

# Formal Proof of Type Preservation of the Dictionary Passing Transform for System F

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# Type Classes in Haskell

## Overloading Equality in Haskell

```
class Eq α where
  eq :: α → α → Bool

instance Eq Nat where
  eq x y = x ≐ y
instance Eq α ⇒ Eq [α] where
  eq [] [] = True
  eq (x : xs) (y : ys) = eq x y && eq xs ys

.. eq 42 0 .. eq [42, 0] [42, 0] ..
```

# Desugaring Type Classes

## Overloading Equality in System $F_0$

```
decl eq in
```

```
inst eq : Nat → Nat → Bool
```

```
  = λx. λy. .. in
```

```
inst eq :  $\forall \alpha. [eq : \alpha \rightarrow \alpha \rightarrow \text{Bool}] \Rightarrow [\alpha] \rightarrow [\alpha] \rightarrow \text{Bool}$ 
```

```
  = λ $\alpha$ . λ(eq :  $\alpha \rightarrow \alpha \rightarrow \text{Bool}$ ). λxs. λys. .. in
```

```
.. eq 42 0 .. eq Nat [42, 0] [42, 0] ..
```

# Dictionary Passing Transform

## Overloading Equality in System $F_0$

```
decl eq in
inst eq : Nat → Nat → Bool
  = λx. λy. .. in
inst eq : ∀α. [eq : α → α → Bool] ⇒ [α] → [α] → Bool
  = Λα. λ(eq : α → α → Bool). λxs. λys. .. in
.. eq 42 0 .. eq Nat [42, 0] [42, 0] ..
```

## System $F_0$ Transformed to System $F$

```
let eq1 : Nat → Nat → Bool
  = λx. λy. .. in
let eq2 : ∀α. (α → α → Bool) → [α] → [α] → Bool
  = Λα. λeq1. λxs. λys. .. in

.. eq1 42 0 .. eq2 Nat eq1 [42, 0] [42, 0] ..
```

# Agda Formalization of System $F_0$

## Syntax Representation in Agda

```
data Term : Sorts → Sort r → Set where
  decl' o' in _      : Term (S ▷ o_s) e_s → Term S e_s
  inst' _ ' = _ ' in _ : Term S o_s → Term S e_s → Term S e_s → Term S e_s
  _ : _              : Term S o_s → Term S  $\tau_s$  → Term S  $c_s$ 
   $\lambda$  _  $\Rightarrow$  _      : Term S  $c_s$  → Term S e_s → Term S e_s
  [ _ ]  $\Rightarrow$  _         : Term S  $c_s$  → Term S  $\tau_s$  → Term S  $\tau_s$ 
  - ...
```

# Agda Formalization of System $F_0$

## Context

```
data Ctx : Sorts → Set where
  ∅ : Ctx []
  _ ► _ : Ctx S → Term S (item-of s) → Ctx (S ▷ s)
  _ ► _ : Ctx S → Cstr S → Ctx S
```

## Constraint Solving

```
data [_] ∈ _ : Cstr S → Ctx S → Set where
  here : [ (' o : τ) ] ∈ (Γ ► (' o : τ))
  under-bind : {l : Term S (item-of s')} →
    [ (' o : τ) ] ∈ Γ → [ (' there o : wk τ) ] ∈ (Γ ► l)
  under-inst : [ c ] ∈ Γ → [ c ] ∈ (Γ ► c')
```

# The Dictionary Passing Transform

# Fun Lemmas on Our Way to Type Preservation



# Type Preservation of the Dictionary Passing Transform