# EE3220 System-on-Chip Design

Lecture Note 3

Embedded System Software Design Basics



#### Overview

- Concurrency
  - How do we make things happen at the right times?
- Software Engineering for Embedded Systems
  - How do we develop working code quickly?

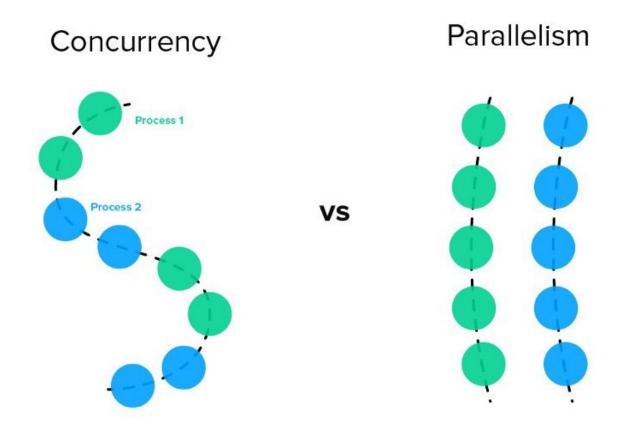


# Concurrency



### Concurrency vs. Parallelism

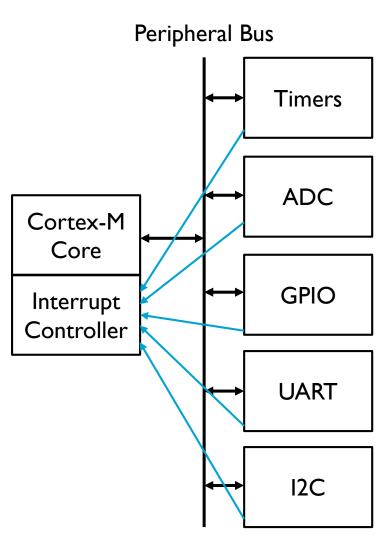
Concurrency is about handling a lot of things (input, output, events) at once.





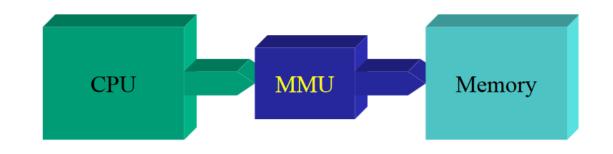
# MCU Hardware & Software for Concurrency

- CPU executes instructions from one or more thread of execution
- Specialized hardware peripherals add dedicated concurrent processing
  - Watchdog timer
  - Analog interfacing
  - Timers
  - Communications with other devices
  - Detecting external signal events
  - LCD driver
- Peripherals use interrupts to notify CPU of events

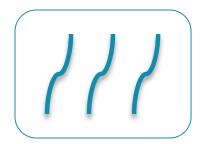




### Process, Task and Thread



- A program is an executable file
- A process is an instance of the program in execution
  - You need memory management unit to enable process independence
- A program is broken into a number of functions running together, called tasks
  - Enable "multi-tasking",
  - Each task is given a priority at creation
  - This priority is used to determine which task runs, a higher one will pre-empt a lower one (a forced context-switching)
- A thread represents a single sequential flow of control
  - A process is usually referred to be "heavy-weight", and expensive in terms of computing resources
  - A thread is usually referred to be "light-weight"
  - A process becomes an overall context in which the threads run

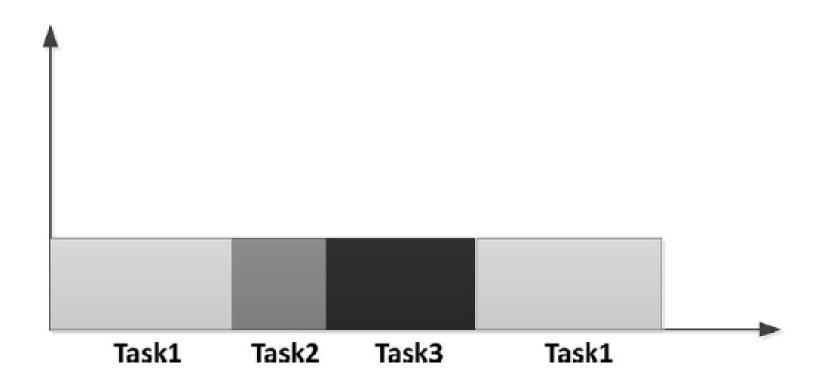


One process Multiple threads



# Task Example

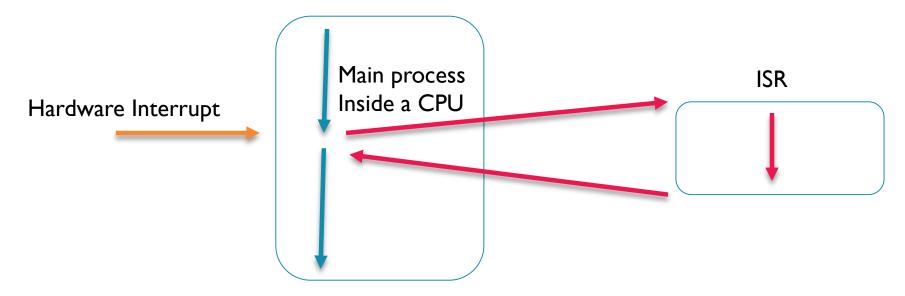
Task I takes the longest time. When it releases the CPU, Task2 starts. Task3 starts after
 Task2 completes. CPU is given back to Task1 after Task3 releases it.





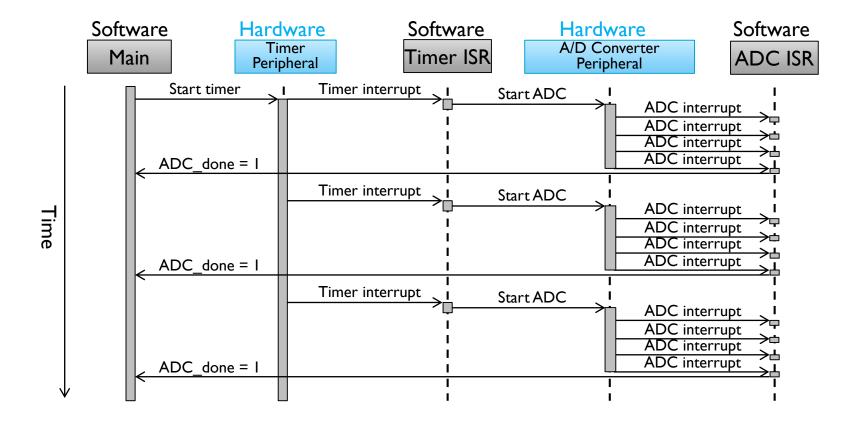
# Interrupt Service Routine (ISR)

- An ISR, also called an interrupt handler, is a software process invoked by an interrupt request from a hardware device.
- This software process handles the request and sends it to the CPU processor, interrupting the active main process. When the ISR is complete, the main process is resumed.





