Datamanagement in the Cloud

1. Introduction

Cloud Computing is:

- SaaS: Software as a Service for end users
- PaaS: Platform aaS for developers
- IaaS: Infrastructure aaS for administrators
- a pay-as-you-go-model

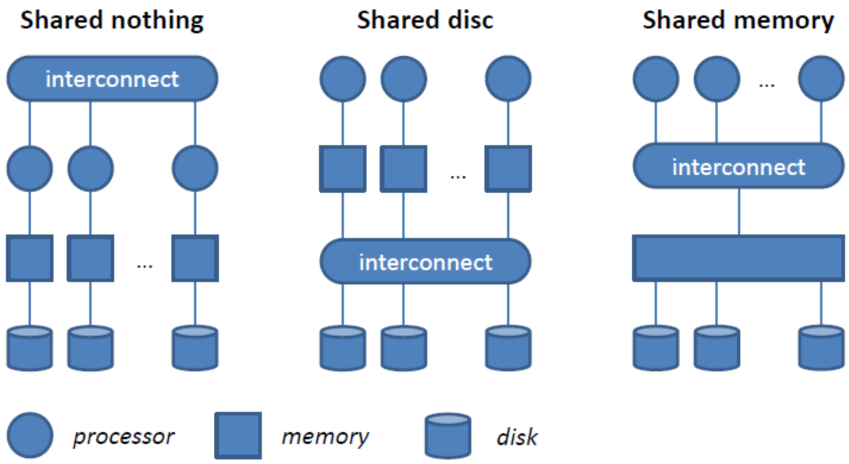
NoSQL

- OLTP (OnLine Transaction Processing)
- OLAP (OnLine Analytical Processing)

Six key features:

1. ability to **horizontally scale simple operation** throughput over many servers
2. ability to **replicate and distribute** (partition) data over many servers
3. simple **call level interface** or protocol
4. **weaker concurrency model** than ACID transactions of most relational DBS
5. efficient use of **distributed indexes and RAM** for data storage
6. ability to **dynamically add new attributes** to data records

1.1. Scaling Databases



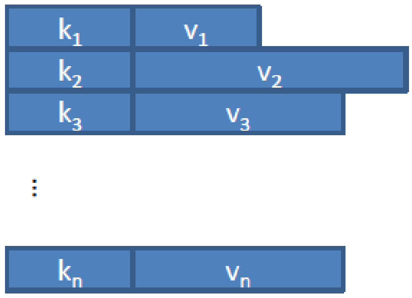
- ***Horizontal:*** Add more nodes
- ***Vertical:*** up-size existing node by adding CPUs, memory
- Move data to where it is needed
- Manage replication for availability and reliability

1.2. Data partitioning

- ***Horizontal:*** distribute groups of tuples of a relation onto different nodes
- ***Vertical:*** distribute groups of columns of a relation onto different nodes
- ***"sharding" -> horizontal***

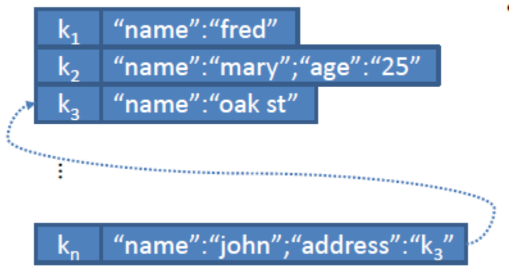
1.3. Data Models

1.3.1. *Key/Value Data Model*



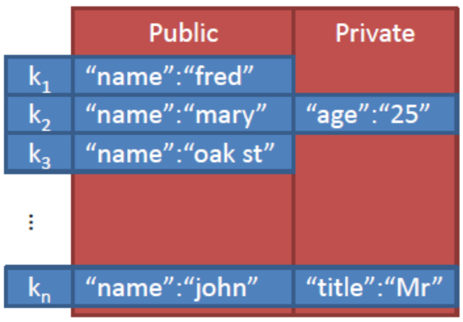
- Interface:
  - put(key, value)
  - get(key): value
- Data Storage:
  - Values (data) are stored based on programmer-defined keys
  - System does not (need to) know about the structure (semantics) of the value
- Queries are expressed in terms of keys
- Indexes are defined over keys:
  - Some systems support secondary indexes over (part of) the value

1.3.2. *Document Data Model*



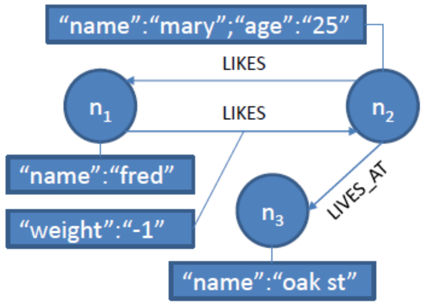
- Interface:
  - set(key, document)
  - get(key): document
  - set(key, name, value)
  - get(key, name): value
- Data storage:
  - Documents (data) is stored based on programmer-defined keys
  - System is aware of the (arbitrary) document structure
  - Support for lists, pointers and nested documents
- Queries expressed in terms of key (or attribute, if index exists)
- Support for key-based indexes and secondary indexes

