# Security of distributed Model Predictive Control under False Data injection

#### Rafael Accácio NOGUEIRA

2022-12-09







https://bit.ly/3g3S6X4











- Electricity Distribution System
- Heat distribution
- Water distribution
- Traffic management (include your problem here)





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"Necessity is the mother of invention"



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- Coupled by constraints
   Technical/ Comfort
- Optimization objectives
   Minimize energy consumption
   Maximize user satisfaction
   Follow a trajectory
- Solution  $\rightarrow$  MPC





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  - Decision variable is the control sequence
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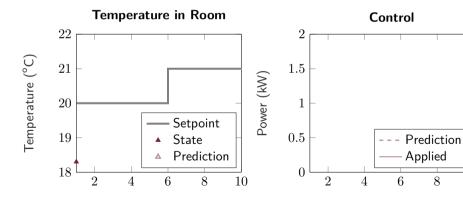
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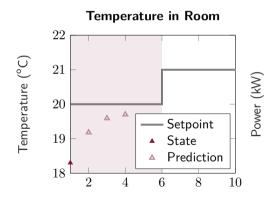


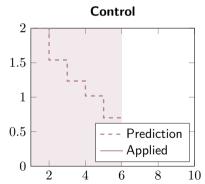


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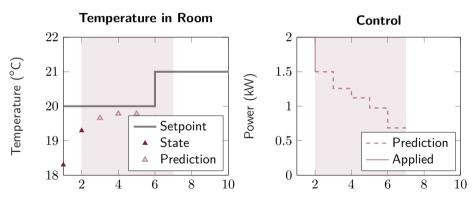






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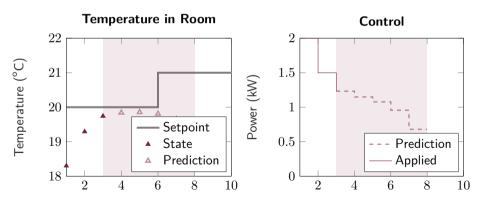
Find optimal control sequence, apply first element





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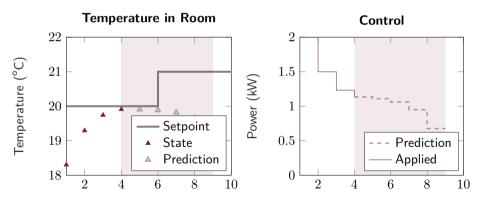
Find optimal control sequence, apply first element, rinse repeat





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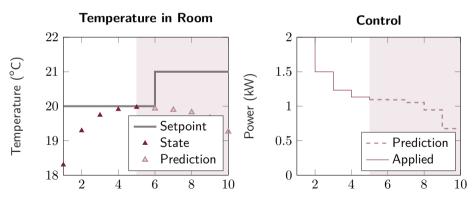
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#### Nothing is perfect

- Problems
  - Complexity of calculation
  - Topology (Geographical distribution)
  - Flexibility (Add/remove parts)
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- Solution: Divide and Conquer (distributed MPC)
  - Break calculation
  - Make systems communicate



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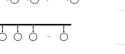




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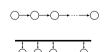






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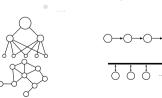








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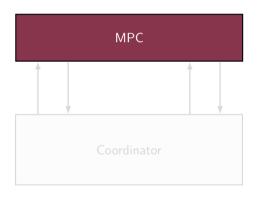






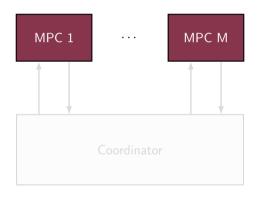






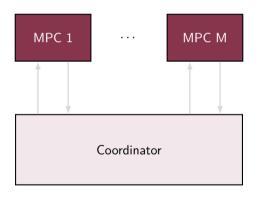
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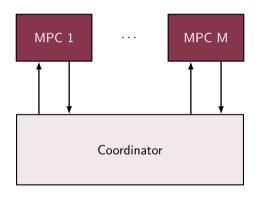
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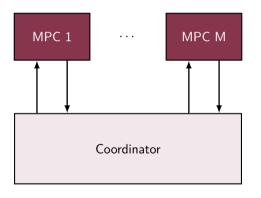
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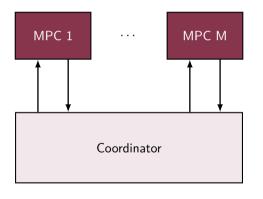
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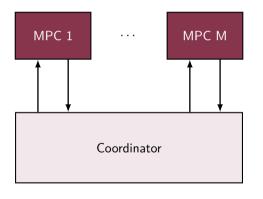




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But what if some agents are ill-intentioned and attack the system?

- How can an agent attack?
- What are the consequences of an attack?
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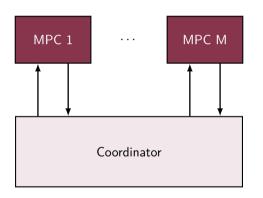


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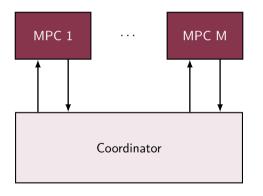


#### Literature



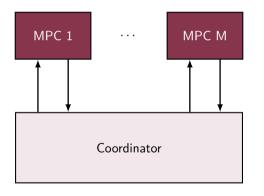
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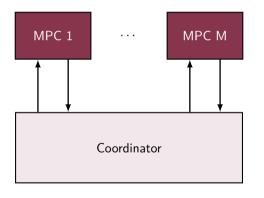
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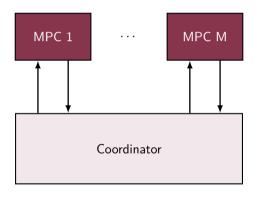
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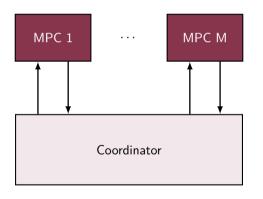




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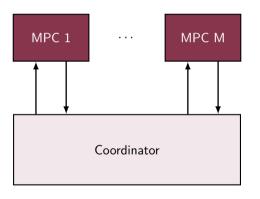