# Concurrent Hardware & Software Operation



 Embedded systems rely on both MCU hardware peripherals and software to get everything done on time



# Pre-emptive vs. Non-preemptive Scheduling

#### Preemptive

- A processor can be preempted to execute another process in the middle of ongoing execution
- CPU utilization is more efficient compared to Non-Preemptive Scheduling
- Preemptive scheduling is prioritized, according to the highest priority process
- Preemptive algorithm has an overhead of switching process from different states

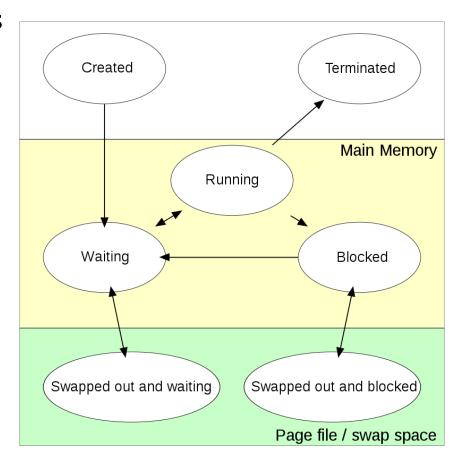
#### Non-Preemptive

- Once the processor has started the execution, it must complete the execution before running the others
- CPU utilization is less efficient compared to preemptive Scheduling
- When a process is being run, a state will be given to this process, and will not be deleted until it finishes the job
- Non-preemptive Scheduling has no overhead of switching the process from different states



#### **Process States**

- Create State When a process is first created, it occupies the "created" or "new" state.
- Ready / Waiting / Idle State A "ready" or "waiting" process has been loaded into main memory and is awaiting execution on a CPU
- Running State A process moves into the running state when it is chosen for execution.
- Terminated State A process may be terminated, either from the "running" state by completing its execution or by explicitly being killed



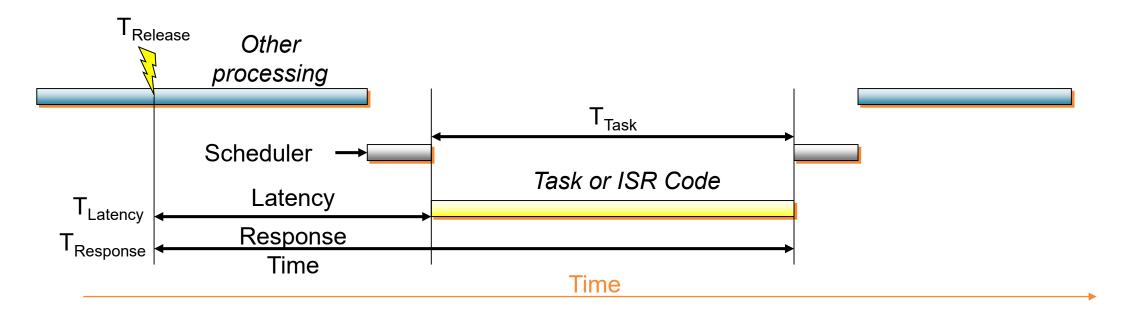


# **CPU** Scheduling

- MCU's Interrupt system provides a basic scheduling approach for CPU
  - "Run this subroutine every time this hardware event occurs"
  - Is adequate for simple systems
- More complex systems need to support multiple concurrent independent threads of execution
  - Use task scheduler to share CPU
  - Different approaches to task scheduling
- How do we make the processor responsive? (How do we make it do the right things at the right times?)
  - If we have more software threads than hardware threads, we need to share the processor.



#### **Definitions**



- T<sub>Release</sub>(i) = Time at which task (or interrupt) i requests service/is released/is ready to run
- T<sub>Latency</sub> (i) = Delay between release and start of service for task i
- T<sub>Response</sub>(i) = Delay between request for service and completion of service for task i
- $T_{Task}(i)$  = Time needed to perform computations for task i
- $T_{ISR}(i) = Time needed to perform interrupt service routine routine routine interrupt service routine rout$



### Scheduling Approaches

- Rely on MCU's hardware interrupt system to run right code
  - Event-triggered scheduling with interrupts
  - Works well for many simple systems
- Use software to schedule CPU's time
  - Static cyclic executive
  - Dynamic priority
    - Without task-level preemption
    - With task-level preemption



# Operating System Scheduling algorithms

- A Process Scheduler schedules different processes to be assigned to the CPU / processor based on particular scheduling algorithms. There are common and popular process scheduling algorithms.
- These algorithms are either non-preemptive or preemptive.
  - First-Come, First-Served (FCFS) Scheduling
  - Shortest-Job-Next (SJN) Scheduling
  - Priority Scheduling
  - Shortest Remaining Time
  - Round Robin(RR) Scheduling
  - Multiple-Level Queues Scheduling
- Cover in more details in the tutorials.



# Event-Triggered Scheduling using Interrupts

- Basic architecture, useful for simple lowpower devices
  - Very little code or time overhead
- Leverages built-in task dispatching of interrupt system
  - Can trigger ISRs with input changes, timer expiration, UART data reception, analog input level crossing comparator threshold
- Function types
  - Main function configures system and then goes to sleep
    - If interrupted, it goes right back to sleep
  - Only interrupts are used for normal program operation

- Example: bike computer
  - Intl: wheel rotation
  - Int2: mode key
  - Int3: clock
  - Output: Liquid Crystal Display



### Bike Computer Functions – Four Handlers

Reset

ISR 1: Wheel rotation

ISR 2: Mode Key ISR 3: Time of Day Timer

```
Configure timer,
inputs and
outputs

cur_time = 0;
rotations = 0;
tenth_miles = 0;

while (1) {
   sleep;
}
```

```
rotations++;
if(rotations>
    R_PER_MILE/10) {
    tenth_miles++;
    rotations = 0;
}
speed =
    circumference /
    (cur_time - prev_time);
compute avg_speed;
prev_time = cur_time;
return from interrupt
```

