



M2M Security

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Outline

- Challenges in M2M Security
- oneM2M Security Solutions
- Access-rights Management
- Appendix. TLS/DTLS





CHALLENGES IN M2M SECURITY

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Uniqueness of M2M Security (1)

- Security is an important function in any communication systems.
- This is especially critical in M2M systems because unlike in cellular systems, the application, service, network and device distribution are not tightly coupled and managed by a single entity such as a network operator.
- For example, in the case of smart meters, the entities involved may include
 - The utility company (M2M application provider)
 - The M2M service provider (MVNO providing M2M service layer)
 - The cellular network provider (providing 3G/4G networks)
 - The meter manufacturer (providing smart meter) and
 - The end user
- Furthermore, limited business relationships exist among these multiple players and thus no mutual trust should be assumed among these players.



Uniqueness of M2M Security (2)

- The economics of M2M do not allow a security provisioning process similar to that of cellular handsets. Security methods developed for M2M should be able to deal with an explosion in the number of devices by manufacturers.
 - A simpler automated process with the capability of efficiently dealing with a large number of devices is desirable.
- Unlike cellular handsets, M2M devices are often unattended and are subject to a higher risk of vandalism and misuse.
 - Security solutions should be designed to protect against any inappropriate use of the M2M devices.

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Examples of Undetected M2M Security Attacks

- Theft of Removable UICC (Universal Integrated Circuit Card) from an M2M device
 - The network cannot tell that the UICC has been installed in a non-M2M device.
- Hijacking through Passive Medium Overhearing with the absence of UICC
 - By overhearing the UATI (Unicast Access Terminal Identifier) and the IP addresses that are assigned to the M2M devices by the network operators, the attacker can generate packet headers that will mislead the network to believe that the packets are from a legitimate client.





Trust Relationships in the M2M

- **Ecosystem**The ecosystem potentially involves M2M application providers, M2M service providers, cellular network providers, device manufacturers and end users.
- Application providers may not have a relationship with device manufacturers.
- However, M2M application providers always maintain a business relationship with end users and M2M service providers; M2M service providers always maintain relationship with network providers.
- It is likely that the device may be used by an end user who has no ownership of the device.
- Examples: (1) smart meters offered by a utility company and (2) OBU (On Board Unit) offered by an M2M application provider for fleet management.





M2M Security Threats

- Customer/M2M Device User
 - Devices access by unauthorized entities
 - Leaking of stored data by eavesdropping at some point of the network
 - End user identity revealed
- M2M Network Provider
 - Threat primarily from inexpensive devices and ease of accessibility of devices
- M2M Service Provider
 - Service availability
 - Data integrity
- **Application Provider**
 - Man-in-the middle attacks to the data transferred from devices to back-end servers.
 - Non authenticated devices.



M2M Security Requirements (1)

- Identification -the process of checking if the identity provided for authentication is valid
- Mutual Authentication the process of validating if the identity supplied in the identification step is associated with a trustworthy credential.
- Authorization responsible for authorizing services and data access to authenticated entities.
- Confidentiality keep secrecy of the data transmitted.

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M2M Security Requirements (2)

- Integrity Guarantee no modification of the original data
- Exclusivity no use of M2M communication modules for non M2M applications.
- Anonymity Keep the sender's identify unrevealed.



Bootstrapping Requirements

- Complex ecosystem: assume no trust relationship between application providers and device manufacturers.
- Scalability: Bootstrapping should be as automated as possible.
- Network authentication first: The M2M service layer is an overlay network over the existing network.
- Flexibility in operator selection: No leakage of key information across operator boundaries.

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Available Security Solutions for Bootstrapping

- Public Key (Asymmetric Key) Solutions
- Smart Card-based Solutions
- Pre-Provisioned Symmetric Key Solutions
- Identity-Based Encryption Solutions





Authenticated key

Symmetric Key vs. Asymmetric Key

- Authenticated key agreements are cryptographic protocols where two or more participants authenticate each other and agree on key material used for securing future communication.
 - Symmetric key protocol
 - require an out-of-band security mechanism to bootstrap a secret key.
 - Asymmetric public-key protocol
 - require certificates and a large-scale Public Key Infrastructure (PKI). ← more flexible

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Public Key Solutions

- Approach: Provisioning a private-public key pair along with a certificate of public key during manufacturing.
- Public Key provided by a CA (Certificate Authority) to encrypt data. Everyone usually knows public key.
- Private Key a secret key known only to a key owner.

Purpose of Key	Key Used	Whose Key Used	
Encrypt data for a recipient	Public key	Receiver	
Sign data	Private key	Sender	
Decrypt data received	Private key	Receiver	
Verify a signature	Public key	Sender	





Public Key Solutions (Cont.)

