FXPERIMENT 7

Experiment No: 7 **Date:** 10/04/2021

Aim: Implementation of Single Source Shortest Path Algo

(Dijkstra's Algo) and estimate its step count

Theory:

Dijkstra's Shortest Path Algorithm

- ➤ Given a graph and a source vertex in the graph, find shortest paths from source to all vertices in the given graph.
- Dijkstra's algorithm is very similar to Prim's algorithm for minimum spanning tree.
- Like Prim's MST, we generate a SPT (shortest path tree) with given source as root.
- ➤ We maintain two sets, one set contains vertices included in shortest path tree, other set includes vertices not yet included in shortest path tree.
- At every step of the algorithm, we find a vertex which is in the other set (set of not yet included) and has a minimum distance from the source.

Time Complexity

- ➤ The time complexity of the above code/algorithm looks O(V^2) as there are two nested while loops.
- If we take a closer look, we can observe that the statements in inner loop are executed O(V+E) times (similar to BFS).
- The inner loop has decreaseKey() operation which takes O(LogV) time.

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➤ So overall time complexity is O(E+V)*O(LogV) which is

$$O((E+V)*LogV) = O(ELogV)$$

- Note that the above code uses Binary Heap for Priority Queue implementation.
- ightharpoonup Time complexity can be reduced to O(E+VLogV) using Fibonacci Heap.
- ➤ The reason is, Fibonacci Heap takes O(1) time for decrease-key operation while Binary Heap takes O(Logn) time.

```
Begin Algorithm Dijkstra (G, s)
    For each vertex v ∈ V
        d [v] ← ∞ // an estimate of the min-weight path from s to v
2
3
4
   d[s] \leftarrow 0
                                                                         Min-heap
   S \leftarrow \mathcal{O} // set of nodes for which we know the min-weight path from s
   Q - V // set of nodes for which we know estimate of min-weight path from s
    While Q \neq \emptyset \leftarrow done |V| times = O(V) time
       u ← EXTRACT-MIN(Q) ← Each extraction takes O(logV) time
9
       S ← S U {u}
       For each vertex v such that (u, v) \in E } done O(E) times totally
10
          If v \in Q and d[v] > d[u] + w(u, v) then
11
                                                       It takes O(logV) time when
12
            d[v] \leftarrow d[u] + w(u, v)
                                                       done once
            Predecessor (v) = u
13
13
          End If
                     Overall Complexity: O(V) + O(V) + O(VlogV) + O(ElogV)
        End For
14
                     Since |V| = \Omega(|E|), the VlogV term is dominated by the
     End While
15
                     ElogV term. Hence, overall complexity = O(|E|*log|V|)
16 End Dijkstra
```

Algorithm

```
Algorithm ShortestPaths(v,cost,dist,n)
// dist[j], 1<jn, is set to the length of the shortest</pre>
// path from vertex v to vertex j in a digraph G with n
// vertices. dist[v] is set to zero. G is represented by its
// cost adjacency matrix cost[1: n. 1:n].
{
      for i:=1 to n do
      { // Initialize S.
             S[i]:= false; dist[i]:= cost[v, i];
      }
      S[v]:= true; dist[v] := 0.0; // Put v in S.
      for num=2 to n - 1 do
      {
      // Determine n - 1 paths from v.
      Choose u from among those vertices not
      in S such that dist[u] is minimum;
      S[u]:= true; // Put u in S.
      for (each w adjacent to u with S[w] = false) do
```

```
// Update distances.

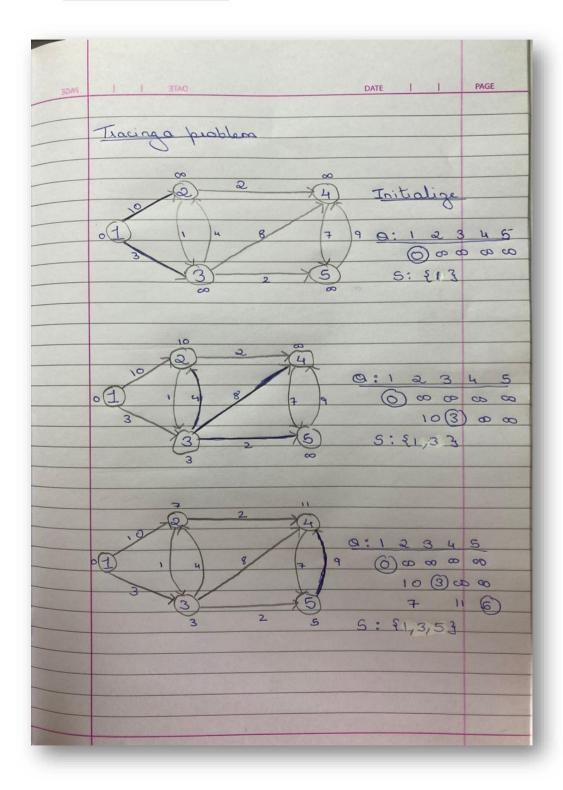
if (dist[w]> dist[u] + cost[u, w])) then

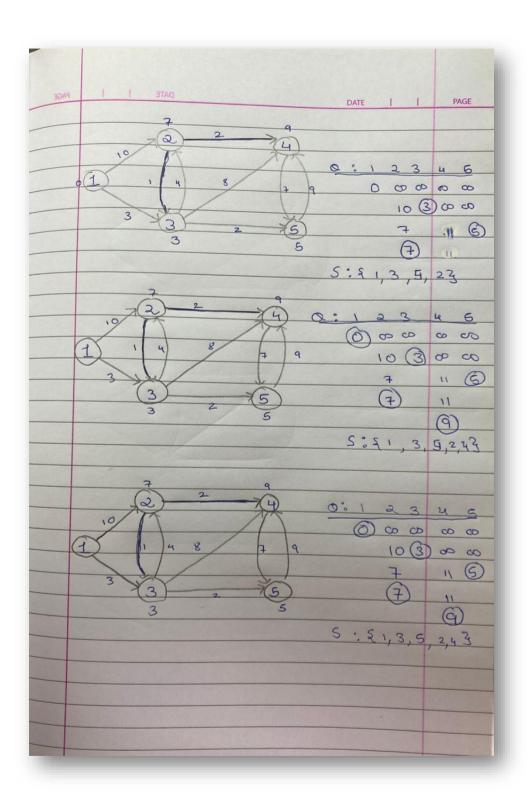
dist[w]: dist[u] + cost[u, w];
}
```

