

Bases de Dados

RELATIONAL MODEL

MAIN EXERCISE

1. REFS relation is an external view containing all the relevant attributes of a bibliographic references' database.

REFS (T, A, I, C, S#, R, S, E, V, N, PI, PF, Y, M).

Referencias¹ (Title, Author, Institution, Email, ISSN#, Journal, Acronym, Editor, Volume, Issue, StartPage, EndPage, Year, Month).

An author belongs to an institution I and has an email C; wrote an article with title T, in a journal R, also known by the acronym S, with the international number S#, edited by E; the article appears in the volume V, number N of the journal, published in year Y and month M, from page PI to page PF. The declared functional dependencies are:

$$\mathbf{D} = \{ \begin{array}{l} S \rightarrow S\#, R, E \\ S\# \rightarrow S \\ S, V, N \rightarrow Y, M \\ S, V, N, PI \rightarrow T, PF \\ A \rightarrow I, C \end{array} \}$$

An article may have several authors who, in turn, can write several articles.

- Is $X = S\# V N A$ a key of Referencias? Justify.
- Obtain a decomposition of Referencias in the third normal form with dependency preservation.
- Compare the decomposition obtained in b) with the initial relation. What are the advantages of the third normal form?
- On the Revista² (S, S#, R, E) relation, the acronym S and the international series number S# are keys and S is the primary key. Through an example, show that, although this relation is in the third normal form, the definition of a primary key is not enough to assure that functional dependencies are not violated. What ability should the DBMS have to assure this?
- Check if the obtained decomposition is in the Boyce-Codd Normal Form (BCNF).

¹ References

² Journal

MORE EXERCISES

- 2) Given the relation $R(A,B,C,D,E)$ and the set of dependencies $F=\{AB\rightarrow C, DE\rightarrow C, B\rightarrow D\}$
- Find all the keys for R .
 - Check if the decomposition $R_1(A,B,C), R_2(C,D,E), R_3(B,D)$ assures a lossless join.
 - Does the conclusion you have reached on b) changes if you add the relation $R_4(A,B,E)$ to the decomposition?
 - Obtain a decomposition in the third normal form.
 - Suppose you decompose the relation $R(A,B,C,D,E)$ in relation $S(A,B,C)$ and other relations. If R functional dependencies are $F=\{A\rightarrow D, BD\rightarrow E, AC\rightarrow E, DE\rightarrow B\}$, what is the minimal basis for S functional dependencies?
- 3) Given the relation $R(C,S,J,D,P,Q,V)$ and the set of functional dependencies $F=\{JP\rightarrow C, SD\rightarrow P, J\rightarrow S\}$
- Find the keys of relation R .
 - The decomposition $R_1(S,D,P), R_2(J,S), R_3(C,J,D,Q,V)$ is in the third normal form? Justify. If it's not, obtain a third normal form decomposition of R .
 - See if the functional dependency $JP\rightarrow C$ is preserved in the decomposition presented in the previous question.
 - For the same relation R but, for the set of functional dependencies $F_1=\{C\rightarrow CSJDPQV, JP\rightarrow C, SD\rightarrow P, J\rightarrow S\}$, obtain a minimal basis for the full set of functional dependencies.
- 4) Consider the relation $R(CPHSAN)$, where C is the course, P is the teacher, H is the hour, S is the room, A is the student and N is the grade, and the following referential integrity constraints:
- Each course has a responsible teacher;
 - In a specific hour, only one course can be at a specific room;
 - A teacher can only be at one room in a certain hour;
 - Each student has only one grade for each course;
 - A student can only be at one room in each moment.
- Based on the previous integrity constraints, what is the set of functional dependencies?
 - What is the key of R ?
 - Decompose R into the third normal form.
 - Check if any of the relations obtained in the previous decomposition is not in BCNF.