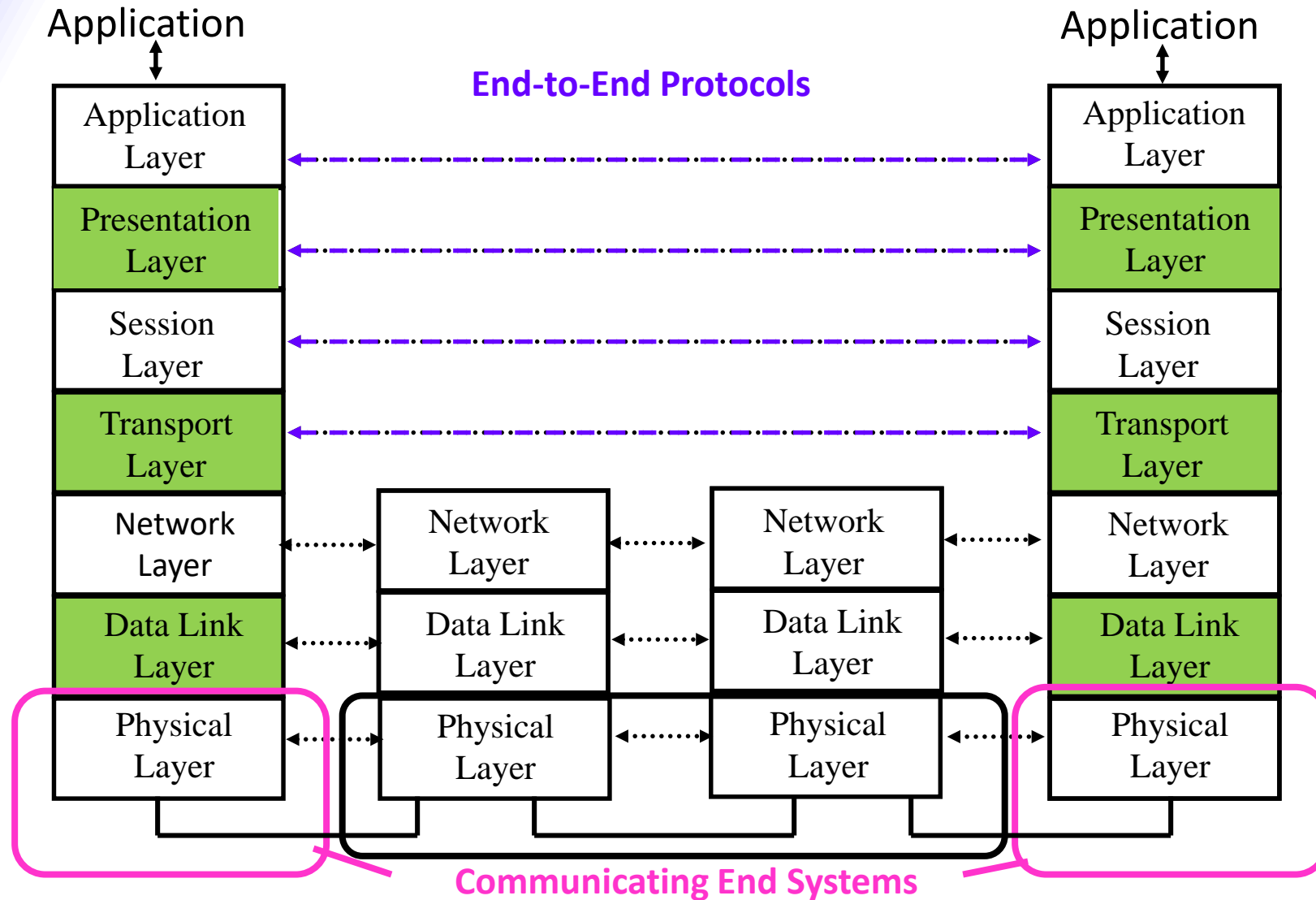


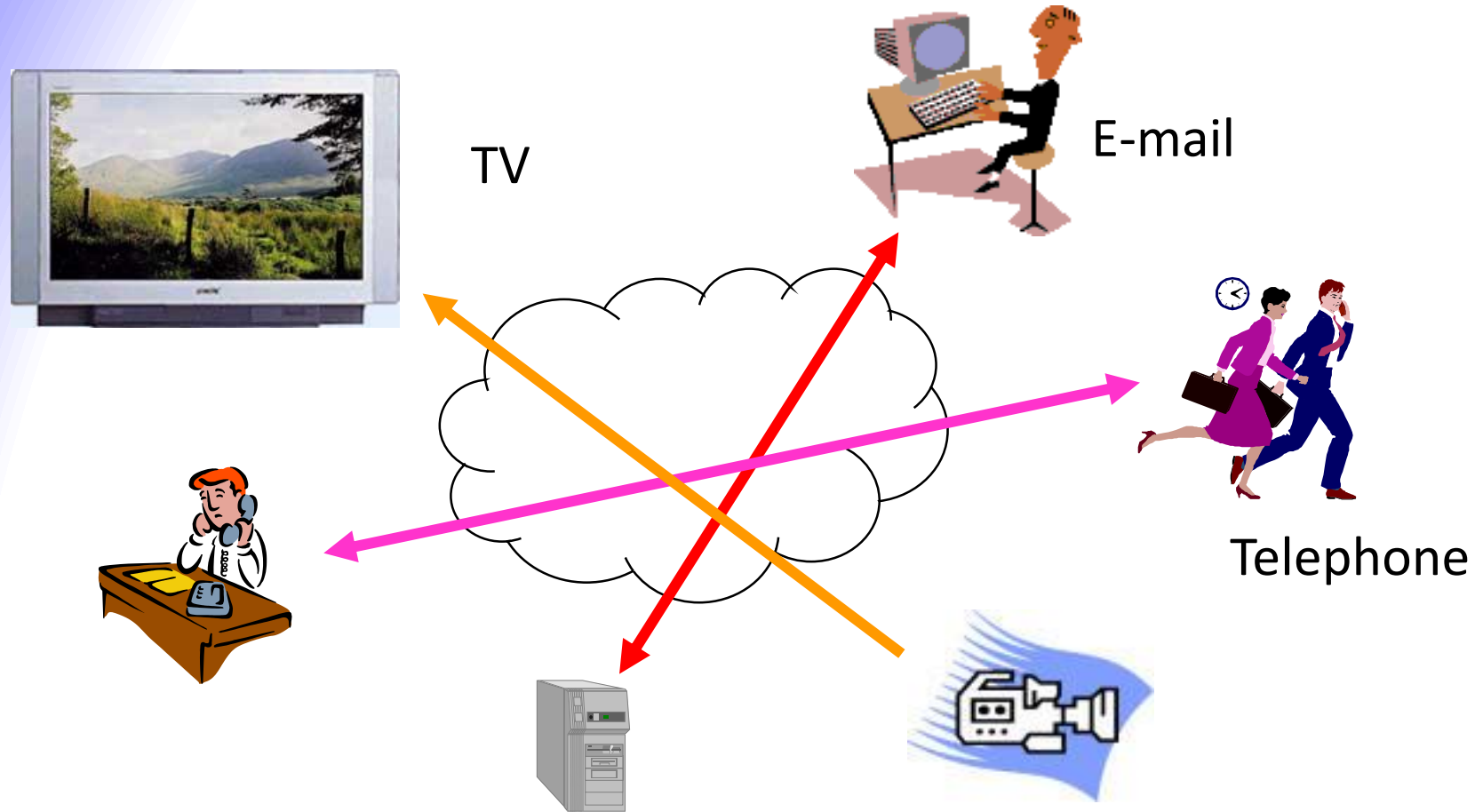
## 6. Data transmission

- \* digital representation
- \* digital and analog transmission
- \* asynchronous and synchronous communications
- \* transmission media

# 7-Layer OSI Reference Model



Digital transmission enables networks to support many services



## 6.1 Digital representation

- Bit: number with value 0 or 1
  - $n$  bits: digital representation for  $0, 1, \dots, 2^n$
  - byte or octet,  $n = 8$
  - computer word,  $n = 16, 32$ , or  $64$
- $n$  bits allows enumeration of  $2^n$  possibilities
  - $n$ -bit field in a header
  - $n$ -bit representation of a voice sample
  - message consisting of  $n$  bits
- *The number of bits required to represent a message is a measure of its information content*
  - more bits  $\rightarrow$  more content

## Block

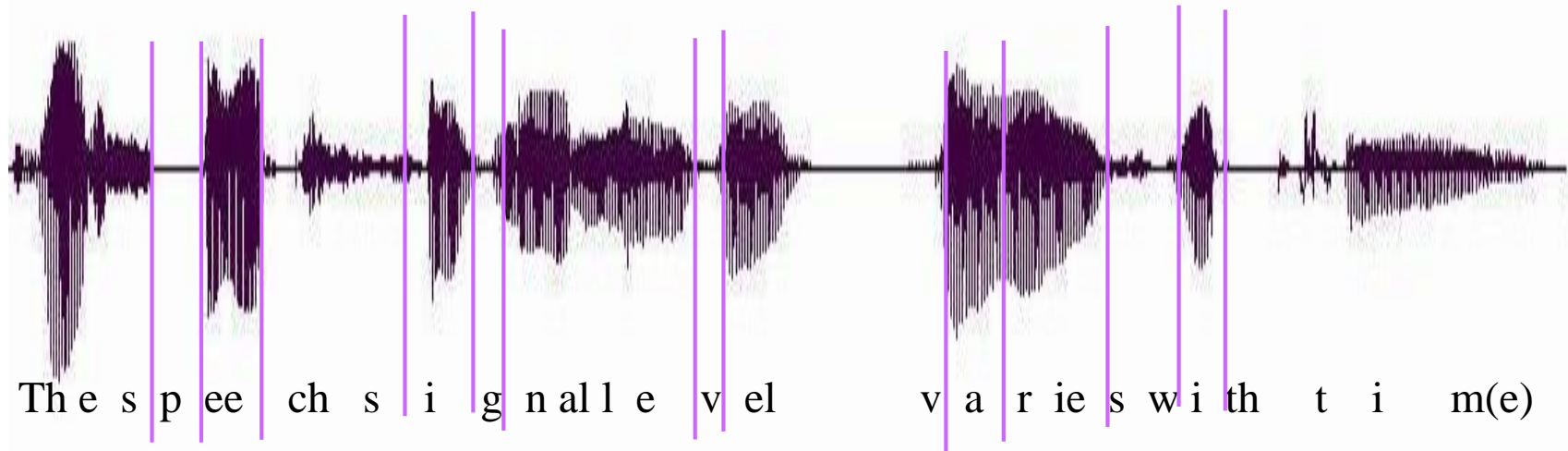
- Information that occurs in a single block
  - text message
  - data file
  - JPEG image
- size = bits / block  
or bytes/block
  - 1 Kbyte =  $2^{10}$  bytes
  - 1 Mbyte =  $2^{20}$  bytes
  - 1 Gbyte =  $2^{30}$  bytes

## Stream

- Information that is produced & transmitted *continuously*
  - real-time voice
  - MP3 audio
  - streaming video (H.261, MPEG-2)
- Bit rate = bits / second (bps)
  - 1 Kbps =  $10^3$  bps
  - 1 Mbps =  $10^6$  bps
  - 1 Gbps =  $10^9$  bps

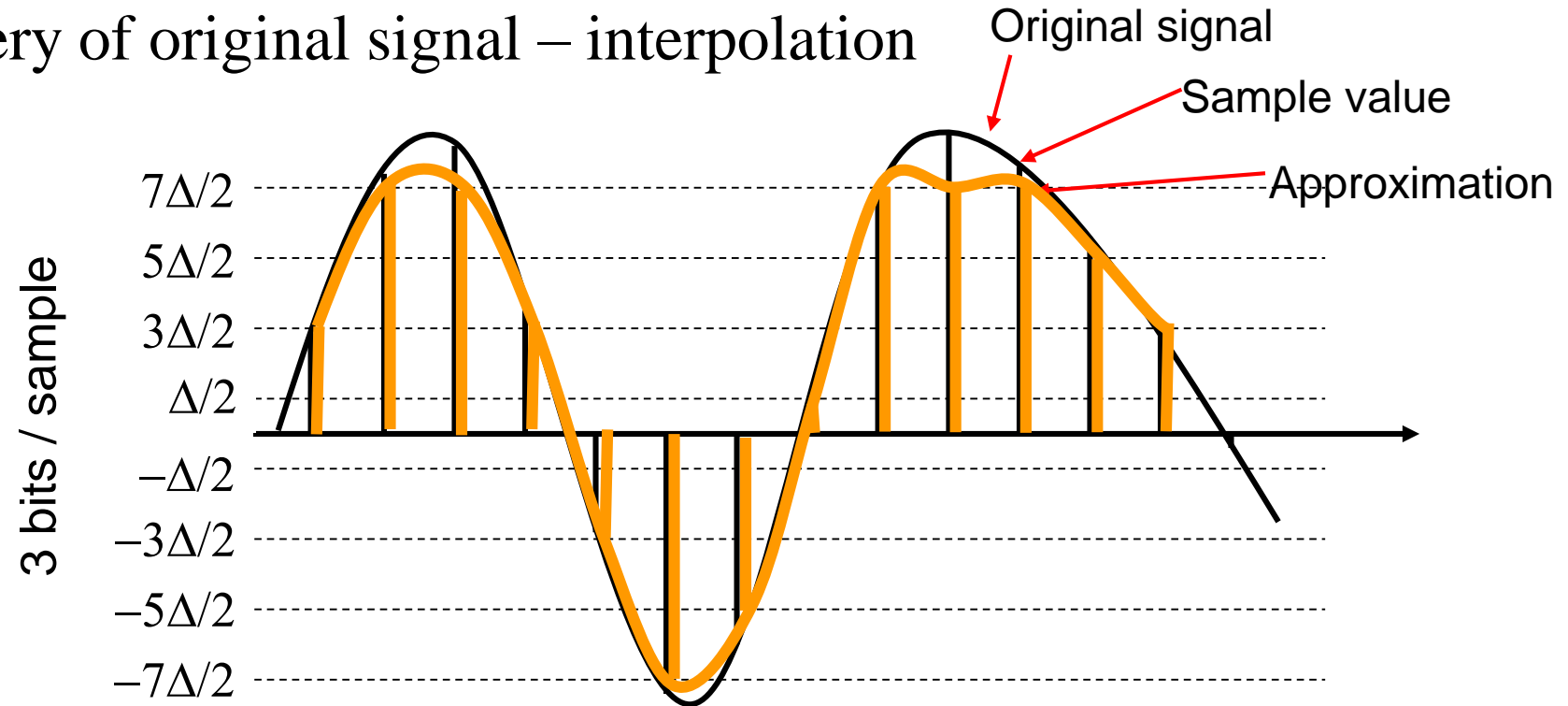
# Stream information

- a real-time voice signal must be digitized & transmitted as it is produced
- analog signal level varies continuously in time



# Digitization of analog signal

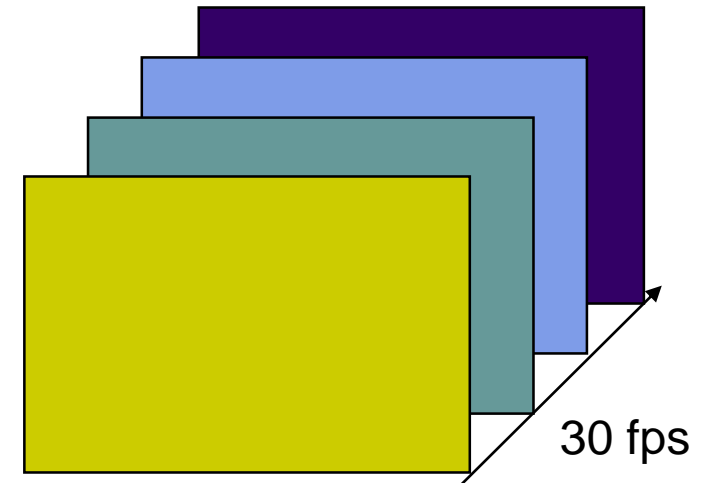
- sample analog signal in time
- find closest approximation – quantization
- recovery of original signal – interpolation



$$R_s = \text{bit rate} = \# \text{ bits/sample} \times \# \text{ samples/second}$$

# Video signal

- sequence of picture frames
  - each picture digitized & compressed
- frame repetition rate
  - 10-30-60 frames/second (fps), depending on quality
- frame resolution
  - small frames for videoconferencing
  - standard frames for conventional broadcast TV
  - HDTV frames

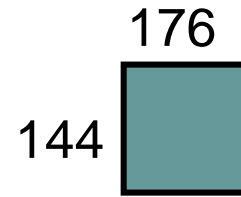


$$\text{Rate} = M \text{ bits/pixel} \times (W \times H) \text{ pixels/frame} \times F \text{ frames/second}$$



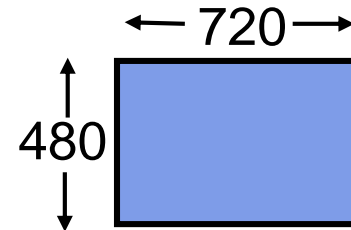
# Video frames

QCIF videoconferencing



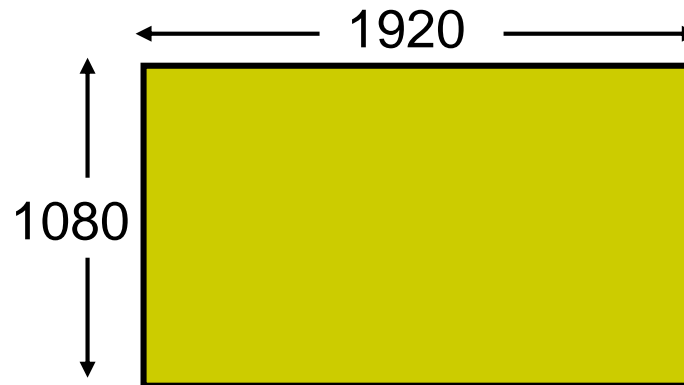
at 30 frames/sec =  
760,000 pixels/sec

Broadcast TV



at 30 frames/sec =  
 $10.4 \times 10^6$  pixels/sec

HDTV



at 30 frames/sec =  
 $62 \times 10^6$  pixels/sec

## Digital video signals

Type	Method	Format	Original	Compressed
Video Conference	H.261	176x144 or 352x288 pix @ 10-30 fps	2-36 Mbps	64-1544 Kbps
Full Motion	MPEG-2	720x480 pix @ 30 fps	249 Mbps	2-6 Mbps
HDTV	MPEG-2	1920x1080 @ 30 fps	1.6 Gbps	19-38 Mbps

# Transmission of stream information

- Constant bit-rate
  - Signals such as digitized telephone voice produce a steady stream: e.g. 64 Kbps
  - Network must support steady transfer of signal, e.g. 64 Kbps circuit
- Variable bit-rate
  - Signals such as digitized video produce a stream that varies in bit rate, e.g. according to motion and detail in a scene
  - Network must support variable transfer rate of signal, e.g. packet switching or rate-smoothing with constant bit-rate circuit

## 6.2 Digital and analog transmission



### Transmitter

- converts information into *signal* suitable for transmission
- injects energy into communications medium or channel
  - telephone converts voice into electric current
  - modem converts bits into tones

### Receiver

- receives energy from medium
- converts received signal into form suitable for delivery to user
  - telephone converts current into voice
  - modem converts tones into bits



### Communication Channel

- Pair of copper wires
- Coaxial cable
- Radio
- Light in optical fiber
- Light in air
- Infrared

### Transmission Impairments

- Signal attenuation
- Signal distortion
- Spurious noise
- Interference from other signals

