Security of distributed Model Predictive Control under False Data injection or How I Learned to Stop and Worry about Everything

Rafael Accácio NOGUEIRA

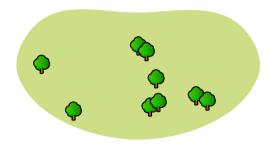
December 12, 2022







https://bit.ly/3g3S6X4























- Electricity Distribution System
- Heat distribution
- Water distribution
- Traffic management (include your problem here)





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Requirements evolve with time



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(include your problem here)





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- Multiple systems interacting
- Coupled by constraints
- Optimization objectives
 - Minimize energy consumptionMaximize user satisfactionFollow a trajectory
- Solution → MPC





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Solution → MPC



- Objective function to optimize
- System's Model (states and inputs)
- Other constraints to respect (QoS, technical restrictions, ...)



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$$\begin{array}{ll} \underset{\boldsymbol{u}[0:N-1|k]}{\operatorname{minimize}} & J(\boldsymbol{x}[0|k],\boldsymbol{u}[0:N-1|k]) \\ & \boldsymbol{x}[\xi|k] = f(\boldsymbol{x}[\xi-1|k],\boldsymbol{u}[\xi-1|k]) \\ \text{subject to} & g_i(\boldsymbol{x}[\xi-1|k],\boldsymbol{u}[\xi-1|k]) \leqslant 0 \\ & h_i(\boldsymbol{x}[\xi-1|k],\boldsymbol{u}[\xi-1|k]) \leqslant 0 \\ & h_i(\boldsymbol{x}[\xi-1|k],\boldsymbol{u}[\xi-1|k]) = 0 \end{array} \right\} \forall i \in \{1,\ldots,n\}$$



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$$\begin{aligned} & \underset{\boldsymbol{u}[0:N-1|k]}{\text{minimize}} & & J(\boldsymbol{x}[0|k], \boldsymbol{u}[0:N-1|k]) \\ & & \boldsymbol{x}[\xi|k] = f(\boldsymbol{x}[\xi-1|k], \boldsymbol{u}[\xi-1|k]) \\ & \text{subject to} & & g_i(\boldsymbol{x}[\xi-1|k], \boldsymbol{u}[\xi-1|k]) \leq 0 \\ & & h_i(\boldsymbol{x}[\xi-1|k], \boldsymbol{u}[\xi-1|k]) = 0 \end{aligned} \right\} \forall i \in \{1, \dots, m\} \\ & & h_i(\boldsymbol{x}[\xi-1|k], \boldsymbol{u}[\xi-1|k]) = 0 \end{aligned}$$



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minimize
$$u[0:N-1|k]$$
 $J(x[0|k], u[0:N-1|k])$ $x[\xi|k] = f(x[\xi-1|k], u[\xi-1|k])$ subject to $g_i(x[\xi-1|k], u[\xi-1|k]) \le 0$ $\forall \xi \in \{1, \dots, N\}$ $\forall i \in \{1, \dots, m\}$ $h_i(x[\xi-1|k], u[\xi-1|k]) = 0$ $\forall j \in \{1, \dots, p\}$



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```
minimize u[0:N-1|k] J(x[0|k], u[0:N-1|k]) x[\xi|k] = f(x[\xi-1|k], u[\xi-1|k]) \} \forall \xi \in \{1, \dots, N\} subject to g_i(x[\xi-1|k], u[\xi-1|k]) \leq 0  h_j(x[\xi-1|k], u[\xi-1|k]) = 0  \forall i \in \{1, \dots, m\} \forall j \in \{1, \dots, p\}
```



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 $\forall \xi \in \{1, \dots, N\}$ $\forall i \in \{1, \dots, m\}$ $\forall j \in \{1, \dots, p\}$



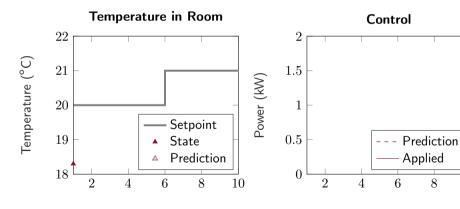
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 $h_j(\boldsymbol{x}[\xi-1|k], \boldsymbol{u}[\xi-1|k]) = 0 \end{cases} \begin{cases} \forall j \in \{1, \dots, m\} \end{cases}$



In a nutshell



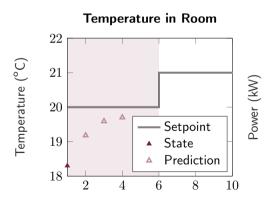


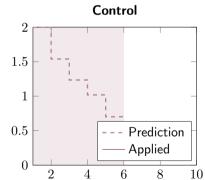
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In a nutshell

Find optimal control sequence

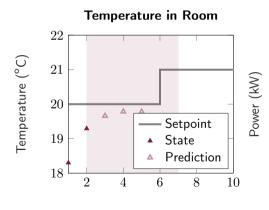


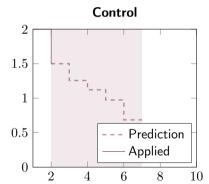




In a nutshell

Find optimal control sequence, apply first element

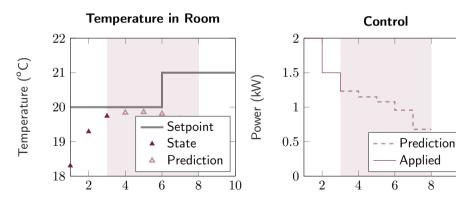






In a nutshell

Find optimal control sequence, apply first element, rinse repeat

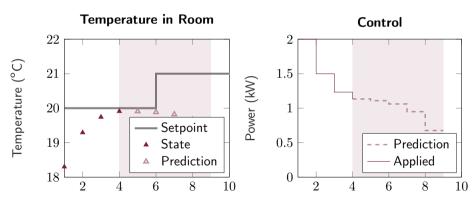




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In a nutshell

Find optimal control sequence, apply first element, rinse repeat \rightarrow Receding Horizon

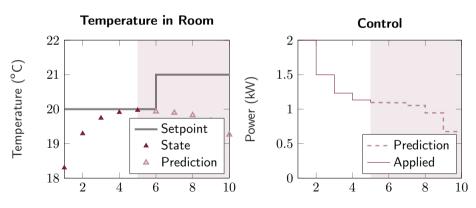




Model Predictive Control

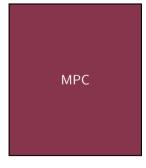
In a nutshell

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- ullet Problem: Complexity depends on N,m,p and sizes of $oldsymbol{x}$ and $oldsymbol{u}$
- Solution: Divide and Conquer¹



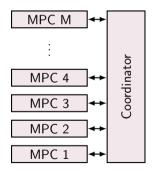


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dMPC

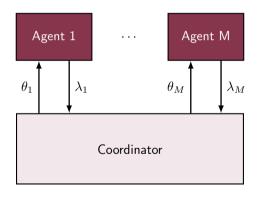


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CentraleSupélec

Optimization Frameworks

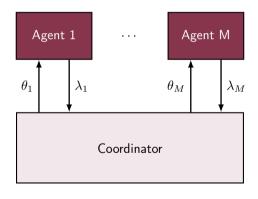


- ullet Agents solve local problems ullet oxdots
- Variables are updated

Convergence



Optimization Frameworks

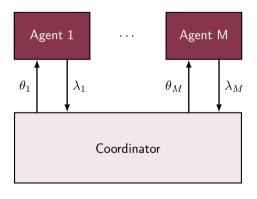


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Convergence



Optimization Frameworks



- Agents solve local problems | Until
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Convergence



Negotiation works if agents comply.

