCmp( const K_CHAR *szStr1_, const K_CHAR *szStr2_ )
00076  {
00077      K_CHAR *szTmp1 = (K_CHAR*) szStr1_;
00078      K_CHAR *szTmp2 = (K_CHAR*) szStr2_;
00079
00080      while (*szTmp1 && *szTmp2)
00081      {
00082          if (*szTmp1++ != *szTmp2++)
00083          {
00084              return 0;
00085          }
00086      }

```

```

00087
00088 // Both terminate at the same length
00089 if (!(*szTmp1) && !(*szTmp2))
00090 {
00091     return 1;
00092 }
00093
00094 return 0;
00095 }
00096
00097 //-----
00098 void DriverList::Init()
00099 {
00100     // Ensure we always have at least one entry - a default in case no match
00101     // is found (/dev/null)
00102     clDevNull.Init();
00103     Add(&clDevNull);
00104 }
00105
00106 //-----
00107 Driver *DriverList::FindByPath( const K_CHAR *m_pcPath )
00108 {
00109     KERNEL_ASSERT( m_pcPath );
00110     Driver *pclTemp = static_cast<Driver*>(m_clDriverList.
    GetHead());
00111
00112     // Iterate through the list of drivers until we find a match, or we
00113     // exhaust our list of installed drivers
00114     while (pclTemp)
00115     {
00116         if(DrvCmp(m_pcPath, pclTemp->GetPath()))
00117         {
00118             return pclTemp;
00119         }
00120         pclTemp = static_cast<Driver*>(pclTemp->GetNext());
00121     }
00122     // No matching driver found - return a pointer to our /dev/null driver
00123     return &clDevNull;
00124 }
00125
00126 #endif

```

15.23 /home/moslevin/Project/R1/kernel/eventflag.cpp File Reference

Event Flag Blocking Object/IPC-Object implementation.

```

#include "mark3cfg.h"
#include "blocking.h"
#include "kernel.h"
#include "thread.h"
#include "eventflag.h"
#include "kernelaware.h"
#include "timerlist.h"

```

Functions

- void [TimedEventFlag_Callback](#) (Thread *pclOwner_, void *pvData_)
TimedEventFlag_Callback.

15.23.1 Detailed Description

Event Flag Blocking Object/IPC-Object implementation.

Definition in file [eventflag.cpp](#).

15.23.2 Function Documentation


```

00076 #endif
00077
00078 // Ensure we're operating in a critical section while we determine
00079 // whether or not we need to block the current thread on this object.
00080 CS_ENTER();
00081
00082 // Check to see whether or not the current mask matches any of the
00083 // desired bits.
00084 g_pstCurrent->SetEventFlagMask(usMask_);
00085
00086 if ((eMode_ == EVENT_FLAG_ALL) || (eMode_ ==
EVENT_FLAG_ALL_CLEAR))
00087 {
00088     // Check to see if the flags in their current state match all of
00089     // the set flags in the event flag group, with this mask.
00090     if ((m_usSetMask & usMask_) == usMask_)
00091     {
00092         bMatch = true;
00093         g_pstCurrent->SetEventFlagMask(usMask_);
00094     }
00095 }
00096 else if ((eMode_ == EVENT_FLAG_ANY) || (eMode_ ==
EVENT_FLAG_ANY_CLEAR))
00097 {
00098     // Check to see if the existing flags match any of the set flags in
00099     // the event flag group with this mask
00100     if (m_usSetMask & usMask_)
00101     {
00102         bMatch = true;
00103         g_pstCurrent->SetEventFlagMask(m_usSetMask & usMask_);
00104     }
00105 }
00106
00107 // We're unable to match this pattern as-is, so we must block.
00108 if (!bMatch)
00109 {
00110     // Reset the current thread's event flag mask & mode
00111     g_pstCurrent->SetEventFlagMask(usMask_);
00112     g_pstCurrent->SetEventFlagMode(eMode_);
00113
00114 #if KERNEL_USE_TIMEOUTS
00115     if (ulTimeMS_)
00116     {
00117         g_pstCurrent->SetExpired(false);
00118         clEventTimer.Init();
00119         clEventTimer.Start(0, ulTimeMS_, TimedEventFlag_Callback, (void*)
this);
00120         bUseTimer = true;
00121     }
00122 #endif
00123
00124 // Add the thread to the object's block-list.
00125 Block(g_pstCurrent);
00126
00127 // Trigger that
00128 bThreadYield = true;
00129 }
00130
00131 // If bThreadYield is set, it means that we've blocked the current thread,
00132 // and must therefore rerun the scheduler to determine what thread to
00133 // switch to.
00134 if (bThreadYield)
00135 {
00136     // Switch threads immediately
00137     Thread::Yield();
00138 }
00139
00140 // Exit the critical section and return back to normal execution
00141 CS_EXIT();
00142
00147 #if KERNEL_USE_TIMEOUTS
00148 if (bUseTimer && bThreadYield)
00149 {
00150     clEventTimer.Stop();
00151 }
00152 #endif
00153
00154 return g_pstCurrent->GetEventFlagMask();
00155 }
00156
00157 //-----
00158 K_USHORT EventFlag::Wait(K_USHORT usMask_,
EventFlagOperation_t eMode_)
00159 {
00160 #if KERNEL_USE_TIMEOUTS
00161     return Wait_i(usMask_, eMode_, 0);
00162 #else

```



```

00163     return Wait_i(usMask_, eMode_);
00164 #endif
00165 }
00166
00167 #if KERNEL_USE_TIMEOUTS
00168 //-----
00169 K_USHORT EventFlag::Wait(K_USHORT usMask_,
    EventFlagOperation_t eMode_, K_ULONG ulTimeMS_)
00170 {
00171     return Wait_i(usMask_, eMode_, ulTimeMS_);
00172 }
00173 #endif
00174
00175 //-----
00176 void EventFlag::Set(K_USHORT usMask_)
00177 {
00178     Thread *pclPrev;
00179     Thread *pclCurrent;
00180     K_BOOL bReschedule = false;
00181     K_USHORT usNewMask;
00182
00183     CS_ENTER();
00184
00185     // Walk through the whole block list, checking to see whether or not
00186     // the current flag set now matches any/all of the masks and modes of
00187     // the threads involved.
00188
00189     m_usSetMask |= usMask_;
00190     usNewMask = m_usSetMask;
00191
00192     // Start at the head of the list, and iterate through until we hit the
00193     // "head" element in the list again. Ensure that we handle the case where
00194     // we remove the first or last elements in the list, or if there's only
00195     // one element in the list.
00196     pclCurrent = static_cast<Thread*>(m_clBlockList.GetHead());
00197
00198     // Do nothing when there are no objects blocking.
00199     if (pclCurrent)
00200     {
00201         // First loop - process every thread in the block-list and check to
00202         // see whether or not the current flags match the event-flag conditions
00203         // on the thread.
00204         do
00205         {
00206             pclPrev = pclCurrent;
00207             pclCurrent = static_cast<Thread*>(pclCurrent->GetNext());
00208
00209             // Read the thread's event mask/mode
00210             K_USHORT usThreadMask = pclPrev->GetEventFlagMask();
00211             EventFlagOperation_t eThreadMode = pclPrev->
GetEventFlagMode();
00212
00213             // For the "any" mode - unblock the blocked threads if one or more bits
00214             // in the thread's bitmask match the object's bitmask
00215             if ((EVENT_FLAG_ANY == eThreadMode) || (
EVENT_FLAG_ANY_CLEAR == eThreadMode))
00216             {
00217                 if (usThreadMask & m_usSetMask)
00218                 {
00219                     pclPrev->SetEventFlagMode(
EVENT_FLAG_PENDING_UNBLOCK);
00220                     pclPrev->SetEventFlagMask(m_usSetMask & usThreadMask);
00221                     bReschedule = true;
00222
00223                     // If the "clear" variant is set, then clear the bits in the mask
00224                     // that caused the thread to unblock.
00225                     if (EVENT_FLAG_ANY_CLEAR == eThreadMode)
00226                     {
00227                         usNewMask &=~ (usThreadMask & usMask_);
00228                     }
00229                 }
00230             }
00231             // For the "all" mode, every set bit in the thread's requested bitmask must
00232             // match the object's flag mask.
00233             else if ((EVENT_FLAG_ALL == eThreadMode) || (
EVENT_FLAG_ALL_CLEAR == eThreadMode))
00234             {
00235                 if ((usThreadMask & m_usSetMask) == usThreadMask)
00236                 {
00237                     pclPrev->SetEventFlagMode(
EVENT_FLAG_PENDING_UNBLOCK);
00238                     pclPrev->SetEventFlagMask(usThreadMask);
00239                     bReschedule = true;
00240
00241                     // If the "clear" variant is set, then clear the bits in the mask
00242                     // that caused the thread to unblock.
00243                     if (EVENT_FLAG_ALL_CLEAR == eThreadMode)

```

```

00244         {
00245             usNewMask &=~ (usThreadMask & usMask_);
00246         }
00247     }
00248 }
00249 }
00250 // To keep looping, ensure that there's something in the list, and
00251 // that the next item isn't the head of the list.
00252 while (pclPrev != m_clBlockList.GetTail());
00253
00254 // Second loop - go through and unblock all of the threads that
00255 // were tagged for unblocking.
00256 pclCurrent = static_cast<Thread*>(m_clBlockList.
GetHead());
00257 K_BOOL bIsTail = false;
00258 do
00259 {
00260     pclPrev = pclCurrent;
00261     pclCurrent = static_cast<Thread*>(pclCurrent->GetNext());
00262
00263     // Check to see if this is the condition to terminate the loop
00264     if (pclPrev == m_clBlockList.GetTail())
00265     {
00266         bIsTail = true;
00267     }
00268
00269     // If the first pass indicated that this thread should be
00270     // unblocked, then unblock the thread
00271     if (pclPrev->GetEventFlagMode() ==
EVENT_FLAG_PENDING_UNBLOCK)
00272     {
00273         Unblock(pclPrev);
00274     }
00275 }
00276 while (!bIsTail);
00277 }
00278
00279 // If we awoke any threads, re-run the scheduler
00280 if (bReschedule)
00281 {
00282     Thread::Yield();
00283 }
00284
00285 // Update the bitmask based on any "clear" operations performed along
00286 // the way
00287 m_usSetMask = usNewMask;
00288
00289 // Restore interrupts - will potentially cause a context switch if a
00290 // thread is unblocked.
00291 CS_EXIT();
00292 }
00293
00294 //-----
00295 void EventFlag::Clear(K_USHORT usMask_)
00296 {
00297     // Just clear the bitfields in the local object.
00298     CS_ENTER();
00299     m_usSetMask &= ~usMask_;
00300     CS_EXIT();
00301 }
00302
00303 //-----
00304 K_USHORT EventFlag::GetMask()
00305 {
00306     // Return the presently held event flag values in this object. Ensure
00307     // we get this within a critical section to guarantee atomicity.
00308     K_USHORT usReturn;
00309     CS_ENTER();
00310     usReturn = m_usSetMask;
00311     CS_EXIT();
00312     return usReturn;
00313 }
00314
00315 #endif // KERNEL_USE_EVENTFLAG

```

15.25 /home/moslevin/Project/R1/kernel/kernel.cpp File Reference

Kernel initialization and startup code.

```
#include "kerneltypes.h"
#include "mark3cfg.h"
#include "kernel.h"
#include "scheduler.h"
#include "thread.h"
#include "threadport.h"
#include "timerlist.h"
#include "message.h"
#include "driver.h"
#include "profile.h"
#include "kernelprofile.h"
#include "tracebuffer.h"
#include "kerneldebug.h"
```

Macros

- `#define __FILE_ID__ KERNEL_CPP`

File ID used in kernel trace calls.

15.25.1 Detailed Description

Kernel initialization and startup code.

Definition in file [kernel.cpp](#).

15.26 kernel.cpp

```

00001 /*-----
00002
00003
00004 |-----|-----|-----|-----|-----|-----|
00005 | \      / | | | \      / | | | \      / | | | \      / | | | \      / | | | \      /
00006 | /      \ | | | /      \ | | | /      \ | | | /      \ | | | /      \ | | | /      \
00007 |-----|-----|-----|-----|-----|-----|
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00021 #include "kerneltypes.h"
00022 #include "mark3cfg.h"
00023
00024 #include "kernel.h"
00025 #include "scheduler.h"
00026 #include "thread.h"
00027 #include "threadport.h"
00028 #include "timerlist.h"
00029 #include "message.h"
00030 #include "driver.h"
00031 #include "profile.h"
00032 #include "kernelprofile.h"
00033 #include "tracebuffer.h"
00034 #include "kerneldebug.h"
00035
00036 K_BOOL Kernel::m_bIsStarted;
00037 K_BOOL Kernel::m_bIsPanic;
00038 panic_func_t Kernel::m_pfPanic;
00039
00040 #if KERNEL_USE_IDLE_FUNC
00041 idle_func_t Kernel::m_pfIdle;
00042 FakeThread_t Kernel::m_clIdle;
00043 #endif
00044
00045 //-----
00046 #if defined __FILE_ID__
00047     #undef __FILE_ID__
00048 #endif

```

```

00049 #define __FILE_ID__      KERNEL_CPP
00050
00051 //-----
00052 void Kernel::Init(void)
00053 {
00054     m_bIsStarted = false;
00055     m_bIsPanic = false;
00056     m_pfPanic = 0;
00057
00058 #if KERNEL_USE_IDLE_FUNC
00059     ((Thread*) &m_clIdle)->InitIdle();
00060     m_pfIdle = 0;
00061 #endif
00062
00063 #if KERNEL_USE_DEBUG
00064     TraceBuffer::Init();
00065 #endif
00066     KERNEL_TRACE( STR_MARK3_INIT );
00067
00068     // Initialize the global kernel data - scheduler, timer-scheduler, and
00069     // the global message pool.
00070     Scheduler::Init();
00071 #if KERNEL_USE_DRIVER
00072     DriverList::Init();
00073 #endif
00074 #if KERNEL_USE_TIMERS
00075     TimerScheduler::Init();
00076 #endif
00077 #if KERNEL_USE_MESSAGE
00078     GlobalMessagePool::Init();
00079 #endif
00080 #if KERNEL_USE_PROFILER
00081     Profiler::Init();
00082 #endif
00083 }
00084
00085 //-----
00086 void Kernel::Start(void)
00087 {
00088     KERNEL_TRACE( STR_THREAD_START );
00089     m_bIsStarted = true;
00090     ThreadPort::StartThreads();
00091     KERNEL_TRACE( STR_START_ERROR );
00092 }
00093
00094 //-----
00095 void Kernel::Panic(K_USHORT usCause_)
00096 {
00097     m_bIsPanic = true;
00098     if (m_pfPanic)
00099     {
00100         m_pfPanic(usCause_);
00101     }
00102     else
00103     {
00104 #if KERNEL_AWARE_SIMULATION
00105         KernelAware::ExitSimulator();
00106 #endif
00107         while(1);
00108     }
00109 }

```

15.27 /home/moslevin/Project/R1/kernel/kernelaware.cpp File Reference

[Kernel](#) aware simulation support.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "kernelaware.h"
#include "threadport.h"

```

Classes

- union [KernelAwareData_t](#)

This structure is used to communicate between the kernel and a kernel- aware host.

Variables

- volatile `K_BOOL g_bIsKernelAware` = false
Will be set to true by a kernel-aware host.
- volatile `K_UCHAR g_ucKACommand`
Kernel-aware simulator command to execute.
- `KernelAwareData_t g_stKADData`
Data structure used to communicate with host.

15.27.1 Detailed Description

Kernel aware simulation support.

Definition in file `kernelaware.cpp`.

15.27.2 Variable Documentation

15.27.2.1 volatile `K_BOOL g_bIsKernelAware` = false

Will be set to true by a kernel-aware host.

Definition at line 71 of file `kernelaware.cpp`.

15.27.2.2 `KernelAwareData_t g_stKADData`

Data structure used to communicate with host.

Definition at line 73 of file `kernelaware.cpp`.

15.28 kernelaware.cpp

```

00001  /*=====
00002
00003  _____
00004  |   /   \   |   /   \   |   /   \   |   /   \   |
00005  |  /     \  |  /     \  |  /     \  |  /     \  |
00006  | /       \ | /       \ | /       \ | /       \ |
00007  |_____|   |_____|   |_____|   |_____|   |
00008
00009  --[Mark3 Realtime Platform]-----
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information
00013  =====*/
00021  #include "kerneltypes.h"
00022  #include "mark3cfg.h"
00023  #include "kernelaware.h"
00024  #include "threadport.h"
00025
00026  #if KERNEL_AWARE_SIMULATION
00027
00028  //-----
00037  typedef union
00038  {
00039      volatile K_USHORT ausBuffer[5];
00040
00044      struct
00045      {
00046          volatile const K_CHAR *szName;
00047      } Profiler;
00052      struct
00053      {
00054          volatile K_USHORT usFile;
00055          volatile K_USHORT usLine;
00056          volatile K_USHORT usCode;
00057          volatile K_USHORT usArg1;
00058          volatile K_USHORT usArg2;

```

```

00059     } Trace;
00064     struct
00065     {
00066         volatile const K_CHAR *szString;
00067     } Print;
00068 } KernelAwareData_t;
00069
00070 //-----
00071 volatile K_BOOL      g_bIsKernelAware = false;
00072 volatile K_UCHAR     g_ucKACommand;
00073 KernelAwareData_t    g_stKADData;
00074
00075 //-----
00076
00077 void KernelAware::ProfileInit(const K_CHAR *szStr_)
00078 {
00079     CS_ENTER();
00080     g_stKADData.Profiler.szName = szStr_;
00081     g_ucKACommand = KA_COMMAND_PROFILE_INIT;
00082     CS_EXIT();
00083 }
00084
00085 //-----
00086 void KernelAware::ProfileStart(void)
00087 {
00088     g_ucKACommand = KA_COMMAND_PROFILE_START;
00089 }
00090
00091 //-----
00092 void KernelAware::ProfileStop(void)
00093 {
00094     g_ucKACommand = KA_COMMAND_PROFILE_STOP;
00095 }
00096
00097 //-----
00098 void KernelAware::ProfileReport(void)
00099 {
00100     g_ucKACommand = KA_COMMAND_PROFILE_REPORT;
00101 }
00102
00103 //-----
00104 void KernelAware::ExitSimulator(void)
00105 {
00106     g_ucKACommand = KA_COMMAND_EXIT_SIMULATOR;
00107 }
00108
00109 //-----
00110 void KernelAware::Trace( K_USHORT usFile_,
00111                         K_USHORT usLine_,
00112                         K_USHORT usCode_ )
00113 {
00114     Trace_i( usFile_, usLine_, usCode_, 0, 0, KA_COMMAND_TRACE_0 );
00115 }
00116
00117 //-----
00118 void KernelAware::Trace( K_USHORT usFile_,
00119                         K_USHORT usLine_,
00120                         K_USHORT usCode_,
00121                         K_USHORT usArg1_ )
00122 {
00123     Trace_i( usFile_, usLine_, usCode_, usArg1_, 0, KA_COMMAND_TRACE_1 );
00124 }
00125
00126 //-----
00127 void KernelAware::Trace( K_USHORT usFile_,
00128                         K_USHORT usLine_,
00129                         K_USHORT usCode_,
00130                         K_USHORT usArg1_,
00131                         K_USHORT usArg2_ )
00132 {
00133     Trace_i( usFile_, usLine_, usCode_, usArg1_, usArg2_,
00134             KA_COMMAND_TRACE_2 );
00135 }
00136
00137 //-----
00138 void KernelAware::Trace_i( K_USHORT usFile_,
00139                           K_USHORT usLine_,
00140                           K_USHORT usCode_,
00141                           K_USHORT usArg1_,
00142                           K_USHORT usArg2_,
00143                           KernelAwareCommand_t eCmd_ )
00144 {
00145     CS_ENTER();
00146     g_stKADData.Trace.usFile = usFile_;
00147     g_stKADData.Trace.usLine = usLine_;
00148     g_stKADData.Trace.usCode = usCode_;
00149     g_stKADData.Trace.usArg1 = usArg1_;

```

```

00149     g_stKADData.Trace.usArg2 = usArg2_;
00150     g_ucKACommand = eCmd_;
00151     CS_EXIT();
00152 }
00153
00154 //-----
00155 void KernelAware::Print(const K_CHAR *szStr_)
00156 {
00157     CS_ENTER();
00158     g_stKADData.Print.szString = szStr_;
00159     g_ucKACommand = KA_COMMAND_PRINT;
00160     CS_EXIT();
00161 }
00162
00163 //-----
00164 K_BOOL KernelAware::IsSimulatorAware(void)
00165 {
00166     return g_bIsKernelAware;
00167 }
00168
00169 #endif

