Security of distributed Model Predictive Control under False Data injection

Rafael Accácio NOGUEIRA

2022-12-07







https://bit.ly/3g3S6X4











- Electricity Distribution System
- Heat distribution
- Water distribution
- Traffic management (include your problem here)





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"Necessity is the mother of invention"



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- Multiple systems interacting
- Coupled by constraints
 Technical/ Comfort
- Optimization objectives
 Minimize energy consumption
 Maximize user satisfaction
 Follow a trajectory
- Solution \rightarrow MPC





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 - Decision variable is the control sequence
 - Objective function to optimize
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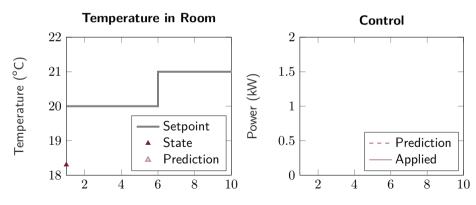
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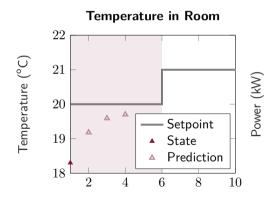
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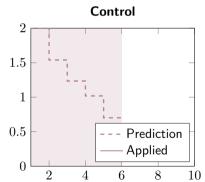




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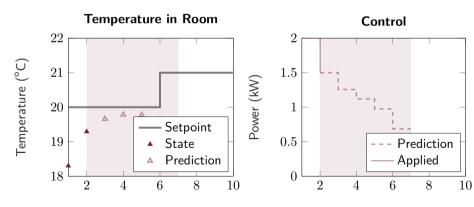






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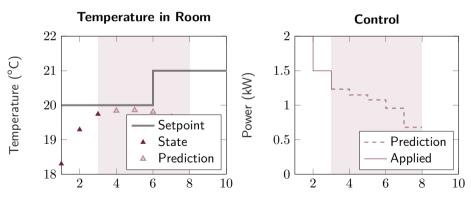
Find optimal control sequence, apply first element





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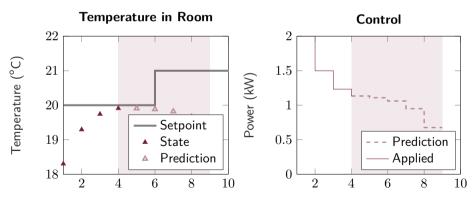
Find optimal control sequence, apply first element, rinse repeat





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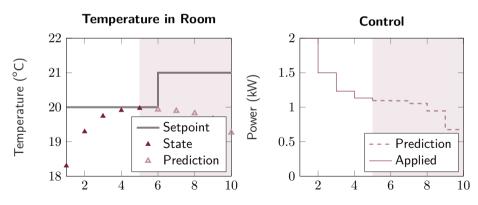
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Nothing is perfect

- Problems
 - Complexity of calculation
 - Topology (Geographical distribution)
 - Flexibility (Add/remove parts)
 - Privacy
- Solution: Divide and Conquer (distributed MPC)
 - Break calculation
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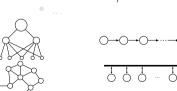








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It is about communication

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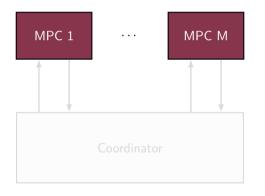






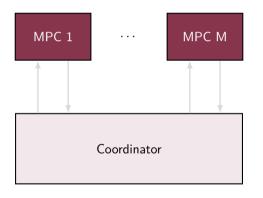






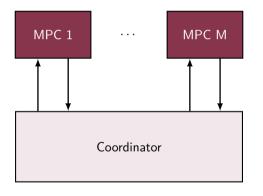
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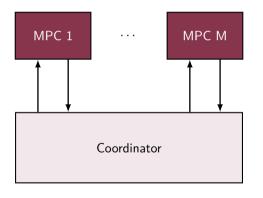
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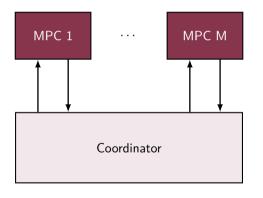
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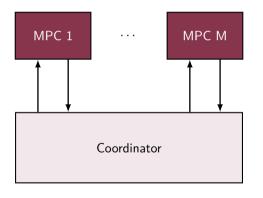




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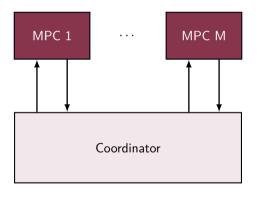




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- How can an agent attack?
- What are the consequences of an attack?
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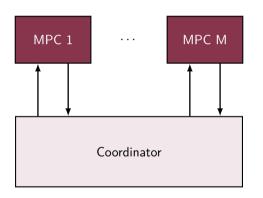


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Literature



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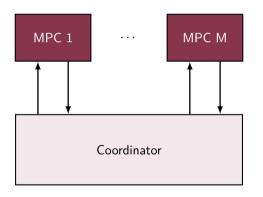
Objective function

Deception Attacks

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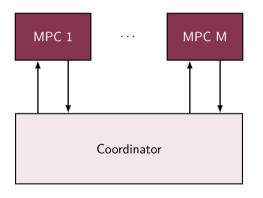
Liar agent





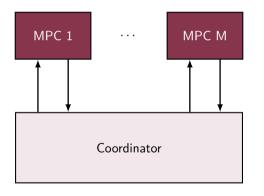
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 - Fake reference
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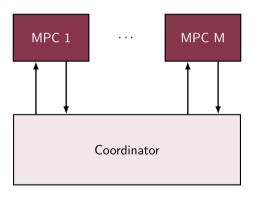
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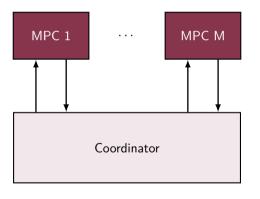




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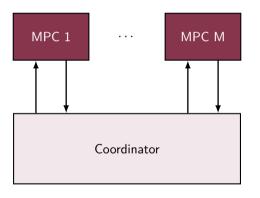
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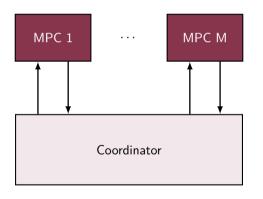
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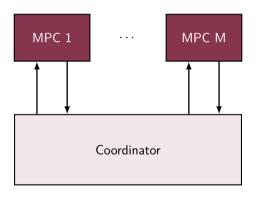


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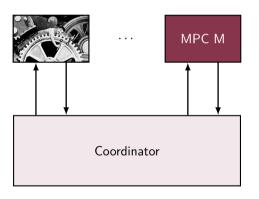


