Competitive Programming Algorithms and Topics

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1	Data Structures		
	1.1	Segment Tree using Pointers	1
	1.2	Range Update Segment Tree	2
	1.3	Range Update Binary Indexed Tree	2
2	Uncategorized		3
	2.1	Longest Increasing Subsequence	3
3	Geometry		3
	3.1	Closest Pair of Points	3
4	Graphs		4
	4.1	Dynammic Connectivity - connected(u,v) query	4

1. Data Structures

1.1. Segment Tree using Pointers

```
/* Segment Tree implementation using pointers */
   /* Can be adapted to Persistent Segment Tree */
   struct node {
    node *left, *right;
     //attributes of node
       //initialize attributes
       left = NULL;
10
       right = NULL;
11
12
13 | void combine(node *ans, node *left, node *right) {
    //combine operation
15 }
16 | void build (node *root, int 1, int r) {
    if(1 == r){
17
1.8
       root->sum = v[1];
19
       return;
20
21
     int m = (1+r) >> 1;
     if(!root->left) root->left = new node();
     if(!root->right) root->right = new node();
23
24
     build(root->left, 1, m);
     build(root->right, m+1, r);
25
     combine(root, root->left, root->right);
26
27 }
28
29 void update(node *root, int 1, int r, int idx, int val) {
    if(1 == r && 1 == idx) {
       //do leaf operation
32
       return;
33
34
     int m = (1+r) >> 1;
35
     if(idx <= m){
       if(!root->left) root->left = new node();
36
37
       update(root->left, l, m, idx, val);
38
39
     else -
40
       if(!root->right) root->right = new node();
41
       update(root->right, m+1, r, idx, val);
42
43
     combine(root, root->left, root->right);
44
   node* query(node *root, int 1, int r, int a, int b) {
47
     if(1 == a && r == b){
48
       return root;
49
50
     int m = (1+r) >> 1;
51
     if(b <= m){
52
       if(!root->left) root->left = new node();
53
       return query(root->left, l, m, a, b);
54
     else if (m < a) {
       if(!root->right) root->right = new node();
57
       return query (root->right, m+1, r, a, b);
58
59
     if(!root->left) root->left = new node();
     if(!root->right) root->right = new node();
     node *left = query(root->left, l,m,a,m);
```

```
node *right = query(root->right, m+1, r, m+1, b);
node *ans = new node();
combine(ans, left, right);
return ans;
}
```

1.2. Range Update Segment Tree

```
/* Range update Segment Tree Implementation */
   /* The first node (ROOT) is defined to 1 (1 - index impl) */
   /* N is the maximum number of elements given by the statement */
   /* Lazy can be inside node structure instead of being another structure */
   #define ROOT 1
   #define N MAX INPUT
a
   struct node {
    //attributes of node
10
11
12
1.3
   node tree[4*N];
14
   node lazy[4*N];
15
16
   node combine (node a, node b) {
17
     node res:
18
     //combine operations
19
     return res;
20
21
   void propagate(int root, int 1 , int r){
22
     //return if there is no update
23
     //update tree using lazy node
24
25
     if(1 != r){
26
       //propagate for left and right child
27
28
     //reset lazy node
29
30
   void range_update(int root, int 1, int r, int a, int b, long long val) {
31
32
    if(l == a && r == b) {
33
       //lazy operation using val
34
       return;
35
36
37
     int m = (1+r)/2;
38
39
     if(b <= m) range update(2*root, 1, m, a, b, val);
40
     else if(m < a) range_update(2*root+1, m+1, r, a, b, val);</pre>
41
     else {
42
       range_update(2*root, 1, m, a, m, val);
43
       range_update(2*root+1, m+1, r, m+1, b, val);
44
45
     propagate(root, l , r);
46
47
     propagate (2*root, 1, m);
48
     propagate (2*root+1, m+1, r);
49
     tree[root] = combine(tree[2*root], tree[2*root+1]);
50
51
52 | node query(int root, int 1, int r, int a, int b) {
53
     propagate(root, l, r);
54
     if(l == a && r == b) return tree[root];
55
     int m = (1+r)/2;
```

