

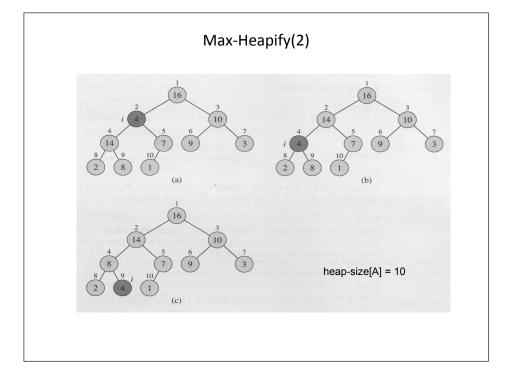
Discussion 4: Week of Sep 26, 2011

#### Outline

- Heaps
- Huffman Encoding
- AVL Trees

## **Heap Review**

- Heapify
  - PercolateUp, TrickleDown
  - O(logN) time for both cases
- BuildHeap
  - O(N) time, proof is non-trivial
- DequeueMax
  - Replace root with last node, then fix down
  - O(logN) time



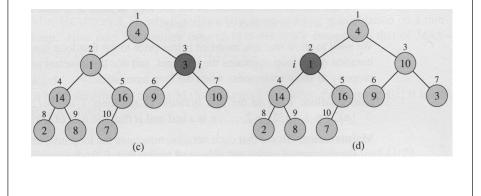
# **Building a heap**

Build-Max-Heap(A)

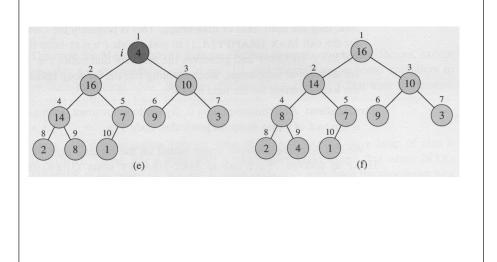
- 1 heap-size[A]  $\leftarrow$  length[A]
- 2 for  $i \leftarrow \lfloor length[A]/2 \rfloor$  downto 1
- 3 do Max-Heapify(A, i)

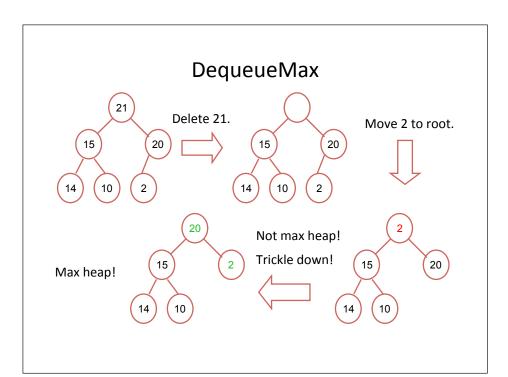
Chapter 6 P.5

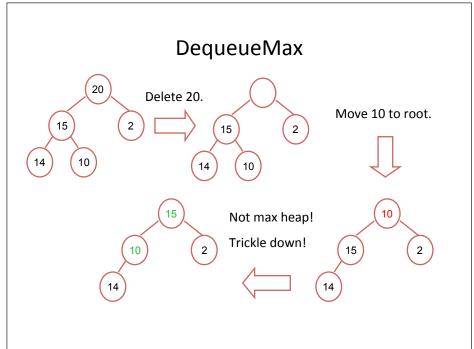
# **Building a heap**

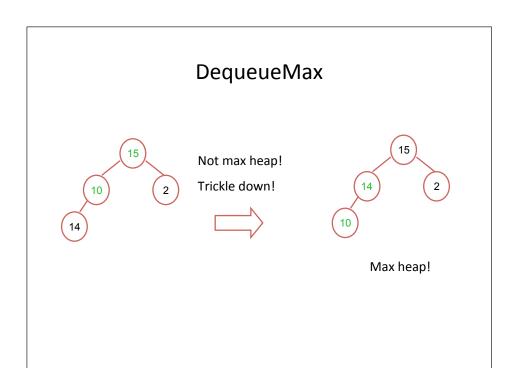


#### **Building a heap**









- Compresses data to use less space.
- Works best on text data, not binary data.
  - Why?

- Compresses data to use less space.
- Works best on text data, not binary data.
  - Huffman encoding uses a character's frequency in order to encode things – the more frequent a character, the better the compression.
  - Binary data looks like: páW.ÈDŽOŸuæ`ç\lß
  - Text data is relatively uniform and predictable.

#### **Huffman Encoding**

Given the text ALIBABA which of the following can be the Huffman prefix code for compressing this text?

A. A=001, B=000, I=01, L=10.

B. A=00, B=01, I=10, L=11.

C. A=0, B=11, I=001, L=000.

D. A=0, B=01, I=100, L=110.

E. A=1, B=01, I=001, L=000.

#### **Huffman Encoding**

- ALIBABA
  - A: 3
  - B: 2
  - -L:1
  - -I:1
- A=001, B=000, I=01, L=10.

- ALIBABA
  - A: 3
  - − B: 2
  - -L:1
  - -I:1
- A=001, B=000, I=01, L=10. No.

#### ALIBABA

- A: 3
- В: 2
- -L:1
- -I:1
- A=00, B=01, I=10, L=11.

#### **Huffman Encoding**

- ALIBABA
  - A: 3
  - B: 2
  - -L:1
  - -I:1
  - A=00, B=01, I=10, L=11. No.

# **Huffman Encoding**

- ALIBABA
  - A: 3
  - В: 2
  - -L:1
  - -I:1
  - A=0, B=11, I=001, L=000.

- ALIBABA
  - A: 3
  - − B: 2
  - -L:1
  - -1:1
  - A=0, B=11, I=001, L=000. Yes.

- ALIBABA
  - A: 3
  - − B: 2
  - -L:1
  - -I:1
  - A=0, B=01, I=100, L=110

## **Huffman Encoding**

- ALIBABA
  - A: 3
  - B: 2
  - -L:1
  - -I:1
  - A=0, B=01, I=100, L=110. Yes.

# **Huffman Encoding**

- ALIBABA
  - A: 3
  - В: 2
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  - -I:1
  - A=1, B=01, I=001, L=000.

- ALIBABA
  - A: 3
  - − B: 2
  - -L:1
  - -I:1
  - A=1, B=01, I=001, L=000. Yes.

#### **AVL Trees**

- A self-balancing binary search tree
- Balance Factor
  - Height(left-sub-tree) Height(right-sub-tree)
- Balance factor of -1, 0, or 1 for each node
- · Rotations needed for maintain balance
  - Single Rotation (LL or RR)
  - Double Rotation(LR or RL)

