

* Find the Way:

IP:

	0	1	2
0	0 0>	1 0>	0
1	<0	0 >	1
2	0	0	0

(left) <0 0> (right)
 turn right
 turn right 0> (down)

* 0 → forward

1 → make it 0,
 turn right,
 more forward.

0> (right) = 0

0 (down) = 1

<0 (left) = 2

^ (top) = 3

OP: (1, 0)

The last cell before
 coming out of matrix

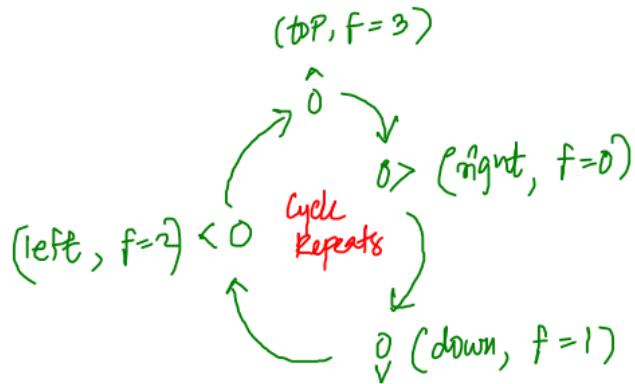
* maintain facing variable
 (0, 1, 2, 3)

(i, j) f=0 (i, j+1) (i, j-1) f=2 (i, j) f=3
 0> → 0> <0 ← <0

0 (i, j)
 ↓ f=1
 0 (i+1, j)

^ 0 (i-1, j)
 ↑ f=3
 ^ 0 (i, j)

How to turn right :



$$f = \emptyset X \neq \emptyset \emptyset X \neq \emptyset \emptyset X 2 \dots$$

