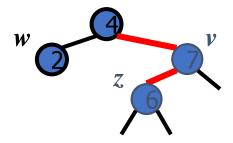
Red-Black Tree Example

Remedying a Double Red

 Consider a double red with child z and parent v, and let w be the sibling of v

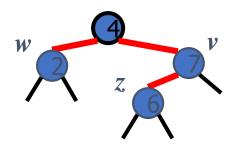
Case 1: *sibling of first red is black*

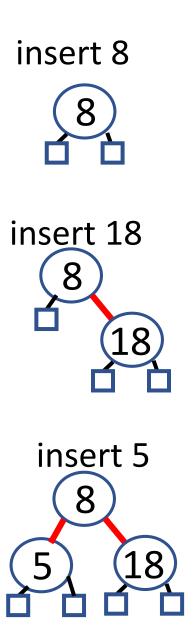
- The double red is an incorrect replacement of a 4-node
- Restructuring: we change the 4-node replacement

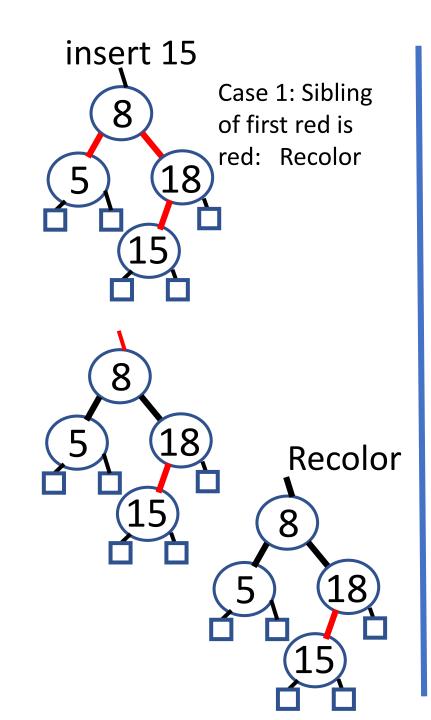


Case 2: sibling of first red is red

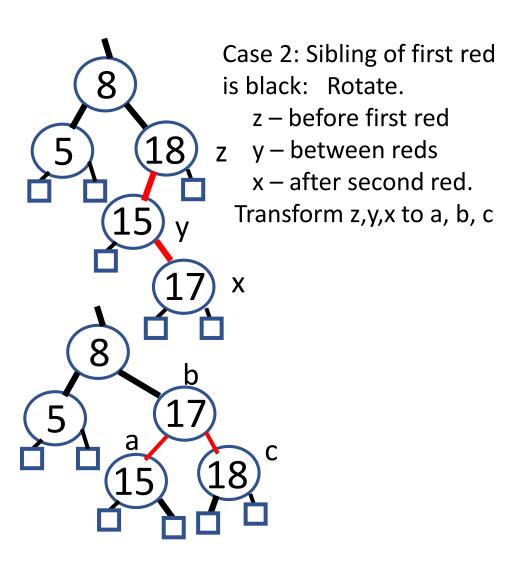
- The double red corresponds to an overflow
- Recoloring: we perform the equivalent of a split

