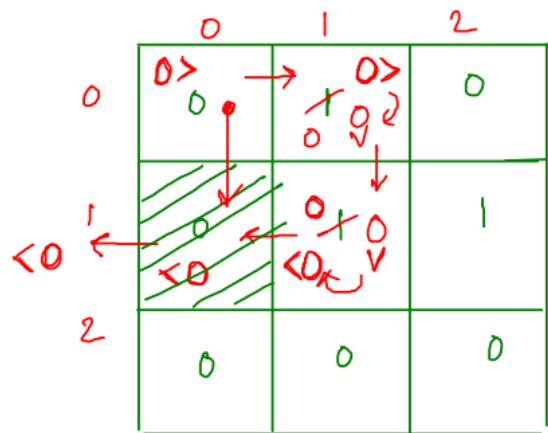


\* Find the way : [Get out of matrix]



op:  $(1, 0)$  [Last cell before coming out of matrix]

\*  $0 \rightarrow$  forward

$1 \rightarrow$  make it 0,  
tron right  
move forward

\* In order to move forward,  
we need to know our facing first, ( maintain facing variable )  
 $(0, 1, 2, 3)$

$(i, j)$      $f=0$      $(i, j+1)$   
 $0 >$      $\longrightarrow$      $0 >$

$< 0$      $f=2$      $< 0$   
 $(i, j-1)$         $(i, j)$

$0$      $(i, j)$

$v$

$f=1$

$0$      $(i+1, j)$

$v$

$0$      $(i-1, j)$

$\uparrow$   
 $f=3$

$0$      $(i, j)$

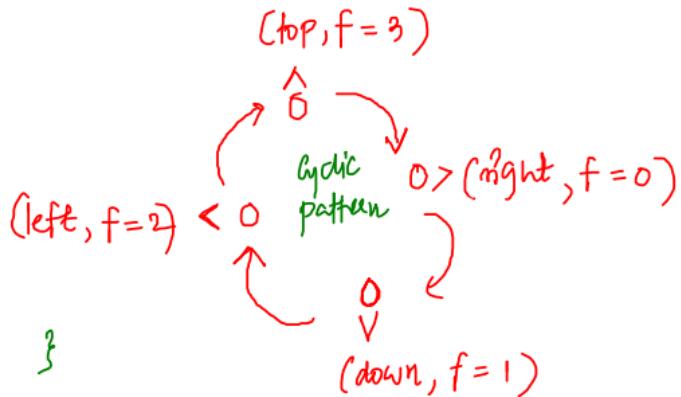
\* How to move forward based on facing .

\* How to know the facing?

Initially, you are at (0,0) facing right

```
let i=0;  
let j=0;  
let facing = 0;  
  
while (true) {
```

\* how to change facing



\* 1<sup>st</sup> know your facing  
then only you can decide  
on how to forward.

$\text{mat}[i][j] = 0 \rightarrow$  no change in facing

$\Rightarrow$  In this facing variable will directly give the direction.

$\text{mat}[i][j] = 1 \rightarrow$  facing will change because you need to turn right.

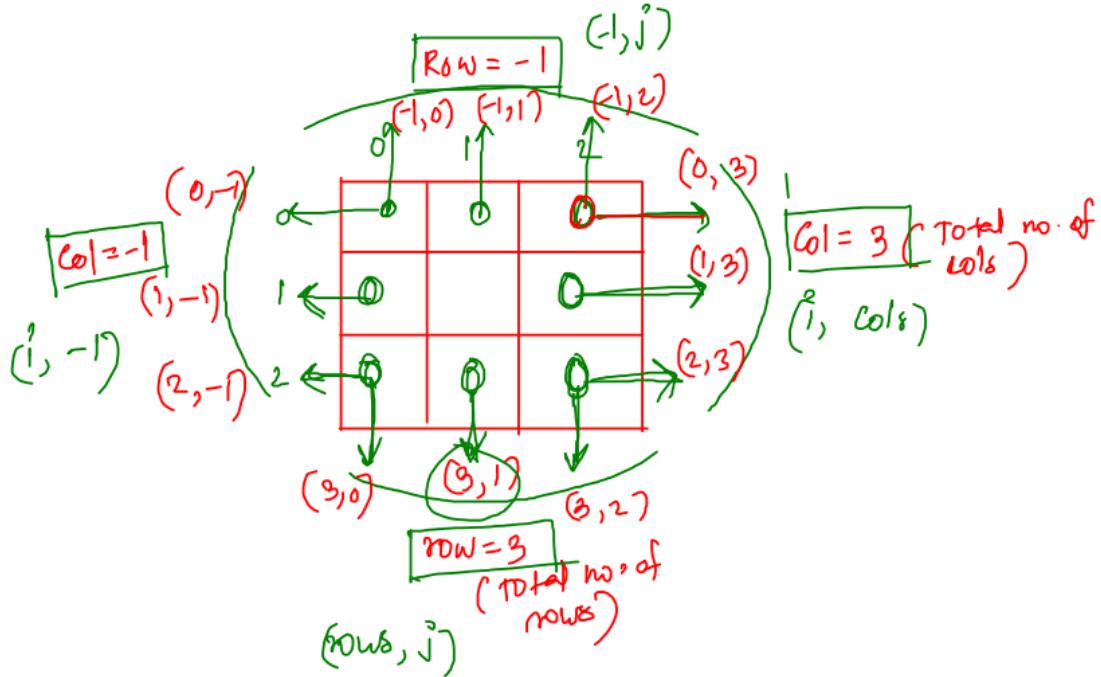
$f = \emptyset X \neq 3 \emptyset X 3 \emptyset X \neq 3 \dots$