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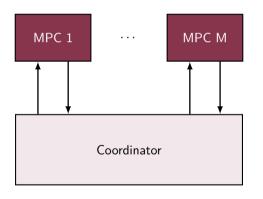


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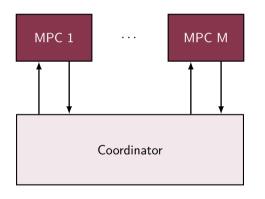
Deception Attacks (Internal change)





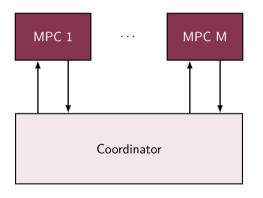
- We are in coordinator's shoes
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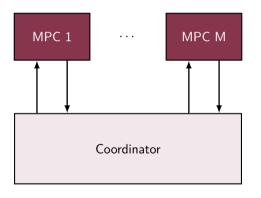
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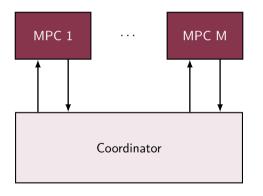
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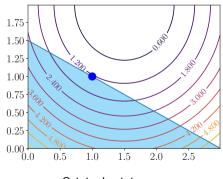
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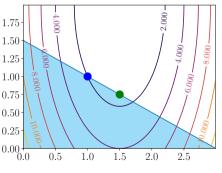
Consequence of an attack

Attack modifies optimization problem

Optimum value is shifted



Original minimum.



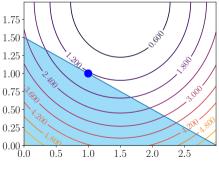
Minimum after attack.



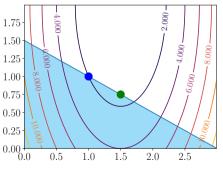
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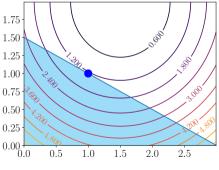


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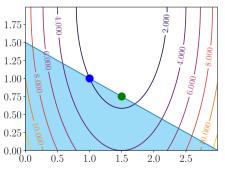


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Original minimum.



Minimum after attack.



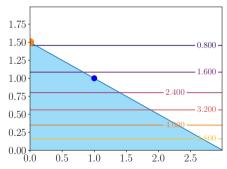
- We can recover by
 - Ignoring attacker
 - Recuperating original behavior (at least trying)



- We can recover by
 - Ignoring attacker
 - Recuperating original behavior (at least trying)



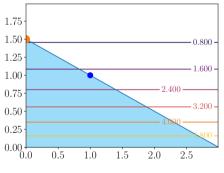
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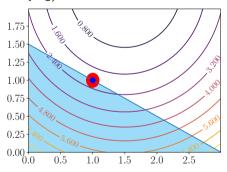
Ignore attacker.



- We can recover by
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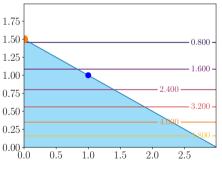
Ignore attacker.



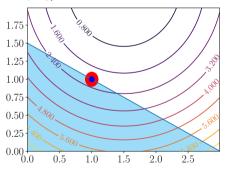
Recover original behavior.



- We can recover by
 - Ignoring attacker
 - Recuperating original behavior (at least trying)



Ignore attacker.



Recover original behavior.



- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
 - Detection/Isolation
 - Mitigation



- Passive (Robust) 1 mode
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 - ① Detection/Isolation
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Attack free

