

## \* Array of Arrays (2D Array) :

```
843 const friends1 = ["Sovan", "Naveen", "Avneet"];
844 const friends2 = ["Sumit", "Nikita", "Deepak", "Kiran"];
845 const friends3 = ["JavaScript", "C++"];
846
847 const allFriends = [friends1, friends2, friends3];
848 console.log(allFriends);
849 console.log(allFriends[0]);
850 console.log(allFriends[1]);
851 console.log(allFriends[2]);
```

## → Matrix Representation →

0	0	1	2
1	"Sovan"	"Naveen"	"Avneet"
2	"Sumit"	"Nikita"	"Deepak"
3	"JavaScript"	"C++"	"Kiran"

↑ rows ← columns →

all Friends =

[ [ "Sovan", "Naveen", "Avneet" ], [ "Sumit", "Nikita", "Deepak", "Kiran" ], [ "JavaScript", "C++" ] ]

① (0)(1)      ② (1)(2)      ③ (2)(0)

Eg: allFriends [1][2] → "Deepak"  
1<sup>st</sup> Row, 2<sup>nd</sup> Col

Eg: allFriends [2][0] → "JS"  
2<sup>nd</sup> Row, 0<sup>th</sup> Col

Eg: allFriends [0][1] → "Naveen"  
0<sup>th</sup> Row, 1<sup>st</sup> Col

allFriends =

	0	1	2	
row 0	"Sovan"	"Naveen"	"Anneet"	(3)
row 1	"Avni"	"Nikita"	"Deepak"	(4)
row 2	"JS"	"C++"		(2)

← columns →

allFriends = [ A(3), A(4), A(2) ] . length = 3

allFriends[0] = 3

\* No. of Rows = 3 = No. of smaller Arrays = no. of elements in allFriends array  
= allFriends.length

\* No. of cols = No. of ele's in a smaller Array

row-0 = allFriends[0].length = 3 } allFriends[row].length

row-1 = allFriends[1].length = 4 }

row-2 = allFriends[2].length = 2 }

[ "Sovan", "Naveen", "Anneet" ].length = 3

Q: Given a  $n \times n$  matrix, print it as follows

rows      cols

IP:  $n = 3$

	0	1	2
0	10	11	12
1	15	13	14
2	16	17	18

$3 \times 3$

Op:

10 11 12

15 13 14

16 17 18

```
for(let r = 0; r < rows; r++) {
```

```
    for(let c = 0; c < cols; c++) {
```

```
        console.log(mat[r][c] + " ");
```

```
}
```

```
console.log();
```

```
}
```

①  $r = 0$

$\rightarrow c = 0 \rightarrow mat[0][0] \rightarrow 10$

$\rightarrow c = 1 \rightarrow mat[0][1] \rightarrow 11$

$\rightarrow c = 2 \rightarrow mat[0][2] \rightarrow 12$

$\rightarrow c = 3 (3 < 3) \times$

②  $r = 1$

$\rightarrow c = 0 \rightarrow mat[1][0] \rightarrow 15$

$\rightarrow c = 1 \rightarrow mat[1][1] \rightarrow 13$

$\rightarrow c = 2 \rightarrow mat[1][2] \rightarrow 14$

$\rightarrow c = 3 (3 < 3) \times$

Q: Create a  $N \times N$  matrix in the following pattern

IP:  $N = 3$

	0	1	2
0	1	2	3
1	1	2	3
2	1	2	3

$3 \times 3$

IP:  $N = 4$

	0	1	2	3
0	1	2	3	4
1	1	2	3	4
2	1	2	3	4
3	1	2	3	4

$4 \times 4$

mat = [ [1, 2, 3], [1, 2, 3], [1, 2, 3] ]

→ To build an entire Matrix,

= first smaller arrays are needed

const arr1 = [1, 2, 3];

const arr2 = [1, 2, 3];

const arr3 = [1, 2, 3];

const mat = [arr1, arr2, arr3];

} build  
smaller arrays

store the  
smaller arr  
in a bigger  
array

```
867 function createMatrix(n) {  
868   const mat = [];  
869   for (let r = 0; r < n; r++) {  
870     const smallArr = [];  
871     for (let i = 0; i < n; i++) {  
872       smallArr.push(i + 1);  
873     }  
874     mat.push(smallArr);  
875   }  
876   return mat;  
877 }
```

