```
General
                              7 Graphs
    Algorithms
                              8 2D Geometry
    Structures
                              9 3D Geometry
    Strings
                              10 Optimization
    Greedy
                              11 Python
    Math
                              12 Additional
     General
run.sh
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
for i in *.in; do
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
done
Header
// use better compiler options
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
// include everything
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
using namespace std;
using namespace __gnu_cxx; // rope
using namespace __gnu_pbds; // tree/trie
// common defines
#define fastio

→ ios_base::sync_with_stdio(0);cin.tie(0);
#define nostacklim rlimit RZ; getrlimit(3,&RZ)
 \rightarrow );RZ.rlim cur=-1;setrlimit(3,\&RZ);
#define DEBUG(v) cerr<<__LINE__ <<": "'<<#v<<" =
\Rightarrow "<<v<<'\n'; #define TIMER

→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
#define i128 __int128
#define u128 unsigned i128
#define ld long double
// global variables
mt19937 rng((uint32_t)chrono::steady

    _clock::now().time_since_epoch().count());
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c - 0';
void read(int& n) {
   char c; n = 0; int s = 1;
   if ((c=getchar_unlocked())=='-') s = -1;
 else n = c - \sqrt{0}:
 while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c -
void read(ld& n) {
 char c; n = 0;

ld m = 0, o = 1; bool d = false; int s = 1;

if ((c=getchar_unlocked())=='-') s = -1;
 else if (c == '.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n') \{ // max - compare = a < b, reset = a < 0 \}
  if (c == '.') d = true;
```

```
else if (d) { m=m*10+c-'0'; o*=0.1; } else n = n * 10 + c - '0':
 n = s * (n + m * o):
void read(double& n) {
ld m; read(m); n = m;
void read(float& n) {
  ld m; read(m); n = m;
void read(string& s) {
 char c; s = ""
 while((c=getchar_unlocked())!=' '&&c!='\n')
bool readline(string& s) {
 char c; s = ""
 while(c=getchar_unlocked()) {
 if (c == '\n') return true;
if (c == EOF) return false;
  s += c:
 return false:
void print(unsigned int n) {
 if (n / 10) print(n / 10);
 putchar unlocked(n % 10 + '0'):
void print(int n) {
 if (n < 0) { putchar_unlocked('-'); n*=-1; }
 print((unsigned int)n);
Additional cout
ostream& operator << (ostream& o, unsigned
\rightarrow __int128 n) {
auto t = n<0 ? -n : n; char b[128], *d =

    end(b);
    do *--d = '0'+t%10, t /= 10; while (t);
    o.rdbuf()->sputn(d,end(b)-d);

 return o:
ostream& operator<<(ostream& o, __int128 n) {
 if (n < 0) return o << '-' << (unsigned
     __int128)n;
 return o << (unsigned int128)n:
ostream& operator<<(ostream& o, __float128 n) {
 return o << (long double)n;
Common Structs
   n-dimension vectors
// Vec<2, int> v(n, m) = arr[n][m]

// Vec<2, int> v(n, m, -1) default init -1

template<int D, typename T>
struct Vec : public vector<Vec<D-1, T>> {
template<typename... Args>
.Vec(int n=0, Args... args) : vector<Vec<D-1,</pre>
    T >> (n, Vec < D-1, T > (args...)) {}
template<typename T>
struct Vec<1, T> : public vector<T> {
   Vec(int n=0, T val=T()) : vector<T>(n, val) {}
    Algorithms
Binary Search
// search for k in [p,n)
template<tvpename T>
int binsearch(T x[], int k, int n, int p = 0) {
for (int i = n; i >= 1; i /= 2)

while (p+i < n && x[p+i] <= k) p += i;

return p; // bool: x[p] == k;
Min/Max Subarray
// min - compare = a > b, reset = a > 0
```

```
// returns {sum, {start, end}}
pair<int, pair<int, int>>
     ContiguousSubarray(int* a, int size,
     bool(*compare)(int, int).
 bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
 → 0, s = 0;
for (int i = 0; i < size; i++) {
cur += a[i];
  if ((*compare)(best, cur)) { best = cur;
 start = s; end = i; }
if ((*reset)(cur)) { cur = 0; s = i + 1; }
 return {best, {start, end}}:
Quickselect
#define OSNE -999999
int partition(int arr[], int 1, int r)
  int x = arr[r], i = 1;
 for (int j = 1; j <= r - 1; j++)
if (arr[j] <= x)
   swap(arr[i++], arr[j]);
  swap(arr[i], arr[r]);
 return i:
 // find k'th smallest element in unsorted array
 \hookrightarrow only if all distinct
 int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - 1 + 1)) return QSNE;
swap(arr[1 + rng() % (r-1+1)], arr[r]);
int pos = partition(arr, 1, r);
 if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return gselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
→ on both axis
pair<int, int> saddleback_search(int** arr, int
 x, int y, int v) {
int i = x-1, j = 0;
while (i >= 0 && j < y) {</pre>
  if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i--: j++;
 return {-1, -1};
Ternary Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a. int b. int (*f)(int)) {
 while (b-a > 4) {
  int m = (a+b)/2;
if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1;
 for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
   a = i;
 return a;
 #define TERNPREC 0.000001
double ternsearch(double a, double b, double
 while (b-a > TERNPREC * 4) {
double m = (a+b)/2;
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
   else b = m + TERNPREC:
 for (double i = a + TERNPREC: i <= b: i +=
 → TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return a;
```

```
Golden Section Search
```

```
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
double goldsection(double a, double b, double

    (*f)(double)) {
    double r = (sqrt(5)-1)/2, eps = 1e-7;
    double x1 = b - r*(b-a), x2 = a + r*(b-a);
    double f1 = f(x1), f2 = f(x2);
}

 while (b-a > eps)
 white (b-a > eps)
if (TERNCOMP(f2,f1)) {
   b = x2; x2 = x1; f2 = f1;
   x1 = b - r*(b-a); f1 = f(x1);
} else {
   a = x1; x1 = x2; f1 = f2;
   x2 = a + r*(b-a); f2 = f(x2);
 return a;
```

3 Structures

Fenwick Tree

```
// Fenwick tree, array of cumulative sums
\hookrightarrow O(log n) updates. O(log n) gets
struct Fenwick {
int n; ll* tree;
void update(int i, int val) {
 while (i <= n) {
  tree[i] += val;
  i += i & (-i);</pre>
Fenwick(int size) {
  tree = new ll[n+1];
 for (int i = 1; i <= n; i++)
tree[i] = 0;
Fenwick(int* arr, int size) : Fenwick(size) {
   for (int i = 0; i < n; i++)
      update(i, arr[i]);</pre>
 ~Fenwick() { delete[] tree; }
11 operator[](int i) {
 if (i < 0 || i > n) return 0;
ll sum = 0;
 while (i>0)
  sum += tree[i];
i -= i & (-i);
 return sum;
11 getRange(int a, int b) { return
    operator[](b) - operator[](a-1); }
```

Hashtable

```
// similar to unordered map, but faster
struct chash {
  const uint64 t C = (11)(2e18 * M_PI) + 71;
  ll operator()(11 x) const { return
     _builtin_bswap64(x*C); }
int main() {
  gp_hash_table<11,int,chash>
\rightarrow hashtable({},{},{},{},{1<<16});
for (int i = 0; i < 100; i++)
. hashtable[i] = 200+i;
.if (hashtable.find(10) != hashtable.end())
. cout << hashtable[10];
```

```
Ordered Set
                                                          for (int i = 0; i < size(arr)-(1<<k)+1; i
 template <typename T>
 using oset = tree<T,null_type,less<T>,rb_tree
                                                           m[k][i] = min(m[k-1][i].
     tag, tree order statistics node update>;
                                                           m[k-1][i+(1<<(k-1))]:
 template <typename T, typename D> using omap = tree <T,D,less <T>,rb_tree
                                                         // min of range [l,r]

→ _tag,tree_order_statistics node update>:

                                                        T query(int 1, int r) {
 int main()
                                                         int k = __lg(r-l+1);
                                                          return \min(m[k][1], m[k][r-(1<< k)+1]);
  oset<int> o_set;
  o set.insert(5); o_set.insert(1);
 \rightarrow o_set.insert(3);
  // get second smallest element
                                                        Trie
  cout << *(o_set.find_by_order(1));</pre>
                                                        typedef trie<string, null_type,
  // number of elements less than k=4
cout << ' ' << o_set.order_of_key(4) << '\n';</pre>

→ trie string access traits<>.
