### Large-Scale and Multi-Structured Databases

# **Document Databases**Data Modelling and Partitioning

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#### Some considerations

A document database could **theoretically** implement a **third normal form schema**.

Tables, as in relational databases, may be "simulated" considering collections with JSON documents with an identical pre-defined structure.

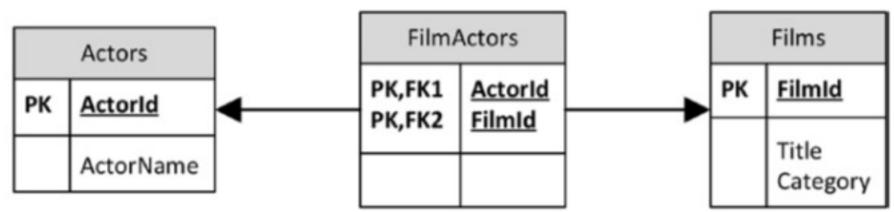


Image extracted from "Guy Harrison, Next Generation Databases, Apress, 2015"







#### JSON Databases: An example

```
{ " id" : 97, "Title" : "BRIDE INTRIGUE",
              "Category" : "Action",
    "Actors" :
      [ { "actorId" : 65, "Name" : "ANGELA HUDSON" } ]
    id": 115, "Title": "CAMPUS REMEMBER",
             "Category": "Action",
                          Actor document
           "actorId": 45, "Name": "REESE KILMER" },
           "actorId": 168, "Name": "WILL WILSON"
[ " id" : 105, "Title" : "BULL SHAWSHANK",
               "Category" : "Action",
    "Actors" :
      [ { "actorId" : 2, "Name" : "NICK WAHLBERG" },
         { "actorId" : 23, "Name" : "SANDRA KILMER" } ]
```

Document databases usually adopts a *reduced number* of collections for modeling data.

**Nested documents** are used for representing **relationships** among the different entities.

Document databases *do not* generally provide *join operations*.

**Programmers like** to have the JSON structure map **closely to the object** structure of their code!!!

Image extracted from "Guy Harrison, Next Generation Databases, Apress, 2015"



#### Data Modeling: Document Embedding

The solution above allows the user to *retrieve* a film and all its actors in a *single operation*.

However, "actors" result to be *duplicated* across multiple documents.

In a complex design this could lead to *issues* and possibly *inconsistencies* if any of the "actor" attributes need to be changed.

Moreover, some JSON databases have some *limitations* of the maximum *dimension* of a single document.

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#### Data Modeling: Document Linking



In the solution above, an array of actor *IDs* has been embedded into the film document.

The IDs can be used to *retrieve* the documents of the actors (in an other collection) who appear in a film.

We are **rolling back** to a relational model!!! Now, at least two collections of document must be defined.







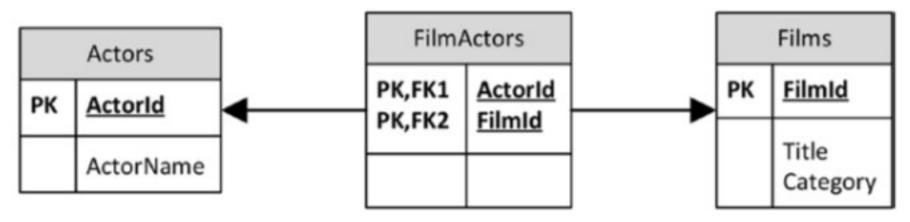
#### Data Modeling: Document Linking

```
{ "_id": 115, "Title": "CAMPUS REMEMBER", "Category": "Action"}

{ "_id": 99, "film": 115, "Actor: ":8, "Role": "Lead actor"}

{ "_id": 8, "Name": "MATTHEW JOHANSSON" }
```

We are *rolling back* to a the third normal form!!!!!









#### Data Modeling: Some Discussions

**Document linking** approaches are usually somewhat an **unnatural** style for a document database.

However, for *some workloads* it may provide the best *balance* between performance and maintainability.

When modeling data for document databases, there is no equivalent of third normal form that defines a "correct" model.

In this context, the *nature of the queries* to be executed *drives* the approach to *model* data.







# Data Modeling: One to Many Relationship

Let consider the *customer* entity which may have associated a list of *address* 

entities.

The basic pattern is that the **one** entity in a one-to-many relation is the **primary document**, and the **many** entities are represented as an array of **embedded documents**.







