

GPU Acceleration of the Material Point Method

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GPGPU for performance enthusiasts

Why would('nt) you?

Drawbacks:

- ▶ Interactivity much easier on CPU, but slow PCI-Bus communication
- ▶ Code is mostly written against GPU architecture
- ▶ A lot of strain on the programmer

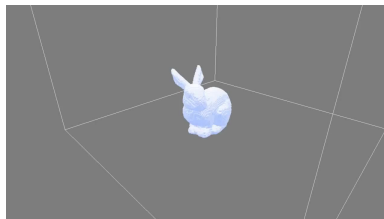
Benefits:

- ▶ Data is already on the GPU for rendering
- ▶ **Higher parallelization acceleration**

A Brief MPM Overview: Do You Want to Build a Snowman?

A short historical summary of MPM:

- ▶ Belongs to family of particle-in-cell(PIC) techniques [EHB57].
- ▶ Initial application to solids [SZS95] → MPM
- ▶ From research to production in *Disney's* animation film *Frozen* [Sto+13].
- ▶ Avalanche research [Gau+18]



Video result of my bachelor thesis on the simulation of snow [Mey15].

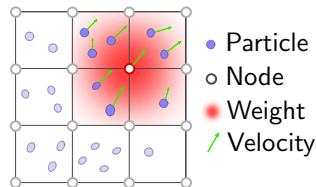
PIC ideas:

- ▶ Particles store all information
- ▶ Eulerian grid as a non-deformable scratchpad

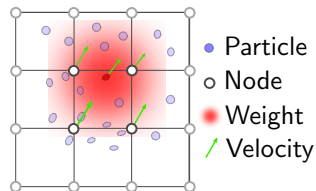
⇒ **meshfree**

Typical PIC/MPM roundtrip:

1. Transfer particle quantities to the grid (*P2G*)
2. Solve discretized governing equations on grid
3. Transfer back (*G2P*)



P2G-transfer \sum_p



G2P-transfer \sum_i

Governing Equations: Conservation of Mass & Momentum

Conservation of mass, continuum assumption holds.

Lagrangian (moving with a particle ${}_0\mathbf{x}$):

$${}_0^t J \rho({}_0\mathbf{x}, t) = \rho({}_0\mathbf{x}, 0). \quad (1)$$

Eulerian (outside observer ${}_t\mathbf{x}$):

$$\frac{\partial}{\partial t} \rho({}_t\mathbf{x}, t) = -\vec{\nabla} \cdot (\rho({}_t\mathbf{x}, t) \mathbf{v}({}_t\mathbf{x}, t)). \quad (2)$$

Lagrangian and Eulerian view measure differently but give same results. Equations are given in the strong form! [\[Jia+16\]](#)[\[Abe12\]](#)

Conservation of momentum:

Lagrangian (moving with a particle ${}_0\mathbf{x}$):

$$\rho({}_0\mathbf{x}, 0) \mathbf{a}({}_0\mathbf{x}, t) = \vec{\nabla} \cdot \mathbf{P}({}_0\mathbf{x}, t) + \mathbf{f}^{\text{body}}({}_0\mathbf{x}, t)_0 J. \quad (3)$$

Eulerian (outside observer ${}_t\mathbf{x}$):

$$\rho({}_t\mathbf{x}, t) \mathbf{a}({}_t\mathbf{x}, t) = \vec{\nabla} \cdot \boldsymbol{\sigma}({}_t\mathbf{x}, t) + \mathbf{f}^{\text{body}}({}_t\mathbf{x}, t) \quad (4)$$

Solving this equation will tell us how the velocity fields $\mathbf{v}({}_t\mathbf{x})$, $\mathbf{v}({}_0\mathbf{x})$ change on the whole domain due to acceleration \mathbf{a} . This is important to advect particles accounting for all forces.

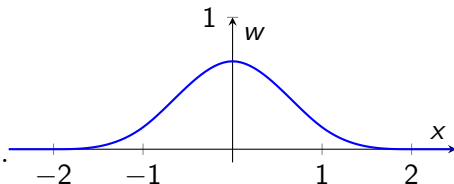
[Jia+16][Abe12]

Discretization of Space and Time

Space Discretization is done in a Galerkin/FEM fashion with grid based interpolants w_i . Here dyadic products of cubic b-splines suffice. w_i should satisfy at least [Gao+17]:

- ▶ Partition of unity: $\sum_i w(\mathbf{x} - \mathbf{x}_i^n) = 1$
- ▶ Identity relation: $\sum_i \mathbf{x}_i w(\mathbf{x} - \mathbf{x}_i^n) = \mathbf{x}$ ($\mathbf{x} = \mathbf{x}_p$)
- ▶ Non-negativity: $w \geq 0$.
- ▶ Limited support.
- ▶ C^1 -continuity.