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Deception Attacks



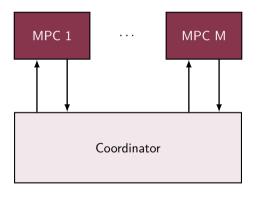
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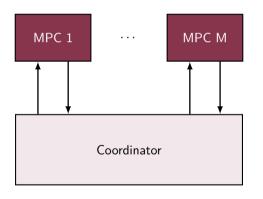
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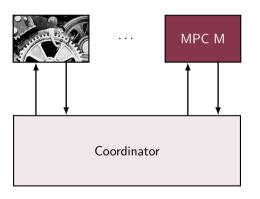


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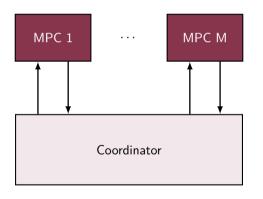


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Liar agent

Deception Attacks (Internal change)

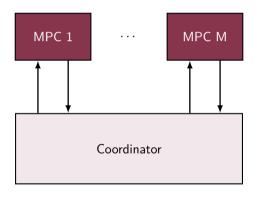




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- What matters is the interface
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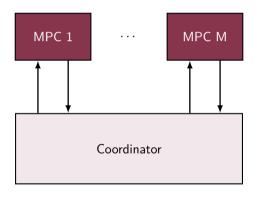






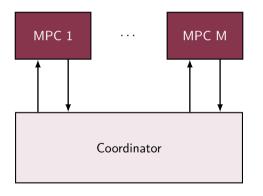
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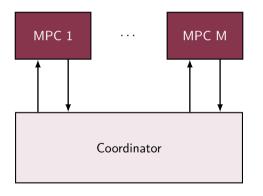
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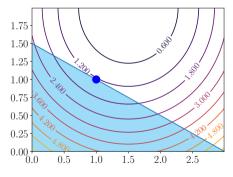


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# Consequence of an attack

- Attack modifies optimization problem
  - Optimum value is shifted

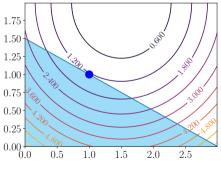


Original minimum.

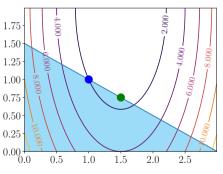


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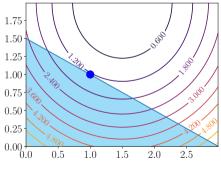


Minimum after attack.

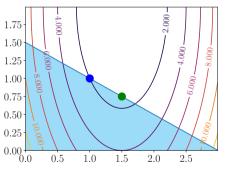


#### Consequence of an attack

- Attack modifies optimization problem
  - Optimum value is shifted



Original minimum.



Minimum after attack.



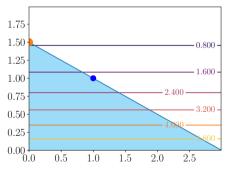
- We can recover by
  - Ignoring attacker
  - Recuperating original behavior (at least trying)



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  - Recuperating original behavior (at least trying)



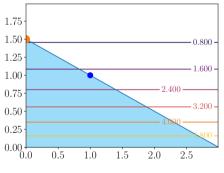
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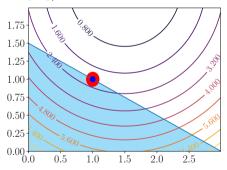
Ignore attacker.



- We can recover by
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  - Recuperating original behavior (at least trying)



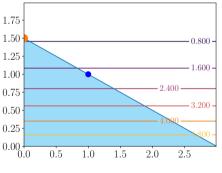
Ignore attacker.



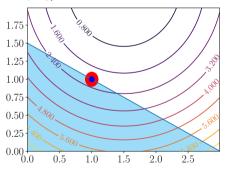
Recover original behavior.



- We can recover by
  - Ignoring attacker
  - Recuperating original behavior (at least trying)



Ignore attacker.



Recover original behavior.



- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
  - Detection/Isolation
  - Mitigation



- Passive (Robust) 1 mode
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- Passive (Robust) 1 mode
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Attack free

- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
  - ① Detection/Isolation
  - Mitigation

Attack free When attack detected



- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
  - ① Detection/Isolation
  - 2 Mitigation

Attack free When attack detected



- Passive (Robust) 1 mode
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```
Attack free
When attack detected
```



- Passive (Robust) 1 mode
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```
Attack free
When attack detected
```



	Decomposition	Resilient/Robust	Detection	Mitigation
[Vel+17a] [Mae+21]	Dual	Robust (Scenario)		NA
[Vel+17b] [Vel+18]	Dual	Robust (f-robust)	NA	NA
[CMI18]	Jacobi-Gauß	-		
[Ana+18] [Ana+19] [Ana+20]	Dual	Resilient	Analyt./Learn.	Disconnect (Robustness)
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### State of art

### Security dMPC

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- Vulnerabilities in distributed MPC based on Primal Decomposition
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- 3 Resilient Primal Decomposition-based dMPC using Artificial Scarcity
- 4 Conclusion
  - and pielded [NBG21] (SysTol'21)
  - o yielded [Nog+22] (NecSys'22)



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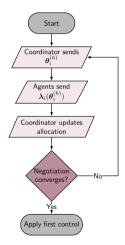


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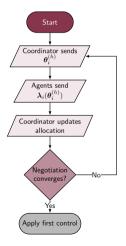


1 Vulnerabilities in distributed MPC based on Primal Decomposition What is the Primal Decomposition? How can an agent attack? Consequences





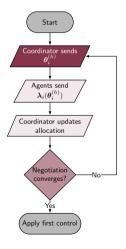








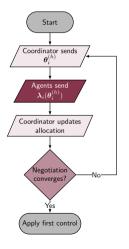
#### or Quantity Decomposition | or Resource Allocation



