Shortest Path Algorithms Verification with Idris

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1 Introduction

Shortest path problems deal with finding the path with minimum distance value between two nodes in a given graph. One variation of shortest path problem is single-source shortest path problem, which focuses on finding the path with minimum distance value from one source to all other vertices within the graph. Dijkstra's and Bellman-Ford are the most renowned single-source shortest path algorithms, and are implemented by software concerning various fields in real-life applications, such as finding the shortest path in road map, or routing path with minimum cost in networks[6][7][8].

Given the importance of Dijkstra's and Bellman-Ford in real-life applications, we are interested in verifying the implementation of both algorithms. We will provide concrete implementations for both algorithms. Based on the specific implementation, we then define functions with precise type signatures which carry out specifications that should hold for the correct implementations of Dijkstra's and Bellman-Ford algorithms, for instance returning the minimum distance value from the source to each node in the graph. Having these functions type checked will then ensure the correctness of our algorithm implementation, and that any program with problematic implementation will fail to compile at the type checking level. Our implementation will use the Idris functional programming language, which embraces powerful tools and features that are siginificantly helpful in program verification.

Contribution

(To be finished.)

The structure of the paper is as follows. Section 2 describes the significance and value of algorithm verification, and reasons of choosing Idris as the language for verifying programs. Section 3 provides some background on Dijkstra's and Bellman-Ford algorithms, follows up by briefly introduction on the Idris functional programming language. Section 4 includes an overview of our verification program, including definition of key concepts, assumptions made by our program, and details on the pseudocode and theoretical proof of Dijkstra's and Bellman-Ford, which serves as important guideline in implementation our verification program. Section 5 covers more details of our verification program, including function type signatures and code of the proof for key lemmas. Section 6 is discussion of our work. Section 7 presents and compares related work, and section 8 gives a breif conclusion.

2 Motivation

Verifying the correctness of programs is important, however in most real-life applications, the correctness of software is never verified directly, rather, it relies on the correctness of the algorithms it implements. This raises an issue concerning the gap between the expected and actual behavior of programs, that theoretical proof of algorithms can never validate the actual behavior of programs. The significance and value of verification, therefore, lies on the fact that it allows us to verify programs themselves rather than the algorithms behind them.

Dijkstra's and Bellman-Ford algorithms are two of the most renowned and widely-applied shortest path algorithms, however existing resource on verifying both algorithms are relatively limited. In this thesis, we offer verifications for the implementations of both algorithm. In additional, we aim to present verification as a programming issue. We want to show that with certain programming languages, verifying the correctness of programs can be achieved with type checking, that if the program's correctness is not guaranteed, then our verification program will fail to be type checked.

Based on the above motivations, the Idris programming language is chosen over other verification tools and proof management systems. Idris is a functional programming language with dependent types, which allows programmers to provide more specification on function's behaviors in its type signature. As we plan to achieve verification with type checking, this feature of Idris can be significantly helpful as often times it is important to establish tight connection between functions and its input data in a verification program. In addition, Idris's compiler-supported interactive editing feature provides precise description of functions' behaviors according to their types, allowing programmer to use types as guidance for writing program, which offers considerable assistance during our implementation. Section 3 covers more backgrounds on the Idris programming language.

3 Background

3.1 Introduction of Idris

Idris is a general-purpose functional programming language with dependent types. Many aspects of Idris is influenced by Haskell and ML. Features of Idris include but not limit to dependent types, `with` rule, `case` expressions, lambda binding, as well as interactive editing.

Data Declaration

As an example of data declaraction in Idris, below shows the definition of natural numbers in Idris standard library:

```
  -- natural number can be either zero(Z) or plus one of another natural
number (S Nat)
  data Nat = Z | S Nat
```

Another syntax similar to that of GADT in Haskell for data declaration is also allowed:

```
-- declaration of List data type in Idris standard library
data List : (elem : Type) -> Type where
  Nil : List elem
```

```
(::) : (x : elem) -> (xs : List elem) -> List elem
```

Dependent Types

Dependent types are types that depend on elements of other types[2]. It allows programmers to specify certain properties of data types explicitly in their type signature. Consider the following definition of a vector data type, where `len` specifies the length of the vector, and `elem` is the element type.

```
-- declaration of Vect data type in Idris standard library
data Vect : (len : Nat) -> (elem : Type) -> Type where
  Nil  : Vect Z elem
  (::) : (x : elem) ->
         (xs : Vect len elem) ->
         Vect (S len) elem
```

The type `Vect len elem` is dependent on the value of type variables `len` and `elem`, specifying that `Vect` is a vector of length `len` containing element of type `elem`. With dependent types, programmers can ensure the behaviors of functions through their type signatures by defining more precise types. Consider the function `concat` below that concantenates two vectors. Given two input vectors V_1 and V_2, the output value of `concat` should be a vector of length $|V_1| + |V_2|$.

```
concat : Vect n Nat -> Vect m Nat -> Vect (n+m) Nat
```

The type siganture of `concat` establishes that the resulting vector of concatenating two vectors of length n and m must be of length $(n + m)$, otherwise `concat` will fail to type check.