\* print column wise :

Eq:

	0	1	2	3
0	1	2	3	4
1	5	6	7	8
2	9	10	11	12
3	13	14	15	16

4x4

Previously,

→ goto each row

→ goto each col

Now,

→ go to each col,

→ go to each row

```
for(let c = 0; c < cols; c++)  
  for(let r = 0; r < rows; r++) {  
    PSW(`mat[r][c] + a = `);  
  }  
}
```

①  $c = 0$

$\rightarrow r = 0, mat[0](0)$

$\rightarrow r = 1, mat[1](0)$

$\rightarrow r = 2, mat[2](0)$

$\rightarrow r = 3, mat[3](0)$

$\rightarrow r = 4 \times$

②  $c = 1$

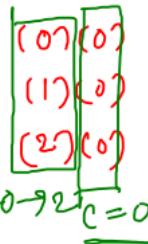
$\rightarrow r = 0 \rightarrow mat[0](1)$

$\rightarrow r = 1 \rightarrow mat[1](1)$

$\rightarrow r = 2 \rightarrow mat[2](1)$

$\rightarrow r = 3 \rightarrow mat[3](1)$

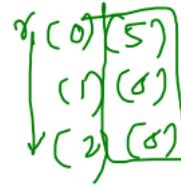
	0	1	2	3	4	5
0	1	2	3	4	5	6
1	7	8	9	10	11	12
2	13	14	15	16	17	18



$r=0 \rightarrow 2$   $c=0$



$0 \leq r < 1$   $c=1$



$r=1$   $0 \leq c < 2$

$\text{Mat}[r][c]$   
 $\rightarrow \text{Mat}[r][c]$

$\text{Mat}[j][i]$   
 $\rightarrow \text{mat}[c][r]$

$0 \rightarrow 5$     $0 \rightarrow 2$

$r = 0 \rightarrow 2$     $c = 0 \rightarrow 5$

} always only follow  $\text{mat}[r][c]$

$\boxed{\text{mat}[4][1]}$

} Never do this

## \* Alternate Matrix Traversal :

Eg:

	0	1	2	3
even	0	1	2	3
odd	1	5	6	7
even	2	9	10	11
odd	3	13	14	15

↓                          4x4

r is even    L → R  
r is odd      R → L

Odd Rows  
(L → R)

1st Row,  
3rd Row,  
5th Row etc...

Even Rows  
(R → L)

2nd Row,  
4th Row  
6th Row etc...

```
1121 function printElementsAlternately(mat, m, n) {  
1122     //Write code here and print output  
1123     for (let r = 0; r < m; r++) {  
1124         if (r % 2 == 0) {  
1125             for (let c = 0; c < n; c++) {  
1126                 process.stdout.write(mat[r][c] + " ");  
1127             }  
1128         } else {  
1129             for (let c = n - 1; c >= 0; c--) {  
1130                 process.stdout.write(mat[r][c] + " ");  
1131             }  
1132         }  
1133     }  
1134 }
```

$r=0 \rightarrow c=\emptyset \times \cancel{\emptyset} \otimes (L \rightarrow R)$

$r=1 \rightarrow c=\cancel{\emptyset} \times \emptyset \otimes (R \rightarrow L)$

\* Transpose of a Matrix :

	0	1	2	3	4
0	1	2	3	4	5
1	6	7	8	9	10
2	11	12	13	14	15
3	16	17	18	19	20
4	21	22	23	24	25

transpose

	0	1	2	3	4
0	1	6	11	16	21
1	2	7	12	17	22
2	3	8	13	18	23
3	4	9	14	19	24
4	5	10	15	20	25

→ Interchanging  
rows and cols.