### Allocation $\theta_i$

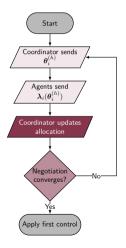


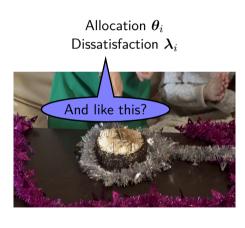




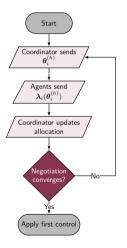


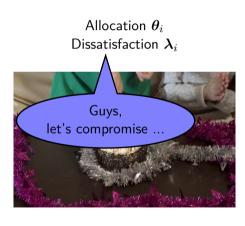




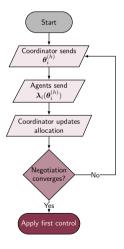












Allocation  $\boldsymbol{\theta}_i$ Dissatisfaction  $\boldsymbol{\lambda}_i$ 





- Objective is sum of local ones
- Constraints couple variables
- lacktriangle Allocate  $heta_i$  for each agent
- They solve local problems and
- $\odot$  Send dual variable  $\lambda_i$
- Allocation is updated (respecting global constraint)

$$egin{array}{ll} & \min & J_i(oldsymbol{x}_i, oldsymbol{u}_i) \ & ext{s. t.} & oldsymbol{h}_i(oldsymbol{x}_i, oldsymbol{u}_i) \leq oldsymbol{ heta}_i : \lambda_i \end{array}$$

$$\boldsymbol{\theta}[k]^{(p+1)} = \text{Proj}^{8}(\boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)})$$



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$$egin{aligned} & \min_{m{u}_1,...,m{u}_M} & \sum_{i\in \mathcal{M}} J_i(m{x}_i,m{u}_i) \ & ext{s.t.} & \sum_{i\in \mathcal{M}} m{h}_i(m{x}_i,m{u}_i) \leq m{u}_{\mathsf{total}} \ & & & \downarrow & \mathsf{For \ each} \ i \in \mathcal{M} \end{aligned}$$

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- $oldsymbol{0}$  Allocate  $oldsymbol{ heta}_i$  for each agent
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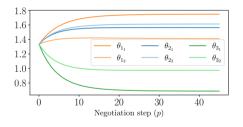
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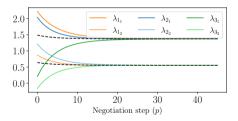
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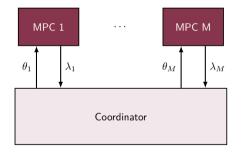
### Example

### Until everybody is equally dissatisfied



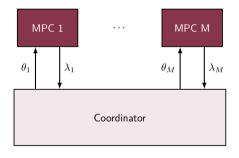






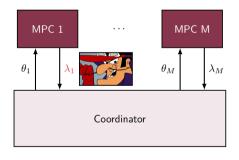
- $\lambda_i$  is the only interface
- ullet  $\lambda_i$  depends on local parameters
- ullet Malicious agent modifies  $oldsymbol{\lambda}_i$





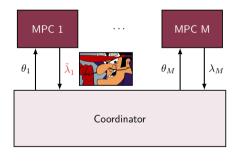
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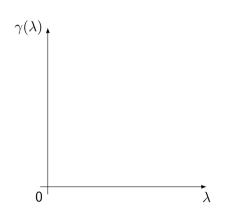


- ullet  $oldsymbol{\lambda}_i$  is the only interface
- ullet  $oldsymbol{\lambda}_i$  depends on local parameters
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$$ilde{oldsymbol{\lambda}}_i = \gamma_i(oldsymbol{\lambda}_i)$$



#### Liar, Liar, Pants of fire

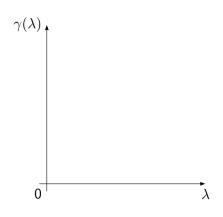


- $\lambda \ge 0$  means dissatisfaction
- $\lambda = 0$  means complete satisfaction

- Attacker satisfied only if it really is
- Attacker is greedy  $\gamma(\lambda) > \lambda$
- Attack is monotonically increasing  $\lambda_b > \lambda_a \rightarrow \gamma(\lambda_b) > \gamma(\lambda_a)$
- Invertible
- If  $\tilde{\lambda}_i = T_i[k]\lambda_i \to \exists T_i[k]^{-1}$



#### Liar, Liar, Pants of fire

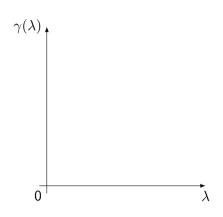


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#### Liar, Liar, Pants of fire

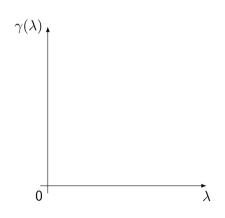


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#### Liar, Liar, Pants of fire



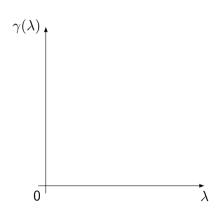
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### Assumptions

Attacker satisfied only if it really is

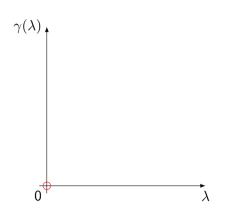
• 
$$\gamma(\lambda) = 0 \to \lambda = 0$$

- Attacker is greedy  $\gamma(\lambda) > \lambda$
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- Invertible
- If  $\tilde{\lambda}_i = T_i[k]\lambda_i \to \exists T_i[k]^{-1}$



#### Liar, Liar, Pants of fire



- $\lambda \ge 0$  means dissatisfaction
- ullet  $\lambda=0$  means complete satisfaction

### Assumptions

Attacker satisfied only if it really is

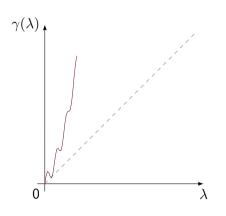
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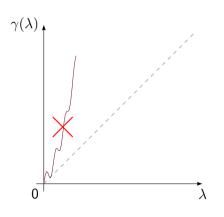
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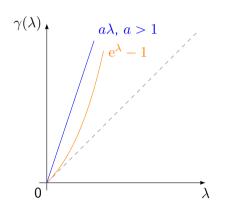
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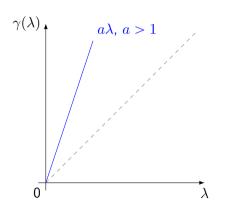
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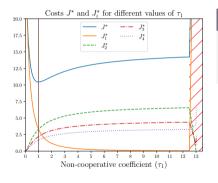


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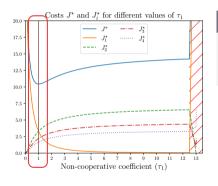
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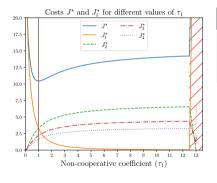
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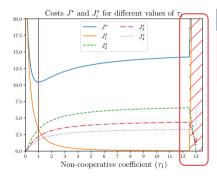
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- Yes! (At least in some cases)



