# **ETH** zürich

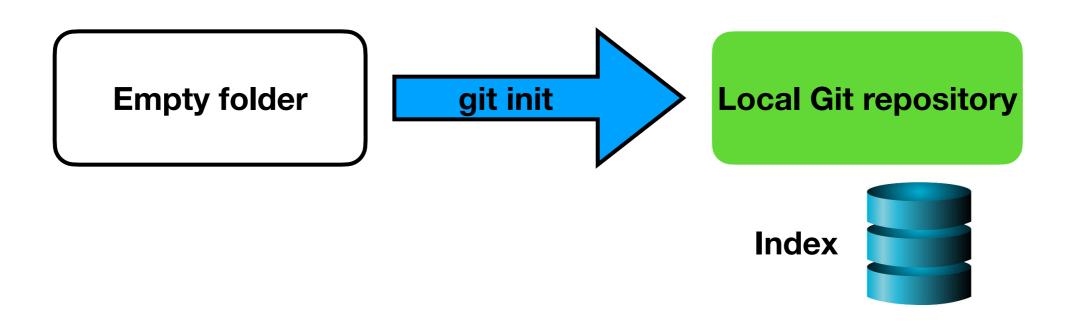


Introduction to Scientific Computation Lecture 1 Fall 2020

Git v2, Complexity, Floating-point arithmetic Numerical stability

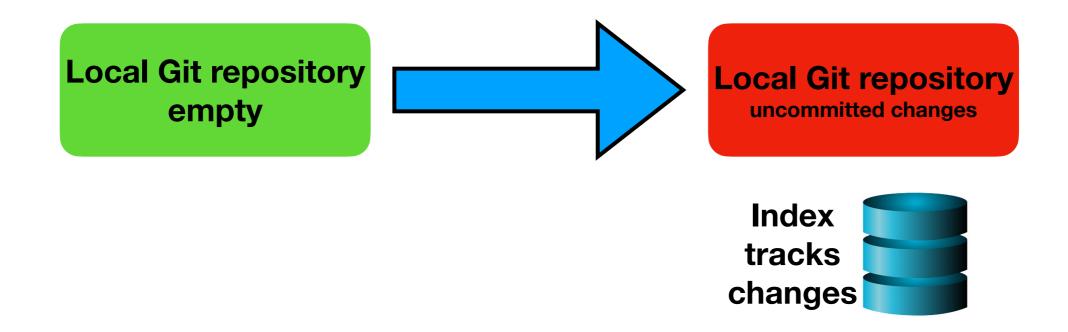


Git v2



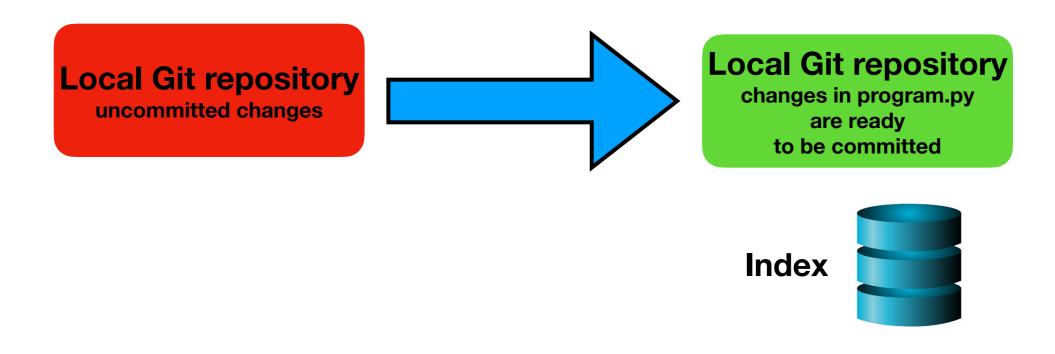


### **Create program.py consisting of K lines**



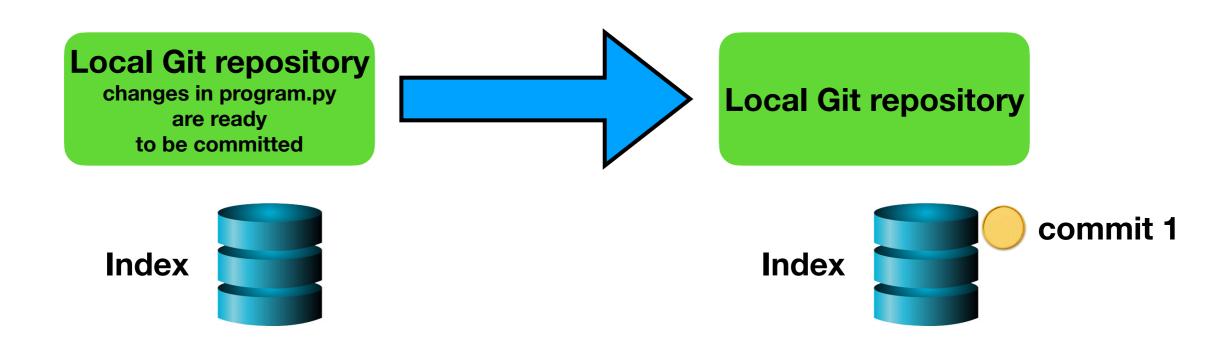


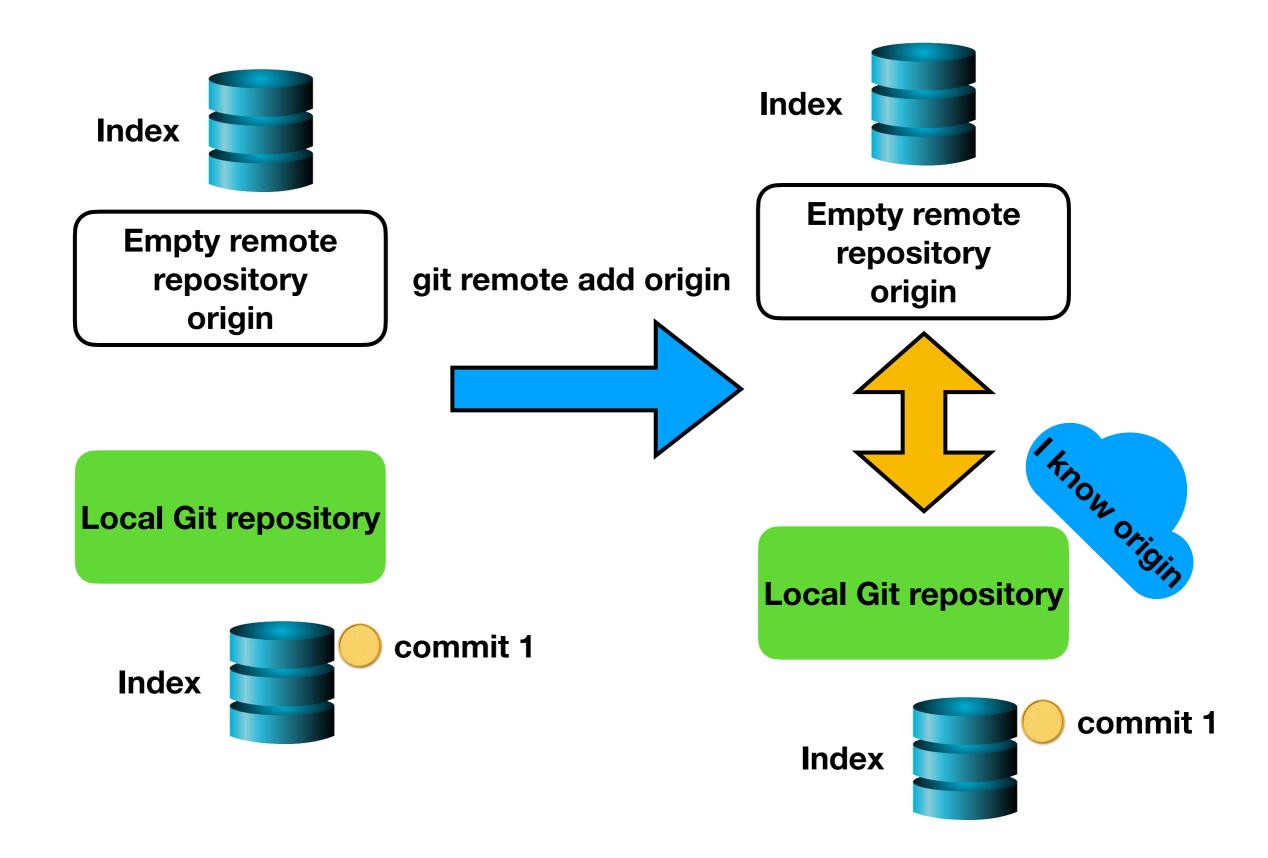
## git add program.py

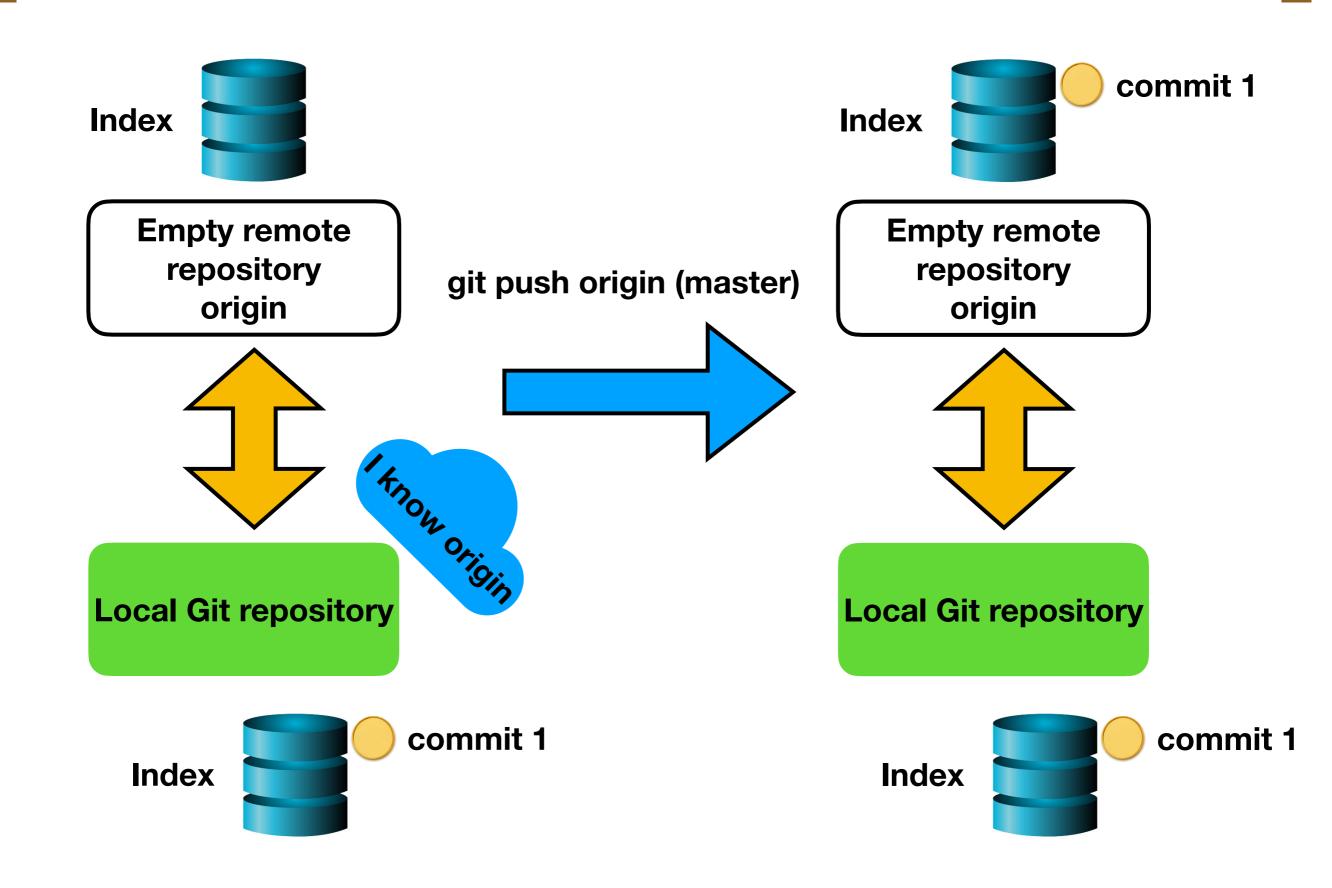




### git commit -m "implemented program"

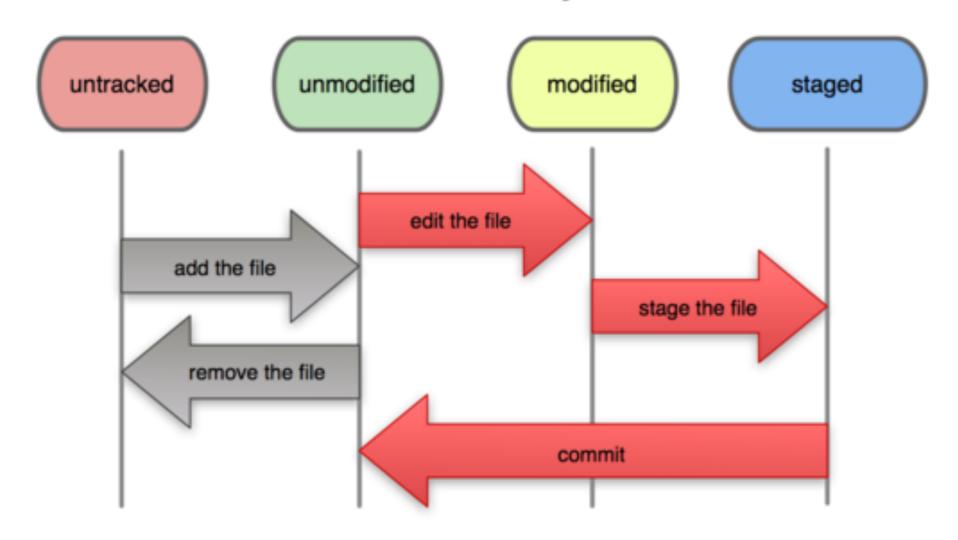