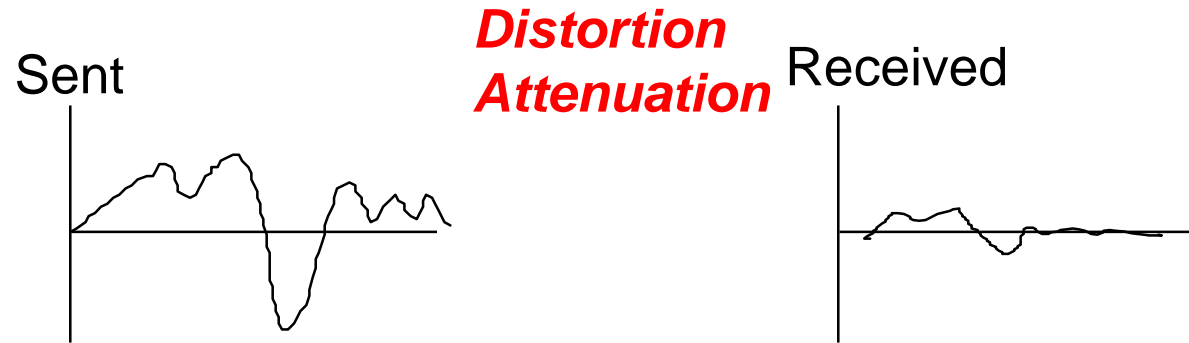
# Analog long-distance communication



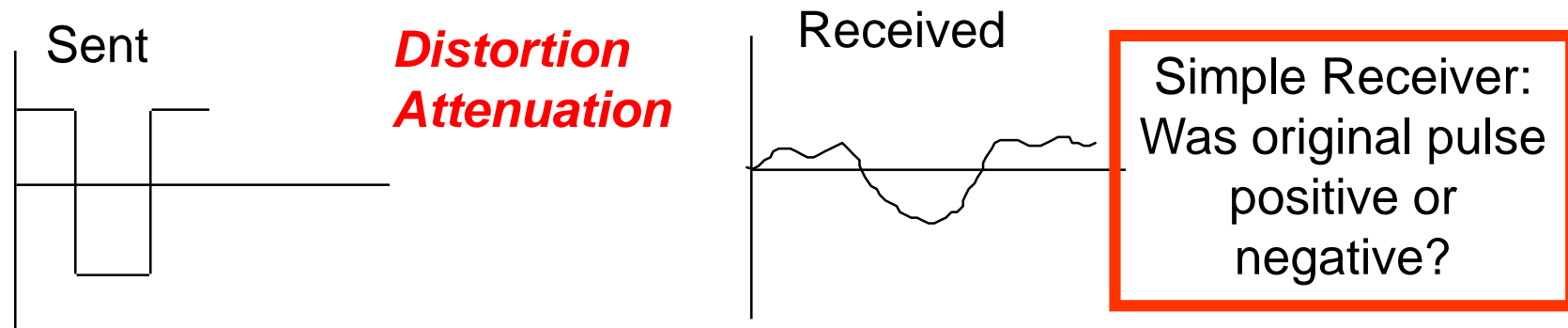
- Each repeater attempts to restore analog signal to its original form
- Restoration is imperfect
  - Distortion is not completely eliminated
  - Noise and interference is only partially removed
- Signal quality decreases with number of repeaters
- Communication is distance-limited
- Analogy: copy a song using a cassette recorder

# Analog vs. digital transmission

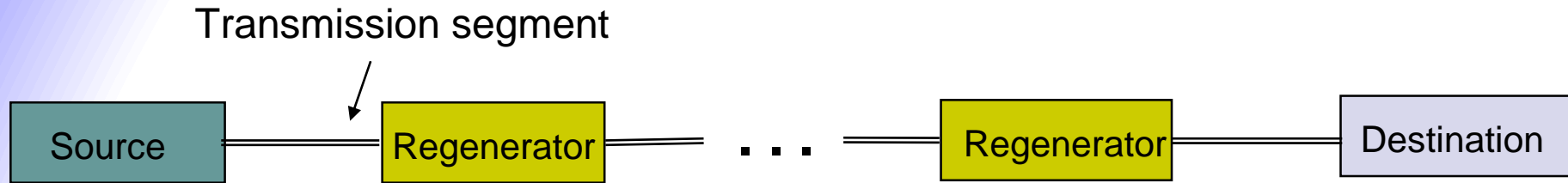
**Analog transmission:** all details must be reproduced accurately



**Digital transmission:** only discrete levels need to be reproduced



# Digital long-distance communication



- Regenerator recovers original data sequence and retransmits on next segment
- Each regeneration is like the first time!
- Can design a transmission system with very low bit error rate
- Analogy: copy an MP3 file
- Communication is possible over very long distances
- Digital systems vs. analog systems
  - Less power, longer distances, lower system cost
  - Monitoring, multiplexing, coding, encryption, protocols...

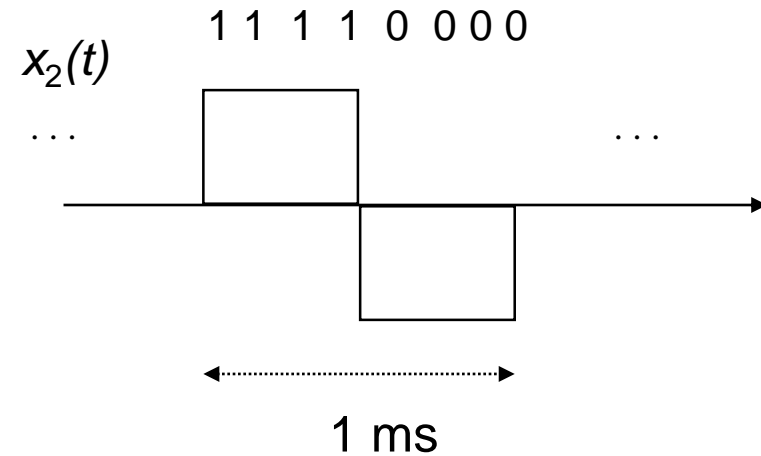
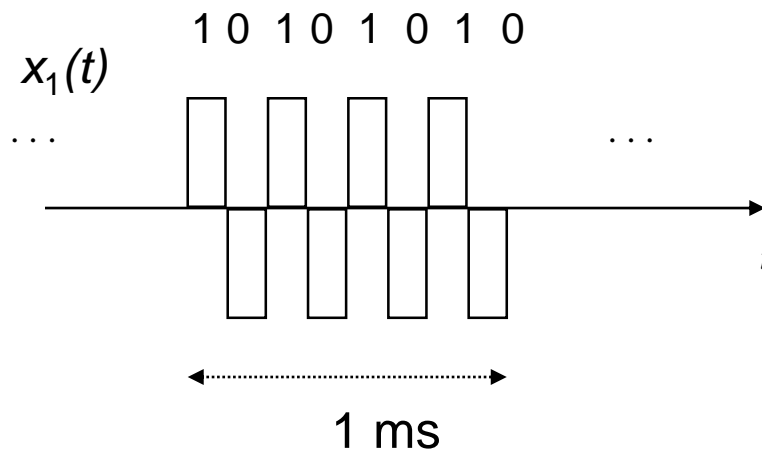


# Digitization of analog signals

1. Sampling: obtain samples of  $x(t)$  at uniformly spaced time intervals
2. Quantization: map each sample into an approximation value of finite precision
  - Pulse Code Modulation (PCM): telephone speech
  - CD audio
3. Compression: to lower bit rate further, apply additional compression method
  - Differential coding: cellular telephone speech
  - Subband coding: MP3 audio

# Sampling rate and bandwidth

- A signal that varies faster needs to be sampled more frequently
- *Bandwidth* measures how fast a signal varies



- What is the bandwidth of a signal?
- How is bandwidth related to sampling rate?

# Periodic signals

- A periodic signal with period  $T$  can be represented as sum of sinusoids using Fourier Series:

$$x(t) = a_0 + a_1 \cos(2\pi f_0 t + \phi_1) + a_2 \cos(2\pi 2f_0 t + \phi_2) + \dots \\ + a_k \cos(2\pi k f_0 t + \phi_k) + \dots$$

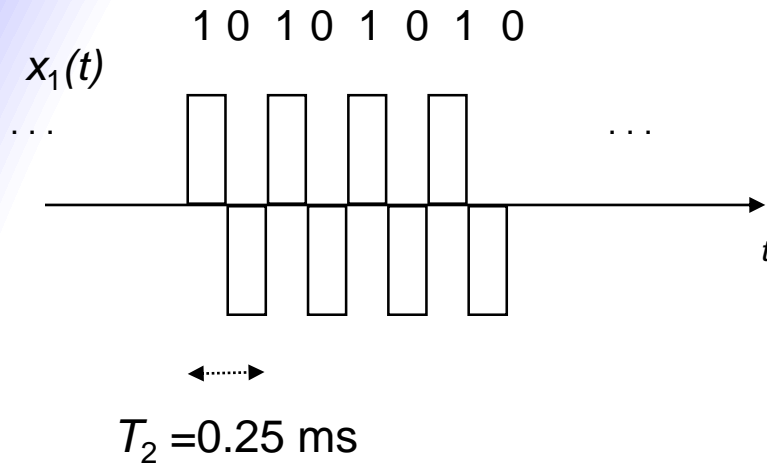
“DC”  
long-term  
average

fundamental  
frequency  $f_0=1/T$   
first harmonic

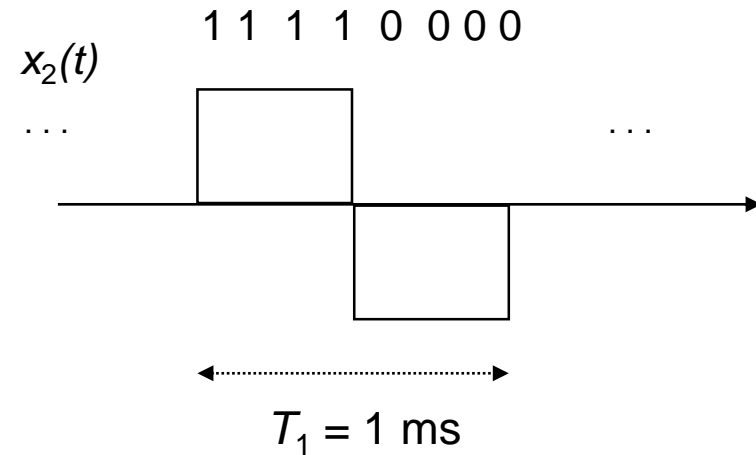
$k$ th harmonic

- $|a_k|$  determines amount of power in  $k$ th harmonic
- Amplitude spectrum  $|a_0|, |a_1|, |a_2|, \dots$

## Example - Fourier series



$$\begin{aligned}
 x_1(t) = & 0 + \frac{4}{\pi} \cos(2\pi 4000t) \\
 & + \frac{4}{3\pi} \cos(2\pi 3(4000)t) \\
 & + \frac{4}{5\pi} \cos(2\pi 5(4000)t) + \dots
 \end{aligned}$$



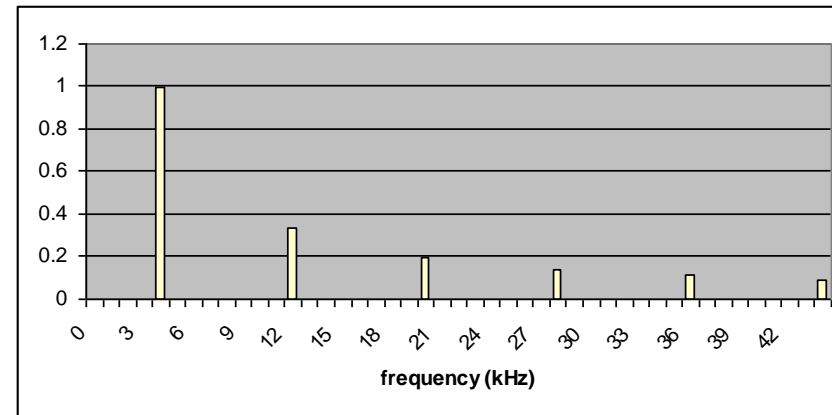
$$\begin{aligned}
 x_2(t) = & 0 + \frac{4}{\pi} \cos(2\pi 1000t) \\
 & + \frac{4}{3\pi} \cos(2\pi 3(1000)t) \\
 & + \frac{4}{5\pi} \cos(2\pi 5(1000)t) + \dots
 \end{aligned}$$

Only odd harmonics have power

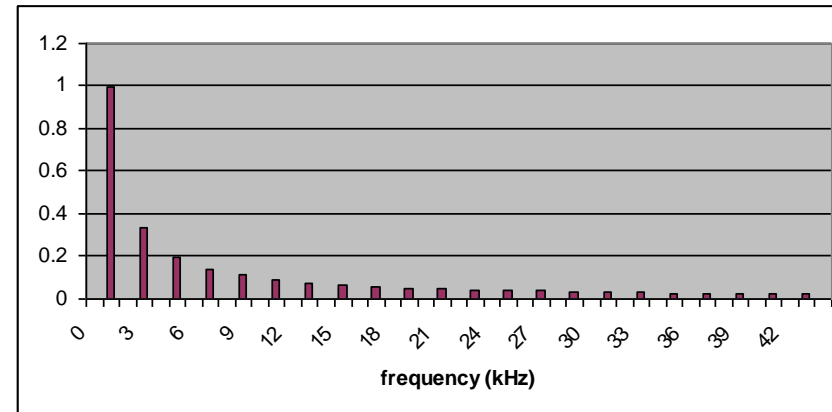
# Spectrum and bandwidth

- Spectrum of a signal: magnitude of amplitudes as a function of frequency
- $x_1(t)$  varies faster in time and has more high frequency content than  $x_2(t)$
- Bandwidth  $W_s$  is defined as range of frequencies where a signal has non-negligible power, e.g. range of band that contains 99% of total signal power

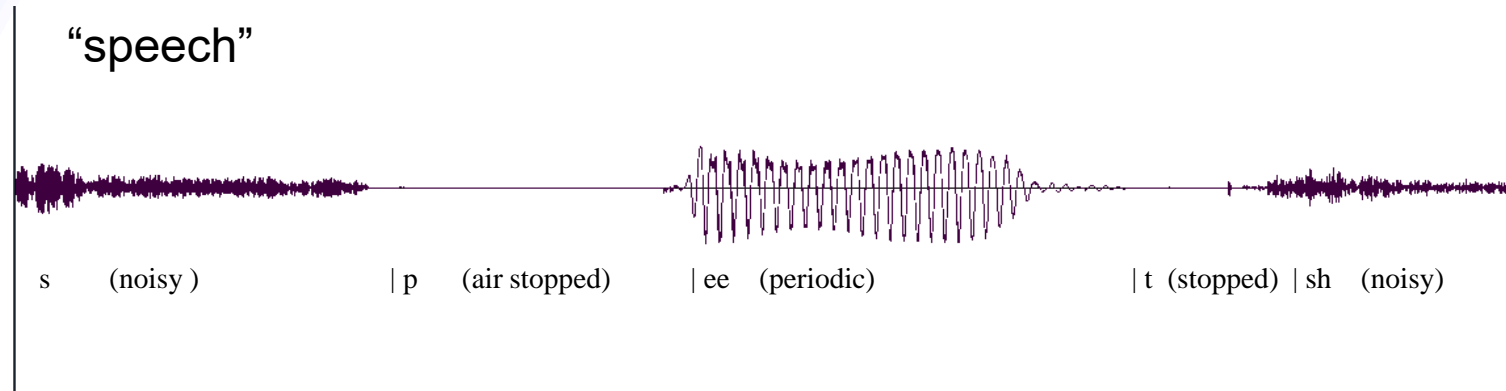
Spectrum of  $x_1(t)$



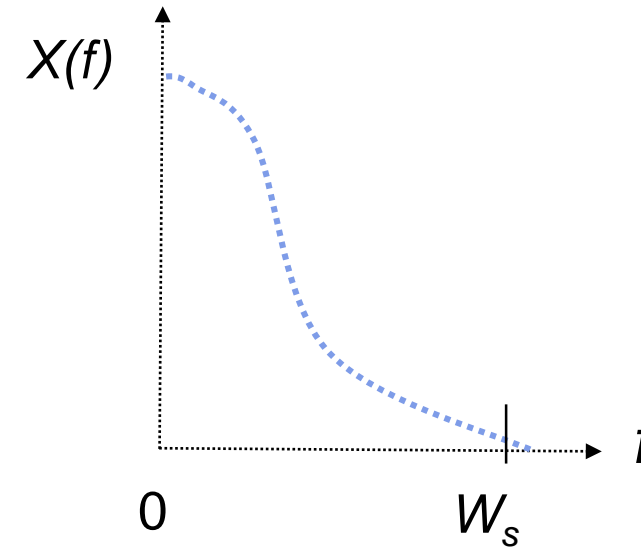
Spectrum of  $x_2(t)$



# Bandwidth of general signals

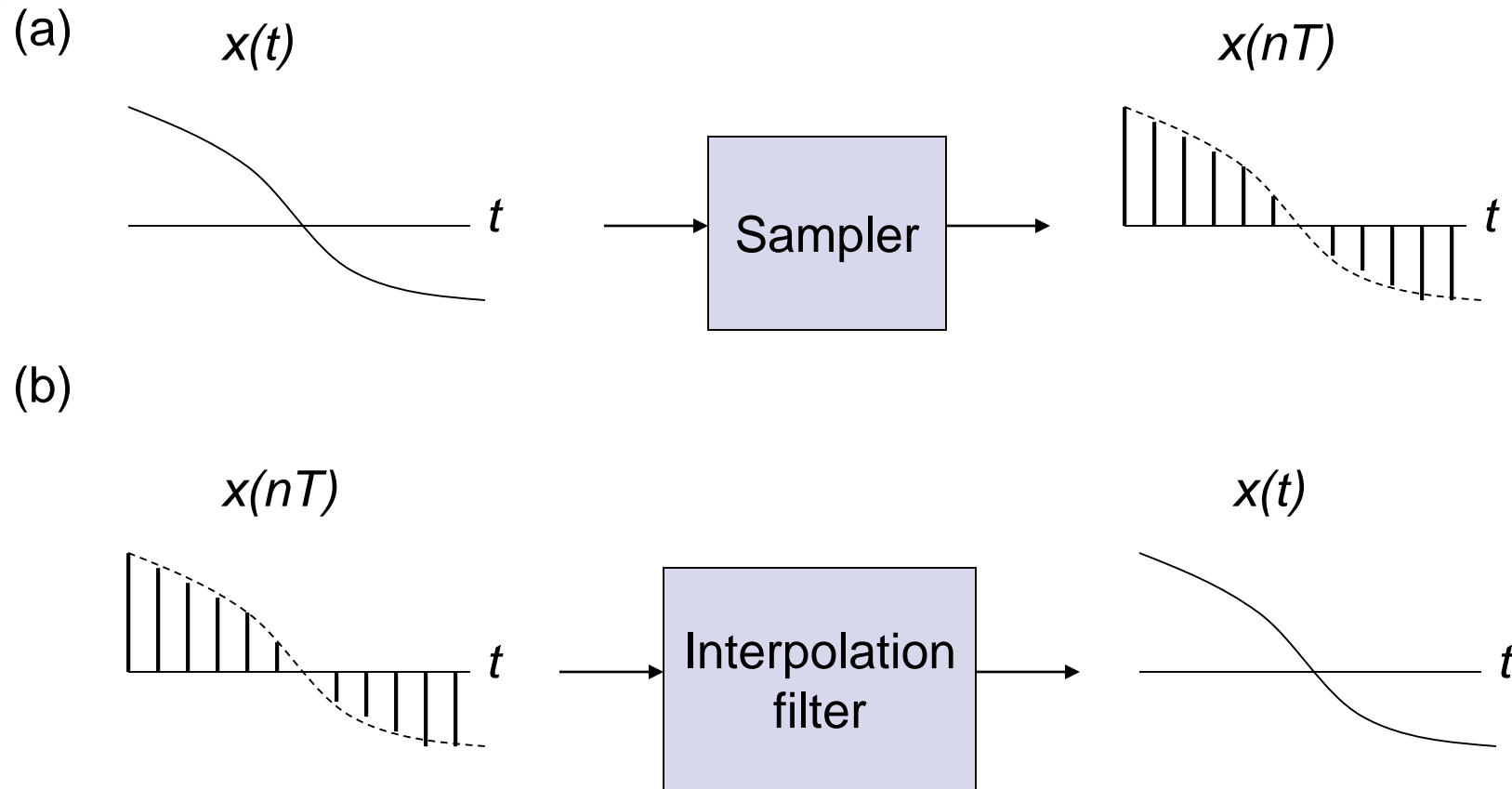


- Not all signals are periodic
  - e.g. voice signal varies according to sound
  - vowels are periodic, “s” is noise-like
- Spectrum of long-term signal
  - averages over many sounds, many speakers
  - involves Fourier transform
- Telephone speech: 4 KHz
- CD Audio: 22 KHz