```

15.29 /home/moslevin/Project/R1/kernel/ksemaphore.cpp File Reference

[Semaphore](#) Blocking-Object Implemenation.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "ksemaphore.h"
#include "blocking.h"
#include "kerneldebug.h"
#include "timerlist.h"

```

Macros

- `#define __FILE_ID__ SEMAPHORE_CPP`
File ID used in kernel trace calls.

Functions

- void [TimedSemaphore_Callback](#) (Thread *pclOwner_, void *pvData_)
TimedSemaphore_Callback.

15.29.1 Detailed Description

[Semaphore](#) Blocking-Object Implemenation.

Definition in file [ksemaphore.cpp](#).

15.29.2 Function Documentation

15.29.2.1 void TimedSemaphore_Callback (Thread * pclOwner_, void * pvData_)

[TimedSemaphore_Callback](#).

This function is called from the timer-expired context to trigger a timeout on this semaphore. This results in the waking of the thread that generated the semaphore pend call that was not completed in time.

Parameters

<i>pclOwner_</i>	Pointer to the thread to wake
<i>pvData_</i>	Pointer to the semaphore object that the thread is blocked on

Definition at line 50 of file [ksemaphore.cpp](#).

15.30 ksemaphore.cpp

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00022 #include "kerneltypes.h"
00023 #include "mark3cfg.h"
00024
00025 #include "ksemaphore.h"
00026 #include "blocking.h"
00027 #include "kerneldebug.h"
00028 //-----
00029 #if defined __FILE_ID__
00030 #undef __FILE_ID__
00031 #endif
00032 #define __FILE_ID__ SEMAPHORE_CPP
00033
00034 #if KERNEL_USE_SEMAPHORE
00035
00036 #if KERNEL_USE_TIMEOUTS
00037 #include "timerlist.h"
00038
00039 //-----
00050 void TimedSemaphore_Callback(Thread *pclOwner_, void *pvData_)
00051 {
00052     Semaphore *pclSemaphore = static_cast<Semaphore*>(pvData_);
00053
00054     // Indicate that the semaphore has expired on the thread
00055     pclOwner_>SetExpired(true);
00056
00057     // Wake up the thread that was blocked on this semaphore.
00058     pclSemaphore->WakeMe(pclOwner_);
00059
00060
00061     if (pclOwner_>GetCurPriority() >= Scheduler::GetCurrentThread
00062         ()->GetCurPriority())
00063     {
00064         Thread::Yield();
00065     }
00066
00067 //-----
00068 void Semaphore::WakeMe(Thread *pclChosenOne_)
00069 {
00070     // Remove from the semaphore waitlist and back to its ready list.
00071     Unblock(pclChosenOne_);
00072 }
00073
00074 #endif // KERNEL_USE_TIMEOUTS
00075
00076 //-----
00077 K_UCHAR Semaphore::WakeNext()
00078 {
00079     Thread *pclChosenOne;
00080
00081     pclChosenOne = m_clBlockList.HighestWaiter();
00082
00083     // Remove from the semaphore waitlist and back to its ready list.
00084     Unblock(pclChosenOne);
00085
00086     // Call a task switch if higher or equal priority thread
00087     if (pclChosenOne->GetCurPriority() >=
00088         Scheduler::GetCurrentThread()->GetCurPriority())
00089     {
00090         return 1;

```



```

00090     }
00091     return 0;
00092 }
00093
00094 //-----
00095 void Semaphore::Init(K_USHORT usInitVal_, K_USHORT usMaxVal_)
00096 {
00097     // Copy the paramters into the object - set the maximum value for this
00098     // semaphore to implement either binary or counting semaphores, and set
00099     // the initial count. Clear the wait list for this object.
00100     m_usValue = usInitVal_;
00101     m_usMaxValue = usMaxVal_;
00102
00103     m_clBlockList.Init();
00104 }
00105
00106 //-----
00107 K_BOOL Semaphore::Post()
00108 {
00109     KERNEL_TRACE_1( STR_SEMAPHORE_POST_1, (K_USHORT)
00110         g_pstCurrent->GetID() );
00111
00112     K_BOOL bThreadWake = 0;
00113     K_BOOL bBail = false;
00114     // Increment the semaphore count - we can mess with threads so ensure this
00115     // is in a critical section. We don't just disable the scheduler since
00116     // we want to be able to do this from within an interrupt context as well.
00117     CS_ENTER();
00118
00119     // If nothing is waiting for the semaphore
00120     if (m_clBlockList.GetHead() == NULL)
00121     {
00122         // Check so see if we've reached the maximum value in the semaphore
00123         if (m_usValue < m_usMaxValue)
00124         {
00125             // Increment the count value
00126             m_usValue++;
00127         }
00128         else
00129         {
00130             // Maximum value has been reached, bail out.
00131             bBail = true;
00132         }
00133     }
00134     else
00135     {
00136         // Otherwise, there are threads waiting for the semaphore to be
00137         // posted, so wake the next one (highest priority goes first).
00138         bThreadWake = WakeNext();
00139     }
00140
00141     CS_EXIT();
00142
00143     // If we weren't able to increment the semaphore count, fail out.
00144     if (bBail)
00145     {
00146         return false;
00147     }
00148
00149     // if bThreadWake was set, it means that a higher-priority thread was
00150     // woken. Trigger a context switch to ensure that this thread gets
00151     // to execute next.
00152     if (bThreadWake)
00153     {
00154         Thread::Yield();
00155     }
00156     return true;
00157 }
00158 //-----
00159 #if KERNEL_USE_TIMEOUTS
00160 K_BOOL Semaphore::Pend_i( K_ULONG ulWaitTimeMS_ )
00161 #else
00162 void Semaphore::Pend_i( void )
00163 #endif
00164 {
00165     KERNEL_TRACE_1( STR_SEMAPHORE_PEND_1, (K_USHORT)
00166         g_pstCurrent->GetID() );
00167
00168     #if KERNEL_USE_TIMEOUTS
00169     Timer clSemTimer;
00170     K_BOOL bUseTimer = false;
00171     #endif
00172
00173     // Once again, messing with thread data - ensure
00174     // we're doing all of these operations from within a thread-safe context.
00175     CS_ENTER();

```

```

00175
00176 // Check to see if we need to take any action based on the semaphore count
00177 if (m_usValue != 0)
00178 {
00179     // The semaphore count is non-zero, we can just decrement the count
00180     // and go along our merry way.
00181     m_usValue--;
00182 }
00183 else
00184 {
00185     // The semaphore count is zero - we need to block the current thread
00186     // and wait until the semaphore is posted from elsewhere.
00187 #if KERNEL_USE_TIMEOUTS
00188     if (ulWaitTimeMS_)
00189     {
00190         g_pstCurrent->SetExpired(false);
00191         clSemTimer.Init();
00192         clSemTimer.Start(0, ulWaitTimeMS_, TimedSemaphore_Callback, (void*)this)
00193     ;
00194         bUseTimer = true;
00195     }
00196 #endif
00197     Block(g_pstCurrent);
00198     // Switch Threads immediately
00199     Thread::Yield();
00200 }
00201
00202 CS_EXIT();
00203
00204 #if KERNEL_USE_TIMEOUTS
00205     if (bUseTimer)
00206     {
00207         clSemTimer.Stop();
00208         return (g_pstCurrent->GetExpired() == 0);
00209     }
00210     return true;
00211 #endif
00212 }
00213
00214 //-----
00215 // Redirect the untimed pend API to the timed pend, with a null timeout.
00216 void Semaphore::Pend()
00217 {
00218     #if KERNEL_USE_TIMEOUTS
00219         Pend_i(0);
00220     #else
00221         Pend_i();
00222     #endif
00223 }
00224
00225 #if KERNEL_USE_TIMEOUTS
00226 //-----
00227 K_BOOL Semaphore::Pend( K_ULONG ulWaitTimeMS_ )
00228 {
00229     return Pend_i( ulWaitTimeMS_ );
00230 }
00231 #endif
00232
00233 //-----
00234 K_USHORT Semaphore::GetCount()
00235 {
00236     K_USHORT usRet;
00237     CS_ENTER();
00238     usRet = m_usValue;
00239     CS_EXIT();
00240     return usRet;
00241 }
00242
00243 #endif

```

15.31 /home/moslevin/Project/R1/kernel/ll.cpp File Reference

Core Linked-List implementation, from which all kernel objects are derived.

```

#include "kerneltypes.h"
#include "kernel.h"
#include "ll.h"
#include "kerneldebug.h"

```



```

00074         Kernel::Panic(PANIC_LIST_UNLINK_FAILED);
00075     }
00076 #endif
00077     node_->prev->next = node_->next;
00078 }
00079 if (node_->next)
00080 {
00081     #if SAFE_UNLINK
00082         if (node_->next->prev != node_)
00083         {
00084             Kernel::Panic(PANIC_LIST_UNLINK_FAILED);
00085         }
00086     #endif
00087     node_->next->prev = node_->prev;
00088 }
00089 if (node_ == m_pstHead)
00090 {
00091     m_pstHead = node_->next;
00092 }
00093 if (node_ == m_pstTail)
00094 {
00095     m_pstTail = node_->prev;
00096 }
00097 node_->ClearNode();
00098 }
00099 }
00100
00101 //-----
00102 void CircularLinkList::Add(LinkListNode *node_)
00103 {
00104     KERNEL_ASSERT( node_ );
00105
00106     // Add a node to the end of the linked list.
00107     if (!m_pstHead)
00108     {
00109         // If the list is empty, initilize the nodes
00110         m_pstHead = node_;
00111         m_pstTail = node_;
00112
00113         m_pstHead->prev = m_pstHead;
00114         m_pstHead->next = m_pstHead;
00115         return;
00116     }
00117
00118     // Move the tail node, and assign it to the new node just passed in
00119     m_pstTail->next = node_;
00120     node_->prev = m_pstTail;
00121     node_->next = m_pstHead;
00122     m_pstTail = node_;
00123     m_pstHead->prev = node_;
00124 }
00125
00126 //-----
00127 void CircularLinkList::Remove(LinkListNode *node_)
00128 {
00129     KERNEL_ASSERT( node_ );
00130
00131     // Check to see if this is the head of the list...
00132     if ((node_ == m_pstHead) && (m_pstHead == m_pstTail))
00133     {
00134         // Clear the head and tail pointers - nothing else left.
00135         m_pstHead = NULL;
00136         m_pstTail = NULL;
00137         return;
00138     }
00139
00140     #if SAFE_UNLINK
00141         // Verify that all nodes are properly connected
00142         if ((node_->prev->next != node_) || (node_->next->prev != node_))
00143         {
00144             Kernel::Panic(PANIC_LIST_UNLINK_FAILED);
00145         }
00146     #endif
00147
00148     // This is a circularly linked list - no need to check for connection,
00149     // just remove the node.
00150     node_->next->prev = node_->prev;
00151     node_->prev->next = node_->next;
00152
00153     if (node_ == m_pstHead)
00154     {
00155         m_pstHead = m_pstHead->next;
00156     }
00157     if (node_ == m_pstTail)
00158     {
00159         m_pstTail = m_pstTail->prev;
00160     }

```

```
00161         node_-->ClearNode();
00162     }
00163
00164     //-----
00165     void CircularLinkedList::PivotForward()
00166     {
00167         if (m_pstHead)
00168         {
00169             m_pstHead = m_pstHead->next;
00170             m_pstTail = m_pstTail->next;
00171         }
00172     }
00173
00174     //-----
00175     void CircularLinkedList::PivotBackward()
00176     {
00177         if (m_pstHead)
00178         {
00179             m_pstHead = m_pstHead->prev;
00180             m_pstTail = m_pstTail->prev;
00181         }
00182     }
```

15.33 /home/moslevin/Project/R1/kernel/message.cpp File Reference

Inter-thread communications via message passing.

```
#include "kerneltypes.h"
#include "mark3cfg.h"
#include "message.h"
#include "threadport.h"
#include "kerneldebug.h"
#include "timerlist.h"
```

Macros

- `#define __FILE_ID__ MESSAGE_CPP`
File ID used in kernel trace calls.

15.33.1 Detailed Description

Inter-thread communications via message passing.

Definition in file [message.cpp](#).

15.34 message.cpp

```

00001 /*-----*
00002
00003      |_____|_____|_____|_____|_____|_____|_____|_____|
00004      |  /  \  /  \  /  \  /  \  /  \  /  \  /  \  /  \  |
00005      | /    \ /    \ /    \ /    \ /    \ /    \ /    \ |
00006      |/_    _/_    _/_    _/_    _/_    _/_    _/_    _|
00007      |_____|_____|_____|_____|_____|_____|_____|_____|
00008
00009  --[Mark3 Realtime Platform]-----*
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information
00013  -----*/
00014
00015  #include "kerneltypes.h"
00016  #include "mark3cfg.h"
00017
00018  #include "message.h"
00019  #include "threadport.h"
00020  #include "kerneldebug.h"
00021
00022

```

```

00029 //-----
00030 #if defined __FILE_ID__
00031 #undef __FILE_ID__
00032 #endif
00033 #define __FILE_ID__ MESSAGE_CPP
00034
00035
00036 #if KERNEL_USE_MESSAGE
00037
00038 #if KERNEL_USE_TIMEOUTS
00039 #include "timerlist.h"
00040 #endif
00041
00042 Message GlobalMessagePool::m_aclMessagePool[
00043     GLOBAL_MESSAGE_POOL_SIZE];
00044 DoubleLinkedList GlobalMessagePool::m_clList;
00045 //-----
00046 void GlobalMessagePool::Init()
00047 {
00048     K_UCHAR i;
00049     GlobalMessagePool::m_clList.Init();
00050     for (i = 0; i < GLOBAL_MESSAGE_POOL_SIZE; i++)
00051     {
00052         GlobalMessagePool::m_aclMessagePool[i].Init();
00053         GlobalMessagePool::m_clList.Add(&(GlobalMessagePool::m_aclMessagePool[i]));
00054     }
00055 }
00056
00057 //-----
00058 void GlobalMessagePool::Push( Message *pclMessage_ )
00059 {
00060     KERNEL_ASSERT( pclMessage_ );
00061
00062     CS_ENTER();
00063
00064     GlobalMessagePool::m_clList.Add(pclMessage_);
00065
00066     CS_EXIT();
00067 }
00068
00069 //-----
00070 Message *GlobalMessagePool::Pop()
00071 {
00072     Message *pclRet;
00073     CS_ENTER();
00074
00075     pclRet = static_cast<Message*>( GlobalMessagePool::m_clList.GetHead() );
00076     if (0 != pclRet)
00077     {
00078         GlobalMessagePool::m_clList.Remove( static_cast<LinkedListNode*>( pclRet ) );
00079     }
00080
00081     CS_EXIT();
00082     return pclRet;
00083 }
00084
00085 //-----
00086 void MessageQueue::Init()
00087 {
00088     m_clSemaphore.Init(0, GLOBAL_MESSAGE_POOL_SIZE);
00089 }
00090
00091 //-----
00092 Message *MessageQueue::Receive()
00093 {
00094     #if KERNEL_USE_TIMEOUTS
00095         return Receive_i(0);
00096     #else
00097         return Receive_i();
00098     #endif
00099 }
00100
00101 //-----
00102 #if KERNEL_USE_TIMEOUTS
00103 Message *MessageQueue::Receive( K_ULONG ulTimeWaitMS_ )
00104 {
00105     return Receive_i( ulTimeWaitMS_ );
00106 }
00107 #endif
00108
00109 //-----
00110 #if KERNEL_USE_TIMEOUTS
00111 Message *MessageQueue::Receive_i( K_ULONG ulTimeWaitMS_ )
00112 #else
00113 Message *MessageQueue::Receive_i( void )
00114 #endif

```

```

00115 {
00116     Message *pclRet;
00117
00118     // Block the current thread on the counting semaphore
00119     #if KERNEL_USE_TIMEOUTS
00120     if (!m_clSemaphore.Pend(ulTimeWaitMS_))
00121     {
00122         return NULL;
00123     }
00124     #else
00125     m_clSemaphore.Pend();
00126     #endif
00127
00128     CS_ENTER();
00129
00130     // Pop the head of the message queue and return it
00131     pclRet = static_cast<Message*>( m_clLinkList.GetHead() );
00132     m_clLinkList.Remove(static_cast<Message*>(pclRet));
00133
00134     CS_EXIT();
00135
00136     return pclRet;
00137 }
00138
00139 //-----
00140 void MessageQueue::Send( Message *pclSrc_ )
00141 {
00142     KERNEL_ASSERT( pclSrc_ );
00143
00144     CS_ENTER();
00145
00146     // Add the message to the head of the linked list
00147     m_clLinkList.Add( pclSrc_ );
00148
00149     // Post the semaphore, waking the blocking thread for the queue.
00150     m_clSemaphore.Post();
00151
00152     CS_EXIT();
00153 }
00154
00155 //-----
00156 K_USHORT MessageQueue::GetCount ()
00157 {
00158     return m_clSemaphore.GetCount ();
00159 }
00160 #endif //KERNEL_USE_MESSAGE

```

15.35 /home/moslevin/Project/R1/kernel/mutex.cpp File Reference

Mutual-exclusion object.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "blocking.h"
#include "mutex.h"
#include "kerneldebug.h"

```

Macros

- `#define __FILE_ID__ MUTEX_CPP`
File ID used in kernel trace calls.

Functions

- void `TimedMutex_Callback (Thread *pclOwner_, void *pvData_)`
TimedMutex_Callback.

15.35.1 Detailed Description

Mutual-exclusion object.

Definition in file [mutex.cpp](#).

15.35.2 Function Documentation

15.35.2.1 void TimedMutex_Callback (Thread * *pclOwner_*, void * *pvData_*)

TimedMutex_Callback.

This function is called from the timer-expired context to trigger a timeout on this mutex. This results in the waking of the thread that generated the mutex claim call that was not completed in time.

Parameters

<i>pclOwner_</i>	Pointer to the thread to wake
<i>pvData_</i>	Pointer to the mutex object that the thread is blocked on

Definition at line 48 of file [mutex.cpp](#).

15.36 mutex.cpp

```

00001 /*=====
00002
00003  _____
00004  |   /   \   |   /   \   |   /   \   |   /   \   |
00005  |  /     \  |  /     \  |  /     \  |  /     \  |
00006  | /       \ | /       \ | /       \ | /       \ |
00007  | \       / | \       / | \       / | \       / |
00008  |  \     /  |  \     /  |  \     /  |  \     /  |
00009  |   \   /   |   \   /   |   \   /   |   \   /   |
00010  |_____/_____|_____/_____|_____/_____|_____/_____|
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00020 #include "kerneltypes.h"
00021 #include "mark3cfg.h"
00022
00023 #include "blocking.h"
00024 #include "mutex.h"
00025 #include "kerneldebug.h"
00026 //-----
00027 #if defined __FILE_ID__
00028     #undef __FILE_ID__
00029 #endif
00030 #define __FILE_ID__      MUTEX_CPP
00031
00032
00033 #if KERNEL_USE_MUTEX
00034
00035 #if KERNEL_USE_TIMEOUTS
00036
00037 //-----
00048 void TimedMutex_Callback(Thread *pclOwner_, void *pvData_)
00049 {
00050     Mutex *pclMutex = static_cast<Mutex*>(pvData_);
00051
00052     // Indicate that the semaphore has expired on the thread
00053     pclOwner_>SetExpired(true);
00054
00055     // Wake up the thread that was blocked on this semaphore.
00056     pclMutex->WakeMe(pclOwner_);
00057
00058     if (pclOwner_>GetCurPriority() >= Scheduler::GetCurrentThread
00059         ()->GetCurPriority())
00059     {
00060         Thread::Yield();
00061     }
00062 }
00063
00064 //-----
00065 void Mutex::WakeMe(Thread *pclOwner_)
00066 {

```



```

00067     // Remove from the semaphore waitlist and back to its ready list.
00068     Unblock(pclOwner_);
00069 }
00070
00071 #endif
00072
00073 //-----
00074 K_UCHAR Mutex::WakeNext ()
00075 {
00076     Thread *pclChosenOne = NULL;
00077
00078     // Get the highest priority waiter thread
00079     pclChosenOne = m_clBlockList.HighestWaiter();
00080
00081     // Unblock the thread
00082     Unblock(pclChosenOne);
00083
00084     // The chosen one now owns the mutex
00085     m_pclOwner = pclChosenOne;
00086
00087     // Signal a context switch if it's a greater than or equal to the current priority
00088     if (pclChosenOne->GetCurPriority() >=
Scheduler::GetCurrentThread()->GetCurPriority())
00089     {
00090         return 1;
00091     }
00092     return 0;
00093 }
00094
00095 //-----
00096 void Mutex::Init ()
00097 {
00098     // Reset the data in the mutex
00099     m_bReady = 1;           // The mutex is free.
00100     m_ucMaxPri = 0;         // Set the maximum priority inheritance state
00101     m_pclOwner = NULL;      // Clear the mutex owner
00102     m_ucRecurse = 0;        // Reset recurse count
00103 }
00104
00105 //-----
00106 #if KERNEL_USE_TIMEOUTS
00107 K_BOOL Mutex::Claim_i(K_ULONG ulWaitTimeMS_)
00108 #else
00109 void Mutex::Claim_i(void)
00110 #endif
00111 {
00112     KERNEL_TRACE_1( STR_MUTEX_CLAIM_1, (K_USHORT)
g_pstCurrent->GetID() );
00113
00114 #if KERNEL_USE_TIMEOUTS
00115     Timer clTimer;
00116     K_BOOL bUseTimer = false;
00117 #endif
00118
00119     // Disable the scheduler while claiming the mutex - we're dealing with all
00120     // sorts of private thread data, can't have a thread switch while messing
00121     // with internal data structures.
00122     Scheduler::SetScheduler(0);
00123
00124     // Check to see if the mutex is claimed or not
00125     if (m_bReady != 0)
00126     {
00127         // Mutex isn't claimed, claim it.
00128         m_bReady = 0;
00129         m_ucRecurse = 0;
00130         m_ucMaxPri = g_pstCurrent->GetPriority();
00131         m_pclOwner = g_pstCurrent;
00132
00133         Scheduler::SetScheduler(1);
00134
00135 #if KERNEL_USE_TIMEOUTS
00136         return true;
00137 #else
00138         return;
00139 #endif
00140     }
00141
00142     // If the mutex is already claimed, check to see if this is the owner thread,
00143     // since we allow the mutex to be claimed recursively.
00144     if (g_pstCurrent == m_pclOwner)
00145     {
00146         // Ensure that we haven't exceeded the maximum recursive-lock count
00147         KERNEL_ASSERT( (m_ucRecurse < 255) );
00148         m_ucRecurse++;
00149
00150         // Increment the lock count and bail
00151         Scheduler::SetScheduler(1);

```

