

## \* Adding 2 arrays :

Eg: 

6	1	2	9
4	3	5	2

 → arr1 ( $n=4$ )  

6	9	7	4
---	---	---	---

 → arr2 ( $m=4$ )

$$\begin{array}{r}
 \textcircled{1} \textcircled{1} \\
 4 3 \textcircled{5} \textcircled{2} \\
 + 6 9 \textcircled{7} \textcircled{4} \\
 \hline
 1 1 3 2 6
 \end{array}$$

$$\begin{aligned}
 5+1 &= 1 \textcircled{1} 2 \\
 c &\leftarrow \text{keep} \\
 1+9+9 &= 1 \textcircled{1} 0 \\
 c &\leftarrow \text{keep} \\
 1+4+6 &= 11
 \end{aligned}$$

```

let i = n-1;
let j = m-1;
let carry = 0;
let res = [];
while (i >= 0 && j >= 0) {

```

```
    const sum = arr1[i] + arr2[j] + carry;
```

```
    res.push(sum % 10);
```

```
    carry = sum / 10;
```

```
i--;
```

```
j--;
```

```
}
```

$$\begin{aligned}
 1+5+7 &= 1 \textcircled{1} 3 \\
 \cancel{\text{carry}} &\rightarrow \text{keep} \\
 \text{sum}/10 &= \textcircled{1} \quad \text{sum}/10 = \textcircled{3}
 \end{aligned}$$

→ carry = 0 or 1,  $c > 1$  is not possible  
max possible addition,

$$\begin{aligned}
 \Rightarrow 9+9+1 &= 1 \textcircled{1} 9 \\
 c &= \textcircled{1} \rightarrow \text{keep}
 \end{aligned}$$

```
let i = n-1;  
let j = m-1;
```

'j' is completed  
but 'i' is  
remaining

```
let carry = 0;  
let res = [];
```

```
while (i >= 0 && j <= 0) {
```

```
    const sum = arr1[i] + arr2[j] + carry;
```

```
    res.push(sum % 10);
```

```
    carry = parseInt(sum / 10);
```

```
    i--;
```

```
    j--;
```

```
}
```

```
while (i >= 0) {
```

```
    const sum = arr1[i] + carry;
```

```
    res.push(sum % 10);
```

```
    carry = parseInt(sum / 10);
```

```
    i--;
```

```
}
```

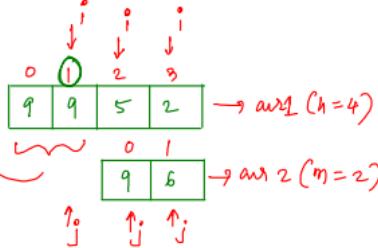
```
if (carry == 1) {
```

```
    res.push(1);
```

```
}
```

```
res = reverse(res);
```

how to  
handle  
remaining  
elements



Carry = 0<sup>1</sup>

① sum = 2 + 6 + 0  
= 8

res = [8]

Carry = 8/10 = 0  
i-- → 2, j-- → 0

= [8, 4]

= [8, 4, 0]

② sum = 5 + 9 + 0  
= 14

= [8, 4, 0, 0]

Carry = 14/10 = 1  
i-- → 1, j-- → -1

③ sum = 9 + 1 = 10

Carry = 10/10 = 1

i--; → 0

④ sum = 9 + 1 = 10

C = 10/10 = 1

i--; → -1

e.g.:

