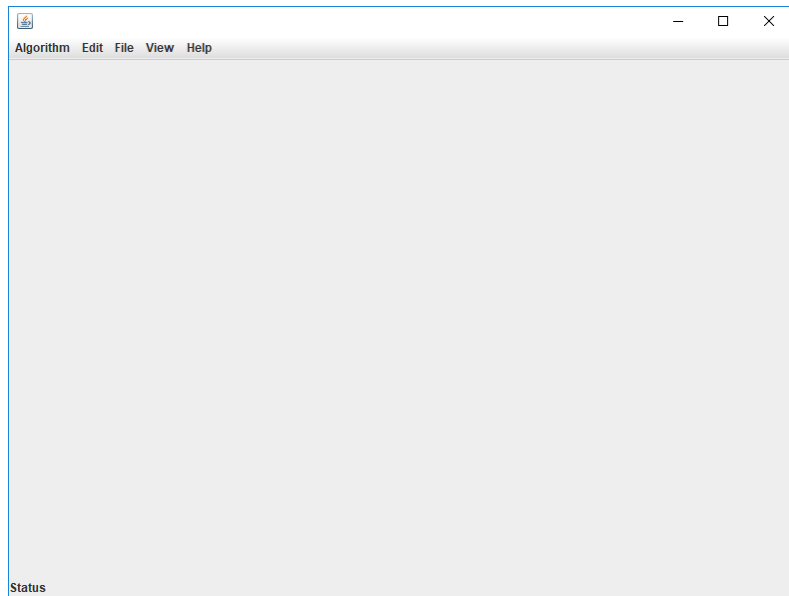


fancy software name

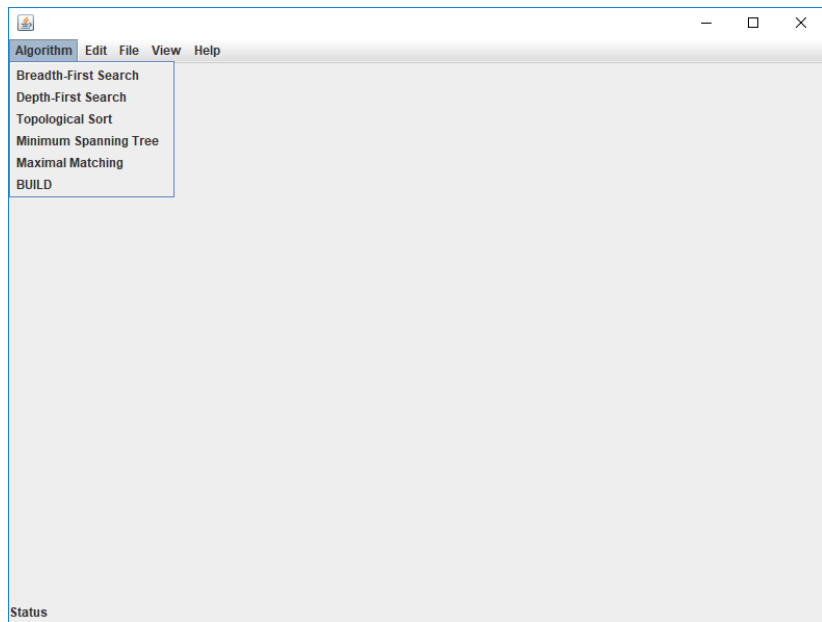
Anica Hoppe, Sonja Türpitz

May 2, 2018

# GUI



# Algorithms



# Breadth/Depth-First Search

Input:  $G = (V, E)$ ,  $s \in V$

$Q, F \leftarrow \emptyset$

$R \leftarrow \{s\}$

▷ visited vertices

**for all**  $(s, v) \in E$  **do**

$Q \leftarrow Q + (s, v)$

**while**  $Q \neq \emptyset$  **do**

$(v, w) \leftarrow \text{choose}(Q)$

▷  $Q$  is queue or stack (BFS/DFS)

**if**  $w \notin R$  or  $v \notin R$  **then**

$R \leftarrow R + w + v$

$F \leftarrow F + (v, w)$

**for all**  $(w, x) \in E$  **do**

$Q \leftarrow Q + (w, x)$

**return**  $(V, F)$

source: lecture graphtheory by Prof. Dr. Marc Hellmuth

# Topological Sort

$L \leftarrow$  empty list ▷ will contain the sorted elements

$S \leftarrow$  set of all vertices with no incoming edge

**while**  $S \neq \emptyset$  **do**

    remove a vertex  $n$  from  $S$

    add  $n$  to *tail* of  $L$

**for all** vertices  $m$  with an edge  $e$  from  $n$  to  $m$  **do**

        remove edge  $e$  from the graph

**if**  $m$  has no other incoming edge **then**

            insert  $m$  into  $S$

**if** graph has edges **then**

**return** error

▷ graph has at least one cycle

**else**

**return**  $L$

source: [https://en.wikipedia.org/wiki/Topological\\_sorting](https://en.wikipedia.org/wiki/Topological_sorting)

