

* Adding two arrays:

0	1	2	3	4
2	1	3	5	4
0	1	2	3	4
3	5	4	2	1

for ($i = n-1 \rightarrow 0$) {

$$\text{res}[i] = \text{arr1}[i] + \text{arr2}[i]$$

}

$$\text{res}[n-1] = \text{arr1}[n-1] + \text{arr2}[n-1]$$

$$= \text{arr1}[4] + \text{arr2}[4]$$

$$= 4 + 1$$

$$\text{arr1}[4] = 5$$

If there carry left after the process,

if ($\text{carry} == 1$) {

add ① to the start of the

0	1	2	3
1	3	2	6
0	1	2	3

0	1	2	3	4
1	3	2	6	
0	1	2	3	4

e.g:

0	1	2	3
4	3	5	2
0	1	2	3
6	9	7	4
10	12	12	6

$\rightarrow \text{arr1}$ ($n=4$)
 $\rightarrow \text{arr2}$

* We do not iterate on arrays, we are actually iterating on some indices.

X according
1st code

let carry = 0;

for (let $i = n-1$; $i >= 0$; $i--$) {

$$\text{const sum} = \text{arr1}[i] + \text{arr2}[i] + \text{carry};$$

$$\text{res}[i] = \text{sum} \% 10;$$

}

$$\text{carry} = \text{sum} / 10;$$

$$\textcircled{1} \quad \text{sum} = 2 + 4 + 0 \\ = 6$$

$$6 \% 10 = \textcircled{6} \rightarrow \text{res} \\ 6 / 10 = \textcircled{0} \rightarrow \text{carry}$$

0	1	2	3
4	3	5	2
6	9	7	4

0	1	2	3
1	3	2	6
0	1	2	3

res :

0	1	2	3
1	3	2	6
0	1	2	3

$$12 \% 10 = \textcircled{2} \rightarrow \text{res} \\ 12 / 10 = \textcircled{1} \rightarrow \text{C}$$

$$\textcircled{5} \quad \text{sum} = 3 + 9 + 1$$

$$= 13$$

$$13 \% 10 = \textcircled{3} \rightarrow \text{res}$$

$$13 / 10 = \textcircled{1} \rightarrow \text{carry}$$

$$\textcircled{4} \quad \text{sum} = 4 + 6 + 1$$

$$= 11$$

$$11 \% 10 = \textcircled{1} \rightarrow \text{res}$$

$$11 / 10 = \textcircled{1} \rightarrow \text{carry}$$

Eg:

0	1	2	9
9	9	5	2

→ arr1 ($n = 4$)

0	1
9	6

→ arr2 ($m = 2$)

$$i = n - 1 \quad \text{carry} = 0$$

$$j = m - 1$$

while ($i >= 0$) {

 let sum;

 if ($j >= 0$) sum = arr1[i] + arr2[j] + carry;

 else sum = arr1[i] + carry;

$$\text{res}[i] = \text{sum} \% 10;$$

$$\text{carry} = \text{sum}/10;$$

$$i--;$$

$$j--;$$

}

$$\begin{array}{r}
 9 \ 9 \ (\textcircled{5}) \ (\textcircled{2}) \\
 + (\textcircled{9}) \ (\textcircled{6}) \\
 \hline
 \end{array}$$

0	1	2	3
9	9	5	2

→ res ($i = 4$)

0	1
9	6

→ arr2 ($m = 2$)

-3	-2	-1
9	6	

→ arr1 ($n = 4$)

