

System programming

# Assemblers

REFERENCES:  
SYSTEM SOFTWARE BY LELAND L. BECK  
SYSTEM SOFTWARE BY SANTANU CHATTOPADHAY  
SYSTEM PROGRAMMING BY SRIMANTAPAL

System programming

## OUTLINE

- > Basic Assembler Functions
- > Machine-dependent Assembler Features
- > Machine-independent Assembler Features
- > Assembler Design Options
- > Appropriate data structures

System programming

## ROLE OF ASSEMBLER

```

graph LR
    SP[Source Program] --> A[Assembler]
    A --> OC[Object Code]
    OC --> L[Linker]
    L --> EC[Executable Code]
    EC --> LO[Loader]
  
```

- > Assembler is the tool to convert an assembly language program into machine language one, understandable by the processor that executes it

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## INTRODUCTION TO ASSEMBLERS

- > Fundamental functions
  - translating mnemonic operation codes to their machine language equivalents
  - assigning machine addresses to symbolic labels
- > Machine dependency
  - different machine instruction formats and codes

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## ASSEMBLERS

- > To Design and implementation of an assembler needs the following:
  - Development of finite specification of the machine language
  - Development of finite specification of the Assembly language
  - Implementation of mapping the that performs
    - > The reading of an assembly language program, statement by statement
    - > The Mapping of each statement into machine language representation and
    - > The generation of machine language module

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## BASIC TASK OF AN ASSEMBLER

- > Noting the external and user defined symbol
  - Passing these symbol to the loader and linker
  - Finding the programs containing the address( or values)
    - > IF not known
  - Generating information for the loader, loading this information into the memory, and placing these values of the symbols in the calling program

The working Principle of an assembly Process

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### LIMITATIONS OF AN ASSEMBLER

- Problems in manual assembly
  - Converting Mnemonic codes
  - Converting decimal numbers
  - Assigning address to program or sub program
  - Counting number of data words
  - Assigning addresses in memory
- General problems in Assembler
  - Designed for particular computers not generic
  - Assembler must be loaded in computer memory...additional memory is required for that

**System programming**

### EXAMPLE 2.1

*Line numbers are not part of the program. They are for reference only.*

5	COPY	START	1000	COPY FILE FROM INPUT TO OUTPUT
10	FIRST	STL	RETADR	SAVE RETURN ADDRESS
15	CLOOP	JSUB	WRREC	READ INPUT RECORD
20		LDA	LENGTH	TEST FOR EOF (LENGTH = 0)
25		COMP	ZERO	
30		JBO	ENDFIL	EXIT IF EOF FOUND
35		JSUB	WRREC	WRITE OUTPUT RECORD
40		J	CLOOP	LOOP
45	ENDFIL	LDA	EOF	INSERT END OF FILE MARKER
50		STA	BUFFER	
55		LDA	THREE	SET LENGTH = 3
60		STA	LENGTH	
65		JSUB	WRREC	WRITE EOF
70		LDL	RETADR	GET RETURN ADDRESS
75		RSUB		RETURN TO CALLER
80	EOF	BYTE	'EOF'	
85	THREE	WORD	3	
90	ZERO	WORD	0	
95	RETADR	RESW	1	
100	LENGTH	RESW	1	LENGTH OF RECORD
105	BUFFER	RESE	4096	4096-BYTE BUFFER AREA

*Forward reference*

*Call subroutine*

*code*

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### Comment line

110				
115				
120				
125	RDREC	LIX	ZERO	CLEAR LOOP COUNTER
130		LDA	ZERO	CLEAR A TO ZERO
135	RLOOP	TD	INPUT	TEST INPUT DEVICE
140		JBO	RLOOP	LOOP UNTIL READY
145		RD	INPUT	READ CHARACTER INTO REGISTER A
150		COMP	ZERO	TEST FOR END OF RECORD ('X'00')
155		JBO	EXIT	EXIT LOOP IF EOR
160		STCH	BUFFER,X	STORE CHARACTER IN BUFFER
165		TIK	MAXLEN	LOOP UNLESS MAX LENGTH
170		JLT	RLOOP	HAS BEEN REACHED
175	EXIT	STX	LENGTH	SAVE RECORD LENGTH
180		RSUB		RETURN TO CALLER
185	INPUT	BYTE	'X'01'	CODE FOR INPUT DEVICE
190	MAXLEN	WORD	4096	
195				

*Indexing mode*

*Hexadecimal number*

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### Subroutine entry point

195				
200				
205				
210	WRREC	LIX	ZERO	CLEAR LOOP COUNTER
215		TD	OUTPUT	TEST OUTPUT DEVICE
220		JBO	WLOOP	LOOP UNTIL READY
225		LDCH	BUFFER,X	GET CHARACTER FROM BUFFER
230		WD	OUTPUT	WRITE CHARACTER
235		TIK	LENGTH	LOOP UNTIL ALL CHARACTERS
240		STL	WLOOP	HAVE BEEN WRITTEN
245		RSUB		RETURN TO CALLER
250	OUTPUT	BYTE	'X'05'	CODE FOR OUTPUT DEVICE
255		END	FIRST	

*Subroutine return point*

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### PURPOSE OF EXAMPLE 2.1 (COPY)

- It is a copy function that reads some records from a specified input device and then copies them to a specified output device
  - Reads a record from the input device (code F1)
  - Copies the record to the output device (code 05)
  - Repeats the above steps until encountering EOF.
  - Then writes EOF to the output device
  - Then call RSUB to return to the caller

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### RDREC AND WRREC

- Data transfer
  - A record is a stream of bytes with a null character (00<sub>16</sub>) at the end.
  - If a record is longer than 4096 bytes, only the first 4096 bytes are copied.
  - EOF is indicated by a zero-length record. (I.e., a byte stream with only a null character.
  - Because the speed of the input and output devices may be different, a buffer is used to temporarily store the record
- Subroutine call and return
  - On line 10, "STL RETADDR" is called to save the return address that is already stored in register L.
  - Otherwise, after calling RD or WR, this COPY cannot return back to its caller.

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## ASSEMBLER DIRECTIVES

- Assembler directives are pseudo instructions
  - ❑ They will not be translated into machine instructions.
  - ❑ They only provide instruction/direction/information to the assembler.
- Basic assembler directives :
  - ❑ **START :**
    - Specify name and starting address for the program
  - ❑ **END :**
    - Indicate the end of the source program, and (optionally) the first executable instruction in the program.

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## ASSEMBLER DIRECTIVES (CONT'D)

- ❑ **BYTE :**
  - Generate character or hexadecimal constant, occupying as many bytes as needed to represent the constant.
- ❑ **WORD :**
  - Generate one-word integer constant
- ❑ **RESB :**
  - Reserve the indicated number of bytes for a data area
- ❑ **RESW :**
  - Reserve the indicated number of words for a data area

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## THE ASSEMBLER'S JOB

- Convert mnemonic operation codes to their machine language codes
- Convert symbolic (e.g., jump labels, variable names) operands to their machine addresses
- Use proper addressing modes and formats to build efficient machine instructions
- Translate data constants into internal machine representations
- Output the object program and provide other information (e.g., for linker and loader)

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## FIGURE 2.1 (PSEUDO CODE)

```

Program copy {
    save return address;
loop: call subroutine RDREC to read one record;
    if length(record)=0 {
        call subroutine WRREC to write EOF;
    } else {
        call subroutine WRREC to write one record;
        goto loop;
    }
    load return address
    return to caller
}
  
```

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## AN EXAMPLE (FIGURE 2.1, CONT.)

```

Subroutine RDREC {
    clear A, X register to 0;
loop: read character from input device to A register
    if not EOR {
        store character into buffer[X];
        X++;
        if X < maximum length
            goto loop;
    }
    store X to length(record);
    return
}
  
```

EOR: character 'x'00'