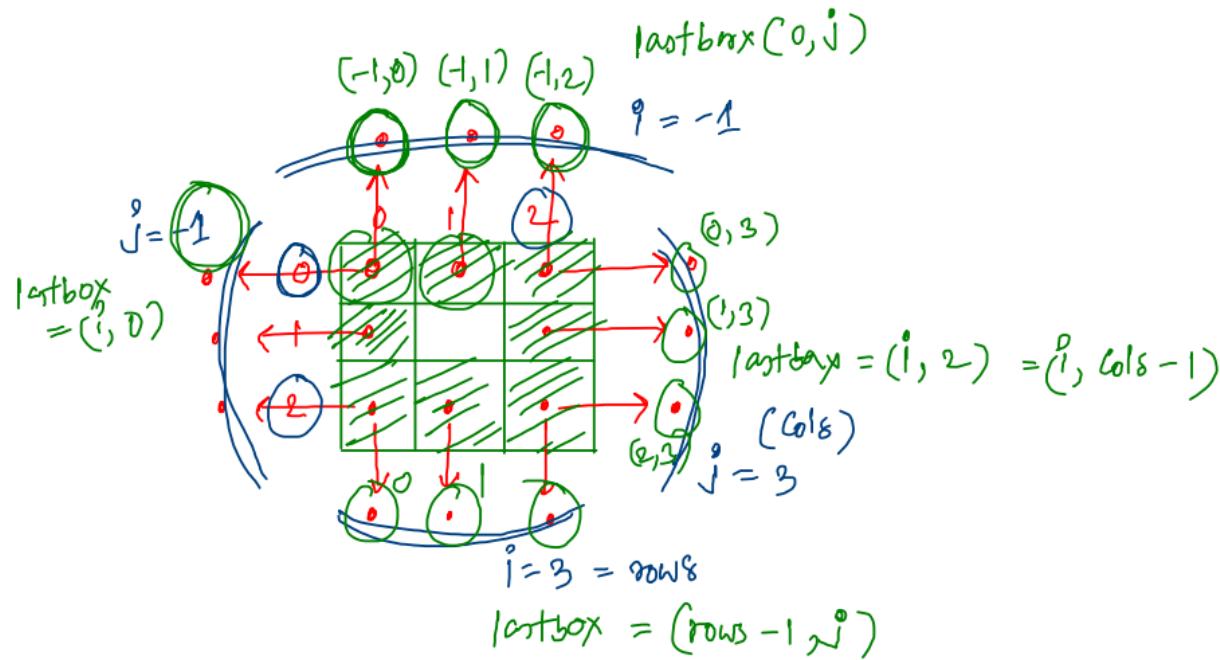
when, $\text{mat}[i][j] = -1$,

a. $\text{mat}[i][j] = 0$

b. turn right, $f = f + 1$, $f = \emptyset X \neq \emptyset \emptyset \overset{\times}{\textcircled{4}} \overset{\times}{\textcircled{5}} \overset{\times}{\textcircled{6}}$

★ $f = (f + 1) \% 4$ $f = \emptyset X \neq \emptyset 0 \quad \overset{\times}{\textcircled{4}} \downarrow \overset{\times}{\textcircled{5}} \downarrow \overset{\times}{\textcircled{6}} \downarrow$

How to check you are out:



4 boundaries

* Maxima / minima :

IP:

0	1	2	3
0	1	2	3
1	4	5	6
2	7	8	9

* ele which minimum
in its Row and
maximum in its col

OP: 7

$$\min R_0 = 1$$

$$\max C_0 = 7$$

$$1 == \min R_0 \text{ if}$$

$$1 == \max C_0$$

$$\boxed{\min R_0 = 1}$$

$$\max C_1 = 8$$

$$8 == \min R_0 \text{ if}$$

$$8 == \max C_1$$

$$\min R_2 = 7$$

$$\max C_2 = 7$$

$$7 == \min R_2 \text{ if}$$

$$7 == \max C_2 \\ (\text{yes})$$

* TC : $O(\text{rows} * \text{cols} * (\text{rows} + \text{cols}))$
 $= O(N^3)$, SC : $O(1)$

for (let i=0; i < rows; i++) {

 for (let j=0; j < cols; j++) {

 const minR = findMinRow
(mat, i);

 const maxC = findMaxCol
(mat, j);

 if (mat[i][j] == minR &&
 mat[i][j] == maxC)

 return mat[i][j];

}

 return -1;

optimise :

3×3 matrix ,

① $i = 0$
 $\rightarrow j=0 \rightarrow \min R_0, \max C_0$
 $\rightarrow j=1 \rightarrow \cancel{\min R_0}, \max C_1$
 $\rightarrow j=2 \rightarrow \cancel{\min R_0}, \max C_2$

② $i = 1$
 $\rightarrow j=0 \rightarrow \min R_1, \max C_0 \times$
 $\rightarrow j=1 \rightarrow \cancel{\min R_1}, \max C_1 \times$
 $\rightarrow j=2 \rightarrow \cancel{\min R_1}, \max C_2 \times$

③ $i = 2$
 $\rightarrow j=0 \rightarrow \min R_2, \max C_0 \times$
 $\rightarrow j=1 \rightarrow \cancel{\min R_2}, \max C_1 \times$
 $\rightarrow j=2 \rightarrow \cancel{\min R_2}, \max C_2 \times$

* we can use the calculated things,
need not calculate multiple times .

① Compute min in each row, and store

$$\min R_0 = 1 \quad \text{If } n \text{ Rows}$$

$$\min R_1 = 4 \rightarrow n \text{ size Arr}$$

$$\min R_2 = 7$$

② Compute max in each col, and store

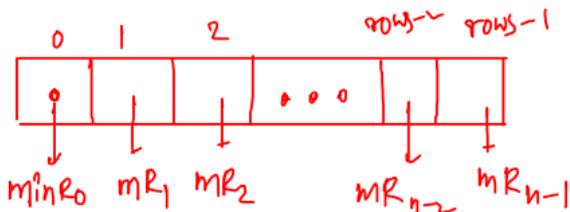
$$\max C_0 = 1 \quad \text{If } m \text{ Col}$$

$$\max C_1 = 8 \rightarrow m \text{ size Arr}$$

$$\max C_2 = 9$$

③ Run the pt discussed approach

How to build table ?



