Security of distributed Model Predictive Control under False Data injection

Rafael Accácio NOGUEIRA

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https://bit.ly/3g3S6X4











- Electricity Distribution System
- Heat distribution
- Water distribution
- Traffic management (include your problem here)





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"Necessity is the mother of invention"



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- Coupled by constraints
 Technical/ Comfort
- Optimization objectives
 Minimize energy consumption
 Maximize user satisfaction
 Follow a trajectory
- Solution \rightarrow MPC





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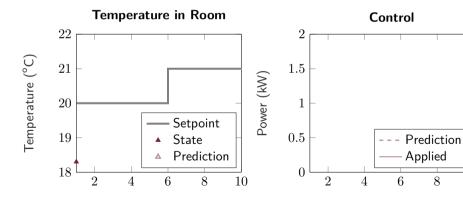
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In a nutshell

Find optimal control sequence

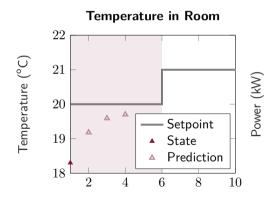


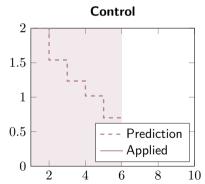


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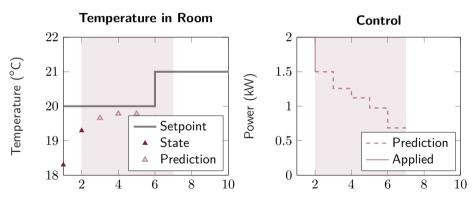






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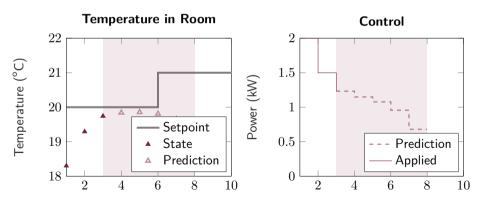
Find optimal control sequence, apply first element





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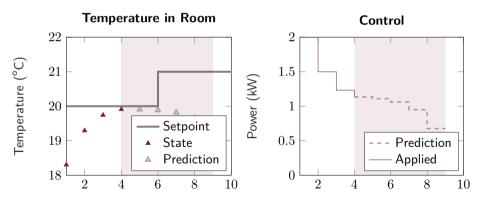
Find optimal control sequence, apply first element, rinse repeat





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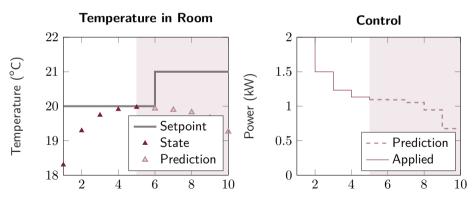
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Nothing is perfect

- Problems
 - Complexity of calculation
 - Topology (Geographical distribution)
 - Flexibility (Add/remove parts)
 - Privacy
- Solution: Divide and Conquer (distributed MPC)
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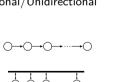








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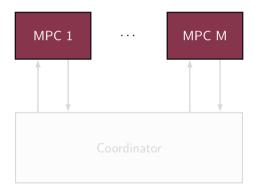






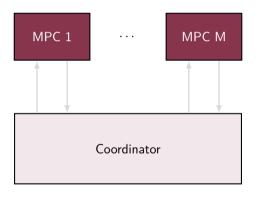






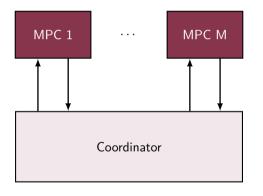
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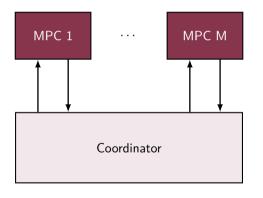
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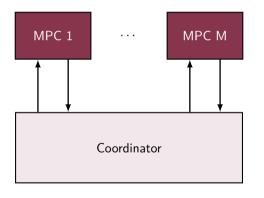
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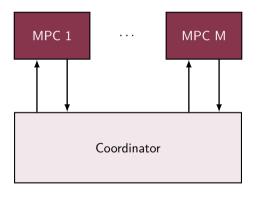




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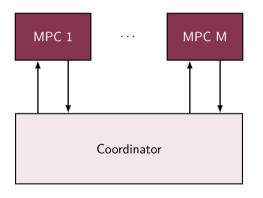




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But what if some agents are ill-intentioned and attack the system?

- How can an agent attack?
- What are the consequences of an attack?
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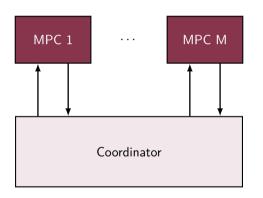


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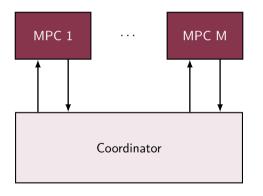


Literature



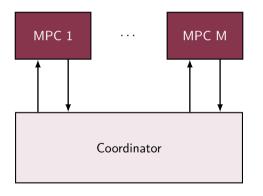
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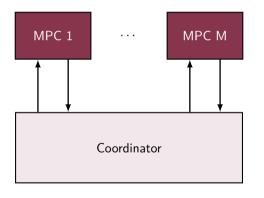
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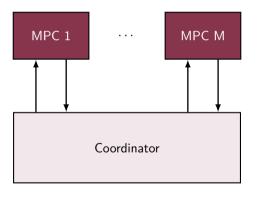
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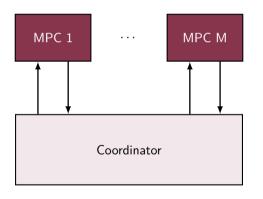




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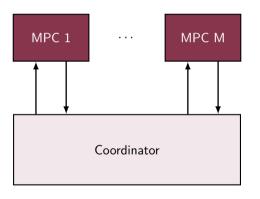


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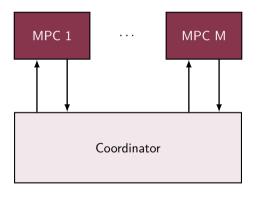
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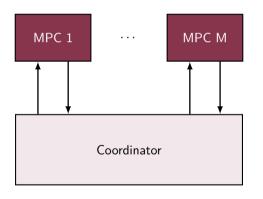
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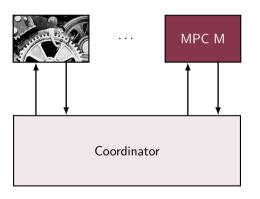


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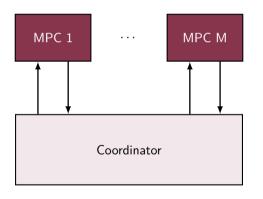


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Deception Attacks (Internal change)

