10/24

## Assignment -1

PAGE NO : / DATE: / / G

what is No SQL? Explain briefly about a great diagram?

No SQL databases are mon relational databases

designed to handle large volume of data, distributed

data starage and high performance requirements.

NO SQL databases are defectedly popular imapplication

that require scalability, quits steration and

efficient handling of by data.

Aggregate data models organize data into "aggregation" which are collections of data directed as single units Aggregates simplify data distribution across Clusters in 10594 databases as the data often needs to be moved on processed together. Too example, on "Order" in an e Commerce system could include customer details ordered items and payment information within one aggregate.

PAGE NO.: 2 DATE: / / This model is especially helpful fordistributed data systems because each aggrugate con manage asa single unit, which supports lastly easier data parlitioning and replication. [ Coustomer regaregate ] [Order Augregate] - - - [order I terms] shipping Address, Payment between J each aggregate can be stared and managed independent making it suitable for large distributed databases where data consistency is crucial MATERIAL PROPERTY AND ADDRESS OF THE PARTY AND

PAGENO : 3 Laplain beriefly about impedance mismately with a next diagram? Empedance mismatch refers to the conflict that arises due to the difference in how data is prepresented in object - oriented programming (oop) and relational databases. data is structured as complex objects, often Containing nexted attributes and hierarchial relationships, while relational databases store data in lables with rows and columns with each table holding flat, tabular data. The impedance mismatch Occurs because: 1. Objects in Code have altributes that con reference other objects, creating complex nexted data structure, but relational databases require data to be normalized and stored across multiple tables

2. Converting these objects to teables

requires on Object - Relational resplany layer, which can lead to performance and consistency essues.

[ Object - Oriented Structure] <--> [ Relational Structure]

(e.g., Hibernate)

In this diagram:

- The Object - Oriented Structure represents

Complex mested relationships, which are more

natural in applications.

- The Relational Structure represents mormolized data in tables which are better suited for storage in relational totalous.

- The ORM layer serves as a bridge, translating between 8 byiets in Code and rows / Columns in the database although it may introduce over head and complexity

white a short make on a) consequences of Aggregate Orientation b) Key value data model () Document data model d) Column family stores e) Greath dutable a) Consequences of Aggregate Orientation · Simplified Distribution - Aggregates are well-sewed for distribution because they are self contained units of data, making a it earlier to portition data across modes in a cluster . Pota Deeple whom - storing related data together con lead to a data desplication if the some data is a part of multiple oggregates · Consistency Challenges - Since oggengales are designed to oberate as a singl untils b) key value Porta Model. · Description - In the key-value model data is story and collection of my - value pairs where each key is conique and make directly to a specific while \* storage as values can be only date type

PAGE NO.: 6 9 Document Poto rodel · Description: Document dolaboses store data in documents, typically using formuly like I son , Blon or XML · Advantages : Supports serve structured data and nested data making it adaptable for Changing data requirements d) Column Formely Stores · Description - Column family stores Or goneige date in Columns ond rows similar to relational tables but more flexible schema definition bote is stored in Collemn families . Advantages : Optimized for queries on specific columns and allows for efficient storage of store data. e) Greaph dotabase · Description: Graph databases use modes, edges and properties to represent and stone data, making them ideal for applications with compar relationships

PAGE NO. 7

Advantages - Efficiently handles highly connected data and complex queries einvolwing reductions what are schemaless dateboses? Explain in details? Schemaless Destabase also known as schema - free one databases that do not requires a predefined schema to define the directure of data.

Unlike traditional relational databases, which require a sor rigid rehema with specified tables, Columns and datatypes.

characteristics of Schemaless Dotaboses

阳

- Thereble Data Structure: Schemaless databases allow each record to have a conseque.

  Structure meoning new fields combe added to documents without attering existing database migrations
  - 2 Rapid development and I terahon & Developers con modify and evalue the database scheme directly in the application code, which accelerates development particularly in Agile environments where regularments Change frequently.

