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      General
run.sh
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
for i in *.in; do
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
done
Header
// use better compiler options
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
// include everything
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
using namespace std;
using namespace __gnu_cxx; // rope
using namespace __gnu_pbds; // tree/trie
// common defines
#define fastio

→ ios_base::sync_with_stdio(0);cin.tie(0);
#define nostacklim rlimit RZ; getrlimit(3,&RZ)
 \rightarrow );RZ.rlim cur=-1;setrlimit(3,\&RZ);
#define DEBUG(v) cerr<<__LINE__<<": "<<#v<<" =
\Rightarrow "<<v<<'\n'; #define TIMER
→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
#define i128 _ int128
#define u128 unsigned i128
#define ld long double
// global variables
mt19937 rng((uint32_t)chrono::steady

    _clock::now().time_since_epoch().count());
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c - '0';
void read(int& n) {
   char c; n = 0; int s = 1;
   if ((c=getchar_unlocked())=='-') s = -1;
 else n = c - \sqrt{0}:
 while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c -
void read(ld& n) {
 char c; n = 0;

.ld m = 0, o = 1; bool d = false; int s = 1;

.if ((c=getchar_unlocked())=='-') s = -1;
 else if (c == '.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n')
  if (c == '.') d = true;
```

```
else if (d) { m=m*10+c-'0'; o*=0.1; } else n = n * 10 + c - '0':
 n = s * (n + m * o):
void read(double& n) {
 ld m; read(m); n = m;
void read(float& n) {
  ld m; read(m); n = m;
 void read(string& s) {
 char c; s = ""
 while((c=getchar_unlocked())!=' '&&c!='\n')
bool readline(string& s) {
 char c; s = ""
 while(c=getchar_unlocked()) {
  if (c == '\n') return true;
if (c == EOF) return false;
  s += c:
 return false:
void print(unsigned int n) {
 if (n / 10) print(n / 10);
 putchar unlocked(n % 10 + '0'):
void print(int n) {
 if (n < 0) { putchar_unlocked('-'); n*=-1; }
 print((unsigned int)n);
Common Structs
   n-dimension vectors
// Vec<2, int> v(n, m) = arr[n][m]

// Vec<2, int> v(n, m, -1) default init -1

template<int D, typename T>
struct Vec : public vector < Vec < D-1. T>> {
 template<typename... Args>
 Vec(int n=0, Args... args) : vector<Vec<D-1,
 \rightarrow T>>(n, Vec<D-1, T>(args...)) {}
};
template<typename T>
struct Vec<1, T> : public vector<T> {
   Vec(int n=0. T val=T()) : vector<T>(n. val) {}
   Algorithms
Binary Search
 // search for k in [p,n)
template<typename T>
int binsearch(T x[], int k, int n, int p = 0) {
 for (int i = n; i >= 1; i /= 2)

while (p+i < n && x[p+i] <= k) p += i;

return p; // bool: x[p] == k;
Min/Max Subarray
   max - compare = a < b, reset = a < 0
 // min - compare = a > b, reset = a > 0
// returns {sum, {start, end}}
pair<int, pair<int, int>>
     ContiguousSubarray(int* a, int size,
     bool(*compare)(int, int),
 bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
 \rightarrow 0, s = 0;
 for (int i = 0: i < size: i++) {
  cur += a[i];
  if ((*compare)(best, cur)) { best = cur;
    start = s; end = i; }
  if ((*reset)(cur)) { cur = 0; s = i + 1; }
 return {best, {start, end}};
Quickselect
 #define QSNE -999999
```

```
int partition(int arr[], int 1, int r)
 int x = arr[r], i = 1;
for (int j = 1; j <= r - 1; j++)
if (arr[j] <= x)
swap(arr[i++], arr[j]);
 swap(arr[i], arr[r]);
 return i:
// find k'th smallest element in unsorted array,
\rightarrow only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - 1 + 1)) return QSNE;
 swap(arr[1 + rng() % (r-1+1)], arr[r]);
 int pos = partition(arr, 1, r);
 if (pos-l==k-1) return arr[pos];
 if (pos-l>k-1) return qselect(arr,l,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
|}
|// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
on both axis
pair<int. int> saddleback search(int** arr. int
\hookrightarrow x, int y, int v) {
int i = x-1, j = 0;
while (i >= 0 && j < y) {
   if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i--: j++;
 return {-1, -1};
Ternary Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a. int b. int (*f)(int)) {
 while (b-a > 4) {
  int m = (a+b)/2;
if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
 . a = i;
return a:
#define TERNPREC 0.000001
double ternsearch (double a, double b, double
     (*f)(double)) {
 while (b-a > TERNPREC * 4) {
    double m = (a+b)/2:
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  > = m;
else b = m + TERNPREC;
 for (double i = a + TERNPREC; i <= b; i +=
    TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
double goldsection(double a, double b, double
 double g(double)) {
  double r = (sqrt(5)-1)/2, eps = 1e-7;
  double x1 = b - r*(b-a), x2 = a + r*(b-a);
  double f1 = f(x1), f2 = f(x2);
 while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
   a = x1; x1 = x2; f1 = f2;
x2 = a + r*(b-a); f2 = f(x2);
 return a:
```

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3 Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums
\hookrightarrow O(\log n) updates, O(\log n) gets
struct Fenwick { int n: ll* tree:
 void update(int i, int val) {
 .++i;
while (i <= n) {
  tree[i] += val;</pre>
   i += i & (-i);
 Fenwick(int size) {
  n = size;
  tree = new ll[n+1];
for (int i = 1; i <= n; i++)
   .tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
 11 operator[](int i) {
  if (i < 0 || i > n) return 0;
  while (i>0)
  sum += tree[i];
i -= i & (-i);
  return sum;
 ll getRange(int a, int b) { return
    operator[](b) - operator[](a-1); }
Hashtable
// similar to unordered map, but faster
struct chash {
    const uint64_t C = (11)(2e18 * M_PI) + 71;
    ll operator()(11 x) const { return
    builtin bswap64(x*C): }
int main() {
  gp_hash_table<11,int,chash>
\rightarrow hashtable({},{},{},{},{1<<16});
for (int i = 0; i < 100; i++)

hashtable[i] = 200+i;

if (hashtable.find(10) != hashtable.end())

cout << hashtable[10];
Ordered Set
template <typename T>
using oset = tree<T,null type,less<T>,rb tree
    tag, tree order statistics node update>;
template <typename T, typename D> using omap = tree<T,D,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
int main()
 oset<int> o_set;
o_set.insert(5); o_set.insert(1);

    o_set.insert(3);
// get second smallest element
 cout << *(o_set.find_by_order(1));</pre>
 // number of elements less than k=4
cout << ' ' << o_set.order_of_key(4) << '\n';</pre>
 // equivalent with ordered map
 omap<int,int> o_map;
o_map[5]=1;o_map[1]=2;o_map[3]=3;
cout << (*(o_map.find_by_order(1))).first;
 cout << ' ' << o_map.order_of_key(4) << '\n';
```

```
Rope
                                                         // print things with prefix "1"
                                                         auto range = trie.prefix_range("1");
// O(\log n) insert, delete, concatenate
                                                         for (auto it = range.first; it !=
int main() {
 // generate rope
                                                         → range.second: it++)
 rope<int> v;
                                                          cout << *it << '
 for (int i = 0; i < 100; i++)
.v.push_back(i);
                                                        Wavelet Tree
 // move range to front
                                                        using iter = vector<int>::iterator;
 rope<int> copy = v.substr(10, 10);
v.erase(10, 10);
                                                        struct WaveletTree {
   Vec<2, int> C; int s;
 v.insert(copy.mutable_begin(), copy);
                                                          // sigma = highest value + 1
                                                         WaveletTree(vector<int>& a. int sigma) :
 // print elements of rope
for (auto it : v) cout << it << "";
                                                            s(sigma), C(sigma*2, 0) {
                                                          build(a.begin(), a.end(), 0, s-1, 1);
                                                         void build(iter b. iter e. int L. int U. int
Segment Tree
                                                          u) {
if (L == U) return
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
                                                           int M = (L+U)/2;
 typedef int T;
                                                           C[u].reserve(e-b+1); C[u].push back(0);
 static constexpr T UNIT = INT_MIN;
                                                          for (auto it = b; it != e; ++it)
  C[u].push_back(C[u].back() + (*it<=M));</pre>
 T f(T a, T b) {
 if (a == UNIT) return b;
if (b == UNIT) return a;
                                                           auto p = stable_partition(b, e, [=](int
                                                            i){return i<=M;});
  return max(a,b);
                                                          build(b, p, L, M, u*2);
 int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
                                                          build(p, e, M+1, U, u*2+1);
                                                          // number of occurrences of x in [0,i)
\rightarrow n(n) {}
                                                         int rank(int x, int i) {
   int L = 0, U = s-1, u = 1, M, r;
   while (L != U) {
 SegmentTree(vector<T> arr)

    SegmentTree(arr.size()) {

 for (int i=0:i<arr.size():i++)
                                                           M = (L+U)/2;
r = C[u][i]; u*=2;

    update(i,arr[i]);

                                                           if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
                                                          return i:
                                                          ^{\prime\prime} number of occurences of x in [l,r)
 T query(int b, int e) { // query [b, e)
                                                         int count(int x, int 1, int r) {
  return rank(x, r) - rank(x, 1);
  Tra = UNIT, rb = UNIT;
  for (b+=n, e+=n; b<=); b/=2, e/=2) {
    if (b % 2) ra = f(ra, s[b++]);
    if (e % 2) rb = f(s[--e], rb);
                                                         // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  return f(ra, rb):
                                                          while (L != U) {
   M = (L+U)/2;
 T get(int p) { return query(p, p+1); }
                                                           ri = C[u][1]; rj = C[u][r]; u*=2;
                                                           if (k \le rj-ri)^{n}l = ri, r = rj, U = M;
Sparse Table
                                                           else k -= řj-rí, l -= ŕi, r -= ŕj,
template < class T> struct SparseTable {
                                                           L = M+1, ++u;
 vector<vector<T>> m;
                                                           return U:
 SparseTable(vector<T> arr) {
  m.push back(arr);
  for (int k = 1: (1<<(k)) <= size(arr): k++)
                                                         // # elements between [x,y] in [l, r)
                                                         mutable int L, U;
  m.push_back(vector<T>(size(arr)-(1<(k)+1));
                                                         int range(int x, int y, int 1, int r) const {
  for (int i = 0; i < size(arr)-(1<<k)+1; i
                                                          if (y < x \text{ or } r <= 1) return 0;
                                                          L = x; U = y;
 [k][i] = min(m[k-1][i],
                                                          return range(1, r, 0, s-1, 1);
\rightarrow m[k-1][i+(1<<(k-1))]:
}
// min of range [l,r]
                                                         int range(int 1, int r, int x, int y, int u)
                                                         → const {
                                                          if (y < L or U < x) return 0;
if (L <= x and y <= U) return r-l;
T query(int 1, int r) {
 int k = _-lg(r-l+1);
                                                          int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return \min(m[k][1], m[k][r-(1<< k)+1]);
                                                          return range(ri, rj, x, M, u*2) + range(1-ri
}
};
                                                            r-rj, M+1, y, u*2+1);
                                                          ^{\prime}// # elements <= x in [l, r]
                                                         int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);
typedef trie<string, null_type,

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>

→ trie_type;

int main() {
                                                             Strings
 // generate trie
 trie_type trie;
                                                        Aho Corasick
 for (int i = 0; i < 20; i++)
                                                           range of alphabet for automata to consider
 trie.insert(to string(i)); // true if new,
                                                           MAXC = 26, OFFC = 'a' if only lowercase
\hookrightarrow false if old
```

```
|const int MAXC = 256:
const int OFFC = 0:
struct aho_corasick {
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
  int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
  for (int i = 0; i < size; i++) {
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
      s[cur].go[c-OFFC] = s.size();
      s.push back(state());
     cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue int> sq;
for (int c = 0; c < MAXC; c++) {
   if (s[0].so[c] != 0) {
      ...s[s[0].so[c]].fail = 0;</pre>
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
   for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
      int failure = s[e].fail;
while (s[failure].go[c] == -1)
        failure = s[failure].fail;
      failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
      for (auto length : s[failure].out)
s[s[e].go[c]].out.insert(length);
      sq.push(s[e].go[c]);
 // list of {start pos, pattern id}
  vector<pair<int, int>> search(string text)
  vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {
  while (s[cur].go[text[i]-OFFC] == -1)
    cur = s[cur].fail;
cur = s[cur].go[text[i]-OFFC];
    if (s[cur].out.size())
    for (auto end : s[cur].out)
. toret.push_back({i - end.first + 1,
     end.second):
  return toret:
Boyer Moore
struct defint { int i = -1; };
vector<int> boyermoore(string txt, string pat)
 vector<int> toret; unordered_map<char, defint>string lcp(string* arr, int n, bool sorted =
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0; i < m; i++) badchar[pat[i]].i
 \rightarrow = i;
int s = 0:
 while (s \leq n - m) {
  int j = m - 1:
```