But what if some agents are ill-intentioned and attack the system?

- What are the consequences of an attack?
- Can we mitigate the effects?



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State of art

	Decomposition	Present vulnerabilities?	Resilient/Robust	Detection	Mitigation
[Vel+17a] [Mae+21]	Dual	Yes	Robust (Scenario)	NA	NA
[Vel+17b] [Vel+18]	Dual	Yes	Robust (f-robust)	NA	NA
[CMI18]	Jacobi-Gauß	Yes	-	-	-
[Ana+18] [Ana+19] [Ana+20]	Dual	Yes	Resilient	Analyt./Learn.	Disconnect (Robustness)
Our	Primal	Yes	Resilient	Active Analyt./Learn.	Data reconstruction



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- 1 Vulnerabilities in distributed MPC based on Primal Decomposition
- Resilient Primal Decomposition-based dMPC for deprived systems
- Resilient Primal Decomposition-based dMPC using Artificial Scarcity



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1 Vulnerabilities in distributed MPC based on Primal Decomposition What is the Primal Decomposition? How can an agent attack? Consequences



Primal Decomposition | Quantity Decomposition | Resource Allocation



Primal Decomposition | Quantity Decomposition | Resource Allocation

minimize
$$J(\boldsymbol{x}[0|k], \boldsymbol{u}[0:N-1|k])$$

$$\boldsymbol{x}[\xi|k] = f(\boldsymbol{x}[\xi-1|k], \boldsymbol{u}[\xi-1|k])$$
subject to $g_i(\boldsymbol{x}[\xi-1|k], \boldsymbol{u}[\xi-1|k]) \leqslant 0$

$$h_j(\boldsymbol{x}[\xi-1|k], \boldsymbol{u}[\xi-1|k]) = 0$$

$$\forall \xi \in \{1, \dots, N\}$$

$$\forall i \in \{1, \dots, m\}$$

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Primal Decomposition | Quantity Decomposition | Resource Allocation

minimize
$$\sum_{i \in \mathcal{M}} \sum_{\xi \in \mathcal{N}} \left[\| \boldsymbol{v}_i[\xi|k] \|_{Q_i}^2 + \| \boldsymbol{u}_i[\xi - 1|k] \|_{R_i}^2 \right]$$
subject to
$$\begin{aligned} \boldsymbol{x}[\xi|k] &= f(\boldsymbol{x}[\xi - 1|k], \boldsymbol{u}[\xi - 1|k]) \\ g_i(\boldsymbol{x}[\xi - 1|k], \boldsymbol{u}[\xi - 1|k]) &\leq 0 \\ h_j(\boldsymbol{x}[\xi - 1|k], \boldsymbol{u}[\xi - 1|k]) &= 0 \end{aligned} \end{aligned}$$

$$\forall \xi \in \{1, \dots, N\}$$

$$\forall i \in \{1, \dots, m\}$$

$$\forall j \in \{1, \dots, p\}$$



Primal Decomposition | Quantity Decomposition | Resource Allocation

Decompose original problem using primal problem $v_i = w_i - x_i$



Primal Decomposition | Quantity Decomposition | Resource Allocation

minimize
$$\sum_{i \in \mathcal{M}} \sum_{\xi \in \mathcal{N}} \left[\| \boldsymbol{v}_i[\xi|k] \|_{Q_i}^2 + \| \boldsymbol{u}_i[\xi - 1|k] \|_{R_i}^2 \right]$$
subject to
$$\begin{aligned} \boldsymbol{x}[\xi|k] &= A \boldsymbol{x}[\xi - 1|k] + B \boldsymbol{u}[\xi - 1|k] \\ g_i(\boldsymbol{x}[\xi - 1|k], \boldsymbol{u}[\xi - 1|k]) &\leq 0 \\ h_j(\boldsymbol{x}[\xi - 1|k], \boldsymbol{u}[\xi - 1|k]) &= 0 \end{aligned} \end{aligned}$$

$$\forall \xi \in \{1, \dots, N\}$$

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Primal Decomposition | Quantity Decomposition | Resource Allocation

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$$\mathbf{u}_{[0:N-1|k]} \sum_{i \in \mathcal{M}} \sum_{\xi \in \mathcal{N}} \left[\| \mathbf{v}_{i}[\xi|k] \|_{Q_{i}}^{2} + \| \mathbf{u}_{i}[\xi-1|k] \|_{R_{i}}^{2} \right]$$
subject to
$$\mathbf{x}[\xi|k] = A_{i}\mathbf{x}[\xi-1|k] + B_{i}\mathbf{u}[\xi-1|k]$$

$$\sum_{i \in \mathcal{M}} \Gamma_{i}\mathbf{u}_{i}[\xi|k] \leqslant \mathbf{u}_{\text{max}}$$

$$\forall \xi \in \mathcal{N}$$



Primal Decomposition | Quantity Decomposition | Resource Allocation

$$\begin{array}{ll} \underset{\boldsymbol{U}_{1}[k], \ldots, \boldsymbol{U}_{M}[k]}{\text{minimize}} & \sum\limits_{i \in \mathcal{M}} \left[\frac{1}{2} \left\| \boldsymbol{U}_{i}[k] \right\|_{H_{i}}^{2} + \boldsymbol{f}_{i}[k]^{T} \boldsymbol{U}_{i}[k] \right] \\ \text{subject to} & \sum\limits_{i \in \mathcal{M}} \left[\bar{\Gamma}_{i} \boldsymbol{U}_{i}[k] \right] \leq \boldsymbol{U}_{\mathsf{max}} \end{array}$$



Primal Decomposition | Quantity Decomposition | Resource Allocation