- A public key certificate allows the exchange of information securely over the Internet using the public key infrastructure (PKI).
- A **certificate** was issued by an official, trusted agency called Certificate Authority (CA) so that a recipient can verify that the certificate is real.
- To verify a certificate is genuine and valid, it is digitally signed by a root certificate belonging to the CA.
- Issues:
 - Need a large-scale PKI to handle billions of devices.
 - Must trust the CA (Trusted Third Party TTP).
- PKI Standards: X.509







Public Key Solutions (Cont.)

Public Key Infrastructure (PKI)

CA Issue Check

Source: http://ei4africa.eu/2013/04/24/ei4africa-boosts-the-deployment-of-public-key-infrastructures-in-africa/

Public key

certificate

User's

private key





Public Key Solutions (Cont.)

- Why not ideal for M2M devices:
 - Public key methods are not supported by existing cellular networks. Existing networks use symmetric key protocols.
 - Cost of PKI is too high one certificate per device means billions of certificates.
 - Safe storage of certificates and private keys is an issue.







Smart Card-based Solutions

- Approach: Devices are equipped with smart cards to provide the necessary key material for use by the M2M service layer.
- Smart cards contain a pre-provisioned secret key that is hard-coded by the smart-card manufacturer for network operators.
- As a result, M2M application and service providers need to trust network operators.
- Issues:
 - Not scalable in cellular M2M deployments that may be comprised of billions of devices.
 - Subscription portability enabled by smart cards has little and no value in the M2M context but introduces the possibility of malicious use of these smart cards in other devices.
 - When switch network operators, all smart cards for an M2M application provider need to be replaced.



Pre-Provisioned Symmetric Keys Solutions

 Approach: Both the M2M device and the service provider are pre-provisioned with the same key for symmetric key operations during device registration.

Issues:

- The M2M service provider and the device owner normally are not involved in the device deployment and distribution.
- Manufacturers have to install the key into the device in the factory and thus they need to provide this information to the service provider.
- Such pre-provisioned keys cannot be used as permanent keys for establishing secure service-layer sessions. However, they can be used for initial mutual authentication between devices and service providers and for use in a secure procedure to generate a new shared permanent key that cannot be derived by untrusted parties such as manufacturers.

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Identity-Based Encryption Solutions

- Recently, Identity-Based Encryption (IBE)
 protocols have been proposed as a viable
 alternative to public-key methods by replacing
 the PKI with a Private-Key Generator (PKG).
- Approach: The public key of a user-device is a mathematical function of the identity that is associated with this key. Thus there is no need to bind an identity of an entity with a public key through the use of **certificates**. The PKG is used to generate the paired private key of the identity-based public key.



Identity-Based Encryption Solutions (Cont.)

- IBE is based on the concept of elliptic curve cryptography (ECC) and can be used for establishing a shared key between two entities.
- Advantages:
 - keys are always valid (no key revocation.)
 - Public keys are derived from identifiers, thus the need for a public key infrastructure is eliminated.
 - Additional information may be encoded into the identifier such as an expiration date for a message.

Drawbacks:

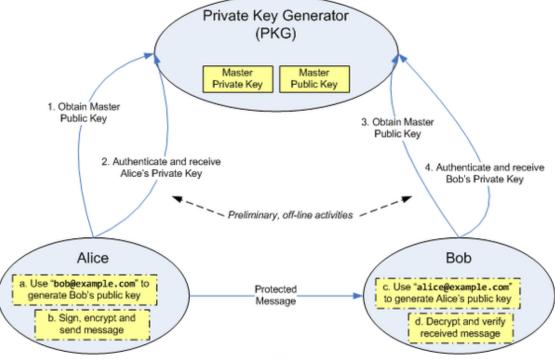
- Key escrow problem: All private keys needed to decrypt encrypted data are held in escrow by PKG.
- Lack of non-repudiation: Because PKG generates private keys for users, it may decrypt and/or sign any message without authorization.
- A secure channel between a user and the PKG is required for transmitting the private key.

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Identity-Based Encryption Solutions (Cont.)



Source: http://en.wikipedia.org/wiki/ID-based_encryption



ONEM2M SECURITY SOLUTIONS

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M2M Service Security Management

- oneM2M Security Architecture
- High Level Flow To Establish M2M Connection
 - M2M Service Layer Security Provisioning
 - M2M Security Association Establishment
- Detailed Flow To Establish M2M Connection
 - M2M Service Layer Security Provisioning
 - M2M Security Association Establishment





oneM2M Security Architecture (1)

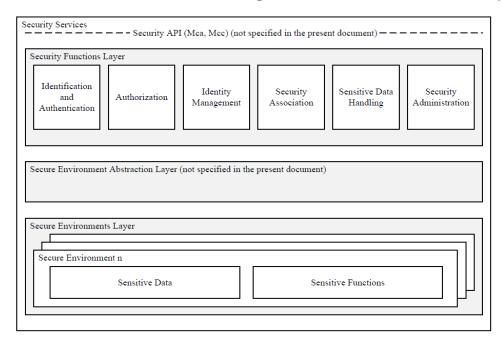


Figure 5.1-1: High level overview of the Security architecture

Source: oneM2M TS-0003-Security-Solutions

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oneM2M Security Architecture (2)

The architecture consists of following layers:

Security Functions layer:

- This layer contains a set of security functions that are exposed at reference point Mca and Mcc.
 - Identification
 - Authentication
 - Authorization
 - Security Association
 - Sensitive Data Handling and
 - Security Administration.