- Passive (Robust) 1 mode
- Active (Resilient) 2 modes
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```
Attack free
When attack detected
```



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Attack free When attack detected



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Attack free
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```



	Decomposition	Resilient/Robust	Detection	Mitigation
[Vel+17a] [Mae+21]	Dual	Robust (Scenario)	NA	NA
$ \begin{array}{c} [Vel + 17b] \\ [Vel + 18] \end{array}$	Dual	Robust (f-robust)	NA	NA
[CMI18]	Jacobi-Gauß	-		
[Ana+18] [Ana+19] [Ana+20]	Dual	Resilient	Analyt./Learn.	Disconnect (Robustness)
Our	Primal	Resilient	Active Analyt./Learn.	Data reconstruction



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State of art

Security dMPC

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- 1 Vulnerabilities in distributed MPC based on Primal Decomposition
- Resilient Primal Decomposition-based dMPC for deprived systems
- Resilient Primal Decomposition-based dMPC using Artificial Scarcity



- 1 Vulnerabilities in distributed MPC based on Primal Decomposition
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- 1 Vulnerabilities in distributed MPC based on Primal Decomposition
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1 Vulnerabilities in distributed MPC based on Primal Decomposition What is the Primal Decomposition? How can an agent attack? Consequences



- Objective is sum of local ones
- Constraints couple variables
- Allocate a part for each agent
- They solve local problems and
- 3 communicate how dissatisfied
- Allocation is updated (respecting global constraint)

$$egin{array}{ll} & \min & J_i(oldsymbol{x}_i, oldsymbol{u}_i) \ & ext{s. t.} & h_i(oldsymbol{x}_i, oldsymbol{u}_i) \leq oldsymbol{ heta}_i : \lambda_i \end{array}$$

$$\boldsymbol{\theta}[k]^{(p+1)} = \operatorname{Proj}^{\mathcal{S}}(\boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)})$$



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$$egin{array}{ll} & \min _{m{u}_1,...,m{u}_M} & \sum_{i \in \mathcal{M}} J_i(m{x}_i,m{u}_i) \ & ext{s.t.} & \sum_{i \in \mathcal{M}} m{h}_i(m{x}_i,m{u}_i) \leq m{u}_{\mathsf{total}} \ & & \downarrow & \mathsf{For each } i \in \mathcal{M} \end{array}$$

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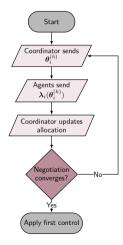
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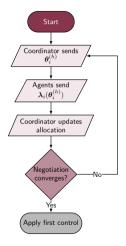
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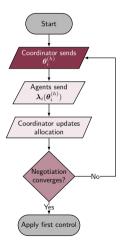








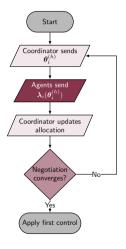




Allocation θ_i

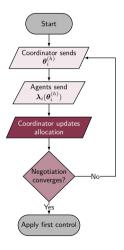


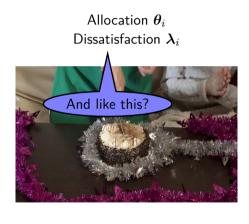




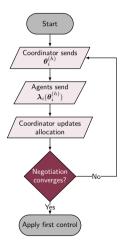






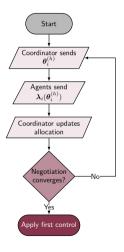










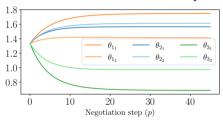


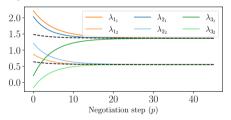
Allocation $oldsymbol{ heta}_i$ Dissatisfaction $oldsymbol{\lambda}_i$



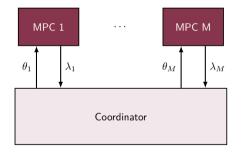


Until everybody is equally dissatisfied



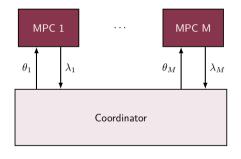






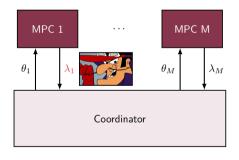
- λ_i is the only interface
- $oldsymbol{\circ}$ $oldsymbol{\lambda}_i$ depends on local parameters
- ullet Malicious agent modifies $oldsymbol{\lambda}_i$





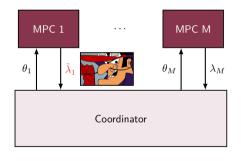
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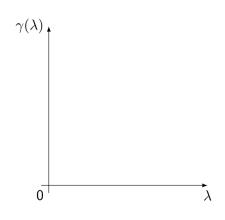


- λ_i is the only interface
- ullet $oldsymbol{\lambda}_i$ depends on local parameters
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$$ilde{oldsymbol{\lambda}}_i = \gamma_i(oldsymbol{\lambda}_i)$$



Liar, Liar, Pants of fire



- $\lambda \ge 0$ means dissatisfaction
- $\lambda = 0$ means complete satisfaction

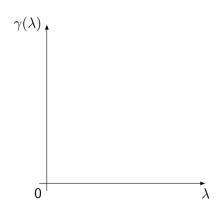
- Attacker is not naïve
- $= \gamma(\lambda) = 0 \to \lambda = 0$
- Attacker is greedy $\gamma(\lambda) > \lambda$
- Really greedy

$$\lambda_b > \lambda_a \to \gamma(\lambda_b) > \gamma(\lambda_a)$$

- Invertible
 - If linear $\rightarrow \exists T^{-1}$



