

COL788: Advanced Topics in Embedded Computing

Lecture 15 – Embedded OS



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Outline

- Previously
 - System Architecture
 - Processor Architecture
 - Memory
- Next
 - Boot Sequence
 - **Operating System**
 - Embedded Linux

Contiki OS

- [Features](#)
 - Memory Allocation
 - Full IP Networking
 - Power Awareness
 - 6Lowpan, RPL, CoAP
 - Dynamic Module Loading
 - Protothreads
 - File System
 - Contiki Shell

CooCox CoOS

- [Features](#)

- Free and open real-time Operating System
- Specially designed for Cortex-M series
- Scalable, minimum system kernel is only 974Bytes
- Adaptive Task Scheduling Algorithm.
- Supports preemptive priority and round-robin
- Interrupt latency is next to 0
- Stack overflow detection option
- Semaphore, Mutex, Flag, Mailbox and Queue for communication & synchronisation
- Supports the platforms of ICCARM, ARMCC, GCC

TinyOS

- [Overview](#)
 - TinyOS is an "operating system" designed for low-power wireless embedded systems. Fundamentally, it is a work scheduler and a collection of drivers for microcontrollers and other ICs commonly used in wireless embedded platforms.

FreeRTOS

- [What is FreeRTOS?](#)
 - FreeRTOS is a class of RTOS that is designed to be small enough to run on a microcontroller
- [What is an RTOS?](#)
 - A part of the operating system called the scheduler is responsible for deciding which program to run when and provides the illusion of simultaneous execution by rapidly switching between each program.
 - The scheduler in a Real Time Operating System (RTOS) is designed to provide a predictable (normally described as *deterministic*) execution pattern.

Embedded Operating System

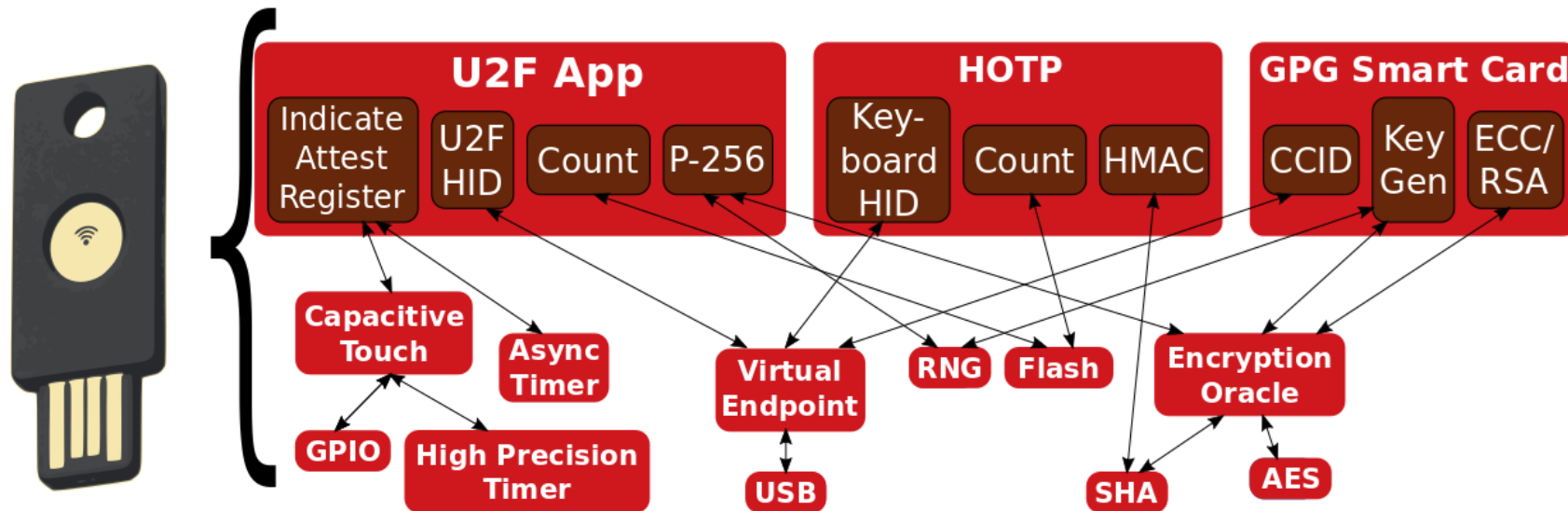
- Some subset of operating system principles, dictated by application
 - Hardware abstraction
 - Block storage versus file system?
 - But not architecture? [What's a runtime?]
 - Resource management
 - Process isolation
 - Job control