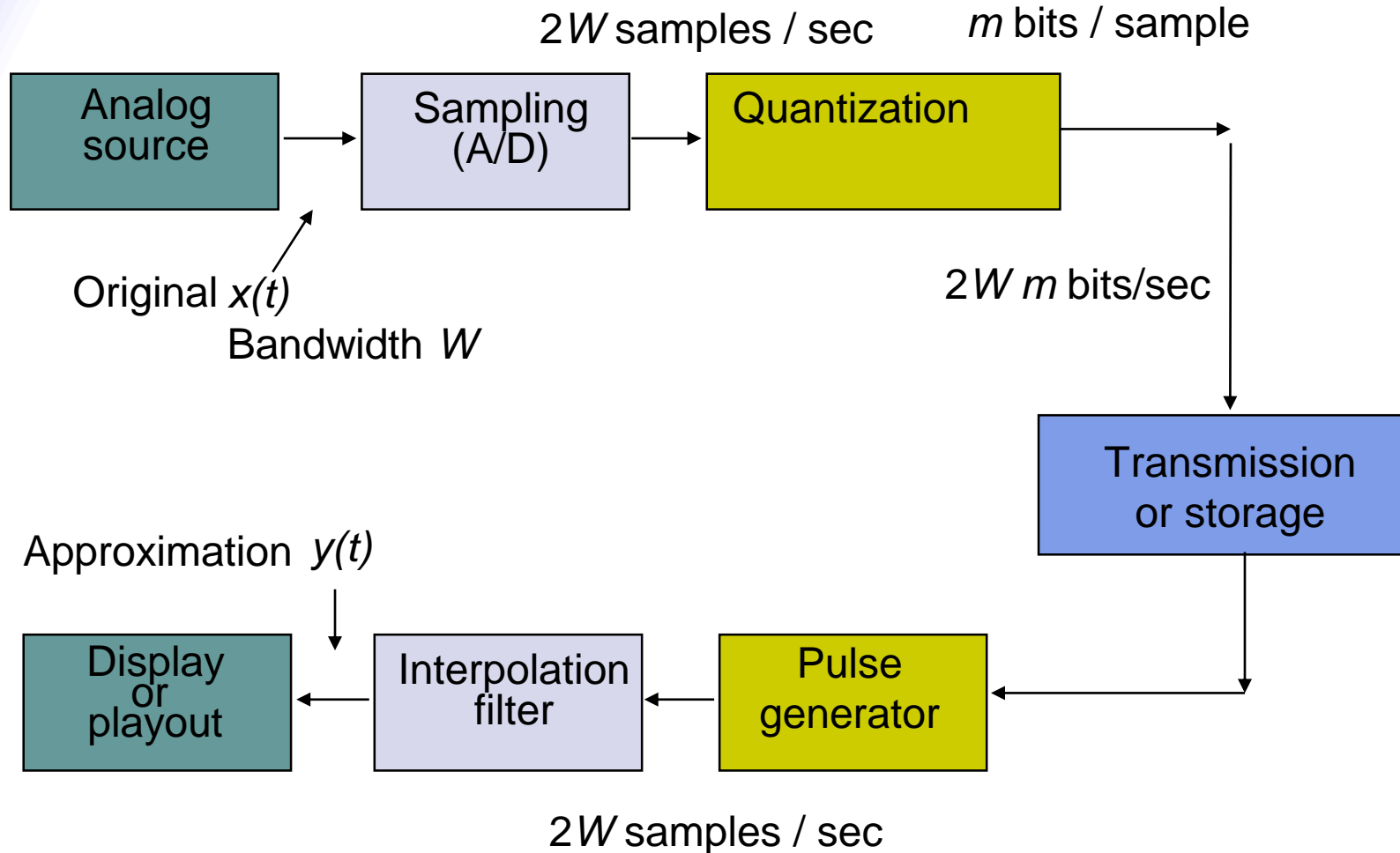


# Sampling theorem

Nyquist: Perfect reconstruction if sampling rate  $1/T > 2W_s$

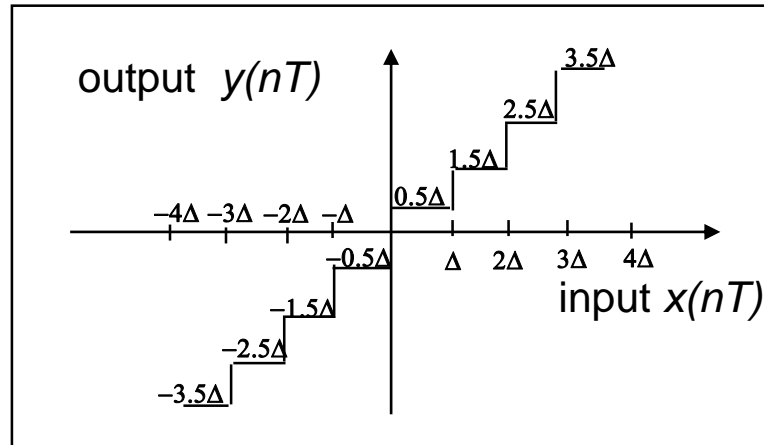


# Digital transmission of analog information



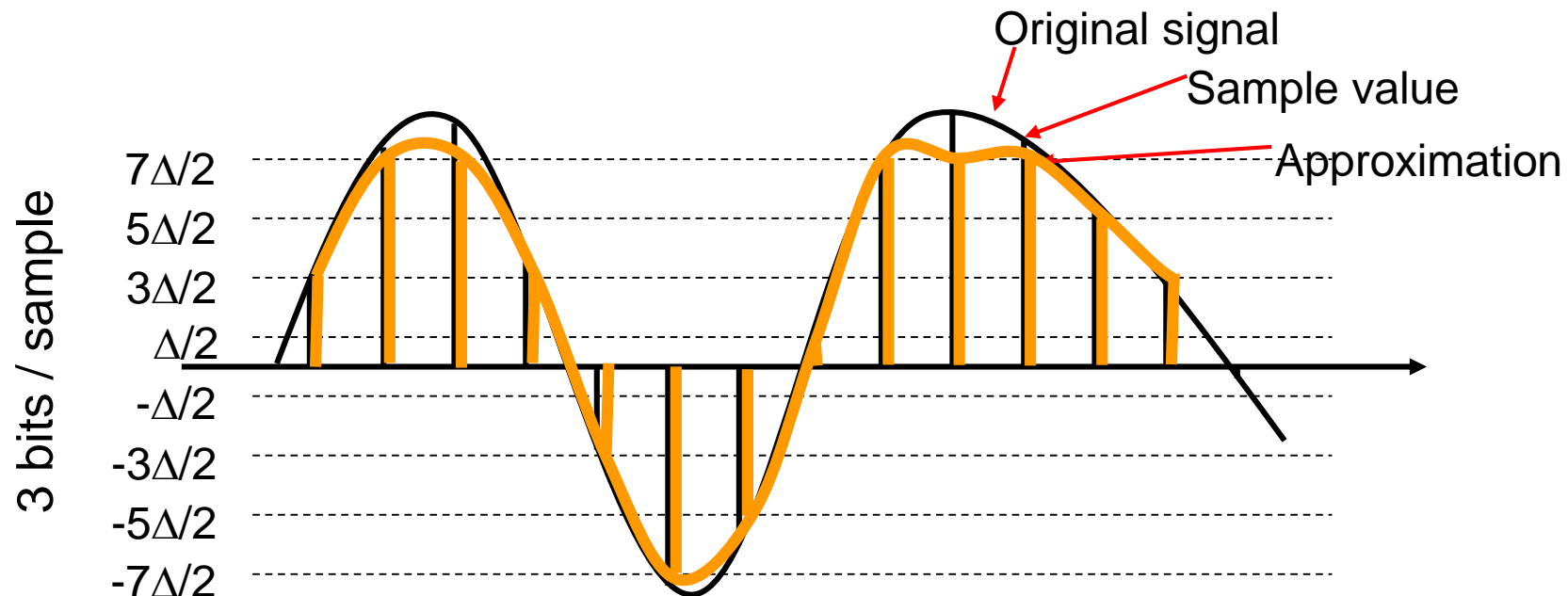


# Quantization of analog samples



Quantizer maps input into closest of  $2^m$  representation values

Quantization error:  
"noise" =  $x(nT) - y(nT)$



# Example: voice and audio

## Telephone voice

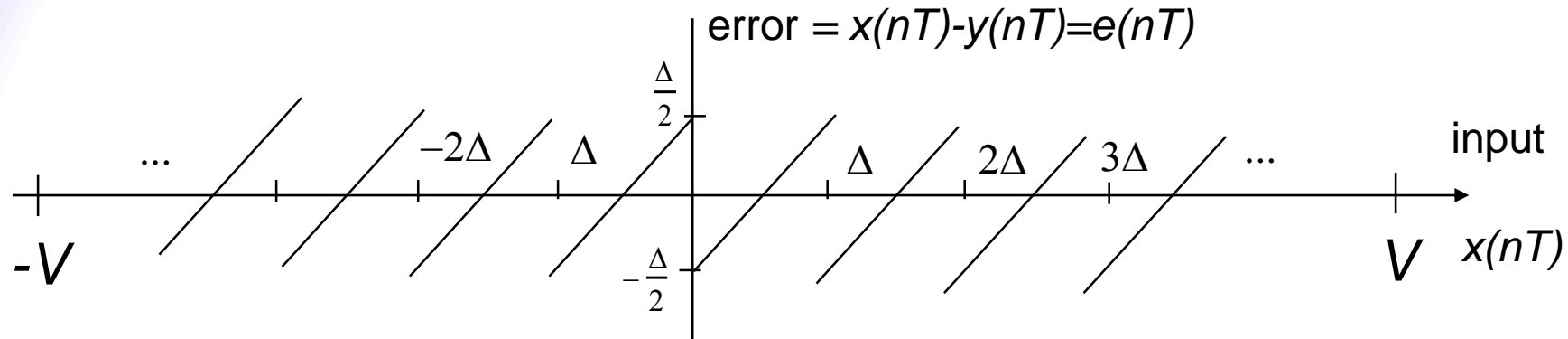
- $W_s = 4 \text{ KHz} \rightarrow 8000 \text{ samples/sec}$
- 8 bits/sample
- $R_s = 8 \times 8000 = 64 \text{ Kbps}$
- Cellular phones use more powerful compression algorithms: 8-12 Kbps

## CD Audio

- $W_s = 22 \text{ KHz} \rightarrow 44,000 \text{ samples/sec}$
- 16 bits/sample
- $R_s = 16 \times 44000 = 704 \text{ Kbps}$  per audio channel
- MP3 uses more powerful compression algorithms: 50 Kbps per audio channel

# Quantizer performance

$M = 2^m$  levels, Dynamic range  $(-V, V)$   $\Delta = 2V/M$



If the number of levels  $M$  is large, then the error is approximately uniformly distributed between  $(-\Delta/2, \Delta/2)$

Average Noise Power = Mean Square Error:

$$\sigma_e^2 = \int_{-\frac{\Delta}{2}}^{\frac{\Delta}{2}} x^2 \frac{1}{\Delta} dx = \frac{\Delta^2}{12}$$

Figure of merit:

Signal-to-Noise Ratio (SNR) = average signal power / average noise power

Let  $\sigma_x^2$  be the signal power, then

$$SNR = \frac{\sigma_x^2}{\Delta^2/12} = \frac{12\sigma_x^2}{4V^2/M^2} = 3 \left(\frac{\sigma_x}{V}\right)^2 M^2 = 3 \left(\frac{\sigma_x}{V}\right)^2 2^{2m}$$

The ratio  $V/\sigma_x \approx 4$

The SNR is usually stated in decibels:

$$SNR \text{ dB} = 10 \log_{10} \sigma_x^2 / \sigma_e^2 = 6m + 10 \log_{10} 3\sigma_x^2 / V^2$$

$$SNR \text{ dB} = 6m - 7.27 \text{ dB} \quad \text{for } V/\sigma_x = 4$$

## Example: telephone speech

$W = 4$  KHz, so Nyquist sampling theorem

$\Rightarrow 2W = 8000$  samples/second

Suppose error requirement = 1% error

$$\text{SNR} = 10 \log(1/.01)^2 = 40 \text{ dB}$$

Assume  $V/\sigma_x = 4$ , then

$$40 \text{ dB} = 6m - 7$$

$$\Rightarrow m = 8 \text{ bits/sample}$$

PCM telephone speech:

$$\text{Bit rate} = 8000 \times 8 \text{ bits/sec} = 64 \text{ Kbps}$$

# Communication channel

- A *physical medium* is an inherent part of a communications system
  - copper wires, radio medium, or optical fiber
- Communications system includes electronic or optical devices that are part of the path followed by a signal
  - equalizers, amplifiers, signal conditioners
- By *communication channel* we refer to the combined end-to-end physical medium and attached devices
- Sometimes we use the term *filter* to refer to a channel especially in the context of a specific mathematical model for the channel



### Signal Bandwidth

- In order to transfer data faster, a signal has to vary more quickly.

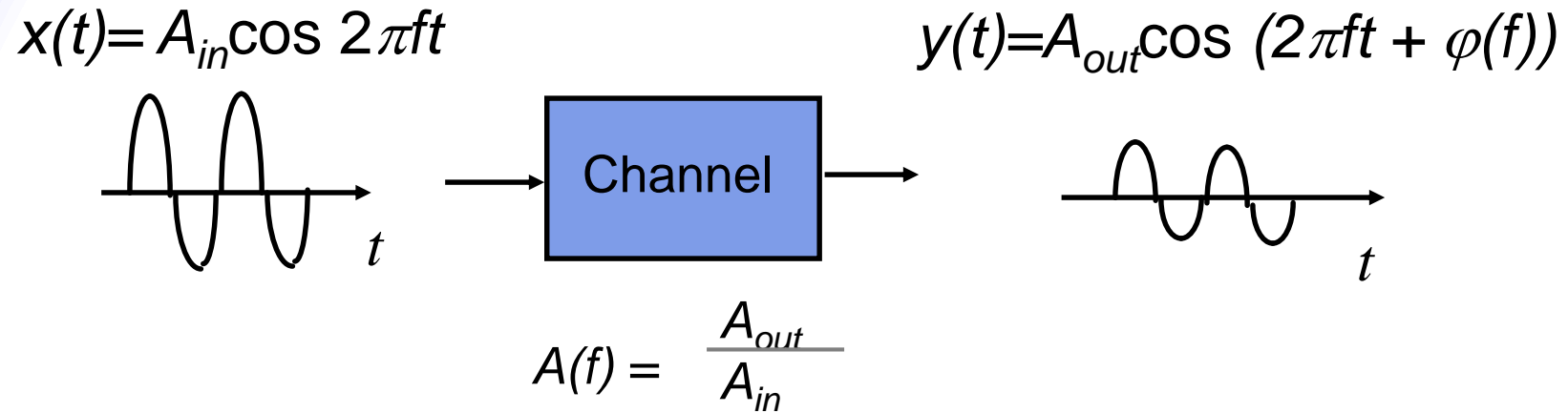
### Channel Bandwidth

- A channel or medium has an inherent limit on how fast the signals it passes can vary
- *Limits how tightly input pulses can be packed*

### Transmission Impairments

- Signal attenuation
- Signal distortion
- Spurious noise
- Interference from other signals
- *Limits accuracy of measurements on received signal*

# Frequency domain channel characterization



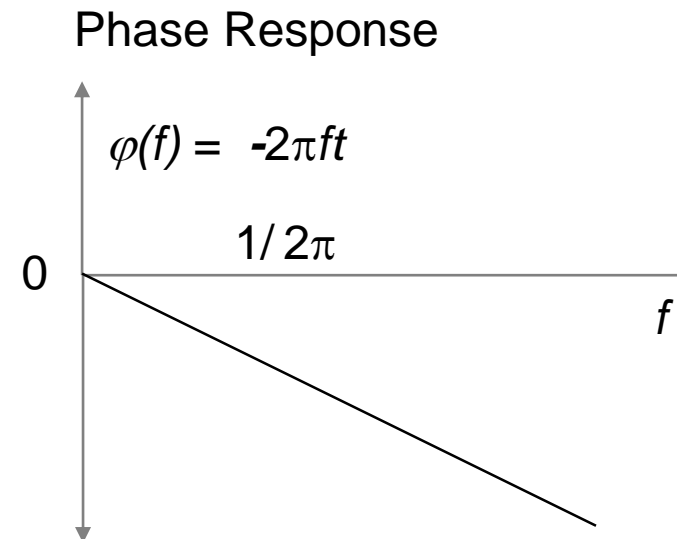
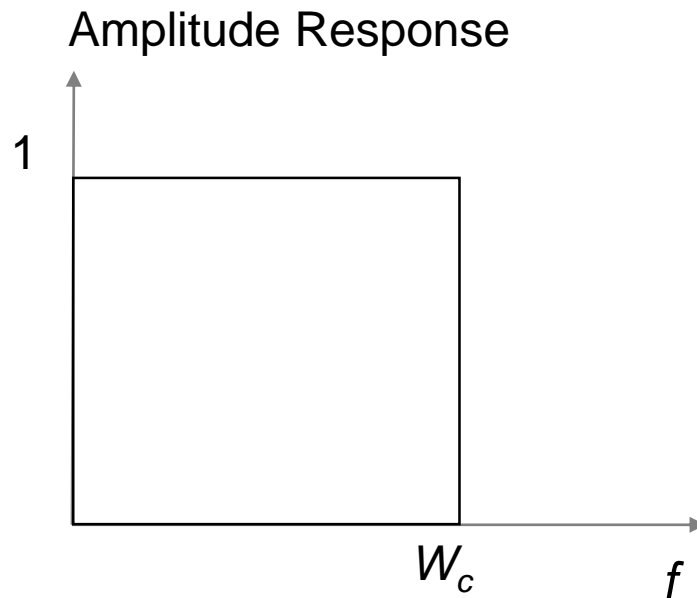
- Apply sinusoidal input at frequency  $f$ 
  - Output is sinusoid at same frequency, but attenuated & phase-shifted
  - Measure amplitude of output sinusoid (of same frequency  $f$ )
  - Calculate amplitude response
    - $A(f)$  = ratio of output amplitude to input amplitude
  - If  $A(f) \approx 1$ , then input signal passes readily
  - If  $A(f) \approx 0$ , then input signal is blocked
- Bandwidth  $W_c$  is range of frequencies passed by channel



