

# Lagrange Multiplier Method

Solve the multi-objective minimization problem:

$$\min_x E1(x), E2(x), \dots, Ek(x) \quad (1)$$

where

$$E_i = 0.5x^T H_i x + x^T f_i \quad (2)$$

and  $E_i$  is deemed “more important” than  $E_{i+1}$  (lexicographical ordering).

## Formulation

$$C_0 = H_0 \quad (3)$$

$$d_0 = f_0 \quad (4)$$

$$C_{i+1} = \begin{bmatrix} H_{i+1} & C_i^T \\ C_i & 0 \end{bmatrix} \quad (5)$$

$$d_i = \begin{bmatrix} f_{i+1} \\ d_i \end{bmatrix} \quad (6)$$

Where  $H_{i+1}$  is padded by zeros to match the size of  $C_i$ , and  $f_{i+1}$  is padded by zeros to match the size of  $d_i$ .

If  $C_{i+1}$  is full rank we solve the equation  $C_{i+1}\vec{z} = \vec{d}_{i+1}$ . The final solution is the first  $n$  elements of  $\vec{z}$ . Otherwise we then find the row space of  $C_{i+1}$  and use that as  $C_{i+1}$ . To find the row space of  $C_{i+1}$  we perform QR decomposition on  $C_{i+1}^T$ .

$$PC_{i+1}^T = QR \quad (7)$$

$$\text{row}(C_{i+1}) = \text{col}(C_{i+1}^T) = \text{col}(Q) \quad (8)$$

From this QR decomposition we can find the rank of  $C_{i+1}$ ,  $r = \text{rank}(C_{i+1})$ , which is the index of the first row of all zeros in  $Q$ . The column space of  $Q$  is spanned by the first  $r$  columns of  $Q$ .

$$C_{i+1} \leftarrow (Q_{:,1:r} R_{1:r,1:r})^T \quad (9)$$

Then we shrink  $d_{i+1}$  to include only the rows corresponding to the row space of  $C_{i+1}$

$$d_{i+1} = (Pd)_{:rx} \quad (10)$$

If there are no more Energies we solve the equation  $C_{i+1}\vec{z} = \vec{d}_{i+1}$ . The final solution is the first  $n$  elements of  $\vec{z}$ .

$$\min_{x, \lambda_1, \lambda_2, \dots, \lambda_{i-1}} \frac{1}{2} x^T H_i x + x^T f_i + 0 \cdot \lambda_1 + 0 \cdot \lambda_2 + \dots + 0 \cdot \lambda_{i-1} \quad (11)$$

such that  $C_{i-1} \cdot [x^T \lambda_1 \lambda_2 \dots \lambda_{i-1}]^T = D_{i-1}$  or

$$y = [x^T \lambda_1 \lambda_2 \dots \lambda_{i-1}]^T \quad (12)$$

$$\min_y \frac{1}{2} y^T A y - y^T B \quad (13)$$

such that  $C_{i-1} \cdot [x^T \lambda_1 \lambda_2 \dots \lambda_{i-1}]^T = D_{i-1}$ .

## Example

For example, on the third iteration we have the following:

$$C_2 = \begin{bmatrix} H_2 & 0 & H_1 & H_0^T \\ 0 & 0 & H_0 & 0 \\ H_1^T & H_0^T & 0 & 0 \\ H_0 & 0 & 0 & 0 \end{bmatrix} \quad (14)$$

where each  $H_i$  and 0 are  $n \times n$  matrices

$$d_2 = \begin{bmatrix} f_2 \\ 0 \\ f_1 \\ f_0 \end{bmatrix} \quad (15)$$

where each  $f_i$  and 0 are  $n \times 1$  vectors.

On the fourth iteration we have the following

$$C_3 = \begin{bmatrix} H_3 & 0 & 0 & 0 & H_2^T & 0 & H_1 & H_0 \\ 0 & 0 & 0 & 0 & 0 & 0 & H_0 & 0 \\ 0 & 0 & 0 & 0 & H_1^T & H_0^T & 0 & 0 \\ 0 & 0 & 0 & 0 & H_0 & 0 & 0 & 0 \\ H_2 & 0 & H_1 & H_0^T & 0 & 0 & 0 & 0 \\ 0 & 0 & H_0 & 0 & 0 & 0 & 0 & 0 \\ H_1^T & H_0^T & 0 & 0 & 0 & 0 & 0 & 0 \\ H_0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} \quad (16)$$

$$d_3 = \begin{bmatrix} f_3 \\ 0 \\ 0 \\ 0 \\ f_2 \\ 0 \\ f_1 \\ f_0 \end{bmatrix} \quad (17)$$

As can be seen the  $C$  matrix doubles in size each iteration. The size of  $C_i$  is  $(n2^i) \times (n2^i)$ .

## Limitations

QR factorization on the  $i$ th set of constraints will be very expensive because the number of non-zeros will be  $O(n2^i)$ .