while  $(j \ge 0 \&\& pat[j] == txt[s + j]) j--;$ 

.if (j < 0) {

```
..toret.push_back(s);
   s += (s + m < n) ? m - badchar[txt[s +
   m]].i : 1:
 .} else
   s += \max(1, i - badchar[txt[s + i]].i):
 return toret:
English Conversion
const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven", "eight", "nine";
const string teens[] ={"ten", "eleven",
    "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen";
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",
    "eighty", "ninety"};
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion", "quintillion", "sextillion",
    "septillion"};
string convert(int num, int carry) {
 if (num < 0) return "negative " +

    convert(-num, 0):

     (num < 10) return ones[num];
(num < 20) return teens[num % 10]
 if (num < 100) return tens[(num / 10) - 2] +
     (num\%10==0?"":"") + ones[num\%10]:
 if (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" + (num%100==0?"":" ") + convert(num % 100,
    0);
 return convert(num / 1000, carry + 1) + " " +
    mags[carry] + " " + convert(num % 1000.
    0):
string convert(int num) {
 return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = 0; i < n + 1; i++)
  next[i] = 0;
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
   j = next[j];
  if (j > 0 | pat[j] == pat[i])
   next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
  if (txt[i] == pat[j]) {
  if (++j == n)
    toret.push_back(i - j + 1);
  } else if (j > 0) {
...j = next[j];
 return toret;
Longest Common Prefix (array)
 // longest common prefix of strings in array
 → false) {
idise; l
if (n == 0) return "";
if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==

    arr[n-1][v])
    r += arr[0][v++];

 return r:
```

```
Longest Common Subsequence
                                                       unsigned int mwb(string s, set<string> dict) { | for (int j = 1; j <= n; j++)</pre>
                                                        int l = s.size();
string lcs(string a, string b) {
                                                        vector<unsigned int> arr(l+1, -1);
 int m = a.length(), n = b.length();
                                                        arr[0] = 0;
for (int i = 0; i < 1; i++) {
 int L[m+1][n+1];
 for (int i = 0; i <= m; i++) {
                                                         if (arr[i] != -1) {
 for (int j = 0; j <= n; j++) {
...if (i == 0 || j == 0) L[i][j] = 0;
...else if (a[i-1] == b[j-1]) L[i][j] =
                                                         for (auto e : díct) {
                                                           int L = e.size();
if (1 >= i + L) {
                                                             bool isGood = true;
\hookrightarrow L[i-1][j-1]+1;
                                                         .....for (int j = 0; isGood && j < L; j++)
.....if (s[i+j] != e[j])
   else L[i][j] = \max(L[i-1][j], L[i][j-1]);
                                                            isGood = false;
if (isGood)
 // return L[m][n]; // length of lcs
                                                            arr[i+L] = min(arr[i]+1, arr[i+L]);
 string out = "";
 int i = m - 1, j = n - 1;
 while (i >= 0 && j >= 0) {
 if (a[i] == b[j]) {
                                                        return arr[1];
   out = a[i--] + out;
                                                       Hashing
  else if (L[i][j+1] > L[i+1][j]) i--;
                                                       #define HASHER 27
  .else j--;
                                                       ull basicHash(string s) {
                                                        ull v = 0:
 return out;
                                                        for (auto c : s) v = (c - 'a' + 1) + v *
                                                        → HASHER;
                                                        return v;
Longest Common Substring
// l is array of palindrome length at that
                                                       const int MAXN = 1000001:
int manacher(string s, int* 1) {
                                                       ull base[MAXN] = {1};
void genBase(int n) {
 int n = s.length() * 2;
                                                        for (int i = 1; i \le n; i++)

base[i] = base[i-1] * HASHER;
 for (int i = 0, j = 0, k; i < n; i += k, j =
 \rightarrow max(j-k, 0)) {
                                                       struct advHash {
  ull v, l; vector<ull> wip;
  while (i >= j \&\& i + j + 1 < n \&\& s[(i-j)/2]
 \Rightarrow = s[(i+j+1)/2]) j++;
                                                        advHash(string& s): v(0) {
  .1[i] = j;
                                                         wip = vector<ull>(s.length()+1);\
 for (k = 1; i >= k && j >= k && l[i-k] !=
                                                         \sin c = c \cos w
   i-k: k++)
                                                         for (int i = 0; i < s.length(); i++)</pre>
  1[i+k] = min(1[i-k], j-k);
                                                          wip[i+1] = (s[i] - 'a' + 1) + wip[i] *
                                                           HASHER:
 return *max_element(1, 1 + n);
                                                         1 = s.length(): v = wip[1]:
                                                        ull del(int pos, int len) {
Cyclic Rotation (Lyndon)
                                                         return v - wip[pos+len]*base[l-pos-len] +
// simple strings = smaller than its nontrivial
                                                           wip[pos]*base[1-pos-len];
\rightarrow suffixes
                                                        ull substr(int pos, int len) {
// lyndon factorization = simple strings
                                                         return del(pos+len, (l-pos-len)) -
\hookrightarrow factorized
 // "abaaba" -> "ab", "aab", "a"
                                                           wip[pos]*base[len];
vector<string> duval(string s) {
 int n = s.length();
                                                        ull replace(int pos, char c) {
                                                         return v - wip[pos+1]*base[l-pos-1] + ((c -
 vector<string> lyndon;
 for (int i = 0; i < n;) {
   int j = i+1, k = i;
   for (; j < n && s[k] <= s[j]; j++)
   if (s[k] < s[j]) k = i;
                                                            'a' + 1) + wip[pos] *
                                                           HASHER) *base[1-pos-1];
                                                        ull replace(int pos, string s) {
                                                         // can't increase total string size
  for (; i <= k; i += j - k)
                                                         ull r = v -
   lyndon.push back(s.substr(i,j-k));
                                                            wip[pos+s.size()]*base[l-pos-s.size()], c
                                                          wip[pos];
 return lyndon;
                                                         for (int i = 0; i < s.size(); i++)
c = (s[i]-'a'+1) + c * HASHER;
}
// lexicographically smallest rotation
                                                          return r + c * base[l-pos-s.size()];
int minRotation(string s) {
 int n = s.length(); s += s;
 auto d = duval(s); int i = 0, a = 0;
while (a + d[i].length() < n) a +=</pre>
                                                       Subsequence Count

    d[i++].length();

                                                        // "banana", "ban" >> 3 (ban, ba..n, b..an)
 while (i && d[i] == d[i-1]) a -=
                                                       ull subsequences(string body, string subs) {

    d[i--].length();

                                                        int m = subs.length(), n = body.length();
                                                        if (m > n) return 0;
 return a;
                                                        ull** arr = new ull*[m+1];
                                                        for (int i = 0; i \le m; i++) arr[i] = new
Minimum Word Boundary
                                                         \rightarrow ull[n+1];
// minimum word boundary
                                                        for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[0][i] = 1;
for (int i = 1; i <= m; i++)
// compose string s using words from dict
// NOTE: can reuse words from dict
```

```
arr[i][j] = arr[i][j-1] + ((body[j-1] ==
   subs[i-1])? arr[i-1][j-1] : 0);
 return arr[m][n];
Suffix Array + LCP
struct SuffixArray {
 vector<int> sa, 1cp;
 SuffixArray(string& s, int lim=256) {
  int n = s.length() + 1, k = 0, a, b;
  vector<int> x(begin(s), end(s)+1), y(n),
 \rightarrow ws(max(n, lim)), rank(n);
  sa = lcp = y;
  iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j)
 \rightarrow 2), lim = p) {
  p = j; iota(begin(y), end(y), n - j);
   for (int i = 0; i < (n); i++)
if (sa[i] >= i)
     .y[p++] = sa[i] - j;
   fill(begin(ws), end(ws), 0);
for (int i = 0; i < (n); i++) ws[x[i]]++;
   for (int i = 1; i < (lim); i++) ws[i] +=
    ws[i - 1];
   for (int i = n; i--;) sa[--ws[x[y[i]]]] =
    v[i];
   swap(x, y); p = 1; x[sa[0]] = 0;
for (int i = 1; i < (n); i++) {
   a = sa[i - 1]; b = sa[i];</pre>
    x[b] = (y[a] = y[b] \&\& y[a + j] = y[b +
    j]) ? p - 1 : p++;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
  for (int i = 0, j; i < n - 1; lcp[rank[i++]]
    = k
  for (k \&\& k--, j = sa[rank[i] - 1];
     s[i + k] == s[j + k]; k++);
 // smallest cyclic shift
 int cyclic() { return sa[0]; }
 // longest repeated substring
 pair<int,int> lrs() {
  int length = -1, index = -1;
  for (int i = 0; i < lcp.size(); i++) {
  if (lcp[i] > length) {
   length = lcp[i];
    index = sa[i]:
  return {index,length};
 }
// count distinct substrings, excluding empty
 int distincts() {
   int n = sa.size() - 1, r = n - sa[0];
   for (int i = 1; i < lcp.size(); i++)</pre>
  r += (n - sa[i]) - lcp[i - 1];
  return r:
 }// count repeated substrings, excluding empty
 int repeateds() {
  .int r' = 0;
  for (int i = 1; i < lcp.size(); i++)
  r += \max(lcp[i] - lcp[i-1], 0);
  return r;
Suffix Tree (Ukkonen's)
struct SuffixTree {
 // n = 2*len+10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
 void ukkadd(int i, int c) { suff:
```

```
if (r[v]<=q) {
  if (q==-1 || c==toi(a[q])) q++; else {
...l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
   p[m] = p[v]; t[m][c] = m+1; t[m][toi(a[q])] = v;
   l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
v=s[p[m]]; q=l[m];
   while (q < r[m]) { v = t[v][toi(a[q])];
   q+=r[v]-l[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
   q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
  fill(r,r+N,(int)(a).size());
 memset(s, 0, sizeof s);
memset(t, -1, sizeof t);
fill(t[1],t[1]+ALPHA,0);
s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
  for(int i=0; i < a. size(); i++)
    ukkadd(i,toi(a[i]));
 // Longest Common Substring between 2 strings
 // returns {length, offset from first string}
 pair<int. int> best:
 int lcs(int_node, int i1, int i2, int olen) {