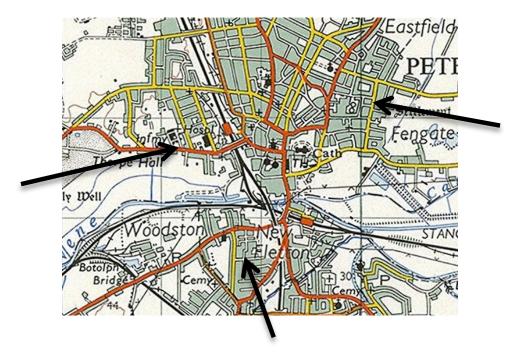
```
mode++;
mode = mode %
   NUM_MODES;
return from interrupt;
```

```
cur_time ++;
lcd refresh--;
if (lcd refresh==0) {
 convert tenth miles
    and display
 convert speed
    and display
  if (mode == 0)
    convert cur time
      and display
  else
    convert avg speed
      and display
  lcd refresh =
     LCD REF PERIOD
```



### A More Complex Application



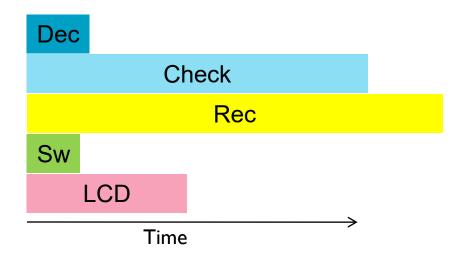


- GPS-based Pothole Alarm and Moving Map
  - Sounds alarm when approaching a pothole
  - Display's vehicle position on LCD
  - Also logs driver's position information
  - Hardware: GPS, user switches, speaker, LCD, flash memory



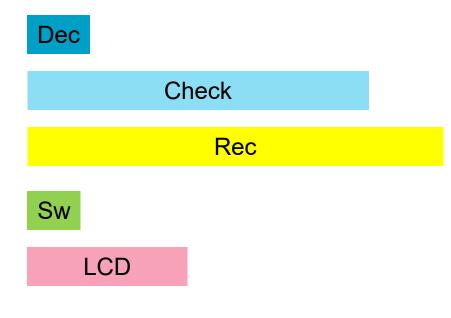
### Application Software Tasks

- Dec: Decode GPS sentence to find current vehicle position.
- Check: Check to see if approaching any pothole locations. Takes longer as the number of potholes in database increases.
- Rec: Record position to flash memory. Takes a long time if erasing a block.
- Sw: Read user input switches. Run 10 times per second
- LCD: Update LCD with map. Run 4 times per second





#### How do we schedule these tasks?



- Task scheduling: Deciding which task should be running now
- Two fundamental questions:
  - Do we run tasks in the same order every time?
    - Yes: Static schedule (cyclic executive, round-robin)
    - No: Dynamic, prioritized schedule
  - Can one task preempt another, or must it wait for completion?
    - Yes: Preemptive
    - No: Non-preemptive (cooperative, run-to-completion)



# Static Schedule (Cyclic Executive)

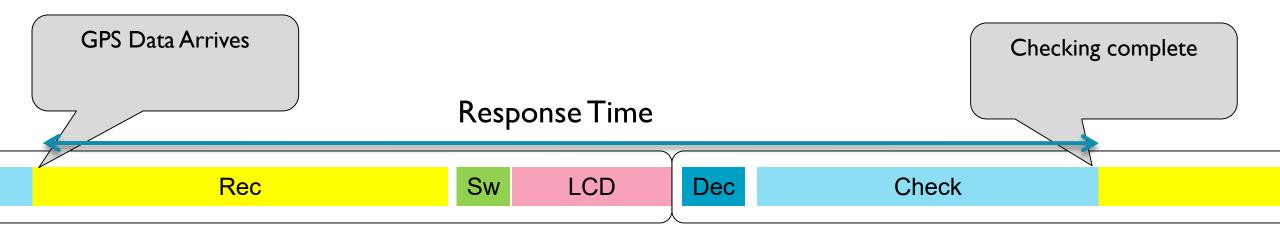
DecCheckRecSwLCDDec

- Pros
  - Very simple
- Cons
  - Always run the same schedule, regardless of changing conditions and relative importance of tasks.
  - All tasks run at same rate. Changing rates requires adding extra calls to the function.
  - Maximum delay is sum of all task run times. Polling/execution rate is I/maximum delay.

```
while (1){
    Dec();
    Check();
    Rec();
    Sw();
    LCD();
}
```



# Static Schedule Example



- What if we receive GPS position right after Rec starts running?
- Delays
  - Have to wait for Rec, Sw, LCD before we start decoding position with Dec.
  - Have to wait for Rec, Sw, LCD, Dec, Check before we know if we are approaching a pothole!

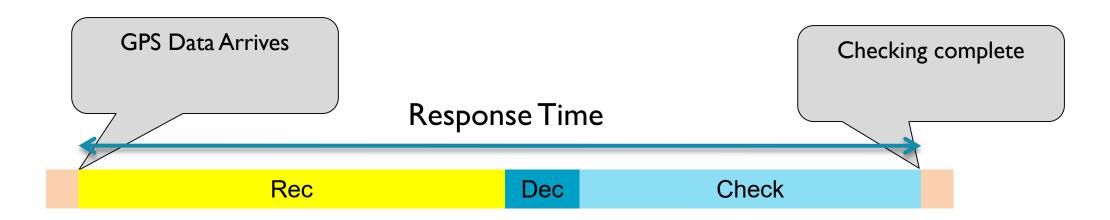


# Dynamic Scheduling

- Allow schedule to be computed on-the-fly
  - Based on importance or something else
  - Simplifies creating multi-rate systems
- Schedule based on importance
  - Prioritization means that less important tasks don't delay more important ones
- How often do we decide what to run?
  - Coarse grain After a task finishes. Called Run-to-Completion (RTC) or non-preemptive
  - Fine grain Any time. Called Preemptive, since one task can preempt another.