Example: USB Authentication Key

Multiple independent applications

No programmability in favor of security

Result: Handful of programmers control software stack

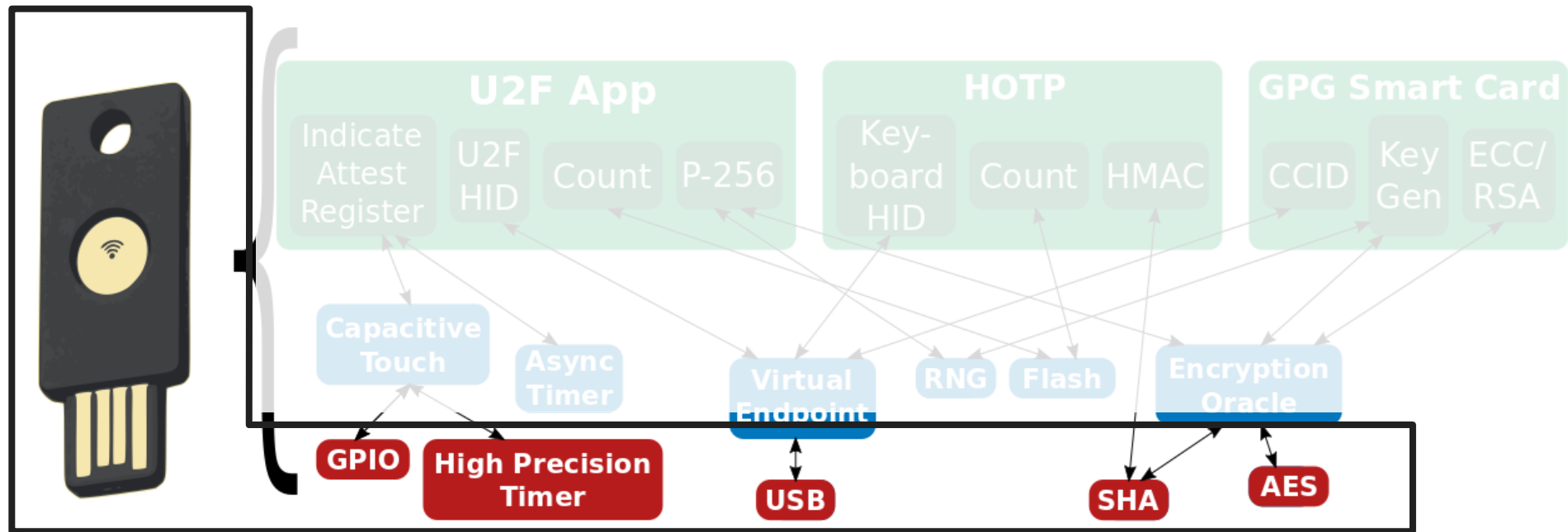


Platform

Build the hardware

Responsible for TCB: core kernel, MCU-specific code

Trusted: complete control over firmware & hardware



