**HDF5 User’s Guide**

**HDF5 Release 1.10**



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# The HDF5 Data Model and File Structure

## Introduction

The Hierarchical Data Format (HDF) implements a model for managing and storing data. The model includes an abstract data model and an abstract storage model (the data format), and libraries to imple­ment the abstract model and to map the storage model to different storage mechanisms. The HDF5 library provides a programming interface to a concrete implementation of the abstract models. The library also implements a model of data transfer, an efficient movement of data from one stored represen­tation to another stored representation. The figure below illustrates the relationships between the mod­els and implementations. This chapter explains these models in detail.

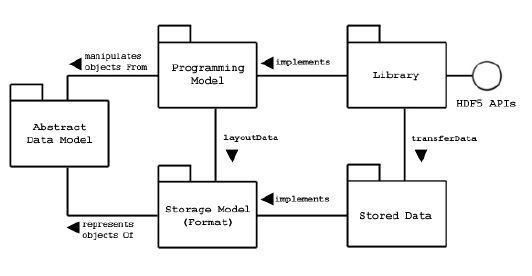


Figure 1-1. HDF5 models and implementations

The *Abstract Data Model* is a conceptual model of data, data types, and data organization. The abstract data model is independent of storage medium or programming environment. The *Storage Model* is a stan­dard representation for the objects of the abstract data model. The [HDF5 File Format Specification](https://portal.hdfgroup.org/display/HDF5/File+Format+Specification) defines the storage model.

The *Programming Model* is a model of the computing environment and includes platforms from small sin­gle systems to large multiprocessors and clusters. The programming model manipulates (instantiates, pop­ulates, and retrieves) objects from the abstract data model.

The *Library* is the concrete implementation of the programming model. The library exports the HDF5 APIs as its interface. In addition to implementing the objects of the abstract data model, the library manages data transfers from one stored form to another. Data transfer examples include reading from disk to mem­ory and writing from memory to disk.

*Stored Data* is the concrete implementation of the storage model. The *Storage Model* is mapped to several storage mechanisms including single disk files, multiple files (family of files), and memory representations.

The HDF5 library is a C module that implements the programming model and abstract data model. The HDF5 library calls the operating system or other storage management software (for example, the MPI/IO Library) to store and retrieve persistent data. The HDF5 library may also link to other software such as fil­ters for compression. The HDF5 library is linked to an application program which may be written in C, C++, Fortran, or Java. The application program implements problem specific algorithms and data structures and calls the HDF5 library to store and retrieve data. The figure below shows the dependencies of these mod­ules.

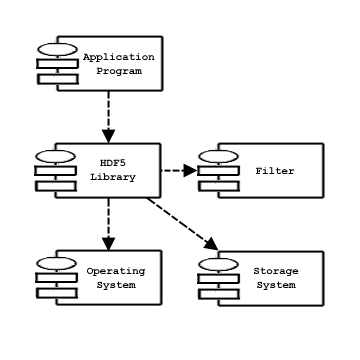


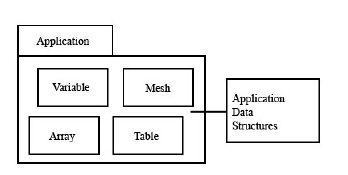
Figure 1-2. The library, the application program, and other modules

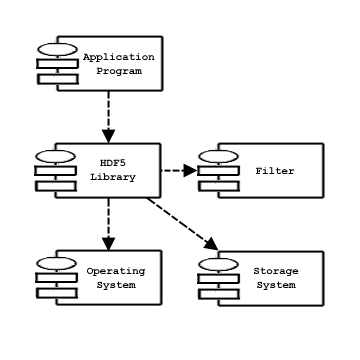
It is important to realize that each of the software components manages data using models and data structures that are appropriate to the component. When data is passed between layers (during storage or retrieval), it is transformed from one representation to another. The figure below suggests some of the kinds of data structures used in the different layers.

The *Application Program* uses data structures that represent the problem and algorithms including vari­ables, tables, arrays, and meshes among other data structures. Depending on its design and function, an application may have quite a few different kinds of data structures and different numbers and sizes of objects.

The *HDF5 Library* implements the objects of the HDF5 abstract data model. Some of these objects include groups, datasets, and attributes. The application program maps the application data structures to a hier­archy of HDF5 objects. Each application will create a mapping best suited to its purposes.

The objects of the HDF5 abstract data model are mapped to the objects of the HDF5 storage model, and stored in a storage medium. The stored objects include header blocks, free lists, data blocks, B-trees, and other objects. Each group or dataset is stored as one or more header and data blocks. See the [HDF5 File Format Specification](https://portal.hdfgroup.org/display/HDF5/File+Format+Specification) for more information on how these objects are organized. The HDF5 library can also use other libraries and modules such as compression.





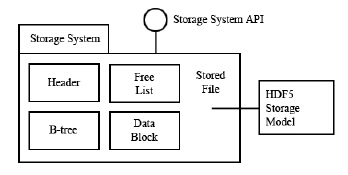


Figure 1-3. Data structures in different layers

The important point to note is that there is not necessarily any simple correspondence between the objects of the application program, the abstract data model, and those of the Format Specification. The organization of the data of application program, and how it is mapped to the HDF5 abstract data model is up to the application developer. The application program only needs to deal with the library and the abstract data model. Most applications need not consider any details of the [HDF5 File Format Specifica­tion](https://portal.hdfgroup.org/display/HDF5/File+Format+Specification) or the details of how objects of abstract data model are translated to and from storage.

## The Abstract Data Model

The abstract data model (ADM) defines concepts for defining and describing complex data stored in files. The ADM is a very general model which is designed to conceptually cover many specific models. Many dif­ferent kinds of data can be mapped to objects of the ADM, and therefore stored and retrieved using HDF5. The ADM is not, however, a model of any particular problem or application domain. Users need to map their data to the concepts of the ADM.

The key concepts include:

•        *Fil*e - a contiguous string of bytes in a computer store (memory, disk, etc.), and the bytes repre­sent zero or more objects of the model

•        *Group* - a collection of objects (including groups)

•        *Dataset* - a multidimensional array of data elements with attributes and other metadata

•        *Dataspace* - a description of the dimensions of a multidimensional array

•        *Datatype* - a description of a specific class of data element including its storage layout as a pattern of bits

•        *Attribute* - a named data value associated with a group, dataset, or named datatype

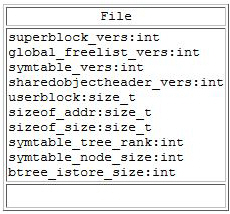
•        *Property List* - a collection of parameters (some permanent and some transient) controlling options in the library

•        *Link* - the way objects are connected

These key concepts are described in more detail below.

### File

Abstractly, an HDF5 file is a container for an organized collection of objects. The objects are groups, data­sets, and other objects as defined below. The objects are organized as a rooted, directed graph. Every HDF5 file has at least one object, the root group. See the figure below. All objects are members of the root group or descendants of the root group.



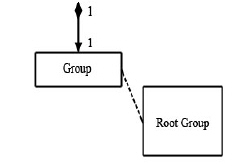


Figure 1-4. The HDF5 file

HDF5 objects have a unique identity within a single HDF5 file and can be accessed only by their names within the hierarchy of the file. HDF5 objects in different files do not necessarily have unique identities, and it is not possible to access a permanent HDF5 object except through a file. For more information, see Section 1.4, *The Structure of an HDF5 File*.

When the file is created, the file creation properties specify settings for the file. The file creation proper­ties include version information and parameters of global data structures. When the file is opened, the file access properties specify settings for the current access to the file. File access properties include parame­ters for storage drivers and parameters for caching and garbage collection. The file creation properties are set permanently for the life of the file, and the file access properties can be changed by closing and reopening the file.

An HDF5 file can be “mounted” as part of another HDF5 file. This is analogous to Unix file system mounts. The root of the mounted file is attached to a group in the mounting file, and all the contents can be accessed as if the mounted file were part of the mounting file.

### Group

An HDF5 group is analogous to a file system directory. Abstractly, a group contains zero or more objects, and every object must be a member of at least one group. The root group is a special case; it may not be a member of any group.

Group membership is actually implemented via link objects. See the figure below. A link object is owned by a group and points to a named object. Each link has a name, and each link points to exactly one object. Each named object has at least one and possibly many links to it.

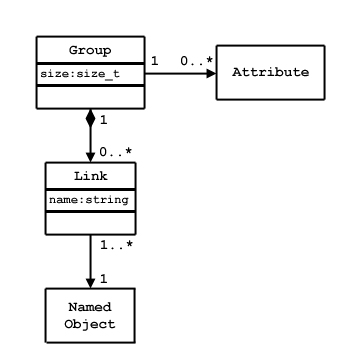


Figure 1-5. Group membership via link objects

There are three classes of named objects: group, dataset, and committed (named) datatype. See the fig­ure below. Each of these objects is the member of at least one group, and this means there is at least one link to it.

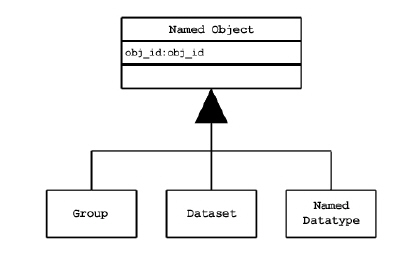


Figure 1-6. Classes of named objects

### Dataset

An HDF5 dataset is a multidimensional (rectangular) array of data elements. See the figure below. The shape of the array (number of dimensions, size of each dimension) is described by the dataspace object (described in the next section below).

A data element is a single unit of data which may be a number, a character, an array of numbers or charac­ters, or a record of heterogeneous data elements. A data element is a set of bits. The layout of the bits is described by the datatype (see below).

The dataspace and datatype are set when the dataset is created, and they cannot be changed for the life of the dataset. The dataset creation properties are set when the dataset is created. The dataset creation properties include the fill value and storage properties such as chunking and compression. These proper­ties cannot be changed after the dataset is created.

The dataset object manages the storage and access to the data. While the data is conceptually a contigu­ous rectangular array, it is physically stored and transferred in different ways depending on the storage properties and the storage mechanism used. The actual storage may be a set of compressed chunks, and the access may be through different storage mechanisms and caches. The dataset maps between the con­ceptual array of elements and the actual stored data.

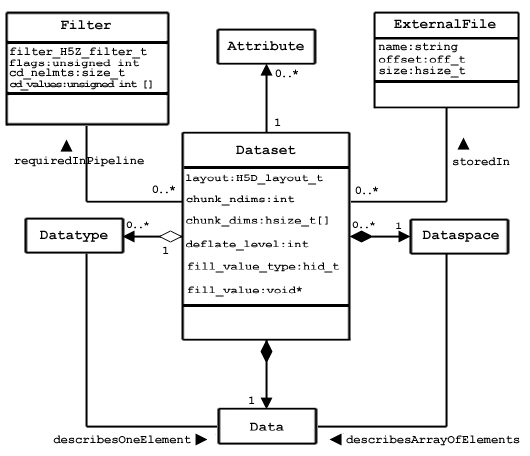


Figure 1-7. The dataset

### Dataspace

The HDF5 dataspace describes the layout of the elements of a multidimensional array. Conceptually, the array is a hyper-rectangle with one to 32 dimensions. HDF5 dataspaces can be extendable. Therefore, each dimension has a current size and a maximum size, and the maximum may be unlimited. The dataspace describes this hyper-rectangle: it is a list of dimensions with the current and maximum (or unlimited) sizes. See the figure below.

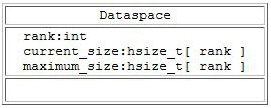


Figure 1-8. The dataspace

Dataspace objects are also used to describe hyperslab selections from a dataset. Any subset of the ele­ments of a dataset can be selected for read or write by specifying a set of hyperslabs. A non-rectangular region can be selected by the union of several (rectangular) dataspaces.

### Datatype

The HDF5 datatype object describes the layout of a single data element. A data element is a single ele­ment of the array; it may be a single number, a character, an array of numbers or carriers, or other data. The datatype object describes the storage layout of this data.

Data types are categorized into 11 classes of datatype. Each class is interpreted according to a set of rules and has a specific set of properties to describe its storage. For instance, floating point numbers have expo­nent position and sizes which are interpreted according to appropriate standards for number representa­tion. Thus, the datatype class tells what the element means, and the datatype describes how it is stored.

The figure below shows the classification of datatypes. Atomic datatypes are indivisible. Each may be a single object such as a number or a string. Composite datatypes are composed of multiple elements of atomic datatypes. In addition to the standard types, users can define additional datatypes such as a 24-bit integer or a 16-bit float.

A dataset or attribute has a single datatype object associated with it. See Figure 7 above. The datatype object may be used in the definition of several objects, but by default, a copy of the datatype object will be private to the dataset.

Optionally, a datatype object can be stored in the HDF5 file. The datatype is linked into a group, and there­fore given a name. A committed datatype (formerly called a named datatype) can be opened and used in any way that a datatype object can be used.

For more information, see Section 6, *HDF5 Datatypes*.

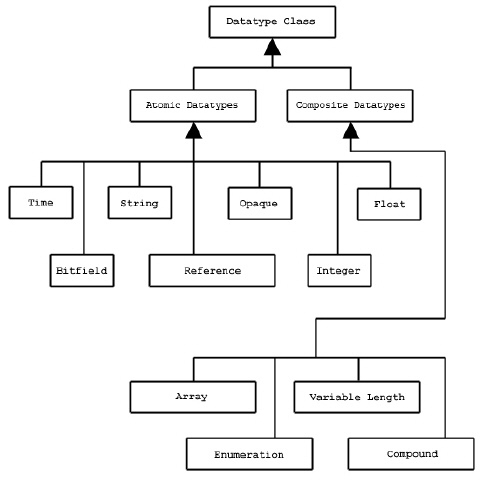


Figure 1-9. Datatype classifications

### Attribute

Any HDF5 named data object (group, dataset, or named datatype) may have zero or more user defined attributes. Attributes are used to document the object. The attributes of an object are stored with the object.

An HDF5 attribute has a name and data. The data portion is similar in structure to a dataset: a dataspace defines the layout of an array of data elements, and a datatype defines the storage layout and interpreta­tion of the elements See the figure below.

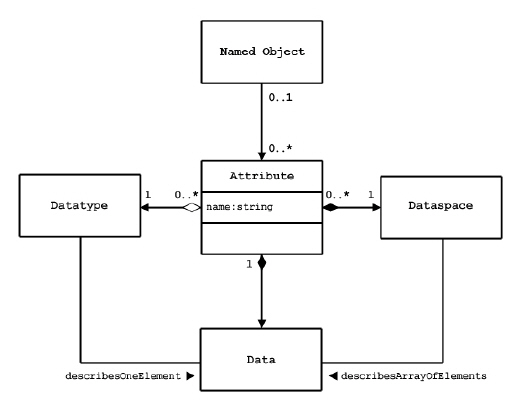


Figure 1-10. Attribute data elements

In fact, an attribute is very similar to a dataset with the following limitations:

•        An attribute can only be accessed via the object

•        Attribute names are significant only within the object

•        An attribute should be a small object

•        The data of an attribute must be read or written in a single access (partial reading or

writing is not allowed)

•        Attributes do not have attributes

Note that the value of an attribute can be an object reference. A shared attribute or an attribute that is a large array can be implemented as a reference to a dataset.

The name, dataspace, and datatype of an attribute are specified when it is created and cannot be changed over the life of the attribute. An attribute can be opened by name, by index, or by iterating through all the attributes of the object.

### Property List

HDF5 has a generic property list object. Each list is a collection of name-value pairs. Each class of property list has a specific set of properties. Each property has an implicit name, a datatype, and a value. See the figure below. A property list object is created and used in ways similar to the other objects of the HDF5 library.

Property Lists are attached to the object in the library, and they can be used by any part of the library. Some properties are permanent (for example, the chunking strategy for a dataset), others are transient (for example, buffer sizes for data transfer). A common use of a Property List is to pass parameters from the calling program to a VFL driver or a module of the pipeline.

Property lists are conceptually similar to attributes. Property lists are information relevant to the behavior of the library while attributes are relevant to the user’s data and application.



Figure 1-11. The property list

Property lists are used to control optional behavior for file creation, file access, dataset creation, dataset transfer (read, write), and file mounting. Some property list classes are shown in the table below. Details of the different property lists are explained in the relevant sections of this document.

|  |  |  |
| --- | --- | --- |
| Table 1-1. Property list classes and their usage | | |
| **Property List Class** | **Used** | **Examples** |
| H5P\_FILE\_CREATE | Properties for file creation. | Set size of user block. |
| H5P\_FILE\_ACCESS | Properties for file access. | Set parameters for VFL driver. An example is MPI I/O. |
| H5P\_DATASET\_CREATE | Properties for dataset cre­ation. | Set chunking, compression, or fill value. |
| H5P\_DATASET\_XFER | Properties for raw data trans­fer (read and write). | Tune buffer sizes or memory management. |
| H5P\_FILE\_MOUNT | Properties for file mounting. |  |

### Link

This section is under construction.

## The HDF5 Storage Model

### The Abstract Storage Model: the HDF5 Format Specification

The [HDF5 File Format Specification](https://portal.hdfgroup.org/display/HDF5/File+Format+Specification) defines how HDF5 objects and data are mapped to a linear address space. The address space is assumed to be a contiguous array of bytes stored on some random access medium.1 The format defines the standard for how the objects of the abstract data model are mapped to linear addresses. The stored representation is self-describing in the sense that the format defines all the information necessary to read and reconstruct the original objects of the abstract data model.

The HDF5 File Format Specification is organized in three parts:

1.      **Level 0**: File signature and super block

2.      **Level 1**: File infrastructure

a.      **Level 1A**: B-link trees and B-tree nodes

b.      **Level 1B**: Group

c.      **Level 1C**: Group entry

d.     **Level 1D**: Local heaps

e.      **Level 1E**: Global heap

f.       **Level 1F**: Free-space index

3.      **Level 2**: Data object

a.      **Level 2A**: Data object headers

b.      **Level 2B**: Shared data object headers

c.      **Level 2C:** Data object data storage

The Level 0 specification defines the header block for the file. Header block elements include a signature, version information, key parameters of the file layout (such as which VFL file drivers are needed), and pointers to the rest of the file. Level 1 defines the data structures used throughout the file: the B-trees, heaps, and groups. Level 2 defines the data structure for storing the data objects and data. In all cases, the data structures are completely specified so that every bit in the file can be faithfully interpreted.

It is important to realize that the structures defined in the HDF5 file format are not the same as the abstract data model: the object headers, heaps, and B-trees of the file specification are not represented in the abstract data model. The format defines a number of objects for managing the storage including header blocks, B-trees, and heaps. The HDF5 File Format Specification defines how the abstract objects (for example, groups and datasets) are represented as headers, B-tree blocks, and other elements.

The HDF5 library implements operations to write HDF5 objects to the linear format and to read from the linear format to create HDF5 objects. It is important to realize that a single HDF5 abstract object is usually stored as several objects. A dataset, for example, might be stored in a header and in one or more data blocks, and these objects might not be contiguous on the hard disk.

### Concrete Storage Model

The HDF5 file format defines an abstract linear address space. This can be implemented in different stor­age media such as a single file or multiple files on disk or in memory. The HDF5 Library defines an open interface called the Virtual File Layer (VFL). The VFL allows different concrete storage models to be selected.

The VFL defines an abstract model, an API for random access storage, and an API to plug in alternative VFL driver modules. The model defines the operations that the VFL driver must and may support, and the plug-in API enables the HDF5 library to recognize the driver and pass it control and data.

A number of VFL drivers have been defined in the HDF5 library. Some work with a single file, and some work with multiple files split in various ways. Some work in serial computing environments, and some work in parallel computing environments. Most work with disk copies of HDF5 files, but one works with a memory copy. These drivers are listed in the “Supported file drivers” table. For more information, see Section 3.11, *Alternate File Storage Layouts and Low-level File Drivers*.

Each driver isolates the details of reading and writing storage so that the rest of the HDF5 library and user program can be almost the same for different storage methods. The exception to this rule is that some VFL drivers need information from the calling application. This information is passed using property lists. For example, the Parallel driver requires certain control information that must be provided by the application.

## The Structure of an HDF5 File

### Overall File Structure

An HDF5 file is organized as a rooted, directed graph. Named data objects are the nodes of the graph, and links are the directed arcs. Each arc of the graph has a name, and the root group has the name “/”. Objects are created and then inserted into the graph with the link operation which creates a named link from a group to the object. For example, the figure below illustrates the structure of an HDF5 file when one data­set is created. An object can be the target of more than one link. The names on the links must be unique within each group, but there may be many links with the same name in different groups. Link names are unambiguous: some ancestor will have a different name, or they are the same object. The graph is navi­gated with path names similar to Unix file systems. An object can be opened with a full path starting at the root group or with a relative path and a starting node (group). Note that all paths are relative to a single HDF5 file. In this sense, an HDF5 file is analogous to a single Unix file system.[2](https://bitbucket.hdfgroup.org/pages/HDFFV/hdf5doc/master/browse/html/UG/HDF5_Users_Guide-Responsive%20HTML5/HDF5_Users_Guide/DataModelAndFileStructure/The_HDF5_Data_Model_and_File_Structure.htm" \l "FNH_1)

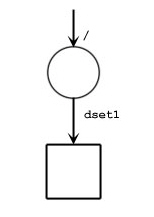
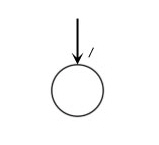


Figure 1-12. An HDF5 file with one dataset

Note: In the figure above are two figures. The top figure represents a newly created file with one group, /. In the bot­tom figure, a dataset called /dset1 has been created.

It is important to note that, just like the Unix file system, HDF5 objects do not have names. The names are associated with paths. An object has a unique (within the file) object identifier, but a single object may have many names because there may be many paths to the same object. An object can be renamed (moved to another group) by adding and deleting links. In this case, the object itself never moves. For that matter, membership in a group has no implication for the physical location of the stored object.

Deleting a link to an object does not necessarily delete the object. The object remains available as long as there is at least one link to it. After all the links to an object are deleted, it can no longer be opened although the storage may or may not be reclaimed.3

It is important to realize that the linking mechanism can be used to construct very complex graphs of objects. For example, it is possible for an object to be shared between several groups and even to have more than one name in the same group. It is also possible for a group to be a member of itself or to be in a “cycle” in the graph. An example of a cycle is where a child is the parent of one of its own ancestors.

### HDF5 Path Names and Navigation

The structure of the file constitutes the name space for the objects in the file. A path name is a string of components separated by ‘/’. Each component is the name of a link or the special character “.” for the cur­rent group. Link names (components) can be any string of ASCII characters not containing ‘/’ (except the string “.” which is reserved). However, users are advised to avoid the use of punctuation and non-printing characters because they may create problems for other software. The figure below gives a BNF grammar for HDF5 path names.

PathName ::= AbsolutePathName | RelativePathName

Separator ::= "/" ["/"]\*

AbsolutePathName ::= Separator [ RelativePathName ]

RelativePathName ::= Component [ Separator RelativePathName ]\*

Component ::=  "." |  Name

Name ::= Character+  -  {"."}

Character ::= {c: c in {{ legal ASCII characters }  -  {'/'}}

Figure 1-13. A BNF grammar for path names

An object can always be addressed by a full or absolute path which would start at the root group. As already noted, a given object can have more than one full path name. An object can also be addressed by a relative path which would start at a group and include the path to the object.

The structure of an HDF5 file is “self-describing.” This means that it is possible to navigate the file to dis­cover all the objects in the file. Basically, the structure is traversed as a graph starting at one node and recursively visiting the nodes of the graph.

### Examples of HDF5 File Structures

The figures below show some possible HDF5 file structures with groups and datasets. The first figure shows the structure of a file with three groups. The second shows a dataset created in “/group1”. The third figure shows the structure after a dataset called dset2 has been added to the root group. The fourth figure shows the structure after another group and dataset have been added.

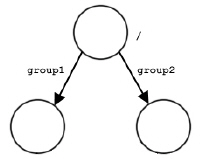


Figure 1-14. An HDF5 file structure with groups

Note: The figure above shows three groups; /group1 and /group2 are members of the root group.

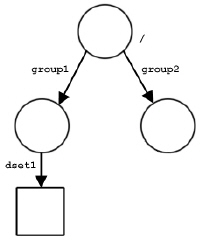


Figure 1-15. An HDF5 file structure with groups and a dataset

Note: The figure above shows that a dataset has been created in /group1: /group1/dset1.

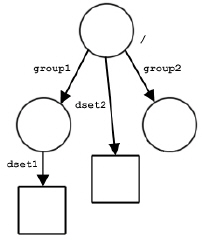


Figure 1-16. An HDF5 file structure with groups and datasets

Note: In the figure above, another dataset has been added as a member of the root group: /dset2.

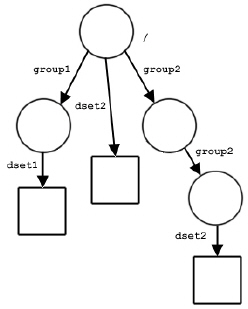


Figure 1-17. Another HDF5 file structure with groups and datasets

Note: In the figure above, another group and dataset have been added reusing object names: /group2/group2/dset2.

1. HDF5 requires random access to the linear address space. For this reason it is not well suited for some data media such as streams.

2.   It could be said that HDF5 extends the organizing concepts of a file system to the internal structure of a single file.

3.   As of HDF5-1.4, the storage used for an object is reclaimed, even if all links are deleted.

# The HDF5 Library and Programming Model

## Introduction

The HDF5 library implements the HDF5 abstract data model and storage model. These models were described in the preceding chapter.

Two major objectives of the HDF5 products are to provide tools that can be used on as many computa­tional platforms as possible (portability), and to provide a reasonably object-oriented data model and pro­gramming interface.

To be as portable as possible, the HDF5 library is implemented in portable C. C is not an object-oriented language, but the library uses several mechanisms and conventions to implement an object model.

One mechanism the HDF5 library uses is to implement the objects as data structures. To refer to an object, the HDF5 library implements its own pointers. These pointers are called identifiers. An identifier is then used to invoke operations on a specific instance of an object. For example, when a group is opened, the API returns a group identifier. This identifier is a reference to that specific group and will be used to invoke future operations on that group. The identifier is valid only within the context it is created and remains valid until it is closed or the file is closed. This mechanism is essentially the same as the mecha­nism that C++ or other object-oriented languages use to refer to objects except that the syntax is C.

Similarly, object-oriented languages collect all the methods for an object in a single name space. An exam­ple is the methods of a C++ class. The C language does not have any such mechanism, but the HDF5 library simulates this through its API naming convention. API function names begin with a common prefix that is related to the class of objects that the function operates on. The table below lists the HDF5 objects and the standard prefixes used by the corresponding HDF5 APIs. For example, functions that operate on data­type objects all have names beginning with H5T.

|  |  |
| --- | --- |
| Table 2-1. The HDF5 API naming scheme | |
| **Prefix** | **Operates on** |
| H5A | Attributes |
| H5D | Datasets |
| H5E | Error reports |
| H5F | Files |
| H5G | Groups |
| H5I | Identifiers |
| H5L | Links |
| H5O | Objects |
| H5P | Property lists |
| H5R | References |
| H5S | Dataspaces |
| H5T | Datatypes |
| H5Z | Filters |

## The HDF5 Programming Model

In this section we introduce the HDF5 programming model by means of a series of short code samples. These samples illustrate a broad selection of common HDF5 tasks. More details are provided in the follow­ing chapters and in the HDF5 Reference Manual.

### Creating an HDF5 File

Before an HDF5 file can be used or referred to in any manner, it must be explicitly created or opened. When the need for access to a file ends, the file must be closed. The example below provides a C code fragment illustrating these steps. In this example, the values for the file creation property list and the file access property list are set to the defaults H5P\_DEFAULT.

hid\_t       file;                 /\* declare file identifier \*/

/\*

\* Create a new file using H5F\_ACC\_TRUNC

\* to truncate and overwrite any file of the same name,

\* default file creation properties, and

\* default file access properties.

\* Then close the file.

\*/

file = H5Fcreate(FILE, H5F\_ACC\_TRUNC, H5P\_DEFAULT, H5P\_DEFAULT);

status = H5Fclose(file);

Code Example 2-1. Creating and closing an HDF5 file

Note: If there is a possibility that a file of the declared name already exists and you wish to open a new file regardless of that possibility, the flag H5F\_ACC\_TRUNC will cause the operation to overwrite the previous file. If the operation should fail in such a circumstance, use the flag H5F\_ACC\_EXCL instead.

### Creating and Initializing a Dataset

The essential objects within a dataset are datatype and dataspace. These are independent objects and are created separately from any dataset to which they may be attached. Hence, creating a dataset requires, at a minimum, the following steps:

1.      Create and initialize a dataspace for the dataset

2.      Define a datatype for the dataset

3.      Create and initialize the dataset

The code in the example below illustrates the execution of these steps.

|  |
| --- |
| hid\_t    dataset, datatype, dataspace;  /\* declare identifiers \*/    /\*  \* Create a dataspace: Describe the size of the array and  \* create the dataspace for a fixed-size dataset.  \*/  dimsf[0] = NX;  dimsf[1] = NY;  dataspace = H5Screate\_simple(RANK, dimsf, NULL); |
| /\*  \* Define a datatype for the data in the dataset.  \* We will store little endian integers.  \*/  datatype = H5Tcopy(H5T\_NATIVE\_INT);  status = H5Tset\_order(datatype, H5T\_ORDER\_LE); |
| /\*  \* Create a new dataset within the file using the defined  \* dataspace and datatype and default dataset creation  \* properties.  \* NOTE: H5T\_NATIVE\_INT can be used as the datatype if  \* conversion to little endian is not needed.  \*/  dataset = H5Dcreate(file, DATASETNAME, datatype, dataspace,  H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT); |
| Code Example 2-2. Create a dataset |

### Closing an Object

An application should close an object such as a datatype, dataspace, or dataset once the object is no lon­ger needed. Since each is an independent object, each must be released (or closed) separately. This action is frequently referred to as releasing the object’s identifier. The code in the example below closes the datatype, dataspace, and dataset that were created in the preceding section.

|  |
| --- |
| H5Tclose(datatype);  H5Dclose(dataset);  H5Sclose(dataspace); |
| Code Example 2-3. Close an object |

There is a long list of HDF5 library items that return a unique identifier when the item is created or opened. Each time that one of these items is opened, a unique identifier is returned. Closing a file does not mean that the groups, datasets, or other open items are also closed. Each opened item must be closed separately.

For more information, see Using Identifiers in the *HDF5 Application Developer’s Guide* under General Topics in HDF5.

#### How Closing a File Effects Other Open Structural Elements

Every structural element in an HDF5 file can be opened, and these elements can be opened more than once. Elements range in size from the entire file down to attributes. When an element is opened, the HDF5 library returns a unique identifier to the application. Every element that is opened must be closed. If an element was opened more than once, each identifier that was returned to the application must be closed. For example, if a dataset was opened twice, both dataset identifiers must be released (closed) before the dataset can be considered closed. Suppose an application has opened a file, a group in the file, and two datasets in the group. In order for the file to be totally closed, the file, group, and datasets must each be closed. Closing the file before the group or the datasets will not affect the state of the group or datasets: the group and datasets will still be open.

There are several exceptions to the above general rule. One is when the H5close function is used. H5close causes a general shutdown of the library: all data is written to disk, all identifiers are closed, and all memory used by the library is cleaned up. Another exception occurs on parallel processing systems. Suppose on a parallel system an application has opened a file, a group in the file, and two datasets in the group. If the application uses the H5Fclose function to close the file, the call will fail with an error. The open group and datasets must be closed before the file can be closed. A third exception is when the file access property list includes the property H5F\_CLOSE\_STRONG. This property closes any open elements when the file is closed with H5Fclose. For more information, see the H5Pset\_fclose\_degree func­tion in the HDF5 Reference Manual.

### Writing or Reading a Dataset to or from a File

Having created the dataset, the actual data can be written with a call to H5Dwrite. See the example below.

|  |
| --- |
| /\*  \* Write the data to the dataset using default transfer properties.  \*/  status = H5Dwrite(dataset, H5T\_NATIVE\_INT, H5S\_ALL, H5S\_ALL,  H5P\_DEFAULT, data); |
| Code Example 2-4. Writing a dataset |

Note that the third and fourth H5Dwrite parameters in the above example describe the dataspaces in memory and in the file, respectively. For now, these are both set to H5S\_ALL which indicates that the entire dataset is to be written. The selection of partial datasets and the use of differing dataspaces in memory and in storage will be discussed later in this chapter and in more detail elsewhere in this guide.

Reading the dataset from storage is similar to writing the dataset to storage. To read an entire dataset, substitute H5Dread for H5Dwrite in the above example.

### Reading and Writing a Portion of a Dataset

The previous section described writing or reading an entire dataset. HDF5 also supports access to portions of a dataset. These parts of datasets are known as selections.

The simplest type of selection is a simple hyperslab. This is an n-dimensional rectangular sub-set of a dataset where n is equal to the dataset’s rank. Other available selections include a more complex hyper­slab with user-defined stride and block size, a list of independent points, or the union of any of these.

The figure below shows several sample selections.

|  |
| --- |
| Pmodel_fig5_a.jpg |
| Pmodel_fig5_b.jpg |
| Pmodel_fig5_c.jpg |
| Pmodel_fig5_d.jpg |
| Pmodel_fig5_e.jpg  Figure 2‑1. Dataset selections |

Note: In the figure above, selections can take the form of a simple hyperslab, a hyperslab with user-defined stride and block, a selection of points, or a union of any of these forms.

Selections and hyperslabs are portions of a dataset. As described above, a simple hyperslab is a rectangu­lar array of data elements with the same rank as the dataset’s dataspace. Thus, a simple hyperslab is a log­ically contiguous collection of points within the dataset.

The more general case of a hyperslab can also be a regular pattern of points or blocks within the dataspace. Four parameters are required to describe a general hyperslab: the starting coordinates, the block size, the stride or space between blocks, and the number of blocks. These parameters are each expressed as a one-dimensional array with length equal to the rank of the dataspace and are described in the table below.

|  |  |
| --- | --- |
| Table 2-2. Hyperslab parameters | |
| **Parameter** | **Definition** |
| *start* | The coordinates of the starting location of the hyperslab in the dataset’s dataspace. |
| *block* | The size of each block to be selected from the dataspace. If the block param­eter is set to NULL, the block size defaults to a single element in each dimen­sion, as if the block array was set to all 1s (all ones). This will result in the selection of a uniformly spaced set of count points starting at start and on the interval defined by stride. |
| *stride* | The number of elements separating the starting point of each element or block to be selected. If the stride parameter is set to NULL, the stride size defaults to 1 (one) in each dimension and no elements are skipped. |
| *count* | The number of elements or blocks to select along each dimension. |

#### Reading Data into a Differently Shaped Memory Block

For maximum flexibility in user applications, a selection in storage can be mapped into a differently-shaped selection in memory. All that is required is that the two selections contain the same number of data elements. In this example, we will first define the selection to be read from the dataset in storage, and then we will define the selection as it will appear in application memory.

Suppose we want to read a 3 x 4 hyperslab from a two-dimensional dataset in a file beginning at the data­set element <1,2>. The first task is to create the dataspace that describes the overall rank and dimensions of the dataset in the file and to specify the position and size of the in-file hyperslab that we are extracting from that dataset. See the code below.

|  |
| --- |
| /\*  \* Define dataset dataspace in file.  \*/  dataspace = H5Dget\_space(dataset);    /\* dataspace identifier \*/  rank      = H5Sget\_simple\_extent\_ndims(dataspace);  status\_n  = H5Sget\_simple\_extent\_dims(dataspace, dims\_out, NULL); |
| /\*  \* Define hyperslab in the dataset.  \*/  offset[0] = 1;  offset[1] = 2;  count[0]  = 3;  count[1]  = 4;  status = H5Sselect\_hyperslab(dataspace, H5S\_SELECT\_SET, offset,  NULL, count, NULL); |

Code Example 2-5. Define the selection to be read from storage

The next task is to define a dataspace in memory. Suppose that we have in memory a three-dimensional 7 x 7 x 3 array into which we wish to read the two-dimensional 3 x 4 hyperslab described above and that we want the memory selection to begin at the element <3,0,0> and reside in the plane of the first two dimen­sions of the array. Since the in-memory dataspace is three-dimensional, we have to describe the in-mem­ory selection as three-dimensional. Since we are keeping the selection in the plane of the first two dimensions of the in-memory dataset, the in-memory selection will be a 3 x 4 x 1 array defined as <3,4,1>.

Notice that we must describe two things: the dimensions of the in-memory array, and the size and posi­tion of the hyperslab that we wish to read in. The code below illustrates how this would be done.

|  |  |
| --- | --- |
| /\*  \* Define memory dataspace.  \*/  dimsm[0] = 7;  dimsm[1] = 7;  dimsm[2] = 3;  memspace = H5Screate\_simple(RANK\_OUT,dimsm,NULL); | |
| /\*  \* Define memory hyperslab.  \*/  offset\_out[0] = 3;  offset\_out[1] = 0;  offset\_out[2] = 0;  count\_out[0]  = 3;  count\_out[1]  = 4;  count\_out[2]  = 1;  status = H5Sselect\_hyperslab(memspace, H5S\_SELECT\_SET, offset\_out, NULL, count\_out, NULL); | |
| Code Example 2-6. Define the memory dataspace and selection |

The hyperslab defined in the code above has the following parameters: start=(3,0,0), count=(3,4,1), stride and block size are NULL.

#### Writing Data into a Differently Shaped Disk Storage Block

Now let’s consider the opposite process of writing a selection from memory to a selection in a dataset in a file. Suppose that the source dataspace in memory is a 50-element, one-dimensional array called vec­tor and that the source selection is a 48-element simple hyperslab that starts at the second element of vector. See the figure below.

|  |
| --- |
| Pmodel_fig2.JPG  Figure 2‑2. A one-dimensional array |

Further suppose that we wish to write this data to the file as a series of 3 x 2-element blocks in a two-dimensional dataset, skipping one row and one column between blocks. Since the source selection con­tains 48 data elements and each block in the destination selection contains 6 data elements, we must define the destination selection with 8 blocks. We will write 2 blocks in the first dimension and 4 in the second. The code below shows how to achieve this objective.

|  |
| --- |
| /\* Select the hyperslab for the dataset in the file, using  \* 3 x 2 blocks, a (4,3) stride, a (2,4) count, and starting  \* at the position (0,1).  \*/ |
| start[0]  = 0; start[1]  = 1;  stride[0] = 4; stride[1] = 3;  count[0]  = 2; count[1]  = 4;  block[0]  = 3; block[1]  = 2;  ret = H5Sselect\_hyperslab(fid, H5S\_SELECT\_SET, start, stride,  count, block); |
| /\*  \* Create dataspace for the first dataset.  \*/  mid1 = H5Screate\_simple(MSPACE1\_RANK, dim1, NULL); |
| /\*  \* Select hyperslab.  \* We will use 48 elements of the vector buffer starting at the  \* second element. Selected elements are 1 2 3 . . . 48  \*/ |
| start[0]  = 1;  stride[0] = 1;  count[0]  = 48;  block[0]  = 1;  ret = H5Sselect\_hyperslab(mid1, H5S\_SELECT\_SET, start, stride,  count, block);    /\*  \* Write selection from the vector buffer to the dataset in the  \* file.  \*  ret = H5Dwrite(dataset, H5T\_NATIVE\_INT, mid1, fid, H5P\_DEFAULT,  vector) |
| Code Example 2-7. The destination selection |

### Getting Information about a Dataset

Although reading is analogous to writing, it is often first necessary to query a file to obtain information about the dataset to be read. For instance, we often need to determine the datatype associated with a dataset, or its dataspace (in other words, rank and dimensions). As illustrated in the code example below, there are several get routines for obtaining this information.

|  |
| --- |
| /\*  \* Get datatype and dataspace identifiers,  \* then query datatype class, order, and size, and  \* then query dataspace rank and dimensions.  \*/    datatype = H5Dget\_type (dataset);  /\* datatype identifier \*/  class = H5Tget\_class (datatype);  if (class == H5T\_INTEGER) printf("Dataset has INTEGER type \n");  order = H5Tget\_order (datatype);  if (order == H5T\_ORDER\_LE) printf("Little endian order \n"); |
| size = H5Tget\_size (datatype);  printf ("Size is %d \n", size);  dataspace = H5Dget\_space (dataset); /\* dataspace identifier \*/    /\* Find rank and retrieve current and maximum dimension  \* sizes.  \*/    rank = H5Sget\_simple\_extent\_dims (dataspace, dims, max\_dims); |
| Code Example 2-8. Routines to get dataset parameters |

### Creating and Defining Compound Datatypes

A compound datatype is a collection of one or more data elements. Each element might be an atomic type, a small array, or another compound datatype.

The provision for nested compound datatypes allows these structures to become quite complex. An HDF5 compound datatype has some similarities to a C struct or a Fortran common block. Though not originally designed with databases in mind, HDF5 compound datatypes are sometimes used in a way that is similar to a database record. Compound datatypes can become either a powerful tool or a complex and difficult-to-debug construct. Reasonable caution is advised.

To create and use a compound datatype, you need to create a datatype with class compound (H5T\_COM­POUND) and specify the total size of the data element in bytes. A compound datatype consists of zero or more uniquely named members. Members can be defined in any order but must occupy non-overlapping regions within the datum. The table below lists the properties of compound datatype members.

|  |  |
| --- | --- |
| Table 2-3. Compound datatype member properties | |
| **Parameter** | **Definition** |
| Index | An index number between zero and N-1, where N is the number of members in the compound. The elements are indexed in the order of their location in the array of bytes. |
| Name | A string that must be unique within the members of the same datatype. |
| Datatype | An HDF5 datatype. |
| Offset | A fixed byte offset which defines the location of the first byte of that member in the compound datatype. |

Properties of the members of a compound datatype are defined when the member is added to the com­pound type. These properties cannot be modified later.

**Defining Compound Datatypes**

Compound datatypes must be built out of other datatypes. To do this, you first create an empty com­pound datatype and specify its total size. Members are then added to the compound datatype in any order.

Each member must have a descriptive name. This is the key used to uniquely identify the member within the compound datatype. A member name in an HDF5 datatype does not necessarily have to be the same as the name of the corresponding member in the C struct in memory although this is often the case. You also do not need to define all the members of the C struct in the HDF5 compound datatype (or vice versa).

Usually a C struct will be defined to hold a data point in memory, and the offsets of the members in mem­ory will be the offsets of the struct members from the beginning of an instance of the struct. The library defines the macro that computes the offset of member m within a struct variable s:

HOFFSET(s,m)

The code below shows an example in which a compound datatype is created to describe complex num­bers whose type is defined by the complex\_t struct.

|  |
| --- |
| Typedef struct {  double re;   /\*real part \*/  double im;   /\*imaginary part \*/  } complex\_t;    complex\_t tmp;  /\*used only to compute offsets \*/  hid\_t complex\_id = H5Tcreate (H5T\_COMPOUND, sizeof tmp);  H5Tinsert (complex\_id, "real", HOFFSET(tmp,re),  H5T\_NATIVE\_DOUBLE);  H5Tinsert (complex\_id, "imaginary", HOFFSET(tmp,im),  H5T\_NATIVE\_DOUBLE); |
| Code Example 2-9. A compound datatype for complex numbers |

### Creating and Writing Extendable Datasets

An extendable dataset is one whose dimensions can grow. One can define an HDF5 dataset to have certain initial dimensions with the capacity to later increase the size of any of the initial dimensions. For example, the figure below shows a  3 x 3 dataset (a) which is later extended to be a 10 x 3 dataset by adding 7 rows (b), and further extended to be a 10 x 5 dataset by adding two columns (c).

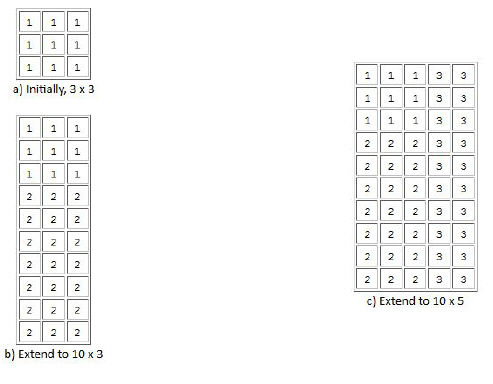


Figure 2‑3. Extending a dataset

HDF5 requires the use of chunking when defining extendable datasets. Chunking makes it possible to extend datasets efficiently without having to reorganize contiguous storage excessively.

To summarize, an extendable dataset requires two conditions:

1.      Define the dataspace of the dataset as unlimited in all dimensions that might eventually be extended

2.      Enable chunking in the dataset creation properties

For example, suppose we wish to create a dataset similar to the one shown in the figure above. We want to start with a 3 x 3 dataset, and then later we will extend it. To do this, go through the steps below.

First, declare the dataspace to have unlimited dimensions. See the code shown below. Note the use of the predefined constant H5S\_UNLIMITED to specify that a dimension is unlimited.

|  |
| --- |
| /\* dataset dimensions at creation time \*/  Hsize\_t dims[2] = {3, 3};  hsize\_t maxdims[2] = {H5S\_UNLIMITED, H5S\_UNLIMITED};    /\*  \* Create the data space with unlimited dimensions.  \*/  dataspace = H5Screate\_simple(RANK, dims, maxdims); |
| Code Example 2-10. Declaring a dataspace with unlimited dimensions |

 Next, set the dataset creation property list to enable chunking. See the code below.

|  |
| --- |
| hid\_t cparms;  hsize\_t chunk\_dims[2] ={2, 5};  /\*  \* Modify dataset creation properties to enable chunking.  \*/  cparms = H5Pcreate (H5P\_DATASET\_CREATE);  status = H5Pset\_chunk(cparms, RANK, chunk\_dims); |
| Code Example 2-11. Enable chunking |

The next step is to create the dataset. See the code below.

|  |
| --- |
| /\*  \* Create a new dataset within the file using cparms  \* creation properties.  \*/  dataset = H5Dcreate(file, DATASETNAME, H5T\_NATIVE\_INT, dataspace,  H5P\_DEFAULT, cparms, H5P\_DEFAULT); |
| Code Example 2-12. Create a dataset |

Finally, when the time comes to extend the size of the dataset, invoke H5Dextend. Extending the dataset along the first dimension by seven rows leaves the dataset with new dimensions of <10,3>. See the code below.

|  |
| --- |
| /\*  \* Extend the dataset. Dataset becomes 10 x 3.  \*/  dims[0] = dims[0] + 7;  size[0] = dims[0];  size[1] = dims[1];  status = H5Dextend (dataset, size); |
| Code Example 2-13. Extend the dataset by seven rows |

### Creating and Working with Groups

Groups provide a mechanism for organizing meaningful and extendable sets of datasets within an HDF5 file. The H5G API provides several routines for working with groups.

#### Creating a Group

With no datatype, dataspace, or storage layout to define, creating a group is considerably simpler than creating a dataset. For example, the following code creates a group called Data in the root group of file.

|  |
| --- |
| /\*  \*  Create a group in the file.  \*/  grp = H5Gcreate(file, "/Data", H5P\_DEFAULT, H5P\_DEFAULT,  H5P\_DEFAULT); |
| Code Example 2-14. Create a group |

A group may be created within another group by providing the absolute name of the group to the H5Gcreate function or by specifying its location. For example, to create the group Data\_new in the group Data, you might use the sequence of calls shown below.

|  |
| --- |
| /\*  \* Create group "Data\_new" in the group "Data" by specifying  \* absolute name of the group.  \*/  grp\_new = H5Gcreate(file, "/Data/Data\_new", H5P\_DEFAULT,  H5P\_DEFAULT, H5P\_DEFAULT);    or    /\*  \* Create group "Data\_new" in the "Data" group.  \*/  grp\_new = H5Gcreate(grp, "Data\_new", H5P\_DEFAULT, H5P\_DEFAULT,  H5P\_DEFAULT); |
| Code Example 2-15. Create a group within a group |

This first parameter of H5Gcreate is a location identifier. file in the first example specifies only the file. grp in the second example specifies a particular group in a particular file. Note that in this instance, the group identifier grp is used as the first parameter in the H5Gcreate call so that the relative name of Data\_new can be used.

The third parameter of H5Gcreate optionally specifies how much file space to reserve to store the names of objects that will be created in this group. If a non-positive value is supplied, the library provides a default size.

Use H5Gclose to close the group and release the group identifier.

#### Creating a Dataset within a Group

As with groups, a dataset can be created in a particular group by specifying either its absolute name in the file or its relative name with respect to that group. The next code excerpt uses the absolute name.

|  |
| --- |
| /\*  \* Create the dataset "Compressed\_Data" in the group Data using  \* the absolute name. The dataset creation property list is  \* modified to use GZIP compression with the compression  \* effort set to 6. Note that compression can be used only when  \* the dataset is chunked.  \*/ |
| dims[0] = 1000;  dims[1] = 20;  cdims[0] = 20;  cdims[1] = 20;  dataspace = H5Screate\_simple(RANK, dims, NULL);  plist = H5Pcreate(H5P\_DATASET\_CREATE);  H5Pset\_chunk(plist, 2, cdims);  H5Pset\_deflate(plist, 6);  dataset = H5Dcreate(file, "/Data/Compressed\_Data",  H5T\_NATIVE\_INT, dataspace, H5P\_DEFAULT, plist, H5P\_DEFAULT); |
| Code Example 2-16. Create a dataset within a group using an absolute name |

 Alternatively, you can first obtain an identifier for the group in which the dataset is to be created, and then create the dataset with a relative name.

|  |
| --- |
| /\*  \* Open the group.  \*/  grp = H5Gopen(file, "Data", H5P\_DEFAULT);    /\*  \* Create the dataset "Compressed\_Data" in the "Data" group  \* by providing a group identifier and a relative dataset  \* name as parameters to the H5Dcreate function.  \*/  dataset = H5Dcreate(grp, "Compressed\_Data", H5T\_NATIVE\_INT,  dataspace, H5P\_DEFAULT, plist, H5P\_DEFAULT); |
| Code Example 2-17. Create a dataset within a group using a relative name |

#### Accessing an Object in a Group

Any object in a group can be accessed by its absolute or relative name. The first code snippet below illus­trates the use of the absolute name to access the dataset Compressed\_Data in the group Data created in the examples above. The second code snippet illustrates the use of the relative name.

|  |
| --- |
| /\*  \* Open the dataset "Compressed\_Data" in the "Data" group.  \*/  dataset = H5Dopen(file, "/Data/Compressed\_Data", H5P\_DEFAULT); |
| Code Example 2-18. Accessing a group using its absolute name |

|  |
| --- |
| /\*  \* Open the group "data" in the file.  \*/  grp  = H5Gopen(file, "Data", H5P\_DEFAULT);    /\*  \* Access the "Compressed\_Data" dataset in the group.  \*/  dataset = H5Dopen(grp, "Compressed\_Data", H5P\_DEFAULT); |
| Code Example 2-19. Accessing a group using its relative name |

### Working with Attributes

An attribute is a small dataset that is attached to a normal dataset or group. Attributes share many of the characteristics of datasets, so the programming model for working with attributes is similar in many ways to the model for working with datasets. The primary differences are that an attribute must be attached to a dataset or a group and sub-setting operations cannot be performed on attributes.

To create an attribute belonging to a particular dataset or group, first create a dataspace for the attribute with the call to H5Screate, and then create the attribute using H5Acreate. For example, the code shown below creates an attribute called “Integer attribute” that is a member of a dataset whose iden­tifier is dataset. The attribute identifier is attr2. H5Awrite then sets the value of the attribute of that of the integer variable point. H5Aclose then releases the attribute identifier.

|  |
| --- |
| Int point = 1; /\* Value of the scalar attribute \*/ |
| /\*  \* Create scalar attribute.  \*/  aid2 = H5Screate(H5S\_SCALAR);  attr2 = H5Acreate(dataset, "Integer attribute", H5T\_NATIVE\_INT,  aid2, H5P\_DEFAULT, H5P\_DEFAULT); |
| /\*  \* Write scalar attribute.  \*/  ret = H5Awrite(attr2, H5T\_NATIVE\_INT, &point); |
| /\*  \* Close attribute dataspace.  \*/  ret = H5Sclose(aid2);    /\*  \* Close attribute.  \*/  ret = H5Aclose(attr2); |
| Code Example 2-20. Create an attribute |

|  |
| --- |
| /\*  \* Attach to the scalar attribute using attribute name, then  \* read and display its value.  \*/  attr = H5Aopen\_by\_name(file\_id, dataset\_name,  "Integer attribute", H5P\_DEFAULT, H5P\_DEFAULT);  ret = H5Aread(attr, H5T\_NATIVE\_INT, &point\_out);  printf("The value of the attribute \"Integer attribute\"  is %d \n", point\_out);  ret = H5Aclose(attr); |
| Code Example 2-21. Read a known attribute |

To read a scalar attribute whose name and datatype are known, first open the attribute using H5Aop­en\_by\_name, and then use H5Aread to get its value. For example, the code shown below reads a scalar attribute called “Integer attribute” whose datatype is a native integer and whose parent dataset has the identifier dataset.

To read an attribute whose characteristics are not known, go through these steps. First, query the file to obtain information about the attribute such as its name, datatype, rank, and dimensions, and then read the attribute. The following code opens an attribute by its index value using H5Aopen\_by\_idx, and then it reads in information about the datatype with H5Aread.

|  |
| --- |
| /\*  \* Attach to the string attribute using its index, then read and  \* display the value.  \*/  attr = H5Aopen\_by\_idx(file\_id, dataset\_name, index\_type,  iter\_order, 2, H5P\_DEFAULT, H5P\_DEFAULT);  atype = H5Tcopy(H5T\_C\_S1);  H5Tset\_size(atype, 4);  ret = H5Aread(attr, atype, string\_out);  printf("The value of the attribute with the index 2 is %s \n",  string\_out); |
| Code Example 2-22. Read an unknown attribute |

In practice, if the characteristics of attributes are not known, the code involved in accessing and process­ing the attribute can be quite complex. For this reason, HDF5 includes a function called H5Aiterate. This function applies a user-supplied function to each of a set of attributes. The user-supplied function can contain the code that interprets, accesses, and processes each attribute.

## The Data Transfer Pipeline

The HDF5 library implements data transfers between different storage locations. At the lowest levels, the HDF5 Library reads and writes blocks of bytes to and from storage using calls to the virtual file layer (VFL) drivers. In addition to this, the HDF5 library manages caches of metadata and a data I/O pipeline. The data I/O pipeline applies compression to data blocks, transforms data elements, and implements selections.

A substantial portion of the HDF5 library’s work is in transferring data from one environment or media to another. This most often involves a transfer between system memory and a storage medium. Data trans­fers are affected by compression, encryption, machine-dependent differences in numerical representa­tion, and other features. So, the bit-by-bit arrangement of a given dataset is often substantially different in the two environments.

Consider the representation on disk of a compressed and encrypted little-endian array as compared to the same array after it has been read from disk, decrypted, decompressed, and loaded into memory on a big-endian system. HDF5 performs all of the operations necessary to make that transition during the I/O pro­cess with many of the operations being handled by the VFL and the data transfer pipeline.

The figure below provides a simplified view of a sample data transfer with four stages. Note that the mod­ules are used only when needed. For example, if the data is not compressed, the compression stage is omitted.

|  |
| --- |
| Pmodel_fig26.JPG |
| Figure 2-4. A data transfer from storage to memory |

For a given I/O request, different combinations of actions may be performed by the pipeline. The library automatically sets up the pipeline and passes data through the processing steps. For example, for a read request (from disk to memory), the library must determine which logical blocks contain the requested data elements and fetch each block into the library’s cache. If the data needs to be decompressed, then the compression algorithm is applied to the block after it is read from disk. If the data is a selection, the selected elements are extracted from the data block after it is decompressed. If the data needs to be transformed (for example, byte swapped), then the data elements are transformed after decompression and selection.

While an application must sometimes set up some elements of the pipeline, use of the pipeline is nor­mally transparent to the user program. The library determines what must be done based on the metadata for the file, the object, and the specific request. An example of when an application might be required to set up some elements in the pipeline is if the application used a custom error-checking algorithm.

In some cases, it is necessary to pass parameters to and from modules in the pipeline or among other parts of the library that are not directly called through the programming API. This is accomplished through the use of dataset transfer and data access property lists.

The VFL provides an interface whereby user applications can add custom modules to the data transfer pipeline. For example, a custom compression algorithm can be used with the HDF5 Library by linking an appropriate module into the pipeline through the VFL. This requires creating an appropriate wrapper for the compression module and registering it with the library with H5Zregister. The algorithm can then be applied to a dataset with an H5Pset\_filter call which will add the algorithm to the selected dataset’s transfer property list.

# The HDF5 File

## Introduction

The purpose of this chapter is to describe how to work with HDF5 data files.

If HDF5 data is to be written to or read from a file, the file must first be explicitly created or opened with the appropriate file driver and access privileges. Once all work with the file is complete, the file must be explicitly closed.

This chapter discusses the following:

* File access modes
* Creating, opening, and closing files
* The use of file creation property lists
* The use of file access property lists
* The use of low-level file drivers

This chapter assumes an understanding of the material presented in the data model chapter. For more information, see "The HDF5 Data Model and File Structure."

## File Access Modes

There are two issues regarding file access:

* What should happen when a new file is created but a file of the same name already exists? Should the create action fail, or should the existing file be overwritten?
* Is a file to be opened with read-only or read-write access?

Four access modes address these concerns. Two of these modes can be used with H5Fcreate, and two modes can be used with H5Fopen.

* H5Fcreate accepts H5F\_ACC\_EXCL or H5F\_ACC\_TRUNC
* H5Fopen accepts H5F\_ACC\_RDONLY or H5F\_ACC\_RDWR

The access modes are described in the table below.

|  |  |
| --- | --- |
| Table 3-1. Access flags and modes | |
| **Access Flag** | **Resulting Access Mode** |
| H5F\_ACC\_EXCL | If the file already exists, H5Fcreate fails. If the file does not exist, it is created and opened with read-write access. (Default) |
| H5F\_ACC\_TRUNC | If the file already exists, the file is opened with read-write access, and new data will overwrite any existing data. If the file does not exist, it is created and opened with read-write access. |
| H5F\_ACC\_RDONLY | An existing file is opened with read-only access. If the file does not exist, H5Fopen fails. (Default) |
| H5F\_ACC\_RDWR | An existing file is opened with read-write access. If the file does not exist, H5Fopen fails. |

By default, H5Fopen opens a file for read-only access; passing H5F\_ACC\_RDWR allows read-write access to the file.

By default, H5Fcreate fails if the file already exists; only passing H5F\_ACC\_TRUNC allows the truncating of an existing file.

## File Creation and File Access Properties

File creation and file access property lists control the more complex aspects of creating and accessing files.

File creation property lists control the characteristics of a file such as the size of the userblock, a user-definable data block; the size of data address parameters; properties of the B-trees that are used to man­age the data in the file; and certain HDF5 Library versioning information.

For more information, see "File Creation Properties." This section has a more detailed discus­sion of file creation properties. If you have no special requirements for these file characteristics, you can simply specify H5P\_DEFAULT for the default file creation property list when a file creation property list is called for.

File access property lists control properties and means of accessing a file such as data alignment charac­teristics, metadata block and cache sizes, data sieve buffer size, garbage collection settings, and parallel I/O. Data alignment, metadata block and cache sizes, and data sieve buffer size are factors in improving I/O performance.

For more information, see "File Access Properties." This section has a more detailed discussion of file access properties. If you have no special requirements for these file access characteristics, you can simply specify H5P\_DEFAULT for the default file access property list when a file access property list is called for.

|  |
| --- |
| UML_FileAndProps.gif |
| Figure 3-1. UML model for an HDF5 file and its property lists |

## Low-level File Drivers

The concept of an HDF5 file is actually rather abstract: the address space for what is normally thought of as an HDF5 file might correspond to any of the following at the storage level:

•        Single file on a standard file system

•        Multiple files on a standard file system

•        Multiple files on a parallel file system

•        Block of memory within an application’s memory space

•        More abstract situations such as virtual files

This HDF5 address space is generally referred to as an HDF5 file regardless of its organization at the stor­age level.

HDF5 accesses a file (the address space) through various types of low-level file drivers. The default HDF5 file storage layout is as an unbuffered permanent file which is a single, contiguous file on local disk. Alter­native layouts are designed to suit the needs of a variety of systems, environments, and applications.

## Programming Model for Files

Programming models for creating, opening, and closing HDF5 files are described in the sub-sections below.

### Creating a New File

The programming model for creating a new HDF5 file can be summarized as follows:

•        Define the file creation property list

•        Define the file access property list

•        Create the file

First, consider the simple case where we use the default values for the property lists. See the example below.

|  |
| --- |
| file\_id = H5Fcreate ("SampleFile.h5", H5F\_ACC\_EXCL,  H5P\_DEFAULT, H5P\_DEFAULT) |
| Code Example 3-1. Creating an HDF5 file using property list defaults |

Note: The example above specifies that H5Fcreate should fail if SampleFile.h5 already exists.

A more complex case is shown in the example below. In this example, we define file creation and access property lists (though we do not assign any properties), specify that H5Fcreate should fail if Sample­File.h5 already exists, and create a new file named SampleFile.h5. The example does not specify a driver, so the default driver, H5FD\_SEC2, will be used.

|  |
| --- |
| fcplist\_id = H5Pcreate (H5P\_FILE\_CREATE)  <...set desired file creation properties...>  faplist\_id = H5Pcreate (H5P\_FILE\_ACCESS)  <...set desired file access properties...>  file\_id = H5Fcreate ("SampleFile.h5", H5F\_ACC\_EXCL,  fcplist\_id, faplist\_id) |
| Code Example 3-2. Creating an HDF5 file using property lists |

 Notes:

A root group is automatically created in a file when the file is first created.