# Dynamic Run-To-Completion (RTC) Schedule

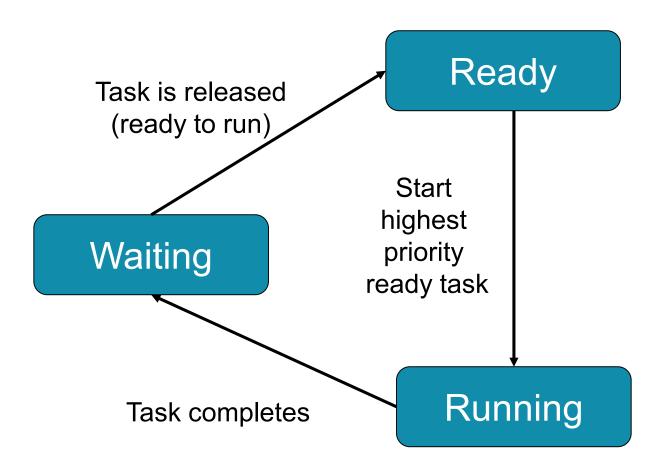


- What if we receive GPS position right after Rec starts running?
- Delays
  - Have to wait for Rec to finish before we start decoding position with Dec.
  - Have to wait for Rec, Dec, Check before we know if we are approaching a pothole



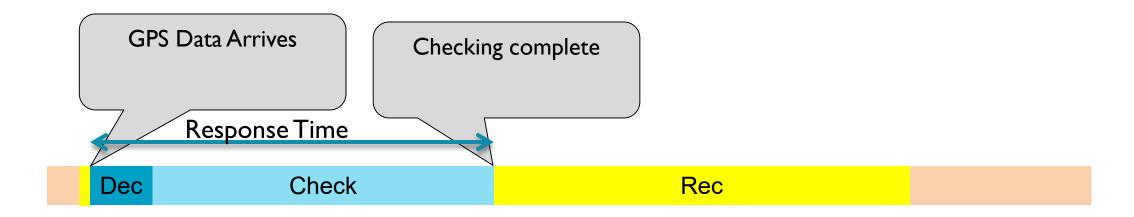
# Task State and Scheduling Rules

- Scheduler chooses among Ready tasks for execution based on priority
- Scheduling Rules
  - If no task is running, scheduler starts the highest priority ready task
  - Once started, a task runs until it completes
  - Tasks then enter waiting state until triggered or released again





# Dynamic Preemptive Schedule

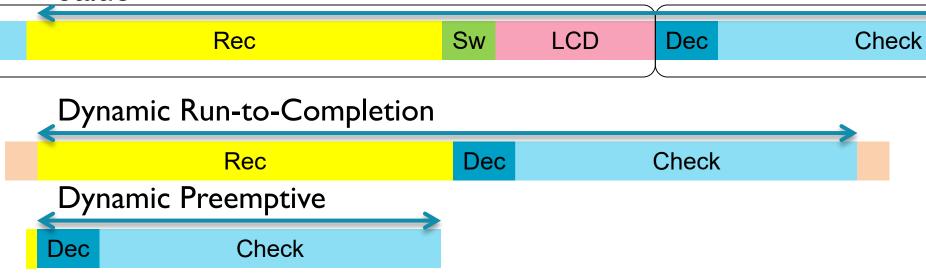


- What if we receive GPS position right after Rec starts running?
- Delays
  - Scheduler switches out Rec so we can start decoding position with Dec immediately
  - Have to wait for Dec, Check to complete before we know if we are approaching a pothole



### Comparison of Response Times

Static



- Pros
  - Preemption offers best response time
    - Can do more processing (support more potholes, or higher vehicle speed)
    - Or can lower processor speed, saving money, power
- Cons
  - Requires more complicated programming, more memory
  - Introduces vulnerability to data race conditions



#### Common Schedulers

- I) Cyclic executive non-preemptive and static
  - Static the scheduling is done at compile time
  - Only one task, like an infinite loop in main()
- 2) Run-to-completion non-preemptive and dynamic
  - Dynamic the scheduling is done at run-time, usually done by the operating system
  - typically have an event queue
  - in strict order of admission by an event loop
- 3) Preemptive and dynamic



# 1) Cyclic Executive with Interrupts

- Two priority levels
  - main code foreground
  - Interrupts background
- Example of a foreground / background system
- Main user code runs in foreground
- Interrupt routines run in background (high priority)
  - Run when triggered
  - Handle most urgent work
  - Set flags to request processing by main loop

```
BOOL DeviceARequest, DeviceBRequest,
DeviceCRequest;
void interrupt HandleDeviceA() {
  /* do A's urgent work */
  DeviceARequest = TRUE;
void main(void) {
  while (TRUE) {
    if (DeviceARequest) {
      FinishDeviceA();
    if (DeviceBRequest) {
      FinishDeviceB();
    if (DeviceCRequest) {
      FinishDeviceC();
```

# 2) Run-To-Completion (Non-Preemptive) Scheduler

- Use a scheduler function to run task functions at the right rates
  - Table stores information per task
    - Period: How many ticks between each task release
    - Release Time: how long until task is ready to run
    - ReadyToRun: task is ready to run immediately
  - Scheduler runs forever, examining schedule table which indicates tasks which are ready to run (have been "released")
  - A periodic timer interrupt triggers an ISR, which updates the schedule table
    - Decrements "time until next release"
    - If this time reaches 0, set that task's Run flag and reload its time with the period
- Follows a "run-to-completion" model
  - A task's execution is not interleaved with any other task
  - Only ISRs can interrupt a task
  - After ISR completes, the previously-running task resumes
- Priority is typically static, so can use a table with highest priority tasks first for a fast, simple scheduler implementation.

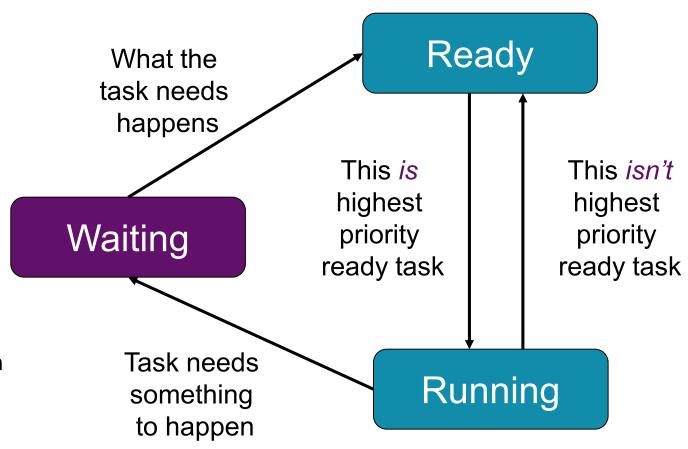
# 3) Preemptive Scheduler

- Task functions need not run to completion, but can be interleaved with each other
  - Simplifies writing software
  - Improves response time
  - Introduces new potential problems
- Worst case response time for highest priority task does not depend on other tasks, only ISRs and scheduler
  - Lower priority tasks depend only on higher priority tasks



