

Discrete Maths

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Information

Rolling a dice

Event A

"it's 2"

$$P(A) = \frac{1}{6} = 0.1666$$

P.S. Dice // Global Terminology.
 $\{1, 4\}$ are termed red outcome.
 $\{2, 3, 5, 6\}$ are termed black outcome.

Event B

"it's red"

$$P(B) = \frac{2}{6} = \frac{1}{3} = 0.3333$$

$$P(A)$$

$$\frac{1}{6} \parallel 0.1666$$

<

$$P(B)$$

$$\frac{2}{6} \parallel 0.3333$$

Less probability

More Information

More probability
Less Information

Probability is inverse of Information.

i.e.

~~Reguly~~ there is no sudden break
at XYZ college,
that is lectures/labs are conducted
100% other than public holidays.

And suddenly somebody told,
"There are no lectures today
(assume it's not ~~public~~ holiday)[↑]

For the real world,
probability of having ~~no~~^{no} lectures on regular
day is a little
the information is high.

Fraud Detection in Data Analytics

Routinely you are using your credit card in certain reason,
but suddenly someone has used the same ~~cc~~ in region Y

there is very less probability / chance but banks consider this as "high value information" to act against fraud.

This discussion shows
any unerring message contains
information and this can be
measured with maths upto a level.

How about formula

$$\text{Information} = -\log_2(\text{Percent})$$

it's 4

$$P(A) \frac{1}{6}$$

$$-\log_2 \frac{1}{6}$$

$$-(\log_2 1 - \log_2 6)$$

$$= -(-\log_2 6) = 2.585$$

it's red $\frac{2}{6}$

$$-\log_2 \frac{1}{3}$$

$$-(\log_2 1 - \log_2 3)$$

$$-(-\log_2 3)$$

$$= 1.585$$

Note that information can take greater than 1 value too.

Event A

it's 4

2.585

Event B

it's 80

>

1.585

[Can information
take
-ve
value?]

Event A has given more information compare to event B and math shows that.

Here, when we used \log base 2, the measurement unit of information has a global name "bit" binary digit.

Note that P is always ≤ 1 , $\log_2 P$ is always nonpositive. $-\log_2 P$ is always non-negative.

Moreover,

the smaller the P , larger the
Quantity $-\log_2 P$ (information).

This is exactly what we want.

Other ways are possible

$1/P$, $1/P^2$, $(-P)$, $\log_P P$, etc.

Entropy (randomness, chaos, disorder
uncertainty, unpredictability, etc)

"The world and especially life
has the highest entropy of them
all"

Entropy of a random variable is
a function which attempts to
characterize the unpredictability/
chaos / disorder/uncertainty of a
random variable.

Event A

fair 6 sided die

Event B

roulette wheel at Casino

36+0 || 36+0+00

which event has more uncertainty?

possible

outcomes are 6

|| Possible outcomes are
36 || ignore, 0/00

Event ————— B has more uncertainty

and hence more
entropy.

"less predictable"⁴

Moreover, entropy is not just about possible outcome but also about their frequency

Event M

Fair die

Event N

Weighted die

i.e.

90% of time its 2
10% 1, 3, 4, 5, 6

less uncertainty

less entropy

"more predictable"

Entropy

$$H(X) = - \sum_{x \in X} P(x) \cdot \log P(x)$$

Compute entropy of a fair coin.

$$P(X=\text{heads}) = \frac{1}{2} \quad P(X=\text{tails}) = \frac{1}{2}$$

H(X)
by calc.
values, ansatz.

$$\begin{aligned} H(P) &= - \sum_{x \in \{\text{heads, tails}\}} P(x) \log P(x) \\ &= - \left[\frac{1}{2} \log \frac{1}{2} + \frac{1}{2} \log \frac{1}{2} \right] \\ &= - \left[-\frac{1}{2} + (-\frac{1}{2}) \right] = -[-1] = 1 \end{aligned}$$