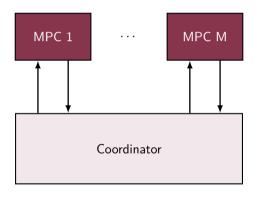




- We are in coordinator's shoes
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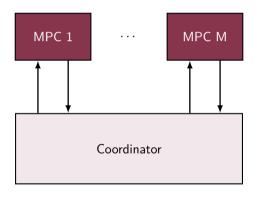






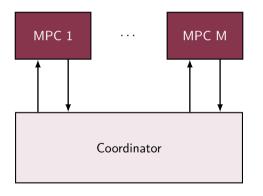
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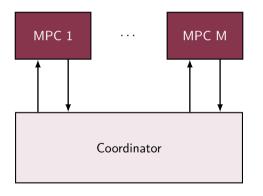
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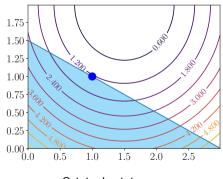
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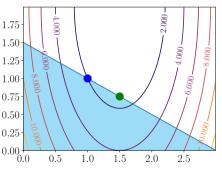
Consequence of an attack

Attack modifies optimization problem

Optimum value is shifted



Original minimum.



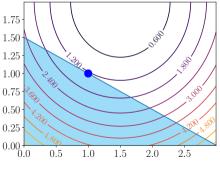
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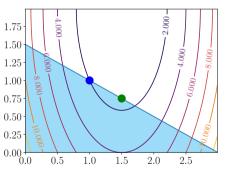
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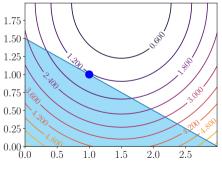


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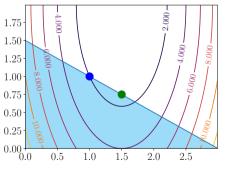


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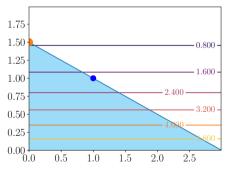
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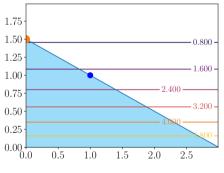
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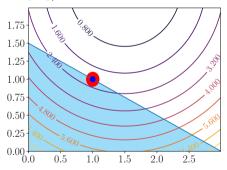
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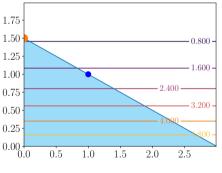
Ignore attacker.



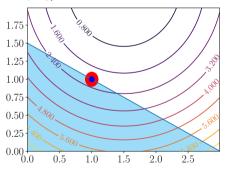
Recover original behavior.



- We can recover by
 - Ignoring attacker
 - Recuperating original behavior (at least trying)



Ignore attacker.



Recover original behavior.



- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
 - Detection/Isolation
 - Mitigation



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- Active (Resilient) 2 modes
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Attack free

- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
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 - Mitigation

Attack free When attack detected



- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
 - ① Detection/Isolation
 - 2 Mitigation

Attack free When attack detected



- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
 - ① Detection/Isolation
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```
Attack free
When attack detected
```



- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
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```
Attack free
When attack detected
```



	Decomposition	Resilient/Robust	Detection	Mitigation
[Vel+17a] [Mae+21]	Dual	Robust (Scenario)	NA	NA
[Vel+17b] [Vel+18]	Dual	Robust (f-robust)	NA	NA
[CMI18]	Jacobi-Gauß	-		
[Ana+18] [Ana+19] [Ana+20]	Dual	Resilient	Analyt./Learn.	Disconnect (Robustness)
Our	Primal	Resilient	Active Analyt./Learn.	Data reconstruction



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State of art

Security dMPC

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- 1 Vulnerabilities in distributed MPC based on Primal Decomposition
- Resilient Primal Decomposition-based dMPC for deprived systems
- Resilient Primal Decomposition-based dMPC using Artificial Scarcity
- 4 Conclusion



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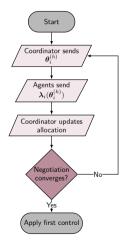


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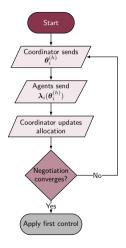


1 Vulnerabilities in distributed MPC based on Primal Decomposition What is the Primal Decomposition? How can an agent attack? Consequences



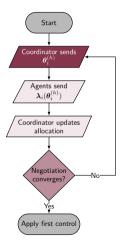








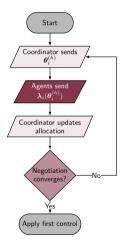




Allocation θ_i

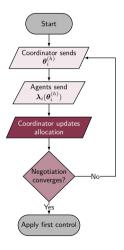


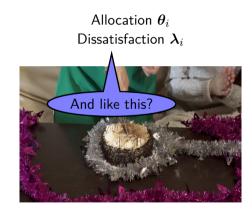




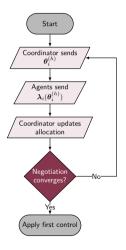






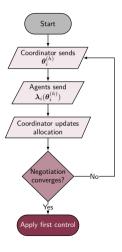












Allocation $oldsymbol{ heta}_i$ Dissatisfaction $oldsymbol{\lambda}_i$





- Objective is sum of local ones
- Constraints couple variables
- $oldsymbol{0}$ Allocate $oldsymbol{ heta}_i$ for each agent
- They solve local problems and
- $oldsymbol{3}$ Send dual variable $oldsymbol{\lambda}_i$
- Allocation is updated (respecting global constraint)

$$egin{array}{ll} & \min & J_i(oldsymbol{x}_i, oldsymbol{u}_i) \ & ext{s. t.} & oldsymbol{h}_i(oldsymbol{x}_i, oldsymbol{u}_i) \leq oldsymbol{ heta}_i + \lambda_i \end{array}$$

$$\boldsymbol{\theta}[k]^{(p+1)} = \operatorname{Proj}^{\mathcal{S}}(\boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)})$$



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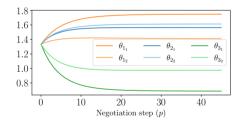
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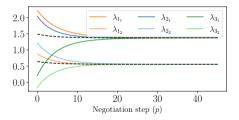
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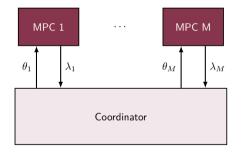


Until everybody is equally dissatisfied



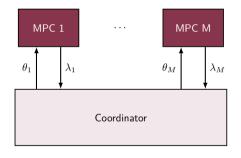






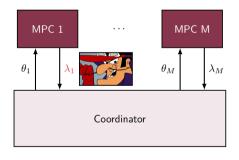
- λ_i is the only interface
- ullet λ_i depends on local parameters
- ullet Malicious agent modifies $oldsymbol{\lambda}_i$