                                                         pat_trie_tag, trie_prefix_search_node_update>
  // equivalent with ordered map
                                                         \rightarrow trie_type;
  omap<int,int> o_map;
                                                        int main() {
  o_map[5]=1;o_map[1]=2;o_map[3]=3;
                                                         // generate trie
  cout << (*(o_map.find_by_order(1))).first;</pre>
                                                        trie_type trie;
  cout << ' ' << o_map.order_of_key(4) << '\n';
                                                        for (int i = 0; i < 20; i++)
                                                         trie.insert(to_string(i)); // true if new,
                                                         \rightarrow false if old
 Rope
                                                         // print things with prefix "1"
 // O(\log n) insert, delete, concatenate
                                                         auto range = trie.prefix range("1");
 int main() {
                                                        for (auto it = range.first; it !=
  // generate rope
  rope<int> v;
                                                           range.second; it++)
  for (int i = 0; i < 100; i++)
                                                         cout << *it << "
  v.push_back(i);
  // move range to front
                                                        Wavelet Tree
  rope<int> copy = v.substr(10, 10);
                                                        using iter = vector<int>::iterator;
  v.erase(10, 10);
                                                        struct WaveletTree {
  v.insert(copy.mutable_begin(), copy);
                                                         Vec<2, int> C; int s;
for (auto it : y)
cout << it << "":
  // print elements of rope
                                                         // sigma = highest value + 1
                                                         WaveletTree(vector<int>& a, int sigma) :
                                                           s(sigma), C(sigma*2, 0) {
                                                         build(a.begin(), a.end(), 0, s-1, 1);
 Segment Tree
 //max(a,b), min(a,b), a+b, a*b, qcd(a,b), a\hat{b}
                                                        void build(iter b, iter e, int L, int U, int
                                                         u) {
if (L == U) return;
 struct SegmentTree {
  typedef int T;
                                                          \underline{int} \widetilde{M} = (\widetilde{L} + \widetilde{U})/2;
  static constexpr T UNIT = INT_MIN;
                                                          C[u].reserve(e-b+1); C[u].push_back(0);
 T f(T a, T b) {
   if (a == UNIT) return b;
   if (b == UNIT) return a;
                                                         for (auto it = b; it != e; ++it)
    C[u].push_back(C[u].back() + (*it<=M));
                                                          auto p = stable_partition(b, e, [=](int
   return max(a,b);
                                                           i) {return i <= M; }):
  int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def).
                                                          build(b, p, L, M, u*2);
                                                          build(p, e, M+1, U, u*2+1);
 → n(n) {}
SegmentTree(vector<T> arr) :
                                                         // number of occurences of x in [0,i)
 int rank(int x, int i) {
                                                         for (int i=0; i < arr.size(); i++)

    update(i,arr[i]);

  void update(int pos, T val) {
                                                          if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
   for (s[pos += n] = val; pos /= 2;)
    s[pos] = f(s[pos * 2], s[pos*2+1]);
                                                         return i:
  T query(int b, int e) { // query [b, e)
   Tra = UNIT, rb = UNIT;
                                                         ^{\prime\prime} number of occurences of x in [l,r)
   for (b+=n, e+=n; b<e; b/=2, e/=2) {
    if (b % 2) ra = f(ra, s[b++]);
    if (e % 2) rb = f(s[--e], rb);
                                                        int count(int x, int 1, int r) {