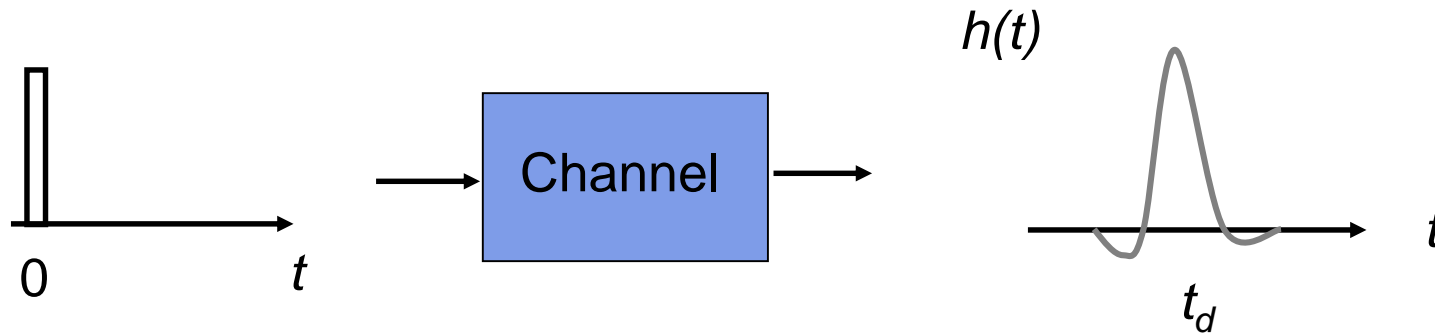
# Ideal Low-Pass Filter

- Ideal filter: all sinusoids with frequency  $f < W_c$  are passed without attenuation and delayed by  $\tau$  seconds; sinusoids at other frequencies are blocked

$$y(t) = A_{in} \cos(2\pi f t - 2\pi f \tau) = A_{in} \cos(2\pi f(t - \tau)) = x(t - \tau)$$



# Time-domain characterization

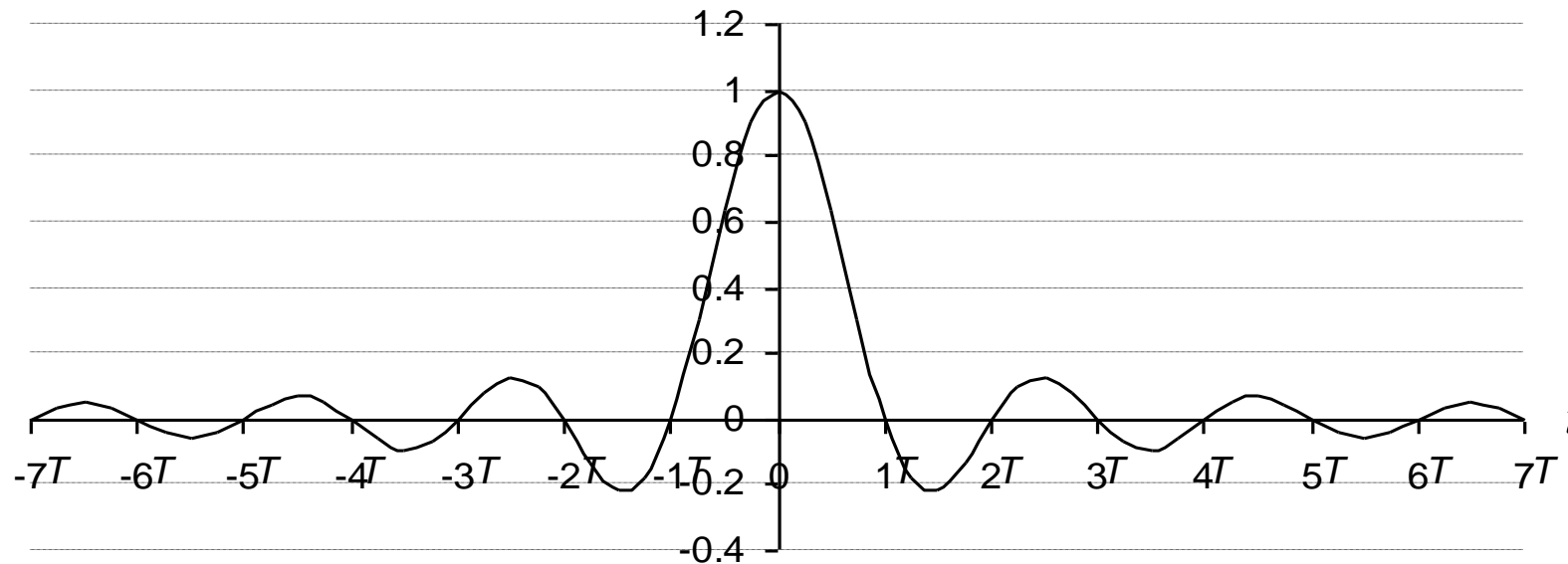


- Time-domain characterization of a channel requires finding the *impulse response*  $h(t)$
- Apply a very narrow pulse to a channel and observe the channel output
  - $h(t)$  typically a delayed pulse with ringing
- Interested in system designs with  $h(t)$  that can be packed closely without interfering with each other

# Nyquist pulse with zero InterSymbol Interference (ISI)

- For channel with ideal low pass amplitude response of bandwidth  $W_c$ , the impulse response is a Nyquist pulse  $h(t)=s(t-\tau)$ , where  $T=1/2W_c$ , and

$$s(t) = \sin(2\pi W_c t) / 2\pi W_c t$$



- $s(t)$  has zero crossings at  $t = kT$ ,  $k = \pm 1, \pm 2, \dots$
- Pulses can be packed every  $T$  seconds with *zero interference*

## Example - composite waveform

Three Nyquist pulses shown separately

- $+s(t)$
- $+s(t-T)$
- $-s(t-2T)$

### Composite waveform

$$r(t) = s(t) + s(t-T) - s(t-2T)$$

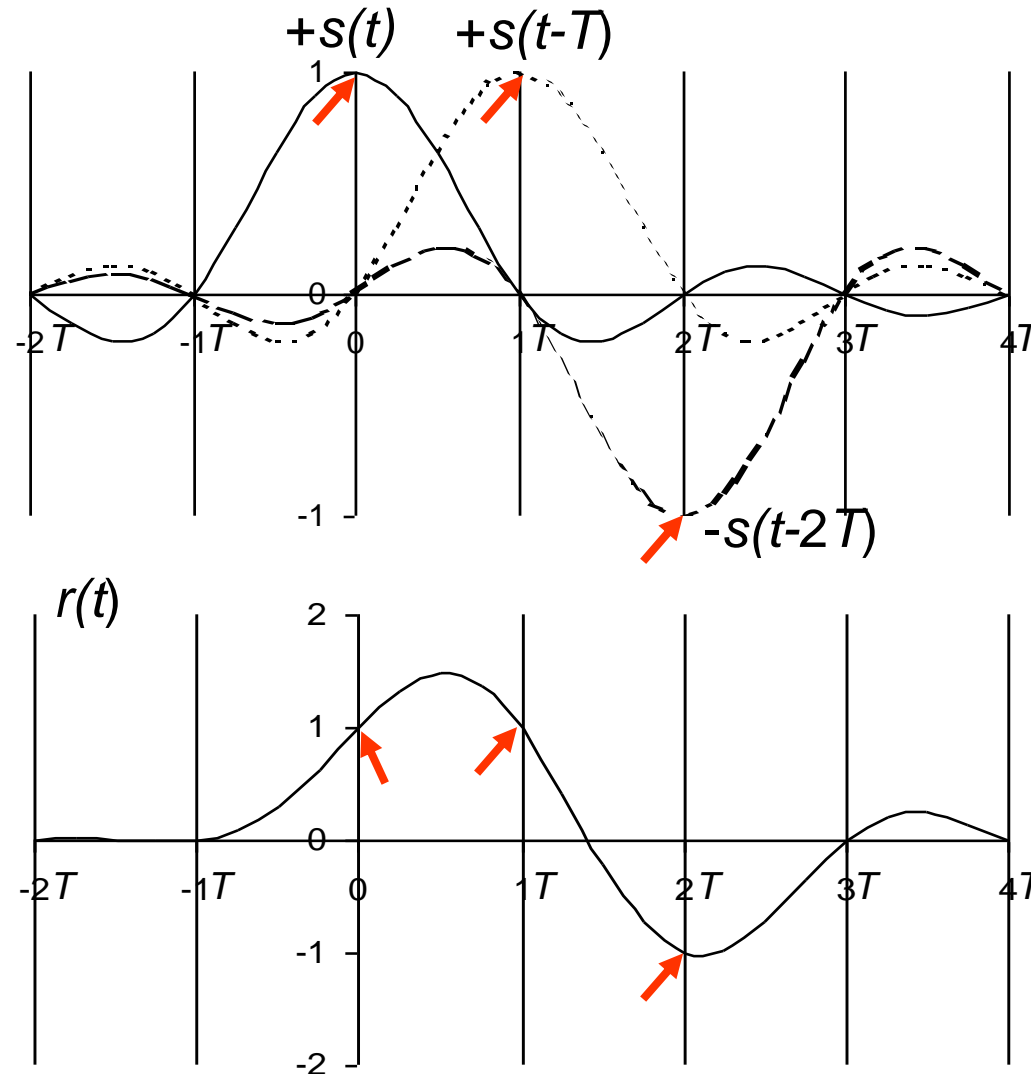
Samples at  $kT$

$$r(0) = s(0) + s(-T) - s(-2T) = +1$$

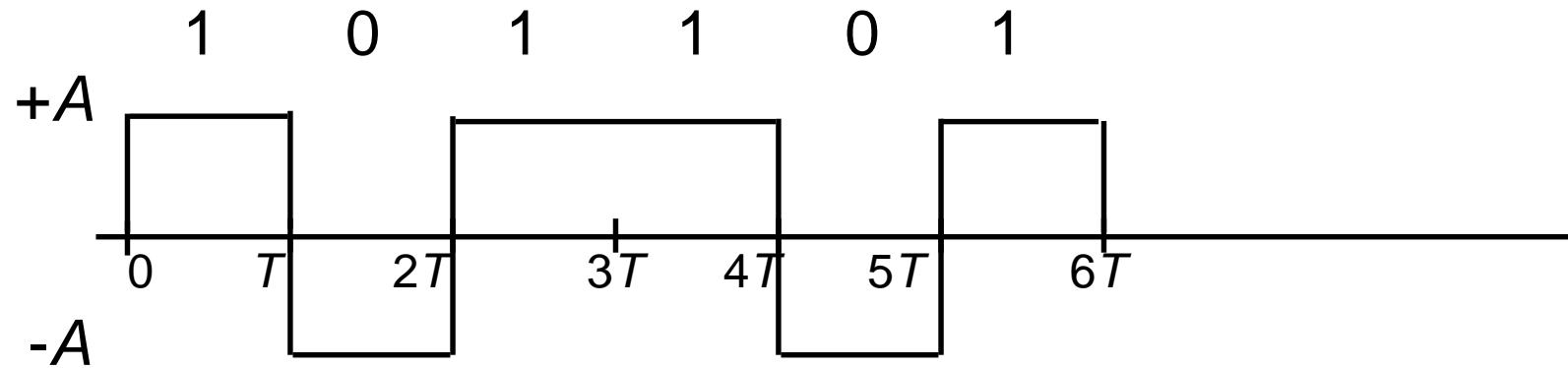
$$r(T) = s(T) + s(0) - s(-T) = +1$$

$$r(2T) = s(2T) + s(T) - s(0) = -1$$

**Zero ISI at sampling times  $kT$**



## Digital binary signal



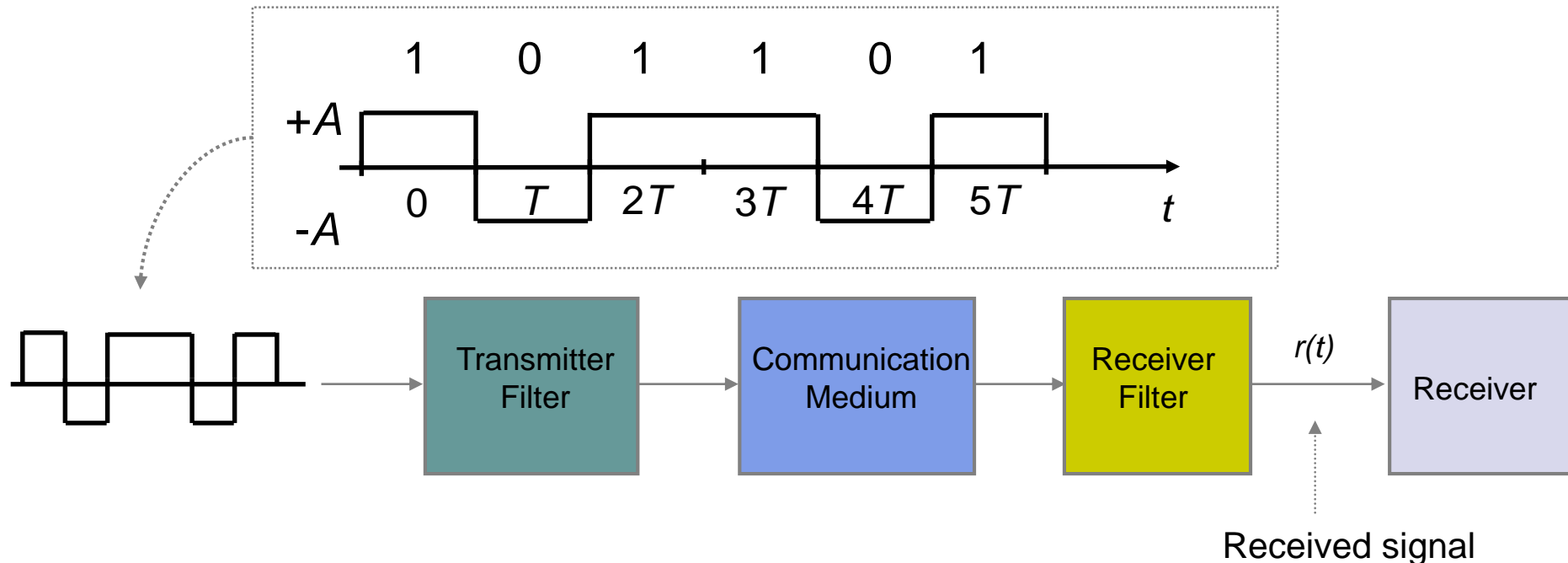
$$\text{Bit rate} = 1 \text{ bit} / T \text{ seconds}$$

For a given communications medium:

- How do we increase transmission speed?
- How do we achieve reliable communications?
- Are there limits to speed and reliability?

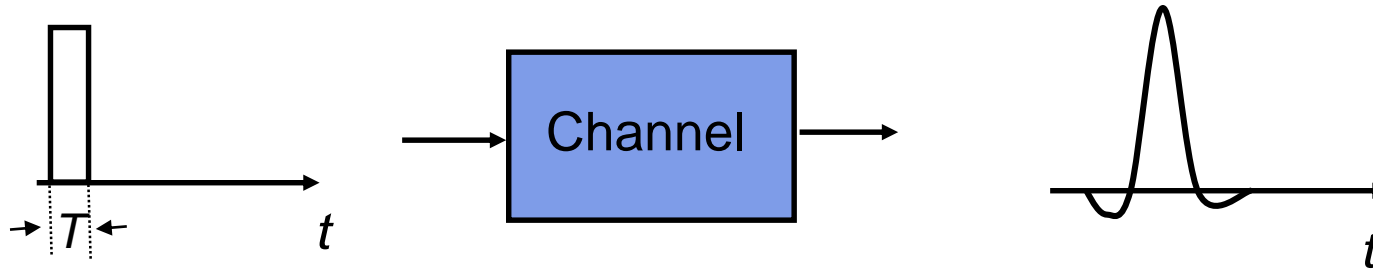
# Signaling with Nyquist pulses

- $p(t)$  pulse at receiver in response to a single input pulse (takes into account pulse shape at input, transmitter and receiver filters, and communications medium)
- $r(t)$  waveform that appears in response to sequence of pulses
- If  $p(t)$  is a Nyquist pulse, then  $r(t)$  has zero intersymbol interference (ISI) when sampled at multiples of  $T$



## Pulse transmission rate

- Objective: maximize pulse rate through a channel, that is, make  $T$  as small as possible



- If input is a narrow pulse, then typical output is a spread-out pulse with ringing
- Question: How frequently can these pulses be transmitted without interfering with each other?
- Answer:  $2 \times W_c$  pulses/second

where  $W_c$  is the bandwidth of the channel

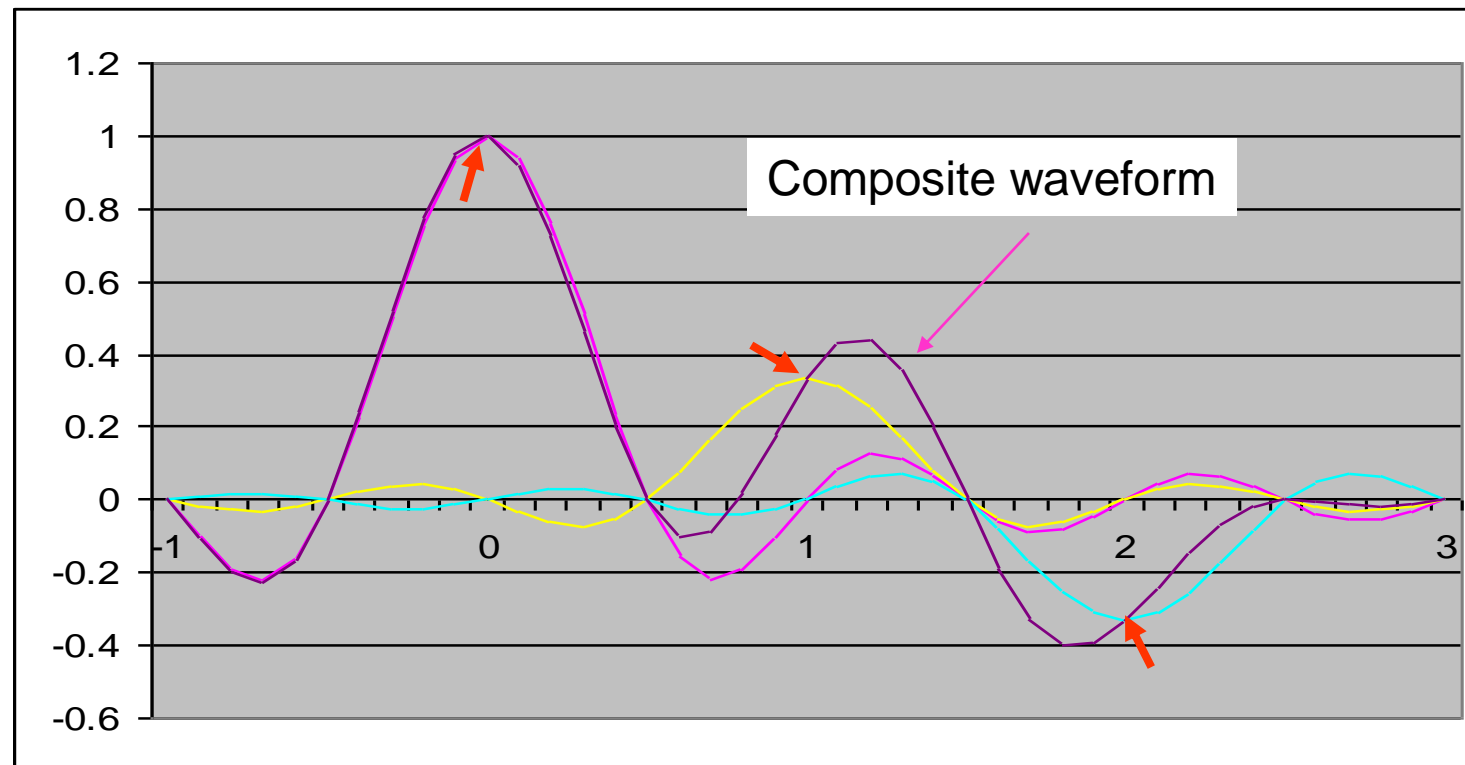
# Multilevel signaling

- Nyquist pulses achieve the maximum signalling rate with zero ISI,  
 $2W_c$  pulses per second or  
 $2W_c \text{ pulses} / W_c \text{ Hz} = 2 \text{ pulses} / \text{Hz}$
- With two signal levels, each pulse carries one bit of information  
Bit rate =  $2W_c$  bits/second
- With  $M = 2^m$  signal levels, each pulse carries  $m$  bits  
Bit rate =  $2W_c$  pulses/sec. \*  $m$  bits/pulse =  $2W_c m$  bps
- *Bit rate can be increased by increasing number of levels*
- *$r(t)$  includes additive noise, that limits number of levels that can be used reliably.*