# Minimum Spanning Tree

```
 $F \leftarrow \emptyset$   
for  $i = 1$  to  $|E|$  do  
    if  $(V, F + e_i)$  is acyclic then  
         $F \leftarrow F + e_i$   
return  $(V, F)$ 
```

source: lecture graphtheory by Prof. Dr. Marc Hellmuth

# Maximal Matching

$M \leftarrow \emptyset$

$E' \leftarrow E$

▷ not visited edges

**while**  $E' \neq \emptyset$  **do**

    choose  $e = (u, v) \in E'$

$M \leftarrow M + (u, v)$

    delete all edges from  $E'$  that are incident to  $u$  or  $v$

**return**  $M$

source: lecture discret optimization by Prof. Dr. Marc Hellmuth,  
own

# BUILD

## BUILD

```
1: INPUT: Set of triples in  $\mathcal{R}$ , leaf set  $\mathcal{L}$ .
2: OUTPUT: A rooted, phylog. tree distinctly leaf-labeled by  $\mathcal{L}$  consistent with
   all rooted triplets in  $\mathcal{R}$ , if one exists; otherwise null.
3: compute  $G(\mathcal{R}, \mathcal{L})$ 
4: compute connected components  $C_1, \dots, C_s$  of  $G(\mathcal{R}, \mathcal{L})$ 
5: if  $s = 1$  and  $|\mathcal{L}| = 1$  then
6:   return tree  $\simeq K_1$ 
7: else if  $s = 1$  and  $|\mathcal{L}| > 1$  then
8:   return null
9: else
10:  for  $i = 1, \dots, s$  do
11:     $T_i = \text{BUILD}(\mathcal{R}|_{V(C_i)}, V(C_i))$ 
12:  end for
13:  if  $T_i \neq \text{null}$  for all  $i = 1, \dots, s$  then
14:    attach all of these trees to a common parent node and let  $T$  be the
    resulting tree; else  $T = \text{null}$ .
15:  end if
16: end if
```



# Requirements

## properties of the graph

- no coloring of edges or vertices
- no weights of edges or vertices  
→ not needed for algorithms
- mainly undirected graphs
- vertex label  
→ for BUILD