Shortening $w_{ip}^n = w(\mathbf{x}_p^n - \mathbf{x}_i^n)$.



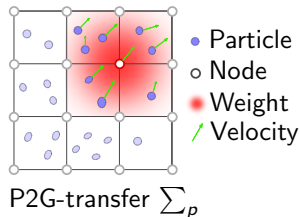
Transfer quantities from particles to grid. Numerical integration where the particles function as quadrature points [SKB08]:

$$m_i = \int_{\Omega} \rho(\mathbf{x}) w_i(\mathbf{x}) d\Omega \approx \sum_p \rho_p w_{ip} V_p \approx \sum_p m_p w_{ip}$$

$$A_i = \int_{\Omega} A(\mathbf{x}) w_i(\mathbf{x}) d\Omega \approx \sum_p A_p w_{ip} V_p.$$

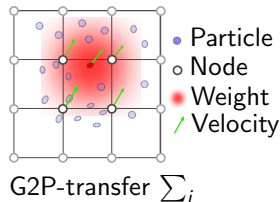
APIC-transfers add a local velocity field \mathbf{C}_p around \mathbf{v}_p :

$$(m\mathbf{v})_i = \sum_p w_{ip} m_p (\mathbf{v}_p + \mathbf{C}_p(\mathbf{x}_i - \mathbf{x}_p))$$



The MPM can be described in the weak formulation (FEM: virtual quantities). This allows us to 'shift' the derivative:

- ▶ $\mathbf{A}_p = \sum_i \mathbf{A}_i w_{ip}$
- ▶ $\nabla \mathbf{A}_p = \sum_i \mathbf{A}_i \nabla w_{ip}^T$
- ▶ $\nabla \cdot \mathbf{A}_p = \sum_i \mathbf{A}_i^T \nabla w_{ip}$



Time discretization with implicit midpoint scheme:

$$v_i^{n+1} - v_i^n = \frac{\Delta t}{m_i^n} f_i^{n+\frac{1}{2}} = \frac{\Delta t}{m_i^n} f \left(t^n + \frac{\Delta t}{2}, \frac{1}{2} y^n + \frac{1}{2} y^{n+1} \right)$$

- ▶ **implicit** requires linear system solve \Rightarrow more stable, larger time steps
- ▶ **midpoint** as it conserves angular momentum

Layout of the data: SoA vs. AoS

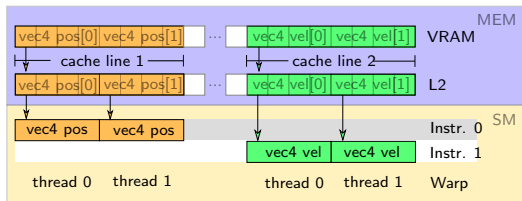
Coalescing:

SoA-Layout

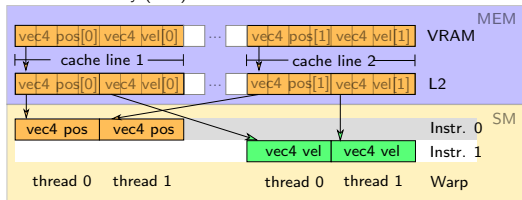
```
struct Container{
    vec4 positions[n];
    vec4 velocities[n];
} Particles;
```

AoS-Layout

```
struct Particle {
    vec4 position;
    vec4 velocity;
} Particles[n];
```



Structures of arrays(SoA)



Arrays of structures(AoS)

Nvidia Nsight[NVI] now offers metrics to identify bottlenecks:

Metric	Description
VRAM SOL%	memory throughput w.r.t. to hardware limit
SM SOL%	instruction throughput
L2 SOL%	L2-cache throughput
Tex SOL%	L1-cache throughput
SM Issue Util.%	amount of cycles an instr. was issued

A simple $\text{map}(y=\text{length}(x))$ shader on 1024×1024 Elements SoA vs. AoS differences:

Layout	$\Delta t_c (\mu s)$	Speedup	VRAM	SM	L2	SM Issue Util.
AoS(1 instr.)	243	-	77.7%	7.3%	30.3%	6.8%
SoA(1 instr.)	120	2.26x	75.4%	14.3%	29.4%	14.0%
AoS(2 instr.)	275	-	61.3%	41.8%	53.8%	48.9%
SoA(2 instr.)	240	1.16x	75.4%	29.4%	20.0%	62.3%

⇒ SoA increases coalescing for non-random access.