Pattern Matching and Totality Checking

Pattern matching is the process of matching values against specific patterns. In Idris, functions are implemented by pattern matching on possible values of inputs. Continuing with the above example of `concat` function that concatenates two vectors, to define `concat`, we need to provide definitions on all possible values of `Vect`, which can either be `Nil`, i.e., a vector of length zero, or a non-empty vector of the pattern `(x :: xs)`.

```
concat : Vect n Nat -> Vect m Nat -> Vect (n+m) Nat
concat Nil v2 = v2
concat v1 Nil = v1
concat (x :: xs) v2 = x :: concat xs v2
```

Functions defined for all possible values of input are total functions, and are guaranteed to produce a result in finite time given well-typed inputs. Partial functions are not total, and hence might crash for some inputs. To secure the termination of programs, every function definition in Idris are checked for totality after type checking. Specifically, Idris decides whether a function terminates based on two aspects: first, function must be defined for all possible inputs; and second, if a function definition includes a recursive call, then there must be an argument that strictly decreases over each recursion, and converges towards a base case. An error or warning will be given for any function that fails totality checking.

3.2 Dijkstra's and Bellman-Ford algorithms

Dijkstra's Algorithm

Dijkstra's algorithm is a greedy algorithm that finds the shortest path from a given source to all other nodes in a directed graph with weighted edges. It was first introduced in 1959 by Edsger Wybe Dijkstra, and it is widely applied in many real-life applications, including shortest path finding in road map, or Internet routing protocols such as the Open Shortest Path First protocol.

Dijkstra's algorithm takes in a directed graph with non-negative edge weights, and computes the shortest path distance from one single source node to all other reachable nodes in the graph. The algorithm maintains a list of unexplored nodes and their distance values to the source node. Initially, the list of unexplored nodes contains all nodes in the input graph, and the distance value of all node are set as infinity except for the source node itself, which is set to zero. The algorithm extracts the node v with minimum distance value from the unexplored list during each iteration, and for each neighbor v' of v, if the path from source to v' via v contributes a smaller distance value, then the distance value of v' is updated.

Bellman-Ford Algorithm

Bellman-Ford algorithm was first introduced by Alfonso Shimbel in 1955, and was published by Richard Bellman and Lester Ford, Jr in 1958 and 1956 respectively. The algorithm solves the issue of calculating the minimum distance value from a single source to all other nodes in a given graph, and different from Dijkstra's algorithm, Bellman-Ford algorithm allows negative edge weights in the input graph, and is capable of detecting the existence of negative cycle(a cycle whose edge weights sum up to a negative value). Applications of Bellman-Ford includes routing protocols such as the Routing Information Protocol.

4 High-Level Contribution

4.1 Dijkstra's Algorithm

4.1.1 Data Structures

Dijkstra's algorithm requires non-negative edge weights and valid input graph, and the data structures in our implementation are designed to ensure these properties of input values. An overview of data structures in our implementation is presented below, and a detailed description is provided under Section 5.
Denote `gsize` as the size of graph, i.e. the number of vertices in a graph. A graph g is defined as a vector containing `gsize` number of adjacent lists, one for each node in the graph, and a node is defined as a data structure carrying a value of type `Fin gsize`. An adjacent list for a node $n \in g$ is defined as a list of tuples $(n', edge_w)$, where the first element n' in each tuple is a neighbor of n in g, and the second element $edge_w$ is the weight of the edge (n, n') in g. To access the adjacent list for a particularly node, the `Fin gsize` type value carried by this node is used to index the graph g. As the graph is defined as a vector of length `gsize`, the definition of node data type ensures that every well-typed node is a valid vertex in the graph, and that each indexing to the graph data structure are guaranteed to be in-bound.
The type of edge weight is user-defined in our implementation. Specifically, we define a `WeightOps`

data type, which carries a user-specified type for the edge weight, along with operators and properties proofs for this type, which includes arithmetic operators, proof of non-negative value, and proof of plus associativity. The definition of `Distance` data type is then parameterized over the user-defined edge weight data type. Since all edge weight are non-negative, the value of `Distance` can only be zero, infinity, or sum of edge weights.

4.1.2 Definition

Our implementation and correctness proof are based on the following definitions of key concepts used in Dijkstra's algorithm.

Definition 4.1. Path
(We adopt the definition of $path$ presented in the `Discrete Mathematics with Applications` book by `SUSANNA S. EPP`.)

A path from node v to w is a finite alternating sequence of adjacent vertices and edges of G, which does not contain any repeated edge or vertex. A path from v to w has the form:

$$ve_0v_0e_1v_2....v_{n-1}e_nw$$

where e_i is an edge in g with endpoints v_{i-1}, v_i. We denote the set of paths from v to w as $path(v, w)$.

Definition 4.2. Prefix of Path
Given a path from node v to w : $p(v, w) = ve_0v_0e_1v_2....v_{n-1}e_nw$, the prefix of this $v - w$ path is defined as the subsequence of $p(v, w)$ that starts with v and ends with some node $w' \in p(v, w)$ (w' is a vertex in the sequence $p(v, w)$).