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## AN EXAMPLE (FIGURE 2.1, CONT.)

```

Subroutine WDREC {
    clear X register to 0;
wloop: get character from buffer[X]
    write character from X to output device
    X++;
    if X < length(record)
        goto wloop;
    return
}
  
```

Line	Loc	Source statement	Object code
5	1000	COPY START 1000	141033
10	1000	FIRST STL RETADR	482039
15	1003	CLOOP JSUB RDREC	001036
20	1006	LDA LENGTH	281030
25	1009	COMP ZERO	301015
30	100C	JBQ ENDFIL	482061
35	100F	JSUB WRREC	3C1003
40	1012	J CLOOP	00102A
45	1015	ENDFIL LDA BOP	0C1039
50	1018	STA BUFFER	00102D
55	101B	LDA THREE	0C1036
60	101E	STA LENGTH	482061
65	1021	JSUB WRREC	081033
70	1024	LDL RETADR	4C0000
75	1027	RSUB	454F46
80	102A	EOF BYTE C' EOF'	000003
85	102D	THREE WORD 3	000000
90	1030	ZERO WORD 0	
95	1033	RETADR RESW 1	
100	1036	LENGTH RESW 1	
105	1039	BUFFER RESB 4096	
110			

110	.	SUBROUTINE TO READ RECORD INTO BUFFER.	
115	.		
120	.		
125	2039	RDREC LDX ZERO	041030
130	203C	LDA ZERO	001030
135	203F	RLOOP TD INPUT	E0205D
140	2042	JBQ RLOOP	30203F
145	2045	RD INPUT	D8205D
150	2048	COMP ZERO	281030
155	204B	JBQ EXIT	302057
160	204E	STCH BUFFER, X	549039
165	2051	TLX MAXLEN	2C205E
170	2054	JLT RLOOP	38203F
175	2057	EXIT STX LENGTH	101036
180	205A	RSUB	4C0000
185	205D	INPUT BYTE X'F1'	F1
190	205E	MAXLEN WORD 4096	001000

195	.	SUBROUTINE TO WRITE RECORD FROM BUFFER	
200	.		
205	.		
210	2061	WRREC LDX ZERO	041030
215	2064	WLOOP TD OUTPUT	E02079
220	2067	JBQ WLOOP	302064
225	206A	LDCH BUFFER, X	509039
230	206D	WD OUTPUT	DC2079
235	2070	TLX LENGTH	2C1036
240	2073	JLT WLOOP	382064
245	2076	RSUB	4C0000
250	2079	OUTPUT BYTE X'05'	05
255		END FIRST	

H   COPY   001000   00107A	Header record:
T   001000   1E   141033   482039   001036   ...	Col. 1 H
T   00101E   15   0C1036   482061   081033   ...	Col. 2-7 Program name
...	Col. 2-7 Starting address of object program (hexadecimal)
T   002079   07   382064   4C0000   05	Col. 8-13 Length of object program in bytes (hexadecimal)
E   001000	Col. 14-19 Length of object program in bytes (hexadecimal)
	Text record:
Col. 1 T	Col. 2-7 Starting address for object code in this record (hexadecimal)
Col. 2-7	Col. 8-9 Length of object code in this record in bytes (hexadecimal)
Col. 10 - 49	Col. 10 - 49 Object code, represented in hexadecimal (10-10+10-10 instructions)
	End record:
Col. 1 E	Col. 2-7 Address of first executable instruction in object program (hexadecimal)
<pre> COPY 00100000107A 00100000107A1036A820390010362810303010134820613C100300102A8C10300102D 001018230C1036A820610236C000003484000003000000 0020391804103600103602050202039A820502810303020573490392C205938203F 0020571C101036A00000F1001000041030E02079202065090390C20792C1036 00207307382064C0000005 001000 </pre>	

Figure 2.3 Object program corresponding to Fig. 2.2.

EXAMPLE OF INSTRUCTION ASSEMBLE			
STCH	BUFFER, X	549039	
8	1	15	
opcode	x	address	
(54) <sub>16</sub>	1 (001) <sub>2</sub>	(039) <sub>16</sub>	m
Forward reference			

Loc	Label	Operator	Operand
1000	FIRST	STL	RETADR
1003	CLOOP	JSUB	RDREC
...	...	...	...
1012	...	J	CLOOP
...	...	...	...
1033	RETADR	RESW	1

## System programming TWO PASS ASSEMBLER

- **Pass 1**
  - ❑ Assign addresses to all statements in the program
  - ❑ Save the values assigned to all **labels** for use in Pass 2
  - ❑ Perform some processing of assembler directives
- **Pass 2**
  - ❑ Assemble instructions
  - ❑ Generate data values defined by BYTE, WORD
  - ❑ Perform processing of assembler directives not done in Pass 1
  - ❑ Write the object program and the assembly listing

## System programming TWO PASS ASSEMBLER

- Read from input line
  - ❑ LABEL, OPCODE, OPERAND

```

graph LR
    SP[Source program] --> P1[Pass 1]
    P1 --> IF[Intermediate file]
    IF --> P2[Pass 2]
    P2 --> OC[Object codes]
    P1 --> OPTAB[OPTAB]
    P1 --> SYMTAB1[SYMTAB]
    P2 --> SYMTAB2[SYMTAB]
  
```

## System programming TWO PASS ASSEMBLER

*READ (Label, opcode, operand)*

*Mnemonic and opcode mappings are referenced from here*

*Label and address mappings enter here*

*Label and address mappings are referenced from here*

## System programming DATA STRUCTURES

- Operation Code Table (OPTAB)
- Symbol Table (SYMTAB)
- Location Counter (LOCCTR)

## System programming OPTAB (OPERATION CODE TABLE)

- **Content**
  - ❑ The mapping between mnemonic and machine code. Also include the instruction format, available addressing modes, and length information.
- **Characteristic**
  - ❑ Static table. The content will never change.
- **Implementation**
  - ❑ Array or hash table. Because the content will never change, we can optimize its search speed.
- In pass 1, OPTAB is used to look up and validate mnemonics in the source program.
- In pass 2, OPTAB is used to translate mnemonics to machine instructions.

## System programming SYMTAB (SYMBOL TABLE)

- **Content**
  - ❑ Include the label name and value (address) for each label in the source program.
  - ❑ Include type and length information (e.g., int64)
  - ❑ With flag to indicate errors (e.g., a symbol defined in two places)
- **Characteristic**
  - ❑ Dynamic table (i.e., symbols may be inserted, deleted, or searched in the table)
- **Implementation**
  - ❑ Hash table can be used to speed up search
  - ❑ Because variable names may be very similar (e.g., LOOP1, LOOP2), the selected hash function must perform well with such non-random keys.
  - ❑ Binary search tree
  - ❑ Linked lists indexed by the first alphabet of the mnemonic

COPY	1000
FIRST	1000
CLOOP	1003
ENDFIL	1015
EOF	1024
THREE	102D
ZERO	1030
RETADR	1033
LENGTH	1036
BUFFER	1039
RDREC	2039



### System programming LOCATION COUNTER (LOCCTR)

- This variable can help in the assignment of addresses.
- It is initialized to the beginning address specified in the START statement.
- After each source statement is processed, the length of the assembled instruction and data area to be generated is added to LOCCTR.
- Thus, when we reach a label in the source program, the current value of LOCCTR gives the address to be associated with that label.