```

00152 #if KERNEL_USE_TIMEOUTS
00153     return true;
00154 #else
00155     return;
00156 #endif
00157 }
00158
00159 // The mutex is claimed already - we have to block now. Move the
00160 // current thread to the list of threads waiting on the mutex.
00161 #if KERNEL_USE_TIMEOUTS
00162 if (ulWaitTimeMS_)
00163 {
00164     g_pstCurrent->SetExpired(false);
00165     clTimer.Init();
00166     clTimer.Start(0, ulWaitTimeMS_, (TimerCallback_t)
TimedMutex_Callback, (void*)this);
00167     bUseTimer = true;
00168 }
00169 #endif
00170 Block(g_pstCurrent);
00171
00172 // Check if priority inheritance is necessary. We do this in order
00173 // to ensure that we don't end up with priority inversions in case
00174 // multiple threads are waiting on the same resource.
00175 if(m_ucMaxPri <= g_pstCurrent->GetPriority())
00176 {
00177     m_ucMaxPri = g_pstCurrent->GetPriority();
00178
00179     Thread *pclTemp = static_cast<Thread*>(m_clBlockList.GetHead());
00180     while(pclTemp)
00181     {
00182         pclTemp->InheritPriority(m_ucMaxPri);
00183         if(pclTemp == static_cast<Thread*>(m_clBlockList.GetTail()))
00184         {
00185             break;
00186         }
00187         pclTemp = static_cast<Thread*>(pclTemp->GetNext());
00188     }
00189     m_pclOwner->InheritPriority(m_ucMaxPri);
00190 }
00191
00192 // Done with thread data -reenable the scheduler
00193 Scheduler::SetScheduler(1);
00194
00195 // Switch threads if this thread acquired the mutex
00196 Thread::Yield();
00197
00198 #if KERNEL_USE_TIMEOUTS
00199 if (bUseTimer)
00200 {
00201     clTimer.Stop();
00202     return (g_pstCurrent->GetExpired() == 0);
00203 }
00204 return true;
00205 #endif
00206 }
00207
00208 //-----
00209 void Mutex::Claim(void)
00210 {
00211 #if KERNEL_USE_TIMEOUTS
00212     Claim_i(0);
00213 #else
00214     Claim_i();
00215 #endif
00216 }
00217
00218 //-----
00219 #if KERNEL_USE_TIMEOUTS
00220 K_BOOL Mutex::Claim(K_ULONG ulWaitTimeMS_)
00221 {
00222     return Claim_i(ulWaitTimeMS_);
00223 }
00224 #endif
00225
00226 //-----
00227 void Mutex::Release()
00228 {
00229     KERNEL_TRACE_1( STR_MUTEX_RELEASE_1, (K_USHORT)
g_pstCurrent->GetID() );
00230
00231     K_BOOL bSchedule = 0;
00232
00233     // Disable the scheduler while we deal with internal data structures.
00234     Scheduler::SetScheduler(0);
00235
00236     // This thread had better be the one that owns the mutex currently...

```

```

00237     KERNEL_ASSERT( (g_pstCurrent == m_pclOwner) );
00238
00239     // If the owner had claimed the lock multiple times, decrease the lock
00240     // count and return immediately.
00241     if (m_ucRecurse)
00242     {
00243         m_ucRecurse--;
00244         Scheduler::SetScheduler(1);
00245         return;
00246     }
00247
00248     // Restore the thread's original priority
00249     if (g_pstCurrent->GetCurPriority() != g_pstCurrent->
GetPriority())
00250     {
00251         g_pstCurrent->SetPriority(g_pstCurrent->
GetPriority());
00252
00253         // In this case, we want to reschedule
00254         bSchedule = 1;
00255     }
00256
00257     // No threads are waiting on this semaphore?
00258     if (m_clBlockList.GetHead() == NULL)
00259     {
00260         // Re-initialize the mutex to its default values
00261         m_bReady = 1;
00262         m_ucMaxPri = 0;
00263         m_pclOwner = NULL;
00264     }
00265     else
00266     {
00267         // Wake the highest priority Thread pending on the mutex
00268         if(WakeNext())
00269         {
00270             // Switch threads if it's higher or equal priority than the current thread
00271             bSchedule = 1;
00272         }
00273     }
00274
00275     // Must enable the scheduler again in order to switch threads.
00276     Scheduler::SetScheduler(1);
00277     if(bSchedule)
00278     {
00279         // Switch threads if a higher-priority thread was woken
00280         Thread::Yield();
00281     }
00282 }
00283
00284 #endif //KERNEL_USE_MUTEX

```

15.37 /home/moslevin/Project/R1/kernel/profile.cpp File Reference

Code profiling utilities.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "profile.h"
#include "kernelprofile.h"
#include "threadport.h"
#include "kerneldebug.h"

```

Macros

- `#define __FILE_ID__ PROFILE_CPP`
File ID used in kernel trace calls.

15.37.1 Detailed Description

Code profiling utilities.

15.39 /home/moslevin/Project/R1/kernel/public/atomic.h File Reference

```
#include "kerneltypes.h"
#include "mark3cfg.h"
#include "threadport.h"
```

Definition in file [atomic.h](#).

[illegible]

For instance, a semaphore Pend operation may result in a call to the Block() method with the currently-executing thread in order to make that thread wait for a semaphore Post. That operation would then invoke the Unblock() method, removing the blocking thread from the semaphore's list, and back into the appropriate thread inside the scheduler.

Care must be taken when implementing blocking objects to ensure that critical sections are used judiciously, otherwise asynchronous events like timers and interrupts could result in non-deterministic and often catastrophic behavior.

Definition in file [blocking.h](#).

15.42 blocking.h

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00047 #ifndef __BLOCKING_H__
00048 #define __BLOCKING_H__
00049
00050 #include "kerneltypes.h"
00051 #include "mark3cfg.h"
00052
00053 #include "ll.h"
00054 #include "threadlist.h"
00055 #include "thread.h"
00056
00057 #if KERNEL_USE_MUTEX || KERNEL_USE_SEMAPHORE || KERNEL_USE_EVENTFLAG
00058
00059 //-----
00065 class BlockingObject
00066 {
00067 protected:
00088     void Block(Thread *pclThread_);
00089
00101     void Unblock(Thread *pclThread_);
00102
00107     ThreadList m_clBlockList;
00108 };
00109
00110 #endif
00111
00112 #endif

```

15.43 /home/moslevin/Project/R1/kernel/public/debugtokens.h File Reference

Hex codes/translation tables used for efficient string tokenization.

Macros

- #define **BLOCKING_CPP** 0x0001 /* SUBSTITUTE="blocking.cpp" */
Source file names start at 0x0000.
- #define **BLOCKING_H** 0x1000 /* SUBSTITUTE="blocking.h" */
Header file names start at 0x1000.
- #define **STR_PANIC** 0x2000 /* SUBSTITUTE="!Panic!" */
Indexed strings start at 0x2000.

15.43.1 Detailed Description

Hex codes/translation tables used for efficient string tokenization. We use this for efficiently encoding strings used for kernel traces, debug prints, etc. The upside - this is really fast and efficient for encoding strings and data. Downside? The tools need to parse this header file in order to convert the enumerated data into actual strings, decoding them.

Definition in file [debugtokens.h](#).

15.44 debugtokens.h

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00025 #ifndef __DEBUG_TOKENS_H__
00026 #define __DEBUG_TOKENS_H__
00027 //-----
00029 #define BLOCKING_CPP      0x0001      /* SUBSTITUTE="blocking.cpp" */
00030 #define DRIVER_CPP       0x0002      /* SUBSTITUTE="driver.cpp" */
00031 #define KERNEL_CPP       0x0003      /* SUBSTITUTE="kernel.cpp" */
00032 #define LL_CPP           0x0004      /* SUBSTITUTE="ll.cpp" */
00033 #define MESSAGE_CPP      0x0005      /* SUBSTITUTE="message.cpp" */
00034 #define MUTEX_CPP        0x0006      /* SUBSTITUTE="mutex.cpp" */
00035 #define PROFILE_CPP      0x0007      /* SUBSTITUTE="profile.cpp" */
00036 #define QUANTUM_CPP      0x0008      /* SUBSTITUTE="quantum.cpp" */
00037 #define SCHEDULER_CPP    0x0009      /* SUBSTITUTE="scheduler.cpp" */
00038 #define SEMAPHORE_CPP    0x000A      /* SUBSTITUTE="semaphore.cpp" */
00039 #define THREAD_CPP       0x000B      /* SUBSTITUTE="thread.cpp" */
00040 #define THREADLIST_CPP   0x000C      /* SUBSTITUTE="threadlist.cpp" */
00041 #define TIMERLIST_CPP    0x000D      /* SUBSTITUTE="timerlist.cpp" */
00042 #define KERNELSWI_CPP    0x000E      /* SUBSTITUTE="kernelswi.cpp" */
00043 #define KERNELTIMER_CPP  0x000F      /* SUBSTITUTE="kerneltimer.cpp" */
00044 #define KPROFILE_CPP     0x0010      /* SUBSTITUTE="kernelprofile.cpp" */
00045 #define THREADPORT_CPP   0x0011      /* SUBSTITUTE="threadport.cpp" */
00046 #define TIMER_CPP        0x0012      /* SUBSTITUTE="timer.cpp" */
00047
00048 //-----
00050 #define BLOCKING_H        0x1000      /* SUBSTITUTE="blocking.h" */
00051 #define DRIVER_H         0x1001      /* SUBSTITUTE="driver.h" */
00052 #define KERNEL_H         0x1002      /* SUBSTITUTE="kernel.h" */
00053 #define KERNELTYPES_H    0x1003      /* SUBSTITUTE="kerneltypes.h" */
00054 #define LL_H             0x1004      /* SUBSTITUTE="ll.h" */
00055 #define MANUAL_H         0x1005      /* SUBSTITUTE="manual.h" */
00056 #define MARK3CFG_H       0x1006      /* SUBSTITUTE="mark3cfg.h" */
00057 #define MESSAGE_H        0x1007      /* SUBSTITUTE="message.h" */
00058 #define MUTEX_H          0x1008      /* SUBSTITUTE="mutex.h" */
00059 #define PROFILE_H        0x1009      /* SUBSTITUTE="profile.h" */
00060 #define PROFILING_RESULTS_H 0x100A   /* SUBSTITUTE="profiling_results.h" */
00061 #define QUANTUM_H        0x100B      /* SUBSTITUTE="quantum.h" */
00062 #define SCHEDULER_H      0x100C      /* SUBSTITUTE="scheduler.h" */
00063 #define SEMAPHORE_H      0x100D      /* SUBSTITUTE="ksemaphore.h" */
00064 #define THREAD_H         0x100E      /* SUBSTITUTE="thread.h" */
00065 #define THREADLIST_H     0x100F      /* SUBSTITUTE="threadlist.h" */
00066 #define TIMERLIST_H      0x1010      /* SUBSTITUTE="timerlist.h" */
00067 #define KERNELSWI_H      0x1011      /* SUBSTITUTE="kernelswi.h" */
00068 #define KERNELTIMER_H    0x1012      /* SUBSTITUTE="kerneltimer.h" */
00069 #define KPROFILE_H       0x1013      /* SUBSTITUTE="kernelprofile.h" */
00070 #define THREADPORT_H     0x1014      /* SUBSTITUTE="threadport.h" */
00071
00072 //-----
00074 #define STR_PANIC          0x2000      /* SUBSTITUTE="!Panic!" */
00075 #define STR_MARK3_INIT     0x2001      /* SUBSTITUTE="Initializing Kernel Objects" */
00076 #define STR_KERNEL_ENTER   0x2002      /* SUBSTITUTE="Starting Kernel" */
00077 #define STR_THREAD_START   0x2003      /* SUBSTITUTE="Switching to First Thread" */
00078 #define STR_START_ERROR    0x2004      /* SUBSTITUTE="Error starting kernel - function should never
return" */
00079 #define STR_THREAD_CREATE  0x2005      /* SUBSTITUTE="Creating Thread" */
00080 #define STR_STACK_SIZE_1   0x2006      /* SUBSTITUTE=" Stack Size: %1" */
00081 #define STR_PRIORITY_1     0x2007      /* SUBSTITUTE=" Priority: %1" */
00082 #define STR_THREAD_ID_1    0x2008      /* SUBSTITUTE=" Thread ID: %1" */

```



```

00083 #define STR_ENTRYPOINT_1      0x2009      /* SUBSTITUTE=" EntryPoint: %1" */
00084 #define STR_CONTEXT_SWITCH_1    0x200A      /* SUBSTITUTE="Context Switch To Thread: %1" */
00085 #define STR_IDLING              0x200B      /* SUBSTITUTE="Idling CPU" */
00086 #define STR_WAKEUP              0x200C      /* SUBSTITUTE="Waking up" */
00087 #define STR_SEMAPHORE_PEND_1     0x200D      /* SUBSTITUTE="Semaphore Pend: %1" */
00088 #define STR_SEMAPHORE_POST_1    0x200E      /* SUBSTITUTE="Semaphore Post: %1" */
00089 #define STR_MUTEX_CLAIM_1       0x200F      /* SUBSTITUTE="Mutex Claim: %1" */
00090 #define STR_MUTEX_RELEASE_1     0x2010      /* SUBSTITUTE="Mutex Release: %1" */
00091 #define STR_THREAD_BLOCK_1      0x2011      /* SUBSTITUTE="Thread %1 Blocked" */
00092 #define STR_THREAD_UNBLOCK_1    0x2012-2015 /* SUBSTITUTE="Thread %1 Unblocked" */
00093 #define STR_ASSERT_FAILED       0x2013      /* SUBSTITUTE="Assertion Failed" */
00094 #define STR_SCHEDULE_1          0x2014      /* SUBSTITUTE="Scheduler chose %1" */
00095 #define STR_THREAD_START_1      0x2015      /* SUBSTITUTE="Thread Start: %1" */
00096 #define STR_THREAD_EXIT_1       0x2016      /* SUBSTITUTE="Thread Exit: %1" */
00097
00098 //-----
00099 #define STR_UNDEFINED           0xFFFF      /* SUBSTITUTE="UNDEFINED" */
00100 #endif

```

15.45 /home/moslevin/Project/R1/kernel/public/driver.h File Reference

[Driver](#) abstraction framework.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "ll.h"

```

Classes

- class [Driver](#)
Base device-driver class used in hardware abstraction.
- class [DriverList](#)
List of [Driver](#) objects used to keep track of all device drivers in the system.

15.45.1 Detailed Description

[Driver](#) abstraction framework.

15.45.2 Intro

This is the basis of the driver framework. In the context of Mark3, drivers don't necessarily have to be based on physical hardware peripherals. They can be used to represent algorithms (such as random number generators), files, or protocol stacks. Unlike FunkOS, where driver IO is protected automatically by a mutex, we do not use this kind of protection - we leave it up to the driver implementor to do what's right in its own context. This also frees up the driver to implement all sorts of other neat stuff, like sending messages to threads associated with the driver. Drivers are implemented as character devices, with the standard array of posix-style accessor methods for reading, writing, and general driver control.

A global driver list is provided as a convenient and minimal "filesystem" structure, in which devices can be accessed by name.

15.45.3 Driver Design

A device driver needs to be able to perform the following operations: -Initialize a peripheral -Start/stop a peripheral -Handle I/O control operations -Perform various read/write operations

At the end of the day, that's pretty much all a device driver has to do, and all of the functionality that needs to be presented to the developer.

We abstract all device drivers using a base-class which implements the following methods: -Start/Open -Stop/Close -Control -Read -Write

A basic driver framework and API can thus be implemented in five function calls - that's it! You could even reduce that further by handling the initialize, start, and stop operations inside the "control" operation.

15.45.4 Driver API

In C++, we can implement this as a class to abstract these event handlers, with virtual void functions in the base class overridden by the inherited objects.

To add and remove device drivers from the global table, we use the following methods:

```
void DriverList::Add( Driver *pclDriver_ );
void DriverList::Remove( Driver *pclDriver_ );
```

`DriverList::Add()/Remove()` takes a single arguments the pointer to the object to operate on.

Once a driver has been added to the table, drivers are opened by NAME using `DriverList::FindByName("/dev/name")`. This function returns a pointer to the specified driver if successful, or to a built in /dev/null device if the path name is invalid. After a driver is open, that pointer is used for all other driver access functions.

This abstraction is incredibly useful any peripheral or service can be accessed through a consistent set of APIs, that make it easy to substitute implementations from one platform to another. Portability is ensured, the overhead is negligible, and it emphasizes the reuse of both driver and application code as separate entities.

Consider a system with drivers for I2C, SPI, and UART peripherals - under our driver framework, an application can initialize these peripherals and write a greeting to each using the same simple API functions for all drivers:

```
pclI2C = DriverList::FindByName("/dev/i2c");
pclUART = DriverList::FindByName("/dev/tty0");
pclSPI = DriverList::FindByName("/dev/spi");

pclI2C->Write(12, "Hello World!");
pclUART->Write(12, "Hello World!");
pclSPI->Write(12, "Hello World!");
```

Definition in file [driver.h](#).

15.46 driver.h

```
00001 /*=====
00002
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00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00105 #include "kerneltypes.h"
00106 #include "mark3cfg.h"
00107
00108 #include "ll.h"
00109
00110 #ifndef __DRIVER_H__
00111 #define __DRIVER_H__
00112
00113 #if KERNEL_USE_DRIVER
00114
00115 class DriverList;
00116 //-----
00121 class Driver : public LinkListNode
00122 {
00123 public:
00129     virtual void Init() = 0;
```

```

00130
00138     virtual K_UCHAR Open() = 0;
00139
00147     virtual K_UCHAR Close() = 0;
00148
00164     virtual K_USHORT Read( K_USHORT usBytes_,
00165                             K_UCHAR *pucData_) = 0;
00166
00183     virtual K_USHORT Write( K_USHORT usBytes_,
00184                             K_UCHAR *pucData_) = 0;
00185
00208     virtual K_USHORT Control( K_USHORT usEvent_,
00209                               void *pvDataIn_,
00210                               K_USHORT usSizeIn_,
00211                               void *pvDataOut_,
00212                               K_USHORT usSizeOut_ ) = 0;
00213
00222     void SetName( const K_CHAR *pcName_ ) { m_pcPath = pcName_; }
00223
00231     const K_CHAR *GetPath() { return m_pcPath; }
00232
00233 private:
00234
00236     const K_CHAR *m_pcPath;
00237 };
00238
00239 //-----
00244 class DriverList
00245 {
00246 public:
00254     static void Init();
00255
00264     static void Add( Driver *pclDriver_ ) { m_clDriverList.
Add(pclDriver_); }
00265
00274     static void Remove( Driver *pclDriver_ ) { m_clDriverList.
Remove(pclDriver_); }
00275
00282     static Driver *FindByPath( const K_CHAR *m_pcPath );
00283
00284 private:
00285
00287     static DoubleLinkedList m_clDriverList;
00288 };
00289
00290 #endif //KERNEL_USE_DRIVER
00291
00292 #endif

```

15.47 /home/moslevin/Project/R1/kernel/public/eventflag.h File Reference

Event Flag Blocking Object/IPC-Object definition.

```

#include "mark3cfg.h"
#include "kernel.h"
#include "kerneltypes.h"
#include "blocking.h"
#include "thread.h"

```

Classes

- class [EventFlag](#)

The [EventFlag](#) class is a blocking object, similar to a semaphore or mutex, commonly used for synchronizing thread execution based on events occurring within the system.

15.47.1 Detailed Description

Event Flag Blocking Object/IPC-Object definition.

Definition in file [eventflag.h](#).

15.48 eventflag.h

```

00001 /*=====
00002
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00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00019 #ifndef __EVENTFLAG_H__
00020 #define __EVENTFLAG_H__
00021
00022 #include "mark3cfg.h"
00023 #include "kernel.h"
00024 #include "kerneltypes.h"
00025 #include "blocking.h"
00026 #include "thread.h"
00027
00028 #if KERNEL_USE_EVENTFLAG
00029
00030 //-----
00046 class EventFlag : public BlockingObject
00047 {
00048 public:
00052     void Init() { m_usSetMask = 0; m_clBlockList.
        Init(); }
00053
00061     K_USHORT Wait(K_USHORT usMask_, EventFlagOperation_t eMode_);
00062
00063 #if KERNEL_USE_TIMEOUTS
00064
00072     K_USHORT Wait(K_USHORT usMask_, EventFlagOperation_t eMode_,
        K_ULONG ulTimeMS_);
00073
00081     void WakeMe(Thread *pclOwner_);
00082
00083 #endif
00084
00090     void Set(K_USHORT usMask_);
00091
00096     void Clear(K_USHORT usMask_);
00097
00102     K_USHORT GetMask();
00103
00104 private:
00105
00106 #if KERNEL_USE_TIMEOUTS
00107
00119     K_USHORT Wait_i(K_USHORT usMask_, EventFlagOperation_t eMode_
        , K_ULONG ulTimeMS_);
00120 #else
00121
00131     K_USHORT Wait_i(K_USHORT usMask_, EventFlagOperation_t eMode_
        );
00132 #endif
00133
00134     K_USHORT m_usSetMask;
00135 };
00136
00137 #endif //KERNEL_USE_EVENTFLAG
00138 #endif //__EVENTFLAG_H__
00139

```

15.49 /home/moslevin/Project/R1/kernel/public/kernel.h File Reference

[Kernel](#) initialization and startup class.

```

#include "mark3cfg.h"
#include "kerneltypes.h"
#include "paniccodes.h"
#include "thread.h"

```



```

00131 #if KERNEL_USE_IDLE_FUNC
00132     static idle_func_t m_pfIdle;
00133     static FakeThread_t m_clIdle;
00134 #endif
00135 };
00136
00137 #endif
00138

```

15.51 /home/moslevin/Project/R1/kernel/public/kernelaware.h File Reference

[Kernel](#) aware simulation support.