Code  
should  
perform  
these operations

$$\left\{ \begin{array}{l} \text{mat}(1)(0) \leftrightarrow \text{mat}(0)(1) \\ \text{mat}(2)(0) \leftrightarrow \text{mat}(0)(2) \\ \text{mat}(2)(1) \leftrightarrow \text{mat}(1)(2) \\ \text{mat}(4)(0) \leftrightarrow \text{mat}(0)(4) \\ \text{mat}(4)(1) \leftrightarrow \text{mat}(1)(4) \end{array} \right.$$

Upper triangular matrix

	0	1	2	3	4
0	.	2	3	4	5
1	6	7	8	9	10
2	11	12	13	14	15
3	16	17	18	19	20
4	21	22	23	24	25

Lower triangular  
matrix

$$\begin{aligned} \text{mat}(3)(0) &\leftrightarrow \text{mat}(0)(3) \\ \text{mat}(3)(1) &\leftrightarrow \text{mat}(1)(3) \\ \text{mat}(3)(2) &\leftrightarrow \text{mat}(2)(3) \\ \text{mat}(4)(2) &\leftrightarrow \text{mat}(2)(4) \\ \text{mat}(4)(3) &\leftrightarrow \text{mat}(3)(4) \end{aligned}$$

	0	1	2	3	4	5
0		1	2	3	4	5
1		6	7	8	9	10
2		11	12	13	14	15
3		16	17	18	19	20
4		21	22	23	24	25
5						

```

for(let r=1; r<rows; r++) {
    for(let c=0; c < r; c++) {
        let temp = mat(r)(c);
        mat(r)(c) = mat(c)(r);
        mat(c)(r) = temp;
    }
}

```

$$\text{row} = 1 \rightarrow c = \cancel{\varnothing} \cancel{\times} \cancel{\cancel{(c < 1)}}$$

$$\text{row} = 2 \rightarrow c = \cancel{\varnothing} \cancel{\times} \cancel{\cancel{(c < 2)}}$$

$$\text{row} = 3 \rightarrow c = \cancel{\varnothing} \cancel{\times} \cancel{\cancel{(c < 3)}}$$

$$\text{row} = 4 \rightarrow c = \cancel{\varnothing} \cancel{\times} \cancel{\cancel{(c < 4)}}$$

$$(r, c) \leftrightarrow (c, r)$$

```

90 function hw1(n) {
91   for (let i = 0; i < n; i++) {
92     for (let j = 0; j < i; j++) {
93       console.log("*");
94     }
95   } → O(N)
96 } → Inner loop - ①
97
98
99 function hw2(n) {           i² ≤ N
100   let i = 1;                i = 1 → √N
101   while (i ** 2 <= n) {    → i² ≤ √N
102     i = i + 1;
103   } → O(√N)   i = 1 → √N
104 }
105
106 function hw3(m, n) {      m = 4
107   while (m != n) {
108     if (m > n) {          n = 1
109       m = m - n;        4₁₀ = 1, m = 3
110     } else {              3₁₀ = 1, m = 2
111       n = n - m;        2₁₀ = 1, m = 1
112     }
113   }                       1₁₀ = 1 ×
114 }
```

$$\begin{aligned}
 m &= 4, n = 1, O(m) \} O(\max(m, n)) \\
 m &= 1, n = 4, O(N) \} (n) O(m+n)
 \end{aligned}$$

```

116 function hw4(n) {
117   let i = 1;
118   while (i < n) {
119     let j = n;
120     while (j > 0) {
121       j = parseInt(j / 2);
122     }
123     i = i * 2;
124   }
125 }

127 function hw5(n) {
128   for (let i = 0; i < parseInt(n / 2); i++) {
129     for (let j = 1; j < n - parseInt(n / 2); j++) {
130       let m = 1;
131       while (m <= n) {
132         m = m * 2;
133       }
134     }
135   }
136 }
```

$i = 1 \xrightarrow{+1} n \ O(\log_2^i)$   
 $j = n \xrightarrow{/2} 1 \ O(\log_2^j)$   
 $i = 1 \xrightarrow{+1} n \ O((\log_2^i)^2)$

$i = 0 \xrightarrow{+1} n/2 \rightarrow O(N)$   
 $j = 1 \xrightarrow{+1} n-n/2 = n/2 \rightarrow O(N)$   
 $m = 1 \xrightarrow{*2} n \ O(\log_2^m)$

$$\begin{aligned}
 &O(n + n + \log_2 n) \\
 &\Rightarrow O(n^2 \log_2 n)
 \end{aligned}$$