Source: oneM2M TS-0003-Security-Solutions





oneM2M Security Architecture (3)

Security Environment Abstraction Layer:

- Includes various security capabilities such as
 - key derivation,
 - data encryption/decryption,
 - o signature generation/verification,
 - security credential read/write from/to the Secure Environments
- The security functions in the Security Functions Layer invoke these functions in order to do the operations related to the Secure Environments.
- In addition, this layer also provides physical access to the Secure Environments. Implementation of this is out of scope of the present document. This layer is not specified in the initial release but is expected to be considered in future releases.

Source: oneM2M TS-0003-Security-Solutions

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oneM2M Security Architecture (4)

Secure Environment (SE) layer:

- This layer contains one or multiple secure environments that provide various security services related to sensitive data storage and sensitive function execution.
- The sensitive data includes SE capability, security keys, local credentials, security policies, identity information, subscription information, and so on.
- The sensitive functions include data encryption, data decryption, and so on.
- Implementation of secure environments is out of scope of the present document.

Source: oneM2M TS-0003-Security-Solutions





Security Function Layer (1)

This layer contains a set of security functions that are exposed at Reference Points Mca and Mcc.

1. Identification

- Check if the identity provided for authentication is valid.
- Once passing the identification, the identified identity will be supplied to authentication process.

2. Authentication

 Validate if the identity supplied in the identification step is associated with a trustworthy credential.

3. Authorization

- Authorize services and data access to authenticated entities according to provisioned Access Control Policies (ACPs) and assigned roles.
- The Authorization function may need to evaluate multiple ACPs in an authorization process.

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Security Function Layer (2)

4. Security Association

 Establish the secure association between two M2M nodes to provide confidentiality and integrity.

5. Sensitive Data Handling

- Protect the sensitive data like local credentials for storage and manipulation.
- Perform security algorithms in cryptographically separated secure environments.

6. Security Administration

- Provide functions to manage the Security functions, resources and attributes.
- Manage sensitive data with their associated identifiers and subscriptions on behalf of other entities.





High Level Flow To Establish M2M Connection

- First, establish connectivity in the underlying Network Services Layer
- Next, Enrollment Phase: Service Layer Security Provisioning
- 3. Next, Operational Phase: Security Association Establishment
- 4. **TLS or DTLS session** established which protects messages being exchanged between adjacent AE/CSE, i.e. hop-by-hop.
- Finally, AEs that need to preserve the privacy of their information exchange from untrusted intermediate nodes may be provisioned to support a direct **security association** between them.

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Objectives in Each Stage

- RSPF (Enrollment Key, Master Credentials)
- SAEF (Security Association)
- Between Applications (Security Association)





Enrollment Phase: Service Layer Security provisioning

- May be performed by
 - a pre-provisioning that can be integrated in the manufacturing or product deployment phase, or
 - by means of a security bootstrapping procedure (i.e. Remote Security Provisioning Framework RSPF) that takes place before the equipment starts actual operation.
- Requires selection of the stakeholder that will provide services through the equipment, especially the M2M Service Provider

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Remote Security Provisioning Frameworks (RSPF)

- Remote Security Provisioning procedures rely on an M2M Enrollment Function (MEF) which can be external to the M2M Service Provider to establish appropriate credentials.
- Provide post-provisioning of the essential information (normally to produce an Enrollment Key, then generate Master Credentials from Enrollment Key) to establish a security association between a Field Domain entity and a chosen M2M Service Provider.





Operational Phase: Security Association Establishment

- M2M services are offered by CSEs to AEs and/or other CSEs.
- To be able to use M2M services offered by one CSE, the AEs and/or CSEs need to be mutually identified and authenticated with that CSE.
- This mutual authentication provides protection from unauthorized access and Denial of Service attacks. It also enables encryption and integrity protection for the exchange of messages across a single Mca, Mcc or Mcc' reference point.
- In addition, communicating AEs that require similar protection for their own information exchanges can be provisioned to apply the same security method to their communications.

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Security Association Establishment Framework (SAEF)

- On the Mca and Mcc reference points, security association establishment between a field domain CSE and an infrastructure CSE is mandatory.
- On the Mcc' reference point, security
 association establishment between IN-CSE and
 IN-CSE is mandatory.
- On the Mca reference point, security association establishment between AEs and the CSE in the field domain is strongly recommended.