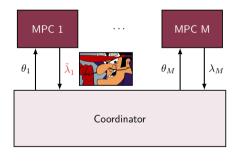
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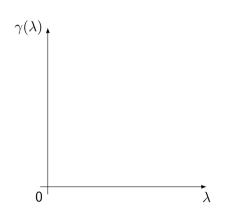


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$$ilde{oldsymbol{\lambda}}_i = \gamma_i(oldsymbol{\lambda}_i)$$



Liar, Liar, Pants of fire

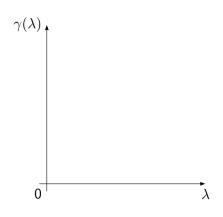


- $\lambda \ge 0$ means dissatisfaction
- $\lambda = 0$ means complete satisfaction

- Attacker satisfied only if it really is
- Attacker is greedy $\gamma(\lambda) > \lambda$
- Attack is monotonically increasing $\lambda_b > \lambda_a \rightarrow \gamma(\lambda_b) > \gamma(\lambda_a)$
- Invertible
- If $\tilde{\lambda}_i = T_i[k]\lambda_i \to \exists T_i[k]^{-1}$



Liar, Liar, Pants of fire

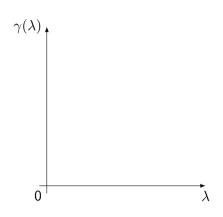


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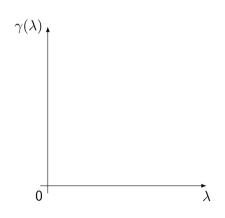


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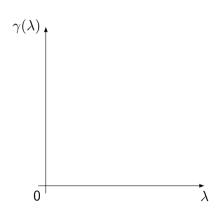
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Liar, Liar, Pants of fire



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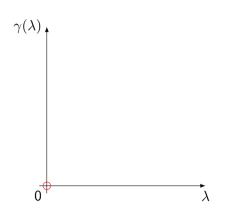
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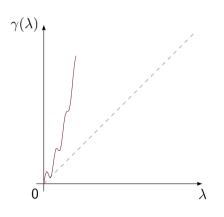
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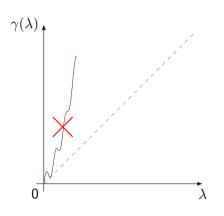
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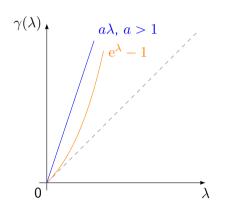
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How does an agent lie?

Liar, Liar, Pants of fire



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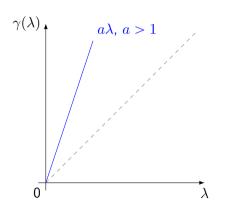
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$$\gamma(\lambda) = 0 \to \lambda = 0$$

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- Attack is monotonically increasing $\lambda_b > \lambda_a \rightarrow \gamma(\lambda_b) > \gamma(\lambda_a)$
- Invertible
- If $\tilde{\lambda}_i = T_i[k]\lambda_i \to \exists T_i[k]^{-1}$



- Agent 1 is non-cooperative
- ullet It uses $ilde{oldsymbol{\lambda}}_1=\gamma_1(oldsymbol{\lambda}_1)= au_1Ioldsymbol{\lambda}_1$
- We can observe 3 things
 - Global minimum when $\tau_1 = 1$
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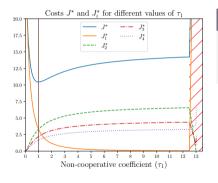


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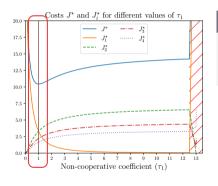
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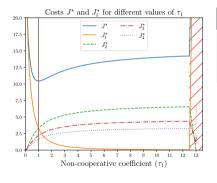
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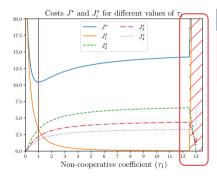
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- Yes! (At least in some cases)



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Outline

Resilient Primal Decomposition-based dMPC for deprived systems
 Analyzing deprived systems
 Building an algorithm
 Applying mechanism



- Unconstrained Solution $\mathring{\boldsymbol{U}}_{i}^{\star}[k]$
- $\bar{\Gamma}_i \mathring{U}_i^{\star}[k] \geq \theta_i[k] \rightarrow \mathsf{Scarcity}$
 - Solution projected onto boundary
 - Same as with equality constraints²

$$\begin{array}{ll} \underset{\boldsymbol{U}_{i}[k]}{\text{minimize}} & \frac{1}{2} \left\| \boldsymbol{U}_{i}[k] \right\|_{H_{i}}^{2} + \boldsymbol{f}_{i}[k]^{T} \boldsymbol{U}_{i}[k] \\ \text{subject to} & \bar{\Gamma}_{i} \boldsymbol{U}_{i}[k] \leq \boldsymbol{\theta}_{i}[k] : \boldsymbol{\lambda}_{i}[k] \end{array}$$



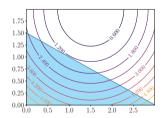
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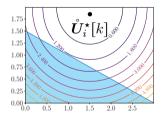
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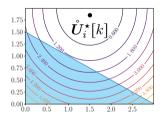
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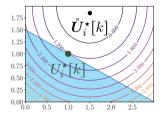




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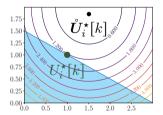




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Analysis

Assumptions

- Quadratic local problems
- Scarcity
- Solution is analytical and affine

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Under attack!

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 - Affine solution

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• Under attack $\rightarrow \lambda_i = T_i[k]\lambda_i$ • Parameters modified

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We know nominal \bar{P}_i

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Rafael Accácio Nogueira

- We estimate $\hat{\tilde{P}}_i[k]$ and $\hat{\tilde{s}}_i[k]$ simultaneously using RLS
- Challenge: Online estimation during negotiation fails
 - Update function couples $heta_i^p$ and $\lambda_i^p o$ low input excitation
- Solution: Send a random³ sequence to increase excitation until convergence.



³A random signal has persistent excitation of any order (

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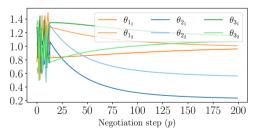
Rafael Accácio Nogueira

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- Challenge: Online estimation during negotiation fails
 - Update function couples $oldsymbol{ heta}_i^p$ and $oldsymbol{\lambda}_i^p o$ low input excitation
- Solution: Send a random³ sequence to increase excitation until convergence.