$$\begin{array}{ll} \underset{\boldsymbol{U}_{1}[k]}{\text{minimize}} & \frac{1}{2} \|\boldsymbol{U}_{1}[k]\|_{H_{1}}^{2} + \boldsymbol{f}_{1}[k]^{T}\boldsymbol{U}_{1}[k] \\ \text{subject to} & \bar{\Gamma}_{1}\boldsymbol{U}_{1}[k] \leq \boldsymbol{\theta}_{1}[k] : \boldsymbol{\lambda}_{1}[k] \\ & \vdots & \boldsymbol{\theta}[k]^{(p+1)} = \operatorname{Proj}^{\mathcal{S}}(\boldsymbol{\theta}[k]^{(p)} + \boldsymbol{\rho}^{(p)}\boldsymbol{\lambda}[k]^{(p)}) \\ \underset{\boldsymbol{U}_{M}[k]}{\text{minimize}} & \frac{1}{2} \|\boldsymbol{U}_{M}[k]\|_{H_{M}}^{2} + \boldsymbol{f}_{M}[k]^{T}\boldsymbol{U}_{M}[k] \\ \text{subject to} & \bar{\Gamma}_{M}\boldsymbol{U}_{M}[k] \leq \boldsymbol{\theta}_{M}[k] : \boldsymbol{\lambda}_{M}[k] \end{array}$$



Primal Decomposition | Quantity Decomposition | Resource Allocation

Decompose original problem using primal problem $\mathcal{S} = \{\boldsymbol{\theta}[k] \mid I_c^M \boldsymbol{\theta}[k] \leq \boldsymbol{U}_{\max} \}$ $\underset{\boldsymbol{U}_1[k]}{\text{minimize}} \quad \frac{1}{2} \|\boldsymbol{U}_1[k]\|_{H_1}^2 + \boldsymbol{f}_1[k]^T \boldsymbol{U}_1[k]$ $\text{subject to} \quad \bar{\Gamma}_1 \boldsymbol{U}_1[k] \leq \boldsymbol{\theta}_1[k] : \boldsymbol{\lambda}_1[k]$ \vdots $\underset{\boldsymbol{U}_M[k]}{\text{minimize}} \quad \frac{1}{2} \|\boldsymbol{U}_M[k]\|_{H_M}^2 + \boldsymbol{f}_M[k]^T \boldsymbol{U}_M[k]$ $\text{subject to} \quad \bar{\Gamma}_M \boldsymbol{U}_M[k] \leq \boldsymbol{\theta}_M[k] : \boldsymbol{\lambda}_M[k]$

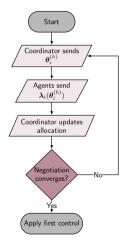


Primal Decomposition | Quantity Decomposition | Resource Allocation

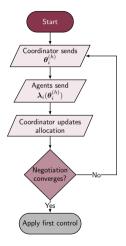
$$S = \{ \boldsymbol{\theta}[k] \mid I_c^M \boldsymbol{\theta}[k] \le \boldsymbol{U}_{\text{max}} \}$$

$$\begin{array}{ll} \underset{\boldsymbol{U}_{1}[k]}{\text{minimize}} & \frac{1}{2} \|\boldsymbol{U}_{1}[k]\|_{H_{1}}^{2} + \boldsymbol{f}_{1}[k]^{T}\boldsymbol{U}_{1}[k] \\ \text{subject to} & \bar{\Gamma}_{1}\boldsymbol{U}_{1}[k] \leq \boldsymbol{\theta}_{1}[k] : \boldsymbol{\lambda}_{1}[k] \\ & \vdots & \boldsymbol{\theta}[k]^{(p+1)} = \operatorname{Proj}^{\mathcal{S}}(\boldsymbol{\theta}[k]^{(p)} + \boldsymbol{\rho}^{(p)}\boldsymbol{\lambda}[k]^{(p)}) \\ \underset{\boldsymbol{U}_{M}[k]}{\text{minimize}} & \frac{1}{2} \|\boldsymbol{U}_{M}[k]\|_{H_{M}}^{2} + \boldsymbol{f}_{M}[k]^{T}\boldsymbol{U}_{M}[k] \\ \text{subject to} & \bar{\Gamma}_{M}\boldsymbol{U}_{M}[k] \leq \boldsymbol{\theta}_{M}[k] : \boldsymbol{\lambda}_{M}[k] \end{array}$$



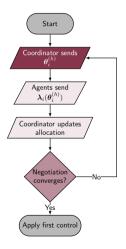






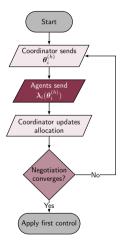






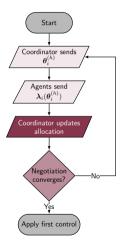






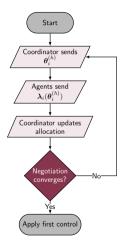






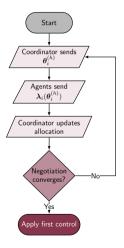






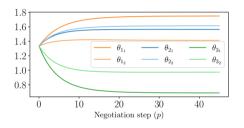


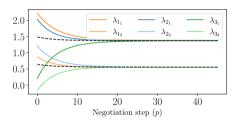














How can a non-cooperative agent attack?

Literature

- [Vel+17a; CMI18] present some kinds of attacks
 - Objective function
 - Selfish Attack
 Fake weights
 - Fake weights
 - Taka constraints
 - Fake constraints
 - Liar agent (use different control)



Literature

- [Vel+17a; CMI18] present some kinds of attacks
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Deception Attack

Fake weightsFake reference

(False Data Injection)

- Fake constraints
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(False Data Injection



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- Deception Attacks (False Data Injection)
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- λ_i depends on the parameters of the system
- Malicious agent sends a different λ_i



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- λ_i depends on the parameters of the system
- Malicious agent sends a different λ_i

$$ilde{oldsymbol{\lambda}}_i = \gamma_i(oldsymbol{\lambda}_i)$$





Let's suppose
$$\gamma_i(\boldsymbol{\lambda}_i) = T_i \boldsymbol{\lambda}_i$$



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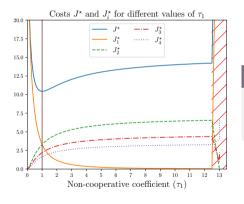


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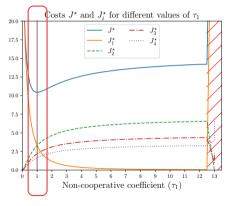
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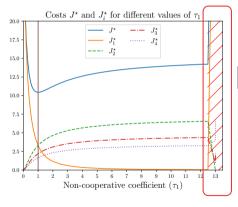
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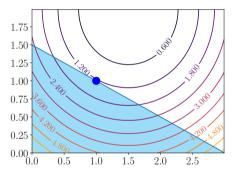
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There are vulnerabilities



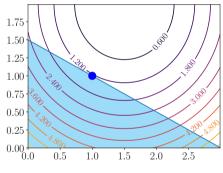
There are vulnerabilities



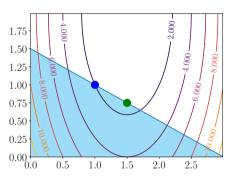
Original minimum.