```

#include "kerneltypes.h"
#include "mark3cfg.h"

```

Classes

- class [KernelAware](#)
The *KernelAware* class.

Enumerations

- enum [KernelAwareCommand_t](#) {
[KA_COMMAND_IDLE](#) = 0, [KA_COMMAND_PROFILE_INIT](#), [KA_COMMAND_PROFILE_START](#), [KA_COMMAND_PROFILE_STOP](#),
[KA_COMMAND_PROFILE_REPORT](#), [KA_COMMAND_EXIT_SIMULATOR](#), [KA_COMMAND_TRACE_0](#), [KA_COMMAND_TRACE_1](#),
[KA_COMMAND_TRACE_2](#), [KA_COMMAND_PRINT](#) }

This enumeration contains a list of supported commands that can be executed to invoke a response from a kernel aware host.

15.51.1 Detailed Description

[Kernel](#) aware simulation support.

Definition in file [kernelaware.h](#).

15.51.2 Enumeration Type Documentation

15.51.2.1 enum KernelAwareCommand_t

This enumeration contains a list of supported commands that can be executed to invoke a response from a kernel aware host.

Enumerator

[KA_COMMAND_IDLE](#) Null command, does nothing.

[KA_COMMAND_PROFILE_INIT](#) Initialize a new profiling session.

[KA_COMMAND_PROFILE_START](#) Begin a profiling sample.

[KA_COMMAND_PROFILE_STOP](#) End a profiling sample.

[KA_COMMAND_PROFILE_REPORT](#) Report current profiling session.

[KA_COMMAND_EXIT_SIMULATOR](#) Terminate the host simulator.

[KA_COMMAND_TRACE_0](#) 0-argument kernel trace

KA_COMMAND_TRACE_1 1-argument kernel trace
KA_COMMAND_TRACE_2 2-argument kernel trace
KA_COMMAND_PRINT Print an arbitrary string of data.

Definition at line 33 of file [kernelaware.h](#).

15.52 kernelaware.h

```

00001 /*=====
00002
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00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00021 #ifndef __KERNEL_AWARE_H__
00022 #define __KERNEL_AWARE_H__
00023
00024 #include "kerneltypes.h"
00025 #include "mark3cfg.h"
00026
00027 #if KERNEL_AWARE_SIMULATION
00028 //-----
00033 typedef enum
00034 {
00035     KA_COMMAND_IDLE = 0,
00036     KA_COMMAND_PROFILE_INIT,
00037     KA_COMMAND_PROFILE_START,
00038     KA_COMMAND_PROFILE_STOP,
00039     KA_COMMAND_PROFILE_REPORT,
00040     KA_COMMAND_EXIT_SIMULATOR,
00041     KA_COMMAND_TRACE_0,
00042     KA_COMMAND_TRACE_1,
00043     KA_COMMAND_TRACE_2,
00044     KA_COMMAND_PRINT
00045 } KernelAwareCommand_t;
00046
00047 //-----
00065 class KernelAware
00066 {
00067 public:
00068     //-----
00079     static void ProfileInit( const K_CHAR *szStr_ );
00080
00081     //-----
00089     static void ProfileStart( void );
00090
00091     //-----
00098     static void ProfileStop( void );
00099
00100     //-----
00108     static void ProfileReport( void );
00109
00110     //-----
00118     static void ExitSimulator( void );
00119
00120     //-----
00128     static void Print( const K_CHAR *szStr_ );
00129
00130     //-----
00141     static void Trace( K_USHORT usFile_,
00142                       K_USHORT usLine_,
00143                       K_USHORT usCode_ );
00144
00145     //-----
00157     static void Trace( K_USHORT usFile_,
00158                       K_USHORT usLine_,
00159                       K_USHORT usCode_,
00160                       K_USHORT usArg1_ );
00161
00162     //-----
00175     static void Trace( K_USHORT usFile_,
00176                       K_USHORT usLine_,
00177                       K_USHORT usCode_,

```

```

00178             K_USHORT usArg1_,
00179             K_USHORT usArg2_);
00180
00181     //-----
00191     static K_BOOL IsSimulatorAware(void);
00192
00193 private:
00194
00195     //-----
00209     static void Trace_i( K_USHORT usFile_,
00210                         K_USHORT usLine_,
00211                         K_USHORT usCode_,
00212                         K_USHORT usArg1_,
00213                         K_USHORT usArg2_,
00214                         KernelAwareCommand_t eCmd_);
00215 };
00216
00217 #endif
00218
00219 #endif

```

15.53 /home/moslevin/Project/R1/kernel/public/kerneldebug.h File Reference

Macros and functions used for assertions, kernel traces, etc.

```

#include "debugtokens.h"
#include "mark3cfg.h"
#include "tracebuffer.h"
#include "kernelaware.h"
#include "paniccodes.h"
#include "kernel.h"

```

Macros

- `#define __FILE_ID__ 0`
Null ID.
- `#define KERNEL_TRACE(x)`
Null *Kernel* Trace Macro.
- `#define KERNEL_TRACE_1(x, arg1)`
Null *Kernel* Trace Macro.
- `#define KERNEL_TRACE_2(x, arg1, arg2)`
Null *Kernel* Trace Macro.
- `#define KERNEL_ASSERT(x)`
Null *Kernel* Assert Macro.

15.53.1 Detailed Description

Macros and functions used for assertions, kernel traces, etc.

Definition in file [kerneldebug.h](#).

15.54 kerneldebug.h

```

00001 /*-----
00002
00003
00004
00005
00006
00007
00008

```



```

00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00020 #ifndef __KERNEL_DEBUG_H__
00021 #define __KERNEL_DEBUG_H__
00022
00023 #include "debugtokens.h"
00024 #include "mark3cfg.h"
00025 #include "tracebuffer.h"
00026 #include "kernelaware.h"
00027 #include "paniccodes.h"
00028 #include "kernel.h"
00029 //-----
00030 #if (KERNEL_USE_DEBUG && !KERNEL_AWARE_SIMULATION)
00031
00032 //-----
00033 #define __FILE_ID__          STR_UNDEFINED
00034
00035 //-----
00036 #define KERNEL_TRACE( x ) \
00037 { \
00038     K_USHORT ausMsg__[5]; \
00039     ausMsg__[0] = 0xACDC; \
00040     ausMsg__[1] = __FILE_ID__; \
00041     ausMsg__[2] = __LINE__; \
00042     ausMsg__[3] = TraceBuffer::Increment(); \
00043     ausMsg__[4] = (K_USHORT)(x); \
00044     TraceBuffer::Write(ausMsg__, 5); \
00045 };
00046
00047 //-----
00048 #define KERNEL_TRACE_1( x, arg1 ) \
00049 { \
00050     K_USHORT ausMsg__[6]; \
00051     ausMsg__[0] = 0xACDC; \
00052     ausMsg__[1] = __FILE_ID__; \
00053     ausMsg__[2] = __LINE__; \
00054     ausMsg__[3] = TraceBuffer::Increment(); \
00055     ausMsg__[4] = (K_USHORT)(x); \
00056     ausMsg__[5] = arg1; \
00057     TraceBuffer::Write(ausMsg__, 6); \
00058 }
00059
00060 //-----
00061 #define KERNEL_TRACE_2( x, arg1, arg2 ) \
00062 { \
00063     K_USHORT ausMsg__[7]; \
00064     ausMsg__[0] = 0xACDC; \
00065     ausMsg__[1] = __FILE_ID__; \
00066     ausMsg__[2] = __LINE__; \
00067     ausMsg__[3] = TraceBuffer::Increment(); \
00068     ausMsg__[4] = (K_USHORT)(x); \
00069     ausMsg__[5] = arg1; \
00070     ausMsg__[6] = arg2; \
00071     TraceBuffer::Write(ausMsg__, 7); \
00072 }
00073
00074 //-----
00075 #define KERNEL_ASSERT( x ) \
00076 { \
00077     if( ( x ) == false ) \
00078     { \
00079         K_USHORT ausMsg__[5]; \
00080         ausMsg__[0] = 0xACDC; \
00081         ausMsg__[1] = __FILE_ID__; \
00082         ausMsg__[2] = __LINE__; \
00083         ausMsg__[3] = TraceBuffer::Increment(); \
00084         ausMsg__[4] = STR_ASSERT_FAILED; \
00085         TraceBuffer::Write(ausMsg__, 5); \
00086         Kernel::Panic(PANIC_ASSERT_FAILED); \
00087     } \
00088 }
00089
00090 #elif (KERNEL_USE_DEBUG && KERNEL_AWARE_SIMULATION)
00091 //-----
00092 #define __FILE_ID__          STR_UNDEFINED
00093
00094 //-----
00095 #define KERNEL_TRACE( x ) \
00096 { \
00097     KernelAware::Trace( __FILE_ID__, __LINE__, x ); \
00098 };
00099
00100 //-----
00101 #define KERNEL_TRACE_1( x, arg1 ) \

```

```

00102 { \
00103     KernelAware::Trace( __FILE_ID__, __LINE__, x, arg1 ); \
00104 }
00105
00106 //-----
00107 #define KERNEL_TRACE_2( x, arg1, arg2 ) \
00108 { \
00109     KernelAware::Trace( __FILE_ID__, __LINE__, x, arg1, arg2 ); \
00110 }
00111
00112 //-----
00113 #define KERNEL_ASSERT( x ) \
00114 { \
00115     if( ( x ) == false ) \
00116     { \
00117         KernelAware::Trace( __FILE_ID__, __LINE__, STR_ASSERT_FAILED ); \
00118         Kernel::Panic( PANIC_ASSERT_FAILED ); \
00119     } \
00120 }
00121
00122 #else
00123 //-----
00124 // Note -- when kernel-debugging is disabled, we still have to define the
00125 // macros to ensure that the expressions compile (albeit, by elimination
00126 // during pre-processing).
00127 //-----
00128 #define __FILE_ID__ 0
00129 //-----
00130 #define KERNEL_TRACE( x )
00131 //-----
00132 #define KERNEL_TRACE_1( x, arg1 )
00133 //-----
00134 #define KERNEL_TRACE_2( x, arg1, arg2 )
00135 //-----
00136 #define KERNEL_ASSERT( x )
00137
00138 #endif // KERNEL_USE_DEBUG
00139
00140 #endif

```

15.55 /home/moslevin/Project/R1/kernel/public/kerneltypes.h File Reference

Basic data type primitives used throughout the OS.

```
#include <stdint.h>
```

Macros

- #define **K_BOOL** uint8_t
Basic boolean data type (true = 1, false = 0)
- #define **K_CHAR** char
The 8-bit signed integer type used by Mark3.
- #define **K_UCHAR** uint8_t
The 8-bit unsigned integer type used by Mark3.
- #define **K_USHORT** uint16_t
The 16-bit unsigned integer type used by Mark3.
- #define **K_SHORT** int16_t
The 16-bit signed integer type used by Mark3.
- #define **K_ULONG** uint32_t
The 32-bit unsigned integer type used by Mark3.
- #define **K_LONG** int32_t
The 32-bit signed integer type used by Mark3.
- #define **K_ADDR** uint32_t
Primitive datatype representing address-size.
- #define **K_WORD** uint32_t
Primitive datatype representing a data word.

Typedefs

- `typedef void(* panic_func_t)(K_USHORT usPanicCode_)`
Function pointer type used to implement kernel-panic handlers.

Enumerations

- enum `EventFlagOperation_t` {
`EVENT_FLAG_ALL`, `EVENT_FLAG_ANY`, `EVENT_FLAG_ALL_CLEAR`, `EVENT_FLAG_ANY_CLEAR`,
`EVENT_FLAG_MODES`, `EVENT_FLAG_PENDING_UNBLOCK` }
- This enumeration describes the different operations supported by the event flag blocking object.*

15.55.1 Detailed Description

Basic data type primitives used throughout the OS.

Definition in file [kerneltypes.h](#).

15.55.2 Enumeration Type Documentation

15.55.2.1 enum EventFlagOperation_t

This enumeration describes the different operations supported by the event flag blocking object.

Enumerator

EVENT_FLAG_ALL Block until all bits in the specified bitmask are set.

EVENT_FLAG_ANY Block until any bits in the specified bitmask are set.

EVENT_FLAG_ALL_CLEAR Block until all bits in the specified bitmask are cleared.

EVENT_FLAG_ANY_CLEAR Block until any bits in the specified bitmask are cleared.

EVENT_FLAG_MODES Count of event-flag modes. Not used by user

EVENT_FLAG_PENDING_UNBLOCK Special code. Not used by user

Definition at line 56 of file kerneltypes.h.

15.56 kerneltypes.h

```
00001 /*=====*/
00002
00003 |_____|_____|_____|_____|_____|_____|_____|_____|_____|_____|
00004 | \   /| \   /| \   /| \   /| \   /| \   /| \   /| \   /|
00005 |  _ _/  _ _/  _ _/  _ _/  _ _/  _ _/  _ _/  _ _/  _ _/  _ _/
00006 |_/___/\_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_/_\
00007 |_____||_____||_____||_____||_____||_____||_____||_____||
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00019 #include <stdint.h>
00020
00021 #ifndef __KERNELTYPES_H__
00022 #define __KERNELTYPES_H__
00023
00024 //-----
00025 #if defined(bool)
00026     #define K_BOOL                bool
00027 #else
00028     #define K_BOOL                uint8_t
00029 #endif
```



```

00022 #ifndef __KSEMAPHORE_H__
00023 #define __KSEMAPHORE_H__
00024
00025 #include "kerneltypes.h"
00026 #include "mark3cfg.h"
00027
00028 #include "blocking.h"
00029 #include "threadlist.h"
00030
00031 #if KERNEL_USE_SEMAPHORE
00032
00033 //-----
00037 class Semaphore : public BlockingObject
00038 {
00039 public:
00049     void Init(K_USHORT usInitVal_, K_USHORT usMaxVal_);
00050
00059     K_BOOL Post();
00060
00067     void Pend();
00068
00069
00081     K_USHORT GetCount();
00082
00083 #if KERNEL_USE_TIMEOUTS
00084
00095     K_BOOL Pend( K_ULONG ulWaitTimeMS_);
00096
00107     void WakeMe(Thread *pclChosenOne_);
00108 #endif
00109
00110 private:
00111
00117     K_UCHAR WakeNext();
00118
00119 #if KERNEL_USE_TIMEOUTS
00120
00128     K_BOOL Pend_i( K_ULONG ulWaitTimeMS_ );
00129 #else
00130
00136     void Pend_i( void );
00137 #endif
00138
00139     K_USHORT m_usValue;
00140     K_USHORT m_usMaxValue;
00141
00142
00143 };
00144
00145 #endif //KERNEL_USE_SEMAPHORE
00146
00147 #endif

```

15.59 /home/moslevin/Project/R1/kernel/public/ll.h File Reference

Core linked-list declarations, used by all kernel list types.

```
#include "kerneltypes.h"
```

Classes

- class [LinkListNode](#)
Basic linked-list node data structure.
- class [LinkList](#)
Abstract-data-type from which all other linked-lists are derived.
- class [DoubleLinkList](#)
Doubly-linked-list data type, inherited from the base [LinkList](#) type.
- class [CircularLinkList](#)
Circular-linked-list data type, inherited from the base [LinkList](#) type.

15.59.1 Detailed Description

Core linked-list declarations, used by all kernel list types. At the heart of RTOS data structures are linked lists. Having a robust and efficient set of linked-list types that we can use as a foundation for building the rest of our kernel types allows us to keep our RTOS code efficient and logically-separated.

So what data types rely on these linked-list classes?

-Threads -ThreadLists -The [Scheduler](#) -Timers, -The [Timer Scheduler](#) -Blocking objects (Semaphores, Mutexes, etc...)

Pretty much everything in the kernel uses these linked lists. By having objects inherit from the base linked-list node type, we're able to leverage the double and circular linked-list classes to manager virtually every object type in the system without duplicating code. These functions are very efficient as well, allowing for very deterministic behavior in our code.

Definition in file [ll.h](#).

15.60 ll.h

```

00001  /*=====
00002
00003  _____
00004  |   \   /   |   \   /   |   \   /   |   \   /   |   \   /   |   \   /   |
00005  |   /   \   |   /   \   |   /   \   |   /   \   |   /   \   |   /   \   |
00006  |_____/_____|_____/_____|_____/_____|_____/_____|_____/_____|_____/_____|
00007
00008
00009  --[Mark3 Realtime Platform]-----
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information
00013  =====*/
00043  #ifndef __LL_H__
00044  #define __LL_H__
00045
00046  #include "kerneltypes.h"
00047
00048  //-----
00049  #ifndef NULL
00050  #define NULL          (0)
00051  #endif
00052
00053  //-----
00059  class LinkList;
00060  class DoubleLinkList;
00061  class CircularLinkList;
00062
00063  //-----
00068  class LinkListNode
00069  {
00070  protected:
00071
00072      LinkListNode *next;
00073      LinkListNode *prev;
00074
00075      LinkListNode() { }
00076
00082      void ClearNode();
00083
00084  public:
00092      LinkListNode *GetNext(void) { return next; }
00093
00101      LinkListNode *GetPrev(void) { return prev; }
00102
00103      friend class LinkList;
00104      friend class DoubleLinkList;
00105      friend class CircularLinkList;
00106  };
00107
00108  //-----
00112  class LinkList
00113  {
00114  protected:
00115      LinkListNode *m_pstHead;
00116      LinkListNode *m_pstTail;
00117
00118  public:

```



```

#include "mark3cfg.h"
#include "kerneltypes.h"
#include "threadport.h"
#include "kernelswi.h"
#include "kerneltimer.h"
#include "kernelprofile.h"
#include "kernel.h"
#include "thread.h"
#include "timerlist.h"
#include "ksemaphore.h"
#include "mutex.h"
#include "eventflag.h"
#include "message.h"
#include "atomic.h"
#include "driver.h"
#include "kernelaware.h"
#include "profile.h"

```

15.63.1 Detailed Description

Single include file given to users of the Mark3 [Kernel](#) API.

Definition in file [mark3.h](#).

15.64 mark3.h

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00021 #ifndef __MARK3_H__
00022 #define __MARK3_H__
00023
00024 #include "mark3cfg.h"
00025 #include "kerneltypes.h"
00026
00027 #include "threadport.h"
00028 #include "kernelswi.h"
00029 #include "kerneltimer.h"
00030 #include "kernelprofile.h"
00031
00032 #include "kernel.h"
00033 #include "thread.h"
00034 #include "timerlist.h"
00035
00036 #include "ksemaphore.h"
00037 #include "mutex.h"
00038 #include "eventflag.h"
00039 #include "message.h"
00040
00041 #include "atomic.h"
00042 #include "driver.h"
00043
00044 #include "kernelaware.h"
00045
00046 #include "profile.h"
00047 #endif

```


15.65 /home/moslevin/Project/R1/kernel/public/mark3cfg.h File Reference

Mark3 [Kernel](#) Configuration.

Macros

- #define [KERNEL_USE_TIMERS](#) (1)
The following options is related to all kernel time-tracking.
- #define [KERNEL_TIMERS_TICKLESS](#) (1)
If you've opted to use the kernel timers module, you have an option as to which timer implementation to use: Tick-based or Tick-less.
- #define [KERNEL_USE_TIMEOUTS](#) (1)
By default, if you opt to enable kernel timers, you also get timeout- enabled versions of the blocking object APIs along with it.
- #define [KERNEL_USE_QUANTUM](#) (1)
Do you want to enable time quanta? This is useful when you want to have tasks in the same priority group share time in a controlled way.
- #define [THREAD_QUANTUM_DEFAULT](#) (4)
This value defines the default thread quantum when [KERNEL_USE_QUANTUM](#) is enabled.
- #define [KERNEL_USE_SEMAPHORE](#) (1)
Do you want the ability to use counting/binary semaphores for thread synchronization? Enabling this features provides fully-blocking semaphores and enables all API functions declared in [semaphore.h](#).
- #define [KERNEL_USE_MUTEX](#) (1)
Do you want the ability to use mutual exclusion semaphores (mutex) for resource/block protection? Enabling this feature provides mutexes, with priority inheritance, as declared in [mutex.h](#).
- #define [KERNEL_USE_EVENTFLAG](#) (1)
Provides additional event-flag based blocking.
- #define [KERNEL_USE_MESSAGE](#) (1)
Enable inter-thread messaging using message queues.
- #define [GLOBAL_MESSAGE_POOL_SIZE](#) (8)
If Messages are enabled, define the size of the default kernel message pool.
- #define [KERNEL_USE_SLEEP](#) (1)
Do you want to be able to set threads to sleep for a specified time? This enables the [Thread::Sleep\(\)](#) API.
- #define [KERNEL_USE_DRIVER](#) (1)
Enabling device drivers provides a posix-like filesystem interface for peripheral device drivers.
- #define [KERNEL_USE_THREADNAME](#) (0)
Provide [Thread](#) method to allow the user to set a name for each thread in the system.
- #define [KERNEL_USE_DYNAMIC_THREADS](#) (1)
Provide extra [Thread](#) methods to allow the application to create (and more importantly destroy) threads at runtime.
- #define [KERNEL_USE_PROFILER](#) (1)
Provides extra classes for profiling the performance of code.
- #define [KERNEL_USE_DEBUG](#) (0)
Provides extra logic for kernel debugging, and instruments the kernel with extra asserts, and kernel trace functionality.
- #define [KERNEL_USE_ATOMIC](#) (0)
Provides support for atomic operations, including addition, subtraction, set, and test-and-set.
- #define [SAFE_UNLINK](#) (0)
"Safe unlinking" performs extra checks on data to make sure that there are no consistencies when performing operations on linked lists.
- #define [KERNEL_AWARE_SIMULATION](#) (1)
Include support for kernel-aware simulation.
- #define [KERNEL_USE_IDLE_FUNC](#) (1)
Enabling this feature removes the necessity for the user to dedicate a complete thread for idle functionality.

15.65.1 Detailed Description

Mark3 [Kernel](#) Configuration. This file is used to configure the kernel for your specific application in order to provide the optimal set of features for a given use case.

Since you only pay the price (code space/RAM) for the features you use, you can usually find a sweet spot between features and resource usage by picking and choosing features a-la-carte. This config file is written in an "interactive" way, in order to minimize confusion about what each option provides, and to make dependencies obvious.

Definition in file [mark3cfg.h](#).

15.65.2 Macro Definition Documentation

15.65.2.1 `#define GLOBAL_MESSAGE_POOL_SIZE (8)`

If Messages are enabled, define the size of the default kernel message pool.

Messages can be manually added to the message pool, but this mechanism is more convenient and automatic. All message queues share their message objects from this global pool to maximize efficiency and simplify data management.

Definition at line 144 of file [mark3cfg.h](#).

15.65.2.2 `#define KERNEL_AWARE_SIMULATION (1)`

Include support for kernel-aware simulation.

Enabling this feature adds advanced profiling, trace, and environment-aware debugging and diagnostic functionality when Mark3-based applications are run on the flavr AVR simulator.

Definition at line 210 of file [mark3cfg.h](#).