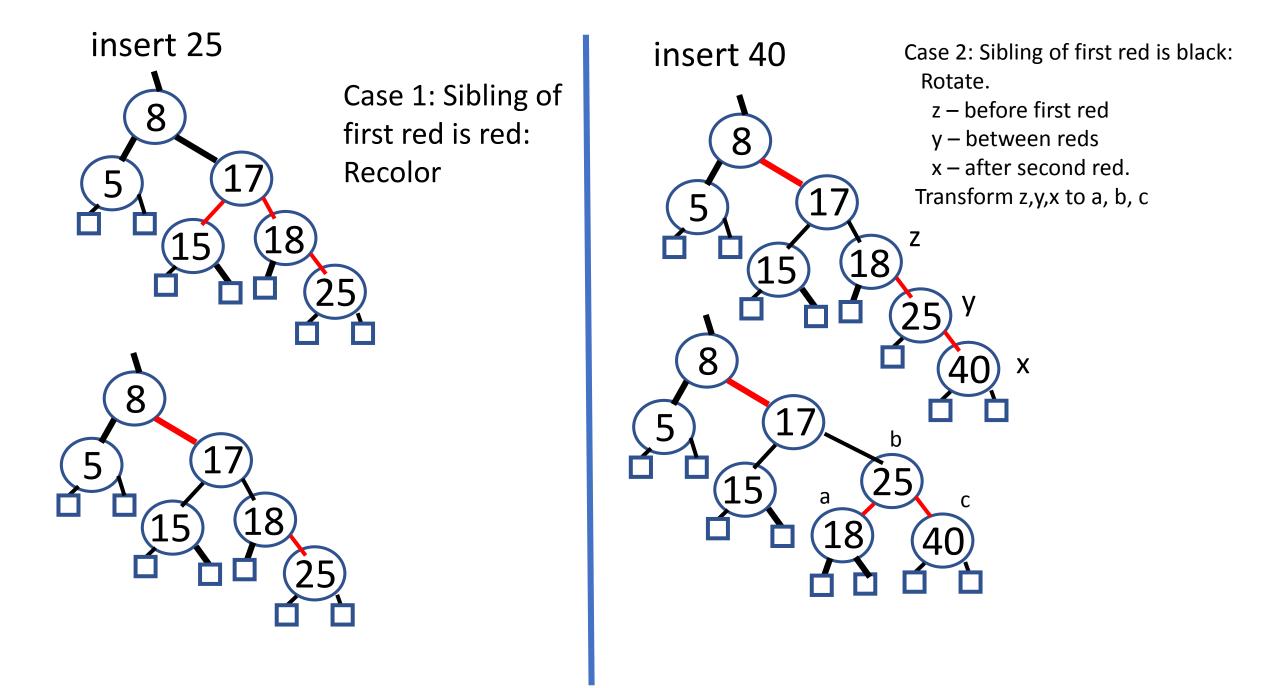


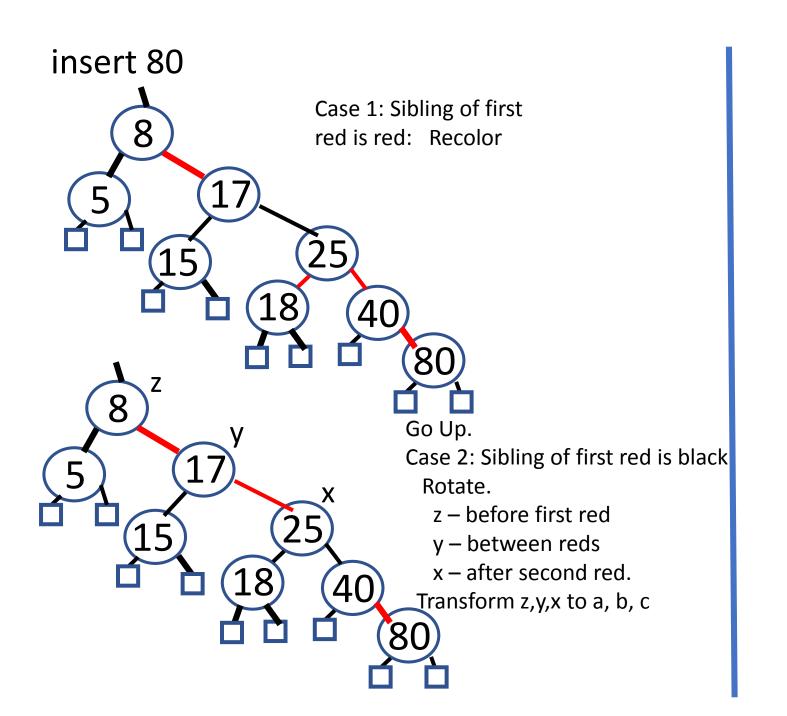


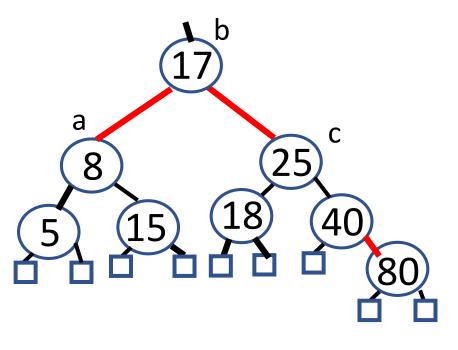


insert 17









Remedying a Double Black

• The algorithm for remedying a double black node w with sibling y considers three cases

Case 1: y is black and has a red child

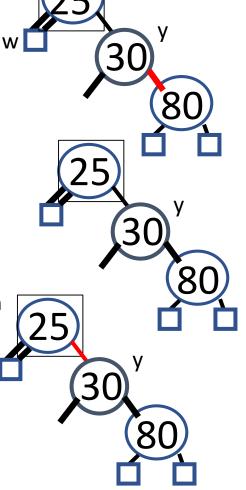
We perform a restructuring, equivalent to a transfer, and we are done