Joint entropy

$$H(E, C) = - \sum_{e \in E} \sum_{c \in C} p(e, c), \log(p(e, c))$$

Example

X Whether it's sunny or rainy

Y Above 70° or below 70°

$$P(\text{sunny, hot}) = \frac{1}{2}$$

$$P(\text{sunny, cool}) = \frac{1}{4}$$

$$P(\text{rainy, hot}) = \frac{1}{4}$$

$$P(\text{rainy, cool}) = 0, \text{ b.s. it's cool when there is thunder or snow.}$$

$$H(x,y) =$$

$$= - \left[\frac{1}{2} \log \frac{1}{2} + \frac{1}{4} \log \frac{1}{4} + \frac{1}{4} \log \frac{1}{4} + 0 \log 0 \right]$$

$$= - \left[-\frac{1}{2} + -\frac{1}{2} + -\frac{1}{2} + 0 \right]$$

$$= \frac{3}{2}$$

How about mutual information

"A quantity that measures a relationship between two random variables that are sampled simultaneously"

On average, how much information is communicated in one random variable about another.

We will prove that mutual information both ways is exactly same. $I(A, B) = I(B, A)$

$$I(x, y) = \sum_{x \in X} \sum_{y \in Y} P(x, y) \cdot \log \left(\frac{P(x, y)}{P(x) \cdot P(y)} \right)$$

Here, $P(X)$ and $P(Y)$
are marginal distributions of X and Y
obtained through the marginalization

process in probability

"Summing out"

i.e.

Independent events

$$P(A \cap B) = P(A) \times P(B)$$

(m. Exclusive)
 $P(A \cup B) = P(A) + P(B)$

Not m. 
 $P(A \cup B) = P(A) + P(B) - P(A \cap B)$

|| Conditional Probs (dependent events)
 $P(A \cap B) = P(B) \cdot P(A|B)$

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Mutual Information

$$I(A, B)$$

Amount of information concerning the occurrence of A that is contained in statement asserting the occurrence of B.

Surprisingly, this is exactly same as

$$I(B, A)$$

Amount of information concerning the occurrence of B that is contained in statement asserting the occurrence of A.

$$I(A;B) = -\log_2 P(A) - (-\log_2 P(A|B))$$

$$= -\log_2 P(A) + \log_2 P(A|B)$$

$$= -\log_2 P(A) + \log_2 \frac{P(A \cap B)}{P(B)}$$

$$= -\log_2 P(A) + \log_2 P(B \cap A) - \log_2 P(B)$$

. / .

$$= -\log_2 P(B) + \log_2 P(B \cap A) - \log_2 P(A)$$

$$= -\log_2 P(B) + \log_2 \frac{P(B \cap A)}{P(A)}$$

$$= -\log_2 P(B) + \log_2 P(B|A)$$

$$= -\log_2 P(B) - (-\log_2 P(B|A))$$

$$= I(B;A)$$

note
that
mutual
information
can
have
-ve
value
to.

$x-y$,
what
if
 y is
greater.

Example

A it's 4

$$P(A) = \frac{1}{6}$$

B, it's not
 $\{1, 4\}$

$$P(B) = \frac{2}{6} = \frac{1}{3}$$

C, even number
 $\{2, 4, 6\}$

$$P(C) = \frac{3}{6} = \frac{1}{2}$$

$$P(A \cap B) = \frac{1}{6} = P(B \cap A)$$

$$P(A \cap C) = \frac{1}{6}$$

$$P(A|B) = \frac{\frac{1}{6}}{\frac{1}{3}} = \frac{1}{6} \times \frac{3}{1} = \frac{1}{2}$$

$$P(A|C) = \frac{\frac{1}{6}}{\frac{1}{2}} = \frac{1}{6} \times \frac{2}{1} = \frac{1}{3}$$