1.3. Range Update Binary Indexed Tree

```
/* Range Update Binary Indexed Tree Implementation */
2 /* Tree is 1 - index */
3 /* Point Update Binary Indexed Tree operations are used as auxiliar */
4 /* N is defined as the maximum number of elements (given by the statement) */
6 #define N MAX INPUT
8 | int bit[2][N+1];
10 | void init(int n) {
    for (int i=1; i<=n; i++) {</pre>
11
12
       bit[0][i] = 0;
13
       bit[1][i] = 0;
14
15 }
16
17
    //auxiliar functions
18
   void update(int *bit, int idx, int val, int n) {
    for(int i = idx; i <= n; i += i&-i) {</pre>
20
21
       bit[i]+=val;
22
23
24
25 | int query(int *bit, int idx) {
     int ans = 0:
2.7
     for(int i=idx; i>0; i -= i&-i) {
28
       ans += bit[i];
29
30
     return ans;
31
32
33
    //end of auxiliar functions
34
35 | void range_update(int 1, int r, int val, int n) {
36
     update(bit[0], 1, val, n);
37
     update(bit[0], r+1, -val, n);
38
     update(bit[1], l, val*(l-1), n);
39
     update(bit[1], r+1, -val*r, n);
40
41
42 int prefix querv(int idx) {
43
     return query(bit[0],idx)*idx - query(bit[1], idx);
44
45
46 | int range_query(int 1, int r){
    return prefix_query(r) - prefix_query(l-1);
47
48 }
```

2. Uncategorized

2.1. Longest Increasing Subsequence

```
/* Longest Increasing Subsequence Implementation */
   /* N is defined as the maximum array size given by the statement */
   #define N MAX_N
4
5
   int v[N];
   int lis[N+1];
8
   void calculate_lis(int n) {
    for(int i=1; i<=n; i++) lis[i] = INT_MAX;</pre>
10
11
     lis[0] = INT_MIN;
12
     for(int i=0; i<n; i++) {</pre>
       int index = lower_bound(lis, lis+n+1, v[i]) - lis;
13
14
15
       lis[index+1] = min(lis[index+1], v[i]);
16
17
```

3. Geometry

3.1. Closest Pair of Points

```
/* Closest Pair of Points Problem Implementation */
   /* Divide and Conquer Approach */
   /* Using the observation of only checking points inside min_dist x min_dist
       from mid */
   /* Binary search boosts search for the right border start point */
   struct vec2 {
7
    long long x, y;
8
10
   bool cmp(vec2 a, vec2 b) {
11
    return a.x < b.x || (a.x == b.x && a.y < b.y);
12
13
14
   pair<vec2, vec2> ans;
15
   long long solve(vector<vec2> &a) {
16
17
18
     long long mid = a[a.size()/2].x;
19
     int n = a.size();
20
21
     vector<vec2> 1:
22
     vector<vec2> r;
23
     int i = 0;
24
     for(; i < a.size()/2; i++) l.push_back(a[i]);</pre>
25
     for(; i < a.size(); i++) r.push_back(a[i]);</pre>
26
     long long d = LLONG_MAX;
27
2.8
29
     if(l.size() > 1) {
30
       d = min(d, solve(1));
31
     } if(r.size() > 1) {
32
       d = min(d, solve(r));
33
34
35
     a.clear();
36
     vector<vec2> 11;
```

```
38
     vector<vec2> rr:
39
40
41
     int j = 0;
42
     i = 0;
43
     for (int k=0; k<n; k++) {</pre>
44
        if(i < 1.size() && j < r.size()){</pre>
45
          if(r[j].y <= l[i].y){
46
            if((r[j].x - mid) * (r[j].x - mid) < d) {
47
              rr.push_back(r[j]);
48
49
            a.push_back(r[j++]);
50
51
          else
            if((l[i].x - mid) * (l[i].x - mid) < d) {
52
53
              ll.push_back(l[i]);
54
55
            a.push_back(l[i++]);
56
57
58
        else if(i < l.size()){
          if((l[i].x - mid) * (l[i].x - mid) < d) {
59
60
            ll.push_back(l[i]);
61
62
          a.push_back(l[i++]);
63
64
65
          if((r[j].x - mid) * (r[j].x - mid) < d) {
66
              rr.push_back(r[j]);
67
68
            a.push_back(r[j++]);
69
70
71
72
      for(int i = 0; i < ll.size(); i++) {</pre>
73
74
        int ini = 0, end = rr.size()-1;
75
       int j;
76
        while(ini < end) {</pre>
77
          j = (ini + end) / 2;
78
          if((rr[j].y - ll[i].y) * (rr[j].y - ll[i].y) > d && rr[j].y < ll[i].y)
79
            ini = j+1;
80
          else end = j;
81
82
        j = ini;
83
84
        for(; j < rr.size(); j++) {</pre>
85
          if((rr[j].y - ll[i].y) * (rr[j].y - ll[i].y) > d) break;
86
          long long cur = (ll[i].x - rr[j].x)*(ll[i].x - rr[j].x)
87
                   +(ll[i].y - rr[j].y)*(ll[i].y - rr[j].y);
88
          if(cur < d) {
            d = cur;
89
90
            long long cur2 = (ans.first.x - ans.second.x) * (ans.first.x -
        ans.second.x)
91
                  +(ans.first.y - ans.second.y) *(ans.first.y - ans.second.y);
            if(cur < cur2)</pre>
92
93
              ans = { ll[i], rr[j] };
94
95
96
97
     return d;
98
```