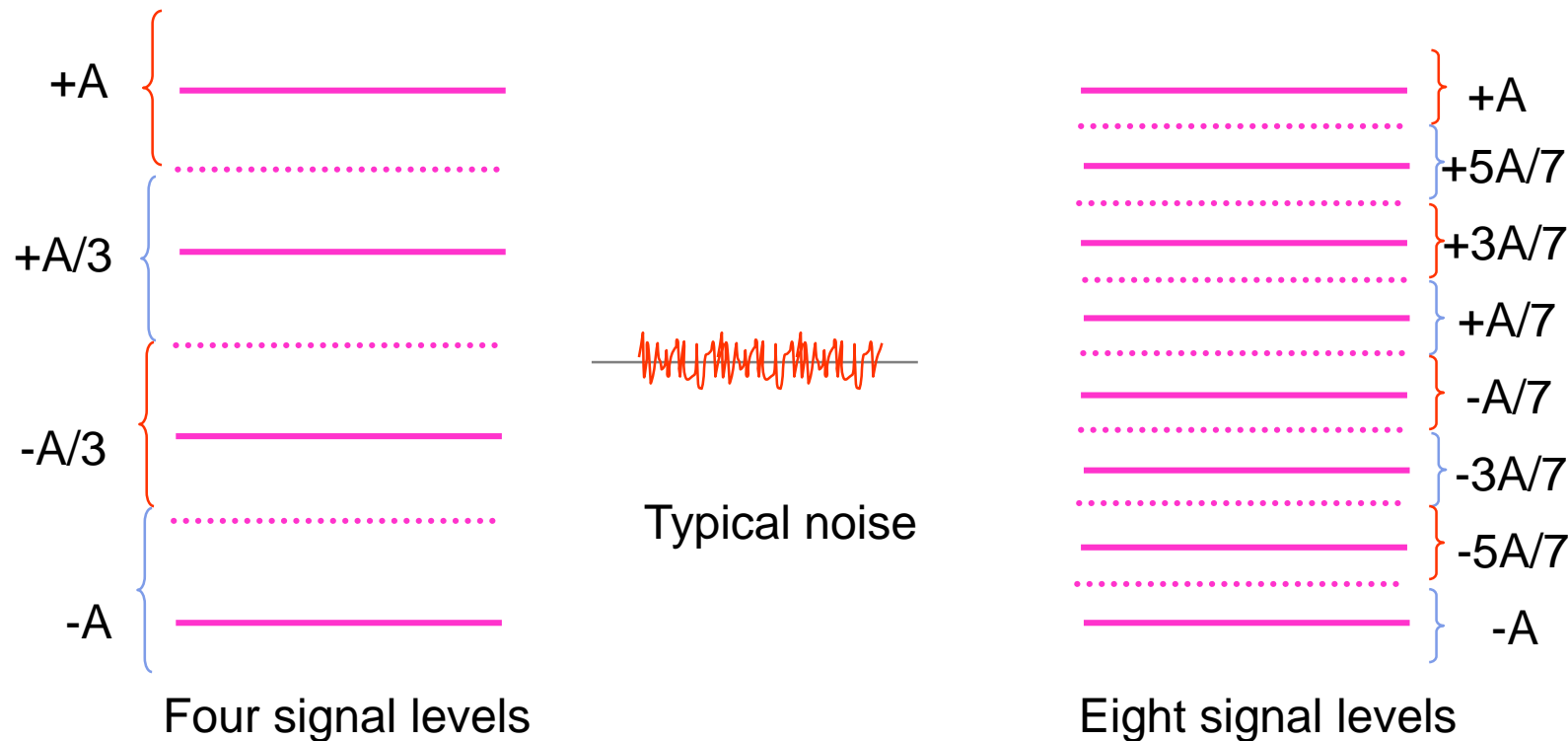
# Example

- Four levels  $\{-1, -1/3, 1/3, +1\}$  for  $\{00, 01, 10, 11\}$
- Waveform for 11, 10, 01 sends  $+1, +1/3, -1/3$
- Zero ISI at sampling instants



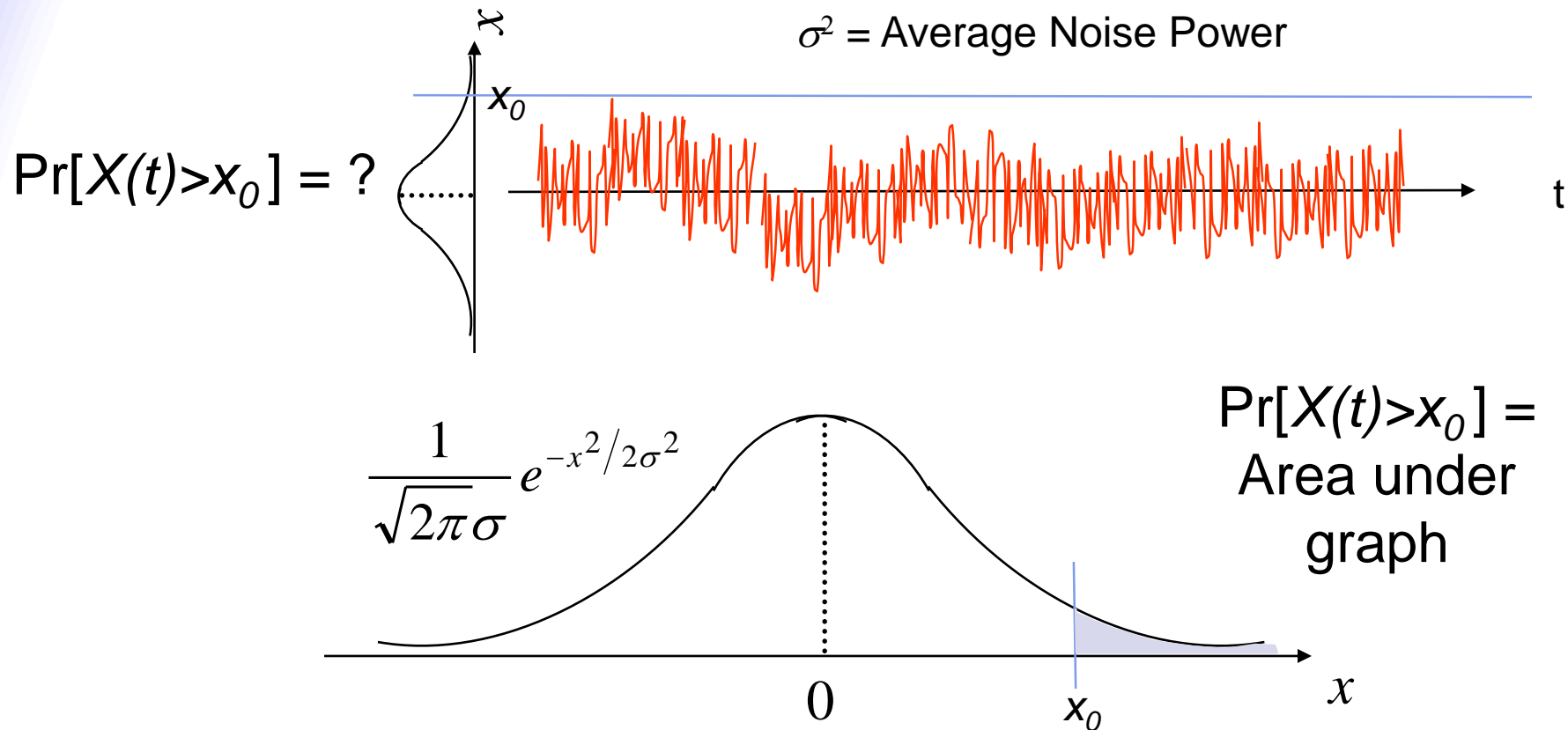
# Noise limits accuracy

- Receiver makes decision based on transmitted pulse level + noise
- Error rate depends on relative value of noise amplitude and spacing between signal levels
- Large (positive or negative) noise values can cause wrong decision
- Noise level below impacts 8-level signaling more than 4-level signaling



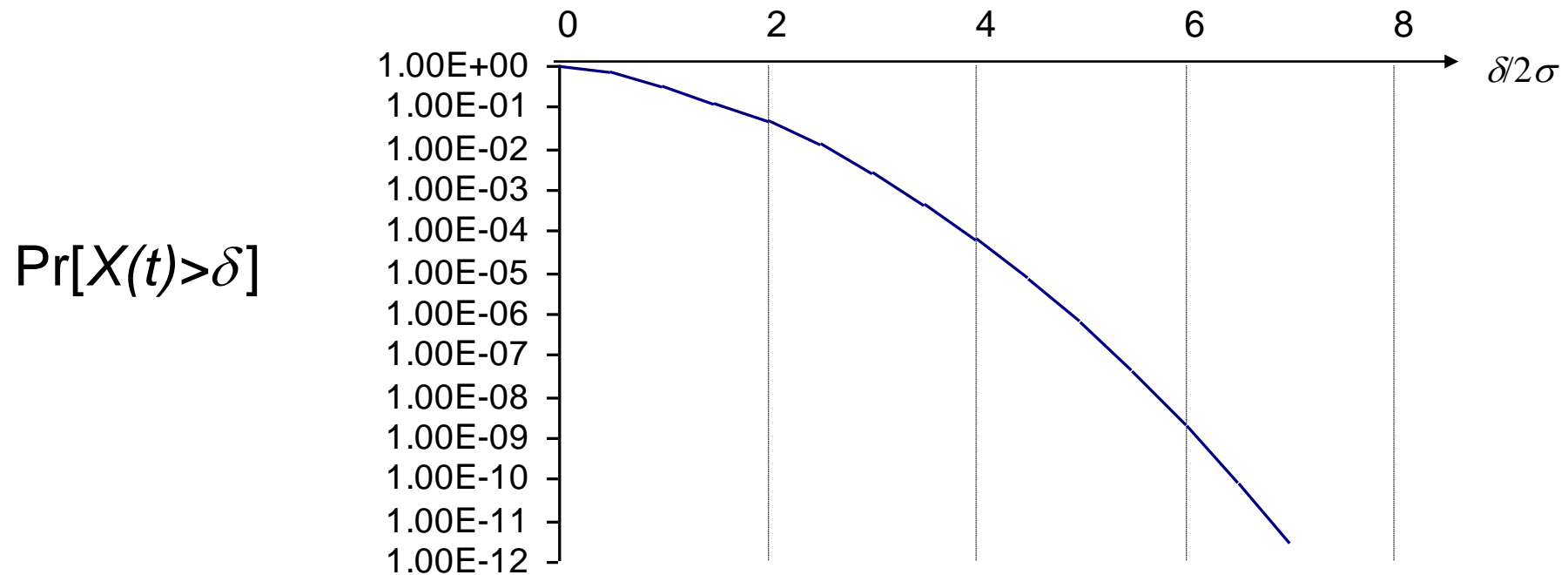
# Noise distribution

- Noise is characterized by probability density of amplitude samples
- Likelihood that certain amplitude occurs
- Thermal electronic noise is inevitable (due to vibrations of electrons)
- Noise distribution is Gaussian (bell-shaped) as below

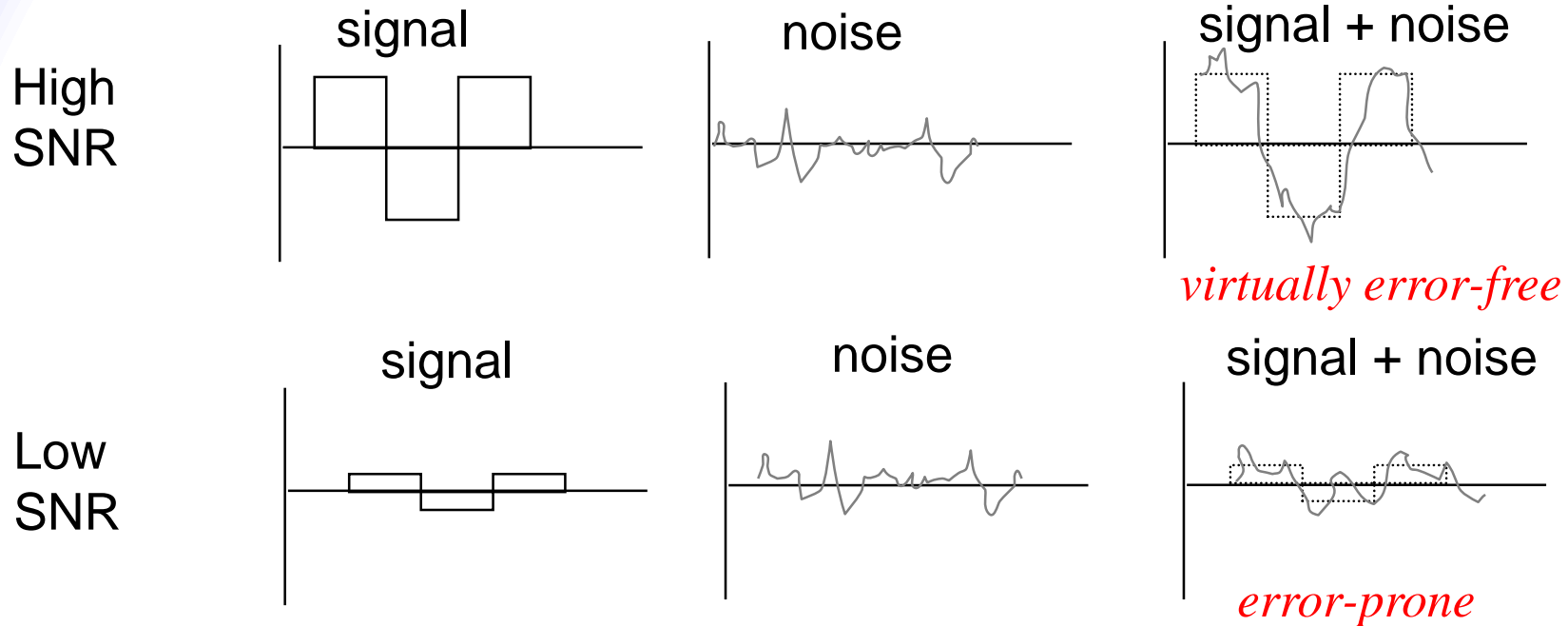


# Probability of error

- Error occurs if noise value exceeds certain magnitude
- Probability of large values drops quickly with Gaussian noise
- Target probability of error achieved by designing system so separation between signal levels is appropriate relative to average noise power



# Channel noise affects reliability



$$\text{SNR} = \frac{\text{Average Signal Power}}{\text{Average Noise Power}}$$

$$\text{SNR (dB)} = 10 \log_{10} \text{SNR}$$

- If transmitted power is limited, then as  $M$  increases spacing between levels decreases
- Presence of noise at receiver causes more frequent errors to occur as  $M$  is increased

### **Shannon channel capacity:**

The maximum reliable transmission rate over an ideal channel with bandwidth  $W$  Hz, with Gaussian distributed noise, and with SNR  $S/N$  is

$$C = W \log_2 ( 1 + S/N ) \text{ bps}$$

- Reliable means error rate can be made arbitrarily small by proper coding

## Example

- Consider a 3 KHz channel with 8-level signaling. Compare bit rate to channel capacity at 20 dB SNR
- 3KHz telephone channel with 8 level signaling  
Bit rate =  $2 \times 3000$  pulses/sec \* 3 bits/pulse = 18 Kbps
- 20 dB SNR means  $10 \log_{10} S/N = 20$   
Implies  $S/N = 100$
- Shannon channel capacity is then  
 $C = 3000 \log (1 + 100) = 19,974$  bps

## Bit rates of digital transmission systems

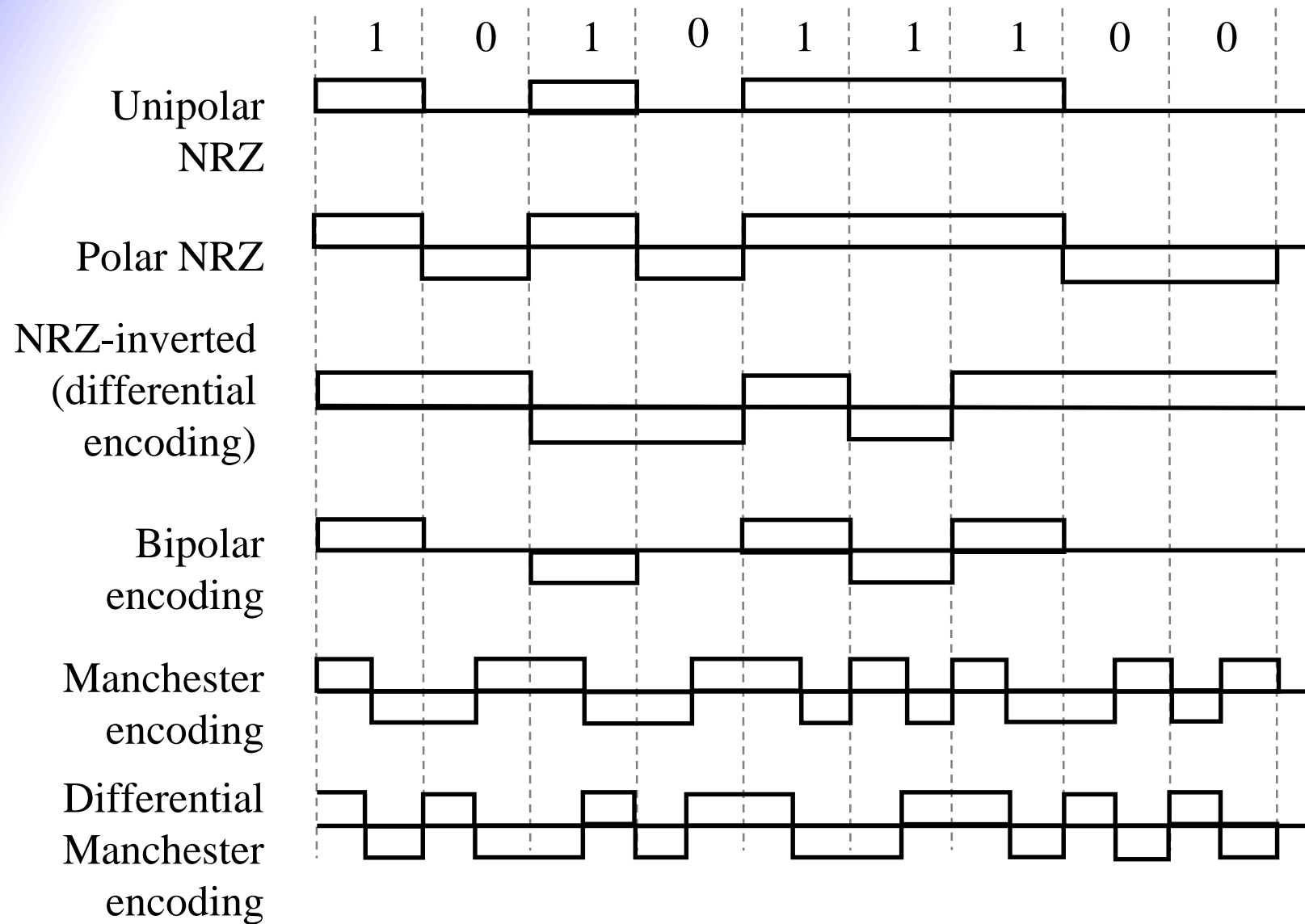
System	Bit Rate	Observations
Telephone twisted pair	33.6-56 Kbps	4 KHz telephone channel
Ethernet twisted pair	10 Mbps, 100 Mbps	100 meters of unshielded twisted copper wire pair
Cable modem	500 Kbps-4 Mbps	Shared CATV return channel
ADSL twisted pair	64-640 Kbps in, 1.536-6.144 Mbps out	Coexists with analog telephone signal
2.4 GHz radio	2-11 Mbps	IEEE 802.11 wireless LAN
28 GHz radio	1.5-45 Mbps	5 Km multipoint radio
Optical fiber	2.5-10 Gbps	1 wavelength
Optical fiber	>1600 Gbps	Many wavelengths