Roles of GBA BSF and MEF (1)

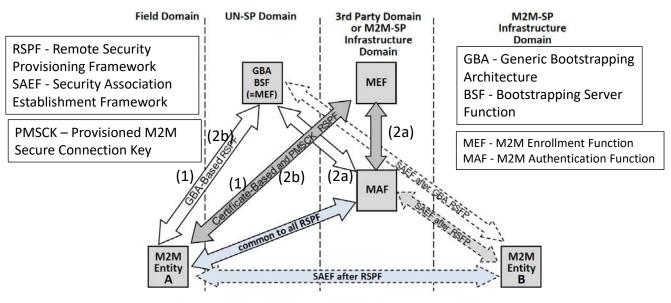


Figure 6.1.2.1-1: Entities involved in Remote Security Provisioning

Source: oneM2M TS-0003-Security-Solutions

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Roles of GBA BSF and MEF (2)

- (1) Three alternatives between M2M Entity A and MEF for service bootstrapping
 - Pre-Provisioned Symmetric Key RSPF (not shown in the diagram)
 - Certificate-Based RSPF
 - GBA-based RSPF
- (2a) & (2b) Derived key by RSPF is distributed to MAF and M2M Entity A for MAF SAEF





SAEF and the Role of MAF (1)

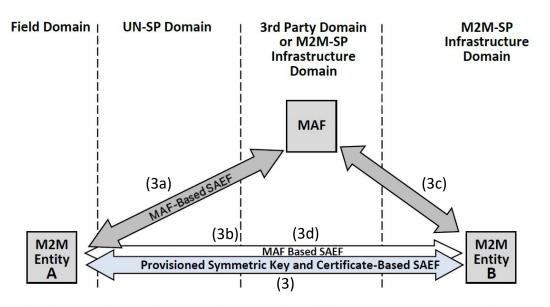


Figure 6.1.2.2.1-1: Entities involved in Security Association Establishment Source: oneM2M TS-0003-Security-Solutions

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SAEF and the Role of MAF (2)

- 3) M2M Entity A establishes security association with M2M Entity B based on
 - Provisioned Symmetric Key SAEF (no MAF involved)
 - Certificate-Based SAEF (no MAF involved)
 - M2M Authentication Function (MAF) SAEF

For MAF SAEF,

- (3a) Entity A handshakes with MAF
- (3b) Entity A sends the derived Kc and KcID to Entity B
- (3c) Entity B handshakes with MAF
- (3d) Entity A and Entity B handshake to set up security association





Security Mechanism defined by oneM2M

	Remote Security Provisioning Framework	Security Association Establishment Framework
1	Pre-provisioned Symmetric Key RSPF	Provisioned Symmetric Key SAEF
Χ	Pre-provisioned Symmetric Key RSPF	Certificate-based SAEF
2	Pre-provisioned Symmetric Key RSPF	MAF-Based SAEF
3	Certificate-based RSPF	Provisioned Symmetric Key SAEF
Χ	Certificate-based RSPF	Certificate-based SAEF
4	Certificate-based RSPF	MAF-based SAEF
5	GBA-based RSPF	Provisioned Symmetric Key SAEF
Χ	GBA-based RSPF	Certificate-based SAEF
6	GBA-based RSPF	MAF-based SAEF
7 8 9		Provisioned Symmetric Key SAEF MAF-based SAEF Certificate-based SAEF

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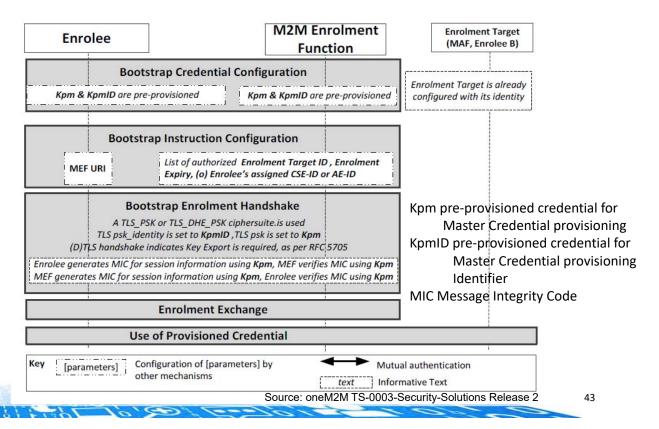