### System programming HOMEWORK #3

SUM	START	4000
FIRST	LDX	ZERO
	LDA	ZERO
LOOP	ADD	TABLE,X
	TIX	COUNT
	JLT	LOOP
	STA	TOTAL
	RSUB	
TABLE	RESW	2000
COUNT	RESW	1
ZERO	WORD	0
TOTAL	RESW	1
	END	FIRST

### System Software Assembler

Pass 1:

```

begin
  read first input line
  if OPCODE = 'START' then
    begin
      save #[OPERAND] as starting address
      initialize LOCCTR to starting address
      write line to intermediate file
      read next input line
    end (if START)
  else
    initialize LOCCTR to 0
    while OPCODE ≠ 'END' do
      begin
        if this is not a comment line then
          begin
            if there is a symbol in the LABEL field then
              begin
                search SYMTAB for LABEL
                if found then
                  set error flag (duplicate symbol)
                else
                  insert (LABEL,LOCCTR) into SYMTAB
              end (if symbol)
            search OPTAB for OPCODE
            if found then
              add 3 (instruction length) to LOCCTR
            else if OPCODE = 'WORD' then
              add 3 to LOCCTR
            else if OPCODE = 'RESW' then
              add 3 * #[OPERAND] to LOCCTR
            else if OPCODE = 'RESEB' then
              add #[OPERAND] to LOCCTR
            else
              begin
                find length of constant in bytes
                add length to LOCCTR
              end (if BYTE)
            else
              set error flag (invalid operation code)
            end (if not a comment)
            write line to intermediate file
          end (while not END)
          write last line to intermediate file
          save (LOCCTR - starting address) as program length
        end (pass 1)
      end
    end
  end

```

Figure 2.4(a) Algorithm for Pass 1 of assembler.

### System Software

Pass 2:

```

begin
  read first input line (from intermediate file)
  if OPCODE = 'START' then
    begin
      write listing line
      read next input line
    end (if START)
  write header record to object program
  initialize first text record
  while OPCODE ≠ 'END' do
    begin
      if this is not a comment line then
        search OPTAB for OPCODE
        if found then
          begin
            if there is a symbol in OPERAND field then
              begin
                search SYMTAB for OPERAND
                if found then
                  store symbol value as operand address
                else
                  begin
                    store 0 as operand address
                    set error flag (undefined symbol)
                  end (if symbol)
                end (if symbol)
                store 0 as operand address
                assemble the object code instruction
                end (if opcode found)
              else if OPCODE = 'BYTE' or 'WORD' then
                convert constant to object code
              if object code will not fit into the current Text record then
                write Text record to object program
                initialize new Text record
              end (if not comment)
              add object code to Text record
            end (if not comment)
            write listing line
            read next input line
          end (while not END)
          write last Text record to object program
          write last listing line
        end (pass 2)
      end
    end
  end

```

Figure 2.4(b) Algorithm for Pass 2 of assembler.

### System programming THE PSEUDO CODE FOR PASS 1

Pass 1:

```

begin
  read first input line
  if OPCODE = 'START' then
    begin
      save #[OPERAND] as starting address
      initialize LOCCTR to starting address
      write line to intermediate file
      read next input line
    end (if START)
  else
    initialize LOCCTR to 0
    while OPCODE ≠ 'END' do
      begin
        if this is not a comment line then
          begin
            if there is a symbol in the LABEL field then
              begin
                search SYMTAB for LABEL
                if found then
                  set error flag (duplicate symbol)
                else
                  insert (LABEL,LOCCTR) into SYMTAB
              end (if symbol)
            search OPTAB for OPCODE
            if found then
              add 3 (instruction length) to LOCCTR
            else if OPCODE = 'WORD' then
              add 3 to LOCCTR
            else if OPCODE = 'RESW' then
              add 3 * #[OPERAND] to LOCCTR
            else if OPCODE = 'RESEB' then
              add #[OPERAND] to LOCCTR
            else
              begin
                find length of constant in bytes
                add length to LOCCTR
              end (if BYTE)
            else
              set error flag (invalid operation code)
            end (if not a comment)
            write line to intermediate file
          end (while not END)
          write last line to intermediate file
          save (LOCCTR - starting address) as program length
        end (pass 1)
      end
    end
  end

```

### System programming

```

search SYMTAB for LABEL
if found then
  set error flag (duplicate symbol)
else
  insert (LABEL,LOCCTR) into SYMTAB
end (if symbol)
search OPTAB for OPCODE
if found then
  add 3 (instruction length) to LOCCTR
else if OPCODE = 'WORD' then
  add 3 to LOCCTR
else if OPCODE = 'RESW' then
  add 3 * #[OPERAND] to LOCCTR
else if OPCODE = 'RESEB' then
  add #[OPERAND] to LOCCTR
else
  begin
    find length of constant in bytes
    add length to LOCCTR
  end (if BYTE)
else
  set error flag (invalid operation code)
end (if not a comment)
write line to intermediate file
end (while not END)
write last line to intermediate file
save (LOCCTR - starting address) as program length
end (pass 1)

```

**System programming**

```

else if OPCODE = 'BYTE' then
begin
    find length of constant in bytes
    add length to LOCCTR
end (if BYTE)
else
    set error flag (invalid operation code)
end (if not a comment)
write line to intermediate file
read next input line
end (while not END)
write last line to intermediate file
save (LOCCTR - starting address) as program length
end (Pass 1)

```

**System programming**

### THE PSEUDO CODE FOR PASS 2

Pass 2:

```

begin
    read first input line (from intermediate file)
    if OPCODE = 'START' then
        begin
            write listing line
            read next input line
        end (if START)
        write Header record to object program
        initialize first Text record
        while OPCODE ≠ 'END' do
            begin
                if this is not a comment line then
                    begin
                        search OPTAB for OPCODE

```

**System programming**

```

if found then
begin
    if there is a symbol in OPERAND field then
        begin
            search SYMTAB for OPERAND
            if found then
                store symbol value as operand address
            else
                begin
                    store 0 as operand address
                    set error flag (undefined symbol)
                end
            end (if symbol)
        else
            store 0 as operand address
            assemble the object code instruction
        end (if opcode found)
    else if OPCODE = 'BYTE' or 'WORD' then
        convert constant to object code

```

**System programming**

```

if object code will not fit into the current Text record then
begin
    write Text record to object program
    initialize new Text record
end
add object code to Text record
end (if not comment)
write listing line
read next input line
end (while not END)
write last Text record to object program
write End record to object program
write last listing line
end (Pass 2)

```

**System programming**

### PASS-1

LOC	Source Statement
1000 COPY	START COPY FILE FROM INPUT TO OUTPUT
1001 FIRST	STL SAVE RETURN ADDRESS
1003 CLOOP	JSUB RDECR READ INPUT RECORD
1006 LDA	LENGTH TEST FOR EOF (LENGTH = 0)
1009 COMP	ZERO
100C JEQ	ENDFIL EXIT IF EOF FOUND
100F JSUB	WRREC WRITE OUTPUT RECORD
1012 J	CLOOP LOOP
1015 ENDFIL	LDA EOF INSERT END OF FILE MARKER
1018 STA	BUFFER
101B LDA	THREE SET LENGTH = 3
101E STA	LENGTH
1021 JSUB	WRREC WRITE EOF
1024 LDH	RETADR GET RETURN ADDRESS
1027 RSUB	RETURN TO CALLER
102A EOF	BYTE C'EOF'
102D THREE	WORD 3
1030 ZERO	WORD 0
1033 RETADR	RESW 1
1036 LENGTH	RESW 1 LENGTH OF RECORD
1039 BUFFER	RESB 4096 4096-BYTE BUFFER AREA
	SUBROUTINE TO READ RECORD INTO BUFFER
2039 RDECR	LDX ZERO CLEAR LOOP COUNTER
	SUBROUTINE TO WRITE RECORD FROM BUFFER
2061 WRREC	LDX ZERO CLEAR LOOP COUNTER
	END FIRST