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### Outline

Resilient Primal Decomposition-based dMPC for deprived systems
 Analyzing deprived systems
 Building an algorithm
 Applying mechanism



- Unconstrained Solution  $\mathring{\boldsymbol{U}}_{i}^{\star}[k]$
- $\bar{\Gamma}_i \mathring{\boldsymbol{U}}_i^{\star}[k] \geq \theta_i[k] \rightarrow \text{Scarcity}$ 
  - Solution projected onto boundary
  - Same as with equality constraints<sup>2</sup>

$$\begin{array}{ll} \underset{\boldsymbol{U}_{i}[k]}{\operatorname{minimize}} & \frac{1}{2} \left\| \boldsymbol{U}_{i}[k] \right\|_{H_{i}}^{2} + \boldsymbol{f}_{i}[k]^{T} \boldsymbol{U}_{i}[k] \\ \text{subject to} & \bar{\Gamma}_{i} \boldsymbol{U}_{i}[k] \leq \boldsymbol{\theta}_{i}[k] : \boldsymbol{\lambda}_{i}[k] \end{array}$$



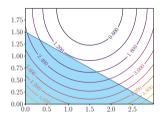
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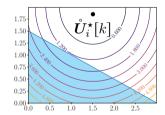
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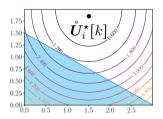
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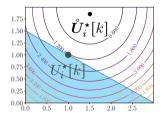
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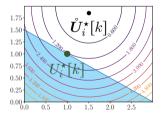
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- Quadratic local problems
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- Solution is analytical and affine

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- ullet Change o Probably an Attack! Let's take advantage of this!



#### Assumption

### We know nominal $\bar{P}_i$

• If we estimate  $\hat{P}_i[k]$  and  $\hat{\tilde{s}}_i[k]$  such as:

$$\tilde{\boldsymbol{\lambda}}_i[k] = -\hat{\tilde{P}}_i[k]\boldsymbol{\theta}_i - \hat{\tilde{\boldsymbol{s}}}_i[k]$$

• If 
$$\left\|\hat{\tilde{P}}_i[k] - \bar{P}_i \right\|_{\scriptscriptstyle E} > \epsilon_P o \mathsf{Attack}$$

• Ok, but how can we estimate  $\widehat{\tilde{P}}_i[k]$ ?



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Rafael Accácio Nogueira

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- Challenge: Online estimation during negotiation fails
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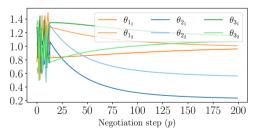
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# Classification of mitigation techniques

- Active (Resilient)
  - Detection/Isolation
  - Mitigation



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- Active (Resilient)
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# Mitigation mechanism

#### Reconstructing $\lambda_i$

- Now, we have  $\hat{\tilde{P}}_i[k]$ 
  - Since  $\tilde{P}_i[k] = T_i[k]\bar{P}_i$
  - We can recover  $T_i[k]^{-1}$

$$\widehat{T_i[k]^{-1}} = P_i \widehat{\tilde{P}}_i[k]^{-1}$$

• Reconstruct  $\lambda_i$ 

$$\hat{\boldsymbol{\lambda}}_i^{ ext{rec}} = -\bar{P}_i \boldsymbol{\theta}_i - \widehat{T_i[k]^{-1}} \hat{\tilde{\boldsymbol{s}}}_i[k]$$

Choose adequate version for coordination

$$oldsymbol{\lambda}_i^{ ext{nod}} = egin{cases} ilde{oldsymbol{\lambda}}_i, & ext{if attack detected} \ ilde{oldsymbol{\lambda}}_i, & ext{otherwise} \end{cases}$$



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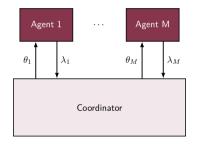
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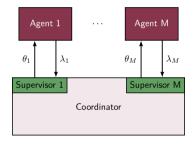




- Supervise exchanges by inquiring the agents
- Estimate how they will behave

- Detect which agents are non-cooperative
- Reconstruct  $\lambda_i$  and use in negotiation

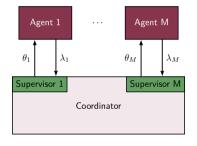




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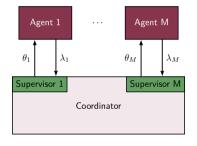




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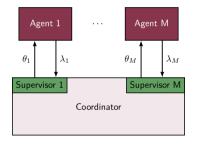




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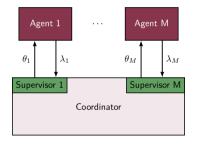




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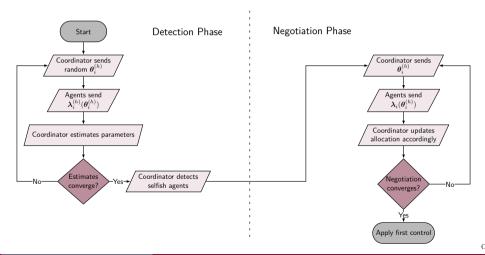




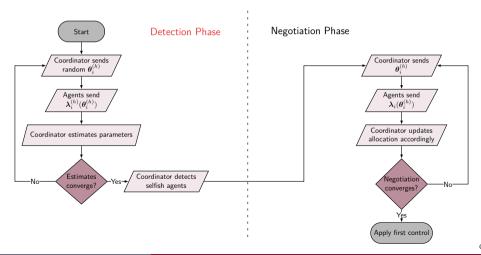
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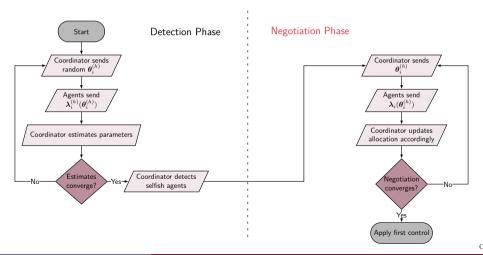




