Faster Shortest Path Computation for Traffic Assignment

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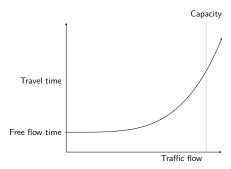
> Department of Engineering Science University of Auckland

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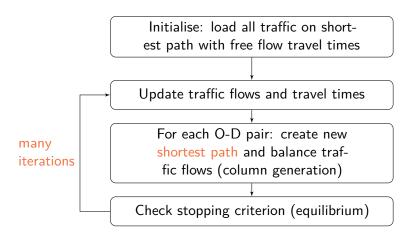
- 1 The traffic assignment problem
- 2 Find faster shortest path algorithms
- 3 Avoiding shortest path calculations
- 4 Conclusion and future work

The traffic assignment problem

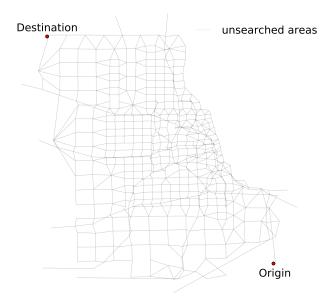
- Assigns traffic to a transportation network
- Used to determine areas of high congestion
- Deals with selecting the best route for vehicles to minimise their travel times (Wardrop equilibrium)
- Path distance is measured by non-linear travel times for capturing congestion effects



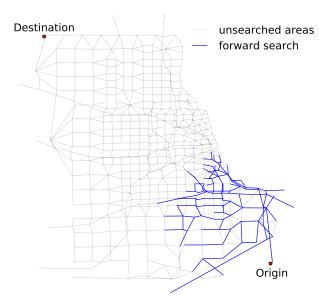
Path equilibration algorithm



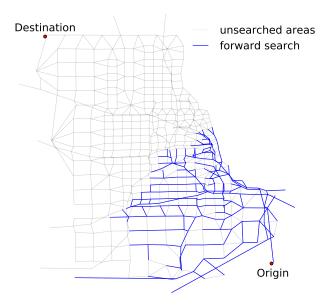
Requires millions of shortest paths to be found



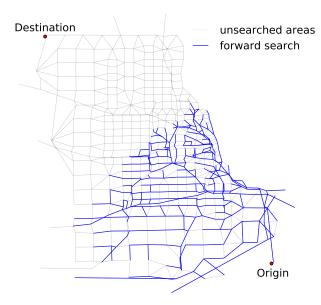
- Chicago Sketch network
- 93,135 O-D pairs
- 546 nodes
- 2,950 arcs



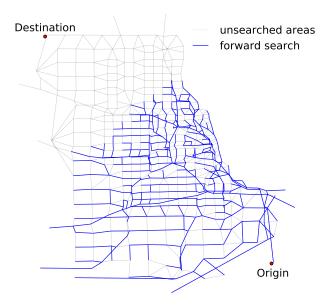
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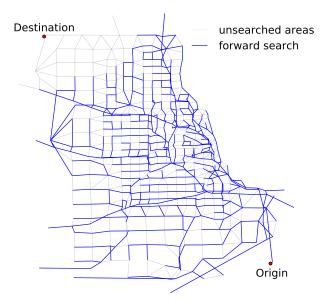
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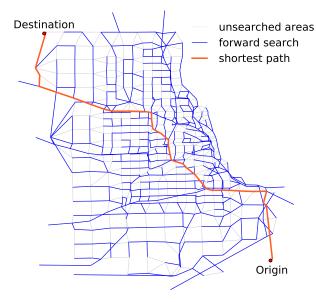
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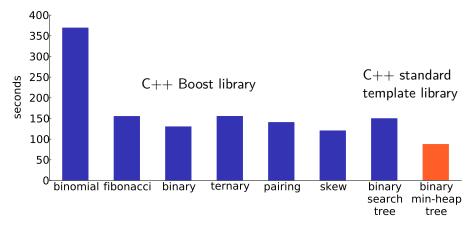
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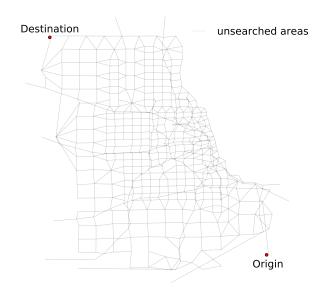


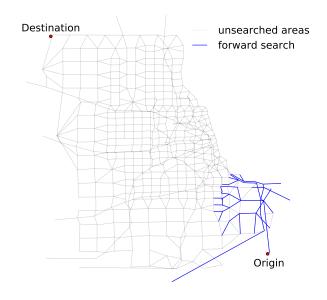
- $\min_{u\in\mathcal{Q}}\left[d_{u}\right]$
- *d_u*: shortest path from origin to *u*
- Q: set of labelled nodes