Definition 4.3. Length of Path
The length of a path $p = ve_0v_0e_1v_2....v_{n-1}e_nw$ is the sum of the weights of all edges in p. We write:

$$length(p) = \sum weight(e_i), \forall e_i \in p.$$

Definition 4.4. Shortest Path
Denote $\Delta(s, v)$ as the shortest path from s to v, and $\delta(v)$ as the length of $\Delta(s, v)$. $\Delta(s, v)$ must fulfills:

$$\Delta(s, v) \in path(s, v)$$
and
$$\forall p' \in path(s, v),\ \delta(v) = length(\Delta(s, v)) \leq length(p')$$

4.1.3 Pseudocode

We denote (u, v) as an edge from node u to v, $weight(u, v)$ as the weight of edge (u, v). Let `gsize` denote the size of the input graph, i.e., the number of nodes in the graph. The type `Graph gsize weight` specifies a graph with `gsize` nodes and edge weight of type `weight`.
Given input graph g and source node s with types:

g : Graph gsize weight
s : Node gsize

We define $unexplored$ as the list of unexplored nodes, and $dist$ as the list storing distance from s to each node $n \in g$

(initially $unexplored$ contains all nodes in graph g)
$unexplored : List(Node\ \ gsize)$
$unexplored = \{v : v \in g\}$

(node value is used to index $dist$, initially distance of all nodes are infinity except the source node)
$dist : List\ \ weight$
$dist[s] = 0, dist[a] = infinity, \forall a \in g, a \neq s$

The Dijkstra's Algorithm runs as follows: Given graph g and source node s, $dist$ stores the distance value from s to all nodes in g calculated by the Dijkstra's algorithm, $dist[v]$ gives the corresponding distance value of v from s. We index $unexplored$ and $dist$ by the number of iterations. Specifically, denote u_i as the node being explored at the i^th iteration, and denote $dist_i$, $unexplored_i$ as the value of distance list and unexplored list at the beginning of the i^{th} iteration. Then during each iteration the Dijkstra's Algorithm calculates $dist, unexplored, explored$ as follows:

$\texttt{choose } u_k \in unexplored_k \texttt{ and } \forall u' \in unexplored_k, dist_k[u_k] \leq dist_k[u']$
$unexplored_{k+1} = unexplored_k - \{u_k\}$
$\texttt{for}(\forall v \in g)\ \{$

$$dist_{k+1}[v] = \begin{cases} min(dist_k[v], (dist_k[u_k] + weight(u_k, v))), & (u_k, v) \in g \\ dist_k[v] & otherwise \end{cases}$$

$\}$

4.1.4 Proof of Correctness

This section provides a theoretical proof for our Dijkstra's implementation, which includes proof of program termination and proof of correct program behavior.

4.1.4.1 Lemmas

Denote $explored$ as the list of nodes in g but not in $unexplored$, i.e., $explored$ stored all nodes whose neighbors have been updated by the algorithm. We index $explored$ by the number of iterations, such that $explored_i$ denotes the value of $explored$ at the beginning of the i^{th} iteration.

Lemma 4.1. Given any two nodes v, w, the prefix of the shortest path $\Delta(v, w)$ is also a shortest path.

Proof. We will prove Lemma 3.1 by contradiction.
Consider any node q in the sequence of $\Delta(v, w)$, we have $\Delta(v, w) = ve_0v_0e_1v_2...v_iqv_j....v_{n-1}e_nw$. Suppose the prefix of $\Delta(v, w)$ from v to q, denote as $p(v, q)$, is not the shortest path from v to q. Then we know $p(v, q) = ve_0v_0e_1v_2...v_iq$ is a path from v to q and $length(p(v, q)) > length(\Delta(v, q))$.
Based on the definition of shortest path, we know:

$$length(\Delta(v, w)) \leq length(p), \forall p \in path(v, w)$$

Fenote the path after the node q as $p(q, w) = qv_j....v_{n-1}e_nw$, since $\Delta(v, w) = ve_0v_0e_1v_2...v_iqv_j....v_{n-1}e_nw$, then $\Delta(v, w) = p(v, q) + p(q, w)$, and that $length(\Delta(v, w)) = length(p(v, q)) + length(p(q, w))$.
Then we have:

$$length(\Delta(v, w)) = length(p(v, q)) + length(p(q, w)) \leq length(p), \forall p \in path(v, w)$$

Since $p(v, q)$ is not the shortest path from v to q by assumption, then based on the definition of shortest path, $length(p(v, q)) < length(\Delta(v, w))$. Hence there exists another $v - w$ path $p'(v, w)$ such that:

$$\begin{aligned}
&p'(v, w) \in path(v, w) \\
&p'(v, w) = \Delta(v, q) + p(q, w) \\
&length(p'(v, w)) = length(\Delta(v, q)) + length(p(q, w)) \\
&\qquad\qquad\qquad\quad < length(p(v, q)) + length(p(q, w)) \\
&\text{i.e. } length(p'(v, w)) < length(\Delta(v, w))
\end{aligned}$$

Hence we have reached a contradiction. Thus by the principle of prove by contradiction, for any the prefix $p(v, q)$ of $\Delta(v, w)$ is the shortest path from v to q. Lemma 3.1 holds. □

Lemma 4.2. After the n^{th} iteration for $n \geq 1$, forall node $v \in explored_{n+1}$, if $dist_{n+1}[v] \neq infinity$, then $dist_{n+1}[v]$ is the length of some $s - v$ path, i.e, $path(s, v) \neq \emptyset$.

Proof. We will prove `Lemma 3.2` by inducting on the number of iterations.
Let P(n) be: After the n^{th} iteration, $n \geq 1$, for all node $v \in g$, if $dist_{n+1}[v] \neq infinity$, then $dist_{n+1}[v]$ is the length of some $s - v$ path.

Base Case : We shall show P(1) holds.
Based on the algorithm, initially $dist_1[s] = 0$ and for all node $v \in g, v \neq s, dist_1[v] = infinity$, then s is the only node whose distance value is not infinity. Based on the definition of path, the path from the source node s to itself is s, $path(s, s) = \{s\}$. Hence P(1) holds.