0	1	2	3
9	9	5	2

→ res ($i = 4$)

$$\begin{aligned}
 \textcircled{1} \quad \text{Sum} &= 9 + 6 + 0 \\
 &= 15
 \end{aligned}$$

$$15 \% 10 = \textcircled{5} \rightarrow \text{res}$$

$$15 / 10 = \textcircled{1} \rightarrow \text{carry}$$

$$\begin{aligned}
 \textcircled{2} \quad \text{Sum} &= 5 + 9 + 0 \\
 &= 14
 \end{aligned}$$

$$14 \% 10 = \textcircled{4} \rightarrow \text{res}$$

$$14 / 10 = \textcircled{1} \rightarrow \text{carry}$$

$$\begin{aligned}
 \textcircled{3} \quad \text{Sum} &= 9 + 1 \\
 &= 10
 \end{aligned}$$

$$10 \% 10 = \textcircled{0} \rightarrow \text{res}$$

$$10 / 10 = \textcircled{1} \rightarrow \text{carry}$$

$$\begin{array}{r}
 10 \\
 10 \\
 \hline
 0
 \end{array}$$

```

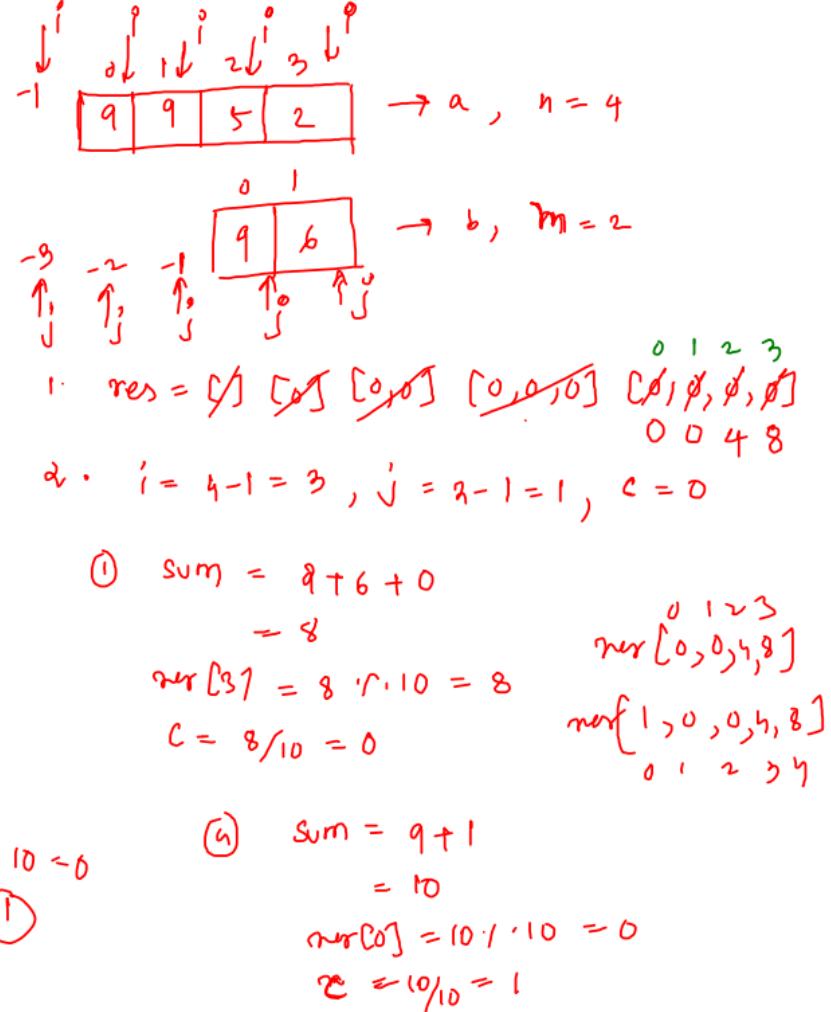
9  // 1. initialize / create empty result array
10 const res = [];
11 for (let i = 0; i < n; i++) {
12   res.push(0);
13 }
14
15 // 2. start adding from last
16 let i = n - 1;
17 let j = m - 1;
18 let carry = 0;
19
20 while (i >= 0) {
21   let sum = j >= 0 ? a[i] + b[j] + carry : a[i] + carry;
22   res[i] = sum % 10;
23   carry = sum / 10;
24
25   i--;
26   j--;
27 }
28
29 // 3. If there is carry left, add 1 to the beginning of array
30 res.unshift(1); if (carry == 1)