# File Status Lifecycle



## (c) <a href="http://git-lectures.github.io">http://git-lectures.github.io</a>

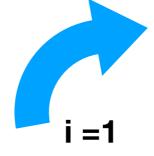




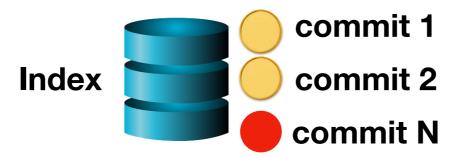
git checkout commit i

**Local Git repository** 

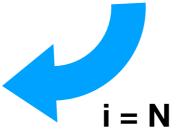


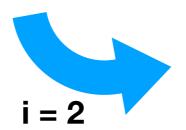






**Local Git repository** 





**Local Git repository** 







**Check here for more examples:** 

http://git-lectures.github.io



One of definitions: it is a measure of computational cost In order to give a more formal definition quite a big formalism should be introduced (see recommended materials for this week). We will operate with a concept.



The notion of complexity is about how good or bad is the algorithm. We want to think in an abstract way, independent of:

- implementation details (Python, C++, matlab)
- hardware (stm32, intel Pentium D, Intel Core i7-8700K)
- etc and whatnot

Let us use constant time operation (elementary operations).

Remember: real runtime doesn't matter (does it ?)

$$3 + 5 =$$

$$\begin{array}{r}
0011 \\
0101 \\
\hline
1000
\end{array}$$

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Remember: real runtime doesn't matter (does it ?)

$$3 + 5 =$$

$$\begin{array}{r}
 0011 \\
 0101 \\
\hline
 1000
\end{array}$$

$$T(n) = c \cdot n$$

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$$T(n)=c\cdot n$$
  $c$  - is a cost for summation of one bits  $c_1$ - for cpu1  $c_2$ - for cpu2



In order to get rid of c we will use **Asymptotic complexity** with big-O notation.