Algorithm writing

- 1. Mark your selected initial node with a current distance of 0 and the rest with infinity.
- 2. Set the non-visited node with the smallest current distance as the current node C.
- 3. For each neighbour N of your current node C:
 - add the current distance of C with the weight of the edge connecting C-N.
 - If it's smaller than the current distance of N, set it as the new current distance of N.
- 4. Mark the current node C as visited.
- 5. If there are non-visited nodes, go to step 2.

Tracing with Example





Program

```
#include<iostream>
using namespace std;
#define V 6
int stepcount=0;
int minDistance(int dist[], bool sptSet[])
{
      int min = INT_MAX;stepcount++;
      int min index;
      for (int v = 0; v < V; v++)
  {
    stepcount++;
    stepcount++;
    if (sptSet[v] == false && dist[v] <= min)</pre>
    {
      min = dist[v], min_index = v;stepcount++;
    }
  }
```

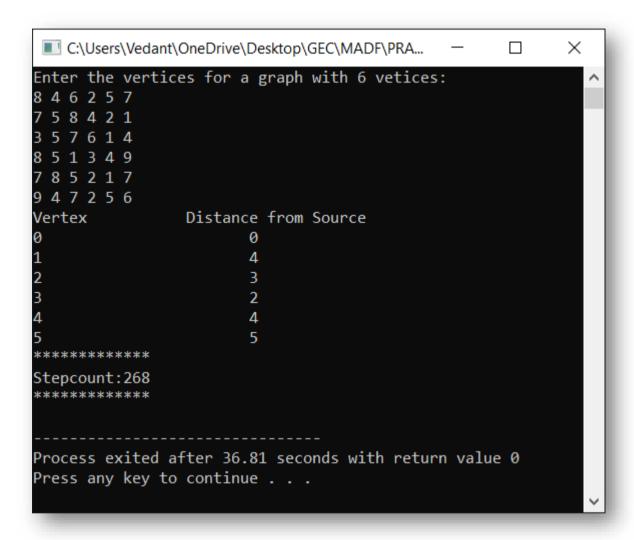
```
stepcount++;
      return min_index;
}
void printSolution(int dist[])
{
  cout<<"Vertex \t\t Distance from Source"<<endl;stepcount++;</pre>
      for (int i = 0; i < V; i++)
  {
      stepcount++;
      cout << i < "\t\t" << dist[i] << endl; step count ++;
  }
}
void dijkstra(int graph[V][V], int src)
{
      int dist[V];
      bool sptSet[V];
```

```
for (int i = 0; i < V; i++)
        {
             stepcount++;
             dist[i] = INT_MAX, sptSet[i] = false;stepcount++;
        }
             dist[src] = 0;stepcount++;
             for (int count = 0; count < V - 1; count++)
             {
           stepcount++;
                   int u = minDistance(dist, sptSet);stepcount++;
                   sptSet[u] = true;stepcount++;
                   for (int v = 0; v < V; v++)
           {
             stepcount++;
             stepcount++;
                          if (!sptSet[v] && graph[u][v] && dist[u] !=
INT_MAX
```

```
&& dist[u] + graph[u][v] < dist[v])
                                 {
                                   dist[v] = dist[u] +
graph[u][v];stepcount++;
                                 }
           }
             }
         stepcount++;
             printSolution(dist);
      }
      int main()
      {
         int graph[V][V];
             cout<<"Enter the vertices for a graph with 6
vetices:";stepcount++;
        for (int i=0;i<V;i++)
         {
           stepcount++;
           for(int j=0;j<V;j++)
```

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Output



Conclusion

- Detailed concept of Single Source Shortest Path Algo (Dijkstra's Algo)was studied successfully.
- > Program using Dijkstra's Algorithm was executed successfully.
- ➤ The step count for the Dijkstra's Algorithm was obtained.