```

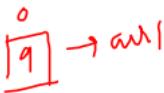
$$\textcircled{2} \quad \text{Sum} = 5 + 9 + 0 \\ = 14$$

$$\text{arr}[2] = 14 \% 10 = 4 \\ c = 14 / 10 = \textcircled{1}$$

$$\textcircled{3} \quad \text{Sum} = 9 + 1 \\ = 10 \\ \text{arr}[1] = 10 \% 10 = 0 \\ c = 10 / 10 = \textcircled{1}$$



* Carry can never be > 1 (it is always $0 \text{ or } 1$)



→ max possible sum of 2 digits
18 → carry = 1

only $n >= m$
calSumUtil(a, b, n, m) {

/* code
} {

user see
this, they
can provide
anything

helped
function

```
calSum(a, b, n, m) {
    if (n >= m) {
        return calSumUtil(a, b, n, m);
    } else {
        return calSumUtil(b, a, m, n);
    }
}
```

* Our previous code assumes $\text{ans1} >= \text{ans2}$
 $(n >= m)$



$$\text{ans} = [0, 1, 8]$$

$$= [1, 4, 8]$$

$$a = [9, 6] \quad n = 2$$
$$b = [9, 9, 5, 2] \quad m = 4$$

9:20 pm - 9:35 pm

BREAK

arr = [^{0 1 2 3}
 1, 1, 2, 3]

arr.unshift(20) → adds an ele to beginning
of array

[^{0 1 2 3 4}
 20, 1, 1, 2, 3]

arr.unshift(4)

[^{0 1 2 3 4 5}
 4, 20, 1, 1, 2, 3]

* Buildings:

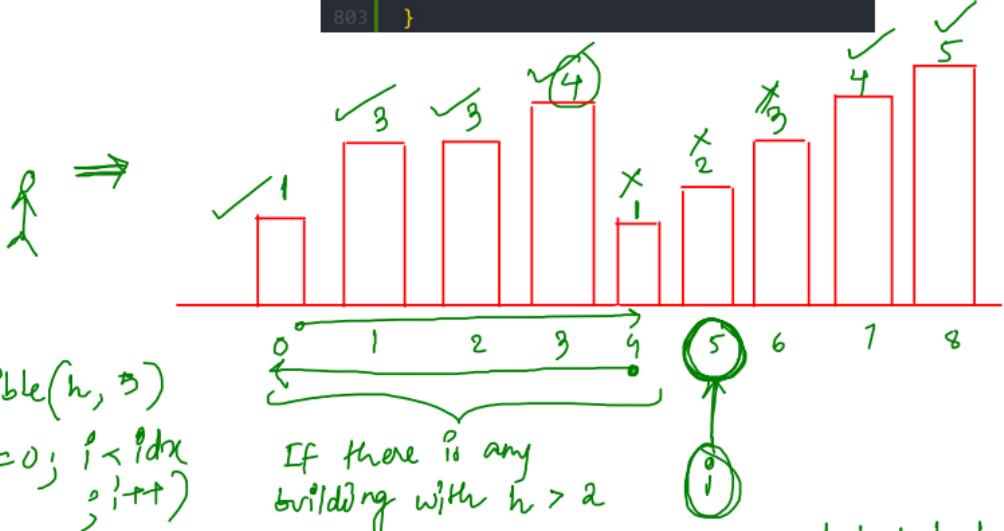
```

796
797
798
799
800
801
802
803
function isVisible(heights, idx) {
  for(let i = idx - 1; i >= 0; i--) {
    if(heights[i] > heights[idx]) {
      return false;
    }
  }
  return true;
}

```

↑
↑: [1, 3, 3, 4, 1, 2, 3, 4, 5]

Op: 6



* go to every element,
if (see) $cnt++$
↓
how to decide?

