# XS Grammar

Syntax for production rules is:

 $A \to B\alpha$ 

The following represent non terminal symbols: - uppercase letters

The following represent terminal symbols: - lowercase words/letters - greek symbols - keywords

## 1. Program

## 1.1. Literals

 $NUM \rightarrow INT \mid FLT$ 

 $VCT \rightarrow vector(NUM, NUM, NUM);$ 

## 1.2. Top Level Statements

 $X \to I X \mid P$ 

 $P \to RL~P \mid FN~P \mid V_{top}~P \mid \epsilon$ 

where

X := XS script

I := Include statement

P := Program

RL := Rule def

FN := Function def

 $V_{top} := Top level variable definition$ 

#### 1.3. Prelude

This consists of all the constants and functions described in

- 1. XS Constant Reference
- 2. XS Function Reference

# 2. Statement

 $\overrightarrow{S \rightarrow V_{decl}} \mid \overrightarrow{V_{def}} \mid \overrightarrow{V_{asgn}} \mid \overrightarrow{IE} \mid \overrightarrow{W} \mid \overrightarrow{F} \mid \overrightarrow{SC} \mid \overrightarrow{R} \mid \overrightarrow{BR} \mid \overrightarrow{CO} \mid \overrightarrow{LBL} \mid \overrightarrow{GT} \mid \overrightarrow{DBG} \mid \overrightarrow{BRPT} \mid \overrightarrow{DP} \mid \overrightarrow{DM} \mid \overrightarrow{DM}$ 

 $\bar{S} \to S \bar{S} \mid \epsilon$ 

 $B \rightarrow \{ \bar{S} \}$ 

 $BS \to B \mid S$ 

where

S := Statement

 $\bar{S} := Statements$ 

B := Bodv

BS := Body or statement

 $V_{decl} := Variable declaration$ 

 $V_{def} := Variable definition$ 

 $V_{asgn} := Variable Assignment$ 

IE := If (Else) statement

W := While loop

F := For loop

SC := Switch case

R := Return statement

BR := Break statement

CO := Continue statement

BRPT := Breakpoint

## 2.1. Top Level Var Def

 $V_{\mathrm{top}} 
ightarrow \mathtt{extern} \ V_{\mathrm{top}} \mid \mathtt{const} \ V_{\mathrm{top}} \mid \mathtt{static} \ V_{\mathrm{top}}$ 

 $V_{\rm top} \to {\rm DTYPE~ID}$  = LIT;

 $DTYPE \rightarrow \mathtt{int} \ | \ \mathtt{float} \ | \ \mathtt{bool} \ | \ \mathtt{string} \ | \ \mathtt{vect}$ 

 $\mathrm{LIT} \rightarrow \mathrm{INT} \mid \mathrm{FLT} \mid \mathrm{STR} \mid \mathrm{VCT} \mid \mathrm{BOOL}$ 

Note: XS currently has bugs with defining top level strings and vectors

where

DTYPE := Datatype

ID := Identifier

LIT := Literal

## 2.2. Var Decl

 $V_{
m decl} 
ightarrow { t static} \ V_{
m decl}$ 

 $V_{
m decl} 
ightarrow {
m DTYPE~ID}$  ;

# 2.3. Var Def

 $V_{\mathrm{def}} \to \mathtt{const}\ V_{\mathrm{def}}$ 

 $V_{def} \rightarrow DTYPE ID = E;$ 

## 2.4. Var Assign

 $V_{asgn} o ID$  = E;

where

E := Expression

### **2.5.** If Else

 $\mathrm{IE} \to \mathtt{if}$  (  $\mathrm{E}$  ) BS ELSE

 $\mathrm{ELSE} 
ightarrow \mathtt{else} \ \mathrm{BS} \mid \epsilon$ 

where

ELSE := Else branch

## 2.6. While

 $W \to \mathtt{while}$  ( E ) BS

## 2.7. For

 $F \to \mbox{for (} V_{\rm asgn} \ {\rm OP_{\rm rel}} \ {\rm INT}$  )  ${\rm BS}$ 

where

 $OP_{rel} := Relational operators$ 

#### 2.8. Switch

 $SC \rightarrow switch (E) \overline{\{CASES\}}$ 

 $\text{CASES} \rightarrow \text{CASE CASES} \mid \text{DEFAULT CASES}_{\text{no default}} \mid \epsilon$ 

 $\mathrm{CASES_{no~default}} \rightarrow \mathrm{CASE~CASES_{no~default}} ~|~ \epsilon$ 

