

Formal Proof of Type Preservation of the Dictionary Passing Transform for System F

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Type Classes in Haskell

Overloading Equality in Haskell

```
class Eq α where
  eq :: α → α → Bool

instance Eq Nat where
  eq x y = x ≐ y
instance Eq α ⇒ Eq [α] where
  eq [] [] = True
  eq (x : xs) (y : ys) = eq x y && eq xs ys

.. eq 42 0 .. eq [42, 0] [42, 0] ..
```

Desugaring Type Classes

Overloading Equality in System F_0

```
decl eq in
```

```
inst eq : Nat → Nat → Bool
```

```
  =  $\lambda x. \lambda y. \dots$  in
```

```
inst eq :  $\forall \alpha. [eq : \alpha \rightarrow \alpha \rightarrow \text{Bool}] \Rightarrow [\alpha] \rightarrow [\alpha] \rightarrow \text{Bool}$ 
```

```
  =  $\Lambda \alpha. \lambda (eq : \alpha \rightarrow \alpha \rightarrow \text{Bool}). \lambda xs. \lambda ys. \dots$  in
```

```
 $\dots$  eq 42 0  $\dots$  eq Nat [42, 0] [42, 0]  $\dots$ 
```

Dictionary Passing Transform

Overloading Equality in System F_0

```
decl eq in
inst eq : Nat → Nat → Bool
  =  $\lambda x. \lambda y. \dots$  in
inst eq :  $\forall \alpha. [eq : \alpha \rightarrow \alpha \rightarrow \text{Bool}] \Rightarrow [\alpha] \rightarrow [\alpha] \rightarrow \text{Bool}$ 
  =  $\Lambda \alpha. \lambda (eq : \alpha \rightarrow \alpha \rightarrow \text{Bool}). \lambda xs. \lambda ys. \dots$  in
.. eq 42 0 .. eq Nat [42, 0] [42, 0] ..
```

System F_0 Transformed to System F

```
let eq1 : Nat → Nat → Bool
  =  $\lambda x. \lambda y. \dots$  in
let eq2 :  $\forall \alpha. (\alpha \rightarrow \alpha \rightarrow \text{Bool}) \rightarrow [\alpha] \rightarrow [\alpha] \rightarrow \text{Bool}$ 
  =  $\Lambda \alpha. \lambda eq_1. \lambda xs. \lambda ys. \dots$  in

.. eq1 42 0 .. eq2 Nat eq1 [42, 0] [42, 0] ..
```

Agda Formalization of System F₀

Syntax Representation in Agda

```
data Term : Sorts → Sort r → Set where
  ' _          : Var S s → Term S s
  tt           : Term S es
  λ'x→ _       : Term (S ▷ es) es → Term S es
  Λ'α→ _       : Term (S ▷ τs) es → Term S es
  λ _ ⇒ _      : Term S cs → Term S es → Term S es
  _ · _        : Term S es → Term S es → Term S es
  _ • _        : Term S es → Term S τs → Term S es
  let 'x = _ in _ : Term S es → Term (S ▷ es) es → Term S es
  decl 'o in _   : Term (S ▷ os) es → Term S es
  inst ' _ ' = _ in _ : Term S os → Term S es → Term S es → Term S es
  _ : _         : Term S os → Term S τs → Term S cs
  'T _         : Term S τs
  _ ⇒ _        : Term S τs → Term S τs → Term S τs
  ∀ 'α _       : Term (S ▷ τs) τs → Term S τs
  [ _ ] ⇒ _    : Term S cs → Term S τs → Term S τs
```

Agda Formalization of System F_0

Context

```
data Ctx : Sorts → Set where
  ∅ : Ctx []
  _ ▶ _ : Ctx S → Stores S s → Ctx (S ▷ s)
  _ ▶ _ : Ctx S → Cstr S → Ctx S
```

Constraint Solving

```
data [_] ∈ _ : Cstr S → Ctx S → Set where
  here : [ (' o : τ) ] ∈ (Γ ▶ (' o : τ))
  under-bind : {ST : Stores S s'} →
    [ (' o : τ) ] ∈ Γ → [ (' there o : wk τ) ] ∈ (Γ ▶ ST)
  under-inst : [ c ] ∈ Γ → [ c ] ∈ (Γ ▶ c')
```

The Dictionary Passing Transform

Fun Lemmas on Our Way to Type Preservation

Type Preservation of the Dictionary Passing Transform