List of Keys

- Kc M2M Secure Connection Key
- KcID M2M Secure Connection Key Identifier
- Ke Enrolment Key
- KeID Enrolment Key Identifier
- Ker Enrolment Re-Authentication Key
- Km Master Credential
- KmID Master Credential Identifier
- Kpm pre-provisioned credential for Master Credential provisioning
- KpmID pre-provisioned credential for Master Credential provisioning Identifier
- Kpsa provisioned credential for M2M Security Association Establishment
- KpsaID provisioned credential for M2M Security Association Establishment Identifier
- Ks temporary Key material referred to in GBA
- Ks..NAF Abbreviation of Ks_(int/ext)_NAF
- Ks_(ext/int)_NAF Derived key in GBA_ME or Derived key in GBA_U which remains on UICC
- Ks_ext_NAF Derived key in GBA_U sent to the ME
- Ks_int_NAF Derived key in GBA_U which remains on UICC
- Ks_NAF Derived key in the ME

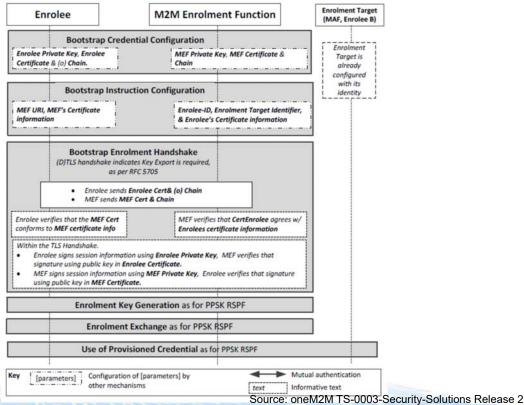
RSP based on Pre-Provisioned Symmetric Key





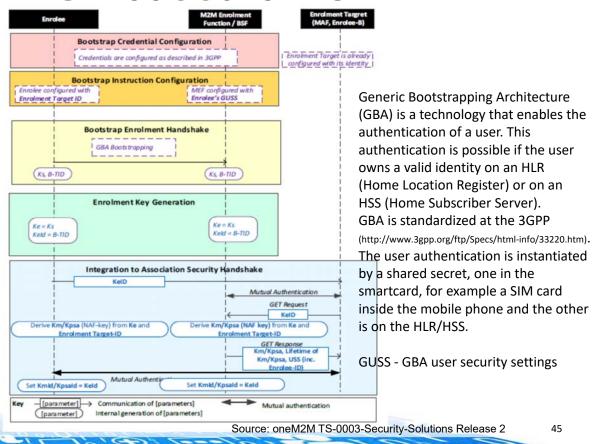
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RSP based on Certificate





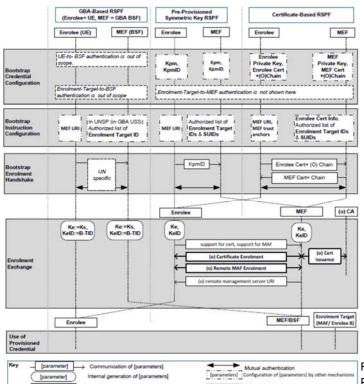






RSP Methods





General Description of RSP

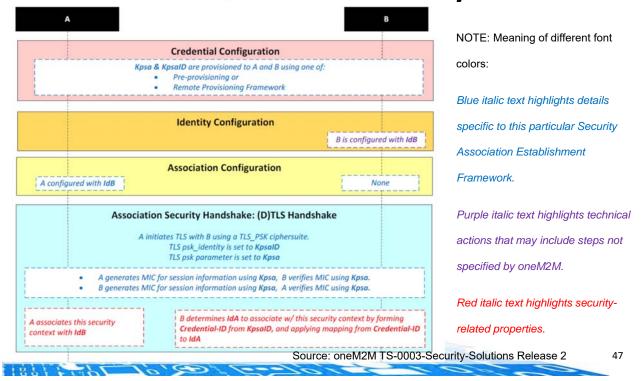
- **Bootstrap Credential Configuration**
- **Bootstrap Instruction Configuration**
- **Bootstrap Enrolment Handshake**
- **Enrolment Key Generation**
- **Enrolment Phase**

Figure 8.3.1.2-1: Overview of the Remote Security Provisioning Frameworks supported by oneM2M