**System programming**

LOC	Source Statement	LOC	Source Statement
1000 COPY	START COPY FILE FROM INPUT TO OUTPUT	1000 COPY	START COPY FILE FROM INPUT TO OUTPUT
1001 FIRST	STL SAVE RETURN ADDRESS	1001 FIRST	STL SAVE RETURN ADDRESS
1003 CLOOP	JSUB RDECR READ INPUT RECORD	1003 CLOOP	JSUB RDECR READ INPUT RECORD
1006 LDA	LENGTH TEST FOR EOF (LENGTH = 0)	1006 LDA	LENGTH TEST FOR EOF (LENGTH = 0)
1009 COMP	ZERO	1009 COMP	ZERO
100C JEQ	ENDFIL EXIT IF EOF FOUND	100C JEQ	ENDFIL EXIT IF EOF FOUND
100F JSUB	WRREC WRITE OUTPUT RECORD	100F JSUB	WRREC WRITE OUTPUT RECORD
1012 J	CLOOP LOOP	1012 J	CLOOP LOOP
1015 ENDFIL	LDA EOF INSERT END OF FILE MARKER	1015 ENDFIL	LDA EOF INSERT END OF FILE MARKER
1018 STA	BUFFER	1018 STA	BUFFER
101B LDA	THREE SET LENGTH = 3	101B LDA	THREE SET LENGTH = 3
101E STA	LENGTH	101E STA	LENGTH
1021 JSUB	WRREC WRITE EOF	1021 JSUB	WRREC WRITE EOF
1024 LDH	RETADR GET RETURN ADDRESS	1024 LDH	RETADR GET RETURN ADDRESS
1027 RSUB	RETURN TO CALLER	1027 RSUB	RETURN TO CALLER
102A EOF	BYTE C'EOF'	102A EOF	BYTE C'EOF'
102D THREE	WORD 3	102D THREE	WORD 3
1030 ZERO	WORD 0	1030 ZERO	WORD 0
1033 RETADR	RESW 1	1033 RETADR	RESW 1
1036 LENGTH	RESW 1 LENGTH OF RECORD	1036 LENGTH	RESW 1 LENGTH OF RECORD
1039 BUFFER	RESB 4096 4096-BYTE BUFFER AREA	1039 BUFFER	RESB 4096 4096-BYTE BUFFER AREA
	SUBROUTINE TO READ RECORD INTO BUFFER		SUBROUTINE TO READ RECORD INTO BUFFER
2039 RDECR	LDX ZERO CLEAR LOOP CO	2039 RDECR	LDX ZERO CLEAR LOOP CO
	SUBROUTINE TO WRITE RECORD FROM BUFF		SUBROUTINE TO WRITE RECORD FROM BUFF
2061 WRREC	LDX ZERO CLEAR LOOP CO	2061 WRREC	LDX ZERO CLEAR LOOP CO
	END FIRST		END FIRST

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**System programming**

## ASSEMBLER DESIGN

- ❑ Step 11: [Replacement opcode] Replace all mnemonic instruction names with their corresponding codes
- ❑ Step 12: [Replacement symbol table] Resolve all symbolic names with assigned values in the table
- ❑ Step 13: [Error if any] Provide error and diagnostic messages if there are invalid mnemonics or unresolved symbolic names
- ❑ Step 14: Provide a machine code program with additional information indicating those addresses which are either relative or absolute, and those values which are designed as literals
- ❑ Step 15: [Termination]
  - Steps 1-10 First pass
  - Steps 11-15 second pass

**System programming**

## ASSEMBLER FEATURES

- Machine Dependent Assembler Features
  - ❑ instruction formats
  - ❑ addressing modes
  - ❑ program relocation
- Machine Independent Assembler Features
  - ❑ literals
  - ❑ symbol-defining statements
  - ❑ Expressions handling in the program
  - ❑ program blocks
  - ❑ control sections and
  - ❑ program linking

**System programming**

## ASSEMBLER DESIGN CRITERIA

- How many times does a translation process look at the source code to resolve the forward references,
  - ❑ i.e., how many passes are required by a translation process?
    - e.g. two pass required two complete pass
- How can one ensure that the translation process will do everything
  - ❑ i.e. symbol table generation, code generation, and program listing on a minimum number of passes and look at the source code with minimum number of times?
- How does one handle forward references?
- How can the total processing time be minimized?

**System programming**

## MACHINE-DEPENDENT ASSEMBLER FEATURES

**System programming**

## INSTRUCTION FORMAT AND ADDRESSING MODE

- SIC/XE
  - ❑ PC-relative or Base-relative addressing: op m
  - ❑ Indirect addressing: op @m
  - ❑ Immediate addressing: op #c
  - ❑ Extended format: +op m
  - ❑ Index addressing: op m,x
  - ❑ register-to-register instructions
  - ❑ larger memory -> multi-programming (program allocation)

**System programming**

## TRANSLATION

- Register translation
  - ❑ register name (A, X, L, B, S, T, F, PC, SW) and their values (0, 1, 2, 3, 4, 5, 6, 8, 9)
  - ❑ preloaded in SYMTAB
- Address translation
  - ❑ Most register-memory instructions use program counter relative or base relative addressing
  - ❑ Format 3: 12-bit address field
    - base-relative: 0~4095
    - pc-relative: -2048~2047
  - ❑ Format 4: 20-bit address field

System programming				
5	COPY	START	0	COPY FILE FROM INPUT TO OUTPUT
10	FIRST	STL	RETADR	SAVE RETURN ADDRESS
12		LDB	#LENGTH	ESTABLISH BASE REGISTER
13		BASE	LENGTH	
15	CLOOP	+JSUB	RDREC	READ INPUT RECORD
20		LDA	LENGTH	TEST FOR EOF (LENGTH = 0)
25		COMP	#0	
30		JBQ	ENDFIL	EXIT IF EOF FOUND
35		+JSUB	WRREC	WRITE OUTPUT RECORD
40		J	CLOOP	LOOP
45	ENDFIL	LDA	EOF	INSERT END OF FILE MARKER
50		STA	BUFFER	
55		LDA	#3	SET LENGTH = 3
60		STA	LENGTH	
65		+JSUB	WRREC	WRITE EOF
70		J	@RETADR	RETURN TO CALLER
80	EOF	BYTE	C'EOF'	
95	RETADR	RESW	1	
100	LENGTH	RESW	1	LENGTH OF RECORD
105	BUFFER	RESB	4096	4096-BYTE BUFFER AREA
110				

System programming				
115				SUBROUTINE TO READ RECORD INTO BUFFER
120				
125	RDREC	CLEAR	X	CLEAR LOOP COUNTER
130		CLEAR	A	CLEAR A TO ZERO
132		CLEAR	S	CLEAR S TO ZERO
133		+LDT	#4096	
135	RLOOP	TD	INPUT	TEST INPUT DEVICE
140		JBQ	RLOOP	LOOP UNTIL READY
145		RD	INPUT	READ CHARACTER INTO REGISTER A
150		COMPR	A,S	TEST FOR END OF RECORD (X'00')
155		JEQ	EXIT	EXIT LOOP IF BOR
160		STCH	BUFFER,X	STORE CHARACTER IN BUFFER
165		TIXR	T	LOOP UNLESS MAX LENGTH
170		JLT	RLOOP	HAS BEEN REACHED
175	EXIT	STX	LENGTH	SAVE RECORD LENGTH
180		RSUB		RETURN TO CALLER
185	INPUT	BYTE	X'F1'	CODE FOR INPUT DEVICE
195				

System programming				
200				SUBROUTINE TO WRITE RECORD FROM BUFFER
205				
210	WRREC	CLEAR	X	CLEAR LOOP COUNTER
212		LDT	LENGTH	
215	WLOOP	TD	OUTPUT	TEST OUTPUT DEVICE
220		JBQ	WLOOP	LOOP UNTIL READY
225		LDCH	BUFFER,X	GET CHARACTER FROM BUFFER
230		WD	OUTPUT	WRITE CHARACTER
235		TIXR	T	LOOP UNTIL ALL CHARACTERS
240		JLT	WLOOP	HAVE BEEN WRITTEN
245		RSUB		RETURN TO CALLER
250	OUTPUT	BYTE	X'05'	CODE FOR OUTPUT DEVICE
255		END	FIRST	