$f = \emptyset X \neq 3 4X$  ( $f = f + 1$   
will not work)

\* Cyclic pattern,  
e.g. helpful

$f = (f + 1) \% 4;$

\* How to check you are out of Matrix



①  $(-1, j)$

Last box  $\Rightarrow (0, j)$

②  $(i, \text{cols})$

Last box  $\Rightarrow (i, \text{cols} - 1)$

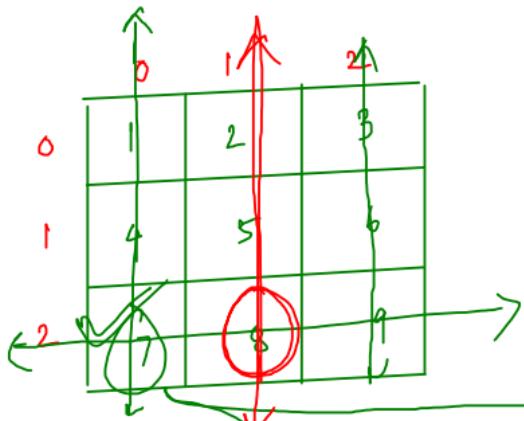
③  $(\text{rows}, j)$

Last box  $\Rightarrow (\text{rows} - 1, j)$

④  $(i, -1)$

Last box  $\Rightarrow (i, 0)$

## \* Maxima / Minima :



$$\min R_0 = 1$$

$$\max C_0 = 7$$

$$\min R_1 = 4$$

$$\max C_1 = 8$$

$$\max C_2 = 9$$

$$\min R_2 = 7$$

\* Find an element which is minimum in its Row and Maximum in its Column

$(i, j)$   $\rightarrow$  belongs to Row  $i$   
 $\rightarrow$  belongs to Col  $j$

1. goto every element  $(i, j)$
2. check  $\text{mat}[i][j] == \min \text{in Row}(i)$   
 if  $\text{mat}[i][j] == \max \text{Col}(j)$

$$\text{mat}(i)(j) = \min \text{Row}(i)$$

if  $\text{mat}(i)(j) == \max \text{Col}(j)$

```

for (let i = 0; i < rows; i++) {
    for (let j = 0; j < cols; j++) {
        const minR = findMinRow(mat, i);
        const maxC = findMaxCol(mat, j);
        if (mat[i][j] == minR || mat[i][j] == maxC) {
            return mat[i][j];
        }
    }
}
return -1;

```

$O(\text{rows} * \text{cols})$

```

function findMinRow(mat, row) {
    /* Code */
}

function findMaxCol(mat, col) {
    /* Code */
}

```

$O(\text{rows})$

$O(\text{cols})$

$$\begin{aligned}
 \text{TC : } O(\text{rows} * \text{cols} + (\text{row} + \text{cols})) &\Rightarrow O(N * N + 2N) \\
 &\Rightarrow O(N^2)
 \end{aligned}$$

## # optimize / Improve :

$3 \times 3$  matrix ,

①  $i=0$   
 $\rightarrow j=0 \rightarrow \min R_0, \max C_0$   
 $\rightarrow j=1 \rightarrow \cancel{\min R_0}, \max C_1$   
 $\rightarrow j=2 \rightarrow \cancel{\min R_0}, \max C_2$

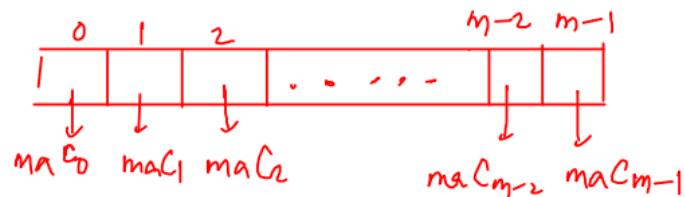
②  $i=1$   
 $\rightarrow j=0 \rightarrow \min R_1, \max C_0$   
 $\rightarrow j=1 \rightarrow \cancel{\min R_1}, \max C_1$   
 $\rightarrow j=2 \rightarrow \cancel{\min R_1}, \max C_2$