15.65.2.3 `#define KERNEL_TIMERS_TICKLESS (1)`

If you've opted to use the kernel timers module, you have an option as to which timer implementation to use: Tick-based or Tick-less.

Tick-based timers provide a "traditional" RTOS timer implementation based on a fixed-frequency timer interrupt. While this provides very accurate, reliable timing, it also means that the CPU is being interrupted far more often than may be necessary (as not all timer ticks result in "real work" being done).

Tick-less timers still rely on a hardware timer interrupt, but uses a dynamic expiry interval to ensure that the interrupt is only called when the next timer expires. This increases the complexity of the timer interrupt handler, but reduces the number and frequency.

Note that the CPU port ([kerneltimer.cpp](#)) must be implemented for the particular timer variant desired.

Definition at line 62 of file [mark3cfg.h](#).

15.65.2.4 `#define KERNEL_USE_ATOMIC (0)`

Provides support for atomic operations, including addition, subtraction, set, and test-and-set.

Add/Sub/Set contain 8, 16, and 32-bit variants.

Definition at line 194 of file [mark3cfg.h](#).

15.65.2.5 `#define KERNEL_USE_DYNAMIC_THREADS (1)`

Provide extra [Thread](#) methods to allow the application to create (and more importantly destroy) threads at runtime.

Useful for designs implementing worker threads, or threads that can be restarted after encountering error conditions.

Definition at line 176 of file [mark3cfg.h](#).

15.65.2.6 #define KERNEL_USE_EVENTFLAG (1)

Provides additional event-flag based blocking.

This relies on an additional per-thread flag-mask to be allocated, which adds 2 bytes to the size of each thread object.

Definition at line 123 of file [mark3cfg.h](#).

15.65.2.7 #define KERNEL_USE_IDLE_FUNC (1)

Enabling this feature removes the necessity for the user to dedicate a complete thread for idle functionality.

This saves a full thread stack, but also requires a bit extra static data. This also adds a slight overhead to the context switch and scheduler, as a special case has to be taken into account.

Definition at line 219 of file [mark3cfg.h](#).

15.65.2.8 #define KERNEL_USE_MESSAGE (1)

Enable inter-thread messaging using message queues.

This is the preferred mechanism for IPC for serious multi-threaded communications; generally anywhere a semaphore or event-flag is insufficient.

Definition at line 131 of file [mark3cfg.h](#).

15.65.2.9 #define KERNEL_USE_PROFILER (1)

Provides extra classes for profiling the performance of code.

Useful for debugging and development, but uses an additional hardware timer.

Definition at line 182 of file [mark3cfg.h](#).

15.65.2.10 #define KERNEL_USE_QUANTUM (1)

Do you want to enable time quanta? This is useful when you want to have tasks in the same priority group share time in a controlled way.

This allows equal tasks to use unequal amounts of the CPU, which is a great way to set up CPU budgets per thread in a round-robin scheduling system. If enabled, you can specify a number of ticks that serves as the default time period (quantum). Unless otherwise specified, every thread in a priority will get the default quantum.

Definition at line 92 of file [mark3cfg.h](#).

15.65.2.11 #define KERNEL_USE_SEMAPHORE (1)

Do you want the ability to use counting/binary semaphores for thread synchronization? Enabling this features provides fully-blocking semaphores and enables all API functions declared in semaphore.h.

If you have to pick one blocking mechanism, this is the one to choose.

Definition at line 109 of file [mark3cfg.h](#).

15.65.2.12 #define KERNEL_USE_THREADNAME (0)

Provide [Thread](#) method to allow the user to set a name for each thread in the system.

Adds a const K_CHAR* pointer to the size of the thread object.

Definition at line 168 of file [mark3cfg.h](#).

15.65.2.13 #define KERNEL_USE_TIMEOUTS (1)

By default, if you opt to enable kernel timers, you also get timeout- enabled versions of the blocking object APIs along with it.

This support comes at a small cost to code size, but a slightly larger cost to realtime performance - as checking for the use of timers in the underlying internal code costs some cycles.

As a result, the option is given to the user here to manually disable these timeout-based APIs if desired by the user for performance and code-size reasons.

Definition at line 77 of file [mark3cfg.h](#).

15.65.2.14 #define KERNEL_USE_TIMERS (1)

The following options is related to all kernel time-tracking.

-timers provide a way for events to be periodically triggered in a lightweight manner. These can be periodic, or one-shot.

-Thread [Quantum](#) (used for round-robin scheduling) is dependent on this module, as is [Thread](#) Sleep functionality.

Definition at line 41 of file [mark3cfg.h](#).

15.65.2.15 #define SAFE_UNLINK (0)

"Safe unlinking" performs extra checks on data to make sure that there are no consistencies when performing operations on linked lists.

This goes beyond pointer checks, adding a layer of structural and metadata validation to help detect system corruption early.

Definition at line 202 of file [mark3cfg.h](#).

15.65.2.16 #define THREAD_QUANTUM_DEFAULT (4)

This value defines the default thread quantum when KERNEL_USE_QUANTUM is enabled.

The thread quantum value is in milliseconds

Definition at line 101 of file [mark3cfg.h](#).

15.66 mark3cfg.h

```

00001  /*=====
00002
00003  |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |
00004  | / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ /
00005  | / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ /
00006  | / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ /
00007  | / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ /
00008  | / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ / \ /
00009  --[Mark3 Realtime Platform]-----
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information

```

```

00013 =====*/
00029 #ifndef __MARK3CFG_H__
00030 #define __MARK3CFG_H__
00031
00041 #define KERNEL_USE_TIMERS (1)
00042
00061 #if KERNEL_USE_TIMERS
00062     #define KERNEL_TIMERS_TICKLESS (1)
00063 #endif
00064
00076 #if KERNEL_USE_TIMERS
00077     #define KERNEL_USE_TIMEOUTS (1)
00078 #else
00079     #define KERNEL_USE_TIMEOUTS (0)
00080 #endif
00081
00091 #if KERNEL_USE_TIMERS
00092     #define KERNEL_USE_QUANTUM (1)
00093 #else
00094     #define KERNEL_USE_QUANTUM (0)
00095 #endif
00096
00101 #define THREAD_QUANTUM_DEFAULT (4)
00102
00109 #define KERNEL_USE_SEMAPHORE (1)
00110
00116 #define KERNEL_USE_MUTEX (1)
00117
00123 #define KERNEL_USE_EVENTFLAG (1)
00124
00130 #if KERNEL_USE_SEMAPHORE
00131     #define KERNEL_USE_MESSAGE (1)
00132 #else
00133     #define KERNEL_USE_MESSAGE (0)
00134 #endif
00135
00143 #if KERNEL_USE_MESSAGE
00144     #define GLOBAL_MESSAGE_POOL_SIZE (8)
00145 #endif
00146
00151 #if KERNEL_USE_TIMERS && KERNEL_USE_SEMAPHORE
00152     #define KERNEL_USE_SLEEP (1)
00153 #else
00154     #define KERNEL_USE_SLEEP (0)
00155 #endif
00156
00161 #define KERNEL_USE_DRIVER (1)
00162
00168 #define KERNEL_USE_THREADNAME (0)
00169
00176 #define KERNEL_USE_DYNAMIC_THREADS (1)
00177
00182 #define KERNEL_USE_PROFILER (1)
00183
00188 #define KERNEL_USE_DEBUG (0)
00189
00194 #define KERNEL_USE_ATOMIC (0)
00195
00202 #define SAFE_UNLINK (0)
00203
00210 #define KERNEL_AWARE_SIMULATION (1)
00211
00219 #define KERNEL_USE_IDLE_FUNC (1)
00220 #endif

```

15.67 /home/moslevin/Project/R1/kernel/public/message.h File Reference

Inter-thread communication via message-passing.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "ll.h"
#include "ksemaphore.h"
#include "timerlist.h"

```

Classes

- class [Message](#)
Class to provide message-based IPC services in the kernel.
- class [GlobalMessagePool](#)
Implements a list of message objects shared between all threads.
- class [MessageQueue](#)
List of messages, used as the channel for sending and receiving messages between threads.

15.67.1 Detailed Description

Inter-thread communication via message-passing. Embedded systems guru Jack Ganssle once said that without a robust form of interprocess communications (IPC), an RTOS is just a toy. Mark3 implements a form of IPC to provide safe and flexible messaging between threads.

Using kernel-managed IPC offers significant benefits over other forms of data sharing (i.e. Global variables) in that it avoids synchronization issues and race conditions common to the practice. Using IPC also enforces a more disciplined coding style that keeps threads decoupled from one another and minimizes global data, preventing careless and hard-to-debug errors.

15.67.2 Using Messages, Queues, and the Global Message Pool

```
// Declare a message queue shared between two threads
MessageQueue my_queue;

int main()
{
    ...
    // Initialize the message queue
    my_queue.init();
    ...
}

void Thread1()
{
    // Example TX thread - sends a message every 10ms
    while(1)
    {
        // Grab a message from the global message pool
        Message *tx_message = GlobalMessagePool::Pop();

        // Set the message data/parameters
        tx_message->SetCode( 1234 );
        tx_message->SetData( NULL );

        // Send the message on the queue.
        my_queue.Send( tx_message );
        Thread::Sleep(10);
    }
}

void Thread2()
{
    while()
    {
        // Blocking receive - wait until we have messages to process
        Message *rx_message = my_queue.Recv();

        // Do something with the message data...

        // Return back into the pool when done
        GlobalMessagePool::Push( rx_message );
    }
}
```

Definition in file [message.h](#).

15.68 message.h

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00080 #ifndef __MESSAGE_H__
00081 #define __MESSAGE_H__
00082
00083 #include "kerneltypes.h"
00084 #include "mark3cfg.h"
00085
00086 #include "ll.h"
00087 #include "ksemaphore.h"
00088
00089 #if KERNEL_USE_MESSAGE
00090
00091 #if KERNEL_USE_TIMEOUTS
00092     #include "timerlist.h"
00093 #endif
00094
00095 //-----
00099 class Message : public LinkListNode
00100 {
00101 public:
00107     void Init() { ClearNode(); m_pvData = NULL; m_usCode = 0; }
00108
00116     void SetData( void *pvData_ ) { m_pvData = pvData_; }
00117
00125     void *GetData() { return m_pvData; }
00126
00134     void SetCode( K_USHORT usCode_ ) { m_usCode = usCode_; }
00135
00143     K_USHORT GetCode() { return m_usCode; }
00144 private:
00145
00147     void *m_pvData;
00148
00150     K_USHORT m_usCode;
00151 };
00152
00153 //-----
00157 class GlobalMessagePool
00158 {
00159 public:
00165     static void Init();
00166
00176     static void Push( Message *pclMessage_ );
00177
00186     static Message *Pop();
00187
00188 private:
00190     static Message m_aclMessagePool[
GLOBAL_MESSAGE_POOL_SIZE];
00191
00193     static DoubleLinkedList m_clList;
00194 };
00195
00196 //-----
00201 class MessageQueue
00202 {
00203 public:
00209     void Init();
00210
00219     Message *Receive();
00220
00221 #if KERNEL_USE_TIMEOUTS
00222
00236     Message *Receive( K_ULONG ulTimeWaitMS_ );
00237 #endif
00238
00247     void Send( Message *pclSrc_ );
00248
00249
00257     K_USHORT GetCount();
00258 private:
00259

```

```

00260 #if KERNEL_USE_TIMEOUTS
00261
00270     Message *Receive_i( K_ULONG ulTimeWaitMS_ );
00271 #else
00272
00279     Message *Receive_i( void );
00280 #endif
00281
00283     Semaphore m_clSemaphore;
00284
00286     DoubleLinkedList m_clLinkList;
00287 };
00288
00289 #endif //KERNEL_USE_MESSAGE
00290
00291 #endif

```

15.69 /home/moslevin/Project/R1/kernel/public/mutex.h File Reference

Mutual exclusion class declaration.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "blocking.h"
#include "timerlist.h"

```

Classes

- class [Mutex](#)

Mutual-exclusion locks, based on [BlockingObject](#).

15.69.1 Detailed Description

Mutual exclusion class declaration. Resource locks are implemented using mutual exclusion semaphores ([Mutex_t](#)). Protected blocks can be placed around any resource that may only be accessed by one thread at a time. If additional threads attempt to access the protected resource, they will be placed in a wait queue until the resource becomes available. When the resource becomes available, the thread with the highest original priority claims the resource and is activated. Priority inheritance is included in the implementation to prevent priority inversion. Always ensure that you claim and release your mutex objects consistently, otherwise you may end up with a deadlock scenario that's hard to debug.

15.69.2 Initializing

Initializing a mutex object by calling:

```
clMutex.Init();
```

15.69.3 Resource protection example

```

clMutex.Claim();
...
<resource protected block>
...
clMutex.Release();

```

Definition in file [mutex.h](#).

15.70 mutex.h

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00050 #ifndef __MUTEX_H_
00051 #define __MUTEX_H_
00052
00053 #include "kerneltypes.h"
00054 #include "mark3cfg.h"
00055
00056 #include "blocking.h"
00057
00058 #if KERNEL_USE_MUTEX
00059
00060 #if KERNEL_USE_TIMEOUTS
00061 #include "timerlist.h"
00062 #endif
00063
00064 //-----
00068 class Mutex : public BlockingObject
00069 {
00070 public:
00077     void Init();
00078
00085     void Claim();
00086
00087 #if KERNEL_USE_TIMEOUTS
00088
00097     K_BOOL Claim(K_ULONG ulWaitTimeMS_);
00098
00111     void WakeMe( Thread *pclOwner_ );
00112
00113 #endif
00114
00121     void Release();
00122
00123 private:
00124
00130     K_UCHAR WakeNext();
00131
00132
00133 #if KERNEL_USE_TIMEOUTS
00134
00142     K_BOOL Claim_i( K_ULONG ulWaitTimeMS_ );
00143 #else
00144
00150     void Claim_i(void);
00151 #endif
00152
00153     K_UCHAR m_ucRecurse;
00154     K_UCHAR m_bReady;
00155     K_UCHAR m_ucMaxPri;
00156     Thread *m_pclOwner;
00157
00158 };
00159
00160 #endif //KERNEL_USE_MUTEX
00161 #endif //__MUTEX_H_
00163

```

15.71 /home/moslevin/Project/R1/kernel/public/paniccodes.h File Reference

Defines the reason codes thrown when a kernel panic occurs.

15.71.1 Detailed Description

Defines the reason codes thrown when a kernel panic occurs.

Definition in file [paniccodes.h](#).

15.72 paniccodes.h

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00020 #ifndef __PANIC_CODES_H
00021 #define __PANIC_CODES_H
00022
00023 #define PANIC_ASSERT_FAILED (1)
00024 #define PANIC_LIST_UNLINK_FAILED (2)
00025 #define PANIC_STACK_SLACK_VIOLATED (3)
00026
00027 #endif // __PANIC_CODES_H
00028

```

15.73 /home/moslevin/Project/R1/kernel/public/profile.h File Reference

High-precision profiling timers.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "ll.h"

```

Classes

- class [ProfileTimer](#)
Profiling timer.

15.73.1 Detailed Description

High-precision profiling timers. Enables the profiling and instrumentation of performance-critical code. Multiple timers can be used simultaneously to enable system-wide performance metrics to be computed in a lightweight manner.

Usage:

```

ProfileTimer clMyTimer;
int i;

clMyTimer.Init();

// Profile the same block of code ten times
for (i = 0; i < 10; i++)
{
    clMyTimer.Start();
    ...
    //Block of code to profile
    ...
    clMyTimer.Stop();
}

```

```

}

// Get the average execution time of all iterations
ulAverageTimer = clMyTimer.GetAverage();

// Get the execution time from the last iteration
ulLastTimer = clMyTimer.GetCurrent();

```

Definition in file [profile.h](#).

15.74 profile.h

```

00001  /*=====
00002
00003  _____
00004  |   \   /   |   \   /   |   \   /   |   \   /   |   \   /   |
00005  |  / \ / \  |  / \ / \  |  / \ / \  |  / \ / \  |  / \ / \  |
00006  | /   /   \ | /   /   \ | /   /   \ | /   /   \ | /   /   \ |
00007  |_____|   |_____|   |_____|   |_____|   |_____|   |
00008
00009  --[Mark3 Realtime Platform]-----
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information
00013  =====*/
00053  #ifndef __PROFILE_H__
00054  #define __PROFILE_H__
00055
00056  #include "kerneltypes.h"
00057  #include "mark3cfg.h"
00058  #include "ll.h"
00059
00060  #if KERNEL_USE_PROFILER
00061
00069  class ProfileTimer
00070  {
00071
00072  public:
00079      void Init();
00080
00087      void Start();
00088
00095      void Stop();
00096
00104      K_ULONG GetAverage();
00105
00114      K_ULONG GetCurrent();
00115
00116  private:
00117
00126      K_ULONG ComputeCurrentTicks(K_USHORT usCount_,
                                K_ULONG ulEpoch_);
00127
00128      K_ULONG m_ulCumulative;
00129      K_ULONG m_ulCurrentIteration;
00130      K_USHORT m_usInitial;
00131      K_ULONG m_ulInitialEpoch;
00132      K_USHORT m_usIterations;
00133      K_UCHAR m_bActive;
00134  };
00135
00136  #endif // KERNEL_USE_PROFILE
00137
00138  #endif

```

15.75 /home/moslevin/Project/R1/kernel/public/quantum.h File Reference

[Thread Quantum](#) declarations for Round-Robin Scheduling.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "thread.h"
#include "timer.h"
#include "timerlist.h"
#include "timerscheduler.h"

```

Classes

- class [Quantum](#)

Static-class used to implement [Thread](#) quantum functionality, which is a key part of round-robin scheduling.

15.75.1 Detailed Description

[Thread Quantum](#) declarations for Round-Robin Scheduling.

Definition in file [quantum.h](#).

15.76 quantum.h

```

00001  /*=====
00002
00003  _____
00004  |   \   /   |   \   /   |   \   /   |   \   /   |   \   /   |
00005  |  / \ / \  |  / \ / \  |  / \ / \  |  / \ / \  |  / \ / \  |
00006  |_/   \_ \_|_/   \_ \_|_/   \_ \_|_/   \_ \_|_/   \_ \_|_/   \_ \_|
00007  |_____|   |_____|   |_____|   |_____|   |_____|   |_____|   |
00008
00009  --[Mark3 Realtime Platform]-----
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information
00013  =====*/
00022  #ifndef __KQUANTUM_H__
00023  #define __KQUANTUM_H__
00024
00025  #include "kerneltypes.h"
00026  #include "mark3cfg.h"
00027
00028  #include "thread.h"
00029  #include "timer.h"
00030  #include "timerlist.h"
00031  #include "timerscheduler.h"
00032
00033  #if KERNEL_USE_QUANTUM
00034  class Timer;
00035
00041  class Quantum
00042  {
00043  public:
00052      static void UpdateTimer();
00053
00060      static void AddThread( Thread *pclThread_ );
00061
00067      static void RemoveThread();
00068
00077      static void SetInTimer(void) { m_bInTimer = true; }
00078
00084      static void ClearInTimer(void) { m_bInTimer = false; }
00085
00086  private:
00098      static void SetTimer( Thread *pclThread_ );
00099
00100      static Timer m_clQuantumTimer;
00101      static K_UCHAR m_bActive;
00102      static K_UCHAR m_bInTimer;
00103  };
00104
00105  #endif //KERNEL_USE_QUANTUM
00106
00107  #endif

```

15.77 /home/moslevin/Project/R1/kernel/public/scheduler.h File Reference

[Thread](#) scheduler function declarations.


```

00013 =====*/
00046 #ifndef __SCHEDULER_H__
00047 #define __SCHEDULER_H__
00048
00049 #include "kerneltypes.h"
00050 #include "thread.h"
00051 #include "threadport.h"
00052
00053 extern volatile Thread *g_pstNext;
00054 extern Thread *g_pstCurrent;
00055
00056 #define NUM_PRIORITIES (8)
00057 //-----
00058
00062 class Scheduler
00063 {
00064 public:
00070     static void Init();
00071
00079     static void Schedule();
00080
00088     static void Add(Thread *pclThread_);
00089
00098     static void Remove(Thread *pclThread_);
00099
00112     static K_BOOL SetScheduler(K_BOOL bEnable_);
00113
00119     static Thread *GetCurrentThread() { return g_pstCurrent; }
00120
00127     static volatile Thread *GetNextThread() { return g_pstNext; }
00128
00137     static ThreadList *GetThreadList(K_UCHAR ucPriority_) { return &
m_aclPriorities[ucPriority_]; }
00138
00145     static ThreadList *GetStopList() { return &m_clStopList; }
00146
00155     static K_UCHAR IsEnabled() { return m_bEnabled; }
00156
00163     static void QueueScheduler() { m_bQueuedSchedule = true; }
00164
00165 private:
00167     static K_BOOL m_bEnabled;
00168
00170     static K_BOOL m_bQueuedSchedule;
00171
00173     static ThreadList m_clStopList;
00174
00176     static ThreadList m_aclPriorities[NUM_PRIORITIES];
00177
00179     static K_UCHAR m_ucPriFlag;
00180 };
00181 #endif
00182

```

15.79 /home/moslevin/Project/R1/kernel/public/thread.h File Reference

Platform independent thread class declarations.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "ll.h"
#include "threadlist.h"
#include "scheduler.h"
#include "threadport.h"
#include "quantum.h"

```

Classes

- class [Thread](#)

Object providing fundamental multitasking support in the kernel.

- struct [FakeThread_t](#)

If the kernel is set up to use an idle function instead of an idle thread, we use a placeholder data structure to "simulate" the effect of having an idle thread in the system.

Typedefs

- typedef void(* [ThreadEntry_t](#))(void *pvArg_)
Function pointer type used for thread entryptoint functions.

Enumerations

- enum [ThreadState_t](#)
Enumeration representing the different states a thread can exist in.

15.79.1 Detailed Description

Platform independent thread class declarations. Threads are an atomic unit of execution, and each instance of the thread class represents an instance of a program running on the processor. The [Thread](#) is the fundamental user-facing object in the kernel - it is what makes multiprocessing possible from application code.

In Mark3, threads each have their own context - consisting of a stack, and all of the registers required to multiplex a processor between multiple threads.