# Data Modelling: Many to Many Relationships

Let consider an example of application in which:

- A student can be enrolled in many courses
- A course can have many students enrolled to it

We can model this situation considering the following two collections:

```
{ courseID: 'C1667',
   title: 'Introduction to Anthropology',
  instructor: 'Dr. Margret Austin',
                                                                   {studentID:'S1837',
  credits: 3,
                                                                     name: 'Brian Nelson',
  enrolledStudents: ['S1837', 'S3737', 'S9825' ...
     'S1847'] },
                                                                     gradYear: 2018,
{ courseID: 'C2873',
                                                                     courses: ['C1667', C2873, 'C3876']},
   title: 'Algorithms and Data Structures',
  instructor: 'Dr. Susan Johnson',
                                                                   {studentID: 'S3737',
   credits: 3,
                                                                     name: 'Yolanda Deltor',
  enrolledStudents: ['S1837', 'S3737', 'S4321', 'S9825'
    ... 'S1847'] },
                                                                            gradYear: 2017,
{ courseID: C3876,
                                                                            courses: [ 'C1667', 'C2873']},
  title: 'Macroeconomics',
  instructor: 'Dr. James Schulen',
  credits: 3,
   enrolledStudents: ['S1837', 'S4321', 'S1470', 'S9825'
    ... 'S1847'] },
```

We have to take care when updating data in this kind of relationship. Indeed, the <u>DBMS will not control</u> the *referential integrity* as in relational DBMSs.

#### Data Modeling: An Example (I)

**Scenario**: Trucks in a company fleet have to transmit location, fuel consumption and other metrics every three minutes to a fleet management data base (one-to-many relationship between a truck and the transmitted details)

We may consider to generate a new document to add to the DB for each data transmission.

```
Let consider a document as follows:

{
    truck_id: 'T87V12',
    time: '08:10:00',
    date: '27-May-2015',
    driver_name: 'Jane Washington',
    fuel_consumption_rate: '14.8 mpg',
    ...
```

At the end of the day, the DB will include 200 new documents for each truck (we consider 20 transmissions per hour, 10 working hours)







#### Data Modeling: An Example (II)

An alternative solution may be to use embedded documents as follows:

```
truck_id: 'T87V12',
   date: '27-May-2015',
   driver_name: 'Jane Washington',
   operational_data:
                  {time : '00:01',
                   fuel_consumption_rate: '14.8 mpg',
                   ...},
                    {time : '00:04',
                   fuel_consumption_rate: '12.2 mpg',
                   ...},
                    {time : '00:07',
                   fuel_consumption_rate: '15.1 mpg',
                   ...},
}
```

Pay attention: we can have a potential *performance problem*!

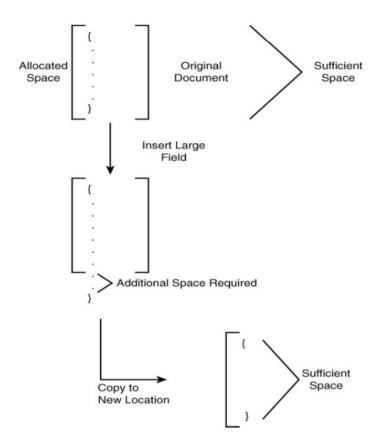






### Planning for Mutable Documents (I)

When a document is created, the DBMS allocates a certain amount of space for the document.



If the document *grows more* than the allocated space, the DBMS has to *relocate* it to another location.

Moreover, the DBMS *must free* the previously allocated space.

The previous steps can adversely affect the system performance.







### Planning for Mutable Documents (II)

A solution for avoiding to move oversized document is to *allocate sufficient* space at the moment in which the document is *created*.

Regarding the previous problem, the following solution may be adopted:

In conclusion, we have to consider the *life cycle* of a document and planning, if possible, the strategies for handling its growing.







#### **Indexing Document Database**

In order to avoid the entire scan of the overall database, DBMSs for document databases (for example MongoDB) allow the definition of *indexes*.

Indexes, like in book indexes, are a **structured set of information** that maps from one attribute to related information.

In general, indexes are *special data structures* that store a small portion of the collection's data set in an *easy to traverse* form.

The index stores the value of a specific field or set of fields, ordered by the value of the field.

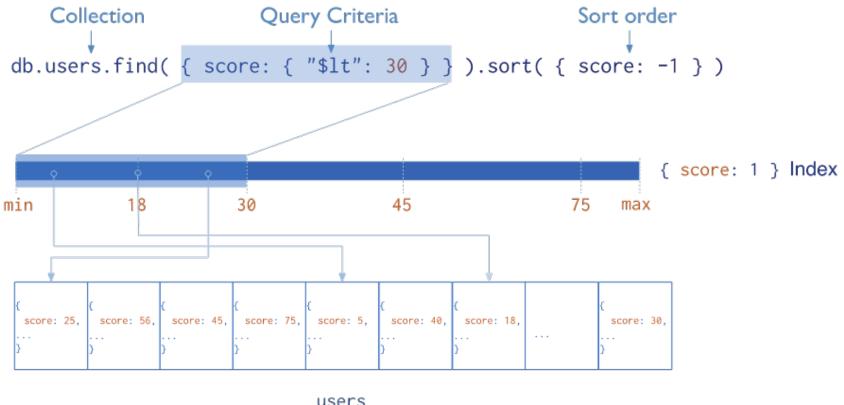
The ordering of the index entries supports *efficient equality matches* and *range-based query operations*.







