## Tutorial 9: Transaction Processing Concepts and Theory CS3402 Database Systems

## **Question 1**

- Which of the following schedules is (conflict) serializable? For each serializable schedule, determine the equivalent serial schedules.
  - a) A:  $r_1(X)$ ;  $r_3(X)$ ;  $w_1(X)$ ;  $r_2(X)$ ;  $w_3(X)$ ;
  - b) B:  $r_1(X)$ ;  $r_3(X)$ ;  $w_3(X)$ ;  $w_1(X)$ ;  $r_2(X)$ ;
  - c) C:  $r_3(X)$ ;  $r_2(X)$ ;  $w_3(X)$ ;  $r_1(X)$ ;  $w_1(X)$ ;

## Question 2

• Consider the following concurrent schedule S. Draw the serialization graph for the schedule. Is it conflict serializable?

T <sub>1</sub>	T <sub>2</sub>	<b>T</b> <sub>3</sub>
	Read(X)	
Write(Y)		
		Read(Y)
	Write(Y)	
Write(X)		
	Commit	
		Write(Z)
Commit		
		Commit

## **Question 3**

- Consider schedules S<sub>1</sub>, S<sub>2</sub> and S<sub>3</sub> below. Determine whether each schedule is strict, cascadeless, recoverable, or nonrecoverable. Determine the strictest recoverability condition that each schedule satisfies.
  - a)  $S_1$ :  $r_1(X)$ ;  $w_1(X)$ ;  $r_2(X)$ ;  $r_1(Y)$ ;  $w_2(X)$ ;  $c_2$ ;  $c_1$ ;
  - b)  $S_2$ :  $r_1(X)$ ;  $w_1(X)$ ;  $r_2(X)$ ;  $r_1(Y)$ ;  $w_2(X)$ ;  $w_1(Y)$ ;  $c_1$ ;  $c_2$ ;
  - c)  $S_3$ :  $r_1(X)$ ;  $w_1(X)$ ;  $w_2(X)$ ;  $w_1(Y)$ ;  $c_1$ ;  $r_2(X)$ ;  $c_2$ ;

Can you change schedule S<sub>3</sub> into a strict schedule?