 if (1[node] <= i1 && i1 < r[node]) return 1;
if (1[node] <= i2 && i2 < r[node]) return 2;
  int mask=0
   len=node?olen+(r[node]-l[node]):0;
  for(int c=0; c<ALPHA; c++) if
   (t[node][c]!=-1)
mask |= lcs(t[node][c], i1, i2, len);
  if (mask==3)
  best=max(best, {len,r[node]-len});
  return mask:
 static pair<int, int> LCS(string s, string t)
 \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2));
 st.lcs(0, s.size(), s.size()+t.size()+1, 0); return st.best;
String Utilities
void lowercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
   ::tolower);
void uppercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
\hookrightarrow ::toupper);
void trim(string &s) {
 s.erase(s.begin(),find_if_not(s.begin(),s
     .end(),[](int c){return
    isspace(c);}));
 s.erase(find_if_not(s.rbegin(),s.rend(),[](int

    c){return isspace(c);}).base(),s.end());

vector<string> split(string& s, char token) {
    vector<string> v; stringstream ss(s);
    for (string e;getline(ss,e,token);)
        v.push_back(e);
    return v;
    Greedy
```

```
Interval Cover
int (L,R) = interval [L,R], in = \{\{l,r\}, index\}\} string a; (L,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}

    vector<pair<pair<double,double>,int>> in) {
```

```
int i = 0; pair<double,int> pos = {L,-1};
                                                         .re %= mo;
    vector<int> a;
     sort(begin(in), end(in));
                                                         re = (re + mo) \% mo;
                                                         return true:
     while (pos.first < R) {
         double cur = pos.first;
while (i < (int)in.size() &&</pre>
                                                        Count Digit Occurences
    in[i].first.first <= cur)</pre>
                                                        /*count(n,d) counts the number of occurences of
                                                        \rightarrow a digit d in the range [0,n]*/
     max(pos,{in[i].first.second,in[i].second})
                                                        ll digit_count(ll n, ll d) {
                                                         ll result = 0:
         if (pos.first == cur) return {};
                                                         while (n != 0) {
   result += ((n%10) == d ? 1 : 0):
         a.push_back(pos.second);
                                                         n /= 10;
     return a;
                                                         return result;
6 Math
                                                        11 count(ll n, ll d) {
                                                         if (n < 10) return (d > 0 && n >= d);
if ((n % 10) != 9) return digit_count(n, d) +
 Catalan Numbers
ull* catalan = new ull[1000000];
count(n-1, d):
                                                         return 10*count(n/10, d) + (n/10) + (d > 0):
                                                        Discrete Logarithm
                                                        int discretelog(int a, int b, int m) {
                                                         11 n = sqrt(m) + 1, an = 1;
                                                        for (ll i = 0; i < n; ++i)
an = (an * a) % m:
 if (catalan[i] >= mod)
catalan[i] -= mod;
                                                         unordered_map<11, 11> vals;
                                                         for (ll q = 0, cur = b; q <= n; q++) {
  vals[cur] = q;
 // TODO: consider binomial coefficient method
                                                          cur = (cur * a) \% m;
Combinatorics (nCr, nPr)
                                                         for (ll p = 1, cur = 1; p \le n; p++) {
                                                          cur = (cur * an) % m;
 // can optimize by precomputing factorials, and
                                                          if (vals.count(cur)) {
    fact[n]/fact[n-r]
                                                           int ans = n * p - vals[cur];
 ull nPr(ull n, ull r) {
                                                           return ans:
 for (ull i = n-r+1; i <= n; i++)
 .v *= i;
return v:
                                                         return -1;
ull nPr(ull n, ull r, ull m) {
                                                        Euler Phi / Totient
 ull v = 1;
for (ull i = n-r+1: i <= n: i++)
                                                        int phi(int n) {
  v = (v * i) \% m;
                                                         int r = n;
                                                        for (int i = 2; i * i <= n; i++) {
    if (n % i == 0) r -= r / i;
    while (n % i == 0) n /= i;
 return v;
ull nCr(ull n, ull r) {
 long double \dot{v} = 1;
 for (ull i = 1; i <= r; i++)
v = v * (n-r+i) /i;
                                                         if (n > 1) r = r / n;
                                                         return r;
 return (ull)(v + 0.001);
                                                        #define n 100000
                                                        ll phi[n+1];
 // requires modulo math
                                                        void computeTotient() {
 // can optimize by precomputing mfac and
                                                         for (int i=1; i<=n; i++) phi[i] = i;
 \hookrightarrow minv-mfac
                                                         for (int p=2; p<=n; p++) {
ull nCr(ull n, ull r, ull m) {
                                                         if (phi[p] == p) {
 return mfac(n, m) * minv(mfac(k, m), m) % m
 \rightarrow minv(mfac(n-k, m), m) % m:
                                                           phi[p] = p-1;
                                                           for (int i = 2*p; i<=n; i += p) phi[i] =
                                                            (phi[i]/p) * (p-1);
Multinomials
ll multinomial(vector<int>& v) {
    ll c = 1, m = v.empty() ? 1 : v[0];
    for(int i = 1; i < v.size(); i++)
        for (int j = 0; j < v[i]; j++)
        c = c * ++m / (j+1);
    }
</pre>
                                                        Factorials
                                                        // digits in factorial
                                                        #define kamenetsky(n) (floor((n * log10(n /
 return c;
                                                         \rightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
                                                        // approximation of factorial #define stirling(n) ((n == 1) ? 1 : sqrt(2 *
 Chinese Remainder Theorem
bool ecrt(l1* r, l1* m, int n, l1% re, l1% mo)
                                                        \hookrightarrow M PI * n) * pow(n / M_E, n))
// natural log of factorial
                                                        #define lfactorial(n) (lgamma(n+1))
  d = egcd(mo, m[i], x, y);
  if ((r[i] - re) % d != 0) return false;

x = (r[i] - re) / d * x % (m[i] / d);

re += x * mo;
                                                        Prime Factorization
                                                        // do not call directly
                                                        ll pollard_rho(ll n, ll s) {
  mo = mo / d * m[i];
                                                        11 x, y;
```

```
x = y = rand() \% (n - 1) + 1;
 int head = 1, tail = 2;
 while (true) {
 x = mult(x, x, n);

x = (x + s) \% n;
 if (x == y) return n;
 11 d = _{gcd(max(x - y, y - x), n);}
 if (1 < d && d < n) return d;
if (++head == tail) y = x, tail <<= 1;
// call for prime factors
void factorize(ll n, vector<ll> &divisor) {
 if (n == 1) return;
 if (isPrime(n)) divisor.push_back(n);
  while (d >= n) d = pollard_rho(n, rand() % (n Fast Fourier Transform
 \rightarrow -1) +1);
 factorize(n / d, divisor);
 factorize(d, divisor);
Factorize Factorials
// NOTE: count distinct divisors of n by
// computing (q1+1)*(q2+1)*...*(qk+1)
// where gi are powers of primes pi dividing n
// use that and this code to solve
→ https://open.kattis.com/problems/divisors
// max power of a prime p dividing n!
// O(log(n))
int legendre(int n, int p) {
 int mx = 0;
while(n>0) n/=p, mx+=n;
 return mx;
bitset<10000> sieve;
vector<int> primes;
// get all primes O(n log n)
// if dealing with small numbers
void genPrimes(int n) {
 sieve[0] = sieve[1] = 1;
 primes.push_back(2);
 for (int i = 3; i <= n; i+=2)
if (i%2 != 0 && !sieve[i]) {
   primes.push_back(i);
   for (int j = i * 3; j \le n; j += i*2)
    sieve[j] = 1;
// make sure you call genPrimes first
// return vector of prime factor powers as
\rightarrow vector v of size pi(n)
// so that v[i] = power of primes[i] dividing
\stackrel{\longrightarrow}{//} \stackrel{n!}{0} (pi(n) * log(n)) where pi(n) is prime
// so basically O(n) since pi(n) = O(n/\log(n))
vector<int> factorize factorial(int n) {
 vector<int> factorization(primes.size(), 0);
 for(int i=0:i<primes.size() && primes[i] <=
 → n:i++) {
 factorization[i] = legendre(n, primes[i]);
 return factorization:
// same thing but for C(n,k)
vector<int> factorize_binom(int n, int k) {
 vector<int> factorization(primes.size(), 0);
 for(int i=0;i<primes.size() && primes[i] <=</pre>
 \rightarrow n;i++) {
  factorization[i] = legendre(n, primes[i])
    legendre(k, primes[i]) - legendre(n-k,
    primes[i]):
return factorization;
Farev Fractions
```

```
// farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
 // length is sum of phi(i) for i = 1 to n
vector<pair<int, int>> farev(int n) {
 int h = 0, k = 1, x = 1, y = 0, r;
vector<pair<int, int>> v;
  v.push back({h, k}):
  r = (n-y)/k;
 y += r*k; x += r*h;

swap(x,h); swap(y,k);

x = -x; y = -y;

} while (k > 1);
 v.push_back({1, 1});
 return v;
const double PI = acos(-1):
void fft(vector<cd>& a, bool invert) {
 int n = a.size();
 for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1;
  for (; j & bit; bit >>= 1) j ^= bit;
  .j ^= bit;
  if (i < j) swap(a[i], a[j]);
 for (int len = 2; len <= n; len <<= 1) {
   double ang = 2 * PI / len * (invert ? -1 :
  cd wlen(cos(ang), sin(ang));
  for (int i = 0; i < n; i += len) {
   cd w(1);
   for (int j = 0; j < len / 2; j++) {
    cd u = a[i+j], v = a[i+j+len/2] * w;
    a[i+i] = u + v:
    a[i+j+len/2] = u - v;
    w *= wlen:
 if (invert)
  for (auto& x : a)
vector<int> fftmult(vector<int> const& a,

    vector<int> const& b) {

 vector<cd> fa(a.begin(), a.end()),

    fb(b.begin(), b.end());

 int n = 1 << (32 - __builtin_clz(a.size() +
 \rightarrow b.size() - 1));
fa.resize(n); fb.resize(n);
 fft(fa, false); fft(fb, false)
 for (int i = 0; i < n; i++) fa[i] *= fb[i]:
 fft(fa, true);
 vector<int> toret(n):
for (int i = 0; i < n; i++) toret[i] =

→ round(fa[i].real());
 return toret;
Greatest Common Denominator
ll egcd(ll a, ll b, ll& x, ll& y) {
   if (b == 0) { x = 1; y = 0; return a; }
   ll gcd = egcd(b, a % b, x, y);
 x = a / b * y;
 swap(x, y);
 return gcd;
Josephus Problem
 // 0-indexed, arbitrary k
int josephus(int n, int k) {
if (n == 1) return 0;
if (k == 1) return n-1;
if (k > n) return (josephus(n-1,k)+k)%n;
 int res = josephus(n-n/k,k)-n\%k;
 return res + ((res<0)?n:res/(k-1)):
```