Estimating $\hat{P}_i[k]$

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Classification of mitigation techniques

- Active (Resilient)
 - Detection/Isolation
 - Mitigation



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Reconstructing λ_i

- Now, we have $\hat{\tilde{P}}_i[k]$
 - Since $\tilde{P}_i[k] = T_i[k]\bar{P}_i$
 - We can recover $T_i[k]^{-1}$

$$\widehat{T_i[k]^{-1}} = P_i \widehat{\tilde{P}}_i[k]^{-1}$$

• Reconstruct λ_i

$$\hat{oldsymbol{\lambda}}_i^{ ext{rec}} = -ar{P}_i oldsymbol{ heta}_i - \widehat{T}_i \widehat{ar{\hat{oldsymbol{s}}}}_i [k]$$

$$oldsymbol{\lambda}_i^{\mathsf{nod}} = egin{cases} ilde{\lambda}_i, & \mathsf{if} \ \mathsf{attack} \ detected \ ilde{\lambda}_i, & \mathsf{otherwise} \end{cases}$$



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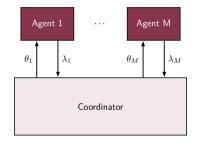
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Complete Mechanism

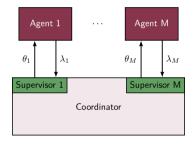


- Supervise exchanges by inquiring the agents
- Estimate how they will behave

Two Phases

- Detect which agents are non-cooperative
- $lue{}$ Reconstruct $oldsymbol{\lambda}_i$ and use in negotiation

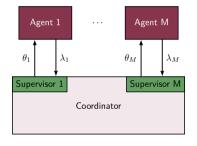




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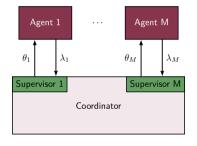




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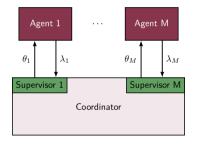




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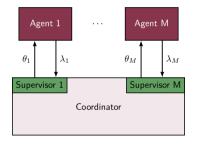




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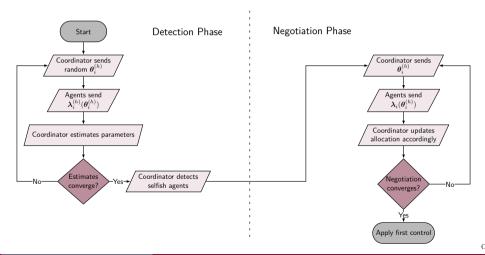




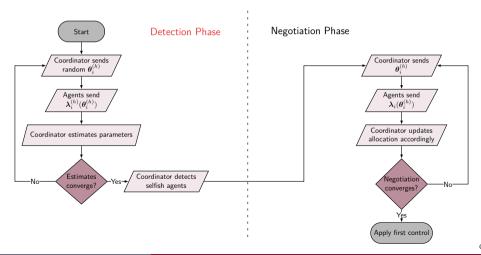
- Supervise exchanges by inquiring the agents
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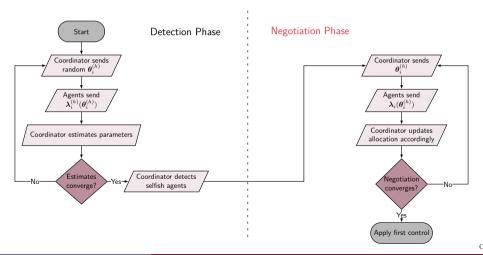