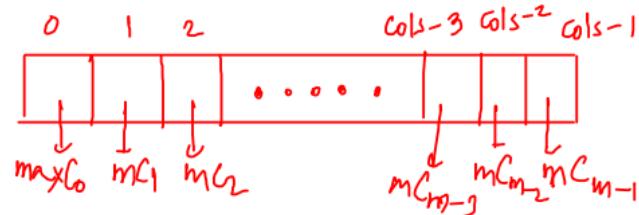
	0	1	2
0	1	(2)	(3)
1	4	5	6
2	7	8	9

PRE COMPUTATION

TECHNIQUE

Calculate once use many times

- ① $i=0$
 $j=0 \rightarrow \min R_0, \max C_0 \rightarrow 7 == 1$
 $j=1 \rightarrow \min R_0, \max C_1 \rightarrow 8 == 2$
 $j=2 \rightarrow \min R_0, \max C_2 \rightarrow 9 == 3$



Time/space:

```

1440 function maximaMinima(mat) {
1441     //Write code here
1442     const rows = mat.length;
1443     const cols = mat[0].length;
1444     const minRow = [];
1445     const maxCol = [];
1446     const extraArr = [0, 1, 2, 3, 4, 5, 6, 7];
1447
1448     for (let r = 0; r < rows; r++) {
1449         const minEle = findMinRow(mat, r);
1450         minRow.push(minEle);
1451     }
1452
1453     for (let c = 0; c < cols; c++) {
1454         const maxEle = findMaxCol(mat, c);
1455         maxCol.push(maxEle);
1456     }
1457
1458     for (let r = 0; r < rows; r++) {
1459         for (let c = 0; c < cols; c++) {
1460             const minR = minRow[r];
1461             const maxC = maxCol[c];
1462             if (minR == mat[r][c] && maxC == mat[r][c]) {
1463                 return mat[r][c];
1464             }
1465         }
1466     }
1467
1468     return -1;
1469 }

```

$O(r \times c)$

$\Rightarrow O(r + c)$

$\Rightarrow O(r \times c)$

$\Rightarrow O(1)$

$\Rightarrow O(r \times c)$

$\Rightarrow O(3 \times r \times c) \Rightarrow O(r \times c) \Rightarrow O(N^2)$

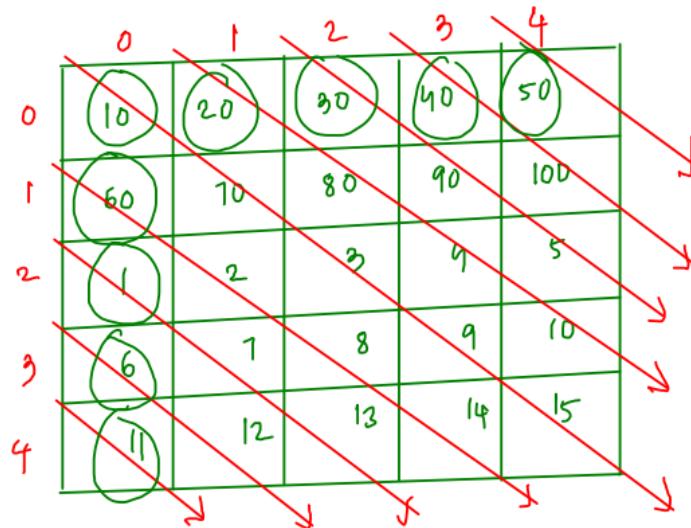
$T.C$

$O(3 \times r \times c) \Rightarrow O(r \times c) \Rightarrow O(N^2)$

$\frac{SC}{r \times c}$

$= O(r + c) = O(2N) \Rightarrow O(N)$

* Diagonal traversal :



① starting point

② follow the diagonal $(+1, +1)$

op: 50, 40, 100, 30, 90, 5, 20, 80, 4, 10,
10, 70, 3, 9, 15, 60, 2, 8, 14, 1, 1, 13,
6, 12, 11