Parallel Reduction & Scan

Assuming an associative $\text{binary_op}(x,y) := x \circ y$, a neutral element e of the binary_op , and an array of values $[a_0, a_1, \dots, a_n]$.

- **Parallel reduction** computes the value:

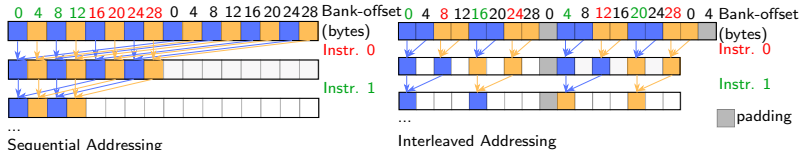
$$r = a_0 \circ a_1 \circ \dots \circ a_n. \quad (5)$$

- (Exclusive) **scan** computes the array:

$$[e, a_0, (a_0 \circ a_1), (a_0 \circ a_1 \circ a_2), \dots, (a_0 \circ a_1 \circ a_2 \circ \dots \circ a_{n-1})]. \quad (6)$$

Here, only shared memory approaches without warp shuffle operations.

Shared Memory Bank Conflicts:



Interleaved addressing causes bank conflicts (Short Scoreboard activity) \Rightarrow padding needed.

Method	Δt_c	Speedup	VRAM	SM	Sel. Warp-Stall Reas.
Interl. no padd.	305	-	23.0%	60.9%	S. Scoreb.(17.2%)
Sequential	141	2.16x	49.8%	37.1%	S. Scoreb.(2.0%)

Table: Parallel reduction on 1024×1024 vectors with $y=\text{length}(x)$ as input.

More elements than thread group size require pyramid schemes.

Sequential work: multiple elements per thread.

- ▶ Memory latency hiding (Long Scoreboard up)
- ▶ Higher reduction factor each dispatch \Rightarrow Less global memory indirections
- ▶ Unrolling loops can help but adds register pressure.

Method	Δt_c	Speedup	VRAM	SM	Sel. Warp-Stall Reas.
Sequential	141	2.16x	49.8%	37.1%	S. Scoreb.(2.0%)
Seq. (2x)	100	3.05x	69.5%	26.2%	L. Scoreb.(80.1%)
Seq. (128x)	98	3.1x	72.9%	16.9%	L. Scoreb.(84.4%)
Seq. (256x)	101	3.0x	66.4%	14.6%	L. Scoreb.(76.9%)

Table: Parallel reduction on 1024×1024 vectors with $y=\text{length}(x)$ as input. Methods have 504, 8, 4 thread groups, respectively. A GTX970 has 13 SMs.

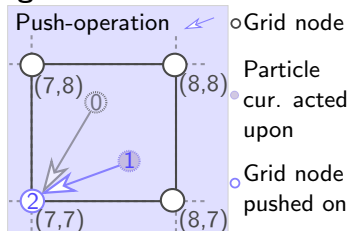
Scan is similar but cannot profit as much from sequential work having to keep multiple elements in register memory.

Binning & Counting Sort: Where Are You?

Grid node does not know its neighboring particles \Rightarrow Binning.

Binning combines nicely with **Counting Sort**:

1. **Binning**: Per node counting.
2. **Scan**: Computes new memory offset for particles.
3. **Reordering**: Give back indexing list or do deep copy.



Sorting can dramatically increase workload performance of subsequent steps for neighboring queries:

1. Deep sorted accesses are now **coalesced**.
2. **Data reuse** due to L2-Cache and/or shared memory.

Double buffer particles to use last sorted state as input for new sorting to profit from item 1 and 2!

Ordering	$\Delta t_c (\mu s)$	Speedup	VRAM	SM	L2	L2-Hit
Random	1,516	-	25.0%	3.4%	9.1%	10.8%
Deep sorted	218	6.95x	75.3%	24.4%	35.0%	37.8%

Table: Order dependency of binning of 1024×1024 randomly positioned particles in a $128 \times 128 \times 128$ grid.