3.2. Convex Hull - Monotone Chain Algorithm

```
/* Convex Hull - Monotone Chain */
   /* Generates Upper and Lower Hull */
   /* It is needed to give array of points ordered by <x,y> */
3
   vector < pair<int, int> > upper, lower;
   vector< pair<int, int> > hull;
   int cross(pair<int, int> & a, pair<int, int> & b, pair<int, int> & c) {
10
       pair<int, int> vecl(b.ff - a.ff, b.ss - a.ss);
11
       pair<int, int> vec2(c.ff - b.ff, c.ss - b.ss);
12
       return vec1.ff*vec2.ss - vec1.ss*vec2.ff;
13
14
15
   void calculate_upper(vector< pair<int, int> > & p) {
       for (int i=0; i<p.size(); i++) {</pre>
16
17
              while(upper.size() >= 2 && cross(upper[upper.size()-2],
       upper.back(), p[i]) >= 0){
18
                   upper.pop_back();
19
20
              upper.push_back(p[i]);
21
22
23
24
   void calculate_lower(vector< pair<int, int> > & p) {
25
     for(int i=0; i<p.size(); i++){
26
              while(lower.size() >= 2 && cross(lower[lower.size()-2],
       lower.back(), p[i]) <= 0){</pre>
27
                   lower.pop_back();
28
29
              lower.push_back(p[i]);
30
31
32
   void merge_hull() {
     for(int i=0; i<upper.size(); i++) hull.push_back(upper[i]);</pre>
35
     for(int i=lower.size()-2; i>0; i--) hull.push_back(lower[i]);
36
```

3.3. Shoelace Formula for Polygon Area

```
/* Shoelace formula */
   /* Calculate area of convex polygon */
   /* Points given in clockwise/counterclockwise order */
3
   int cross(pair<int, int> & a, pair<int, int> & b) {
     return a.ff*b.ss - a.ss*b.ff;
7
8
   int shoelace(vector< pair<int, int> > & p) {
    int area = 0;
10
11
     for(int i=0; i<hull.size(); i++){</pre>
12
       area += cross(hull[i], hull[(i+1)%hull.size()]);
13
14
    return abs(area/2);
15
```