(how can you say
the building is
hidden or not)

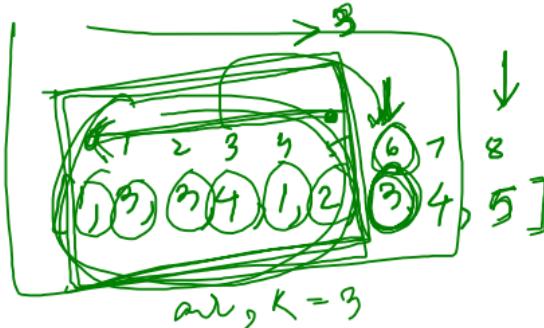
↓ ↓ ↓ ↓ ↓ ↓ ↓
0 1 2 3 4 5 6 7 8
[1, 3, 3, 4, 1, 2, 3, 4, 5]
✓ ✓ ✓ ✓ X X ✓

$$cnt = \emptyset \times 2 \oplus 4 \times 6$$

```

31 v function isVisible(heights, idx) {
32 v   for (let i = idx - 1; i >= 0; i--) {
33 v     if (heights[i] > heights[idx]) {
34       return false;
35     }
36   }
37   return true;
38 }
39
40 v function countVisibleRoofs(heights) {
41   // Write your code here
42   const n = heights.length;
43   let cnt = 0;
44 v   for (let i = 0; i < n; i++) {
45 v     if (isVisible(heights, i)) {
46       cnt++;
47     }
48   }
49
50   return cnt;
51 }

```



$i=1$
 $4 > 5$ (2, 3, 7, 6)
 $i=6$ $3 > 5$ (1, 2, 7, 8)
 $i=5$ $2 > 4$ (1, 2, 5)
 $i=4$ $1 > 5$
 $i=3$ $4 > 5$

\Rightarrow Given an array check whether
 is there any element greater than
 some k

$(heights, \textcircled{idx}) \rightarrow (h, b)$

① $i=3$

$h[3] > h[6]$

$(4) > (3)$
blocked

I can't see
6th \textcircled{idx}
building

① $i=5$

$h[5] > h[6] (2 > 3)$

② $i=4$

$h[4] > h[6] (1 > 3)$

Improve :

```
for(going to every ele) {
    isVisible() → ( iterating on
                    left part)
```

}

i = 0, isVisible() → 0

i = 1, isVisible() → 1

i = 2, isVisible() → 2

i = 3, 3

i = 4, 4

:

:

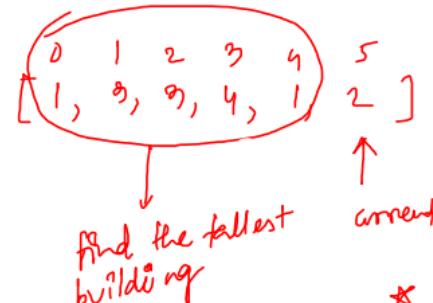
i = n-1, n-1

$$0 + 1 + 2 + 3 + 4 + \dots + n-1$$

$$\Rightarrow \frac{n(n-1)}{2} \Rightarrow \frac{n^2 - n}{2}$$

$\approx n^2$ instruction
operations

running maximum :



If tallest > Current
+ you cannot see

- * If tallest building itself cannot block me
- then no other building on left will block me.



- * Using RM, there is no need of calculating max again & gain

Q : [1, 4, 2, 9, 6, 5]

$$\textcircled{1} \quad -\infty > 1 \rightarrow \text{no } \checkmark$$

$$\textcircled{2} \quad 1 > 4 \rightarrow \text{no } \checkmark$$

$$\textcircled{3} \quad 4 > 2 \rightarrow \text{yes } \times$$

$$\textcircled{4} \quad 4 > 9 \rightarrow \text{yes } \times$$

$$\textcircled{5} \quad 9 > 6 \rightarrow \text{NO } \checkmark$$

$$\textcircled{6} \quad 6 > 5 \rightarrow \text{yes } \times$$

runningMax = $-\infty$ \rightarrow This will tell you the maxafe on left part

$1 > -\infty$	\times
$4 > 1$	\times
$6 > 4$	\checkmark

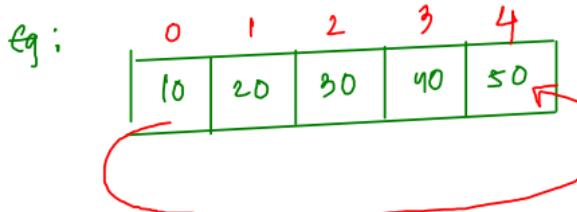
cut = $\emptyset \times \checkmark \times \times$