# Line coding

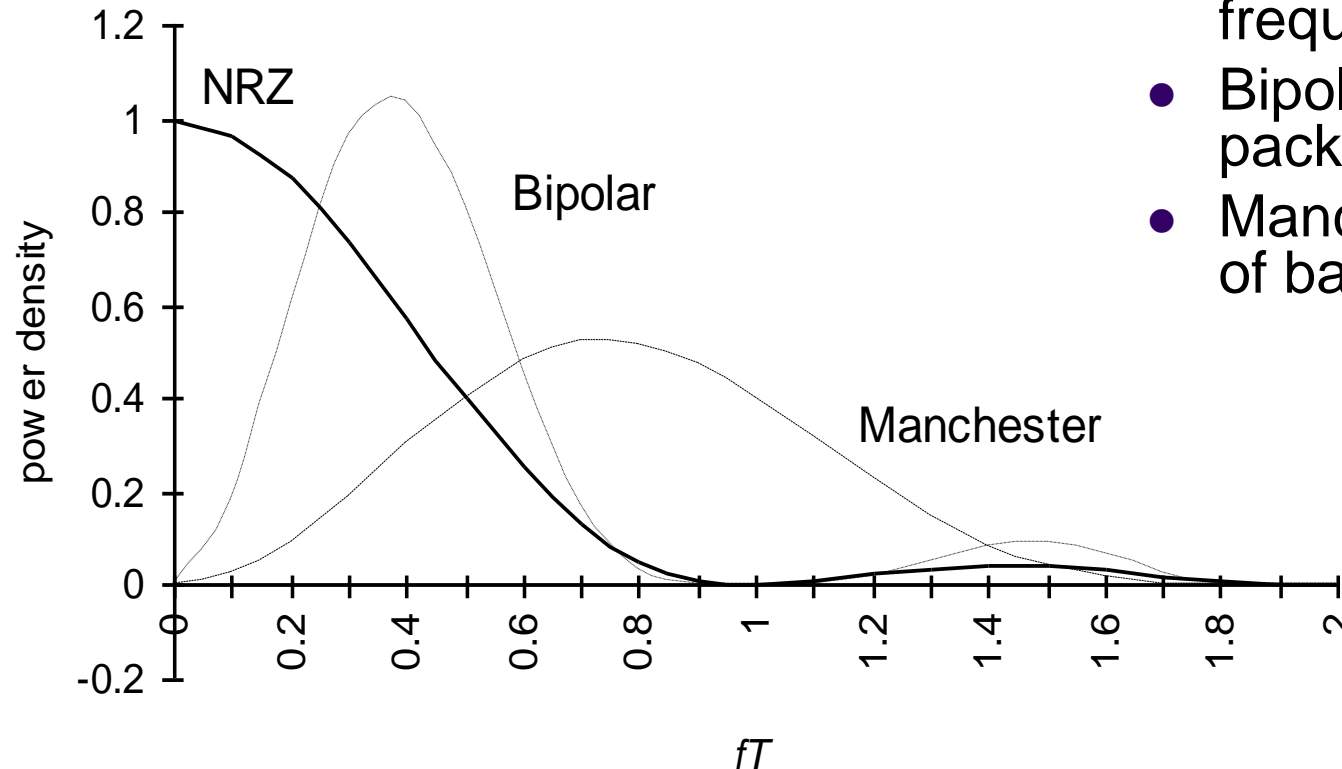
- Mapping of binary information sequence into the digital signal that enters the channel
  - Example - “1” maps to +A square pulse; “0” to −A pulse
- Line code selected to meet system requirements:
  - *Transmitted power*: Power consumption = \$
  - *Bit timing*: Transitions in signal help timing recovery
  - *Bandwidth efficiency*: Excessive transitions wastes bandwidth
  - *Low frequency content*: Some channels block low frequencies
    - Waveform should not have low-frequency content
  - *Error detection*: Ability to detect errors helps
  - *Complexity/cost*: Is code implementable in chip at high speed?

# Examples

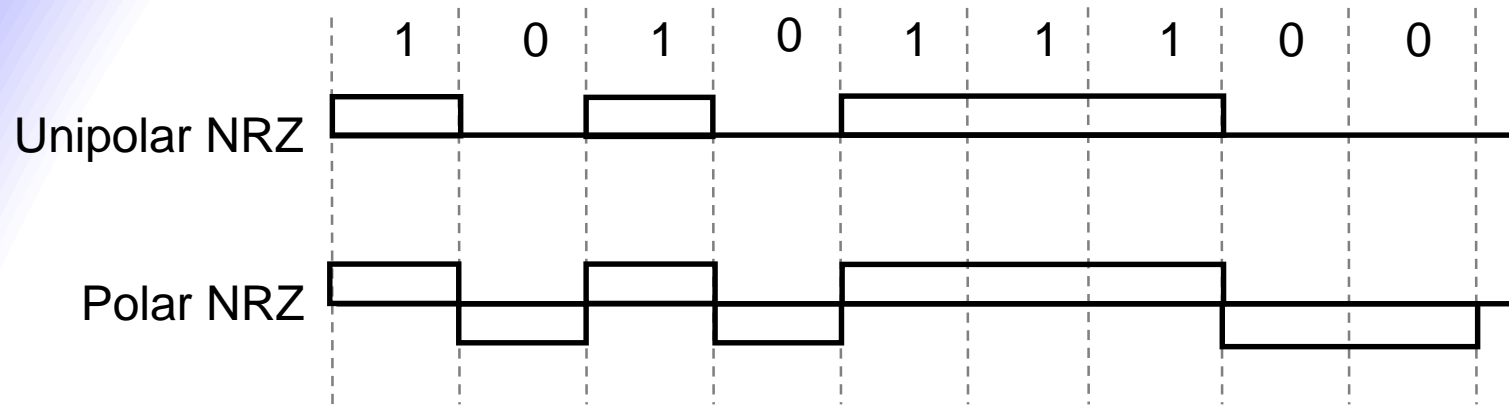


# Spectrum of line codes

- Assume 1s and 0s independent and equiprobable



- NRZ has high content at low frequencies
- Bipolar tightly packed around  $T/2$
- Manchester wasteful of bandwidth

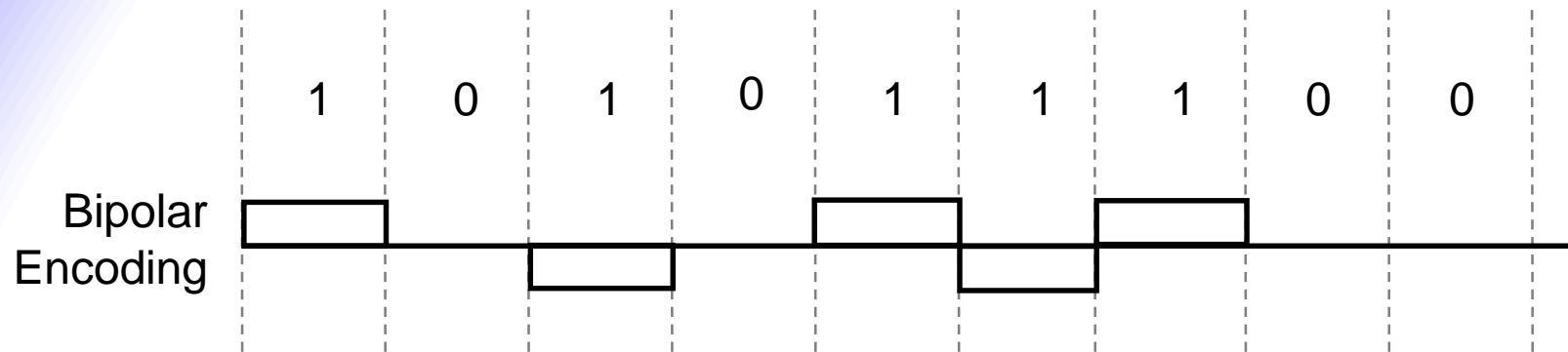


### Unipolar NRZ

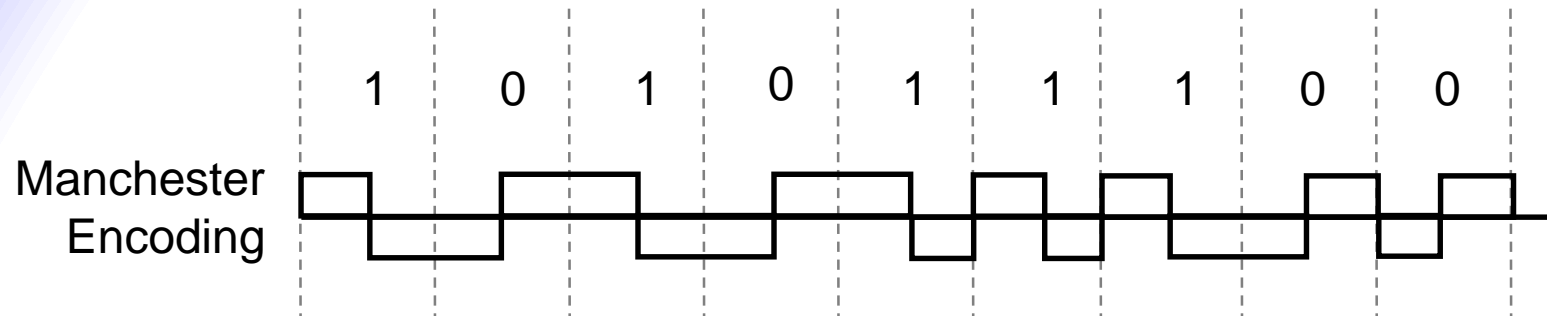
- “1” maps to +A pulse
- “0” maps to no pulse
- High Average Power  
 $0.5 \cdot A^2 + 0.5 \cdot 0^2 = A^2/2$
- Long strings of A or 0
  - Poor timing
  - Low-frequency content
- Simple

### Polar NRZ

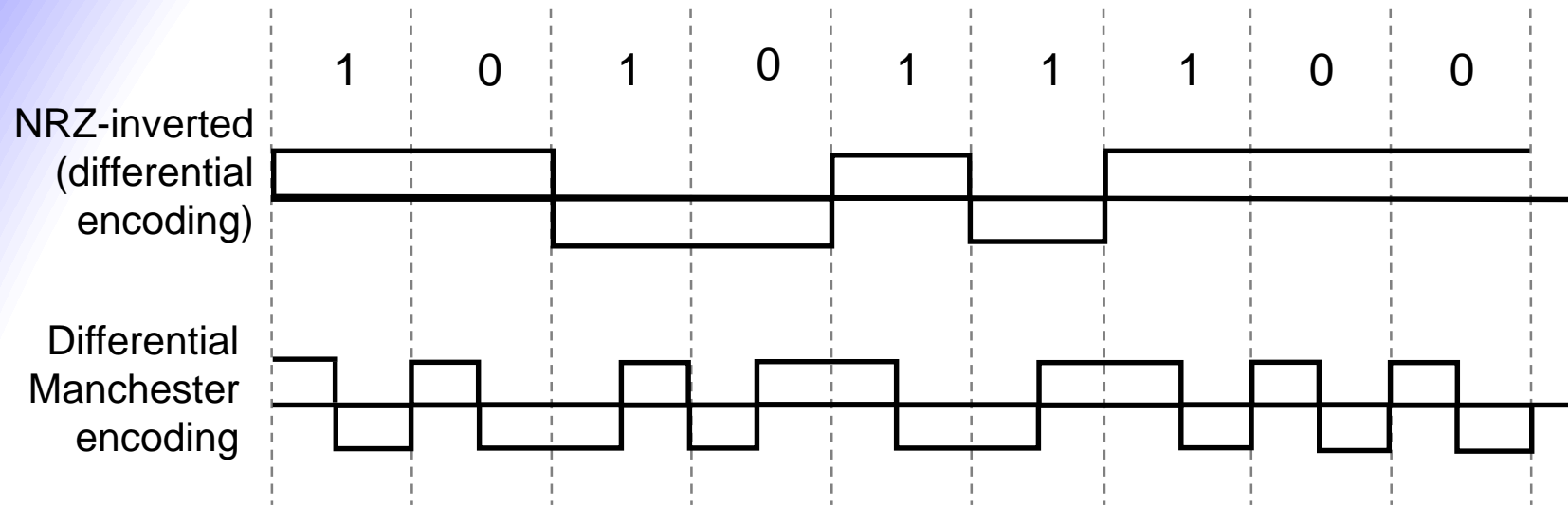
- “1” maps to +A/2 pulse
- “0” maps to -A/2 pulse
- Better Average Power  
 $0.5 \cdot (A/2)^2 + 0.5 \cdot (-A/2)^2 = A^2/4$
- Long strings of +A/2 or -A/2
  - Poor timing
  - Low-frequency content
- Simple



- Three signal levels:  $\{-A, 0, +A\}$
- “1” maps to  $+A$  or  $-A$  in alternation
- “0” maps to no pulse
  - Every +pulse matched by -pulse so little content at low frequencies
- String of 1s produces a square wave
  - Spectrum centered at  $T/2$
- Long string of 0s causes receiver to lose synchronization
- Zero-substitution codes



- “1” maps into  $A/2$  first  $T/2$ ,  $-A/2$  last  $T/2$
- “0” maps into  $-A/2$  first  $T/2$ ,  $A/2$  last  $T/2$
- Every interval has transition in middle
  - Timing recovery easy
  - Uses double the minimum bandwidth
- Simple to implement
- Used in 10-Mbps Ethernet and other LAN standards
- $mBnB$  line code
- Maps block of  $m$  bits into  $n$  bits
- Manchester code is 1B2B code
- 4B5B code used in FDDI LAN
- 8B10B code used in Gigabit Ethernet
- 64B66B code used in 10G Ethernet

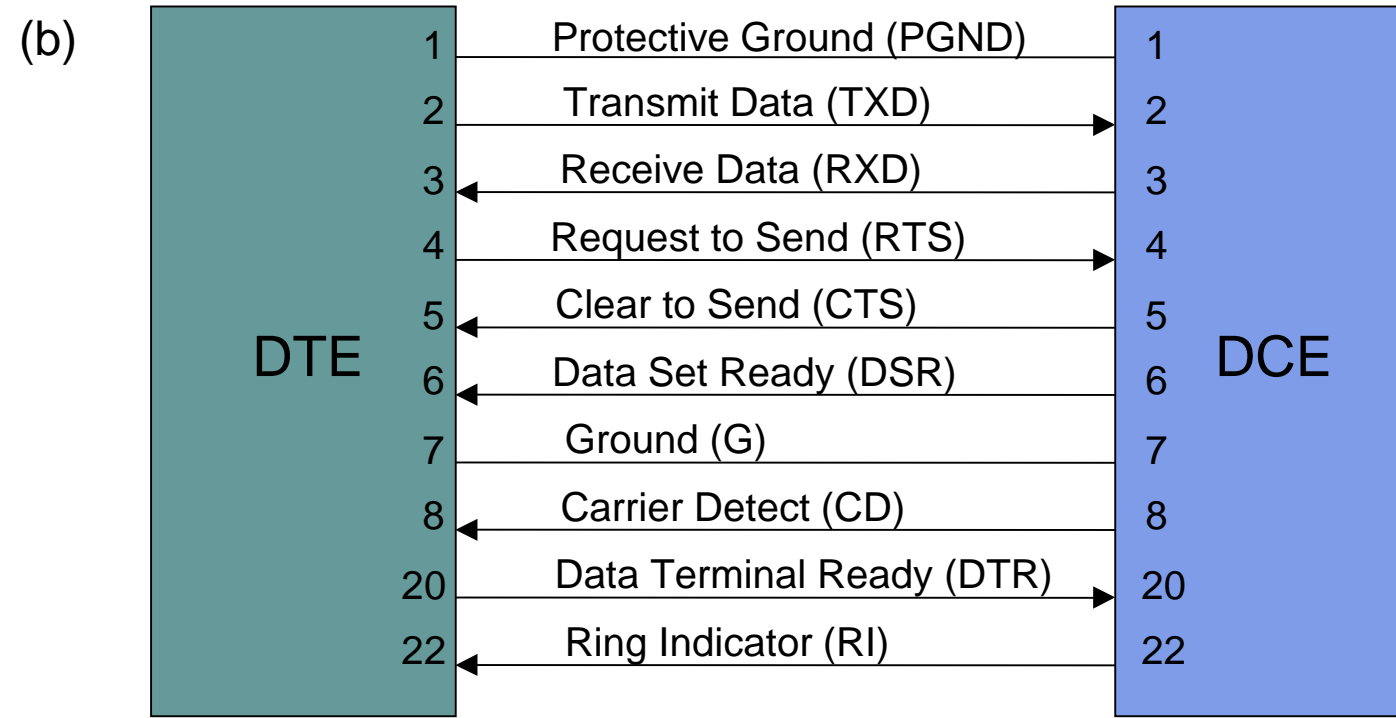
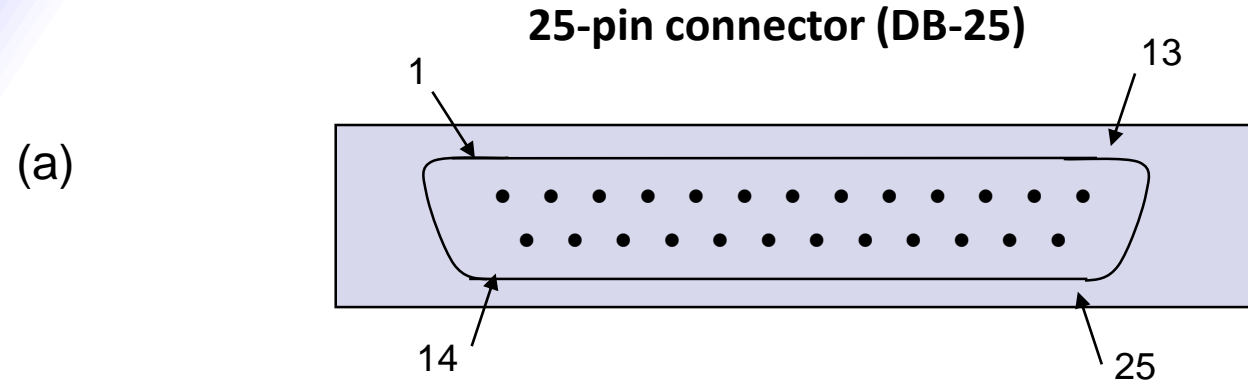