File property lists, once defined, can be reused when another file is created within the same application.

3.5.2. Opening an Existing File

The programming model for opening an existing HDF5 file can be summarized as follows:

•        Define or modify the file access property list including a low-level file driver (optional)

•        Open the file

The code in the example below shows how to open an existing file with read-only access.

|  |
| --- |
| faplist\_id = H5Pcreate (H5P\_FILE\_ACCESS)  status = H5Pset\_fapl\_stdio (faplist\_id)  file\_id = H5Fopen ("SampleFile.h5", H5F\_ACC\_RDONLY,  faplist\_id) |
| Code Example 3-3. Opening an HDF5 file |

### Closing a File

The programming model for closing an HDF5 file is very simple:

•        Close file

We close SampleFile.h5 with the code in the example below.

|  |
| --- |
| status = H5Fclose (file\_id) |
| Code Example 3-4. Closing an HDF5 file |

Note that H5Fclose flushes all unwritten data to storage and that file\_id is the identifier returned for SampleFile.h5 by H5Fopen.

More comprehensive discussions regarding all of these steps are provided below.

## Using h5dump to View a File

h5dump is a command-line utility that is included in the HDF5 distribution. This program provides a straight-forward means of inspecting the contents of an HDF5 file. You can use h5dump to verify that a program is generating the intended HDF5 file. h5dump displays ASCII output formatted according to the HDF5 DDL grammar.

The following h5dump command will display the contents of SampleFile.h5:

h5dump SampleFile.h5

If no datasets or groups have been created in and no data has been written to the file, the output will look something like the following:

HDF5 "SampleFile.h5" {

GROUP "/" {

}

}

Note that the root group, indicated above by /, was automatically created when the file was created.

h5dump is described on the [Tools](https://portal.hdfgroup.org/display/HDF5/h5dump) page under [Libraries and Tools Reference](https://portal.hdfgroup.org/display/HDF5/Libraries+and+Tools+Reference). The HDF5 DDL grammar is described in the document [DDL in BNF for HDF5](https://portal.hdfgroup.org/display/HDF5/DDL+in+BNF+for+HDF5).

## File Function Summaries

General library functions and macros (H5), file functions (H5F), file related property list functions (H5P), and file driver functions (H5P) are listed below.

|  |  |
| --- | --- |
| Function Listing 3-1. General library functions and macros (H5) | |
| **C Function**  **Fortran Function** | **Purpose** |
| H5check\_version  h5check\_version\_f | Verifies that HDF5 library versions are consis­tent. |
| H5close  h5close\_f | Flushes all data to disk, closes all open identi­fiers, and cleans up memory. |
| H5dont\_atexit  h5dont\_atexit\_f | Instructs the library not to install the atexit cleanup routine. |
| H5garbage\_collect  h5garbage\_collect\_f | Garbage collects on all free-lists of all types. |
| H5get\_libversion  h5get\_libversion\_f | Returns the HDF library release number. |
| H5open  h5open\_f | Initializes the HDF5 library. |
| H5set\_free\_list\_limits  h5set\_free\_list\_limits\_f | Sets free-list size limits. |
| H5\_VERSION\_GE  (no Fortran subroutine) | Determines whether the version of the library being used is greater than or equal to the specified version. |
| H5\_VERSION\_LE  (no Fortran subroutine) | Determines whether the version of the library being used is less than or equal to the speci­fied version. |

|  |  |
| --- | --- |
| Function Listing 3-2. File functions (H5F) | |
| **C Function**  **Fortran Function** | **Purpose** |
| H5Fclear\_elink\_file\_cache  (no Fortran subroutine) | Clears the external link open file cache for a file. |
| H5Fclose  h5fclose\_f | Closes HDF5 file. |
| H5Fcreate  h5fcreate\_f | Creates new HDF5 file. |
| H5Fflush  h5fflush\_f | Flushes data to HDF5 file on storage medium. |
| H5Fget\_access\_plist  h5fget\_access\_plist\_f | Returns a file access property list identifier. |
| H5Fget\_create\_plist  h5fget\_create\_plist\_f | Returns a file creation property list identifier. |
| H5Fget\_file\_image  h5fget\_file\_image\_f | Retrieves a copy of the image of an existing, open file. |
| H5Fget\_filesize  h5fget\_filesize\_f | Returns the size of an HDF5 file. |
| H5Fget\_freespace  h5fget\_freespace\_f | Returns the amount of free space in a file. |
| H5Fget\_info  (no Fortran subroutine) | Returns global information for a file. |
| H5Fget\_intent  (no Fortran subroutine) | Determines the read/write or read-only status of a file. |
| H5Fget\_mdc\_config  (no Fortran subroutine) | Obtain current metadata cache configuration for target file. |
| H5Fget\_mdc\_hit\_rate  (no Fortran subroutine) | Obtain target file’s metadata cache hit rate. |
| H5Fget\_mdc\_size  (no Fortran subroutine) | Obtain current metadata cache size data for specified file. |
| H5Fget\_mpi\_atomicity  h5fget\_mpi\_atomicity\_f | Retrieves the atomicity mode in use. |
| H5Fget\_name  h5fget\_name\_f | Retrieves the name of the file to which the object belongs. |
| H5Fget\_obj\_count  h5fget\_obj\_count\_f | Returns the number of open object identifiers for an open file. |
| H5Fget\_obj\_ids  h5fget\_obj\_ids\_f | Returns a list of open object identifiers. |
| H5Fget\_vfd\_handle  (no Fortran subroutine) | Returns pointer to the file handle from the virtual file driver. |
| H5Fis\_hdf5  h5fis\_hdf5\_f | Determines whether a file is in the HDF5 for­mat. |
| H5Fmount  h5fmount\_f | Mounts a file. |
| H5Fopen  h5fopen\_f | Opens existing HDF5 file. |
| H5Freopen  h5freopen\_f | Returns a new identifier for a previously-opened HDF5 file. |
| H5Freset\_mdc\_hit\_rate\_stats  (no Fortran subroutine) | Reset hit rate statistics counters for the target file. |
| H5Fset\_mdc\_config  (no Fortran subroutine) | Use to configure metadata cache of target file. |
| H5Fset\_mpi\_atomicity  h5fset\_mpi\_atomicity\_f | Use to set the MPI atomicity mode. |
| H5Funmount  h5funmount\_f | Unmounts a file. |

|  |  |
| --- | --- |
| Function Listing 3-3. File creation property list functions (H5P) | |
| **C Function**  **Fortran Function** | **Purpose** |
| H5Pset/get\_userblock  h5pset/get\_userblock\_f | Sets/retrieves size of userblock. |
| H5Pset/get\_sizes  h5pset/get\_sizes\_f | Sets/retrieves byte size of offsets and lengths used to address objects in HDF5 file. |
| H5Pset/get\_sym\_k  h5pset/get\_sym\_k\_f | Sets/retrieves size of parameters used to con­trol symbol table nodes. |
| H5Pset/get\_istore\_k  h5pset/get\_istore\_k\_f | Sets/retrieves size of parameter used to con­trol B-trees for indexing chunked datasets. |
| H5Pget\_file\_image  h5pget\_file\_image\_f | Retrieves a copy of the file image designated as the initial content and structure of a file. |
| H5Pset\_file\_image  h5pset\_file\_image\_f | Sets an initial file image in a memory buffer. |
| H5Pset\_shared\_mesg\_nindexes  h5pset\_shared\_mesg\_nindexes\_f | Sets number of shared object header mes­sage indexes. |
| H5Pget\_shared\_mesg\_nindexes  (no Fortran subroutine) | Retrieves number of shared object header message indexes in file creation property list. |
| H5Pset\_shared\_mesg\_index  h5pset\_shared\_mesg\_index\_f | Configures the specified shared object header message index. |
| H5Pget\_shared\_mesg\_index  (no Fortran subroutine) | Retrieves the configuration settings for a shared message index. |
| H5Pset\_shared\_mesg\_phase\_change  (no Fortran subroutine) | Sets shared object header message storage phase change thresholds. |
| H5Pget\_shared\_mesg\_phase\_change  (no Fortran subroutine) | Retrieves shared object header message phase change information. |
| H5Pget\_version  h5pget\_version\_f | Retrieves version information for various objects for file creation property list. |

|  |  |
| --- | --- |
| Function Listing 3-4. File access property list functions (H5P) | |
| **C Function**  **Fortran Function** | **Purpose** |
| H5Pset/get\_alignment  h5pset/get\_alignment\_f | Sets/retrieves alignment properties. |
| H5Pset/get\_cache  h5pset/get\_cache\_f | Sets/retrieves metadata cache and raw data chunk cache parameters. |
| H5Pset/get\_elink\_file\_cache\_size  (no Fortran subroutine) | Sets/retrieves the size of the external link open file cache from the specified file access property list. |
| H5Pset/get\_fclose\_degree  h5pset/get\_fclose\_degree\_f | Sets/retrieves file close degree property. |
| H5Pset/get\_gc\_references  h5pset/get\_gc\_references\_f | Sets/retrieves garbage collecting references flag. |
| H5Pset\_family\_offset  h5pset\_family\_offset\_f | Sets offset property for low-level access to a file in a family of files. |
| H5Pget\_family\_offset  (no Fortran subroutine) | Retrieves a data offset from the file access property list. |
| H5Pset/get\_meta\_block\_size  h5pset/get\_meta\_block\_size\_f | Sets the minimum metadata block size or retrieves the current metadata block size set­ting. |
| H5Pset\_mdc\_config  (no Fortran subroutine) | Set the initial metadata cache configuration in the indicated File Access Property List to the supplied value. |
| H5Pget\_mdc\_config  (no Fortran subroutine) | Get the current initial metadata cache config­uration from the indicated File Access Prop­erty List. |
| H5Pset/get\_sieve\_buf\_size  h5pset/get\_sieve\_buf\_size\_f | Sets/retrieves maximum size of data sieve buffer. |
| H5Pset\_libver\_bounds  h5pset\_libver\_bounds\_f | Sets bounds on library versions, and indirectly format versions, to be used when creating objects. |
| H5Pget\_libver\_bounds  (no Fortran subroutine) | Retrieves library version bounds settings that indirectly control the format versions used when creating objects. |
| H5Pset\_small\_data\_block\_size  h5pset\_small\_data\_block\_size\_f | Sets the size of a contiguous block reserved for small data. |
| H5Pget\_small\_data\_block\_size  h5pget\_small\_data\_block\_size\_f | Retrieves the current small data block size setting. |

|  |  |
| --- | --- |
| Function Listing 3-5. File driver functions (H5P) | |
| **C Function**  **Fortran Function** | **Purpose** |
| H5Pset\_driver  (no Fortran subroutine) | Sets a file driver. |
| H5Pget\_driver  h5pget\_driver\_f | Returns the identifier for the driver used to create a file. |
| H5Pget\_driver\_info  (no Fortran subroutine) | Returns a pointer to file driver information. |
| H5Pset/get\_fapl\_core  h5pset/get\_fapl\_core\_f | Sets the driver for buffered memory files (in RAM) or retrieves information regarding the driver. |
| H5Pset\_fapl\_direct  h5pset\_fapl\_direct\_f | Sets up use of the direct I/O driver. |
| H5Pget\_fapl\_direct  h5pget\_fapl\_direct\_f | Retrieves the direct I/O driver settings. |
| H5Pset/get\_fapl\_family  h5pset/get\_fapl\_family\_f | Sets driver for file families, designed for sys­tems that do not support files larger than 2 gigabytes, or retrieves information regarding driver. |
| H5Pset\_fapl\_log  (no Fortran subroutine) | Sets logging driver. |
| H5Pset/get\_fapl\_mpio  h5pset/get\_fapl\_mpio\_f | Sets driver for files on parallel file systems (MPI I/O) or retrieves information regarding the driver. |
| H5Pset\_fapl\_mpiposix  h5pset\_fapl\_mpiposix\_f | No longer available. |
| H5Pget\_fapl\_mpiposix  h5pget\_fapl\_mpiposix\_f | No longer available. |
| H5Pset/get\_fapl\_multi  h5pset/get\_fapl\_multi\_f | Sets driver for multiple files, separating cate­gories of metadata and raw data, or retrieves information regarding driver. |
| H5Pset\_fapl\_sec2  h5pset\_fapl\_sec2\_f | Sets driver for unbuffered permanent files or retrieves information regarding driver. |
| H5Pset\_fapl\_split  h5pset\_fapl\_split\_f | Sets driver for split files, a limited case of mul­tiple files with one metadata file and one raw data file. |
| H5Pset\_fapl\_stdio  H5Pset\_fapl\_stdio\_f | Sets driver for buffered permanent files. |
| H5Pset\_fapl\_windows  (no Fortran subroutine) | Sets the Windows I/O driver. |
| H5Pset\_multi\_type  (no Fortran subroutine) | Specifies type of data to be accessed via the MULTI driver enabling more direct access. |
| H5Pget\_multi\_type  (no Fortran subroutine) | Retrieves type of data property for MULTI driver. |

## Creating or Opening an HDF5 File

This section describes in more detail how to create and how to open files.

New HDF5 files are created and opened with H5Fcreate; existing files are opened with H5Fopen. Both functions return an object identifier which must eventually be released by calling H5Fclose.

To create a new file, call H5Fcreate:

hid\_t H5Fcreate (const char \*name, unsigned flags, hid\_t fcpl\_id,     hid\_t fapl\_id)

H5Fcreate creates a new file named name in the current directory. The file is opened with read and write access; if the H5F\_ACC\_TRUNC flag is set, any pre-existing file of the same name in the same directory is truncated. If H5F\_ACC\_TRUNC is not set or H5F\_ACC\_EXCL is set and if a file of the same name exists, H5Fcreate will fail.

The new file is created with the properties specified in the property lists fcpl\_id and fapl\_id. fcpl is short for file creation property list. fapl is short for file access property list. Specifying H5P\_DEFAULT for either the creation or access property list will use the library’s default creation or access properties.

If H5Fcreate successfully creates the file, it returns a file identifier for the new file. This identifier will be used by the application any time an object identifier, an OID, for the file is required. Once the application has finished working with a file, the identifier should be released and the file closed with H5Fclose.

To open an existing file, call H5Fopen:

hid\_t H5Fopen (const char \*name, unsigned flags, hid\_t fapl\_id)

H5Fopen opens an existing file with read-write access if H5F\_ACC\_RDWR is set and read-only access if H5F\_ACC\_RDONLY is set.

fapl\_id is the file access property list identifier. Alternatively, H5P\_DEFAULT indicates that the applica­tion relies on the default I/O access parameters. Creating and changing access property lists is docu­mented further below.

A file can be opened more than once via multiple H5Fopen calls. Each such call returns a unique file iden­tifier and the file can be accessed through any of these file identifiers as long as they remain valid. Each of these file identifiers must be released by calling H5Fclose when it is no longer needed.

For more information, see "File Access Modes."

For more information, see "File Property Lists."

## Closing an HDF5 File

H5Fclose both closes a file and releases the file identifier returned by H5Fopen or H5Fcreate. H5F­close must be called when an application is done working with a file; while the HDF5 Library makes every effort to maintain file integrity, failure to call H5Fclose may result in the file being abandoned in an incomplete or corrupted state.

To close a file, call H5Fclose:

herr\_t H5Fclose (hid\_t file\_id)

This function releases resources associated with an open file. After closing a file, the file identifier, file\_id, cannot be used again as it will be undefined.

H5Fclose fulfills three purposes: to ensure that the file is left in an uncorrupted state, to ensure that all data has been written to the file, and to release resources. Use [H5Fflush](http://www.hdfgroup.org/HDF5/doc/RM/RM_H5F.html#File-Flush) if you wish to ensure that all data has been written to the file but it is premature to close it.

Note regarding serial mode behavior: When H5Fclose is called in serial mode, it closes the file and termi­nates new access to it, but it does not terminate access to objects that remain individually open within the file. That is, if H5Fclose is called for a file but one or more objects within the file remain open, those objects will remain accessible until they are individually closed. To illustrate, assume that a file, fileA, contains a dataset, data\_setA, and that both are open when H5Fclose is called for fileA. data\_setA will remain open and accessible, including writable, until it is explicitly closed. The file will be automati­cally and finally closed once all objects within it have been closed.

Note regarding parallel mode behavior: Once H5Fclose has been called in parallel mode, access is no longer available to any object within the file.

## File Property Lists

Additional information regarding file structure and access are passed to H5Fcreate and H5Fopen through property list objects. Property lists provide a portable and extensible method of modifying file properties via simple API functions. There are two kinds of file-related property lists:

•        File creation property lists

•        File access property lists

In the following sub-sections, we discuss only one file creation property, userblock size, in detail as a model for the user. Other file creation and file access properties are mentioned and defined briefly, but the model is not expanded for each; complete syntax, parameter, and usage information for every prop­erty list function is provided in the "H5P: Property List Interface" section of the HDF5 Reference Manual. For more information, see "Properties and Property Lists in HDF5."

### Creating a Property List

If you do not wish to rely on the default file creation and access properties, you must first create a prop­erty list with H5Pcreate.

hid\_t H5Pcreate (hid\_t cls\_id)

type is the type of property list being created. In this case, the appropriate values are H5P\_FILE\_CRE­ATE for a file creation property list and H5P\_FILE\_ACCESS for a file access property list.

Thus, the following calls create a file creation property list and a file access property list with identifiers fcpl\_id and fapl\_id, respectively:

fcpl\_id = H5Pcreate (H5P\_FILE\_CREATE)

fapl\_id = H5Pcreate (H5P\_FILE\_ACCESS)

Once the property lists have been created, the properties themselves can be modified via the functions described in the following sub-sections.

### File Creation Properties

File creation property lists control the file metadata, which is maintained in the superblock of the file. These properties are used only when a file is first created.

#### Userblock Size

herr\_t H5Pset\_userblock (hid\_t plist, hsize\_t size)

herr\_t H5Pget\_userblock (hid\_t plist, hsize\_t \*size)

The userblock is a fixed-length block of data located at the beginning of the file and is ignored by the HDF5 library. This block is specifically set aside for any data or information that developers determine to be use­ful to their applications but that will not be used by the HDF5 library. The size of the userblock is defined in bytes and may be set to any power of two with a minimum size of 512 bytes. In other words, userblocks might be 512, 1024, or 2048 bytes in size.

This property is set with H5Pset\_userblock and queried via H5Pget\_userblock. For example, if an application needed a 4K userblock, then the following function call could be used:

status = H5Pset\_userblock(fcpl\_id, 4096)

The property list could later be queried with

status = H5Pget\_userblock(fcpl\_id, size)

and the value 4096 would be returned in the parameter size.

Other properties, described below, are set and queried in exactly the same manner. Syntax and usage are detailed in the "H5P: Property List Interface" section of the HDF5 Reference Manual.

#### Offset and Length Sizes

This property specifies the number of bytes used to store the offset and length of objects in the HDF5 file. Values of 2, 4, and 8 bytes are currently supported to accommodate 16-bit, 32-bit, and 64-bit file address spaces.

These properties are set and queried via H5Pset\_sizes and H5Pget\_sizes.

#### Symbol Table Parameters

The size of symbol table B-trees can be controlled by setting the 1/2-rank and 1/2-node size parameters of the B-tree.

These properties are set and queried via H5Pset\_sym\_k and H5Pget\_sym\_k.

#### Indexed Storage Parameters

The size of indexed storage B-trees can be controlled by setting the 1/2-rank and 1/2-node size parameters of the B-tree.

These properties are set and queried via H5Pset\_istore\_k and H5Pget\_istore\_k.

#### Version Information

Various objects in an HDF5 file may over time appear in different versions. The HDF5 Library keeps track of the version of each object in the file.

Version information is retrieved via H5Pget\_version.

### File Access Properties

This section discusses file access properties that are not related to the low-level file drivers. File drivers are discussed separately later in this chapter. For more information, see "Alternate File Storage Layouts and Low-level File Drivers."

File access property lists control various aspects of file I/O and structure.

#### Data Alignment

Sometimes file access is faster if certain data elements are aligned in a specific manner. This can be con­trolled by setting alignment properties via the H5Pset\_alignment function. There are two values involved:

•        A threshold value

•        An alignment interval

Any allocation request at least as large as the threshold will be aligned on an address that is a multiple of the alignment interval.

#### Metadata Block Allocation Size

Metadata typically exists as very small chunks of data; storing metadata elements in a file without block­ing them can result in hundreds or thousands of very small data elements in the file. This can result in a highly fragmented file and seriously impede I/O. By blocking metadata elements, these small elements can be grouped in larger sets, thus alleviating both problems.

H5Pset\_meta\_block\_size sets the minimum size in bytes of metadata block allocations. H5Pget\_meta\_block\_size retrieves the current minimum metadata block allocation size.

#### Metadata Cache

Metadata and raw data I/O speed are often governed by the size and frequency of disk reads and writes. In many cases, the speed can be substantially improved by the use of an appropriate cache.

H5Pset\_cache sets the minimum cache size for both metadata and raw data and a preemption value for raw data chunks. H5Pget\_cache retrieves the current values.

#### Data Sieve Buffer Size

Data sieve buffering is used by certain file drivers to speed data I/O and is most commonly when working with dataset hyperslabs. For example, using a buffer large enough to hold several pieces of a dataset as it is read in for hyperslab selections will boost performance noticeably.

H5Pset\_sieve\_buf\_size sets the maximum size in bytes of the data sieve buffer. H5Pget\_sieve\_buf\_size retrieves the current maximum size of the data sieve buffer.

#### Garbage Collection References

Dataset region references and other reference types use space in an HDF5 file’s global heap. If garbage collection is on (1) and the user passes in an uninitialized value in a reference structure, the heap might become corrupted. When garbage collection is off (0), however, and the user re-uses a reference, the pre­vious heap block will be orphaned and not returned to the free heap space. When garbage collection is on, the user must initialize the reference structures to 0 or risk heap corruption.

H5Pset\_gc\_references sets the garbage collecting references flag.

## Alternate File Storage Layouts and Low-level File Drivers

The concept of an HDF5 file is actually rather abstract: the address space for what is normally thought of as an HDF5 file might correspond to any of the following:

•        Single file on standard file system

•        Multiple files on standard file system

•        Multiple files on parallel file system

•        Block of memory within application’s memory space

•        More abstract situations such as virtual files

This HDF5 address space is generally referred to as an HDF5 file regardless of its organization at the stor­age level.

HDF5 employs an extremely flexible mechanism called the virtual file layer, or VFL, for file I/O. A full understanding of the VFL is only necessary if you plan to write your own drivers (see "Virtual File Layer" and "List of VFL Functions" in the HDF5 Technical Notes). For our purposes here, it is sufficient to know that the low-level drivers used for file I/O reside in the VFL, as illustrated in the following figure. Note that H5FD\_STREAM is not available with 1.8.x and later versions of the library.

|  |
| --- |
| VFL_Drivers.jpg |
| Figure 3-2. I/O path from application to VFL and low-level drivers to storage |

As mentioned above, HDF5 applications access HDF5 files through various low-level file drivers. The default driver for that layout is the POSIX driver (also known as the SEC2 driver), H5FD\_SEC2. Alternative layouts and drivers are designed to suit the needs of a variety of systems, environments, and applications. The drivers are listed in the table below.

|  |  |  |  |
| --- | --- | --- | --- |
| Table 3-2. Supported file drivers | | | |
| **Driver Name** | **Driver**  **Identifier** | **Description** | **Related API** |
| POSIX | H5FD\_SEC2 | This driver uses POSIX file-system functions like read and write to perform I/O to a single, permanent file on local disk with no system buffering. This driver is POSIX-compliant and is the default file driver for all systems. | H5Pset\_fapl\_sec2 |
| Direct | H5FD\_DIRECT | This is the H5FD\_SEC2 driver except data is written to or read from the file synchronously without being cached by the system. | H5Pset\_fapl\_direct |
| Log | H5FD\_LOG | This is the H5FD\_SEC2 driver with logging capabilities. | H5Pset\_fapl\_log |
| Windows | H5FD\_WINDOWS | This driver was modified in HDF5-1.8.8 to be a wrapper of the POSIX driver, H5FD\_SEC2. This change should not affect user applications. | H5Pset\_fapl\_windows |
| STDIO | H5FD\_STDIO | This driver uses func­tions from the standard C stdio.h to perform I/O to a single, perma­nent file on local disk with additional system buffering. | H5Pset\_fapl\_stdio |
| Memory | H5FD\_CORE | With this driver, an application can work with a file in memory for faster reads and writes. File contents are kept in memory until the file is closed. At clos­ing, the memory version of the file can be writ­ten back to disk or abandoned. | H5Pset\_fapl\_core |
| Family | H5FD\_FAMILY | With this driver, the HDF5 file’s address space is partitioned into pieces and sent to sepa­rate storage files using an underlying driver of the user’s choice. This driver is for systems that do not support files larger than 2 gigabytes. | H5Pset\_fapl\_family |
| Multi | H5FD\_MULTI | With this driver, data can be stored in multi­ple files according to the type of the data. I/O might work better if data is stored in sepa­rate files based on the type of data. The Split driver is a special case of this driver. | H5Pset\_fapl\_multi |
| Split | H5FD\_SPLIT | This file driver splits a file into two parts. One part stores metadata, and the other part stores raw data. This splitting a file into two parts is a limited case of the Multi driver. | H5Pset\_fapl\_split |
| Parallel | H5FD\_MPIO | This is the standard HDF5 file driver for par­allel file systems. This driver uses the MPI standard for both com­munication and file I/O. | H5Pset\_fapl\_mpio |
| Parallel POSIX | H5FD\_MPIPOSIX | This driver is no longer available. |  |
| Stream | H5FD\_STREAM | This driver is no longer available. |  |

For more information, see the HDF5 Reference Manual entries for the function calls shown in the column on the right in the table above.

Note that the low-level file drivers manage alternative file storage layouts. Dataset storage layouts (chunk­ing, compression, and external dataset storage) are managed independently of file storage layouts.

If an application requires a special-purpose low-level driver, the VFL provides a public API for creating one. For more information on how to create a driver, see “Virtual File Layer” and “List of VFL Functions” in the HDF5 Technical Notes.

### Identifying the Previously-used File Driver

When creating a new HDF5 file, no history exists, so the file driver must be specified if it is to be other than the default.

When opening existing files, however, the application may need to determine which low-level driver was used to create the file. The function H5Pget\_driver is used for this purpose. See the example below.

|  |
| --- |
| hid\_t H5Pget\_driver (hid\_t fapl\_id) |
| Code Example 3-5. Identifying a driver |

H5Pget\_driver returns a constant identifying the low-level driver for the access property list fapl\_id. For example, if the file was created with the POSIX (aka SEC2) driver, H5Pget\_driver returns H5F­D\_SEC2.

If the application opens an HDF5 file without both determining the driver used to create the file and set­ting up the use of that driver, the HDF5 Library will examine the superblock and the driver definition block to identify the driver. See the [HDF5 File Format Specification](https://portal.hdfgroup.org/display/HDF5/File+Format+Specification) for detailed descriptions of the superblock and the driver definition block.

### The POSIX (aka SEC2) Driver

The POSIX driver, H5FD\_SEC2, uses functions from section 2 of the POSIX manual to access unbuffered files stored on a local file system. This driver is also known as the SEC2 driver. The HDF5 Library buffers metadata regardless of the low-level driver, but using this driver prevents data from being buffered again by the lowest layers of the library.

The function H5Pset\_fapl\_sec2 sets the file access properties to use the POSIX driver. See the example below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_sec2 (hid\_t fapl\_id) |
| Code Example 3-6. Using the POSIX, aka SEC2, driver |

Any previously-defined driver properties are erased from the property list.

Additional parameters may be added to this function in the future. Since there are no additional variable settings associated with the POSIX driver, there is no H5Pget\_fapl\_sec2 function.

### The Direct Driver

The Direct driver, H5FD\_DIRECT, functions like the POSIX driver except that data is written to or read from the file synchronously without being cached by the system.

The functions H5Pset\_fapl\_direct and H5Pget\_fapl\_direct are used to manage file access prop­erties. See the example below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_direct( hid\_t fapl\_id, size\_t alignment,  size\_t block\_size, size\_t cbuf\_size )  herr\_t H5Pget\_fapl\_direct( hid\_t fapl\_id, size\_t \*alignment,  size\_t \*block\_size, size\_t \*cbuf\_size ) |
| Code Example 3-7. Using the Direct driver |

H5Pset\_fapl\_direct sets the file access properties to use the Direct driver; any previously defined driver properties are erased from the property list. H5Pget\_fapl\_direct retrieves the file access prop­erties used with the Direct driver. fapl\_id is the file access property list identifier. alignment is the memory alignment boundary. block\_size is the file system block size. cbuf\_size is the copy buffer size.

Additional parameters may be added to this function in the future.

### The Log Driver

The Log driver, H5FD\_LOG, is designed for situations where it is necessary to log file access activity.

The function H5Pset\_fapl\_log is used to manage logging properties. See the example below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_log (hid\_t fapl\_id, const char \*logfile,  unsigned int flags, size\_t buf\_size) |
| Code Example 3-8. Logging file access |

H5Pset\_fapl\_log sets the file access property list to use the Log driver. File access characteristics are identical to access via the POSIX driver. Any previously defined driver properties are erased from the prop­erty list.

Log records are written to the file logfile.

The logging levels set with the verbosity parameter are shown in the table below.

|  |  |
| --- | --- |
| Table 3-3. Logging levels | |
| **Level** | **Comments** |
| 0 | Performs no logging. |
| 1 | Records where writes and reads occur in the file. |
| 2 | Records where writes and reads occur in the file and what kind of data is writ­ten at each location. This includes raw data or any of several types of metadata (object headers, superblock, B-tree data, local headers, or global headers). |

There is no H5Pget\_fapl\_log function.

Additional parameters may be added to this function in the future.

### The Windows Driver

The Windows driver, H5FD\_WINDOWS, was modified in HDF5-1.8.8 to be a wrapper of the POSIX driver, H5FD\_SEC2. In other words, if the Windows drivers is used, any file I/O will instead use the functionality of the POSIX driver. This change should be transparent to all user applications. The Windows driver used to be the default driver for Windows systems. The POSIX driver is now the default.

The function H5Pset\_fapl\_windows sets the file access properties to use the Windows driver. See the example below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_windows (hid\_t fapl\_id) |
| Code Example 3-9. Using the Windows driver |

Any previously-defined driver properties are erased from the property list.

Additional parameters may be added to this function in the future. Since there are no additional variable settings associated with the POSIX driver, there is no H5Pget\_fapl\_windows function.

### The STDIO Driver

The STDIO driver, H5FD\_STDIO, accesses permanent files in a local file system like the POSIX driver does. The STDIO driver also has an additional layer of buffering beneath the HDF5 Library.

The function H5Pset\_fapl\_stdio sets the file access properties to use the STDIO driver. See the exam­ple below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_stdio (hid\_t fapl\_id) |
| Code Example 3-10. Using the STDIO driver |

Any previously defined driver properties are erased from the property list.

Additional parameters may be added to this function in the future. Since there are no additional variable settings associated with the STDIO driver, there is no H5Pget\_fapl\_stdio function.

### The Memory (aka Core) Driver

There are several situations in which it is reasonable, sometimes even required, to maintain a file entirely in system memory. You might want to do so if, for example, either of the following conditions apply:

•        Performance requirements are so stringent that disk latency is a limiting factor

•        You are working with small, temporary files that will not be retained and, thus,

need not be writ­ten to storage media

The Memory driver, H5FD\_CORE, provides a mechanism for creating and managing such in-memory files. The functions H5Pset\_fapl\_core and H5Pget\_fapl\_core manage file access properties. See the example below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_core (hid\_t access\_properties,  size\_t block\_size, hbool\_t backing\_store)  herr\_t H5Pget\_fapl\_core (hid\_t access\_properties,  size\_t \*block\_size), hbool\_t \*backing\_store) |
| Code Example 3-11. Managing file access for in-memory files |

H5Pset\_fapl\_core sets the file access property list to use the Memory driver; any previously defined driver properties are erased from the property list.

Memory for the file will always be allocated in units of the specified block\_size.

The backing\_store Boolean flag is set when the in-memory file is created. backing\_store indicates whether to write the file contents to disk when the file is closed. If backing\_store is set to 1 (TRUE), the file contents are flushed to a file with the same name as the in-memory file when the file is closed or access to the file is terminated in memory. If backing\_store is set to 0 (FALSE), the file is not saved.

The application is allowed to open an existing file with the H5FD\_CORE driver. While using H5Fopen to open an existing file, if backing\_store is set to 1 and the flag for H5Fopen is set to H5F\_ACC\_RDWR, changes to the file contents will be saved to the file when the file is closed. If backing\_store is set to 0 and the flag for H5Fopen is set to H5F\_ACC\_RDWR, changes to the file contents will be lost when the file is closed. If the flag for H5Fopen is set to H5F\_ACC\_RDONLY, no change to the file will be allowed either in memory or on file.

If the file access property list is set to use the Memory driver, H5Pget\_fapl\_core will return block\_­size and backing\_store with the relevant file access property settings.

Note the following important points regarding in-memory files:

•        Local temporary files are created and accessed directly from memory without ever

being written to disk

•        Total file size must not exceed the available virtual memory

•        Only one HDF5 file identifier can be opened for the file, the identifier returned by

H5Fcreate or H5Fopen

•        The changes to the file will be discarded when access is terminated unless

backing\_store is set to 1

Additional parameters may be added to these functions in the future.

See the "HDF5 File Image Operations" section for information on more advanced usage of the Memory file driver, and see the "Modified Region Writes" section for information on how to set write operations so that only modified regions are written to storage.

### The Family Driver

HDF5 files can become quite large, and this can create problems on systems that do not support files larger than 2 gigabytes. The HDF5 file family mechanism is designed to solve the problems this creates by splitting the HDF5 file address space across several smaller files. This structure does not affect how meta­data and raw data are stored: they are mixed in the address space just as they would be in a single, contig­uous file.

HDF5 applications access a family of files via the Family driver, H5FD\_FAMILY. The functions H5Pset\_­fapl\_family and H5Pget\_fapl\_family are used to manage file family properties. See the example below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_family (hid\_t fapl\_id,  hsize\_t memb\_size, hid\_t member\_properties)    herr\_t H5Pget\_fapl\_family (hid\_t fapl\_id,  hsize\_t \*memb\_size, hid\_t \*member\_properties) |
| Code Example 3-12. Managing file family properties |

Each member of the family is the same logical size though the size and disk storage reported by file system listing tools may be substantially smaller. Examples of file system listing tools are ’ls -l’ on a Unix sys­tem or the detailed folder listing on an Apple Macintosh or Microsoft Windows system. The name passed to H5Fcreate or H5Fopen should include a printf(3c)-style integer format specifier which will be replaced with the family member number. The first family member is numbered zero (0).

H5Pset\_fapl\_family sets the access properties to use the Family driver; any previously defined driver properties are erased from the property list. member\_properties will serve as the file access property list for each member of the file family. memb\_size specifies the logical size, in bytes, of each family mem­ber. memb\_size is used only when creating a new file or truncating an existing file; otherwise the mem­ber size is determined by the size of the first member of the family being opened. Note: If the size of the off\_t type is four bytes, the maximum family member size is usually 2^31-1 because the byte at offset 2,147,483,647 is generally inaccessible.

H5Pget\_fapl\_family is used to retrieve file family properties. If the file access property list is set to use the Family driver, member\_properties will be returned with a pointer to a copy of the appropriate member access property list. If memb\_size is non-null, it will contain the logical size, in bytes, of family members.

Additional parameters may be added to these functions in the future.

#### 3.11.8.1. Unix Tools and an HDF5 Utility

It occasionally becomes necessary to repartition a file family. A command-line utility for this purpose, h5repart, is distributed with the HDF5 library.

h5repart [-v] [-b block\_size[suffix]] [-m member\_size[suffix]] source destination

h5repart repartitions an HDF5 file by copying the source file or file family to the destination file or file family, preserving holes in the underlying UNIX files. Families are used for the source and/or destination if the name includes a printf-style integer format such as %d. The -v switch prints input and output file names on the standard error stream for progress monitoring, -b sets the I/O block size (the default is 1KB), and -m sets the output member size if the destination is a family name (the default is 1GB). block\_size and member\_size may be suffixed with the letters g, m, or k for GB, MB, or KB respectively.

The h5repart utility is described on the [Tools](https://portal.hdfgroup.org/display/HDF5/Tools) page of the HDF5 Reference Manual.

An existing HDF5 file can be split into a family of files by running the file through split(1) on a UNIX sys­tem and numbering the output files. However, the HDF5 Library is lazy about extending the size of family members, so a valid file cannot generally be created by concatenation of the family members.

Splitting the file and rejoining the segments by concatenation (split(1) and cat(1) on UNIX systems) does not generate files with holes; holes are preserved only through the use of h5repart.

### The Multi Driver

In some circumstances, it is useful to separate metadata from raw data and some types of metadata from other types of metadata. Situations that would benefit from use of the Multi driver include the following:

•        In networked situations where the small metadata files can be kept on local disks but

larger raw data files must be stored on remote media

•        In cases where the raw data is extremely large

•         In situations requiring frequent access to metadata held in RAM while the raw data

can be effi­ciently held on disk

In either case, access to the metadata is substantially easier with the smaller, and possibly more localized, metadata files. This often results in improved application performance.

The Multi driver, H5FD\_MULTI, provides a mechanism for segregating raw data and different types of metadata into multiple files. The functions H5Pset\_fapl\_multi and H5Pget\_fapl\_multi are used to manage access properties for these multiple files. See the example below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_multi (hid\_t fapl\_id,  const H5FD\_mem\_t \*memb\_map,  const hid\_t \*memb\_fapl,  const char \* const \*memb\_name,  const haddr\_t \*memb\_addr,  hbool\_t relax)  herr\_t H5Pget\_fapl\_multi (hid\_t fapl\_id,  const H5FD\_mem\_t \*memb\_map,  const hid\_t \*memb\_fapl,  const char \*\*memb\_name,  const haddr\_t \*memb\_addr,  hbool\_t \*relax) |
| Code Example 3-13. Managing access properties for multiple files |

H5Pset\_fapl\_multi sets the file access properties to use the Multi driver; any previously defined driver properties are erased from the property list. With the Multi driver invoked, the application will provide a base name to H5Fopen or H5Fcreate. The files will be named by that base name as modified by the rule indicated in memb\_name. File access will be governed by the file access property list memb\_properties.

See H5Pset\_fapl\_multi and H5Pget\_fapl\_multi in the HDF5 Reference Manual for descriptions of these functions and their usage.

Additional parameters may be added to these functions in the future.

### The Split Driver

The Split driver, H5FD\_SPLIT, is a limited case of the Multi driver where only two files are created. One file holds metadata, and the other file holds raw data.

The function H5Pset\_fapl\_split is used to manage Split file access properties. See the example below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_split (hid\_t access\_properties,  const char \*meta\_extension, hid\_t meta\_properties,  const char \*raw\_extension, hid\_t raw\_properties) |
| Code Example 3-14. Managing access properties for split files |

H5Pset\_fapl\_split sets the file access properties to use the Split driver; any previously defined driver properties are erased from the property list.

With the Split driver invoked, the application will provide a base file name such as file\_name to H5F­create or H5Fopen. The metadata and raw data files in storage will then be named file\_name.meta\_extension and file\_name.raw\_extension, respectively. For example, if meta\_extension is defined as .meta and raw\_extension is defined as .raw, the final filenames will be file\_name.meta and file\_name.raw.

Each file can have its own file access property list. This allows the creative use of other low-level file driv­ers. For instance, the metadata file can be held in RAM and accessed via the Memory driver while the raw data file is stored on disk and accessed via the POSIX driver. Metadata file access will be governed by the file access property list in meta\_properties. Raw data file access will be governed by the file access property list in raw\_properties.

Additional parameters may be added to these functions in the future. Since there are no additional vari­able settings associated with the Split driver, there is no H5Pget\_fapl\_split function.

### The Parallel Driver

Parallel environments require a parallel low-level driver. HDF5’s default driver for parallel systems is called the Parallel driver, H5FD\_MPIO. This driver uses the MPI standard for both communication and file I/O.

The functions H5Pset\_fapl\_mpio and H5Pget\_fapl\_mpio are used to manage file access properties for the H5FD\_MPIO driver. See the example below.

|  |
| --- |
| herr\_t H5Pset\_fapl\_mpio (hid\_t fapl\_id, MPI\_Comm comm, MPI\_info info)  herr\_t H5Pget\_fapl\_mpio (hid\_t fapl\_id, MPI\_Comm \*comm, MPI\_info \*info) |
| Code Example 3-15. Managing parallel file access properties |

The file access properties managed by H5Pset\_fapl\_mpio and retrieved by H5Pget\_fapl\_mpio are the MPI communicator, comm, and the MPI info object, info. comm and info are used for file open. info is an information object much like an HDF5 property list. Both are defined in MPI\_FILE\_OPEN of MPI-2.

The communicator and the info object are saved in the file access property list fapl\_id. fapl\_id can then be passed to MPI\_FILE\_OPEN to create and/or open the file.

H5Pset\_fapl\_mpio and H5Pget\_fapl\_mpio are available only in the parallel HDF5 Library and are not collective functions. The Parallel driver is available only in the parallel HDF5 Library.

Additional parameters may be added to these functions in the future.

## Code Examples for Opening and Closing Files

### Example Using the H5F\_ACC\_TRUNC Flag

The following example uses the H5F\_ACC\_TRUNC flag when it creates a new file. The default file creation and file access properties are also used. Using H5F\_ACC\_TRUNC means the function will look for an exist­ing file with the name specified by the function. In this case, that name is FILE. If the function does not find an existing file, it will create one. If it does find an existing file, it will empty the file in preparation for a new set of data. The identifier for the "new" file will be passed back to the application program. For more information, see "File Access Modes."

|  |
| --- |
| hid\_t file; /\* identifier \*/    /\* Create a new file using H5F\_ACC\_TRUNC access, default  \* file creation properties, and default file access  \*/ properties.  file = H5Fcreate(FILE, H5F\_ACC\_TRUNC, H5P\_DEFAULT,  H5P\_DEFAULT);    /\* Close the file. \*/  status = H5Fclose(file); |
| Code Example 3-16. Creating a file with default creation and access properties |

### Example with the File Creation Property List

The example below shows how to create a file with 64-bit object offsets and lengths.

|  |
| --- |
| hid\_t create\_plist;  hid\_t file\_id;  create\_plist = H5Pcreate(H5P\_FILE\_CREATE);  H5Pset\_sizes(create\_plist, 8, 8);  file\_id = H5Fcreate(“test.h5”, H5F\_ACC\_TRUNC, create\_plist, H5P\_DEFAULT);  .  .  .  H5Fclose(file\_id); |
| Code Example 3-17. Creating a file with 64-bit offsets |

### Example with the File Access Property List

This example shows how to open an existing file for independent datasets access by MPI parallel I/O:

|  |
| --- |
| hid\_t access\_plist;  hid\_t file\_id;  access\_plist = H5Pcreate(H5P\_FILE\_ACCESS);  H5Pset\_fapl\_mpi(access\_plist, MPI\_COMM\_WORLD,  MPI\_INFO\_NULL);    /\* H5Fopen must be called collectively \*/  file\_id = H5Fopen(“test.h5”, H5F\_ACC\_RDWR, access\_plist);  .  .  .  /\* H5Fclose must be called collectively \*/  H5Fclose(file\_id); |
| Code Example 3-18. Opening an existing file for parallel I/O |

## Working with Multiple HDF5 Files

Multiple HDF5 files can be associated so that the files can be worked with as though all the information is in a single HDF5 file. A temporary association can be set up by means of the H5Fmount function. A perma­nent association can be set up by means of the external link function H5Lcreate\_external.

The purpose of this section is to describe what happens when the H5Fmount function is used to mount one file on another.

When a file is mounted on another, the mounted file is mounted at a group, and the root group of the mounted file takes the place of that group until the mounted file is unmounted or until the files are closed.

The figure below shows two files before one is mounted on the other. File1 has two groups and three datasets. The group that is the target of the A link has links, Z and Y, to two of the datasets. The group that is the target of the B link has a link, W, to the other dataset. File2 has three groups and three datasets. The groups in File2 are the targets of the AA, BB, and CC links. The datasets in File2 are the targets of the ZZ, YY, and WW links.

|  |
| --- |
| Files_fig3.JPG |
| Figure 3-3. Two separate files |

The figure below shows the two files after File2 has been mounted File1 at the group that is the target of the B link.

|  |
| --- |
| Files_fig4.JPG |
| Figure 3-4. File2 mounted on File1 |

Note: In the figure above, the dataset that is the target of the W link is not shown. That dataset is masked by the mounted file.

If a file is mounted on a group that has members, those members are hidden until the mounted file is unmounted. There are two ways around this if you need to work with a group member. One is to mount the file on an empty group. Another is to open the group member before you mount the file. Opening the group member will return an identifier that you can use to locate the group member.

The example below shows how H5Fmount might be used to mount File2 onto File1.

|  |
| --- |
| status = H5Fmount(loc\_id, "/B", child\_id, plist\_id) |
| Code Example 3-19. Using H5Fmount |

Note: In the code example above, loc\_id is the file identifier for File1, /B is the link path to the group where File2 is mounted, child\_id is the file identifier for File2, and plist\_id is a property list identifier.

For more information, see "HDF5 Groups." See the entries for H5Fmount, H5Funmount, and H5Lcreate\_external in the HDF5 Reference Manual.

# HDF5 Groups

## Introduction

### The Group Object

Abstractly, an HDF5 group contains zero or more objects and every object must be a member of at least one group. The root group, the sole exception, may not belong to any group.

|  |
| --- |
| groups_fig4.JPG |
| Figure 4-4. Abstract model of the HDF5 group object |

Group membership is actually implemented via link objects. See the figure above. A link object is owned by a group and points to a named object. Each link has a name, and each link points to exactly one object. Each named object has at least one and possibly many links to it.

There are three classes of named objects: group, dataset, and committed datatype (formerly called named datatype). See the figure below. Each of these objects is the member of at least one group, which means there is at least one link to it.

|  |
| --- |
| groups_fig5.JPG |
| Figure 4-5. Classes of named objects |

The primary operations on a group are to add and remove members and to discover member objects. These abstract operations, as listed in the figure below, are implemented in the H5G APIs. For more information, see "Group Function Summaries."

To add and delete members of a group, links from the group to existing objects in the file are created and deleted with the link and unlink operations. When a new named object is created, the HDF5 Library executes the link operation in the background immediately after creating the object (in other words, a new object is added as a member of the group in which it is created without further user intervention).

Given the name of an object, the get\_object\_info method retrieves a description of the object, including the number of references to it. The iterate method iterates through the members of the group, returning the name and type of each object.

|  |
| --- |
| Groups_fig6.JPG |
| Figure 4-6. The group object |

Every HDF5 file has a single root group, with the name /. The root group is identical to any other HDF5 group, except:

•        The root group is automatically created when the HDF5 file is created (H5Fcreate).

•        The root group has no parent, but by convention has a reference count of 1.

•        The root group cannot be deleted (in other words, unlinked)!

### The Hierarchy of Data Objects

An HDF5 file is organized as a rooted, directed graph using HDF5 group objects. The named data objects are the nodes of the graph, and the links are the directed arcs. Each arc of the graph has a name, with the special name / reserved for the root group. New objects are created and then inserted into the graph with a link operation that is automatically executed by the library; existing objects are inserted into the graph with a link operation explicitly called by the user, which creates a named link from a group to the object.

An object can be the target of more than one link.

The names on the links must be unique within each group, but there may be many links with the same name in different groups. These are unambiguous, because some ancestor must have a different name, or else they are the same object. The graph is navigated with path names, analogous to Unix file systems. For more information, see "HDF5 Path Names." An object can be opened with a full path starting at the root group, or with a relative path and a starting point. That starting point is always a group, though it may be the current working group, another specified group, or the root group of the file. Note that all paths are relative to a single HDF5 file. In this sense, an HDF5 file is analogous to a single UNIX file system.[1](https://bitbucket.hdfgroup.org/pages/HDFFV/hdf5doc/master/browse/html/UG/HDF5_Users_Guide-Responsive%20HTML5/HDF5_Users_Guide/Groups/HDF5_Groups.htm#FNH_3)

It is important to note that, just like the UNIX file system, HDF5 objects do not have names, the names are associated with paths. An object has an object identifier that is unique within the file, but a single object may have many names because there may be many paths to the same object. An object can be renamed, or moved to another group, by adding and deleting links. In this case, the object itself never moves. For that matter, membership in a group has no implication for the physical location of the stored object.

Deleting a link to an object does not necessarily delete the object. The object remains available as long as there is at least one link to it. After all links to an object are deleted, it can no longer be opened, and the storage may be reclaimed.

It is also important to realize that the linking mechanism can be used to construct very complex graphs of objects. For example, it is possible for an object to be shared between several groups and even to have more than one name in the same group. It is also possible for a group to be a member of itself, or to create other cycles in the graph, such as in the case where a child group is linked to one of its ancestors.

HDF5 also has soft links similar to UNIX soft links. A soft link is an object that has a name and a path name for the target object. The soft link can be followed to open the target of the link just like a regular or hard link. The differences are that the hard link cannot be created if the target object does not exist and it always points to the same object. A soft link can be created with any path name, whether or not the object exists; it may or may not, therefore, be possible to follow a soft link. Furthermore, a soft link’s target object may be changed.

### HDF5 Path Names

The structure of the HDF5 file constitutes the name space for the objects in the file. A path name is a string of components separated by slashes (/). Each component is the name of a hard or soft link which points to an object in the file. The slash not only separates the components, but indicates their hierarchical relation­ship; the component indicated by the link name following a slash is a always a member of the component indicated by the link name preceding that slash.

The first component in the path name may be any of the following:

•        The special character dot (., a period), indicating the current group

•        The special character slash (/), indicating the root group

•        Any member of the current group

Component link names may be any string of ASCII characters not containing a slash or a dot (/ and ., which are reserved as noted above). However, users are advised to avoid the use of punctuation and non-printing characters, as they may create problems for other software. The figure below provides a BNF grammar for HDF5 path names.

|  |
| --- |
| PathName ::= AbsolutePathName | RelativePathName  Separator ::= "/" ["/"]\*  AbsolutePathName ::= Separator [ RelativePathName ]  RelativePathName ::= Component [ Separator RelativePathName ]\*  Component ::=  "." |  Characters  Characters ::= Character+   -  { "." }  Character ::= {c:  c Î { { legal ASCII characters } - {'/'} } |
| Figure 4-7. A BNF grammar for HDF5 path names |

An object can always be addressed by a either a full or absolute path name, starting at the root group, or by a relative path name, starting in a known location such as the current working group. As noted else­where, a given object may have multiple full and relative path names.

Consider, for example, the file illustrated in the figure below. Dataset1 can be identified by either of these absolute path names:

/GroupA/Dataset1

/GroupA/GroupB/GroupC/Dataset1

Since an HDF5 file is a directed graph structure, and is therefore not limited to a strict tree structure, and since this illustrated file includes the sort of circular reference that a directed graph enables, Dataset1 can also be identified by this absolute path name:

/GroupA/GroupB/GroupC/GroupA/Dataset1

Alternatively, if the current working location is GroupB, Dataset1 can be identified by either of these rel­ative path names:

GroupC/Dataset1

GroupC/GroupA/Dataset1

Note that relative path names in HDF5 do not employ the ../ notation, the UNIX notation indicating a parent directory, to indicate a parent group.

|  |
| --- |
| Group_fig2_800001.jpg |
| Figure 4-8. A file with a circular reference |

### Group Implementations in HDF5

The original HDF5 group implementation provided a single indexed structure for link storage. A new group implementation, as of HDF5 Release 1.8.0, enables more efficient compact storage for very small groups, improved link indexing for large groups, and other advanced features.

•        The original indexed format remains the default. Links are stored in a B-tree in the

group’s local heap.

•        Groups created in the new compact-or-indexed format, the implementation introduced

with Release 1.8.0, can be tuned for performance, switching between the compact and

indexed for­mats at thresholds set in the user application.

* The compact format will conserve file space and processing overhead when

working with small groups and is particularly valuable when a group contains

no links. Links are stored as a list of messages in the group’s header.

* The indexed format will yield improved performance when working with large

groups. A large group may contain thousands to millions of members. Links

are stored in a fractal heap and indexed with an improved B-tree.

•        The new implementation also enables the use of link names consisting of non-ASCII

character sets (see H5Pset\_char\_encoding) and is required for all link types other than

hard or soft links; the link types other than hard or soft links are external links and

user-defined links (see the H5L APIs).

The original group structure and the newer structures are not directly interoperable. By default, a group will be created in the original indexed format. An existing group can be changed to a compact-or-indexed format if the need arises; there is no capability to change back. As stated above, once in the compact-or-indexed format, a group can switch between compact and indexed as needed.

Groups will be initially created in the compact-or-indexed format only when one or more of the following conditions is met:

•        The low version bound value of the library version bounds property has been set to

Release 1.8.0 or later in the file access property list (see H5Pset\_libver\_bounds).

Currently, that would require an H5Pset\_libver\_bounds call with the low parameter

set to H5F\_LIBVER\_LATEST.

When this property is set for an HDF5 file, all objects in the file will be created using

the latest available format; no effort will be made to create a file that can be read by

older libraries.

•        The creation order tracking property, H5P\_CRT\_ORDER\_TRACKED, has been set

in the group cre­ation property list (see H5Pset\_link\_creation\_order).

An existing group, currently in the original indexed format, will be converted to the compact-or-indexed format upon the occurrence of any of the following events:

•        An external or user-defined link is inserted into the group.

•        A link named with a string composed of non-ASCII characters is inserted into the

group.

The compact-or-indexed format offers performance improvements that will be most notable at the extremes (for example, in groups with zero members and in groups with tens of thousands of members). But measurable differences may sometimes appear at a threshold as low as eight group members. Since these performance thresholds and criteria differ from application to application, tunable settings are pro­vided to govern the switch between the compact and indexed formats (see H5Pset\_link\_phase\_change). Optimal thresholds will depend on the application and the operating environment.

Future versions of HDF5 will retain the ability to create, read, write, and manipulate all groups stored in either the original indexed format or the compact-or-indexed format.

## Description of the Group Object

### The Group Object

Abstractly, an HDF5 group contains zero or more objects and every object must be a member of at least one group. The root group, the sole exception, may not belong to any group.

|  |
| --- |
| groups_fig4.JPG |
| Figure 4-4. Abstract model of the HDF5 group object |

Group membership is actually implemented via link objects. See the figure above. A link object is owned by a group and points to a named object. Each link has a name, and each link points to exactly one object. Each named object has at least one and possibly many links to it.

There are three classes of named objects: group, dataset, and committed datatype (formerly called named datatype). See the figure below. Each of these objects is the member of at least one group, which means there is at least one link to it.

|  |
| --- |
| groups_fig5.JPG |
| Figure 4-5. Classes of named objects |

The primary operations on a group are to add and remove members and to discover member objects. These abstract operations, as listed in the figure below, are implemented in the H5G APIs. For more information, see "Group Function Summaries."

To add and delete members of a group, links from the group to existing objects in the file are created and deleted with the link and unlink operations. When a new named object is created, the HDF5 Library executes the link operation in the background immediately after creating the object (in other words, a new object is added as a member of the group in which it is created without further user intervention).

Given the name of an object, the get\_object\_info method retrieves a description of the object, including the number of references to it. The iterate method iterates through the members of the group, returning the name and type of each object.

|  |
| --- |
| Groups_fig6.JPG |
| Figure 4-6. The group object |

Every HDF5 file has a single root group, with the name /. The root group is identical to any other HDF5 group, except:

•        The root group is automatically created when the HDF5 file is created (H5Fcreate).

•        The root group has no parent, but by convention has a reference count of 1.

•        The root group cannot be deleted (in other words, unlinked)!

### The Hierarchy of Data Objects

An HDF5 file is organized as a rooted, directed graph using HDF5 group objects. The named data objects are the nodes of the graph, and the links are the directed arcs. Each arc of the graph has a name, with the special name / reserved for the root group. New objects are created and then inserted into the graph with a link operation that is automatically executed by the library; existing objects are inserted into the graph with a link operation explicitly called by the user, which creates a named link from a group to the object.

An object can be the target of more than one link.

The names on the links must be unique within each group, but there may be many links with the same name in different groups. These are unambiguous, because some ancestor must have a different name, or else they are the same object. The graph is navigated with path names, analogous to Unix file systems. For more information, see "HDF5 Path Names." An object can be opened with a full path starting at the root group, or with a relative path and a starting point. That starting point is always a group, though it may be the current working group, another specified group, or the root group of the file. Note that all paths are relative to a single HDF5 file. In this sense, an HDF5 file is analogous to a single UNIX file system.[1](https://bitbucket.hdfgroup.org/pages/HDFFV/hdf5doc/master/browse/html/UG/HDF5_Users_Guide-Responsive%20HTML5/HDF5_Users_Guide/Groups/HDF5_Groups.htm" \l "FNH_3)

It is important to note that, just like the UNIX file system, HDF5 objects do not have names, the names are associated with paths. An object has an object identifier that is unique within the file, but a single object may have many names because there may be many paths to the same object. An object can be renamed, or moved to another group, by adding and deleting links. In this case, the object itself never moves. For that matter, membership in a group has no implication for the physical location of the stored object.

Deleting a link to an object does not necessarily delete the object. The object remains available as long as there is at least one link to it. After all links to an object are deleted, it can no longer be opened, and the storage may be reclaimed.

It is also important to realize that the linking mechanism can be used to construct very complex graphs of objects. For example, it is possible for an object to be shared between several groups and even to have more than one name in the same group. It is also possible for a group to be a member of itself, or to create other cycles in the graph, such as in the case where a child group is linked to one of its ancestors.

HDF5 also has soft links similar to UNIX soft links. A soft link is an object that has a name and a path name for the target object. The soft link can be followed to open the target of the link just like a regular or hard link. The differences are that the hard link cannot be created if the target object does not exist and it always points to the same object. A soft link can be created with any path name, whether or not the object exists; it may or may not, therefore, be possible to follow a soft link. Furthermore, a soft link’s target object may be changed.

### HDF5 Path Names

The structure of the HDF5 file constitutes the name space for the objects in the file. A path name is a string of components separated by slashes (/). Each component is the name of a hard or soft link which points to an object in the file. The slash not only separates the components, but indicates their hierarchical relation­ship; the component indicated by the link name following a slash is a always a member of the component indicated by the link name preceding that slash.

The first component in the path name may be any of the following:

•        The special character dot (., a period), indicating the current group

•        The special character slash (/), indicating the root group

•        Any member of the current group

Component link names may be any string of ASCII characters not containing a slash or a dot (/ and ., which are reserved as noted above). However, users are advised to avoid the use of punctuation and non-printing characters, as they may create problems for other software. The figure below provides a BNF grammar for HDF5 path names.

|  |
| --- |
| PathName ::= AbsolutePathName | RelativePathName  Separator ::= "/" ["/"]\*  AbsolutePathName ::= Separator [ RelativePathName ]  RelativePathName ::= Component [ Separator RelativePathName ]\*  Component ::=  "." |  Characters  Characters ::= Character+   -  { "." }  Character ::= {c:  c Î { { legal ASCII characters } - {'/'} } |
| Figure 4-7. A BNF grammar for HDF5 path names |

An object can always be addressed by a either a full or absolute path name, starting at the root group, or by a relative path name, starting in a known location such as the current working group. As noted else­where, a given object may have multiple full and relative path names.

Consider, for example, the file illustrated in the figure below. Dataset1 can be identified by either of these absolute path names:

/GroupA/Dataset1

/GroupA/GroupB/GroupC/Dataset1

Since an HDF5 file is a directed graph structure, and is therefore not limited to a strict tree structure, and since this illustrated file includes the sort of circular reference that a directed graph enables, Dataset1 can also be identified by this absolute path name:

/GroupA/GroupB/GroupC/GroupA/Dataset1

Alternatively, if the current working location is GroupB, Dataset1 can be identified by either of these rel­ative path names:

GroupC/Dataset1

GroupC/GroupA/Dataset1

Note that relative path names in HDF5 do not employ the ../ notation, the UNIX notation indicating a parent directory, to indicate a parent group.

|  |
| --- |
| Group_fig2_800001.jpg |
| Figure 4-8. A file with a circular reference |

### Group Implementations in HDF5

The original HDF5 group implementation provided a single indexed structure for link storage. A new group implementation, as of HDF5 Release 1.8.0, enables more efficient compact storage for very small groups, improved link indexing for large groups, and other advanced features.

•        The original indexed format remains the default. Links are stored in a B-tree in the

group’s local heap.

•        Groups created in the new compact-or-indexed format, the implementation introduced

with Release 1.8.0, can be tuned for performance, switching between the compact and

indexed for­mats at thresholds set in the user application.

* The compact format will conserve file space and processing overhead when

working with small groups and is particularly valuable when a group contains

no links. Links are stored as a list of messages in the group’s header.

* The indexed format will yield improved performance when working with large

groups. A large group may contain thousands to millions of members. Links

are stored in a fractal heap and indexed with an improved B-tree.

•        The new implementation also enables the use of link names consisting of non-ASCII

character sets (see H5Pset\_char\_encoding) and is required for all link types other than

hard or soft links; the link types other than hard or soft links are external links and

user-defined links (see the H5L APIs).