Inductive Hypothesis : Suppose $\forall i, 1 \leq i \leq k$, P(i) holds. That is, after the i^{th} iteration, $1 \leq i \leq k$, for all nodes $v \in g$, if $dist_{i+1}[v] \neq infinity$, then $dist_{n+1}[v]$ is the length of some $s - v$ path.

Inductive Step : We shall show P(k+1) holds.
For node u_{k+1} being explored during the $(k+1)^{th}$ iteration, based on the algorithm, $dist_{k+1}[u_{k+1}]$ is calculated as:

$$dist_{k+2}[u_{k+1}] = \begin{cases} min(dist_{k+1}[u_{k+1}], dist_{k+1}[u_{k+1}] + weight(u_{k+1}, u_{k+1})), & (u_{k+1}, u_{k+1}) \in g \\ dist_{k+1}[u_{k+1}] & otherwise \end{cases} \Bigg\}$$

Since the distance value from u_{k+1} to itself is 0, then $dist_{k+2}[u_{k+1}] = dist_{k+1}[u_{k+1}]$, and that $dist_{k+2}[u_{k+1}]$ and $dist_{k+1}[u_{k+1}]$ are the length of the same $s - u_{k+1}$ path if there exists one.
If $dist_{k+2}[u_{k+1}] \neq infinity$, then $dist_{k+1}[u_{k+1}] = dist_{k+2}[u_{k+1}] \neq infinity$. Since $k \leq k$ and $dist_{k+1}[u_{k+1}] \neq infinity$, then based on the inductive hypothesis, $dist_{k+1}[u_{k+1}]$ is the length of some $s - u_{k+1}$ path, and hence $dist_{k+2}[u_{k+1}]$ is the length of some $s - u_{k+1}$ path.
Then for all node $v \in g$ other than u_{k+1}, there are two cases: (1) $(u_{k+1}, v) \in g$; (2) u_{k+1} does not have an edge to v. We will prove P(k+1) holds in both cases separately.

Case (1): $(u_{k+1}, v) \in g$
Based on the algorithm, as $(u_{k+1}, v) \in g$, $dist_{k+2}[v] = min(dist_{k+1}[v], dist_{k+1}[u_{k+1}]+weight(u_{k+1}, v))$.

- If $dist_{k+1}[v] < dist_{k+1}[u_{k+1}] + weight(u_{k+1}, v)$, then $dist_{k+2}[v] = dist_{k+1}[v]$. Then if $dist_{k+2}[v] \neq infinity$, we have $dist_{k+1}[v] \neq infinity$, and that $dist_{k+2}[v]$ and $dist_{k+1}[v]$ are the length of the same $s-v$ path if there exists one. Since $dist_{k+1}[v] \neq infinity$, the inductive hypothesis implies that $dist_{k+1}[v]$ is the length of some $s - v$ path, hence $dist_{k+2}[v]$ is the length of some $s - v$ path. P(k+1) holds.

- If $dist_{k+1}[v] \geq dist_{k+1}[u_{k+1}]+weight(u_{k+1}, v)$, then $dist_{k+2}[v] = dist_{k+1}[u_{k+1}]+weight(w, v)$. If $dist_{k+2}[v] \neq infinity$, then it follows that $dist_{k+1}[u_{k+1}] = dist_{k+2}[v]-weight(u_{k+1}, v) \neq infinity$. Then the inductive hypothesis implies that $dist_{k+1}[u_{k+1}]$ must be the length of some $s - u_{k+1}$ path, denote as $p(s, u_{k+1})$. Since there is an edge $(u_{k+1}, v) \in g$, then $dist_{k+2}[v] = dist_{k+1}[u_{k+1}] + weight(u_{k+1}, v)$ must be the length of the $s - v$ path through u_{k+1}. P(k+1) holds.

Hence P(k+1) holds under under `Case(1)`.

Case (2): u_{k+1} **does not have an edge to** v
Under this case, our algorithm indicates that $dist_{k+2}[v] = dist_{k+1}[v]$, and that $dist_{k+1}[v]$ and $dist_{k+2}[v]$ are the length of the same $s - v$ path if there exists one. If $dist_{k+1}[v] = dist_{k+2}[v] \neq infinity$, then based on the inductive hypothesis, $dist_{k+1}[v]$ is the length of some $s-v$ path, and hence $dist_{k+2}[v]$ is the length of some $s - v$ path. P(k+1) holds under `Case(2)`.

We have proved P(k+1) holds for u_{k+1} and both cases for all nodes $v \in g$ other than u_{k+1}. Hence by the principle of prove by induction, P(n) holds. Thus `Lemma 3.2` holds. □

Lemma 4.3. For any node $v \in g$, if after the i^{th} iteration, $dist_{i+1}[v] = \delta(v)$, then for each proceeding j^{th} iteration, $j > i$, $dist_{j+1}[v] = dist_{i+1}[v] = \delta(v)$.