The MPM Specific Transfers

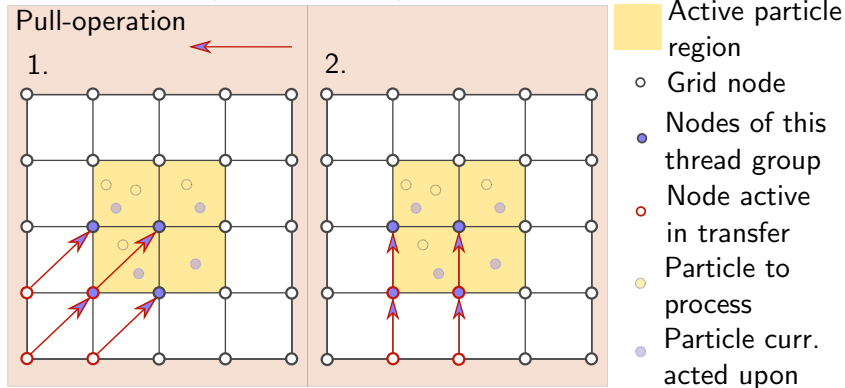
All MPM operations belong to one of those parallelization schemes:

- ▶ 1 thread : 1 particle: $\square_p = \square_p \circ \square_p \circ \dots \circ \square_p$.
- ▶ 1 thread : 1 node: $\square_i = \square_i \circ \square_i \circ \dots \circ \square_i$.
- ▶ G2P-transfer: $\square_p = \sum_i \square_i \circ \square_{ip}$.
- ▶ P2G-transfer: $\square_i = \sum_p \square_p \circ \square_{ip}$.

MPM-Transfers are executed **multiple times per physical frame** with varying numbers of variables and mathematical operations.

⇒ Preprocessing steps only need to be done **once per physical frame**. Sorting already introduced as one of these.

G2P-Transfer: $\square_p = \sum_i \square_i \circ \square_{ip}$.



Pull-Operations read!

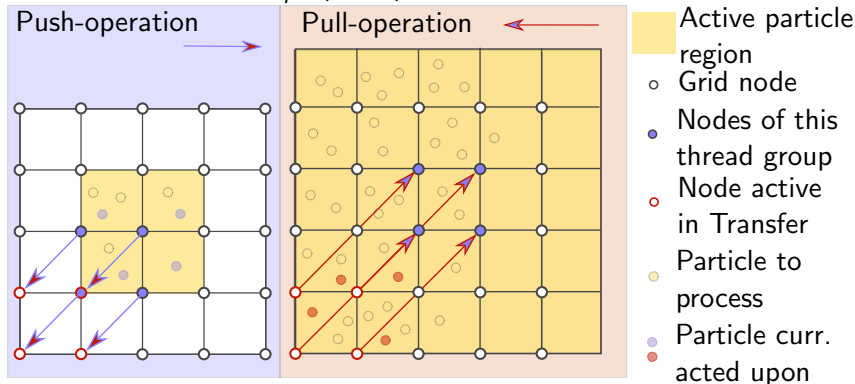
Similar to filter or stencil operations on the GPU:

- ▶ 1 node : 1 thread, split grid into blocks corr. to thread groups.
- ▶ Interpolation function however dependent on particle position.
- ▶ Needs to be rerun for every particle in the cell.

Typical setup of transfers:

1. Initialize shared memory (pull: with nodes from global memory).
2. Perform transfers.
3. Write back to global memory (push: global atomics since writes on halo).

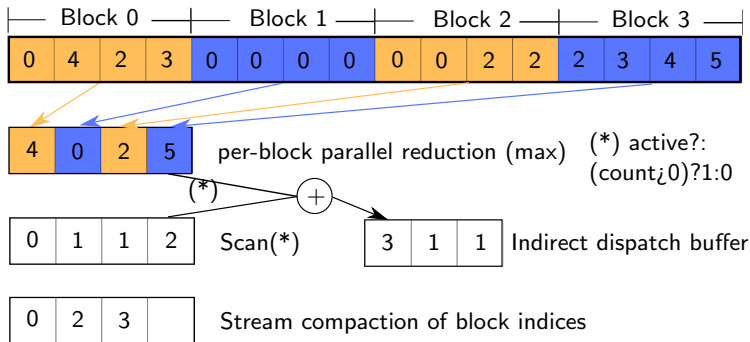
P2G-transfer: $\square_i = \sum_p \square_p \circ \square_{ip}$.



Pull-Operations read, Push-operations write!
 \Rightarrow Race conditions.

Typically simulation domain(grid) much bigger than simulation model. Filtering inactive blocks as a preprocess improves performance.

- ▶ Block is active if any cell counter is active.
- ▶ Cell counter is active if it has at least one particle.