The [Thread](#) class inherits directly from the [LinkListNode](#) class to facilitate efficient thread management using Double, or Double-Circular linked lists.

Definition in file [thread.h](#).

15.80 thread.h

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00035 #ifndef __THREAD_H__
00036 #define __THREAD_H__
00037
00038 #include "kerneltypes.h"
00039 #include "mark3cfg.h"
00040
00041 #include "ll.h"
00042 #include "threadlist.h"
00043 #include "scheduler.h"
00044 #include "threadport.h"
00045 #include "quantum.h"
00046
00047 //-----
00051 typedef void (*ThreadEntry_t)(void *pvArg_);
00052
00053 //-----
00057 typedef enum
00058 {
00059     THREAD_STATE_EXIT = 0,
00060     THREAD_STATE_READY,
00061     THREAD_STATE_BLOCKED,
00062     THREAD_STATE_STOP,
00063 } ThreadState_t;
00064
00067 //-----
00071 class Thread : public LinkListNode
00072 {
00073 public:
00093     void Init(K_WORD *paucStack_,
00094              K_USHORT usStackSize_,

```

```

00095         K_UCHAR ucPriority_,
00096         ThreadEntry_t pfEntryPoint_,
00097         void *pvArg_ );
00098
00106     void Start();
00107
00108
00115     void Stop();
00116
00117 #if KERNEL_USE_THREADNAME
00118
00127     void SetName(const K_CHAR *szName_) { m_szName = szName_; }
00128
00135     const K_CHAR* GetName() { return m_szName; }
00136 #endif
00137
00146     ThreadList *GetOwner(void) { return m_pclOwner; }
00147
00155     ThreadList *GetCurrent(void) { return m_pclCurrent; }
00156
00165     K_UCHAR GetPriority(void) { return m_ucPriority; }
00166
00174     K_UCHAR GetCurPriority(void) { return m_ucCurPriority; }
00175
00176 #if KERNEL_USE_QUANTUM
00177
00184     void SetQuantum( K_USHORT usQuantum_ ) { m_usQuantum = usQuantum_; }
00185
00193     K_USHORT GetQuantum(void) { return m_usQuantum; }
00194 #endif
00195
00203     void SetCurrent( ThreadList *pclNewList_ ) {
00204         m_pclCurrent = pclNewList_; }
00205
00212     void SetOwner( ThreadList *pclNewList_ ) { m_pclOwner = pclNewList_; }
00213
00214
00227     void SetPriority(K_UCHAR ucPriority_);
00228
00238     void InheritPriority(K_UCHAR ucPriority_);
00239
00240 #if KERNEL_USE_DYNAMIC_THREADS
00241
00252     void Exit();
00253 #endif
00254
00255 #if KERNEL_USE_SLEEP
00256
00264     static void Sleep(K_ULONG ulTimeMs_);
00265
00274     static void USleep(K_ULONG ulTimeUs_);
00275 #endif
00276
00284     static void Yield(void);
00285
00293     void SetID( K_UCHAR ucID_ ) { m_ucThreadID = ucID_; }
00294
00302     K_UCHAR GetID() { return m_ucThreadID; }
00303
00304
00317     K_USHORT GetStackSlack();
00318
00319 #if KERNEL_USE_EVENTFLAG
00320
00327     K_USHORT GetEventFlagMask() { return m_usFlagMask; }
00328
00333     void SetEventFlagMask(K_USHORT usMask_) {
00334         m_usFlagMask = usMask_; }
00335
00340     void SetEventFlagMode(EventFlagOperation_t eMode_ ) {
00341         m_eFlagMode = eMode_; }
00342
00346     EventFlagOperation_t GetEventFlagMode() { return
00347         m_eFlagMode; }
00348 #endif
00349 #if KERNEL_USE_TIMEOUTS || KERNEL_USE_SLEEP
00350
00353     Timer *GetTimer();
00354 #endif
00355 #if KERNEL_USE_TIMEOUTS
00356
00364     void SetExpired( K_BOOL bExpired_ );
00365
00372     K_BOOL GetExpired();
00373 #endif
00374

```



```

00375 #if KERNEL_USE_IDLE_FUNC
00376
00381     void InitIdle();
00382 #endif
00383
00390     ThreadState_t GetState()          { return
m_eState; }
00391
00399     void SetState( ThreadState_t eState_ ) { m_eState = eState_; }
00400
00401     friend class ThreadPort;
00402
00403 private:
00411     static void ContextSwitchSWI(void);
00412
00417     void SetPriorityBase(K_UCHAR ucPriority_);
00418
00420     K_WORD *m_pwStackTop;
00421
00423     K_WORD *m_pwStack;
00424
00426     K_UCHAR m_ucThreadID;
00427
00429     K_UCHAR m_ucPriority;
00430
00432     K_UCHAR m_ucCurPriority;
00433
00435     ThreadState_t m_eState;
00436
00437 #if KERNEL_USE_THREADNAME
00438     const K_CHAR *m_szName;
00440 #endif
00441
00443     K_USHORT m_usStackSize;
00444
00446     ThreadList *m_pclCurrent;
00447
00449     ThreadList *m_pclOwner;
00450
00452     ThreadEntry_t m_pfEntryPoint;
00453
00455     void *m_pvArg;
00456
00457 #if KERNEL_USE_QUANTUM
00458     K_USHORT m_usQuantum;
00460 #endif
00461
00462 #if KERNEL_USE_EVENTFLAG
00463     K_USHORT m_usFlagMask;
00465
00467     EventFlagOperation_t m_eFlagMode;
00468 #endif
00469
00470 #if KERNEL_USE_TIMEOUTS || KERNEL_USE_SLEEP
00471     Timer m_clTimer;
00473 #endif
00474 #if KERNEL_USE_TIMEOUTS
00475     K_BOOL m_bExpired;
00477 #endif
00478
00479 };
00480
00481 #if KERNEL_USE_IDLE_FUNC
00482 //-----
00494 typedef struct
00495 {
00496     LinkListNode *next;
00497     LinkListNode *prev;
00498
00500     K_WORD *m_pwStackTop;
00501
00503     K_WORD *m_pwStack;
00504
00506     K_UCHAR m_ucThreadID;
00507
00509     K_UCHAR m_ucPriority;
00510
00512     K_UCHAR m_ucCurPriority;
00513
00515     ThreadState_t m_eState;
00516
00517 #if KERNEL_USE_THREADNAME
00518     const K_CHAR *m_szName;
00520 #endif
00521
00522 } FakeThread_t;
00523 #endif

```

```
00524
00525 #endif
```

15.81 /home/moslevin/Project/R1/kernel/public/threadlist.h File Reference

[Thread](#) linked-list declarations.

```
#include "kerneltypes.h"
#include "ll.h"
```

Classes

- class [ThreadList](#)

This class is used for building thread-management facilities, such as schedulers, and blocking objects.

15.81.1 Detailed Description

[Thread](#) linked-list declarations.

Definition in file [threadlist.h](#).

15.82 threadlist.h

```
00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00022 #ifndef __THREADLIST_H__
00023 #define __THREADLIST_H__
00024
00025 #include "kerneltypes.h"
00026 #include "ll.h"
00027
00028 class Thread;
00029
00034 class ThreadList : public CircularLinkedList
00035 {
00036 public:
00040 ThreadList() { m_ucPriority = 0; m_pucFlag = NULL; }
00041
00049 void SetPriority(K_UCHAR ucPriority_);
00050
00059 void SetFlagPointer(K_UCHAR *pucFlag_);
00060
00068 void Add(LinkListNode *node_);
00069
00083 void Add(LinkListNode *node_, K_UCHAR *pucFlag_,
K_UCHAR ucPriority_);
00084
00092 void Remove(LinkListNode *node_);
00093
00101 Thread *HighestWaiter();
00102 private:
00103
00105 K_UCHAR m_ucPriority;
00106
00108 K_UCHAR *m_pucFlag;
00109 };
00110
00111 #endif
```

00112

15.83 /home/moslevin/Project/R1/kernel/public/timer.h File Reference

[Timer](#) object declarations.

```
#include "kerneltypes.h"
#include "mark3cfg.h"
#include "ll.h"
```

Classes

- class [Timer](#)

[Timer](#) - an event-driven execution context based on a specified time interval.

Macros

- #define [TIMERLIST_FLAG_ONE_SHOT](#) (0x01)
[Timer](#) is one-shot.
- #define [TIMERLIST_FLAG_ACTIVE](#) (0x02)
[Timer](#) is currently active.
- #define [TIMERLIST_FLAG_CALLBACK](#) (0x04)
[Timer](#) is pending a callback.
- #define [TIMERLIST_FLAG_EXPIRED](#) (0x08)
[Timer](#) is actually expired.
- #define [MAX_TIMER_TICKS](#) (0x7FFFFFFF)
Maximum value to set.
- #define [MIN_TICKS](#) (3)
The minimum tick value to set.

Typedefs

- typedef void(* [TimerCallback_t](#))([Thread](#) *pclOwner_, void *pvData_)
This type defines the callback function type for timer events.

15.83.1 Detailed Description

[Timer](#) object declarations.

Definition in file [timer.h](#).

15.83.2 Macro Definition Documentation

15.83.2.1 #define [TIMERLIST_FLAG_EXPIRED](#) (0x08)

[Timer](#) is actually expired.

Definition at line 36 of file [timer.h](#).

15.83.3 Typedef Documentation

15.83.3.1 typedef void(* TimerCallback_t)(Thread *pclOwner_, void *pvData_)

This type defines the callback function type for timer events.

Since these are called from an interrupt context, they do not operate from within a thread or object context directly – as a result, the context must be manually passed into the calls.

pclOwner_ is a pointer to the thread that owns the timer pvData_ is a pointer to some data or object that needs to know about the timer's expiry from within the timer interrupt context.

Definition at line 94 of file [timer.h](#).

15.84 timer.h

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00021 #ifndef __TIMER_H__
00022 #define __TIMER_H__
00023
00024 #include "kerneltypes.h"
00025 #include "mark3cfg.h"
00026
00027 #include "ll.h"
00028
00029 #if KERNEL_USE_TIMERS
00030 class Thread;
00031
00032 //-----
00033 #define TIMERLIST_FLAG_ONE_SHOT (0x01)
00034 #define TIMERLIST_FLAG_ACTIVE (0x02)
00035 #define TIMERLIST_FLAG_CALLBACK (0x04)
00036 #define TIMERLIST_FLAG_EXPIRED (0x08)
00037
00038 //-----
00039 #if KERNEL_TIMERS_TICKLESS
00040
00041 //-----
00042 #define MAX_TIMER_TICKS (0x7FFFFFFF)
00043
00044 //-----
00045 /*
00046     Ugly macros to support a wide resolution of delays.
00047     Given a 16-bit timer @ 16MHz & 256 cycle prescaler, this gives us...
00048     Max time, SECONDS_TO_TICKS: 68719s
00049     Max time, MSECONDS_TO_TICKS: 6871.9s
00050     Max time, USECONDS_TO_TICKS: 6.8719s
00051
00052     ...With a 16us tick resolution.
00053
00054     Depending on the system frequency and timer resolution, you may want to
00055     customize these values to suit your system more appropriately.
00056 */
00057 //-----
00058 #define SECONDS_TO_TICKS(x) (((K_ULONG)x) * TIMER_FREQ)
00059 #define MSECONDS_TO_TICKS(x) (((((K_ULONG)x) * (TIMER_FREQ/100)) + 5) / 10))
00060 #define USECONDS_TO_TICKS(x) (((((K_ULONG)x) * TIMER_FREQ) + 50000) / 1000000))
00061
00062 //-----
00063 #define MIN_TICKS (3)
00064
00065 //-----
00066 #else
00067 //-----
00068 // Tick-based timers, assuming 1khz tick rate
00069 #define MAX_TIMER_TICKS (0x7FFFFFFF)
00070

```

```

00071 //-----
00072 // add time because we don't know how far in an epoch we are when a call is made.
00073 #define SECONDS_TO_TICKS(x)      ((K_ULONG)(x) * 1000) + 1)
00074 #define MSECONDS_TO_TICKS(x)     ((K_ULONG)(x + 1))
00075 #define USECONDS_TO_TICKS(x)     ((K_ULONG)(x + 999)) / 1000)
00076
00077 //-----
00078 #define MIN_TICKS                  (1)
00079 //-----
00080
00081 #endif // KERNEL_TIMERS_TICKLESS
00082
00083 //-----
00094 typedef void (*TimerCallback_t)(Thread *pclOwner_, void *pvData_);
00095
00096 //-----
00097 class TimerList;
00098 class TimerScheduler;
00099 class Quantum;
00100 class Timer : public LinkListNode
00101 {
00102 public:
00103     Timer() { }
00104
00105     void Init() { ClearNode(); m_ulInterval = 0;
00106                 m_ulTimerTolerance = 0; m_ulTimeLeft = 0;
00107                 m_ucFlags = 0; }
00108
00109     void Start( K_BOOL bRepeat_, K_ULONG ulIntervalMs_,
00110                TimerCallback_t pfCallback_, void *pvData_ );
00111
00112     void Start( K_BOOL bRepeat_, K_ULONG ulIntervalMs_,
00113                K_ULONG ulToleranceMs_, TimerCallback_t pfCallback_, void *pvData_ );
00114
00115     void Stop();
00116
00117     void SetFlags( K_UCHAR ucFlags_ ) { m_ucFlags = ucFlags_; }
00118
00119     void SetCallback( TimerCallback_t pfCallback_ ) {
00120         m_pfCallback = pfCallback_; }
00121
00122     void SetData( void *pvData_ ) { m_pvData = pvData_; }
00123
00124     void SetOwner( Thread *pclOwner_ ) { m_pclOwner = pclOwner_; }
00125
00126     void SetIntervalTicks(K_ULONG ulTicks_);
00127
00128     void SetIntervalSeconds(K_ULONG ulSeconds_);
00129
00130     K_ULONG GetInterval() { return m_ulInterval; }
00131
00132     void SetIntervalMSeconds(K_ULONG ulMSeconds_);
00133
00134     void SetIntervalUSeconds(K_ULONG ulUSeconds_);
00135
00136     void SetTolerance(K_ULONG ulTicks_);
00137 private:
00138     friend class TimerList;
00139
00140     K_UCHAR m_ucFlags;
00141
00142     TimerCallback_t m_pfCallback;
00143
00144     K_ULONG m_ulInterval;
00145
00146     K_ULONG m_ulTimeLeft;
00147
00148     K_ULONG m_ulTimerTolerance;
00149
00150     Thread *m_pclOwner;
00151
00152     void *m_pvData;
00153 };
00154
00155 #endif // KERNEL_USE_TIMERS
00156 #endif

```

15.85 /home/moslevin/Project/R1/kernel/public/timerlist.h File Reference

[Timer](#) list declarations.

```
#include "kerneltypes.h"
#include "mark3cfg.h"
#include "timer.h"
```

Classes

- class [TimerList](#)

[TimerList](#) class - a doubly-linked-list of timer objects.

15.85.1 Detailed Description

[Timer](#) list declarations. These classes implements a linked list of timer objects attached to the global kernel timer scheduler.

Definition in file [timerlist.h](#).

15.86 timerlist.h

```
00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00024 #ifndef __TIMERLIST_H__
00025 #define __TIMERLIST_H__
00026
00027 #include "kerneltypes.h"
00028 #include "mark3cfg.h"
00029
00030 #include "timer.h"
00031 #if KERNEL_USE_TIMERS
00032
00033 //-----
00037 class TimerList : public DoubleLinkedList
00038 {
00039 public:
00046     void Init();
00047
00055     void Add(Timer *pclListNode_);
00056
00064     void Remove(Timer *pclListNode_);
00065
00072     void Process();
00073
00074 private:
00076     K_ULONG m_ulNextWakeup;
00077
00079     K_UCHAR m_bTimerActive;
00080 };
00081
00082 #endif // KERNEL_USE_TIMERS
00083
00084 #endif
```

Timer scheduler declarations.

Classes

- "Static" Class used to interface a global [TimerList](#) with the rest of the kernel.*

Timer scheduler declarations.

Definition in file [timerscheduler.h](#).

```

00001  /*-----
00002
00003  |-----|-----|-----|-----|-----|-----|
00004  |   \   /   |   \   /   |   \   /   |   \   /   |   \   /   |
00005  |   /   \   |   /   \   |   /   \   |   /   \   |   /   \   |
00006  |-----|-----|-----|-----|-----|-----|
00007  |-----|-----|-----|-----|-----|-----|
00008
00009  --[Mark3 Realtime Platform]-----
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information
00013  =====*/
00021  #ifndef __TIMERSCHEDULER_H__
00022  #define __TIMERSCHEDULER_H__
00023
00024  #include "kerneltypes.h"
00025  #include "mark3cfg.h"
00026
00027  #include "ll.h"
00028  #include "timer.h"
00029  #include "timerlist.h"
00030
00031  #if KERNEL_USE_TIMERS
00032
00033  //-----
00038  class TimerScheduler
00039  {
00040  public:
00047      static void Init() { m_clTimerList.Init(); }
00048
00057      static void Add(Timer *pclListNode_)
00058          {m_clTimerList.Add(pclListNode_); }
00059
00068      static void Remove(Timer *pclListNode_)
00069          {m_clTimerList.Remove(pclListNode_); }
00070
00079      static void Process() {m_clTimerList.Process();}
00080  private:
00081
00083      static TimerList m_clTimerList;
00084  };
00085
00086  #endif //KERNEL_USE_TIMERS
00087
00088  #endif //__TIMERSCHEDULER_H__
00089

```

15.89 /home/moslevin/Project/R1/kernel/public/tracebuffer.h File Reference

[Kernel](#) trace buffer class declaration.

```
#include "kerneltypes.h"
#include "mark3cfg.h"
#include "writebuf16.h"
```

15.89.1 Detailed Description

[Kernel](#) trace buffer class declaration. Global kernel trace-buffer. Used to instrument the kernel with lightweight encoded print statements. If something goes wrong, the tracebuffer can be examined for debugging purposes. Also, subsets of kernel trace information can be extracted and analyzed to provide information about runtime performance, thread-scheduling, and other nifty things in real-time.

Definition in file [tracebuffer.h](#).

15.90 tracebuffer.h

```
00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00024 #ifndef __TRACEBUFFER_H__
00025 #define __TRACEBUFFER_H__
00026
00027 #include "kerneltypes.h"
00028 #include "mark3cfg.h"
00029 #include "writebuf16.h"
00030
00031 #if KERNEL_USE_DEBUG && !KERNEL_AWARE_SIMULATION
00032
00033 #define TRACE_BUFFER_SIZE          (16)
00034
00038 class TraceBuffer
00039 {
00040 public:
00046     static void Init();
00047
00055     static K_USHORT Increment();
00056
00065     static void Write( K_USHORT *pusData_, K_USHORT usSize_ );
00066
00075     void SetCallback( WriteBufferCallback pfCallback_ )
00076     { m_clBuffer.SetCallback( pfCallback_ ); }
00077 private:
00078
00079     static WriteBuffer16 m_clBuffer;
00080     static volatile K_USHORT m_usIndex;
00081     static K_USHORT m_ausBuffer[ (TRACE_BUFFER_SIZE / sizeof( K_USHORT )) ];
00082 };
00083
00084 #endif //KERNEL_USE_DEBUG
00085
00086 #endif
```

15.91 /home/moslevin/Project/R1/kernel/public/writebuf16.h File Reference

Thread-safe circular buffer implementation with 16-bit elements.


```
#include "kerneltypes.h"
#include "mark3cfg.h"
```

15.91.1 Detailed Description

Thread-safe circular buffer implementation with 16-bit elements.

Definition in file [writebuf16.h](#).