③  $i=2$   
 $\rightarrow j=0 \rightarrow \min R_2, \max C_0$   
 $\rightarrow j=1 \rightarrow \cancel{\min R_2}, \max C_1$   
 $\rightarrow j=2 \rightarrow \cancel{\min R_2}, \max C_2$

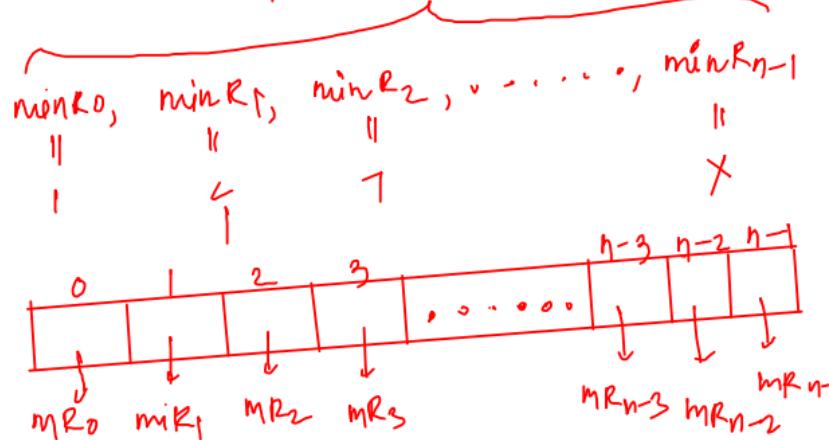
\* we are calculating always which is redundant / duplicate , calculate once and store it somewhere .

	0	1	2
0	1	2	3
1	4	5	6
2	1	8	9

② Compute all the maximas  
in every row ( $m \times n$ )



① Compute all the minimas  
in every row,  $n$  rows



findMinRow(mat, z)



$minR[z]$

```

1205 function maximaMinima(mat) {
1206   const rows = mat.length;
1207   const cols = mat[0].length;
1208
1209   // 1. Precompute, calculate minR, maxC beforehand
1210   const minR = [];
1211   const maxC = [];
1212   for (let r = 0; r < rows; r++) {  $O(\text{rows} + \text{cols})$ 
1213     const minEle = findMinRow(mat, r);
1214     minR.push(minEle);
1215   }
1216
1217   for (let c = 0; c < cols; c++) {  $O(\text{cols} + \text{rows})$ 
1218     const maxEle = findMaxCol(mat, c);
1219     maxC.push(maxEle);
1220   }
1221
1222   for (let i = 0; i < rows; i++) {
1223     for (let j = 0; j < cols; j++) {
1224       if (mat[i][j] == minR[i] && mat[i][j] == maxC[j]) {
1225         return mat[i][j];
1226       }
1227     }
1228   }
1229   return -1;
1230 }

```

findMinRow  $\rightarrow O(\text{cols})$   
 findMaxCol  $\rightarrow O(\text{rows})$

\* Precomputation,  
 (when you repetitive or redundant  
 calculation, use this technique)

$O(3 * \text{rows} * \text{cols})$

$\Rightarrow O(3N^2)$

$\Rightarrow O(N^2)$  : TC

SC:  $O(\text{rows} + \text{cols}) \Rightarrow O(2N) \Rightarrow O(N)$ .

$\downarrow$   
 $\downarrow$   
 minR[] maxC[]

## \* Special Matrix:

	0	1	2
0	0	0	2
1	0	2	0
2	3	0	1

① If the ele belongs  
to diag / anti-diag  $r \neq c$

② all others ele = 0

\* given(r, c)

how do you know if it is  
a diag / anti diag ele?

$\Rightarrow$  In diag,  $(r == c)$   
 $(0,0), (1,1), (2,2)$

In anti-diag,  $(r + c == n - 1)$   
 $c = n - r - 1$

```
for (let r = 0; r < rows; r++) {
```

```
    for (let c = 0; c < cols; c++) {
```

```
        if (r == c || c == n - r - 1) {
```

```
            if (mat[r][c] != 0) {
```

```
                return false;
```

```
}
```

```
else {
```

```
    if (mat[r][c] != 0) {
```

```
        return false;
```

```
    }
```

```
return true;
```

\* Upper and Lower triangular Sum :