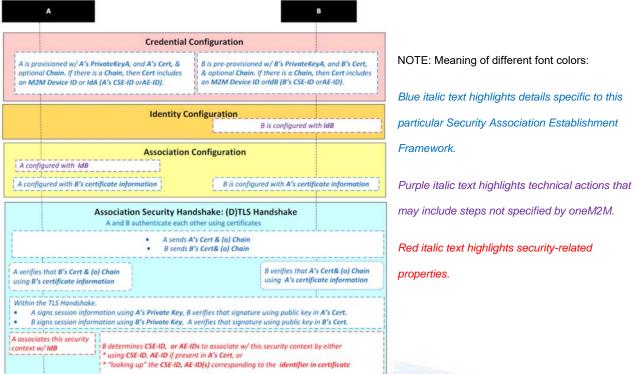
SAE Based on Provisioned Symmetric Key







SAE Based on Certificate

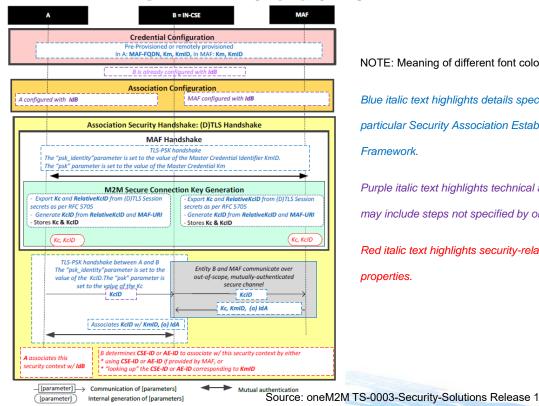


Source: oneM2M TS-0003-Security-Solutions Release 2





SAE Based on MAF



NOTE: Meaning of different font colors:

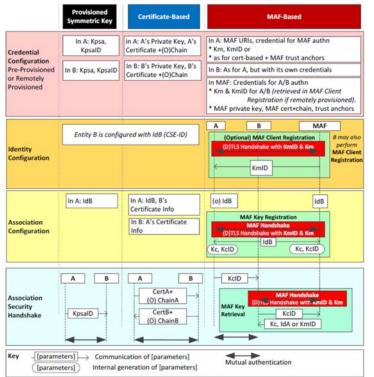
Blue italic text highlights details specific to this particular Security Association Establishment Framework.

Purple italic text highlights technical actions that may include steps not specified by oneM2M.

Red italic text highlights security-related properties.



SAE Methods



General description of SAEF

- **Credential Configuration**
- **Identity Configuration**
- **Association Configuration**
- Association Security Handshake

Figure 8.2.1-1: Overview of the Security Association Establishment Frameworks supported by oneM2M

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ACCESS-RIGHTS MANAGEMENT

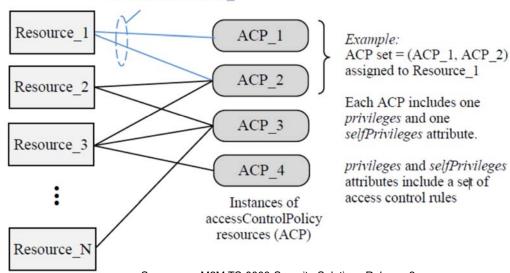
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Relation between Resource Instances and Access Control Policies

List of IDs in accessControlPolicyIDs attribute of Resource 1



Source: oneM2M TS-0003-Security-Solutions Release 2





selfPrivileges vs. Privileges

- Access requests to ACP's itself are evaluated against the selfPrivileges attribute of that ACP.
- Access requests to instances of all other resource types, are evaluated against the privileges attributes of the ACP set associated with the targeted resource.







Parameters indicated in the request message

Parameter	Description	Mandatory/O ptional	Usage in access control mechanism
to	URI of target resource	M	Selection of accessControlPolicy associated with the target resource
fr	Identifier representing the originator of the request	M (see Note)	Evaluated against accessControlOriginators in privileges and selfPrivileges attributes
role	Role of the originator	0	Evaluated against accessControlOriginators in privileges and selfPrivileges attributes
ор	Requested operation	M	Evaluated against accessControlOperations in privileges and selfPrivileges attributes
fc	filterUsage condition tag in Filter criteria	0	Differentiation between Retrieve and Discovery operations
Tokens	ESData-protected Tokens	0	Contains authorization information (e.g. Role-IDs) to be used in the decision for the request
Token IDs	tokenIDs or Local-Token-ID	0	Identifies Tokens containing authorization information (e.g. Role-IDs) to be used in the decision for the request

Source: oneM2M TS-0003-Security-Solutions Release 2



Condition for Permit Granted

- For requests to <accessControlPolicy>
 resource type, authorization is granted if the
 request is evaluated to "Permit" for at least
 one selfPrivileges attribute.
- For other resource types, authorization is granted if the request is evaluated to "Permit" for at least one privileges attribute.



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Access Control Rules

The set of access control rules is denoted as acrs.