- ▶ Block and halo are always target of shared memory operations
⇒ P2G-pull low occupancy.
- ▶ **Batching** multiple particles can increase performance due to hiding synchronization, unroll!
- ▶ Transfers respect shared memory bank conflicts fully.
- ▶ **Warp divergence** for varying cell counts.

Method	$\Delta t_c (\mu s)$	Speedup	VRAM	L2	SM
global	44,442	-	4.6 %	34.4%	7.7%
global sorted	20,484	2.21x	7.0 %	44.0%	16.1%
P2G-sync	2,595	17.47x	5.9%	7.6%	67.0%

P2G-transfers of one uniformly million particles with 4 particles per cell with random velocities between $v_x, v_y, v_z \in [-1.0; 1.0]$ in a $128 \times 128 \times 128$ grid. Block size is (8,4,4).

A comparison to [Gao+18]

Simultaneously being worked on. Largely same decision making:

	Me	[Gao+18]
Sort	Count/Histogram for each var.	Count/Histogram sel. variables
Filtering domain	Filter-op.	Sparse Grid structure
Transfers	Shared mem. only	Warp-shuffle op.

Warp shuffle allows for fast parallel segmented reduction of cells of varying counts.

⇒ Solves warp divergence and shared memory issues mostly, thread groups now correspond to particles. Faster for varying counts.

Thank you for your attention!

- ▶ **Bachelor thesis:** <https://github.com/MeyerFabian/snow>
- ▶ **Master thesis:** <https://github.com/MeyerFabian/msc>
- ▶ **Presentation:** <https://github.com/MeyerFabian/msc/pres>
- ▶ **Code** (right now): <https://github.com/mpm-msc/snow>

Questions?



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The Pretty Strong but Mathematically Weak Formulation

Weak Formulation (or Principle of Virtual Work):

Dot product equations with arbitrarily 'test functions' \mathbf{q} and apply divergence theorem:

$$\int_{\Omega^0} {}_0\mathbf{q} \cdot \left[({}_0\rho_0)({}_0\mathbf{a}) - {}_0\mathbf{f}^{\text{body}t} {}_0J \right] d_0\mathbf{x} = \int_{\partial\Omega^{t^n}} {}_t\mathbf{q} \cdot \boldsymbol{\sigma} d_t\mathbf{A} - \int_{\Omega^{t^n}} \nabla_t \mathbf{q} : \boldsymbol{\sigma} d_t\mathbf{x}. \quad (7)$$

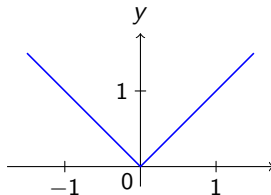
A strong solution is also a solution to the weak formulation. Leave out body forces (like gravity) and boundary condition (e.g. collisions) for now:

$$\int_{\Omega^0} {}_0\mathbf{q} \cdot ({}_0\rho_0)({}_0\mathbf{a}) d_0\mathbf{x} = \int_{\Omega^{t^n}} \nabla_t \mathbf{q} : \boldsymbol{\sigma} d_t\mathbf{x}. \quad (8)$$

Weak Derivative:

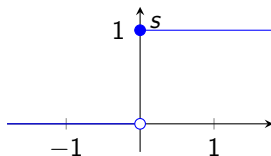
$y = |t|$ has weak derivative:

$$v = \begin{cases} -1, & \text{if } t < 0 \\ c, & \text{if } t = 0 \\ 1, & \text{if } t > 0 \end{cases}$$



Heaviside step function has no weak derivative:

$$s = \begin{cases} 0, & \text{if } t < 0 \\ 1, & \text{if } t \geq 0 \end{cases}$$



Allows for point loads, material discontinuities and more. [\[Bat06\]](#)

Space Discretization is done in a Galerkin/FEM fashion with grid based interpolants w_i with limited support. Here dyadic products

$$w_i(\mathbf{x}) = w(\mathbf{x} - \mathbf{x}_i) = w\left(\frac{1}{h}(\mathbf{x} - \mathbf{x}_i)\right) = w\left(\frac{1}{h}(x - x_i)\right)w\left(\frac{1}{h}(y - y_i)\right)w\left(\frac{1}{h}(z - z_i)\right) \quad (9)$$

of cubic b-splines suffice:

$$w(x) = \begin{cases} \frac{1}{2}|x|^3 - |x|^2 + \frac{2}{3} & 0 \leq |x| < 1 \\ \frac{1}{6}(2 - |x|)^3 & 1 \leq |x| < 2 \\ 0 & 2 \leq |x| \end{cases} \quad (10)$$

