General Morroach to Jolvina
General Approach to Solving Optimization problems using Dynamic Programmin
1. Characterize the structure of an opt. Solution
2- Recursively define the value of an opt.
3- Compute the value of an opt solution in a bottom up fashion
4 Contract are at sell a +1
4- Construct an opt. Sol. from computed intomation
4- Construct an opt. Sol. from Computed information
4- Construct an opt. Sol. from Computed information
4- Construct an opt. Sol. from Computed information.
4- Construct an opt. Sol. from Computed information.
4- Construct an opt. Sal. from Computed information
4- Construct an opt. Sal. from Computed information

Problem Statement
_ We have I resource
" n requests labeled 1 to n
_ Each request has start time Si,
finish time fi, and
- weight w:
Goal: Select a subset 5 = {1.n}
of mutually compatible interval
- of markey companies in in reveals
of mutually compatible intervals So as to Maximize \(\sum_{i \in S} \)

Case 1 if it is, value of the opt. So! =

wi + value of the opt. So! for

the sub problem that consists

only of compatible requests with i

Case 2 if it isn't, value of the opt. So! =

value of the opt. So! without job!

Sort requests in order of non-decreasing finish time.

fi (fz (fm

Define P(j) for an interval j to be the largest index i < j such that interval is j are disjoint.

Def Let O; denote the opt. Solution to the problem consisting of requests {1j} Let OPT(j) denote the value of O;
Def Let Oj denote the opt solution to the problem consisting of requests [1j] Let OPT(j) denote the value of Oj

Solution:
Compute - opt (j) if j = 0 then return 0
Singula - Spi (7)
- 1 1 = 0 Then
- seturn o
else





Memoization

Store the value of Compute-opt in a
globally accessible place the first

time we compute it. Then simply

use this precomputed value in

place of all future recursive

callo.

M-Compute-opt (j)

if j=0 then

return 0

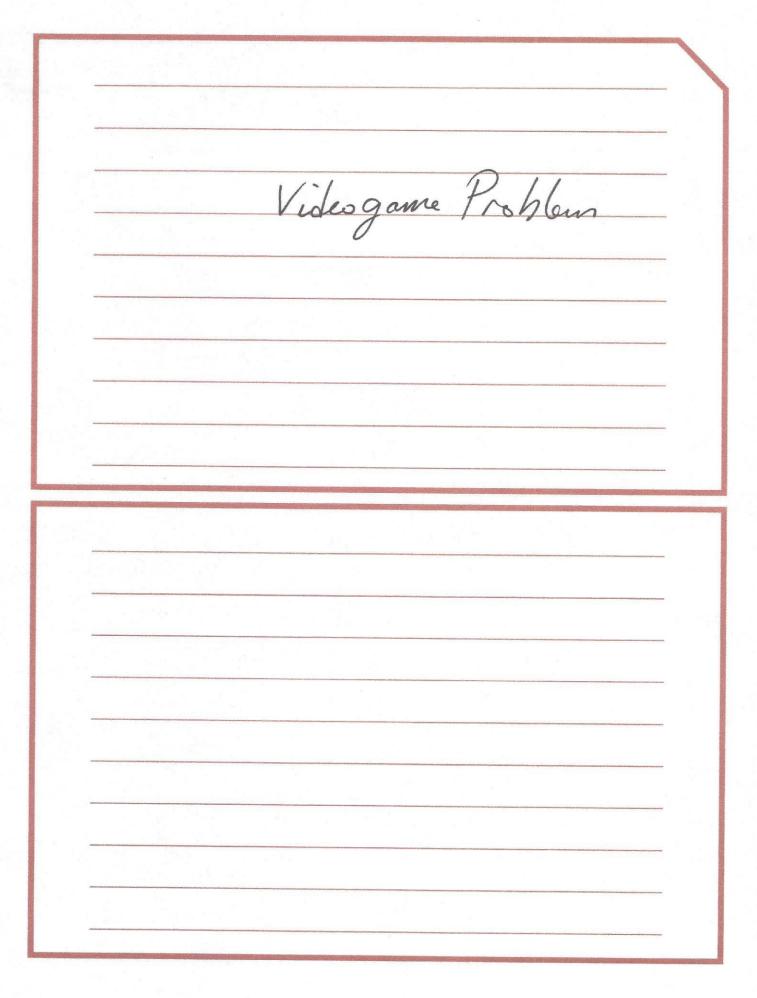
else if M(j) is not empty then





Find-Solution
if j>0 then if wj + M[p(j)] > M[j-1] then
- autput j toge ther us the results
output j together us the results J Find-Solution (P(j))
else
end if end if
Final-Solutions (j-1)
end if











0-1 knapsack &	
Subset Sum	

Problem Statement
A Single resource
- Requests {1n} each take time w:
- reguess (n) each take time wi
to process
- Can schedule jobs at any time between 0 to W
- between 0 to W
Objective: To schedule jobs such
that we massimize the machine's
- utilization
Sandsh

OPT(i,w) = value of the opt. Solution using a subset of the items {1. i? with Max. allowed weight w.

if $n \notin O$, Then OPT(n, w) =if $n \in O$, Then OPT(n, w) =

Of $\omega(\omega_i)$, then $OPT(i,\omega) = OPT(i-1,\omega)$ else, $OPT(i,\omega) = Max(OPT(i-1,\omega),$ $\omega_i + OPT(i-1,\omega-\omega_i)$

Subset-sum (n, w)	
	1
array $M[0,w]=0$ for each $w=0$	to W
for i=1 to n	
for w=o to W	
use recurrence formula	
to compute M[i, w]	
endfor	
endfor	
Return M[n,w]	



Dendo-polynomial time
An algorithm runs in seculo-polynomial time if its runing time is a polynomial in the numeric value of the input
of the imput

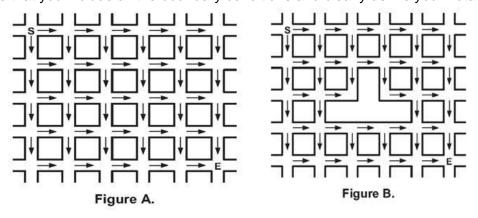
An algorithm runs in polynomial
time if its running time is a
polynomial in the lagth of the
input (or output).



Discussion 6

- **1.** You are to compute the <u>total number of ways</u> to make a change for a given amount m. Assume that we have an unlimited supply of coins and all denominations are sorted in ascending order: $1 = d_1 < d_2 < ... < d_n$. Formulate the solution to this problem as a dynamic programming problem.
- **2.** Graduate students get a lot of free food at various events. Suppose you have a schedule of the next *n* days marked with those days when you get a free dinner, and those days on which you must acquire dinner on your own. On any given day you can buy dinner at the cafeteria for \$3. Alternatively, you can purchase one week's groceries for \$10, which will provide dinner for each day that week (that day and the six that follow). However, because you don't have a fridge, the groceries will go bad after seven days (including the day of purchase) and any leftovers must be discarded. Due to your very busy schedule, these are your only two options for dinner each night. Your goal is to eat dinner every night while minimizing the money you spend on food.
- **3.** You are in Downtown of a city and all the streets are one-way streets. You can only go east (right) on the east-west (left-right) streets, and you can only go south (down) on the north-south (up-down) streets. This is called a Manhattan walk.
- a) In Figure A below, how many unique ways are there to go from the intersection marked S (coordinate (0,0)) to the intersection marked E (coordinate (n,m))?

 Formulate the solution to this problem as a dynamic programming problem. Please make sure that you include all the boundary conditions and clearly define your notations you use.



b) Repeat this process with Figure B; be wary of dead ends.