The original group structure and the newer structures are not directly interoperable. By default, a group will be created in the original indexed format. An existing group can be changed to a compact-or-indexed format if the need arises; there is no capability to change back. As stated above, once in the compact-or-indexed format, a group can switch between compact and indexed as needed.

Groups will be initially created in the compact-or-indexed format only when one or more of the following conditions is met:

•        The low version bound value of the library version bounds property has been set to

Release 1.8.0 or later in the file access property list (see H5Pset\_libver\_bounds).

Currently, that would require an H5Pset\_libver\_bounds call with the low parameter

set to H5F\_LIBVER\_LATEST.

When this property is set for an HDF5 file, all objects in the file will be created using

the latest available format; no effort will be made to create a file that can be read by

older libraries.

•        The creation order tracking property, H5P\_CRT\_ORDER\_TRACKED, has been set

in the group cre­ation property list (see H5Pset\_link\_creation\_order).

An existing group, currently in the original indexed format, will be converted to the compact-or-indexed format upon the occurrence of any of the following events:

•        An external or user-defined link is inserted into the group.

•        A link named with a string composed of non-ASCII characters is inserted into the

group.

The compact-or-indexed format offers performance improvements that will be most notable at the extremes (for example, in groups with zero members and in groups with tens of thousands of members). But measurable differences may sometimes appear at a threshold as low as eight group members. Since these performance thresholds and criteria differ from application to application, tunable settings are pro­vided to govern the switch between the compact and indexed formats (see H5Pset\_link\_phase\_change). Optimal thresholds will depend on the application and the operating environment.

Future versions of HDF5 will retain the ability to create, read, write, and manipulate all groups stored in either the original indexed format or the compact-or-indexed format.

## Using h5dump

You can use h5dump, the command-line utility distributed with HDF5, to examine a file for purposes either of determining where to create an object within an HDF5 file or to verify that you have created an object in the intended place.

In the case of the new group created later in this chapter, the following h5dump command will display the contents of FileA.h5:

h5dump FileA.h5

For more information, see "Creating a Group."

Assuming that the discussed objects, GroupA and GroupB are the only objects that exist in FileA.h5, the output will look something like the following:

HDF5 "FileA.h5" {

GROUP "/" {

GROUP GroupA {

GROUP GroupB {

}

}

}

}

h5dump is described on the “HDF5 Tools” page of the HDF5 Reference Manual.

The HDF5 DDL grammar is described in the document DDL in BNF for HDF5.

## Group Function Summaries

Functions that can be used with groups (H5G functions) and property list functions that can used with groups (H5P functions) are listed below. A number of group functions have been deprecated. Most of these have become link (H5L) or object (H5O) functions. These replacement functions are also listed below.

|  |  |
| --- | --- |
| Function Listing 4-1. Group functions (H5G) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Gcreate  h5gcreate\_f | Creates a new empty group and gives it a name. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Gcreate\_anon  h5gcreate\_anon\_f | Creates a new empty group without linking it into the file structure. |
| H5Gopen  h5gopen\_f | Opens an existing group for modification and returns a group identifier for that group. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Gclose  h5gclose\_f | Closes the specified group. |
| H5Gget\_create\_plist  h5gget\_create\_plist\_f | Gets a group creation property list identifier. |
| H5Gget\_info  h5gget\_info\_f | Retrieves information about a group. Use instead of H5Gget\_num\_objs. |
| H5Gget\_info\_by\_idx  h5gget\_info\_by\_idx\_f | Retrieves information about a group accord­ing to the group’s position within an index. |
| H5Gget\_info\_by\_name  h5gget\_info\_by\_name\_f | Retrieves information about a group. |
| (no C function)  h5gget\_obj\_info\_idx\_f | Returns name and type of the group member identified by its index. Use with the h5gn\_members\_f function. h5gget\_ob­j\_info\_idx\_f and h5gn\_members\_f are the Fortran equivalent of the C function H5Literate. |
| (no C function)  h5gn\_members\_f | Returns the number of group members. Use with the h5gget\_obj\_info\_idx\_f func­tion. |

|  |  |
| --- | --- |
| Function Listing 4-2. Link (H5L) and object (H5O) functions | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Lcreate\_hard  h5lcreate\_hard\_f | Creates a hard link to an object. Replaces H5Glink and H5Glink2. |
| H5Lcreate\_soft  h5lcreate\_soft\_f | Creates a soft link to an object. Replaces H5Glink and H5Glink2. |
| H5Lcreate\_external  h5lcreate\_external\_f | Creates a soft link to an object in a different file. Replaces H5Glink and H5Glink2. |
| H5Lcreate\_ud  (no Fortran subroutine) | Creates a link of a user-defined type. |
| H5Lget\_val  (no Fortran subroutine) | Returns the value of a symbolic link. Replaces H5Gget\_linkval. |
| H5Literate  h5literate\_f | Iterates through links in a group. Replaces H5Giterate. See also H5Ovisit and H5Lvisit. |
| H5Literate\_by\_name  h5literate\_by\_name\_f | Iterates through links in a group. |
| H5Lvisit  (no Fortran subroutine) | Recursively visits all links starting from a spec­ified group. |
| H5Ovisit  h5ovisit\_f | Recursively visits all objects accessible from a specified object. |
| H5Lget\_info  h5lget\_info\_f | Returns information about a link. Replaces H5Gget\_objinfo. |
| H5Oget\_info  (no Fortran subroutine) | Retrieves the metadata for an object specified by an identifier. Replaces H5Gget\_objinfo. |
| H5Lget\_name\_by\_idx  h5lget\_name\_by\_idx\_f | Retrieves name of the nth link in a group, according to the order within a specified field or index. Replaces H5Gget\_ob­jname\_by\_idx. |
| H5Oget\_info\_by\_idx  (no Fortran subroutine) | Retrieves the metadata for an object, identi­fying the object by an index position. Replaces H5Gget\_objtype\_by\_idx. |
| H5Oget\_info\_by\_name  h5oget\_info\_by\_name\_f | Retrieves the metadata for an object, identi­fying the object by location and relative name. |
| H5Oset\_comment  (no Fortran subroutine) | Sets the comment for specified object. Replaces H5Gset\_comment. |
| H5Oget\_comment  (no Fortran subroutine) | Gets the comment for specified object. Replaces H5Gget\_comment. |
| H5Ldelete  h5ldelete\_f | Removes a link from a group. Replaces H5Gunlink. |
| H5Lmove  h5lmove\_f | Renames a link within an HDF5 file. Replaces H5Gmove and H5Gmove2. |

|  |  |
| --- | --- |
| Function Listing 4-3. Group creation property list functions (H5P) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Pall\_filters\_avail  (no Fortran subroutine) | Verifies that all required filters are available. |
| H5Pget\_filter  h5pget\_filter\_f | Returns information about a filter in a pipe­line. The C function is a macro: see “API Com­patibility Macros in HDF5.” |
| H5Pget\_filter\_by\_id  h5pget\_filter\_by\_id\_f | Returns information about the specified filter. The C function is a macro: see “API Compati­bility Macros in HDF5.” |
| H5Pget\_nfilters  h5pget\_nfilters\_f | Returns the number of filters in the pipeline. |
| H5Pmodify\_filter  h5pmodify\_filter\_f | Modifies a filter in the filter pipeline. |
| H5Premove\_filter  h5premove\_filter\_f | Deletes one or more filters in the filter pipe­line. |
| H5Pset\_deflate  h5pset\_deflate\_f | Sets the deflate (GNU gzip) compression method and compression level. |
| H5Pset\_filter  h5pset\_filter\_f | Adds a filter to the filter pipeline. |
| H5Pset\_fletcher32  h5pset\_fletcher32\_f | Sets up use of the Fletcher32 checksum filter. |
| H5Pset\_fletcher32  h5pset\_fletcher32\_f | Sets up use of the Fletcher32 checksum filter. |
| H5Pset\_link\_phase\_change  h5pset\_link\_phase\_change\_f | Sets the parameters for conversion between compact and dense groups. |
| H5Pget\_link\_phase\_change  h5pget\_link\_phase\_change\_f | Queries the settings for conversion between compact and dense groups. |
| H5Pset\_est\_link\_info  h5pset\_est\_link\_info\_f | Sets estimated number of links and length of link names in a group. |
| H5Pget\_est\_link\_info  h5pget\_est\_link\_info\_f | Queries data required to estimate required local heap or object header size. |
| H5Pset\_nlinks  h5pset\_nlinks\_f | Sets maximum number of soft or user-defined link traversals. |
| H5Pget\_nlinks  h5pget\_nlinks\_f | Retrieves the maximum number of link tra­versals. |
| H5Pset\_link\_creation\_order  h5pset\_link\_creation\_order\_f | Sets creation order tracking and indexing for links in a group. |
| H5Pget\_link\_creation\_order  h5pget\_link\_creation\_order\_f | Queries whether link creation order is tracked and/or indexed in a group. |
| H5Pset\_create\_intermediate\_group  h5pset\_create\_inter\_group\_f | Specifies in the property list whether to cre­ate missing intermediate groups. |
| H5Pget\_create\_intermediate\_group  (no Fortran subroutine) | Determines whether the property is set to enable creating missing intermediate groups. |
| H5Pset\_char\_encoding  h5pset\_char\_encoding\_f | Sets the character encoding used to encode a string. Use to set ASCII or UTF-8 character encoding for object names. |
| H5Pget\_char\_encoding  h5pget\_char\_encoding\_f | Retrieves the character encoding used to cre­ate a string. |

|  |  |
| --- | --- |
| Function Listing 4-4. Other external link functions | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Pset/get\_elink\_file\_cache\_size  (no Fortran subroutine) | Sets/retrieves the size of the external link open file cache from the specified file access property list. |
| H5Fclear\_elink\_file\_cache  (no Fortran subroutine) | Clears the external link open file cache for a file. |

## Programming Model for Groups

The programming model for working with groups is as follows:

1.      Create a new group or open an existing one.

2.      Perform the desired operations on the group.

•        Create new objects in the group.

•        Insert existing objects as group members.

•        Delete existing members.

•        Open and close member objects.

•        Access information regarding member objects.

•        Iterate across group members.

•        Manipulate links.

3.      Terminate access to the group (Close the group).

### Creating a Group

To create a group, use H5Gcreate, specifying the location and the path of the new group. The location is the identifier of the file or the group in a file with respect to which the new group is to be identified. The path is a string that provides either an absolute path or a relative path to the new group. For more information, see "HDF5 Path Names." A path that begins with a slash (/) is an absolute path indi­cating that it locates the new group from the root group of the HDF5 file. A path that begins with any other character is a relative path. When the location is a file, a relative path is a path from that file’s root group; when the location is a group, a relative path is a path from that group.

The sample code in the example below creates three groups. The group Data is created in the root direc­tory; two groups are then created in /Data, one with absolute path, the other with a relative path.

|  |
| --- |
| hid\_t file;  file = H5Fopen(....);    group = H5Gcreate(file, "/Data", H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT);  group\_new1 = H5Gcreate(file, "/Data/Data\_new1", H5P\_DEFAULT, H5P\_DEFAULT,  H5P\_DEFAULT);  group\_new2 = H5Gcreate(group, "Data\_new2", H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT); |
| Code Example 4-1. Creating three new groups |

The third H5Gcreate parameter optionally specifies how much file space to reserve to store the names that will appear in this group. If a non-positive value is supplied, a default size is chosen.

### Opening a Group and Accessing an Object in that Group

Though it is not always necessary, it is often useful to explicitly open a group when working with objects in that group. Using the file created in the example above, the example below illustrates the use of a previ­ously-acquired file identifier and a path relative to that file to open the group Data.

Any object in a group can be also accessed by its absolute or relative path. To open an object using a rela­tive path, an application must first open the group or file on which that relative path is based. To open an object using an absolute path, the application can use any location identifier in the same file as the target object; the file identifier is commonly used, but object identifier for any object in that file will work. Both of these approaches are illustrated in the example below.

Using the file created in the examples above, the example below provides sample code illustrating the use of both relative and absolute paths to access an HDF5 data object. The first sequence (two function calls) uses a previously-acquired file identifier to open the group Data, and then uses the returned group iden­tifier and a relative path to open the dataset CData. The second approach (one function call) uses the same previously-acquired file identifier and an absolute path to open the same dataset.

|  |
| --- |
| group = H5Gopen(file, "Data", H5P\_DEFAULT);  dataset1 = H5Dopen(group, "CData", H5P\_DEFAULT);    dataset2 = H5Dopen(file, "/Data/CData", H5P\_DEFAULT); |
| Code Example 4-2. Open a dataset with relative and absolute paths |

### Creating a Dataset in a Specific Group

Any dataset must be created in a particular group. As with groups, a dataset may be created in a particular group by specifying its absolute path or a relative path. The example below illustrates both approaches to creating a dataset in the group /Data.

|  |
| --- |
| dataspace = H5Screate\_simple(RANK, dims, NULL);  dataset1 = H5Dcreate(file, "/Data/CData", H5T\_NATIVE\_INT,  dataspace, H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT);    group = H5Gopen(file, "Data", H5P\_DEFAULT);  dataset2 = H5Dcreate(group, "Cdata2", H5T\_NATIVE\_INT,  dataspace, H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT); |
| Code Example 4-3. Create a dataset with absolute and relative paths |

### Closing a Group

To ensure the integrity of HDF5 objects and to release system resources, an application should always call the appropriate close function when it is through working with an HDF5 object. In the case of groups, H5Gclose ends access to the group and releases any resources the HDF5 library has maintained in sup­port of that access, including the group identifier.

As illustrated in the example below, all that is required for an H5Gclose call is the group identifier acquired when the group was opened; there are no relative versus absolute path considerations.

|  |
| --- |
| herr\_t status;  status = H5Gclose(group); |
| Code Example 4-4. Close a group |

A non-negative return value indicates that the group was successfully closed and the resources released; a negative return value indicates that the attempt to close the group or release resources failed.

### Creating Links

As previously mentioned, every object is created in a specific group. Once created, an object can be made a member of additional groups by means of links created with one of the H5Lcreate\_\* functions.

A link is, in effect, a path by which the target object can be accessed; it therefore has a name which func­tions as a single path component. A link can be removed with an H5Ldelete call, effectively removing the target object from the group that contained the link (assuming, of course, that the removed link was the only link to the target object in the group).

#### Hard Links

There are two kinds of links, hard links and symbolic links. Hard links are reference counted; symbolic links are not. When an object is created, a hard link is automatically created. An object can be deleted from the file by removing all the hard links to it.

Working with the file from the previous examples, the code in the example below illustrates the creation of a hard link, named Data\_link, in the root group, /, to the group Data. Once that link is created, the dataset Cdata can be accessed via either of two absolute paths, /Data/Cdata or /Data\_Link/Cdata.

|  |
| --- |
| status = H5Lcreate\_hard(Data\_loc\_id, "Data", DataLink\_loc\_id,  "Data\_link", H5P\_DEFAULT, H5P\_DEFAULT)    dataset1 = H5Dopen(file, "/Data\_link/CData", H5P\_DEFAULT);  dataset2 = H5Dopen(file, "/Data/CData", H5P\_DEFAULT); |
| Code Example 4-5. Create a hard link |

The example below shows example code to delete a link, deleting the hard link Data from the root group. The group /Data and its members are still in the file, but they can no longer be accessed via a path using the component /Data.

|  |
| --- |
| status = H5Ldelete(Data\_loc\_id, "Data", H5P\_DEFAULT);    dataset1 = H5Dopen(file, "/Data\_link/CData", H5P\_DEFAULT);  /\* This call should succeed; all path components  \* still exist  \*/  dataset2 = H5Dopen(file, "/Data/CData", H5P\_DEFAULT);  /\* This call will fail; the path component '/Data'  \* has been deleted.  \*/ |
| Code Example 4-6. Delete a link |

When the last hard link to an object is deleted, the object is no longer accessible. H5Ldelete will not pre­vent you from deleting the last link to an object. To see if an object has only one link, use the H5Oget\_info function. If the value of the rc (reference count) field in the is greater than 1, then the link can be deleted without making the object inaccessible.

The example below shows H5Oget\_info to the group originally called Data.

|  |
| --- |
| status = H5Oget\_info(Data\_loc\_id, object\_info); |
| Code Example 4-7. Finding the number of links to an object |

It is possible to delete the last hard link to an object and not make the object inaccessible. Suppose your application opens a dataset, and then deletes the last hard link to the dataset. While the dataset is open, your application still has a connection to the dataset. If your application creates a hard link to the dataset before it closes the dataset, then the dataset will still be accessible.

#### Symbolic Links

Symbolic links are objects that assign a name in a group to a path. Notably, the target object is determined only when the symbolic link is accessed, and may, in fact, not exist. Symbolic links are not reference counted, so there may be zero, one, or more symbolic links to an object.

The major types of symbolic links are soft links and external links. Soft links are symbolic links within an HDF5 file and are created with the H5Lcreate\_soft function. Symbolic links to objects located in exter­nal files, in other words external links, can be created with the H5Lcreate\_external function. Symbolic links are removed with the H5Ldelete function.

The example below shows the creating two soft links to the group /Data.

|  |
| --- |
| status = H5Lcreate\_soft(path\_to\_target, link\_loc\_id, "Soft2",  H5P\_DEFAULT, H5P\_DEFAULT);  status = H5Lcreate\_soft(path\_to\_target, link\_loc\_id, "Soft3",  H5P\_DEFAULT, H5P\_DEFAULT);    dataset = H5Dopen(file, "/Soft2/CData", H5P\_DEFAULT); |
| Code Example 4-8. Create a soft link |

With the soft links defined in the example above, the dataset CData in the group /Data can now be opened with any of the names /Data/CData, /Soft2/CData, or /Soft3/CData.

In release 1.8.7, a cache was added to hold the names of files accessed via external links. The size of this cache can be changed to help improve performance. For more information, see the entry in the HDF5 Ref­erence Manual for the H5Pset\_elink\_file\_cache\_size function call.

#### Note Regarding Hard Links and Soft Links

Note that an object’s existence in a file is governed by the presence of at least one hard link to that object. If the last hard link to an object is removed, the object is removed from the file and any remaining soft link becomes a dangling link, a link whose target object does not exist.

#### Moving or Renaming Objects, and a Warning

An object can be renamed by changing the name of a link to it with H5Lmove. This has the same effect as creating a new link with the new name and deleting the link with the old name.

Exercise caution in the use of H5Lmove and H5Ldelete as these functions each include a step that unlinks a pointer to an HDF5 object. If the link that is removed is on the only path leading to an HDF5 object, that object will become permanently inaccessible in the file.

##### Scenario 1: Removing the Last Link

To avoid removing the last link to an object or otherwise making an object inaccessible, use the H5Oget\_info function. Make sure that the value of the reference count field (rc) is greater than 1.

##### Scenario 2: Moving a Link that Isolates an Object

Consider the following example: assume that the group group2 can only be accessed via the following path, where top\_group is a member of the file’s root group:

/top\_group/group1/group2/

Using H5Lmove, top\_group is renamed to be a member ofgroup2. At this point, since top\_group was the only route from the root group to group1, there is no longer a path by which one can access group1, group2, or any member datasets. And since top\_group is now a member of group2, top\_group itself and any member datasets have thereby also become inaccessible.

#### Mounting a File

An external link is a permanent connection between two files. A temporary connection can be set up with the H5Fmount function. [For more information, see "The HDF5 File."](https://bitbucket.hdfgroup.org/pages/HDFFV/hdf5doc/master/browse/html/UG/HDF5_Users_Guide-Responsive%20HTML5/HDF5_Users_Guide/TheFile/The_HDF5_File.htm#XREF_TheFileChapter)For more information, see the[H5Fmount](http://www.hdfgroup.org/HDF5/doc/RM/RM_H5F.html#File-Mount) function in the [HDF5 Reference Manual](http://www.hdfgroup.org/HDF5/doc/RM/RM_H5Front.html).

### Discovering Information about Objects

There is often a need to retrieve information about a particular object. The H5Lget\_info and H5Oget\_info functions fill this niche by returning a description of the object or link in an H5L\_info\_t or H5O\_info\_t structure.

### Discovering Objects in a Group

To examine all the objects or links in a group, use the H5Literate or H5Ovisit functions to examine the objects, and use the H5Lvisit function to examine the links. H5Literate is useful both with a single group and in an iterative process that examines an entire file or section of a file (such as the contents of a group or the contents of all the groups that are members of that group) and acts on objects as they are encountered. H5Ovisit recursively visits all objects accessible from a specified object. H5Lvisit recur­sively visits all the links starting from a specified group.

### Discovering All of the Objects in the File

The structure of an HDF5 file is self-describing, meaning that an application can navigate an HDF5 file to discover and understand all the objects it contains. This is an iterative process wherein the structure is tra­versed as a graph, starting at one node and recursively visiting linked nodes. To explore the entire file, the traversal should start at the root group.

## Examples of File Structures

This section presents several samples of HDF5 file structures.

|  |  |
| --- | --- |
| groups_fig27_a.JPG | groups_fig27_b.JPG |
| a) The file contains three groups: the root group, /group1, and /group2. | b) The dataset dset1 (or /group1/dset1) is created in /group1. |
|  |  |
| groups_fig27_aa.JPG | groups_fig27_bb.JPG |
| c) A link named dset2 to the same dataset is created in /group2. | d) The link from /group1 to dset1 is removed. The dataset is still in the file, but can be accessed only as /group2/dset2. |
| Figure 4-9. Some file structures | |

The figure above shows examples of the structure of a file with three groups and one dataset. The file in part a contains three groups: the root group and two member groups. In part b, the dataset dset1 has been created in /group1. In part c, a link named dset2 from /group2 to the dataset has been added. Note that there is only one copy of the dataset; there are two links to it and it can be accessed either as /group1/dset1 or as /group2/dset2.

The figure in part d above illustrates that one of the two links to the dataset can be deleted. In this case, the link from /group1 has been removed. The dataset itself has not been deleted; it is still in the file but can only be accessed as /group1/dset2.

|  |  |
| --- | --- |
| groups_fig28_a.JPG | groups_fig28_b.JPG |
| a) dset1 has two names: /group2/dset1 and /group1/GXX/dset1. | b) dset1 again has two names: /group1/dset1 and /group1/dset2. |
|  |  |
| groups_fig28_c.JPG | groups_fig28_d.JPG |
| c) dset1 has three names: /group1/dset1, /group2/dset2, and /group1/GXX/dset2. | d) dset1 has an infinite number of available path names. |
| Figure 4-10. More sample file structures | |

The figure above illustrates loops in an HDF5 file structure. The file in part a contains three groups and a dataset; group2 is a member of the root group and of the root group’s other member group, group1. group2 thus can be accessed by either of two paths: /group2 or /group1/GXX. Similarly, the dataset can be accessed either as /group2/dset1 or as /group1/GXX/dset1.

Part b illustrates a different case: the dataset is a member of a single group but with two links, or names, in that group. In this case, the dataset again has two names, /group1/dset1 and /group1/dset2.

In part c, the dataset dset1 is a member of two groups, one of which can be accessed by either of two names. The dataset thus has three path names: /group1/dset1, /group2/dset2, and /group1/GXX/dset2.

And in part d, two of the groups are members of each other and the dataset is a member of both groups. In this case, there are an infinite number of paths to the dataset because GXX and GYY can be traversed any number of times on the way from the root group, /, to the dataset. This can yield a path name such as /group1/GXX/GYY/GXX/GYY/GXX/dset2.

|  |  |
| --- | --- |
| groups_fig29_a.JPG | groups_fig29_b.JPG |
| a) The file contains only hard links. | b) A soft link is added from group2 to  /group1/dset1. |
|  |  |
| groups_fig29_c.JPG | groups_fig29_d.JPG |
| c) A soft link named dset3 is added with a tar­get that does not yet exist. | d) The target of the soft link is created or linked. |
| Figure 4-11. Hard and soft links | |

The figure above takes us into the realm of soft links. The original file, in part a, contains only three hard links. In part b, a soft link named dset2 from group2 to /group1/dset1 has been created, making this dataset accessible as /group2/dset2.

In part c, another soft link has been created in group2. But this time the soft link, dset3, points to a tar­get object that does not yet exist. That target object, dset, has been added in part d and is now accessible as either /group2/dset or /group2/dset3.

1.   It could be said that HDF5 extends the organizing concepts of a file system to the internal structure of a single file.

# HDF5 Datasets

## Introduction

An HDF5 dataset is an object composed of a collection of data elements, or raw data, and metadata that stores a description of the data elements, data layout, and all other information necessary to write, read, and interpret the stored data. From the viewpoint of the application the raw data is stored as a one-dimensional or multi-dimensional array of elements (the raw data), those elements can be any of several numerical or character types, small arrays, or even compound types similar to C structs. The dataset object may have attribute objects. See the figure below.

|  |
| --- |
| Dsets_fig1.JPG |
| Figure 5-1. Application view of a dataset |

A dataset object is stored in a file in two parts: a header and a data array. The header contains information that is needed to interpret the array portion of the dataset, as well as metadata (or pointers to metadata) that describes or annotates the dataset. Header information includes the name of the object, its dimen­sionality, its number-type, information about how the data itself is stored on disk (the storage layout), and other information used by the library to speed up access to the dataset or maintain the file’s integrity.

The HDF5 dataset interface, comprising the H5D functions, provides a mechanism for managing HDF5 datasets including the transfer of data between memory and disk and the description of dataset proper­ties.

A dataset is used by other HDF5 APIs, either by name or by an identifier. For more information, see “Using Identifiers.”

### Link/Unlink

A dataset can be added to a group with one of the H5Lcreate calls, and deleted from a group with H5Ldelete. The link and unlink operations use the name of an object, which may be a dataset. The data­set does not have to open to be linked or unlinked.

### Object Reference

A dataset may be the target of an object reference. The object reference is created by H5Rcreate with the name of an object which may be a dataset and the reference type H5R\_OBJECT. The dataset does not have to be open to create a reference to it.

An object reference may also refer to a region (selection) of a dataset. The reference is created with H5Rcreate and a reference type of H5R\_DATASET\_REGION.

An object reference can be accessed by a call to H5Rdereference. When the reference is to a dataset or dataset region, the H5Rdeference call returns an identifier to the dataset just as if H5Dopen has been called.

### Adding Attributes

A dataset may have user-defined attributes which are created with H5Acreate and accessed through the H5A API. To create an attribute for a dataset, the dataset must be open, and the identifier is passed to H5Acreate. The attributes of a dataset are discovered and opened using H5Aopen\_name, H5Aop­en\_idx, or H5Aiterate; these functions use the identifier of the dataset. An attribute can be deleted with H5Adelete which also uses the identifier of the dataset.

## Dataset Function Summaries

Functions that can be used with datasets (H5D functions) and property list functions that can used with datasets (H5P functions) are listed below.

|  |  |
| --- | --- |
| Function Listing 5-1. Dataset functions (H5D) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Dcreate  h5dcreate\_f | Creates a dataset at the specified location. The C function is a macro: see “[API Compati­bility Macros in HDF5](http://www.hdfgroup.org/HDF5/doc/RM/APICompatMacros.html).” |
| H5Dcreate\_anon  h5dcreate\_anon\_f | Creates a dataset in a file without linking it into the file structure. |
| H5Dopen  h5dopen\_f | Opens an existing dataset. The C function is a macro: see “[API Compatibility Macros in HDF5](http://www.hdfgroup.org/HDF5/doc/RM/APICompatMacros.html).” |
| H5Dclose  h5dclose\_f | Closes the specified dataset. |
| H5Dget\_space  h5dget\_space\_f | Returns an identifier for a copy of the dataspace for a dataset. |
| H5Dget\_space\_status  h5dget\_space\_status\_f | Determines whether space has been allo­cated for a dataset. |
| H5Dget\_type  h5dget\_type\_f | Returns an identifier for a copy of the data­type for a dataset. |
| H5Dget\_create\_plist  h5dget\_create\_plist\_f | Returns an identifier for a copy of the dataset creation property list for a dataset. |
| H5Dget\_access\_plist  (no Fortran subroutine) | Returns the dataset access property list asso­ciated with a dataset. |
| H5Dget\_offset  h5dget\_offset\_f | Returns the dataset address in a file. |
| H5Dget\_storage\_size  h5dget\_storage\_size\_f | Returns the amount of storage required for a dataset. |
| H5Dvlen\_get\_buf\_size  h5dvlen\_get\_max\_len\_f | Determines the number of bytes required to store variable-length (VL) data. |
| H5Dvlen\_reclaim  h5dvlen\_reclaim\_f | Reclaims VL datatype memory buffers. |
| H5Dread  h5dread\_f | Reads raw data from a dataset into a buffer. |
| H5Dwrite  h5dwrite\_f | Writes raw data from a buffer to a dataset. |
| H5Diterate  (no Fortran subroutine) | Iterates over all selected elements in a dataspace. |
| H5Dgather  (no Fortran subroutine) | Gathers data from a selection within a mem­ory buffer. |
| H5Dscatter  (no Fortran subroutine) | Scatters data into a selection within a mem­ory buffer. |
| H5Dfill  h5dfill\_f | Fills dataspace elements with a fill value in a memory buffer. |
| H5Dset\_extent  h5dset\_extent\_f | Changes the sizes of a dataset’s dimensions. |

|  |  |
| --- | --- |
| Function Listing 5-2. Dataset creation property list functions (H5P) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Pset\_layout  h5pset\_layout\_f | Sets the type of storage used to store the raw data for a dataset. |
| H5Pget\_layout  h5pget\_layout\_f | Returns the layout of the raw data for a data­set. |
| H5Pset\_chunk  h5pset\_chunk\_f | Sets the size of the chunks used to store a chunked layout dataset. |
| H5Pget\_chunk  h5pget\_chunk\_f | Retrieves the size of chunks for the raw data of a chunked layout dataset. |
| H5Pset\_deflate  h5pset\_deflate\_f | Sets compression method and compression level. |
| H5Pset\_fill\_value  h5pset\_fill\_value\_f | Sets the fill value for a dataset. |
| H5Pget\_fill\_value  h5pget\_fill\_value\_f | Retrieves a dataset fill value. |
| H5Pfill\_value\_defined  (no Fortran subroutine) | Determines whether the fill value is defined. |
| H5Pset\_fill\_time  h5pset\_fill\_time\_f | Sets the time when fill values are written to a dataset. |
| H5Pget\_fill\_time  h5pget\_fill\_time\_f | Retrieves the time when fill value are written to a dataset. |
| H5Pset\_alloc\_time  h5pset\_alloc\_time\_f | Sets the timing for storage space allocation. |
| H5Pget\_alloc\_time  h5pget\_alloc\_time\_f | Retrieves the timing for storage space alloca­tion. |
| H5Pset\_filter  h5pset\_filter\_f | Adds a filter to the filter pipeline. |
| H5Pall\_filters\_avail  (no Fortran subroutine) | Verifies that all required filters are available. |
| H5Pget\_nfilters  h5pget\_nfilters\_f | Returns the number of filters in the pipeline. |
| H5Pget\_filter  h5pget\_filter\_f | Returns information about a filter in a pipe­line. The C function is a macro: see “API Com­patibility Macros in HDF5.” |
| H5Pget\_filter\_by\_id  h5pget\_filter\_by\_id\_f | Returns information about the specified filter. The C function is a macro: see “API Compati­bility Macros in HDF5.” |
| H5Pmodify\_filter  h5pmodify\_filter\_f | Modifies a filter in the filter pipeline. |
| H5Premove\_filter  h5premove\_filter\_f | Deletes one or more filters in the filter pipe­line. |
| H5Pset\_fletcher32  h5pset\_fletcher32\_f | Sets up use of the Fletcher32 checksum filter. |
| H5Pset\_nbit  h5pset\_nbit\_f | Sets up use of the n-bit filter. |
| H5Pset\_scaleoffset  h5pset\_scaleoffset\_f | Sets up use of the scale-offset filter. |
| H5Pset\_shuffle  h5pset\_shuffle\_f | Sets up use of the shuffle filter. |
| H5Pset\_szip  h5pset\_szip\_f | Sets up use of the Szip compression filter. |
| H5Pset\_external  h5pset\_external\_f | Adds an external file to the list of external files. |
| H5Pget\_external\_count  h5pget\_external\_count\_f | Returns the number of external files for a dataset. |
| H5Pget\_external  h5pget\_external\_f | Returns information about an external file. |
| H5Pset\_char\_encoding  h5pset\_char\_encoding\_f | Sets the character encoding used to encode a string. Use to set ASCII or UTF-8 character encoding for object names. |
| H5Pget\_char\_encoding  h5pget\_char\_encoding\_f | Retrieves the character encoding used to cre­ate a string. |

|  |  |
| --- | --- |
| Function Listing 5-3. Dataset access property list functions (H5P) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Pset\_buffer  h5pset\_buffer\_f | Sets type conversion and background buffers. |
| H5Pget\_buffer  h5pget\_buffer\_f | Reads buffer settings. |
| H5Pset\_chunk\_cache  h5pset\_chunk\_cache\_f | Sets the raw data chunk cache parameters. |
| H5Pget\_chunk\_cache  h5pget\_chunk\_cache\_f | Retrieves the raw data chunk cache parame­ters. |
| H5Pset\_edc\_check  h5pset\_edc\_check\_f | Sets whether to enable error-detection when reading a dataset. |
| H5Pget\_edc\_check  h5pget\_edc\_check\_f | Determines whether error-detection is enabled for dataset reads. |
| H5Pset\_filter\_callback  (no Fortran subroutine) | Sets user-defined filter callback function. |
| H5Pset\_data\_transform  h5pset\_data\_transform\_f | Sets a data transform expression. |
| H5Pget\_data\_transform  h5pget\_data\_transform\_f | Retrieves a data transform expression. |
| H5Pset\_type\_conv\_cb  (no Fortran subroutine) | Sets user-defined datatype conversion call­back function. |
| H5Pget\_type\_conv\_cb  (no Fortran subroutine) | Gets user-defined datatype conversion call­back function. |
| H5Pset\_hyper\_vector\_size  h5pset\_hyper\_vector\_size\_f | Sets number of I/O vectors to be read/written in hyperslab I/O. |
| H5Pget\_hyper\_vector\_size  h5pget\_hyper\_vector\_size\_f | Retrieves number of I/O vectors to be read/written in hyperslab I/O. |
| H5Pset\_btree\_ratios  h5pset\_btree\_ratios\_f | Sets B-tree split ratios for a dataset transfer property list. |
| H5Pget\_btree\_ratios  h5pget\_btree\_ratios\_f | Gets B-tree split ratios for a dataset transfer property list. |
| H5Pset\_vlen\_mem\_manager  (no Fortran subroutine) | Sets the memory manager for variable-length datatype allocation in H5Dread and H5Dv­len\_reclaim. |
| H5Pget\_vlen\_mem\_manager  (no Fortran subroutine) | Gets the memory manager for variable-length datatype allocation in H5Dread and H5Dv­len\_reclaim. |
| H5Pset\_dxpl\_mpio  h5pset\_dxpl\_mpio\_f | Sets data transfer mode. |
| H5Pget\_dxpl\_mpio  h5pget\_dxpl\_mpio\_f | Returns the data transfer mode. |
| H5Pset\_dxpl\_mpio\_chunk\_opt  (no Fortran subroutine) | Sets a flag specifying linked-chunk I/O or multi-chunk I/O. |
| H5Pset\_dxpl\_mpio\_chunk\_opt\_num  (no Fortran subroutine) | Sets a numeric threshold for linked-chunk I/O. |
| H5Pset\_dxpl\_mpio\_chunk\_opt\_ratio  (no Fortran subroutine) | Sets a ratio threshold for collective I/O. |
| H5Pset\_dxpl\_mpio\_collective\_opt  (no Fortran subroutine) | Sets a flag governing the use of independent versus collective I/O. |
| H5Pset\_multi\_type  (no Fortran subroutine) | Sets the type of data property for the MULTI driver. |
| H5Pget\_multi\_type  (no Fortran subroutine) | Retrieves the type of data property for the MULTI driver. |
| H5Pset\_small\_data\_block\_size  h5pset\_small\_data\_block\_size\_f | Sets the size of a contiguous block reserved for small data. |
| H5Pget\_small\_data\_block\_size  h5pget\_small\_data\_block\_size\_f | Retrieves the current small data block size setting. |

## Programming Model for Datasets

This section explains the programming model for datasets.

### General Model

The programming model for using a dataset has three main phases:

•        Obtain access to the dataset

•        Operate on the dataset using the dataset identifier returned at access

•        Release the dataset

These three phases or steps are described in more detail below the figure.

A dataset may be opened several times and operations performed with several different identifiers to the same dataset. All the operations affect the dataset although the calling program must synchronize if nec­essary to serialize accesses.

Note that the dataset remains open until every identifier is closed. The figure below shows the basic sequence of operations.

|  |
| --- |
| Dsets_fig2.JPG |
| Figure 5-2. Dataset programming sequence |

Creation and data access operations may have optional parameters which are set with property lists. The general programming model is:

•        Create property list of appropriate class (dataset create, dataset transfer)

•        Set properties as needed; each type of property has its own format and datatype

•        Pass the property list as a parameter of the API call

The steps below describe the programming phases or steps for using a dataset.

#### Step 1. Obtain Access

A new dataset is created by a call to H5Dcreate. If successful, the call returns an identifier for the newly created dataset.

Access to an existing dataset is obtained by a call to H5Dopen. This call returns an identifier for the existing dataset.

An object reference may be dereferenced to obtain an identifier to the dataset it points to.

In each of these cases, the successful call returns an identifier to the dataset. The identifier is used in sub­sequent operations until the dataset is closed.

#### Step 2. Operate on the Dataset

The dataset identifier can be used to write and read data to the dataset, to query and set properties, and to perform other operations such as adding attributes, linking in groups, and creating references.

The dataset identifier can be used for any number of operations until the dataset is closed.

#### Step 3. Close the Dataset

When all operations are completed, the dataset identifier should be closed. This releases the dataset.

After the identifier is closed, it cannot be used for further operations.

### Create Dataset

A dataset is created and initialized with a call to H5Dcreate. The dataset create operation sets permanent properties of the dataset:

•        Name

•        Dataspace

•        Datatype

•        Storage properties

These properties cannot be changed for the life of the dataset, although the dataspace may be expanded up to its maximum dimensions.

**Name**

A dataset name is a sequence of alphanumeric ASCII characters. The full name would include a tracing of the group hierarchy from the root group of the file. An example is /rootGroup/groupA/subgroup23/dataset1. The local name or relative name within the lowest-level group containing the dataset would include none of the group hierarchy. An example is Dataset1.

**Dataspace**

The dataspace of a dataset defines the number of dimensions and the size of each dimension. The dataspace defines the number of dimensions, and the maximum dimension sizes and current size of each dimension. The maximum dimension size can be a fixed value or the constant H5D\_UNLIMITED, in which case the actual dimension size can be changed with calls to H5Dset\_extent, up to the maximum set with the maxdims parameter in the H5Screate\_simple call that established the dataset’s original dimen­sions. The maximum dimension size is set when the dataset is created and cannot be changed.

**Datatype**

Raw data has a datatype which describes the layout of the raw data stored in the file. The datatype is set when the dataset is created and can never be changed. When data is transferred to and from the dataset, the HDF5 library will assure that the data is transformed to and from the stored format.

**Storage Properties**

Storage properties of the dataset are set when it is created. The required inputs table below shows the categories of storage properties. The storage properties cannot be changed after the dataset is created.

**Filters**

When a dataset is created, optional filters are specified. The filters are added to the data transfer pipeline when data is read or written. The standard library includes filters to implement compression, data shuf­fling, and error detection code. Additional user-defined filters may also be used.

The required filters are stored as part of the dataset, and the list may not be changed after the dataset is created. The HDF5 library automatically applies the filters whenever data is transferred.

**Summary**

A newly created dataset has no attributes and no data values. The dimensions, datatype, storage proper­ties, and selected filters are set. The table below lists the required inputs, and the second table below lists the optional inputs.

|  |  |
| --- | --- |
| Table 5-1. Required inputs | |
| **Required Inputs** | **Description** |
| Dataspace | The shape of the array. |
| Datatype | The layout of the stored elements. |
| Name | The name of the dataset in the group. |

|  |  |
| --- | --- |
| Table 5-2. Optional inputs | |
| **Optional Inputs** | **Description** |
| Storage Layout | How the data is organized in the file including chunking. |
| Fill Value | The behavior and value for uninitialized data. |
| External Storage | Option to store the raw data in an external file. |
| Filters | Select optional filters to be applied. One of the filters that might be applied is compression. |

**Example**

To create a new dataset, go through the following general steps:

•        Set dataset characteristics (optional where default settings are acceptable)

•        Datatype

•        Dataspace

•        Dataset creation property list

•        Create the dataset

•        Close the datatype, dataspace, and property list (as necessary)

•        Close the dataset

Example 1 below shows example code to create an empty dataset. The dataspace is 7 x 8, and the data­type is a big-endian integer. The dataset is created with the name “dset1” and is a member of the root group, “/”.

|  |
| --- |
| hid\_t    dataset, datatype, dataspace;    /\*  \* Create dataspace: Describe the size of the array and  \* create the dataspace for fixed-size dataset.  \*/ |
| dimsf[0] = 7;  dimsf[1] = 8;  dataspace = H5Screate\_simple(2, dimsf, NULL);  /\*  \* Define datatype for the data in the file.  \* For this example, store little-endian integer numbers.  \*/ |
| datatype = H5Tcopy(H5T\_NATIVE\_INT);  status = H5Tset\_order(datatype, H5T\_ORDER\_LE);  /\*  \* Create a new dataset within the file using defined  \* dataspace and datatype. No properties are set.  \*/  dataset = H5Dcreate(file, "/dset", datatype, dataspace,  H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT); |
| H5Dclose(dataset);  H5Sclose(dataspace);  H5Tclose(datatype); |
| Code Example 5-1. Create an empty dataset |

Example 2, below, shows example code to create a similar dataset with a fill value of ‘-1’. This code has the same steps as in the example above, but uses a non-default property list. A file creation property list is cre­ated, and then the fill value is set to the desired value. Then the property list is passed to the H5Dcreate call.

|  |
| --- |
| hid\_t    dataset, datatype, dataspace;  hid\_t plist;  /\* property list \*/  int fillval = -1;  dimsf[0] = 7;  dimsf[1] = 8;  dataspace = H5Screate\_simple(2, dimsf, NULL); |
| datatype = H5Tcopy(H5T\_NATIVE\_INT);  status = H5Tset\_order(datatype, H5T\_ORDER\_LE);    /\*  \* Example of Dataset Creation property list: set fill value  \* to '-1'  \*/ |
| plist = H5Pcreate(H5P\_DATASET\_CREATE);  status = H5Pset\_fill\_value(plist, datatype, &fillval);    /\* Same as above, but use the property list \*/  dataset = H5Dcreate(file, "/dset", datatype, dataspace,  H5P\_DEFAULT, plist, H5P\_DEFAULT); |
| H5Dclose(dataset);  H5Sclose(dataspace);  H5Tclose(datatype);  H5Pclose(plist); |
| Code Example 5-2. Create a dataset with fill value set |

 After this code is executed, the dataset has been created and written to the file. The data array is uninitial­ized. Depending on the storage strategy and fill value options that have been selected, some or all of the space may be allocated in the file, and fill values may be written in the file.

### Data Transfer Operations on a Dataset

Data is transferred between memory and the raw data array of the dataset through H5Dwrite and H5Dread operations. A data transfer has the following basic steps:

1.      Allocate and initialize memory space as needed

2.      Define the datatype of the memory elements

3.      Define the elements to be transferred (a selection, or all the elements)

4.      Set data transfer properties (including parameters for filters or file drivers) as needed

5.      Call the H5D API

Note that the location of the data in the file, the datatype of the data in the file, the storage properties, and the filters do not need to be specified because these are stored as a permanent part of the dataset. A selection of elements from the dataspace is specified; the selected elements may be the whole dataspace.

The figure below shows a diagram of a write operation which transfers a data array from memory to a dataset in the file (usually on disk). A read operation has similar parameters with the data flowing the other direction.

|  |
| --- |
| Dsets_fig5.JPG |
| Figure 5-3. A write operation |

**Memory Space**

The calling program must allocate sufficient memory to store the data elements to be transferred. For a write (from memory to the file), the memory must be initialized with the data to be written to the file. For a read, the memory must be large enough to store the elements that will be read. The amount of storage needed can be computed from the memory datatype (which defines the size of each data element) and the number of elements in the selection.

**Memory Datatype**

The memory layout of a single data element is specified by the memory datatype. This specifies the size, alignment, and byte order of the element as well as the datatype class. Note that the memory datatype must be the same datatype class as the file, but may have different byte order and other properties. The HDF5 Library automatically transforms data elements between the source and destination layouts. For more information, see "HDF5 Datatypes."

For a write, the memory datatype defines the layout of the data to be written; an example is IEEE floating-point numbers in native byte order. If the file datatype (defined when the dataset is created) is different but compatible, the HDF5 Library will transform each data element when it is written. For example, if the file byte order is different than the native byte order, the HDF5 library will swap the bytes.

For a read, the memory datatype defines the desired layout of the data to be read. This must be compati­ble with the file datatype, but should generally use native formats such as byte orders. The HDF5 library will transform each data element as it is read.

**Selection**

The data transfer will transfer some or all of the elements of the dataset depending on the dataspace selection. The selection has two dataspace objects: one for the source, and one for the destination. These objects describe which elements of the dataspace to be transferred. Some (partial I/O) or all of the data may be transferred. Partial I/O is defined by defining hyperslabs or lists of elements in a dataspace object.

The dataspace selection for the source defines the indices of the elements to be read or written. The two selections must define the same number of points, but the order and layout may be different. The HDF5 Library automatically selects and distributes the elements according to the selections. It might, for exam­ple, perform a scatter-gather or sub-set of the data.

**Data Transfer Properties**

For some data transfers, additional parameters should be set using the transfer property list. The table below lists the categories of transfer properties. These properties set parameters for the HDF5 Library and may be used to pass parameters for optional filters and file drivers. For example, transfer properties are used to select independent or collective operation when using MPI-I/O.

|  |  |
| --- | --- |
| Table 5-3. Categories of transfer properties | |
| **Properties** | **Description** |
| Library parameters | Internal caches, buffers, B-Trees, etc. |
| Memory management | Variable-length memory management, data overwrite |
| File driver management | Parameters for file drivers |
| Filter management | Parameters for filters |

**Data Transfer Operation (Read or Write)**

The data transfer is done by calling H5Dread or H5Dwrite with the parameters described above. The HDF5 Library constructs the required pipeline, which will scatter-gather, transform datatypes, apply the requested filters, and use the correct file driver.

During the data transfer, the transformations and filters are applied to each element of the data in the required order until all the data is transferred.

**Summary**

To perform a data transfer, it is necessary to allocate and initialize memory, describe the source and desti­nation, set required and optional transfer properties, and call the H5D API.

**Examples**

The basic procedure to write to a dataset is the following:

•        Open the dataset.

•        Set the dataset dataspace for the write (optional if dataspace is H5S\_SELECT\_ALL).

•        Write data.

•        Close the datatype, dataspace, and property list (as necessary).

•        Close the dataset.

Example 3 below shows example code to write a 4 x 6 array of integers. In the example, the data is initial­ized in the memory array dset\_data. The dataset has already been created in the file, so it is opened with H5Dopen.

The data is written with H5Dwrite. The arguments are the dataset identifier, the memory datatype (H5T\_NATIVE\_INT), the memory and file selections (H5S\_ALL in this case: the whole array), and the default (empty) property list. The last argument is the data to be transferred.

|  |
| --- |
| hid\_t       file\_id, dataset\_id;  /\* identifiers \*/  herr\_t      status;  int         i, j, dset\_data[4][6]; |
| /\* Initialize the dataset. \*/  for (i = 0; i < 4; i++)     for (j = 0; j < 6; j++)  dset\_data[i][j] = i \* 6 + j + 1; |
| /\* Open an existing file. \*/  file\_id = H5Fopen("dset.h5", H5F\_ACC\_RDWR, H5P\_DEFAULT);    /\* Open an existing dataset. \*/  dataset\_id = H5Dopen(file\_id, "/dset", H5P\_DEFAULT); |
| /\* Write the entire dataset, using 'dset\_data':  memory type is 'native int'  write the entire dataspace to the entire dataspace,  no transfer properties,  \*/  status = H5Dwrite(dataset\_id, H5T\_NATIVE\_INT, H5S\_ALL,  H5S\_ALL, H5P\_DEFAULT, dset\_data);    status = H5Dclose(dataset\_id); |
| Code Example 5-3. Write an array of integers |

Example 4 below shows a similar write except for setting a non-default value for the transfer buffer. The code is the same as Example 3, but a transfer property list is created, and the desired buffer size is set. The H5Dwrite function has the same arguments, but uses the property list to set the buffer.

|  |
| --- |
| hid\_t       file\_id, dataset\_id;  hid\_t       xferplist;  herr\_t      status;  int         i, j, dset\_data[4][6]; |
| file\_id = H5Fopen("dset.h5", H5F\_ACC\_RDWR, H5P\_DEFAULT);    dataset\_id = H5Dopen(file\_id, "/dset", H5P\_DEFAULT);    /\*  \* Example: set type conversion buffer to 64MB  \*/ |
| xferplist = H5Pcreate(H5P\_DATASET\_XFER);  status = H5Pset\_buffer( xferplist, 64 \* 1024 \*1024, NULL, NULL);    /\* Write the entire dataset, using 'dset\_data':  memory type is 'native int'  write the entire dataspace to the entire dataspace,  set the buffer size with the property list,  \*/ |
| status = H5Dwrite(dataset\_id, H5T\_NATIVE\_INT, H5S\_ALL,  H5S\_ALL, xferplist, dset\_data);    status = H5Dclose(dataset\_id); |
| Code Example 5-4. Write an array using a property list |

 The basic procedure to read from a dataset is the following:

•        Define the memory dataspace of the read (optional if dataspace is

H5S\_SELECT\_ALL).

•        Open the dataset.

•        Get the dataset dataspace (if using H5S\_SELECT\_ALL above).

Else define dataset dataspace of read.

•        Define the memory datatype (optional).

•        Define the memory buffer.

•        Open the dataset.

•        Read data.

•        Close the datatype, dataspace, and property list (as necessary).

•        Close the dataset.

The example below shows code that reads a 4 x 6 array of integers from a dataset called “dset1”. First, the dataset is opened. The H5Dread call has parameters:

•        The dataset identifier (from H5Dopen)

•        The memory datatype (H5T\_NATVE\_INT)

•        The memory and file dataspace (H5S\_ALL, the whole array)

•        A default (empty) property list

•        The memory to be filled

|  |
| --- |
| hid\_t       file\_id, dataset\_id;  herr\_t      status;  int         i, j, dset\_data[4][6];    /\* Open an existing file. \*/  file\_id = H5Fopen("dset.h5", H5F\_ACC\_RDWR, H5P\_DEFAULT); |
| /\* Open an existing dataset. \*/  dataset\_id = H5Dopen(file\_id, "/dset", H5P\_DEFAULT);    /\* read the entire dataset, into 'dset\_data':  memory type is 'native int'  read the entire dataspace to the entire dataspace,  no transfer properties, |
| \*/  status = H5Dread(dataset\_id, H5T\_NATIVE\_INT, H5S\_ALL,  H5S\_ALL, H5P\_DEFAULT, dset\_data);    status = H5Dclose(dataset\_id); |
| Code Example 5-5. Read an array from a dataset |

### Retrieve the Properties of a Dataset

The functions listed below allow the user to retrieve information regarding a dataset including the data­type, the dataspace, the dataset creation property list, and the total stored size of the data.

|  |  |
| --- | --- |
| Function Listing 5-4. Retrieve dataset information | |
| **Query Function** | **Description** |
| H5Dget\_space | Retrieve the dataspace of the dataset as stored in the file. |
| H5Dget\_type | Retrieve the datatype of the dataset as stored in the file. |
| H5Dget\_create\_plist | Retrieve the dataset creation properties. |
| H5Dget\_storage\_size | Retrieve the total bytes for all the data of the dataset. |
| H5Dvlen\_get\_buf\_size | Retrieve the total bytes for all the variable-length data of the dataset. |

The example below illustrates how to retrieve dataset information.

|  |
| --- |
| hid\_t       file\_id, dataset\_id;  hid\_t       dspace\_id, dtype\_id, plist\_id;  herr\_t      status; |
| /\* Open an existing file. \*/  file\_id = H5Fopen("dset.h5", H5F\_ACC\_RDWR, H5P\_DEFAULT);    /\* Open an existing dataset. \*/  dataset\_id = H5Dopen(file\_id, "/dset", H5P\_DEFAULT); |
| dspace\_id = H5Dget\_space(dataset\_id);  dtype\_id = H5Dget\_type(dataset\_id);  plist\_id = H5Dget\_create\_plist(dataset\_id);    /\* use the objects to discover the properties of the dataset \*/    status = H5Dclose(dataset\_id); |
| Code Example 5-6. Retrieve dataset |

## Data Transfer

The HDF5 library implements data transfers through a pipeline which implements data transformations (according to the datatype and selections), chunking (as requested), and I/O operations using different mechanisms (file drivers). The pipeline is automatically configured by the HDF5 library. Metadata is stored in the file so that the correct pipeline can be constructed to retrieve the data. In addition, optional filters such as compression may be added to the standard pipeline.

The figure below illustrates data layouts for different layers of an application using HDF5. The application data is organized as a multidimensional array of elements. The HDF5 format specification defines the stored layout of the data and metadata. The storage layout properties define the organization of the abstract data. This data is written and read to and from some storage medium.

|  |
| --- |
| Dsets_fig9.JPG |
| Figure 5-4. Data layouts in an application |

The last stage of a write (and first stage of a read) is managed by an HDF5 file driver module. The virtual file layer of the HDF5 Library implements a standard interface to alternative I/O methods, including mem­ory (AKA “core”) files, single serial file I/O, multiple file I/O, and parallel I/O. The file driver maps a simple abstract HDF5 file to the specific access methods.

The raw data of an HDF5 dataset is conceived to be a multidimensional array of data elements. This array may be stored in the file according to several storage strategies:

•        Contiguous

•        Chunked

•        Compact

The storage strategy does not affect data access methods except that certain operations may be more or less efficient depending on the storage strategy and the access patterns.

Overall, the data transfer operations (H5Dread and H5Dwrite) work identically for any storage method, for any file driver, and for any filters and transformations. The HDF5 library automatically manages the data transfer process. In some cases, transfer properties should or must be used to pass additional param­eters such as MPI/IO directives when using the parallel file driver.

### The Data Pipeline

When data is written or read to or from an HDF5 file, the HDF5 library passes the data through a sequence of processing steps which are known as the HDF5 data pipeline. This data pipeline performs operations on the data in memory such as byte swapping, alignment, scatter-gather, and hyperslab selections. The HDF5 library automatically determines which operations are needed and manages the organization of memory operations such as extracting selected elements from a data block. The data pipeline modules operate on data buffers: each module processes a buffer and passes the transformed buffer to the next stage.

The table below lists the stages of the data pipeline. The figure below the table shows the order of pro­cessing during a read or write.

|  |  |
| --- | --- |
| Table 5-4. Stages of the data pipeline | |
| **Layers** | **Description** |
| I/O initiation | Initiation of HDF5 I/O activities (H5Dwrite and H5Dread) in a user’s application program. |
| Memory hyperslab opera­tion | Data is scattered to (for read), or gathered from (for write) the application’s memory buffer (bypassed if no datatype conversion is needed). |
| Datatype conversion | Datatype is converted if it is different between memory and stor­age (bypassed if no datatype conversion is needed). |
| File hyperslab operation | Data is gathered from (for read), or scattered to (for write) to file space in memory (bypassed if no datatype conversion is needed). |
| Filter pipeline | Data is processed by filters when it passes. Data can be modified and restored here (bypassed if no datatype conversion is needed, no filter is enabled, or dataset is not chunked). |
| Virtual File Layer | Facilitate easy plug-in file drivers such as MPIO or POSIX I/O. |
| Actual I/O | Actual file driver used by the library such as MPIO or STDIO. |

|  |
| --- |
| Dsets_fig10.JPG |
| Figure 5-5. The processing order in the data pipeline |

The HDF5 library automatically applies the stages as needed.

When the memory dataspace selection is other than the whole dataspace, the memory hyperslab stage scatters/gathers the data elements between the application memory (described by the selection) and a contiguous memory buffer for the pipeline. On a write, this is a gather operation; on a read, this is a scat­ter operation.

When the memory datatype is different from the file datatype, the datatype conversion stage transforms each data element. For example, if data is written from 32-bit big-endian memory, and the file datatype is 32-bit little-endian, the datatype conversion stage will swap the bytes of every element. Similarly, when data is read from the file to native memory, byte swapping will be applied automatically when needed.

The file hyperslab stage is similar to the memory hyperslab stage, but is managing the arrangement of the elements according to the dataspace selection. When data is read, data elements are gathered from the data blocks from the file to fill the contiguous buffers which are then processed by the pipeline. When data is read, the elements from a buffer are scattered to the data blocks of the file.

### Data Pipeline Filters

In addition to the standard pipeline, optional stages, called filters, can be inserted in the pipeline. The standard distribution includes optional filters to implement compression and error checking. User applica­tions may add custom filters as well.

The HDF5 library distribution includes or employs several optional filters. These are listed in the table below. The filters are applied in the pipeline between the virtual file layer and the file hyperslab operation. See the figure above. The application can use any number of filters in any order.

|  |  |
| --- | --- |
| Table 5-5. Data pipeline filters | |
| **Filter** | **Description** |
| gzip compression | Data compression using zlib. |
| Szip compression | Data compression using the Szip library. See The HDF Group website for more information regarding the Szip filter. |
| N-bit compression | Data compression using an algorithm specialized for n-bit datatypes. |
| Scale-offset compression | Data compression using a “scale and offset” algorithm. |
| Shuffling | To improve compression performance, data is regrouped by its byte position in the data unit. In other words, the 1st, 2nd, 3rd, and 4th bytes of integers are stored together respectively. |
| Fletcher32 | Fletcher32 checksum for error-detection. |

Filters may be used only for chunked data and are applied to chunks of data between the file hyperslab stage and the virtual file layer. At this stage in the pipeline, the data is organized as fixed-size blocks of ele­ments, and the filter stage processes each chunk separately.

Filters are selected by dataset creation properties, and some behavior may be controlled by data transfer properties. The library determines what filters must be applied and applies them in the order in which they were set by the application. That is, if an application calls H5Pset\_shuffle and then H5Pset\_de­flate when creating a dataset’s creation property list, the library will apply the shuffle filter first and then the deflate filter.

For more information, see "Using the N-bit Filter." For more information, see "Using the Scale-offset Filter."

### File Drivers

I/O is performed by the HDF5 virtual file layer. The file driver interface writes and reads blocks of data; each driver module implements the interface using different I/O mechanisms. The table below lists the file drivers currently supported. Note that the I/O mechanisms are separated from the pipeline processing: the pipeline and filter operations are identical no matter what data access mechanism is used.

|  |  |
| --- | --- |
| Table 5-6. I/O file drivers | |
| **File Driver** | **Description** |
| H5FD\_CORE | Store in memory (optional backing store to disk file). |
| H5FD\_FAMILY | Store in a set of files. |
| H5FD\_LOG | Store in logging file. |
| H5FD\_MPIO | Store using MPI/IO. |
| H5FD\_MULTI | Store in multiple files. There are several options to control layout. |
| H5FD\_SEC2 | Serial I/O to file using Unix “section 2” functions. |
| H5FD\_STDIO | Serial I/O to file using Unix “stdio” functions. |

Each file driver writes/reads contiguous blocks of bytes from a logically contiguous address space. The file driver is responsible for managing the details of the different physical storage methods.

In serial environments, everything above the virtual file layer tends to work identically no matter what storage method is used.

Some options may have substantially different performance depending on the file driver that is used. In particular, multi-file and parallel I/O may perform considerably differently from serial drivers depending on chunking and other settings.

### Data Transfer Properties to Manage the Pipeline

Data transfer properties set optional parameters that control parts of the data pipeline. The function list­ing below shows transfer properties that control the behavior of the library.

|  |  |
| --- | --- |
| Function Listing 5-5. Data transfer property list functions | |
| **C Function** | **Purpose** |
| H5Pset\_buffer | Maximum size for the type conversion buffer and the back­ground buffer. May also supply pointers to application-allo­cated buffers. |
| H5Pset\_hyper\_cache | Whether to cache hyperslab blocks during I/O. |
| H5Pset\_btree\_ratios | Set the B-tree split ratios for a dataset transfer property list. The split ratios determine what percent of children go in the first node when a node splits. |

Some filters and file drivers require or use additional parameters from the application program. These can be passed in the data transfer property list. The table below shows file driver property list functions.

|  |  |
| --- | --- |
| Function Listing 5-6. File driver property list functions | |
| **C Function** | **Purpose** |
| H5Pset\_dxpl\_mpio | Control the MPI I/O transfer mode (independent or collective) during data I/O operations. |
| H5Pset\_small\_data\_block\_size | Reserves blocks of size bytes for the contiguous storage of the raw data portion of small datasets. The HDF5 Library then writes the raw data from small datasets to this reserved space which reduces unnecessary discontinuities within blocks of metadata and improves I/O performance. |
| H5Pset\_edc\_check | Disable/enable EDC checking for read. When selected, EDC is always written. |

The transfer properties are set in a property list which is passed as a parameter of the H5Dread or H5Dwrite call. The transfer properties are passed to each pipeline stage. Each stage may use or ignore any property in the list. In short, there is one property list that contains all the properties.