Liar, Liar, Pants of fire



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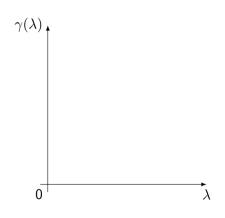
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Liar, Liar, Pants of fire

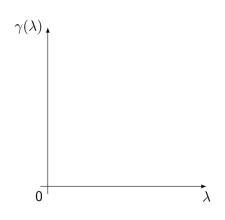


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Liar, Liar, Pants of fire



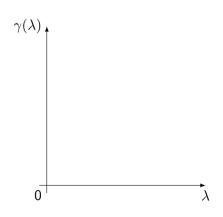
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Liar, Liar, Pants of fire



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Assumptions

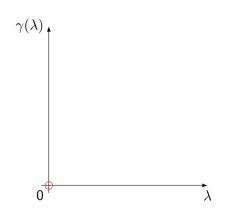
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Liar, Liar, Pants of fire



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Assumptions

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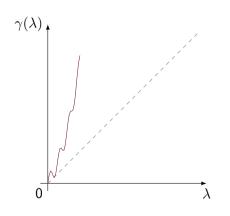
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• If linear $\rightarrow \exists T^{-1}$



Liar, Liar, Pants of fire



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Assumptions

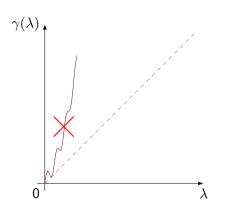
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$$\gamma(\lambda) = 0 \to \lambda = 0$$

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Liar, Liar, Pants of fire



- $\lambda \geqslant 0$ means dissatisfaction
- $\lambda = 0$ means complete satisfaction

Assumptions

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$$\gamma(\lambda) = 0 \to \lambda = 0$$

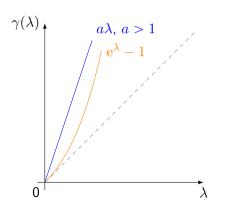
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Liar, Liar, Pants of fire



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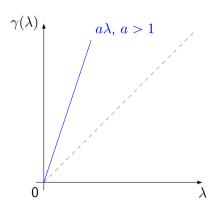
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How does an agent lie?

Liar, Liar, Pants of fire



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Assumptions

Attacker is not naïve

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$$\gamma(\lambda) = 0 \to \lambda = 0$$

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- Agent 1 is non-cooperative
- It uses $ilde{\lambda}_1 = \gamma_1(\lambda_1) = au_1 I \lambda_1$
- We can observe 3 things
 - Global minimum when $\tau_1 = 1$
 - Agent 1 benefits if τ_1 increases (inverse otherwise)
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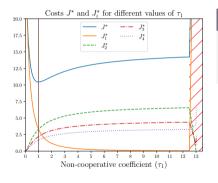


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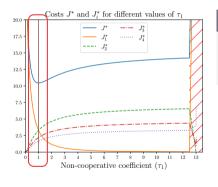
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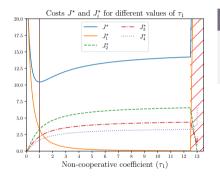
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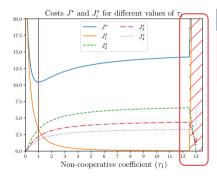
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- Yes! (At least in some cases)
- Let's start slowly



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Outline

Resilient Primal Decomposition-based dMPC for deprived systems
 Analyzing deprived systems
 Building an algorithm
 Applying mechanism



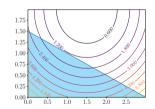
- Unconstrained Solution $\hat{m{U}}_i^{\star}[k]$
- All constraints active = Scarcity
 Solution projected onto boundary

$$\begin{array}{ll} \underset{U_{i}[k]}{\text{minimize}} & \frac{1}{2} \left\| U_{i}[k] \right\|_{H_{i}}^{2} + f_{i}[k]^{T} U_{i}[k] \\ \text{subject to} & \bar{\Gamma}_{i} U_{i}[k] \leq \theta_{i}[k] : \lambda_{i}[k] \end{array}$$



- Unconstrained Solution $\mathring{m{U}}_i^{\star}[k]$
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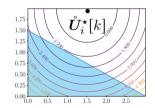
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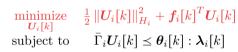
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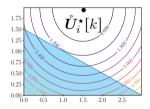
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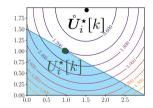




TL;DR: Systems where all constraints are active

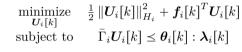
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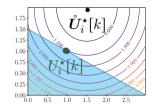
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Analysis

Assumptions

- Quadratic local problems
- Scarcity

