Separating Use and Reuse to Improve Both

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Subtyping and Subclassing

- Subtyping without subclassing; easy: Java interfaces
- Subclassing without subtyping; hard:

```
class A{ int ma(){ return Utils.m(this); } }
class Utils{ static int m(A a){..} }
class B extends A{ int mb(){return this.ma();} }
..
new B().mb();
```

This-leaking problem: `this` written in class `A` is of type `A`. when `B` inherits the code of `A`, code `Uitls.m(this)` would pass a `this` as an `A`. To be sound, $B \leq A$ must hold.

Three different approaches separating subtyping and subclassing:

- DeepFJig
- TraitRecordJ
- PackageTemplate

Widely different approaches, similar core ideas that we synthesize and improve: *separating code from use and code for reuse*

We synthesize the best from these 3 approaches.
Traits! (inspired from the original approach, not Scala/Rust traits).

- **Traits** for Code Reuse: traits can be reused to produce more code.
- **Classes** for Code Use: classes can be instantiated.
- **Types**: classes and interface names can be used as a types.
- **Exact types**: all classes are final; inheritance by reusing traits.
- **Sub types**: interfaces are the only way to induce subtyping.
- **Not a type**: trait names can not be used as types.

  `This` can be used to refer to the eventual class of the code.
- **State**: No explicit constructors/fields; state is implicitly defined.
- **Composition**: traits code can be composed to obtain more code.

This leaking in our approach

```
IA ={ interface
  method int ma ()
  }
Utils ={
  static method int m ( IA a ){ return ...}
  }
ta ={ implements IA
  method int ma (){ return Utils.m ( this ); }
  }
A = Use ta
B = Use ta
```

`this` written in trait `ta` is of type `This` and `IA`. when `B` inherits the code of `ta`, code `Uitls.m(this)` would pass a `this` as an `IA`. In this way, even if `B` and `A` share the same code, there are not in a subtype relation. (Note that `ta` is not a valid typename)

Trait composition

The composed code contains all members from the used traits. Methods
with same name and type signature are summed into a single one. At
most one of those summed methods can have a body, which will be
propagated into the result.

```
t1={
  method String hello();
  method String helloWorld(){return hello()+"␣World";}
  }

t2={ method String hello(){return "Hi";} }

t3= Use t1,t2
//-- flatten to --------------------------
t3= {
  method String hello(){return "Hi";}
  method String helloWorld(){return hello()+"␣World";}
  }
```

How to make this process work with constructors and fields while achieving the following goals:

- managing fields in a way that borrows the elegance of summing methods;
- actually initializing objects, leaving no `null` fields;
- making it easy to add new fields;
- allowing self instantiation: a trait method can instantiate the class using it.

Abstract State Operations

Idea: use abstract methods as getters, setters and factory methods.
Conventionally, an abstract class is a class with some abstract method.
Here, a class whose abstract methods can be seen as state operations is
a concrete coherent class.
For example Java

```
class A{
  int x;
  A(int x){this.x=x;}
  int x(){return this.x;}
  void x(int that){this.x=that;}
  }
```

42_μ

```
A={
  static method This of(Int x)
  method Int x()
  method Void x(Int that)
  }
```

Coherent class

- A class with no abstract methods is coherent (just like Java `Math`, for example). Such classes have no instances and are only useful for calling static methods.
- A class with a single abstract `static` method returning `This` and with parameters $T_1\ x_1, \ldots, T_n\ x_n$ is coherent if all the other abstract methods can be seen as *abstract state operations* over one of $x_1, \ldots, x_n$. That is:
  - A method $T_i\ x_i()$ is interpreted as an abstract state method: a *getter* for x_i.
  - A method `void` $x_i(T_i$ `that`) is a *setter* for x_i.

static factory method:

```
Point p=Point.of(3,4)
```

getters:

```
p.x()==3,    p.y()==4
```

summing points:

```
p.sum(Point.of(10,20))
```
 equivalent to `Point.of(13,24)`

NOTE: sum operation creates a new point

If no code reuse is desired, it is easy to use Java to encode such `Point` class. However, we would like to define operations `sum`, `mul`, `div` and `sub` independently and compose them to create classes with the operations we want, so that it is easy to have points with `sum` and `mul`, points with just `sum` or points with just `mul`.

For example, we expect `PointSumMul.sum(PointSumMul) : PointSumMul`.

This breaks subtyping: `PointSumMul ⊄ PointSum`.

```
p= { //point trait with abstract state
  method int x() //getter for field x
  method int y() //getter for field y
  static method This of(int x,int y) //factory method
  }

Point= Use p //this is a concrete coherent class
...
.. Point.of(3,7).x() ..//valid code
```

Abstract state

```
p={
  method int x()
  method int y()
  static method This of(int x,int y)
  }

pointSum= Use p, {
  method This sum(This that){
    return This.of(this.x()+that.x(),this.y()+that.y());
    }}

pointMul= Use p, {
  method This mul(This that){
    return This.of(this.x()*that.x(),this.y()*that.y());
    }}

pointDiv= ...

MyPoint= Use pointSum,pointMul,pointDiv,...
```

Encoding the point example above, with the 4 arithmetic operations, and instantiating all the 16 possible permutations as separate classes:

Language	Lines of code	members	classes/traits
Java7	115	50	16
Classless Java	82	34	16
Scala	81	40	21
42_μ	32	7	21

State extensibility

```
p={
  method int x()
  method int y()
  static method This of(int x,int y)
  }
pointSum= Use p, {
  method This sum(This that){
    return This.of(this.x()+that.x(),this.y()+that.y());
    }}

colored= { method Color color() }

CPoint= Use pointSum,colored,{
  static method This of(int x,int y){
    return This.of(x,y,Color.of(/*red*/));
    }
  static method This of(int x,int y,Color color)
  }
```

Done in this way, `CPoint.sum` would return red points.

State extensibility, with withers

A wither is like a field setter, but creates a new copy of the object

```
p= {
  method int x()
  method int y()
  method This withX(int that)
  method This withY(int that)
  method This merge(This that)
  }
pointSum= Use p, {
  method This sum(This that){
    return this.merge(that)
      .withX(this.x()+that.x()).withY(this.y()+that.y());
  }}
colored= {
  method Color color()
  method This withColor(Color that)
  method This merge(This that){
    return this.withColor(this.color().mix(that.color()));
  }}
CPoint= Use pointSum,colored,{
  static method This of(int x, int y, Color color)}
```

Independent extensibility

```
flavored = {
  method Flavor flavor ()
  method This withFlavor ( Flavor that )
  method This merge ( This that ){
    return this.withFlavor ( that.flavor ());
  }}
FCPoint = Use //aliasing conflicting implementations
  colored [ super merge as m1 ],
  flavored [ super merge as m2 ],
  pointSum ,{
    method This merge ( This that ){
      return this.m1 ( that ).m2 ( that );
    }
    static method This of (
      int x , int y , Color color , Flavor flavor )
    }
```

More in the paper

- Transparent handling of nested classes
- `This` type generalized for family polymorphism
- A very natural encoding of the Expression Problem
- Simple formalization of our language

Thanks

Questions?