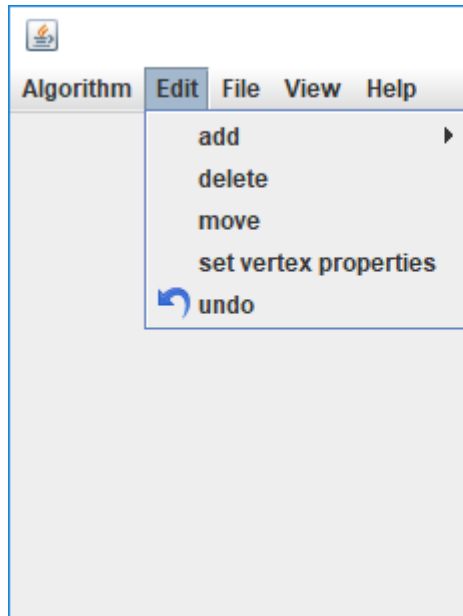
~> as easy as possible

# Requirements

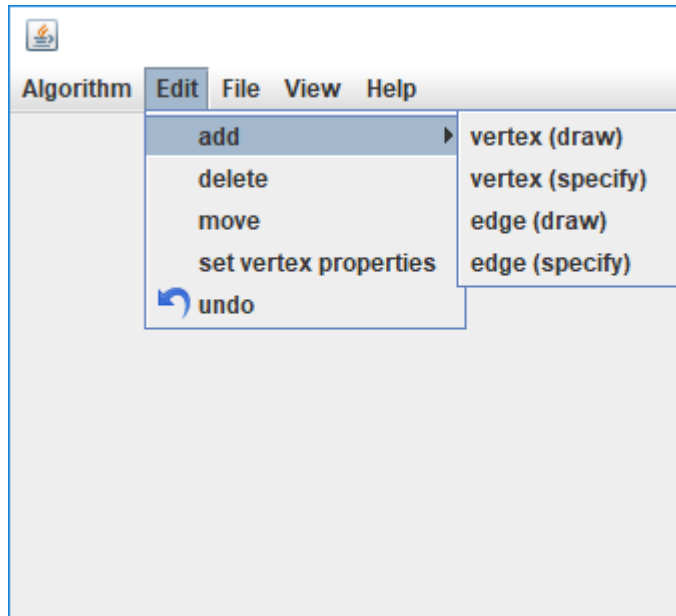
## format of textfile

- |                         |   |
|-------------------------|---|
| 1                       | 'u' or 'd' (undirected or directed)       |
| 2                       | Integer $n$ (number of vertices)          |
| 3                       | Integer $m$ (number of edges)             |
| $4 - (n + 4)$           | vertex list (name, label) for each vertex |
| $(n + 5) - (m + n + 5)$ | edge list (u,v) for each edge             |