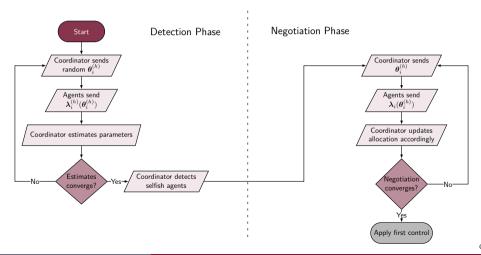




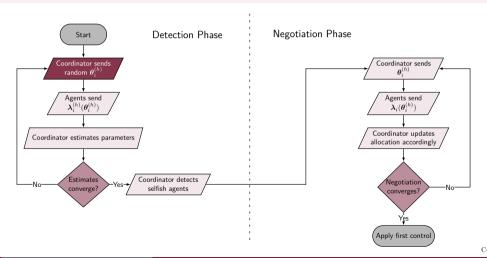


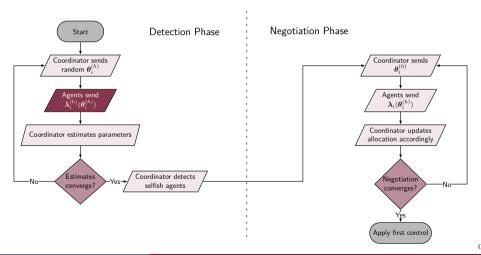




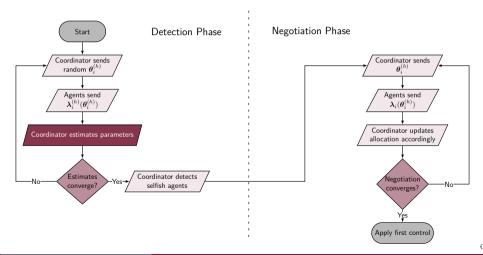




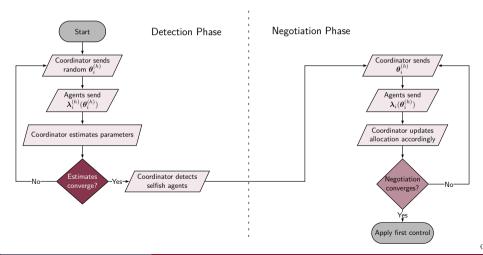




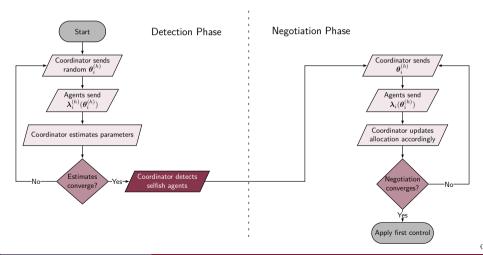




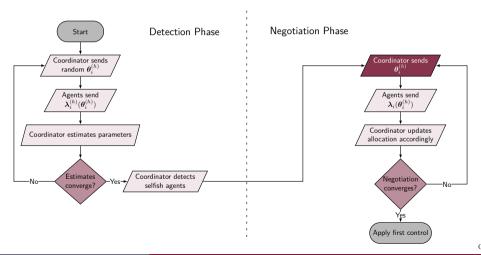


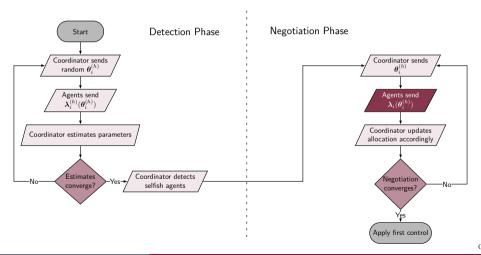




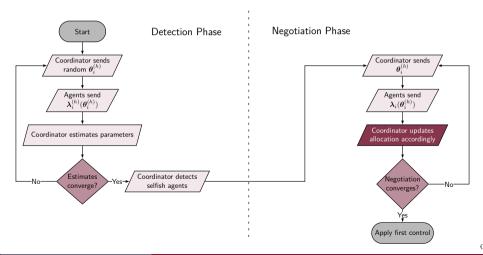




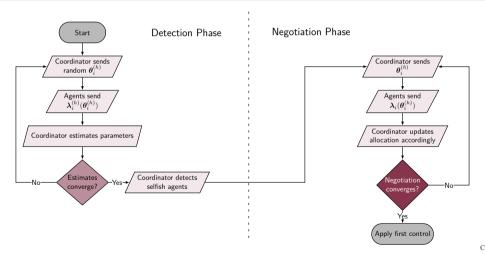


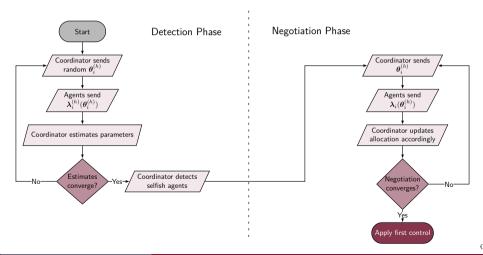




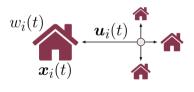






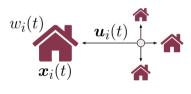






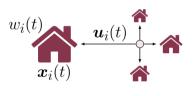
- Houses modeled using 3R-2C (monozone)
- Not enough power
- Period of 5h
- 5 Section 105
 - Agent Laborta (dMPC)
 - Agent I cheats (RPdMPC-DS)





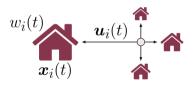
- Houses modeled using 3R-2C (monozone)
- Not enough power
- Period of 5h
- 3 scenarios
 - N Nominal
 - Agent I cheats (dMPC)
 - S Agent I cheats (RPdMPC-DS)





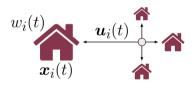
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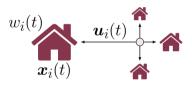
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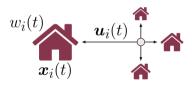
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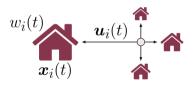
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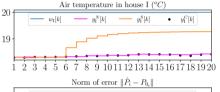


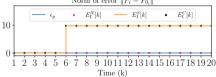


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Temporal





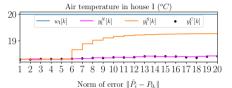
Temperature in house I. Error $E_I(k)$.

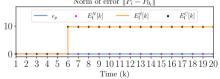
Nominal, S Selflish, C Corrected



Applying mechanism

Temporal



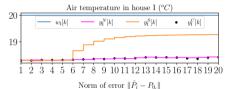


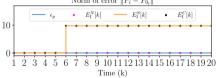
Temperature in house I. Error $E_I(k)$.

- Agent starts cheating in k=6
- S Agent increases its comfort
- Restablish behavior close to §



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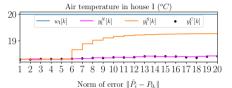


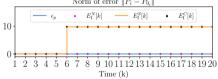
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Temporal



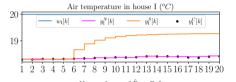


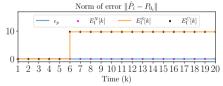
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Costs

Objective functions J_i (Normalized error %)

Agent	Selfish	Corrected
ı	-36.3	0.503
Ш	21.671	-0.547
Ш	17.387	-0.004
IV	17.626	-0.09
Global	3.526	0.016



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Outline

Resilient Primal Decomposition-based dMPC using Artificial Scarcity Relaxing some assumptions Adapting the algorithm Applying mechanism



Relaxing scarcity assumption

- Systems are not completely deprived
 - We can't change our constraints to equality ones anymore
 - Nor use the simpler update equation

minimize
$$\frac{1}{U_i[k]} \| \boldsymbol{U}_i[k] \|_{H_i}^2 + \boldsymbol{f}_i[k]^T \boldsymbol{U}_i[k]$$
subject to
$$\bar{\Gamma}_i \boldsymbol{U}_i[k] \leq \boldsymbol{\theta}_i[k] : \boldsymbol{\lambda}_i[k]$$

$$\boldsymbol{\theta}[k]^{(p+1)} = \operatorname{Proj}^{8}(\boldsymbol{\theta}[k]^{(p)} + \boldsymbol{\rho}^{(p)} \boldsymbol{\lambda}[k]^{(p)})$$



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$$\begin{array}{ll} \underset{\boldsymbol{U}_{i}[k]}{\text{minimize}} & \frac{1}{2} \|\boldsymbol{U}_{i}[k]\|_{H_{i}}^{2} + \boldsymbol{f}_{i}[k]^{T} \boldsymbol{U}_{i}[k] \\ \text{subject to} & \bar{\Gamma}_{i} \boldsymbol{U}_{i}[k] \leq \boldsymbol{\theta}_{i}[k] : \boldsymbol{\lambda}_{i}[k] \end{array}$$

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Solution for $\lambda_i[k]$

Instead of having one single affine solution

$$\boldsymbol{\lambda}_i[k] = -P_i\boldsymbol{\theta}_i[k] - \boldsymbol{s}_i[k]$$

Now, we may have multiple (Piecewise affine function)

$$\lambda_i[k] = \begin{cases} -P_i^{(0)}\theta_i[k] - s_i^{(0)}[k], & \text{if } \theta_i[k] \in \mathcal{R}_{\lambda_i}^0 \\ \vdots & \vdots \\ -P_i^{(2^{n_{\text{ineq}}}-1)}\theta_i[k] - s_i^{(2^{n_{\text{ineq}}}-1)}[k], & \text{if } \theta_i[k] \in \mathcal{R}_{\lambda}^{2^{n_{\text{ineq}}}-1} \end{cases}$$



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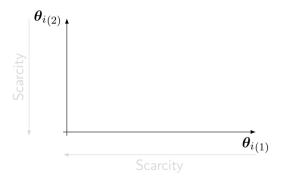
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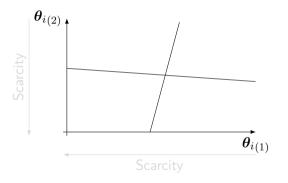


Solution for $\lambda_i[k]$ (Continued)





Solution for $\lambda_i[k]$ (Continued)

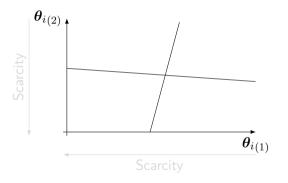


Separation surfaces depend on state and local parameters.

Unknown by the coordinator.