generate  $0 \le a/b \le 1$  ordered,  $b \le n$ 

```
// fast case if k=2, traditional josephus
                                                           return tetraloop(a,b,m) % m;
int josephus(int n) {
return 2*(n-(1<<(32-builtin clz(n)-1)));
                                                           Matrix
                                                          template<typename T>
Least Common Multiple
                                                          struct Mat : public Vec<2, T> {
                                                           int w. h:
#define lcm(a,b) ((a*b)/__gcd(a,b))
                                                           Mat(int x, int y) : Vec<2, T>(x, y), w(x),
Modulo Operations
                                                           \rightarrow h(v) {}
#define MOD 1000000007
                                                           static Mat<T> identity(int n) { Mat<T> m(n,n)
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0))
                                                               for (int i=0;i<n;i++) m[i][i] = 1; return
#define mult(a,b,m) ((ull)a*b%m)
                                                           #define msub(a, b, m) (a-b+((a < b)?m:0)) 11 mpow(11 b, 11 e, 11 m) {
                                                            for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
(*this)[i][j] += m[i][j];
 11 x = 1:
 while (e > 0) {
    if (e % 2) x = (x * b) % m;
    b = (b * b) % m;
                                                            return *this;
  e /= 2;
                                                           Mat<T>& operator-=(const Mat<T>& m) {
                                                            for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
(*this)[i][j] -= m[i][j];
 return x % m:
ull mfac(ull n, ull m) {
  ull f = 1;
                                                             return *this;
for (int i = n; i > 1; i--)

f = (f * i) % m;

return f;
                                                           Mat<T> operator*(const Mat<T>& m) {
                                                             Mat < T > z(w,m.h);
                                                            for (int i = 0; i < w; i++)

for (int j = 0; j < h; j++)

for (int k = 0; k < m.h; k++)

...z[i][k] += (*this)[i][j] * m[j][k];
// if m is not guaranteed to be prime
ll minv(ll b, ll m) {
ll x = 0, y = 0;
if (egcd(b, m, x, y) != 1) return -1;
                                                               return z:
 return (x % m + m) % m:
                                                           Mat<T> operator+(const Mat<T>& m) { Mat<T>
                                                              a=*this; return a+=m; }
11 mdiv_compmod(int a, int b, int m) {
  if (_gcd(b, m) != 1) return -1;
                                                           Mat<T> operator-(const Mat<T>& m) { Mat<T>

→ a=*this; return a-=m; }
Mat<T>& operator*=(const Mat<T>& m) { return
 return mult(a, minv(b, m), m);
                                                           → *this = (*this)*m; }
\frac{1}{1} if m is prime (like 10^{9}+7)
11 mdiv_primemod (int a, int b, int m) {
                                                           Mat<T> power(int n) {
                                                            Mat<T> a = Mat<T>::identity(w),m=*this;
 return mult(a, mpow(b, m-2, m), m);
                                                             for (;n;n/=2,m*=m) if (n\&1) a *= m;
   tonelli \ shanks = sqrt(n) \ % \ m, \ m \ is \ prime
11 legendre(ll a, ll m){
 if (a % m==0) return 0;
if (m == 2) return 1;
return mpow(a,(m-1)/2,m);
                                                          Matrix Exponentiation
                                                             F(n) = c \lceil 0 \rceil * F(n-1) + c \lceil 1 \rceil * F(n-2) + \dots
                                                          // b is the base cases of same length c
ll matrix_exponentiation(ll n, vector<ll> c,
11 msqrt(ll n, ll m) {
ll s = \_builtin\_ctzll(m-1), q = (m-111)>>s,

  vector<ll> b) {
  if (nth < b.size()) return b[nth-1];
}</pre>
\rightarrow z = rand()%(m-1)+1;
 if (m == 2) return 1;
if (s == 1) return mpow(n,(m+1)/411,m);
                                                           Mat<11> a(c.size(), c.size()); ll s = 0;
for (int i = 0; i < c.size(); i++) a[i][0] =
 while (legendre(z,m)!=m-1) z = rand()\%(m-1)+1;
                                                           c[i];
for (int i = 0; i < c.size() - 1; i++)</pre>
11 c = mpow(z,q,m), r = mpow(n,(q+1)/2,m), t
\rightarrow = mpow(n,q,m), M = s;
                                                           \rightarrow a[i][i+1] = 1;
a = a.power(nth - c.size());
for (int i = 0; i < c.size(); i++)
                                                           s += a[i][0] * b[i];
return s;
  11 b = c;
 for (int'j = 0; j < M-i-1; j++) b = (b * b) %
 r = r * b % m; c = b * b % m; t = t * c % m;
                                                          Matrix Subarray Sums
\hookrightarrow M = i:
                                                           template<class T> struct MatrixSum {
 return r;
                                                           MatrixSum(Vec<2, T>& v) {
                                                            p = Vec<2,T>(v.size()+1, v[0].size()+1);
Modulo Tetration
                                                             for (int i = 0; i < v.size(); i++)
11 tetraloop(ll a, ll b, ll m) {
                                                             for (int j = 0; j < v[0].size(); j++)
p[i+1][j+1] = v[i][j] + p[i][j+1] +
 if(b == 0 | a == 1) return 1;
 11 w = tetraloop(a,b-1,phi(m)), r = 1;
                                                              p[i+1][i] - p[i][i]:
 for (;w;w/=2) {
 if (w&1)
                                                           T sum(int u, int l, int d, int r) {
    return p[d][r] - p[d][l] - p[u][r] + p[u][l];
  r *= a: if (r >= m) r -= (r/m-1)*m:
  a *= a; if (a >= m) a -= (a/m-1)*m;
 return r:
                                                          Mobius Function
                                                                                                                       ull mod = mpow(a, temp, n);
int tetration(int a, int b, int m) {
  if (a == 0 || m == 1) return ((b+1)&1)%m;
                                                          const int MAXN = 10000000;
                                                          // mu[n] = 0 iff n has no square factors
                                                                                                                       while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
```

```
|// 1 = even number prime factors, -1 = odd
short mu[MAXN] = \{0,1\};
void mobius(){
  for (int i = 1; i < MAXN; i++)
    if (mu[i])</pre>
   for (int' j = i + i; j < MAXN; j += i)
     mu[j] -= mu[i];
Nimber Arithmetic
 #define nimAdd(a,b) ((a)^(b))
ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
    M[0]^{(M[0] <<16)}, M[1]^{(M[1] <<8)},
     M[2]^{(M[2] << 4)}, M[3]^{(M[3] << 2)},
    M[4]^(M[4]<<1);
  if (i-- == 0) return a&b;
int k=1<<i:
