

CHAPTER 5

The Relational Data Model and Relational Database Constraints

Chapter Outline

- Relational Model Concepts
- Relational Model Constraints and Relational Database Schemas
- Update Operations and Dealing with Constraint Violations

Relational Model Concepts

- The relational Model of Data is based on the concept of a Relation
 - The strength of the relational approach to data management comes from the formal foundation provided by the theory of relations
- We review the essentials of the formal relational model in this chapter
- In practice, there is a standard model based on SQL this
 is described in Chapters 6 and 7 as a language
- Note: There are several important differences between the formal model and the practical model, as we shall see

Relational Model Concepts

- A Relation is a mathematical concept based on the ideas of sets
- The model was first proposed by Dr. E.F. Codd of IBM Research in 1970 in the following paper:
 - "A Relational Model for Large Shared Data Banks," Communications of the ACM, June 1970
- The above paper caused a major revolution in the field of database management and earned Dr.
 Codd the coveted ACM Turing Award

Informal Definitions

- Informally, a relation looks like a table of values.
- A relation typically contains a set of rows.
- The data elements in each row represent certain facts that correspond to a real-world entity or relationship
 - In the formal model, rows are called tuples
- Each column has a column header that gives an indication of the meaning of the data items in that column
 - In the formal model, the column header is called an attribute name (or just attribute)

Example of a Relation

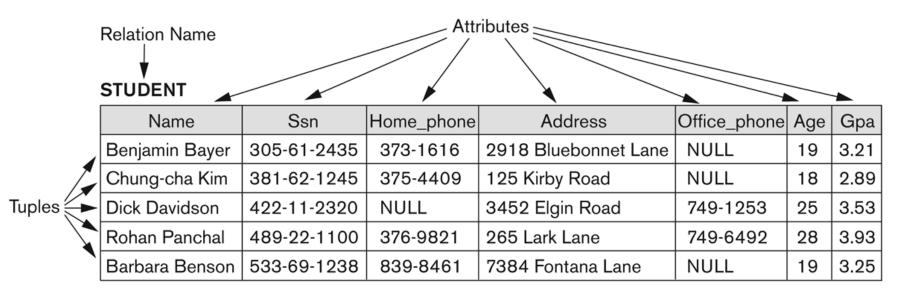


Figure 5.1

The attributes and tuples of a relation STUDENT.

Informal Definitions

- Key of a Relation:
 - Each row has a value of a data item (or set of items)
 that uniquely identifies that row in the table
 - Called the key
 - In the STUDENT table, SSN is the key
 - Sometimes row-ids or sequential numbers are assigned as keys to identify the rows in a table
 - Called artificial key or surrogate key

Formal Definitions - Schema

- The **Schema** (or description) of a Relation:
 - Denoted by R(A1, A2,An)
 - R is the name of the relation
 - The attributes of the relation are A1, A2, ..., An
- Example:
 - CUSTOMER (Cust-id, Cust-name, Address, Phone#)
 - CUSTOMER is the relation name
 - Defined over the four attributes: Cust-id, Cust-name, Address, Phone#
- Each attribute has a domain or a set of valid values.
 - For example, the domain of Cust-id is 6 digit numbers.

Formal Definitions - Tuple

- A tuple is an ordered set of values (enclosed in angled brackets '< ... >')
- Each value is derived from an appropriate domain.
- A row in the CUSTOMER relation is a 4-tuple and would consist of four values, for example:
 - <632895, "John Smith", "101 Main St. Atlanta, GA 30332", "(404) 894-2000">
 - This is called a 4-tuple as it has 4 values
 - A tuple (row) in the CUSTOMER relation.
- A relation is a set of such tuples (rows)

Formal Definitions - Domain

- A domain has a logical definition:
 - Example: "USA_phone_numbers" are the set of 10 digit phone numbers valid in the U.S.
- A domain also has a data-type or a format defined for it.
 - The USA_phone_numbers may have a format: (ddd)ddd-dddd where each d is a decimal digit.
 - Dates have various formats such as year, month, date formatted as yyyy-mm-dd, or as dd mm,yyyy etc.
- The attribute name designates the role played by a domain in a relation:
 - Used to interpret the meaning of the data elements corresponding to that attribute
 - Example: The domain Date may be used to define two attributes named "Invoice-date" and "Payment-date" with different meanings

Formal Definitions - State

- The relation state is a subset of the Cartesian product of the domains of its attributes
 - each domain contains the set of all possible values the attribute can take.
- Example: attribute Cust-name is defined over the domain of character strings of maximum length
 25
 - dom(Cust-name) is varchar(25)
- The role these strings play in the CUSTOMER relation is that of the name of a customer.