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$$\lambda_i[k] = -P_i \boldsymbol{\theta}_i[k] - \boldsymbol{s}_i[k]$$

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- P_i is time invariant
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 $ldsymbol{lap}$ see here



Under attack!

- Normal behavior
 - Affine solution

$$\lambda \cdot [k] = -P \cdot \theta \cdot [k] - s \cdot [k]$$

- Under attack $\to \tilde{\lambda}_i = T_i[k]\lambda_i$
 - Parameters modified

$$\boldsymbol{\theta}^{(p+1)} = \tilde{\mathcal{A}}_{\theta}[k]\boldsymbol{\theta}^{(p)} + \tilde{\mathcal{B}}_{\theta}[k]$$

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We know nominal \bar{P}_i

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 - Update function couples θ_i^p and $\lambda_i^p \to \text{low input excitation}$
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Rafael Accácio Nogueira

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Classification of mitigation techniques

- Active (Resilient)
 - Detection/Isolation
 - Mitigation



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Mitigation mechanism

Reconstructing λ_i

- We now have $\hat{\tilde{P}}_i[k]$
 - Since $\tilde{P}_i[k] = T_i[k]\bar{P}_i$
 - We can recover $T_i[k]^{-1}$

$$\widehat{T_i[k]^{-1}} = P_i \widehat{\tilde{P}}_i[k]^{-1}$$

• Reconstruct λ_i

$$\hat{oldsymbol{\lambda}}_i^{ ext{rec}} = -ar{P}_i oldsymbol{ heta}_i - \widehat{T}_i \widehat{ar{ar{s}}}_i [k]$$

Choose adequate version for coordination

$$oldsymbol{\lambda}_i^{ ext{pod}} = egin{cases} ilde{oldsymbol{\lambda}}_i, & ext{if attack detected} \ ilde{oldsymbol{\lambda}}_i, & ext{otherwise} \end{cases}$$



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Mitigation mechanism

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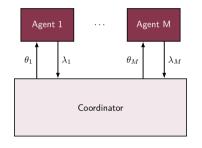
• Reconstruct λ_i

$$\overset{\scriptscriptstyle\mathsf{rec}}{\pmb{\lambda}}_i = -ar{P}_i \pmb{\theta}_i - \widehat{T}_i \widehat{\pmb{[}k]}^{-1} \widehat{\hat{\pmb{s}}}_i [k]$$

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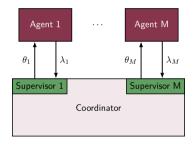




- Supervise exchanges by inquiring the agents
- Estimate how they will behave

- Detect which agents are non-cooperative
- @ Reconstruct $oldsymbol{\lambda}_i$ and use in negotiation

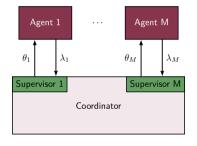




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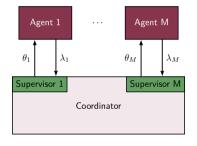




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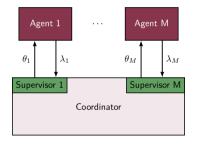




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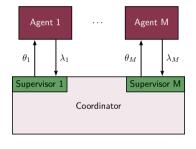




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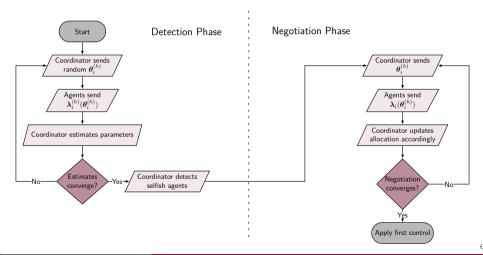




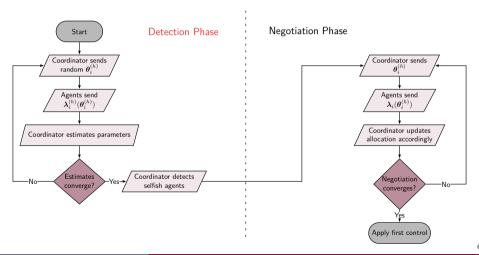
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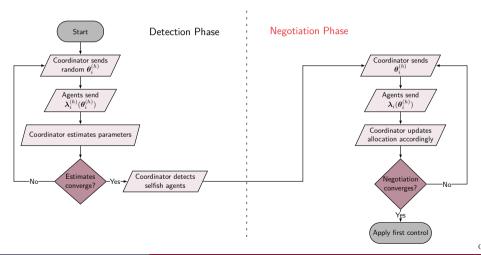
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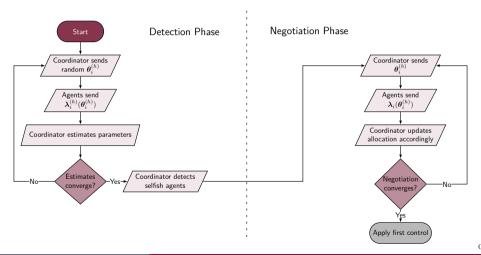




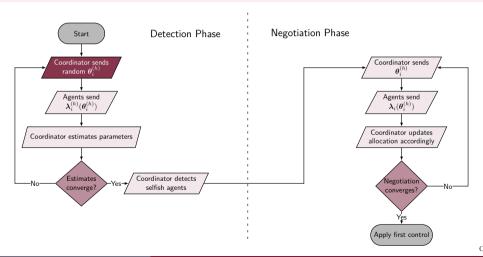




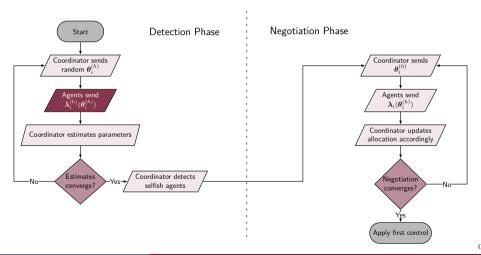




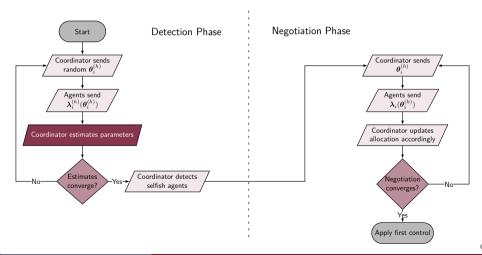


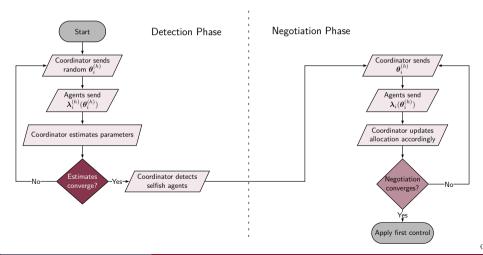




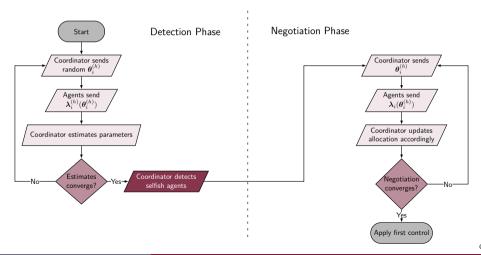




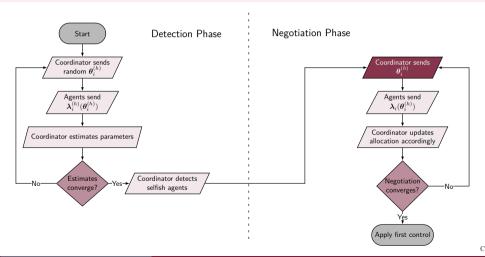


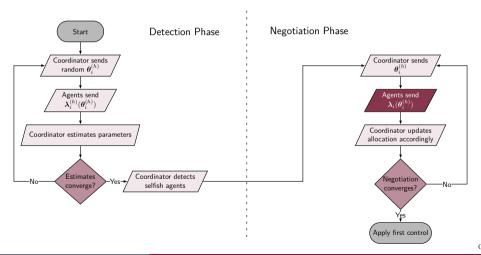




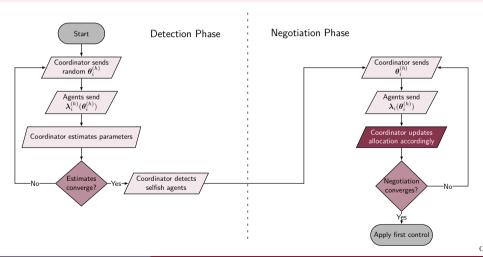




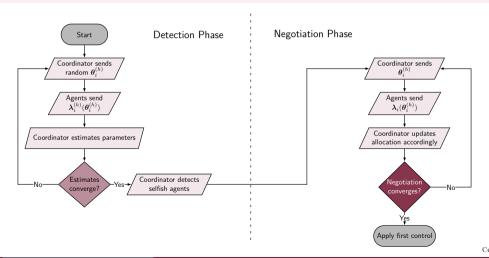


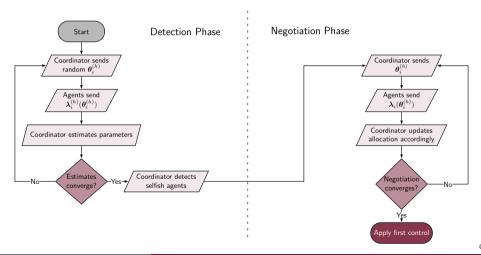




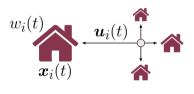






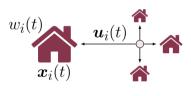






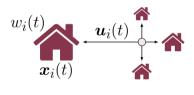
- Houses modeled using 3R-2C (monozone)
- Not enough power
- Period of 5h
- 3 scenarios
 - Nominal
 - Agent I cheats (dMPC)
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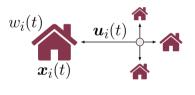
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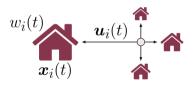
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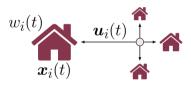
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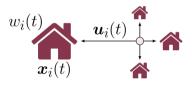
