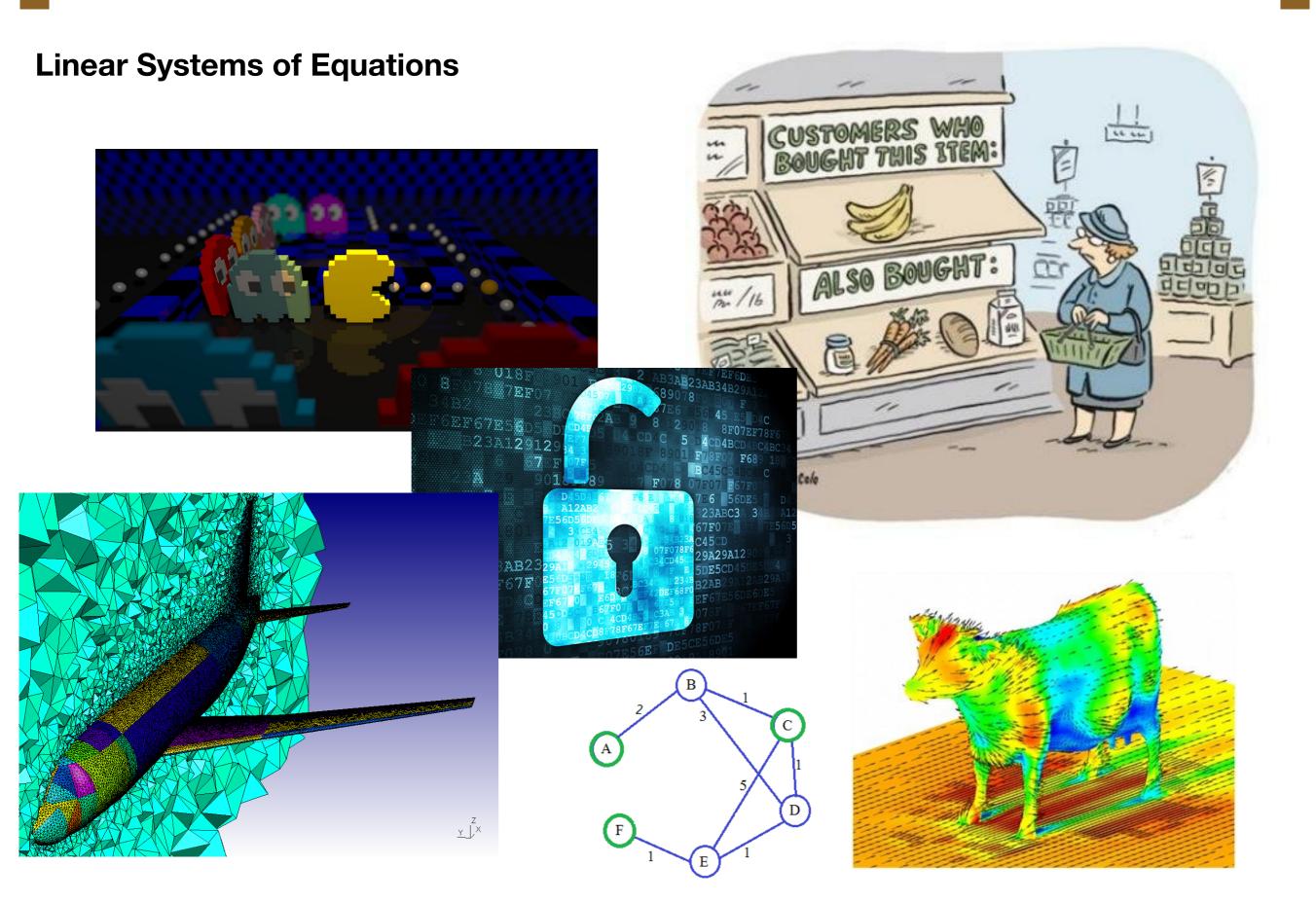


Introduction to Scientific Computation Lecture 2 Fall 2019

Liners Systems of Equations, SVD MatVec operations



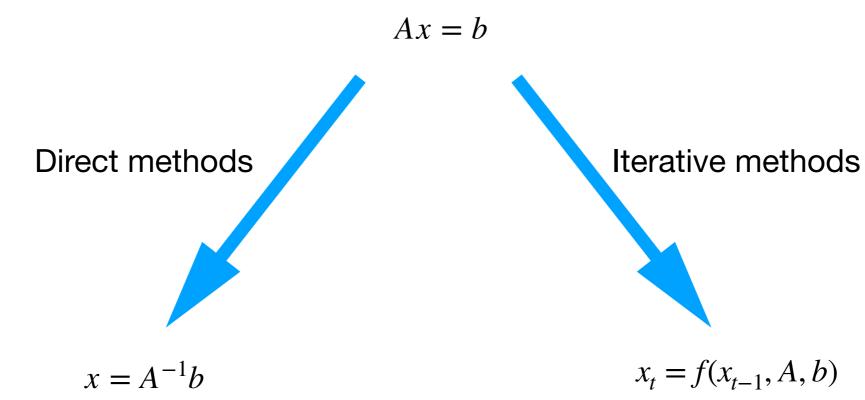


How to solve this?