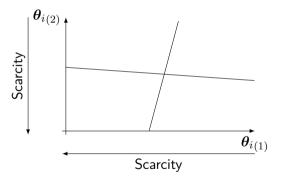


Solution for $\lambda_i[k]$ (Continued)



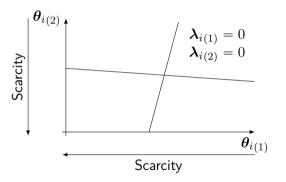


Solution for $\lambda_i[k]$ (Continued)



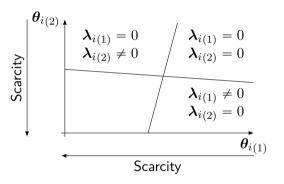


Solution for $\lambda_i[k]$ (Continued)



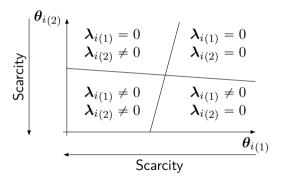


Solution for $\lambda_i[k]$ (Continued)



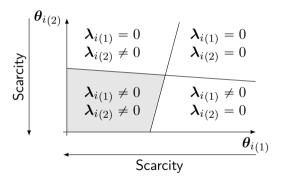


Solution for $\lambda_i[k]$ (Continued)





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 Scarcity Sparsity

All constraints active
$$-P_i^{(0)}\theta_i[k] - s_i^{(0)}[k] \qquad \rightarrow \qquad -P_i\theta_i[k] - s_i[k]$$
 None constraints active
$$-P_i^{\left(2^{n_{\mathsf{ineq}}}-1\right)}\theta_i[k] - s_i^{\left(2^{n_{\mathsf{ineq}}}-1\right)}[k] \qquad \rightarrow \qquad \mathbf{0}$$



$$\boldsymbol{\lambda}_{i}[k] = \begin{cases} -P_{i}^{(0)}\boldsymbol{\theta}_{i}[k] - \boldsymbol{s}_{i}^{(0)}[k], & \text{if } \boldsymbol{\theta}_{i}[k] \in \mathcal{R}_{\boldsymbol{\lambda}_{i}}^{0} \\ \vdots & \vdots \\ -P_{i}^{(2^{n_{\text{ineq}}}-1)}\boldsymbol{\theta}_{i}[k] - \boldsymbol{s}_{i}^{(2^{n_{\text{ineq}}}-1)}[k], & \text{if } \boldsymbol{\theta}_{i}[k] \in \mathcal{R}_{\boldsymbol{\lambda}_{i}}^{2^{n_{\text{ineq}}}-1} \end{cases}$$
 Scarcity

All constraints active
$$-P_i^{(0)}\theta_i[k] - s_i^{(0)}[k] \qquad \rightarrow \qquad -P_i\theta_i[k] - s_i[k]$$
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 Scarcity Sparsity

All constraints active
$$\begin{array}{ccc} -P_i^{(0)} \boldsymbol{\theta_i}[k] - \boldsymbol{s_i^{(0)}}[k] & \rightarrow & -P_i \boldsymbol{\theta_i}[k] - \boldsymbol{s_i}[k] \\ \text{None constraints active} & -P_i^{\left(2^{n_{\text{ineq}}}-1\right)} \boldsymbol{\theta_i}[k] - \boldsymbol{s_i^{\left(2^{n_{\text{ineq}}}-1\right)}}[k] & \rightarrow & 0 \end{array}$$



$$\boldsymbol{\lambda}_{i}[k] = \begin{cases} -P_{i}^{(0)}\boldsymbol{\theta}_{i}[k] - \boldsymbol{s}_{i}^{(0)}[k], & \text{if } \boldsymbol{\theta}_{i}[k] \in \mathcal{R}_{\boldsymbol{\lambda}_{i}}^{0} \\ \vdots & \vdots \\ -P_{i}^{(2^{n_{\text{ineq}}}-1)}\boldsymbol{\theta}_{i}[k] - \boldsymbol{s}_{i}^{(2^{n_{\text{ineq}}}-1)}[k], & \text{if } \boldsymbol{\theta}_{i}[k] \in \mathcal{R}_{\boldsymbol{\lambda}_{i}}^{2^{n_{\text{ineq}}}-1} \end{cases}$$
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All constraints active
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 None constraints active
$$-P_i^{\left(2^{n_{\text{ineq}}}-1\right)} \boldsymbol{\theta}_i[k] - \boldsymbol{s}_i^{\left(2^{n_{\text{ineq}}}-1\right)}[k] \quad \rightarrow \quad \boldsymbol{0}$$



Under attack!

$$\tilde{\boldsymbol{\lambda}}_i[k] = T_i[k]\boldsymbol{\lambda}_k$$

$$\tilde{\pmb{\lambda}}_i[k] = \begin{cases} -\widetilde{P_i}^{(0)} \pmb{\theta}_i[k] - \widetilde{\pmb{s}_i}^{(0)}[k], & \text{if } \pmb{\theta}_i[k] \in \mathbb{R}^0 \\ \vdots & \vdots \\ -\widetilde{P_i}^{(2^{n_{\text{ineq}}}-1)} \pmb{\theta}_i[k] - \widetilde{\pmb{s}_i}^{(2^{n_{\text{ineq}}}-1)}[k], & \text{if } \pmb{\theta}_i[k] \in \mathbb{R}^{2^{n_{\text{ineq}}}-1}_{\pmb{\lambda}_i} \end{cases}$$

- ullet If we can estimate $\widetilde{P}_i^{\,(0)}$ we can use same strategy than before
- ullet Problem: We don't known in which region $oldsymbol{ heta}_i$ is
- Solution: Let's force it using Artificial Scarcity



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Who is it? Who is it?

Assumption

We known a point $\overset{\circ}{\theta}_i$ which activates all constraints⁴

$$\theta_{i(2)}$$

$$\lambda_{i(1)} = 0$$

$$\lambda_{i(2)} \neq 0$$

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CentraleSupélec

⁴If we have local constraints, we suppose this point respects then

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CentraleSupélec

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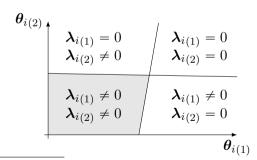
Rafael Accácio Nogueira

Artificial Scarcity

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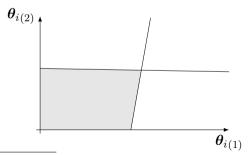


CentraleSupélec

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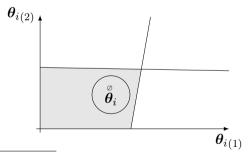
CentraleSupéleo

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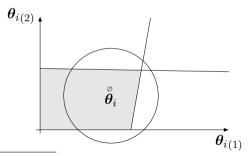


CentraleSupélec

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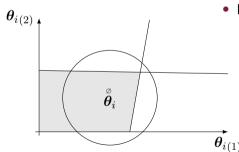
CentraleSupélec

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- How to known the radius?
 - We don't.
 - Let's estimate $\widehat{\widetilde{P}}_i^{(0)}[k]$ nonetheless