                                                         return rank(x, r) - rank(x, 1);
                                                        // kth smallest in [l, r)
int kth(int k, int l, int r) const {
  int L = 0, U = s-1, u = 1, M, ri, rj;
  while (L.!= U) {
   return f(ra, rb);
  .T get(int p) { return query(p, p+1); }
                                                          M = (L+U)/2;
ri = C[u][1]; rj = C[u][r]; u*=2;
 Sparse Table
                                                           if (k \le rj-ri)^{\circ}l = ri, r = rj, U = M;
 template<class T> struct SparseTable {
                                                           else k -= rj-ri, l -= ri, r -= rj,
  vector<vector<T>> m;
                                                          L = M+1, ++u:
  SparseTable(vector<T> arr) {
   m.push back(arr):
                                                         return U:
   for (int k = 1; (1 << (k)) <= size(arr); k++) {|}
    m.push back(vector<T>(size(arr)-(1<<k)+1)); \ \ // # elements between [x,y] in [l,r)
```

```
mutable int L. U:
 int range(int x, int y, int 1, int r) const {
  if (y < x \text{ or } r \le 1) return 0;
  L = x; U = y;
return range(1, r, 0, s-1, 1);
 int range(int 1, int r, int x, int y, int u)
 - const {
   if (y < L or U < x) return 0;
   if (L <= x and y <= U) return r-1;</pre>
  int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return range(ri, rj, x, M, u*2) + range(1-ri,
    r-rj, M+1, y, u+2+1);
  ^{\prime}// # elements <= x in [l, r)
 int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);
     Strings
Aho Corasick
// range of alphabet for automata to consider
// MAXC = 26, OFFC = 'a' if only lowercase const int MAXC = 256; const int OFFC = 0;
struct aho_corasick {
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
  véctor<state> s;
 int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
  for (int i = 0; i < size; i++) {
   int cur = 0;
   for (int c : arr[i]) {
    if (s[cur].go[c-OFFC] == -1) {
      s[cur].go[c-OFFC] = s.size();
      s.push_back(state());
     cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
     s[0].go[c] = 0;
  queue int> sq;
for (int c = 0; c < MAXC; c++) {
   if (s[0].so[c] != 0) {
      ...s[s[0].so[c]].fail = 0;</pre>
     sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
   for (int c = 0; c < MAXC; c++) {
  if (s[e].go[c] != -1) {
     int failure = s[e].fail;
      while (s[failure].go[c] == -1)
      failure = s[failure].fail;
failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
      for (auto length : s[failure].out)
       s[s[e].go[c]].out.insert(length);
      sq.push(s[e].go[c]);
 // list of {start pos, pattern id}
 vector<pair<int, int>> search(string text)
  vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {
   while (s[cur].go[text[i]-OFFC] == -1)
```

```
..cur = s[cur].fail;
   cur = s[cur].go[text[i]-OFFC];
   if (s[cur].out.size())
    for (auto end : s[cur].out)
toret.push_back({i - end.first + 1,
    end.second});
  return toret:
Boyer Moore
struct defint { int i = -1: }:
vector<int> bovermoore(string txt, string pat)
vector<int> toret; unordered_map<char, defint>
→ badchar;
 int m = pat.size(), n = txt.size();
 for (int i = 0; i < m; i++) badchar[pat[i]].i
 \rightarrow = i;
 int s = 0:
 while (s <= n - m) {
  int j = m - 1;
  while (j \ge 0 \&\& pat[j] == txt[s + j]) j--;
  if (j < 0) {
   toret.push back(s);
   s += (s + m < n) ? m - badchar[txt[s +
\rightarrow m]].i : 1;
 .} else
   s += max(1, j - badchar[txt[s + j]].i);
 return toret;
English Conversion
const string ones[] = {"", "one", "two",
   "three", "four", "five", "six", "seven", "eight", "nine";
const string teens[] ={"ten", "eleven",
   "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen"};
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy".
"eighty", "ninety");
const string mags[] = {"thousand", "million",
    "billion", "trillion", "quadrillion",
    "quintillion", "sextillion",
    "septillion"}:
string convert(int num, int carry) {
if (num < 0) return "negative " +
    convert(-num, 0);
if (num < 10) return ones[num];
if (num < 20) return teens[num % 10];
if (num < 100) return tens[(num / 10) - 2] +
    (num%10==0?"":" ") + ones[num % 10];
(num < 1000) return ones[num / 100] -
     (num/100==0?"":" ") + "hundred" +
     (num%100==0?"":" ") + convert(num % 100.
    0):
return convert(num / 1000, carry + 1) + " " +
    mags[carrv] + " " + convert(num % 1000.