 $\mathrm{CASE} \to \mathtt{case}\;\mathrm{LIT}\;\colon\mathrm{B}$ 

 $DEFAULT \rightarrow \texttt{default} \; : \; B$ 

#### where

CASES := Optional case statements with one optional default case

 $CASES_{no, default} := Optional case statements only$ 

CASE := case statement

DEFAULT := default statement

#### 2.9. Functions

 $\overline{FN} o \mathtt{extern} \ FN \mid \mathtt{mutable} \ FN$ 

 $\text{FN} \to \text{RTYPE ID (ARGS}_{\text{formal}}) \text{B}$ 

 $\mathrm{RTYPE} \to \mathtt{void} \mid \mathrm{DTYPE}$ 

 $ARGS_{formal} \rightarrow \overline{ARG \mid ARG, ARGS}_{formal}$ 

 $ARG \rightarrow DTYPE ID = LIT \mid \epsilon$ 

where

RTYPE := Return type

 $ARGS_c := Formal Arguments$ 

ARG := Argument

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2.10. Return
\mathrm{R} 	o \mathtt{return} ( \mathrm{E} );
\mathrm{R} 	o \mathtt{return};
2.11. Rules
\mathrm{RL} 	o \mathtt{rule} \ \mathrm{ID} \ \mathrm{RPS} \ \mathrm{B}
RPS \rightarrow RP RPS \mid \epsilon
\mathrm{RP} \rightarrow \mathrm{RP}_{\mathrm{act}} \mid \mathrm{RP}_{\mathrm{grp}} \mid \mathrm{RP}_{\mathrm{sfq}} \mid \mathrm{RP}_{\mathrm{xfq}} \mid \mathrm{RP}_{\mathrm{rim}} \mid \mathrm{RP}_{\mathrm{pty}}
\mathrm{RP}_\mathrm{act} \to \mathtt{active} | inactive
\mathrm{RP}_{\mathrm{grp}} 	o \mathtt{group} \; \mathrm{ID}
\mathrm{RP}_{\mathrm{sfq}} 	o \mathtt{minInterval} \ \mathrm{INT} \ | \ \mathtt{highFrequency}
\mathrm{RP}_{\mathrm{xfq}} 	o \mathtt{maxInterval} \ \mathrm{INT}
\mathrm{RP}_{\mathrm{rim}} 	o \mathtt{runImmediately}
\mathrm{RP}_{\mathrm{pty}} 	o \mathtt{priority} \; \mathrm{INT}
Note: only one of each parameter can be present in a rule def, this needs to be a linting time check,
unfortunately its not possible to bake it into the grammar
2.12. Postfix
\mathrm{DP} \to \mathrm{ID}\text{++}
\mathrm{DM} \to \mathrm{ID}\text{--}
2.13. Include
I \rightarrow \text{include } STR;
2.14. Break
\mathrm{BR} 	o \mathtt{break};
2.15. Continue
\mathrm{CO} \to \mathtt{continue};
```

2.16. Label Def LBL  $\rightarrow$  label ID;

#### 2.17. Goto

 $GT \rightarrow goto ID;$ 

## 2.18. Function Call (Statement)

FNCS  $\rightarrow$  FNC;

#### 2.19. Debug

 $DBG \rightarrow dbg \ ID;$ 

Note: I don't know what this does in XS, its valid syntax though.

## 2.20. Breakpoint

 $\mathrm{BRPT} \to \mathtt{breakpoint};$ 

Note: This will pause XS execution. I don't know if its possible to resume execution/if this keyword is useful.

#### 2.21. Class

 $CLS \rightarrow class ID \{ MEM_VARS \};$ 

MEM\_VARS  $\rightarrow$  DTYPE ID = E; MEM\_VARS |  $\epsilon$ 

Note: I don't know how to use classes in XS, its valid syntax though. The furthest I've gotten is declaring a class variable: ClsName obj;. Initialising it or accessing member values doesn't seem possible.

#### 2.22. Docstring

Todo

## 3. Expression

 $E7 \rightarrow LIT \mid ID \mid P \mid FNC$ 

 $\rm E6 \rightarrow \rm E7 \mid \rm E6 * \rm E7 \mid \rm E6$  / E7 | E6 % E7

 $E5 \rightarrow E6 \mid E5 + E6 \mid E5 - E6$ 

 $E4 \rightarrow E5 \mid E4 \triangleleft E5 \mid E4 \triangleright E5 \mid E4 \triangleright E5 \mid E4 \triangleleft E5$ 

 $E3 \rightarrow E4 \mid E3 == E4 \mid E3 != E4$ 

 $\mathrm{E2} \rightarrow \mathrm{E3} \mid \mathrm{E2}$  &&  $\mathrm{E3}$ 

 $E1 \rightarrow E2 \mid E1 \mid \mid E2$ 

 $E \to E1$ 

### 3.1. Parenthesis

 $P \rightarrow (E)$ 

# 3.2. Function Call (Expression)

 $\overline{FNC} \to \overline{ID(ARGS_{actual})}$ 

 $ARGS_{actual} \rightarrow ARGS' \mid \epsilon$ 

 $ARGS' \rightarrow E \mid E, ARGS'$ 

# where

 $ARGS_{actual} := Actual Arguments$ 

# 4. The Full Grammar