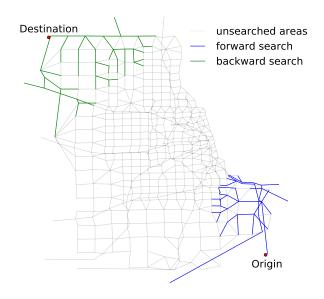
Priority queues

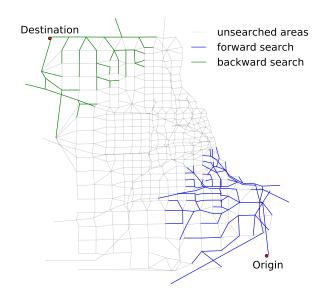
- Priority queue a data structure for storing the searched nodes in priority order so the next location to search can be found easily
- O(1) extract minimum, $O(\log(n))$ insert

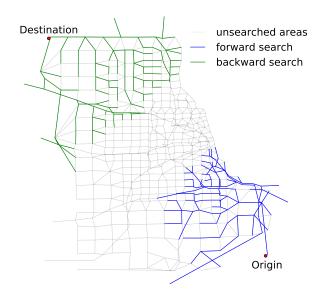


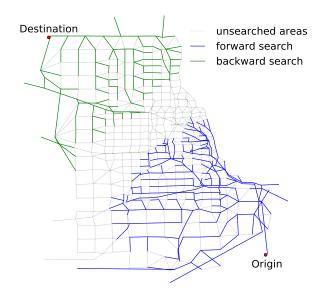


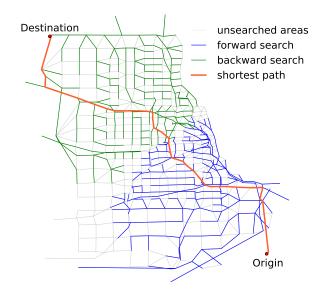


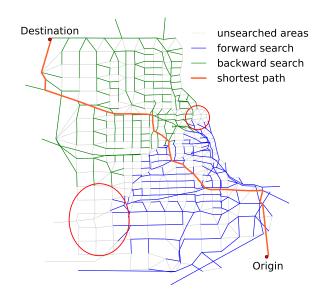


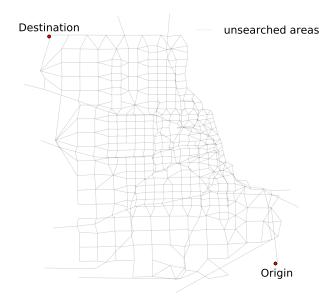




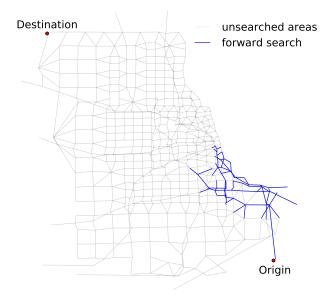




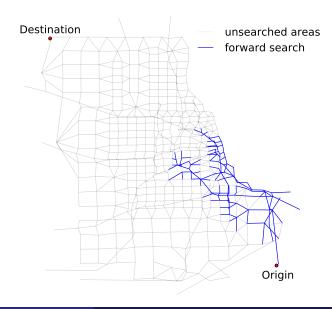




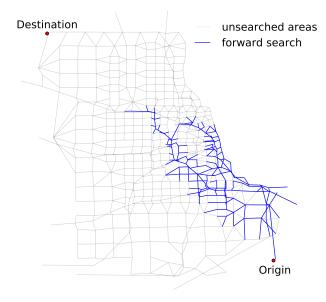
- $\min_{u \in \mathcal{Q}} [d_u + h_u]$
- *d_u*: shortest path from origin to *u*
- h_u: shortest path estimate from u to destination
- Q: set of labelled nodes



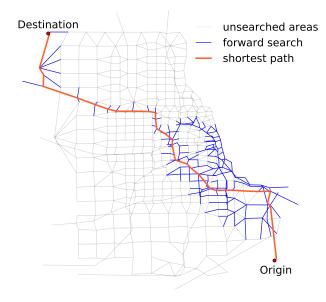
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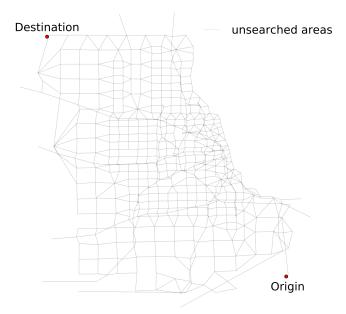
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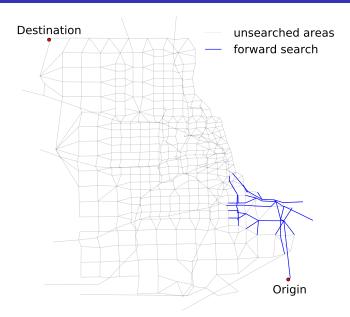