- Errors in some systems cause transposition in polarity,  $+A$  become  $-A$  and vice versa
  - All subsequent bits in Polar NRZ coding would be in error
- Differential line coding provides robustness to this type of error
- “1” mapped into transition in signal level
- “0” mapped into no transition in signal level
- Same spectrum as NRZ
- Errors occur in pairs
- Also used with Manchester coding

## 6.3 Asynchronous and synchronous communications

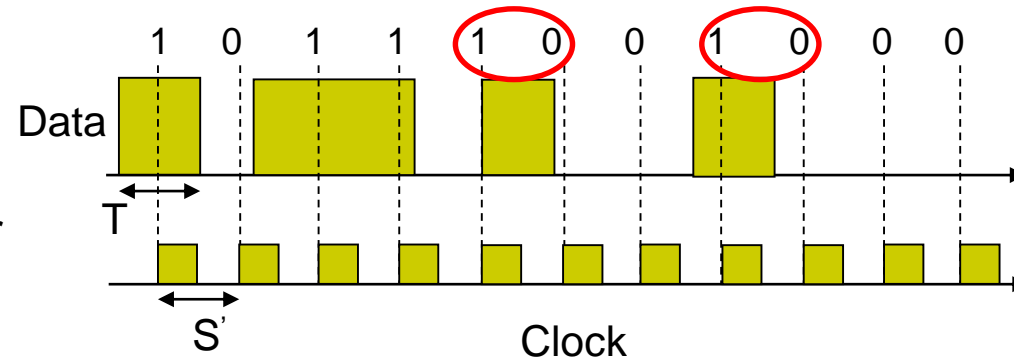
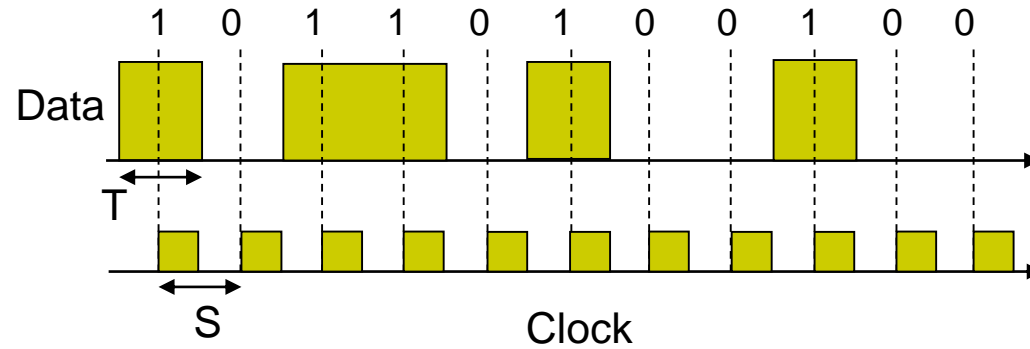
- Recommended Standard (RS) 232
- serial line interface between computer and modem or similar device
- Data Terminal Equipment (DTE): computer
- Data Communications Equipment (DCE): modem
- mechanical and electrical specification



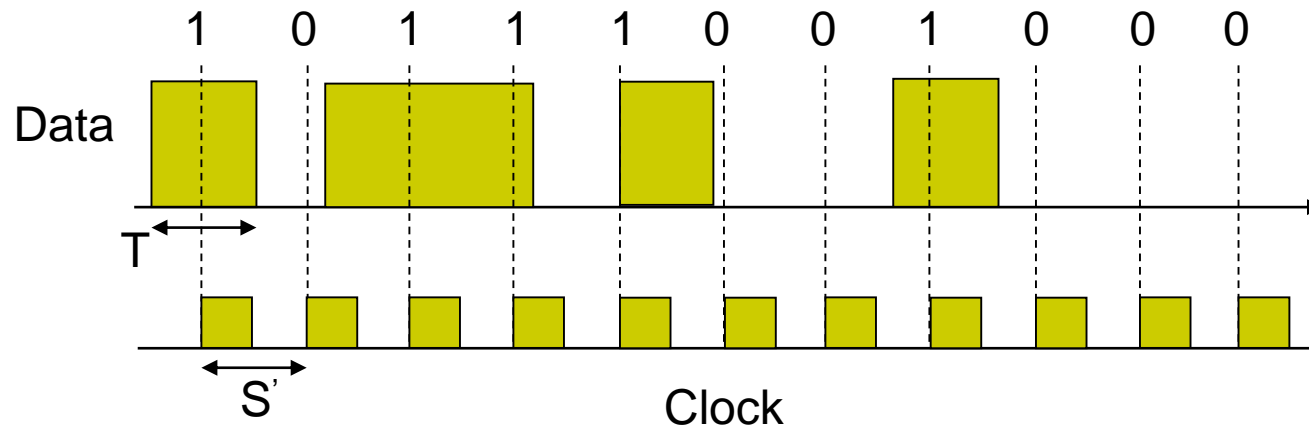


# Synchronization

- Synchronization of clocks in transmitters and receivers.
  - clock drift causes a loss of synchronization
- Example: assume '1' and '0' are represented by V volts and 0 volts respectively
  - Correct reception
  - Incorrect reception due to incorrect clock (slower clock)

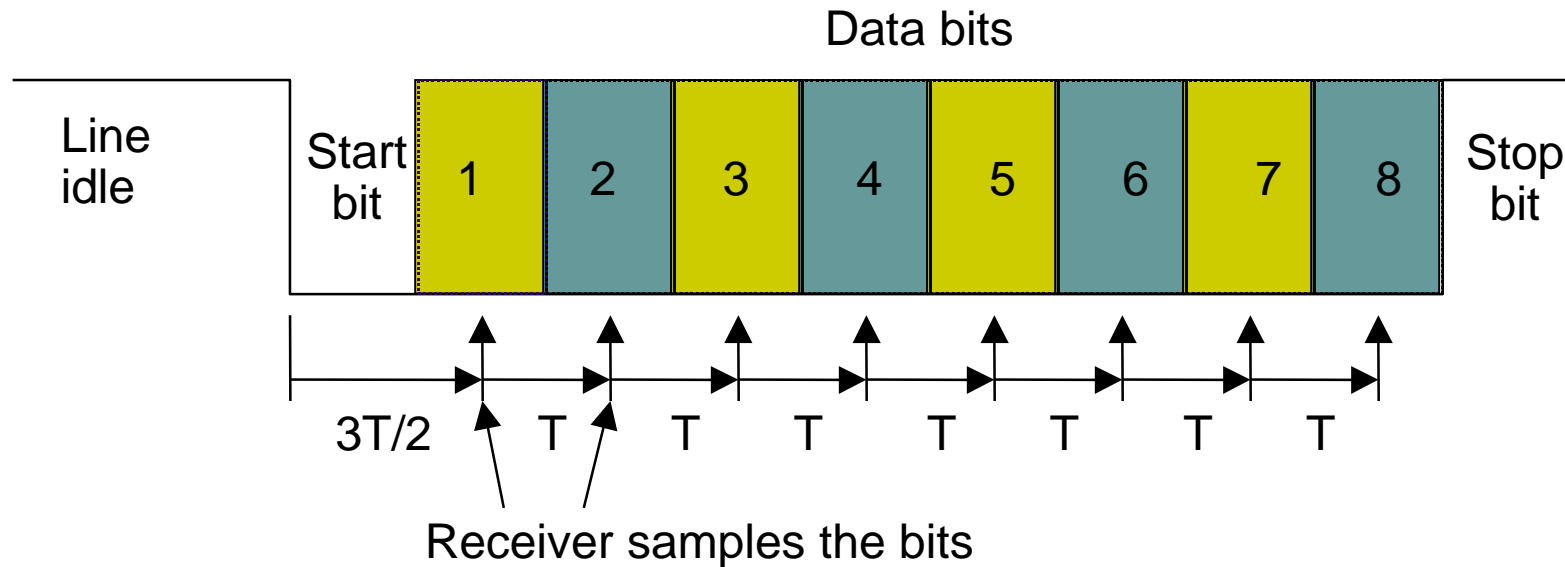


- Incorrect reception
- How to avoid a loss of synchronization?
  - Asynchronous transmission
  - Synchronous transmission



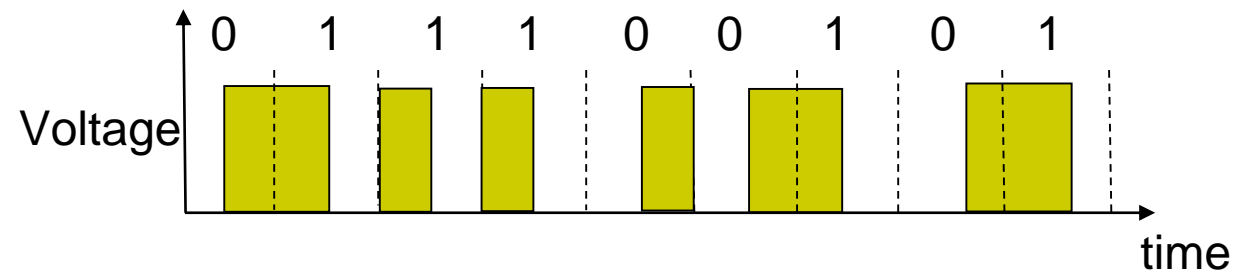
# Asynchronous transmission

- Avoids synchronization loss by specifying a short maximum length for the bit sequences and resetting the clock in the beginning of each bit sequence.

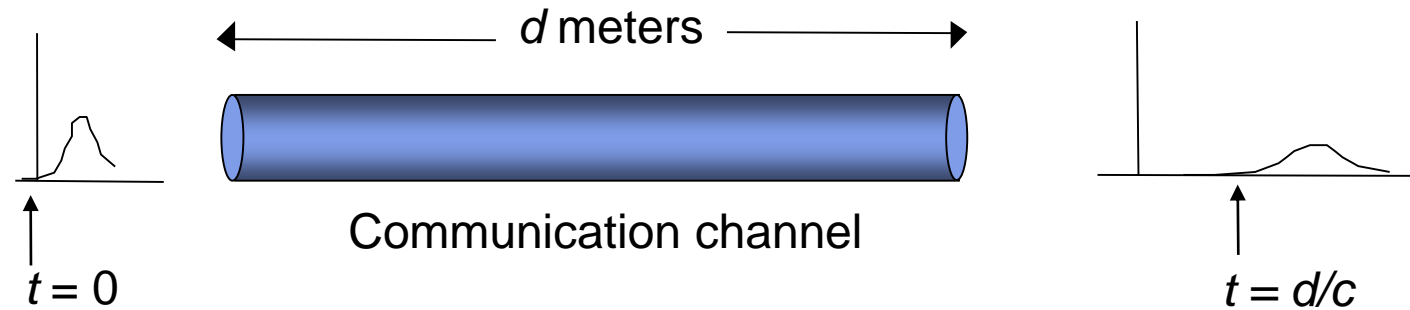


# Synchronous transmission

- Sequence contains data + clock information (line coding)
  - i.e. Manchester encoding, self-synchronizing codes, is used.
- R transition for  $R$  bits per second transmission

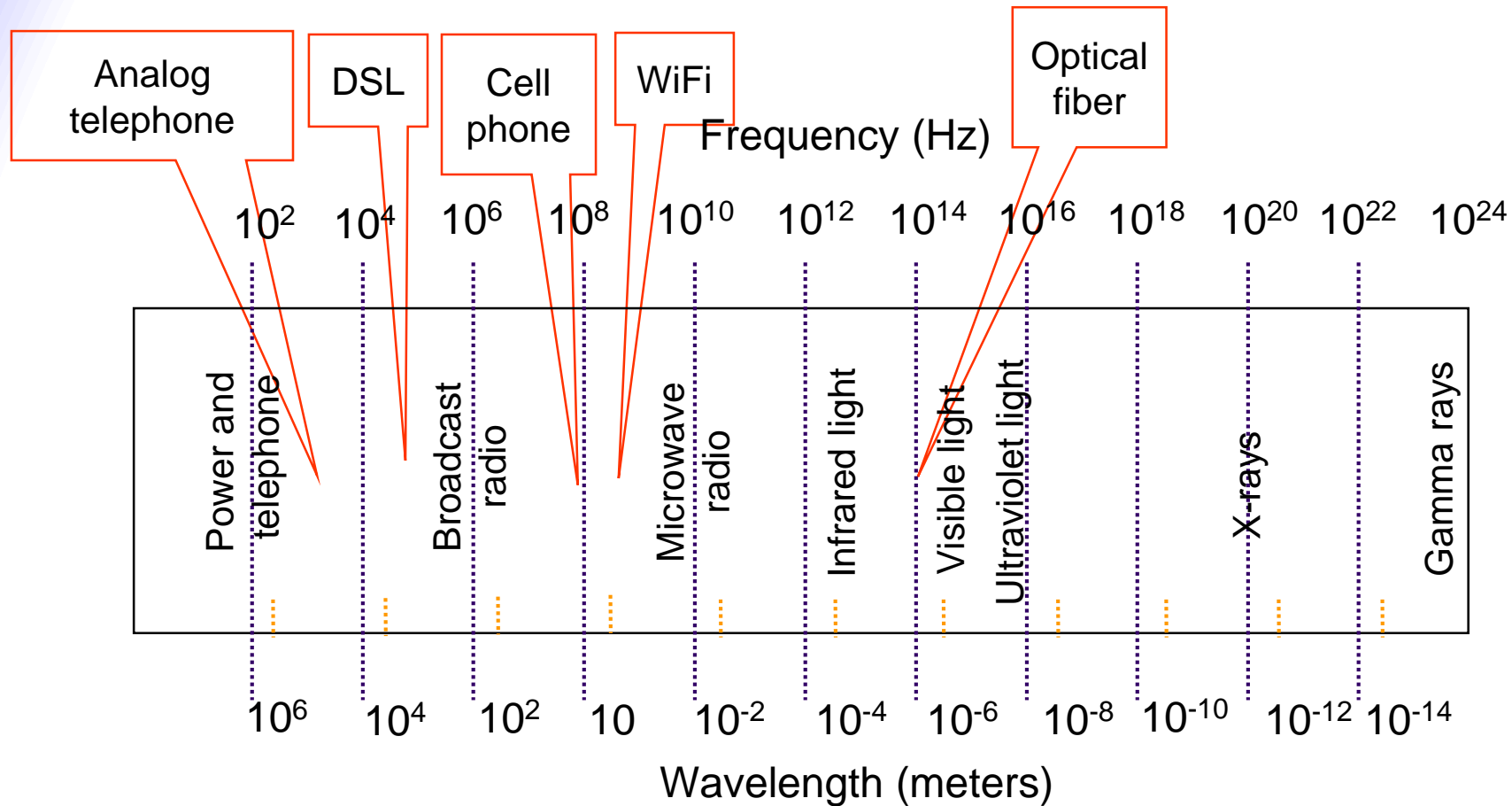


## 6.4 Transmission media



- Information bearing capacity
  - Amplitude response and bandwidth
    - depend on distance
  - Susceptibility to noise and interference
    - error rates and SNRs
- Propagation speed of signal
  - $c = 3 \times 10^8$  m/s in vacuum
  - $n = c/\sqrt{e}$  speed of light in medium where  $e > 1$  is the dielectric constant of the medium
  - $n = 2.3 \times 10^8$  m/s in copper wire;  $n = 2.0 \times 10^8$  m/s in optical fiber

- frequency of communication signals



## Wireless Media

- Signal energy propagates in space, limited directionality
- Interference possible, so spectrum regulated
- Limited bandwidth
- Simple infrastructure: antennas & transmitters
- No physical connection between network & user
- Users can move

## Wired Media

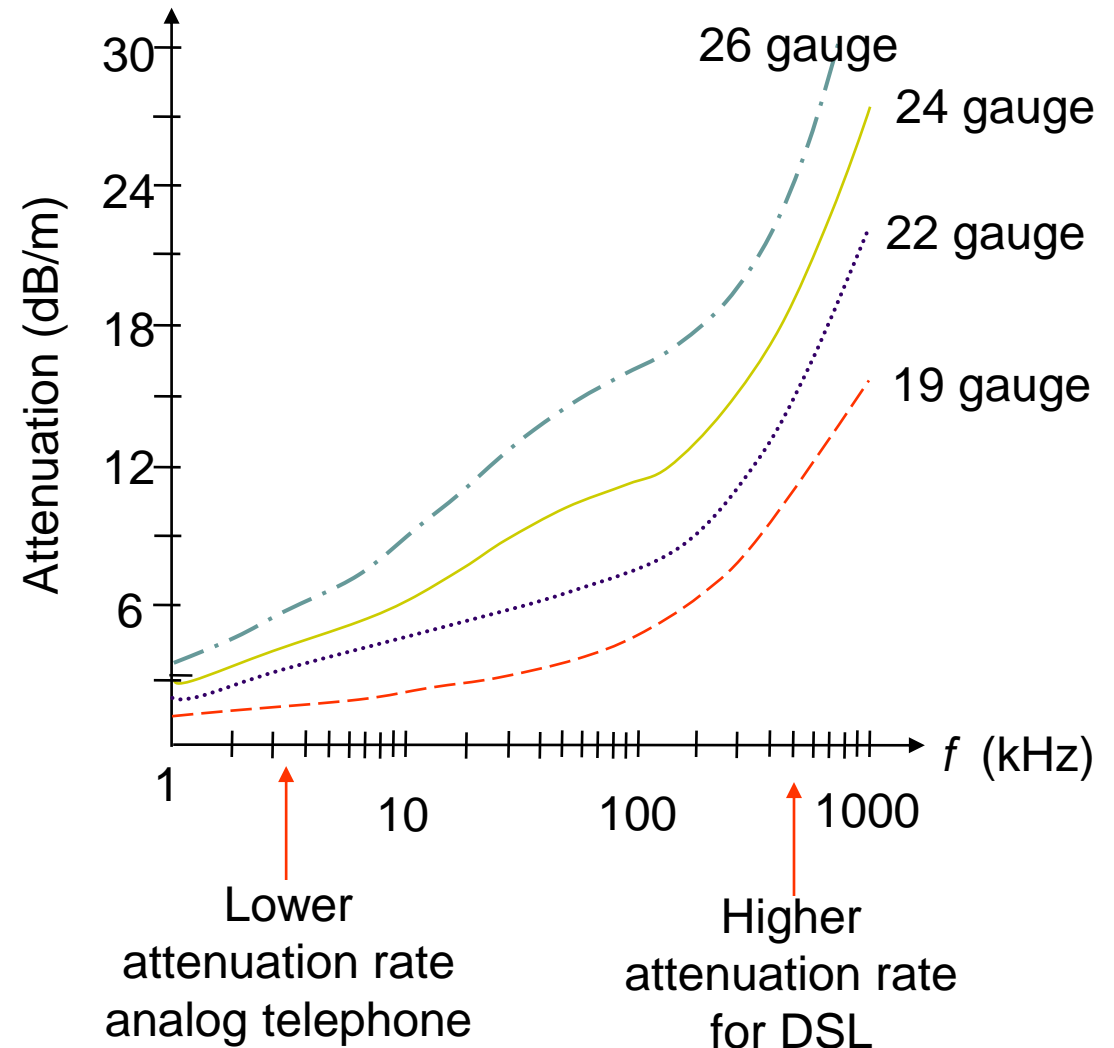
- Signal energy contained & guided within medium
- Spectrum can be re-used in separate media (wires or cables), more scalable
- Extremely high bandwidth
- Complex infrastructure: ducts, conduits, poles, right-of-way