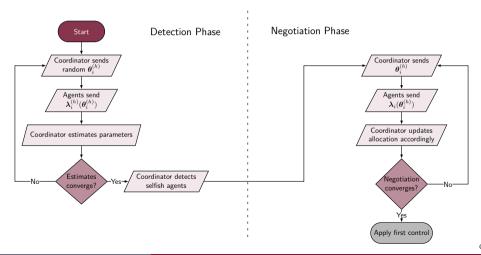




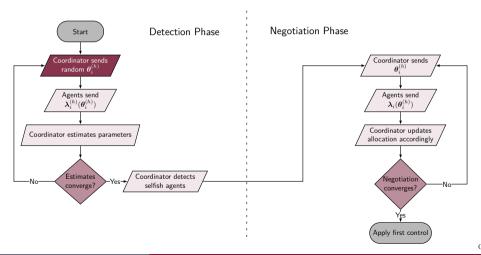


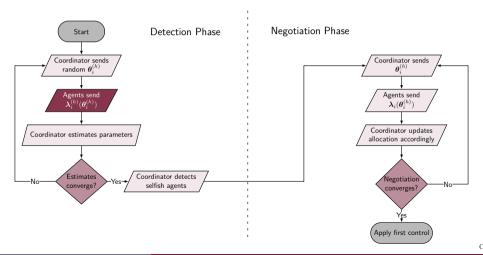




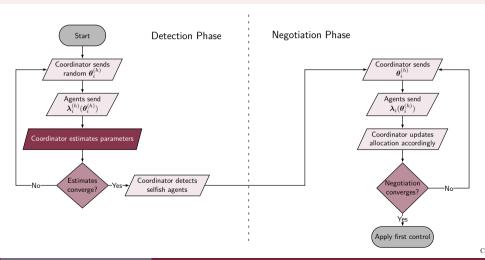




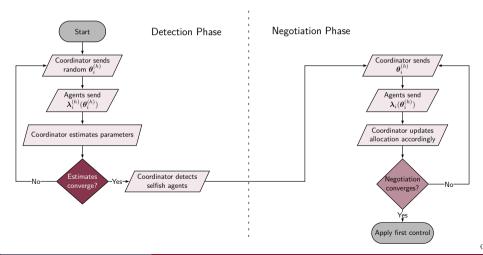




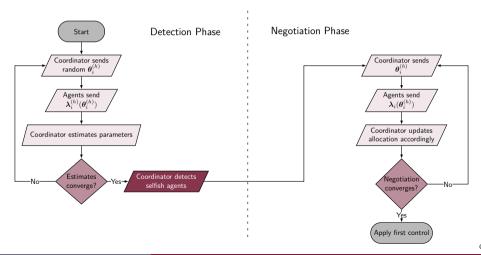




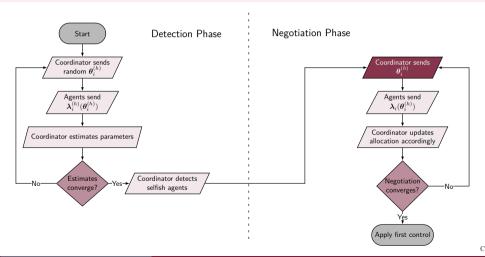


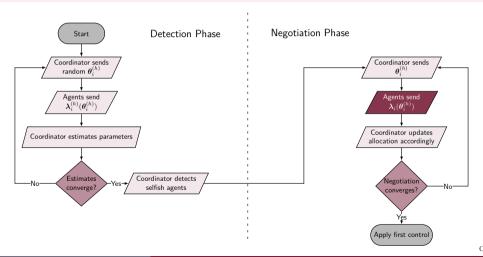




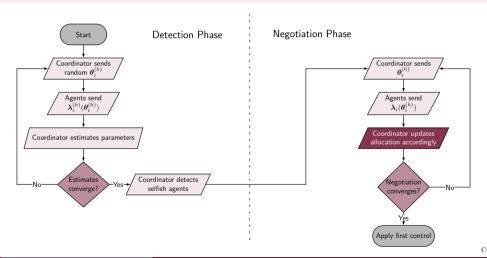


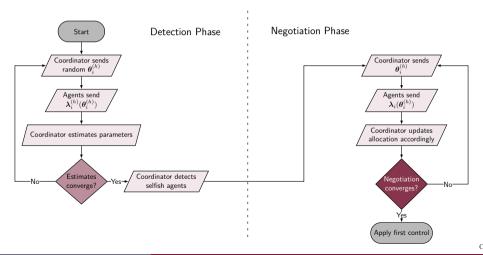




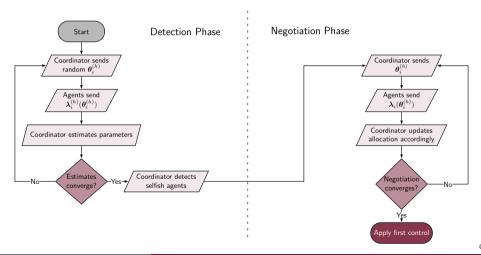




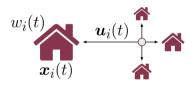






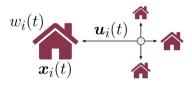






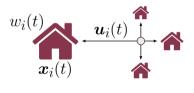
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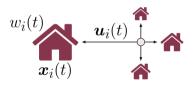
### District Heating Network (4 Houses)

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Applying mechanism

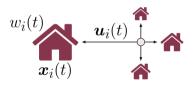
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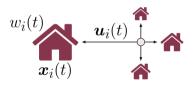
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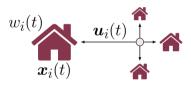
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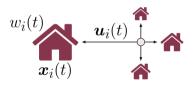
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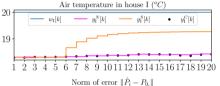
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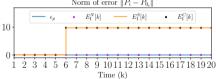


Applying mechanism

### Results

### Temporal

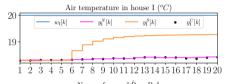


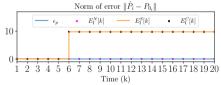


Temperature in house I. Error  $E_I(k)$ .