  ull s=nimMul(a,b,i), m=M[5-i],
     t=nimMul(((a^(a>>k))&m)|(s\&~m),
     ((b^{(b)}) \& m) | (m \& (\sim m >> 1)) << k, i):
  return ((s^t)&m)<<k|((s^(t>>k))&m);
Permutation
//c = array \ size, \ n = nth \ perm, \ return \ index
vector<int> gen permutation(int c, int n) {
 vector<int> idx(c), per(c), fac(c); int i;
 for (i = 0; i < c; i++) idx[i] = i;
 for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i;
 for (i = c - 1; i >= 0; i-
per[c-i-1] = idx[fac[i]],
  idx.erase(idx.begin() + fac[i]);
 return per;
 // get what nth permutation of vector
int get permutation(vector<int>& v) {
 int use = 0, i = 1, r = 0;
 for (int e : v) {
   r = r * i++ + __builtin_popcount(use &
 → -(1<<e));</pre>
  use |= 1 << e:
 return r;
Permutation (string/multiset)
string freq2str(vector<int>& v) {
 string s;
 for (int i = 0: i < v.size(): i++)
  for (int j = 0; j < v[i]; j++)
 s += (char)(i + 'A');
return s:
// nth perm of multiset, n is 0-indexed
string gen_permutation(string s, ll n) {
 vector<int> freq(26, 0);
 for (auto e : s) freq[e - 'A']++;
 for (int i = 0; i < 26; i++) if (freq[i] > 0)
  freq[i]--; ll v = multinomial(freq);
  if (n < v) return (char)(i+'A') +
 → gen_permutation(freq2str(freq), n);
  freq[i]++; n -= v;
 return "":
Miller-Rabin Primality Test
// Miller-Rabin primality test - O(10 log^3 n)
bool isPrime(ull n) {
   if (n < 2) return false;
   if (n == 2) return true;
if (n % 2 == 0) return false;
ull s = n - 1;
while (s % 2 == 0) s /= 2;
 for (int i = 0; i < 10; i++) {
  ull temp = s;
ull a = rand() % (n - 1) + 1;
```

```
mod = mult(mod, mod, n);
    temp *= 2:
  if (mod!=n-1&&temp%2==0) return false:
 return true:
Sieve of Eratosthenes
bitset<100000001> sieve;
// generate sieve - O(n log n)
void genSieve(int n) {
    sieve[0] = sieve[1] = 1;
    for (ull i = 3; i * i < n; i += 2)
        if (!sieve[i])</pre>
   for (ull j = i * 3; j <= n; j += i * 2)
...sieve[j] = 1;
 .
// query sieve after it's generated - O(1)
bool querySieve(int n) {
 return n == 2 || (n % 2 != 0 && !sieve[n]);
 Compile-time Prime Sieve
const int MAXN = 100000;
template<int N>
struct Sieve {
  bool sieve[N];
  constexpr Sieve() : sieve() {
  sieve[0] = sieve[1] = 1;
  for (int i = 2; i * i < N; i++)
if (!sieve[i])
     for (int j = i * 2; j < N; j += i)
...sieve[j] = 1;
bool isPrime(int n) {
   static constexpr Sieve<MAXN> s;
 return !s.sieve[n]:
Simpson's / Approximate Integrals
    integrate f from a to b, k iterations
// threefacts of from a to , k to the factor of k ((b-a)/2k) ^2/2 ((b-a)/2k) ^2/2 ((b-a)/2k) ^2/2 ((b-a)/2k) for (a,b)/2 ((b-a)/2k) for (a,b)/2 ((b-a)/2k) is a function "double func (double (a,b)/2).
double Simpsons (double a. double b. int k.
 double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;
for (int i = 0; i < k; i++)
. t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
 \leftrightarrow (*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
 Common Equations Solvers
 // ax^2 + bx + c = 0, find x
vector<double> solveEq(double a, double b,
 double c) {
.vector<double> r;
.double z = b * b - 4 * a * c;
 if (z == 0)
  r.push_back(-b/(2*a));
 else if (z > 0) {
  r.push back((sgrt(z)-b)/(2*a)):
  r.push_back((sqrt(z)+b)/(2*a));
 return r;
 \frac{1}{2} / ax^3 + bx^2 + cx + d = 0, find x
vector < double > solveEq (double a, double b,

    double c, double d) {
    vector < double > res;
}

 long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
 \rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z \le 0) {
  theta = a\cos(r/sqrt(q*q*q));
```

```
res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push back(sq*cos((theta+2.0*PI)/3.0) -
    a1/3.0):
 res.push_back(sq*cos((theta+4.0*PI)/3.0)
\rightarrow a1/3.0);
 élse {
  res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
 res[0] = (res[0] + q / res[0]) *
\leftrightarrow ((r<0)?1:-1) - a1 / 3.0:
 return res:
// linear diophantine equation ax + by = c,
    find x and y
// infinite solutions of form x+k*b/g, y-k*a/g bool solveEq(ll a, ll b, ll c, ll &x, ll &y, ll
 g = \overline{egcd(abs(a), abs(b), x, y)};
 if (c % g) return false;
x *= c / g * ((a < 0) ? -1 : 1);

y *= c / g * ((b < 0) ? -1 : 1);

return true;
^{\prime\prime}/ m = # equations, n = # variables, a[m][n+1]
\rightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
\rightarrow a[i][n+1]
// find a solution of some kind to linear
\rightarrow equation
const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a >
→ -eps); }
vector <double > solveEq(double **a, int m, int
\underset{int}{\hookrightarrow} n) { = 0;
 for (int i = 0; i < n; i++) {
  for (int j = cur; j < m; j++) {
  if (!zero(a[j][i])) {
    if (j != cur) swap(a[j], a[cur]);
     for (int sat = 0; sat < m; sat++) {
  if (sat == cur) continue;</pre>
  cuouple num = a[sat][i] / a[cur][i]
for (int sot = 0; sot <= n; sot++)
          a[sat][sot] -= a[cur][sot] * num;
}
cur++;</pre>
       double num = a[sat][i] / a[cur][i];
    break;
 for (int j = cur; j < m; j++)
 if (!zero(a[j][n])) return vector<double>();
 vector<double > ans(n,0);
 for (int i = 0, sat = 0; i < n; i++)
if (sat < m && !zero(a[sat][i]))
 ans[i] = a[sat][n] / a[sat++][i];
return ans;
}
// solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 is no
⇒ solutions

// `alls` is whether to find all solutions, or
\hookrightarrow any
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&

→ b, Vec<1, double>& x, bool alls=false) {
 .<u>int</u> n = A.size(), m = x.size(), rank = 0, br,
vector<int> col(m); iota(begin(col), end(col);
 for(int i = 0; i < n; i++) {
  double v, bv = 0;</pre>
  for(int r = i; r < n; r++)
  for(int c = i; c < n; c++)
    if ((v = fabs(A[r][c])) > bv)
        br = r, bc = c, bv = v;
    if (bv <= eps) {
for(int j = i; j < n; j++)
...if (fabs(b[j]) > eps)
```

```
swap(b[i], b[br]);
swap(col[i], col[bc]);
  for(int j = 0; j < n; j++)
  swap(A[j][i], A[j][bc]);
bv = 1.0 / A[i][i];</pre>
  for(int j = (alls)?0:i+1; j < n; j++) {
   if (j != i) {
    double fac = A[j][i] * bv;
    b[j] -= fac * b[i];
    for(int k = i+1; k < m; k++)
A[j][k] -= fac*A[i][k];</pre>
  rank++:
 if (alls) for (int i = 0; i < m; i++) x[i] =
 for (int i = rank; i--;) {
  bool isGood = true;
  if (alls)
   for (int j = rank; isGood && j < m; j++)
...if (fabs(A[i][j]) > eps)
      isGood = false;
  b[i] /= A[i][i];
  if (isGood) x[col[i]] = b[i];
if (!alls)
  for(int j = 0; j < i; j++)
b[j] -= A[j][i] * b[i];
 return rank;
Gravcode Conversions
ull graycode2ull(ull n) {
  ull i = 0;
 for (; n; n = n >> 1) i ^= n;
ull ull2graycode(ull n) {
return n ^ (n >> 1):
Date Utilities
// handles -4799-01-01 to 1465001-12-31
int date2int(int y, int m, int d){
 return 1461*(y+4800+(m-14)/12)/4+367*(m-2-(m
    -14)/12*12)/12-3*((y+4900+(m-14)/12)/100)
   /4+d-32075;
pair<int,pair<int,int>> int2date(int x){
 int n,i,j;
 n=4*x/146097;
 x=(146097*n+3)/4
 i=(4000*(x+1))/1461001:
 x = 1461 * i/4 - 31;
 j=80*x/2447;
 return \{100*(n-49)+i+j/11, \{j+2-12*(j/11), \}
    x-2447*i/80};
int dayOfWeek(int y, int m, int d){ //0=sunday
  static int cal[]={0,3,2,5,0,3,5,1,4,6,2,4};
 y-=m<3;
return (y+y/4-y/100+y/400+cal[m-1]+d)\%7;
Unix/Epoch Time
   O-indexed month/time, 1-indexed day
/// minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
```