    0):
string convert(int num) {
return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
    vector<int> toret;
 int m = txt.length(), n = pat.length();
int next[n + 1];
for (int i = 0; i < n + 1; i++)
  next[i] = 0;</pre>
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
 while (j > 0 && pat[j] != pat[i])
```

```
j = next[j];
                                                                                           vector<string> duval(string s) {
   if (j > 0 | pat[j] == pat[i])
                                                                                             int n = s.length();
                                                                                             vector<string> lyndon;
    next[i + 1] = i + 1;
                                                                                            for (int i = 0; i < n;) {
    int j = i+1, k = i;
    for (; j < n && s[k] <= s[j]; j++)
    if (s[k] < s[j]) k = i;
  for (int i = 0, j = 0; i < m; i++) {
  if (txt[i] == pat[j]) {
  ..if (++i == n)
                                                                                              else k++;
for (; i <= k; i += j - k)
     toret.push_back(i - j + 1);
   } else if (j > 0) {
                                                                                                lyndon.push_back(s.substr(i,j-k));
    j = next[j];
                                                                                            return lyndon;
                                                                                            // lexicographically smallest rotation
 return toret;
                                                                                           int minRotation(string s) {
                                                                                            int n = s.length(); s += s;
Longest Common Prefix (array)
                                                                                            auto d = duval(s); int i = 0, a = 0;
while (a + d[i].length() < n) a +=</pre>
// longest common prefix of strings in array
string lcp(string* arr, int n, bool sorted =
                                                                                                  d[i++].length();
                                                                                            while (i && d[i] == d[i-1]) a -=
 → false) {
if (n == 0) return "":
                                                                                            → d[i--].length();
 if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
                                                                                            return a:
 while (v < arr[0].length() && arr[0][v] ==
Minimum Word Boundary
                                                                                            // minimum word boundary
                                                                                            // compose string s using words from dict
                                                                                            // NOTE: can reuse words from dict
                                                                                           unsigned int mwb(string s, set<string> dict) {
Longest Common Subsequence
                                                                                            int l = s.size();
vector<unsigned int> arr(l+1, -1);
string lcs(string a, string b) {
 int m = a.length(), n = b.length();
                                                                                            arr[0] = 0;
for (int i = 0; i < 1; i++) {
   if (arr[i] != -1) {
 int L[m+1][n+1]:
 for (int i = 0; i <= m; i++) {
    for (int j = 0; j <= n; j++) {
        if (i == 0 || j == 0) L[i][j] = 0;
        else if (a[i-1] == b[j-1]) L[i][j] =
                                                                                                for (auto e : dict) {
                                                                                                  int L = e.size();
if (l >= i + L) {
   bool isGood = true;

    L[i-1][j-1]+1;

                                                                                              ....for (int j = 0; isGood && j < L; j++)
.....if (s[i+j] != e[j])
     else L[i][j] = \max(L[i-1][j], L[i][j-1]);
                                                                                           in (lsGood)
in arr[i+L] = min(arr[i]+1, arr[i+L]);
in arr[i+L] = min(arr[i]+1, arr[i+L]);
in arr [i+L] = min(arr[i]+1, arr[i]+1, ar
 // return L[m][n]; // length of lcs
 string out = ""
  int i = m - 1, j = n - 1;
 while (i >= 0 && j >= 0) {
 if (a[i] == b[j]) {
                                                                                            return arr[1]:
   out = a[i--] + out;
                                                                                           Hashing