 $acrs = \{ acr(1), acr(2), ..., acr(k), ..., acr(K) \}$ Where

 $acr(k) = \{acr(k)_accessControlOriginators, acr(k)_accessControlOperations\}$ OR

 $acr(k) = \{acr(k)_accessControlOriginators, acr(k)_accessControlOperations, acr(k)_accessControlContexts\}$



ControlOriginators

- The accessControlOriginators parameter comprises a list of CSE-IDs and/or AE-IDs of any format defined in oneM2M TS-0001.
- It is allowed to include the wildcard characters, e.g. "*", into the URI string of CSE-ID and AE-ID at any level.







accessControlOperations

- The accessControlOperations parameter comprises a list of admissible operations which can be any subset of the following elements: Create, Request, Update, Delete, Discover, and Notify.
- While Create, Request, Update, Delete, and Notify operation are explicitly indicated in the op parameter of a request message.
- the Discovery operation is indicated by op = retrieve in combination with the provisioning of fc and Disrestype parameters in the request message.





accessControlContexts

Parameters of accessControlContexts

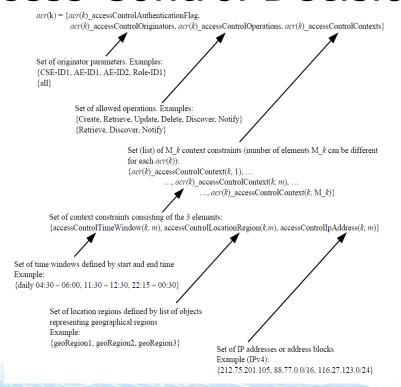
Parameter	Usage Description	Mandatory/Optional	Formats
accessControlTimeWindow	Set of Time Windows	0	List of time intervals where access
	that can be authorized		can be granted in extended
			crontab format
accessControlLocationRegion	Set of Location Regions that can be authorized	0	Latitude/longitude coordinates, and a radius defining a circular region around the coordinates Country code
accessControllpAddress	Set of IPv4 and IPv6 addresses that can be authorized	0	IPv4: dotted-decimal notation with CIDR suffix IPv6: colon separated groups of hexadecimal digits with CIDR suffix

Source: oneM2M TS-0003-Security-Solutions Release 2





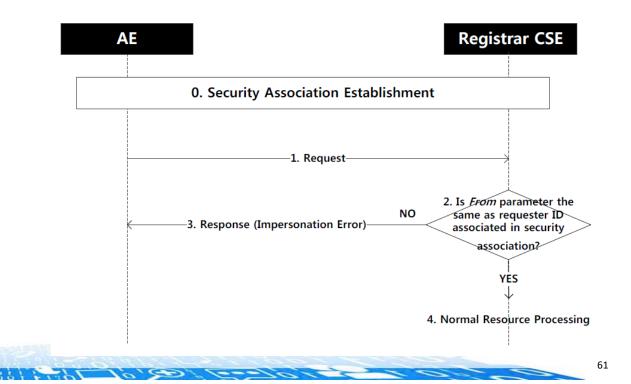
Access Control Decision







AE Impersonation Prevention









- D. Boswarthick, O. Elloumi, O. Hersent (Editors), M2M Communications: A Systems Approach, Wiley, 2012
- 2. 3GPP TS 24.302 Access to the 3GPP Evolved Packet Core (EPC) via non-3GPP access networks
- 3. 3GPP TS 33.102 3G Security Architecture
- 4. 3GPP TS 33.402 System Architecture Evolution (SAE); Security aspects of non-3GPP accesses
- 5. ETSI TS 102 690 Machine-to-Machine communications (M2M); Functional Architecture
- 6. ETSI TS 102 921 Machine-to-Machine communications (M2M); mla, dla and mld interfaces
- 7. oneM2M TS-0001: Functional Architecture
- 8. oneM2M TS-0003: Security-Solutions
- 9. RFC 5246 The Transport Layer Security (TLS) Protocol Version 1.2
- 10. RFC 6347 Datagram Transport Layer Security Version 1.2





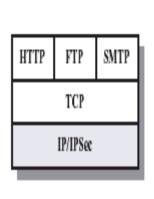
Appendix. TLS/DTLS Security Protocols

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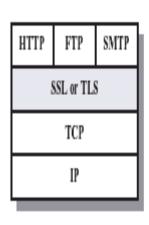




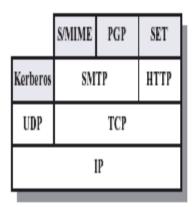
Layers of Security



(a) Network Level



(b) Transport Level



(c) Application Level





History of TLS

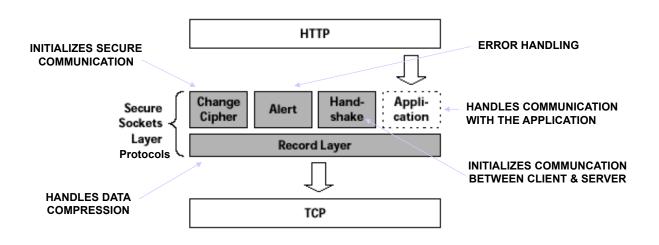
- Evolved from SSL
 - Unreleased v1 (Netscape)
 - Flawed-but-useful v2
 - Version 3 from scratch
 - Standard TLS1.0
 - o SSL3.0 with minor tweaks, hence Version field is 3.1
- Defined in RFC2246 (1999),
 http://www.ietf.org/rfc/rfc2246.txt
- Open-source implementation at http://www.openssl.org/





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TLS Protocol Architecture







TLS Protocol Details

TLS consists of two protocols

- Handshake protocol
 - Use public-key cryptography (or pre-shared key) to establish a shared secret key between the client and the server
- Record protocol
 - Use the secret key established in the handshake protocol to protect communication between the client and the server
- We will focus on the handshake protocol using public-key cryptography.