4. Graphs

4.1. Dynammic Connectivity - connected(u,v) query

```
1 /* Dynammic Connectivity Implementation */
```

```
2 | /* Uses Divide and Conquer Offline approach */
 3 /* Able to answer if two vertex <u,v> are connected */
   /* No multi-edges allowed */
5 /* DSU + Rollback is used to backtrack merges */
6 | /* N is defined as the maximum graph size given by input */
   #define N MAX INPUT
10 | int uf[N];
11 | int sz[N];
13 | struct event{
    int op, u, v, l, r;
15
     event() {}
16
    event(int o, int a, int b, int x, int y) : op(o), u(a), v(b), l(x), r(y) {}
17
1.8
19 map< pair<int, int>, int > edge_to_l;
20 | stack< pair<int*,int> > hist;
21 | vector<event> events;
23 int init(int n) {
    for(int i=0; i<=n; i++) {
25
       uf[i] = i;
26
       sz[i] = 1;
27
28
29
30
   int find(int u) {
31
     if(uf[u] == u) return u;
32
     else return find(uf[u]);
33
34
35 void merge(int u, int v) {
36
     int a = find(u);
37
     int b = find(v);
38
     if(a == b) return;
39
     if(sz[a] < sz[b]){
40
       hist.push(make_pair(&uf[a], uf[a]));
41
       uf[a] = b;
42
       hist.push(make_pair(&sz[b], sz[b]));
       sz[b] += sz[a];
43
44
45
46
       hist.push(make pair(&uf[b], uf[b]));
47
       hist.push(make_pair(&sz[a], sz[a]));
48
       uf[b] = a;
49
       sz[a] += sz[b];
50
51 }
52
53 int snap() {
54
    return hist.size();
55 }
56
57 void rollback(int t) {
     while(hist.size() > t){
59
       pair<int*, int> aux = hist.top();
60
       hist.pop();
61
        *aux.first = aux.second;
62
63
64
65 void solve(int 1, int r){
66 | if(1 == r){
```

```
if(events[1].op == 2){
 68
           if(find(events[1].u) == find(events[1].v)) cout << "YES" << endl;</pre>
 69
           else cout << "NO" << endl:
 70
 71
        return;
 72
 73
 74
      int m = (1+r)/2;
 75
      //doing for [L,m]
      int t = snap();
 76
 77
      for (int i=1; i<=r; i++) {</pre>
 78
        if(events[i].op == 0 || events[i].op == 1){
 79
          if(events[i].1 <= 1 && m <= events[i].r) merge(events[i].u,</pre>
         events[i].v);
 81
      solve(1, m);
 82
      rollback(t);
 83
 84
 85
      //doing for [m+1, R]
 86
      t = snap();
      for(int i=1; i<=r; i++){
 88
        if(events[i].op == 0 || events[i].op == 1){
 89
          if(events[i].1 <= m+1 && r <= events[i].r) merge(events[i].u,</pre>
         events[i].v);
 91
 92
      solve (m+1, r);
 93
      rollback(t):
 94
 95
 96
    void offline_process() {
 97
      int n, q;
      cin >> n >> q; //number of vertex and queries
 98
      init(n);
 99
100
      for(int i=0; i<q; i++) {
101
         string op;
102
         int u, v;
103
         cin >> op >> u >> v; //add, remove or query for u, v
        if(u > v) swap(u,v);
104
105
         if(op == "add") {
106
           events.push_back(event(0, u, v, i, -1));
107
           edge_to_l[make_pair(u,v)] = i;
108
109
         else if(op == "rem"){
110
           int l = edge_to_l[make_pair(u,v)];
111
           events.push_back(event(1, u, v, 1, i));
           events[l].r = i;
112
113
114
         else if(op == "conn"){
115
           events.push_back(event(2, u, v, -1, -1));
116
117
      for(int i=0; i<q; i++) {</pre>
118
119
        if(events[i].op == 0){
120
          if(events[i].r == -1){
121
             events[i].r = events.size();
122
             events.push_back(event(1, events[i].u, events[i].v, events[i].l,
         events[i].r));
123
124
125
126
```