$$P(B|A) = \frac{\frac{1}{6}}{\frac{1}{6}} = 1$$

$$P(C|A) = \frac{\frac{1}{6}}{\frac{1}{6}} = 1$$

$$\begin{aligned}
 I(A, B) &= -\lg p(A) + \lg p(A|B) \\
 &= -\lg \frac{1}{6} + \lg \frac{1}{2} \\
 &= 2.585 - 1 \\
 &= 1.585 \text{ bits}
 \end{aligned}$$

$$\begin{aligned}
 I(B, A) &= -\lg p(B) + \lg p(B|A) \\
 &= -\lg \left(\frac{1}{3}\right) + \lg 1 \\
 &= 1.585 + 0 \\
 &= 1.585
 \end{aligned}$$

$$\begin{aligned}
 I(A, C) &= -\log P(A) + \log P(A|_C) \\
 &= -\log \frac{1}{6} + \log \frac{1}{3} \\
 &= 2.585 - 1.585 \\
 &= 1 \text{ bit}
 \end{aligned}$$

$$\begin{aligned}
 I(C, A) &= -\log P(C) + \log P(C|_A) \\
 &= -\log \frac{1}{2} + \log 1 \\
 &= -(-1) + 0 \\
 &= 1 \text{ bit}
 \end{aligned}$$

Note that M.I. is symmetric measure.

Example

Consider the problem of estimating the likelihood that there will be 1-hour examination when the professor is scheduled to go out of town.

[observe probability and information, mutual info values carefully]

x_1 : Professor out of town and exam ~~conducted~~

x_2 : Prof out of town and ~~exam not conducted~~

x_3 : Prof in town and exam conducted ??

x_4 : Prof in town and ~~exam not conducted~~

Let $S = \{x_1, x_2, x_3, x_4\}$

Let A denote the event that an exam is conducted.

B Professor is out of town

$$P(A) = \frac{1}{2} + \frac{3}{16}$$

$$P(A/B) = \frac{\frac{1}{2}}{\frac{1}{2} + \frac{1}{16}} = \frac{8}{9}$$

Let C denote the event that
the professor is in town

$$P(A/C) = \frac{\frac{3}{16}}{\frac{3}{16} + \frac{1}{4}} = \frac{3}{7}$$

	P ^{out} of Town	P ⁱⁿ of Town	Exam Conducted
$P(x_1)$	$\frac{1}{2}$	Yes	Yes
$P(x_2)$	$\frac{1}{16}$	Yes	Given
$P(x_3)$	$\frac{3}{16}$	No	Yes
$P(x_4)$	$\frac{1}{4}$	No	No

Information that exam will be given / already given

$$-\lg p(A) = -\lg \frac{11}{16} = -(\lg 11 - \lg 16) = 3 \cdot 4.6 + 4 = 0.54 \text{ bits}$$

mutual

Information about the fact

"professor is out of town" on the fact
that exam will be given

$$\begin{aligned} I(A, B) &= -\lg \frac{11}{16} + \lg \frac{8}{9} \\ &= -3.46 + 4 + 3 - 3.17 \\ &= 0.37 \text{ bits} \end{aligned}$$

$$I(A, C) = -\lg \frac{11}{16} + \lg \frac{3}{7}$$

$$= -3.46 + 4 + 1.58 - 2.81 = -0.69 \text{ bits}$$

Logically, if the prof in town he shall conduct lecture but
exam was conducted is surprising shown by -ve MI.

Applications

Communication Theory

Transmitter ————— received

f. P.D.F, Probability density function
(continuous)

CDF Cumulative density function

$$CDF(x) = \text{Prob}(X \leq x)$$

PMF Probability mass function
(discrete)

probability distribution function ?

- funfair rings and products play
- shoot balloons
- what shall be MRP of a product?
- why casino / lottery company will always have profit?
- MCQ
- Bernoulli trials
- Movies
 - Focus 2015
 - Jesly and Merge go large
- Teen Patti
- 21 tablet
- Travelling Salesman Prob. (TSP)