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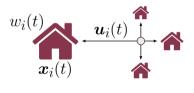
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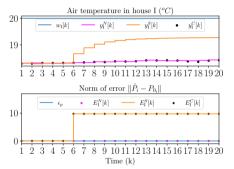




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Temporal

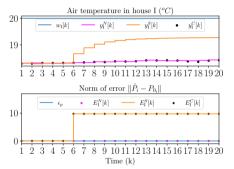


Temperature in house I. Error $E_I(k)$.

Nominal, S Selflish, C Corrected



Temporal

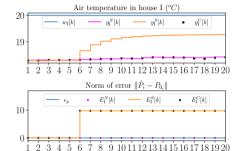


Temperature in house I. Error $E_I(k)$.

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Temporal



Temperature close to reference

Applying mechanism

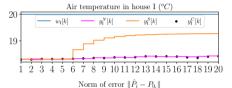
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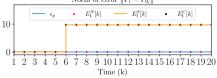
Time (k)

Nominal. S Selflish. Corrected



Temporal





Temperature in house I. Error $E_I(k)$.

Nominal, S Selflish, C Corrected

- S Temperature close to reference



Costs

Objective functions J_i (% error)

Agent	Scenario N	Scenario S	Scenario C
1	299.5	190.8 (-36.3)	301.0 (0.0)
П	192.4	234.1 (21.7)	191.4 (-0.5)
Ш	305.9	359.1 (17.4)	305.9 (-0.0)
IV	297.5	349.9 (17.6)	297.2 (-0.1)
Global	1095.3	1133.9 (3.5)	1095.5 (0.0)



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Outline

Resilient Primal Decomposition-based dMPC using Artificial Scarcity Relaxing some assumptions Adapting the algorithm Results



- Let's relax the scarcity assumption
- And add some local constraints
- Similarly we have the local problems and update

$$\begin{array}{ll} \underset{\boldsymbol{U}[k]}{\text{minimize}} & \frac{1}{2} \|\boldsymbol{U}[k]\|_{H}^{2} + \boldsymbol{f}[k]^{T} \boldsymbol{U}[k] \\ \text{subject to} & \bar{\Gamma} \boldsymbol{U}[k] \leq \boldsymbol{U}_{\text{max}} \\ & U[k] \in \mathcal{U} \end{array}$$

minimize
$$\frac{1}{U_i[k]} \| \boldsymbol{U}_i[k] \|_{H_i}^2 + \boldsymbol{f}_i[k]^T \boldsymbol{U}_i[k]$$
subject to
$$\bar{\Gamma}_i \boldsymbol{U}_i[k] \leq \boldsymbol{\theta}_i[k] : \boldsymbol{\lambda}_i[k]$$

$$\boldsymbol{U}_i[k] \in \boldsymbol{\mathcal{U}}_i$$

$$\boldsymbol{\theta}[k]^{(p+1)} = \operatorname{Proj}^{\mathbb{S}}(\boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)})$$



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$$egin{aligned} & \min & \min_{oldsymbol{U}[k]} & rac{1}{2} \left\| oldsymbol{U}[k]
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$$\begin{array}{ll} \underset{\boldsymbol{U}_{i}[k]}{\text{minimize}} & \frac{1}{2} \|\boldsymbol{U}_{i}[k]\|_{H_{i}}^{2} + \boldsymbol{f}_{i}[k]^{T} \boldsymbol{U}_{i}[k] \\ \text{subject to} & \bar{\Gamma}_{i} \boldsymbol{U}_{i}[k] \leq \boldsymbol{\theta}_{i}[k] : \boldsymbol{\lambda}_{i}[k] \\ & \boldsymbol{U}_{i}[k] \in \boldsymbol{\mathcal{U}}_{i} \end{array}$$

$$\boldsymbol{\theta}[k]^{(p+1)} = \text{Proj}^{\mathbb{S}}(\boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)})$$



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 $\boldsymbol{U}[k] \in \mathcal{U}$

minimize
$$\frac{1}{U_i[k]} \| \boldsymbol{U}_i[k] \|_{H_i}^2 + \boldsymbol{f}_i[k]^T \boldsymbol{U}_i[k]$$
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$$\boldsymbol{\theta}[k]^{(p+1)} = \operatorname{Proj}^{S}(\boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)})$$



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$$\frac{1}{2} \| \boldsymbol{U}[k] \|_H^2 + \boldsymbol{f}[k]^T \boldsymbol{U}[k]$$
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Solution for $\lambda_i[k]$

Instead of having

$$\lambda_i[k] = -P_i \boldsymbol{\theta}_i[k] - \boldsymbol{s}_i[k]$$

Now we have

For n_{ineq} constraints $o 2^{n_{\mathsf{ineq}}}$ permutations



Solution for $\lambda_i[k]$

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Solution for $\lambda_i[k]$

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$$\boldsymbol{\lambda}_i[k] = -P_i \boldsymbol{\theta}_i[k] - \boldsymbol{s}_i[k]$$

Now we have