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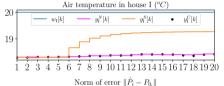


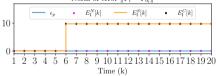
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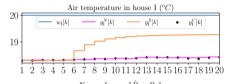


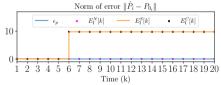
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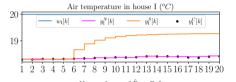


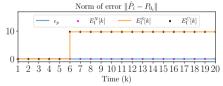
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### Costs

Objective functions  $J_i$  (Normalized error %)

Agent	Selfish	Corrected
ı	-36.3	0.503
Ш	21.671	-0.547
Ш	17.387	-0.004
IV	17.626	-0.09
Global	3.526	0.016



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### Outline

Resilient Primal Decomposition-based dMPC using Artificial Scarcity Relaxing some assumptions Adapting the algorithm Applying mechanism



- Systems are not completely deprived
  - We can't change our constraints to equality ones anymore
  - Nor use the simpler update equation

$$\begin{array}{ll}
\text{minimize} & \frac{1}{2} \| \boldsymbol{U}_i[k] \|_{H_i}^2 + \boldsymbol{f}_i[k]^T \boldsymbol{U}_i[k] \\
\text{subject to} & \bar{\Gamma}_i \boldsymbol{U}_i[k] \leq \boldsymbol{\theta}_i[k] : \boldsymbol{\lambda}_i[k] \\
\boldsymbol{\theta}[k]^{(p+1)} = \operatorname{Proj}^{\mathbb{S}}(\boldsymbol{\theta}[k]^{(p)} + \boldsymbol{\rho}^{(p)} \boldsymbol{\lambda}[k]^{(p)})
\end{array}$$



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$$\boldsymbol{\theta}[k]^{(p+1)} = \text{Proj}^{\mathcal{S}}(\boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)})$$



### Solution for $\lambda_i[k]$

### Instead of having one single affine solution

$$\boldsymbol{\lambda}_i[k] = -P_i\boldsymbol{\theta}_i[k] - \boldsymbol{s}_i[k]$$

Now, we may have multiple (Piecewise affine function)

$$\boldsymbol{\lambda}_i[k] = \begin{cases} -P_i^{(0)}\boldsymbol{\theta}_i[k] - s_i^{(0)}[k], & \text{if } \boldsymbol{\theta}_i[k] \in \mathcal{R}_{\lambda_i}^0 \\ \vdots & \vdots \\ -P_i^{(2^{n_{\text{ineq}}}-1)}\boldsymbol{\theta}_i[k] - s_i^{(2^{n_{\text{ineq}}}-1)}[k], & \text{if } \boldsymbol{\theta}_i[k] \in \mathcal{R}_{\lambda}^{2^{n_{\text{ineq}}}-1} \end{cases}$$



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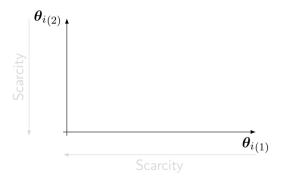
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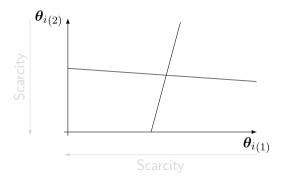


### Solution for $\lambda_i[k]$ (Continued)





#### Solution for $\lambda_i[k]$ (Continued)

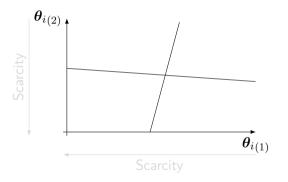


Separation surfaces depend on state and local parameters.

Unknown by the coordinator.

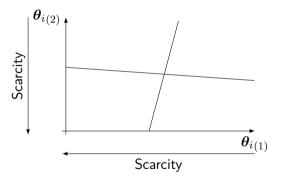


#### Solution for $\lambda_i[k]$ (Continued)



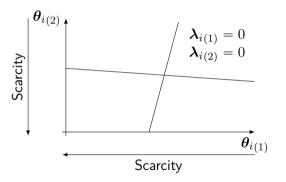


#### Solution for $\lambda_i[k]$ (Continued)



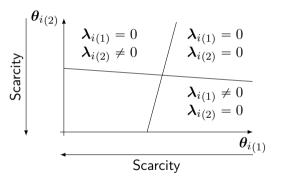


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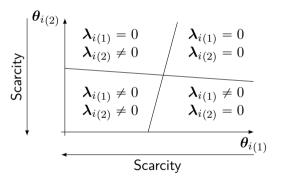


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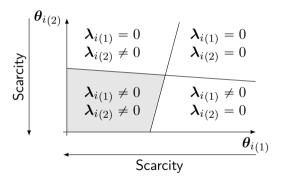


#### Solution for $\lambda_i[k]$ (Continued)





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 Scarcity Sparsity

All constraints active 
$$-P_i^{(0)}\theta_i[k] - s_i^{(0)}[k] \qquad \rightarrow \qquad -P_i\theta_i[k] - s_i[k]$$
 None constraints active 
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Who is it? Who is it?

### Assumption

We know a point  $\overset{\circ}{\theta}_i$  which activates all constraints

$$\theta_{i(2)}$$

$$\lambda_{i(1)} = 0$$

$$\lambda_{i(2)} \neq 0$$

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CentraleSupélec

<sup>4</sup> If we have local constraints, we suppose this point respects them

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CentraleSupélec

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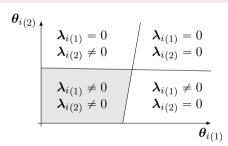
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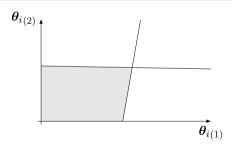


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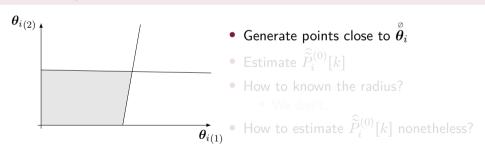
Rafael Accácio Nogueira

### Artificial Scarcity

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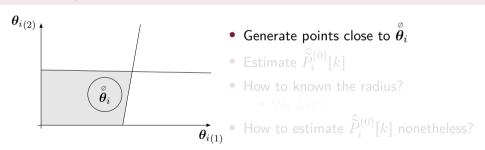
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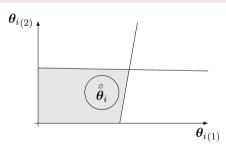


