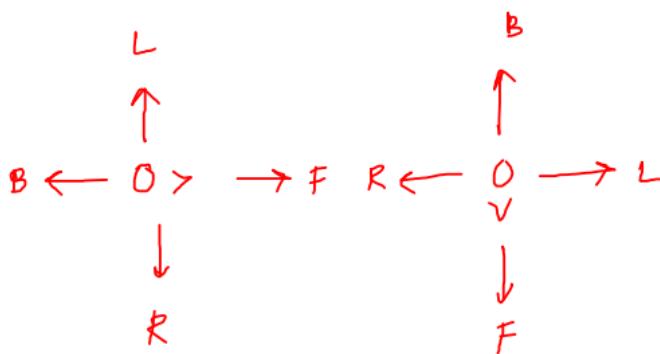


* Find the way:

	0	1	2
0	0 >	0 X 0 >	0
1	0 ←	0 X 0 ↓	1
2	0	0	0



* 0 - forward

1 - make it 0,
turn right,
move forward

Let,

$$0 > (\text{right}) = 0$$

$$0 \downarrow (\text{down}) = 1$$

$$<0 (\text{left}) = 2$$

$$\hat{0} (\text{top}) = 3$$

* facing

$$(i, j) \begin{matrix} f=0 \\ 0 > \end{matrix} (i, j+1) \quad (i, j-1) \begin{matrix} f=2 \\ <0 \end{matrix} (i, j)$$

$$\begin{matrix} (i, j) & 0 \\ \downarrow & \hat{0} \\ (i+1, j) & 0 \end{matrix} \quad \begin{matrix} \hat{0} & (i-1, j) \\ \uparrow & f=3 \\ (i, j) & 0 \end{matrix}$$

* How move forward
looks like in various
scenarios?

```

if (facing == 0) {
    j++;
}

```

```

} else if (facing == 1) {
    i++;
}

```

```

} else if (facing == 2) {
    j--;
}

```

```

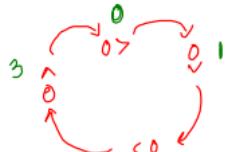
} else { // 0
    i--;
}
}

```

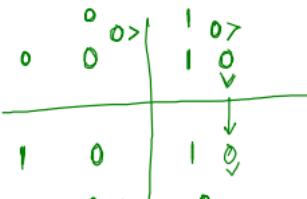
$0 > = 0$
 $0 < = 1 \swarrow$ turn right
 * How turn right looks like?

$< 0 = 2$
 $0 > = 1 \circlearrowleft$
 $f = \frac{0 X X Z}{(cyclic pattern)} \frac{0 X Z}{\dots} \frac{0 X Z}{\dots}$