15.92 writebuf16.h

```
00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00020 #ifndef __WRITEBUF16_H__
00021 #define __WRITEBUF16_H__
00022
00023 #include "kerneltypes.h"
00024 #include "mark3cfg.h"
00025
00026 #if KERNEL_USE_DEBUG && !KERNEL_AWARE_SIMULATION
00027
00032 typedef void (*WriteBufferCallback)( K_USHORT *pusData_, K_USHORT usSize_ );
00033
00040 class WriteBuffer16
00041 {
00042 public:
00053     void SetBuffers( K_USHORT *pusData_, K_USHORT usSize_ )
00054     {
00055         m_pusData = pusData_;
00056         m_usSize = usSize_;
00057         m_usHead = 0;
00058         m_usTail = 0;
00059     }
00060
00072     void SetCallback( WriteBufferCallback pfCallback_ )
00073     { m_pfCallback = pfCallback_; }
00074
00083     void WriteData( K_USHORT *pusBuf_, K_USHORT usLen_ );
00084
00094     void WriteVector( K_USHORT **ppusBuf_, K_USHORT *pusLen_,
00095                     K_UCHAR ucCount_);
00096 private:
00097     K_USHORT *m_pusData;
00098
00099     volatile K_USHORT m_usSize;
00100     volatile K_USHORT m_usHead;
00101     volatile K_USHORT m_usTail;
00102
00103     WriteBufferCallback m_pfCallback;
00104 };
00105 #endif
00106
00107 #endif
```

15.93 /home/moslevin/Project/R1/kernel/quantum.cpp File Reference

[Thread Quantum](#) Implementation for Round-Robin Scheduling.


```

00027 #include "quantum.h"
00028 #include "kerneldebug.h"
00029 #include "kernelaware.h"
00030 //-----
00031 #if defined __FILE_ID__
00032     #undef __FILE_ID__
00033 #endif
00034 #define __FILE_ID__      QUANTUM_CPP
00035
00036 #if KERNEL_USE_QUANTUM
00037
00038 //-----
00039 static volatile K_BOOL bAddQuantumTimer; // Indicates that a timer add is pending
00040
00041 //-----
00042 Timer Quantum::m_clQuantumTimer; // The global timernodelist_t object
00043 K_UCHAR Quantum::m_bActive;
00044 K_UCHAR Quantum::m_bInTimer;
00045 //-----
00056 static void QuantumCallback(Thread *pclThread_, void *pvData_)
00057 {
00058     // Validate thread pointer, check that source/destination match (it's
00059     // in its real priority list). Also check that this thread was part of
00060     // the highest-running priority level.
00061     if (pclThread_>GetPriority() >= Scheduler::GetCurrentThread()->
        GetPriority())
00062     {
00063         if (pclThread_>GetCurrent()->GetHead() != pclThread_>
            GetCurrent()->GetTail() )
00064         {
00065             bAddQuantumTimer = true;
00066             pclThread_>GetCurrent()->PivotForward();
00067         }
00068     }
00069 }
00070
00071 //-----
00072 void Quantum::SetTimer(Thread *pclThread_)
00073 {
00074     m_clQuantumTimer.SetIntervalMSeconds(pclThread_>
        GetQuantum());
00075     m_clQuantumTimer.SetFlags(TIMERLIST_FLAG_ONE_SHOT);
00076     m_clQuantumTimer.SetData(NULL);
00077     m_clQuantumTimer.SetCallback((TimerCallback_t)
        QuantumCallback);
00078     m_clQuantumTimer.SetOwner(pclThread_);
00079 }
00080
00081 //-----
00082 void Quantum::AddThread(Thread *pclThread_)
00083 {
00084     if (m_bActive
00085     #if KERNEL_USE_IDLE_FUNC
00086         || (pclThread_ == Kernel::GetIdleThread())
00087     #endif
00088     )
00089     {
00090         return;
00091     }
00092
00093     // If this is called from the timer callback, queue a timer add...
00094     if (m_bInTimer)
00095     {
00096         bAddQuantumTimer = true;
00097         return;
00098     }
00099
00100     // If this isn't the only thread in the list.
00101     if ( pclThread_>GetCurrent()->GetHead() !=
00102         pclThread_>GetCurrent()->GetTail() )
00103     {
00104         Quantum::SetTimer(pclThread_);
00105         TimerScheduler::Add(&m_clQuantumTimer);
00106         m_bActive = 1;
00107     }
00108 }
00109
00110 //-----
00111 void Quantum::RemoveThread(void)
00112 {
00113     if (!m_bActive)
00114     {
00115         return;
00116     }
00117
00118     // Cancel the current timer
00119     TimerScheduler::Remove(&m_clQuantumTimer);

```

```

00120     m_bActive = 0;
00121 }
00122
00123 //-----
00124 void Quantum::UpdateTimer(void)
00125 {
00126     // If we have to re-add the quantum timer (more than 2 threads at the
00127     // high-priority level...)
00128     if (bAddQuantumTimer)
00129     {
00130         // Trigger a thread yield - this will also re-schedule the
00131         // thread *and* reset the round-robin scheduler.
00132         Thread::Yield();
00133         bAddQuantumTimer = false;
00134     }
00135 }
00136
00137 #endif //KERNEL_USE_QUANTUM

```

15.95 /home/moslevin/Project/R1/kernel/scheduler.cpp File Reference

Strict-Priority + Round-Robin thread scheduler implementation.

```

#include "kerneltypes.h"
#include "ll.h"
#include "scheduler.h"
#include "thread.h"
#include "threadport.h"
#include "kernel.h"
#include "kerneldebug.h"

```

Macros

- `#define __FILE_ID__ SCHEDULER_CPP`
File ID used in kernel trace calls.

Variables

- volatile `Thread * g_pstNext`
Pointer to the currently-chosen next-running thread.
- `Thread * g_pstCurrent`
Pointer to the currently-running thread.
- static const `K_UCHAR aucCLZ [16] = {255,0,1,1,2,2,2,3,3,3,3,3,3,3}`
This implements a 4-bit "Count-leading-zeros" operation using a RAM-based lookup table.

15.95.1 Detailed Description

Strict-Priority + Round-Robin thread scheduler implementation.

Definition in file [scheduler.cpp](#).

15.95.2 Variable Documentation

15.95.2.1 `const K_UCHAR aucCLZ[16] = {255,0,1,1,2,2,2,3,3,3,3,3,3,3}` `[static]`

This implements a 4-bit "Count-leading-zeros" operation using a RAM-based lookup table.

It is used to efficiently perform a CLZ operation under the assumption that a native CLZ instruction is unavailable. This table is further optimized to provide a 0xFF result in the event that the index value is itself zero, allowing us to quickly identify whether or not subsequent 4-bit LUT operations are required to complete the scheduling process.

Definition at line 56 of file scheduler.cpp.

15.96 scheduler.cpp

```

00001 /*
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00022 #include "kerneltypes.h"
00023 #include "ll.h"
00024 #include "scheduler.h"
00025 #include "thread.h"
00026 #include "threadport.h"
00027 #include "kernel.h"
00028 #include "kerneldebug.h"
00029 //-----
00030 #if defined __FILE_ID__
00031     #undef __FILE_ID__
00032 #endif
00033 #define __FILE_ID__ SCHEDULER_CPP
00034
00035 //-----
00036 volatile Thread *g_pstNext;
00037 Thread *g_pstCurrent;
00038
00039 //-----
00040 K_BOOL Scheduler::m_bEnabled;
00041 K_BOOL Scheduler::m_bQueuedSchedule;
00042
00043 ThreadList Scheduler::m_clStopList;
00044 ThreadList Scheduler::m_aclPriorities[
    NUM_PRIORITIES];
00045 K_UCHAR Scheduler::m_ucPriFlag;
00046
00047 //-----
00056 static const K_UCHAR aucCLZ[16] = {255,0,1,1,2,2,2,2,3,3,3,3,3,3,3,3};
00057
00058 //-----
00059 void Scheduler::Init()
00060 {
00061     m_ucPriFlag = 0;
00062     for (int i = 0; i < NUM_PRIORITIES; i++)
00063     {
00064         m_aclPriorities[i].SetPriority(i);
00065         m_aclPriorities[i].SetFlagPointer(&
m_ucPriFlag);
00066     }
00067     m_bQueuedSchedule = false;
00068 }
00069
00070 //-----
00071 void Scheduler::Schedule()
00072 {
00073     K_UCHAR ucPri = 0;
00074
00075     // Figure out what priority level has ready tasks (8 priorities max)
00076     // To do this, we apply our current active-thread bitmap (m_ucPriFlag)
00077     // and perform a CLZ on the upper four bits. If no tasks are found
00078     // in the higher priority bits, search the lower priority bits. This
00079     // also assumes that we always have the idle thread ready-to-run in
00080     // priority level zero.
00081     ucPri = aucCLZ[m_ucPriFlag >> 4];
00082     if (ucPri == 0xFF)
00083     {
00084         ucPri = aucCLZ[m_ucPriFlag & 0x0F];
00085     }
00086     else
00087     {
00088         ucPri += 4;

```

```

00089     }
00090
00091     #if KERNEL_USE_IDLE_FUNC
00092     if (ucPri == 0xFF)
00093     {
00094         // There aren't any active threads at all - set g_pstNext to IDLE
00095         g_pstNext = Kernel::GetIdleThread();
00096     }
00097     else
00098     #endif
00099     {
00100         // Get the thread node at this priority.
00101         g_pstNext = (Thread*)( m_aclPriorities[ucPri].GetHead() );
00102     }
00103     KERNEL_TRACE_1( STR_SCHEDULE_1, (K_USHORT)((Thread*)g_pstNext)->GetID() );
00104
00105 }
00106
00107 //-----
00108 void Scheduler::Add(Thread *pclThread_)
00109 {
00110     m_aclPriorities[pclThread_->GetPriority()].Add(pclThread_);
00111 }
00112
00113 //-----
00114 void Scheduler::Remove(Thread *pclThread_)
00115 {
00116     m_aclPriorities[pclThread_->GetPriority()].Remove(pclThread_);
00117 }
00118
00119 //-----
00120 K_BOOL Scheduler::SetScheduler(K_BOOL bEnable_)
00121 {
00122     K_BOOL bRet ;
00123     CS_ENTER();
00124     bRet = m_bEnabled;
00125     m_bEnabled = bEnable_;
00126     // If there was a queued scheduler event, dequeue and trigger an
00127     // immediate Yield
00128     if (m_bEnabled && m_bQueuedSchedule)
00129     {
00130         m_bQueuedSchedule = false;
00131         Thread::Yield();
00132     }
00133     CS_EXIT();
00134     return bRet;
00135 }

```

15.97 /home/moslevin/Project/R1/kernel/thread.cpp File Reference

Platform-Independent thread class Definition.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "thread.h"
#include "scheduler.h"
#include "kernelswi.h"
#include "timerlist.h"
#include "ksemaphore.h"
#include "quantum.h"
#include "kernel.h"
#include "kerneldebug.h"

```

Macros

- #define `__FILE_ID__` `THREAD_CPP`

File ID used in kernel trace calls.

Functions

- static void [ThreadSleepCallback](#) (Thread *pclOwner_, void *pvData_)

This callback is used to wake up a thread once the interval has expired.

15.97.1 Detailed Description

Platform-Independent thread class Definition.

Definition in file [thread.cpp](#).

15.98 thread.cpp

```

00001  /*=====
00002
00003  _____
00004  |   \   |   |   |   |   |   |   |   |   |
00005  |  / \  |   |   |   |   |   |   |   |   |
00006  | /   \ |   |   |   |   |   |   |   |   |
00007  |_____|   |   |   |   |   |   |   |   |
00008
00009  --[Mark3 Realtime Platform]-----
00010
00011  Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012  See license.txt for more information
00013  =====*/
00022  #include "kerneltypes.h"
00023  #include "mark3cfg.h"
00024
00025  #include "thread.h"
00026  #include "scheduler.h"
00027  #include "kernelswi.h"
00028  #include "timerlist.h"
00029  #include "ksemaphore.h"
00030  #include "quantum.h"
00031  #include "kernel.h"
00032  #include "kerneldebug.h"
00033
00034  //-----
00035  #if defined __FILE_ID__
00036      #undef __FILE_ID__
00037  #endif
00038  #define __FILE_ID__      THREAD_CPP
00039
00040  //-----
00041  void Thread::Init( K_WORD *pwStack_,
00042                   K_USHORT usStackSize_,
00043                   K_UCHAR ucPriority_,
00044                   ThreadEntry_t pfEntryPoint_,
00045                   void *pvArg_ )
00046  {
00047      static K_UCHAR ucThreadID = 0;
00048
00049      KERNEL_ASSERT( pwStack_ );
00050      KERNEL_ASSERT( pfEntryPoint_ );
00051
00052      ClearNode();
00053
00054      m_ucThreadID = ucThreadID++;
00055
00056      KERNEL_TRACE_1( STR_STACK_SIZE_1, usStackSize_ );
00057      KERNEL_TRACE_1( STR_PRIORITY_1, (K_UCHAR)ucPriority_ );
00058      KERNEL_TRACE_1( STR_THREAD_ID_1, (K_USHORT)
00059      m_ucThreadID );
00059      KERNEL_TRACE_1( STR_ENTRYPOINT_1, (K_USHORT)pfEntryPoint_ );
00060
00061      // Initialize the thread parameters to their initial values.
00062      m_pwStack = pwStack_;
00063      m_pwStackTop = TOP_OF_STACK(pwStack_, usStackSize_);
00064
00065      m_usStackSize = usStackSize_;
00066
00067      #if KERNEL_USE_QUANTUM
00068          m_usQuantum = THREAD_QUANTUM_DEFAULT;
00069      #endif
00070
00071      m_ucPriority = ucPriority_ ;
00072      m_ucCurPriority = m_ucPriority;

```

```

00073     m_pfEntryPoint = pfEntryPoint_;
00074     m_pvArg = pvArg_;
00075     m_eState = THREAD_STATE_STOP;
00076
00077     #if KERNEL_USE_THREADNAME
00078         m_szName = NULL;
00079     #endif
00080     #if KERNEL_USE_TIMERS
00081         m_clTimer.Init();
00082     #endif
00083
00084     // Call CPU-specific stack initialization
00085     ThreadPort::InitStack(this);
00086
00087     // Add to the global "stop" list.
00088     CS_ENTER();
00089     m_pclOwner = Scheduler::GetThreadList(
m_ucPriority);
00090     m_pclCurrent = Scheduler::GetStopList();
00091     m_pclCurrent->Add(this);
00092     CS_EXIT();
00093 }
00094
00095 //-----
00096 void Thread::Start(void)
00097 {
00098     // Remove the thread from the scheduler's "stopped" list, and add it
00099     // to the scheduler's ready list at the proper priority.
00100     KERNEL_TRACE_1( STR_THREAD_START_1, (K_USHORT)
m_ucThreadID );
00101
00102     CS_ENTER();
00103     Scheduler::GetStopList()->Remove(this);
00104     Scheduler::Add(this);
00105     m_pclOwner = Scheduler::GetThreadList(
m_ucPriority);
00106     m_pclCurrent = m_pclOwner;
00107     m_eState = THREAD_STATE_READY;
00108
00109     #if KERNEL_USE_QUANTUM
00110         if (GetCurPriority() >= Scheduler::GetCurrentThread()->
GetCurPriority())
00111         {
00112             // Deal with the thread Quantum
00113             Quantum::RemoveThread();
00114             Quantum::AddThread(this);
00115         }
00116     #endif
00117
00118     if (Kernel::IsStarted())
00119     {
00120         if (GetCurPriority() >= Scheduler::GetCurrentThread()->
GetCurPriority())
00121         {
00122             Thread::Yield();
00123         }
00124     }
00125     CS_EXIT();
00126 }
00127
00128 //-----
00129 void Thread::Stop()
00130 {
00131     K_BOOL bReschedule = 0;
00132
00133     CS_ENTER();
00134
00135     // If a thread is attempting to stop itself, ensure we call the scheduler
00136     if (this == Scheduler::GetCurrentThread())
00137     {
00138         bReschedule = true;
00139     }
00140
00141     // Add this thread to the stop-list (removing it from active scheduling)
00142     // Remove the thread from scheduling
00143     if (m_eState == THREAD_STATE_READY)
00144     {
00145         Scheduler::Remove(this);
00146     }
00147     else if (m_eState == THREAD_STATE_BLOCKED)
00148     {
00149         m_pclCurrent->Remove(this);
00150     }
00151
00152     m_pclOwner = Scheduler::GetStopList();
00153     m_pclCurrent = m_pclOwner;
00154     m_pclOwner->Add(this);

```



```

00155     m_eState = THREAD_STATE_STOP;
00156
00157 #if KERNEL_USE_TIMERS
00158     // Just to be safe - attempt to remove the thread's timer
00159     // from the timer-scheduler (does no harm if it isn't
00160     // in the timer-list)
00161     TimerScheduler::Remove(&m_clTimer);
00162 #endif
00163
00164     CS_EXIT();
00165
00166     if (bReschedule)
00167     {
00168         Thread::Yield();
00169     }
00170 }
00171
00172 #if KERNEL_USE_DYNAMIC_THREADS
00173 //-----
00174 void Thread::Exit()
00175 {
00176     K_BOOL bReschedule = 0;
00177
00178     KERNEL_TRACE_1( STR_THREAD_EXIT_1, m_ucThreadID );
00179
00180     CS_ENTER();
00181
00182     // If this thread is the actively-running thread, make sure we run the
00183     // scheduler again.
00184     if (this == Scheduler::GetCurrentThread())
00185     {
00186         bReschedule = 1;
00187     }
00188
00189     // Remove the thread from scheduling
00190     if (m_eState == THREAD_STATE_READY)
00191     {
00192         Scheduler::Remove(this);
00193     }
00194     else if (m_eState == THREAD_STATE_BLOCKED)
00195     {
00196         m_pclCurrent->Remove(this);
00197     }
00198
00199     m_pclCurrent = 0;
00200     m_pclOwner = 0;
00201     m_eState = THREAD_STATE_EXIT;
00202
00203     // We've removed the thread from scheduling, but interrupts might
00204     // trigger checks against this thread's currently priority before
00205     // we get around to scheduling new threads. As a result, set the
00206     // priority to idle to ensure that we always wind up scheduling
00207     // new threads.
00208     m_ucCurPriority = 0;
00209     m_ucPriority = 0;
00210
00211 #if KERNEL_USE_TIMERS
00212     // Just to be safe - attempt to remove the thread's timer
00213     // from the timer-scheduler (does no harm if it isn't
00214     // in the timer-list)
00215     TimerScheduler::Remove(&m_clTimer);
00216 #endif
00217
00218     CS_EXIT();
00219
00220     if (bReschedule)
00221     {
00222         // Choose a new "next" thread if we must
00223         Thread::Yield();
00224     }
00225 }
00226 #endif
00227
00228 #if KERNEL_USE_SLEEP
00229 //-----
00230 static void ThreadSleepCallback( Thread *pclOwner_, void *pvData_ )
00231 {
00232     Semaphore *pclSemaphore = static_cast<Semaphore*>(pvData_);
00233     // Post the semaphore, which will wake the sleeping thread.
00234     pclSemaphore->Post();
00235 }
00236
00237 //-----
00238 void Thread::Sleep(K_ULONG ulTimeMs_)
00239 {
00240     Semaphore clSemaphore;
00241     Timer *pclTimer = g_pstCurrent->GetTimer();

```

```

00243
00244 // Create a semaphore that this thread will block on
00245 clSemaphore.Init(0, 1);
00246
00247 // Create a one-shot timer that will call a callback that posts the
00248 // semaphore, waking our thread.
00249 pclTimer->Init();
00250 pclTimer->SetIntervalMSeconds(ulTimeMs_);
00251 pclTimer->SetCallback(ThreadSleepCallback);
00252 pclTimer->SetData((void*)&clSemaphore);
00253 pclTimer->SetFlags(TIMERLIST_FLAG_ONE_SHOT);
00254
00255 // Add the new timer to the timer scheduler, and block the thread
00256 TimerScheduler::Add(pclTimer);
00257 clSemaphore.Pend();
00258 }
00259
00260 //-----
00261 void Thread::USleep(K_ULONG ulTimeUs_)
00262 {
00263     Semaphore clSemaphore;
00264     Timer *pclTimer = g_pstCurrent->GetTimer();
00265
00266     // Create a semaphore that this thread will block on
00267     clSemaphore.Init(0, 1);
00268
00269     // Create a one-shot timer that will call a callback that posts the
00270     // semaphore, waking our thread.
00271     pclTimer->Init();
00272     pclTimer->SetIntervalUSEconds(ulTimeUs_);
00273     pclTimer->SetCallback(ThreadSleepCallback);
00274     pclTimer->SetData((void*)&clSemaphore);
00275     pclTimer->SetFlags(TIMERLIST_FLAG_ONE_SHOT);
00276
00277     // Add the new timer to the timer scheduler, and block the thread
00278     TimerScheduler::Add(pclTimer);
00279     clSemaphore.Pend();
00280 }
00281 #endif // KERNEL_USE_SLEEP
00282
00283 //-----
00284 K_USHORT Thread::GetStackSlack()
00285 {
00286     K_USHORT usCount = 0;
00287
00288     CS_ENTER();
00289
00290     for (usCount = 0; usCount < m_usStackSize; usCount++)
00291     {
00292         if (m_pwStack[usCount] != 0xFF)
00293         {
00294             break;
00295         }
00296     }
00297
00298     CS_EXIT();
00299
00300     return usCount;
00301 }
00302
00303 //-----
00304 void Thread::Yield()
00305 {
00306     CS_ENTER();
00307     // Run the scheduler
00308     if (Scheduler::IsEnabled())
00309     {
00310         Scheduler::Schedule();
00311
00312         // Only switch contexts if the new task is different than the old task
00313         if (Scheduler::GetCurrentThread() !=
00314             Scheduler::GetNextThread())
00315         {
00316             #if KERNEL_USE_QUANTUM
00317                 // new thread scheduled. Stop current quantum timer (if it exists),
00318                 // and restart it for the new thread (if required).
00319                 Quantum::RemoveThread();
00320                 Quantum::AddThread((Thread*)g_pstNext);
00321             #endif
00322             Thread::ContextSwitchSWI();
00323         }
00324     }
00325     else
00326     {
00327         Scheduler::QueueScheduler();
00328     }
00329 }

```

```

00330     CS_EXIT();
00331 }
00332
00333 //-----
00334 void Thread::SetPriorityBase(K_UCHAR ucPriority_)
00335 {
00336     GetCurrent()->Remove(this);
00337
00338     SetCurrent(Scheduler::GetThreadList(
m_ucPriority));
00339
00340     GetCurrent()->Add(this);
00341 }
00342
00343 //-----
00344 void Thread::SetPriority(K_UCHAR ucPriority_)
00345 {
00346     K_BOOL bSchedule = 0;
00347
00348     CS_ENTER();
00349     // If this is the currently running thread, it's a good idea to reschedule
00350     // Or, if the new priority is a higher priority than the current thread's.
00351     if ((g_pstCurrent == this) || (ucPriority_ > g_pstCurrent->
GetPriority()))
00352     {
00353         bSchedule = 1;
00354     }
00355     Scheduler::Remove(this);
00356     CS_EXIT();
00357
00358     m_ucCurPriority = ucPriority_;
00359     m_ucPriority = ucPriority_;
00360
00361     CS_ENTER();
00362     Scheduler::Add(this);
00363     CS_EXIT();
00364
00365     if (bSchedule)
00366     {
00367         if (Scheduler::IsEnabled())
00368         {
00369             CS_ENTER();
00370             Scheduler::Schedule();
00371             #if KERNEL_USE_QUANTUM
00372             // new thread scheduled. Stop current quantum timer (if it exists),
00373             // and restart it for the new thread (if required).
00374             Quantum::RemoveThread();
00375             Quantum::AddThread((Thread*)g_pstNext);
00376             #endif
00377             CS_EXIT();
00378             Thread::ContextSwitchSWI();
00379         }
00380         else
00381         {
00382             Scheduler::QueueScheduler();
00383         }
00384     }
00385 }
00386
00387 //-----
00388 void Thread::InheritPriority(K_UCHAR ucPriority_)
00389 {
00390     SetOwner(Scheduler::GetThreadList(ucPriority_));
00391     m_ucCurPriority = ucPriority_;
00392 }
00393
00394 //-----
00395 void Thread::ContextSwitchSWI()
00396 {
00397     // Call the context switch interrupt if the scheduler is enabled.
00398     if (Scheduler::IsEnabled() == 1)
00399     {
00400         KERNEL_TRACE_1( STR_CONTEXT_SWITCH_1, (K_USHORT)((
Thread*)g_pstNext)->GetID() );
00401         KernelSWI::Trigger();
00402     }
00403 }
00404
00405 #if KERNEL_USE_TIMEOUTS
00406 //-----
00407 Timer *Thread::GetTimer() { return &
m_clTimer; }
00408
00409 //-----
00410 void Thread::SetExpired( K_BOOL bExpired_ ) {
m_bExpired = bExpired_; }
00411

```



```

00030 #endif
00031 #define __FILE_ID__      THREADLIST_CPP
00032
00033 //-----
00034 void ThreadList::SetPriority(K_UCHAR ucPriority_)
00035 {
00036     m_ucPriority = ucPriority_;
00037 }
00038
00039 //-----
00040 void ThreadList::SetFlagPointer( K_UCHAR *pucFlag_)
00041 {
00042     m_pucFlag = pucFlag_;
00043 }
00044
00045 //-----
00046 void ThreadList::Add(LinkListNode *node_) {
00047     CircularLinkList::Add(node_);
00048     CircularLinkList::PivotForward();
00049
00050     // We've specified a bitmap for this threadlist
00051     if (m_pucFlag)
00052     {
00053         // Set the flag for this priority level
00054         *m_pucFlag |= (1 << m_ucPriority);
00055     }
00056 }
00057
00058 //-----
00059 void ThreadList::Add(LinkListNode *node_, K_UCHAR *pucFlag_,
K_UCHAR ucPriority_) {
00060     // Set the threadlist's priority level, flag pointer, and then add the
00061     // thread to the threadlist
00062     SetPriority(ucPriority_);
00063     SetFlagPointer(pucFlag_);
00064     Add(node_);
00065 }
00066
00067 //-----
00068 void ThreadList::Remove(LinkListNode *node_) {
00069     // Remove the thread from the list
00070     CircularLinkList::Remove(node_);
00071
00072     // If the list is empty...
00073     if (!m_pstHead)
00074     {
00075         // Clear the bit in the bitmap at this priority level
00076         if (m_pucFlag)
00077         {
00078             *m_pucFlag &= ~(1 << m_ucPriority);
00079         }
00080     }
00081 }
00082
00083 //-----
00084 Thread *ThreadList::HighestWaiter()
00085 {
00086     Thread *pclTemp = static_cast<Thread*>(GetHead());
00087     Thread *pclChosen = pclTemp;
00088
00089     K_UCHAR ucMaxPri = 0;
00090
00091     // Go through the list, return the highest-priority thread in this list.
00092     while(1)
00093     {
00094         // Compare against current max-priority thread
00095         if (pclTemp->GetPriority() >= ucMaxPri)
00096         {
00097             ucMaxPri = pclTemp->GetPriority();
00098             pclChosen = pclTemp;
00099         }
00100
00101         // Break out if this is the last thread in the list
00102         if (pclTemp == static_cast<Thread*>(GetTail()))
00103         {
00104             break;
00105         }
00106
00107         pclTemp = static_cast<Thread*>(pclTemp->GetNext());
00108     }
00109     return pclChosen;
00110 }

```

15.101 /home/moslevin/Project/R1/kernel/timer.cpp File Reference

Timer implementations.

