# **Recursion - III Complete**

#### Permutation

To print all the permutations of a string.

Idea: for each character s[i] in the given string, we add a character in the ans string and then solve s.substr(0,i) + s.substr(i+1)

Sample Input:

**ABC** 

Sample Output:

**ABC** 

**ACB** 

**BAC** 

BCA

CAB

CBA

Time Complexity: O(N\*2<sup>n</sup>) Space Complexity: O(2<sup>n</sup>)

```
void permutation(string s, string ans) {

if (s.length() == 0) {
   cout << ans << endl;
   return;
}

for (int i = 0; i < s.length(); i++) {
   char ch = s[i];
   string ros = s.substr(0, i) + s.substr(i + 1);

   permutation(ros, ans + ch);
}</pre>
```

## Tiling problem

Find the number of ways to tile the floor with 1x2 and 1x1 tiles.

Idea: Tile[i] = Tile[i-1] (1x1) + Tile[i-2](1x2)

Time Complexity: O(2<sup>n</sup>)

Space Complexity: O(2<sup>n</sup>) //Memory is used for call stack as well

```
int tilingWays(int n) {
    if (n == 0) {
        return 0;
    }
    if (n == 1) {
        return 1;
    }
    return tilingWays(n - 1) + tilingWays(n - 2);
}
```

## Knapsack [IMP]

Given n items, each item has a certain value and weight. We have to maximize the value of the objects we can accommodate in a bag of weight W.

Idea: For each item, we have two choices, either to include it or exclude it.

Time Complexity: O(2<sup>n</sup>)

Space Complexity: O(2<sup>n</sup>) //space for call stack

# **Friends Pairing Problem**

There are n friends, we have to find all the pairings possible. Each person can be paired with only one person or does not pair with anyone.

Idea: we have two options, i'th friend does not get paired or we have n-1 options to pair it with anyone.

Time Complexity: O(n)
Space Complexity: O(n)

```
int friendsPairing(int n) {
    if (n == 0 || n == 1 || n == 2) {
        return n;
    }
    return friendsPairing(n - 1) + friendsPairing(n - 2) * (n - 1);
}
```

#### CountPaths

```
Find the number of ways to reach e from s.
```

Idea:

We have 6 ways to go forward (1,2,3,4,5,6).

At the starting point s,

Current answer = countPath(s+1,e) + countPath(s+2,e) + countPath(s+3,e) + countPath(s+4,e) + countPath(s+5,e) + countPath(s+6,e)

Time Complexity: O(2<sup>n</sup>)

Space Complexity: O(2<sup>n</sup>)

```
int countPath(int s, int e) {
    if (s == e) {
        return 1;
    }
    if (s > e) {
        return 0;
    }
    int count = 0;
    for (int i = 1; i <= 6; i++) {
        count += countPath(s + i, e);
    }
    return count;
}</pre>
```

#### CountPathMaze

Given a 2D grid, find the number of ways to reach (n-1, n-1).

You can go to (i,j) from (i-1,j) and (i,j-1).

Time Complexity: O(2<sup>n</sup>)

Space Complexity: O(2<sup>n</sup>)

```
int countPathMaze(int n, int i, int j) {
    if (i == n - 1 && j == n - 1) {
        return 1;
    }
    if (i >= n || j >= n) {
        return 0;
    }

    return countPathMaze(n, i + 1, j) +
        countPathMaze(n, i, j + 1);
}
```