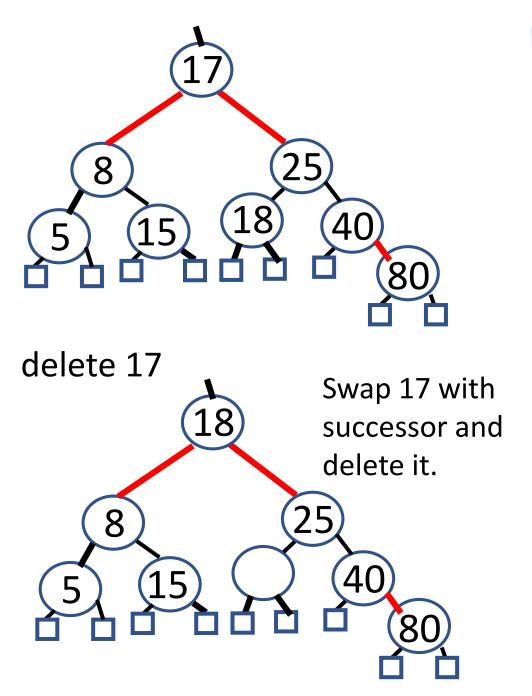
Case 2: y is black and its children are both black

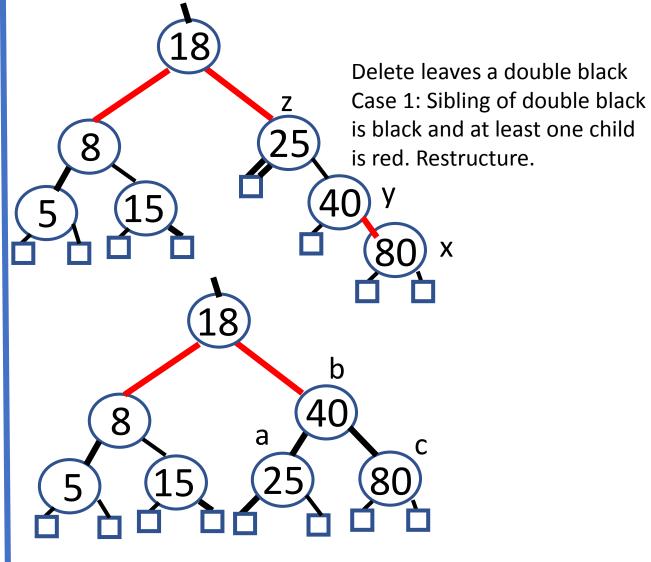
 We perform a recoloring, equivalent to a fusion, which may propagate up the double black violation

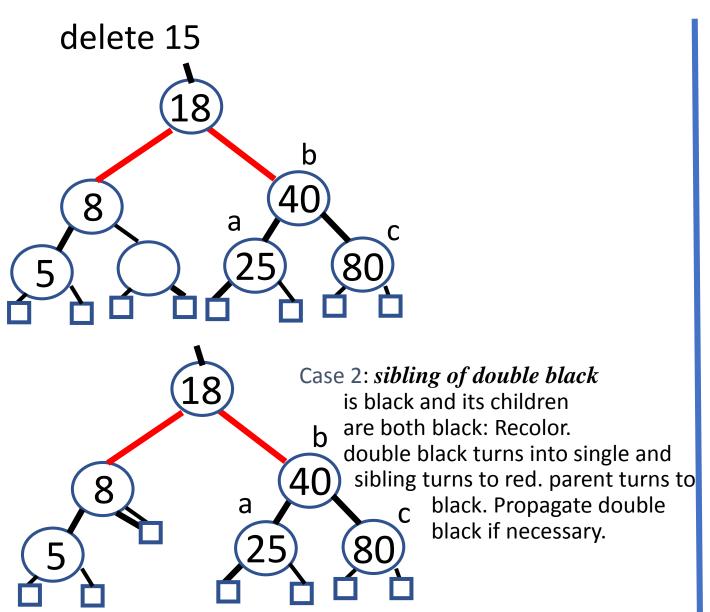
Case 3: y is red

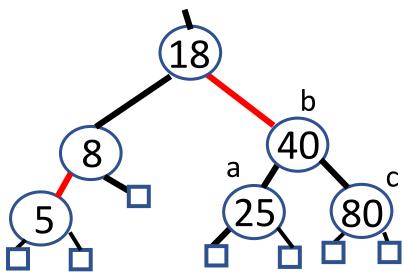
- We perform an adjustment, equivalent to choosing a different representation of a 3-node, after which either Case 1 or Case 2 applies
- Deletion in a red-black tree takes $O(\log n)$ time











delete 5