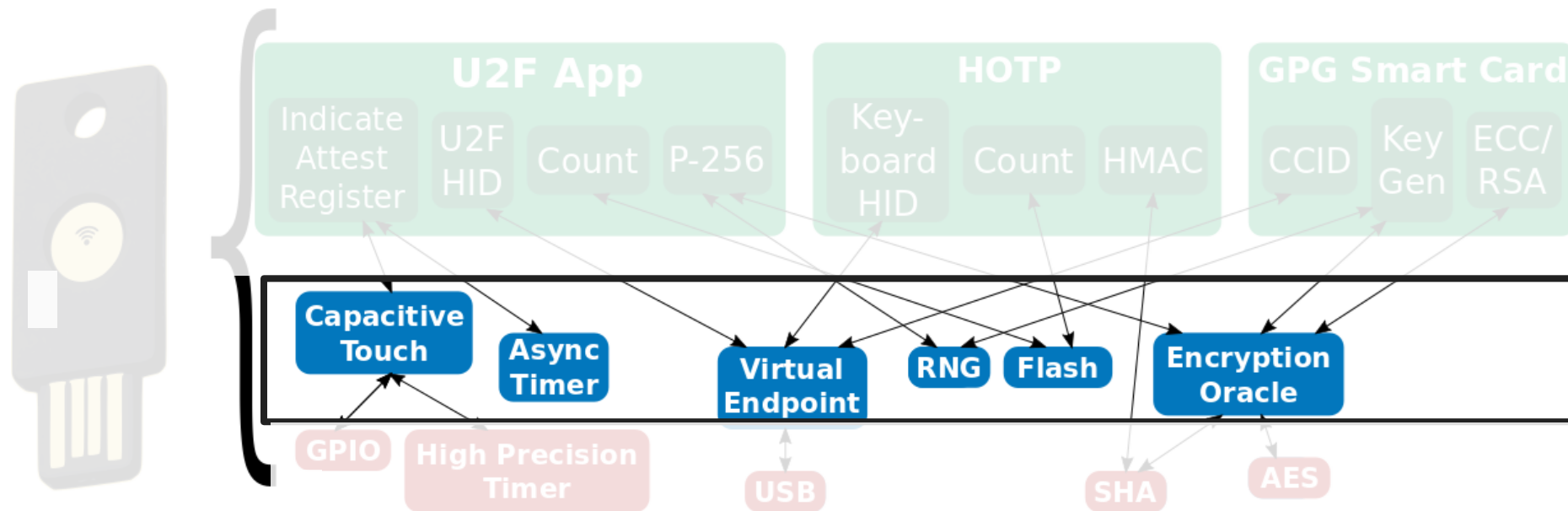
Goal: possible to correctly extend TCB

OS Services

Most OS services come from community

Device drivers, networking protocols, timers...

Platform provider can audit but won't catch all bugs



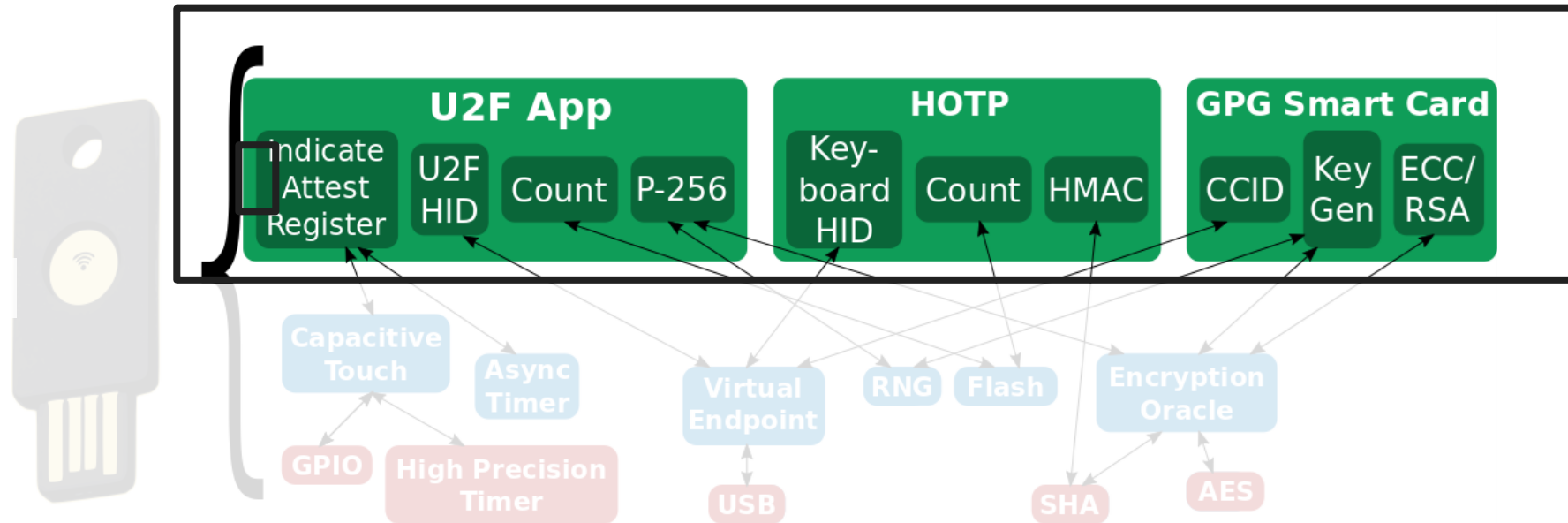
Goal: protect kernel from safety violations