There are vulnerabilities



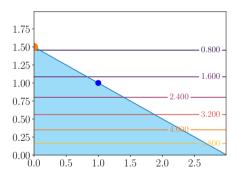
Original minimum.



Minimum after attack.

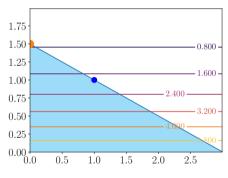




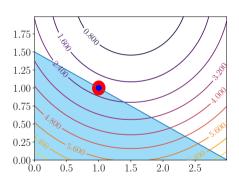


Ignore attacker.



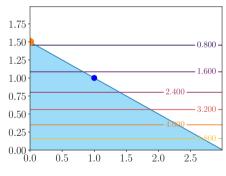


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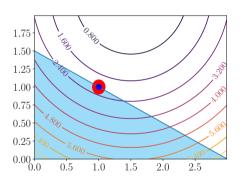


Recover original behavior.





Ignore attacker.



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- Recover original behavior
 - Invert effects of function $\gamma_i(\lambda_i)$
- Is $\gamma_i(\lambda_i)$ invertible?
 - Not necessarily, but let's reason



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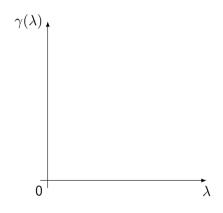
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Unidimensional Case



- $\lambda \ge 0$ means dissatisfaction
- $\lambda = 0$ means complete satisfaction

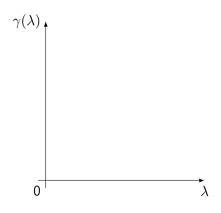
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$$\gamma(\lambda) = 0 \Leftrightarrow \lambda = 0$$

- Attacker is greedy $\gamma(\lambda) > \lambda$
- But not smart

$$\lambda_b > \lambda_a \to \gamma(\lambda_b) > \gamma(\lambda_a)$$



Unidimensional Case



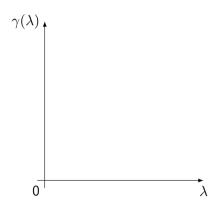
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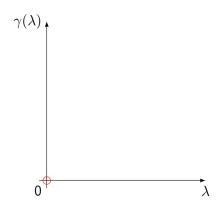
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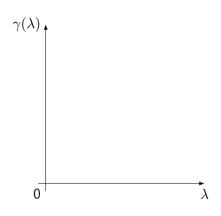
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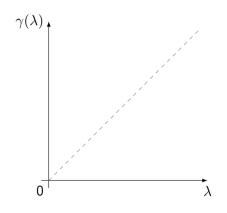
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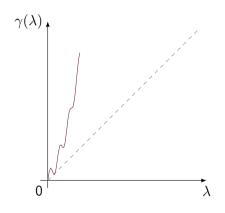
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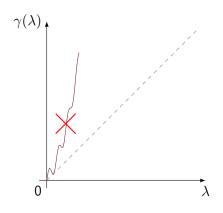
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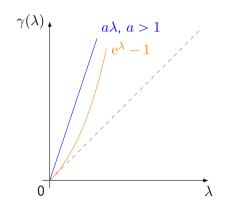
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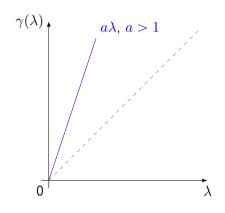
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- $T\lambda = 0 \Leftrightarrow \lambda = 0 \to \exists T^{-1}$
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Is $\gamma_i(\boldsymbol{\lambda}_i)$ invertible?

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Linear Multidimensional Case

- $\gamma(\lambda) = T\lambda$
- $T\lambda = 0 \Leftrightarrow \lambda = 0 \to \exists T^{-1}$
- Each agent could choose different T_i each time $\to T_i[k]$



Consequences

- Recover original behavior
 - Invert effects of function $\gamma_i(\lambda_i) o \mathsf{Estimate}\ T_i[k]^{-1}$
- But how? Analyzing the system



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Outline

Resilient Primal Decomposition-based dMPC for deprived systems
Analyzing deprived systems
Building an algorithm
Applying mechanism



minimize
$$\frac{1}{U_i[k]} \| \boldsymbol{U}_i[k] \|_{H_i}^2 + \boldsymbol{f}_i[k]^T \boldsymbol{U}_i[k]$$

subject to $\bar{\Gamma}_i \boldsymbol{U}_i[k] \leq \boldsymbol{\theta}_i[k] : \boldsymbol{\lambda}_i[k]$

- Unconstrained solution $\mathring{m{U}}_i^{\star}[k] = -H_i^{-1} f_i[k]$
- Deprived if constraints are active → Γ̄_iŮ_i*[k] > θ_i[k], ∀k
 Solution projected onto boundaries (equality constraints)



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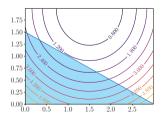
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$$\begin{array}{ll}
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\mathbf{U}_i[k] & \bar{\Gamma}_i \boldsymbol{U}_i[k] \leq \boldsymbol{\theta}_i[k] : \boldsymbol{\lambda}_i[k]
\end{array}$$
subject to

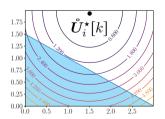
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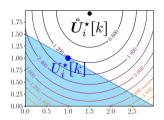
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- Scarcity → Competition → Consensus/Compromise (or cheating 📠)



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Analyzing deprived systems

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Analysis

- We can transform inequality constraints into equality ones²
- Solution is analytical and trivial.