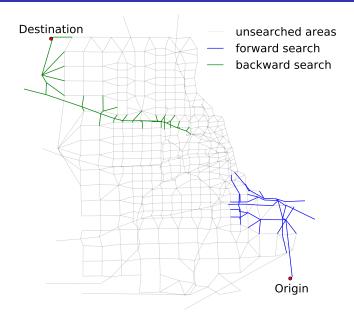
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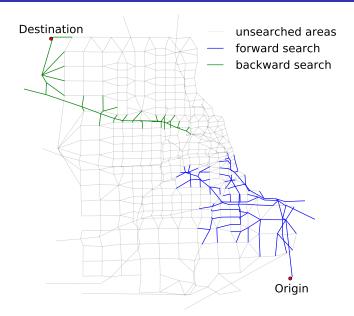


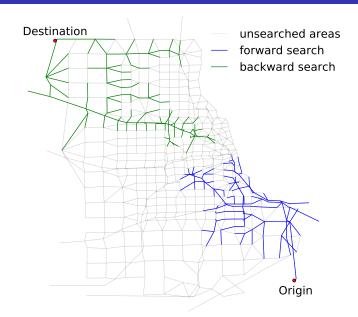
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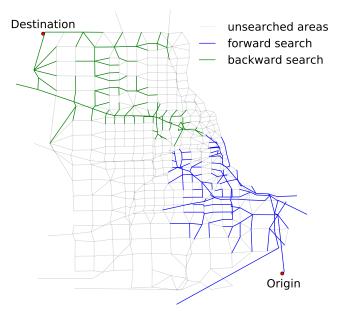


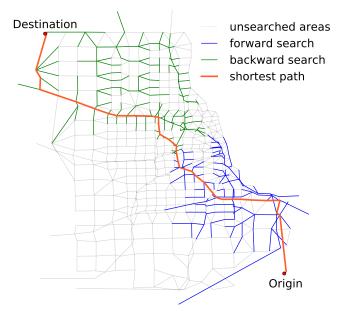




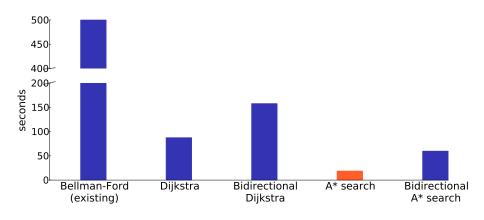




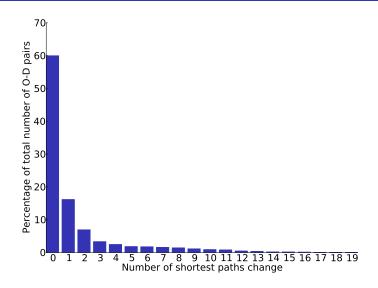




Shortest path algorithm results

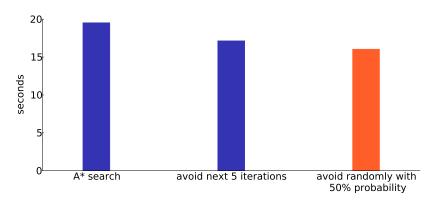


Shortest paths change between iterations for Path Equilibration



Avoid shortest path calculations

- 1 avoid the next few iterations if the shortest paths of the previous two iterations are identical
- 2 randomly avoid the next shortest path calculation in the hope that the shortest path of previous and current iteration are identical



Conclusion and future work

Conclusion

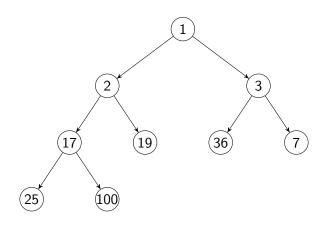
- Best performance: A* search algorithm using min-heap tree with random avoiding strategy
- 30 times faster than the existing implemented Bellman-Ford algorithm
- Bidirectional algorithms are worse compared to the unidirectional ones

Future work

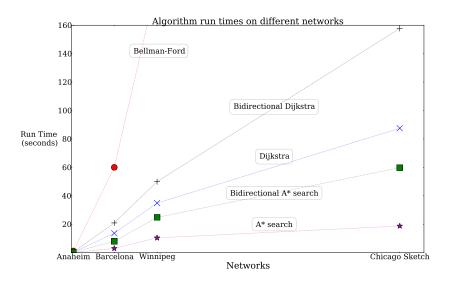
- Pre-processing: A* search with landmarks
- Multi-thread on GPU
- Test the avoiding strategies on other algorithms that solve the traffic assignment problem

Appendix

Binary min-heap tree



Shortest path algorithm results on different networks



Non-linear travel time function