```
#include "kerneltypes.h"
#include "mark3cfg.h"
#include "timer.h"
#include "timerlist.h"
#include "timerscheduler.h"
#include "kerneltimer.h"
#include "threadport.h"
#include "kerneldebug.h"
#include "quantum.h"
```

Macros

- `#define __FILE_ID__ TIMER_CPP`
File ID used in kernel trace calls.

15.101.1 Detailed Description

Timer implementations.

Definition in file [timer.cpp](#).

15.102 timer.cpp

```
00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00022 #include "kerneltypes.h"
00023 #include "mark3cfg.h"
00024
00025 #include "timer.h"
00026 #include "timerlist.h"
00027 #include "timerscheduler.h"
00028 #include "kerneltimer.h"
00029 #include "threadport.h"
00030 #include "kerneldebug.h"
00031 #include "quantum.h"
00032
00033 //-----
00034 #if defined __FILE_ID__
00035     #undef __FILE_ID__
00036 #endif
00037 #define __FILE_ID__      TIMER_CPP
00038
00039 #if KERNEL_USE_TIMERS
00040
00041 //-----
00042 void Timer::Start( K_BOOL bRepeat_, K_ULONG ulIntervalMs_,
00043     TimerCallback_t pfCallback_, void *pvData_ )
00044 {
00045     SetIntervalMSeconds(ulIntervalMs_);
00046     m_ulTimerTolerance = 0;
00047     m_pfCallback = pfCallback_;
00048     m_pvData = pvData_;
00049     if (!bRepeat_)
```

```

00049     {
00050         m_ucFlags = TIMERLIST_FLAG_ONE_SHOT;
00051     }
00052     else
00053     {
00054         m_ucFlags = 0;
00055     }
00056     m_pclOwner = Scheduler::GetCurrentThread();
00057     TimerScheduler::Add(this);
00058 }
00059
00060 //-----
00061 void Timer::Start( K_BOOL bRepeat_, K_ULONG ulIntervalMs_,
00062                  K_ULONG ulToleranceMs_, TimerCallback_t pfCallback_, void *pvData_ )
00063 {
00064     m_ulTimerTolerance = MSECONDS_TO_TICKS(ulToleranceMs_);
00065     Start(bRepeat_, ulIntervalMs_, pfCallback_, pvData_);
00066 }
00067 //-----
00068 void Timer::Stop()
00069 {
00070     TimerScheduler::Remove(this);
00071 }
00072
00073 //-----
00074 void Timer::SetIntervalTicks( K_ULONG ulTicks_ )
00075 {
00076     m_ulInterval = ulTicks_;
00077 }
00078
00079 //-----
00081 //-----
00082 void Timer::SetIntervalSeconds( K_ULONG ulSeconds_ )
00083 {
00084     m_ulInterval = SECONDS_TO_TICKS(ulSeconds_);
00085 }
00086
00087 //-----
00088 void Timer::SetIntervalMSeconds( K_ULONG ulMSeconds_ )
00089 {
00090     m_ulInterval = MSECONDS_TO_TICKS(ulMSeconds_);
00091 }
00092
00093 //-----
00094 void Timer::SetIntervalUSeconds( K_ULONG ulUSeconds_ )
00095 {
00096     m_ulInterval = USECONDS_TO_TICKS(ulUSeconds_);
00097 }
00098
00099 //-----
00100 void Timer::SetTolerance(K_ULONG ulTicks_)
00101 {
00102     m_ulTimerTolerance = ulTicks_;
00103 }
00104
00105 #endif
00106

```

15.103 /home/moslevin/Project/R1/kernel/timerlist.cpp File Reference

Implements timer list processing algorithms, responsible for all timer tick and expiry logic.

```

#include "kerneltypes.h"
#include "mark3cfg.h"
#include "timerlist.h"
#include "kerneltimer.h"
#include "threadport.h"
#include "kerneldebug.h"
#include "quantum.h"

```

Macros

- #define `__FILE_ID__` TIMERLIST_CPP

File ID used in kernel trace calls.

15.103.1 Detailed Description

Implements timer list processing algorithms, responsible for all timer tick and expiry logic.

Definition in file [timerlist.cpp](#).

15.104 timerlist.cpp

```

00001 /*=====
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00023 #include "kerneltypes.h"
00024 #include "mark3cfg.h"
00025
00026 #include "timerlist.h"
00027 #include "kerneltimer.h"
00028 #include "threadport.h"
00029 #include "kerneldebug.h"
00030 #include "quantum.h"
00031
00032 //-----
00033 #if defined __FILE_ID__
00034     #undef __FILE_ID__
00035 #endif
00036 #define __FILE_ID__      TIMERLIST_CPP
00037
00038 #if KERNEL_USE_TIMERS
00039
00040 //-----
00041 TimerList TimerScheduler::m_clTimerList;
00042
00043 //-----
00044 void TimerList::Init(void)
00045 {
00046     m_bTimerActive = 0;
00047     m_ulNextWakeup = 0;
00048 }
00049
00050 //-----
00051 void TimerList::Add(Timer *pclListNode_)
00052 {
00053     #if KERNEL_TIMERS_TICKLESS
00054         K_BOOL bStart = 0;
00055         K_LONG lDelta;
00056     #endif
00057
00058     CS_ENTER();
00059
00060     #if KERNEL_TIMERS_TICKLESS
00061         if (GetHead() == NULL)
00062         {
00063             bStart = 1;
00064         }
00065     #endif
00066
00067     pclListNode_>ClearNode();
00068     DoubleLinkedList::Add(pclListNode_);
00069
00070     // Set the initial timer value
00071     pclListNode_>m_ulTimeLeft = pclListNode_>m_ulInterval;
00072
00073     #if KERNEL_TIMERS_TICKLESS
00074         if (!bStart)
00075         {
00076             // If the new interval is less than the amount of time remaining...
00077             lDelta = KernelTimer::TimeToExpiry() - pclListNode_>
m_ulInterval;

```



```

00078
00079         if (lDelta > 0)
00080         {
00081             // Set the new expiry time on the timer.
00082             m_ulNextWakeup = KernelTimer::SubtractExpiry((
K_ULONG)lDelta);
00083         }
00084     }
00085     else
00086     {
00087         m_ulNextWakeup = pclListNode->m_ulInterval;
00088         KernelTimer::SetExpiry(m_ulNextWakeup);
00089         KernelTimer::Start();
00090     }
00091 #endif
00092
00093     // Set the timer as active.
00094     pclListNode->m_ucFlags |= TIMERLIST_FLAG_ACTIVE;
00095     CS_EXIT();
00096 }
00097
00098 //-----
00099 void TimerList::Remove(Timer *pclLinkListNode_)
00100 {
00101     CS_ENTER();
00102
00103     DoubleLinkedList::Remove(pclLinkListNode_);
00104
00105     #if KERNEL_TIMERS_TICKLESS
00106     if (this->GetHead() == NULL)
00107     {
00108         KernelTimer::Stop();
00109     }
00110 #endif
00111
00112     CS_EXIT();
00113 }
00114
00115 //-----
00116 void TimerList::Process(void)
00117 {
00118     #if KERNEL_TIMERS_TICKLESS
00119     K_ULONG ulNewExpiry;
00120     K_ULONG ulOvertime;
00121     K_BOOL bContinue;
00122 #endif
00123
00124     Timer *pclNode;
00125     Timer *pclPrev;
00126
00127     #if KERNEL_USE_QUANTUM
00128     Quantum::SetInTimer();
00129 #endif
00130     #if KERNEL_TIMERS_TICKLESS
00131     // Clear the timer and its expiry time - keep it running though
00132     KernelTimer::ClearExpiry();
00133     do
00134     {
00135 #endif
00136         pclNode = static_cast<Timer*>(GetHead());
00137         pclPrev = NULL;
00138
00139         #if KERNEL_TIMERS_TICKLESS
00140         bContinue = 0;
00141         ulNewExpiry = MAX_TIMER_TICKS;
00142 #endif
00143
00144         // Subtract the elapsed time interval from each active timer.
00145         while (pclNode)
00146         {
00147             // Active timers only...
00148             if (pclNode->m_ucFlags & TIMERLIST_FLAG_ACTIVE)
00149             {
00150                 // Did the timer expire?
00151                 #if KERNEL_TIMERS_TICKLESS
00152                 if (pclNode->m_ulTimeLeft <= m_ulNextWakeup)
00153                 #else
00154                 pclNode->m_ulTimeLeft--;
00155                 if (0 == pclNode->m_ulTimeLeft)
00156                 #endif
00157                 {
00158                     // Yes - set the "callback" flag - we'll execute the callbacks later
00159                     pclNode->m_ucFlags |= TIMERLIST_FLAG_CALLBACK;
00160
00161                     if (pclNode->m_ucFlags & TIMERLIST_FLAG_ONE_SHOT)
00162                     {
00163                         // If this was a one-shot timer, deactivate the timer.

```

```

00164         pclNode->m_ucFlags |= TIMERLIST_FLAG_EXPIRED;
00165         pclNode->m_ucFlags &= ~TIMERLIST_FLAG_ACTIVE;

00166     }
00167     else
00168     {
00169         // Reset the interval timer.
00170         // I think we're good though...
00171         pclNode->m_ulTimeLeft = pclNode->
m_ulInterval;
00172
00173     #if KERNEL_TIMERS_TICKLESS
00174         // If the time remaining (plus the length of the tolerance interval)
00175         // is less than the next expiry interval, set the next expiry interval.
00176         K_ULONG ulTmp = pclNode->m_ulTimeLeft + pclNode->
m_ulTimerTolerance;
00177
00178         if (ulTmp < ulNewExpiry)
00179         {
00180             ulNewExpiry = ulTmp;
00181         }
00182     #endif
00183     }
00184     }
00185     #if KERNEL_TIMERS_TICKLESS
00186     else
00187     {
00188         // Not expiring, but determine how K_LONG to run the next timer interval for.
00189         pclNode->m_ulTimeLeft -= m_ulNextWakeup;
00190         if (pclNode->m_ulTimeLeft < ulNewExpiry)
00191         {
00192             ulNewExpiry = pclNode->m_ulTimeLeft;
00193         }
00194     }
00195     #endif
00196     }
00197     pclNode = static_cast<Timer*>(pclNode->GetNext());
00198 }
00199
00200 // Process the expired timers callbacks.
00201 pclNode = static_cast<Timer*>(GetHead());
00202 while (pclNode)
00203 {
00204     pclPrev = NULL;
00205
00206     // If the timer expired, run the callbacks now.
00207     if (pclNode->m_ucFlags & TIMERLIST_FLAG_CALLBACK)
00208     {
00209         // Run the callback. these callbacks must be very fast...
00210         pclNode->m_pfCallback( pclNode->m_pclOwner, pclNode->
m_pvData );
00211         pclNode->m_ucFlags &= ~TIMERLIST_FLAG_CALLBACK;
00212
00213         // If this was a one-shot timer, let's remove it.
00214         if (pclNode->m_ucFlags & TIMERLIST_FLAG_ONE_SHOT)
00215         {
00216             pclPrev = pclNode;
00217         }
00218     }
00219     pclNode = static_cast<Timer*>(pclNode->GetNext());
00220
00221     // Remove one-shot-timers
00222     if (pclPrev)
00223     {
00224         Remove(pclPrev);
00225     }
00226 }
00227
00228 #if KERNEL_TIMERS_TICKLESS
00229 // Check to see how much time has elapsed since the time we
00230 // acknowledged the interrupt...
00231 ulOvertime = KernelTimer::GetOvertime();
00232
00233 if( ulOvertime >= ulNewExpiry ) {
00234     m_ulNextWakeup = ulOvertime;
00235     bContinue = 1;
00236 }
00237
00238 // If it's taken longer to go through this loop than would take us to
00239 // the next expiry, re-run the timing loop
00240 } while (bContinue);
00241
00242 // This timer elapsed, but there's nothing more to do...
00243 // Turn the timer off.
00244 if (ulNewExpiry >= MAX_TIMER_TICKS)
00245 {

```

```
00248         KernelTimer::Stop();
00249     }
00250     else
00251     {
00252         // Update the timer with the new "Next Wakeup" value, plus whatever
00253         // overtime has accumulated since the last time we called this handler
00254
00255         m_ulNextWakeup = KernelTimer::SetExpiry(ulNewExpiry +
ulOvertime);
00256     }
00257 #endif
00258 #if KERNEL_USE_QUANTUM
00259     Quantum::ClearInTimer();
00260 #endif
00261 }
00262
00263
00264 #endif //KERNEL_USE_TIMERS
```

15.105 /home/moslevin/Project/R1/kernel/tracebuffer.cpp File Reference

Kernel trace buffer class definition.

```
#include "kerneltypes.h"
#include "tracebuffer.h"
#include "mark3cfg.h"
#include "writebuf16.h"
#include "kerneldebug.h"
```

15.105.1 Detailed Description

Kernel trace buffer class definition.

Definition in file [tracebuffer.cpp](#).

15.106 tracebuffer.cpp

```

00001 /*
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 =====*/
00014 #include "kerneltypes.h"
00015 #include "tracebuffer.h"
00016 #include "mark3cfg.h"
00017 #include "writebuf16.h"
00018 #include "kerneldebug.h"
00019
00020
00021 #if KERNEL_USE_DEBUG && !KERNEL_AWARE_SIMULATION
00022 //-----
00023 WriteBuffer16 TraceBuffer::m_clBuffer;
00024 volatile K_USHORT TraceBuffer::m_usIndex;
00025 K_USHORT TraceBuffer::m_auBuffer[ (TRACE_BUFFER_SIZE/sizeof(K_USHORT)) ];
00026
00027 //-----
00028 void TraceBuffer::Init()
00029 {
00030     m_clBuffer.SetBuffers(m_auBuffer, TRACE_BUFFER_SIZE/sizeof(K_USHORT));
00031     m_usIndex = 0;
00032 }
00033
00034 //-----
00035 K_USHORT TraceBuffer::Increment()

```

```

00040 {
00041     return m_usIndex++;
00042 }
00043
00044 //-----
00045 void TraceBuffer::Write( K_USHORT *pusData_, K_USHORT usSize_ )
00046 {
00047     // Pipe the data directly to the circular buffer
00048     m_clBuffer.WriteData(pusData_, usSize_);
00049 }
00050
00051 #endif
00052

```

15.107 /home/moslevin/Project/R1/kernel/writebuf16.cpp File Reference

16 bit circular buffer implementation with callbacks.

```

#include "kerneltypes.h"
#include "writebuf16.h"
#include "kerneldebug.h"
#include "threadport.h"

```

15.107.1 Detailed Description

16 bit circular buffer implementation with callbacks.

Definition in file [writebuf16.cpp](#).

15.108 writebuf16.cpp

```

00001 /*-----
00002
00003
00004
00005
00006
00007
00008
00009 --[Mark3 Realtime Platform]-----
00010
00011 Copyright (c) 2012-2015 Funkenstein Software Consulting, all rights reserved.
00012 See license.txt for more information
00013 -----*/
00020 #include "kerneltypes.h"
00021 #include "writebuf16.h"
00022 #include "kerneldebug.h"
00023 #include "threadport.h"
00024
00025 #if KERNEL_USE_DEBUG && !KERNEL_AWARE_SIMULATION
00026
00027 //-----
00028 void WriteBuffer16::WriteData( K_USHORT *pusBuf_, K_USHORT usLen_ )
00029 {
00030     K_USHORT *apusBuf[1];
00031     K_USHORT ausLen[1];
00032
00033     apusBuf[0] = pusBuf_;
00034     ausLen[0] = usLen_;
00035
00036     WriteVector( apusBuf, ausLen, 1 );
00037 }
00038
00039 //-----
00040 void WriteBuffer16::WriteVector( K_USHORT **ppusBuf_, K_USHORT *pusLen_,
00041     K_UCHAR ucCount_ )
00042 {
00043     K_USHORT usTempHead;
00044     K_UCHAR i;
00045     K_UCHAR j;
00046     K_USHORT usTotalLen = 0;
00047     K_BOOL bCallback = false;

```

```

00047     K_BOOL bRollover = false;
00048     // Update the head pointer synchronously, using a small
00049     // critical section in order to provide thread safety without
00050     // compromising on responsiveness by adding lots of extra
00051     // interrupt latency.
00052
00053     CS_ENTER();
00054
00055     usTempHead = m_usHead;
00056     {
00057         for (i = 0; i < ucCount_; i++)
00058         {
00059             usTotalLen += pusLen_[i];
00060         }
00061         m_usHead = (usTempHead + usTotalLen) % m_usSize;
00062     }
00063     CS_EXIT();
00064
00065     // Call the callback if we cross the 50% mark or rollover
00066     if (m_usHead < usTempHead)
00067     {
00068         if (m_pfCallback)
00069         {
00070             bCallback = true;
00071             bRollover = true;
00072         }
00073     }
00074     else if ((usTempHead < (m_usSize >> 1)) && (m_usHead >= (m_usSize >> 1)))
00075     {
00076         // Only trigger the callback if it's non-null
00077         if (m_pfCallback)
00078         {
00079             bCallback = true;
00080         }
00081     }
00082
00083     // Are we going to roll-over?
00084     for (j = 0; j < ucCount_; j++)
00085     {
00086         K_USHORT usSegmentLength = pusLen_[j];
00087         if (usSegmentLength + usTempHead >= m_usSize)
00088         {
00089             // We need to two-part this... First part: before the rollover
00090             K_USHORT usTempLen;
00091             K_USHORT *pusTmp = &m_pusData[ usTempHead ];
00092             K_USHORT *pusSrc = ppusBuf_[j];
00093             usTempLen = m_usSize - usTempHead;
00094             for (i = 0; i < usTempLen; i++)
00095             {
00096                 *pusTmp++ = *pusSrc++;
00097             }
00098
00099             // Second part: after the rollover
00100             usTempLen = usSegmentLength - usTempLen;
00101             pusTmp = m_pusData;
00102             for (i = 0; i < usTempLen; i++)
00103             {
00104                 *pusTmp++ = *pusSrc++;
00105             }
00106         }
00107         else
00108         {
00109             // No rollover - do the copy all at once.
00110             K_USHORT *pusSrc = ppusBuf_[j];
00111             K_USHORT *pusTmp = &m_pusData[ usTempHead ];
00112             for (K_USHORT i = 0; i < usSegmentLength; i++)
00113             {
00114                 *pusTmp++ = *pusSrc++;
00115             }
00116         }
00117     }
00118
00119     // Call the callback if necessary
00120     if (bCallback)
00121     {
00122         if (bRollover)
00123         {
00124             // Rollover - process the back-half of the buffer
00125             m_pfCallback( &m_pusData[ m_usSize >> 1 ], m_usSize >> 1 );
00126         }
00127         else
00128         {
00129             // 50% point - process the front-half of the buffer
00130             m_pfCallback( m_pusData, m_usSize >> 1 );
00131         }
00132     }
00133 }

```

```
00134 }  
00135  
00136 #endif
```

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