→ hour, int minute, int second) {

struct tm t; time\_t epoch;
t.tm\_year = year - 1900; t.tm\_mon = month;
t.tm\_mday = day; t.tm\_hour = hour;

t.tm\_min = minute; t.tm\_sec = second; t.tm\_isdst = 0; // 1 = daylights savings

return -1:

swap(A[i], A[br]);

break:

```
epoch = mktime(&t);
 return (ull)epoch;
vector<int> toDate(ull epoch) {
 time t e=epoch; struct tm t=*localtime(&e);
 return {t.tm_year+1900,t.tm_mon,t.tm_mday,t_
    .tm hour.t.tm min.t.tm sec}:
int getWeekday(ull epoch) {
 time_t e=epoch; struct tm t=*localtime(&e);
 return t.tm wday: // 0-6. 0 = sunday
int getDayofYear(ull epoch) {
 time_t e=epoch; struct tm t=*localtime(&e);
 return t.tm_yday; // 0-365
const int months[] =
→ {31,28,31,30,31,30,31,30,31,30,31};
bool validDate(int year, int month, int day) {
    bool leap = !(year%(year%25?4:16));
    if (month >= 12) return false;
    return day <= months[month] + (leap &&
    month == 1):
Theorems and Formulae
```

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1)) $\frac{1)+!(n-2)), !n=n! \sum_{i=0}^{n} \frac{(-1)^{i}}{i!}, !n=\left[\frac{n!}{e}\right]}{\text{In a partially ordered set, a chain is a subset of}}$ 

elements that are all comparable to eachother. An antichain is a subset where no two are comparable.

imal antichain equals the size of a minimal of world add\_edge(Vec<2, edge> &graph, int u, int w, bool directed=true) { chain cover of a partially ordered set S. The graph[u].push\_back({u,v,w}); width of S is the maximum size of an antichain if (!directed) graph[v].push\_back({v,u,w}); in S, which is equal to the minimum number  $|_{\text{vector} < \text{int}}^{\text{f}} \rangle$  dijkstra(Vec<2, edge> &graph, int of chains needed to cover S, or the minimum of chains such that all elements are in vector(int) D(graph.size(), inf); at least one chain.

Rosser's Theorem states the nth prime pq.push({src,src,0}); number is greater than n \* ln(n) for n > 1.

Nicomachi's Theorem states  $1^3 + 2^3 + ... +$  $n^3 = (1 + 2 + ... + n)^2$  and is equivalent to  $(n^{\frac{n+1}{2}})^2$ .

Lagrange's Four Square Theorem states every natural number is the sum of the squares of four non-negative integers. This is a special case of the Fermat Polygonal Number **Theorem** where every positive integer is a sum of at most n s-gonal numbers. The nth Eulerian Path s-gonal number  $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$ 

```
7 Graphs
struct edge {
int u,v,w;
 edge (int u,int v,int w) : u(u),v(v),w(w) {}
 edge (): u(0), v(0), w(0) {}
bool operator < (const edge &e1, const edge
```

 $\rightarrow$  &e2) { return e1.w < e2.w; }

```
|bool operator > (const edge &e1, const edge
struct subset {
int p, rank, sz;
 subset(int p) : p(p), rank(0), sz(1) {}
subset() : p(0), rank(0), sz(0) {}
 void make set(int p) { p= p, rank=0, sz=1; }
\mathbf{BFS}
// adjacency list named 'graph'
int visited[MAX]:
int parent[MAX];
int vc = 0;
vector<int> bfs(int start, int end) {
 visited[start] = vc:
 parent[start] = -1;
 queue<int> q
 q.push(start);
 while (!q.empty()) {
 int v = q.front(); q.pop();
for (auto e : graph[v]) {
   if (visited[e] != vc) {
    visited[e] = vc;
    q.push(e);
    parent[e] = v;
    if (e == end) goto DONE;
 // path reconstruction
 if (visited[end] != vc) return {};
vector<int> path;
 for (int v = end: v = -1: v = parent[v])
 path.push back(v);
 return path;
Diikstra's
```