5. Math and Number Theory

5.1. Binomial Coefficient DP

```
/* Dynammic Programming for Binomial Coefficient Calculation */
/* Using Stiefel Rule C(n, k) = C(n-1, k) + C(n-1, k-1) */

int binomial(int n ,int k) {
   int c[n+10][k + 10];
   memset(c, 0 , sizeof c);
   c[0][0] = 1;
   for(int i = 1;i<=n;i++) {
      for(int j = min(i, k);j>0;j--) {
        c[i][j] = c[i-1][j] + c[i-1][j-1];
    }
}
return c[n][k];
}
```

5.2. Erathostenes Sieve + Logn Prime Factorization

```
/* Erasthostenes Sieve Implementation */
   /* Calculate primes from 2 to N */
   /* lf[i] stores the lowest prime factor of i(logn factorization) */
5 | bitset<N> prime;
6 | int lf[N];
8 void run sieve(int n) {
9 for(int i=0; i<=n; i++) lf[i] = i;</pre>
10 prime.set();
    prime[0] = false:
    prime[1] = false;
13
     for (int p = 2; p*p <= n; p++) {
14
       if(prime[p]){
15
         for (int i=p*p; i<=n; i+=p) {</pre>
16
           prime[i] = false;
           lf[i] = min(lf[i], p);
17
18
19
20
21 }
```

5.3. Matrix Exponentiation

```
/* Matrix Exponentiation Implementation */
   typedef vector< vector<int> > Matrix;
   Matrix operator * (const Matrix & a, const Matrix & b) {
     Matrix c(a.size(), vector<int>(b[0].size()));
     for(int i = 0; i<a.size(); i++){</pre>
        for(int j = 0; j < b[0].size(); j++) {</pre>
          for(int k = 0; k<b.size(); k++){</pre>
10
            c[i][j] += (a[i][k]*b[k][j]);
11
12
13
14
     return c;
15 }
16
17 Matrix exp(Matrix & a, int k) {
18 | if(k == 1) return a;
```

6. String Algorithms

6.1. KMP Failure Function + String Matching

```
/* Knuth - Morris - Pratt Algorithm */
    /* Failure Function for String Matching */
3
   int pi[N];
   string s, t;
   void prefix(int n) {
     pi[0] = 0;
     for(int i = 1; i < n; i++) {
10
11
       pi[i] = pi[i-1];
12
       while(pi[i] > 0 \& \& t[i] != t[pi[i]]) pi[i] = pi[pi[i]-1];
13
       if(t[i] == t[pi[i]]) pi[i]++;
14
15
16
17
   void matching (int n) {
18
    int j = 0;
     for(int i=0; i<n; i++) {</pre>
19
20
       while(j > 0 \&\& s[i] != t[j]) j = pi[j-1];
       if(s[i] == t[j]) j++;
21
22
       if(j == t.size()){
         cout << "match in " << j-t.size()+1 << endl;</pre>
23
24
          j = pi[j-1];
25
26
27
```

6.2. Z-Function

```
/* Z-function */
   /* Calculate the size K of the largest substring which is a prefix */
   vector<int> z;
   void make(string s) {
     int n = s.size();
     z.resize(n);
9
     z[0] = 0;
1.0
     int 1 = 0, r = 0;
11
     for(int i=1; i<n; i++) {</pre>
12
       if(i > r){
13
         1 = i;
14
         r = i;
15
       z[i] = min(z[i-1], r-i+1);
16
17
       while(i + z[i] < n && s[i + z[i]] == s[z[i]]) z[i]++;
18
       if(i + z[i] > r) {
19
         l = i;
20
         r = i + z[i] - 1;
21
22
23
```