0	1
9	6

$\rightarrow \text{arr1} (n=2)$

i is completed  
but j is remaining

0	1	2	3
9	9	5	2

$\rightarrow \text{arr2} (m=4)$

$\rightarrow i, j$  run together  
 $\rightarrow j$  is alone

\*  $i^{\circ}$  can be alone

In previous

or  $j^{\circ}$  can be alone

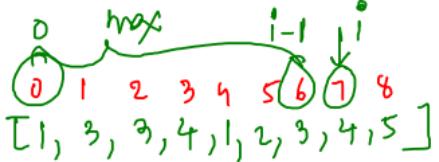
$\rightarrow i, j$  run together

any of these can be possible

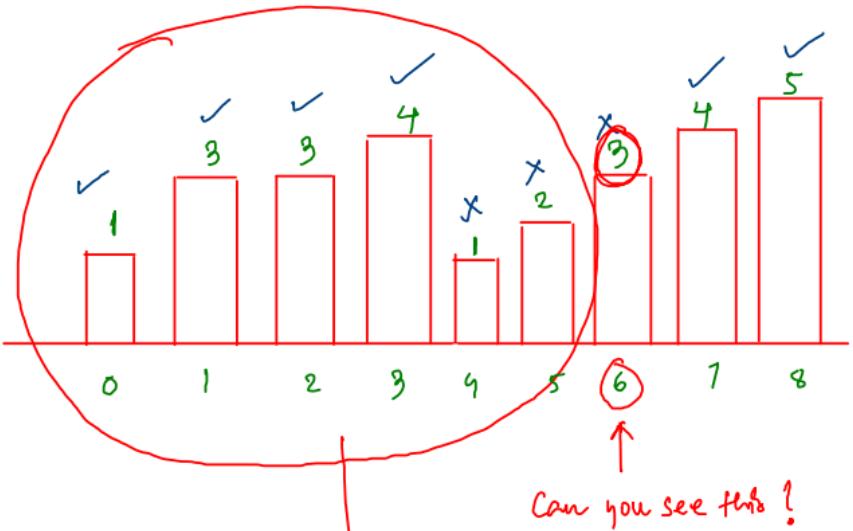
$\rightarrow i$  is alone

$\rightarrow$  handle ' $j$ ' in the same way as you handled ' $i$ '.

## Buildings :



$$09 = 6$$



\* How to decide whether you can see current building or not?

fallest building  
= 4

Can you see this?

$\text{currH} < \text{fallest of left par}$

⇒ This means (3) will be hidden due to tallest (4) on the left

→ only count when  $\text{currH} \geq \text{fallest on left}$

0 1 2 3 4 5 6 7 8  
[1, 3, 3, 4, 1, 2, 3, 4, 5]

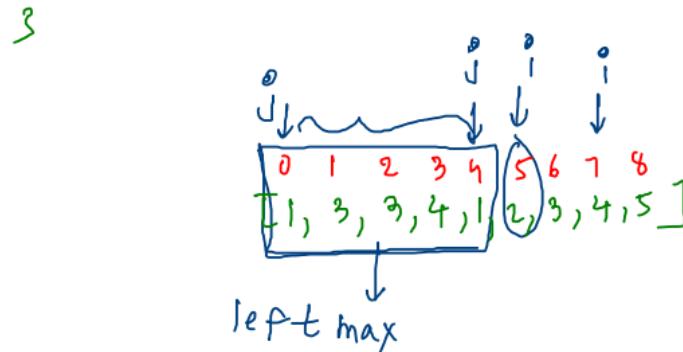
$i=0 \rightarrow 1 \geq -\infty \rightarrow \text{out} = 1$   
 $i=1 \rightarrow 3 \geq 1 \rightarrow \text{out} = 2$   
 $i=2 \rightarrow 3 \geq 3 \rightarrow \text{out} = 3$   
 $i=3 \rightarrow 4 \geq 3 \rightarrow \text{out} = 4$   
 $i=4 \rightarrow 1 \geq 4 \rightarrow \text{hidden}$   
 $i=5 \rightarrow 2 \geq 4 \rightarrow \text{hidden}$   
 $i=6 \rightarrow 3 \geq 4 \rightarrow \text{hidden}$   
 $i=7 \rightarrow 4 \geq 4 \rightarrow \text{out} = 5$   
 $i=8 \rightarrow 5 \geq 4 \rightarrow \text{out} = 6$

```

let cut = 0;
for(let i=0; i<n; i++) {
    if (arr[i] >= tallest on left [0...i-1]) {
        cut++;
    }
}

```

how to get  
this?



```

let leftMax = -∞
for(let j=0; j<i; j++) {
    leftMax = Math.max(leftMax,
                        arr[j]);
}