## const int inf = 20000001; // change as needed Dilworth's theorem states the size of a max // use add\_edge(..., true) for digraphs priority\_queue < edge, vector < edge >, → greater<edge>> pq; D[src]=0; while(!pq.empty()) { edge e = pq.top(); pq.pop(); int v = e.v; for(int i=0;i<graph[v].size();i++) {</pre> int u = graph[v][i].v; if(D[v] + graph[v][i].w < D[u]) { D[u] = D[v] + graph[v][i].w; pq.push({src,u,D[u]}); return D:

```
#define edge_list vector<edge>
#define adj sets vector<set<int>>
struct EulerPathGraph {
adj_sets graph; // actually indexes incident
edge_list edges; int n; vector<int> indeg;
EulerPathGraph(int n): n(n) {
graph = *(new adj_sets(n, set<int>()));
}
 indeg = *(new vector<int>(n,0));
```

```
void add_edge(int u, int v) {
                                                       |vector<edge> mst(Vec<2, edge> graph, 11
                                                        graph[u].insert(edges.size());
  indeg[v]++;
  edges.push back(edge(u,v,0));

    greater<edge>> pq;
vector<edge> MST;

 bool eulerian_path(vector<int> &circuit) {
                                                        bitset<20001> marked; // change size as needed
  if(edges.size()==0) return false;
                                                        marked[0] = 1:
  stack<int> st;
int a[] = {-1, -1};
                                                        for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
  for(int v=0; v<n; v++)
  if(indeg[v]!=graph[v].size()) {
    bool b = indeg[v] > graph[v].size();
                                                        → pq.size()!=0) {
                                                         edge e = pq.top(); pq.pop();
                                                          int u = e.u, v = e.w;
if(marked[u] && marked[v]) continue;
    if (abs(((int)indeg[v])-((int)graph[v])
    .size())) > 1) return
                                                         else if(marked[u]) swap(u, v);
for(edge_ep : graph[u]) pq.push(ep);
    false;
if (a[b] != -1) return false;
                                                          marked[u] = 1:
   a[b] = v;
                                                          MST.push_back(e);
                                                          total += e.w;
  int s = (a[0]!=-1 \&\& a[1]!=-1 ? a[0] :
                                                        return MST:
\rightarrow (a[0]==-1 && a[1]==-1 ? edges[0].u : -1));
  if(s==-1) return false;
  while(!st.empty() || !graph[s].empty()) {
                                                        Union Find
  if (graph[s].empty()) {
                                                       int uf find(subset* s, int i) {
   if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
   struct line { point a, b; };
   struct circle { point c; double r; };
    circuit.push back(s): s = st.top():
                                                        return s[i].p;
   st.pop(); }
   else {
                                                       void uf_union(subset* s, int x, int y) {
    int w = edges[*graph[s].begin()].v;
                                                        int xp = uf_find(s, x), yp = uf_find(s, y);
    graph[s].erase(graph[s].begin());
                                                        if (s[xp].rank > s[yp].rank) s[yp].p = xp,
    st.push(s); s = w:
                                                           s[xp].sz += s[yp].sz;
                                                        else if (s[xp].rank < s[yp].rank) s[xp].p =
  circuit.push back(s):
                                                           vp. s[vp].sz += s[xp].sz:
  return circuit.size()-1==edges.size();
                                                        else s[yp].\bar{p} = xp, s[xp].rank++, s[xp].sz +=
                                                           s[yp].sz;
                                                       void uf_size(subset *s, int i) {
Floyd Warshall
                                                        return s[uf_find(s, i)].sz;
const ll inf = 1LL << 62;
#define FOR(i,n) for (int i = 0; i < n; i++)
                                                       Bipartite Graph
void floydWarshall(Vec<2, 11>& m) {
int n = m.size();
FOR(i,n) m[i][i] = min(m[i][i], OLL);
FOR(k,n) FOR(i,n) FOR(j,n) if (m[i][k] != inf
                                                        A bipartite graph has "left" and "right" set of
                                                       \rightarrow nodes
Every edge has an endpoint in each set (L/R)
A matching is a subset of all edges
  auto newDist = max(m[i][k] + m[k][j], -inf);
                                                       Such that each vertex is an endpoint
  m[i][j] = min(m[i][j], newDist);
                                                       Of at most one edge in the subset
                                                       sgrt(V)*E time
 FOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n)
                                                        tested on "piano lessons"
 if (m[i][k] != inf && m[k][j] != inf)
                                                       sourced from
\hookrightarrow m[i][j] = -inf;
                                                            https://codeforces.com/blog/entry/58048
                                                        #define MAXNODES 1001
Bellman Ford
                                                       bitset<MAXNODES> V:
const int inf = 20000001;
vector<1l> bellman_ford(vector<edge> edges, int
                                                       bool match(int node, Vec<2,int> &G, vector<int>

    &R, vector<int> &L) {
    if (V[node]) return false;
}

    src, int V) {
    vector<ll> D(V,inf);
                                                        V[node] = 1:
D[src] = 0;
for (int i=1;i<=V-1;i++)
                                                        for(auto vec : G[node]) {
   if (R[vec] == -1 || match(R[vec], G, R, L))
 for (edge e : edges)
 if (D[e.u] != inf && D[e.u] + e.w < D[e.v])
...D[e.v] = D[e.u] + e.w;
                                                          L[node] = vec; R[vec] = node; return true;
 // detect negative cycles: *typically* 2 is as
\hookrightarrow good as V-1 for this
                                                        return false;
 for (int i=1:i<=V-1:i++)
 for (edge e : edges)
  -of cage : eages/
if (D[e.u] != inf && D[e.u] + e.w < D[e.v])
.D[e.v] = -inf;
                                                        vector<pair<int, int>> bipartite match(Vec<2,
                                                        int> &G, int m) {
vector<int> L(G.size(), -1), R(m, -1);
return D;
                                                        V.reset();
                                                        bool running = true;
                                                        while (running) {
Minimum Spanning Tree
                                                         running = false;
// returns vector of edges in the mst
                                                         V.reset();
for (int i=0;i<L.size();i++)
if (L[i] == -1)
// graph[i] = vector of edges incident to
    vertex i
\nearrow places total weight of the mst in rak{S}total
                                                            running |= match(i, G, R, L);
// if returned vector has size != n-1, there is
```

```
for (int i = 0; i < L.size(); ++i)
  if(L[i]!=-1) ret.push_back({i, L[i]});</pre>
 return ret;
2D Grid Shortcut
#define inbound(x,n) (0<=xx < n)
|#define fordir(x,y,n,m) for(auto[dx,dy]:dir)i
    (inbound(x+dx,n)&&inbound(y+dy,m))
const pair<int,int> dir[] =
\hookrightarrow {{1,0},{0,1},{-1,0},{0,-1}};
     2D Geometry
#define point complex<double>
#define EPS 0.0000001
#define sq(a) ((a)*(a))
#define cb(a) ((a)*(a)*(a))
double dot(point a, point b) { return

    real(coni(a)*b): }

double cross (point a, point b) { return

    imag(conj(a)*b); }

struct line { point a, b; };
struct segment { point a, b; };
struct triangle { point a, b, c; };
struct rectangle { point tl, br; };
struct convex_polygon {
  vector<point> points;
 convex_polygon(vector<point> points) :
 → points(points) {}
 convex_polygon(triangle a) {
 points.push_back(a.a); points.push_back(a.b);
 → points.push_back(a.c);
 convex polygon(rectangle a) {
 points.push_back(a.tl);
    points.push back({real(a.tl),
    imag(a.br)}):
  points.push_back(a.br);
    points.push back({real(a.br),
    imag(a.tl)});
struct polygon {
 vector<point> points;
 polygon(vector point points) :
 → points(points) {}
 polygon(triangle a) {
 points.push_back(a.a); points.push back(a.b);
    points.push_back(a.c);
 polygon(rectangle a) {
  points.push_back(a.tl);
    points.push_back({real(a.tl),
    imag(a.br)});
  points.push_back(a.br);
    points.push_back({real(a.br),
    imag(a.tl)}):
 polygon(convex_polygon a) {
  for (point v : a.points)
  points.push_back(v);
   triangle methods
double area_heron(double a, double b, double
 \rightarrow c) {
if (a < b) swap(a, b);
 if (a < c) swap(a, c);
 if (b < c) swap(b, c);
 if (a > b + c) return -1:
return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
    /16.0);
```

vector<pair<int.int>> ret:

```
// segment methods
double lengthsq(segment a) { return
    sq(real(a.a) - real(a.b)) + sq(imag(a.a) -
    imag(a.b)): }
double length(segment a) { return

    sqrt(lengthsq(a)); }

   circle methods
double circumference(circle a) { return 2 * a.r

→ * M PI; }

double area(circle a) { return sq(a.r) * M_PI;
\overset{	o}{\rightarrow} } // rectangle methods
double width(rectangle a) { return

→ abs(real(a.br) - real(a.tl)); }

double height (rectangle a) { return

→ abs(imag(a.br) - real(a.tl)); }

double diagonal (rectangle a) { return

    sqrt(sq(width(a)) + sq(height(a))); }

double area (rectangle a) { return width(a) >
→ height(a); }
double perimeter(rectangle a) { return 2 *
    (width(a) + height(a)); }
   check if `a` fit's inside `b
// swap equalities to exclude tight fits
bool doesfitInside(rectangle a, rectangle b) {
  int x = width(a), w = width(b), y = height(a),
\rightarrow h = height(b):
if (x > y) swap(x, y);
if (w > h) swap(w, h);
 if (w < x) return false;
 if (v <= h) return true:
double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
return sq(a) <= sq(b) + sq(c);</pre>
// polygon methoas
|// negative area = CCW, positive = CW
   polygon methods
double area(polygon a) {
 double area = 0.0; int n = a.points.size();
for (int i = 0, j = 1; i < n; i++, j = (j +
\rightarrow 1) % n)
 area += (real(a.points[j]-a.points[i]))*
return area / 2.0:
// get both unsigned area and centroid
pair<double, point> area_centroid(polygon a) {
 int n = a.points.size():
 double area = 0:
 point c(0, 0);
for (int i = n - 1, j = 0; j < n; i = j++) {
    double v = cross(a.points[i], a.points[j]) /
 c += (a.points[i] + a.points[j]) * (v / 3);
 c /= area;
return {area, c};
Intersection
// -1 coincide, 0 parallel, 1 intersection
int intersection(line a, line b, point& p) {
if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
 p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
\rightarrow -a.a, b.b - b.a) * (b - a) + a;
 if (abs(cross(a.b - a.a. a.b - b.a)) > EPS)

→ return 0:

return -1;
// area of intersection
double intersection(circle a, circle b) {
double d = abs(a.c - b.c);
if (d <= b.r - a.r) return area(a);</pre>
 if (d <= a.r - b.r) return area(b);
 if (d \ge a.r + b.r) return 0;
```

```
double alpha = acos((sq(a.r) + sq(d) -
 \rightarrow sq(b.r)) / (2 * a.r * d));
 double beta = acos((sq(b.r) + sq(d) - sq(a.r))
 \rightarrow / (2 * b.r * d));
 return sq(a.r) * (alpha - 0.5 * sin(2 *
        alpha)) + sq(b.r) * (beta - 0.5 * sin(2 *
}
// -1 outside, 0 inside, 1 tangent, 2
int intersection (circle a, circle b,

    vector<point>& inter) {

  double d2 = norm(b.c - a.c), rS = a.r + b.r,
 \rightarrow rD = a.r - b.r;
if (d2 > sq(rS)) return -1;
  if (d2 < sq(rD)) return 0;
  double ca = 0.5 * (1 + rS * rD / d2):
 point z = point(ca, sqrt(sq(a.r) / d2 -
  inter.push back(a.c + (b.c - a.c) * z);
  if (abs(imag(z)) > EPS) inter.push_back(a.c +
 \rightarrow (b.c - a.c) * conj(z));
 return inter.size();
 // points of intersection
vector<point> intersection(line a, circle c) {
 vector<point> intersection of the section of the se
  point m = \dot{a}.b * real(c.c / a.b);
  double d2 = norm(m - c.c);
  if (d2 > sq(c.r)) return 0;
  double 1 = sqrt((sq(c.r) - d2) / norm(a.b));
  inter.push_back(a.a + m + 1 * a.b);
  if (abs(1) > EPS) inter.push_back(a.a + m - 1
 \rightarrow * a.b);
 return inter;
 // area of intersection
double intersection(rectangle a, rectangle b) {
 double x1 = max(real(a.tl), real(b.tl)), y1 =

→ max(imag(a.tl), imag(b.tl));
 double x2 = min(real(a.br), real(b.br)), y2 =

→ min(imag(a.br), imag(b.br));
 return (x2 <= x1 || y2 <= y1) ? 0 :
       (x2-x1)*(y2-y1);
 Convex Hull
bool cmp(point a, point b) {
 if (abs(real(a) - real(b)) > EPS) return

→ real(a) < real(b);
</pre>
 if (abs(imag(a) - imag(b)) > EPS) return

    imag(a) < imag(b);
</pre>
 return false:
convex_polygon convexhull(polygon a) {
  sort(a.points.begin(), a.points.end(), cmp);
  vector<point> lower, upper;
  for (int i = 0; i < a.points.size(); i++) {</pre>
   while (lower.size() >= 2 &&
        cross(lower.back() - lower[lower.size()
      2], a.points[i] - lower.back()) < EPS)
     .lower.pop_back();
    while (upper size() >= 2 &&
        cross(upper.back() - upper[upper.size()
      2], a.points[i] - upper.back()) > -EPS)
      upper.pop_back();
    lower.push_back(a.points[i]);
    upper.push_back(a.points[i]);
  lower.insert(lower.end(), upper.rbegin() + 1,

    upper.rend());

 return convex_polygon(lower);
Maximum Colinear Points
```