$\hat{0} = 3$
 $< 0 = 2$
 $\hat{0} = 3 \curvearrowright 0 > = 0$



* facing = (facing + mat[1][1]) % 4



facing = 0

$$(0+0) \% 4 = 0$$

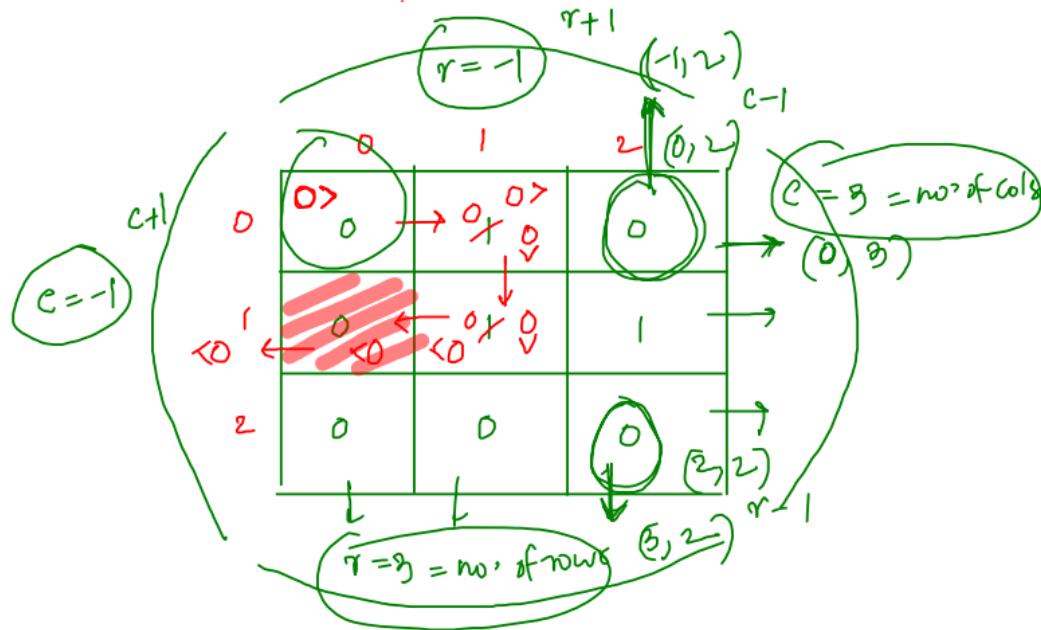
$$(\textcircled{0}+1) \% 4 = 1 \% 4 = \textcircled{1}$$

$< 0 \leftarrow < 0 \leftarrow < 0$

$$\textcircled{1} \% 4 = \textcircled{2}$$

$$\textcircled{0} + 0 \% 4 = \textcircled{2}$$

* how to decide exit?



```

1234     while (true) {
1235         facing = (facing + matrix[i][j]) % 4;
1236
1237         if (matrix[i][j] == 1) {
1238             matrix[i][j] = 0;
1239         }
1240
1241         if (facing == 0) {
1242             j++;
1243         } else if (facing == 1) {
1244             i++;
1245         } else if (facing == 2) {
1246             j--;
1247         } else {
1248             i--;
1249         }
1250
1251         if (j == cols) {
1252             // exit from right
1253             return [i, j - 1];
1254         } else if (i == rows) {
1255             // exit from bottom
1256             return [i - 1, j];
1257         } else if (j == -1) {
1258             // exit from left
1259             return [i, j + 1];
1260         } else if (i == -1) {
1261             // exit from top
1262             return [i + 1, j];
1263         }
1264     }
1265 }
```

0	1	2	
0	0 ↗	0 ↘ 0 ↓	0
1	0 ↙	0 ↗ 0 ↓	1
2	0	0	0

$$\begin{aligned}
 i &= \emptyset & 1 \\
 j &= \emptyset & \times \emptyset (-1) \\
 \text{facing} &= \emptyset \times 2
 \end{aligned}$$

9:20pm BREAK
 $\simeq 9:35\text{pm}$

$$\begin{aligned}
 \textcircled{1} \quad \text{facing} &= (0+0)\%4 \\
 &= 0
 \end{aligned}$$

$$j++ \Rightarrow j = \emptyset 1$$

$$\begin{aligned}
 \textcircled{2} \quad \text{facing} &= (1+1)\%4 \\
 &= 2
 \end{aligned}$$

$$j-- \Rightarrow j = \emptyset 0$$

$$\begin{aligned}
 \textcircled{3} \quad \text{facing} &= (0+1)\%4 \\
 &= 1
 \end{aligned}$$

$$i++ \Rightarrow i = \emptyset 1$$

$$\begin{aligned}
 \textcircled{4} \quad \text{facing} &= (2+0)\%4 \\
 &= 2
 \end{aligned}$$

$$j-- \Rightarrow j = \emptyset -1$$

$$[i, j+1] \Rightarrow [1, -1+1] \Rightarrow [1, 0]$$

* Maxima Minima :

	0	1	2
0	1	2	3
1	4	5	6
2	7	8	9

* find an element which is minimum in its Row and maximum in its Col.

