

COL352 HOMEWORK 1

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3)To show that:-

Show that regular languages are closed under the repeat operation, where repeat operation on a language L is given by

$$\text{repeat}(L) = \{l_1l_1l_1l_2\dots l_k \mid s \in L\}$$

Sol:

Let $A_1 \Rightarrow (Q_1, \Sigma, q_0^1, F_1, \delta_1)$ be the automata accepting L .

Let A be a finite state automata then $(Q, \Sigma, q_0, F, \delta)$ such that

$$Q \Rightarrow Q_1$$

$$q_0 \Rightarrow q_0^1$$

$$F \Rightarrow F_1$$

$$\delta(q, a) \Rightarrow \{q, \delta_1(q, a)\} \text{ where } a \in \Sigma.$$

4) Design an algorithm that takes as input the descriptions of two DFAs, D1 and D2, and determines whether they recognize the same language.

Sol: $L1 \Rightarrow L(D1)$, $L2 \Rightarrow L(D2)$

Lets say $L3 \Rightarrow L2 \text{ dash} \Rightarrow L(D2 \text{ dash})$

$L4 \Rightarrow L1 \cap L3$.

If $L4 \Rightarrow \{ \emptyset \}$ then D1 and D2 recognize the same language .

$L4 \Rightarrow L(D) \Rightarrow \{ \emptyset \}$

if D has no path from starting state to final state

5) To prove that :-

Show that if A is regular, then so is A^R . In other

words, regular languages are closed under the reverse operation

sol: Let $B \Rightarrow (Q_1, \Sigma, q_0^1, F_1, \delta_1)$ be the automata accepting A .

Let $B^R \Rightarrow (Q_2, \Sigma, q_0^2, F_2, \delta_2)$ Such that

$$Q_2 \Rightarrow Q_1$$

$$Q_0^2 \Rightarrow q \mid q \in F_1$$

$$F_2 \Rightarrow q_0^1$$

$$\delta_2(q, a) \Rightarrow \delta_1^{-1}(q, a)$$

correctness of proof:

$w \in \Sigma^*$, w is accepted by B^R if w^R is accepted by B .

2)To prove:

All NFAs recognize the class of regular languages

Sol: All NFA $M (Q, \Sigma, q_0, F, \delta)$ that accepts x belongs to Σ^* , and every possible state M could be taking input x is a state from F .

There are 2 steps to prove that –NFAs recognize of regular languages.

1)every regular language is recognized by some /all NFA

2) every all/NFA recognizes a regular language

Step1)

Let L be the any regular language and M be the DFA that recognizes L . for every input string there is exactly one possible state q belong Q that M could be after reading x .

So, if M is shown as all NFA then it accept x iff (if and only if) q belongs F , which happens iff x belongs to L . from the above every regular language is recognized by some/all NFA

Step2)

Let A be the language recognized by N , then proof is accordingly given below to prove A is regular.

- Constructing a DFA that recognizes $M = (Q, \Sigma, q_0, F, \delta)$ that recognizes A .

For any subset $S \subseteq Q$, $E(S)$ is the set of all states q belongs Q that can reach from S .

Therefore every all NFA recognizes a regular language from step1 and step2 all NFA recognizes the class of regular language

1)construct an NFA that accepts strings that don't have all the characters from Σ . Below I have explained the NFA and drawn the construction of NFA of l_1, l_2, \dots, l_k