Line	Loc	Source statement		Object code
5	0000	COPY	START 0	
10	0000	FIRST	STL RETADR	17202D
12	0003		LDB #LENGTH	69202D
13			BASE LENGTH	
15	0006	CLOOP	+JSUB RDREC	4B101036
20	000A		LDA LENGTH	032026
25	000D		COMP #0	290000
30	0010		JBQ ENDFIL	332007
35	0013		+JSUB WRREC	4B10105D
40	0017		J CLOOP	3F2FEC
45	001A	ENDFIL	LDA EOF	032010
50	001D		STA BUFFER	0F2016
55	0020		LDA #3	010003
60	0023		STA LENGTH	0F200D
65	0026		+JSUB WRREC	4B10105D
70	002A		J @RETADR	3E2003
80	002D	EOF	BYTE C'EOF'	454F46
95	0030	RETADR	RESW 1	
100	0033	LENGTH	RESW 1	
105	0036	BUFFER	RESB 4096	

System programming				
110				
115				SUBROUTINE TO READ RECORD INTO BUFFER
120				
125	1036	RDREC	CLEAR X	B410
130	1038		CLEAR A	B400
132	103A		CLEAR S	B440
133	103C		+LDT #4096	75101000
135	1040	RLOOP	TD INPUT	E32019
140	1043		JBQ RLOOP	332FFA
145	1046		RD INPUT	DB2013
150	1049		COMPR A,S	A004
155	104B		JEQ EXIT	332008
160	104E		STCH BUFFER,X	57C003
165	1051		TIXR T	B850
170	1053		JLT RLOOP	3B2FEA
175	1056	EXIT	STX LENGTH	134000
180	1059		RSUB	4F0000
185	105C	INPUT	BYTE X'F1'	F1
195				

System programming				
195				
200				SUBROUTINE TO WRITE RECORD FROM BUFFER
205				
210	105D	WRREC	CLEAR X	B410
212	105F		LDT LENGTH	774000
215	1062	WLOOP	TD OUTPUT	E32011
220	1065		JBQ WLOOP	332FFA
225	1068		LDCH BUFFER,X	53C003
230	106B		WD OUTPUT	DF2008
235	106E		TIXR T	B850
240	1070		JLT WLOOP	3B2FEF
245	1073		RSUB	4F0000
250	1076	OUTPUT	BYTE X'05'	05
255			END FIRST	

### System programming PC-RELATIVE ADDRESSING MODES

➤ PC-relative

□ 10 0000 FIRST STL RETADR 17202D

op(6)	n	I	x	b	p	c	disp(12)
(14) <sub>16</sub>	1	1	0	0	1	0	(02D) <sub>16</sub>

➤ displacement= RETADR - PC = 30-3 = 2D

□ 40 0017 J CLOOP 3F2FEC

op(6)	n	I	x	b	p	c	disp(12)
(FEC) <sub>16</sub>	1	1	0	0	1	0	(3C) <sub>16</sub>

➤ displacement= CLOOP-PC= 6 - 1A= -14= FEC

### System programming BASE-RELATIVE ADDRESSING MODES

➤ Base-relative

□ base register is under the control of the programmer

□ 12 LDB #LENGTH

□ 13 BASE LENGTH

□ 160 104E STCH BUFFER, X 57C003

op(6)	n	I	x	b	p	c	disp(12)
(54) <sub>16</sub>	1	1	1	1	0	0	(003) <sub>16</sub>
(54) <sub>16</sub>	1	1	1	0	1	0	0036-1051= -101B <sub>16</sub>

➤ displacement= BUFFER - B = 0036 - 0033 = 3

□ NOBASE is used to inform the assembler that the contents of the base register no longer be relied upon for addressing

### System programming IMMEDIATE ADDRESS TRANSLATION

➤ Immediate addressing

□ 55 0020 LDA #3 010003

op(6)	n	I	x	b	p	c	disp(12)
(00) <sub>16</sub>	0	1	0	0	0	0	(003) <sub>16</sub>

□ 133 103C +LDT #4096 75101000

op(6)	n	I	x	b	p	c	disp(20)
(74) <sub>16</sub>	0	1	0	0	0	1	(01000) <sub>16</sub>

### System programming IMMEDIATE ADDRESS TRANSLATION (CONT.)

➤ Immediate addressing

□ 12 0003 LDB #LENGTH 69202D

op(6)	n	I	x	b	p	c	disp(12)
(68) <sub>16</sub>	0	1	0	0	1	0	(02D) <sub>16</sub>
(68) <sub>16</sub>	0	1	0	0	0	0	(033) <sub>16</sub>

690033

➤ the immediate operand is the symbol LENGTH

➤ the address of this symbol LENGTH is loaded into register B

➤ LENGTH=0033=PC+displacement=0006+02D

➤ if immediate mode is specified, the target address becomes the operand

### System programming INDIRECT ADDRESS TRANSLATION

➤ Indirect addressing

□ target addressing is computed as usual (PC-relative or BASE-relative)

□ only the n bit is set to 1

□ 70 002A J @RETADR 3E2003

op(6)	n	I	x	b	p	c	disp(12)
(3C) <sub>16</sub>	1	0	0	0	1	0	(003) <sub>16</sub>

➤ TA=RETADR=0030

➤ TA= (PC) +disp=002D+0003

### System programming PROGRAM RELOCATION

➤ Example Fig. 2.1

□ Absolute program, starting address 1000

e.g. 55 101B LDA THREE 00102D

□ Relocate the program to 2000

e.g. 55 101B LDA THREE 00202D

□ Each Absolute address should be modified

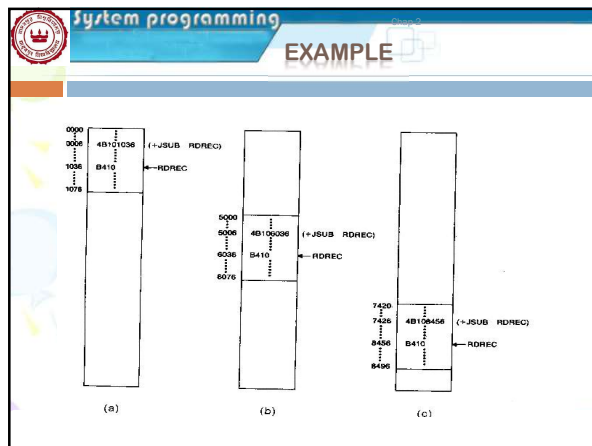
➤ Example Fig. 2.5:

□ Except for absolute address, the rest of the instructions need not be modified

➤ not a memory address (immediate addressing)

➤ PC-relative, Base-relative

□ The only parts of the program that require modification at load time are those that specify direct addresses



**System programming**

### RELOCATABLE PROGRAM

- Modification record
  - Col 1 M
  - Col 2-7 Starting location of the address field to be modified, relative to the beginning of the program
  - Col 8-9 length of the address field to be modified, in half-bytes

**System programming**

### OBJECT CODE

```

H0COPY 000000001077
T0000001D17202D69202D4B1010360320262900003320074B10105D3F2FEC032010
T00001D130F20160100030F200D4B10105D3E2003454F46
T0010361D3410B400B44075101000E32019332FFAD2013A00433200857C0038850
T0010531D382FEA1340004F0000F18410774000E32011332FFA53C003D72008B850
T001070073B2FEF4F000005
M00000705
M00001405
M00002705
R0000000
  
```

End of Sec 2-2

**System programming**

### MACHINE-INDEPENDENT ASSEMBLER FEATURES

- Literals
- Symbol Defining Statement
- Expressions
- Program Blocks
- Control Sections and Program Linking

**System programming**

### LITERALS

- Design idea
  - Let programmers to be able to write the value of a constant operand as a part of the instruction that uses it.
  - This avoids having to define the constant elsewhere in the program and make up a label for it.
- Example
  - e.g. 45 001A ENDFIL LDA =C'EOF' 032010
  - 93 LTORG
  - 002D \* =C'EOF' 454F46
  - e.g. 215 1062 WLOOP TD =X'05' E32011