* running stream
Technique
(max/min/
sum/prod)

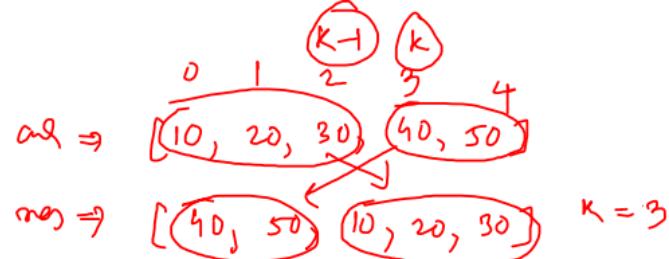
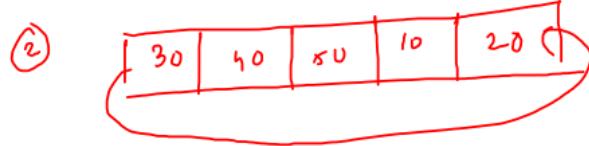
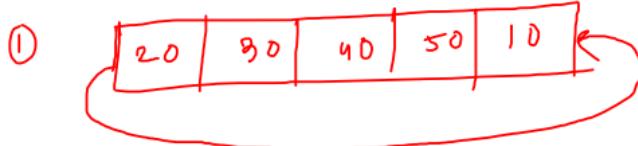
```

819 function countVisibleRoofs(heights) {
820     // Write your code here
821     const n = heights.length;
822     let cnt = 0;
823     let runningMax = -Infinity;
824     for (let i = 0; i < n; i++) {
825         if (runningMax <= heights[i]) {
826             cnt++;
827             runningMax = heights[i];
828         }
829     }
830     return cnt;
831 }
832 }
```

* Rotate Array :



$k_r = 3$



① ($k \rightarrow n-1$) (push to res) ($3 \rightarrow 4$)

② ($0 \rightarrow k-1$) (push to res) ($0 \rightarrow 2$)

($res = []$), $\star k = k_r \cdot n$

for(let $i = k$; $i < n$; $i++$) {

 res.push(arr[i]);
}

for(let $i = 0$; $i < k$; $i++$) {

 res.push(arr[i]);
}

$res = []$

= [10]

= [40, 50]

= [40, 50, 10]

= [40, 50, 10, 20]

= [40, 50, 10, 20, 30]

* what if $K \geq n$?

$$K = 100;$$

$$n=5 [10, 20, 30, 40, 50]$$

for ($K \rightarrow n-1$)

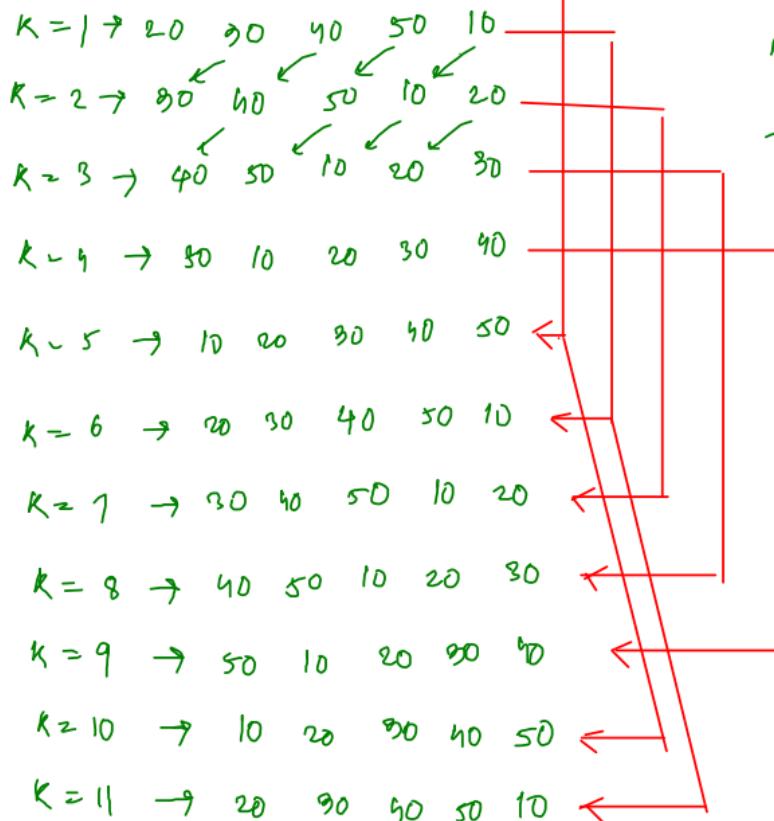
$$\Rightarrow (100 \rightarrow 4) \times$$