# Task State and Scheduling Rules

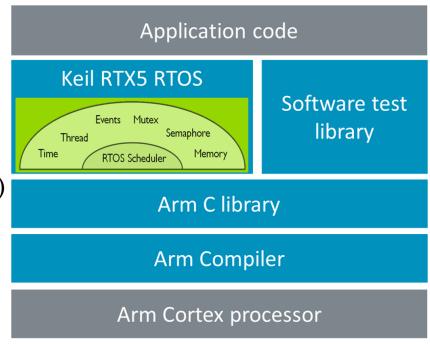
- Scheduler chooses among Ready tasks for execution based on priority
- Scheduling Rules
  - A task's activities may lead it to waiting (blocked)
  - A waiting task never gets the CPU. It must be signaled by an ISR or another task.
  - Only the scheduler moves tasks between ready and running





#### What's an RTOS?

- What does Real-Time mean?
  - Can calculate and guarantee the maximum response time for each task and interrupt service routine
  - This "bounding" of response times allows use in hard-real-time systems (which have deadlines which must be met)
- What's in the RTOS
  - Task Scheduler
    - Preemptive, prioritized to minimize response times
    - Interrupt support
  - Core Integrated RTOS services
    - Inter-process communication and synchronization (safe data sharing)
    - Time management
  - Optional Integrated RTOS services
    - I/O abstractions?
    - memory management?
    - file system?
    - networking support?
    - GUI??



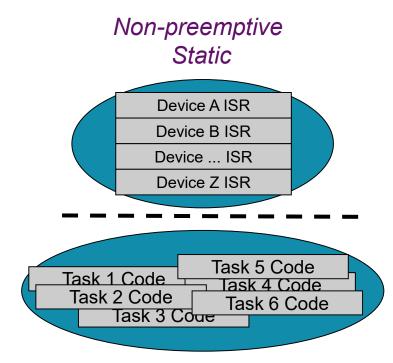


### Definition of Embedded System

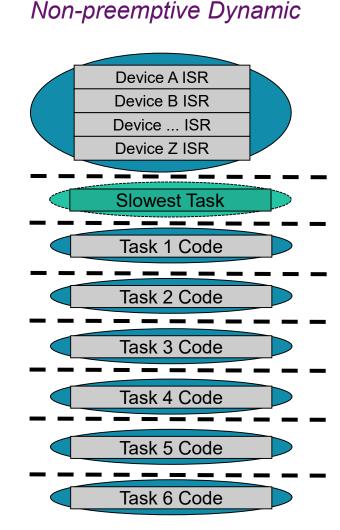
- An embedded system is a special-purpose computer system designed to perform one or a few dedicated functions
- sometimes with real-time computing constraints
- It is usually embedded as part of a complete device including hardware and mechanical parts
  - As shown in Wikipedia
- In real-time systems, the correct behaviour of a system depends, not only on the values of results that are produced, but also on the time at which they are produced
  - John Stankovic. Misconceptions about real-time computing IEEE Computer,
     October 1988



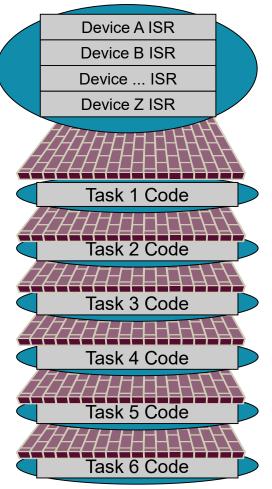
### Comparison of Timing Dependence



 Code can be delayed by everything at same level (in oval) or above



#### Preemptive Dynamic





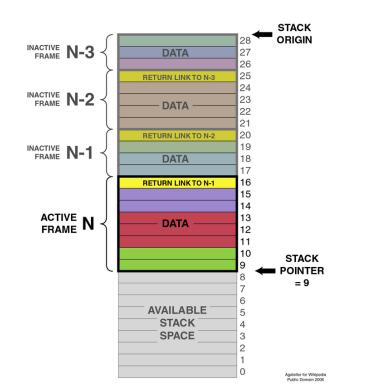
# Recall: Memory Structure

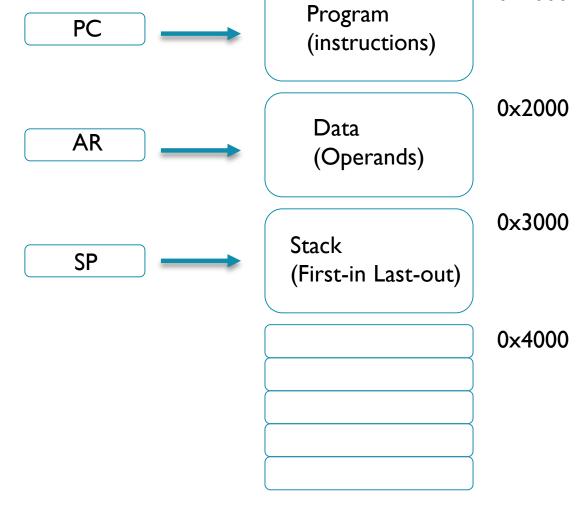
Memory with instruction, data, stack segments

PC = Program Counter

AR = Address Register

SP = Stack Pointer

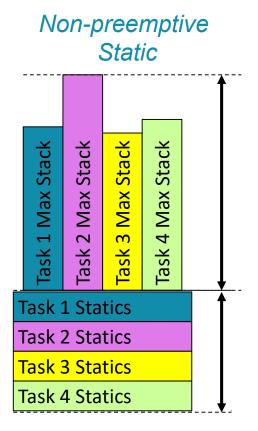


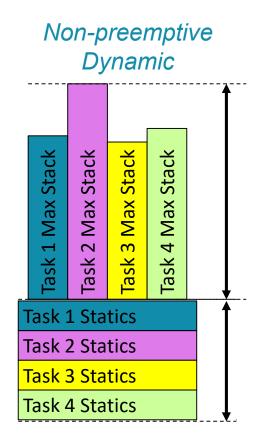




 $0 \times 1000$ 

## Comparison of RAM Requirements





- Preemption requires space for each stack generally
- Need space for all static variables (including globals)

Max  $\mathcal{C}$ **Preemptive Dynamic** Stack Max 7 Stack Max Task Task 1 Statics Task 2 Statics Task 3 Statics Task 4 Statics



# Software Engineering For Embedded Systems



# Good Enough Software, Soon Enough

- How do we make software correct enough without going bankrupt?
  - Need to be able to develop (and test) software efficiently
- Follow a good plan
  - Start with customer requirements
  - Design architectures to define the building blocks of the systems (tasks, modules, etc.)
  - Add missing requirements
    - Fault detection, management and logging
    - Real-time issues
    - Compliance to a firmware standards manual
    - Fail-safes
  - Create detailed design
  - Implement the code, following a good development process
    - Perform frequent design and code reviews
    - Perform frequent testing (unit and system testing, preferably automated)
    - Use revision control to manage changes
  - Perform post-mortems to improve development process



#### What happens when the plan meets reality?

- We want a robust plan which considers likely risks
  - What if the code turns out to be a lot more complex than we expected?
  - What if there is a bug in our code (or a library)?
  - What if the system doesn't have enough memory or throughput?
  - What if the system is too expensive?
  - What if the lead developer quits?
  - What if the lead developer is incompetent, lazy, or both (and won't quit!)?
  - What if the rest of the team gets sick?
  - What if the customer adds new requirements?
  - What if the customer wants the product two months early?
- Successful software engineering depends on balancing many factors, many of which are non-technical!