#### An Example of Index





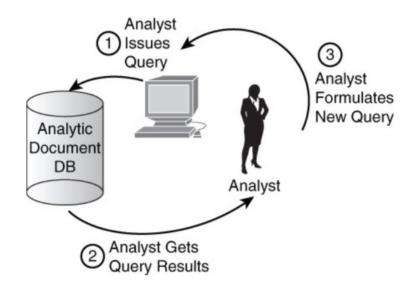






#### Read-Heavy Applications

In the figure we show the classical scheme of a read-heavy application (business intelligence and analytics applications):



In this kind of applications, the use **of several indexes** allows the user to quickly access to the database. For example, indexes can be defined for easily retrieve documents describing objects related to a specific **geographic region** or to a specific **type**.







#### Write-Heavy Applications

The example of the truck information transmission (each three minutes) is a typical write-heavy application.

The *higher* the number of *indexes* adopted the *higher* the amount of *time* required for closing a write operation.

Indeed, all the *indexes* must be *updated* (and created at the beginning).

Reducing the number of indexes, allow us to obtain systems with *fast write* operation responses. On the other hand, we have to accept to deal with *slow read operations*.

In conclusion, the number and the type of indexes to adopt must be identified as a *trade-off* solution.

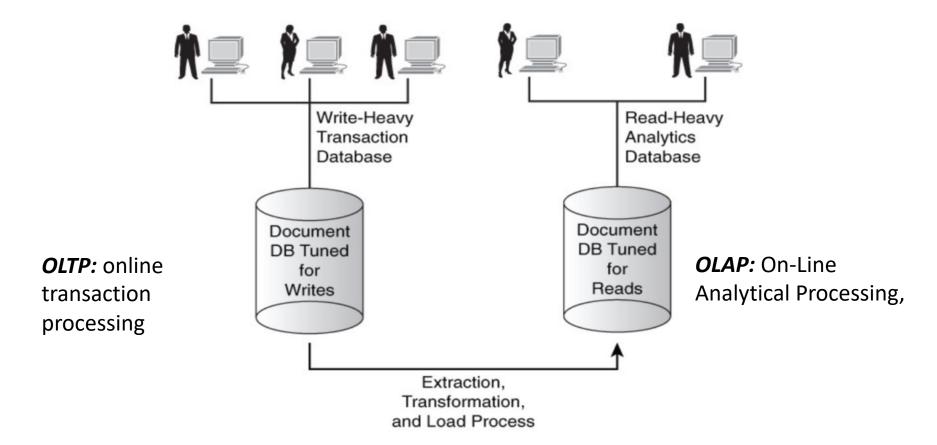






#### **Transactions Processing Systems**

These systems are designed for fast write operation and targeted reads, as shown in the figure below:









### Sharding

- Sharding or horizontal partitioning is the process of dividing data into blocks or chunks.
- Each block, labeled as shard, is deployed on a specific node (server) of a cluster.
- Each node can contain only one shard.
- In case of data replication, a shard can be hosted by more than one node.

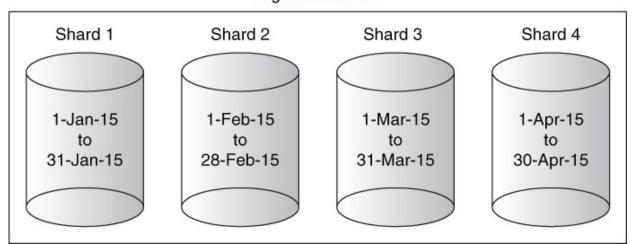


Image extracted from: "Dan Sullivan, NoSQL For Mere Mortals, Addison-Wesley, 2015"





#### Advantages of Sharding

- Allows handling heavy loads and the increase of system users.
- Data may be *easily distributed* on a variable number of servers that may be *added* or *removed* by request.
- Cheaper than vertical scaling (adding ram and disks, upgrading CPUs to a single server).
- Combined with *replications*, ensures a *high availability* of the system and *fast* responses.







#### **Shard Keys**

To implement sharding, document database designers have to select a **shard key** and a **partitioning method**.

A shard key is **one or more** fields, that exist **in all documents** in a collection, that is used to separate documents.

**Examples** of shard keys may be: Unique **document ID**, **Name**, **Date**, such as creation date, **Category** or type, Geographical region.

Actually, *any atomic field* in a document may be chosen as a shard key.

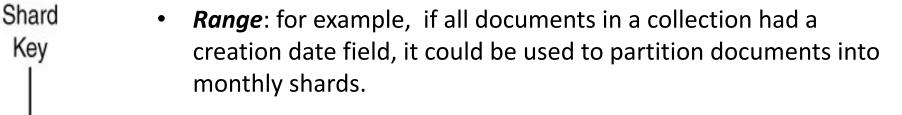






#### Partition Algorithms

There are three main categories of partition algorithms, based on:



• *Hashing*: a hash function can be used to determine where to place a document. *Consistent hashing* may be also used.

Algorithm

Shard N

• List: for example, let imagine a product database

• **List**: for example, let imagine a product database with several types (electronics, appliances, household goods, books, and clothes). These product types could be used as a shard key to allocate documents across five different servers.



Partitioning





#### Suggested Readings

Chapter 4 of the book "Guy Harrison, Next Generation Databases, Apress, 2015".

Chapters 7,8 of the book "Dan Sullivan, NoSQL For Mere Mortals, Addison-Wesley, 2015"