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- ullet Generate points close to  $\stackrel{\circ}{ heta}_i$
- Estimate  $\widehat{\widetilde{P}}_i^{(0)}[k]$ 
  - How to known the radius?
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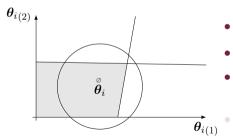
CentraleSupélec

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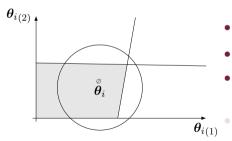
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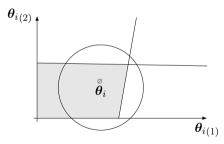
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- Iterative method to estimate parameters of multimodal models<sup>5</sup>
- We give multiple observations  $m{ heta}_i^o[k]$  and  $ilde{m{\lambda}}_i^o[k]$
- At each step we calculate
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- At the end we have
  - Parameters with associated region index
  - Observations with associated region index
- We consult the index associated to  $\overset{\circ}{ heta}_i$
- We recover the associated parameter, i.e.,  $\widehat{\tilde{P}}_i^{(0)}[k]$

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<sup>&</sup>lt;sup>5</sup>Such as our PWA function using some tricks

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CentraleSupélec

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#### Same same, but different

#### Assumption

### We know nominal $ar{P}_i{}^{(0)}$

Detection

$$\left\|\widehat{\widetilde{P}}_{i}^{(0)}[k] - \bar{P}_{i}^{(0)}\right\|_{F} \geqslant \epsilon_{P_{i}^{(0)}}$$

$$\widehat{T_i[k]^{-1}} = \overline{P_i}^{(0)} \widehat{\widetilde{P}_i}^{(0)}[k]^{-1}.$$

$$\hat{oldsymbol{\lambda}}_i = \widehat{T_i[k]^{-1}} ilde{oldsymbol{\lambda}}_i.$$



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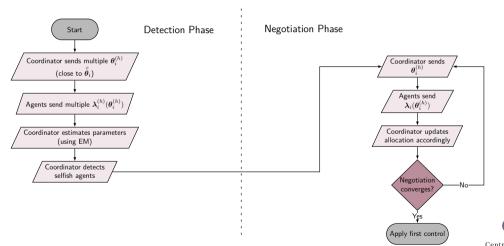
Detection

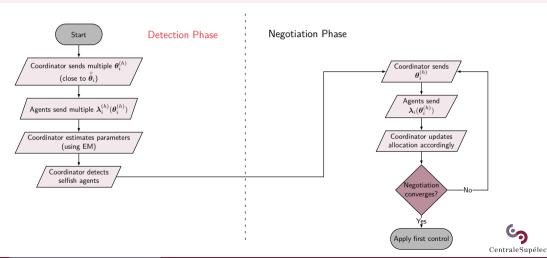
$$\left\| \hat{\tilde{P}}_{i}^{(0)}[k] - \bar{P}_{i}^{(0)} \right\|_{F} \geqslant \epsilon_{P_{i}^{(0)}}$$

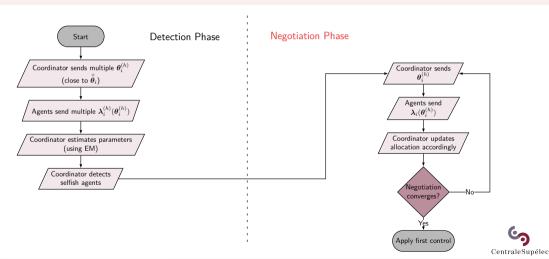
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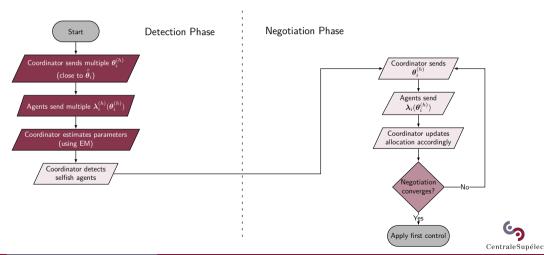
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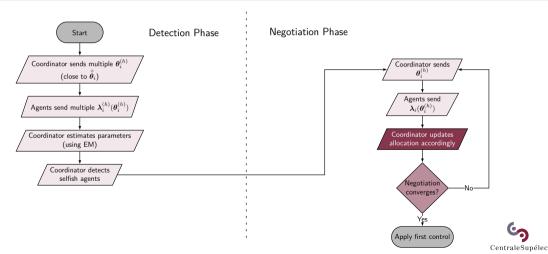




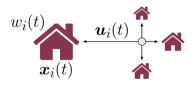








### Example



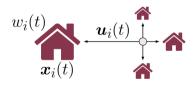
### District Heating Network (4 Houses)

- Houses modeled using 3R-2C
- Not enough power
- Period of 5h  $(T_s = 0.25h)$ 
  - 3 scenarios
    - Nominal
    - Agent I cheats (dMPC)
    - S Agent I cheats (RPdMPC-AS)



Applying mechanism

### Example



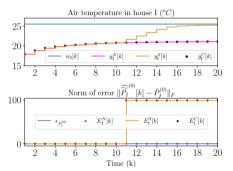
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### Results

### Temporal



Temperature in house I. Error  $E_I(k)$ .



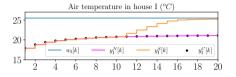


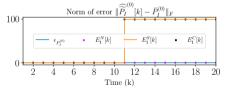




### Results

#### Temporal





Temperature in house I. Error  $E_I(k)$ .

Nominal, S Selflish Corrected





### Results

#### Costs

Objective functions  $J_i$  (Normalized error %)

Agent	Selfish	Corrected
ı	-36.489	-4.12e - 05
Ш	35.813	1.74e - 05
Ш	29.225	2.14e - 05
IV	37.541	1.73e - 05
Global	10.689	-6e - 07



### Too good to be true!

### It's a kind of magic!

- Unfortunately EM is not magic
  - Slow convergence
  - Dependency on initialization
    - No guarantees of achieving global optimal
- Some "solutions":
  - Force some parameters to converge faster (case dependant)
  - Run multiple times with different initialization and pick best
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## Outline



- How can an agent attack? ✓
  - Attacker can change the communication to receive more ressources.
- What are the consequences of an attack? ✓
  - Suboptimality and maybe instability
- Can we mitigate the effects?
  - Yes! By exploring the scarcity of the systems!