### Storage Strategies

The raw data is conceptually a multi-dimensional array of elements that is stored as a contiguous array of bytes. The data may be physically stored in the file in several ways. The table below lists the storage strat­egies for a dataset.

|  |  |
| --- | --- |
| Table 5-7. Dataset storage strategies | |
| **Storage Strategy** | **Description** |
| Contiguous | The dataset is stored as one continuous array of bytes. |
| Chunked | The dataset is stored as fixed-size chunks. |
| Compact | A small dataset is stored in the metadata header. |

The different storage strategies do not affect the data transfer operations of the dataset: reads and writes work the same for any storage strategy.

These strategies are described in the following sections.

**Contiguous**

A contiguous dataset is stored in the file as a header and a single continuous array of bytes. See the figure below. In the case of a multi-dimensional array, the data is serialized in row major order. By default, data is stored contiguously.

|  |
| --- |
| Dsets_fig12.JPG |
| Figure 5-6. Contiguous data storage |

 Contiguous storage is the simplest model. It has several limitations. First, the dataset must be a fixed-size: it is not possible to extend the limit of the dataset or to have unlimited dimensions. In other words, if the number of dimensions of the array might change over time, then chunking storage must be used instead of contiguous. Second, because data is passed through the pipeline as fixed-size blocks, compression and other filters cannot be used with contiguous data.

**Chunked**

The data of a dataset may be stored as fixed-size chunks. See the figure below. A chunk is a hyper-rectan­gle of any shape. When a dataset is chunked, each chunk is read or written as a single I/O operation, and individually passed from stage to stage of the data pipeline.

|  |
| --- |
| Dsets_fig13.JPG |
| Figure 5-7. Chunked data storage |

Chunks may be any size and shape that fits in the dataspace of the dataset. For example, a three dimen­sional dataspace can be chunked as 3-D cubes, 2-D planes, or 1-D lines. The chunks may extend beyond the size of the dataspace. For example, a 3 x 3 dataset might by chunked in 2 x 2 chunks. Sufficient chunks will be allocated to store the array, and any extra space will not be accessible. So, to store the 3 x 3 array, four 2 x 2 chunks would be allocated with 5 unused elements stored.

Chunked datasets can be unlimited in any direction and can be compressed or filtered.

Since the data is read or written by chunks, chunking can have a dramatic effect on performance by opti­mizing what is read and written. Note, too, that for specific access patterns such as parallel I/O, decompo­sition into chunks can have a large impact on performance.

Two restrictions have been placed on chunk shape and size:

•        The rank of a chunk must be less than or equal to the rank of the dataset

•        Chunk size cannot exceed the size of a fixed-size dataset; for example, a dataset consisting of a 5 x 4 fixed-size array cannot be defined with 10 x 10 chunks

**Compact**

For contiguous and chunked storage, the dataset header information and data are stored in two (or more) blocks. Therefore, at least two I/O operations are required to access the data: one to access the header, and one (or more) to access data. For a small dataset, this is considerable overhead.

A small dataset may be stored in a continuous array of bytes in the header block using the compact stor­age option. This dataset can be read entirely in one operation which retrieves the header and data. The dataset must fit in the header. This may vary depending on the metadata that is stored. In general, a com­pact dataset should be approximately 30 KB or less total size. See the figure below.

|  |
| --- |
| Dsets_fig14.JPG |
| Figure 5-8. Compact data storage |

### Partial I/O Sub-setting and Hyperslabs

Data transfers can write or read some of the data elements of the dataset. This is controlled by specifying two selections: one for the source and one for the destination. Selections are specified by creating a dataspace with selections.

Selections may be a union of hyperslabs or a list of points. A hyperslab is a contiguous hyper-rectangle from the dataspace. Selected fields of a compound datatype may be read or written. In this case, the selection is controlled by the memory and file datatypes.

Summary of procedure:

1.      Open the dataset

2.      Define the memory datatype

3.      Define the memory dataspace selection and file dataspace selection

4.      Transfer data (H5Dread or H5Dwrite)

For more information, see "HDF5 Dataspaces and Partial I/O."

## Allocation of Space in the File

When a dataset is created, space is allocated in the file for its header and initial data. The amount of space allocated when the dataset is created depends on the storage properties. When the dataset is modified (data is written, attributes added, or other changes), additional storage may be allocated if necessary.

|  |  |
| --- | --- |
| Table 5-8. Initial dataset size | |
| **Object** | **Size** |
| Header | Variable, but typically around 256 bytes at the creation of a simple dataset with a simple datatype. |
| Data | Size of the data array (number of elements x size of element). Space allocated in the file depends on the storage strategy and the allocation strategy. |

**Header**

A dataset header consists of one or more header messages containing persistent metadata describing var­ious aspects of the dataset. These records are defined in the HDF5 File Format Specification. The amount of storage required for the metadata depends on the metadata to be stored. The table below summarizes the metadata.

|  |  |
| --- | --- |
| Table 5-9. Metadata storage sizes | |
| **Header Information** | **Approximate Storage Size** |
| Datatype (required) | Bytes or more. Depends on type. |
| Dataspace (required) | Bytes or more. Depends on number of dimensions and hsize\_t. |
| Layout (required) | Points to the stored data. Bytes or more. Depends on hsize\_t and number of dimensions. |
| Filters | Depends on the number of filters. The size of the filter message depends on the name and data that will be passed. |

The header blocks also store the name and values of attributes, so the total storage depends on the num­ber and size of the attributes.

In addition, the dataset must have at least one link, including a name, which is stored in the file and in the group it is linked from.

The different storage strategies determine when and how much space is allocated for the data array. See the discussion of fill values below for a detailed explanation of the storage allocation.

**Contiguous Storage**

For a continuous storage option, the data is stored in a single, contiguous block in the file. The data is nominally a fixed-size, (number of elements x size of element). The figure below shows an example of a two dimensional array stored as a contiguous dataset.

Depending on the fill value properties, the space may be allocated when the dataset is created or when first written (default), and filled with fill values if specified. For parallel I/O, by default the space is allo­cated when the dataset is created.

|  |
| --- |
| Dsets_fig15.JPG |
| Figure 5-9. A two dimensional array stored as a contiguous dataset |

**Chunked Storage**

For chunked storage, the data is stored in one or more chunks. Each chunk is a continuous block in the file, but chunks are not necessarily stored contiguously. Each chunk has the same size. The data array has the same nominal size as a contiguous array (number of elements x size of element), but the storage is allo­cated in chunks, so the total size in the file can be larger than the nominal size of the array. See the figure below.

If a fill value is defined, each chunk will be filled with the fill value. Chunks must be allocated when data is written, but they may be allocated when the file is created, as the file expands, or when data is written.

For serial I/O, by default chunks are allocated incrementally, as data is written to the chunk. For a sparse dataset, chunks are allocated only for the parts of the dataset that are written. In this case, if the dataset is extended, no storage is allocated.

For parallel I/O, by default chunks are allocated when the dataset is created or extended with fill values written to the chunk.

In either case, the default can be changed using fill value properties. For example, using serial I/O, the properties can select to allocate chunks when the dataset is created.

|  |
| --- |
| Dsets_fig16.JPG |
| Figure 5-10. A two dimensional array stored in chunks |

**Changing Dataset Dimensions**

H5Dset\_extent is used to change the current dimensions of the dataset within the limits of the dataspace. Each dimension can be extended up to its maximum or unlimited. Extending the dataspace may or may not allocate space in the file and may or may not write fill values, if they are defined. See the example code below.

The dimensions of the dataset can also be reduced. If the sizes specified are smaller than the dataset’s cur­rent dimension sizes, H5Dset\_extent will reduce the dataset’s dimension sizes to the specified values. It is the user’s responsibility to ensure that valuable data is not lost; H5Dset\_extent does not check.

|  |
| --- |
| hid\_t       file\_id, dataset\_id;  Herr\_t      status;  size\_t      newdims[2];    /\* Open an existing file. \*/  file\_id = H5Fopen("dset.h5", H5F\_ACC\_RDWR, H5P\_DEFAULT); |
| /\* Open an existing dataset. \*/  dataset\_id = H5Dopen(file\_id, "/dset", H5P\_DEFAULT);    /\* Example:  dataset is 2 x 3, each dimension is UNLIMITED \*/  /\* extend to 2 x 7 \*/  newdims[0] = 2;  newdims[1] = 7; |
| status = H5Dset\_extent(dataset\_id, newdims);    /\* dataset is now 2 x 7 \*/    status = H5Dclose(dataset\_id); |
| Code Example 5-7. Using H5Dset\_extent to increase the size of a dataset |

### Storage Allocation in the File: Early, Incremental, Late

The HDF5 Library implements several strategies for when storage is allocated if and when it is filled with fill values for elements not yet written by the user. Different strategies are recommended for different storage layouts and file drivers. In particular, a parallel program needs storage allocated during a collective call (for example, create or extend), while serial programs may benefit from delaying the allocation until the data is written.

Two file creation properties control when to allocate space, when to write the fill value, and the actual fill value to write.

**When to Allocate Space**

The table below shows the options for when data is allocated in the file. Early allocation is done during the dataset create call. Certain file drivers (especially MPI-I/O and MPI-POSIX) require space to be allocated when a dataset is created, so all processors will have the correct view of the data.

|  |  |
| --- | --- |
| Table 5-10. File storage allocation options | |
| **Strategy** | **Description** |
| Early | Allocate storage for the dataset immediately when the dataset is cre­ated. |
| Late | Defer allocating space for storing the dataset until the dataset is written. |
| Incremental | Defer allocating space for storing each chunk until the chunk is written. |
| Default | Use the strategy (Early, Late, or Incremental) for the storage method and access method. This is the recommended strategy. |

Late allocation is done at the time of the first write to dataset. Space for the whole dataset is allocated at the first write.

Incremental allocation (chunks only) is done at the time of the first write to the chunk. Chunks that have never been written are not allocated in the file. In a sparsely populated dataset, this option allocates chunks only where data is actually written.

The “Default” property selects the option recommended as appropriate for the storage method and access method. The defaults are shown in the table below. Note that Early allocation is recommended for all Parallel I/O, while other options are recommended as the default for serial I/O cases.

|  |  |  |
| --- | --- | --- |
| Table 5-11. Default storage options | | |
| **Storage Type** | **Serial I/O** | **Parallel I/O** |
| Contiguous | Late | Early |
| Chunked | Incremental | Early |
| Compact | Early | Early |

**When to Write the Fill Value**

The second property is when to write the fill value. The possible values are “Never” and “Allocation”. The table below shows these options.

|  |  |
| --- | --- |
| Table 5-12. When to write fill values | |
| **When** | **Description** |
| Never | Fill value will never be written. |
| Allocation | Fill value is written when space is allocated. (Default for chunked and contigu­ous data storage.) |

**What Fill Value to Write**

The third property is the fill value to write. The table below shows the values. By default, the data is filled with zeros. The application may choose no fill value (Undefined). In this case, uninitialized data may have random values. The application may define a fill value of an appropriate type. For more information, see "Fill Values."

|  |  |
| --- | --- |
| Table 5-13. Fill values to write | |
| **What to Write** | **Description** |
| Default | By default, the library fills allocated space with zeros. |
| Undefined | Allocated space is filled with random values. |
| User-defined | The application specifies the fill value. |

Together these three properties control the library’s behavior. The table below summarizes the possibili­ties during the dataset create-write-close cycle.

|  |  |  |  |
| --- | --- | --- | --- |
| Table 5-14. Storage allocation and fill summary | | | |
| **When to allocate space** | **When to write fill value** | **What fill value to write** | **Library create-write-close behavior** |
| Early | Never | - | Library allocates space when dataset is cre­ated, but never writes a fill value to dataset. A read of unwritten data returns undefined val­ues. |
| Late | Never | - | Library allocates space when dataset is writ­ten to, but never writes a fill value to the dataset. A read of unwritten data returns undefined values. |
| Incremental | Never | - | Library allocates space when a dataset or chunk (whichever is the smallest unit of space) is written to, but it never writes a fill value to a dataset or a chunk. A read of unwritten data returns undefined values. |
| - | Allocation | Undefined | Error on creating the dataset. The dataset is not created. |
| Early | Allocation | Default or User-defined | Allocate space for the dataset when the data­set is created. Write the fill value (default or user-defined) to the entire dataset when the dataset is created. |
| Late | Allocation | Default or User-defined | Allocate space for the dataset when the appli­cation first writes data values to the dataset. Write the fill value to the entire dataset before writing application data values. |
| Incremental | Allocation | Default or User-defined | Allocate space for the dataset when the appli­cation first writes data values to the dataset or chunk (whichever is the smallest unit of space). Write the fill value to the entire data­set or chunk before writing application data values. |

During the H5Dread function call, the library behavior depends on whether space has been allocated, whether the fill value has been written to storage, how the fill value is defined, and when to write the fill value. The table below summarizes the different behaviors.

|  |  |  |  |
| --- | --- | --- | --- |
| Table 5-15. H5Dread summary | | | |
| **Is space allocated in the file?** | **What is the fill value?** | **When to write the fill value?** | **Library read behavior** |
| No | Undefined | <<any>> | Error. Cannot create this dataset. |
| No | Default or User-defined | <<any>> | Fill the memory buffer with the fill value. |
| Yes | Undefined | <<any>> | Return data from storage (dataset). Trash is possible if the application has not written data to the portion of the dataset being read. |
| Yes | Default or User-defined | Never | Return data from storage (dataset). Trash is possible if the application has not written data to the portion of the dataset being read. |
| Yes | Default or User-defined | Allocation | Return data from storage (dataset). |

There are two cases to consider depending on whether the space in the file has been allocated before the read or not. When space has not yet been allocated and if a fill value is defined, the memory buffer will be filled with the fill values and returned. In other words, no data has been read from the disk. If space has been allocated, the values are returned from the stored data. The unwritten elements will be filled accord­ing to the fill value.

### Deleting a Dataset from a File and Reclaiming Space

HDF5 does not at this time provide an easy mechanism to remove a dataset from a file or to reclaim the storage space occupied by a deleted object.

Removing a dataset and reclaiming the space it used can be done with the H5Ldelete function and the h5repack utility program. With the H5Ldelete function, links to a dataset can be removed from the file structure. After all the links have been removed, the dataset becomes inaccessible to any application and is effectively removed from the file. The way to recover the space occupied by an unlinked dataset is to write all of the objects of the file into a new file. Any unlinked object is inaccessible to the application and will not be included in the new file. Writing objects to a new file can be done with a custom program or with the h5repack utility program.

For more information, see "HDF5 Groups."

### Releasing Memory Resources

The system resources required for HDF5 objects such as datasets, datatypes, and dataspaces should be released once access to the object is no longer needed. This is accomplished via the appropriate close function. This is not unique to datasets but a general requirement when working with the HDF5 Library; failure to close objects will result in resource leaks.

In the case where a dataset is created or data has been transferred, there are several objects that must be closed. These objects include datasets, datatypes, dataspaces, and property lists.

The application program must free any memory variables and buffers it allocates. When accessing data from the file, the amount of memory required can be determined by calculating the size of the memory datatype and the number of elements in the memory selection.

Variable-length data are organized in two or more areas of memory. For more information, see "Variable-length Datatypes." When writing data, the application creates an array of vl\_info\_t which contains pointers to the elements. The elements might be, for example, strings. In the file, the variable-length data is stored in two parts: a heap with the variable-length values of the data elements and an array of vl\_info\_t elements. When the data is read, the amount of memory required for the heap can be determined with the H5Dget\_vlen\_buf\_size call.

The data transfer property may be used to set a custom memory manager for allocating variable-length data for a H5Dread. This is set with the H5Pset\_vlen\_mem\_manager call.

To free the memory for variable-length data, it is necessary to visit each element, free the variable-length data, and reset the element. The application must free the memory it has allocated. For memory allocated by the HDF5 Library during a read, the H5Dvlen\_reclaim function can be used to perform this opera­tion.

### External Storage Properties

The external storage format allows data to be stored across a set of non-HDF5 files. A set of segments (off­sets and sizes) in one or more files is defined as an external file list, or EFL, and the contiguous logical addresses of the data storage are mapped onto these segments. Currently, only the H5D\_CONTIGUOUS storage format allows external storage. External storage is enabled by a dataset creation property. The table below shows the API.

|  |  |
| --- | --- |
| Table 5-16. External storage API | |
| **Function** | **Description** |
| herr\_t H5Pset\_external (hid\_t plist, const char \*name, off\_t offset, hsize\_t size) | This function adds a new segment to the end of the external file list of the specified dataset creation property list. The segment begins a byte offset of file name and continues for size bytes. The space represented by this segment is adjacent to the space already represented by the external file list. The last segment in a file list may have the size H5F\_UNLIMITED, in which case the external file may be of unlim­ited size and no more files can be added to the external files list. |
| int H5Pget\_external\_count (hid\_t plist) | Calling this function returns the number of segments in an external file list. If the dataset creation property list has no external data, then zero is returned. |
| herr\_t H5Pget\_external (hid\_t plist, int idx, size\_t name\_size, char \*name, off\_t \*offset, hsize\_t \*size) | This is the counterpart for the H5Pset\_ex­ternal() function. Given a dataset creation property list and a zero-based index into that list, the file name, byte offset, and segment size are returned through non-null argu­ments. At most name\_size characters are copied into the name argument which is not null terminated if the file name is longer than the supplied name buffer (this is similar to strncpy()). |

The figure below shows an example of how a contiguous, one-dimensional dataset is partitioned into three parts and each of those parts is stored in a segment of an external file. The top rectangle represents the logical address space of the dataset while the bottom rectangle represents an external file.

|  |
| --- |
| Dsets_fig19.JPG |
| Figure 5-11. External file storage |

The example below shows code that defines the external storage for the example. Note that the segments are defined in order of the logical addresses they represent, not their order within the external file. It would also have been possible to put the segments in separate files. Care should be taken when setting up segments in a single file since the library does not automatically check for segments that overlap.

|  |
| --- |
| Plist = H5Pcreate (H5P\_DATASET\_CREATE);  H5Pset\_external (plist, "velocity.data", 3000, 1000);  H5Pset\_external (plist, "velocity.data", 0, 2500);  H5Pset\_external (plist, "velocity.data", 4500, 1500); |
| Code Example 5-8. External storage |

The figure below shows an example of how a contiguous, two-dimensional dataset is partitioned into three parts and each of those parts is stored in a separate external file. The top rectangle represents the logical address space of the dataset while the bottom rectangles represent external files.

|  |
| --- |
| Dsets_fig20.jpg |
| Figure 5-12. Partitioning a 2-D dataset for external storage |

The example below shows code for the partitioning described above. In this example, the library maps the multi-dimensional array onto a linear address space as defined by the HDF5 format specification, and then maps that address space into the segments defined in the external file list.

|  |
| --- |
| plist = H5Pcreate (H5P\_DATASET\_CREATE);  H5Pset\_external (plist, "scan1.data", 0, 24);  H5Pset\_external (plist, "scan2.data", 0, 24);  H5Pset\_external (plist, "scan3.data", 0, 16); |
| Code Example 5-9. Partitioning a 2-D dataset for external storage |

 The segments of an external file can exist beyond the end of the (external) file. The library reads that part of a segment as zeros. When writing to a segment that exists beyond the end of a file, the external file is automatically extended. Using this feature, one can create a segment (or set of segments) which is larger than the current size of the dataset. This allows the dataset to be extended at a future time (provided the dataspace also allows the extension).

All referenced external data files must exist before performing raw data I/O on the dataset. This is nor­mally not a problem since those files are being managed directly by the application or indirectly through some other library. However, if the file is transferred from its original context, care must be taken to assure that all the external files are accessible in the new location.

## Using HDF5 Filters

This section describes in detail how to use the n-bit and scale-offset filters.

### Using the N-bit Filter

N-bit data has n significant bits, where n may not correspond to a precise number of bytes. On the other hand, computing systems and applications universally, or nearly so, run most efficiently when manipulat­ing data as whole bytes or multiple bytes.

Consider the case of 12-bit integer data. In memory, that data will be handled in at least 2 bytes, or 16 bits, and on some platforms in 4 or even 8 bytes. The size of such a dataset can be significantly reduced when written to disk if the unused bits are stripped out.

The n-bit filter is provided for this purpose, packing n-bit data on output by stripping off all unused bits and unpacking on input, restoring the extra bits required by the computational processor.

**N-bit Datatype**

An n-bit datatype is a datatype of n significant bits. Unless it is packed, an n-bit datatype is presented as an n-bit bitfield within a larger-sized value. For example, a 12-bit datatype might be presented as a 12-bit field in a 16-bit, or 2-byte, value.

Currently, the datatype classes of n-bit datatype or n-bit field of a compound datatype or an array data­type are limited to integer or floating-point.

The HDF5 user can create an n-bit datatype through a series of function calls. For example, the follow­ing calls create a 16-bit datatype that is stored in a 32-bit value with a 4-bit offset:

hid\_t nbit\_datatype = H5Tcopy(H5T\_STD\_I32LE);

H5Tset\_precision(nbit\_datatype, 16);

H5Tset\_offset(nbit\_datatype, 4);

In memory, one value of the above example n-bit datatype would be stored on a little-endian machine as follows:

|  |  |  |  |
| --- | --- | --- | --- |
| byte 3 | byte 2 | byte 1 | byte 0 |
| ???????? | ????SPPP | PPPPPPPP | PPPP???? |

*Note: Key: S - sign bit, P - significant bit, ? - padding bit. Sign bit is included in signed integer datatype precision.*

**N-bit Filter**

When data of an n-bit datatype is stored on disk using the n-bit filter, the filter packs the data by stripping off the padding bits; only the significant bits are retained and stored. The values on disk will appear as fol­lows:

|  |  |  |
| --- | --- | --- |
| 1st value | 2nd value |  |
| SPPPPPPP PPPPPPPP | SPPPPPPP PPPPPPPP | ... |

*Note: Key: S - sign bit, P - significant bit, ? - padding bit. Sign bit is included in signed integer datatype precision.*

The n-bit filter can be used effectively for compressing data of an n-bit datatype, including arrays and the n-bit fields of compound datatypes. The filter supports complex situations where a compound datatype contains member(s) of a compound datatype or an array datatype has a compound datatype as the base type.

At present, the n-bit filter supports all datatypes. For datatypes of class time, string, opaque, reference, ENUM, and variable-length, the n-bit filter acts as a no-op which is short for no operation. For conve­nience, the rest of this section refers to such datatypes as no-op datatypes.

As is the case with all HDF5 filters, an application using the n-bit filter must store data with chunked stor­age.

**How Does the N-bit Filter Work?**

The n-bit filter always compresses and decompresses according to dataset properties supplied by the HDF5 library in the datatype, dataspace, or dataset creation property list.

The dataset datatype refers to how data is stored in an HDF5 file while the memory datatype refers to how data is stored in memory. The HDF5 library will do datatype conversion when writing data in memory to the dataset or reading data from the dataset to memory if the memory datatype differs from the dataset datatype. Datatype conversion is performed by HDF5 library before n-bit compression and after n-bit decompression.

The following sub-sections examine the common cases:

•        N-bit integer conversions

•        N-bit floating-point conversions

**N-bit Integer Conversions**

Integer data with a dataset of integer datatype of less than full precision and a memory datatype of H5T\_NATIVE\_INT, provides the simplest application of the n-bit filter.

The precision of H5T\_NATIVE\_INT is 8 multiplied by sizeof(int). This value, the size of an int in bytes, differs from platform to platform; we assume a value of 4 for the following illustration. We further assume the memory byte order to be little-endian.

In memory, therefore, the precision of H5T\_NATIVE\_INT is 32 and the offset is 0. One value of H5T\_NA­TIVE\_INT is laid out in memory as follows:

|  |
| --- |
| Dsets_NbitInteger1.JPG |
| Figure 5-13. H5T\_NATIVE\_INT in memory |

*Note: Key: S - sign bit, P - significant bit, ? - padding bit. Sign bit is included in signed integer datatype precision.*

Suppose the dataset datatype has a precision of 16 and an offset of 4. After HDF5 converts values from the memory datatype to the dataset datatype, it passes something like the following to the n-bit filter for compression:

|  |
| --- |
| Dsets_NbitInteger2.JPG |
| Figure 5-14. Passed to the n-bit filter |

*Note: Key: S - sign bit, P - significant bit, ? - padding bit. Sign bit is included in signed integer datatype precision.*

Notice that only the specified 16 bits (15 significant bits and the sign bit) are retained in the conversion. All other significant bits of the memory datatype are discarded because the dataset datatype calls for only 16 bits of precision. After n-bit compression, none of these discarded bits, known as padding bits will be stored on disk.

**N-bit Floating-point Conversions**

Things get more complicated in the case of a floating-point dataset datatype class. This sub-section pro­vides an example that illustrates the conversion from a memory datatype of H5T\_NATIVE\_FLOAT to a dataset datatype of class floating-point.

As before, let the H5T\_NATIVE\_FLOAT be 4 bytes long, and let the memory byte order be little-endian. Per the IEEE standard, one value of H5T\_NATIVE\_FLOAT is laid out in memory as follows:

|  |
| --- |
| Dsets_NbitFloating1.JPG |
| Figure 5-15. H5T\_NATIVE\_FLOAT in memory |

*Note: Key: S - sign bit, E - exponent bit, M - mantissa bit, ? - padding bit. Sign bit is included in floating-point datatype precision.*

Suppose the dataset datatype has a precision of 20, offset of 7, mantissa size of 13, mantissa position of 7, exponent size of 6, exponent position of 20, and sign position of 26. For more information, see "Definition of Datatypes."

After HDF5 converts values from the memory datatype to the dataset datatype, it passes something like the following to the n-bit filter for compression:

|  |
| --- |
| Dsets_NbitFloating2.JPG |
| Figure 5-16. Passed to the n-bit filter |

*Note: Key: S - sign bit, E - exponent bit, M - mantissa bit, ? - padding bit. Sign bit is included in floating-point datatype precision.*

The sign bit and truncated mantissa bits are not changed during datatype conversion by the HDF5 library. On the other hand, the conversion of the 8-bit exponent to a 6-bit exponent is a little tricky:

The bias for the new exponent in the n-bit datatype is:

2(n-1)-1

The following formula is used for this exponent conversion:

exp8 - (2(8-1)-1) = exp6 - (2(6-1)-1) = actual exponent value

where exp8 is the stored decimal value as represented by the 8-bit exponent, and exp6 is the stored decimal value as represented by the 6-bit exponent.

In this example, caution must be taken to ensure that, after conversion, the actual exponent value is within the range that can be represented by a 6-bit exponent. For example, an 8-bit exponent can repre­sent values from -127 to 128 while a 6-bit exponent can represent values only from -31 to 32.

**N-bit Filt****er Behavior**

The n-bit filter was designed to treat the incoming data byte by byte at the lowest level. The purpose was to make the n-bit filter as generic as possible so that no pointer cast related to the datatype is needed.

Bitwise operations are employed for packing and unpacking at the byte level.

Recursive function calls are used to treat compound and array datatypes.

**N-bit Compression**

The main idea of n-bit compression is to use a loop to compress each data element in a chunk. Depending on the datatype of each element, the n-bit filter will call one of four functions. Each of these functions per­forms one of the following tasks:

•        Compress a data element of a no-op datatype

•        Compress a data element of an atomic datatype

•        Compress a data element of a compound datatype

•        Compress a data element of an array datatype

No-op datatypes: The n-bit filter does not actually compress no-op datatypes. Rather, it copies the data buffer of the no-op datatype from the non-compressed buffer to the proper location in the compressed buffer; the compressed buffer has no holes. The term “compress” is used here simply to distinguish this function from the function that performs the reverse operation during decompression.

Atomic datatypes: The n-bit filter will find the bytes where significant bits are located and try to compress these bytes, one byte at a time, using a loop. At this level, the filter needs the following information:

•        The byte offset of the beginning of the current data element with respect to the

beginning of the input data buffer

•        Datatype size, precision, offset, and byte order

The n-bit filter compresses from the most significant byte containing significant bits to the least significant byte. For big-endian data, therefore, the loop index progresses from smaller to larger while for little-endian, the loop index progresses from larger to smaller.

In the extreme case of when the n-bit datatype has full precision, this function copies the content of the entire non-compressed datatype to the compressed output buffer.

Compound datatypes: The n-bit filter will compress each data member of the compound datatype. If the member datatype is of an integer or floating-point datatype, the n-bit filter will call the function described above. If the member datatype is of a no-op datatype, the filter will call the function described above. If the member datatype is of a compound datatype, the filter will make a recursive call to itself. If the mem­ber datatype is of an array datatype, the filter will call the function described below.

Array datatypes: The n-bit filter will use a loop to compress each array element in the array. If the base datatype of array element is of an integer or floating-point datatype, the n-bit filter will call the function described above. If the base datatype is of a no-op datatype, the filter will call the function described above. If the base datatype is of a compound datatype, the filter will call the function described above. If the member datatype is of an array datatype, the filter will make a recursive call of itself.

**N-bit Decompression**

The n-bit decompression algorithm is very similar to n-bit compression. The only difference is that at the byte level, compression packs out all padding bits and stores only significant bits into a continuous buffer (unsigned char) while decompression unpacks significant bits and inserts padding bits (zeros) at the proper positions to recover the data bytes as they existed before compression.

**Storing N-bit Parameters to Array cd\_value[]**

All of the information, or parameters, required by the n-bit filter are gathered and stored in the array cd\_values[] by the private function H5Z\_set\_local\_nbit and are passed to another private func­tion, H5Z\_filter\_nbit, by the HDF5 Library.

These parameters are as follows:

•        Parameters related to the datatype

•        The number of elements within the chunk

•        A flag indicating whether compression is needed

The first and second parameters can be obtained using the HDF5 dataspace and datatype interface calls.

A compound datatype can have members of array or compound datatype. An array datatype’s base data­type can be a complex compound datatype. Recursive calls are required to set parameters for these com­plex situations.

Before setting the parameters, the number of parameters should be calculated to dynamically allocate the array cd\_values[], which will be passed to the HDF5 Library. This also requires recursive calls.

For an atomic datatype (integer or floating-point), parameters that will be stored include the datatype’s size, endianness, precision, and offset.

For a no-op datatype, only the size is required.

For a compound datatype, parameters that will be stored include the datatype’s total size and number of members. For each member, its member offset needs to be stored. Other parameters for members will depend on the respective datatype class.

For an array datatype, the total size parameter should be stored. Other parameters for the array’s base type depend on the base type’s datatype class.

Further, to correctly retrieve the parameter for use of n-bit compression or decompression later, parame­ters for distinguishing between datatype classes should be stored.

**Implementa****tion**

Three filter callback functions were written for the n-bit filter:

•        H5Z\_can\_apply\_nbit

•        H5Z\_set\_local\_nbit

•        H5Z\_filter\_nbit

These functions are called internally by the HDF5 library. A number of utility functions were written for the function H5Z\_set\_local\_nbit. Compression and decompression functions were written and are called by function H5Z\_filter\_nbit. All these functions are included in the file H5Znbit.c.

The public function H5Pset\_nbit is called by the application to set up the use of the n-bit filter. This function is included in the file H5Pdcpl.c. The application does not need to supply any parameters.

**How N-bit Parameters are Stored**

A scheme of storing parameters required by the n-bit filter in the array cd\_values[] was developed uti­lizing recursive function calls.

Four private utility functions were written for storing the parameters associated with atomic (integer or floating-point), no-op, array, and compound datatypes:

•        H5Z\_set\_parms\_atomic

•        H5Z\_set\_parms\_array

•        H5Z\_set\_parms\_nooptype

•  H5Z\_set\_parms\_compound

The scheme is briefly described below.

First, assign a numeric code for datatype class atomic (integer or float), no-op, array, and compound data­type. The code is stored before other datatype related parameters are stored.

The first three parameters of cd\_values[] are reserved for:

1.      The number of valid entries in the array cd\_values[]

2.      A flag indicating whether compression is needed

3.      The number of elements in the chunk

Throughout the balance of this explanation, i represents the index of cd\_values[].

In the function H5Z\_set\_local\_nbit:

1.      i = 2

2.      Get the number of elements in the chunk and store in cd\_value[i]; increment i

3.      Get the class of the datatype:

•        For an integer or floating-point datatype, call H5Z\_set\_parms\_atomic

•        For an array datatype, call H5Z\_set\_parms\_array

•        For a compound datatype, call H5Z\_set\_parms\_compound

•        For none of the above, call H5Z\_set\_parms\_noopdatatype

4.      Store i in cd\_value[0] and flag in cd\_values[1]

In the function H5Z\_set\_parms\_atomic:

1.      Store the assigned numeric code for the atomic datatype in cd\_value[i]; increment i

2.      Get the size of the atomic datatype and store in cd\_value[i]; increment i

3.      Get the order of the atomic datatype and store in cd\_value[i]; increment i

4.      Get the precision of the atomic datatype and store in cd\_value[i]; increment i

5.      Get the offset of the atomic datatype and store in cd\_value[i]; increment i

6.      Determine the need to do compression at this point

In the function H5Z\_set\_parms\_nooptype:

1.      Store the assigned numeric code for the no-op datatype in cd\_value[i]; increment i

2.      Get the size of the no-op datatype and store in cd\_value[i]; increment i

In the function H5Z\_set\_parms\_array:

1.      Store the assigned numeric code for the array datatype in cd\_value[i]; increment i

2.      Get the size of the array datatype and store in cd\_value[i]; increment i

3.      Get the class of the array’s base datatype.

•        For an integer or floating-point datatype, call H5Z\_set\_parms\_atomic

•        For an array datatype, call H5Z\_set\_parms\_array

•        For a compound datatype, call H5Z\_set\_parms\_compound

•        If none of the above, call H5Z\_set\_parms\_noopdatatype

In the function H5Z\_set\_parms\_compound:

1.      Store the assigned numeric code for the compound datatype in cd\_value[i]; increment i

2.      Get the size of the compound datatype and store in cd\_value[i]; increment i

3.      Get the number of members and store in cd\_values[i]; increment i

4.      For each member

•        Get the member offset and store in cd\_values[i]; increment i

•        Get the class of the member datatype

•        For an integer or floating-point datatype, call H5Z\_set\_parms\_atomic

•        For an array datatype, call H5Z\_set\_parms\_array

•        For a compound datatype, call H5Z\_set\_parms\_compound

•        If none of the above, call H5Z\_set\_parms\_noopdatatype

**N-bit Compression and Decompression Functions**

The n-bit compression and decompression functions above are called by the private HDF5 function H5Z\_­filter\_nbit. The compress and decompress functions retrieve the n-bit parameters from cd\_val­ues[] as it was passed by H5Z\_filter\_nbit. Parameters are retrieved in exactly the same order in which they are stored and lower-level compression and decompression functions for different datatype classes are called.

N-bit compression is not implemented in place. Due to the difficulty of calculating actual output buffer size after compression, the same space as that of the input buffer is allocated for the output buffer as passed to the compression function. However, the size of the output buffer passed by reference to the compression function will be changed (smaller) after the compression is complete.

**Usage Examples**

The following code example illustrates the use of the n-bit filter for writing and reading n-bit integer data.

|  |
| --- |
| #include "hdf5.h"  #include "stdlib.h"  #include "math.h"  #define H5FILE\_NAME  "nbit\_test\_int.h5"  #define DATASET\_NAME "nbit\_int"  #define NX 200  #define NY 300  #define CH\_NX 10  #define CH\_NY 15 |
| int main(void)  {     hid\_t   file, dataspace, dataset, datatype, mem\_datatype, dset\_create\_props;     hsize\_t dims[2], chunk\_size[2];     int     orig\_data[NX][NY];     int     new\_data[NX][NY];     int     i, j;     size\_t  precision, offset; |
| /\* Define dataset datatype (integer), and set precision,     \* offset     \*/     datatype = H5Tcopy(H5T\_NATIVE\_INT);     precision = 17; /\* precision includes sign bit \*/     if(H5Tset\_precision(datatype,precision)<0) {        printf("Error: fail to set precision\n");        return -1;     } |
| offset = 4;     if(H5Tset\_offset(datatype,offset)<0) {        printf("Error: fail to set offset\n");        return -1;     } |
| /\* Copy to memory datatype \*/     mem\_datatype = H5Tcopy(datatype); |
| /\* Set order of dataset datatype \*/     if(H5Tset\_order(datatype, H5T\_ORDER\_BE)<0) {        printf("Error: fail to set endianness\n");        return -1;     } |
| /\* Initialize data buffer with random data within correct     \* range corresponding to the memory datatype's precision     \* and offset.     \*/     for (i=0; i < NX; i++)        for (j=0; j < NY; j++)           orig\_data[i][j] = rand() % (int)pow(2, precision-1)                 <<offset; |
| /\* Describe the size of the array. \*/     dims[0] = NX;     dims[1] = NY;     if((dataspace = H5Screate\_simple (2, dims, NULL))<0) {        printf("Error: fail to create dataspace\n");        return -1;     } |
| /\*     \* Create a new file using read/write access, default file     \* creation properties, and default file access properties.     \*/ |
| if((file = H5Fcreate (H5FILE\_NAME, H5F\_ACC\_TRUNC, H5P\_DEFAULT, H5P\_DEFAULT))<0) {        printf("Error: fail to create file\n");        return -1;     } |
| /\*     \* Set the dataset creation property list to specify that     \* the raw data is to be partitioned into 10 x 15 element     \* chunks and that each chunk is to be compressed.     \*/     chunk\_size[0] = CH\_NX;     chunk\_size[1] = CH\_NY; |
| if((dset\_create\_props = H5Pcreate (H5P\_DATASET\_CREATE))<0) {        printf("Error: fail to create dataset property\n");        return -1;     }     if(H5Pset\_chunk (dset\_create\_props, 2, chunk\_size)<0) {        printf("Error: fail to set chunk\n");        return -1;     } |
| /\*     \* Set parameters for n-bit compression; check the description     \* of the H5Pset\_nbit function in the HDF5 Reference Manual     \* for more information.     \*/ |
| if(H5Pset\_nbit (dset\_create\_props)<0) {        printf("Error: fail to set nbit filter\n");        return -1;     } |
| /\*     \* Create a new dataset within the file.  The datatype     \* and dataspace describe the data on disk, which may     \* be different from the format used in the application's     \* memory.     \*/ |
| if((dataset = H5Dcreate(file, DATASET\_NAME, datatype,           dataspace, H5P\_DEFAULT,           dset\_create\_props, H5P\_DEFAULT))<0) {        printf("Error: fail to create dataset\n");        return -1;     } |
| /\*     \* Write the array to the file. The datatype and dataspace     \* describe the format of the data in the 'orig\_data' buffer.     \* The raw data is translated to the format required on disk,     \* as defined above. We use default raw data transfer     \* properties.     \*/ |
| if(H5Dwrite (dataset, mem\_datatype, H5S\_ALL, H5S\_ALL,           H5P\_DEFAULT, orig\_data)<0) {        printf("Error: fail to write to dataset\n");        return -1;     } |
| H5Dclose (dataset);       if((dataset = H5Dopen(file, DATASET\_NAME, H5P\_DEFAULT))<0) {        printf("Error: fail to open dataset\n");        return -1;     } |
| /\*     \* Read the array. This is similar to writing data,     \* except the data flows in the opposite direction.     \* Note: Decompression is automatic.     \*/ |
| if(H5Dread (dataset, mem\_datatype, H5S\_ALL, H5S\_ALL,           H5P\_DEFAULT, new\_data)<0) {        printf("Error: fail to read from dataset\n");        return -1;     } |
| H5Tclose (datatype);     H5Tclose (mem\_datatype);     H5Dclose (dataset); |
| H5Sclose (dataspace);     H5Pclose (dset\_create\_props);     H5Fclose (file);       return 0;  } |
| Code Example 5-10. N-bit compression for integer data |

*Note: The code example above illustrates the use of the n-bit filter for writing and reading n-bit integer data.*

The following code example illustrates the use of the n-bit filter for writing and reading n-bit floating-point data.

|  |
| --- |
| #include "hdf5.h"  #define H5FILE\_NAME  "nbit\_test\_float.h5"  #define DATASET\_NAME "nbit\_float"  #define NX 2  #define NY 5  #define CH\_NX 2  #define CH\_NY 5 |
| int main(void)  {     hid\_t   file, dataspace, dataset, datatype, dset\_create\_props;     hsize\_t dims[2], chunk\_size[2];     /\* orig\_data[] are initialized to be within the range that     \* can be represented by dataset datatype (no precision     \* loss during datatype conversion)     \*/ |
| float   orig\_data[NX][NY] = {{188384.00, 19.103516,     -1.0831790e9, -84.242188, 5.2045898}, {-49140.000,     2350.2500, -3.2110596e-1,    6.4998865e-5, -0.0000000}};     float   new\_data[NX][NY];     size\_t  precision, offset; |
| /\* Define single-precision floating-point type for dataset     \*---------------------------------------------------------------     \* size=4 byte, precision=20 bits, offset=7 bits,     \* mantissa size=13 bits, mantissa position=7, |
| \* exponent size=6 bits, exponent position=20,     \* exponent bias=31.     \* It can be illustrated in little-endian order as:     \* (S - sign bit, E - exponent bit, M - mantissa bit,     \*  ? - padding bit)     \* |
| \*      3            2            1         0     \* ?????SEE EEEEMMMM MMMMMMMM M???????     \* |
| \* To create a new floating-point type, the following     \* properties must be set in the order of     \* set fields -> set offset -> set precision -> set size. |
| \* All these properties must be set before the type can     \* function. Other properties can be set anytime. Derived     \* type size cannot be expanded bigger than original size     \* but can be decreased. There should be no holes     \* among the significant bits. Exponent bias usually     \* is set 2^(n-1)-1, where n is the exponent size.     \*---------------------------------------------------------------\*/ |
| datatype = H5Tcopy(H5T\_IEEE\_F32BE);     if(H5Tset\_fields(datatype, 26, 20, 6, 7, 13)<0) {        printf("Error: fail to set fields\n");        return -1;     } |
| offset = 7;     if(H5Tset\_offset(datatype,offset)<0) {        printf("Error: fail to set offset\n");        return -1;     } |
| precision = 20;     if(H5Tset\_precision(datatype,precision)<0) {        printf("Error: fail to set precision\n");        return -1;     } |
| if(H5Tset\_size(datatype, 4)<0) {        printf("Error: fail to set size\n");        return -1;     }     if(H5Tset\_ebias(datatype, 31)<0) {        printf("Error: fail to set exponent bias\n");        return -1;     } |
| /\* Describe the size of the array. \*/     dims[0] = NX;     dims[1] = NY;     if((dataspace = H5Screate\_simple (2, dims, NULL))<0) {        printf("Error: fail to create dataspace\n");        return -1;     } |
| /\*     \* Create a new file using read/write access, default file     \* creation properties, and default file access properties.     \*/ |
| if((file = H5Fcreate (H5FILE\_NAME, H5F\_ACC\_TRUNC,           H5P\_DEFAULT, H5P\_DEFAULT))<0) {        printf("Error: fail to create file\n");        return -1;     } |
| /\*     \* Set the dataset creation property list to specify that     \* the raw data is to be partitioned into 2 x 5 element     \* chunks and that each chunk is to be compressed.     \*/ |
| chunk\_size[0] = CH\_NX;     chunk\_size[1] = CH\_NY;     if((dset\_create\_props = H5Pcreate (H5P\_DATASET\_CREATE))<0) {        printf("Error: fail to create dataset property\n");        return -1;     } |
| if(H5Pset\_chunk (dset\_create\_props, 2, chunk\_size)<0) {        printf("Error: fail to set chunk\n");        return -1;     } |
| /\*     \* Set parameters for n-bit compression; check the description     \* of the H5Pset\_nbit function in the HDF5 Reference Manual     \* for more information.     \*/     if(H5Pset\_nbit (dset\_create\_props)<0) {        printf("Error: fail to set nbit filter\n");        return -1;     } |
| /\*     \* Create a new dataset within the file. The datatype     \* and dataspace describe the data on disk, which may     \* be different from the format used in the application's     \* memory.     \*/ |
| if((dataset = H5Dcreate(file, DATASET\_NAME, datatype,           dataspace, H5P\_DEFAULT,           dset\_create\_plists, H5P\_DEFAULT))<0) {        printf("Error: fail to create dataset\n");        return -1;     } |
| /\*     \* Write the array to the file. The datatype and dataspace     \* describe the format of the data in the 'orig\_data' buffer.     \* The raw data is translated to the format required on disk,     \* as defined above. We use default raw data transfer     \* properties.     \*/ |
| if(H5Dwrite (dataset, H5T\_NATIVE\_FLOAT, H5S\_ALL, H5S\_ALL,           H5P\_DEFAULT, orig\_data)<0) {        printf("Error: fail to write to dataset\n");        return -1;     } |
| H5Dclose (dataset);       if((dataset = H5Dopen(file, DATASET\_NAME, H5P\_DEFAULT))<0) {        printf("Error: fail to open dataset\n");        return -1;     } |
| /\*     \* Read the array. This is similar to writing data,     \* except the data flows in the opposite direction.     \* Note: Decompression is automatic.     \*/ |
| if(H5Dread (dataset, H5T\_NATIVE\_FLOAT, H5S\_ALL, H5S\_ALL,           H5P\_DEFAULT, new\_data)<0) {        printf("Error: fail to read from dataset\n");        return -1;     } |
| H5Tclose (datatype);     H5Dclose (dataset);     H5Sclose (dataspace);     H5Pclose (dset\_create\_props);     H5Fclose (file);       return 0;  } |
| Code Example 5-11. N-bit compression for floating-point data |

*Note: The code example above illustrates the use of the n-bit filter for writing and reading n-bit floating-point data.*

**Limita****tions**

Because the array cd\_values[] has to fit into an object header message of 64K, the n-bit filter has an upper limit on the number of n-bit parameters that can be stored in it. To be conservative, a maximum of 4K is allowed for the number of parameters.

The n-bit filter currently only compresses n-bit datatypes or fields derived from integer or floating-point datatypes. The n-bit filter assumes padding bits of zero. This may not be true since the HDF5 user can set padding bit to be zero, one, or leave the background alone. However, it is expected the n-bit filter will be modified to adjust to such situations.

The n-bit filter does not have a way to handle the situation where the fill value of a dataset is defined and the fill value is not of an n-bit datatype although the dataset datatype is.

### Using the Scale-offset Filter

Generally speaking, scale-offset compression performs a scale and/or offset operation on each data value and truncates the resulting value to a minimum number of bits (minimum-bits) before storing it.

The current scale-offset filter supports integer and floating-point datatypes only. For the floating-point datatype, float and double are supported, but long double is not supported.

Integer data compression uses a straight-forward algorithm. Floating-point data compression adopts the GRiB data packing mechanism which offers two alternate methods: a fixed minimum-bits method, and a variable minimum-bits method. Currently, only the variable minimum-bits method is implemented.

Like other I/O filters supported by the HDF5 library, applications using the scale-offset filter must store data with chunked storage.

Integer type: The minimum-bits of integer data can be determined by the filter. For example, if the maxi­mum value of data to be compressed is 7065 and the minimum value is 2970. Then the “span” of dataset values is equal to (max-min+1), which is 4676. If no fill value is defined for the dataset, the minimum-bits is: ceiling(log2(span)) = 12. With fill value set, the minimum-bits is: ceiling(log2(span+1)) = 13.

HDF5 users can also set the minimum-bits. However, if the user gives a minimum-bits that is less than that calculated by the filter, the compression will be lossy.

Floating-point type: The basic idea of the scale-offset filter for the floating-point type is to transform the data by some kind of scaling to integer data, and then to follow the procedure of the scale-offset filter for the integer type to do the data compression. Due to the data transformation from floating-point to inte­ger, the scale-offset filter is lossy in nature.

Two methods of scaling the floating-point data are used: the so-called D-scaling and E-scaling. D-scaling is more straightforward and easy to understand. For HDF5 1.8 release, only the D-scaling method has been implemented.

**Design**

Before the filter does any real work, it needs to gather some information from the HDF5 Library through API calls. The parameters the filter needs are:

•        The minimum-bits of the data value

•        The number of data elements in the chunk

•        The datatype class, size, sign (only for integer type), byte order, and fill value if

defined

Size and sign are needed to determine what kind of pointer cast to use when retrieving values from the data buffer.

The pipeline of the filter can be divided into four parts: (1)pre-compression; (2)compression; (3)decom­pression; (4)post-decompression.

Depending on whether a fill value is defined or not, the filter will handle pre-compression and post-decompression differently.

The scale-offset filter only needs the memory byte order, size of datatype, and minimum-bits for compres­sion and decompression.

Since decompression has no access to the original data, the minimum-bits and the minimum value need to be stored with the compressed data for decompression and post-decompression.

**Integer Type**

Pre-compression: During pre-compression minimum-bits is calculated if it is not set by the user. For more information on how minimum-bits are calculated, see section 6.1. “The N-bit Filter.”

If the fill value is defined, finding the maximum and minimum values should ignore the data element whose value is equal to the fill value.

If no fill value is defined, the value of each data element is subtracted by the minimum value during this stage.

If the fill value is defined, the fill value is assigned to the maximum value. In this way minimum-bits can represent a data element whose value is equal to the fill value and subtracts the minimum value from a data element whose value is not equal to the fill value.

The fill value (if defined), the number of elements in a chunk, the class of the datatype, the size of the datatype, the memory order of the datatype, and other similar elements will be stored in the HDF5 object header for the post-decompression usage.

After pre-compression, all values are non-negative and are within the range that can be stored by mini­mum-bits.

Compression: All modified data values after pre-compression are packed together into the compressed data buffer. The number of bits for each data value decreases from the number of bits of integer (32 for most platforms) to minimum-bits. The value of minimum-bits and the minimum value are added to the data buffer and the whole buffer is sent back to the library. In this way, the number of bits for each modi­fied value is no more than the size of minimum-bits.

Decompression: In this stage, the number of bits for each data value is resumed from minimum-bits to the number of bits of integer.

Post-decompression: For the post-decompression stage, the filter does the opposite of what it does during pre-compression except that it does not calculate the minimum-bits or the minimum value. These values were saved during compression and can be retrieved through the resumed data buffer. If no fill value is defined, the filter adds the minimum value back to each data element.

If the fill value is defined, the filter assigns the fill value to the data element whose value is equal to the maximum value that minimum-bits can represent and adds the minimum value back to each data element whose value is not equal to the maximum value that minimum-bits can represent.

**Floating-poi****nt Type**

The filter will do data transformation from floating-point type to integer type and then handle the data by using the procedure for handling the integer data inside the filter. Insignificant bits of floating-point data will be cut off during data transformation, so this filter is a lossy compression method.

There are two scaling methods: D-scaling and E-scaling. The HDF5 1.8 release only supports D-scaling. D-scaling is short for decimal scaling. E-scaling should be similar conceptually. In order to transform data from floating-point to integer, a scale factor is introduced. The minimum value will be calculated. Each data element value will subtract the minimum value. The modified data will be multiplied by 10 (Decimal) to the power of scale\_factor, and only the integer part will be kept and manipulated through the rou­tines for the integer type of the filter during pre-compression and compression. Integer data will be divided by 10 to the power of scale\_factor to transform back to floating-point data during decompres­sion and post-decompression. Each data element value will then add the minimum value, and the floating-point data are resumed. However, the resumed data will lose some insignificant bits compared with the original value.

For example, the following floating-point data are manipulated by the filter, and the D-scaling factor is 2.

{104.561, 99.459, 100.545, 105.644}

The minimum value is 99.459, each data element subtracts 99.459, the modified data is

{5.102, 0, 1.086, 6.185}

Since the D-scaling factor is 2, all floating-point data will be multiplied by 10^2 with this result:

{510.2, 0, 108.6, 618.5}

The digit after decimal point will be rounded off, and then the set looks like:

{510, 0, 109, 619}

After decompression, each value will be divided by 10^2 and will be added to the offset 99.459.

The floating-point data becomes

{104.559, 99.459, 100.549, 105.649}.

The relative error for each value should be no more than 5\* (10^(D-scaling factor +1)). D-scaling some­times is also referred as a variable minimum-bits method since for different datasets the minimum-bits to represent the same decimal precision will vary. The data value is modified to 2 to power of scale\_fac­tor for E-scaling. E-scaling is also called fixed-bits method since for different datasets the minimum-bits will always be fixed to the scale factor of E-scaling. Currently, HDF5 ONLY supports the D-scaling (variable minimum-bits) method.

**Implementation**

The scale-offset filter implementation was written and included in the file H5Zscaleoffset.c. Function H5Pset\_scaleoffset was written and included in the file “H5Pdcpl.c”. The HDF5 user can supply minimum-bits by calling function H5Pset\_scaleoffset.

The scale-offset filter was implemented based on the design outlined in this section. However, the follow­ing factors need to be considered:

1.      The filter needs the appropriate cast pointer whenever it needs to retrieve data values.

2.      The HDF5 Library passes to the filter the to-be-compressed data in the format of the dataset data­type, and the filter passes back the decompressed data in the same format. If a fill value is defined, it is also in dataset datatype format. For example, if the byte order of the dataset data­type is different from that of the memory datatype of the platform, compression or decompres­sion performs an endianness conversion of data buffer. Moreover, it should be aware that memory byte order can be different during compression and decompression.

3.      The difference of endianness and datatype between file and memory should be considered when saving and retrieval of minimum-bits, minimum value, and fill value.

4.      If the user sets the minimum-bits to full precision of the datatype, no operation is needed at the filter side. If the full precision is a result of calculation by the filter, then the minimum-bits needs to be saved for decompression but no compression or decompression is needed (only a copy of the input buffer is needed).

5.      If by calculation of the filter, the minimum-bits is equal to zero, special handling is needed. Since it means all values are the same, no compression or decompression is needed. But the minimum-bits and minimum value still need to be saved during compression.

6.      For floating-point data, the minimum value of the dataset should be calculated at first. Each data element value will then subtract the minimum value to obtain the “offset” data. The offset data will then follow the steps outlined above in the discussion of floating-point types to do data trans­formation to integer and rounding. For more information, see "Floating-point Type."

**Usage Examples**

The following code example illustrates the use of the scale-offset filter for writing and reading integer data.

|  |
| --- |
| #include "hdf5.h"  #include "stdlib.h"  #define H5FILE\_NAME  "scaleoffset\_test\_int.h5"  #define DATASET\_NAME "scaleoffset\_int"  #define NX 200  #define NY 300  #define CH\_NX 10  #define CH\_NY 15 |
| int main(void)  {     hid\_t   file, dataspace, dataset, datatype, dset\_create\_props;     hsize\_t dims[2], chunk\_size[2];     int     orig\_data[NX][NY];     int     new\_data[NX][NY];     int     i, j, fill\_val; |
| /\* Define dataset datatype \*/     datatype = H5Tcopy(H5T\_NATIVE\_INT);       /\* Initiliaze data buffer \*/     for (i=0; i < NX; i++)        for (j=0; j < NY; j++)           orig\_data[i][j] = rand() % 10000; |
| /\* Describe the size of the array. \*/     dims[0] = NX;     dims[1] = NY;     if((dataspace = H5Screate\_simple (2, dims, NULL))<0) {        printf("Error: fail to create dataspace\n");        return -1;     } |
| /\*     \* Create a new file using read/write access, default file     \* creation properties, and default file access properties.     \*/     if((file = H5Fcreate (H5FILE\_NAME, H5F\_ACC\_TRUNC,           H5P\_DEFAULT, H5P\_DEFAULT))<0) {        printf("Error: fail to create file\n");        return -1;     } |
| /\*     \* Set the dataset creation property list to specify that     \* the raw data is to be partitioned into 10 x 15 element     \* chunks and that each chunk is to be compressed.     \*/     chunk\_size[0] = CH\_NX;     chunk\_size[1] = CH\_NY; |
| if((dset\_create\_props = H5Pcreate (H5P\_DATASET\_CREATE))<0) {        printf("Error: fail to create dataset property\n");        return -1;     }     if(H5Pset\_chunk (dset\_create\_props, 2, chunk\_size)<0) {        printf("Error: fail to set chunk\n");        return -1;     } |
| /\* Set the fill value of dataset \*/     fill\_val = 10000;     if (H5Pset\_fill\_value(dset\_create\_props, H5T\_NATIVE\_INT,           &fill\_val)<0) {        printf("Error: can not set fill value for dataset\n");        return -1;     } |
| /\*     \* Set parameters for scale-offset compression. Check the     \* description of the H5Pset\_scaleoffset function in the     \* HDF5 Reference Manual for more information [3].     \*/ |
| if(H5Pset\_scaleoffset (dset\_create\_props, H5Z\_SO\_INT,           H5Z\_SO\_INT\_MINIMUMBITS\_DEFAULT)<0) {        printf("Error: fail to set scaleoffset filter\n");        return -1;     } |
| /\*     \* Create a new dataset within the file. The datatype     \* and dataspace describe the data on disk, which may     \* or may not be different from the format used in the     \* application's memory. The link creation and     \* dataset access property list parameters are passed     \* with default values.     \*/ |
| if((dataset = H5Dcreate (file, DATASET\_NAME, datatype,           dataspace, H5P\_DEFAULT,           dset\_create\_props, H5P\_DEFAULT))<0) {        printf("Error: fail to create dataset\n");        return -1;     } |
| /\*     \* Write the array to the file. The datatype and dataspace     \* describe the format of the data in the 'orig\_data' buffer.     \* We use default raw data transfer properties.     \*/ |
| if(H5Dwrite (dataset, H5T\_NATIVE\_INT, H5S\_ALL, H5S\_ALL,           H5P\_DEFAULT, orig\_data)<0) {        printf("Error: fail to write to dataset\n");        return -1;     }       H5Dclose (dataset); |
| if((dataset = H5Dopen(file, DATASET\_NAME, H5P\_DEFAULT))<0) {        printf("Error: fail to open dataset\n");        return -1;     } |
| /\*     \* Read the array. This is similar to writing data,     \* except the data flows in the opposite direction.     \* Note: Decompression is automatic.     \*/ |
| if(H5Dread (dataset, H5T\_NATIVE\_INT, H5S\_ALL, H5S\_ALL,           H5P\_DEFAULT, new\_data)<0) {        printf("Error: fail to read from dataset\n");        return -1;     } |
| H5Tclose (datatype);     H5Dclose (dataset);     H5Sclose (dataspace);     H5Pclose (dset\_create\_props);     H5Fclose (file);       return 0;  } |
| Code Example 5-12. Scale-offset compression integer data |

*Note: The code example above illustrates the use of the scale-offset filter for writing and reading integer data.*

The following code example illustrates the use of the scale-offset filter (set for variable minimum-bits method) for writing and reading floating-point data.

|  |
| --- |
| #include "hdf5.h"  #include "stdlib.h"  #define H5FILE\_NAME  "scaleoffset\_test\_float\_Dscale.h5"  #define DATASET\_NAME "scaleoffset\_float\_Dscale"  #define NX 200  #define NY 300  #define CH\_NX 10  #define CH\_NY 15 |
| int main(void)  {     hid\_t   file, dataspace, dataset, datatype, dset\_create\_props;     hsize\_t dims[2], chunk\_size[2];     float   orig\_data[NX][NY];     float   new\_data[NX][NY];     float   fill\_val;     int     i, j; |
| /\* Define dataset datatype \*/     datatype = H5Tcopy(H5T\_NATIVE\_FLOAT);       /\* Initiliaze data buffer \*/     for (i=0; i < NX; i++)        for (j=0; j < NY; j++)           orig\_data[i][j] = (rand() % 10000) / 1000.0; |
| /\* Describe the size of the array. \*/     dims[0] = NX;     dims[1] = NY;     if((dataspace = H5Screate\_simple (2, dims, NULL))<0) {        printf("Error: fail to create dataspace\n");        return -1;     } |
| /\*     \* Create a new file using read/write access, default file     \* creation properties, and default file access properties.     \*/ |
| if((file = H5Fcreate (H5FILE\_NAME, H5F\_ACC\_TRUNC,           H5P\_DEFAULT, H5P\_DEFAULT))<0) {        printf("Error: fail to create file\n");        return -1;     } |
| /\*     \* Set the dataset creation property list to specify that     \* the raw data is to be partitioned into 10 x 15 element     \* chunks and that each chunk is to be compressed.     \*/     chunk\_size[0] = CH\_NX;     chunk\_size[1] = CH\_NY; |
| if((dset\_create\_props = H5Pcreate (H5P\_DATASET\_CREATE))<0) {        printf("Error: fail to create dataset property\n");        return -1;     }     if(H5Pset\_chunk (dset\_create\_props, 2, chunk\_size)<0) {        printf("Error: fail to set chunk\n");        return -1;     } |
| /\* Set the fill value of dataset \*/     fill\_val = 10000.0;     if (H5Pset\_fill\_value(dset\_create\_props, H5T\_NATIVE\_FLOAT,           &fill\_val)<0) {        printf("Error: can not set fill value for dataset\n");        return -1;     } |
| /\*     \* Set parameters for scale-offset compression; use variable     \* minimum-bits method, set decimal scale factor to 3. Check     \* the description of the H5Pset\_scaleoffset function in the     \* HDF5 Reference Manual for more information [3].     \*/ |
| if(H5Pset\_scaleoffset (dset\_create\_props, H5Z\_SO\_FLOAT\_DSCALE,           3)<0)    {        printf("Error: fail to set scaleoffset filter\n");        return -1;     } |
| /\*     \* Create a new dataset within the file. The datatype     \* and dataspace describe the data on disk, which may     \* or may not be different from the format used in the     \* application's memory.     \*/ |
| if((dataset = H5Dcreate (file, DATASET\_NAME, datatype,           dataspace, H5P\_DEFAULT,           dset\_create\_props, H5P\_DEFAULT))<0) {        printf("Error: fail to create dataset\n");        return -1; |
| } |
| /\*     \* Write the array to the file. The datatype and dataspace     \* describe the format of the data in the 'orig\_data' buffer.     \* We use default raw data transfer properties.     \*/ |
| if(H5Dwrite (dataset, H5T\_NATIVE\_FLOAT, H5S\_ALL, H5S\_ALL,           H5P\_DEFAULT, orig\_data)<0) {        printf("Error: fail to write to dataset\n");        return -1;     }       H5Dclose (dataset); |
| if((dataset = H5Dopen(file, DATASET\_NAME, H5P\_DEFAULT))<0) {        printf("Error: fail to open dataset\n");        return -1;     } |
| /\*     \* Read the array. This is similar to writing data,     \* except the data flows in the opposite direction.     \* Note: Decompression is automatic.     \*/ |
| if(H5Dread (dataset, H5T\_NATIVE\_FLOAT, H5S\_ALL, H5S\_ALL,           H5P\_DEFAULT, new\_data)<0) {        printf("Error: fail to read from dataset\n");        return -1;     } |
| H5Tclose (datatype);     H5Dclose (dataset);     H5Sclose (dataspace);     H5Pclose (dset\_create\_props);     H5Fclose (file);       return 0;  } |
| Code Example 5-13. Scale-offset compression floating-point data |

*Note: The code example above illustrates the use of the scale-offset filter for writing and reading floating-point data.*

**Limitations**

For floating-point data handling, there are some algorithmic limitations to the GRiB data packing mecha­nism:

1.      Both the E-scaling and D-scaling methods are lossy compression

2.      For the D-scaling method, since data values have been rounded to integer values (positive) before truncating to the minimum-bits, their range is limited by the maximum value that can be repre­sented by the corresponding unsigned integer type (the same size as that of the floating-point type)

**Suggestions**

The following are some suggestions for using the filter for floating-point data:

1.      It is better to convert the units of data so that the units are within certain common range (for example, 1200m to 1.2km)

2.      If data values to be compressed are very near to zero, it is strongly recommended that the user sets the fill value away from zero (for example, a large positive number); if the user does nothing, the HDF5 library will set the fill value to zero, and this may cause undesirable compression results

3.      Users are not encouraged to use a very large decimal scale factor (for example, 100) for the D-scaling method; this can cause the filter not to ignore the fill value when finding maximum and minimum values, and they will get a much larger minimum-bits (poor compression)

### Using the Szip Filter

See The HDF Group website for further information regarding the Szip filter.