**System programming**

### LITERALS VS. IMMEDIATE OPERANDS

- Immediate Operands
  - The operand value is assembled as part of the machine instruction
  - e.g. 55 0020 LDA #3 010003
- Literals
  - The assembler generates the specified value as a constant at some other memory location
  - e.g. 45 001A ENDFILLDA =C'EOF' 032010
- Compare (Fig. 2.6)
  - e.g. 45 001A ENDFIL LDA EOF 032010
  - 80 002D EOF BYTE C'EOF'454F46

**System programming**

## LITERAL - IMPLEMENTATION (1/3)

- **Literal pools**
  - ❑ Normally literals are placed into a pool at the end of the program
    - see Fig. 2.10 (END statement)
  - ❑ In some cases, it is desirable to place literals into a pool at some other location in the object program
    - assembler directive `LTORG`
    - reason: keep the literal operand close to the instruction

**System programming**

## LITERAL - IMPLEMENTATION (2/3)

- **Duplicate literals**
  - ❑ e.g. 215      1062    WLOOP    TD    =X'05'
  - ❑ e.g. 230      106B            WD    =X'05'
  - ❑ The assemblers should recognize duplicate literals and store only one copy of the specified data value
    - **Comparison of the defining expression**
      - \* Same literal name with different value, e.g. `LOCCTR=*`
    - **Comparison of the generated data value**
      - \* The benefits of using generate data value are usually not great enough to justify the additional complexity in the assembler

**System programming**

## LITERAL - IMPLEMENTATION (3/3)

- **LITTAB**
  - ❑ literal name, the operand value and length, the address assigned to the operand
- **Pass 1**
  - ❑ build LITTAB with literal name, operand value and length, leaving the address unassigned
  - ❑ when `LTORG` statement is encountered, assign an address to each literal not yet assigned an address
- **Pass 2**
  - ❑ search LITTAB for each literal operand encountered
  - ❑ generate data values using `BYTE` or `WORD` statements
  - ❑ generate modification record for literals that represent an address in the program

**System programming**

## SYMBOL-DEFINING STATEMENTS

- **Labels on instructions or data areas**
  - ❑ the value of such a label is the address assigned to the statement
- **Defining symbols**
  - ❑ symbol `EQU` value
  - ❑ value can be: \* constant, \* other symbol, \* expression
  - ❑ making the source program easier to understand
  - ❑ no forward reference

**System programming**

## SYMBOL-DEFINING STATEMENTS

- **Example 1**
  - ❑ `MAXLEN EQU 4096`
  - ❑ `+LDT #MAXLEN`      `+LDT #4096`
- **Example 2 (Many general purpose registers)**
  - ❑ `BASE EQU R1`
  - ❑ `COUNT EQU R2`
  - ❑ `INDEX EQU R3`
- **Example 3**
  - ❑ `MAXLEN EQU BUFEND-BUFFER`

**System programming**

## ORG (ORIGIN)

- Indirectly assign values to symbols
- Reset the location counter to the specified value
  - `ORG value`
- Value can be: \* constant, \* other symbol, \* expression
- No forward reference
- **Example**
  - ❑ `SYMBOL: 6bytes`
  - ❑ `VALUE: 1word`
  - ❑ `FLAGS: 2bytes`
  - ❑ `LDA VALUE, X`

	SYMBOL	VALUE	FLAGS
STAB			
(100 entries)			



### System programming

## ORG EXAMPLE

- Using EQU statements
  - STAB RESB 1100
  - SYMBOLEQU STAB
  - VALUE EQU STAB+6
  - FLAG EQU STAB+9
- Using ORG statements
  - STAB RESB 1100
  - ORG STAB
  - SYMBOL RESB 6
  - VALUE RESW 1
  - FLAGS RESB 2
  - ORG STAB+1100

### System programming

## EXPRESSIONS

- Expressions can be classified as absolute expressions or relative expressions
  - MAXLEN EQU BUFEND-BUFFER
  - BUFEND and BUFFER both are relative terms, representing addresses within the program
  - However the expression BUFEND-BUFFER represents an absolute value
- When relative terms are paired with opposite signs, the dependency on the program starting address is canceled out; the result is an absolute value

### System programming

## SYMTAB

- None of the relative terms may enter into a multiplication or division operation
- Errors:
  - BUFEND+BUFFER
  - 100-BUFFER
  - 3\*BUFFER
- The type of an expression
  - keep track of the types of all symbols defined in the program

Symbol	Type	Value
RETADR	R	30
BUFFER	R	36
BUFEND	R	1036
MAXLEN	A	1000

### System programming

## EXAMPLE 2.9

SYMTAB		LITTAB			
Name	Value				
COPY	0	C'EOF'	454F46	3	002D
FIRST	0	X'05'	05	1	1076
CLOOP	6				
ENDFIL	1A				
RETADR	30				
LENGTH	33				
BUFFER	36				
BUFEND	1036				
MAXLEN	1000				
RDREC	1036				
RLOOP	1040				
EXIT	1056				
INPUT	105C				
WREC	105D				
WLOOP	1062				

### System programming

## PROGRAM BLOCKS

- Program blocks
  - refer to segments of code that are rearranged within a single object program unit
  - USE [blockname]
  - Default block
  - Example: [Figure 2.11](#)
  - Each program block may actually contain several separate segments of the source program

### System programming

## EXAMPLE 2.11

Line	Source statement	Comment
10	COPY	START 0
15	FIRST	SETUP
20	CLOOP	LENGTH
25		ENDFIL
30		RETADR
35		LENGTH
40		BUFFER
45		BUFEND
50		MAXLEN
55		RDREC
60		RLOOP
65		EXIT
70		INPUT
75		WREC
80		WLOOP
85		END

Figure 2.11 Example of a program with multiple program blocks.

### System programming

## PROGRAM BLOCKS - IMPLEMENTATION

**Pass 1**

- each program block has a separate location counter
- each label is assigned an address that is relative to the start of the block that contains it
- at the end of Pass 1, the latest value of the location counter for each block indicates the length of that block
- the assembler can then assign to each block a starting address in the object program

**Pass 2**

- The address of each symbol can be computed by adding the assigned block starting address and the relative address of the symbol to that block

### Figure 2.12 (a)

Line	Loc/Block	Source statement	Object code
5	0000 0	COPY START 0	172063
10	0000 0	FIRST STL RETADR	4B2021
15	0003 0	CLOOP JSUB RDRRC	032060
20	0006 0	LDA LENGTH	290000
25	0009 0	COMP #0	332006
30	000C 0	JEQ ENDFIL	4B203B
35	000F 0	JSUB WRRRC	032055
40	0012 0	J CLOOP	0F2056
45	0015 0	ENDFIL LDA =C'EOF'	010003
50	0018 0	STA BUFFER	0F2048
55	001B 0	LDA #3	4B2029
60	001E 0	STA LENGTH	3E203F
65	0021 0	JSUB WRRRC	
70	0024 0	J RETADR	
75	0027 1	USE CDATA	
80	0030 1	RESW 1	
85	0033 1	LENGTH RESW 1	
90	0036 1	USE CBLKS	
95	0039 2	RESB 4096	
100	0042 2	BUFFER EQU *	
105	0045 2	BUFEND EQU *	
110	0048 2	MAXLEN EQU BUFEND-BUFFER	
115		SUBROUTINE TO READ RECORD INTO BUFFER	

There is no block number for MAXLEN. This is because MAXLEN is an absolute symbol.