If we solve for λ;

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subject to $\bar{\Gamma}_i \boldsymbol{U}_i[k] \geqslant \boldsymbol{\theta}_i[k] : \boldsymbol{\lambda}_i[k]$

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Rafael Accácio Nogueira

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$$S = \{ \boldsymbol{\theta}[k] \mid I_c^M \boldsymbol{\theta}[k] \leqslant \boldsymbol{U}_{\text{max}} \}$$

$$\boldsymbol{\theta}[k]^{(p+1)} = \operatorname{Proj}^{\mathcal{S}}(\boldsymbol{\theta}[k]^{(p)} + \rho^{(p)}\boldsymbol{\lambda}[k]^{(p)})$$



$$\mathcal{S} = \{ \boldsymbol{\theta}[k] \mid I_c^M \boldsymbol{\theta}[k] = \boldsymbol{U}_{\max} \}$$

$$\boldsymbol{\theta}_{i}^{(p+1)} = \boldsymbol{\theta}_{i}^{(p)} + \rho^{(p)} \left(\boldsymbol{\lambda}_{i}^{(p)} - \frac{1}{M} \sum_{j=1}^{M} \boldsymbol{\lambda}_{j}^{(p)} \right), \forall i \in \mathcal{M}$$



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$$oldsymbol{\lambda}_i[k] = -P_ioldsymbol{ heta}_i[k] - oldsymbol{s}_i[k]$$
 $oldsymbol{\mathcal{S}} = \{oldsymbol{ heta}[k] \mid I_c^Moldsymbol{ heta}[k] = oldsymbol{U}_{ ext{max}}\}$
 $oldsymbol{ heta}_i^{(p+1)} = oldsymbol{ heta}_i^{(p)} +
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ight), orall i \in \mathcal{M}$



$$\lambda_i[k] = -P_i \boldsymbol{\theta}_i[k] - \boldsymbol{s}_i[k]$$

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$$\boldsymbol{\theta}_{i}^{(p+1)} = \boldsymbol{\theta}_{i}^{(p)} + \rho^{(p)} \left(\boldsymbol{\lambda}_{i}^{(p)} - \frac{1}{M} \sum_{j=1}^{M} \boldsymbol{\lambda}_{j}^{(p)} \right), \forall i \in \mathcal{M}$$

$$oldsymbol{ heta}^{(p+1)} = \mathcal{A}_{ heta} oldsymbol{ heta}^{(p)} + \mathcal{B}_{ heta}[k]$$
) see here



Under attack!

Normal behavior

$$\lambda_i[k] = -P_i \boldsymbol{\theta}_i[k] - \boldsymbol{s}_i[k],$$

Under attack

$$\hat{\lambda}_i = T_i[k]\lambda_i = -T_iP_i\theta_i[k] - T_is_i[k]$$

$$\tilde{\lambda}_i = -\tilde{P}_i[k]\theta_i[k] - \tilde{s}_i[k]$$



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Analyzing deprived systems

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Under attack

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Analyzing deprived systems

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$$\tilde{\boldsymbol{\lambda}}_{i} = T_{i}[k]\boldsymbol{\lambda}_{i} = -T_{i}P_{i}\boldsymbol{\theta}_{i}[k] - T_{i}\boldsymbol{s}_{i}[k]$$
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$$egin{align} ilde{m{\lambda}}_i &= T_i[k] m{\lambda}_i = -T_i P_i m{ heta}_i[k] - T_i m{s}_i[k] \ & \\ ilde{m{\lambda}}_i &= - ilde{P}_i[k] m{ heta}_i[k] - ilde{m{s}}_i[k] \ & \\ m{ heta}^{(p+1)} &= ilde{m{\mathcal{A}}}_{m{ heta}}[k] m{ heta}^{(p)} + ilde{\mathcal{B}}_{m{ heta}}[k] \end{split}$$



Assumption

We know nominal \bar{P}_i

Assumption

Attacker chooses
$$\tilde{\lambda}_i = \gamma_i(\lambda_i) = T_i(k)\lambda_i$$

 $-T_i(k)P_i\theta_i - T_i(k)s_i(k) \rightarrow -\tilde{P}_i\theta_i - \tilde{s}_i(k)$

• We can estimate \hat{P}_i and $\hat{\tilde{s}}_i(k)$ such as:

$$\widetilde{\boldsymbol{\lambda}}_i = \gamma_i(\boldsymbol{\lambda}_i(\boldsymbol{\theta}_i)) = -\widehat{\widetilde{P}}_i(k)\boldsymbol{\theta}_i - \widehat{\widetilde{\boldsymbol{s}}}_i(k)$$

• If
$$\hat{\tilde{P}}_i(k) \neq \bar{P}_i \rightarrow \mathsf{Attack}$$

Rafael Accácio Nogueira

CentraleSupélec

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CentraleSupélec

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Assumption

Attacker chooses
$$\tilde{\boldsymbol{\lambda}}_i = \gamma_i(\boldsymbol{\lambda}_i) = T_i(k)\boldsymbol{\lambda}_i - T_i(k)P_i\boldsymbol{\theta}_i - T_i(k)\boldsymbol{s}_i(k) \rightarrow -\tilde{P}_i\boldsymbol{\theta}_i - \tilde{s}_i(k)$$

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Rafael Accácio Nogueira



¹Using Recursive Least Squares

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- Error $E_i(k) = \|\hat{\tilde{P}}_i(k) \bar{P}_i\|_F$
- ullet Create threshold ϵ_P
- Indicator $d_i \in \{0, 1\}$ detects the attack in agent i.