$$\min R_0 = 1 \quad \max R_1 = 4 \\ \max C_0 = 7 \quad \max C_1 = 8$$

$$\min R_2 = 7 \\ \max C_0 = 1$$

```
for (let i=0; i<rows; i++) { → O(rows)
    for (let j=0; j<cols; j++) { → O(cols)
        const minR = minRow(mat, i); → Ocols
        const maxC = maxCol(mat, j); → O(rows)
        if (mat[i][j] == minR & &
            mat[i][j] == maxC) {
            return mat[i][j];
        }
    }
}
* TC : (rows * cols + (cols+rows))
: O(N * N * (N+N))
: O(2N2) : O(N3)
```

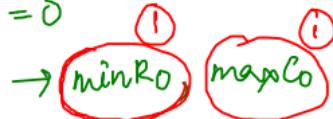
SC: O(1)

optimize / Improve :

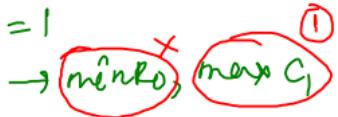
3 rows → find min in each row } store it
8 cols → find max in each col }

$i = 0$

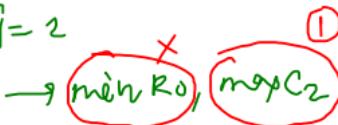
$\rightarrow j = 0$



$\rightarrow j = 1$



$\rightarrow j = 2$

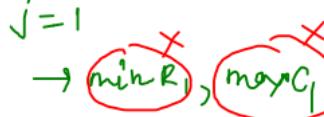


$i = 1$

$\rightarrow j = 0$



$\rightarrow j = 1$



$\rightarrow j = 2$



$i = 2$

$\rightarrow j = 0$



$\rightarrow j = 1$



$\rightarrow j = 2$



* calculate once and store it somewhere so

that you reuse it instead of calculating again & again.

matrix

	0	1	2
0	1	2	3
1	4	5	6
2	7	8	9

rows
cols
 $n \times m$
 3×3

m_1 variables

$$\left\{ \begin{array}{l} \max C_0 = -\infty \times 4 + 1 \\ \max C_1 = -\infty \times 5 + 8 \\ \max C_2 = -\infty \times 8 + 9 \\ \vdots \\ \max C_{m-1} \end{array} \right.$$

* preComputation
(calculating once, using it again & again)

	$\max R_0$	$\max C_1$	$\max C_2$	$\max C_{m-1}$	
0	-10	-10	-10	... 0 ...	-10

X X S
A S S
1 8 9

n variables

$$\left\{ \begin{array}{l} \min R_0 = \infty \mid 1 \\ \min R_1 = \infty \mid 4 \\ \min R_2 = \infty \mid 7 \\ \vdots \\ \min R_{n-1} \end{array} \right.$$

	$\min R_0$	$\min R_1$	$\min R_2$		$\min R_{n-1}$
0	∞	∞	∞	... 0 ...	∞

1 4 7

* TC : precomputation + logic

* SC : $(\text{rows} + \text{cols}) \uparrow$
: $O(N)$

$$: O(\text{rows} + \text{cols} + (\text{rows} \times \text{cols})) + O(\text{rows} \times \text{cols})$$

$$: O(\text{rows} \times \text{cols}) : O(N^2) \downarrow$$

Imp methods / techniques / pattern ?

flag variable \rightarrow store in flag \rightarrow break \rightarrow out of loop $\xrightarrow{\text{yes}}$ NO

digit iteration \rightarrow $\%10$ (last digit), $/10$ (remove last digit)

two pointer \rightarrow reverse an array

rotation (cyclic) \rightarrow use $\%$

generating pair, triplets, subarray \rightarrow template (code structure)

running stream \rightarrow max/min/sum/prod at every point/idx of array

precomputation \rightarrow store in your mind use later

}

① what did you like the most from lectures?

② what did you hate the most overall?

③ one / two words about the module?