Applications

Implement end-user functionality

“Third-party” developers: unknown to platform provider

Potentially *malicious*



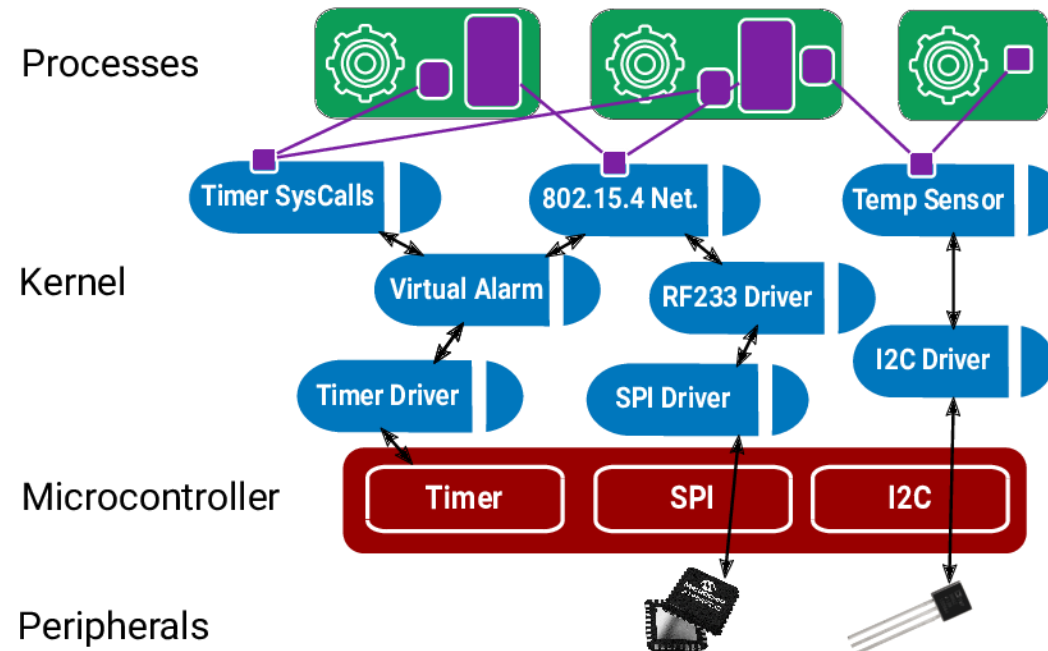
Goal: end-users can install 3rd-party apps

A Microcontroller OS

Processes: Use the Memory Protection Unit

Capsules: Type-safe Rust API for *safe* driver development

Grants: Bind dynamic kernel resources to process lifetime



What's Next?

- Lecture 16
 - September 15, Thursday, 12 pm – 1 pm