1.3.3. *Column Family Data Model*



- Interface:
  - define(family)
  - insert(family, key, columns)
  - get(family, key): columns
- Data storage:
  - <name, value, timestamp> triples (so-called columns) are stored based on a column family and key; a column family is similar to a document
  - System is aware of (arbitrary) structure of column family
  - System uses column family information to replicate and distribute data
- Queries are expressed based on key and columns family
- Secondary indexes per column family are typically supported

Graph Data Model



- Interface:
  - create: id
  - get(id)
  - connect(id1, id2): id
  - addAttribute(id, name, value)
  - getAttribute(id, name): value
- Data Storage:
  - Data is stored in terms of nodes and (typed) edges
  - Both nodes and edges can have (arbitrary) attributes
- Queries are expressed based on system ids (if no indexes exists)
- Secondary indexes for nodes and edges are supported
  - Retrieve nodes by attributes and edges by type, start and/or end node, and/or attributes

1.4. Consistency Models

Recap ACID:

- **Atomicity:** "all or nothing" the whole transaction or no transaction are committed
- **Consistency:** transactions never observe or result in inconsistent data
- **Isolation:** transactions are not aware of concurrent transactions
- **Durability:** once committed, the state of a transaction is permanent

1.4.1. CAP Theorem

- Three properties are desirable and expected from a real-world shared-data systems
  - **C:** data consistency
  - **A:** availability
  - **P:** tolerance of network partition
- Only two of these properties can be satisfied by a system at any given time

Criticism

- Asymmetry of CAP properties:
  - Consistency is a property of the system in general
  - Availability is a property of the system only when there is a partition
- There are not three different choices
  - In practice, CA and CP are indistinguishable, since A is only sacrificed when there is a partition
  - Used as an excuse to not bother with consistency
- Other costs to consistency:
  - Overhead of synchronization schemes
  - Latency

1.4.2. PACELC

- If there is a partition **P**, choose availability **A** or consistency **C**
- Else **E** choose latency **L** or consistency **C**

1.4.3. Strong vs. Weak Consistency

Strong consistency

- After an update is committed, each subsequent access will return the update value

Weak consistency

- A number of conditions might need to be met before the updated value is returned
- **Inconsistency window:** period between update and the point in time when every access is guaranteed to return the updated value

1.4.4. Eventual Consistency

- Specific form of weak consistency
- "If no new updates are made, eventually all accesses will return the last updated values"
- In the absence of failures, the maximum size of the **inconsistency window** can be determined based on:
    - Communication delays
    - System load
    - Number of replicas
    - …

Models of Eventual Consistency

- Causal consistency
    - If A communicated to B that it has updates a value, as subsequent access by B will return the updated value, and a write is guaranteed to supersede the earlier write
    - Access by C that has no causal relationship to A is subject to normal eventual consistency rules
- Read-your-writes consistency
    - After updating a value, a process will always read the updated value and never see an older value
- Session consistency
    - Data is accessed in a session where read-your-writes is guaranteed
    - Guarantees do not span over sessions
- Monotonic read consistency
    - If a process has seen a particular value, any subsequent access will never return any pervious value
- Monotonic Write consistency
    - System guarantees to serialize the writes of one process

- Properties can be combined
  - E.g. monotonic read + session-level consistency
  - E.g. monotonic reads + read-your-own-writes

Configurations

- Definitions:
  - **N:** number of nodes that store a replica
  - **W:** number of replicas that need to acknowledge a write operation
  - **R:** number of replicas that are accessed for a read operation
- W+R> N
  - E.g. **synchronous replication** (N=2, W=2, R =1)
  - Write set and read set always overlap
  - Strong consistency can be guaranteed through **quorum protocols**
  - Risk of reduced availability: in basic quorum protocols, operations fail if fewer than the required number of nodes respond, due to node failure
- R=1, W=N
  - Optimized for **read access:** single read will return a value
  - Write operation involves all nodes and risks not to succeed
- R=N, W=1
  - Optimized for **write access:** write operation involves only one node and relies on lazy (epidemic) technique to update other replicas
  - Read operation involves all nodes and returns "latest" value
  - Durability is not guaranteed in presence of failures
- W < (N+1)/2
  - Risk of conflicting writes
- W+R <= N
  - **Weak/eventual consistency**

BASE = **B**asically **A**vailable, **S**oft state, **E**ventual Consistency