# HDF5 Datatypes

## Introduction and Definitions

An HDF5 dataset is an array of data elements, arranged according to the specifications of the dataspace. In general, a data element is the smallest addressable unit of storage in the HDF5 file. (Compound datatypes are the exception to this rule.) The HDF5 datatype defines the storage format for a single data element. See the figure below.

The model for HDF5 attributes is extremely similar to datasets: an attribute has a dataspace and a data­type, as shown in the figure below. The information in this chapter applies to both datasets and attributes.

|  |
| --- |
| Dtypes_fig1.JPG |
| Figure 6-1. Datatypes, dataspaces, and datasets |

Abstractly, each data element within the dataset is a sequence of bits, interpreted as a single value from a set of values (for example, a number or a character). For a given datatype, there is a standard or conven­tion for representing the values as bits, and when the bits are represented in a particular storage the bits are laid out in a specific storage scheme such as 8-bit bytes with a specific ordering and alignment of bytes within the storage array.

HDF5 datatypes implement a flexible, extensible, and portable mechanism for specifying and discovering the storage layout of the data elements, determining how to interpret the elements (for example, as float­ing point numbers), and for transferring data from different compatible layouts.

An HDF5 datatype describes one specific layout of bits. A dataset has a single datatype which applies to every data element. When a dataset is created, the storage datatype is defined. After the dataset or attri­bute is created, the datatype cannot be changed.

•        The datatype describes the storage layout of a single data element

•        All elements of the dataset must have the same type

•        The datatype of a dataset is immutable

When data is transferred (for example, a read or write), each end point of the transfer has a datatype, which describes the correct storage for the elements. The source and destination may have different (but compatible) layouts, in which case the data elements are automatically transformed during the transfer.

HDF5 datatypes describe commonly used binary formats for numbers (integers and floating point) and characters (ASCII). A given computing architecture and programming language supports certain number and character representations. For example, a computer may support 8-, 16-, 32-, and 64-bit signed inte­gers, stored in memory in little-endian byte order. These would presumably correspond to the C program­ming language types ‘char’, ‘short’, ‘int’, and ‘long’.

When reading and writing from memory, the HDF5 library must know the appropriate datatype that describes the architecture specific layout. The HDF5 library provides the platform independent ‘NATIVE’ types, which are mapped to an appropriate datatype for each platform. So the type ‘H5T\_NATIVE\_INT’ is an alias for the appropriate descriptor for each platform.

Data in memory has a datatype:

•        The storage layout in memory is architecture-specific

•        The HDF5 ‘NATIVE’ types are predefined aliases for the architecture-specific

memory layout

•        The memory datatype need not be the same as the stored datatype of the dataset

In addition to numbers and characters, an HDF5 datatype can describe more abstract classes of types including enumerations, strings, bit strings, and references (pointers to objects in the HDF5 file). HDF5 supports several classes of composite datatypes which are combinations of one or more other datatypes. In addition to the standard predefined datatypes, users can define new datatypes within the datatype classes.

The HDF5 datatype model is very general and flexible:

•        For common simple purposes, only predefined types will be needed

•        Datatypes can be combined to create complex structured datatypes

•        If needed, users can define custom atomic datatypes

•        Committed datatypes can be shared by datasets or attributes

## HDF5 Datatype Model

The HDF5 library implements an object-oriented model of datatypes. HDF5 datatypes are organized as a logical set of base types, or datatype classes. Each datatype class defines a format for representing logical values as a sequence of bits. For example the H5T\_INTEGER class is a format for representing twos com­plement integers of various sizes.

A datatype class is defined as a set of one or more datatype properties. A datatype property is a property of the bit string. The datatype properties are defined by the logical model of the datatype class. For exam­ple, the integer class (twos complement integers) has properties such as “signed or unsigned”, “length”, and “byte-order”. The float class (IEEE floating point numbers) has these properties, plus “exponent bits”, “exponent sign”, etc.

A datatype is derived from one datatype class: a given datatype has a specific value for the datatype prop­erties defined by the class. For example, for 32-bit signed integers, stored big-endian, the HDF5 datatype is a sub-type of integer with the properties set to signed=1, size=4 (bytes), and byte-order=BE.

The HDF5 datatype API (H5T functions) provides methods to create datatypes of different datatype classes, to set the datatype properties of a new datatype, and to discover the datatype properties of an existing datatype.

The datatype for a dataset is stored in the HDF5 file as part of the metadata for the dataset.

A datatype can be shared by more than one dataset in the file if the datatype is saved to the file with a name. This shareable datatype is known as a committed datatype. In the past, this kind of datatype was called a named datatype.

When transferring data (for example, a read or write), the data elements of the source and destination storage must have compatible types. As a general rule, data elements with the same datatype class are compatible while elements from different datatype classes are not compatible. When transferring data of one datatype to another compatible datatype, the HDF5 Library uses the datatype properties of the source and destination to automatically transform each data element. For example, when reading from data stored as 32-bit signed integers, big-endian into 32-bit signed integers, little-endian, the HDF5 Library will automatically swap the bytes.

Thus, data transfer operations (H5Dread, H5Dwrite, H5Aread, H5Awrite) require a datatype for both the source and the destination.

|  |
| --- |
| Dtypes_fig2.JPG |
| Figure 6-2. The datatype model |

The HDF5 library defines a set of predefined datatypes, corresponding to commonly used storage for­mats, such as twos complement integers, IEEE Floating point numbers, etc., 4- and 8-byte sizes, big-endian and little-endian byte orders. In addition, a user can derive types with custom values for the properties. For example, a user program may create a datatype to describe a 6-bit integer, or a 600-bit floating point number.

In addition to atomic datatypes, the HDF5 library supports composite datatypes. A composite datatype is an aggregation of one or more datatypes. Each class of composite datatypes has properties that describe the organization of the composite datatype. See the figure below. Composite datatypes include:

•        Compound datatypes: structured records

•        Array: a multidimensional array of a datatype

•        Variable-length: a one-dimensional array of a datatype

|  |
| --- |
| Dtypes_fig3.JPG |
| Figure 6-3. Composite datatypes |

### Datatype Classes and Properties

The figure below shows the HDF5 datatype classes. Each class is defined to have a set of properties which describe the layout of the data element and the interpretation of the bits. The table below lists the prop­erties for the datatype classes.

|  |
| --- |
| Dtypes_fig4.JPG |
| Figure 6-4. Datatype classes |

|  |  |  |  |
| --- | --- | --- | --- |
| Table 6-1. Datatype classes and their properties | | | |
| **Class** | **Description** | **Properties** | **Notes** |
| Integer | Twos complement integers | Size (bytes), precision (bits), offset (bits), pad, byte order, signed/unsigned |  |
| Float | Floating Point numbers | Size (bytes), precision (bits), offset (bits), pad, byte order, sign position, exponent position, expo­nent size (bits), exponent sign, exponent bias, man­tissa position, mantissa (size) bits, mantissa sign, mantissa normalization, internal padding | See IEEE 754 for a defini­tion of these properties. These properties describe non-IEEE 754 floating point formats as well. |
| Character | Array of 1-byte character encoding | Size (characters), Charac­ter set, byte order, pad/no pad, pad character | Currently, ASCII and UTF-8 are supported. |
| Bitfield | String of bits | Size (bytes), precision (bits), offset (bits), pad, byte order | A sequence of bit values packed into one or more bytes. |
| Opaque | Uninterpreted data | Size (bytes), precision (bits), offset (bits), pad, byte order, tag | A sequence of bytes, stored and retrieved as a block. The ‘tag’ is a string that can be used to label the value. |
| Enumeration | A list of discrete values, with sym­bolic names in the form of strings. | Number of elements, ele­ment names, element val­ues | Enumeration is a list of pairs (name, value). The name is a string; the value is an unsigned integer. |
| Reference | Reference to object or region within the HDF5 file |  | See the Reference API, H5R |
| Array | Array (1-4 dimen­sions) of data ele­ments | Number of dimensions, dimension sizes, base datatype | The array is accessed atomically: no selection or sub-setting. |
| Variable-length | A variable-length 1-dimensional array of data ele­ments | Current size, base type |  |
| Compound | A Datatype of a sequence of Data­types | Number of members, member names, member types, member offset, member class, member size, byte order |  |

### Predefined Datatypes

The HDF5 library predefines a modest number of commonly used datatypes. These types have standard symbolic names of the form H5T\_arch\_base where arch is an architecture name and base is a pro­gramming type name (Table 2). New types can be derived from the predefined types by copying the pre­defined type (see H5Tcopy()) and then modifying the result.

The base name of most types consists of a letter to indicate the class (Table 3), a precision in bits, and an indication of the byte order (Table 4).

Table 5 shows examples of predefined datatypes. The full list can be found in the “HDF5 Predefined Data­types” section of the HDF5 Reference Manual.

|  |  |
| --- | --- |
| Table 6-2. Architectures used in predefined datatypes | |
| **Architecture Name** | **Description** |
| IEEE | IEEE-754 standard floating point types in various byte orders. |
| STD | This is an architecture that contains semi-standard datatypes like signed two’s complement integers, unsigned integers, and bitfields in various byte orders. |
| C  FORTRAN | Types which are specific to the C or Fortran programming languages are defined in these architectures. For instance, H5T\_C\_S1 defines a base string type with null termination which can be used to derive string types of other lengths. |
| NATIVE | This architecture contains C-like datatypes for the machine on which the library was compiled. The types were actually defined by running the H5detect program when the library was compiled. In order to be portable, applications should almost always use this architecture to describe things in memory. |
| CRAY | Cray architectures. These are word-addressable, big-endian systems with non-IEEE floating point. |
| INTEL | All Intel and compatible CPU’s including 80286, 80386, 80486, Pen­tium, Pentium-Pro, and Pentium-II. These are little-endian systems with IEEE floating-point. |
| MIPS | All MIPS CPU’s commonly used in SGI systems. These are big-endian systems with IEEE floating-point. |
| ALPHA | All DEC Alpha CPU’s, little-endian systems with IEEE floating-point. |

|  |  |
| --- | --- |
| Table 6-3. Base types | |
| **B** | **Bitfield** |
| F | Floating point |
| I | Signed integer |
| R | References |
| S | Character string |
| U | Unsigned integer |

|  |  |
| --- | --- |
| Table 6-4. Byte order | |
| **BE** | **Big-endian** |
| LE | Little-endian |

|  |  |
| --- | --- |
| Table 6-5. Some predefined datatypes | |
| **Example** | **Description** |
| H5T\_IEEE\_F64LE | Eight-byte, little-endian, IEEE floating-point |
| H5T\_IEEE\_F32BE | Four-byte, big-endian, IEEE floating point |
| H5T\_STD\_I32LE | Four-byte, little-endian, signed two’s complement integer |
| H5T\_STD\_U16BE | Two-byte, big-endian, unsigned integer |
| H5T\_C\_S1 | One-byte, null-terminated string of eight-bit characters |
| H5T\_INTEL\_B64 | Eight-byte bit field on an Intel CPU |
| H5T\_CRAY\_F64 | Eight-byte Cray floating point |
| H5T\_STD\_ROBJ | Reference to an entire object in a file |

The HDF5 library predefines a set of NATIVE datatypes which are similar to C type names. The native types are set to be an alias for the appropriate HDF5 datatype for each platform. For example, H5T\_NA­TIVE\_INT corresponds to a C int type. On an Intel based PC, this type is the same as H5T\_STD\_I32LE, while on a MIPS system this would be equivalent to H5T\_STD\_I32BE. Table 6 shows examples of NATIVE types and corresponding C types for a common 32-bit workstation.

|  |  |
| --- | --- |
| Table 6-6. Native and 32-bit C datatypes | |
| **Example** | **Corresponding C Type** |
| H5T\_NATIVE\_CHAR | char |
| H5T\_NATIVE\_SCHAR | signed char |
| H5T\_NATIVE\_UCHAR | unsigned char |
| H5T\_NATIVE\_SHORT | short |
| H5T\_NATIVE\_USHORT | unsigned short |
| H5T\_NATIVE\_INT | int |
| H5T\_NATIVE\_UINT | unsigned |
| H5T\_NATIVE\_LONG | long |
| H5T\_NATIVE\_ULONG | unsigned long |
| H5T\_NATIVE\_LLONG | long long |
| H5T\_NATIVE\_ULLONG | unsigned long long |
| H5T\_NATIVE\_FLOAT | float |
| H5T\_NATIVE\_DOUBLE | double |
| H5T\_NATIVE\_LDOUBLE | long double |
| H5T\_NATIVE\_HSIZE | hsize\_t |
| H5T\_NATIVE\_HSSIZE | hssize\_t |
| H5T\_NATIVE\_HERR | herr\_t |
| H5T\_NATIVE\_HBOOL | hbool\_t |
| H5T\_NATIVE\_B8 | 8-bit unsigned integer or 8-bit buffer in memory |
| H5T\_NATIVE\_B16 | 16-bit unsigned integer or 16-bit buffer in memory |
| H5T\_NATIVE\_B32 | 32-bit unsigned integer or 32-bit buffer in memory |
| H5T\_NATIVE\_B64 | 64-bit unsigned integer or 64-bit buffer in memory |

## How Datatypes are Used

### The Datatype Object and the HDF5 Datatype API

The HDF5 library manages datatypes as objects. The HDF5 datatype API manipulates the datatype objects through C function calls. New datatypes can be created from scratch or copied from existing datatypes. When a datatype is no longer needed its resources should be released by calling H5Tclose().

The datatype object is used in several roles in the HDF5 data model and library. Essentially, a datatype is used whenever the format of data elements is needed. There are four major uses of datatypes in the HDF5 library: at dataset creation, during data transfers, when discovering the contents of a file, and for specify­ing user-defined datatypes. See the table below.

|  |  |
| --- | --- |
| Table 6-7. Datatype uses | |
| **Use** | **Description** |
| Dataset creation | The datatype of the data elements must be declared when the dataset is created. |
| Data transfer | The datatype (format) of the data elements must be defined for both the source and destination. |
| Discovery | The datatype of a dataset can be interrogated to retrieve a complete description of the storage layout. |
| Creating user-defined datatypes | Users can define their own datatypes by creating datatype objects and setting their properties. |

### Dataset Creation

All the data elements of a dataset have the same datatype. When a dataset is created, the datatype for the data elements must be specified. The datatype of a dataset can never be changed. The example below shows the use of a datatype to create a dataset called “/dset”. In this example, the dataset will be stored as 32-bit signed integers in big-endian order.

|  |
| --- |
| hid\_t dt;  dt = H5Tcopy(H5T\_STD\_I32BE);  dataset\_id = H5Dcreate(file\_id, “/dset”, dt, dataspace\_id,  H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT); |
| Code Example 6-1. Using a datatype to create a dataset |

### Data Transfer (Read and Write)

Probably the most common use of datatypes is to write or read data from a dataset or attribute. In these operations, each data element is transferred from the source to the destination (possibly rearranging the order of the elements). Since the source and destination do not need to be identical (in other words, one is disk and the other is memory), the transfer requires both the format of the source element and the des­tination element. Therefore, data transfers use two datatype objects, for the source and destination.

When data is written, the source is memory and the destination is disk (file). The memory datatype describes the format of the data element in the machine memory, and the file datatype describes the desired format of the data element on disk. Similarly, when reading, the source datatype describes the format of the data element on disk, and the destination datatype describes the format in memory.

In the most common cases, the file datatype is the datatype specified when the dataset was created, and the memory datatype should be the appropriate NATIVE type.

The examples below show samples of writing data to and reading data from a dataset. The data in mem­ory is declared C type ‘int’, and the datatype H5T\_NATIVE\_INT corresponds to this type. The datatype of the dataset should be of datatype class H5T\_INTEGER.

|  |
| --- |
| int  dset\_data[DATA\_SIZE];    status = H5Dwrite(dataset\_id, H5T\_NATIVE\_INT, H5S\_ALL, H5S\_ALL,  H5P\_DEFAULT, dset\_data); |
| Code Example 6-2. Writing to a dataset |

|  |
| --- |
| int dset\_data[DATA\_SIZE];    status = H5Dread(dataset\_id, H5T\_NATIVE\_INT, H5S\_ALL, H5S\_ALL,  H5P\_DEFAULT, dset\_data); |
| Code Example 6-3. Reading from a dataset |

### Discovery of Data Format

The HDF5 Library enables a program to determine the datatype class and properties for any datatype. In order to discover the storage format of data in a dataset, the datatype is obtained, and the properties are determined by queries to the datatype object. The example below shows code that analyzes the datatype for an integer and prints out a description of its storage properties (byte order, signed, size).

|  |
| --- |
| switch (H5Tget\_class(type)) {  case H5T\_INTEGER:  ord = H5Tget\_order(type);  sgn = H5Tget\_sign(type);  printf(“Integer ByteOrder= ”);  switch (ord) { |
| case H5T\_ORDER\_LE:  printf(“LE”);  break;  case H5T\_ORDER\_BE:  printf(“BE”);  break;  } |
| printf(“ Sign= ”);  switch (sgn) {  case H5T\_SGN\_NONE:  printf(“false”);  break;  case H5T\_SGN\_2:  printf(“true”);  break;  } |
| printf(“ Size= ”);  sz = H5Tget\_size(type);  printf(“%d”, sz);  printf(“\n”);  break; |
| Code Example 6-4. Discovering datatype properties |

### Creating and Using User-defined Datatypes

Most programs will primarily use the predefined datatypes described above, possibly in composite data­types such as compound or array datatypes. However, the HDF5 datatype model is extremely general; a user program can define a great variety of atomic datatypes (storage layouts). In particular, the datatype properties can define signed and unsigned integers of any size and byte order, and floating point numbers with different formats, size, and byte order. The HDF5 datatype API provides methods to set these proper­ties.

User-defined types can be used to define the layout of data in memory; examples might match some platform specific number format or application defined bit-field. The user-defined type can also describe data in the file such as an application-defined format. The user-defined types can be translated to and from standard types of the same class, as described above.

## Datatype (H5T) Function Summaries

Functions that can be used with datatypes (H5T functions) and property list functions that can be used with datatypes (H5P functions) are listed below.

|  |  |
| --- | --- |
| Function Listing 6-1. General datatype operations | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Tcreate  h5tcreate\_f | Creates a new datatype. |
| H5Topen  h5topen\_f | Opens a committed datatype. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Tcommit  h5tcommit\_f | Commits a transient datatype to a file. The datatype is now a committed datatype. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Tcommit\_anon  h5tcommit\_anon\_f | Commits a transient datatype to a file. The datatype is now a committed datatype, but it is not linked into the file structure. |
| H5Tcommitted  h5tcommitted\_f | Determines whether a datatype is a commit­ted or a transient type. |
| H5Tcopy  h5tcopy\_f | Copies an existing datatype. |
| H5Tequal  h5tequal\_f | Determines whether two datatype identifiers refer to the same datatype. |
| H5Tlock  (no Fortran subroutine) | Locks a datatype. |
| H5Tget\_class  h5tget\_class\_f | Returns the datatype class identifier. |
| H5Tget\_create\_plist  h5tget\_create\_plist\_f | Returns a copy of a datatype creation prop­erty list. |
| H5Tget\_size  h5tget\_size\_f | Returns the size of a datatype. |
| H5Tget\_super  h5tget\_super\_f | Returns the base datatype from which a data­type is derived. |
| H5Tget\_native\_type  h5tget\_native\_type\_f | Returns the native datatype of a specified datatype. |
| H5Tdetect\_class  (no Fortran subroutine) | Determines whether a datatype is of the given datatype class. |
| H5Tget\_order  h5tget\_order\_f | Returns the byte order of a datatype. |
| H5Tset\_order  h5tset\_order\_f | Sets the byte ordering of a datatype. |
| H5Tdecode  h5tdecode\_f | Decode a binary object description of data­type and return a new object identifier. |
| H5Tencode  h5tencode | Encode a datatype object description into a binary buffer. |
| H5Tclose  h5tclose\_f | Releases a datatype. |

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| Function Listing 6-2. Conversion functions | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Tconvert  h5tconvert\_f | Converts data between specified datatypes. |
| H5Tcompiler\_conv  h5tcompiler\_conv\_f | Check whether the library’s default conver­sion is hard conversion. |
| H5Tfind  (no Fortran subroutine) | Finds a conversion function. |
| H5Tregister  (no Fortran subroutine) | Registers a conversion function. |
| H5Tunregister  (no Fortran subroutine) | Removes a conversion function from all con­version paths. |

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| --- | --- |
| Function Listing 6-3. Atomic datatype properties | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Tset\_size  h5tset\_size\_f | Sets the total size for an atomic datatype. |
| H5Tget\_precision  h5tget\_precision\_f | Returns the precision of an atomic datatype. |
| H5Tset\_precision  h5tset\_precision\_f | Sets the precision of an atomic datatype. |
| H5Tget\_offset  h5tget\_offset\_f | Retrieves the bit offset of the first significant bit. |
| H5Tset\_offset  h5tset\_offset\_f | Sets the bit offset of the first significant bit. |
| H5Tget\_pad  h5tget\_pad\_f | Retrieves the padding type of the least and most-significant bit padding. |
| H5Tset\_pad  h5tset\_pad\_f | Sets the least and most-significant bits pad­ding types. |
| H5Tget\_sign  h5tget\_sign\_f | Retrieves the sign type for an integer type. |
| H5Tset\_sign  h5tset\_sign\_f | Sets the sign property for an integer type. |
| H5Tget\_fields  h5tget\_fields\_f | Retrieves floating point datatype bit field information. |
| H5Tset\_fields  h5tset\_fields\_f | Sets locations and sizes of floating point bit fields. |
| H5Tget\_ebias  h5tget\_ebias\_f | Retrieves the exponent bias of a floating-point type. |
| H5Tset\_ebias  h5tset\_ebias\_f | Sets the exponent bias of a floating-point type. |
| H5Tget\_norm  h5tget\_norm\_f | Retrieves mantissa normalization of a float­ing-point datatype. |
| H5Tset\_norm  h5tset\_norm\_f | Sets the mantissa normalization of a floating-point datatype. |
| H5Tget\_inpad  h5tget\_inpad\_f | Retrieves the internal padding type for unused bits in floating-point datatypes. |
| H5Tset\_inpad  h5tset\_inpad\_f | Fills unused internal floating point bits. |
| H5Tget\_cset  h5tget\_cset\_f | Retrieves the character set type of a string datatype. |
| H5Tset\_cset  h5tset\_cset\_f | Sets character set to be used. |
| H5Tget\_strpad  h5tget\_strpad\_f | Retrieves the storage mechanism for a string datatype. |
| H5Tset\_strpad  h5tset\_strpad\_f | Defines the storage mechanism for character strings. |

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| Function Listing 6-4. Enumeration datatypes | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Tenum\_create  h5tenum\_create\_f | Creates a new enumeration datatype. |
| H5Tenum\_insert  h5tenum\_insert\_f | Inserts a new enumeration datatype member. |
| H5Tenum\_nameof  h5tenum\_nameof\_f | Returns the symbol name corresponding to a specified member of an enumeration data­type. |
| H5Tenum\_valueof  h5tenum\_valueof\_f | Returns the value corresponding to a speci­fied member of an enumeration datatype. |
| H5Tget\_member\_value  h5tget\_member\_value\_f | Returns the value of an enumeration data­type member. |
| H5Tget\_nmembers  h5tget\_nmembers\_f | Retrieves the number of elements in a com­pound or enumeration datatype. |
| H5Tget\_member\_name  h5tget\_member\_name\_f | Retrieves the name of a compound or enu­meration datatype member. |
| H5Tget\_member\_index  (no Fortran subroutine) | Retrieves the index of a compound or enu­meration datatype member. |

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| Function Listing 6-5. Compound datatype properties | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Tget\_nmembers  h5tget\_nmembers\_f | Retrieves the number of elements in a com­pound or enumeration datatype. |
| H5Tget\_member\_class  h5tget\_member\_class\_f | Returns datatype class of compound datatype member. |
| H5Tget\_member\_name  h5tget\_member\_name\_f | Retrieves the name of a compound or enu­meration datatype member. |
| H5Tget\_member\_index  h5tget\_member\_index\_f | Retrieves the index of a compound or enu­meration datatype member. |
| H5Tget\_member\_offset  h5tget\_member\_offset\_f | Retrieves the offset of a field of a compound datatype. |
| H5Tget\_member\_type  h5tget\_member\_type\_f | Returns the datatype of the specified mem­ber. |
| H5Tinsert  h5tinsert\_f | Adds a new member to a compound data­type. |
| H5Tpack  h5tpack\_f | Recursively removes padding from within a compound datatype. |

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| Function Listing 6-6. Array datatypes | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Tarray\_create  h5tarray\_create\_f | Creates an array datatype object. The C func­tion is a macro: see “API Compatibility Macros in HDF5.” |
| H5Tget\_array\_ndims  h5tget\_array\_ndims\_f | Returns the rank of an array datatype. |
| H5Tget\_array\_dims  h5tget\_array\_dims\_f | Returns sizes of array dimensions and dimen­sion permutations. The C function is a macro: see “API Compatibility Macros in HDF5.” |

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| --- | --- |
| Function Listing 6-7. Variable-length datatypes | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Tvlen\_create  h5tvlen\_create\_f | Creates a new variable-length datatype. |
| H5Tis\_variable\_str  h5tis\_variable\_str\_f | Determines whether datatype is a variable-length string. |

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| --- | --- |
| Function Listing 6-8. Opaque datatypes | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Tset\_tag  h5tset\_tag\_f | Tags an opaque datatype. |
| H5Tget\_tag  h5tget\_tag\_f | Gets the tag associated with an opaque data­type. |

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| Function Listing 6-9. Conversions between datatype and text | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5LTtext\_to\_dtype  (no Fortran subroutine) | Creates a datatype from a text description. |
| H5LTdtype\_to\_text  (no Fortran subroutine) | Generates a text description of a datatype. |

|  |  |
| --- | --- |
| Function Listing 6-10. Datatype creation property list functions (H5P) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Pset\_char\_encoding  h5pset\_char\_encoding\_f | Sets the character encoding used to encode a string. Use to set ASCII or UTF-8 character encoding for object names. |
| H5Pget\_char\_encoding  h5pget\_char\_encoding\_f | Retrieves the character encoding used to cre­ate a string. |

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| Function Listing 6-11. Datatype access property list functions (H5P) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Pset\_type\_conv\_cb  (no Fortran subroutine) | Sets user-defined datatype conversion call­back function. |
| H5Pget\_type\_conv\_cb  (no Fortran subroutine) | Gets user-defined datatype conversion call­back function. |

## Programming Model for Datatypes

The HDF5 Library implements an object-oriented model of datatypes. HDF5 datatypes are organized as a logical set of base types, or datatype classes. The HDF5 Library manages datatypes as objects. The HDF5 datatype API manipulates the datatype objects through C function calls. The figure below shows the abstract view of the datatype object. The table below shows the methods (C functions) that operate on datatype objects. New datatypes can be created from scratch or copied from existing datatypes.

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| --- |
| Dtypes_fig5.JPG |
| Figure 6-5. The datatype object |

|  |  |
| --- | --- |
| Table 6-8. General operations on datatype objects | |
| **API Function** | **Description** |
| hid\_t H5Tcreate (H5T\_class\_t class, size\_t size) | Create a new datatype object of datatype class class. The following datatype classes are supported with this function:  •        H5T\_COMPOUND  •        H5T\_OPAQUE  •        H5T\_ENUM  Other datatypes are created with H5Tcopy(). |
| hid\_t H5Tcopy (hid\_t type) | Obtain a modifiable transient datatype which is a copy of type. If type is a data­set identifier then the type returned is a modifiable transient copy of the datatype of the specified dataset. |
| hid\_t H5Topen (hid\_t location, const char \*name, H5P\_DEFAULT) | Open a committed datatype. The commit­ted datatype returned by this function is read-only. |
| htri\_t H5Tequal (hid\_t type1, hid\_t type2) | Determines if two types are equal. |
| herr\_t H5Tclose (hid\_t type) | Releases resources associated with a data­type obtained from H5Tcopy, H5Topen, or H5Tcreate. It is illegal to close an immutable transient datatype (for exam­ple, predefined types). |
| herr\_t H5Tcommit (hid\_t location, const char \*name, hid\_t type, H5P\_DE­FAULT, H5P\_DEFAULT, H5P\_DEFAULT) | Commit a transient datatype (not immutable) to a file to become a commit­ted datatype. Committed datatypes can be shared. |
| htri\_t H5Tcommitted (hid\_t type) | Test whether the datatype is transient or committed (named). |
| herr\_t H5Tlock (hid\_t type) | Make a transient datatype immutable (read-only and not closable). Predefined types are locked. |

In order to use a datatype, the object must be created (H5Tcreate), or a reference obtained by cloning from an existing type (H5Tcopy), or opened (H5Topen). In addition, a reference to the datatype of a data­set or attribute can be obtained with H5Dget\_type or H5Aget\_type. For composite datatypes a refer­ence to the datatype for members or base types can be obtained (H5Tget\_member\_type, H5Tget\_super). When the datatype object is no longer needed, the reference is discarded with H5Tclose.

Two datatype objects can be tested to see if they are the same with H5Tequal. This function returns true if the two datatype references refer to the same datatype object. However, if two datatype objects define equivalent datatypes (the same datatype class and datatype properties), they will not be considered ‘equal’.

A datatype can be written to the file as a first class object (H5Tcommit). This is a committed datatype and can be used in the same way as any other datatype.

### Discovery of Datatype Properties

Any HDF5 datatype object can be queried to discover all of its datatype properties. For each datatype class, there are a set of API functions to retrieve the datatype properties for this class.

#### 6.5.1.1. Properties of Atomic Datatypes

Table 9 lists the functions to discover the properties of atomic datatypes. Table 10 lists the queries rele­vant to specific numeric types. Table 11 gives the properties for atomic string datatype, and Table 12 gives the property of the opaque datatype.

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| --- | --- |
| Table 6-9. Functions to discover properties of atomic datatypes | |
| **Functions** | **Description** |
| H5T\_class\_t H5Tget\_class (hid\_t type) | The datatype class: H5T\_INTEGER, H5T\_FLOAT, H5T\_STRING, H5T\_BIT­FIELD, H5T\_OPAQUE, H5T\_COMPOUND, H5T\_REFERENCE, H5T\_ENUM, H5T\_VLEN, H5T\_ARRAY |
| size\_t H5Tget\_size (hid\_t type) | The total size of the element in bytes, including padding which may appear on either side of the actual value. |
| H5T\_order\_t H5Tget\_order (hid\_t type) | The byte order describes how the bytes of the datatype are laid out in memory. If the lowest memory address contains the least significant byte of the datum then it is said to be little-endian or H5T\_ORDER\_LE. If the bytes are in the opposite order then they are said to be big-endian or H5T\_ORDER\_BE. |
| size\_t H5Tget\_precision (hid\_t type) | The precision property identifies the number of significant bits of a datatype and the offset property (defined below) identifies its location. Some datatypes occupy more bytes than what is needed to store the value. For instance, a short on a Cray is 32 significant bits in an eight-byte field. |
| int H5Tget\_offset (hid\_t type) | The offset property defines the bit loca­tion of the least significant bit of a bit field whose length is precision. |
| herr\_t H5Tget\_pad (hid\_t type, H5T\_pad\_t \*lsb, H5T\_pad\_t \*msb) | Padding is the bits of a data element which are not significant as defined by the precision and offset properties. Pad­ding in the low-numbered bits is lsb pad­ding and padding in the high-numbered bits is msb padding. Padding bits can be set to zero (H5T\_PAD\_ZERO) or one (H5T\_PAD\_ONE). |

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| --- | --- |
| Table 6-10. Functions to discover properties of atomic numeric datatypes | |
| **Functions** | **Description** |
| H5T\_sign\_t H5Tget\_sign (hid\_t type) | (INTEGER) Integer data can be signed two’s complement (H5T\_SGN\_2) or unsigned (H5T\_SGN\_NONE). |
| herr\_t H5Tget\_fields (hid\_t type, size\_t \*spos, size\_t \*epos, size\_t \*esize, size\_t \*mpos, size\_t \*msize) | (FLOAT) A floating-point data element has bit fields which are the exponent and mantissa as well as a mantissa sign bit. These properties define the location (bit position of least significant bit of the field) and size (in bits) of each field. The sign bit is always of length one and none of the fields are allowed to overlap. |
| size\_t H5Tget\_ebias (hid\_t type) | (FLOAT) The exponent is stored as a non-negative value which is ebias larger than the true exponent. |
| H5T\_norm\_t H5Tget\_norm (hid\_t type) | (FLOAT) This property describes the nor­malization method of the mantissa.  •        H5T\_NORM\_MSBSET: the mantissa is shifted left (if non-zero) until the first bit after the radix point is set and the exponent is adjusted accordingly. All bits of the man­tissa after the radix point are stored.  •        H5T\_NORM\_IMPLIED: the man­tissa is shifted left \ (if non-zero) until the first bit after the radix point is set and the exponent is adjusted accordingly. The first bit after the radix point is not stored since it’s always set.  •        H5T\_NORM\_NONE: the fractional part of the mantissa is stored without normalizing it. |
| H5T\_pad\_t H5Tget\_inpad (hid\_t type) | (FLOAT) If any internal bits (that is, bits between the sign bit, the mantissa field, and the exponent field but within the pre­cision field) are unused, then they will be filled according to the value of this prop­erty. The padding can be: H5T\_PAD\_­NONE, H5T\_PAD\_ZERO, or H5T\_PAD\_ONE. |

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| Table 6-11. Functions to discover properties of atomic string datatypes | |
| **Functions** | **Description** |
| H5T\_cset\_t H5Tget\_cset (hid\_t type) | Two character sets are currently sup­ported: ASCII (H5T\_CSET\_ASCII) and UTF-8 (H5T\_CSET\_UTF8). |
| H5T\_str\_t H5Tget\_strpad (hid\_t type) | The string datatype has a fixed length, but the string may be shorter than the length. This property defines the storage mecha­nism for the left over bytes. The options are: H5T\_STR\_NULLTERM, H5T\_STR\_NULLPAD, or H5T\_STR\_SPACEPAD. |

|  |  |
| --- | --- |
| Table 6-12. Functions to discover properties of atomic opaque datatypes | |
| **Functions** | **Description** |
| char \*H5Tget\_tag(hid\_t type\_id) | A user-defined string. |

#### 6.5.1.2. Properties of Composite Datatypes

The composite datatype classes can also be analyzed to discover their datatype properties and the data­types that are members or base types of the composite datatype. The member or base type can, in turn, be analyzed. The table below lists the functions that can access the datatype properties of the different composite datatypes.

|  |  |
| --- | --- |
| Table 6-13. Functions to discover properties of composite datatypes | |
| **Functions** | **Description** |
| int H5Tget\_nmembers(hid\_t type\_id) | (COMPOUND) The number of fields in the compound datatype. |
| H5T\_class\_t H5Tget\_member\_class (hid\_t cdtype\_id, unsigned member\_no) | (COMPOUND) The datatype class of com­pound datatype member member\_no. |
| char \* H5Tget\_member\_name (hid\_t type\_id, unsigned field\_idx) | (COMPOUND) The name of field field\_idx of a compound datatype. |
| size\_t H5Tget\_member\_offset (hid\_t type\_id, unsigned memb\_no) | (COMPOUND) The byte offset of the beginning of a field within a compound datatype. |
| hid\_t H5Tget\_member\_type (hid\_t type\_id, unsigned field\_idx) | (COMPOUND) The datatype of the speci­fied member. |
| int H5Tget\_array\_ndims (hid\_t adtype\_id) | (ARRAY) The number of dimensions (rank) of the array datatype object. |
| int H5Tget\_array\_dims (hid\_t adtype\_id, hsize\_t \*dims[]) | (ARRAY) The sizes of the dimensions and the dimension permutations of the array datatype object. |
| hid\_t H5Tget\_super(hid\_t type) | (ARRAY, VL, ENUM) The base datatype from which the datatype type is derived. |
| herr\_t H5Tenum\_nameof(hid\_t type void \*value, char \*name, size\_t size) | (ENUM) The symbol name that corre­sponds to the specified value of the enu­meration datatype. |
| herr\_t H5Tenum\_valueof(hid\_t type char \*name, void \*value) | (ENUM) The value that corresponds to the specified name of the enumeration datatype. |
| herr\_t H5Tget\_member\_value (hid\_t type unsigned memb\_no, void \*value) | (ENUM) The value of the enumeration datatype member memb\_no. |

### Definition of Datatypes

The HDF5 library enables user programs to create and modify datatypes. The essential steps are:

1. Create a new datatype object of a specific composite datatype class, or copy an

existing atomic datatype object

2.      Set properties of the datatype object

3.      Use the datatype object

4.      Close the datatype object

To create a user-defined atomic datatype, the procedure is to clone a predefined datatype of the appropri­ate datatype class (H5Tcopy), and then set the datatype properties appropriate to the datatype class. The table below shows how to create a datatype to describe a 1024-bit unsigned integer.

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| --- |
| hid\_t new\_type = H5Tcopy (H5T\_NATIVE\_INT);  H5Tset\_precision(new\_type, 1024);  H5Tset\_sign(new\_type, H5T\_SGN\_NONE); |
| Code Example 6-5. Create a new datatype |

Composite datatypes are created with a specific API call for each datatype class. The table below shows the creation method for each datatype class. A newly created datatype cannot be used until the datatype properties are set. For example, a newly created compound datatype has no members and cannot be used.

|  |  |
| --- | --- |
| Table 6-14. Functions to create each datatype class | |
| **Datatype Class** | **Function to Create** |
| COMPOUND | H5Tcreate |
| OPAQUE | H5Tcreate |
| ENUM | H5Tenum\_create |
| ARRAY | H5Tarray\_create |
| VL | H5Tvlen\_create |

Once the datatype is created and the datatype properties set, the datatype object can be used.

Predefined datatypes are defined by the library during initialization using the same mechanisms as described here. Each predefined datatype is locked (H5Tlock), so that it cannot be changed or destroyed. User-defined datatypes may also be locked using H5Tlock.

#### 6.5.2.1. User-defined Atomic Datatypes

Table 15 summarizes the API methods that set properties of atomic types. Table 16 shows properties spe­cific to numeric types, Table 17 shows properties specific to the string datatype class. Note that offset, pad, etc. do not apply to strings. Table 18 shows the specific property of the OPAQUE datatype class.

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| --- | --- |
| Table 6-15. API methods that set properties of atomic datatypes | |
| **Functions** | **Description** |
| herr\_t H5Tset\_size (hid\_t type, size\_t size) | Set the total size of the element in bytes. This includes padding which may appear on either side of the actual value. If this property is reset to a smaller value which would cause the significant part of the data to extend beyond the edge of the datatype, then the offset property is dec­remented a bit at a time. If the offset reaches zero and the significant part of the data still extends beyond the edge of the datatype then the precision property is decremented a bit at a time. Decreasing the size of a datatype may fail if the H5T\_FLOAT bit fields would extend beyond the significant part of the type. |
| herr\_t H5Tset\_order (hid\_t type, H5T\_order\_t order) | Set the byte order to little-endian (H5T\_ORDER\_LE) or big-endian (H5T\_ORDER\_BE). |
| herr\_t H5Tset\_precision (hid\_t type, size\_t precision) | Set the number of significant bits of a datatype. The offset property (defined below) identifies its location. The size property defined above represents the entire size (in bytes) of the datatype. If the precision is decreased then padding bits are inserted on the MSB side of the signif­icant bits (this will fail for H5T\_FLOAT types if it results in the sign, mantissa, or exponent bit field extending beyond the edge of the significant bit field). On the other hand, if the precision is increased so that it “hangs over” the edge of the total size then the offset property is decre­mented a bit at a time. If the offset reaches zero and the significant bits still hang over the edge, then the total size is increased a byte at a time. |
| herr\_t H5Tset\_offset (hid\_t type, size\_t offset) | Set the bit location of the least significant bit of a bit field whose length is preci­sion. The bits of the entire data are num­bered beginning at zero at the least significant bit of the least significant byte (the byte at the lowest memory address for a little-endian type or the byte at the highest address for a big-endian type). The offset property defines the bit loca­tion of the least significant bit of a bit field whose length is precision. If the offset is increased so the significant bits “hang over” the edge of the datum, then the size property is automatically incre­mented. |
| herr\_t H5Tset\_pad (hid\_t type, H5T\_pad\_t lsb, H5T\_pad\_t msb) | Set the padding to zeros (H5T\_PAD\_ZERO) or ones (H5T\_PAD\_ONE). Padding is the bits of a data element which are not sig­nificant as defined by the precision and offset properties. Padding in the low-numbered bits is lsb padding and pad­ding in the high-numbered bits is msb padding. |

|  |  |
| --- | --- |
| Table 6-16. API methods that set properties of numeric datatypes | |
| **Functions** | **Description** |
| herr\_t H5Tset\_sign (hid\_t type, H5T\_sign\_t sign) | (INTEGER) Integer data can be signed two’s complement (H5T\_SGN\_2) or unsigned (H5T\_SGN\_NONE). |
| herr\_t H5Tset\_fields (hid\_t type, size\_t spos, size\_t epos, size\_t esize, size\_t mpos, size\_t msize) | (FLOAT) Set the properties define the location (bit position of least significant bit of the field) and size (in bits) of each field. The sign bit is always of length one and none of the fields are allowed to overlap. |
| herr\_t H5Tset\_ebias (hid\_t type, size\_t ebias) | (FLOAT) The exponent is stored as a non-negative value which is ebias larger than the true exponent. |
| herr\_t H5Tset\_norm (hid\_t type, H5T\_norm\_t norm) | (FLOAT) This property describes the nor­malization method of the mantissa.  •        H5T\_NORM\_MSBSET: the mantissa is shifted left (if non-zero) until the first bit after the radix point is set and the exponent is adjusted accordingly. All bits of the man­tissa after the radix point are stored.  •        H5T\_NORM\_IMPLIED: the man­tissa is shifted left (if non-zero) until the first bit after the radix point is set and the exponent is adjusted accordingly. The first bit after the radix point is not stored since it is always set.  •        H5T\_NORM\_NONE: the fractional part of the mantissa is stored without normalizing it. |
| herr\_t H5Tset\_inpad (hid\_t type, H5T\_pad\_t inpad) | (FLOAT) If any internal bits (that is, bits between the sign bit, the mantissa field, and the exponent field but within the pre­cision field) are unused, then they will be filled according to the value of this prop­erty. The padding can be: H5T\_PAD\_­NONE, H5T\_PAD\_ZERO or H5T\_PAD\_ONE. |

|  |  |
| --- | --- |
| Table 6-17. API methods that set properties of string datatypes | |
| **Functions** | **Description** |
| herr\_t H5Tset\_size (hid\_t type, size\_t size) | Set the length of the string, in bytes. The precision is automatically set to 8\*size. |
| herr\_t H5Tset\_precision (hid\_t type, size\_t precision) | The precision must be a multiple of 8. |
| herr\_t H5Tset\_cset (hid\_t type\_id, H5T\_cset\_t cset) | Two character sets are currently sup­ported: ASCII (H5T\_CSET\_ASCII) and UTF-8 (H5T\_CSET\_UTF8). |
| herr\_t H5Tset\_strpad (hid\_t type\_id, H5T\_str\_t strpad) | The string datatype has a fixed length, but the string may be shorter than the length. This property defines the storage mecha­nism for the left over bytes. The method used to store character strings differs with the programming language:   * C usually null terminates strings * Fortran left-justifies and space-pads strings   Valid string padding values, as passed in the parameter strpad, are as follows:   * H5T\_STR\_NULLTERM: Null termi­nate (as C does) * H5T\_STR\_NULLPAD: Pad with zeros * H5T\_STR\_SPACEPAD: Pad with spaces (as FORTRAN does) |

|  |  |
| --- | --- |
| Table 6-18. API methods that set properties of opaque datatypes | |
| **Functions** | **Description** |
| herr\_t H5Tset\_tag (hid\_t type\_id const char \*tag) | Tags the opaque datatype type\_id with an ASCII identifier tag. |

**Examples**

The example below shows how to create a 128-bit little-endian signed integer type. Increasing the preci­sion of a type automatically increases the total size. Note that the proper procedure is to begin from a type of the intended datatype class which in this case is a NATIVE INT.

|  |
| --- |
| hid\_t new\_type = H5Tcopy (H5T\_NATIVE\_INT);  H5Tset\_precision (new\_type, 128);  H5Tset\_order (new\_type, H5T\_ORDER\_LE); |
| Code Example 6-6. Create a new 128-bit little-endian signed integer datatype |

The figure below shows the storage layout as the type is defined. The H5Tcopy creates a datatype that is the same as H5T\_NATIVE\_INT. In this example, suppose this is a 32-bit big-endian number (Figure a). The precision is set to 128 bits, which automatically extends the size to 8 bytes (Figure b). Finally, the byte order is set to little-endian (Figure c).

|  |
| --- |
| Dtypes_fig6.JPG |
| Figure 6-6. The storage layout for a new 128-bit little-endian signed integer datatype |

The significant bits of a data element can be offset from the beginning of the memory for that element by an amount of padding. The offset property specifies the number of bits of padding that appear to the “right of” the value. The table and figure below show how a 32-bit unsigned integer with 16-bits of preci­sion having the value 0x1122 will be laid out in memory.

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| Table 6-19. Memory Layout for a 32-bit unsigned integer | | | | |
| **Byte Position** | **Big-Endian Offset=0** | **Big-Endian Offset=16** | **Little-Endian Offset=0** | **Little-Endian Offset=16** |
| 0: | [pad] | [0x11] | [0x22] | [pad] |
| 1: | [pad] | [0x22] | [0x11] | [pad] |
| 2: | [0x11] | [pad] | [pad] | [0x22] |
| 3: | [0x22] | [pad] | [pad] | [0x11] |

|  |
| --- |
| Dtypes_fig7.JPG |
| Figure 6-7. Memory Layout for a 32-bit unsigned integer |

If the offset is incremented then the total size is incremented also if necessary to prevent significant bits of the value from hanging over the edge of the datatype.

The bits of the entire data are numbered beginning at zero at the least significant bit of the least signifi­cant byte (the byte at the lowest memory address for a little-endian type or the byte at the highest address for a big-endian type). The offset property defines the bit location of the least significant bit of a bit field whose length is precision. If the offset is increased so the significant bits “hang over” the edge of the datum, then the size property is automatically incremented.

To illustrate the properties of the integer datatype class, the example below shows how to create a user-defined datatype that describes a 24-bit signed integer that starts on the third bit of a 32-bit word. The datatype is specialized from a 32-bit integer, the precision is set to 24 bits, and the offset is set to 3.

|  |
| --- |
| hid\_t dt;    dt = H5Tcopy(H5T\_SDT\_I32LE);    H5Tset\_precision(dt, 24);  H5Tset\_offset(dt,3);  H5Tset\_pad(dt, H5T\_PAD\_ZERO, H5T\_PAD\_ONE); |
| Code Example 6-7. A user-defined datatype with a 24-bit signed integer |

The figure below shows the storage layout for a data element. Note that the unused bits in the offset will be set to zero and the unused bits at the end will be set to one, as specified in the H5Tset\_pad call.

|  |
| --- |
| Dtypes_fig8.JPG |
| Figure 6-8. A user-defined integer datatype with a range of -1,048,583 to 1,048,584 |

To illustrate a user-defined floating point number, the example below shows how to create a 24-bit float­ing point number that starts 5 bits into a 4 byte word. The floating point number is defined to have a man­tissa of 19 bits (bits 5-23), an exponent of 3 bits (25-27), and the sign bit is bit 28. (Note that this is an illustration of what can be done and is not necessarily a floating point format that a user would require.)

|  |
| --- |
| hid\_t dt;    dt = H5Tcopy(H5T\_IEEE\_F32LE);    H5Tset\_precision(dt, 24);  H5Tset\_fields (dt, 28, 25, 3, 5, 19);  H5Tset\_pad(dt, H5T\_PAD\_ZERO, H5T\_PAD\_ONE);  H5Tset\_inpad(dt, H5T\_PAD\_ZERO); |
| Code Example 6-8. A user-defined 24-bit floating point datatype |

|  |
| --- |
| Dtypes_fig9.JPG |
| Figure 6-9. A user-defined floating point datatype |

The figure above shows the storage layout of a data element for this datatype. Note that there is an unused bit (24) between the mantissa and the exponent. This bit is filled with the inpad value which in this case is 0.

The sign bit is always of length one and none of the fields are allowed to overlap. When expanding a float­ing-point type one should set the precision first; when decreasing the size one should set the field posi­tions and sizes first.

#### 6.5.2.2. Composite Datatypes

All composite datatypes must be user-defined; there are no predefined composite datatypes.

#### 6.5.2.2.1. Compound Datatypes

The subsections below describe how to create a compound datatype and how to write and read data of a compound datatype.

**Defining Compound Datatypes**

Compound datatypes are conceptually similar to a C struct or Fortran derived types. The compound data­type defines a contiguous sequence of bytes, which are formatted using one up to 2^16 datatypes (mem­bers). A compound datatype may have any number of members, in any order, and the members may have any datatype, including compound. Thus, complex nested compound datatypes can be created. The total size of the compound datatype is greater than or equal to the sum of the size of its members, up to a max­imum of 2^32 bytes. HDF5 does not support datatypes with distinguished records or the equivalent of C unions or Fortran EQUIVALENCE statements.

Usually a C struct or Fortran derived type will be defined to hold a data point in memory, and the offsets of the members in memory will be the offsets of the struct members from the beginning of an instance of the struct. The HDF5 C library provides a macro HOFFSET (s,m) to calculate the member’s offset. The HDF5 Fortran applications have to calculate offsets by using sizes of members datatypes and by taking in consideration the order of members in the Fortran derived type.

HOFFSET(s,m)

This macro computes the offset of member m within a struct s

offsetof(s,m)

This macro defined in stddef.h does exactly the same thing as the HOFFSET() macro.

Note for Fortran users: Offsets of Fortran structure members correspond to the offsets within a packed datatype (see explanation below) stored in an HDF5 file.

Each member of a compound datatype must have a descriptive name which is the key used to uniquely identify the member within the compound datatype. A member name in an HDF5 datatype does not nec­essarily have to be the same as the name of the member in the C struct or Fortran derived type, although this is often the case. Nor does one need to define all members of the C struct or Fortran derived type in the HDF5 compound datatype (or vice versa).

Unlike atomic datatypes which are derived from other atomic datatypes, compound datatypes are created from scratch. First, one creates an empty compound datatype and specifies its total size. Then members are added to the compound datatype in any order. Each member type is inserted at a designated offset. Each member has a name which is the key used to uniquely identify the member within the compound datatype.

The example below shows a way of creating an HDF5 C compound datatype to describe a complex num­ber. This is a structure with two components, “real” and “imaginary”, and each component is a double. An equivalent C struct whose type is defined by the complex\_t struct is shown.

|  |
| --- |
| typedef struct {     double re;   /\*real part\*/     double im;   /\*imaginary part\*/  } complex\_t;    hid\_t complex\_id = H5Tcreate (H5T\_COMPOUND, sizeof (complex\_t));  H5Tinsert (complex\_id, “real”, HOFFSET(complex\_t,re),  H5T\_NATIVE\_DOUBLE);  H5Tinsert (complex\_id, “imaginary”, HOFFSET(complex\_t,im),  H5T\_NATIVE\_DOUBLE); |
| Code Example 6-9. A compound datatype for complex numbers in C |

The example below shows a way of creating an HDF5 Fortran compound datatype to describe a complex number. This is a Fortran derived type with two components, “real” and “imaginary”, and each component is DOUBLE PRECISION. An equivalent Fortran TYPE whose type is defined by the TYPE complex\_t is shown.

|  |
| --- |
| TYPE complex\_t     DOUBLE PRECISION re   ! real part     DOUBLE PRECISION im;  ! imaginary part  END TYPE complex\_t    CALL h5tget\_size\_f(H5T\_NATIVE\_DOUBLE, re\_size, error)  CALL h5tget\_size\_f(H5T\_NATIVE\_DOUBLE, im\_size, error)  complex\_t\_size = re\_size + im\_size  CALL h5tcreate\_f(H5T\_COMPOUND\_F, complex\_t\_size, type\_id)  offset = 0  CALL h5tinsert\_f(type\_id, “real”, offset, H5T\_NATIVE\_DOUBLE, error)  offset = offset + re\_size  CALL h5tinsert\_f(type\_id, “imaginary”, offset, H5T\_NATIVE\_DOUBLE, error) |
| Code Example 6-10. A compound datatype for complex numbers in Fortran |

Important Note: The compound datatype is created with a size sufficient to hold all its members. In the C example above, the size of the C struct and the HOFFSET macro are used as a convenient mechanism to determine the appropriate size and offset. Alternatively, the size and offset could be manually deter­mined: the size can be set to 16 with “real” at offset 0 and “imaginary” at offset 8. However, different plat­forms and compilers have different sizes for “double” and may have alignment restrictions which require additional padding within the structure. It is much more portable to use the HOFFSET macro which assures that the values will be correct for any platform.

The figure below shows how the compound datatype would be laid out assuming that NATIVE\_DOUBLE are 64-bit numbers and that there are no alignment requirements. The total size of the compound data­type will be 16 bytes, the “real” component will start at byte 0, and “imaginary” will start at byte 8.

|  |
| --- |
| Dtypes_fig10.JPG |
| Figure 6-10. Layout of a compound datatype |

 The members of a compound datatype may be any HDF5 datatype including the compound, array, and variable-length (VL) types. The figure and example below show the memory layout and code which cre­ates a compound datatype composed of two complex values, and each complex value is also a compound datatype as in the figure above.

|  |
| --- |
| Dtypes_fig11.JPG |
| Figure 6-11. Layout of a compound datatype nested in a compound datatype |

|  |
| --- |
| typedef struct {     complex\_t x;     complex\_t y;  } surf\_t;    hid\_t complex\_id, surf\_id; /\*hdf5 datatypes\*/ |
| complex\_id = H5Tcreate (H5T\_COMPOUND, sizeof(complex\_t));  H5Tinsert (complex\_id, “re”, HOFFSET(complex\_t,re),  H5T\_NATIVE\_DOUBLE);  H5Tinsert (complex\_id, “im”, HOFFSET(complex\_t,im),  H5T\_NATIVE\_DOUBLE); |
| surf\_id = H5Tcreate (H5T\_COMPOUND, sizeof(surf\_t));  H5Tinsert (surf\_id, “x”, HOFFSET(surf\_t,x), complex\_id);  H5Tinsert (surf\_id, “y”, HOFFSET(surf\_t,y), complex\_id); |
| Code Example 6-11. Code for a compound datatype nested in a compound datatype |

 Note that a similar result could be accomplished by creating a compound datatype and inserting four fields. See the figure below. This results in the same layout as the figure above. The difference would be how the fields are addressed. In the first case, the real part of ‘y’ is called ‘y.re’; in the second case it is ‘y-re’.

|  |
| --- |
| typedef struct {     complex\_t x;     complex\_t y;  } surf\_t; |
| hid\_t surf\_id = H5Tcreate (H5T\_COMPOUND, sizeof(surf\_t));  H5Tinsert (surf\_id, “x-re”, HOFFSET(surf\_t,x.re),H5T\_NATIVE\_DOUBLE);  H5Tinsert (surf\_id, “x-im”, HOFFSET(surf\_t,x.im),H5T\_NATIVE\_DOUBLE);  H5Tinsert (surf\_id, “y-re”, HOFFSET(surf\_t,y.re),H5T\_NATIVE\_DOUBLE);  H5Tinsert (surf\_id, “y-im”, HOFFSET(surf\_t,y.im),H5T\_NATIVE\_DOUBLE); |
| Code Example 6-12. Another compound datatype nested in a compound datatype |

 The members of a compound datatype do not always fill all the bytes. The HOFFSET macro assures that the members will be laid out according to the requirements of the platform and language. The example below shows an example of a C struct which requires extra bytes of padding on many platforms. The sec­ond element, ‘b’, is a 1-byte character followed by an 8 byte double, ‘c’. On many systems, the 8-byte value must be stored on a 4- or 8-byte boundary. This requires the struct to be larger than the sum of the size of its elements.

In the example below, sizeof and HOFFSET are used to assure that the members are inserted at the cor­rect offset to match the memory conventions of the platform. The figure below shows how this data ele­ment would be stored in memory, assuming the double must start on a 4-byte boundary. Notice the extra bytes between ‘b’ and ‘c’.

|  |
| --- |
| typedef struct s1\_t {     int    a;     char   b;     double c;  } s1\_t;    s1\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(s1\_t));  H5Tinsert(s1\_tid, “a\_name”, HOFFSET(s1\_t, a), H5T\_NATIVE\_INT);  H5Tinsert(s1\_tid, “b\_name”, HOFFSET(s1\_t, b), H5T\_NATIVE\_CHAR);  H5Tinsert(s1\_tid, “c\_name”, HOFFSET(s1\_t, c), H5T\_NATIVE\_DOUBLE); |
| Code Example 6-13. A compound datatype that requires padding |

|  |
| --- |
| Dtypes_fig23.JPG |
| Figure 6-12. Memory layout of a compound datatype that requires padding |

However, data stored on disk does not require alignment, so unaligned versions of compound data struc­tures can be created to improve space efficiency on disk. These unaligned compound datatypes can be created by computing offsets by hand to eliminate inter-member padding, or the members can be packed by calling H5Tpack (which modifies a datatype directly, so it is usually preceded by a call to H5Tcopy).

The example below shows how to create a disk version of the compound datatype from the figure above in order to store data on disk in as compact a form as possible. Packed compound datatypes should gener­ally not be used to describe memory as they may violate alignment constraints for the architecture being used. Note also that using a packed datatype for disk storage may involve a higher data conversion cost.

|  |
| --- |
| hid\_t s2\_tid = H5Tcopy (s1\_tid);  H5Tpack (s2\_tid); |
| Code Example 6-14. Create a packed compound datatype in C |

 The example below shows the sequence of Fortran calls to create a packed compound datatype. An HDF5 Fortran compound datatype never describes a compound datatype in memory and compound data is ALWAYS written by fields as described in the next section. Therefore packing is not needed unless the off­set of each consecutive member is not equal to the sum of the sizes of the previous members.

|  |
| --- |
| CALL h5tcopy\_f(s1\_id, s2\_id, error)  CALL h5tpack\_f(s2\_id, error) |
| Code Example 6-15. Create a packed compound datatype in Fortran |

**Creating and Writing Datasets with Compound Datatypes**

Creating datasets with compound datatypes is similar to creating datasets with any other HDF5 datatypes. But writing and reading may be different since datasets that have compound datatypes can be written or read by a field (member) or subsets of fields (members). The compound datatype is the only composite datatype that supports “sub-setting” by the elements the datatype is built from.