120	0027 0	RDRRC	USE	X	B410
125	0027 0		CLEAR	A	B400
130	0029 0		CLEAR	S	B440
132	002B 0		+LDT	#MAXLEN	75101000
133	002D 0				
135	0031 0	RLOOP	TD	INPUT	E32038
140	0034 0		JEQ	RLOOP	332FFA
145	0037 0		RD	INPUT	D82032
150	003A 0		COMPR	A,S	A008
155	003C 0		JEQ	EXIT	332008
160	003F 0		STCH	BUFFER,X	57A02F
165	0042 0		TTRR	T	B850
170	0044 0		JLT	RLOOP	3B2FEA
175	0047 0	EXIT	STX	LENGTH	13201F
180	004A 0		RSUB		4F0000
183	004C 0		USE	CDATA	
185	004E 0		BYTE	X'F1'	F1
195					
200					
205					
208	004D 0		USE		
210	004D 0	WRRRC	CLEAR	X	B410
212	004F 0		LDT	LENGTH	772017
215	0052 0	WLOOP	LD	=X'05'	E3201B
220	0055 0		JEQ	WLOOP	332FFA
225	0058 0		LCH	BUFFER,X	53A016
230	005B 0		WD	=X'05'	DF2012
235	005E 0		TTRR	T	B850
240	0060 0		JLT	WLOOP	3B2FEF
245	0063 0		RSUB		4F0000
252	0067 1		USE	CDATA	
253	0067 1		LDRG		
255	006A 1		=C'EOF'		454F46
	006A 1		=X'05'		05
			END	FIRST	

0063+3

### System programming

## FIGURE 2.12

Each source line is given a relative address assigned and a block number

Block name	Block number	Address	Length
(default)	0	0000	0066
CDATA	1	0066	000B
CBLKS	2	0071	1000

For absolute symbol, there is no block number

Example

- line 107
- 20 0006 0 LDA LENGTH 032060
- LENGTH=(Block 1)+0003= 0066+0003= 0069
- LOCCTR=(Block 0)+0009= 0009

### System programming

## PROGRAM READABILITY

**Program readability**

- No extended format instructions on lines 15, 35, 65
- No needs for base relative addressing (line 13, 14)
- LTORG is used to make sure the literals are placed ahead of any large data areas (line 253)

**Object code**

- It is not necessary to physically rearrange the generated code in the object program
- see Fig. 2.13, Fig. 2.14

### System programming

## Figure 2.13

```

HCOPY 000000001071
T0000001E172063AB20210320602900003320064B203B3F2FEE0320550F2056010003
T00001E090F20484B20293E203F
T00000271DB410B400B44075101000E32038332FFADB203A00433200857A02FB850
T0000044093B2FEA13201F4F0000
T000006C01F1
T00004D19B410772017E3201B332FFA53A016DF2012B8503B2FEF4F0000
T000006D04454F4605
E0000000

```

Figure 2.14

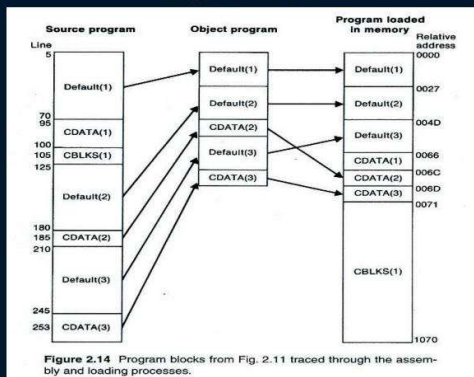


Figure 2.14 Program blocks from Fig. 2.11 traced through the assembly and loading processes.

### System programming CONTROL SECTIONS AND PROGRAM LINKING

- Control Sections
  - are most often used for subroutines or other logical subdivisions of a program
  - the programmer can assemble, load, and manipulate each of these control sections separately
  - instruction in one control section may need to refer to instructions or data located in another section
  - because of this, there should be some means for linking control sections together
  - Fig. 2.15, 2.16

### System programming EXTERNAL DEFINITION AND REFERENCES

- External definition
  - EXTDEF name [, name]
  - EXTDEF names symbols that are defined in this control section and may be used by other sections
- External reference
  - EXTREF name [,name]
  - EXTREF names symbols that are used in this control section and are defined elsewhere
- Example
  - 15 0003 CLOOP +JSUB RDREC 4B100000
  - 160 0017 +STCH BUFFER,X 57900000
  - 190 0028 MAXLEN WORD BUFEND-BUFFER 000000

### System programming IMPLEMENTATION

- The assembler must include information in the object program that will cause the loader to insert proper values where they are required
- Define record
  - Col. 1 D
  - Col. 2-7 Name of external symbol defined in this control section
  - Col. 8-13 Relative address within this control section (hexadecimal)
  - Col.14-73 Repeat information in Col. 2-13 for other external symbols
- Refer record
  - Col. 1 R
  - Col. 2-7 Name of external symbol referred to in this control section
  - Col. 8-73 Name of other external reference symbols

### System programming MODIFICATION RECORD

- Modification record
  - Col. 1 M
  - Col. 2-7 Starting address of the field to be modified (hexadecimal)
  - Col. 8-9 Length of the field to be modified, in half-bytes (hexadecimal)
  - Col.11-16 External symbol whose value is to be added to or subtracted from the indicated field
  - Note: control section name is automatically an external symbol, i.e. it is available for use in Modification records.
- Example
  - Figure 2.17
  - M00000405+RDREC
  - M00000705+COPY

### System programming EXTERNAL REFERENCES IN EXPRESSION

- Earlier definitions
  - required all of the relative terms be paired in an expression (an absolute expression), or that all except one be paired (a relative expression)
- New restriction
  - Both terms in each pair must be relative within the same control section
  - Ex: BUFEND-BUFFER
  - Ex: RDREC-COPY
- In general, the assembler cannot determine whether or not the expression is legal at assembly time. This work will be handled by a linking loader.

System programming

One-pass assemblers  
Multi-pass assemblers  
Two-pass assembler with overlay structure

## ASSEMBLER DESIGN OPTIONS

System programming

## TWO-PASS ASSEMBLER WITH OVERLAY STRUCTURE

- For small memory
  - ❑ pass 1 and pass 2 are never required at the same time
  - ❑ three segments
    - root: driver program and shared tables and subroutines
    - pass 1
    - pass 2
  - ❑ tree structure
  - ❑ overlay program

System programming

## ONE-PASS ASSEMBLERS

- Main problem
  - ❑ forward references
    - data items
    - labels on instructions
- Solution
  - ❑ data items: require all such areas be defined before they are referenced
  - ❑ labels on instructions: no good solution
  - It is possible, although inconvenient, to do so for data items.
  - Forward jump to instruction items cannot be easily eliminated.
    - Insert (label, address to be modified) to SYMTAB
    - Usually, address to be modified is stored in a linked-list

System programming

## FORWARD REFERENCE IN ONE-PASS ASSEMBLER

- For any symbol that has not yet been defined
  1. omit the address translation
  2. insert the symbol into SYMTAB, and mark this symbol undefined
  3. the address that refers to the undefined symbol is added to a list of forward references associated with the symbol table entry
  4. when the definition for a symbol is encountered, the proper address for the symbol is then inserted into any instructions previous generated according to the forward reference list

Name	Type	location	Size	Section id	is-global

Name	Offset(hex) to be corrected

System programming

## ONE-PASS ASSEMBLERS

- Two types of one-pass assembler
  - ❑ load-and-go
    - produces object code directly in memory for immediate execution
  - ❑ the other
    - produces usual kind of object code for later execution

System programming

## LOAD-AND-GO ASSEMBLER

- Characteristics
  - ❑ Useful for program development and testing
  - ❑ Avoids the overhead of writing the object program out and reading it back
  - ❑ Both one-pass and two-pass assemblers can be designed as load-and-go.
  - ❑ However one-pass also avoids the over head of an additional pass over the source program
  - ❑ For a load-and-go assembler, the actual address must be known at assembly time, we can use an absolute program

### System programming LOAD-AND-GO ASSEMBLER (CONT.)

- At the end of the program
  - any SYMTAB entries that are still marked with \* indicate undefined symbols
  - search SYMTAB for the symbol named in the END statement and jump to this location to begin execution
- The actual starting address must be specified at assembly time
- Example
  - Figure 2.18, 2.19