   else if (L[i][j+1] > L[i+1][j]) i--;
                                                                                           #define HASHER 27
   else j--;
                                                                                           ull basicHash(string s) {
                                                                                            ull v = 0;
 return out;
                                                                                            for (auto c : s) v = (c - 'a' + 1) + v *
                                                                                             → HASHER;
Longest Common Substring
                                                                                            return v:
// l is array of palindrome length at that
                                                                                           const int MAXN = 1000001;
ull base[MAXN] = {1};
int manacher(string s, int* 1) {
                                                                                           void genBase(int n) {
 int n = s.length() * 2;
                                                                                            for (int i = 1; i <= n; i++)
base[i] = base[i-1] * HASHER;
 for (int i = 0, j = 0, k; i < n; i += k, j =
 \rightarrow max(j-k, 0)) {
                                                                                           struct advHash {
  ull v, l; vector<ull> wip;
  while (i \ge j \&\& i + j + 1 < n \&\& s[(i-j)/2]
 \Rightarrow == s[(i+j+1)/2]) j++;
                                                                                            advHash(string& s): v(0) {
   .1[i] = j;
                                                                                              wip = vector <ull>(s.length()+1);
  for (k = 1; i >= k \&\& j >= k \&\& l[i-k] !=
                                                                                               wip[0] = 0;
      i-k: k++)
                                                                                              for (int i = 0; i < s.length(); i++)
wip[i+1] = (s[i] - 'a' + 1) + wip[i] *
    1[i+k] = min(1[i-k], j-k);
                                                                                                HASHER;
 return *max_element(1, 1 + n);
                                                                                              l = s.length(); v = wip[l];
                                                                                             ull del(int pos, int len) {
Cyclic Rotation (Lyndon)
                                                                                              return v - wip[pos+len]*base[l-pos-len] +
// simple strings = smaller than its nontrivial
                                                                                                  wip[pos]*base[l-pos-len];
// lyndon factorization = simple strings
                                                                                             ull substr(int pos, int len) {
     factorized
                                                                                             return del(pos+len, (1-pos-len)) -
// "abaaba" -> "ab", "aab", "a"
```

→ wip[pos]*base[len]:

```
ull replace(int pos, char c) {
 return v - wip[pos+1]*base[l-pos-1] + ((c -
     'a' + 1) + wip[pos] *
ull replace(int pos, string s) {
  // can't increase total string size
    wip[pos+s.size()]*base[l-pos-s.size()], c =
   wip[pos];
 for (int i = 0; i < s.size(); i++)
.c = (s[i]-'a'+1) + c * HASHER;
  return r + c * base[1-pos-s.size()]:
Subsequence Count
// "banana", "ban" >> 3 (ban, ba..n, b..an)
ull subsequences(string body, string subs) {
 int m = subs.length(), n = body.length();
 if (m > n) return 0;
 ull** arr = new ull*[m+1];
for (int i = 0; i <= m; i++) arr[i] = new
 \rightarrow ull[n+1];
for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[0][i] = 1;
for (int i = 1; i <= m; i++)
 for (int j = 1; j <= n; j++)
...arr[i][j] = arr[i][j-1] + ((body[j-1] ==
 \rightarrow subs[i-1])? arr[i-1][i-1] : 0):
return arr[m][n];
Suffix Array + LCP
struct SuffixArray {
 vector<int> sa, lcp;
 SuffixArray(string& s, int lim=256) {
int n = s.length() + 1, k = 0, a, b;
  vector<int> x(begin(s), end(s)+1), y(n),
 \rightarrow ws(max(n, lim)), rank(n);
  sa = lcp = y;
iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = \max(1, j *
 \rightarrow 2), lim = p) {
   p = j; iota(begin(y), end(y), n - j);
   for (int i = 0; i < (n); i++)
    if (sa[i] >= j)
     v[p++] = sa[i] -
   fill(begin(ws), end(ws), 0);
   for (int i = 0; i < (n); i++) ws[x[i]]++;
for (int i = 1; i < (lim); i++) ws[i] +=
   ws[i - 1];
   for (int i = n: i--:) sa[--ws[x[v[i]]]] =
 y[i];
  swap(x, y); p = 1; x[sa[0]] = 0;
for (int i = 1; i < (n); i++) {
   a = sa[i - 1]; b = sa[i];</pre>
    x[b] = (y[a] = y[b] \&\& y[a + j] = y[b + j]
    j]) ? p - 1 : p++;
 for (int i = 1: i < (n): i++) rank[sa[i]] =