- $d_i = 1$ if $E_i(k) > \epsilon_P$, 0 otherwise



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Estimating $\hat{\tilde{P}}_i[k]$

- ullet We estimate \hat{P}_i and $\hat{ ilde{s}}_i(k)$ simultaneously using Recursive Least Squares
- Problem: Estimation during negotiation fails
 - Consecutive λ_i^p and θ_i^p are linearly dependent \rightarrow low input excitation
- Solution: Send sequence of random values of θ_i until estimates converge



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Building an algorithm

Estimating $\widehat{\widetilde{P}}_i[k]$

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• Main idea: Reconstruct λ_i and use in negotiation

Assumption

We suppose $\tilde{\lambda}_i = \mathbf{0}$ only if $\lambda_i = \mathbf{0}$, which implies $T_i(k)$ invertible

• Estimate the inverse of $T_i(k)$

$$\widehat{T_i(k)^{-1}} = \bar{P}_i \widehat{\tilde{P}_i}(k)^{-1}$$

$$\lambda_{i \text{rec}} = \widehat{T_i(k)}^{-1} \tilde{\lambda}_i = -\bar{P}_i \theta_i - \widehat{T_i(k)}^{-1} \hat{\tilde{s}}_i(k)$$



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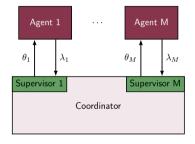
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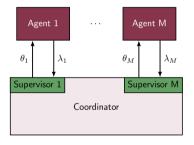
$$\boldsymbol{\lambda}_{i \text{rec}} = \widehat{T_i(k)^{-1}} \tilde{\boldsymbol{\lambda}}_i = -\bar{P}_i \boldsymbol{\theta}_i - \widehat{T_i(k)^{-1}} \hat{\tilde{\boldsymbol{s}}}_i(k)$$





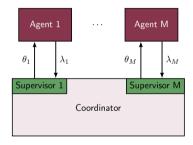
- Detect which agents are non-cooperative
- 2 Reconstruct λ_i and use in negotiation





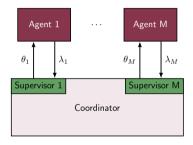
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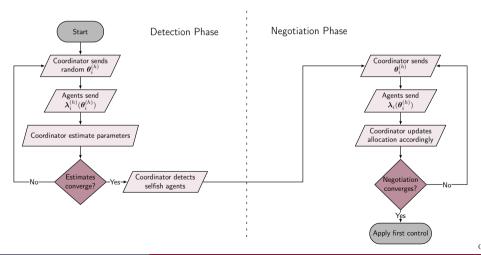


- Detect which agents are non-cooperative
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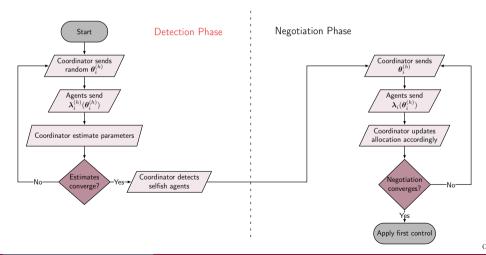
Complete algorithm

RPdMPC-DS

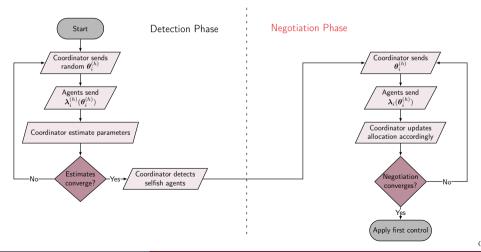


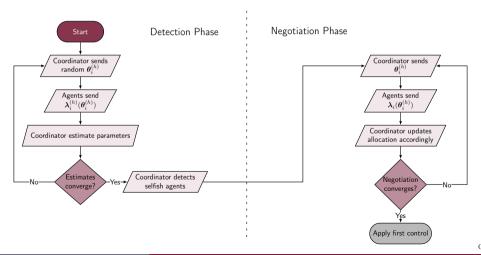
Complete algorithm

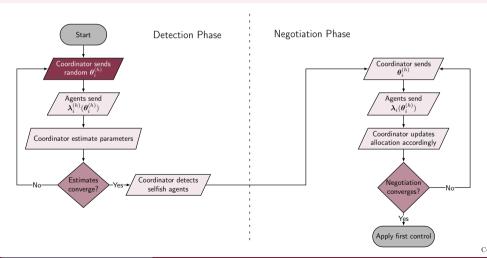
RPdMPC-DS

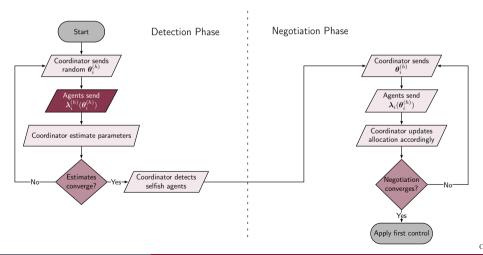


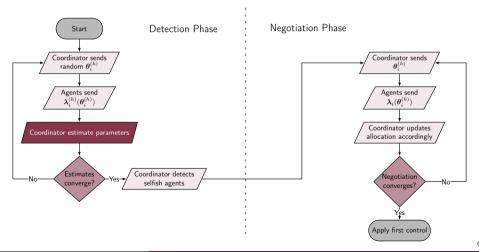


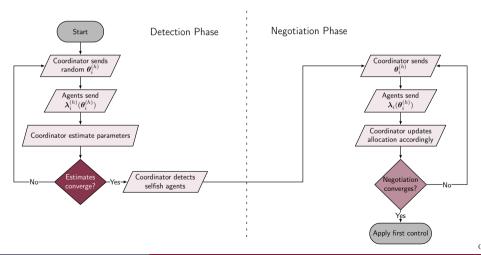


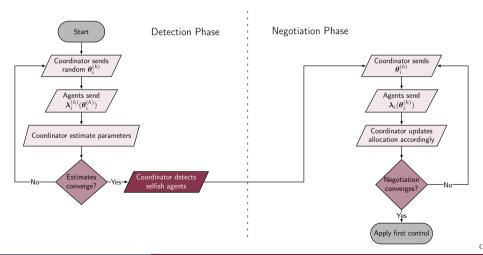


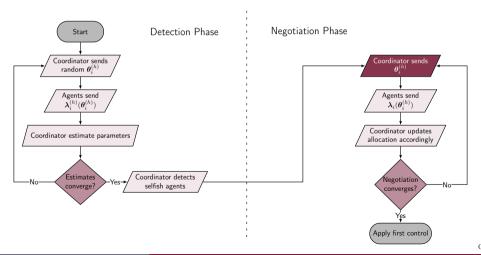


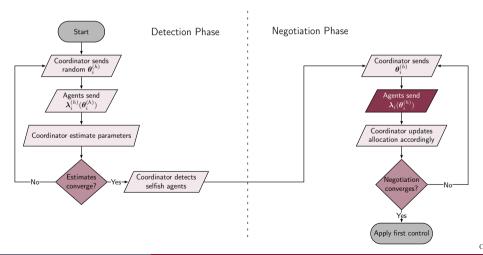




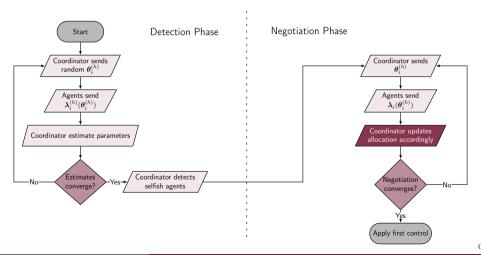


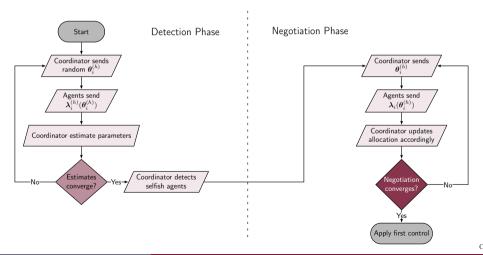


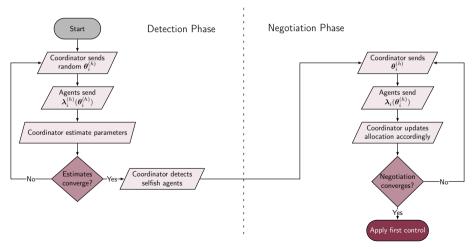














- 4 distinct rooms modeled using 3R-2C
- Initial temperature under 20°C
- Not enough power to achieve setpoint $\left(\sum_{i=1}^4 m{u}_i(k) \leqslant 4 \mathrm{kW} \right)$
- Simulated for a period of 5h
- ZOH $T_s = 0.25h$
- 3 scenarios
 - Nominal
 - 2 Agent I non cooperative from k>6 with T=4*
 - 3 Similar but with secure algorithm



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Temporal

- N Nominal
- S Selflish behavior
- C selfish behavior with Correction



Temporal

- N Nominal
- S Selflish behavior
- C selfish behavior with Correction



Table 1: Comparison of costs J_i^N and J_G^N

Agent	Nominal	Selfish	Selfish + correction
1	103	64	104
П	73	91	73
Ш	100	123	101
IV	132	154	131
Global	408	442	409



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Agent	Nominal	Selfish	Selfish + correction
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Outline

Resilient Primal Decomposition-based dMPC using Artificial Scarcity Relaxing some assumptions Adapting the algorithm Results



Analysing

$$\boldsymbol{\lambda}_{i}[k] = \begin{cases} -P_{i}^{(0)}\boldsymbol{\theta}_{i}[k] - \boldsymbol{s}_{i}^{(0)}[k], & \text{if } \boldsymbol{\theta}_{i}[k] \in \mathcal{R}_{\boldsymbol{\lambda}_{i}}^{n} \\ \vdots & \vdots & \ddots \\ -P_{i}^{\left(2^{n_{\mathsf{ineq}}}-1\right)}\boldsymbol{\theta}_{i}[k] - \boldsymbol{s}_{i}^{\left(2^{n_{\mathsf{ineq}}}-1\right)}[k], & \text{if } \boldsymbol{\theta}_{i}[k] \in \mathcal{R}_{\boldsymbol{\lambda}_{i}}^{2^{n_{\mathsf{ineq}}}-1} \end{cases}$$
(1)



Artificial Scarcity

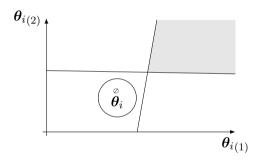


Figure 1: Ball $\mathcal{B}(\overset{\circ}{\theta}_i, r)$.