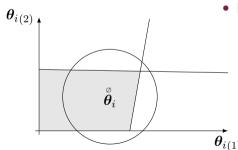


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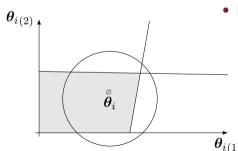
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- Iterative method to estimate parameters of multimodal models⁵
- We give multiple observations $m{ heta}_i^o[k]$ and $ilde{m{\lambda}}_i^o[k]$
- At each step we calculate
 - lacktriangle the probability of each $(\widetilde{P}_i^{(n)}[k],\widehat{\hat{s}}_i^{(n)}[k])$ having generated each $ilde{\lambda}_i^o[k]$
 - mew estimates $(\widetilde{P}_i^{(n)}[k],\widehat{s}_i^{(n)}[k])$ based on the probabilities
- At the end we have
 - Parameters with associated region index
 - Observations with associated region index
- We consult the index associated to $\overset{\circ}{ heta}_i$
- We recover the associated parameter, i.e., $\widehat{\tilde{P}}_i^{(0)}[k]$

CentraleSupélec

⁵Such as our PWA function using some tricks

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CentraleSupélec

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CentraleSupélec

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 - $\qquad \qquad \textbf{(b)} \ \ \, \text{the probability of each} \ \, (\widehat{\widetilde{P}}_i^{(n)}[k],\widehat{\widetilde{s}}_i^{(n)}[k]) \ \, \text{having generated each} \ \, \widetilde{\lambda}_i^o[k]$
 - Moreover new estimates $(\widetilde{P}_i^{(n)}[k],\widehat{\widetilde{s}}_i^{(n)}[k])$ based on the probabilities
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CentraleSupélec

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CentraleSupélec

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CentraleSupélec

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CentraleSupélec

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 - Observations with associated region index
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CentraleSupélec

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CentraleSupélec

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CentraleSupélec

- Iterative method to estimate parameters of multimodal models⁵
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- At the end we have

- Parameters with associated region index
- Observations with associated region index
- We consult the index associated to $\tilde{\theta}_i$
- \bullet We recover the associated parameter, i.e., $\widehat{\widetilde{P}}_{i}^{(0)}\lceil k \rceil$

CentraleSupélec

Same same, but different

Assumption

We know nominal $ar{P}_i{}^{(0)}$

Detection

$$\left\|\widehat{\widetilde{P}}_{i}^{(0)}[k] - \bar{P}_{i}^{(0)}\right\|_{F} \geqslant \epsilon_{P_{i}^{(0)}}$$

$$\widehat{T_i[k]^{-1}} = \bar{P}_i^{(0)} \widehat{\tilde{P}}_i^{(0)}[k]^{-1}.$$

$$\hat{oldsymbol{\lambda}}_i = \widehat{T_i[k]^{-1}} \tilde{oldsymbol{\lambda}}_i.$$



Same same, but different

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Detection

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Detection

$$\left\| \hat{\tilde{P}}_{i}^{(0)}[k] - \bar{P}_{i}^{(0)} \right\|_{F} \ge \epsilon_{P_{i}^{(0)}}$$

$$\widehat{T_i[k]^{-1}} = \overline{P_i}^{(0)} \widehat{P_i}^{(0)}[k]^{-1}.$$

$$\hat{oldsymbol{\lambda}}_i = \widehat{T_i[k]^{-1}} \tilde{oldsymbol{\lambda}}_i.$$



Same same, but different

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We know nominal $\bar{P}_i^{(0)}$

Detection

 $\left\| \hat{\tilde{P}}_{i}^{(0)}[k] - \bar{P}_{i}^{(0)} \right\|_{F} \ge \epsilon_{P_{i}^{(0)}}$

$$\widehat{T_i[k]^{-1}} = \bar{P}_i^{(0)} \widehat{\hat{P}}_i^{(0)}[k]^{-1}.$$

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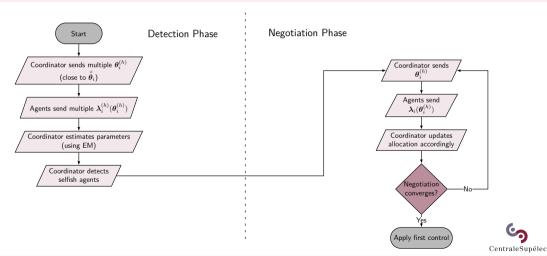
Detection

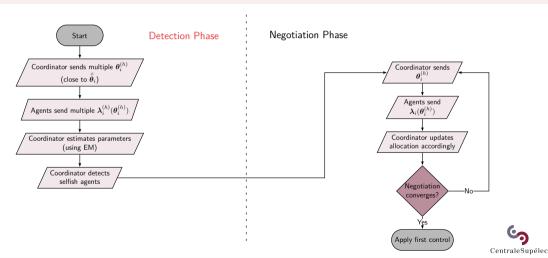
$$\left\| \hat{\tilde{P}}_{i}^{(0)}[k] - \bar{P}_{i}^{(0)} \right\|_{F} \ge \epsilon_{P_{i}^{(0)}}$$

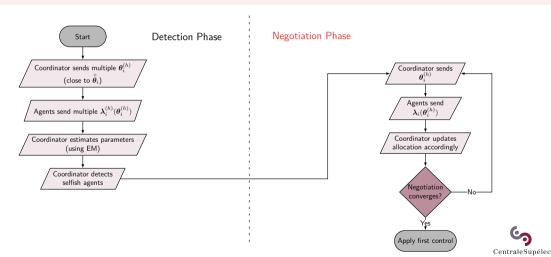
$$\widehat{T_i[k]^{-1}} = \bar{P}_i^{(0)} \widehat{\hat{P}}_i^{(0)}[k]^{-1}.$$

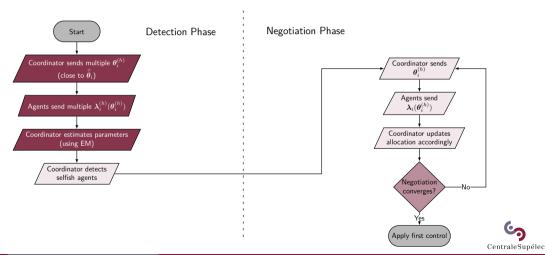
$$\hat{\boldsymbol{\lambda}}_i = \widehat{T_i[k]^{-1}} \tilde{\boldsymbol{\lambda}}_i.$$

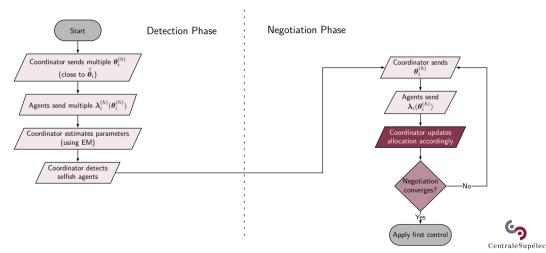




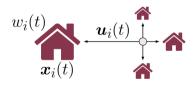








Example



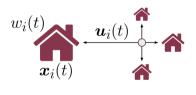
District Heating Network (4 Houses)

- Houses modeled using 3R-2C
- Not enough power
- Period of 5h $(T_s = 0.25h)$
 - 3 scenarios
 - Nominal
 - Agent I cheats (dMPC)
 - S Agent I cheats (RPdMPC-AS)



Applying mechanism

Example



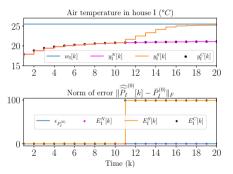
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Results

Temporal



Temperature in house I. Error $E_I(k)$.



