

# SECURE VOTING MECHANISM WITH BLIND SIGNATURES

Vote concatenated with random number  $x \leq n$   
 $\Rightarrow (opt \parallel x)$   
 Voter calculates hash of his vote:  $hash(opt \parallel x) = m$   
 Voter calculates a random number  $r$  coprime to  $n_2$ .

Voter to ballot box

a. His vote:  $(opt \parallel x)$   
 b. Signed hash of vote  
 $= (m^{sd}) \% n_2 = SH$

Ballot box first  
 decrypts signed  
 hash of vote with  
 signing authority's  
 public key.  
 $= (SH^{se}) \% n_2 = m$

Next compares  
 $hash(opt \parallel x) == m$

Then vote opt is  
 counted.

Voter

Voter id : ID  
 Voter's Public key :  $vpubk = ve, n_1$   
 Voter's private key:  $vpvk = vd, n_1$   
 Voter vote option number: opt

Voter multiplies signed blinded hash  
 with  $r$  inverse  $= r^{-1}$   
 $= ((m^{sd}) \times r \times \text{inverse}(r)) \% n_2$   
 $= (m^{sd}) \% n_2 \Rightarrow$  effectively hash of vote  
 encrypted with signing authority's pvt key  
 which is basically his sign.

Voter to signing authority

Voter id : ID  
 Block for authentication  $= (hash(ID)^{vd}) \% n_1 = BA$   
 (hash(ID) encrypted with voter private key)  
 Blinded hash  $m' = (m \times (r^{se})) \% n_2$

Signing authority to voter

Signs the blinded hash:  $(m'^{sd}) \% n_2$   
 $= (m^{sd}) \times (r^{se \times sd}) \% n_2$   
 $= (m^{sd}) \times (r^1) \% n_2 = (m^{sd}) \times r \% n_2$   
 ( $m'$  encrypted with voters private key)

Signing authority authenticates  
 the voter first by checking if  
 $hash(ID) == (BA^{ve}) \% n_1$

Signing Authority

Signing Authority Public key:  $spubk = se, n_2$   
 Signing authority Private key:  $spvk = sd, n_2$

Ballot Box