The example below shows a C example of creating and writing a dataset with a compound datatype.

|  |
| --- |
| typedef struct s1\_t {     int a;     float b;     double c;  } s1\_t; |
| s1\_t data[LENGTH];    /\* Initialize data \*/  for (i = 0; i < LENGTH; i++) {     data[i].a = i;     data[i].b = i\*i;     data[i].c = 1./(i+1);  ... |
| s1\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(s1\_t));  H5Tinsert(s1\_tid, “a\_name”, HOFFSET(s1\_t, a),  H5T\_NATIVE\_INT);  H5Tinsert(s1\_tid, “b\_name”, HOFFSET(s1\_t, b),  H5T\_NATIVE\_FLOAT);  H5Tinsert(s1\_tid, “c\_name”, HOFFSET(s1\_t, c),  H5T\_NATIVE\_DOUBLE);  ... |
| dataset\_id = H5Dcreate(file\_id, “SDScompound.h5”, s1\_t,  space\_id, H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT);  H5Dwrite (dataset\_id, s1\_tid, H5S\_ALL, H5S\_ALL,  H5P\_DEFAULT, data); |
| Code Example 6-16. Create and write a dataset with a compound datatype in C |

 The example below shows the content of the file written on a little-endian machine.

|  |
| --- |
| HDF5 “SDScompound.h5” {  GROUP “/” {     DATASET “ArrayOfStructures” {        DATATYPE  H5T\_COMPOUND {           H5T\_STD\_I32LE “a\_name”;           H5T\_IEEE\_F32LE “b\_name”;           H5T\_IEEE\_F64LE “c\_name”; |
| }        DATASPACE  SIMPLE { ( 3 ) / ( 3 ) }        DATA {        (0): {                 0,                 0,                 1           }, |
| (1): {                 1,                 1,                 0.5           }, |
| (2): {                 2,                 4,                 0.333333           }        }     }  }  } |
| Code Example 6-17. Create and write a little-endian dataset with a compound datatype in C |

It is not necessary to write the whole data at once. Datasets with compound datatypes can be written by field or by subsets of fields. In order to do this one has to remember to set the transfer property of the dataset using the H5Pset\_preserve call and to define the memory datatype that corresponds to a field. The example below shows how float and double fields are written to the dataset.

|  |
| --- |
| typedef struct sb\_t {     float b;     double c;  } sb\_t; |
| typedef struct sc\_t {     float b;     double c;  } sc\_t;  sb\_t data1[LENGTH];  sc\_t data2[LENGTH]; |
| /\* Initialize data \*/  for (i = 0; i < LENGTH; i++) {        data1.b = i\*i;        data2.c = 1./(i+1);  }  ...  /\* Create dataset as in example 15 \*/  ... |
| /\* Create memory datatypes corresponding to float \*/  /\* and double datatype fields \*/    sb\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(sb\_t));  H5Tinsert(sb\_tid, “b\_name”, HOFFSET(sb\_t, b), H5T\_NATIVE\_FLOAT);  sc\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(sc\_t));  H5Tinsert(sc\_tid, “c\_name”, HOFFSET(sc\_t, c), H5T\_NATIVE\_DOUBLE);  ... |
| /\* Set transfer property \*/  xfer\_id = H5Pcreate(H5P\_DATASET\_XFER);  H5Pset\_preserve(xfer\_id, 1);  H5Dwrite (dataset\_id, sb\_tid, H5S\_ALL, H5S\_ALL, xfer\_id, data1);  H5Dwrite (dataset\_id, sc\_tid, H5S\_ALL, H5S\_ALL, xfer\_id, data2); |
| Code Example 6-18. Writing floats and doubles to a dataset |

The figure below shows the content of the file written on a little-endian machine. Only float and double fields are written. The default fill value is used to initialize the unwritten integer field.

|  |
| --- |
| HDF5 “SDScompound.h5” {  GROUP “/” {     DATASET “ArrayOfStructures” {        DATATYPE  H5T\_COMPOUND {           H5T\_STD\_I32LE “a\_name”;           H5T\_IEEE\_F32LE “b\_name”;           H5T\_IEEE\_F64LE “c\_name”;        } |
| DATASPACE  SIMPLE { ( 3 ) / ( 3 ) }        DATA {        (0): {                 0,                 0,                 1              }, |
| (1): {                 0,                 1,                 0.5              }, |
| (2): {                 0,                 4,                 0.333333 |
| }           }        }  }  } |
| Code Example 6-19. Writing floats and doubles to a dataset on a little-endian system |

The example below contains a Fortran example that creates and writes a dataset with a compound data­type. As this example illustrates, writing and reading compound datatypes in Fortran is always done by fields. The content of the written file is the same as shown in the example above.

|  |
| --- |
| ! One cannot write an array of a derived datatype in  ! Fortran.  TYPE s1\_t     INTEGER a     REAL b     DOUBLE PRECISION c  END TYPE s1\_t |
| TYPE(s1\_t) d(LENGTH)  ! Therefore, the following code initializes an array  ! corresponding to each field in the derived datatype  ! and writes those arrays to the dataset    INTEGER, DIMENSION(LENGTH) :: a  REAL, DIMENSION(LENGTH) :: b  DOUBLE PRECISION, DIMENSION(LENGTH) :: c |
| ! Initialize data     do i = 1, LENGTH        a(i) = i-1        b(i) = (i-1) \* (i-1)        c(i) = 1./i     enddo    ... |
| ! Set dataset transfer property to preserve partially  ! initialized fields during write/read to/from dataset  ! with compound datatype.  !  CALL h5pcreate\_f(H5P\_DATASET\_XFER\_F, plist\_id, error)  CALL h5pset\_preserve\_f(plist\_id, .TRUE., error)  ...  ! |
| ! Create compound datatype.  !  ! First calculate total size by calculating sizes of  ! each member  !  CALL h5tget\_size\_f(H5T\_NATIVE\_INTEGER, type\_sizei, error)  CALL h5tget\_size\_f(H5T\_NATIVE\_REAL, type\_sizer, error)  CALL h5tget\_size\_f(H5T\_NATIVE\_DOUBLE, type\_sized, error)  type\_size = type\_sizei + type\_sizer + type\_sized  CALL h5tcreate\_f(H5T\_COMPOUND\_F, type\_size, dtype\_id, error) |
| !  ! Insert members  !  ! |
| ! INTEGER member  !  offset = 0  CALL h5tinsert\_f(dtype\_id, “a\_name”, offset,  H5T\_NATIVE\_INTEGER, error)  !  ! REAL member  !  offset = offset + type\_sizei  CALL h5tinsert\_f(dtype\_id, “b\_name”, offset, H5T\_NATIVE\_REAL,  error)  ! |
| ! DOUBLE PRECISION member  !  offset = offset + type\_sizer  CALL h5tinsert\_f(dtype\_id, “c\_name”, offset,  H5T\_NATIVE\_DOUBLE, error)    !  ! Create the dataset with compound datatype.  ! |
| CALL h5dcreate\_f(file\_id, dsetname, dtype\_id, dspace\_id, &  dset\_id, error, H5P\_DEFAULT\_F, H5P\_DEFAULT\_F,  H5P\_DEFAULT\_F)  !  ...  ! Create memory types. We have to create a compound  ! datatype for each member we want to write.  ! |
| CALL h5tcreate\_f(H5T\_COMPOUND\_F, type\_sizei, dt1\_id, error)  offset = 0  CALL h5tinsert\_f(dt1\_id, “a\_name”, offset,  H5T\_NATIVE\_INTEGER, error)  !  CALL h5tcreate\_f(H5T\_COMPOUND\_F, type\_sizer, dt2\_id, error)  offset = 0 |
| CALL h5tinsert\_f(dt2\_id, “b\_name”, offset, H5T\_NATIVE\_REAL,  error)  !  CALL h5tcreate\_f(H5T\_COMPOUND\_F, type\_sized, dt3\_id, error)  offset = 0  CALL h5tinsert\_f(dt3\_id, “c\_name”, offset, H5T\_NATIVE\_DOUBLE,  error) |
| !  ! Write data by fields in the datatype. Fields order  ! is not important.  !  CALL h5dwrite\_f(dset\_id, dt3\_id, c, data\_dims, error,  xfer\_prp = plist\_id)  CALL h5dwrite\_f(dset\_id, dt2\_id, b, data\_dims, error,  xfer\_prp = plist\_id)  CALL h5dwrite\_f(dset\_id, dt1\_id, a, data\_dims, error,  xfer\_prp = plist\_id) |
| Code Example 6-20. Create and write a dataset with a compound datatype in Fortran |

**Reading Datasets with Compound Datatypes**

Reading datasets with compound datatypes may be a challenge. For general applications there is no way to know a priori the corresponding C structure. Also, C structures cannot be allocated on the fly during dis­covery of the dataset’s datatype. For general C, C++, Fortran and Java application the following steps will be required to read and to interpret data from the dataset with compound datatype:

1.      Get the identifier of the compound datatype in the file with the H5Dget\_type call

2.      Find the number of the compound datatype members with the H5Tget\_nmembers call

3.      Iterate through compound datatype members

* Get member class with the H5Tget\_member\_class call
* Get member name with the H5Tget\_member\_name call
* Check class type against predefined classes
  + H5T\_INTEGER
  + H5T\_FLOAT
  + H5T\_STRING
  + H5T\_BITFIELD
  + H5T\_OPAQUE
  + H5T\_COMPOUND
  + H5T\_REFERENCE
  + H5T\_ENUM
  + H5T\_VLEN
  + H5T\_ARRAY
* If class is H5T\_COMPOUND, then go to step 2 and repeat all steps under step 3. If class is not H5T\_COMPOUND, then a member is of an atomic class and can be read to a corresponding buf­fer after discovering all necessary information specific to each atomic type (for example, size of the integer or floats, super class for enumerated and array datatype, and its sizes)

The examples below show how to read a dataset with a known compound datatype.

The first example below shows the steps needed to read data of a known structure. First, build a memory datatype the same way it was built when the dataset was created, and then second use the datatype in an H5Dread call.

|  |
| --- |
| typedef struct s1\_t {     int a;     float b;     double c;  } s1\_t; |
| s1\_t \*data;    ...  s1\_tid = H5Tcreate(H5T\_COMPOUND, sizeof(s1\_t));  H5Tinsert(s1\_tid, “a\_name”, HOFFSET(s1\_t, a),  H5T\_NATIVE\_INT);  H5Tinsert(s1\_tid, “b\_name”, HOFFSET(s1\_t, b),  H5T\_NATIVE\_FLOAT);  H5Tinsert(s1\_tid, “c\_name”, HOFFSET(s1\_t, c),  H5T\_NATIVE\_DOUBLE);  ... |
| dataset\_id = H5Dopen(file\_id, “SDScompound.h5”,  H5P\_DEFAULT);  ...  data = (s1\_t \*) malloc (sizeof(s1\_t)\*LENGTH);  H5Dread(dataset\_id, s1\_tid, H5S\_ALL, H5S\_ALL,  H5P\_DEFAULT, data); |
| Code Example 6-21. Read a dataset using a memory datatype |

Instead of building a memory datatype, the application could use the H5Tget\_native\_type function. See the example below.

|  |
| --- |
| typedef struct s1\_t {     int a;     float b;     double c;  } s1\_t; |
| s1\_t \*data;  hid\_t file\_s1\_t, mem\_s1\_t;  ...  dataset\_id = H5Dopen(file\_id, “SDScompound.h5”, H5P\_DEFAULT);  /\* Discover datatype in the file \*/  file\_s1\_t  = H5Dget\_type(dataset\_id); |
| /\* Find corresponding memory datatype \*/  mem\_s1\_t   = H5Tget\_native\_type(file\_s1\_t, H5T\_DIR\_DEFAULT);  ...  data = (s1\_t \*) malloc (sizeof(s1\_t)\*LENGTH);  H5Dread (dataset\_id, mem\_s1\_tid, H5S\_ALL, H5S\_ALL, H5P\_DEFAULT, data); |
| Code Example 6-22. Read a dataset using H5Tget\_native\_type |

 The example below shows how to read just one float member of a compound datatype.

|  |
| --- |
| typedef struct s1\_t {     float b;  } sf\_t;    sf\_t \*data;    ...  sf\_tid = H5Tcreate(H5T\_COMPOUND, sizeof(sf\_t));  H5Tinsert(s1\_tid, “b\_name”, HOFFSET(sf\_t, b), H5T\_NATIVE\_FLOAT);  ...  dataset\_id = H5Dopen(file\_id, “SDScompound.h5”, H5P\_DEFAULT);  ...  data = (sf\_t \*) malloc (sizeof(sf\_t)\*LENGTH);  H5Dread(dataset\_id, sf\_tid, H5S\_ALL, H5S\_ALL, H5P\_DEFAULT, data); |
| Code Example 6-23. Read one floating point member of a compound datatype |

The example below shows how to read float and double members of a compound datatype into a struc­ture that has those fields in a different order. Please notice that H5Tinsert calls can be used in an order different from the order of the structure’s members.

|  |
| --- |
| typedef struct s1\_t {     double c;     float b;  } sdf\_t;    sdf\_t \*data;    ...  sdf\_tid = H5Tcreate(H5T\_COMPOUND, sizeof(sdf\_t));  H5Tinsert(sdf\_tid, “b\_name”, HOFFSET(sdf\_t, b), H5T\_NATIVE\_FLOAT);  H5Tinsert(sdf\_tid, “c\_name”, HOFFSET(sdf\_t, c), H5T\_NATIVE\_DOUBLE);  ...  dataset\_id = H5Dopen(file\_id, “SDScompound.h5”, H5P\_DEFAULT);  ...  data = (sdf\_t \*) malloc (sizeof(sdf\_t)\*LENGTH);  H5Dread(dataset\_id, sdf\_tid, H5S\_ALL, H5S\_ALL, H5P\_DEFAULT, data); |
| Code Example 6-24. Read float and double members of a compound datatype |

##### 6.5.2.2.2. Array

Many scientific datasets have multiple measurements for each point in a space. There are several natural ways to represent this data, depending on the variables and how they are used in computation. See the table and the figure below.

|  |  |  |
| --- | --- | --- |
| Table 6-20. Representing data with multiple measurements | | |
| **Storage Strategy** | **Stored as** | **Remarks** |
| Multiple planes | Several datasets with identical dataspaces | This is optimal when variables are accessed individually, or when often uses only selected variables. |
| Additional dimen­sion | One dataset, the last “dimension” is a vec­tor of variables | This can give good performance, although selecting only a few variables may be slow. This may not reflect the science. |
| Record with multi­ple values | One dataset with compound datatype | This enables the variables to be read all together or selected. Also handles “vectors” of heterogeneous data. |
| Vector or Tensor value | One dataset, each data element is a small array of values. | This uses the same amount of space as the previous two, and may represent the science model better. |

|  |  |
| --- | --- |
| Dtypes_fig26_pic1of4.JPG | Dtypes_fig26_pic2of4.JPG |
| Dtypes_fig26_pic3of4.JPG | Dtypes_fig26_pic4of4.JPG |
| Figure 6-13. Representing data with multiple measurements | |

The HDF5 H5T\_ARRAY datatype defines the data element to be a homogeneous, multi-dimensional array. See Figure 13d above. The elements of the array can be any HDF5 datatype (including compound and array), and the size of the datatype is the total size of the array. A dataset of array datatype cannot be sub­divided for I/O within the data element: the entire array of the data element must be transferred. If the data elements need to be accessed separately, for example, by plane, then the array datatype should not be used. The table below shows advantages and disadvantages of various storage methods.

|  |  |  |
| --- | --- | --- |
| Table 6-21. Storage method advantages and disadvantages | | |
| **Method** | **Advantages** | **Disadvantages** |
| a) Multiple Datasets | Easy to access each plane, can select any plane(s) | Less efficient to access a ‘col­umn’ through the planes |
| b) N+1 Dimension | All access patterns supported | Must be homogeneous data­type    The added dimension may not make sense in the scientific model |
| c) Compound Datatype | Can be heterogeneous datatype | Planes must be named, selec­tion is by plane    Not a natural representation for a matrix |
| d) Array | A natural representation for vector or tensor data | Cannot access elements sepa­rately (no access by plane) |

An array datatype may be multi-dimensional with 1 to H5S\_MAX\_RANK (the maximum rank of a dataset is currently 32) dimensions. The dimensions can be any size greater than 0, but unlimited dimensions are not supported (although the datatype can be a variable-length datatype).

An array datatype is created with the H5Tarray\_create call, which specifies the number of dimensions, the size of each dimension, and the base type of the array. The array datatype can then be used in any way that any datatype object is used. The example below shows the creation of a datatype that is a two-dimensional array of native integers, and this is then used to create a dataset. Note that the dataset can be a dataspace that is any number and size of dimensions. The figure below shows the layout in memory assuming that the native integers are 4 bytes. Each data element has 6 elements, for a total of 24 bytes.

|  |
| --- |
| hid\_t file, dataset;  hid\_t datatype, dataspace;  hsize\_t adims[] = {3, 2};    datatype = H5Tarray\_create(H5T\_NATIVE\_INT, 2, adims,  NULL);    dataset = H5Dcreate(file, datasetname, datatype,  dataspace, H5P\_DEFAULT, H5P\_DEFAULT,  H5P\_DEFAULT); |
| Code Example 6-25. Create a two-dimensional array datatype |

|  |
| --- |
| Dtypes_fig28.JPG |
| Figure 6-14. Memory layout of a two-dimensional array datatype |

##### 6.5.2.2.3. Variable-length Datatypes

A variable-length (VL) datatype is a one-dimensional sequence of a datatype which are not fixed in length from one dataset location to another. In other words, each data element may have a different number of members. Variable-length datatypes cannot be divided; the entire data element must be transferred.

VL datatypes are useful to the scientific community in many different ways, possibly including:

•        Ragged arrays: Multi-dimensional ragged arrays can be implemented with the last (fastest chang­ing) dimension being ragged by using a VL datatype as the type of the element stored.

•        Fractal arrays: A nested VL datatype can be used to implement ragged arrays of ragged arrays, to whatever nesting depth is required for the user.

•        Polygon lists: A common storage requirement is to efficiently store arrays of polygons with differ­ent numbers of vertices. A VL datatype can be used to efficiently and succinctly describe an array of polygons with different numbers of vertices.

•        Character strings: Perhaps the most common use of VL datatypes will be to store C-like VL charac­ter strings in dataset elements or as attributes of objects.

•        Indices (for example, of objects within a file): An array of VL object references could be used as an index to all the objects in a file which contain a particular sequence of dataset values.

•        Object Tracking: An array of VL dataset region references can be used as a method of tracking objects or features appearing in a sequence of datasets.

A VL datatype is created by calling H5Tvlen\_create which specifies the base datatype. The first example below shows an example of code that creates a VL datatype of unsigned integers. Each data element is a one-dimensional array of zero or more members and is stored in the hvl\_t structure. See the second example below.

|  |
| --- |
| tid1 = H5Tvlen\_create (H5T\_NATIVE\_UINT);    dataset=H5Dcreate(fid1, “Dataset1”, tid1, sid1,  H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT); |
| Code Example 6-26. Create a variable-length datatype of unsigned integers |

|  |
| --- |
| typedef struct  {     size\_t len; /\* Length of VL data \*/                       /\*(in base type units) \*/     void \*p; /\* Pointer to VL data \*/  } hvl\_t; |
| Code Example 6-27. Data element storage for members of the VL datatype |

The first example below shows how the VL data is written. For each of the 10 data elements, a length and data buffer must be allocated. Below the two examples is a figure that shows how the data is laid out in memory.

An analogous procedure must be used to read the data. See the second example below. An appropriate array of vl\_t must be allocated, and the data read. It is then traversed one data element at a time. The H5Dvlen\_reclaim call frees the data buffer for the buffer. With each element possibly being of different sequence lengths for a dataset with a VL datatype, the memory for the VL datatype must be dynamically allocated. Currently there are two methods of managing the memory for VL datatypes: the standard C malloc/free memory allocation routines or a method of calling user-defined memory management rou­tines to allocate or free memory (set with H5Pset\_vlen\_mem\_manager). Since the memory allocated when reading (or writing) may be complicated to release, the H5Dvlen\_reclaim function is provided to traverse a memory buffer and free the VL datatype information without leaking memory.

|  |
| --- |
| hvl\_t wdata[10];           /\* Information to write \*/    /\* Allocate and initialize VL data to write \*/  for(i=0; i < 10; i++) {     wdata[i].p = malloc((i+1)\*sizeof(unsigned int));     wdata[i].len = i+1;     for(j=0; j<(i+1); j++)        ((unsigned int \*)wdata[i].p)[j]=i\*10+j;  }    ret=H5Dwrite(dataset, tid1, H5S\_ALL, H5S\_ALL, H5P\_DEFAULT, wdata); |
| Code Example 6-28. Write VL data |

|  |
| --- |
| hvl\_t rdata[SPACE1\_DIM1];  ret=H5Dread(dataset, tid1, H5S\_ALL, H5S\_ALL, xfer\_pid, rdata);    for(i=0; i<SPACE1\_DIM1; i++) {     printf(“%d: len %d ”,rdata[i].len);     for(j=0; j<rdata[i].len; j++) {        printf(“ value: %u\n”,((unsigned int \*)rdata[i].p)[j]);     }  }  ret=H5Dvlen\_reclaim(tid1, sid1, xfer\_pid, rdata); |
| Code Example 6-29. Read VL data |

|  |
| --- |
| Dtypes_fig33.JPG |
| Figure 6-15. Memory layout of a VL datatype |

The user program must carefully manage these relatively complex data structures. The H5Dvlen\_re­claim function performs a standard traversal, freeing all the data. This function analyzes the datatype and dataspace objects, and visits each VL data element, recursing through nested types. By default, the system free is called for the pointer in each vl\_t. Obviously, this call assumes that all of this memory was allocated with the system malloc.

The user program may specify custom memory manager routines, one for allocating and one for freeing. These may be set with the H5Pvlen\_mem\_manager, and must have the following prototypes:

* typedef void \*(\*H5MM\_allocate\_t)(size\_t size, void \*info);
* typedef void (\*H5MM\_free\_t)(void \*mem, void \*free\_info);

The utility function H5Dget\_vlen\_buf\_size checks the number of bytes required to store the VL data from the dataset. This function analyzes the datatype and dataspace object to visit all the VL data ele­ments, to determine the number of bytes required to store the data for the in the destination storage (memory). The size value is adjusted for data conversion and alignment in the destination.

## Other Non-numeric Datatypes

Several datatype classes define special types of objects.

### Strings

Text data is represented by arrays of characters, called strings. Many programming languages support dif­ferent conventions for storing strings, which may be fixed or variable-length, and may have different rules for padding unused storage. HDF5 can represent strings in several ways. See the figure below.

|  |
| --- |
| The strings to store are “Four score” and “lazy programmers.” |
| a) H5T\_NATIVE\_CHAR: The dataset is a one-dimensional array with 29 elements, and each element is a single character. |
| Dtypes_fig16a.JPG |
| b) Fixed-length string: The dataset is a one-dimensional array with two elements, and each element is 20 characters. |
| Dtypes_fig16b.JPG |
| c) Variable-length string: The dataset is a one-dimensional array with two elements, and each ele­ment is a variable-length string. This is the same result when stored as a fixed-length string except that the first element of the array will need only 11 bytes for storage instead of 20. |
| Dtypes_fig16c.JPG |
| Dtypes_fig16d.JPG |
| Figure 6-16. A string stored as one-character elements in a one-dimensional array |

First, a dataset may have a dataset with datatype H5T\_NATIVE\_CHAR with each character of the string as an element of the dataset. This will store an unstructured block of text data, but gives little indication of any structure in the text. See item a in the figure above.

A second alternative is to store the data using the datatype class H5T\_STRING with each element a fixed length. See item b in the figure above. In this approach, each element might be a word or a sentence, addressed by the dataspace. The dataset reserves space for the specified number of characters, although some strings may be shorter. This approach is simple and usually is fast to access, but can waste storage space if the length of the Strings varies.

A third alternative is to use a variable-length datatype. See item c in the figure above. This can be done using the standard mechanisms described above. The program would use vl\_t structures to write and read the data.

A fourth alternative is to use a special feature of the string datatype class to set the size of the datatype to H5T\_VARIABLE. See item c in the figure above. The example below shows a declaration of a datatype of type H5T\_C\_S1 which is set to H5T\_VARIABLE. The HDF5 Library automatically translates between this and the vl\_t structure. Note: the H5T\_VARIABLE size can only be used with string datatypes.

|  |
| --- |
| tid1 = H5Tcopy (H5T\_C\_S1);  ret = H5Tset\_size (tid1, H5T\_VARIABLE); |
| Code Example 6-30. Set the string datatype size to H5T\_VARIABLE |

Variable-length strings can be read into C strings (in other words, pointers to zero terminated arrays of char). See the example below.

|  |
| --- |
| char \*rdata[SPACE1\_DIM1];    ret=H5Dread(dataset, tid1, H5S\_ALL, H5S\_ALL, xfer\_pid, rdata);    for(i=0; i<SPACE1\_DIM1; i++) {  printf(“%d: len: %d, str is: %s\n”, i, strlen(rdata[i]),  rdata[i]);  }    ret=H5Dvlen\_reclaim(tid1, sid1, xfer\_pid, rdata); |
| Code Example 6-31. Read variable-length strings into C strings |

### Reference

In HDF5, objects (groups, datasets, and committed datatypes) are usually accessed by name. There is another way to access stored objects - by reference. There are two reference datatypes: object reference and region reference. Object reference objects are created with H5Rcreate and other calls (cross refer­ence). These objects can be stored and retrieved in a dataset as elements with reference datatype. The first example below shows an example of code that creates references to four objects, and then writes the array of object references to a dataset. The second example below shows a dataset of datatype reference being read and one of the reference objects being dereferenced to obtain an object pointer.

In order to store references to regions of a dataset, the datatype should be H5T\_REGION\_OBJ. Note that a data element must be either an object reference or a region reference: these are different types and cannot be mixed within a single array.

A reference datatype cannot be divided for I/O: an element is read or written completely.

|  |
| --- |
| dataset= H5Dcreate (fid1, “Dataset3”, H5T\_STD\_REF\_OBJ, sid1, H5P\_DEFAULT,  H5P\_DEFAULT, H5P\_DEFAULT);     /\* Create reference to dataset \*/     ret = H5Rcreate(&wbuf[0], fid1,“/Group1/Dataset1”, H5R\_OBJECT, -1); |
| /\* Create reference to dataset \*/     ret = H5Rcreate(&wbuf[1], fid1, “/Group1/Dataset2”, H5R\_OBJECT, -1);     /\* Create reference to group \*/     ret = H5Rcreate(&wbuf[2], fid1, “/Group1”, H5R\_OBJECT, -1); |
| /\* Create reference to committed datatype \*/     ret = H5Rcreate(&wbuf[3], fid1, “/Group1/Datatype1”,           H5R\_OBJECT, -1);       /\* Write selection to disk \*/    ret=H5Dwrite(dataset, H5T\_STD\_REF\_OBJ, H5S\_ALL, H5S\_ALL, H5P\_DEFAULT, wbuf); |
| Code Example 6-32. Create object references and write to a dataset |

|  |
| --- |
| rbuf = malloc(sizeof(hobj\_ref\_t)\*SPACE1\_DIM1);    /\* Read selection from disk \*/  ret=H5Dread(dataset, H5T\_STD\_REF\_OBJ, H5S\_ALL, H5S\_ALL,  H5P\_DEFAULT, rbuf);    /\* Open dataset object \*/  dset2 = H5Rdereference(dataset, H5R\_OBJECT, &rbuf[0]); |
| Code Example 6-33. Read a dataset with a reference datatype |

### ENUM

The enum datatype implements a set of (name, value) pairs, similar to C/C++ enum. The values are cur­rently limited to native integer datatypes. Each name can be the name of only one value, and each value can have only one name.

The data elements of the ENUMERATION are stored according to the datatype. An example would be as an array of integers. The example below shows an example of how to create an enumeration with five ele­ments. The elements map symbolic names to 2-byte integers. See the table below.

|  |
| --- |
| hid\_t hdf\_en\_colors;  short val;  hdf\_en\_colors = H5Tcreate(H5T\_ENUM, sizeof(short));     H5Tenum\_insert(hdf\_en\_colors, “RED”, (val=0,&val));     H5Tenum\_insert(hdf\_en\_colors, “GREEN”, (val=1,&val));     H5Tenum\_insert(hdf\_en\_colors, “BLUE”, (val=2,&val));     H5Tenum\_insert(hdf\_en\_colors, “WHITE”, (val=3,&val));     H5Tenum\_insert(hdf\_en\_colors, “BLACK”, (val=4,&val));       H5Dcreate(fileid, datasetname, hdf\_en\_colors, spaceid,           H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT); |
| Code Example 6-34. Create an enumeration with five elements |

|  |  |
| --- | --- |
| Table 6-22. An enumeration with five elements | |
| **Name** | **Value** |
| RED | 0 |
| GREEN | 1 |
| BLUE | 2 |
| WHITE | 3 |
| BLACK | 4 |

The figure below shows how an array of eight values might be stored. Conceptually, the array is an array of symbolic names [BLACK, RED, WHITE, BLUE, ...]. See item a in the figure below. These are stored as the val­ues and are short integers. So, the first 2 bytes are the value associated with “BLACK”, which is the number 4, and so on. See item b in the figure below.

|  |
| --- |
| a) Logical data to be written - eight elements |
| Dtypes_fig17a.JPG |
| Dtypes_fig40.JPG |
| b) The storage layout. Total size of the array is 16 bytes, 2 bytes per element. |
| Figure 6-17. Storing an enum array |

The order that members are inserted into an enumeration type is unimportant; the important part is the associations between the symbol names and the values. Thus, two enumeration datatypes will be consid­ered equal if and only if both types have the same symbol/value associations and both have equal under­lying integer datatypes. Type equality is tested with the H5Tequal function.

If a particular architecture type is required, a little-endian or big-endian datatype for example, use a native integer datatype as the ENUM base datatype and use H5Tconvert on values as they are read from or written to a dataset.

### Opaque

In some cases, a user may have data objects that should be stored and retrieved as blobs with no attempt to interpret them. For example, an application might wish to store an array of encrypted certificates which are 100 bytes long.

While an arbitrary block of data may always be stored as bytes, characters, integers, or whatever, this might mislead programs about the meaning of the data. The opaque datatype defines data elements which are uninterpreted by HDF5. The opaque data may be labeled with H5Tset\_tag with a string that might be used by an application. For example, the encrypted certificates might have a tag to indicate the encryption and the certificate standard.

### Bitfield

Some data is represented as bits, where the number of bits is not an integral byte and the bits are not nec­essarily interpreted as a standard type. Some examples might include readings from machine registers (for example, switch positions), a cloud mask, or data structures with several small integers that should be store in a single byte.

This data could be stored as integers, strings, or enumerations. However, these storage methods would likely result in considerable wasted space. For example, storing a cloud mask with one byte per value would use up to eight times the space of a packed array of bits.

The HDF5 bitfield datatype class defines a data element that is a contiguous sequence of bits, which are stored on disk in a packed array. The programming model is the same as for unsigned integers: the data­type object is created by copying a predefined datatype, and then the precision, offset, and padding are set.

While the use of the bitfield datatype will reduce storage space substantially, there will still be wasted space if the bitfield as a whole does not match the 1-, 2-, 4-, or 8-byte unit in which it is written. The remaining unused space can be removed by applying the N-bit filter to the dataset containing the bitfield data. For more information, see "Using the N-bit Filter."

## Fill Values

The “fill value” for a dataset is the specification of the default value assigned to data elements that have not yet been written. In the case of a dataset with an atomic datatype, the fill value is a single value of the appropriate datatype, such as ‘0’ or ‘-1.0’. In the case of a dataset with a composite datatype, the fill value is a single data element of the appropriate type. For example, for an array or compound datatype, the fill value is a single data element with values for all the component elements of the array or compound data­type.

The fill value is set (permanently) when the dataset is created. The fill value is set in the dataset creation properties in the H5Dcreate call. Note that the H5Dcreate call must also include the datatype of the dataset, and the value provided for the fill value will be interpreted as a single element of this datatype. The example below shows code which creates a dataset of integers with fill value -1. Any unwritten data elements will be set to -1.

|  |
| --- |
| hid\_t       plist\_id;  int filler;    filler = -1;  plist\_id = H5Pcreate(H5P\_DATASET\_CREATE);  H5Pset\_fill\_value(plist\_id, H5T\_NATIVE\_INT, &filler);    /\* Create the dataset with fill value ‘-1’. \*/  dataset\_id = H5Dcreate(file\_id, “/dset”, H5T\_STD\_I32BE,  dataspace\_id, H5P\_DEFAULT, plist\_  id, H5P\_DEFAULT); |
| Code Example 6-35. Create a dataset with a fill value of -1 |

|  |
| --- |
| typedef struct s1\_t {     int a;     char b;     double c;  } s1\_t;  s1\_t filler; |
| s1\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(s1\_t));  H5Tinsert(s1\_tid, “a\_name”, HOFFSET(s1\_t, a), H5T\_NATIVE\_INT);  H5Tinsert(s1\_tid, “b\_name”, HOFFSET(s1\_t, b), H5T\_NATIVE\_CHAR);  H5Tinsert(s1\_tid, “c\_name”, HOFFSET(s1\_t, c), H5T\_NATIVE\_DOUBLE);    filler.a = -1;  filler.b = ‘\*’;  filler.c = -2.0; |
| plist\_id = H5Pcreate(H5P\_DATASET\_CREATE);  H5Pset\_fill\_value(plist\_id, s1\_tid, &filler);    /\* Create the dataset with fill value \*/  /\* (-1, ‘\*’, -2.0). \*/  dataset = H5Dcreate(file, datasetname, s1\_tid, space, H5P\_DEFAULT, plist\_id, H5P\_DEFAULT); |
| Code Example 6-36. Create a fill value for a compound datatype |

The figure above shows how to create a fill value for a compound datatype. The procedure is the same as the previous example except the filler must be a structure with the correct fields. Each field is initialized to the desired fill value.

The fill value for a dataset can be retrieved by reading the dataset creation properties of the dataset and then by reading the fill value with H5Pget\_fill\_value. The data will be read into memory using the storage layout specified by the datatype. This transfer will convert data in the same way as H5Dread. The example below shows how to get the fill value from the dataset created in the example "Create a dataset with a fill value of -1".

|  |
| --- |
| hid\_t plist2;  int filler;    dataset\_id = H5Dopen(file\_id, “/dset”, H5P\_DEFAULT);  plist2 = H5Dget\_create\_plist(dataset\_id);    H5Pget\_fill\_value(plist2, H5T\_NATIVE\_INT, &filler);    /\* filler has the fill value, ‘-1’ \*/ |
| Code Example 6-37. Retrieve a fill value |

A similar procedure is followed for any datatype. The example below shows how to read the fill value for the compound datatype created in an example above. Note that the program must pass an element large enough to hold a fill value of the datatype indicated by the argument to H5Pget\_fill\_value. Also, the program must understand the datatype in order to interpret its components. This may be difficult to determine without knowledge of the application that created the dataset.

|  |
| --- |
| char \*       fillbuf;  int sz;  dataset = H5Dopen( file, DATASETNAME, H5P\_DEFAULT);    s1\_tid = H5Dget\_type(dataset); |
| sz = H5Tget\_size(s1\_tid);    fillbuf = (char \*)malloc(sz);    plist\_id = H5Dget\_create\_plist(dataset); |
| H5Pget\_fill\_value(plist\_id, s1\_tid, fillbuf);    printf(“filler.a: %d\n”,((s1\_t \*) fillbuf)->a);  printf(“filler.b: %c\n”,((s1\_t \*) fillbuf)->b);  printf(“filler.c: %f\n”,((s1\_t \*) fillbuf)->c); |
| Code Example 6-38. Read the fill value for a compound datatype |

## Complex Combinations of Datatypes

Several composite datatype classes define collections of other datatypes, including other composite data­types. In general, a datatype can be nested to any depth, with any combination of datatypes.

For example, a compound datatype can have members that are other compound datatypes, arrays, VL datatypes. An array can be an array of array, an array of compound, or an array of VL. And a VL datatype can be a variable-length array of compound, array, or VL datatypes.

These complicated combinations of datatypes form a logical tree, with a single root datatype, and leaves which must be atomic datatypes (predefined or user-defined). The figure below shows an example of a logical tree describing a compound datatype constructed from different datatypes.

Recall that the datatype is a description of the layout of storage. The complicated compound datatype is constructed from component datatypes, each of which describes the layout of part of the storage. Any datatype can be used as a component of a compound datatype, with the following restrictions:

1.      No byte can be part of more than one component datatype (in other words, the fields cannot overlap within the compound datatype)

2.      The total size of the components must be less than or equal to the total size of the compound datatype

These restrictions are essentially the rules for C structures and similar record types familiar from program­ming languages. Multiple typing, such as a C union, is not allowed in HDF5 datatypes.

|  |
| --- |
| Dtypes_fig45.JPG |
| Figure 6-18. A compound datatype built with different datatypes |

### Creating a Complicated Compound Datatype

To construct a complicated compound datatype, each component is constructed, and then added to the enclosing datatype description. The example below shows how to create a compound datatype with four members:

* “T1”, a compound datatype with three members
* “T2”, a compound datatype with two members
* “T3”, a one-dimensional array of integers
* “T4”, a string

Below the example code is a figure that shows this datatype as a logical tree. The output of the h5dump utility is shown in the example below the figure.

Each datatype is created as a separate datatype object. Figure 20 below shows the storage layout for the four individual datatypes. Then the datatypes are inserted into the outer datatype at an appropriate off­set. Figure 21 below shows the resulting storage layout. The combined record is 89 bytes long.

The Dataset is created using the combined compound datatype. The dataset is declared to be a 4 by 3 array of compound data. Each data element is an instance of the 89-byte compound datatype. Figure 22 below shows the layout of the dataset, and expands one of the elements to show the relative position of the component data elements.

Each data element is a compound datatype, which can be written or read as a record, or each field may be read or written individually. The first field (“T1”) is itself a compound datatype with three fields (“T1.a”, “T1.b”, and “T1.c”). “T1” can be read or written as a record, or individual fields can be accessed. Similarly, the second filed is a compound datatype with two fields (“T2.f1”, “T2.f2”).

The third field (“T3”) is an array datatype. Thus, “T3” should be accessed as an array of 40 integers. Array data can only be read or written as a single element, so all 40 integers must be read or written to the third field. The fourth field (“T4”) is a single string of length 25.

|  |
| --- |
| typedef struct s1\_t {     int a;     char b;     double c;  } s1\_t; |
| typedef struct s2\_t {     float f1;     float f2;  } s2\_t;  hid\_t      s1\_tid, s2\_tid, s3\_tid, s4\_tid, s5\_tid; |
| /\* Create a datatype for s1 \*/  s1\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(s1\_t));  H5Tinsert(s1\_tid, “a\_name”, HOFFSET(s1\_t, a), H5T\_NATIVE\_INT);  H5Tinsert(s1\_tid, “b\_name”, HOFFSET(s1\_t, b), H5T\_NATIVE\_CHAR);  H5Tinsert(s1\_tid, “c\_name”, HOFFSET(s1\_t, c), H5T\_NATIVE\_DOUBLE); |
| /\* Create a datatype for s2. \*.  s2\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(s2\_t));  H5Tinsert(s2\_tid, “f1”, HOFFSET(s2\_t, f1), H5T\_NATIVE\_FLOAT);  H5Tinsert(s2\_tid, “f2”, HOFFSET(s2\_t, f2), H5T\_NATIVE\_FLOAT); |
| /\* Create a datatype for an Array of integers \*/  s3\_tid = H5Tarray\_create(H5T\_NATIVE\_INT, RANK, dim);    /\* Create a datatype for a String of 25 characters \*/  s4\_tid = H5Tcopy(H5T\_C\_S1);  H5Tset\_size(s4\_tid, 25); |
| /\*  \* Create a compound datatype composed of one of each of  \* these types. The total size is the sum of the size of  \* each.  \*/    sz = H5Tget\_size(s1\_tid) + H5Tget\_size(s2\_tid) +           H5Tget\_size(s3\_tid) + H5Tget\_size(s4\_tid); |
| s5\_tid = H5Tcreate (H5T\_COMPOUND, sz);    /\* Insert the component types at the appropriate \*/  \* offsets.  \*/ |
| H5Tinsert(s5\_tid, “T1”, 0, s1\_tid);  H5Tinsert(s5\_tid, “T2”, sizeof(s1\_t), s2\_tid);  H5Tinsert(s5\_tid, “T3”, sizeof(s1\_t)+sizeof(s2\_t), s3\_tid);  H5Tinsert(s5\_tid, “T4”, (sizeof(s1\_t) +sizeof(s2\_t)+ H5Tget\_size(s3\_tid)), s4\_tid); |
| /\*  \* Create the dataset with this datatype.  \*/  dataset = H5Dcreate(file, DATASETNAME, s5\_tid, space,  H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT); |
| Code Example 6-39. Create a compound datatype with four members |

|  |
| --- |
| Dtypes_fig47.JPG |
| Figure 6-19. Logical tree for the compound datatype with four members |

|  |
| --- |
| DATATYPE  H5T\_COMPOUND {     H5T\_COMPOUND {        H5T\_STD\_I32LE “a\_name”;        H5T\_STD\_I8LE “b\_name”;        H5T\_IEEE\_F64LE “c\_name”;     } “T1”;     H5T\_COMPOUND {        H5T\_IEEE\_F32LE “f1”;        H5T\_IEEE\_F32LE “f2”;     } “T2”; |
| H5T\_ARRAY { [10] H5T\_STD\_I32LE } “T3”;     H5T\_STRING {        STRSIZE 25;        STRPAD H5T\_STR\_NULLTERM;        CSET H5T\_CSET\_ASCII;        CTYPE H5T\_C\_S1;     } “T4”;  } |
| Code Example 6-40. Output from h5dump for the compound datatype |

|  |
| --- |
| a) Compound type ‘s1\_t’, size 16 bytes. |
| Dtypes_fig20a.JPG |
| b) Compound type ‘s2\_t’, size 8 bytes. |
| Dtypes_fig20b.JPG |
| c) Array type ‘s3\_tid’, 40 integers, total size 40 bytes. |
| Dtypes_fig20c.JPG |
| d) String type ‘s4\_tid’, size 25 bytes. |
| Dtypes_fig20d.JPG |
| Figure 6-20. The storage layout for the four member datatypes |

|  |
| --- |
| Dtypes_fig50.JPG |
| Figure 6-21. The storage layout of the combined four members |

|  |
| --- |
| a) A 4 x 3 array of Compound Datatype |
| dtypes_fig51new.JPG |
| b) Element [1,1] expanded |
| Figure 6-22. The layout of the dataset |

### Analyzing and Navigating a Compound Datatype

A complicated compound datatype can be analyzed piece by piece to discover the exact storage layout. In the example above, the outer datatype is analyzed to discover that it is a compound datatype with four members. Each member is analyzed in turn to construct a complete map of the storage layout.

The example below shows an example of code that partially analyzes a nested compound datatype. The name and overall offset and size of the component datatype is discovered, and then its type is analyzed depending on the datatype class. Through this method, the complete storage layout can be discovered.

|  |
| --- |
| s1\_tid = H5Dget\_type(dataset); |
| if (H5Tget\_class(s1\_tid) == H5T\_COMPOUND) {     printf(“COMPOUND DATATYPE {\n”);     sz = H5Tget\_size(s1\_tid);     nmemb = H5Tget\_nmembers(s1\_tid);     printf(“ %d bytes\n”,sz);     printf(“ %d members\n”,nmemb); |
| for (i =0; i < nmemb; i++) {        s2\_tid = H5Tget\_member\_type(s1\_tid, i);        if (H5Tget\_class(s2\_tid) == H5T\_COMPOUND) {           /\* recursively analyze the nested type. \*/ |
| } else if (H5Tget\_class(s2\_tid) == H5T\_ARRAY) {           sz2 = H5Tget\_size(s2\_tid);           printf(“  %s: NESTED ARRAY DATATYPE offset %d size %d                    {\n”,              H5Tget\_member\_name(s1\_tid, i),              H5Tget\_member\_offset(s1\_tid, i),                    sz2);              H5Tget\_array\_dims(s2\_tid, dim);                    s3\_tid = H5Tget\_super(s2\_tid);                    /\* Etc., analyze the base type of the array \*/ |
| } else {           /\* analyze a simple type \*/           printf(“ %s: type code %d offset %d size %d\n”,              H5Tget\_member\_name(s1\_tid, i),              H5Tget\_class(s2\_tid),              H5Tget\_member\_offset(s1\_tid, i),              H5Tget\_size(s2\_tid));        }        /\* and so on…. \*/ |
| Code Example 6-41. Analyzing a compound datatype and its members |

## Life Cycle of the Datatype Object

Application programs access HDF5 datatypes through identifiers. Identifiers are obtained by creating a new datatype or by copying or opening an existing datatype. The identifier can be used until it is closed or until the library shuts down. See items a and b in the figure below. By default, a datatype is transient, and it disappears when it is closed.

When a dataset or attribute is created (H5Dcreate or H5Acreate), its datatype is stored in the HDF5 file as part of the dataset or attribute object. See item c in the figure below. Once an object created, its data­type cannot be changed or deleted. The datatype can be accessed by calling H5Dget\_type, H5Aget\_­type, H5Tget\_super, or H5Tget\_member\_type. See item d in the figure below. These calls return an identifier to a transient copy of the datatype of the dataset or attribute unless the datatype is a committed datatype.

Note that when an object is created, the stored datatype is a copy of the transient datatype. If two objects are created with the same datatype, the information is stored in each object with the same effect as if two different datatypes were created and used.

A transient datatype can be stored using H5Tcommit in the HDF5 file as an independent, named object, called a committed datatype. Committed datatypes were formerly known as named datatypes. See item e in the figure below. Subsequently, when a committed datatype is opened with H5Topen (item f), or is obtained with H5Tget\_type or similar call (item k), the return is an identifier to a transient copy of the stored datatype. The identifier can be used in the same way as other datatype identifiers except that the committed datatype cannot be modified. When a committed datatype is copied with H5Tcopy, the return is a new, modifiable, transient datatype object (item f).

When an object is created using a committed datatype (H5Dcreate, H5Acreate), the stored datatype is used without copying it to the object. See item j in the figure below. In this case, if multiple objects are cre­ated using the same committed datatype, they all share the exact same datatype object. This saves space and makes clear that the datatype is shared. Note that a committed datatype can be shared by objects within the same HDF5 file, but not by objects in other files. For more information on copying committed datatypes to other HDF5 files, see the “Copying Committed Datatypes with H5Ocopy” topic in the “Addi­tional Resources” chapter.

A committed datatype can be deleted from the file by calling H5Ldelete which replaces H5Gunlink. See item i in the figure below. If one or more objects are still using the datatype, the committed datatype can­not be accessed with H5Topen, but will not be removed from the file until it is no longer used. H5Tget\_­type and similar calls will return a transient copy of the datatype.

|  |
| --- |
| Dtypes_fig53.JPG |
| Figure 6-23. Life cycle of a datatype |

Transient datatypes are initially modifiable. Note that when a datatype is copied or when it is written to the file (when an object is created) or the datatype is used to create a composite datatype, a copy of the current state of the datatype is used. If the datatype is then modified, the changes have no effect on data­sets, attributes, or datatypes that have already been created. See the figure below.

A transient datatype can be made read-only (H5Tlock). Note that the datatype is still transient, and oth­erwise does not change. A datatype that is immutable is read-only but cannot be closed except when the entire library is closed. The predefined types such as H5T\_NATIVE\_INT are immutable transient types.

|  |
| --- |
| Dtypes_fig54.JPG |
| Figure 6-24. Transient datatype states: modifiable, read-only, and immutable |

To create two or more datasets that share a common datatype, first commit the datatype, and then use that datatype to create the datasets. See the example below.

|  |
| --- |
| hid\_t t1 = ...some transient type...;  H5Tcommit (file, “shared\_type”, t1, H5P\_DEFAULT, H5P\_DEFAULT,  H5P\_DEFAULT);  hid\_t dset1 = H5Dcreate (file, “dset1”, t1, space, H5P\_DEFAULT,  H5P\_DEFAULT, H5P\_DEFAULT);  hid\_t dset2 = H5Dcreate (file, “dset2”, t1, space, H5P\_DEFAULT,  H5P\_DEFAULT, H5P\_DEFAULT); |
| hid\_t dset1 = H5Dopen (file, “dset1”, H5P\_DEFAULT);  hid\_t t2 = H5Dget\_type (dset1);  hid\_t dset3 = H5Dcreate (file, “dset3”, t2, space, H5P\_DEFAULT,  H5P\_DEFAULT, H5P\_DEFAULT);  hid\_t dset4 = H5Dcreate (file, “dset4”, t2, space, H5P\_DEFAULT,  H5P\_DEFAULT, H5P\_DEFAULT); |
| Code Example 6-42. Create a shareable datatype |

|  |  |
| --- | --- |
| Table 6-23. Datatype APIs | |
| **Function** | **Description** |
| hid\_t H5Topen (hid\_t location, const char \*name) | A committed datatype can be opened by calling this function, which returns a data­type identifier. The identifier should even­tually be released by calling H5Tclose() to release resources. The committed data­type returned by this function is read-only or a negative value is returned for failure. The location is either a file or group iden­tifier. |
| herr\_t H5Tcommit (hid\_t location, const char \*name, hid\_t type, H5P\_DE­FAULT, H5P\_DEFAULT, H5P\_DEFAULT) | A transient datatype (not immutable) can be written to a file and turned into a com­mitted datatype by calling this function. The location is either a file or group iden­tifier and when combined with name refers to a new committed datatype. |
| htri\_t H5Tcommitted (hid\_t type) | A type can be queried to determine if it is a committed type or a transient type. If this function returns a positive value then the type is committed. Datasets which return committed datatypes with H5Dget\_type() are able to share the datatype with other datasets in the same file. |

## Data Transfer: Datatype Conversion and Selection

When data is transferred (write or read), the storage layout of the data elements may be different. For example, an integer might be stored on disk in big-endian byte order and read into memory with little-endian byte order. In this case, each data element will be transformed by the HDF5 Library during the data transfer.

The conversion of data elements is controlled by specifying the datatype of the source and specifying the intended datatype of the destination. The storage format on disk is the datatype specified when the data­set is created. The datatype of memory must be specified in the library call.

In order to be convertible, the datatype of the source and destination must have the same datatype class (with the exception of enumeration type). Thus, integers can be converted to other integers, and floats to other floats, but integers cannot (yet) be converted to floats. For each atomic datatype class, the possible conversions are defined. An enumeration datatype can be converted to an integer or a floating-point num­ber datatype.

Basically, any datatype can be converted to another datatype of the same datatype class. The HDF5 Library automatically converts all properties. If the destination is too small to hold the source value then an over­flow or underflow exception occurs. If a handler is defined with the H5Pset\_type\_conv\_cb function, it will be called. Otherwise, a default action will be performed. The table below summarizes the default actions.

|  |  |  |
| --- | --- | --- |
| Table 6-24. Default actions for datatype conversion exceptions | | |
| **Datatype Class** | **Possible Exceptions** | **Default Action** |
| Integer | Size, offset, pad |  |
| Float | Size, offset, pad, ebits |  |
| String | Size | Truncates, zero terminate if required. |
| Enumeration | No field | All bits set |

For example, when reading data from a dataset, the source datatype is the datatype set when the dataset was created, and the destination datatype is the description of the storage layout in memory. The destina­tion datatype must be specified in the H5Dread call. The example below shows an example of reading a dataset of 32-bit integers. The figure below the example shows the data transformation that is performed.

|  |
| --- |
| /\* Stored as H5T\_STD\_BE32 \*/  /\* Use the native memory order in the destination \*/  mem\_type\_id = H5Tcopy(H5T\_NATIVE\_INT);  status = H5Dread(dataset\_id, mem\_type\_id, mem\_space\_id,  file\_space\_id,  xfer\_plist\_id,  buf ); |
| Code Example 6-43. Specify the destination datatype with H5Dread |

|  |
| --- |
| Source Datatype: H5T\_STD\_BE32 |
| Dtypes_fig57a.JPG |
| .... |
| dtypes_fig57_arrowWithText.JPG |
| Destination Datatype: H5T\_STD\_LE32 |
| Dtypes_fig57b.JPG |
| .... |
| Figure 6-25. Layout of a datatype conversion |

One thing to note in the example above is the use of the predefined native datatype H5T\_NATIVE\_INT. Recall that in this example, the data was stored as a 4-bytes in big-endian order. The application wants to read this data into an array of integers in memory. Depending on the system, the storage layout of mem­ory might be either big or little-endian, so the data may need to be transformed on some platforms and not on others. The H5T\_NATIVE\_INT type is set by the HDF5 Library to be the correct type to describe the storage layout of the memory on the system. Thus, the code in the example above will work correctly on any platform, performing a transformation when needed.

There are predefined native types for most atomic datatypes, and these can be combined in composite datatypes. In general, the predefined native datatypes should always be used for data stored in memory.

|  |
| --- |
| **Storage Properties** |
| Predefined native datatypes describe the storage properties of memory. |

For composite datatypes, the component atomic datatypes will be converted. For a variable-length data­type, the source and destination must have compatible base datatypes. For a fixed-size string datatype, the length and padding of the strings will be converted. Variable-length strings are converted as variable-length datatypes.

For an array datatype, the source and destination must have the same rank and dimensions, and the base datatype must be compatible. For example an array datatype of 4 x 3 32-bit big-endian integers can be transferred to an array datatype of 4 x 3 little-endian integers, but not to a 3 x 4 array.

For an enumeration datatype, data elements are converted by matching the symbol names of the source and destination datatype. The figure below shows an example of how two enumerations with the same names and different values would be converted. The value ‘2’ in the source dataset would be converted to ‘0x0004’ in the destination.

If the source data stream contains values which are not in the domain of the conversion map then an over­flow exception is raised within the library.

|  |
| --- |
| Dtypes_fig58.JPG |
| Figure 6-26. An enum datatype conversion |

The library also allows conversion from enumeration to a numeric datatype. A numeric datatype is either an integer or a floating-point number. This conversion can simplify the application program because the base type for an enumeration datatype is an integer datatype. The application program can read the data from a dataset of enumeration datatype in file into a memory buffer of numeric datatype. And it can write enumeration data from memory into a dataset of numeric datatype in file, too.

For compound datatypes, each field of the source and destination datatype is converted according to its type. The name of the fields must be the same in the source and the destination in order for the data to be converted.

The example below shows the compound datatypes shows sample code to create a compound datatype with the fields aligned on word boundaries (s1\_tid) and with the fields packed (s2\_tid). The former is suit­able as a description of the storage layout in memory, the latter would give a more compact store on disk. These types can be used for transferring data, with s2\_tid used to create the dataset, and s1\_tid used as the memory datatype.

|  |
| --- |
| typedef struct s1\_t {     int a;     char b;     double c;  } s1\_t;    s1\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(s1\_t));    H5Tinsert(s1\_tid, “a\_name”, HOFFSET(s1\_t, a), H5T\_NATIVE\_INT);  H5Tinsert(s1\_tid, “b\_name”, HOFFSET(s1\_t, b), H5T\_NATIVE\_CHAR);  H5Tinsert(s1\_tid, “c\_name”, HOFFSET(s1\_t, c), H5T\_NATIVE\_DOUBLE);    s2\_tid = H5Tcopy(s1\_tid);  H5Tpack(s2\_tid); |
| Code Example 6-44. Create an aligned and packed compound datatype |

When the data is transferred, the fields within each data element will be aligned according to the datatype specification. The figure below shows how one data element would be aligned in memory and on disk. Note that the size and byte order of the elements might also be converted during the transfer.

It is also possible to transfer some of the fields of compound datatypes. Based on the example above, the example below shows a compound datatype that selects the first and third fields of the s1\_tid. The sec­ond datatype can be used as the memory datatype, in which case data is read from or written to these two fields, while skipping the middle field. The second figure below shows the layout for two data elements.

|  |
| --- |
| Dtypes_fig60.JPG |
| Figure 6-27. Alignment of a compound datatype |

|  |
| --- |
| typedef struct s1\_t {     int a;     char b;     double c;  } s1\_t;    typedef struct s2\_t {   /\* two fields from s1\_t \*/     int a;     double c;  } s2\_t; |
| s1\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(s1\_t));    H5Tinsert(s1\_tid, “a\_name”, HOFFSET(s1\_t, a), H5T\_NATIVE\_INT);  H5Tinsert(s1\_tid, “b\_name”, HOFFSET(s1\_t, b), H5T\_NATIVE\_CHAR);  H5Tinsert(s1\_tid, “c\_name”, HOFFSET(s1\_t, c), H5T\_NATIVE\_DOUBLE);    s2\_tid = H5Tcreate (H5T\_COMPOUND, sizeof(s2\_t));    H5Tinsert(s1\_tid, “a\_name”, HOFFSET(s2\_t, a), H5T\_NATIVE\_INT);  H5Tinsert(s1\_tid, “c\_name”, HOFFSET(s2\_t, c), H5T\_NATIVE\_DOUBLE); |
| Code Example 6-45. Transfer some fields of a compound datatype |

|  |
| --- |
| Dtypes_fig62.JPG |
| Figure 6-28. Layout when an element is skipped |

## Text Descriptions of Datatypes: Conversion to and from

HDF5 provides a means for generating a portable and human-readable text description of a datatype and for generating a datatype from such a text description. This capability is particularly useful for creating complex datatypes in a single step, for creating a text description of a datatype for debugging purposes, and for creating a portable datatype definition that can then be used to recreate the datatype on many platforms or in other applications.

These tasks are handled by two functions provided in the HDF5 Lite high-level library:

* H5LTtext\_to\_dtype         Creates an HDF5 datatype in a single step.
* H5LTdtype\_to\_text         Translates an HDF5 datatype into a text description.

Note that this functionality requires that the HDF5 High-Level Library (H5LT) be installed.

While H5LTtext\_to\_dtype can be used to generate any sort of datatype, it is particularly useful for complex datatypes.

H5LTdtype\_to\_text is most likely to be used in two sorts of situations: when a datatype must be closely examined for debugging purpose or to create a portable text description of the datatype that can then be used to recreate the datatype on other platforms or in other applications.

These two functions work for all valid HDF5 datatypes except time, bitfield, and reference datatypes.

The currently supported text format used by H5LTtext\_to\_dtype and H5LTdtype\_to\_text is the data description language (DDL) and conforms to the HDF5 DDL. The portion of the HDF5 DDL that defines HDF5 datatypes appears below.

|  |
| --- |
| <datatype> ::= <atomic\_type> | <compound\_type> |                       <array\_type> | <variable\_length\_type>    <atomic\_type> ::= <integer> | <float> | <time> |                          <string> | <bitfield> | <opaque> |                          <reference> | <enum> |
| <integer> ::= H5T\_STD\_I8BE | H5T\_STD\_I8LE |                       H5T\_STD\_I16BE | H5T\_STD\_I16LE |                       H5T\_STD\_I32BE | H5T\_STD\_I32LE |                       H5T\_STD\_I64BE | H5T\_STD\_I64LE |                       H5T\_STD\_U8BE | H5T\_STD\_U8LE |                       H5T\_STD\_U16BE | H5T\_STD\_U16LE |                       H5T\_STD\_U32BE | H5T\_STD\_U32LE |                       H5T\_STD\_U64BE | H5T\_STD\_U64LE |                       H5T\_NATIVE\_CHAR | H5T\_NATIVE\_UCHAR |                       H5T\_NATIVE\_SHORT | H5T\_NATIVE\_USHORT |                       H5T\_NATIVE\_INT | H5T\_NATIVE\_UINT |                       H5T\_NATIVE\_LONG | H5T\_NATIVE\_ULONG |                       H5T\_NATIVE\_LLONG | H5T\_NATIVE\_ULLONG |
| <float> ::= H5T\_IEEE\_F32BE | H5T\_IEEE\_F32LE |                 H5T\_IEEE\_F64BE | H5T\_IEEE\_F64LE |                 H5T\_NATIVE\_FLOAT | H5T\_NATIVE\_DOUBLE |                 H5T\_NATIVE\_LDOUBLE    <time> ::= TBD |
| <string> ::= H5T\_STRING { STRSIZE <strsize> ;     STRPAD <strpad> ;     CSET <cset> ;     CTYPE <ctype> ;}    <strsize> ::= <int\_value> | H5T\_VARIABLE  <strpad> ::= H5T\_STR\_NULLTERM | H5T\_STR\_NULLPAD |                    H5T\_STR\_SPACEPAD  <cset> ::= H5T\_CSET\_ASCII | H5T\_CSET\_UTF8  <ctype> ::= H5T\_C\_S1 | H5T\_FORTRAN\_S1 |
| <bitfield> ::= TBD    <opaque> ::= H5T\_OPAQUE { OPQ\_SIZE <opq\_size>;  OPQ\_TAG <opq\_tag>; }  opq\_size ::= <int\_value>  opq\_tag ::= "<string>"    <reference> ::= Not supported |
| <compound\_type> ::= H5T\_COMPOUND { <member\_type\_def>+ }  <member\_type\_def> ::= <datatype> <field\_name> <offset>opt;  <field\_name> ::= "<identifier>"  <offset> ::= : <int\_value>    <variable\_length\_type> ::= H5T\_VLEN { <datatype> } |
| <array\_type> ::= H5T\_ARRAY { <dim\_sizes> <datatype> }  <dim\_sizes> ::= [<dimsize>] | [<dimsize>] <dim\_sizes>  <dimsize> ::= <int\_value>    <enum> ::= H5T\_ENUM { <enum\_base\_type>; <enum\_def>+ }  <enum\_base\_type> ::= <integer>  // Currently enums can only hold integer type data, but  // they may be expanded in the future to hold any  // datatype  <enum\_def> ::= <enum\_symbol> <enum\_val>;  <enum\_symbol> ::= "<identifier>"  <enum\_val> ::= <int\_value> |
| Code Example 6-46. The definition of HDF5 datatypes from the HDF5 DDL |

The definitions of opaque and compound datatype above are revised for HDF5 Release 1.8. In Release 1.6.5. and earlier, they were defined as follows:

|  |
| --- |
| <opaque> ::= H5T\_OPAQUE { <identifier> }      <compound\_type> ::= H5T\_COMPOUND { <member\_type\_def>+ }  <member\_type\_def> ::= <datatype> <field\_name> ;  <field\_name> ::= <identifier> |
| Code Example 6-47. Old definitions of the opaque and compound datatypes |

##### Examples

The code sample below illustrates the use of H5LTtext\_to\_dtype to generate a variable-length string datatype.

|  |
| --- |
| hid\_t   dtype;  if((dtype = H5LTtext\_to\_dtype(“H5T\_STRING {                                                        STRSIZE H5T\_VARIABLE;                                                        STRPAD H5T\_STR\_NULLPAD;                                                        CSET H5T\_CSET\_ASCII;                                                        CTYPE H5T\_C\_S1;                                                        }”, H5LT\_DDL))<0)  goto out; |
| Code Example 6-48. Creating a variable-length string datatype from a text description |

The code sample below illustrates the use of H5LTtext\_to\_dtype to generate a complex array data­type.

|  |
| --- |
| hid\_t   dtype;  if((dtype = H5LTtext\_to\_dtype(“H5T\_ARRAY { [5][7][13] H5T\_ARRAY                                                        { [17][19] H5T\_COMPOUND                                                           {                                                              H5T\_STD\_I8BE                                                              \“arr\_compound\_1\”;                                                              H5T\_STD\_I32BE                                                              \“arr\_compound\_2\”;                                                           }                                                        }                                                     }”, H5LT\_DDL))<0)  goto out; |
| Code Example 6-49. Creating a complex array datatype from a text description |

# HDF5 Dataspaces and Partial I/O

## Introduction

The HDF5 dataspace is a required component of an HDF5 dataset or attribute definition. The dataspace defines the size and shape of the dataset or attribute raw data. In other words, a dataspace defines the number of dimensions and the size of each dimension of the multidimensional array in which the raw data is represented. The dataspace must be defined when the dataset or attribute is created.