$$Ax = b$$



#### How to solve this?



Exact solution, except for rounding errors, fixed number of computational steps

Often work better in practice, due to care about numerical stability, iteratively improve solution estimate



### **Direct methods**

If all linear equations are uncoupled:



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$$x_i = \frac{b_i}{a_{ii}}$$

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Can we diagonalize every matrix?

$$P^{-1}AP$$

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Can we diagonalize every matrix?

Keyword: Jordan Normal Form

$$P^{-1}AP A = \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix}$$



### **Upper triangular** matrix:

$$A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ & a_{22} & \ddots & \vdots \\ & & \ddots & \vdots \\ & & & a_{nn} \end{bmatrix}$$

$$a_{ij} = 0, \forall i > j$$

$$a_{nn}x_n = b_n, \implies x_n = \frac{b_n}{a_{nn}}$$

$$a_{n-1,n-1}x_{n-1,n-1} = b_{n-1} - a_{n-1,n}x_n, \implies x_{n-1} = \frac{b_{n-1} - a_{n-1,n}x_n}{a_{n-1,n-1}}$$

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Quiz: What is the final complexity?

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Quiz: What is the final complexity?

**Answer**:  $O(N^2)$ 



#### **Gaussian Elimination**

Exact solution, except for rounding errors

$$A = LU$$

$$Ly = b$$

$$Ux = y$$

$$PA = LU$$

Not unique

Problems when A is singular

Solution is LUP - additional permutation matrix which aims to maximize absolute values of diagonal elements

Finding LU (LUP) has complexity  $O(N^3)$ 



**Special case: SPD matrices** 

**Symmetric positive definite**: if  $A = A^T$  and  $x^T A x > 0, \forall x \neq 0$ 

$$A = LL^*$$

Two times less memory than LU

### **Conditioning**

Relative error: 
$$\frac{||x - \hat{x}||}{||x||}$$

$$r = b - A\hat{x} = Ax - A\hat{x} = A(x - \hat{x})$$

$$x - \hat{x} = A^{-1}r \implies ||x - \hat{x}|| = ||A^{-1}r|| \le ||A^{-1}|| ||r||$$

$$\frac{||x - \hat{x}||}{||x||} \le ||A^{-1}|| ||r|| \frac{||A||}{||b||}$$

Note:

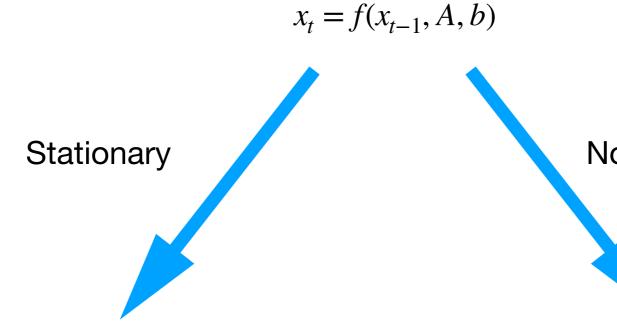
$$A\hat{x} = b + \Delta b \implies ||\Delta b|| = ||r||$$

Condition number:  $k(A) = ||A||||A^{-1}||$ 

$$\frac{||x-\hat{x}||}{||x||} \le k(A) \frac{||\Delta b||}{||b||}$$
 the relative error in the solution is bounded by the condition number of the matrix A times the relative error in the residual.



#### **Iterative methods**



- Jacobi
- Gauss-Seidel

 $Mx_t = Nx_{t-1} + b, A = M - N$ 

• SOR

Non stationary

- Gradient descent
- Conjugate gradients (for SPD)



### **Stationary methods**

$$Mx_t = Nx_{t-1} + b, A = M - N$$

We need initial guess for the solution:  $x_0$ 

Error is:  $e_0 = x - x_0$ 

$$r_0 = Ae_0 = b - Ax_0$$

$$x = x_0 + e_0 = x_0 + A^{-1}r_0$$

Set  $M\widetilde{e} = r_k$ 

Then 
$$x_t = x_{t-1} + \widetilde{e}$$

We do not know the exact solution (if we know, why do we care?) and cannot compute