Proof. We will prove `Lemma 3.3` by induction on the number iterations after the i^{th} iteration.
Let P(n) be: For any node $v \in g$, if after the i^{th} iteration, $dist_{i+1}[v] = \delta(v)$, then for the $(i+n)^{th}$ iteration, $n \geq 1$, $dist_{i+n+1}[v] = dist_{i+1}[v] = \delta(v)$

Base Case : We shall show P(1) holds.
During the $(i+1)^{th}$ iteration, suppose u_{i+1} is the node being explored, then $dist_{i+2}[v]$ is calculated as:

$$dist_{i+2}[v] = \begin{Bmatrix} min(dist_{i+1}[v], dist_{i+1}[u_{i+1}] + weight(u_{i+1}, v)), & (u_{i+1}, v)) \in g \\ dist_{i+1}[v] & otherwise \end{Bmatrix}$$

If $(u_{i+1}, v)) \in g$, then if $dist_{i+1}[u_{i+1}]$ is the length of some $s - u_{i+1}$ path, then $(dist_{i+1}[u_{i+1}] + weight(u_{i+1}, v))$ is the length of some $s - v$ path. Since $dist_{i+1}[v] = \delta(v)$, then based on the definition of shortest path, $dist_{i+1}[v] \leq dist_{i+1}[u_{i+1}] + weight(u_{i+1}, v)$, and hence $dist_{i+2}[v] = dist_{i+1}[v] = \delta(v)$.
If u_{i+1} does not have an edge to v, then $dist_{i+2}[v] = dist_{i+1}[v] = \delta(v)$.
Hence in either cases, $dist_{i+2}[v] = dist_{i+1}[v] = \delta(v)$. P(1) holds.

Inductive Hypothesis : Suppose P(k) holds, that is, if after the i^{th} iteration, $dist_{i+1}[v] = \delta(v)$, then for the $(i+k)^{th}$ iteration, $n \geq 1$, $dist_{i+k+1}[v] = dist_{i+1}[v] = \delta(v)$.

Inductive Step : We shall show P(k+1) holds.
For the node u_{i+k+1} being explored during the $(i+k+1)^{th}$ iteration, there are two cases: (1) $(u_{i+k+1}, v) \in g$; (2) u_{i+k+1} does not have an edge to v. We will show that P(k+1) holds under both cases separately.
Case 1: $(u_{i+k+1}, v) \in g$
If u_{i+k+1} has an edge to v, then based on the algorithm, for $dist_{i+k+2}[v]$, we have:

$$dist_{i+k+2}[v] = min(dist_{i+k+1}[v], dist_{i+k+1}[u_{i+k+1}] + weight(u_{i+k+1}, v))$$

Since based on our inductive hypothesis, $dist_{i+k+1}[v] = dist_{i+1}[v] = \delta(v)$, then if $dist_{i+k+1}[u_{i+k+1}]$ is the length of some $s - u_{i+k+1}$ path, then $(dist_{i+k+1}[u_{i+1}] + weight(u_{i+k+1}, v))$ is the length of some $s - v$ path, and hence $dist_{i+k+1}[v] = \delta(v) \leq (dist_{i+k+1}[u_{i+1}] + weight(u_{i+k+1}, v))$. Then:

$$\begin{aligned} dist_{i+k+2}[v] &= min(dist_{i+k+1}[v], dist_{i+k+1}[u_{i+k+1}] + weight(u_{i+k+1}, v)) \\ &= dist_{i+k+1}[v] \\ &= dist_{i+1}[v] = \delta(v) \end{aligned}$$

P(k+1) holds under `Case 1`.

Case 2: u_{i+k+1} **does not have an edge to** v
Since u_{i+k+1} does not have an edge to v, then $dist_{i+k+2}[v] = dist_{i+k+1}[v]$. Based on the inductive hypothesis, $dist_{i+k+1}[v] = dist_{i+1}[v] = \delta(v)$. then $dist_{i+k+2}[v] = dist_{i+1}[v] = \delta(v)$. P(k+1) holds for `Case (2)`.
Thus P(k+1) holds. By the principle of prove by induction, P(n) holds. `Lemma 3.3` proved.

□

Lemma 4.4. Assume g is a connected graph, that the source node s has a path to every node in g. After the n^{th} iteration of the algorithm for $n \geq 1$, forall node $v \in explored_{n+1}$, we have:

1. $\delta(v) \leq \delta(v')$, $\forall v' \in unexplored_{n+1}$.
2. $dist_{n+1}[v] = \delta(v)$

Proof. We will prove `Lemma 3.4` by inducting on the number of iterations.
Let P(n) be: After the n^{th} iteration of the algroithm for $n \geq 1$, forall node $w \in explored_{n+1}$: (1) $\delta(w) \leq \delta(w')$, $\forall w' \in unexplored_{n+1}$; (2) $dist_{n+1}[w] = \delta(w)$.

Base Case : We shall show P(1) holds
Based on the algorithm, during the first iteration, the node with minimum distance value is the source node s with $dist_1[s] = 0$. Hence during the first iteration, only s is removed from $unexplored_1$ and added to $explored_2$. Since all edge weights are non-negative, then the shortest distance value from s to s is indeed 0, hence $dist_2[s] = 0 = \delta(s)$ and $\delta(s) \leq \delta(v')$, $\forall v' \in unexplored_2$.
P(1) holds.