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- Insights from the analysis of the solutions of the optimization problems
  - We found some parameters that are constant when there is no cheating
  - The same parameters change when system is attacked
- Exploiting the solution, we find how to invert the cheating function
  - Straightforward if system is scarce
  - If not scarce we try to force it artifically
    - \* However solution is PWA and we need special estimation method in
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- Partial/incremental reconstruction of cheating matrix
- Study of robustness + noise
- Resilient strategy with soft constraints
- Recursive EM (or alternative)
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### Questions? Comments?

Repository https://github.com/Accacio/thesis



Contact rafael.accacio.nogueira@gmail.com



# For Further Reading I



K.J. Åström and B. Wittenmark. <u>Adaptive Control</u>. Addison-Wesley series in electrical and computer engineering: Control engineering. Addison-Wesley, 1989. ISBN: 9780201097207. DOI: 10.1007/978-3-662-08546-2\\_24.



José M Maestre, Rudy R Negenborn, et al.

<u>Distributed Model Predictive Control made easy.</u> Vol. 69. Springer, 2014. ISBN: 978-94-007-7005-8.



Wicak Ananduta et al. "Resilient Distributed Model Predictive Control for Energy Management of Interconnected Microgrids". In:

Optimal Control Applications and Methods 41.1 (2020),
pp. 146–169. DOI: 10.1002/oca.2534. URL: https://onlinelibrary.wiley.com/doi/pdf/10.1002/oca.2534.



# For Further Reading II



José M. Maestre et al. "Scenario-Based Defense Mechanism Against Vulnerabilities in Lagrange-Based Dmpc". In: Control Eng Pract 114 (2021), p. 104879. ISSN: 0967-0661. DOI: 10.1016/j.conengprac.2021.104879.



Rafael Accácio Nogueira et al. "Expectation-Maximization Based Defense Mechanism for Distributed Model Predictive Control". In: <a href="IFAC-PapersOnLine">IFAC-PapersOnLine</a> 55.13 (2022). 9th IFAC Conference on Networked Systems NECSYS 2022, pp. 73–78. ISSN: 2405-8963. DOI: 10.1016/j.ifacol.2022.07.238.



Pablo Velarde et al. "Vulnerabilities in Lagrange-Based Distributed Model Predictive Control". In:

Optimal Control Applications and Methods 39.2 (Sept. 2018), pp. 601–621. DOI: 10.1002/oca.2368.



# For Further Reading III



Wicak Ananduta et al. "Resilient Distributed Energy Management for Systems of Interconnected Microgrids". In: 2018 IEEE Conference on Decision and Control (CDC). 2018, pp. 3159–3164. DOI: 10.1109/CDC.2018.8619548.



Wicak Ananduta et al. "A Resilient Approach for Distributed MPC-Based Economic Dispatch in Interconnected Microgrids". In: 2019 18th European Control Conference (ECC). 2019, pp. 691–696. DOI: 10.23919/ECC.2019.8796208.



P. Chanfreut, J. M. Maestre, and H. Ishii. "Vulnerabilities in Distributed Model Predictive Control based on Jacobi-Gauss Decomposition". In: 2018 European Control Conference (ECC). June 2018, pp. 2587–2592. DOI: 10.23919/ECC.2018.8550239.



# For Further Reading IV



Rafael Accácio Nogueira, Romain Bourdais, and Hervé Guéguen. "Detection and Mitigation of Corrupted Information in Distributed Model Predictive Control Based on Resource Allocation". In: 2021 5th Conference on Control and Fault-Tolerant Systems (SysTol). 2021, pp. 329–334. DOI: 10.1109/SysTol52990.2021.9595927.



Pablo Velarde et al. "Scenario-based defense mechanism for distributed model predictive control". In: 2017 IEEE 56th Annual Conference on Decision and Control (CDC). IEEE. Dec. 2017, pp. 6171–6176. DOI: 10.1109/CDC.2017.8264590.



# For Further Reading V



Pablo Velarde et al. "Vulnerabilities in Lagrange-Based DMPC in the Context of Cyber-Security". In:

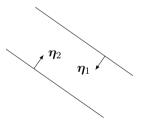
2017 IEEE International Conference on Autonomic Computing (ICAC). July 2017, pp. 215–220. DOI: 10.1109/ICAC.2017.53.



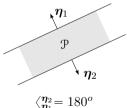
## Conditions



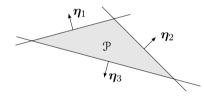
One way to ensure this, is to make the original constraints to form a cone.



No intersection



 $\langle \eta_1^2 = 180^o$ 



A 3-sided polyhedron.



# $\theta$ dynamics

**√** back

$$\boldsymbol{\theta}^{(p+1)} = \mathcal{A}_{\boldsymbol{\theta}} \boldsymbol{\theta}^{(p)} + \mathcal{B}_{\boldsymbol{\theta}}[k]$$

where

$$\mathcal{A}_{\theta} = \begin{bmatrix} I - \frac{M-1}{M} \rho^{(p)} P_{1} & \frac{1}{M} \rho^{(p)} P_{2} & \dots & \frac{1}{M} \rho^{(p)} P_{M} \\ \frac{1}{M} \rho^{(p)} P_{1} & I - \frac{M-1}{M} \rho^{(p)} P_{2} & \dots & \frac{1}{M} \rho^{(p)} P_{M} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{1}{M} \rho^{(p)} P_{1} & \frac{1}{M} \rho^{(p)} P_{2} & \dots & I - \frac{M-1}{M} \rho^{(p)} P_{M} \end{bmatrix}$$

$$\mathcal{B}_{\theta}[k] = \begin{bmatrix} -\frac{M-1}{M} \rho^{(p)} s_{1}[k] + \frac{1}{M} \rho^{(p)} s_{2}[k] \cdots - \frac{1}{M} \rho^{(p)} s_{M}[k] \\ \frac{1}{M} \rho^{(p)} s_{1}[k] - \frac{M-1}{M} \rho^{(p)} s_{2}[k] \cdots - \frac{1}{M} \rho^{(p)} s_{M}[k] \\ \vdots & \vdots \\ \frac{1}{M} \rho^{(p)} s_{1}[k] + \frac{1}{M} \rho^{(p)} s_{2}[k] \cdots - \frac{M-1}{M} \rho^{(p)} s_{M}[k] \end{bmatrix}$$