#### Risk Reduction

- Plan to the work to accommodate risks
- Identify likely risks up front
  - Historical problem areas
  - New implementation technologies
  - New product features
  - New product line
- Severity of risk is combination of likelihood and impact of failure



#### Software Lifecycle Concepts

- Coding is the most visible part of a software development process but is not the only one
- Before we can code, we must know
  - What must the code do? Requirements specification
  - How will the code be structured? Design specification
    - (only at this point can we start writing code)
- How will we know if the code works? Test plan / Test cases
  - Best performed when defining requirements
- The software will likely be enhanced over time Extensive downstream modification and maintenance!
  - Corrections, adaptations, enhancements & preventive maintenance



#### Requirements

- Ganssle's Reason #5 for why embedded projects fail: Vague Requirements
  - http://www.ganssle.com/articles/jackstoptenlist.htm
- Types of requirements
  - Functional what the system needs to do
  - Nonfunctional emergent system behaviors such as response time, reliability, energy efficiency, safety, etc.
  - Constraints limit design choices
- Representations
  - Text Liable to be incomplete, bloated, ambiguous, even contradictory
  - Diagrams (state charts, flow charts, message sequence charts)
  - Concise
  - Can often be used as design documents
- Traceability
  - Each requirement should be verifiable with a test
- Stability
  - Requirements churn leads to inefficiency and often "recency" problem (most recent requirement change is assumed to be most important)



# Design Before Coding

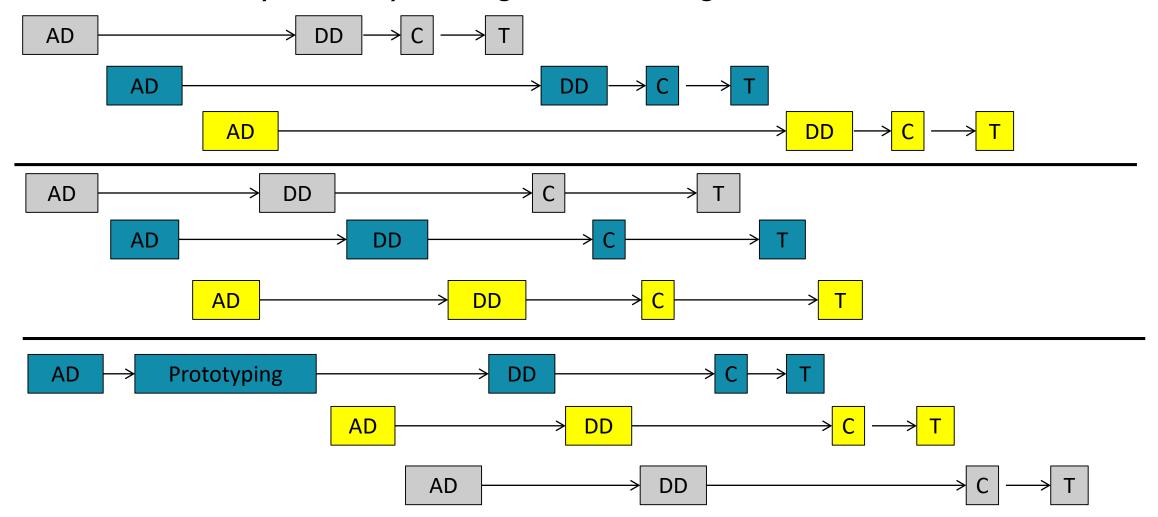


- Ganssle's reason #9: Starting coding too soon
- Underestimating the complexity of the needed software is a very common risk
- Writing code locks you in to specific implementations
  - Starting too early may paint you into a corner
- Benefits of designing system before coding
  - Get early insight into system's complexity, allowing more accurate effort estimation and scheduling
  - Can use design diagrams rather than code to discuss what system should do and how. Ganssle's reason #7: Bad Science
  - Can use design diagrams in documentation to simplify code maintenance and reduce risks of staff turnover



#### Design Before Coding

How much of the system do you design before coding?



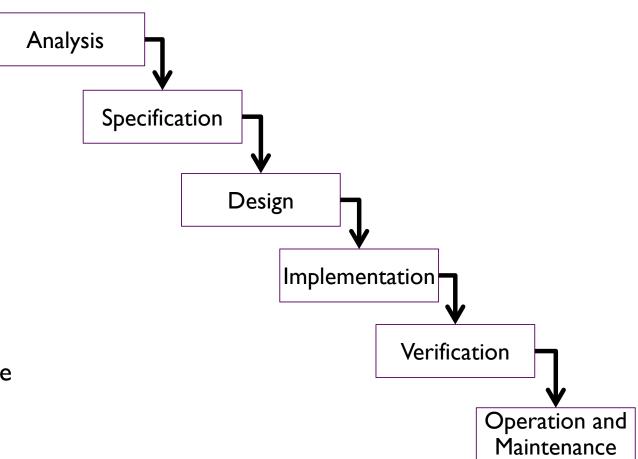
#### Development Models

- How do we schedule these pieces?
- Consider amount of development risk
  - New MCU?
  - Exceptional requirements (throughput, power, safety certification, etc.)
  - New product?
  - New customer?
  - Changing requirements?
- Choose model based on risk
  - Low: Can create detailed plan. Big-up-front design, waterfall
  - High: Use iterative or agile development method, spiral. Prototype high-risk parts first