$$\boldsymbol{\lambda}_i[k] = \begin{cases} -P_i^{(0)} \boldsymbol{\theta}_i[k] - \boldsymbol{s}_i^{(0)}[k], & \text{if } \boldsymbol{\theta}_i[k] \in \mathcal{R}_{\boldsymbol{\lambda}_i}^n \\ \vdots & \vdots \\ -P_i^{\left(2^{n_{\mathsf{ineq}}}-1\right)} \boldsymbol{\theta}_i[k] - \boldsymbol{s}_i^{\left(2^{n_{\mathsf{ineq}}}-1\right)}[k], & \text{if } \boldsymbol{\theta}_i[k] \in \mathcal{R}_{\boldsymbol{\lambda}_i}^{2^{n_{\mathsf{ineq}}}-1} \end{cases} \quad \text{Increasingly Sparse}$$



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$$\boldsymbol{\lambda}_i[k] = \begin{cases} -P_i^{(0)}\boldsymbol{\theta}_i[k] - \boldsymbol{s}_i^{(0)}[k], & \text{if } \boldsymbol{\theta}_i[k] \in \mathcal{R}_{\boldsymbol{\lambda}_i}^n \\ \vdots & \vdots \\ -P_i^{\left(2^{n_{\mathsf{ineq}}}-1\right)}\boldsymbol{\theta}_i[k] - \boldsymbol{s}_i^{\left(2^{n_{\mathsf{ineq}}}-1\right)}[k], & \text{if } \boldsymbol{\theta}_i[k] \in \mathcal{R}_{\boldsymbol{\lambda}_i}^{2^{n_{\mathsf{ineq}}}-1} \end{cases} \text{ Increasingly Sparse}$$



Solution for $\lambda_i[k]$

Instead of having

$$\boldsymbol{\lambda}_i[k] = -P_i \boldsymbol{\theta}_i[k] - \boldsymbol{s}_i[k]$$

Now we have

$$\boldsymbol{\lambda}_i[k] = \begin{cases} -P_i^{(0)}\boldsymbol{\theta}_i[k] - \boldsymbol{s}_i^{(0)}[k], & \text{if } \boldsymbol{\theta}_i[k] \in \mathcal{R}_{\boldsymbol{\lambda}_i}^n \\ \vdots & \vdots \\ -P_i^{\left(2^{n_{\mathsf{ineq}}}-1\right)}\boldsymbol{\theta}_i[k] - \boldsymbol{s}_i^{\left(2^{n_{\mathsf{ineq}}}-1\right)}[k], & \text{if } \boldsymbol{\theta}_i[k] \in \mathcal{R}_{\boldsymbol{\lambda}_i}^{2^{n_{\mathsf{ineq}}}-1} \end{cases} \text{ Increasingly Sparse}$$

For n_{ineq} constraints $\to 2^{n_{\mathsf{ineq}}}$ permutations



Solution for $\lambda_i[k]$

Instead of having

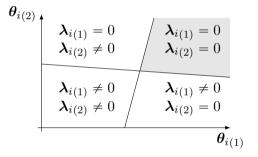
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Solution for $\lambda_i[k]$ (Continued)



Two constraints partitioning θ_i solution space.



Negotiation

$$\operatorname{Proj}^{\mathcal{S}}(\boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)}) = \underset{\boldsymbol{x}}{\operatorname{argmin}} \left\{ \begin{array}{ll} \underset{\boldsymbol{U}[k]}{\operatorname{minimize}} & \left\| \boldsymbol{x} - \boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)} \right\| \\ \underset{\boldsymbol{x}}{\operatorname{subject to}} & I_{c}^{M}\boldsymbol{x} \leq \boldsymbol{U}_{\max} : \boldsymbol{\mu} \end{array} \right\}$$



Negotiation (Continued)

$$\boldsymbol{\theta}[k]^{(p+1)} = \begin{cases} \boldsymbol{x}_0 + I_c^{M^{(0)}} \left[-P_{\boldsymbol{\mu}}^{(0)} \boldsymbol{U}_{max} + \boldsymbol{s}_{\boldsymbol{\mu}}^{(0)}[k] \right], & \text{if } \boldsymbol{x}_0 \in \mathcal{R}_{\boldsymbol{\mu}}^0 \\ \vdots & & \vdots \\ \boldsymbol{x}_0, & \text{if } \boldsymbol{x}_0 \in \mathcal{R}_{\boldsymbol{\mu}}^{2^c - 1} \end{cases}$$

where
$$\boldsymbol{x}_0 = \boldsymbol{\theta}[k]^{(p)} + \rho^{(p)} \boldsymbol{\lambda}[k]^{(p)}$$



Conclusion

$$\bar{\epsilon} = \underbrace{2^c} \times \underbrace{2^{n_{\mathrm{ineq}}} \times \cdots \times 2^{n_{\mathrm{ineq}}}}_{M \times \mathrm{regions in \ each} \ \pmb{\lambda}_i} = 2^{c + M n_{\mathrm{ineq}}}$$



Ideal

$$\widehat{T_i[k]^{-1}} = \bar{P}_i \widehat{\tilde{P}}_i[k]^{-1}$$

• $P_i^{(0)}$ only invertible

$$\widehat{T_i[k]^{-1}} = \bar{P_i}^{(0)} \widehat{\tilde{P}_i}^{(0)}[k]^{-1}$$

But how to force scarcity? Artificial Scarcity



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• But how to force scarcity? Artificial Scarcity



Who is it? Who is it?





Who is it? Who is it?





Who is it? Who is it?





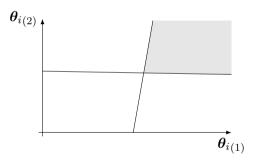
Who is it? Who is it?

$$\begin{array}{c|c} \boldsymbol{\theta}_{i(2)} \\ \boldsymbol{\lambda}_{i(1)} = 0 \\ \boldsymbol{\lambda}_{i(2)} \neq 0 \end{array} \begin{array}{c} \boldsymbol{\lambda}_{i(1)} = 0 \\ \boldsymbol{\lambda}_{i(2)} = 0 \end{array}$$

$$\begin{array}{c|c} \boldsymbol{\lambda}_{i(1)} \neq 0 \\ \boldsymbol{\lambda}_{i(1)} \neq 0 \\ \boldsymbol{\lambda}_{i(2)} \neq 0 \end{array} \begin{array}{c} \boldsymbol{\lambda}_{i(1)} \neq 0 \\ \boldsymbol{\lambda}_{i(2)} = 0 \end{array}$$

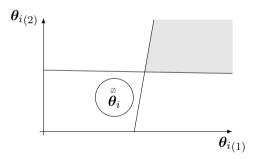


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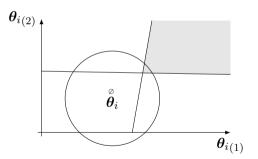


