



Performance of MPI Sends of Non-Contiguous Data Victor Eijkhout eijkhout@tacc.utexas.edu IPDPS / SNACS 2020



Question

MPI derived data types are very convenient; do they carry a performance penalty?



What are MPI derived types, and why

MPI send buffers are contiguous and single datatype; derived types describe non-contiguous or heterogeneous data

Convenience:

defining and using a derived type is much easier than allocate buffer, copy data, send, deallocate

```
MPI_Datatype newtype;
MPI_Type_something( .... description ... &newtype );
MPI_Type_commit( newtype );
MPI_Send( buffer, 1, newtype, dest,tag,comm );
MPI_Type_free(&newtype);
```

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Simple case

Send strided data;

```
float memory[2*N], buffer[N];
for (int i=0; i<N; i++)
  buffer[i] = memory[2*i];
MPI_Send( buffer, N, MPI_FLOAT, dest, tag, comm );</pre>
```

Simple performance model

Sending N words means

- read from memory to buffer;
- send buffer, probably overlapped with read
- \Rightarrow time O(N).

Strided data:

- Read 2N from memory;
- write back, probably overlapped;
- Read N from memory to buffer,
- send buffer, probably overlapped
- \Rightarrow time O(3N).

Expectation: strided send 1/3 performance of contiguous send.



Can we go faster than 1/3 peak?

Yes, but only with hardware support for direct streaming from non-contiguous locations.

M. Li, H. Subramoni, K. Hamidouche, X. Lu, and D. K. Panda. High performance mpi datatype support with user-mode memory registration: Challenges, designs, and benefits. In 2015 IEEE International Conference on Cluster Computing, pages 226–235, Sept 2015.

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Why would performance be lower than 1/3 peak?

- Construction of the datatype
- Index calculation overhead during copy

but in practice mostly:

allocation of internal MPI buffers.

Big problem: an application can allocate a buffer and reuse, but MPI needs to allocate, free, re-allocate, free, re-re-allocate,...



Simple use of derived types

Type vector and type subarray: same performance, so probably no overhead in index calculations. lower than 1/3 performance: internal buffer allocation

Force MPI to maintain a buffer

- Buffered sends: 1/3 performance as predicted
- Persistend sends: lower performance, reason unclear.

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One-sided communication

Could have the same performance as other non-buffer schemes.

In practice:

much worse performance for small messages (overhead), sometimes for large messages (reason unclear)



Packing

Use MPI_Pack to copy elements: one function call per element, so very slow

Pack one derived data: in practice attains 1/3 peak performance because reused buffer in user space.



Findings

- Performance often as expected; lower than hardware peak by 1/3
- Disappointments: persistend sends and one-sided performed worse than expected.
 - Exception: one-sided on the new Intel UCX layer seems improved.
- The processor is important next to the network:
 Stampede2 has Knights Landing and Skylake processors on the same network
 - The lower scalar performance of KNL is especially noticable on small messages.

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MPI derived types

Recommendations

 Use buffered sends if the calculation of buffer space is doable and not excessive.

(Note: total space of all simultaneously outstanding MPI_Bsend calls.)

Otherwise pack derived type.



MPI derived types

Repository of code and results

https://github.com/TACC/mpipacking

