## **Tutorial 9:**

Question 1: Let X=aabbacab and Y=baabcbb. Find the shortest common super-sequence for X and Y. (Backtracking process is required.)

## Solution:

	Y j=0	b j=1	a j=2	а	b	С	b	b
X i=0	0	← 1	← 2	← 3	← 4	← 5	← 6	← 7
a i=1	↑1	↑2	<b>₹</b> 2	₹3	↑4	↑5	16	<b>↑</b> 7
a i=2	↑2	↑3	₹3	₹3	↑4	<b>↑</b> 5	↑6	↑7
b	↑3	<b>^</b> 3	↑4	↑4	<b>~ 4</b>	← 5	٦6	<b>5</b> 7
b	↑4	<b>~ 4</b>	<b>↑</b> 5	↑5	₹ 5	← 6	<b>5</b> 6	<b>5</b> 7
а	<b>↑</b> 5	↑5	₹ 5	٨6	<b>↑</b> 6	← 7	↑7	18 ↑
С	<b>16</b>	16	<b>16</b>	↑7	↑7	<b>5</b> 7	18	↑9
а	↑7	↑7	<b>5</b> 7	<b>5</b> 7	↑8	18	↑9	↑ 10
b	18	₹8	18	18	<b>5</b> 8	↑9	<b>5</b> 9	₹ 10

Backtracking:  $b \to a \to c \to a \to b \to c \to b \to a \to a \to b$ 

 $SCS = \{baabcbacab\}$ 

## Question 2:

**Input:** An array A[1..n] of n integers (positive or negative).

**Task:** Use dynamic method to find a non-empty interval [i, j] such that A[i]+A[i+1]+...+A[j] is maximized.

Example: Given an array: -1, 2 -3, 4, 5, -1

The sum of interval [1,1]=-1, [1,2]=-1+2=1, [3, 5]=-3+4+5=6.

Hint: Let d[i] be the cost of the max sum of intervals ending at position i.

That is, d[i]=max {sum[1,i], sum[2, i], ..., sum[I,i]}.

Find recursive equation and use it to design a DP algorithm.

The final solution is the subinterval with the maximal d value.

In this example, d[1]=sum of [1,1]=-1, d[2]=sum of [2,2]=2. d[3]=sum of [2,3]=-1, d[4]=sum of [4,4]=4.

Answer:

$$d(i) = \begin{cases} A[i] & \text{if } i = 1\\ d[i-1] + A[i] & \text{if } d[i-1] > 0\\ A[i] & \text{if } d[i-1] \le 0 \end{cases}$$

Alg:

Phase 1

d(1):=A[i]
For i=2 to n do:
If d(i-1)>0,
 d(i)=d(i-1)+A[i], B[i]=1,

/\* containing A[i] and optimal interval ending at i-1. Otherwise, d(i)=A[i], B[i]=0,

/\* the optimal interval ending at i contains only A[i]. //\* B for backtracking.

Phase 2: Find j with the maximal d value. (It can also be done in phase 1)

Phase 3: Backtracking:

i=j, while(i>1 & B(i)=1)

j=j-1,

The optimal interval is [A[i],..., A[j]].