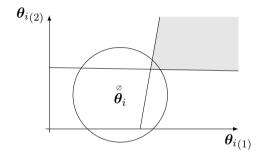


Figure 2: Ball $\mathcal{B}(\overset{\circ}{m{ heta}}_i,r)$ traversing zones.



Expectation Maximization

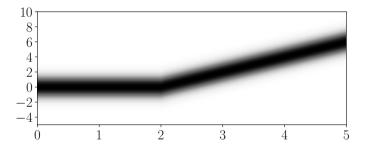


Figure 3: Gaussian Mixture for a 1D PWA function with 2 modes.



Expectation Maximization

Algorithm

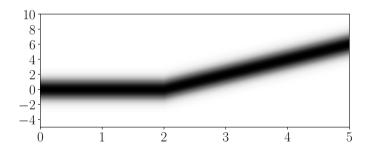
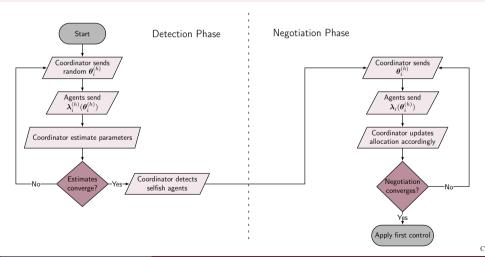
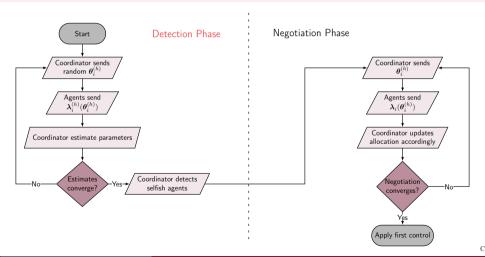
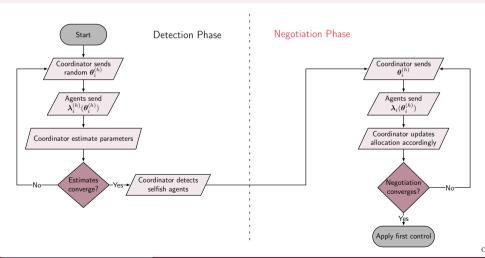


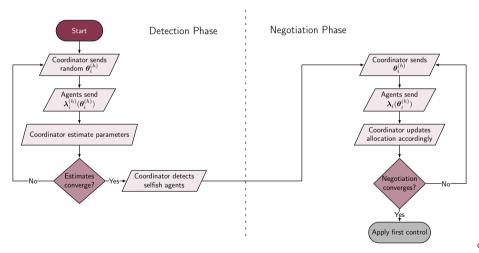
Figure 4: Gaussian Mixture for a 1D PWA function with 2 modes.

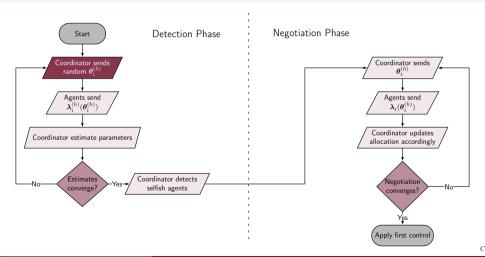


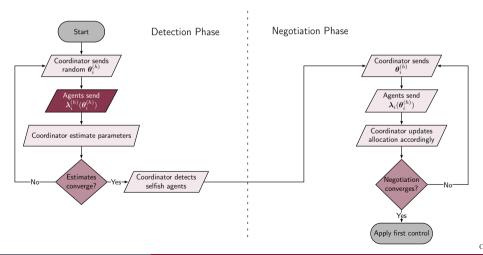


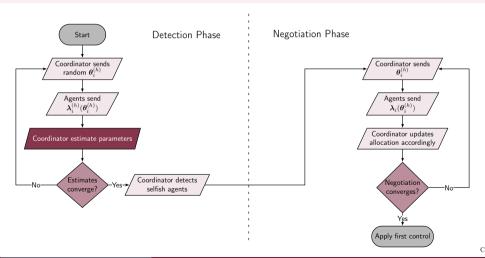


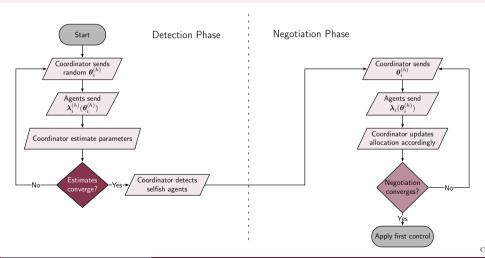


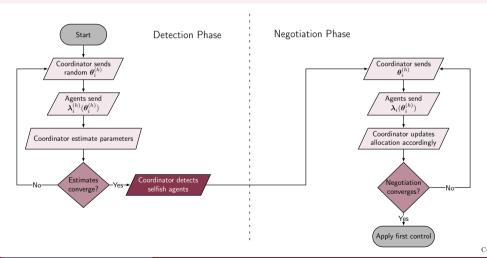


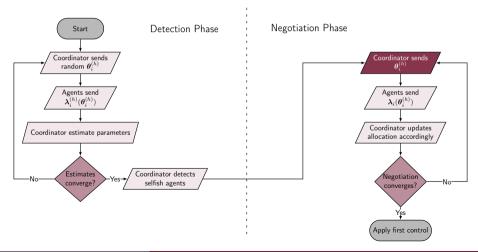


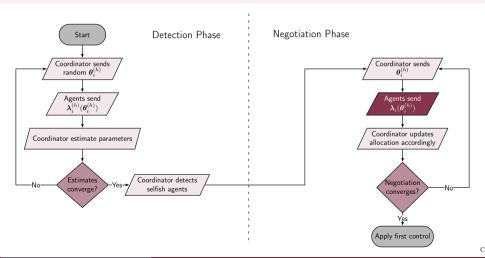


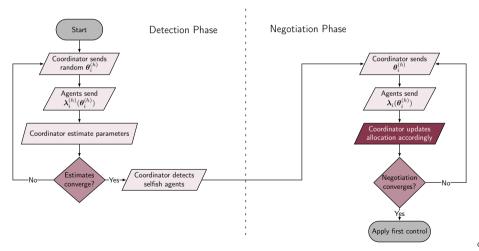


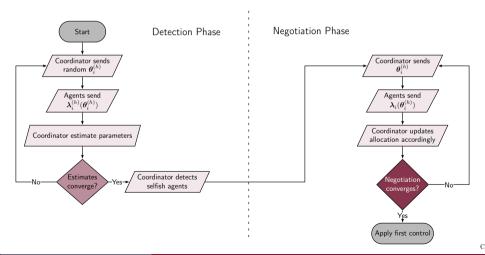


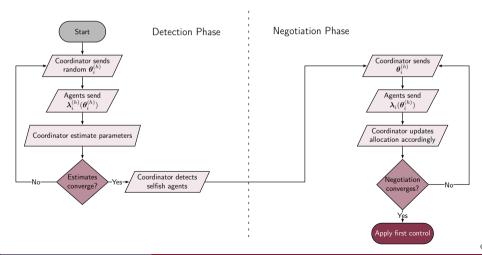














Thank you!

Repository https://github.com/Accacio/thesis



Contact rafael.accacio.nogueira@gmail.com



For Further Reading I



José M Maestre, Rudy R Negenborn, et al. *Distributed Model Predictive Control made easy*. Vol. 69. Springer, 2014. ISBN: 978-94-007-7005-8.



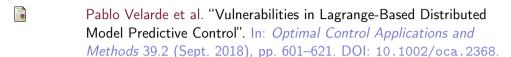
Wicak Ananduta et al. "Resilient Distributed Model Predictive Control for Energy Management of Interconnected Microgrids". In: Optimal Control Applications and Methods 41.1 (2020), pp. 146–169. DOI: 10.1002/oca.2534. URL: https://onlinelibrary.wiley.com/doi/pdf/10.1002/oca.2534.



José M. Maestre et al. "Scenario-Based Defense Mechanism Against Vulnerabilities in Lagrange-Based Dmpc". In: Control Eng Pract 114 (2021), p. 104879. ISSN: 0967-0661. DOI: 10.1016/j.conengprac.2021.104879.



For Further Reading II



Wicak Ananduta et al. "Resilient Distributed Energy Management for Systems of Interconnected Microgrids". In: 2018 IEEE Conference on Decision and Control (CDC). 2018, pp. 3159–3164. DOI: 10.1109/CDC.2018.8619548.