The dataspace is also used during dataset I/O operations, defining the elements of the dataset that partic­ipate in the I/O operation.

This chapter explains the dataspace object and its use in dataset and attribute creation and data transfer. It also describes selection operations on a dataspace used to implement sub-setting, sub-sampling, and scatter-gather access to datasets.

## Dataspace (H5S) Function Summaries

This section provides a reference list of dataspace functions, the H5S APIs, with brief descriptions. The functions are presented in the following categories:

* Dataspace management functions
* Dataspace query functions
* Dataspace selection functions: hyperslabs
* Dataspace selection functions: points

The rest of the chapter will provide examples and explanations of how to use these functions.

|  |  |
| --- | --- |
| Function Listing 7-1. Dataspace management functions | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Screate  h5screate\_f | Creates a new dataspace of a specified type. |
| H5Scopy  h5scopy\_f | Creates an exact copy of a dataspace. |
| H5Sclose  h5sclose\_f | Releases and terminates access to a dataspace. |
| H5Sdecode  h5sdecode\_f | Decode a binary object description of a dataspace and return a new object identifier. |
| H5Sencode  h5sencode | Encode a dataspace object description into a binary buffer. |
| H5Screate\_simple  h5screate\_simple\_f | Creates a new simple dataspace and opens it for access. |
| H5Sis\_simple  h5sis\_simple\_f | Determines whether a dataspace is a simple dataspace. |
| H5Sextent\_copy  h5sextent\_copy\_f | Copies the extent of a dataspace. |
| H5Sextent\_equal  h5sextent\_equal\_f | Determines whether two dataspace extents are equal. |
| H5Sset\_extent\_simple  h5sset\_extent\_simple\_f | Sets or resets the size of an existing dataspace. |
| H5Sset\_extent\_none  h5sset\_extent\_none\_f | Removes the extent from a dataspace. |

|  |  |
| --- | --- |
| Function Listing 7-2. Dataspace query functions | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Sget\_simple\_extent\_dims  h5sget\_simple\_extent\_dims\_f | Retrieves dataspace dimension size and maxi­mum size. |
| H5Sget\_simple\_extent\_ndims  h5sget\_simple\_extent\_ndims\_f | Determines the dimensionality of a dataspace. |
| H5Sget\_simple\_extent\_npoints  h5sget\_simple\_extent\_npoints\_f | Determines the number of elements in a dataspace. |
| H5Sget\_simple\_extent\_type  h5sget\_simple\_extent\_type\_f | Determine the current class of a dataspace. |

|  |  |
| --- | --- |
| Function Listing 7-3. Dataspace selection functions: hyperslabs | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Soffset\_simple  h5soffset\_simple\_f | Sets the offset of a simple dataspace. |
| H5Sget\_select\_type  h5sget\_select\_type\_f | Determines the type of the dataspace selec­tion. |
| H5Sget\_select\_hyper\_nblocks  h5sget\_select\_hyper\_nblocks\_f | Get number of hyperslab blocks. |
| H5Sget\_select\_hyper\_blocklist  h5sget\_select\_hyper\_blocklist\_f | Gets the list of hyperslab blocks currently selected. |
| H5Sget\_select\_bounds  h5sget\_select\_bounds\_f | Gets the bounding box containing the current selection. |
| H5Sselect\_all  h5sselect\_all\_f | Selects the entire dataspace. |
| H5Sselect\_none  h5sselect\_none\_f | Resets the selection region to include no ele­ments. |
| H5Sselect\_valid  h5sselect\_valid\_f | Verifies that the selection is within the extent of the dataspace. |
| H5Sselect\_hyperslab  h5sselect\_hyperslab\_f | Selects a hyperslab region to add to the cur­rent selected region. |

|  |  |
| --- | --- |
| Function Listing 7-4. Dataspace selection functions: points | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Sget\_select\_npoints  h5sget\_select\_npoints\_f | Determines the number of elements in a dataspace selection. |
| H5Sget\_select\_elem\_npoints  h5sget\_select\_elem\_npoints\_f | Gets the number of element points in the cur­rent selection. |
| H5Sget\_select\_elem\_pointlist  h5sget\_select\_elem\_pointlist\_f | Gets the list of element points currently selected. |
| H5Sselect\_elements  h5sselect\_elements\_f | Selects array elements to be included in the selection for a dataspace. |

## Definition of Dataspace Objects and the Dataspace Programming Model

This section introduces the notion of the HDF5 dataspace object and a programming model for creating and working with dataspaces.

### Dataspace Objects

An HDF5 dataspace is a required component of an HDF5 dataset or attribute. A dataspace defines the size and the shape of a dataset’s or an attribute’s raw data. Currently, HDF5 supports the following types of the dataspaces:

* Scalar dataspaces
* Simple dataspaces
* Null dataspaces

A scalar dataspace, H5S\_SCALAR, represents just one element, a scalar. Note that the datatype of this one element may be very complex; example would be a compound structure with members being of any allowed HDF5 datatype, including multidimensional arrays, strings, and nested compound structures. By convention, the rank of a scalar dataspace is always 0 (zero); think of it geometrically as a single, dimen­sionless point, though that point may be complex.

A simple dataspace, H5S\_SIMPLE, is a multidimensional array of elements. The dimensionality of the dataspace (or the rank of the array) is fixed and is defined at creation time. The size of each dimension can grow during the life time of the dataspace from the current size up to the maximum size. Both the current size and the maximum size are specified at creation time. The sizes of dimensions at any particular time in the life of a dataspace are called the current dimensions, or the dataspace extent. They can be queried along with the maximum sizes.

A null dataspace, H5S\_NULL, contains no data elements. Note that no selections can be applied to a null dataset as there is nothing to select.

As shown in the UML diagram in the figure below, an HDF5 simple dataspace object has three attributes: the rank or number of dimensions; the current sizes, expressed as an array of length rank with each ele­ment of the array denoting the current size of the corresponding dimension; and the maximum sizes, expressed as an array of length rank with each element of the array denoting the maximum size of the corresponding dimension.

|  |
| --- |
| Dspace_fig1new.JPG |
| Figure 7-1. A simple dataspace |

Note: A simple dataspace is defined by its rank, the current size of each dimension, and the maximum size of each dimension.

The size of a current dimension cannot be greater than the maximum size, which can be unlimited, speci­fied as H5S\_UNLIMITED. Note that while the HDF5 file format and library impose no maximum size on an unlimited dimension, practically speaking its size will always be limited to the biggest integer available on the particular system being used.

Dataspace rank is restricted to 32, the standard limit in C on the rank of an array, in the current implemen­tation of the HDF5 library. The HDF5 file format, on the other hand, allows any rank up to the maximum integer value on the system, so the library restriction can be raised in the future if higher dimensionality is required.

Note that most of the time Fortran applications calling HDF5 will work with dataspaces of rank less than or equal to seven, since seven is the maximum number of dimensions in a Fortran array. But dataspace rank is not limited to seven for Fortran applications.

The current dimensions of a dataspace, also referred to as the dataspace extent, define the bounding box for dataset elements that can participate in I/O operations.

### Dataspace Programming Model

The programming model for creating and working with HDF5 dataspaces can be summarized as follows:

1. Create a dataspace
2. Use the dataspace to create a dataset in the file or to describe a data array in memory
3. Modify the dataspace to define dataset elements that will participate in I/O operations
4. Use the modified dataspace while reading/writing dataset raw data or to create a region refer­ence
5. Close the dataspace when no longer needed

The rest of this section will address steps 1, 2, and 5 of the programming model; steps 3 and 4 will be dis­cussed in later sections of this chapter.

#### 7.3.2.1. Creating a Dataspace

A dataspace can be created by calling the H5Screate function (h5screate\_f in Fortran). Since the defi­nition of a simple dataspace requires the specification of dimensionality (or rank) and initial and maximum dimension sizes, the HDF5 Library provides a convenience API, H5Screate\_simple (h5screate\_sim­ple\_f) to create a simple dataspace in one step.

The following examples illustrate the usage of these APIs.

#### 7.3.2.2. Creating a Scalar Dataspace

A scalar dataspace is created with the H5Screate or the h5screate\_f function.

In C:

hid\_t space\_id;

. . .

space\_id = H5Screate(H5S\_SCALAR);

In Fortran:

INTEGER(HID\_T) :: space\_id

. . .

CALL h5screate\_f(H5S\_SCALAR\_F, space\_id, error)

As mentioned above, the dataspace will contain only one element. Scalar dataspaces are used more often for describing attributes that have just one value. For example, the attribute temperature with the value Celsius is used to indicate that the dataset with this attribute stores temperature values using the Cel­sius scale.

#### 7.3.2.3. Creating a Null Dataspace

A null dataspace is created with the H5Screate or the h5screate\_f function.

In C:

hid\_t space\_id;

. . .

space\_id = H5Screate(H5S\_NULL);

In Fortran:

    (H5S\_NULL not yet implemented in Fortran.)

INTEGER(HID\_T) :: space\_id

. . .

CALL h5screate\_f(H5S\_NULL\_F, space\_id, error)

As mentioned above, the dataspace will contain no elements.

#### 7.3.2.4. Creating a Simple Dataspace

Let’s assume that an application wants to store a two-dimensional array of data, A(20,100). During the life of the application, the first dimension of the array can grow up to 30; there is no restriction on the size of the second dimension. The following steps are used to declare a dataspace for the dataset in which the array data will be stored.

In C:

hid\_t space\_id;

int rank = 2;

hsize\_t current\_dims[2] = {20, 100};

hsize\_t max\_dims[2] = {30, H5S\_UNLIMITED};

. . .

space\_id = H5Screate(H5S\_SIMPLE);

H5Sset\_extent\_simple(space\_id,rank,current\_dims,max\_dims);

In Fortran:

INTEGER(HID\_T) :: space\_id

INTEGER :: rank = 2

INTEGER(HSIZE\_T) :: current dims = /( 20, 100)/

INTEGER(HSIZE\_T) :: max\_dims = /(30, H5S\_UNLIMITED\_F)/

INTEGER error

. . .

CALL h5screate\_f(H5S\_SIMPLE\_F, space\_id, error)

CALL h5sset\_extent\_simple\_f(space\_id, rank, current\_dims, max\_dims, error)

Alternatively, the convenience APIs H5Screate\_simple/h5screate\_simple\_f can replace the H5Screate/h5screate\_f and H5Sset\_extent\_simple/h5sset\_extent\_simple\_f calls.

In C:

space\_id = H5Screate\_simple(rank, current\_dims, max\_dims);

In Fortran:

CALL h5screate\_simple\_f(rank, current\_dims, space\_id, error, max\_dims)

In this example, a dataspace with current dimensions of 20 by 100 is created. The first dimension can be extended only up to 30. The second dimension, however, is declared unlimited; it can be extended up to the largest available integer value on the system.

Note that when there is a difference between the current dimensions and the maximum dimensions of an array, then chunking storage must be used. In other words, if the number of dimensions may change over the life of the dataset, then chunking must be used. If the array dimensions are fixed (if the number of cur­rent dimensions is equal to the maximum number of dimensions when the dataset is created), then con­tiguous storage can be used. For more information, see "Data Transfer."

Maximum dimensions can be the same as current dimensions. In such a case, the sizes of dimensions can­not be changed during the life of the dataspace object. In C, NULL can be used to indicate to the H5Scre­ate\_simple and H5Sset\_extent\_simple functions that the maximum sizes of all dimensions are the same as the current sizes. In Fortran, the maximum size parameter is optional for h5screate\_simple\_f and can be omitted when the sizes are the same.

In C:

space\_id = H5Screate\_simple(rank, current\_dims, NULL);

In Fortran:

CALL h5screate\_f(rank, current\_dims, space\_id, error)

The created dataspace will have current and maximum dimensions of 20 and 100 correspondingly, and the sizes of those dimensions cannot be changed.

#### 7.3.2.5. C versus Fortran Dataspaces

Dataspace dimensions are numbered from 1 to rank. HDF5 uses C storage conventions, assuming that the last listed dimension is the fastest-changing dimension and the first-listed dimension is the slowest chang­ing. The HDF5 file format storage layout specification adheres to the C convention and the HDF5 Library adheres to the same convention when storing dataspace dimensions in the file. This affects how C pro­grams and tools interpret data written from Fortran programs and vice versa. The example below illus­trates the issue.

When a Fortran application describes a dataspace to store an array as A(20,100), it specifies the value of the first dimension to be 20 and the second to be 100. Since Fortran stores data by columns, the first-listed dimension with the value 20 is the fastest-changing dimension and the last-listed dimension with the value 100 is the slowest-changing. In order to adhere to the HDF5 storage convention, the HDF5 Fortran wrapper transposes dimensions, so the first dimension becomes the last. The dataspace dimensions stored in the file will be 100,20 instead of 20,100 in order to correctly describe the Fortran data that is stored in 100 columns, each containing 20 elements.

When a Fortran application reads the data back, the HDF5 Fortran wrapper transposes the dimensions once more, returning the first dimension to be 20 and the second to be 100, describing correctly the sizes of the array that should be used to read data in the Fortran array A(20,100).

When a C application reads data back, the dimensions will come out as 100 and 20, correctly describing the size of the array to read data into, since the data was written as 100 records of 20 elements each. Therefore C tools such as h5dump and h5ls always display transposed dimensions and values for the data written by a Fortran application.

Consider the following simple example of equivalent C 3 x 5 and Fortran 5 x 3 arrays. As illustrated in the figure below, a C application will store a 3 x 5 2-dimensional array as three 5-element rows. In order to store the same data in the same order, a Fortran application must view the array as a 5 x 3 array with three 5-element columns. The dataspace of this dataset, as written from Fortran, will therefore be described as 5 x 3 in the application but stored and described in the file according to the C convention as a 3 x 5 array. This ensures that C and Fortran applications will always read the data in the order in which it was written. The HDF5 Fortran interface handles this transposition automatically.

In C (from h5\_write.c):

#define NX     3                      /\* dataset dimensions \*/

#define NY     5

. . .

int         data[NX][NY];          /\* data to write \*/

. . .

/\*

\* Data  and output buffer initialization.

\*/

for (j = 0; j < NX; j++) {

         for (i = 0; i < NY; i++)

            data[j][i] = i + 1 + j\*NY;

}

/\*

\*  1  2  3  4  5

\*  6  7  8  9 10

\* 11 12 13 14 15

\*/

. . .

dims[0] = NX;

dims[1] = NY;

dataspace = H5Screate\_simple(RANK, dims, NULL);

For more information, see "h5\_write.c."

In Fortran (from h5\_write.f90):

INTEGER, PARAMETER :: NX = 3

INTEGER, PARAMETER :: NY = 5

. . .

INTEGER(HSIZE\_T), DIMENSION(2) :: dims = (/3,5/) ! Dataset dimensions

---

INTEGER     ::    data(NX,NY)

. . .

!

! Initialize data

!

         do i = 1, NX

            do j = 1, NY

            data(i,j) = j + (i-1)\*NY

          enddo

enddo

!

! Data

!

!  1  2  3  4  5

!  6  7  8  9 10

! 11 12 13 14 15

. . .

CALL h5screate\_simple\_f(rank, dims, dspace\_id, error)

For more information, see "h5\_write.f90."

In Fortran (from h5\_write\_tr.f90):

INTEGER, PARAMETER :: NX = 3

INTEGER, PARAMETER :: NY = 5

. . .

INTEGER(HSIZE\_T), DIMENSION(2) :: dims = (/NY, NX/) ! Dataset dimensions

. . .

!

! Initialize data

!

do i = 1, NY

         do j = 1, NX

            data(i,j) = i + (j-1)\*NY

         enddo

enddo

!

! Data

!

!  1  6  11

!  2  7  12

!  3  8  13

!  4  9  14

!  5 10  15

. . .

CALL h5screate\_simple\_f(rank, dims, dspace\_id, error)

For more information, see "h5\_write\_tr.f90."

|  |
| --- |
| A dataset stored by a C program in a 3 x 5 array: |
| Dspace_fig2a.JPG |
| The same dataset stored by a Fortran program in a 5 x 3 array: |
| Dspace_fig2b.JPG |
| The first dataset above as written to an HDF5 file from C or the second dataset above as written from Fortran: |
| Dspace_fig2c.JPG |
| The first dataset above as written to an HDF5 file from Fortran: |
| Dspace_fig2d.JPG |
| Figure 7-2. Comparing C and Fortran dataspaces |

Note: The HDF5 library stores arrays along the fastest-changing dimension. This approach is often referred to as being “in C order.” C, C++, and Java work with arrays in row-major order. In other words, the row, or the last dimen­sion, is the fastest-changing dimension. Fortran, on the other hand, handles arrays in column-major order making the column, or the first dimension, the fastest-changing dimension. Therefore, the C and Fortran arrays illustrated in the top portion of this figure are stored identically in an HDF5 file. This ensures that data written by any language can be meaningfully read, interpreted, and manipulated by any other.

#### 7.3.2.6. Finding Dataspace Characteristics

The HDF5 library provides several APIs designed to query the characteristics of a dataspace.

The function H5Sis\_simple (h5sis\_simple\_f) returns information about the type of a dataspace. This function is rarely used and currently supports only simple and scalar dataspaces.

To find out the dimensionality, or rank, of a dataspace, use H5Sget\_simple\_extent\_ndims (h5sget\_­simple\_extent\_ndims\_f). H5Sget\_simple\_extent\_dims can also be used to find out the rank. See the example below. If both functions return 0 for the value of rank, then the dataspace is scalar.

To query the sizes of the current and maximum dimensions, use H5Sget\_simple\_extent\_dims (h5sget\_simple\_extent\_dims\_f).

The following example illustrates querying the rank and dimensions of a dataspace using these functions.

In C:

hid\_t space\_id;

int rank;

hsize\_t  \*current\_dims;

hsize\_t  \*max\_dims;

---------

rank=H5Sget\_simple\_extent\_ndims(space\_id);

(or rank=H5Sget\_simple\_extent\_dims(space\_id, NULL, NULL);)

current\_dims= (hsize\_t)malloc(rank\*sizeof(hsize\_t));

max\_dims=(hsize\_t)malloc(rank\*sizeof(hsize\_t));

H5Sget\_simple\_extent\_dims(space\_id, current\_dims, max\_dims);

Print values here for the previous example

## Dataspaces and Data Transfer

Read and write operations transfer data between an HDF5 file on disk and in memory. The shape that the array data takes in the file and in memory may be the same, but HDF5 also allows users the ability to rep­resent data in memory in a different shape than in the file. If the shape of an array in the file and in mem­ory will be the same, then the same dataspace definition can be used for both. If the shape of an array in memory needs to be different than the shape in the file, then the dataspace definition for the shape of the array in memory can be changed. During a read operation, the array will be read into the different shape in memory, and during a write operation, the array will be written to the file in the shape specified by the dataspace in the file. The only qualification is that the number of elements read or written must be the same in both the source and the destination dataspaces.

Item a in the figure below shows a simple example of a read operation in which the data is stored as a 3 by 4 array in the file (item b) on disk, but the program wants it to be a 4 by 3 array in memory. This is accom­plished by setting the memory dataspace to describe the desired memory layout, as in item c. The read operation reads the data in the file array into the memory array.

|  |
| --- |
| Dspace_fig4.JPG |
| Figure 7-3. Data layout before and after a read operation |

|  |
| --- |
| Dspace_fig5.JPG |
| Figure 7-4. Moving data from disk to memory |

Both the source and destination are stored as contiguous blocks of storage with the elements in the order specified by the dataspace. The figure above shows one way the elements might be organized. In item a, the elements are stored as 3 blocks of 4 elements. The destination is an array of 12 elements in memory (see item c). As the figure suggests, the transfer reads the disk blocks into a memory buffer (see item b), and then writes the elements to the correct locations in memory. A similar process occurs in reverse when data is written to disk.

### Data Selection

In addition to rearranging data, the transfer may select the data elements from the source and destina­tion.

Data selection is implemented by creating a dataspace object that describes the selected elements (within the hyper rectangle) rather than the whole array. Two dataspace objects with selections can be used in data transfers to read selected elements from the source and write selected elements to the destination. When data is transferred using the dataspace object, only the selected elements will be transferred.

This can be used to implement partial I/O, including:

* Sub-setting - reading part of a large dataset
* Sampling - reading selected elements (for example, every second element) of a dataset
* Scatter-gather - read non-contiguous elements into contiguous locations (gather) or read contigu­ous elements into non-contiguous locations (scatter) or both

To use selections, the following steps are followed:

1. Get or define the dataspace for the source and destination
2. Specify one or more selections for source and destination dataspaces
3. Transfer data using the dataspaces with selections

A selection is created by applying one or more selections to a dataspace. A selection may override any other selections (H5T\_SELECT\_SET) or may be “Ored” with previous selections on the same dataspace (H5T\_SELECT\_OR). In the latter case, the resulting selection is the union of the selection and all previ­ously selected selections. Arbitrary sets of points from a dataspace can be selected by specifying an appro­priate set of selections.

Two selections are used in data transfer, so the source and destination must be compatible, as described below.

There are two forms of selection, hyperslab and point. A selection must be either a point selection or a set of hyperslab selections. Selections cannot be mixed.

The definition of a selection within a dataspace, not the data in the selection, cannot be saved to the file unless the selection definition is saved as a region reference. For more information, see "References to Dataset Regions."

#### 7.4.1.1. Hyperslab Selection

A hyperslab is a selection of elements from a hyper rectangle. An HDF5 hyperslab is a rectangular pattern defined by four arrays. The four arrays are summarized in the table below.

The offset defines the origin of the hyperslab in the original dataspace.

The stride is the number of elements to increment between selected elements. A stride of ‘1’ is every ele­ment, a stride of ‘2’ is every second element, etc. Note that there may be a different stride for each dimen­sion of the dataspace. The default stride is 1.

The count is the number of elements in the hyperslab selection. When the stride is 1, the selection is a hyper rectangle with a corner at the offset and size count[0] by count[1] by.... When stride is greater than one, the hyperslab bounded by the offset and the corners defined by stride[n] \* count[n].

|  |  |
| --- | --- |
| Table 7-1. Hyperslab elements | |
| **Parameter** | **Description** |
| Offset | The starting location for the hyperslab. |
| Stride | The number of elements to separate each element or block to be selected. |
| Count | The number of elements or blocks to select along each dimension. |
| Block | The size of the block selected from the dataspace. |

The block is a count on the number of repetitions of the hyperslab. The default block size is ‘1’, which is one hyperslab. A block of 2 would be two hyperslabs in that dimension, with the second starting at off­set[n]+ (count[n] \* stride[n]) + 1.

A hyperslab can be used to access a sub-set of a large dataset. The figure below shows an example of a hyperslab that reads a rectangle from the middle of a larger two dimensional array. The destination is the same shape as the source.

|  |
| --- |
| Dspace_fig6.JPG |
| Figure 7-5. Access a sub-set of data with a hyperslab |

Hyperslabs can be combined to select complex regions of the source and destination. The figure below shows an example of a transfer from one non-rectangular region into another non-rectangular region. The source is defined as the union of two hyperslabs, and the destination is the union of three hyperslabs.

|  |
| --- |
| Dspace_fig7.JPG |
| Figure 7-6. Build complex regions with hyperslab unions |

Hyperslabs may also be used to collect or scatter data from regular patterns. The figure below shows an example where the source is a repeating pattern of blocks, and the destination is a single, one dimensional array.

|  |
| --- |
| Dspace_fig8.JPG |
| Figure 7-7. Use hyperslabs to combine or disperse data |

#### 7.4.1.2. Select Points

The second type of selection is an array of points such as coordinates. Essentially, this selection is a list of all the points to include. The figure below shows an example of a transfer of seven elements from a two dimensional dataspace to a three dimensional dataspace using a point selection to specify the points.

|  |
| --- |
| Dspace_fig9.JPG |
| Figure 7-8. Point selection |

#### 7.4.1.3. Rules for Defining Selections

A selection must have the same number of dimensions (rank) as the dataspace it is applied to, although it may select from only a small region such as a plane from a 3D dataspace. Selections do not affect the extent of the dataspace, the selection may be larger than the dataspace. The boundaries of selections are reconciled with the extent at the time of the data transfer.

#### 7.4.1.4. Data Transfer with Selections

A data transfer (read or write) with selections is the same as any read or write, except the source and des­tination dataspace have compatible selections.

During the data transfer, the following steps are executed by the library:

* The source and destination dataspaces are checked to assure that the selections are compatible.
* Each selection must be within the current extent of the dataspace. A selection may be defined to extend outside the current extent of the dataspace, but the dataspace cannot be accessed if the selection is not valid at the time of the access.
* The total number of points selected in the source and destination must be the same. Note that the dimensionality of the source and destination can be different (for example, the source could be 2D, the destination 1D or 3D), and the shape can be different, but the num­ber of elements selected must be the same.
* The data is transferred, element by element.

Selections have an iteration order for the points selected, which can be any permutation of the dimen­sions involved (defaulting to ‘C’ array order) or a specific order for the selected points, for selections com­posed of single array elements with H5Sselect\_elements.

The elements of the selections are transferred in row-major, or C order. That is, it is assumed that the first dimension varies slowest, the second next slowest, and so forth. For hyperslab selections, the order can be any permutation of the dimensions involved (defaulting to ‘C’ array order). When multiple hyperslabs are combined, the hyperslabs are coalesced into contiguous reads and writes.

In the case of point selections, the points are read and written in the order specified.

### Programming Model

#### Selecting Hyperslabs

Suppose we want to read a 3x4 hyperslab from a dataset in a file beginning at the element <1,2> in the dataset, and read it into a 7 x 7 x 3 array in memory. See the figure below. In order to do this, we must cre­ate a dataspace that describes the overall rank and dimensions of the dataset in the file as well as the posi­tion and size of the hyperslab that we are extracting from that dataset.

|  |
| --- |
| Dspace_fig10.JPG |
| Figure 7-9. Selecting a hyperslab |

The code in the first example below illustrates the selection of the hyperslab in the file dataspace. The sec­ond example below shows the definition of the destination dataspace in memory. Since the in-memory dataspace has three dimensions, the hyperslab is an array with three dimensions with the last dimension being 1: <3,4,1>. The third example below shows the read using the source and destination dataspaces with selections.

|  |
| --- |
| /\*  \* get the file dataspace.  \*/  dataspace = H5Dget\_space(dataset); /\* dataspace \*/                                                     /\* identifier \*/    /\*  \* Define hyperslab in the dataset.  \*/ |
| offset[0] = 1;  offset[1] = 2;  count[0]  = 3;  count[1]  = 4;  status = H5Sselect\_hyperslab(dataspace, H5S\_SELECT\_SET, offset, NULL, count, NULL); |
| Code Example 7-1. Selecting a hyperslab |

|  |
| --- |
| /\*  \* Define memory dataspace.  \*/  dimsm[0] = 7;  dimsm[1] = 7;  dimsm[2] = 3;  memspace = H5Screate\_simple(3,dimsm,NULL); |
| /\*  \* Define memory hyperslab.  \*/  offset\_out[0] = 3;  offset\_out[1] = 0;  offset\_out[2] = 0;  count\_out[0]  = 3;  count\_out[1]  = 4;  count\_out[2]  = 1;  status = H5Sselect\_hyperslab(memspace, H5S\_SELECT\_SET, offset\_out, NULL, count\_out, NULL); |
| Code Example 7-2. Defining the destination memory |

|  |
| --- |
| ret = H5Dread(dataset, H5T\_NATIVE\_INT, memspace, dataspace, H5P\_DEFAULT, data); |
| Code Example 7-3. A sample read specifying source and destination dataspaces |

#### 7.4.2.2. Example with Strides and Blocks

Consider an 8 x 12 dataspace into which we want to write eight 3 x 2 blocks in a two dimensional array from a source dataspace in memory that is a 50-element one dimensional array. See the figure below.

|  |
| --- |
| Dspace_fig14.JPG |
| Figure 7-10. Write from a one dimensional array to a two dimensional array |

The example below shows code to write 48 elements from the one dimensional array to the file dataset starting with the second element in vector. The destination hyperslab has the following parameters: off­set=(0,1), stride=(4,3), count=(2,4), block=(3,2). The source has the parameters: offset=(1), stride=(1), count=(48), block=(1). After these operations, the file dataspace will have the values shown in item b in the figure above. Notice that the values are inserted in the file dataset in row-major order.

|  |
| --- |
| /\* Select hyperslab for the dataset in the file, using  \* 3 x 2 blocks, (4,3) stride (2,4) count starting at  \* the position (0,1).  \*/ |
| offset[0]  = 0; offset[1]  = 1;  stride[0] = 4; stride[1] = 3;  count[0]  = 2; count[1]  = 4;  block[0]  = 3; block[1]  = 2;  ret = H5Sselect\_hyperslab(fid, H5S\_SELECT\_SET, offset, stride, count, block); |
| /\*  \* Create dataspace for the first dataset.  \*/  mid1 = H5Screate\_simple(MSPACE1\_RANK, dim1, NULL);    /\*  \* Select hyperslab.  \* We will use 48 elements of the vector buffer starting  \* at the second element. Selected elements are  \* 1 2 3 . . . 48  \*/ |
| offset[0]  = 1;  stride[0] = 1;  count[0]  = 48;  block[0]  = 1;  ret = H5Sselect\_hyperslab(mid1, H5S\_SELECT\_SET, offset, stride, count, block);    /\*  \* Write selection from the vector buffer to the dataset  \* in the file.  \*/  ret = H5Dwrite(dataset, H5T\_NATIVE\_INT, midd1, fid, H5P\_DEFAULT, vector) |
| Code Example 7-4. Write from a one dimensional array to a two dimensional array |

#### 7.4.2.3. Selecting a Union of Hyperslabs

The HDF5 library allows the user to select a union of hyperslabs and write or read the selection into another selection. The shapes of the two selections may differ, but the number of elements must be equal.

|  |
| --- |
| Dspace_fig16a.jpg |
| Dspace_fig16b.jpg |
| Dspace_fig16c.jpg |
| Figure 7-11. Transferring hyperslab unions |

The figure above shows the transfer of a selection that is two overlapping hyperslabs from the dataset into a union of hyperslabs in the memory dataset. Note that the destination dataset has a different shape from the source dataset. Similarly, the selection in the memory dataset could have a different shape than the selected union of hyperslabs in the original file. For simplicity, the selection is that same shape at the des­tination.

To implement this transfer, it is necessary to:

1. Get the source dataspace
2. Define one hyperslab selection for the source
3. Define a second hyperslab selection, unioned with the first
4. Get the destination dataspace
5. Define one hyperslab selection for the destination
6. Define a second hyperslab seletion, unioned with the first
7. Execute the data transfer (H5Dread or H5Dwrite) using the source and destination dataspaces

The example below shows example code to create the selections for the source dataspace (the file). The first hyperslab is size 3 x 4 and the left upper corner at the position (1,2). The hyperslab is a simple rectan­gle, so the stride and block are 1. The second hyperslab is 6 x 5 at the position (2,4). The second selection is a union with the first hyperslab (H5S\_SELECT\_OR).

|  |
| --- |
| fid = H5Dget\_space(dataset);    /\*  \* Select first hyperslab for the dataset in the file.  \*/ |
| offset[0] = 1; offset[1] = 2;  block[0] = 1; block[1] = 1;  stride[0] = 1; stride[1] = 1;  count[0]  = 3; count[1]  = 4;  ret = H5Sselect\_hyperslab(fid, H5S\_SELECT\_SET, offset, stride, count, block); |
| /\*  \* Add second selected hyperslab to the selection.  \*/ |
| offset[0] = 2; offset[1] = 4;  block[0] = 1; block[1] = 1;  stride[0] = 1; stride[1] = 1;  count[0]  = 6; count[1]  = 5;  ret = H5Sselect\_hyperslab(fid, H5S\_SELECT\_OR, offset, stride, count, block); |
| Code Example 7-5. Select source hyperslabs |

The example below shows example code to create the selection for the destination in memory. The steps are similar. In this example, the hyperslabs are the same shape, but located in different positions in the dataspace. The first hyperslab is 3 x 4 and starts at (0,0), and the second is 6 x 5 and starts at (1,2).

Finally, the H5Dread call transfers the selected data from the file dataspace to the selection in memory.

In this example, the source and destination selections are two overlapping rectangles. In general, any number of rectangles can be OR’ed, and they do not have to be contiguous. The order of the selections does not matter, but the first should use H5S\_SELECT\_SET; subsequent selections are unioned using H5S\_SELECT\_OR.

It is important to emphasize that the source and destination do not have to be the same shape (or number of rectangles). As long as the two selections have the same number of elements, the data can be trans­ferred.

|  |
| --- |
| /\*  \* Create memory dataspace.  \*/  mid = H5Screate\_simple(MSPACE\_RANK, mdim, NULL); |
| /\*  \* Select two hyperslabs in memory. Hyperslabs has the  \* same size and shape as the selected hyperslabs for  \* the file dataspace.  \*/ |
| offset[0] = 0; offset[1] = 0;  block[0] = 1; block[1] = 1;  stride[0] = 1; stride[1] = 1;  count[0]  = 3; count[1]  = 4;  ret = H5Sselect\_hyperslab(mid, H5S\_SELECT\_SET, offset, stride, count, block); |
| offset[0] = 1; offset[1] = 2;  block[0] = 1; block[1] = 1;  stride[0] = 1; stride[1] = 1;  count[0]  = 6; count[1]  = 5;  ret = H5Sselect\_hyperslab(mid, H5S\_SELECT\_OR, offset, stride, count, block);    ret = H5Dread(dataset, H5T\_NATIVE\_INT, mid, fid, H5P\_DEFAULT, matrix\_out); |
| Code Example 7-6. Select destination hyperslabs |

#### 7.4.2.4. Selecting a List of Independent Points

It is also possible to specify a list of elements to read or write using the function H5Sselect\_elements. The procedure is similar to hyperslab selections.

1. Get the source dataspace
2. Set the selected points
3. Get the destination dataspace
4. Set the selected points
5. Transfer the data using the source and destination dataspaces

The figure below shows an example where four values are to be written to four separate points in a two dimensional dataspace. The source dataspace is a one dimensional array with the values 53, 59, 61, 67. The destination dataspace is an 8 x 12 array. The elements are to be written to the points (0,0), (3,3), (3,5), and (5,6). In this example, the source does not require a selection. The example below the figure shows example code to implement this transfer.

A point selection lists the exact points to be transferred and the order they will be transferred. The source and destination are required to have the same number of elements. A point selection can be used with a hyperslab (for example, the source could be a point selection and the destination a hyperslab, or vice versa), so long as the number of elements selected are the same.

|  |
| --- |
| Dspace_fig19a.jpg |
| Dspace_fig19b.jpg |
| Dspace_fig19c.jpg |
| Figure 7-12. Write data to separate points |

|  |
| --- |
| hsize\_t dim2[] = {4};  int values[] = {53, 59, 61, 67};    /\* Array to store selected points from the  \* file dataspace  \*/  hssize\_t coord[4][2]; |
| /\*  \* Create dataspace for the second dataset.  \*/ |
| mid2 = H5Screate\_simple(1, dim2, NULL);    /\*  \* Select sequence of NPOINTS points in the file  \* dataspace.  \*/ |
| coord[0][0] = 0; coord[0][1] = 0;  coord[1][0] = 3; coord[1][1] = 3;  coord[2][0] = 3; coord[2][1] = 5;  coord[3][0] = 5; coord[3][1] = 6;    ret = H5Sselect\_elements(fid, H5S\_SELECT\_SET, NPOINTS, (const hssize\_t \*\*)coord);    ret = H5Dwrite(dataset, H5T\_NATIVE\_INT, mid2, fid, H5P\_DEFAULT, values); |
| Code Example 7-7. Write data to separate points |

#### 7.4.2.5. Combinations of Selections

Selections are a very flexible mechanism for reorganizing data during a data transfer. With different com­binations of dataspaces and selections, it is possible to implement many kinds of data transfers including sub-setting, sampling, and reorganizing the data. The table below gives some example combinations of source and destination, and the operations they implement.

|  |  |  |
| --- | --- | --- |
| Table 7-2. Selection operations | | |
| **Source** | **Destination** | **Operation** |
| All | All | Copy whole array |
| All | All (different shape) | Copy and reorganize array |
| Hyperslab | All | Sub-set |
| Hyperslab | Hyperslab (same shape) | Selection |
| Hyperslab | Hyperslab (different shape) | Select and rearrange |
| Hyperslab with stride or block | All or hyperslab with stride 1 | Sub-sample, scatter |
| Hyperslab | Points | Scatter |
| Points | Hyperslab or all | Gather |
| Points | Points (same) | Selection |
| Points | Points (different) | Reorder points |

## Dataspace Selection Operations and Data Transfer

This section is under construction.

## References to Dataset Regions

Another use of selections is to store a reference to a region of a dataset. An HDF5 object reference object is a pointer to an object (dataset, group, or committed datatype) in the file. A selection can be used to cre­ate a pointer to a set of selected elements of a dataset, called a region reference. The selection can be either a point selection or a hyperslab selection.

A region reference is an object maintained by the HDF5 Library. The region reference can be stored in a dataset or attribute, and then read. The dataset or attribute is defined to have the special datatype, H5T\_STD\_REF\_DSETREG.

To discover the elements and/or read the data, the region reference can be dereferenced. The H5Rde­frerence call returns an identifier for the dataset, and then the selected dataspace can be retrieved with H5Rget\_select call. The selected dataspace can be used to read the selected data elements.

For more information, see "Reference."

### Example Uses for Region References

Region references are used to implement stored pointers to data within a dataset. For example, features in a large dataset might be indexed by a table. See the figure below. This table could be stored as an HDF5 dataset with a compound datatype, for example, with a field for the name of the feature and a region ref­erence to point to the feature in the dataset. See the second figure below.

|  |
| --- |
| Dspace_fig21.JPG |
| Figure 7-13. Features indexed by a table |

|  |
| --- |
| Dspace_fig22.JPG |
| Figure 7-14. Storing the table with a compound datatype |

### Creating References to Regions

To create a region reference:

1. Create or open the dataset that contains the region
2. Get the dataspace for the dataset
3. Define a selection that specifies the region
4. Create a region reference using the dataset and dataspace with selection
5. Write the region reference(s) to the desired dataset or attribute

The figure below shows a diagram of a file with three datasets. Dataset D1 and D2 are two dimensional arrays of integers. Dataset R1 is a one dimensional array of references to regions in D1 and D2. The regions can be any valid selection of the dataspace of the target dataset.

|  |
| --- |
| Dspace_fig23.JPG |
| Figure 7-15. A file with three datasets |

Note: In the figure above, R1 is a 1 D array of region pointers; each pointer refers to a selection in one dataset.

The example below shows code to create the array of region references. The references are created in an array of type hdset\_reg\_ref\_t. Each region is defined as a selection on the dataspace of the dataset, and a reference is created using H5Rcreate(). The call to H5Rcreate() specifies the file, dataset, and the dataspace with selection.

|  |
| --- |
| /\* create an array of 4 region references \*/  hdset\_reg\_ref\_t ref[4];  /\*  \* Create a reference to the first hyperslab in the first  \* Dataset.  \*/ |
| offset[0] = 1; offset[1] = 1;  count[0] = 3; count[1] = 2;  status =  H5Sselect\_hyperslab(space\_id, H5S\_SELECT\_SET,  offset, NULL, count, NULL);  status = H5Rcreate(&ref[0], file\_id, "D1", H5R\_DATASET\_REGION, space\_id); |
| /\*  \* The second reference is to a union of hyperslabs in  \* the first Dataset  \*/ |
| offset[0] = 5;  offset[1] = 3;  count[0] = 1; count[1] = 4;  status = H5Sselect\_none(space\_id);  status = H5Sselect\_hyperslab(space\_id, H5S\_SELECT\_SET, offset, NULL, count, NULL);  offset[0] = 6;   offset[1] = 5;  count[0] = 1;  count[1] = 2; |
| status = H5Sselect\_hyperslab(space\_id, H5S\_SELECT\_OR, offset, NULL, count, NULL);  status = H5Rcreate(&ref[1], file\_id, "D1", H5R\_DATASET\_REGION, space\_id); |
| /\*  \* the fourth reference is to a selection of points in  \* the first Dataset  \*/ |
| status = H5Sselect\_none(space\_id);  coord[0][0] = 4; coord[0][1] = 4;  coord[1][0] = 2; coord[1][1] = 6;  coord[2][0] = 3; coord[2][1] = 7;  coord[3][0] = 1; coord[3][1] = 5;  coord[4][0] = 5; coord[4][1] = 8;  status = H5Sselect\_elements(space\_id, H5S\_SELECT\_SET, num\_points, (const hssize\_t \*\*)coord);  status = H5Rcreate(&ref[3], file\_id, "D1", H5R\_DATASET\_REGION, space\_id); |
| /\*  \* the third reference is to a hyperslab in the second  \* Dataset  \*/  offset[0] = 0;  offset[1] = 0;  count[0] = 4; count[1] = 6;  status = H5Sselect\_hyperslab(space\_id2, H5S\_SELECT\_SET, offset, NULL, count, NULL);  status = H5Rcreate(&ref[2], file\_id, "D2", H5R\_DATASET\_REGION, space\_id2); |
| Code Example 7-8. Create an array of region references |

When all the references are created, the array of references is written to the dataset R1. The dataset is declared to have datatype H5T\_STD\_REF\_DSETREG. See the example below.

|  |
| --- |
| Hsize\_t dimsr[1];  dimsr[0] = 4;  /\*  \* Dataset with references.  \*/ |
| spacer\_id = H5Screate\_simple(1, dimsr, NULL);  dsetr\_id = H5Dcreate(file\_id, "R1", H5T\_STD\_REF\_DSETREG, spacer\_id, H5P\_DEFAULT,  H5P\_DEFAULT, H5P\_DEFAULT); |
| /\*  \* Write dataset with the references.  \*/  status = H5Dwrite(dsetr\_id, H5T\_STD\_REF\_DSETREG, H5S\_ALL, H5S\_ALL, H5P\_DEFAULT, ref); |
| Code Example 7-9. Write the array of references to a dataset |

When creating region references, the following rules are enforced.

* The selection must be a valid selection for the target dataset, just as when transferring data
* The dataset must exist in the file when the reference is created (H5Rcreate)
* The target dataset must be in the same file as the stored reference

### Reading References to Regions

To retrieve data from a region reference, the reference must be read from the file, and then the data can be retrieved. The steps are:

1. Open the dataset or attribute containing the reference objects
2. Read the reference object(s)
3. For each region reference, get the dataset (H5R\_dereference) and dataspace (H5Rget\_space)
4. Use the dataspace and datatype to discover what space is needed to store the data, allocate the correct storage and create a dataspace and datatype to define the memory data layout

The example below shows code to read an array of region references from a dataset, and then read the data from the first selected region. Note that the region reference has information that records the data­set (within the file) and the selection on the dataspace of the dataset. After dereferencing the regions ref­erence, the datatype, number of points, and some aspects of the selection can be discovered. (For a union of hyperslabs, it may not be possible to determine the exact set of hyperslabs that has been combined.) The table below the code example shows the inquiry functions.

When reading data from a region reference, the following rules are enforced:

* The target dataset must be present and accessible in the file
* The selection must be a valid selection for the dataset

|  |
| --- |
| dsetr\_id = H5Dopen (file\_id, "R1", H5P\_DEFAULT);    status = H5Dread(dsetr\_id, H5T\_STD\_REF\_DSETREG, H5S\_ALL, H5S\_ALL, H5P\_DEFAULT,  ref\_out); |
| /\*  \* Dereference the first reference.  \*   1) get the dataset (H5Rdereference)  \*   2) get the selected dataspace (H5Rget\_region)  \*/ |
| dsetv\_id = H5Rdereference(dsetr\_id, H5R\_DATASET\_REGION, &ref\_out[0]);  space\_id = H5Rget\_region(dsetr\_id, H5R\_DATASET\_REGION, &ref\_out[0]);    /\*  \*  Discover how many points and shape of the data  \*/  ndims = H5Sget\_simple\_extent\_ndims(space\_id); |
| H5Sget\_simple\_extent\_dims(space\_id,dimsx,NULL);    /\*  \* Read and display hyperslab selection from the dataset.  \*/  dimsy[0] = H5Sget\_select\_npoints(space\_id);  spacex\_id = H5Screate\_simple(1, dimsy, NULL); |
| status = H5Dread(dsetv\_id, H5T\_NATIVE\_INT, H5S\_ALL,  space\_id, H5P\_DEFAULT, data\_out);  printf("Selected hyperslab: ");  for (i = 0; i < 8; i++)  {     printf("\n");     for (j = 0; j < 10; j++)        printf("%d ", data\_out[i][j]);  }  printf("\n"); |
| Code Example 7-10. Read an array of region references; read from the first selection |

|  |  |
| --- | --- |
| Table 7-3. The inquiry functions | |
| **Function** | **Information** |
| H5Sget\_select\_npoints | The number of elements in the selection (hyperslab or point selection). |
| H5Sget\_select\_bounds | The bounding box that encloses the selected points (hyperslab or point selection). |
| H5Sget\_select\_hyper\_nblocks | The number of blocks in the selection. |
| H5Sget\_select\_hyper\_blocklist | A list of the blocks in the selection. |
| H5Sget\_select\_elem\_npoints | The number of points in the selection. |
| H5Sget\_select\_elem\_pointlist | The points. |

## Sample Programs

This section contains the full programs from which several of the code examples in this chapter were derived. The h5dump output from the program’s output file immediately follows each program.

### h5\_write.c

----------

#include "hdf5.h"

#define H5FILE\_NAME        "SDS.h5"

#define DATASETNAME "C Matrix"

#define NX     3                      /\* dataset dimensions \*/

#define NY     5

#define RANK   2

int

main (void)

{

   hid\_t       file, dataset;         /\* file and dataset identifiers \*/

   hid\_t       datatype, dataspace;   /\* identifiers \*/

   hsize\_t     dims[2];               /\* dataset dimensions \*/

   herr\_t      status;

   int         data[NX][NY];          /\* data to write \*/

   int         i, j;

   /\*

   \* Data  and output buffer initialization.

   \*/

   for (j = 0; j < NX; j++) {

      for (i = 0; i < NY; i++)

         data[j][i] = i + 1 + j\*NY;

   }

   /\*

   \*  1  2  3  4  5

   \*  6  7  8  9 10

   \* 11 12 13 14 15

   \*/

   /\*

   \* Create a new file using H5F\_ACC\_TRUNC access,

   \* default file creation properties, and default file

   \* access properties.

   \*/

   file = H5Fcreate(H5FILE\_NAME, H5F\_ACC\_TRUNC, H5P\_DEFAULT, H5P\_DEFAULT);

   /\*

   \* Describe the size of the array and create the data space for fixed

   \* size dataset.

   \*/

   dims[0] = NX;

   dims[1] = NY;

   dataspace = H5Screate\_simple(RANK, dims, NULL);

   /\*

   \* Create a new dataset within the file using defined dataspace and

   \* datatype and default dataset creation properties.

   \*/

   dataset = H5Dcreate(file, DATASETNAME, H5T\_NATIVE\_INT, dataspace,

                           H5P\_DEFAULT, H5P\_DEFAULT, H5P\_DEFAULT);

   /\*

   \* Write the data to the dataset using default transfer properties.

   \*/

   status = H5Dwrite(dataset, H5T\_NATIVE\_INT, H5S\_ALL, H5S\_ALL,

                        H5P\_DEFAULT, data);

   /\*

   \* Close/release resources.

   \*/

   H5Sclose(dataspace);

   H5Dclose(dataset);

   H5Fclose(file);

   return 0;

}

SDS.out

-------

HDF5 "SDS.h5" {

GROUP "/" {

   DATASET "C Matrix" {

      DATATYPE  H5T\_STD\_I32BE

      DATASPACE  SIMPLE { ( 3, 5 ) / ( 3, 5 ) }

      DATA {

         1, 2, 3, 4, 5,

         6, 7, 8, 9, 10,

         11, 12, 13, 14, 15

      }

   }

}

}

### h5\_write.f90

   PROGRAM DSETEXAMPLE

   USE HDF5 ! This module contains all necessary modules

   IMPLICIT NONE

   CHARACTER(LEN=7), PARAMETER :: filename = "SDSf.h5" ! File name

   CHARACTER(LEN=14), PARAMETER :: dsetname = "Fortran Matrix" ! Dataset name

   INTEGER, PARAMETER :: NX = 3

   INTEGER, PARAMETER :: NY = 5

   INTEGER(HID\_T) :: file\_id       ! File identifier

   INTEGER(HID\_T) :: dset\_id       ! Dataset identifier

   INTEGER(HID\_T) :: dspace\_id     ! Dataspace identifier

   INTEGER(HSIZE\_T), DIMENSION(2) :: dims = (/3,5/) ! Dataset dimensions

   INTEGER     ::    rank = 2                       ! Dataset rank

   INTEGER     ::    data(NX,NY)

   INTEGER     ::   error ! Error flag

   INTEGER     :: i, j

   !

   ! Initialize data

   !

      do i = 1, NX

         do j = 1, NY

            data(i,j) = j + (i-1)\*NY

         enddo

      enddo

   !

   ! Data

   !

   !  1  2  3  4  5

   !  6  7  8  9 10

   ! 11 12 13 14 15

   !

   ! Initialize FORTRAN interface.

   !

   CALL h5open\_f(error)

   !

   ! Create a new file using default properties.

   !

   CALL h5fcreate\_f(filename, H5F\_ACC\_TRUNC\_F, file\_id, error)

   !

   ! Create the dataspace.

   !

   CALL h5screate\_simple\_f(rank, dims, dspace\_id, error)

   !

   ! Create and write dataset using default properties.

   !

   CALL h5dcreate\_f(file\_id, dsetname, H5T\_NATIVE\_INTEGER, dspace\_id, &

                        dset\_id, error, H5P\_DEFAULT\_F, H5P\_DEFAULT\_F, &

                        H5P\_DEFAULT\_F)

   CALL h5dwrite\_f(dset\_id, H5T\_NATIVE\_INTEGER, data, dims, error)

   !

   ! End access to the dataset and release resources used by it.

   !

   CALL h5dclose\_f(dset\_id, error)

   !

   ! Terminate access to the data space.

   !

   CALL h5sclose\_f(dspace\_id, error)

   !

   ! Close the file.

   !

   CALL h5fclose\_f(file\_id, error)

   !

   ! Close FORTRAN interface.

   !

   CALL h5close\_f(error)

   END PROGRAM DSETEXAMPLE

SDSf.out

--------

HDF5 "SDSf.h5" {

GROUP "/" {

   DATASET "Fortran Matrix" {

      DATATYPE  H5T\_STD\_I32BE

      DATASPACE  SIMPLE { ( 5, 3 ) / ( 5, 3 ) }

      DATA {

         1, 6, 11,

         2, 7, 12,

         3, 8, 13,

         4, 9, 14,

         5, 10, 15

      }

   }

}

}

### h5\_write\_tr.f90

   PROGRAM DSETEXAMPLE

   USE HDF5 ! This module contains all necessary modules

   IMPLICIT NONE

   CHARACTER(LEN=10), PARAMETER :: filename = "SDSf\_tr.h5" ! File name

   CHARACTER(LEN=24), PARAMETER :: dsetname = "Fortran Transpose Matrix"

                                                         ! Dataset name

   INTEGER, PARAMETER :: NX = 3

   INTEGER, PARAMETER :: NY = 5

   INTEGER(HID\_T) :: file\_id       ! File identifier

   INTEGER(HID\_T) :: dset\_id       ! Dataset identifier

   INTEGER(HID\_T) :: dspace\_id     ! Dataspace identifier

   INTEGER(HSIZE\_T), DIMENSION(2) :: dims = (/NY, NX/) ! Dataset dimensions

   INTEGER     ::    rank = 2                       ! Dataset rank

   INTEGER     ::    data(NY,NX)

   INTEGER     ::   error ! Error flag

   INTEGER     :: i, j

   !

   ! Initialize data

   !

      do i = 1, NY

         do j = 1, NX

            data(i,j) = i + (j-1)\*NY

         enddo

      enddo

   !

   ! Data

   !

   !  1  6  11

   !  2  7  12

   !  3  8  13

   !  4  9  14

   !  5 10  15

   !

   ! Initialize FORTRAN interface.

   !

   CALL h5open\_f(error)

   !

   ! Create a new file using default properties.

   !

   CALL h5fcreate\_f(filename, H5F\_ACC\_TRUNC\_F, file\_id, error)

   !

   ! Create the dataspace.

   !

   CALL h5screate\_simple\_f(rank, dims, dspace\_id, error)

   !

   ! Create and write dataset using default properties.

   !

   CALL h5dcreate\_f(file\_id, dsetname, H5T\_NATIVE\_INTEGER, dspace\_id, &

                        dset\_id, error, H5P\_DEFAULT\_F, H5P\_DEFAULT\_F, &

                        H5P\_DEFAULT\_F)

   CALL h5dwrite\_f(dset\_id, H5T\_NATIVE\_INTEGER, data, dims, error)

   !

   ! End access to the dataset and release resources used by it.

   !

   CALL h5dclose\_f(dset\_id, error)

   !

   ! Terminate access to the data space.

   !

   CALL h5sclose\_f(dspace\_id, error)

   !

   ! Close the file.

   !

   CALL h5fclose\_f(file\_id, error)

   !

   ! Close FORTRAN interface.

   !

   CALL h5close\_f(error)

   END PROGRAM DSETEXAMPLE

SDSf\_tr.out

-----------

HDF5 "SDSf\_tr.h5" {

GROUP "/" {

   DATASET "Fortran Transpose Matrix" {

      DATATYPE  H5T\_STD\_I32LE

      DATASPACE  SIMPLE { ( 3, 5 ) / ( 3, 5 ) }

      DATA {

            1, 2, 3, 4, 5,

            6, 7, 8, 9, 10,

            11, 12, 13, 14, 15

      }

   }

}

}

# HDF5 Attributes

## Introduction

An HDF5 attribute is a small metadata object describing the nature and/or intended usage of a primary data object. A primary data object may be a dataset, group, or committed datatype.

Attributes are assumed to be very small as data objects go, so storing them as standard HDF5 datasets would be quite inefficient. HDF5 attributes are therefore managed through a special attributes interface, H5A, which is designed to easily attach attributes to primary data objects as small datasets containing metadata information and to minimize storage requirements.

Consider, as examples of the simplest case, a set of laboratory readings taken under known temperature and pressure conditions of 18.0 degrees Celsius and 0.5 atmospheres, respectively. The temperature and pressure stored as attributes of the dataset could be described as the following name/value pairs:

temp=18.0

pressure=0.5

While HDF5 attributes are not standard HDF5 datasets, they have much in common:

* An attribute has a user-defined dataspace and the included metadata has a user-assigned data­type
* Metadata can be of any valid HDF5 datatype
* Attributes are addressed by name

But there are some very important differences:

* There is no provision for special storage such as compression or chunking
* There is no partial I/O or sub-setting capability for attribute data
* Attributes cannot be shared
* Attributes cannot have attributes
* Being small, an attribute is stored in the object header of the object it describes and is thus attached directly to that object

The “Special Issues” section describes how to handle attributes that are large in size and how to handle large numbers of attributes. For more information, see "Special Issues."

This chapter discusses or lists the following:

* The HDF5 attributes programming model
* H5A function summaries
* Working with HDF5 attributes
  + The structure of an attribute
  + Creating, writing, and reading attributes
  + Accessing attributes by name or index
  + Obtaining information regarding an object’s attributes
  + Iterating across an object’s attributes
  + Deleting an attribute
  + Closing attributes
* Special issues regarding attributes

In the following discussions, attributes are generally attached to datasets. Attributes attached to other pri­mary data objects such as groups or committed datatypes are handled in exactly the same manner.

## Programming Model for Attributes

The figure below shows the UML model for an HDF5 attribute and its associated dataspace and datatype.

|  |
| --- |
| UML_Attribute.jpg |
| Figure 8-1. The UML model for an HDF5 attribute |

Creating an attribute is similar to creating a dataset. To create an attribute, the application must specify the object to which the attribute is attached, the datatype and dataspace of the attribute data, and the attribute creation property list.

The following steps are required to create and write an HDF5 attribute:

1. Obtain the object identifier for the attribute’s primary data object
2. Define the characteristics of the attribute and specify the attribute creation property list

* Define the datatype
* Define the dataspace
* Specify the attribute creation property list

1. Create the attribute
2. Write the attribute data (optional)
3. Close the attribute (and datatype, dataspace, and attribute creation property list, if necessary)
4. Close the primary data object (if appropriate)

### To Open and Read or Write an Existing Attribute

The following steps are required to open and read/write an existing attribute. Since HDF5 attributes allow no partial I/O, you need specify only the attribute and the attribute’s memory datatype to read it:

1. Obtain the object identifier for the attribute’s primary data object
2. Obtain the attribute’s name or index
3. Open the attribute

* Get attribute dataspace and datatype (optional)

1. Specify the attribute’s memory type
2. Read and/or write the attribute data
3. Close the attribute
4. Close the primary data object (if appropriate)

## Attribute (H5A) Function Summaries

Functions that can be used with attributes (H5A functions) and functions that can be used with property lists (H5P functions) are listed below.

|  |  |
| --- | --- |
| Function Listing 8-1. Attribute functions (H5A) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Acreate  h5acreate\_f | Creates a dataset as an attribute of another group, dataset, or committed datatype. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Acreate\_by\_name  h5acreate\_by\_name\_f | Creates an attribute attached to a specified object. |
| H5Aexists  h5aexists\_f | Determines whether an attribute with a given name exists on an object. |
| H5Aexists\_by\_name  h5aexists\_by\_name\_f | Determines whether an attribute with a given name exists on an object. |
| H5Aclose  h5aclose\_f | Closes the specified attribute. |
| H5Adelete  h5adelete\_f | Deletes an attribute. |
| H5Adelete\_by\_idx  h5adelete\_by\_idx\_f | Deletes an attribute from an object according to index order. |
| H5Adelete\_by\_name  h5adelete\_by\_name\_f | Removes an attribute from a specified loca­tion. |
| H5Aget\_create\_plist  h5aget\_create\_plist\_f | Gets an attribute creation property list identi­fier. |
| H5Aget\_info  h5aget\_info\_f | Retrieves attribute information by attribute identifier. |
| H5Aget\_info\_by\_idx  h5aget\_info\_by\_idx\_f | Retrieves attribute information by attribute index position. |
| H5Aget\_info\_by\_name  h5aget\_info\_by\_name\_f | Retrieves attribute information by attribute name. |
| H5Aget\_name  h5aget\_name\_f | Gets an attribute name. |
| H5Aget\_name\_by\_idx  h5aget\_name\_by\_idx\_f | Gets an attribute name by attribute index position. |
| H5Aget\_space  h5aget\_space\_f | Gets a copy of the dataspace for an attribute. |
| H5Aget\_storage\_size  h5aget\_storage\_size\_f | Returns the amount of storage required for an attribute. |
| H5Aget\_type  h5aget\_type\_f | Gets an attribute datatype. |
| H5Aiterate  (no Fortran subroutine) | Calls a user’s function for each attribute attached to a data object. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Aiterate\_by\_name  (no Fortran subroutine) | Calls user-defined function for each attribute on an object. |
| H5Aopen  h5aopen\_f | Opens an attribute for an object specified by object identifier and attribute name. |
| H5Aopen\_by\_idx  h5aopen\_by\_idx\_f | Opens an existing attribute that is attached to an object specified by location and name. |
| H5Aopen\_by\_name  h5aopen\_by\_name\_f | Opens an attribute for an object by object name and attribute name. |
| H5Aread  h5aread\_f | Reads an attribute. |
| H5Arename  h5arename\_f | Renames an attribute. |
| H5Arename\_by\_name  h5arename\_by\_name\_f | Renames an attribute. |
| H5Awrite  H5awrite\_f | Writes an attribute. |

|  |  |
| --- | --- |
| Function Listing 8-2. Attribute creation property list functions (H5P) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Pset\_char\_encoding  h5pset\_char\_encoding\_f | Sets the character encoding used to encode a string. Use to set ASCII or UTF-8 character encoding for object names. |
| H5Pget\_char\_encoding  h5pget\_char\_encoding\_f | Retrieves the character encoding used to cre­ate a string. |
| H5Pget\_attr\_creation\_order  h5pget\_attr\_creation\_order\_f | Retrieves tracking and indexing settings for attribute creation order. |
| H5Pget\_attr\_phase\_change  h5pget\_attr\_phase\_change\_f | Retrieves attribute storage phase change thresholds. |
| H5Pset\_attr\_creation\_order  h5pget\_attr\_creation\_order\_f | Sets tracking and indexing of attribute cre­ation order. |
| H5Pset\_attr\_phase\_change  h5pset\_attr\_phase\_change\_f | Sets attribute storage phase change thresh­olds. |

## Working with Attributes

### The Structure of an Attribute

An attribute has two parts: name and value(s).