Formal Definitions - Summary

- Formally,
 - Given R(A1, A2,, An)
 - $r(R) \subset dom(A1) \times dom(A2) \times \times dom(An)$
- R(A1, A2, ..., An) is the **schema** of the relation
- R is the name of the relation
- A1, A2, ..., An are the attributes of the relation
- r(R): a specific state (or "value" or "population") of relation R – this is a set of tuples (rows)
 - r(R) = {t1, t2, ..., tn} where each ti is an n-tuple
 - ti = <v1, v2, ..., vn> where each vj element-of dom(Aj)

Formal Definitions - Example

- Let R(A1, A2) be a relation schema:
 - Let dom(A1) = {0,1}
 - Let dom(A2) = {a,b,c}
- Then: dom(A1) X dom(A2) is all possible combinations:

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{<0,a>, <0,b>, <0,c>, <1,a>, <1,b>, <1,c>}
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- The relation state r(R) ⊂ dom(A1) X dom(A2)
- For example: r(R) could be {<0,a>, <0,b>, <1,c>}
 - this is one possible state (or "population" or "extension") r of the relation R, defined over A1 and A2.
 - It has three 2-tuples: <0,a> , <0,b> , <1,c>

Definition Summary

<u>Informal Terms</u>	Formal Terms
Table	Relation
Column Header	Attribute
All possible Column Values	Domain
Row	Tuple
Table Definition	Schema of a Relation
Populated Table	State of the Relation

Example – A relation STUDENT

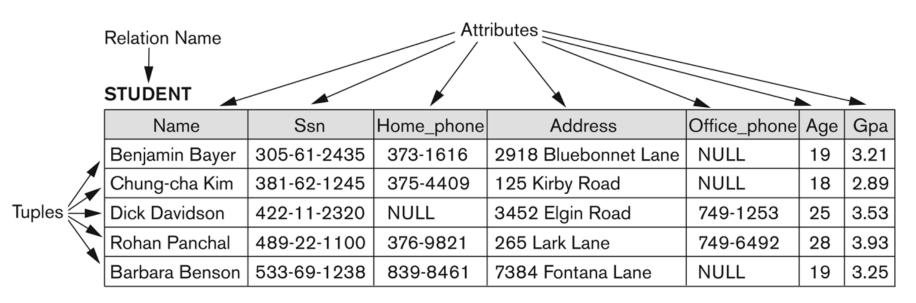


Figure 5.1
The attributes and tuples of a relation STUDENT.

Characteristics Of Relations

- Ordering of tuples in a relation r(R):
 - The tuples are not considered to be ordered, even though they appear to be in the tabular form.
- Ordering of attributes in a relation schema R (and of values within each tuple):
 - We will consider the attributes in R(A1, A2, ..., An) and the values in t=<v1, v2, ..., vn> to be ordered.
 - (However, a more general alternative definition of relation does not require this ordering. It includes both the name and the value for each of the attributes).
 - Example: t= { <name, "John" >, <SSN, 123456789> }
 - This representation may be called as "self-describing".

Same state as previous Figure (but with different order of tuples)

Figure 5.2

The relation STUDENT from Figure 5.1 with a different order of tuples.

STUDENT

Name	Ssn	Home_phone	Address	Office_phone	Age	Gpa
Dick Davidson	422-11-2320	NULL	3452 Elgin Road	749-1253	25	3.53
Barbara Benson	533-69-1238	839-8461	7384 Fontana Lane	NULL	19	3.25
Rohan Panchal	489-22-1100	376-9821	265 Lark Lane	749-6492	28	3.93
Chung-cha Kim	381-62-1245	375-4409	125 Kirby Road	NULL	18	2.89
Benjamin Bayer	305-61-2435	373-1616	2918 Bluebonnet Lane	NULL	19	3.21

Characteristics Of Relations

- Values in a tuple:
 - All values are considered atomic (indivisible).
 - Each value in a tuple must be from the domain of the attribute for that column
 - If tuple t = <v1, v2, ..., vn> is a tuple (row) in the relation state r of R(A1, A2, ..., An)
 - Then each *vi* must be a value from *dom(Ai)*
 - A special **null** value is used to represent values that are unknown or not available or inapplicable in certain tuples.

Characteristics Of Relations

Notation:

- We refer to component values of a tuple t by:
 - t[Ai] or t.Ai
 - This is the value vi of attribute Ai for tuple t
- Similarly, t[Au, Av, ..., Aw] refers to the subtuple of t containing the values of attributes Au, Av, ..., Aw, respectively in t