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Who is it? Who is it?





Expectation Maximization

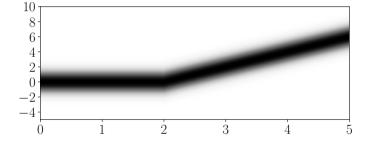


Figure 1: Gaussian Mixture for a 1D PWA function with 2 modes.



Expectation Maximization

Algorithm

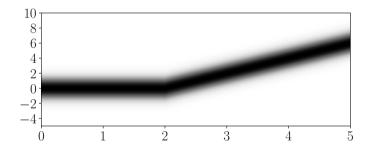


Figure 2: Gaussian Mixture for a 1D PWA function with 2 modes.



• Error
$$E_i^{(0)}[k] = \left\| \widehat{\tilde{P}}_i^{(0)}[k] - \bar{P}_i^{(0)} \right\|_F$$

- \bullet Create threshold $\epsilon_{P_i^{(0)}}$
- Indicator $\mathfrak{d}_i \in \{0,1\}$ detects the attack in agent i.

$$\bullet \ \mathfrak{d}_{i}^{(0)} = \mathbb{1}_{\{E_{i}^{(0)}[k] \geqslant \epsilon_{P_{i}^{(0)}}\}}$$



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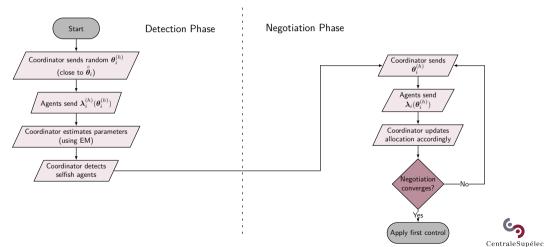
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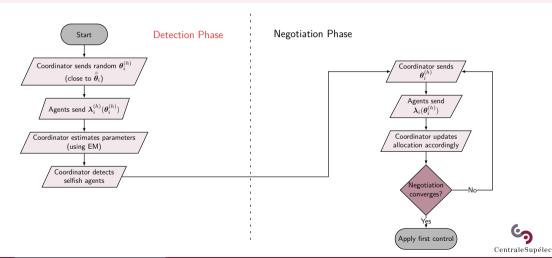


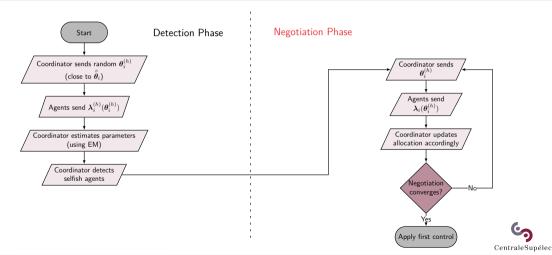
Rafael Accácio Nogueira

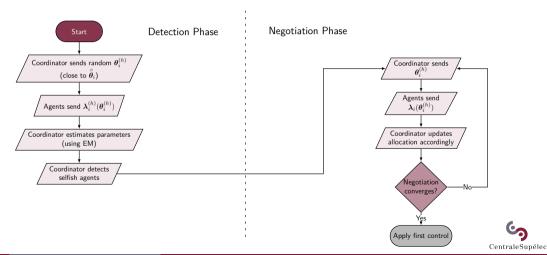
Complete algorithm

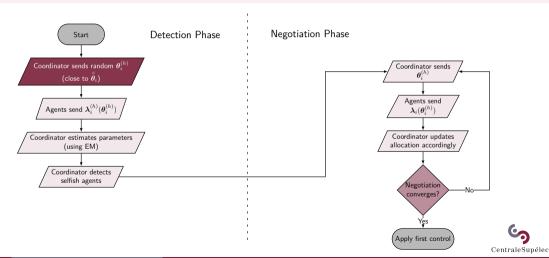
RPdMPC-AS Refaire

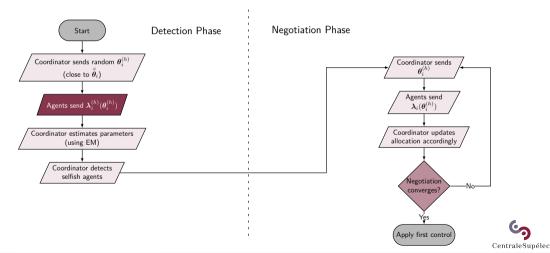


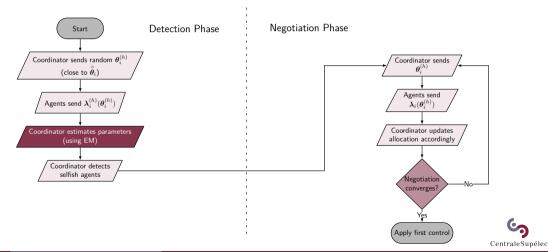


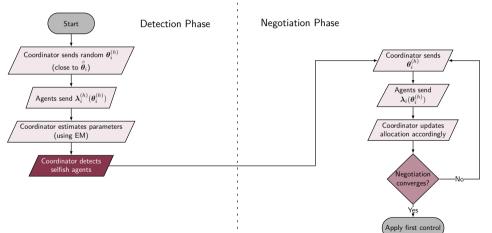


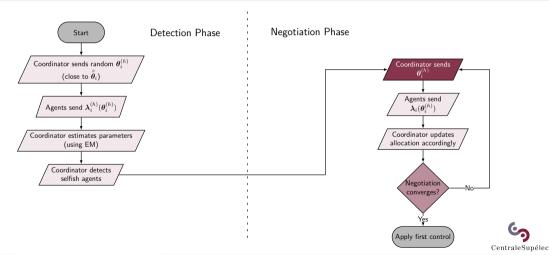


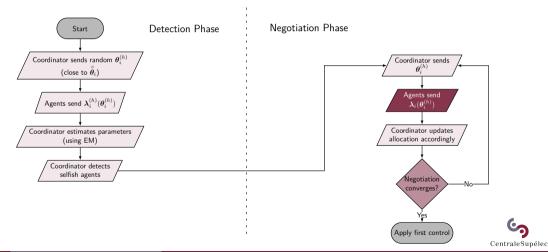


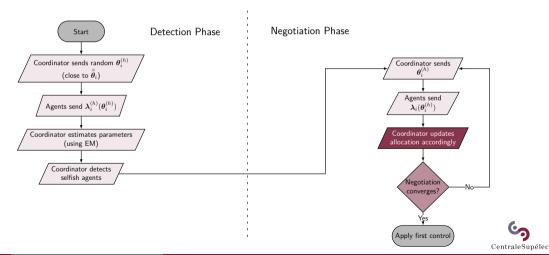


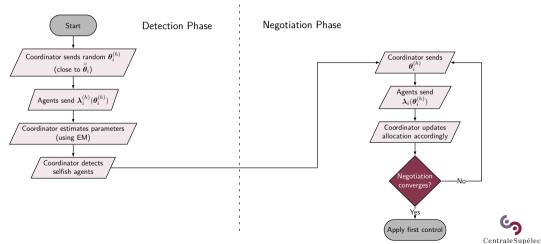


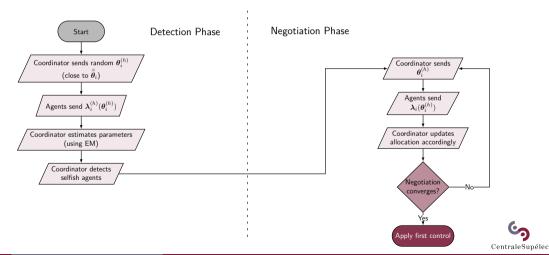


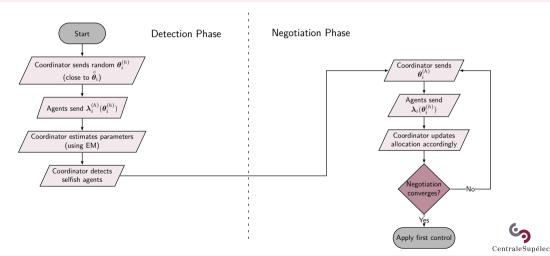




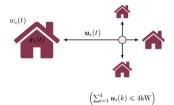








Example



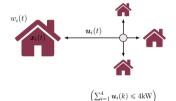
District Heating Network (4 Houses)

- Houses modeled using 3R-2C
- Not enough power
- Period of 5h $(T_s = 0.25h)$
- 3 scenarios
 - Nominal
 - Agent I cheats (dMPC)
 - S Agent I cheats (RPdMPC-AS)

$$T_I = \begin{bmatrix} 14.43288267 & 0 & 0 & 0 & 0 \\ 0 & 13.4590903 & 0 & 0 & 0 \\ 0 & 0 & 6.93065061 & 0 & 0 \\ 0 & 0 & 0 & 3.4447393 \end{bmatrix}$$



Example



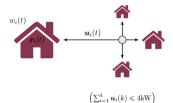
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Example



District Heating Network (4 Houses)

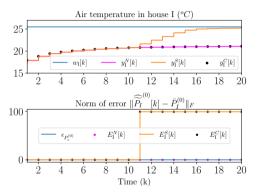
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$$\bullet \quad T_I = \left[\begin{array}{cccc} 14.43288267 & 0. & 0. & 0. \\ 0. & 13.4590903 & 0. & 0. \\ 0. & 0. & 6.93065061 & 0. \\ 0. & 0. & 0. & 3.4447393 \end{array} \right]$$



Results

Temporal

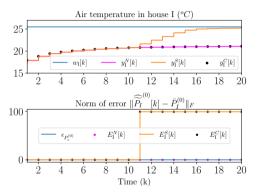


Temperature in house I and the variable $E_I(k)$ for different scenarios.



Results

Temporal

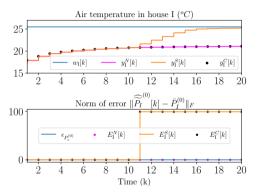


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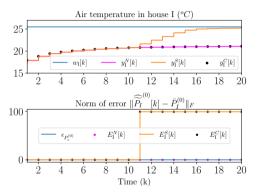


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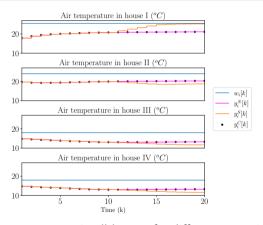
Temporal



Temperature in house I and the variable $E_I(k)$ for different scenarios.



Temporal (Continued)



Air temperature in all houses for different scenarios.





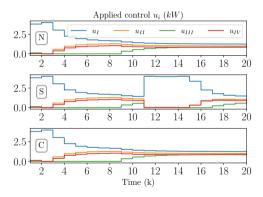








Control



Control applied in all houses for different scenarios.



Costs

Objective functions J_i (% error).

Agent	Scenario N	Scenario S	Scenario C
ı	19868.2	12618.5 (-36.5)	19868.2 (-0.0)
II	13784.5	18721.1 (35.8)	13784.5 (0.0)
Ш	17276.0	22324.9 (29.2)	17276.1 (0.0)
IV	10086.0	13872.4 (37.5)	$10086.0\ (0.0)$
Global	61014.7	67536.9 (10.7)	61014.7 (-0.0)



- Vulnerabilities of Primal decomposition dMPC
- Resilient strategy for 2 kinds of systems
 - Deprived systems (where demand is greater than total resources)
 - Systems with possible artificial scarcity (sensible optimal demand)



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- Partial reconstruction of cheating matrix
- Resilient strategy with soft constraints
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- ...



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- ...



Thank you!

 ${\begin{tabular}{l} Repository\\ https://github.com/Accacio/thesis\\ \end{tabular}}$



Contact rafael.accacio.nogueira@gmail.com



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