0	1	2	3
0	1	2	3
1	5	6	7
2	9	10	11
3	13	14	15
			16

upper

lower

lower :  $(r, c) \Rightarrow (r \geq c) \checkmark$

$(0, 0)$

$(1, 0) \quad (1, 1)$

$(2, 0) \quad (2, 1) \quad (2, 2)$

$(3, 0) \quad (3, 1) \quad (3, 2) \quad (3, 3)$

upper :  $(r, c) \Rightarrow (r \leq c) \checkmark$

$(0, 0) \quad (0, 1) \quad (0, 2) \quad (0, 3)$

$(1, 1) \quad (1, 2) \quad (1, 3)$

$(2, 2) \quad (2, 3)$

$(3, 3)$

## \* Spiral Matrix 2:

	0	1	2	3	4
0	10	20	30	40	50
1	60	70	80	90	100
2	110	120	130	140	150
3	160	170	180	190	200
4	210	220	230	240	250

minR → ↑ minC      ↑ maxC → maxR

① leftwall      res[ ]

```
for (let r = minR; r <= maxR; r++) {
  psw ( mat[r][minC] );
}
minC++;
↓ store in arr
res.push( mat[r][minC] )
```

② bottom wall

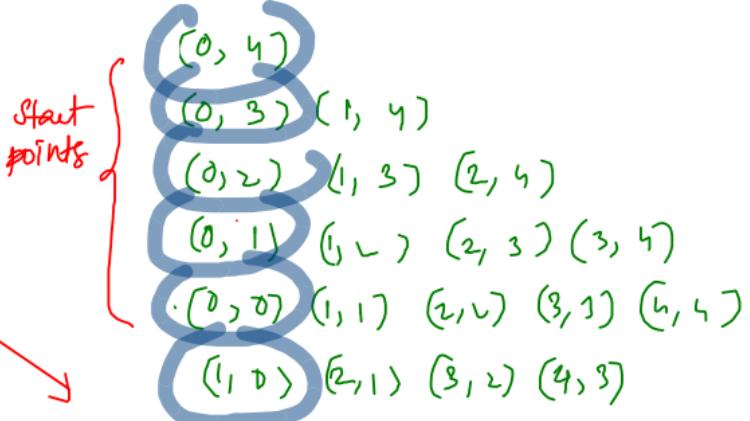
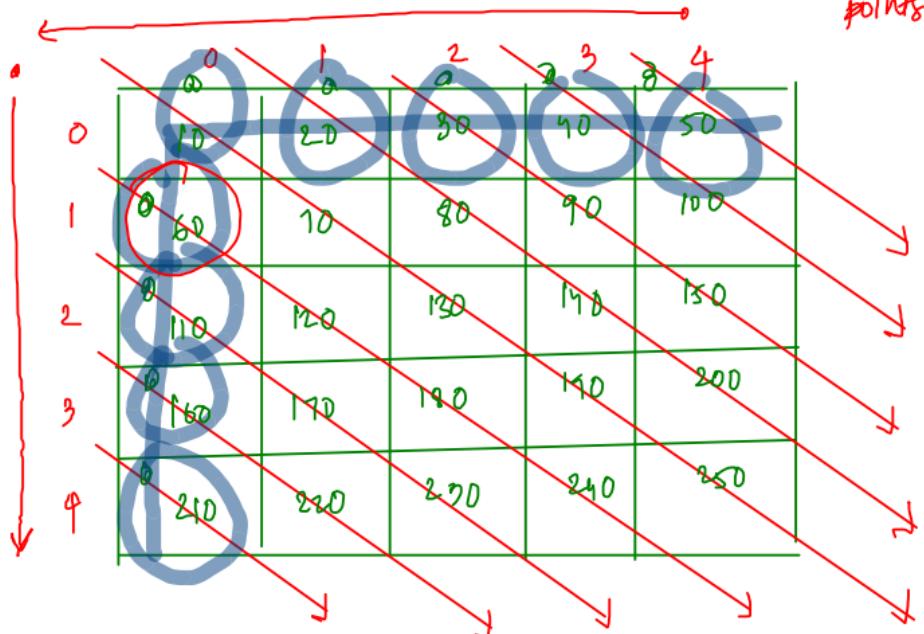
```
for (let c = minC; c <= maxC; c++) {
  psw ( mat[minR][c] );
}
```

maxR--;

③ rightwall

④ top wall

\* Diagonal Traversal Different!



(0, 1) (+1,+1)

→ (1, 2) (+1,+1)

• style  
diagonal traversal  
→ (2, 3) (+1,+1)  
→ (3, 4)

(1, 0) → (2, 1) → (3, 2) → (4, 3)  
(+1,+1) → (+1,+1) → (+1,+1) → (+1,+1)

\* find the starting points and print every diagonal with starting point