## &points) { if(points.size() <= 2) return points.size();</pre> int best = 0; unordered\_map<Slope, int> counter; for(int i=0:i<points.size():i++) { for(int j=i+1; j<points.size(); j++) {</pre> Slope slope(points[i].second-points[j] .second,points[i].first-points[j].first); best = max(best, ++counter[slope]+1): if(i != points.size()-1) counter.clear(): return best; 3D Geometry struct point3d { double x, y, z; point3d operator+(point3d a) const { return $\rightarrow$ {x+a.x, y+a.y, z+a.z}; } point3d operator\*(double a) const { return $\rightarrow$ {x\*a, y\*a, z\*a}; } point3d operator-() const { return {-x, -y, point3d operator-(point3d a) const { return $\rightarrow$ \*this + -a; } point3d operator/(double a) const { return \*this \* (1/a); } double norm() { return x\*x + y\*y + z\*z; } double abs() { return sqrt(norm()); } point3d normalize() { return \*this / this->abs(); } double dot(point3d a, point3d b) { return $\rightarrow$ a.x\*b.x + a.y\*b.y + a.z\*b.z; } point3d cross(point3d a, point3d b) { return $\{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,$ a.x\*b.y - a.y\*b.x; } struct line3d { point3d a, b; }; struct plane { double a, b, c, d; } // a\*x + $\rightarrow b*y + c*z + d = 0$ struct sphere { point3d c; double r; }; #define sq(a) ((a)\*(a)) #define $c\bar{b}(a)$ ((a)\*(a)\*(a))

|const ll range = 10000:

11 p, q;

if(qP==0) {

p = 1, q = 0;return:

p = pP, q = qP;

namespace std {

template<>

range;

→ M PI: }

unsigned infinity (1.0)

11 g = \_\_gcd(pP, qP); pP /= g, qP /= g; if(qP\_< 0) pP\_\*= -1, qP \*= -1;

struct hash<Slope> { // typical

rectangular/lattice hash

// n points in [-range, range]

// compute the largest colinear subset

return other.p == p && other.q == q;

Slope(ll pP=0, ll qP=0) {

struct Slope { // a rational number with

```
double volume(circle a) { return 4.0/3.0 *
                                                              \hookrightarrow cb(a.r) * M PI; }
                                                              10
                                                                     Optimization
                                                             Snoob
                                                              // SameNumberOfOneBits, next permutation
                                                             int snoob(int a) {
                                                              int b = a & -a, c = a + b;
return c | ((a ^ c) >> 2) / b;
                                                              // example usage
                                                             int main() {
   char l1[] = {'1', '2', '3', '4', '
   char l2[] = {'a', 'b', 'c', 'd'};
   int d1 = 5, d2 = 4;
   // prints 12345abad, 1234a5bcd, ...
bool operator == (const Slope &other) const {
                                                               int^*min = (1 < < d1) - 1, max = min < < d2;
                                                               for (int i = min; i < max; i = snoob(i)) {
   int p1 = 0, p2 = 0, v = i;
   while (p1 < d1 || p2 < d2) {
                                                                 cout \langle \langle (v \& 1) ? 11[p1++] : 12[p2++] \rangle;
                                                                 v /= 2;
  size_t operator() (const Slope &r) const {
   return (2*range+1) * (r.p + range) + r.q +
                                                                cout << '\n':
                                                             Powers
                                                             bool isPowerOf2(ll a) {
                                                              return a > 0 \&\& !(a \& a-1);
int max_colinear_points(vector<pair<11,11>>
                                                             bool isPowerOf3(11 a) {
   return a>0&&!(12157665459056928801u11%a);
                                                             bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
                                                               return abs(x-round(x)) < 0.00000000001;
                                                             11 Python
                                                              Recursion Limit Removal (Basic)
                                                             import sys
                                                             sys.setrecursionlimit(10**6)
                                                              Recursion Limit Removal (Advanced)
```

```
# @bootstrap over recursive function
# replace 'return' with 'yield'
# for when sys method does not work
from types import GeneratorType
def bootstrap(f, stack=[]):
 def wrappedfunc(*args, **kwargs):
  if stack:
return f(*args, **kwargs)
  else:
   to = f(*args, **kwargs)
   while True:
if type(to) is GeneratorType:
      stack.append(to)
     to = next(to)
      stack.pop()
     if not stack:
   to = stack[-1].send(to)
 return wrappedfunc
# EXAMPLE recursive fibonacci
@bootstrap
 if (n < 2):
yield n
 yield (yield f(n-1)) + (yield f(n-2))
```

## Python 3 Compatibility

```
import sys from future import division, print function O(n\sqrt{n})
                                                         if sys.version_info[0] < 3:
#define cb(a) ((a)*(a)*(a))
double surface(circle a) { return 4 * sq(a.r) * from future_builtins import ascii, filter,

→ hex, map, oct, zip
```

```
12 Additional
Judge Speed
   kattis: 0.50s
 / codeforces: 0.421s
// atcoder: 0.455s
#include <bits/stdc++.h>
using namespace std;
int v = 1e9/2, p = 1;
int main() {
  for (int i = 1; i <= v; i++) p *= i;</pre>
cout << p;
Judge Pre-Contest Checks
    int128 and float128 support?
does extra or missing whitespace cause WA?
-documentation up to date?
-printer usage available and functional?
// each case tests a different fail condition
// try them before contests to see error codes
```

struct g { int arr[1000000]; g(){}};

// O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE

if (n == 3) while(1) putchar\_unlocked('a');

vector<g> a;

if (n == 4) assert(0);

return n + judge(n + 1);

if (n == 5) 0 / 0; if (n == 6) \*(int\*)(0) = 0;

## GCC Builtin Docs // 128-bit integer \_\_int128 a; unsigned \_\_int128 b; // 128-bit float // minor improvements over long double \_\_float128 c; // log2 floor \_\_lg(n); // number of 1 bits // can add ll like popcountll for long longs \_\_builtin\_popcount(n); // number of trailing zeroes \_\_builtin\_ctz(n); // number of leading zeroes \_\_builtin\_clz(n); // 1-indexed least significant 1 bit

```
__builtin_ffs(n);
// parity of number
__builtin_parity(n);
Limits
                       \pm 2147483647 \mid \pm 2^{31} - 1 \mid 10^9
int
                                             \frac{1}{2}<sup>32</sup> - 1|10<sup>9</sup>
                          4294967295
uint
        \pm 9223372036854775807 | \pm 2^{63} - 1 | 10^{18}
                                               \overline{2}^{64} - \overline{1}|10^{19}
         18446744073709551615
ull
```

```
|1128| \pm 170141183460469231...| \pm 2^{127} - 1|10^{38}
 |\underline{u128}| \ \ 340282366920938463... | \ \ \bar{2}^{128} - 1| 10^{38} 
Complexity classes input size (per second):
O(n^n) or O(n!)
```

```
n < 10
O(2^n)
                                         n < 30
O(n^3)
                                       n < 1000
O(n^2)
                                     n < 30000
                                        n < 10^6
O(n \log n)
                                        n < 10^7
O(n)
                                        n < 10^9
```