TLS Handshake Phase (1)

ClientHello ServerHello ServerCertificate + CertifcateRequest ClientCertificate ServerHelloDone

ClientHello

- List of the ciphersuites* and compression methods
- A random number

ServerHello

- A choice of the ciphersuite it can support
- A random number (different from client's one)
- Session ID

ServerCertificate

- Name and public key of the server
- Signed by a CA

ClientCertificate

- Name and public key of the server
- Signed by a CA

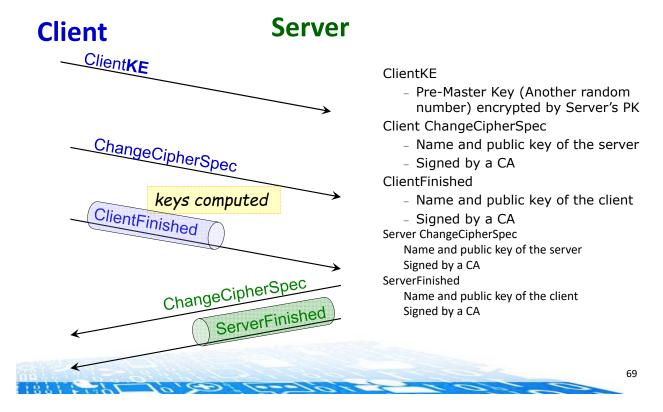
ServerHelloDone

 Ready to establish a mutual secret key for use in further communications





TLS Handshake Phase (2)







Key Computation

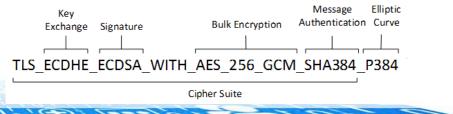
- Generation of Master Key from Clients' random number, Server's random number and Pre-Master Key
- Generation of Key Material from Master Key
- Generation of Symmetric Keys and other secret values from Key Material



Suito シスプラを支票和時代課題

*Cipher Suite

- A cipher suite is a combination of a key exchange algorithm, a bulk encryption algorithm, a message authentication code (MAC) algorithm, and a pseudorandom function (PRF).
- The key exchange algorithm, e.g. ECDHE_ECDSA, is used to determine if and how the client and server will authenticate during the handshake.
- The bulk encryption algorithm, e.g. AES_128_GCM, is used to encrypt the message stream. It also includes the key size and the lengths of explicit and implicit initialization vectors (cryptographic nonces).
- The message authentication code algorithm, e.g. SHA256, is used to create the message digest, a cryptographic hash of each block of the message stream.
- The pseudorandom function, e.g. P384, using the MAC algorithm's hash function, is used to create the master secret, a 48-byte secret shared between the two peers in the connection. The master secret is used as a source of entropy when creating session keys, such as the one used to create the MAC.



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DTLS

- Datagram Transport Layer Security (DTLS) is defined in RFC 4347.
- TLS has been designed for reliable transport protocols, that is it expects no lost or reordered messages from the transport layer.
- If it detects that a message is lost or out of order, it reasonably assumes an attack and drops the connection.
- DTLS is designed to provide reliable transport protocols over unreliable transport protocols where message losses and reordering are very likely.
- It is designed to tolerate unreliability and not creating any security issues while doing so.





DTLS Handshake Phase

Client		Server		
ClientHello			Flight 1	
		HelloVerifyRequest	Flight 2	
ClientHello			Flight 3	
	•	ServerHello ServerCertificate CertificateRequest ServerKeyExchange ServerHelloDone	Flight 4	
ClientCertificate ClientKeyExchange CertificateVerify ChangeCipherSpec Finished	→		Flight 5	
	-	ChangeCipherSpec Finished	Flight 6	73





Considerations of Unreliable Transport – DTLS Flight & Timer

- With an unreliable transport UDP, DTLS has to ensure the reliability of handshake messages itself. Therefore, it needs a timer to retransmit lost messages.
- For increased efficiency, DTLS does not use a timer for every message, but for bundles of messages, called fights.
- A fight contains all messages before the sending side changes. For every fight sent a timer is started, and if there is no response until the timer expires, the entire fight will be retransmitted.





Other Issues of DTLS

- DTLS uses an additional handshake message, called HelloVerifyRequest to prevent DOS attacks.
- DTLS has to provide its own fragmentation mechanism.
- DTLS has to deal with reordered messages, which can likely occur with unreliable transports.
- DTLS has its own Replay Check.