Inductive Hypothesis : Suppose P(i) is true for all $1 \leq i \leq k$. That is, after the i^{th} iteration forall $1 < i \leq k$, forall node $w \in explored_{i+1}$: (1) $\delta(w) \leq \delta(w')$, $\forall w' \in unexplored_{i+1}$; (2) $dist_{i+1}[w] = \delta(w)$;

Inductive Step : We shall show P(k+1) holds. That is, forall node $w \in explored_{k+2}$, (1) $\delta(w) \leq \delta(w')$, $\forall w' \in unexplored_{k+2}$; (2) $dist_{k+2}[w] = \delta(w)$;
Suppose u_{k+1} is the node added into $explored$ during the $(k+1)^{th}$ iteration, then $explored_{k+2} = explored_{k+1} \cup \{u_{k+1}\}$. We will show that (1) and (2) holds for all nodes in $explored_{k+1}$ in `Part (a)`, and `Part (b)` proves (1) and (2) holds for u_{k+1}, so that (1) and (2) holds forall nodes in $explored_{k+2}$.

- **Part(a)**: `WTP: After the` $(k+1)^{th}$ `iteration,` $\forall w \in explored_{k+1}$`, (a.1)` $\delta(w) \leq \delta(w')$, $\forall w' \in unexplored_{k+2}$`; (a.2)` $dist_{i+1}[w] = \delta(w)$

  Consider each node $q \in (explored_{k+1} \cap explored_{k+2}) = explored_{k+1}$, q must be explored before the $(k+1)^{th}$ iteration. Suppose q is explored during the i^{th} iteration for some $i < k+1$, then based on our inductive hypothesis, $dist_{i+1}[q] = \delta(q)$, and $\delta(q) \leq \delta(q'), \forall q' \in unexplored_{i+1}$.
  Proof of (a.1): Based on the algorithm, for each iteration, the algorithm explores exactly one node and never revisits any explored nodes. For each node $q \in explored_{k+1}$ mentioned above, since q is explored before the $(k+1)^{th}$ iteration, then $unexplored_{k+1} \subseteq unexplored_{i+1}$. Since $\delta(q) \leq \delta(q'), \forall q' \in unexplored_{i+1}$, and $unexplored_{i+1}$ includes all

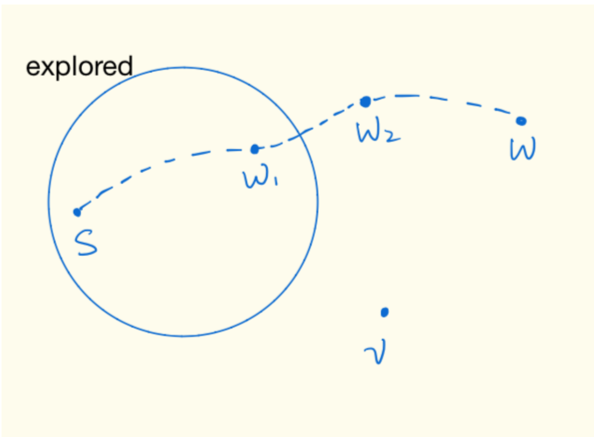
node in $unexplored_{k+1}$, then $\delta(q) \leq \delta(q'), \forall q' \in unexplored_{k+1}$. (1) holds for $explored_{k+1}$.
Proof of (a.2): For all proceeding j^{th} iterations, $j > i$, suppose node q'' is the node being explored for the j^{th} iteration, then the value of $dist_{j+1}[q]$ is calculated as:

$$dist_{j+1}[q] = \begin{cases} min(dist_j[q], dist_j[q''] + weight(q'', q)), & (q'', q) \in g \\ dist_j[q] & otherwise \end{cases}$$

Since $dist_{i+1}[q] = \delta(q) \leq length(p)$ for all path p from s to q, then for each proceeding j^{th} iteration after the i^{th} iteration, there does not exists such q'' such that $dist_j[q''] + weight(q'', q) < \delta(q) = dist_{i+1}[q]$. Hence $dist_{j+1}[q] = \delta(q) = dist_{i+1}[q], \forall j > i$. Since $k + 1 > i$, then for all $q \in S$, $dist_{k+1} = \delta(q) = dist_{i+1}[q]$. (2) holds for $explored_{k+1}$.

Hence we have proved that both (1) and (2) holds for all nodes in $explored_{k+1}$.

- **Part(b)**: After the $(k+1)^{th}$ iteration, (1) and (2) holds for u_{k+1}.
  We want to show: **(b.1)** $\delta(u_{k+1}) \leq \delta(v'), \forall v' \in unexplored_{k+2}$; and **(b.2)** $dist_{k+1}[u_{k+1}] = \delta(u_{k+1})$.
  Proof of (b.1): $\delta(u_{k+1}) \leq \delta(v')$**,** $\forall v' \in unexplored_{k+2}$
  We will prove (b.1) by contradiction. Suppose there exists $w \in unexplored_{k+2}$, such that $\delta(u_{k+1}) > \delta(w)$.
  The assumption states that source s has a path to every node in g, then $dist_{k+1}[u_{k+1}] \neq infinity$. Thus `Lemma 3.2` implies that $dist_{k+1}[u_{k+1}]$ is the length of some $s - v$ path. Based on the definition of shortest path, $\delta(u_{k+1}) \leq dist_{k+1}[u_{k+1}]$. Since $\delta(u_{k+1}) > \delta(w)$ and $\delta(u_{k+1}) \leq dist_{k+1}[u_{k+1}]$, then we have $\delta(w) < dist_{k+1}[u_{k+1}]$([NE1]).
  Consider the shortest path $\Delta(s, w)$ from s to w, $\delta(w) = length(\Delta(s, w))$. Since $w \notin explored_{k+2}$, then there must exists some node in $\Delta(s, w)$ that are not in $explored_{k+2}$. Suppose the first node along $\Delta(s, w)$ that is not in the $explored_{k+2}$ list is w_2, and the node right before w_2 in the s to w_2 subpath is w_1, thus $w_1 \in explored_{k+2}$. The image below illustrates this construction:



Denote the subpath from s to w_1 in $\Delta(s, w)$ as $p(s, w_1)$, subpath from s to w_2 in $\Delta(s, w)$ as $p(s, w_2)$, and subpath w_2 to w as $p(w_2, w)$. Based on `Definition 2.2 Prefix of Path`, $p(s, w_1)$ is a prefix of $\Delta(s, w)$. Since $p(s, w_1)$ is the prefix of the shortest $s - w$ path,

then based on `Lemma 3.1`, $p(s, w_1)$ is the shortest path from s to w_1, $\Delta(s, w_1) = p(s, w_1)$, $length(p(s, w_1)) = \delta(w_1)$.
Similarly, since $p(s, w_2) = p(s, w_1) + (w_1, w_2)$, then $p(s, w_2)$ is a prefix of $\Delta(s, w)$, and hence `Lemma 3.1` implies that $p(s, w_2)$ is the shortest path from s to w_2. Then we have:

$$\begin{aligned}\Delta(s, w_2) &= p(s, w_2) = p(s, w_1) + (w_1, w_2)\\ \delta(w_2) &= length(\Delta(s, w_2))\\ &= length(p(s, w_2))\\ &= length(p(s, w_1)) + weight(w_1, w_2)\\ &= \delta(w_1) + weight(w_1, w_2) \text{ ([E1])}\end{aligned}$$

For $\Delta(s, w)$ we have:

$$\begin{aligned}\delta(w) &= length(p_w)\\ &= length(p(s, w_1)) + weight(w_1, w_2) + length(p(w_2, w))\\ &= \delta(w_1) + weight(w_1, w_2) + length(p(w_2, w))\end{aligned}$$

Since all edge weights are positive, then:

$$\delta(w_2) = \delta(w_1) + weight(w_1, w_2) \leq \delta(w) \text{ ([E2])}$$

Since $w_1 \in explored_{k+2}$, there are two cases to consider: $w_1 = u_{k+1}$ and $w_1 \neq u_{k+1}$. We will prove P(k+1) under both cases below.

`Case 1:` $w_1 = u_{k+1}$
When $w_1 = u_{k+1}$, then substitude w_1 by u_{k+1} in [E2], we have:

$$\delta(w_1) + weight(w_1, w_2) = \delta(u_{k+1}) + weight(u_{k+1}, w_2) \leq delta(w)$$
$$\text{i.e. } \delta(u_{k+1}) \leq \delta(w)$$

which contradicts with our assumption that $\delta(u_{k+1}) > \delta(w)$. Hence by the principle of prove by contradiction, $\delta(u_{k+1}) < \delta(w)$. (1) holds for P(k+1).

`Case 2:` $w_1 \neq u_{k+1}$
Since $w_1 \in explored_{k+2}$ and $w_1 \neq u_{k+1}$, w_1 is explored before the $(k+1)^{th}$ iteration. i.e., $w_1 \in explored_{k+1}$. Suppose w_1 is being explored during the i^{th} iteration, $i < k+1$, then based on the algorithm, the value of $dist_{i+1}[w_1]$ is calculated as:

$$dist_{i+1}[w_1] = \begin{Bmatrix} min(dist_i[w_1], dist_i[w_1] + weight(w_1, w_1)), & (w_1, w_1) \in g \\ dist_i[w_1] & otherwise \end{Bmatrix}$$

Thus $dist_{i+1}[w_1] = dist_i[w_1]$. Since the inductive hypothesis implies that $dist_{i+1}[w_1] = \delta(w_1)$, then $dist_i[w_1] = \delta(w_1)$.
Since w_1 has an edge to w_2, then $dist_{i+1}[w_2]$ must have been updated according as follows:

$$
\begin{aligned}
dist_{i+1}[w_2] &= min(dist_i[w_2], dist_i[w_1] + weight(w_1 + w_2)) \\
&= min(dist_i[w_2], \delta(w_1) + weight(w_1 + w_2))
\end{aligned}
$$

Based on [E1] we know that $\delta(w_2) = \delta(w_1)+weight(w_1+w_2)$, then $dist_{i+1}[w_2] = min(dist_i[w_2], \delta(w_2))$. If $dist_i[w_2] = infinity$, then $dist_{i+1}[w_2] = min(dist_i[w_2], \delta(w_2)) = \delta(w_2)$. If $dist_i[w_2] \neq infinity$, then based on `Lemma 3.2`, $dist_i[w_2]$ is the length of some $s - w_2$ path. Since $\delta(w_2) \leq length(p), \forall p \in path(s, w_2)$, then $dist_{i+1}[w_2] = min(dist_i[w_2], \delta(w_2)) = \delta(w_2)$. Hence in either cases, we conclude that $dist_{i+1}[w_2] = \delta(w_2)$.
Since $dist_{i+1}[w_2] = \delta(w_2)$ and $i < k + 1$, then based on `Lemma 3.3`, we have $dist_{k+1}[w_2] = dist_{i+1} = \delta(w_2)$. Based on [E2], $\delta(w_2) < \delta(w)$, then $dist_{k+1}[w_2] < \delta(w)$ [NE2]. Combining with [NE1], we have:

$$
\begin{aligned}
&\delta(w) < dist_{k+1}[u_{k+1}] \text{ (from [NE1])} \\
&dist_{k+1}[w_2] < \delta(w)
\end{aligned}
$$