Wicak Ananduta et al. "A Resilient Approach for Distributed MPC-Based Economic Dispatch in Interconnected Microgrids". In: 2019 18th European Control Conference (ECC). 2019, pp. 691–696. DOI: 10.23919/ECC.2019.8796208.



For Further Reading III



P. Chanfreut, J. M. Maestre, and H. Ishii. "Vulnerabilities in Distributed Model Predictive Control based on Jacobi-Gauss Decomposition". In: 2018 European Control Conference (ECC). June 2018, pp. 2587–2592. DOI: 10.23919/ECC.2018.8550239.



Pablo Velarde et al. "Scenario-based defense mechanism for distributed model predictive control". In: 2017 IEEE 56th Annual Conference on Decision and Control (CDC). IEEE. Dec. 2017, pp. 6171–6176. DOI: 10.1109/CDC.2017.8264590.



Pablo Velarde et al. "Vulnerabilities in Lagrange-Based DMPC in the Context of Cyber-Security". In: 2017 IEEE International Conference on Autonomic Computing (ICAC). July 2017, pp. 215–220. DOI: 10.1109/ICAC.2017.53.



Conditions

√ back

One way to ensure this, is to make the original constraint (??) to have at most as many rows as columns, i.e., $\# u_{\text{max}} \leq n_u$, although it may be a little restrictive.



θ dynamics

√ back

$$\boldsymbol{\theta}^{(p+1)} = \mathcal{A}_{\theta} \boldsymbol{\theta}^{(p)} + \mathcal{B}_{\theta}[k]$$

where

$$\mathcal{A}_{\theta} = \begin{bmatrix} I - \frac{M-1}{M} \rho^{(p)} P_{1} & \frac{1}{M} \rho^{(p)} P_{2} & \dots & \frac{1}{M} \rho^{(p)} P_{M} \\ \frac{1}{M} \rho^{(p)} P_{1} & I - \frac{M-1}{M} \rho^{(p)} P_{2} & \dots & \frac{1}{M} \rho^{(p)} P_{M} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{1}{M} \rho^{(p)} P_{1} & \frac{1}{M} \rho^{(p)} P_{2} & \dots & I - \frac{M-1}{M} \rho^{(p)} P_{M} \end{bmatrix}$$

$$\mathcal{B}_{\theta}[k] = \begin{bmatrix} -\frac{M-1}{M} \rho^{(p)} s_{1}[k] + \frac{1}{M} \rho^{(p)} s_{2}[k] \dots - \frac{1}{M} \rho^{(p)} s_{M}[k] \\ \frac{1}{M} \rho^{(p)} s_{1}[k] - \frac{M-1}{M} \rho^{(p)} s_{2}[k] \dots - \frac{1}{M} \rho^{(p)} s_{M}[k] \\ \vdots & \vdots \\ \frac{1}{M} \rho^{(p)} s_{1}[k] + \frac{1}{M} \rho^{(p)} s_{2}[k] \dots - \frac{M-1}{M} \rho^{(p)} s_{M}[k] \end{bmatrix}$$