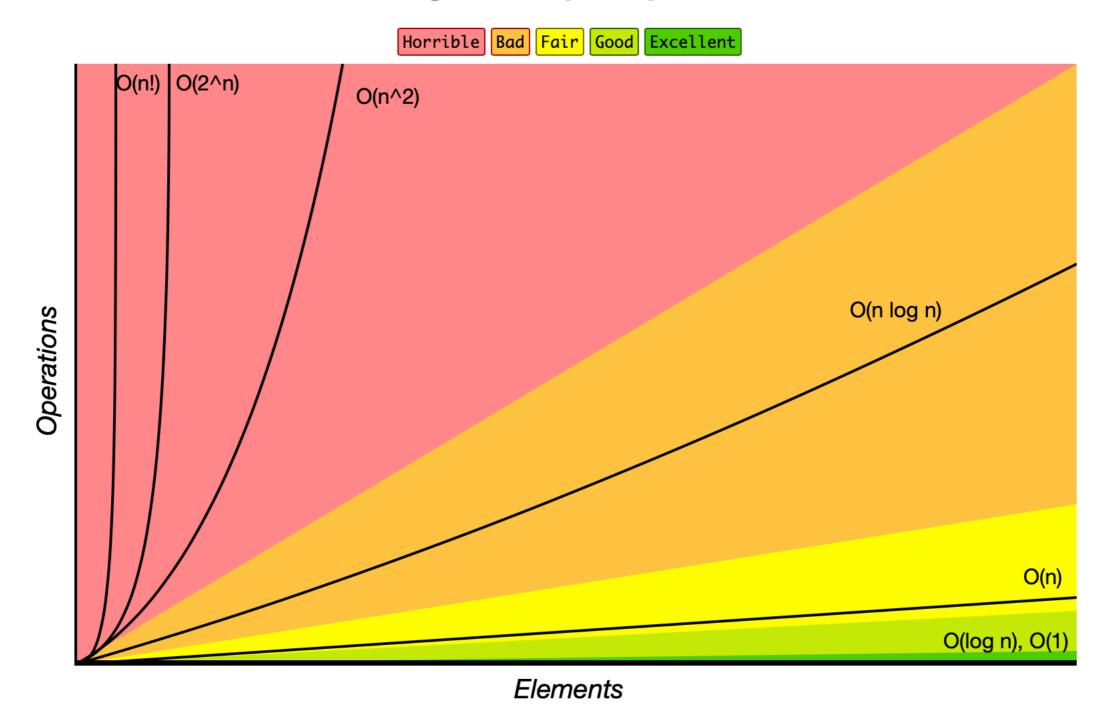
In order to get rid of c we will use **Asymptotic complexity** with big-O notation.

<u>Definition:</u> for any monotonic functions f(n) and g(n) from positive integers to the positive integers f(n) = O(g(n)) when there exists constants c > 0 and  $n_0 > 0$  such that

$$f(n) \le c \cdot g(n)$$
, for all  $n \ge n_0$ 

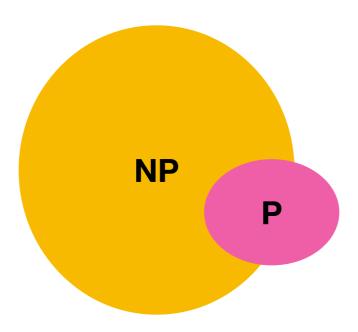


## **Big-O Complexity Chart**



(c) http://www.bigocheatsheet.com





???



#### Representation of numbers

#### **Fixed point**

The most straightforward format for the representation of real numbers is **fixed point** representation, a.k.a **Qm.n** format.

**Qm.n** number is in the range  $[-2^m, 2^m - 2^{-n}]$  with the resolution  $2^{-n}$  requires m + n + 1 bits for storage.



#### **Representation of numbers**

### **Floating point**

We are mostly interested in **floating point numbers** represented as e.g.

$$1.2345 = \underbrace{12345}_{\text{significand base}} \times \underbrace{10}_{-4}$$

**IEEE 754** 

# **ETH** zürich

#### Single and double precision

$$X = \pm 1.\overline{b_1 b_2 \dots b_K} \cdot 2^e$$

$$X \in [X - \Delta X, X + \Delta X],$$

#### **Single**

- **Double**

#### **Absolute accuracy**

$$\Delta X \le \frac{1}{2} 2^{-K} * 2^e \le |X| \cdot 2^{-K-1}$$

#### **Relative accuracy**

$$\frac{\Delta X}{X} \le 2^{-K-1}$$

The **relative accuracy** of single precision is  $10^{-7} - 10^{-8}$ , while for double precision is  $10^{-14} - 10^{-16}$ .

A float32 takes 4 bytes, float64, or double precision, takes 8 bytes.

These are the only two floating point-types supported in hardware.

You should use double precision in CSE and single on GPU/Data Science.

Format	Total bits	Significand bits	Exponent bits	Smallest number	Largest number
Single precision	32	23 + 1 sign	8	ca. 1.2 · 10 <sup>-38</sup>	ca. 3.4 · 10 <sup>38</sup>
Double precision	64	52 + 1 sign	11	ca. 5.0 · 10-324	ca. 1.8 · 10 <sup>308</sup>



#### **Machine zero exists**

```
import numpy as np

n = 1.0
d = np.inf
i = 0

while d > 0:
    d = n - n / 2
    n = n / 2
    i += 1

(1 / 2)**(i - 2)
5e-324
```

```
import numpy as np

n = 1.0
d = np.inf
i = 0

while d > 0:
    d = n - n / 2
    n = n / 2
    i += 1

(1 / 2)**(i - 1)
```

0.0



# See code