HDF5 attributes are sometimes discussed as name/value pairs in the form name=value.

An attribute’s name is a null-terminated ASCII or UTF-8 character string. Each attribute attached to an object has a unique name.

The value portion of the attribute contains one or more data elements of the same datatype.

HDF5 attributes have all the characteristics of HDF5 datasets except that there is no partial I/O capability. In other words, attributes can be written and read only in full with no sub-setting.

### Creating, Writing, and Reading Attributes

If attributes are used in an HDF5 file, these functions will be employed: H5Acreate, H5Awrite, and H5Aread. H5Acreate and H5Awrite are used together to place the attribute in the file. If an attribute is to be used and is not currently in memory, H5Aread generally comes into play usually in concert with one each of the H5Aget\_\* and H5Aopen\_\* functions.

To create an attribute, call H5Acreate:

hid\_t H5Acreate (hid\_t loc\_id, const char \*name,

            hid\_t type\_id, hid\_t space\_id, hid\_t create\_plist,

            hid\_t access\_plist)

loc\_id identifies the object (dataset, group, or committed datatype) to which the attribute is to be attached. name, type\_id, space\_id, and create\_plist convey, respectively, the attribute’s name, datatype, dataspace, and attribute creation property list. The attribute’s name must be locally unique: it must be unique within the context of the object to which it is attached.

H5Acreate creates the attribute in memory. The attribute does not exist in the file until H5Awrite writes it there.

To write or read an attribute, call H5Awrite or H5Aread, respectively:

herr\_t H5Awrite (hid\_t attr\_id, hid\_t mem\_type\_id, const void \*buf)

herr\_t H5Aread (hid\_t attr\_id, hid\_t mem\_type\_id, void \*buf)

attr\_id identifies the attribute while mem\_type\_id identifies the in-memory datatype of the attribute data.

H5Awrite writes the attribute data from the buffer buf to the file. H5Aread reads attribute data from the file into buf.

The HDF5 Library converts the metadata between the in-memory datatype, mem\_type\_id, and the in-file datatype, defined when the attribute was created, without user intervention.

### Accessing Attributes by Name or Index

Attributes can be accessed by name or index value. The use of an index value makes it possible to iterate through all of the attributes associated with a given object.

To access an attribute by its name, use the H5Aopen\_by\_name function. H5Aopen\_by\_name returns an attribute identifier that can then be used by any function that must access an attribute such as H5Aread.Use the function H5Aget\_name to determine an attribute’s name.

To access an attribute by its index value, use the H5Aopen\_by\_idx function. To determine an attribute index value when it is not already known, use the H5Oget\_info function. H5Aopen\_by\_idx is generally used in the course of opening several attributes for later access. Use H5Aiterate if the intent is to per­form the same operation on every attribute attached to an object.

### Obtaining Information Regarding an Object’s Attributes

In the course of working with HDF5 attributes, one may need to obtain any of several pieces of informa­tion:

* An attribute name
* The dataspace of an attribute
* The datatype of an attribute
* The number of attributes attached to an object

**To obtain an attribute’s name**, call H5Aget\_name with an attribute identifier, attr\_id:

ssize\_t H5Aget\_name (hid\_t attr\_id, size\_t buf\_size, char \*buf)

As with other attribute functions, attr\_id identifies the attribute; buf\_size defines the size of the buf­fer; and buf is the buffer to which the attribute’s name will be read.

If the length of the attribute name, and hence the value required for buf\_size, is unknown, a first call to H5Aget\_name will return that size. If the value of buf\_size used in that first call is too small, the name will simply be truncated in buf. A second H5Aget\_name call can then be used to retrieve the name in an appropriately-sized buffer.

**To determine the dataspace or datatype of an attribute**, call H5Aget\_space or H5Aget\_type, respec­tively:

hid\_t H5Aget\_space (hid\_t attr\_id)

hid\_t H5Aget\_type (hid\_t attr\_id)

H5Aget\_space returns the dataspace identifier for the attribute attr\_id.

H5Aget\_type returns the datatype identifier for the attribute attr\_id.

**To determine the number of attributes attached to an object**, use the H5Oget\_info function. The func­tion signature is below.

herr\_t H5Oget\_info( hid\_t object\_id, H5O\_info\_t \*object\_info  )

The number of attributes will be returned in the object\_info buffer. This is generally the preferred first step in determining attribute index values. If the call returns N, the attributes attached to the object object\_id have index values of 0 through N-1.

### Iterating across an Object’s Attributes

It is sometimes useful to be able to perform the identical operation across all of the attributes attached to an object. At the simplest level, you might just want to open each attribute. At a higher level, you might wish to perform a rather complex operation on each attribute as you iterate across the set.

To iterate an operation across the attributes attached to an object, one must make a series of calls to H5Aiterate:

herr\_t H5Aiterate (hid\_t obj\_id, H5\_index\_t index\_type,

            H5\_iter\_order\_t order, hsize\_t \*n, H5A\_operator2\_t op,

            void \*op\_data)

H5Aiterate successively marches across all of the attributes attached to the object specified in loc\_id, performing the operation(s) specified in op\_func with the data specified in op\_data on each attribute.

When H5Aiterate is called, index contains the index of the attribute to be accessed in this call. When H5Aiterate returns, index will contain the index of the next attribute. If the returned index is the null pointer, then all attributes have been processed, and the iterative process is complete.

op\_func is a user-defined operation that adheres to the H5A\_operator\_t prototype. This prototype and certain requirements imposed on the operator’s behavior are described in the H5Aiterate entry in the HDF5 Reference Manual.

op\_data is also user-defined to meet the requirements of op\_func. Beyond providing a parameter with which to pass this data, HDF5 provides no tools for its management and imposes no restrictions.

### Deleting an Attribute

Once an attribute has outlived its usefulness or is no longer appropriate, it may become necessary to delete it.

To delete an attribute, call H5Adelete:

herr\_t H5Adelete (hid\_t loc\_id, const char \*name)

H5Adelete removes the attribute name from the group, dataset, or committed datatype specified in loc\_id.

H5Adelete must not be called if there are any open attribute identifiers on the object loc\_id. Such a call can cause the internal attribute indexes to change; future writes to an open attribute would then pro­duce unintended results.

### Closing an Attribute

As is the case with all HDF5 objects, once access to an attribute it is no longer needed, that attribute must be closed. It is best practice to close it as soon as practicable; it is mandatory that it be closed prior to the H5close call closing the HDF5 Library.

To close an attribute, call H5Aclose:

herr\_t H5Aclose (hid\_t attr\_id)

H5Aclose closes the specified attribute by terminating access to its identifier, attr\_id.

## Special Issues

Some special issues for attributes are discussed below.

**Large Numbers of Attributes Stored in Dense Attribute Storage**

The dense attribute storage scheme was added in version 1.8 so that datasets, groups, and committed datatypes that have large numbers of attributes could be processed more quickly.

Attributes start out being stored in an object's header. This is known as compact storage. For more information, see "Storage Strategies."

As the number of attributes grows, attribute-related performance slows. To improve performance, dense attribute storage can be initiated with the H5Pset\_attr\_phase\_change function. See the HDF5 Refer­ence Manual for more information.

When dense attribute storage is enabled, a threshold is defined for the number of attributes kept in com­pact storage. When the number is exceeded, the library moves all of the attributes into dense storage at another location. The library handles the movement of attributes and the pointers between the locations automatically. If some of the attributes are deleted so that the number falls below the threshold, then the attributes are moved back to compact storage by the library.

The improvements in performance from using dense attribute storage are the result of holding attributes in a heap and indexing the heap with a B-tree.

Note that there are some disadvantages to using dense attribute storage. One is that this is a new feature. Datasets, groups, and committed datatypes that use dense storage cannot be read by applications built with earlier versions of the library. Another disadvantage is that attributes in dense storage cannot be compressed.

**Large Attributes Stored in Dense Attribute Storage**

We generally consider the maximum size of an attribute to be 64K bytes. The library has two ways of stor­ing attributes larger than 64K bytes: in dense attribute storage or in a separate dataset. Using dense attri­bute storage is described in this section, and storing in a separate dataset is described in the next section.

To use dense attribute storage to store large attributes, set the number of attributes that will be stored in compact storage to 0 with the H5Pset\_attr\_phase\_change function. This will force all attributes to be put into dense attribute storage and will avoid the 64KB size limitation for a single attribute in compact attribute storage.

The example code below illustrates how to create a large attribute that will be kept in dense storage.

|  |
| --- |
| /\*  \* Test use of dense attribute  \*/ |
| #define N 82000000  #include "hdf5.h"  #include <stdio.h>  #include <stdlib.h>    int main(){ |
| hid\_t fid, gid, sid, aid, gpid, fpid;  hsize\_t dims[] = {N};  double \*buf;  int i;  herr\_t status; |
| buf = (double \*) malloc(sizeof(double) \* N);  for (i=0; i <N; i++) { buf[i] = -100.0; }  fpid = H5Pcreate (H5P\_FILE\_ACCESS);  status = H5Pset\_libver\_bounds (fpid, H5F\_LIBVER\_LATEST, H5F\_LIBVER\_LATEST);  fid = H5Fcreate("adense.h5", H5F\_ACC\_TRUNC, H5P\_DEFAULT, fpid);  gpid = H5Pcreate (H5P\_GROUP\_CREATE);  status = H5Pset\_attr\_phase\_change (gpid, 0, 0); |
| gid = H5Gcreate(fid, "testgrp", H5P\_DEFAULT, gpid, H5P\_DEFAULT);  sid = H5Screate\_simple(1, dims, NULL);    aid = H5Acreate(gid, "bar", H5T\_NATIVE\_DOUBLE, sid, H5P\_DEFAULT, H5P\_DEFAULT);  status = H5Awrite(aid, H5T\_NATIVE\_DOUBLE, buf); |
| /\* If you remove these two lines, it doesn't crash \*/  status = H5Aclose(aid);  status = H5Pclose (gpid);  status = H5Pclose (fpid);  status = H5Gclose(gid);  status = H5Fclose (fid);    return 0;  } |
| Code Example 8-1. Create a large attribute in dense storage |

**Large Attributes Stored in a Separate Dataset**

In addition to dense attribute storage (see above), a large attribute can be stored in a separate dataset. In the figure below, DatasetA holds an attribute that is too large for the object header in Dataset1. By putting a pointer to DatasetA as an attribute in Dataset1, the attribute becomes available to those working with Dataset1.

This way of handling large attributes can be used in situations where backward compatibility is important and where compression is important. Applications built with versions before 1.8.x can read large attri­butes stored in separate datasets. Datasets can be compressed while attributes cannot.

|  |
| --- |
| Shared_Attribute.jpg |
| Figure 8-2. A large or shared HDF5 attribute and its associated dataset(s) |

Note: In the figure above, DatasetA is an attribute of Dataset1 that is too large to store in Dataset1's header. Data­setA is associated with Dataset1 by means of an object reference pointer attached as an attribute to Dataset1. The attribute in DatasetA can be shared among multiple datasets by means of additional object reference pointers attached to additional datasets.

**Shared Attributes**

Attributes written and managed through the H5A interface cannot be shared. If shared attributes are required, they must be handled in the manner described above for large attributes and illustrated in the figure above.

**Attribute Names**

While any ASCII or UTF-8 character may be used in the name given to an attribute, it is usually wise to avoid the following kinds of characters:

* Commonly used separators or delimiters such as slash, backslash, colon, and semi-colon (\, /, :, ;)
* Escape characters
* Wild cards such as asterisk and question mark (\*, ?)

NULL can be used within a name, but HDF5 names are terminated with a NULL: whatever comes after the NULL will be ignored by HDF5.

The use of ASCII or UTF-8 characters is determined by the character encoding property. See H5Pset\_char\_encoding in the HDF5 Reference Manual.

**No Special I/O or Storage**

HDF5 attributes have all the characteristics of HDF5 datasets except the following:

* Attributes are written and read only in full: there is no provision for partial I/O or sub-setting
* No special storage capability is provided for attributes: there is no compression or chunking, and attributes are not extendable

# HDF5 Error Handling

## Introduction

The HDF5 library provides an error reporting mechanism for both the library itself and for user application programs. It can trace errors through function stack and error information like file name, function name, line number, and error description.

"Basic Error Handling Operations" discusses the basic error concepts such as error stack, error record, and error message and describes the related API functions. These concepts and func­tions are sufficient for application programs to trace errors inside the HDF5 Library.

“Advanced Error Handling Operations” talks about the advanced concepts of error class and error stack handle and talks about the related functions. With these concepts and functions, an application library or program using the HDF5 library can have its own error report blended with HDF5’s error report.

Starting with Release 1.8, we have a new set of Error Handling API functions. For the purpose of backward compatibility with version 1.6 and before, we still keep the old API functions, H5Epush, H5Eprint, H5Ewalk, H5Eclear, H5Eget\_auto, H5Eset\_auto. These functions do not have the error stack as parameter. The library allows them to operate on the default error stack. Users do not have to change their code to catch up with the new Error API but are encouraged to do so.

The old API is similar to functionality discussed in "Basic Error Handling Operations". The functionality discussed in "Advanced Error Handling Operations", the ability of allowing applications to add their own error records, is the new design for the Error Handling API.

## Programming Model for Error Handling

This section is under construction.

## Error Handling (H5E) Function Summaries

Functions that can be used to handle errors (H5E functions) are listed below.

|  |  |
| --- | --- |
| Function Listing 9-1. Error handling functions (H5E) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Eauto\_is\_v2  (no Fortran subroutine) | Determines the type of error stack. |
| H5Eclear  h5eclear\_f | Clears the error stack for the current thread. The C function is a macro: see “API Compati­bility Macros in HDF5.” |
| H5Eclear\_stack  (no Fortran subroutine) | Clears the error stack for the current thread. |
| H5Eclose\_msg  (no Fortran subroutine) | Closes an error message identifier. |
| H5Eclose\_stack  (no Fortran subroutine) | Closes object handle for error stack. |
| H5Ecreate\_msg  (no Fortran subroutine) | Add major error message to an error class. |
| H5Eget\_auto  h5eget\_auto\_f | Returns the current settings for the automatic error stack traversal function and its data. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Eget\_class\_name  (no Fortran subroutine) | Retrieves error class name. |
| H5Eget\_current\_stack  (no Fortran subroutine) | Registers the current error stack. |
| H5Eget\_msg  (no Fortran subroutine) | Retrieves an error message. |
| H5Eget\_num  (no Fortran subroutine) | Retrieves the number of error messages in an error stack. |
| H5Epop  (no Fortran subroutine) | Deletes specified number of error messages from the error stack. |
| H5Eprint  h5eprint\_f | Prints the error stack in a default manner. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Epush  (no Fortran subroutine) | Pushes new error record onto error stack. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Eregister\_class  (no Fortran subroutine) | Registers a client library or application pro­gram to the HDF5 error API. |
| H5Eset\_auto  h5eset\_auto\_f | Turns automatic error printing on or off. The C function is a macro: see “API Compatibility Macros in HDF5.” |
| H5Eset\_current\_stack  (no Fortran subroutine) | Replaces the current error stack. |
| H5Eunregister\_class  (no Fortran subroutine) | Removes an error class. |
| H5Ewalk  (no Fortran subroutine) | Walks the error stack for the current thread, calling a specified function. The C function is a macro: see “API Compatibility Macros in HDF5.” |

## Basic Error Handling Operations

Let us first try to understand the error stack. An error stack is a collection of error records.   Error records can be pushed onto or popped off the error stack. By default, when an error occurs deep within the HDF5 Library, an error record is pushed onto an error stack and that function returns a failure indication. Its caller detects the failure, pushes another record onto the stack, and returns a failure indication. This con­tinues until the API function called by the application returns a failure indication. The next API function being called will reset the error stack. All HDF5 Library error records belong to the same error class. For more information, see "Advanced Error Handling Operations."

### Error Stack and Error Message

In normal circumstances, an error causes the stack to be printed on the standard error stream automati­cally. This automatic error stack is the library’s default stack. For all the functions in this section, whenever an error stack ID is needed as a parameter, H5E\_DEFAULT can be used to indicate the library’s default stack. The first error record of the error stack, number #000, is produced by the API function itself and is usually sufficient to indicate to the application what went wrong.

**Example: An Error Report**

If an application calls H5Tclose on a predefined datatype, then the message in the example below is printed on the standard error stream. This is a simple error that has only one component, the API func­tion; other errors may have many components.

|  |
| --- |
| HDF5-DIAG: Error detected in HDF5 (1.6.4) thread 0.     #000: H5T.c line 462 in H5Tclose(): predefined datatype        major: Function argument        minor: Bad value |
| Code Example 9-1. An error report |

In the example above, we can see that an error record has a major message and a minor message. A major message generally indicates where the error happens. The location can be a dataset or a dataspace, for example. A minor message explains further details of the error. An example is “unable to open file”. Another specific detail about the error can be found at the end of the first line of each error record. This error description is usually added by the library designer to tell what exactly goes wrong. In the example above, the “predefined datatype” is an error description.

### Print and Clear an Error Stack

Besides the automatic error report, the error stack can also be printed and cleared by the functions H5Eprint() and H5Eclear\_stack(). If an application wishes to make explicit calls to H5Eprint() to print the error stack, the automatic printing should be turned off to prevent error messages from being displayed twice (see H5Eset\_auto() below).

**To print an error stack:**

herr\_t H5Eprint(hid\_t error\_stack, FILE \* stream)

This function prints the error stack specified by error\_stack on the specified stream, stream. If the error stack is empty, a one-line message will be printed. The following is an example of such a message. This message would be generated if the error was in the HDF5 library.

HDF5-DIAG: Error detected in HDF5 Library version: 1.5.62 thread 0.

**To clear an error stack:**

herr\_t H5Eclear\_stack(hid\_t error\_stack)

The H5Eclear\_stack function shown above clears the error stack specified by error\_stack. H5E\_DE­FAULT can be passed in to clear the current error stack. The current stack is also cleared whenever an API function is called; there are certain exceptions to this rule such as H5Eprint().

### Mute Error Stack

Sometimes an application calls a function for the sake of its return value, fully expecting the function to fail; sometimes the application wants to call H5Eprint() explicitly. In these situations, it would be mis­leading if an error message were still automatically printed. Using the H5Eset\_auto() function can con­trol the automatic printing of error messages.

To enable or disable automatic printing of errors:

herr\_t H5Eset\_auto(hid\_t error\_stack, H5E\_auto\_t func, void \*client\_data)

The H5Eset\_auto function can be used to turns on or off the automatic printing of errors for the error stack specified by error\_stack. When turned on (non-null func pointer), any API function which returns an error indication will first call func, passing it client\_data as an argument. When the library is first initialized the auto printing function is set to H5Eprint() (cast appropriately) and client\_data is the standard error stream pointer, stderr.

**To see the current settings:**

herr\_t H5Eget\_auto(hid\_t error\_stack, H5E\_auto\_t \* func, void \*\*client\_data)

The function above returns the current settings for the automatic error stack traversal function, func, and its data, client\_data. If either or both of the arguments are null, then the value is not returned.

**Example: Error Control**

An application can temporarily turn off error messages while “probing” a function. See the example below.

|  |
| --- |
| /\* Save old error handler \*/  H5E\_auto2\_t  oldfunc;  void \*old\_client\_data;    H5Eget\_auto(error\_stack, &old\_func, &old\_client\_data); |
| /\* Turn off error handling \*/  H5Eset\_auto(error\_stack, NULL, NULL);    /\* Probe. Likely to fail, but that’s okay \*/  status = H5Fopen (......);    /\* Restore previous error handler \*/  H5Eset\_auto(error\_stack, old\_func, old\_client\_data); |
| Code Example 9-2. Turn off error messages while probing a function |

 Or automatic printing can be disabled altogether and error messages can be explicitly printed.

|  |
| --- |
| /\* Turn off error handling permanently \*/  H5Eset\_auto(error\_stack, NULL, NULL);    /\* If failure, print error message \*/  if (H5Fopen (....)<0) {     H5Eprint(H5E\_DEFAULT, stderr);     exit (1);  } |
| Code Example 9-3. Disable automatic printing and explicitly print error messages |

### Customized Printing of an Error Stack

Applications are allowed to define an automatic error traversal function other than the default H5Eprint(). For instance, one can define a function that prints a simple, one-line error message to the standard error stream and then exits. The first example below defines a such a function. The second exam­ple below installs the function as the error handler.

|  |
| --- |
| herr\_t  my\_hdf5\_error\_handler(void \*unused)  {     fprintf (stderr, “An HDF5 error was detected. Bye.\n”);     exit (1);  } |
| Code Example 9-4. Defining a function to print a simple error message |

|  |
| --- |
| H5Eset\_auto(H5E\_DEFAULT, my\_hdf5\_error\_handler, NULL); |
| Code Example 9-5. The user-defined error handler |

### Walk through the Error Stack

The H5Eprint() function is actually just a wrapper around the more complex H5Ewalk() function which traverses an error stack and calls a user-defined function for each member of the stack. The exam­ple below shows how H5Ewalk is used.

herr\_t H5Ewalk(hid\_t err\_stack, H5E\_direction\_t direction,  H5E\_walk\_t func, void \*client\_data)

The error stack err\_stack is traversed and func is called for each member of the stack. Its arguments are an integer sequence number beginning at zero (regardless of direction) and the client\_data pointer. If direction is H5E\_WALK\_UPWARD, then traversal begins at the inner-most function that detected the error and concludes with the API function. Use H5E\_WALK\_DOWNWARD for the opposite order.

### Traverse an Error Stack with a Callback Function

An error stack traversal callback function takes three arguments: n is a sequence number beginning at zero for each traversal, eptr is a pointer to an error stack member, and client\_data is the same pointer used in the example above passed to H5Ewalk(). See the example below.

typedef herr\_t (\*H5E\_walk\_t)(unsigned n, H5E\_error2\_t \*eptr,

            void \*client\_data)

The H5E\_error2\_t structure is shown below.

typedef struct {

            hid\_t cls\_id;

            hid\_t maj\_num;

            hid\_t min\_num;

            unsigned line;

            const char \*func\_name;

            const char \*file\_name;

            const char \*desc;

} H5E\_error2\_t;

The maj\_num and min\_num are major and minor error IDs, func\_name is the name of the function where the error was detected, file\_name and line locate the error within the HDF5 Library source code, and desc points to a description of the error.

**Example: Callback Function**

The following example shows a user-defined callback function.

|  |
| --- |
| #define MSG\_SIZE       64    herr\_t  custom\_print\_cb(unsigned n, const H5E\_error2\_t \*err\_desc,  void\* client\_data) |
| {     FILE               \*stream  = (FILE \*)client\_data;     char               maj[MSG\_SIZE];     char               min[MSG\_SIZE];     char               cls[MSG\_SIZE];     const int               indent = 4; |
| /\* Get descriptions for the major and minor error     \* numbers     \*/     if(H5Eget\_class\_name(err\_desc->cls\_id, cls, MSG\_SIZE)<0)           TEST\_ERROR;       if(H5Eget\_msg(err\_desc->maj\_num, NULL, maj, MSG\_SIZE)<0)           TEST\_ERROR; |
| if(H5Eget\_msg(err\_desc->min\_num, NULL, min, MSG\_SIZE)<0)           TEST\_ERROR;       fprintf (stream, “%\*serror #%03d: %s in %s():              line %u\n”,              indent, “”, n, err\_desc->file\_name,              err\_desc->func\_name, err\_desc->line); |
| fprintf (stream, “%\*sclass: %s\n”, indent\*2, “”, cls);     fprintf (stream, “%\*smajor: %s\n”, indent\*2, “”, maj);     fprintf (stream, “%\*sminor: %s\n”, indent\*2, “”, min);       return 0;    error:     return -1;  } |
| Code Example 9-6. A user-defined callback function |

**Programming Note for C++ Developers Using C Functions**

If a C routine that takes a function pointer as an argument is called from within C++ code, the C routine should be returned from normally.

Examples of this kind of routine include callbacks such as H5Pset\_elink\_cb and H5Pset\_type\_con­v\_cb and functions such as H5Tconvert and H5Ewalk2.

Exiting the routine in its normal fashion allows the HDF5 C Library to clean up its work properly. In other words, if the C++ application jumps out of the routine back to the C++ “catch” statement, the library is not given the opportunity to close any temporary data structures that were set up when the routine was called. The C++ application should save some state as the routine is started so that any problem that occurs might be diagnosed.

## Advanced Error Handling Operations

The section above, see "Basic Error Handling Operations", discusses the basic error handling operations of the library. In that section, all the error records on the error stack are from the library itself. In this section, we are going to introduce the operations that allow an application program to push its own error records onto the error stack once it declares an error class of its own through the HDF5 Error API.

**Example: An Error Report**

An error report shows both the library’s error record and the application’s error records. See the example below.

|  |
| --- |
| Error Test-DIAG: Error detected in Error Program (1.0)           thread 8192:     #000: ../../hdf5/test/error\_test.c line 468 in main():           Error test failed        major: Error in test        minor: Error in subroutine |
| #001: ../../hdf5/test/error\_test.c line 150 in           test\_error(): H5Dwrite failed as supposed to        major: Error in IO        minor: Error in H5Dwrite |
| HDF5-DIAG: Error detected in HDF5 (1.7.5) thread 8192:     #002: ../../hdf5/src/H5Dio.c line 420 in H5Dwrite():           not a dataset        major: Invalid arguments to routine        minor: Inappropriate type |
| Code Example 9-7. An error report |

 In the line above error record #002 in the example above, the starting phrase is HDF5. This is the error class name of the HDF5 library. All of the library’s error messages (major and minor) are in this default error class. The Error Test in the beginning of the line above error record #000 is the name of the application’s error class. The first two error records, #000 and #001, are from application’s error class.

By definition, an error class is a group of major and minor error messages for a library (the HDF5 library or an application library built on top of the HDF5 library) or an application program. The error class can be registered for a library or program through the HDF5 Error API. Major and minor messages can be defined in an error class. An application will have object handles for the error class and for major and minor mes­sages for further operation. See the example below.

|  |
| --- |
| #define MSG\_SIZE       64    herr\_t  custom\_print\_cb(unsigned n, const H5E\_error2\_t \*err\_desc,  void\* client\_data)  { |
| FILE                        \*stream  = (FILE \*)client\_data;     char                        maj[MSG\_SIZE];     char                        min[MSG\_SIZE];     char                        cls[MSG\_SIZE];     const int                        indent = 4; |
| /\* Get descriptions for the major and minor error     \* numbers     \*/     if(H5Eget\_class\_name(err\_desc->cls\_id, cls, MSG\_SIZE)<0)        TEST\_ERROR;       if(H5Eget\_msg(err\_desc->maj\_num, NULL, maj, MSG\_SIZE)<0)        TEST\_ERROR; |
| if(H5Eget\_msg(err\_desc->min\_num, NULL, min, MSG\_SIZE)<0)        TEST\_ERROR;       fprintf (stream, “%\*serror #%03d: %s in %s():           line %u\n”,           indent, “”, n, err\_desc->file\_name,           err\_desc->func\_name, err\_desc->line); |
| fprintf (stream, “%\*sclass: %s\n”, indent\*2, “”, cls);     fprintf (stream, “%\*smajor: %s\n”, indent\*2, “”, maj);     fprintf (stream, “%\*sminor: %s\n”, indent\*2, “”, min);       return 0;    error:     return -1;  } |
| Code Example 9-8. Defining an error class |

### More Error API Functions

The Error API has functions that can be used to register or unregister an error class, to create or close error messages, and to query an error class or error message. These functions are illustrated below.

**To register an error class:**

hid\_t H5Eregister\_class(const char\* cls\_name, const char\* lib\_name,

            const char\* version)

This function registers an error class with the HDF5 Library so that the application library or program can report errors together with the HDF5 Library.

**To add an error message to an error class:**

hid\_t H5Ecreate\_msg(hid\_t class, H5E\_type\_t msg\_type, const char\* mesg)

This function adds an error message to an error class defined by an application library or program. The error message can be either major or minor which is indicated by parameter msg\_type.

**To get the name of an error class:**

ssize\_t H5Eget\_class\_name(hid\_t class\_id, char\* name, size\_t size)

This function retrieves the name of the error class specified by the class ID.

**To retrieve an error message:**

ssize\_t H5Eget\_msg(hid\_t mesg\_id, H5E\_type\_t\* mesg\_type, char\* mesg,

            size\_t size)

This function retrieves the error message including its length and type.

**To close an error message:**

herr\_t H5Eclose\_msg(hid\_t mesg\_id)

This function closes an error message.

**To remove an error class:**

herr\_t H5Eunregister\_class(hid\_t class\_id)

This function removes an error class from the Error API.

**Example: Error Class and its Message**

The example below shows how an application creates an error class and error messages.

|  |
| --- |
| /\* Create an error class \*/  class\_id = H5Eregister\_class(ERR\_CLS\_NAME, PROG\_NAME,  PROG\_VERS);    /\* Retrieve class name \*/  H5Eget\_class\_name(class\_id, cls\_name, cls\_size); |
| /\* Create a major error message in the class \*/  maj\_id = H5Ecreate\_msg(class\_id, H5E\_MAJOR, “... ...”);    /\* Create a minor error message in the class \*/  min\_id = H5Ecreate\_msg(class\_id, H5E\_MINOR, “... ...”); |
| Code Example 9-9. Create an error class and error messages |

The example below shows how an application closes error messages and unregisters the error class.

|  |
| --- |
| H5Eclose\_msg(maj\_id);  H5Eclose\_msg(min\_id);  H5Eunregister\_class(class\_id); |
| Code Example 9-10. Closing error messages and unregistering the error class |

### Pushing an Application Error Message onto Error Stack

An application can push error records onto or pop error records off of the error stack just as the library does internally. An error stack can be registered, and an object handle can be returned to the application so that the application can manipulate a registered error stack.

**To register the current stack:**

hid\_t H5Eget\_current\_stack(void)

This function registers the current error stack, returns an object handle, and clears the current error stack. An empty error stack will also be assigned an ID.

**To replace the current error stack with another:**

herr\_t H5Eset\_current\_stack(hid\_t error\_stack)

This function replaces the current error stack with another error stack specified by error\_stack and clears the current error stack. The object handle error\_stack is closed after this function call.

**To push a new error record to the error stack:**

herr\_t H5Epush(hid\_t error\_stack, const char\* file, const char\* func,

            unsigned line, hid\_t cls\_id, hid\_t major\_id, hid\_t minor\_id,

            const char\* desc, ... )

This function pushes a new error record onto the error stack for the current thread.

**To delete some error messages:**

herr\_t H5Epop(hid\_t error\_stack, size\_t count)

This function deletes some error messages from the error stack.

**To retrieve the number of error records:**

int H5Eget\_num(hid\_t error\_stack)

This function retrieves the number of error records from an error stack.

**To clear the error stack:**

herr\_t H5Eclear\_stack(hid\_t error\_stack)

This function clears the error stack.

**To close the object handle for an error stack:**

herr\_t H5Eclose\_stack(hid\_t error\_stack)

This function closes the object handle for an error stack and releases its resources.

**Example: Working with an Error Stack**

The example below shows how an application pushes an error record onto the default error stack.

|  |
| --- |
| /\* Make call to HDF5 I/O routine \*/  if((dset\_id=H5Dopen(file\_id, dset\_name, access\_plist))<0)  {     /\* Push client error onto error stack \*/ |
| H5Epush(H5E\_DEFAULT,\_\_FILE\_\_,FUNC,\_\_LINE\_\_,cls\_id,        CLIENT\_ERR\_MAJ\_IO,CLIENT\_ERR\_MINOR\_OPEN,        “H5Dopen failed”);       /\* Indicate error occurred in function \*/     return(0);  } |
| Code Example 9-11. Pushing an error message to an error stack |

The example below shows how an application registers the current error stack and creates an object han­dle to avoid another HDF5 function from clearing the error stack.

|  |
| --- |
| if(H5Dwrite(dset\_id, mem\_type\_id, mem\_space\_id,        file\_space\_id, dset\_xfer\_plist\_id, buf)<0)  { |
| /\* Push client error onto error stack \*/     H5Epush(H5E\_DEFAULT,\_\_FILE\_\_,FUNC,\_\_LINE\_\_,cls\_id,        CLIENT\_ERR\_MAJ\_IO,CLIENT\_ERR\_MINOR\_HDF5,        “H5Dwrite failed”); |
| /\* Preserve the error stack by assigning an object     \* handle to it     \*/     error\_stack = H5Eget\_current\_stack();       /\* Close dataset \*/     H5Dclose(dset\_id); |
| /\* Replace the current error stack with the     \* preserved one     \*/     H5Eset\_current\_stack(error\_stack);       Return(0);  } |
| Code Example 9-12. Registering the error stack |

# Properties and Property Lists in HDF5

## Introduction

HDF5 properties and property lists make it possible to shape or modify an HDF5 file, group, dataset, attri­bute, committed datatype, or even an I/O stream, in a number of ways. For example, you can do any of the following:

* Customize the storage layout of a file to suit a project or task.
* Create a chunked dataset.
* Apply compression or filters to raw data.
* Use either ASCII or UTF-8 character encodings.
* Create missing groups on the fly.
* Switch between serial and parallel I/O.
* Create consistency within a single file or across an international project.

Some properties enable an HDF5 application to take advantage of the capabilities of a specific computing environment while others make a file more compact; some speed the reading or writing of data while oth­ers enable more record-keeping at a per-object level. HDF5 offers nearly one hundred specific properties that can be used in literally thousands of combinations to maximize the usability of HDF5-stored data.

At the most basic level, a property list is a collection of properties, represented by name/value pairs that can be passed to various HDF5 functions, usually modifying default settings. A property list inherits a set of properties and values from a property list class. But that statement hardly provides a complete picture; in the rest of this section and in the next section, “Property List Classes, Property Lists, and Properties”, we will discuss these things in much more detail. After reading that material, the reader should have a reason­ably complete understanding of how properties and property lists can be used in HDF5 applications.

|  |
| --- |
| PropListEcosystem.png |
| Figure 10-1. The HDF5 property environment |

 The remaining sections in this chapter discuss the following topics:

•        What are properties, property lists, and property list classes?

•        Property list programming model

•        Generic property functions

•        Summary listings of property list functions

•        Additional resources

The discussions and function listings in this chapter focus on general property operations, object and link properties, and related functions.

File, group, dataset, datatype, and attribute properties are discussed in the chapters devoted to those fea­tures, where that information will be most convenient to users. For example, "HDF5 Datasets" discusses dataset creation property lists and functions, dataset access property lists and func­tions, and dataset transfer property lists and functions. This chapter does not duplicate those discussions.

Generic property operations are an advanced feature and are beyond the scope of this guide.

This chapter assumes an understanding of the following chapters of this HDF5 User’s Guide:

* "The HDF5 Data Model and File Structure"
* "The HDF5 Library and Programming Model"

## Property List Classes, Property Lists, and Properties

HDF5 property lists and the property list interface H5P provide a mechanism for storing characteristics of objects in an HDF5 file and economically passing them around in an HDF5 application. In this capacity, property lists significantly reduce the burden of additional function parameters throughout the HDF5 API. Another advantage of property lists is that features can often be added to HDF5 by adding only property list functions to the API; this is particularly true when all other requirements of the feature can be accom­plished internally to the library.

For instance, a file creation operation needs to know several things about a file, such as the size of the userblock or the sizes of various file data structures. Bundling this information as a property list simplifies the interface by reducing the number of parameters to the function H5Fcreate.

As illustrated in the figure above ("The HDF5 property environment"), the HDF5 property environment is a three-level hierarchy:

* Property list classes
* Property lists
* Properties

The following subsections discuss property list classes, property lists, and properties in more detail.

### Property List Classes

A property list class defines the roles that property lists of that class can play. Each class includes all prop­erties that are valid for that class with each property set to its default value. HDF5 offers a property lists class for each of the following situations.

|  |  |  |
| --- | --- | --- |
| Table 10-1. Property list classes in HDF5 | | |
| **Property List Class** |  | **For further discussion** |
| File creation (FCPL)  File access (FAPL)  File mount (FMPL) | H5P\_FILE\_CREATE  H5P\_FILE\_ACCESS  H5P\_FILE\_MOUNT | See various sections of "The HDF5 File".  Used only as H5P\_DEFAULT. For more information, see "File Mount Properties". |
| Object creation (OCPL)  Object copy (OCPYPL) | H5P\_OBJECT\_CREATE  H5P\_OBJECT\_COPY | See the table of "Object property functions (H5P)". |
| Group creation (GCPL)  Group access (GAPL) | H5P\_GROUP\_CREATE  H5P\_GROUP\_ACCESS | See "Programming Model for Groups". |
| Link creation (LCPL)  Link access (LAPL) | H5P\_LINK\_CREATE  H5P\_LINK\_ACCESS | See examples in "Programming Model for Properties and Property Lists" and the table of "Link creation property functions (H5P)". |
| Dataset creation (DCPL)  Dataset access (DAPL)  Dataset transfer (DXPL) | H5P\_DATASET\_CREATE  H5P\_DATASET\_ACCESS  H5P\_DATASET\_XFER | See "Programming Model for Datasets". |
| Datatype creation (TCPL)  Datatype access (TAPL) | H5P\_DATATYPE\_CREATE  H5P\_DATATYPE\_ACCESS | See various sections of "HDF5 Datatypes". |
| String creation (STRCPL) | H5P\_STRING\_CREATE | See "Programming Model for Datasets" and "Programming Model for Datatypes". |
| Attribute creation (ACPL) | H5P\_ATTRIBUTE\_CREATE | See "Working with Attributes". |

Note: In the table above, the abbreviations to the right of each property list class name in this table are widely used in both HDF5 programmer documentation and HDF5 source code. For example, FCPL is file creation property list, OCPL is object creation property list, OCPYPL is object copy property list, and STRCPL is string creation property list. These abbreviations may appear in either uppercase or lowercase.

The “HDF5 property list class inheritance hierarchy” figure, immediately following, illustrates the inheri­tance hierarchy of HDF5’s property list classes. Properties are defined at the root of the HDF5 property environment (“Property List Class Root” in the figure below). Property list classes then inherit properties from that root, either directly or indirectly through a parent class. In every case, a property list class inher­its only the properties relevant to its role. For example, the object creation property list class (OCPL) inher­its all properties that are relevant to the creation of any object while the group creation property list class (GCPL) inherits only those properties that are relevant to group creation.

|  |
| --- |
| PropListClassInheritance.png |
| Figure 10-2. HDF5 property list class inheritance hierarchy |

Note: In the figure above, property list classes displayed in black are directly accessible through the programming interface; the root of the property environment and the STRCPL and OCPL property list classes, in gray above, are not user-accessible. The red empty set symbol indicates that the file mount property list class (FMPL) is an empty class; that is, it has no settable properties. For more information, see "File Mount Properties". Abbreviations used in this figure are defined in the preceding table, “Property list classes in HDF5”.

### Property Lists

A property list is a collection of related properties that are used together in specific circumstances. A new property list created from a property list class inherits the properties of the property list class and each property’s default value. A fresh dataset creation property list, for example, includes all of the HDF5 prop­erties relevant to the creation of a new dataset.

Property lists are implemented as containers holding a collection of name/value pairs. Each pair specifies a property name and a value for the property. A property list usually contains information for one to many properties.

HDF5’s default property values are designed to be reasonable for general use cases. Therefore, an applica­tion can often use a property list without modification. On the other hand, adjusting property list settings is a routine action and there are many reasons for an application to do so.

A new property list may either be derived from a property list class or copied from an existing property list. When a property list is created from a property list class, it contains all the properties that are relevant to the class, with each property set to its default value. A new property list created by copying an existing property list will contain the same properties and property values as the original property list. In either case, the property values can be changed as needed through the HDF5 API.

Property lists can be freely reused to create consistency. For example, a single set of file, group, and data­set creation property lists might be created at the beginning of a project and used to create hundreds, thousands, even millions, of consistent files, file structures, and datasets over the project’s life. When such consistency is important to a project, this is an economical means of providing it.

### Properties

A property is the basic element of the property list hierarchy. HDF5 offers nearly one hundred properties controlling things ranging from file access rights, to the storage layout of a dataset, through optimizing the use of a parallel computing environment.

Further examples include the following:

|  |  |  |
| --- | --- | --- |
| **Purpose** | **Examples** | **Property List** |
| Specify the driver to be used to open a file | A POSIX driver or an MPI IO driver | FAPL |
| Specify filters to be applied to a data­set | Gzip compression or checksum evalu­ation | DCPL |
| Specify whether to record key times associated with an object | Creation time and/or last-modified time | OCPL |
| Specify the access mode for a file opened via an external link | Read-only or read-write | LAPL |

Each property is initialized with a default value. For each property, there are one or more dedicated H5Pset\_\* calls that can be used to change that value.

**Creation, access, and transfer properties:**

Properties fall into one of several major categories: creation properties, access properties, and transfer properties.

Creation properties control permanent object characteristics. These characteristics must be established when an object is created, cannot change through the life of the object (they are immutable), and the property setting usually has a permanent presence in the file.

Examples of creation properties include:

* Whether a dataset is stored in a compact, contiguous, or chunked layout

The default for this dataset creation property (H5Pset\_layout) is that a dataset is stored in a contiguous block. This works well for datasets with a known size limit that will fit easily in system memory.

A chunked layout is important if a dataset is to be compressed, to enable extending the dataset’s size, or to enable caching during I/O.

A compact layout is suitable only for very small datasets because the raw data is stored in the object header.

* Creation of intermediate groups when adding an object to an HDF5 file

This link creation property (H5Pset\_create\_intermediate\_group) enables an application to add an object in a file without having to know that the group or group hierarchy containing that object already exists. With this property set, HDF5 automatically creates missing groups. If this property is not set, an application must verify that each group in the path exists, and create those that do not, before creating the new object; if any group is missing, the create operation will fail.

* Whether an HDF5 file is a single file or a set of tightly related files that form a virtual HDF5 file

Certain file creation properties enable the application to select one of several file layouts. Exam­ples of the available layouts include a standard POSIX-compliant layout (H5Pset\_fapl\_sec2), a family of files (H5Pset\_fapl\_family), and a split file layout that separates raw data and meta­data into separate files (H5Pset\_fapl\_split). These and other file layout options are discussed in "Alternate File Storage Layouts and Low-level File Drivers".

* To enable error detection when creating a dataset

In settings where data integrity is vulnerable, it may be desirable to set checksumming when data­sets are created (H5Pset\_fletcher32). A subsequent application will then have a means to ver­ify data integrity when reading the dataset.

Access properties control transient object characteristics. These characteristics may change with the cir­cumstances under which an object is accessed.

Examples of access properties include:

* The driver used to open a file

For example, a file might be created with the MPI I/O driver (H5Pset\_fapl\_mpio) during high-speed data acquisition in a parallel computing environment. The same file might later be analyzed in a serial computing environment with I/O access handled through the serial POSIX driver (H5Pset\_fapl\_sec2).

* Optimization settings in specialized environments

Optimizations differ across computing environments and according to the needs of the task being performed, so are transient by nature.

Transfer properties apply only to datasets and control transient aspects of data I/O. These characteristics may change with the circumstances under which data is accessed.

Examples of dataset transfer properties include:

* To enable error detection when reading a dataset

If checksumming has been set on a dataset (with H5Pset\_fletcher32, in the dataset creation property list), an application reading that dataset can choose whether check for data integrity (H5Pset\_edc\_check).

* Various properties to optimize chunked data I/O on parallel computing systems

HDF5 provides several properties for tuning I/O of chunked datasets in a parallel computing envi­ronment (H5Pset\_dxpl\_mpio\_chunk\_opt, H5Pset\_dxpl\_mpio\_chunk\_opt\_num, H5Pset\_dxpl\_mpio\_chunk\_opt\_ratio, and H5Pget\_mpio\_actual\_chunk\_opt\_mode).

Optimal settings differ due to the characteristics of a computing environment and due to an appli­cation’s data access patterns; even when working with the same file, these settings might change for every application and every platform.

## Programming Model for Properties and Property Lists

The programming model for HDF5 property lists is actually quite simple:

1.      Create a property list.

2.      Modify the property list, if required.

3.      Use the property list.

4.      Close the property list.

There are nuances, of course, but that is the basic process.

In some cases, you will not have to define property lists at all. If the default property settings are sufficient for your application, you can tell HDF5 to use the default property list.

The following sections first discuss the use of default property lists, then each step of the programming model, and finally a few less frequently used property list operations.

### Using Default Property Lists

Default property lists can simplify many routine HDF5 tasks because you do not always have to create every property list you use.

An application that would be well-served by HDF5’s default property settings can use the default property lists simply by substituting the value H5P\_DEFAULT for a property list identifier. HDF5 will then apply the default property list for the appropriate property list class.

For example, the function H5Dcreate2 calls for a link creation property list, a dataset creation property list, and a dataset access property list. If the default properties are suitable for a dataset, this call can be made as

dset\_id = H5Dcreate2( loc\_id, name, dtype\_id, space\_id, H5P\_DEFAULT, H5P\_DEFAULT,

H5P\_DEFAULT );

HDF5 will then apply the default link creation, dataset creation, and dataset access property lists correctly.

Of course, you would not want to do this without considering where it is appropriate, as there may be unforeseen consequences. Consider, for example, the use of chunked datasets. Optimal chunking is quite dependent on the makeup of the dataset and the most common access patterns, both of which must be taken into account in setting up the size and shape of chunks.

### Basic Steps of the Programming Model

The steps of the property list programming model are described in the sub-sections below.

#### Create a Property List

A new property list can be created either as an instance of a property list class or by copying an existing property list. Consider the following examples. A new dataset creation property list is first created “from scratch” with H5Pcreate. A second dataset creation property list is then created by copying the first one with H5Pcopy.

dcplA\_id = H5Pcreate (H5P\_DATASET\_CREATE);

The new dataset creation property list is created as an instance of the property list class H5P\_­DATASET\_CREATE.

The new dataset creation  property list’s identifier is returned in dcplA\_id and the property list is initialized with default dataset creation property values.

A list of valid classes appears in the table "Property list classes in HDF5".

dcplB\_id = H5Pcopy (dcplA\_id);

A new dataset creation property list, dcplB\_id, is created as a copy of dcplA\_id and is initial­ized with dataset creation property values currently in dcplA\_id.

At this point, dcplA\_id and dcplB\_id are identical; they will both contain any modified property values that were changed in dcplA\_id before dcplB\_id was created. They may, however, diverge as additional property values are reset in each.

While we are creating property lists, let’s create a link creation property list; we will need this property list when the new dataset is linked into the file below:

lcplAB\_id = H5Pcreate (H5P\_LINK\_CREATE);

#### 10.3.2.2. Change Property Values

This section describes how to set property values.

Later in this section, the dataset creation property lists dcplA\_id and dcplB\_id created in the section above will be used respectively to create chunked and contiguous datasets. To set this up, we must set the layout property in each property list. The following example sets dcplA\_id for chunked datasets and dcplB\_id for contiguous datasets:

error = H5Pset\_layout (dcplA\_id, H5D\_CHUNKED);

error = H5Pset\_layout (dcplB\_id, H5D\_CONTIGUOUS);

Since dcplA\_id specifies a chunked layout, we must also set the number of dimensions and the size of the chunks. The example below specifies that datasets created with dcplA\_id will be 3-dimensional and that the chunk size will be 100 in each dimension:

error = H5Pset\_chunk (dcplA\_id, 3, [100,100,100]);

These datasets will be created with UTF-8 encoded names. To accomplish that, the following example sets the character encoding property in the link creation property list to create link names with UTF-8 encod­ing:

error = H5Pset\_char\_encoding (lcplAB\_id, H5T\_CSET\_UTF8);

dcplA\_id can now be used to create chunked datasets and dcplB\_id to create contiguous datasets. And with the use of lcplAB\_id, they will be created with UTF-8 encoded names.

#### Use the Property List

Once the required property lists have been created, they can be used to control various HDF5 processes. For illustration, consider dataset creation.

Assume that the datatype dtypeAB and the dataspaces dspaceA and dspaceB have been defined and that the location identifier locAB\_id specifies the group AB in the current HDF5 file. We have already cre­ated the required link creation and dataset creation property lists.  For the sake of illustration, we assume that the default dataset access property list meets our application requirements. The following calls would create the datasets dsetA and dsetB in the group AB. The raw data in dsetA will be contiguous while dsetB raw data will be chunked; both datasets will have UTF-8 encoded link names:

dsetA\_id = H5Dcreate2( locAB\_id, dsetA, dtypeAB, dspaceA\_id,

            lcplAB\_id, dcplA\_id, H5P\_DEFAULT );

dsetB\_id = H5Dcreate2( locAB\_id, dsetB, dtypeAB, dspaceB\_id,

            lcplAB\_id, dcplB\_id, H5P\_DEFAULT );

#### Close the Property List

Generally, creating or opening anything in an HDF5 file results in an HDF5 identifier. These identifiers are of HDF5 type hid\_t and include things like file identifiers, often expressed as file\_id; dataset identifi­ers, dset\_id; and property list identifiers, plist\_id. To reduce the risk of memory leaks, all of these identifiers must be closed once they are no longer needed.

Property list identifiers are no exception to this rule, and H5Pclose is used for this purpose. The calls immediately following would close the property lists created and used in the examples above.

error = H5Pclose (dcplA\_id);

error = H5Pclose (dcplB\_id);

error = H5Pclose (lcplAB\_id);

### Additional Property List Operations

A few property list operations fall outside of the programming model described above. This section describes those operations.

#### Query the Class of an Existing Property List

Occasionally an application will have a property list but not know the corresponding property list class. A call such as in the following example will retrieve the unknown class of a known property list:

PList\_Class = H5Pget\_class (dcplA\_id);

Upon this function’s return, PList\_Class will contain the value H5P\_DATASET\_CREATE indicating that dcplA\_id is a dataset creation property list.

#### Determine Current Creation Property List Settings in an Existing Object

After a file has been created, another application may work on the file without knowing how the creation properties for the file were set up. Retrieving these property values is often unnecessary; HDF5 can read the data and knows how to deal with any properties it encounters.

But sometimes an application must do something that requires knowing the creation property settings. HDF5 makes the acquisition of this information fairly straight-forward; for each property setting call, H5Pset\_\*, there is a corresponding H5Pget\_\* call to retrieve the property’s current setting.

Consider the following examples which illustrate the determination of dataset layout and chunking set­tings:

The application must first identify the creation property list with the appropriate get creation property list call. There is one such call for each kind of object.

H5Dget\_create\_plist will return a property list identifier for the creation property list that was used to create the dataset. Call it DCPL1\_id.

H5Pset\_layout sets a dataset’s layout to be compact, contiguous, or chunked.

H5Pget\_layout called with DCPL1\_id will return the dataset’s layout, either H5D\_­COMPACT, H5D\_CONTIGUOUS, or H5D\_CHUNKED.

H5Pset\_chunk sets the rank of a dataset, that is the number of dimensions it will have, and the maximum size of each dimension.

H5Pget\_chunk, also called with DCPL1\_id, will return the rank of the dataset and the maximum size of each dimension.

If a creation property value has not been explicitly set, these H5Pget\_ calls will return the property’s default value.

#### Determine Access Property Settings

Access property settings are quite different from creation properties. Since access property settings are not retained in an HDF5 file or object, there is normally no knowledge of the settings that were used in the past. On the other hand, since access properties do not affect characteristics of the file or object, this is not normally an issue. For more information, see "Access and Creation Property Exceptions."

One circumstance under which an application might need to determine access property settings might be when a file or object is already open but the application does not know the property list settings. In that case, the application can use the appropriate get access property list call to retrieve a property list identi­fier. For example, if the dataset dsetA from the earlier examples is still open, the following call would return an identifier for the dataset access property list in use:

dsetA\_dacpl\_id = H5Dget\_access\_plist( dsetA\_id );

The application could then use the returned property list identifier to analyze the property settings.

## Generic Properties Interface and User-defined Properties

HDF5’s generic property interface provides tools for managing the entire property hierarchy and for the creation and management of user-defined property lists and properties. This interface also makes it possi­ble for an application or a driver to create, modify, and manage custom properties, property lists, and property list classes. A comprehensive list of functions for this interface appears under “Generic Property Operations (Advanced)” in the “H5P: Property List Interface” section of the HDF5 Reference Manual.

Further discussion of HDF5’s generic property interface and user-defined properties and property lists is beyond the scope of this document.

## Property List Function Summaries

General property functions, generic property functions and macros, property functions that are used with multiple types of objects, and object and link property functions are listed below.

Property list functions that apply to a specific type of object are listed in the chapter that discusses that object. For example, the Datasets chapter has two property list function listings: one for dataset creation property list functions and one for dataset access property list functions. As has been stated, this chapter is not intended to describe every property list function.

|  |  |
| --- | --- |
| Function Listing 10-1. General property list functions (H5P) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Pcreate  h5pcreate\_f | Creates a new property list as an instance of a specified parent property list class. |
| H5Pcopy  h5pcopy\_f | Creates a new property list by copying the specified existing property list. |
| H5Pget\_class  h5pget\_class\_f | Retrieves the parent property list class of the specified property list. |
| H5Pclose  h5pclose\_f | Closes the specified property list. |

Object property functions can be used with several kinds of objects.

|  |  |
| --- | --- |
| Function Listing 10-2. Object property functions (H5P) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| **Object Creation Properties** |  |
| H5Pget\_attr\_creation\_order  h5pget\_attr\_creation\_order\_f | Retrieves tracking and indexing settings for attribute creation order. |
| H5Pget\_attr\_phase\_change  h5pget\_attr\_phase\_change\_f | Retrieves attribute storage phase change thresholds. |
| H5Pget\_obj\_track\_times  h5pget\_obj\_track\_times\_f | Determines whether times associated with an object are being recorded. |
| H5Pset\_attr\_creation\_order  h5pset\_attr\_creation\_order\_f | Sets tracking and indexing of attribute creation order. |
| H5Pset\_attr\_phase\_change  h5pset\_attr\_phase\_change\_f | Sets attribute storage phase change thresholds. |
| H5Pset\_obj\_track\_times  h5pset\_obj\_track\_times\_f | Sets the recording of times associated with an object. |
| **Object Copy Properties** |  |
| H5Padd\_merge\_committed\_dtype\_path  (no Fortran subroutine) | Adds a path to the list of paths that will be searched in the destination file for a matching committed datatype. |
| H5Pfree\_merge\_committed\_dtype\_paths  (no Fortran subroutine) | Clears the list of paths stored in an object copy property list. |
| H5Pget\_copy\_object  h5pget\_copy\_object\_f | Retrieves the properties to be used when an object is copied. |
| H5Pget\_mcdt\_search\_cb  (no Fortran subroutine) | Retrieves the callback function from the specified object copy property list. |
| H5Pset\_copy\_object  h5pset\_copy\_object\_f | Sets the properties to be used when an object is copied. |
| H5Pset\_mcdt\_search\_cb  (no Fortran subroutine) | Sets the callback function that H5Ocopy will invoke before searching the entire destination file for a matching committed datatype. |

The following table lists link creation properties. Since the creation of a link is almost always a step in the creation of an object, these properties may also be set in group creation property lists, dataset creation property lists, datatype creation property lists, and the more generic object creation property lists. Some are also applicable to the attribute creation property lists.

|  |  |
| --- | --- |
| Function Listing 10-3. Link creation property functions (H5P) | |
| **C Function**  **Fortran Subroutine** | **Purpose** |
| H5Pget\_char\_encoding  h5pget\_char\_encoding\_f | Queries the character encoding used to encode link or attribute names.  Note: Use with link, object, dataset, datatype, group, or attribute creation property lists. |
| H5Pset\_char\_encoding  h5pset\_char\_encoding\_f | Sets the character encoding used to encode link and attribute names.  Note: Use with link, object, dataset, datatype, group, or attribute creation property lists. |
| H5Pget\_create\_intermediate\_group  h5pget\_create\_intermediate\_group\_f | Queries setting for creation of intermedi­ate groups.  Note: Use with link creation property lists, which in turn can be used in the create call for any dataset, datatype, or group. |
| H5Pset\_create\_intermediate\_group  h5pset\_create\_intermediate\_group\_f | Specifies whether to create intermediate groups when they do not already exist.  Note: Use with link creation property lists, which in turn can be used in the create call for any dataset, datatype, or group. |

Note: In the function listing above, the properties can be used with any of the indicated property lists.

## Additional Property List Resources

Property lists are ubiquitous in an HDF5 environment and are therefore discussed in many places in HDF5 documentation. The following sections and listings in the HDF5 User’s Guide are of particular interest:

* In the “HDF5 Data Model and File Structure” chapter, see "Property List".
* In the “HDF5 File” chapter, see the following sections and listings:
  + "File Creation and File Access Properties"
  + "File Property Lists"
  + "Example with the File Creation Property List"
  + "Example with the File Access Property List"
  + "File creation property list functions (H5P)"
  + "File access property list functions (H5P)"
  + "File driver functions (H5P)"
* In the “HDF5 Attributes” chapter, see "Attribute creation property list functions (H5P)".
* In the “HDF5 Groups” chapter, see "Group creation property list functions (H5P)".
* Property lists are discussed throughout "HDF5 Datasets".

All property list functions are described in the “H5P: Property List Interface” section of the HDF5 Reference Manual. The function index at the top of the page provides a categorized listing grouped by property list class. Those classes are listed below:

* *File creation properties*
* *File access properties*
* *Group creation properties*
* *Dataset creation properties*
* *Dataset access properties*
* *Dataset transfer properties*
* *Link creation properties*
* *Link access properties*
* *Object creation properties*
* *Object copy properties*

Additional categories not related to the class structure are as follows:

* *General property list operations*
* *Generic property list functions*

The general property functions can be used with any property list; the generic property functions constitute an advanced feature.

The in-memory file image feature of HDF5 uses property lists in a manner that differs substantially from their use elsewhere in HDF5. Those who plan to use in-memory file images must study “File Image Opera­tions” (PDF) in the Advanced Topics in HDF5 collection.

## Notes

**File Mount Prope****rti****es**

While the file mount property list class H5P\_FILE\_MOUNT is a valid HDF5 property list class, no file mount properties are defined by the HDF5 Library. References to a file mount property list should always be expressed as H5P\_DEFAULT, meaning the default file mount property list.

**Access and Creation Property Excepti****o****ns**

There are a small number of exceptions to the rule that creation properties are always retained in a file or object and access properties are never retained.

The following properties are file access properties but they are not transient; they have permanent and different effects on a file. They could be validly classified as file creation properties as they must be set at creation time to properly create the file. But they are access properties because they must also be set when a file is reopened to properly access the file.

|  |  |
| --- | --- |
| **Property** | **Related function** |
| Family file driver | H5Pset\_fapl\_family |
| Split file driver | H5Pset\_fapl\_split |
| Core file driver | H5Pset\_fapl\_core |

The following is a link creation property, but it is not relevant after an object has been created and is not retained in the file or object.

|  |  |
| --- | --- |
| **Property** | **Related function** |
| Create missing intermediate groups | H5Pset\_create\_intermediate\_groups |

# Additional Resources

These documents provide additional information for the use and tuning of specific HDF5 features.

|  |  |
| --- | --- |
| Table 11-1. Additional resources | |
| **Document** | **Comments** |
| *HDF5 Examples* | Code examples by API. |
| *Chunking in HDF5* | Structuring the use of chunking and tun­ing it for performance. |
| *Using the Direct Chunk Write Function* | Describes another way that chunks can be written to datasets. |
| *Copying Committed Datatypes with H5Ocopy* | Describes how to copy to another file a dataset that uses a committed datatype or an object with an attribute that uses a committed datatype so that the commit­ted datatype in the destination file can be used by multiple objects. |
| *Metadata Caching in HDF5* | Managing the HDF5 metadata cache and tuning it for performance. |
| *HDF5 Dynamically Loaded Filters* | Describes how an HDF5 application can apply a filter that is not registered with the HDF5 Library. |
| *HDF5 File Image Operations* | Describes how to work with HDF5 files in memory. Disk I/O is not required when file images are opened, created, read from, or written to. |
| *Modified Region Writes* | Describes how to set write operations for in-memory files so that only modified regions are written to storage. Available when the Core (Memory) VFD is used. |
| *Using Identifiers* | Describes how identifiers behave and how they should be treated. |
| *Using UTF-8 Encoding in HDF5 Applications* | Describes the use of UTF-8 Unicode char­acter encodings in HDF5 applications. |
| *Freeing Memory Allocated by the HDF5 Library* | Describes how inconsistent memory man­agement can cause heap corruption or resource leaks and possible solutions. |
| *HDF5 Glossary* | A glossary of terms. |