```

```

863 function countVisibleRoofs(heights) {
864     // Write your code here
865     const n = heights.length;
866     let cnt = 0;
867
868     for (let i = 0; i < n; i++) {
869         let leftMax = -Infinity;
870         for (let j = 0; j < i; j++) {
871             leftMax = Math.max(leftMax, heights[j]);
872
873             if (heights[i] >= leftMax) {
874                 cnt++;
875             }
876
877         }
878
879         return cnt;
880     }

```

$\eta = 5$

$0 \quad 1 \quad 2 \quad 3 \quad 4$ $[a, b, c, d, e]$	$1 + 2 + 3 + 4 = \text{total}$ $= \frac{n(n-1)}{2}$
--	--

$\approx n^2$  iterations

↓

$n$  iterations  
(Optimization needed)

①  $i = 0$       ②  $i = 1$   
 $j = 0 \times$        $(j = 0), \times$

③  $i = 2$   
 $(j = 0, 1), \times$

④  $i = 3$   
 $(j = 0, 1, 2), \times$

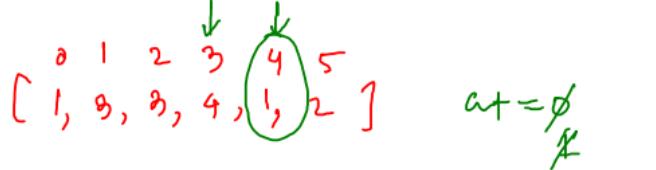
⑤  $i = 4$   
 $(j = 0, 1, 2, 3), \times$

## # Running stream / Carry forward :

- keep track of max element until the current index
- here running max is indicating the leftMax

\* runningMin, runningSum,  
runningProduct, ↪  
runningMax

### Applications



$$\text{runningMax} = -\infty \neq 3$$

$$\textcircled{1} \quad i = 0,$$

$$1 \geq -\infty \rightarrow \text{yes}$$

$$rm = \max(1, -\infty) = 1$$

$$\textcircled{4} \quad i = 3,$$

$$4 > 3 \rightarrow \text{yes}$$

$$rm = \max(4, 3) = 4$$

$$\textcircled{2} \quad i = 1,$$

$$3 \geq 1 \rightarrow \text{yes}$$

$$rm = \max(1, 3) = 3$$

$$\textcircled{5} \quad i = 4$$

$$1 > 4 \rightarrow \text{no}$$

$$rm = \max(4, 1) = 4$$

$$\textcircled{3} \quad i = 2,$$

$$3 > 3 \rightarrow \text{yes}$$

$$rm = \max(3, 3) = 3$$

$$\textcircled{6} \quad i = 5$$

$$4 > 4 \rightarrow \text{no}$$

$$rm = \max(4, 2) = 4$$

## \* Rotate Array : [In place]

e.g:

0	1	2	3	4
10	20	30	40	50

①

0	1	2	3	4
20	30	40	50	10

$k = 1$

②

0	1	2	3	4
30	40	50	10	20

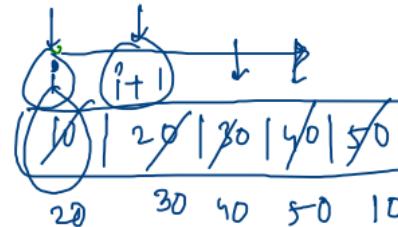
$k = 2$

③

0	1	2	3	4
40	50	10	20	30

$k = 3$

#



$x = 10$

$arr[?] = arr[i+1]$

repeat this  $k$  times

```

902 |   for (let time = 0; time < k; time++) {
903 |     const firstEle = arr[0];
904 |     for (let i = 0; i < n - 1; i++) {
905 |       arr[i] = arr[i + 1];
906 |     }
907 |     arr[n - 1] = firstEle;
908 |   }

```

①  $time = 0$

$\rightarrow n$  iterations

③  $time = 2$

$\rightarrow n$  iterations

②  $time = 1$

$\rightarrow n$  iterations

:

$time = k - 1$   
 $\rightarrow n$  iterations

$\Rightarrow$  total =  $k + n$  # optimised to  $\Theta(n)$

#2:

$k = 0,$	10	20	30	40	50	10
$k = 1,$	20	30	40	50	10	20
$k = 2,$	30	40	50	10	20	30
$k = 3,$	40	50	10	20	30	40
$k = 4,$	50	10	20	30	40	50
$k = 5,$	10	20	30	40	50	10
$k = 6,$	20	30	40	50	10	20
$k = 7,$	30	40	50	10	20	30
$k = 8,$	40	50	10	20	30	40
$k = 9,$	50	10	20	30	40	50
$k = 10,$	10	20	30	40	50	10
$k = 11,$	20	30	40	50	10	20

\* Cycle  $\Rightarrow 4 \% 5$  (help)  $k \% 5$   
 $k \% 4$

$k = 0, 5, 10, 15, 20, 25, \dots \rightarrow 0$

$k = 1, 6, 11, 16, 21, 26, \dots \rightarrow 1$

$k = 2, 7, 12, 17, 22, 27, \dots \rightarrow 2$

$k = 3, 8, 13, 18, 23, \dots \rightarrow 3$

$k = 4, 9, 14, 19, 24, 29, \dots \rightarrow 4$

$\Rightarrow$  given  $k$ , answer for  $k = k \% n$

let  $k = 9298$

$$\rightarrow 9298 \% 5 = 3$$

( $K = 3 \rightarrow 40 \ 50 \ 10 \ 20 \ 30$ )

$arr = [10, 20, 30, 40, 50], k = 3$

① Reverse entire array,

$[50, 40, 30, 20, 10]$

② reverse first  $n-k$  elements

$$n-k = 5-3 = 2$$

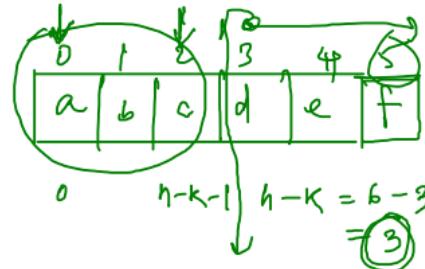
$\begin{matrix} 0 & 1 & 2 & 3 & 4 \\ 40 & 50 & 30 & 20 & 10 \end{matrix}$

③ reverse last  $k$  elements

$\begin{matrix} 40 & 50 & 10 & 20 & 30 \\ 0 & 1 & 2 & 3 & 4 \end{matrix}$

function reverseArr( $arr, s, e$ ) {

    /\*  
      code  
    \*/  
     $\downarrow$   
     $\{$



function rotateArr( $arr, k$ ) {

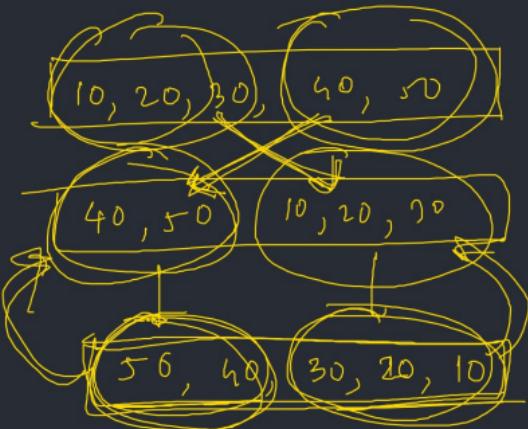
$$K = k \% n;$$

reverseArr( $arr, 0, n-1$ );

reverseArr( $arr, 0, n-k-1$ );

reverseArr( $arr, n-k, n-1$ );

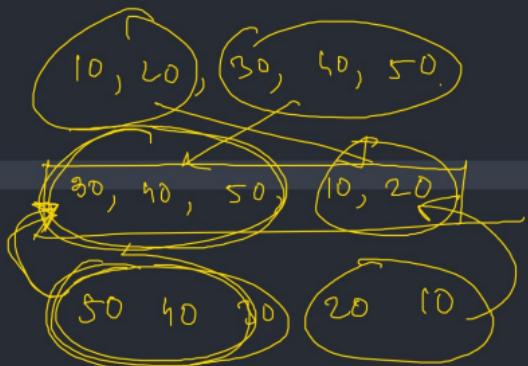
$\}$



$K = 3$

$\leftarrow$   
left

$n - K$



$K = 2$

$\leftarrow$   
left  
 $K$

- \* running stream (v. imp)
- \* cycle  $\Rightarrow$  v. helps
- \* examples are imp to correct the code