### System programming EXAMPLE PROGRAM FIGURE 2.18

Line	Loc	Source statement	Object code
0	1000	COPY START 1000	
1	1003	EOF BYTE C'BOF'	454F46
2	1006	THREE WORD 3	000003
3	1009	ZERO WORD 0	000000
4	100B	RETADR RESW 1	
5	100C	LENGTH RESW 1	
6	100F	BUFFER RESB 4096	
9			
10	200F	FIRST STL RETADR	141009
15	2012	CLOOP JSUB RDREC	48203D
20	2015	LDA LENGTH	00100C
25	2018	COMP ZERO	281006
30	201B	JRQ ENDFIL	302024
35	201E	JSUB WRREC	482062
40	2021	C CLOOP	302012
45	2024	LDA EOF	001000
50	2027	STA BUFFER	0C100F
55	202A	LDA THREE	001003
60	202D	STA LENGTH	0C100C
65	2030	JSUB WRREC	482062
70	2033	LLL RETADR	081009
75	2036	RSUB	4C0000
110			

### System programming OBJECT CODE IN MEMORY AND SYMTAB FOR LOAD-AND-GO

After scanning line 40 of the program in Fig. 2.18

Memory address	Contents	Symbol	Value
1000	454F4600 00030000 00xxxxxx xxxxxxxx	LENGTH	100C
1010	xxxxxx xx xxxxxx xxxxxxxx xxxxxxxx	RDREC	*
*	*	THREE	1003
*	*	ZERO	1006
2000	xxxxxx xxxxxxxx xxxxxxxx xxxxxxxx	WRREC	*
2010	100946 — 00100C 28100630	EOF	1000
2020	— 302012	ENDFIL	*
*	*	RETADR	1009
*	*	BUFFER	100F
*	*	CLOOP	2012
*	*	FIRST	200F

Diagram showing memory addresses and their contents, with arrows pointing to the symbol table entries. For example, the 'RDREC' symbol points to address 2013, and 'WRREC' points to address 201F.

### System programming OBJECT CODE IN MEMORY AND SYMTAB FOR LOAD-AND-GO

After scanning line 160 of the program in Fig. 2.18

Memory address	Contents	Symbol	Value
1000	454F4600 00030000 00xxxxxx xxxxxxxx	LENGTH	100C
1010	xxxxxx xxxxxxxx xxxxxxxx xxxxxxxx	RDREC	203D
*	*	THREE	1003
*	*	ZERO	1006
2000	xxxxxx xxxxxxxx xxxxxxxx xxxxxxxx	WRREC	*
2010	10094600 3001000C 28100630 202448 —	EOF	1000
2020	— 302012 0010000C 100F0010 — 302024	ENDFIL	2024
2030	48 — 00 10094600 00F10010 00041006	RETADR	1009
2040	00100600 20393020 43082039 28100630	BUFFER	100F
2050	— 480 0F	CLOOP	2012
*	*	FIRST	200F
*	*	MAXLEN	203A
*	*	INPUT	2039
*	*	EXIT	*
*	*	FLCOP	2043

Diagram showing memory addresses and their contents, with arrows pointing to the symbol table entries. For example, the 'RDREC' symbol points to address 203D, and 'WRREC' points to address 201F.

### System programming PRODUCING OBJECT CODE

- When external working-storage devices are not available or too slow (for the intermediate file between the two passes)
- Solution:
  - When definition of a symbol is encountered, the assembler must generate another Tex record with the correct operand address
  - The loader is used to complete forward references that could not be handled by the assembler
  - The object program records must be kept in their original order when they are presented to the loader
- Example: Figure 2.20

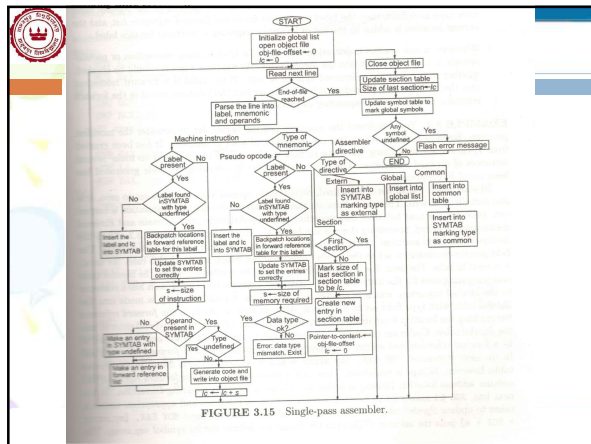
### System programming ONE-PASS WITH OUTPUT FILE

```

H0COPY 00100000107A
T00100009454F46000003000000
T00200F151410094800000100C2810063000004800003C2012
T00201C022024
T002024190010000C100F0010030C100C480000810094C0000F1001000
T00201302203D
T00203D1E041006001006E02039302043D8203928100630000054900F2C203A382043
T0020502205B
T0020580710100C4C000005
T00201F022062
T002031022062
T00206218041006E0206130206550900FDC20612C100C3820654C0000
E00200F

```





System programming

MULTI-PASS  
ASSEMBLERS

- For a two pass assembler, forward references in symbol definition are not allowed:

ALPHA	EQU	BETA
BETA	EQU	DELTA
DELTA	RESW	1

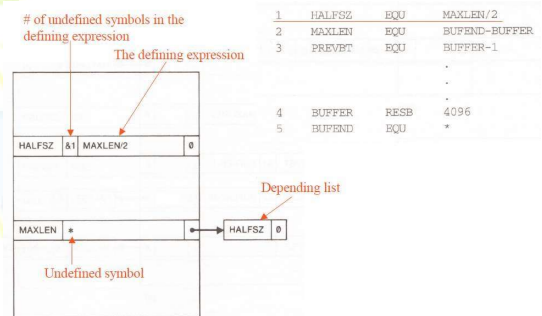
- Symbol definition must be completed in pass 1.
- Prohibiting forward references in symbol definition is not a serious inconvenience.
  - Forward references tend to create difficulty for a person reading the program.

## System programming

- For a forward reference in symbol definition, we store in the SYMTAB:
  - The symbol name
  - The defining expression
  - The number of undefined symbols in the defining expression
- The undefined symbol (marked with a flag \*) associated with a list of symbols depend on this undefined symbol.
- When a symbol is defined, we can recursively evaluate the symbol expressions depending on the newly defined symbol.

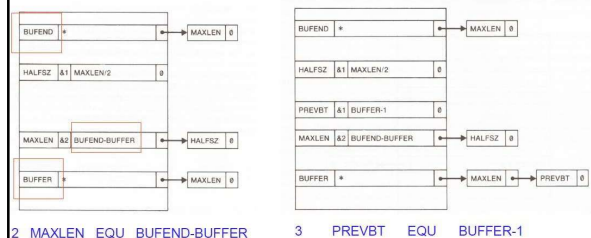
System programming

MULTI-PASS EXAMPLE,  
FIGURE 2.21



System programming

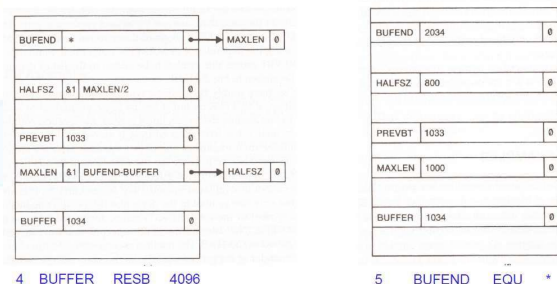
MULTI-PASS EXAMPLE:  
FIGURE 2.21



2 MAXLEN EQU BUFEND-BUFFER

```
3  PREVB  EQU  BUFFER-1
```

System programming



```
4  BUFFER  RESB  4096
```

5 BUFEND EQU