6.3. Rolling Hash

```
/* Rolling Hash Implementation */
   /* Uses 1-indexed string */
   long long BASE = 137
   long long PRIME = (int)1e9+7;
   long long hash[N+1];
   long long base[N+1];
   long long invBase[N+1];
1.0
11 long long expo(long long a, long long k) {
12 | if(k == 0) return 1LL;
13
     else if(k == 1) return a;
14
     long long aux = expo(a, k/2);
     aux %= PRIME;
15
16
     aux *= aux;
     aux %= PRIME;
17
     if(k%2) aux *= a;
19
     aux %= PRIME;
20
     return aux;
21
22
23
   void calculate(string a) {
     base[0] = 1;
24
25
     invBase[0] = 1;
26
     hash[0] = 0;
     for(int i=1; i<=a.size(); i++) {</pre>
       hash[i] += BASE*hash[i-1] + a[i-1];
28
2.9
       hash[i] % = PRIME;
30
       base[i] = base[i-1] *BASE;
31
       base[i] %= PRIME;
32
       invBase[i] = expo(base[i], PRIME-2);
33
34
35
   long long range_hash(int i, int j) {
37
     return ((h[j] - h[i-1])*invBase[j-i+1])%PRIME;
38
```

6.4. Suffix Array + Linear Sort

```
/* Suffix Array using Counting Sort Implementation */
   /* rnk is inverse of sa arrav */
   /* aux arrays are needed for sorting step */
4 /* inverse sorting (using rotating arrays and blocks of power of 2) */
   /* rmg data structure needed for calculating lcp of two non adjacent
        suffixes sorted */
7 | int rnk[N], tmp[N], sa[N], sa_aux[N], lcp[N];
8 int block=0, n;
9 string s;
11 | bool suffixcmp(int i, int j) {
12 | if(rnk[i] != rnk[j]) return rnk[i] < rnk[j];</pre>
    i+=block, j+=block;
15
16
     return rnk[i] < rnk[j];</pre>
17 }
1.8
19 void suffixSort(int MAX_VAL) {
    for(int i=0; i<=MAX_VAL; i++) tmp[i] = 0;</pre>
```

```
for(int i=0; i<n; i++) tmp[rnk[i]]++;</pre>
     for(int i=1; i<=MAX VAL; i++) tmp[i] += tmp[i-1];</pre>
22
23
     for(int i = n-1; i>=0; i--) {
          int aux = sa[i]-block;
24
25
          aux%=n;
26
         if(aux < 0) aux+=n;
27
          sa_aux[--tmp[rnk[aux]]] = aux;
28
29
     for(int i=0; i<n; i++) sa[i] = sa_aux[i];</pre>
30
     tmp[0] = 0;
     for(int i=1; i<n; i++) tmp[i] = tmp[i-1] + suffixcmp(sa[i-1], sa[i]);</pre>
31
32
     for(int i=0; i<n; i++) rnk[sa[i]] = tmp[i];</pre>
33
34
35 void build_sa() {
     s+='\0';
36
37
     n++;
38
     for(int i=0; i<n; i++) {
39
       sa[i] = i;
40
       rnk[i] = s[i];
41
       tmp[i] = 0;
42
43
     suffixSort (256):
44
     block = 1;
45
     while (tmp[n-1] != n-1) {
46
       suffixSort(tmp[n-1]);
47
       block*=2;
48
     for(int i=0; i<n-1; i++) sa[i] = sa[i+1];</pre>
49
50
     n--;
51
     tmp[0] = 0;
     for(int i=1; i < n; i++) tmp[i] = tmp[i-1] + suffixcmp(sa[i-1], sa[i]);</pre>
52
53
     for(int i=0; i<n; i++) rnk[sa[i]] = tmp[i];</pre>
54
55
   void calculate lcp(){
56
57
     int last = 0;
     for (int i=0; i<n; i++) {</pre>
58
59
       if(rnk[i] == n-1) continue;
       int x = rnk[i];
60
61
       lcp[x] = max(0, last-1);
       while (sa[x] + lcp[x] < n \&\& sa[x+1] + lcp[x] < n \&\& s[sa[x]+lcp[x]] ==
62
        s[sa[x+1]+lcp[x]]){
         lcp[x]++;
64
65
       last = lcp[x];
66
67
68
69 int lcp(int x, int y) {
    if(x == y) return n - x;
70
71
    if(rnk[x] > rnk[y]) swap(x,y);
72
    return rmq(rnk[x], rnk[y]-1); //segtree
73
```