PAGE NO.: 8 DATE: / / Advantages of Schema less Patakuses 1) Ad aptability - At application requirements change developers can add, nemone, or modify fields without restructeuring the database allowing anich adaptation. 2) Efficiency for nodern Applications money modern applications handle large Volume of data that do not fit well into riged schemas. 3.> Speed of Deployment; without the need for database migrations, capitates to data models Can be deployed faster, reducing downtone and accelerating feature delivery

PAGE NO.: 9

(5)

what are distribution models? Explain two paths

Refers to strategies used to distribute data ocross multiple modes or servers within a network. Distribution improves performance sudshifty, availability and fault tolerance in large scale applications. In destributed database systems, totally can be managed across several to servers, ensuring that If one mode fails, others can continue to provide access to data.

Two path of bata Distribution

1. Republication:

"Vescription - Reflication involves

Copying data across multiple modes, ensuring
that data is available on more than one services.

This approach emprous data availability and
foult holorome because if one service fails
the data still be accessed from another server

12 Shouling; · Description - Shorting is a form of horizontal fortitioning where data isdivided into smaller distinct justs ( called shords ), each shord is stored on a st separate services. Foot somes (made) holds only a subaset of the data rather than a copy of the entire dataset Benefits: Shording allows for distributed write operations as each short con handle write request indepenatly This helps with realability because as data grows, new modes con added to store additional shords

wouthe a short motes on a) single server [6] 6) Combining shorting and Repaliation distribution data model a) Single Server · Description : Ina single server setup, all data rasides on a osignagle mochine. This is the simplest database Configuration where the somes handles all read and werete operations · Advortages : A single server setup is easy to so up monoge od reason about, making it ideal for applications with minim data storage and orcers needs. It is also eliminates the complete of distributed systems, such as network latericy and data consistency. b) Cornbining Shording and replication Risteribution Data model Description + Combining and replication involves for tetioning the dala into multiple shords and neplecting each should across several nodes

is suring that each short has multiple copies for redunday. · Advantages - Scalability: Shording allows the sigstem to handle large datasets and distribute write operations across modes, reducing the lood on ony single mode. - Fault Toleronce :- Replication ensures data redundancy, so if one made fails, the data constill be accessed from another replica of the short - High Availability : By suplicating should, need and write operations concordince even if some nodes are down, increasing she one all availability of the database.

PAGE NO.: 13 [] Explain about Updale and Read Consistency with on example? Update and Read Consistency in Distributed Ratabases In distributed databases, consistency ensures that data is synchronized across multiple modes so that weers see the same data sugardless of where or when they access it 1. Update consistency ? Direction : Update Consistency scotlers to the quarantee that when data is updated in one location this update will be proposed consistency across all nodes where the duta es replicated. This helps arroid conflicts, Example ? suppose duro cisers, Alice and Bob, buy to reserve the last available moon in a hotel at the same lime without update consistency both users night see the room as available and attempt to book it simultaneously, resulting in a conflict.

PAGE NO.: KO DATE:

## 2 - Read Consistency -

Percription: Read Consistency, also known as read your - writes consistency ensures that when a user writes or updates data they can immediately see the change whom reading it.

a blog and then refresher the fage, read consistency ensures she immediately sees has comment, even if it's stored on multiple modes, without read Consistency, there might not see her own comment immediately which can be confusing.

PAGE NO.: (5

Explain about CAP Theorem

18

the CAP Theorem also known as Brewer's Theorem was formulated by computer scientist Esic Brewer and states that is impossible for a distributed data system to achieve all three of the following properties simultaneously:

1. Consistency:

every read operations returns the most recent white or on evers. Exentially, all modes in the rystem see the same data and the same time.

tor example, if Alice apostis here profile that a consistent system would ensure that oneyone viewing Alice sprofile profile immediate sees the updated photo.

2. Availability means that every request some sectives a response, even if some modes are down. In an available system the septem always responds to read and white requests.



experientle, it a mode in a lestrates segten tails, on available system usual growth the request strongs at may not graffect the most great data

o partition technique:

o partition technique:

o partition technique emplies that

the system continues to Operate even

if there is a loss of communication

between modes. A partition-tolerent system

hondles exercises where modes com't

communicate with each other by ensuring

that the system continues functioning

despite network failure

network partition between two data centers, the systems should Continue openty endependently on each side of the fortistion.