- Attenuation varies with media
  - Dependence on distance of central importance
- Wired media has exponential dependence
  - Received power at  $d$  meters proportional to  $10^{-kd}$
  - Attenuation in dB =  $k d$ , where  $k$  is dB/meter
- Wireless media has logarithmic dependence
  - Received power at  $d$  meters proportional to  $d^{-n}$
  - Attenuation in dB =  $n \log d$ , where  $n$  is path loss exponent;  $n=2$  in free space
  - Signal level maintained for much longer distances
  - Space communications possible

# Twisted pair

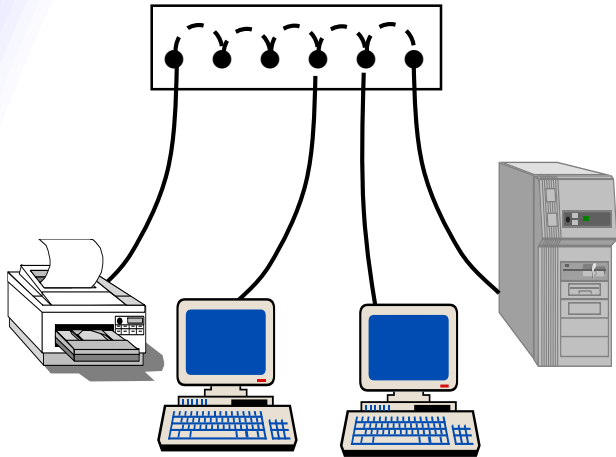
- Two insulated copper wires arranged in a regular spiral pattern to minimize interference
- Various thicknesses, e.g. 0.016 inch (24 gauge)
- Low cost
- Telephone subscriber loop from customer to CO
- Old trunk plant connecting telephone COs
- Intra-building telephone from wiring closet to desktop
- In old installations, loading coils added to improve quality in 3 KHz band, but more attenuation at higher frequencies



Standard	Data Rate	Distance
T-1	1.544 Mbps	18,000 feet, 5.5 Km
DS2	6.312 Mbps	12,000 feet, 3.7 Km
1/4 STS-1	12.960 Mbps	4500 feet, 1.4 Km
1/2 STS-1	25.920 Mbps	3000 feet, 0.9 Km
STS-1	51.840 Mbps	1000 feet, 300 m

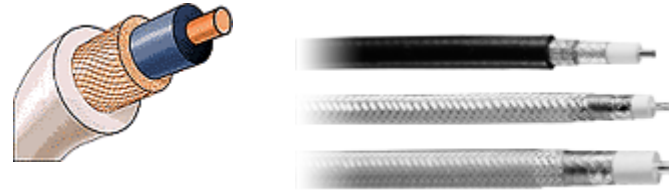
- Twisted pairs can provide high bit rates at short distances
- Asymmetric Digital Subscriber Loop (ADSL)
  - High-speed Internet Access
  - Lower 3 KHz for voice
  - Upper band for data
  - 64 Kbps inbound
  - 640 Kbps outbound
- Much higher rates possible at shorter distances
  - Strategy for telephone companies is to bring fiber close to home and then twisted pair
  - Higher-speed access + video

# Ethernet LAN

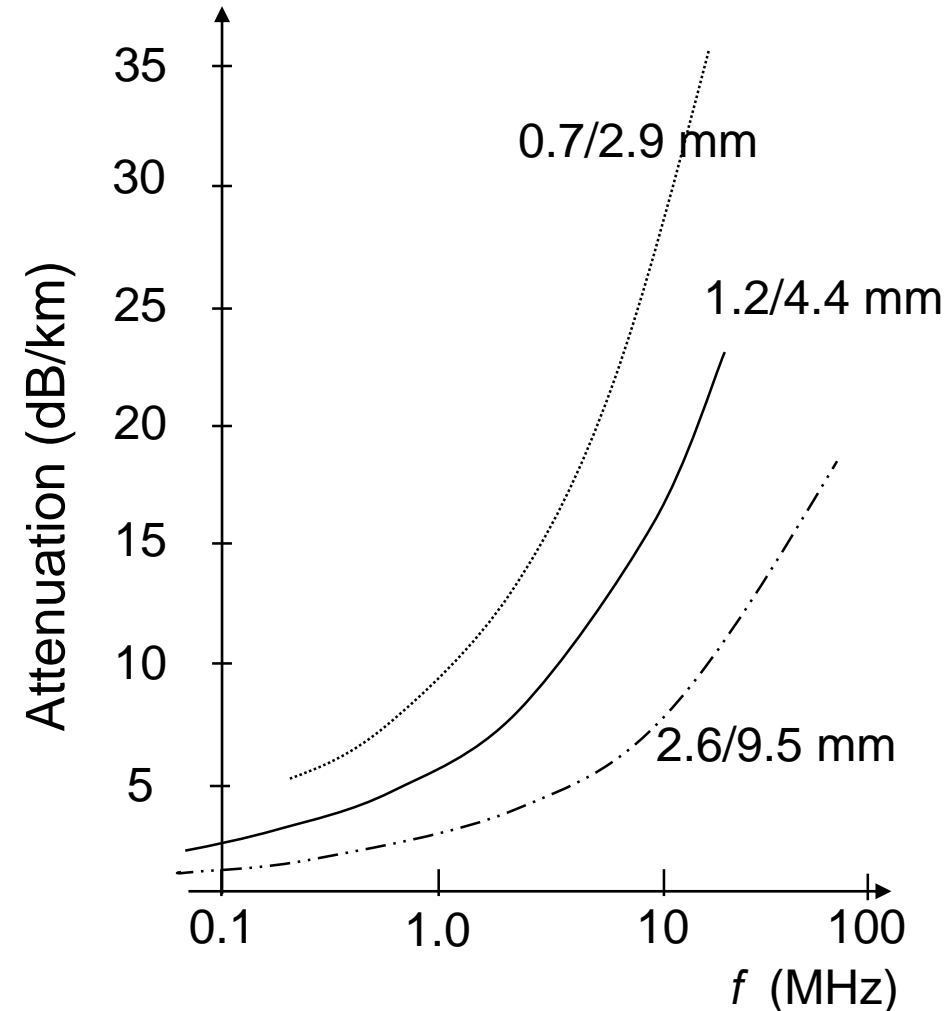


- Category 3 unshielded twisted pair (UTP): ordinary telephone wires
- Category 5 UTP: tighter twisting to improve signal quality
- Shielded twisted pair (STP): to minimize interference; costly
- 10BASE-T Ethernet
  - 10 Mbps, Baseband, Twisted pair
  - Two Cat3 pairs
  - Manchester coding, 100 meters
- 100BASE-T4 *Fast* Ethernet
  - 100 Mbps, Baseband, Twisted pair
  - Four Cat3 pairs
  - Three pairs for one direction at-a-time
  - 100/3 Mbps per pair;
  - 3B6T line code, 100 meters
- Cat5 and STP provide other options

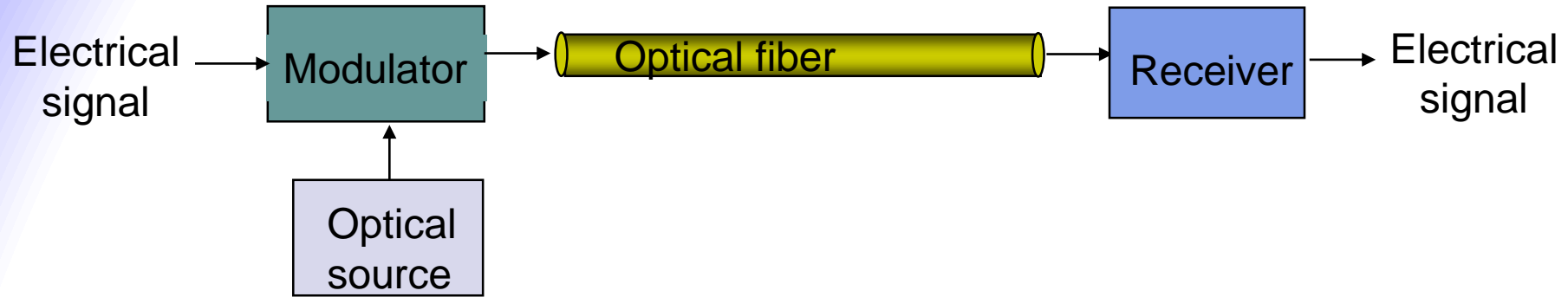
# Coaxial cable



- Cylindrical braided outer conductor surrounds insulated inner wire conductor
- High interference immunity
- Higher bandwidth than twisted pair
- Hundreds of MHz
- Cable TV distribution
- Long distance telephone transmission
- Original Ethernet LAN medium

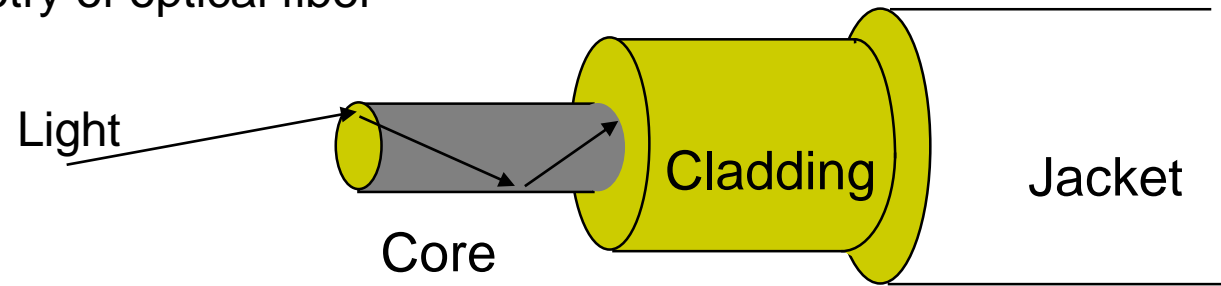


# Optical fiber

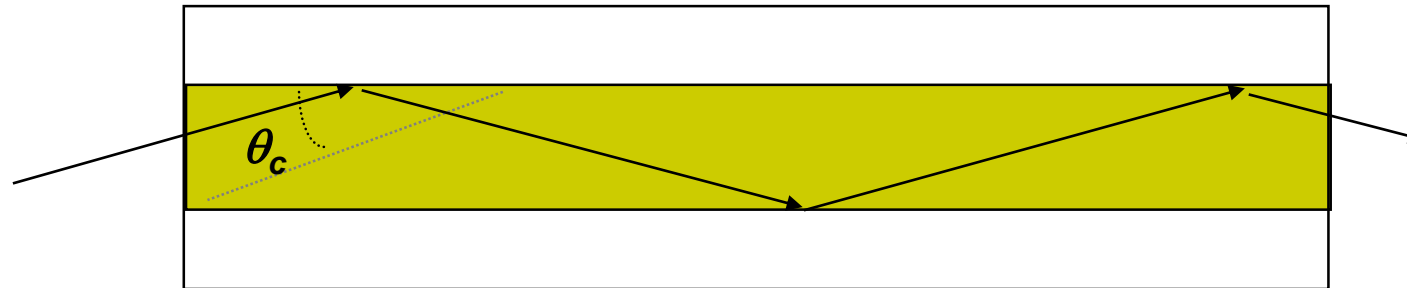


- Light sources (lasers, LEDs) generate pulses of light that are transmitted on optical fiber
  - Very long distances (>1000 Km)
  - Very high speeds (>40 Gbps/wavelength)
  - Nearly error-free (BER of  $10^{-15}$ )
- Profound influence on network architecture
  - Dominates long distance transmission
  - Distance less of a cost factor in communications
  - Plentiful bandwidth for new services

### Geometry of optical fiber

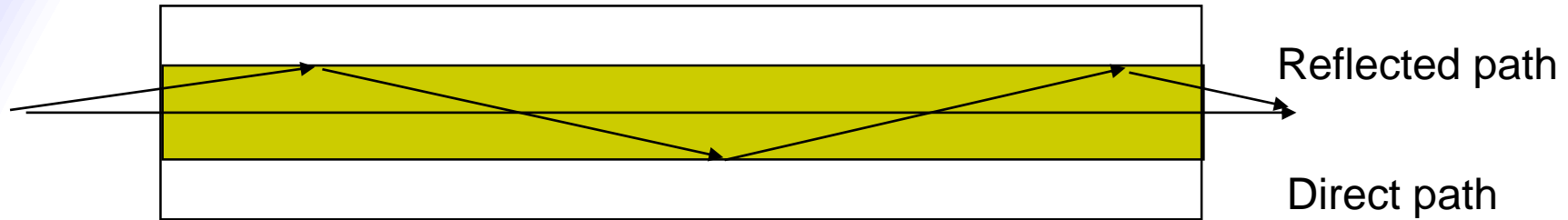


### Total Internal Reflection in optical fiber

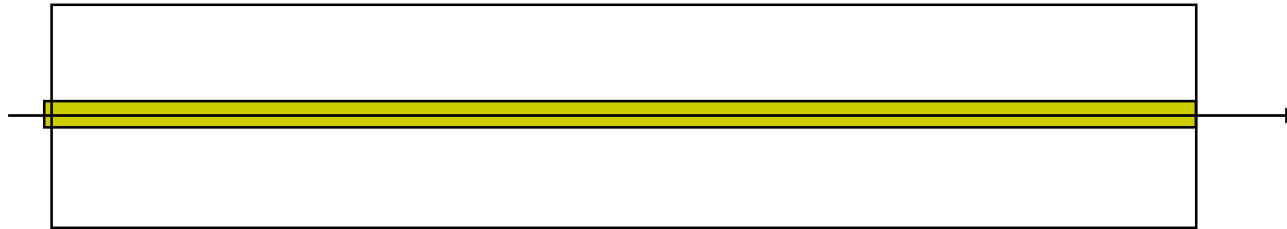


- Very fine glass cylindrical core surrounded by concentric layer of glass (cladding)
- Core has higher index of refraction than cladding
- Light rays incident at less than critical angle  $\theta_c$  is completely reflected back into the core

Multimode fiber: multiple rays follow different paths



Single-mode fiber: only direct path propagates in fiber

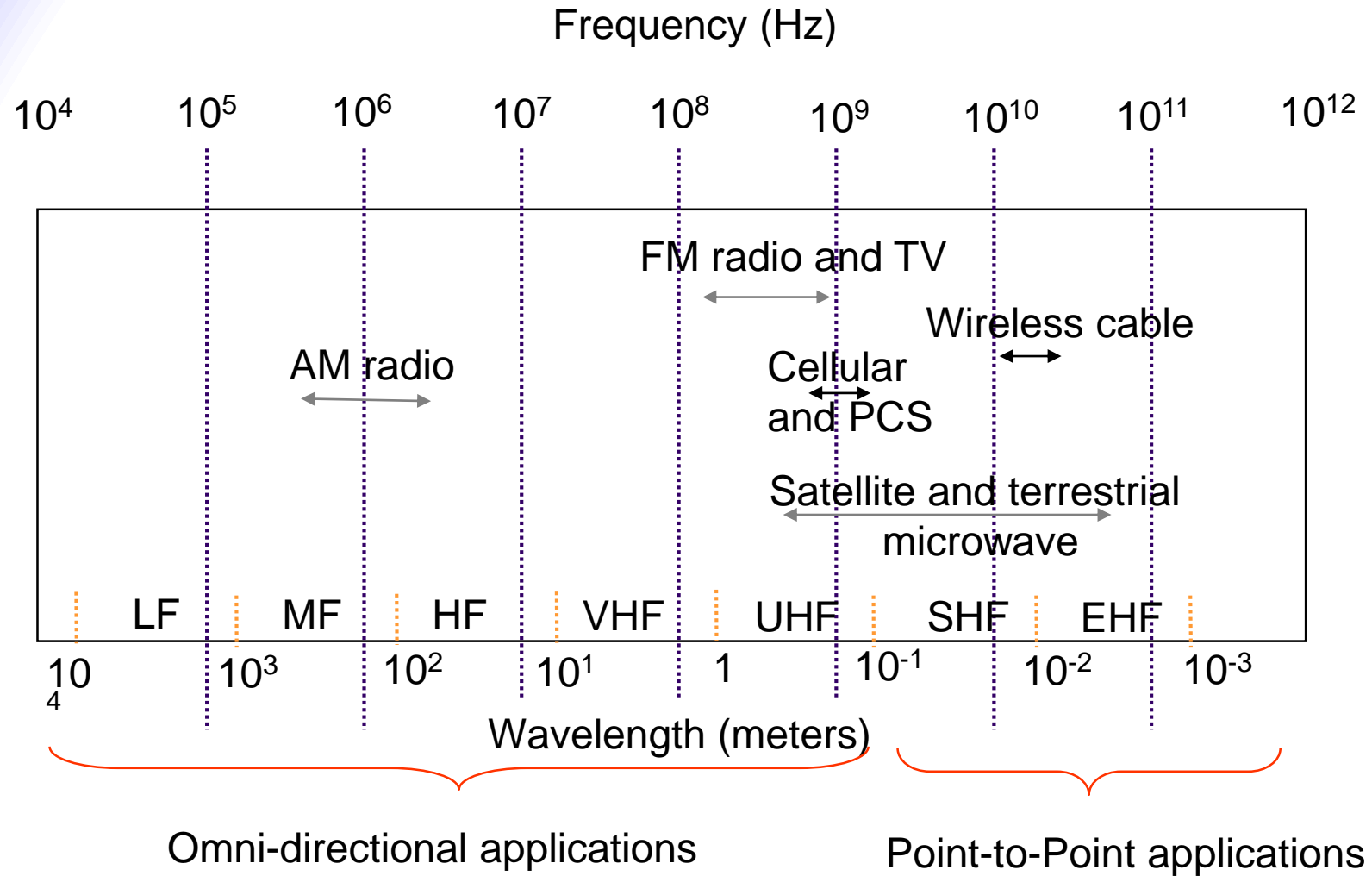


- Multimode: thicker core, shorter reach
  - Rays on different paths interfere causing dispersion and limiting bit rate
- Single mode: very thin core supports only one mode (path)
  - More expensive lasers, but achieves very high speeds



# Radio transmission

- Radio signals: antenna transmits sinusoidal signal (“carrier”) that radiates in air/space
- Information embedded in carrier signal using modulation, e.g. Quadrature Amplitude Modulation (QAM)
- Communications without tethering
  - cellular phones, satellite transmissions, wireless LANs
- Multipath propagation causes fading
- Interference from other users
- Spectrum regulated by national and international regulatory organizations



# Examples

## Cellular Phone

- Allocated spectrum
- First generation:
  - 800, 900 MHz
  - Initially analog voice
- Second generation:
  - 1800-1900 MHz
  - Digital voice, messaging

## Wireless LAN

- Unlicensed ISM spectrum
  - Industrial, Scientific, Medical
  - 902-928 MHz, 2.400-2.4835 GHz, 5.725-5.850 GHz
- IEEE 802.11 LAN standard
  - 11-54 Mbps

## Satellite Communications

- Geostationary satellite @ 36,000 Km above equator
- Relays microwave signals from uplink frequency to downlink frequency
- Long distance telephone
- Satellite TV broadcast

## Chapter Summary

- ◆ digital representation of information
- ◆ digital transmission – digitization of analog signal, characterization of communication channel, limitations, line coding
- ◆ asynchronous and synchronous communications
- ◆ transmission media

# Reference

Chapter 3, Communication  
Networks: Fundamental Concepts  
and Key Architectures