#### Waterfall (Idealized)

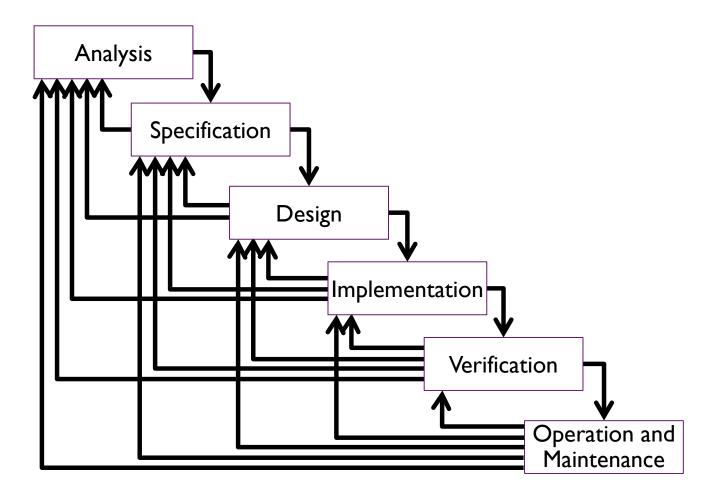
- Plan the work, and then work the plan
- BUFD: Big Up-Front Design
- Model implies that we and the customers know
  - All of the requirements up front
  - All of the interactions between components, etc.
  - How long it will take to write the software and debug it





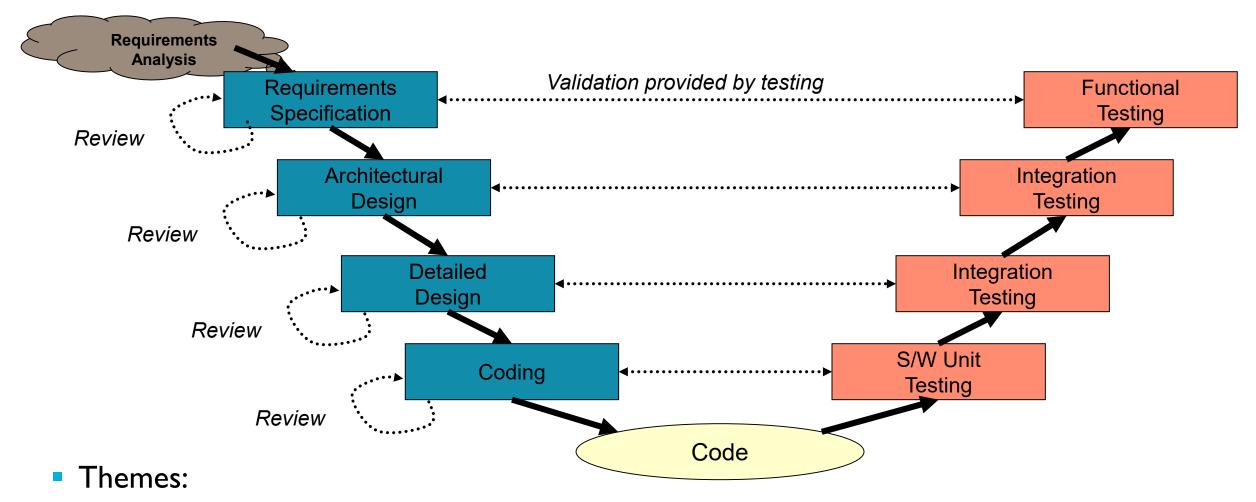
## Waterfall (As Implemented)

 Reality: We are not omniscient, so there is plenty of backtracking





#### V Model Overview



- Link front and back ends of life-cycle for efficiency
- Provide "traceability" to ensure nothing falls through the cracks



#### 1. Requirements Specification and Validation Plan

- Result of Requirements Analysis
- Should contain:
  - Introduction with goals and objectives of system
  - Description of problem to solve
  - Functional description
    - provides a "processing narrative" per function
    - lists and justifies design constraints
    - explains performance requirements
  - Behavioral description shows how system reacts to internal or external events and situations
    - State-based behavior
    - General control flow
    - General data flow
  - Validation criteria
    - tell us how we can decide that a system is acceptable. (Are we done yet?)
    - is the foundation for a validation test plan
  - Bibliography and Appendix refer to all documents related to project and provide supplementary information



# 2. Architectural (High-Level) Design

- Architecture defines the structure of the system
  - Components
  - Externally visible properties of components
  - Relationships among components
- Architecture is a representation which lets the designer...
  - Analyze the design's effectiveness in meeting requirements
  - Consider alternative architectures early
  - Reduce down-stream implementation risks
- Architecture matters because...
  - It's small and simple enough to fit into a single person's brain (as opposed to comprehending the entire program's source code)
  - It gives stakeholders a way to describe and therefore discuss the system



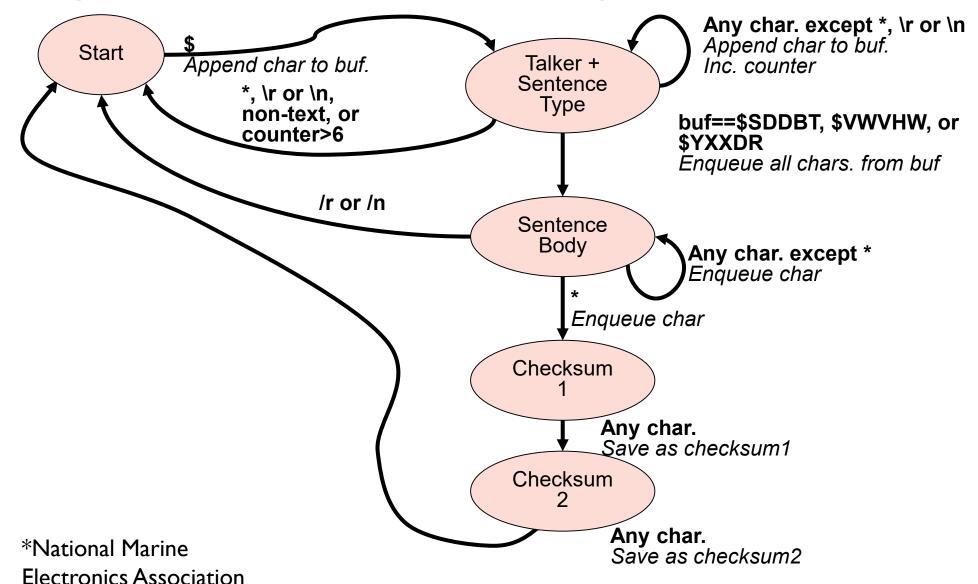
## 3. Detailed Design



- Describe aspects of how system behaves
  - Flow charts for control or data
  - State machine diagram
  - Event sequences
- Graphical representations very helpful
  - Can provide clear, single-page visualization of what system component should do
- Unified Modeling Language (UML)
  - Provides many types of diagrams
  - Some are useful for embedded system design to describe structure or behavior