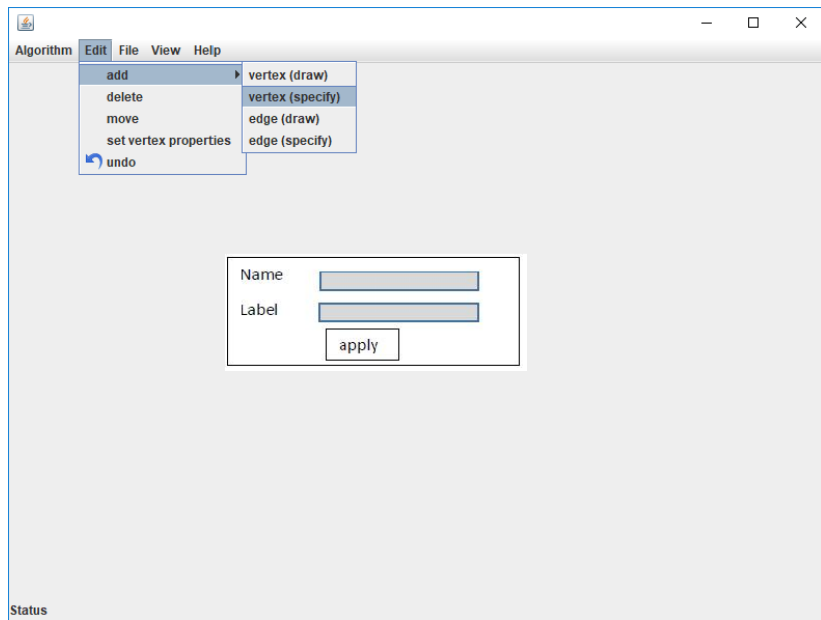
# Menues



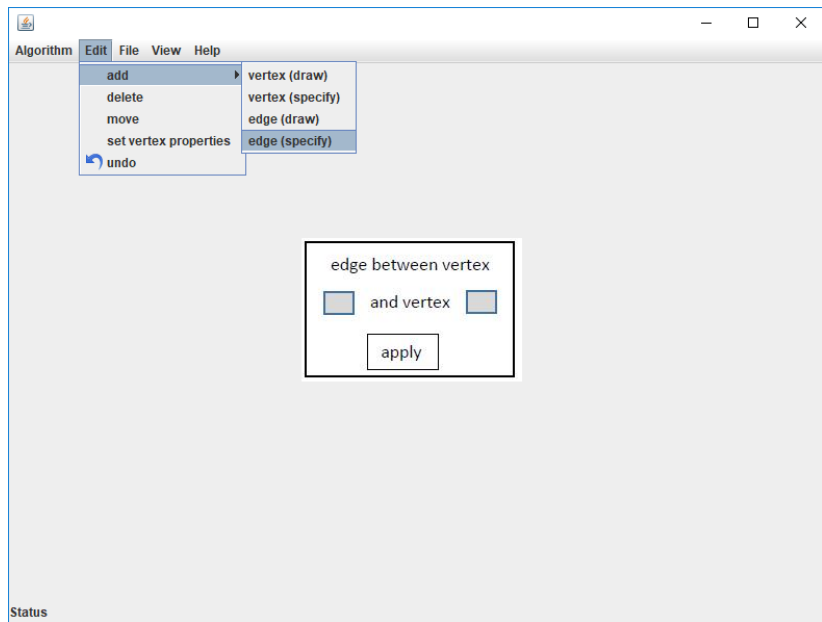
# Menues



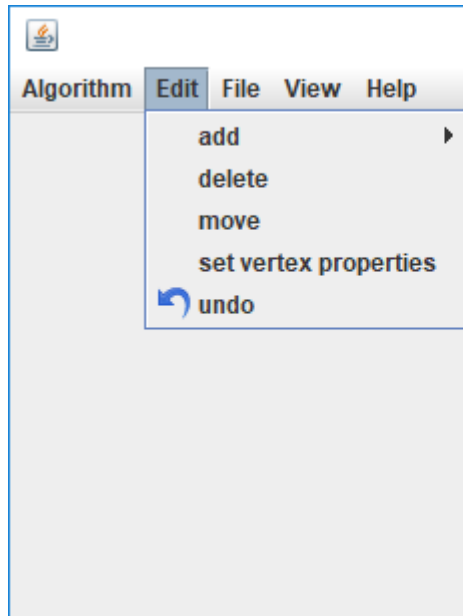
# Menus



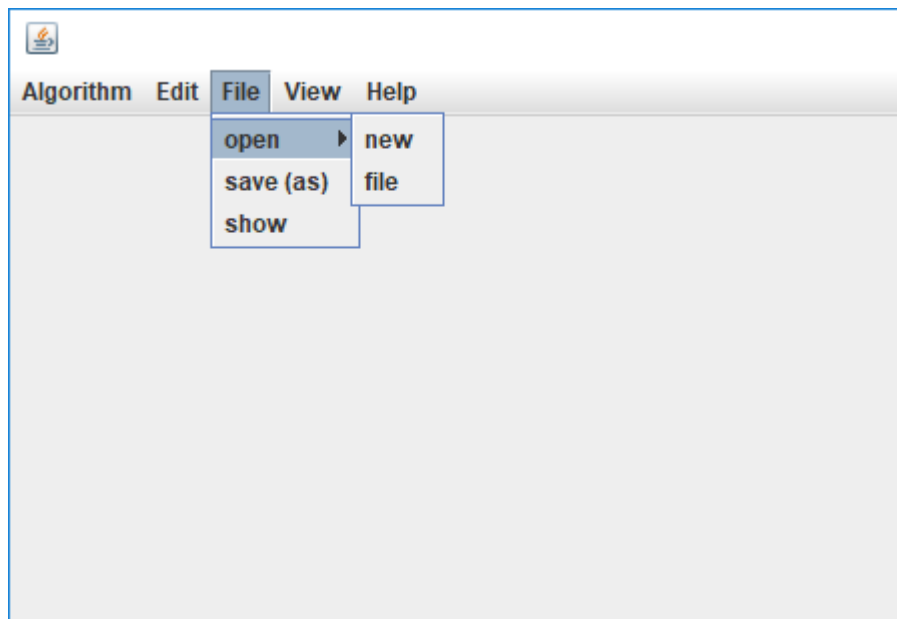
# Menus



# Menues



# Menues





# Menues

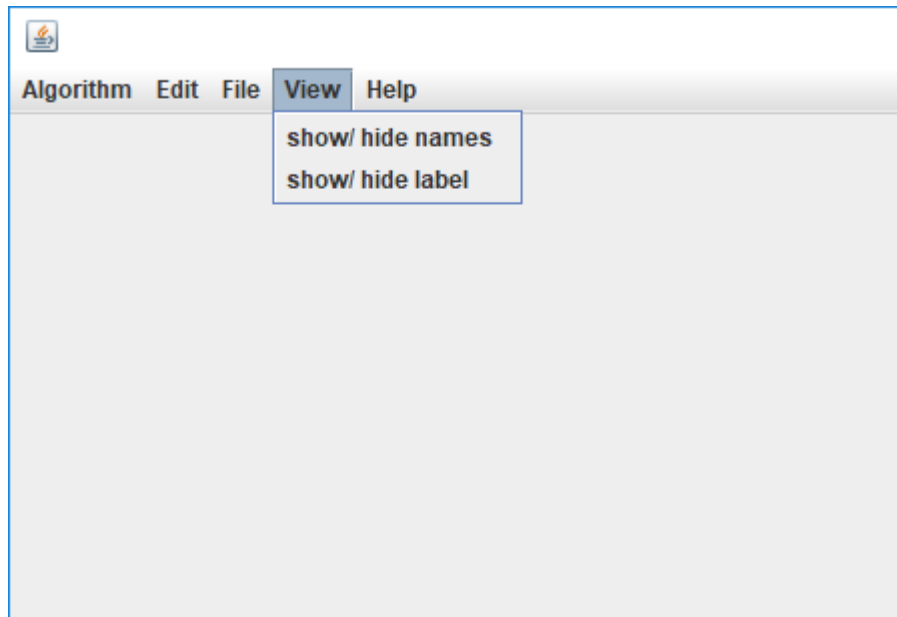
$G=(V,E)$

$V=...$

apply

cancel

# Menus



Thank you for your attention!