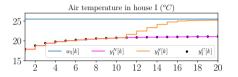


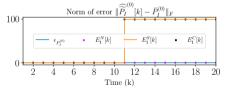


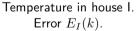


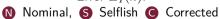
Results

Temporal













Results

Costs

Objective functions J_i (Normalized error %)

Agent	Selfish	Corrected
1	-36.489	-0.0
П	35.813	0.0
Ш	29.225	0.0
IV	37.541	0.0
Global	10.689	-0.0



- Unfortunately EM is not magic
 - Slow convergence
 - Dependency on initialization
 - No guarantees of achieving global optimal
- Some "solutions":
 - Force some parameters to converge faster (case dependant)
 - Run multiple times with different initialization and pick best
 - Associate with other methods of the same family



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Too good to be true!

It's a kind of magic!

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Outline



- How can an agent attack? ✓
 - Attacker can change the communication to receive more ressources.
- What are the consequences of an attack? ✓
 - Suboptimality and maybe instability
- Can we mitigate the effects?
 - Yes! By exploring the scarcity of the systems!



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- Insights from the analysis of the solutions of the optimization problems
 - We found some parameters that are constant when there is no cheating
 - The same parameters change when system is attacked
- Exploiting the solution, we find how to invert the cheating function
 - Straightforward if system is scarce
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- Study of robustness + noise
- Resilient strategy with soft constraints
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Thank you!

 ${\begin{tabular}{l} Repository\\ https://github.com/Accacio/thesis\\ \end{tabular}}$



Contact rafael.accacio.nogueira@gmail.com



For Further Reading I



K.J. Åström and B. Wittenmark. <u>Adaptive Control</u>. Addison-Wesley series in electrical and computer engineering: Control engineering. Addison-Wesley, 1989. ISBN: 9780201097207. DOI: 10.1007/978-3-662-08546-2\ 24.



José M Maestre, Rudy R Negenborn, et al.

<u>Distributed Model Predictive Control made easy.</u> Vol. 69. Springer, 2014. ISBN: 978-94-007-7005-8.



Wicak Ananduta et al. "Resilient Distributed Model Predictive Control for Energy Management of Interconnected Microgrids". In: Optimal Control Applications and Methods 41.1 (2020), pp. 146–169. DOI: 10.1002/oca.2534. URL: https://onlinelibrary.wiley.com/doi/pdf/10.1002/oca.2534.



For Further Reading II



José M. Maestre et al. "Scenario-Based Defense Mechanism Against Vulnerabilities in Lagrange-Based Dmpc". In: Control Eng Pract 114 (2021), p. 104879. ISSN: 0967-0661. DOI: 10.1016/j.conengprac.2021.104879.



Pablo Velarde et al. "Vulnerabilities in Lagrange-Based Distributed Model Predictive Control". In:

Optimal Control Applications and Methods 39.2 (Sept. 2018), pp. 601–621. DOI: 10.1002/oca.2368.



Wicak Ananduta et al. "Resilient Distributed Energy Management for Systems of Interconnected Microgrids". In:

2018 IEEE Conference on Decision and Control (CDC). 2018, pp. 3159–3164. DOI: 10.1109/CDC.2018.8619548.



For Further Reading III



Wicak Ananduta et al. "A Resilient Approach for Distributed MPC-Based Economic Dispatch in Interconnected Microgrids". In: 2019 18th European Control Conference (ECC). 2019, pp. 691–696. DOI: 10.23919/ECC.2019.8796208.



P. Chanfreut, J. M. Maestre, and H. Ishii. "Vulnerabilities in Distributed Model Predictive Control based on Jacobi-Gauss Decomposition". In: 2018 European Control Conference (ECC). June 2018, pp. 2587–2592. DOI: 10.23919/ECC.2018.8550239.



Pablo Velarde et al. "Scenario-based defense mechanism for distributed model predictive control". In: 2017 IEEE 56th Annual Conference on Decision and Control (CDC). IEEE. Dec. 2017, pp. 6171–6176. DOI: 10.1109/CDC.2017.8264590.

For Further Reading IV



Pablo Velarde et al. "Vulnerabilities in Lagrange-Based DMPC in the Context of Cyber-Security". In:

2017 IEEE International Conference on Autonomic Computing (ICAC). July 2017, pp. 215–220. DOI: 10.1109/ICAC.2017.53.



Conditions

♦ back

One way to ensure this, is to make the original constraint (??) to have at most as many rows as columns, i.e., $\#u_{\max} \leq n_u$, although it may be a little restrictive.



θ dynamics

√ back

$$\boldsymbol{\theta}^{(p+1)} = \mathcal{A}_{\boldsymbol{\theta}} \boldsymbol{\theta}^{(p)} + \mathcal{B}_{\boldsymbol{\theta}}[k]$$

where

$$\mathcal{A}_{\theta} = \begin{bmatrix} I - \frac{M-1}{M} \rho^{(p)} P_{1} & \frac{1}{M} \rho^{(p)} P_{2} & \dots & \frac{1}{M} \rho^{(p)} P_{M} \\ \frac{1}{M} \rho^{(p)} P_{1} & I - \frac{M-1}{M} \rho^{(p)} P_{2} & \dots & \frac{1}{M} \rho^{(p)} P_{M} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{1}{M} \rho^{(p)} P_{1} & \frac{1}{M} \rho^{(p)} P_{2} & \dots & I - \frac{M-1}{M} \rho^{(p)} P_{M} \end{bmatrix}$$

$$\mathcal{B}_{\theta}[k] = \begin{bmatrix} -\frac{M-1}{M} \rho^{(p)} s_{1}[k] + \frac{1}{M} \rho^{(p)} s_{2}[k] \cdots - \frac{1}{M} \rho^{(p)} s_{M}[k] \\ \frac{1}{M} \rho^{(p)} s_{1}[k] - \frac{M-1}{M} \rho^{(p)} s_{2}[k] \cdots - \frac{1}{M} \rho^{(p)} s_{M}[k] \\ \vdots & \vdots \\ \frac{1}{M} \rho^{(p)} s_{1}[k] + \frac{1}{M} \rho^{(p)} s_{2}[k] \cdots - \frac{M-1}{M} \rho^{(p)} s_{M}[k] \end{bmatrix}$$