Hence $dist_{k+1}[w_2] < dist_{k+1}[u_{k+1}]$ [NE2].
Based on our assumption, at the beginning of the $(k+1)^{th}$ generation, $u_{k+1}, w_2 \notin explored_{k+1}$ and u_{k+1} is selected by the algorithm, then we must have $dist_{k+1}[w_2] \geq dist_{k+1}[u_{k+1}]$, which contradicts with [NE2]. Hence by the principle of prove by contradiction, there does not exsist $w \in unexplored_{k+2}$, such that $\delta(u_{k+1}) > \delta(w)$, i.e. $\delta(u_{k+1}) \leq \delta(w), \forall w \in unexplored_{k+2}$. Hence (b.1) holds.

Proof of (b.2): After the $(k + 1)^{th}$ iteration, $dist_{k+1}[u_{k+1}] = \delta(u_{k+1})$
We will prove this by contradiction.
Suppose $dist_{k+1}[u_{k+1}]$ is the length of some path p from s to u_{k+1}. Assume the shortest path from s to u_{k+1} is some path different from p, i.e. $\Delta(s, u_{k+1}) \neq p$, $\delta(u_{k+1}) \leq dist_{k+1}[u_{k+1}]$([NE3]). Suppose v' is the node just before u_{k+1} in $\Delta(s, u_{k+1})$.

$$
\delta(u_{k+1}) = \delta(v') + weight(v', u_{k+1}) < dist_{k+1}[u_{k+1}]
$$

Since all edge weights are non-negative, then: $\delta(v') \leq \delta(u_{k+1})$

Based on (a.1) and (b.1), after the $(k + 1)^{th}$ iteration, for all nodes $q \in unexplored_{k+2}$, $\delta(q) \geq \delta(u_{k+1})$, and $\delta(v') \leq \delta(u_{k+1})$, then v' cannot be in $unexplored_{k+2}$. Since $unexplored_{k+1} = unexplored_{k+2} \cup u_{k+1}$, then $v' \notin unexplored_{k+1}$. Hence at the beginning of the $(k + 1)^{th}$ iteration, v' is already explored. Since v' is explored before the $(k + 1)^{th}$ iteration and v' has an edge to u_{k+1}, then the algorithm must have considered $(\delta(v') + weight(v', u_{k+1}))$ against $dist_{k+1}[u_{k+1}]$ and chose $min((\delta(v') + weight(v', u_{k+1})), dist_{k+1}[u_{k+1}])$, which is $dist_{k+1}[u_{k+1}]$. Thus $dist_{k+1}[u_{k+1}] \leq (\delta(v')+weight(v', u_{k+1}))$, i.e. $dist_{k+1}[u_{k+1}] \leq \delta(u_{k+1})$, which contradicts with our assumption [NE3]. Hence by the principle of prove by contradiction, $dist_{k+1}[u_{k+1}] = \delta(u_{k+1})$. (b.2) holds.

Since we have proved both (1) and (2) forall nodes in $explored_{k+1}$ after the $(k + 1)^{th}$ iteration, P(k+1) holds. Then by the principle of prove by induction, `Lemma 3.4` holds. □

4.1.4.2 Proof of Termination

Proof. The inner for loop is guaranteed to terminate as the algorithm goes through each adjacent node exactly once. As the size of list `unexplored` decreases by one during each iteration of the while loop, the algorithm is guaranteed to terminate. □

4.1.4.3 Prove of Correctness

Proof. By applying `Lemma 3.4` to the last iteration of the algorithm, we obtained that for all nodes n in the explored list, $dist[n]$ is indeed the shortest path distance value from source s to n, hence Dijkstra's algorithm indeed calculates the shortest path distance value from the source s to each node $n \in g$. □

5 Low-Level Contribution

6 Discussion

7 Related Work

Existing work on verifying Dijkstra's algorithm is relatively limited, and few resources are found for the verification of Bellman-Ford algorithm. Robin Mange and Jonathan Kuhn demonstrates an implementation that verifies Dijkstra's algorithm with the Jahob verification system in their report on efficient proving of Java implementations[3]. Although few resource has been found on the concrete implementation of this work, the report illustrates that as Jahob allows programmers to provide specification of their function's behaviors in high-level logic(HOL), program verification can be reduced to the problem of the validity of HOL formulas.

Klasen et. al. from the University of Koblenz and Landau present a concrete implementation of Dijkstra's algorithm in Java and its verification with the KeY system[1]. The verification process involves specifying the behavior of each function with preconditions, postconditions, and invariants, and the KeY system checks a function as correct with respect to its specifications if the postconditions hold after execution. Jean-Christophe Filliâtre, a senior researcher from the National Center for Scientific Research(CNRS), offers an implementation of Dijkstra's algorithm along with its verification in Why3, a deductive program verification platform that relies on external theorem provers[4][5]. Both verifications above are largely dependent on theorem proving systems. Unlike Filliâtre and Klasen et. al., our work relies on a significantly smaller trusted code base, indicating that considerable amount of proofs will be presented explicitly in our implementation rather than provided by external theorem provers.

8 Conclusion

References

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Appendix

Statutory Declaration