We can compute

Preconditioner: M

$$A = D + E + F$$

Jacobi: M = D

Gauss-Seidel: M = D + E

SOR:  $M = D + \frac{1}{\omega}E$ 

$$A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{bmatrix} = \begin{bmatrix} 0 & 0 & 0 & 0 \\ a_{21} & 0 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & 0 \end{bmatrix} E$$

$$+ \begin{bmatrix} 0 & a_{12} & \dots & a_{1n} \\ 0 & 0 & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 0 \end{bmatrix} F$$

$$+ \begin{bmatrix} a_{11} & 0 & \dots & 0 \\ 0 & a_{22} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & a_{nn} \end{bmatrix} D$$



## **Nonstationary methods**

$$x_t = x_{t-1} + \alpha_t p_t$$



### **Gradient descent**

$$x_t = x_{t-1} + \alpha_t p_t$$

$$p_k = r_k$$

or

$$M\widetilde{e} = r_t$$

$$M = \alpha_t I$$

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### **Conjugate gradient (SPD only)**

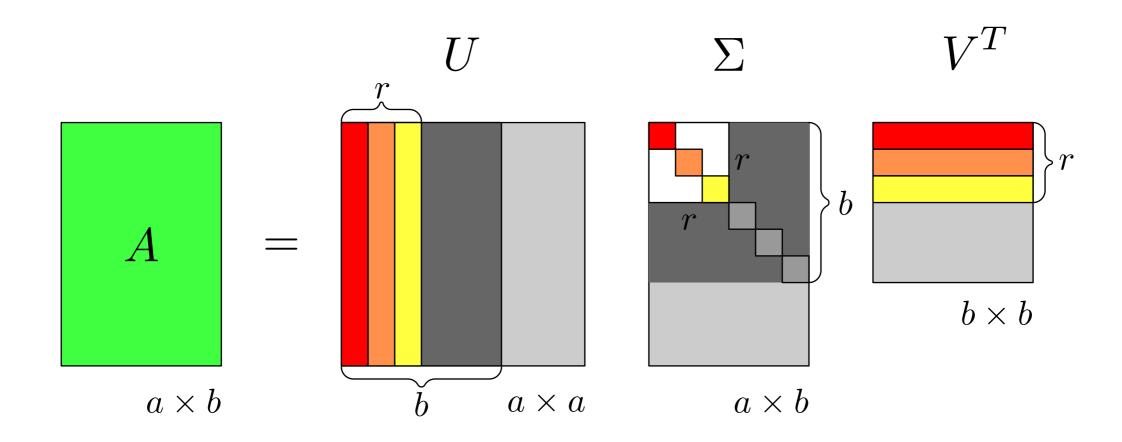
$$x_t = x_{t-1} + \alpha_t p_t$$

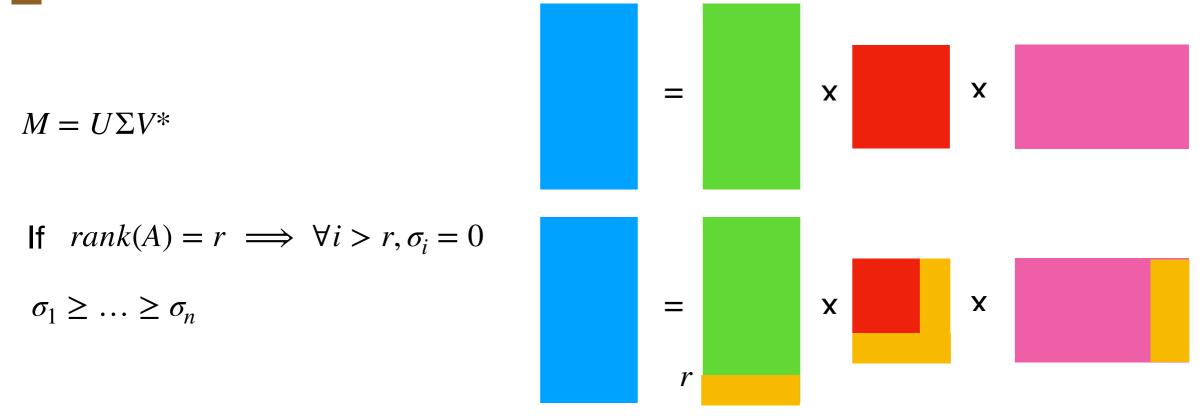
$$p_t = r_t - \sum_{i \le t} \frac{p_i^T A r_t}{p_i^T A p_i} p_i$$

$$\alpha_t = \frac{p_t^T r_t}{p_t^T A p_t}$$



## SVD





Eckart-Young theorem: best low-rank approximation can be computed with SVD

$$\min_{rank(B)=r} ||A - B||_2 = ||A - A_r||_2 = \sigma_{r+1}$$

r - rank approximation is setting  $\forall i > r, \sigma_i = 0$