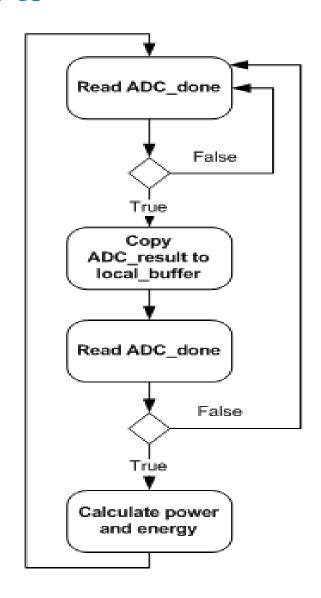


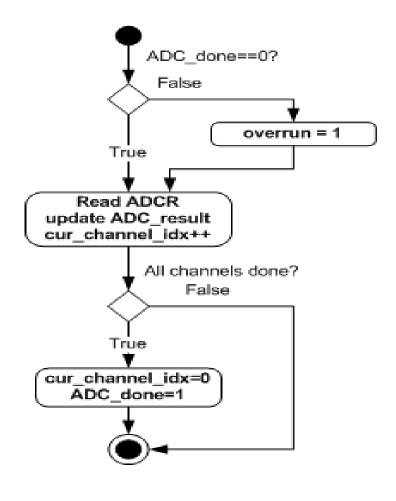
# Example: State Machine for Parsing NMEA-0183\*





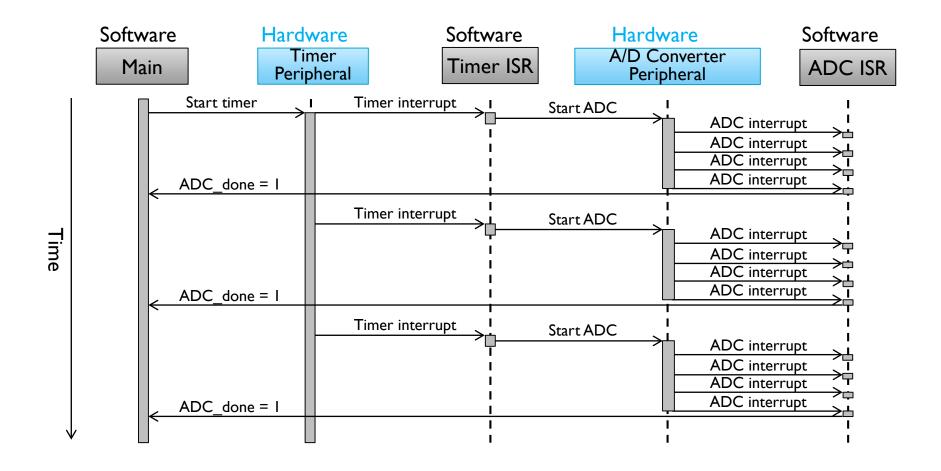
#### **Flowcharts**







## Sequence of Interactions between Components





## 4. Coding and Code Inspections

- Coding driven directly by Detailed Design Specification
- Use a version control system while developing the code
- Follow a coding standard
  - Eliminate stylistic variations which make understanding code more difficult
  - Avoid known questionable practices
  - Spell out best practices to make them easier to follow
- Perform code reviews
- Test effectively
  - Automation
  - Regression testing



#### Peer Code Review

- Inspect the code before testing it
- Extensive positive industry results from code inspections
  - IBM removed 82% of bugs
  - 9 hours saved by finding each defect
  - For AT&T quality rose by 1000% and productivity by 14%



- 80% of the errors detected by HP's inspections were unlikely to be caught by testing
- HP, Shell Research, Bell Northern, AT&T: inspections 20-30x more efficient than testing





## 5. Software Testing

- Testing IS NOT "the process of verifying the program works correctly"
  - The program probably won't work correctly in all possible cases
    - Professional programmers have I-3 bugs per 100 lines of code after it is "done"
  - Testers shouldn't try to prove the program works correctly (impossible)
    - If you want and expect your program to work, you'll unconsciously miss failure because human beings are inherently biased
- The purpose of testing is to find problems quickly
  - Does the software violate the specifications?
  - Does the software violate unstated requirements?
- The purpose of finding problems is to fix the ones which matter
  - Fix the most important problems, as there isn't enough time to fix all of them
  - The Pareto Principle defines "the vital few, the trivial many"
    - Bugs are uneven in frequency a vital few contribute the majority of the program failures. Fix these first.



## Approaches to Testing

#### Incremental Testing

- Code a function and then test it (module/unit/element testing)
- Then test a few working functions together (integration testing)
  - Continue enlarging the scope of tests as you write new functions
- Incremental testing requires extra code for the test harness
  - A driver function calls the function to be tested
  - A stub function might be needed to simulate a function called by the function under test, and which returns or modifies data.
  - The test harness can automate the testing of individual functions to detect later bugs

#### Big Bang Testing

- Code up all of the functions to create the system
- Test the complete system
  - Plug and pray



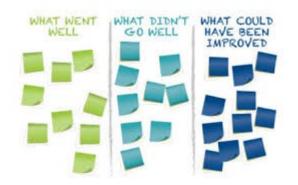
# Why Test Incrementally?



- Finding out what failed is much easier
  - With Big Bang, since no function has been thoroughly tested, most probably have bugs
  - Now the question is "Which bug in which module causes the failure I see?"
  - Errors in one module can make it difficult to test another module
    - Errors in fundamental modules (e.g. kernel) can appear as bugs in other many other dependent modules
- Less finger pointing = happier SW development team
  - It's clear who made the mistake, and it's clear who needs to fix it
- Better automation
  - Drivers and stubs initially require time to develop, but save time for future testing



# 6. Perform Project Retrospectives



- Goals improve your engineering processes
  - Extract all useful information learned from the just-completed project provide "virtual experience" to others
  - Provide positive non-confrontational feedback
  - Document problems and solutions clearly and concisely for future use
- Basic rule: problems need solutions
- Often small changes improve performance, but are easy to forget



## Summary - Embedded Software Design Basics

#### Concurrency

- How do we make things happen at the right times?
- Process, Task, Thread
- Preemptive vs. Non-preemptive
- Scheduling methods
- Memory structure

#### Software Engineering for Embedded Systems

- How do we develop working code quickly?
- Requirements
- Development model Waterfall
- Code review
- Sufficient testing