for ($0 \rightarrow K-1$)

$$\Rightarrow (0 \rightarrow 99) \checkmark$$

0	1	2	3	4
10	20	30	40	50

$$n=5$$



$$\begin{array}{c}
 \stackrel{+5}{\curvearrowright} \stackrel{+6}{\curvearrowright} \stackrel{+6}{\curvearrowright} \stackrel{+5}{\curvearrowright} \\
 K = \textcircled{0} \xrightarrow[5]{10, 15, 20} 100 \\
 K = \textcircled{1} 6, 11, 16, 21 \\
 K = \textcircled{2} 1, 12, 17 \\
 K = \textcircled{3} 8, 13, 18 \\
 K = \textcircled{4} 9, 14, 19
 \end{array}$$

$$\begin{array}{l}
 K = 100 \\
 \boxed{K = K \cdot n} \\
 = 100 \cdot 5 \\
 = \textcircled{0}
 \end{array}$$

$$\begin{array}{l}
 K = 163 \\
 = 163 \cdot 5 \\
 = \textcircled{3}
 \end{array}$$

Improve :

1st code was taking ~~O(n)~~ time, we have to do $\Theta(n)$ place (in original array)

arr = [10, 20, 30, 40, 50] $n=5$
 $k=3$

① reverse the entire array

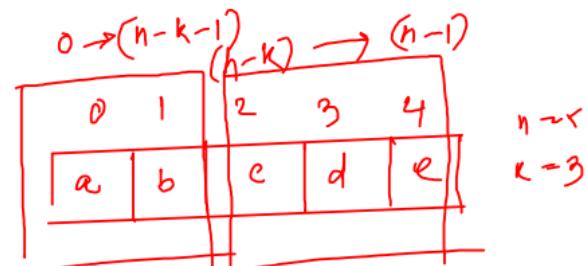
[50, 40, 30, 20, 10]

② reverse 1st $n-k$ elements ($5-3=2$)

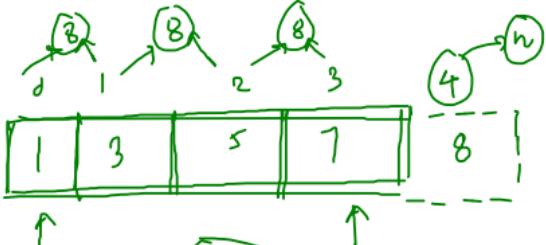
[40, 50, 30, 20, 10]

③ reverse last k elements (3)

[40, 50, 10, 20, 30]



1st $n-k$
↓
1st (5-3)
1st (2)



$i=0$ $1 < 8$ and $8 < 3$
 $i=1$ $3 < 8$ and $8 < 5$
 $i=2$ $5 < 8$ and $8 < 7$

$\left. \begin{array}{l} \\ \\ \end{array} \right\} \text{return } i+1;$



$i>=4$

$3>2>1$

$5>2>1$

$i>=0$

4

~~if ($b \leq a[0]$) return 0;~~ X

~~if ($b > a[n-1]$) return n;~~

for ($i=0$; $i < n-1$; $i+1$) {

~~if ($b \leq a[i+1]$) {~~

~~return $i+1$;~~

}

~~}~~ if ($a[i] \geq b$) {

~~return i ;~~

}