 for (int i = 0, j; i < n - 1; lcp[rank[i++]]
\rightarrow = k)
 for (k && k--, j = sa[rank[i] - 1];
... s[i + k] == s[j + k]; k++);
 // smallest cyclic shift
 int cyclic() { return sa[0]; }
 // longest repeated substring
 pair<int,int> lrs() {
  int length = -1, index = -1;
  for (int i = 0; i < lcp.size(); i++) {
  if (lcp[i] > length) {
   length = lcp[i];
    index = sa[i];
```

```
return {index,length};
// count distinct substrings, excluding empty
int distincts() {
    int n = sa.size() - 1, r = n - sa[0];
    for (int i = 1; i < lcp.size(); i++)
 r += (n - sa[i]) - lcp[i - 1];
 return r:
// count repeated substrings, excluding empty
int repeateds() {
int r = 0;

for (int i = 1; i < lcp.size(); i++)

.r += max(lcp[i] - lcp[i-1], 0);
 return r:
```

Suffix Tree (Ukkonen's)

```
struct SuffixTree {
// n = 2*len+10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
\rightarrow t[N][ALPHA],1[N],r[N],p[N],s[N],v=0,q=0,m=2;
string a;
void ukkadd(int i, int c) { suff:
 if (r[v]<=q) {
  if (t[v][c]==-1) { t[v][c]=m; l[m]=i; p[m++]=v; v=s[v]; q=r[v]; goto suff; }
   v=t[v][c]; q=l[v];
 while (q < r[m]) { v = t[v][toi(a[q])];
   q+=r[v]-1[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
q=r[v]-(q-r[m]); m+=2; goto suff;
SuffixTree(string a) : a(a) {
 fill(r,r+N,(int)(a).size());
 memset(s, 0, sizeof s);

memset(t, -1, sizeof t);

fill(t[1],t[1]+ALPHA,0);

s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
 for(int i=0; i < a. size(); i++)
    ukkadd(i.toi(a[i])):
// Longest Common Substring between 2 strings
 // returns {length, offset from first string}
pair<int, int> best;
int lcs(int node, int i1, int i2, int olen) {
   if (1[node] <= i1 && i1 < r[node]) return 1;
   if (1[node] <= i2 && i2 < r[node]) return 2;
 int mask=0,
  len=node?olen+(r[node]-l[node]):0;
 for(int c=0; c<ALPHA; c++) if
   (t[node][c]!=-1)
mask |= lcs(t[node][c], i1, i2, len);
  if (mask==3)
→ best=max(best, {len,r[node]-len});
 return mask;
static pair<int, int> LCS(string s, string t)
 SuffixTree
 \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2)):
 st.lcs(0, s.size(), s.size()+t.size()+1, 0); return st.best;
```

```
long double v = 1;
String Utilities
                                                            for (ull i = 1; i <= r; i++)
void lowercase(string& s) {
                                                             v = v * (n-r+i) /i;
 transform(s.begin(), s.end(), s.begin(),
                                                            return (ull)(v + 0.001);
    ::tolower):
                                                               requires modulo math
void uppercase(string& s) {
                                                            // can optimize by precomputing mfac and
 transform(s.begin(), s.end(), s.begin(),
 ull nCr(ull n, ull r, ull m) {
                                                            return mfac(n, m) * minv(mfac(k, m), m) % m *
void trim(string &s) {
                                                                minv(mfac(n-k, m), m) \% m;
 s.erase(s.begin(),find_if_not(s.begin(),s
      .end(), [](int c){return
                                                            Multinomials
    isspace(c);}));
                                                           11 multinomial(vector<int>& v) {
11 c = 1, m = v.empty() ? 1 : v[0];
  for(int i = 1; i < v.size(); i++)</pre>
 s.erase(find_if_not(s.rbegin(),s.rend(),[](int
    c){return isspace(c):}).base().s.end()):
                                                             for (int j = 0; j < v[i]; j++)
...c = c * ++m / (j+1);
vector<string> split(string& s, char token) {
     vector<string> v; stringstream ss(s);
                                                            return c:
     for (string e;getline(ss,e,token);)
          v.push back(e);
                                                            Chinese Remainder Theorem
                                                            bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo)
                                                            \vec{1}1 x, y, d; mo = m[0]; re = r[0];
      Greedy
                                                            for (int i = 1: i < n: i++) {
                                                             if ([i] - r, i, i, i+, i, d = egcd(mo, m[i], x, y);
if ((r[i] - re) % d != 0) return false;
x = (r[i] - re) / d * x % (m[i] / d);
re += x * mo;
Interval Cover
// L,R = interval [L,R], in = {{l,r}, index}

// does not handle case where L == R

vector<int> intervalCover(double L, double R
                                                              mo = mo / d * m[i];
 → vector<pair<pair<double,double>,int>> in)
                                                              re %= mo;
     int i = 0; pair < double, int > pos = {L,-1};

    vector<int> a;
    sort(begin(in), end(in));

                                                            re = (re + mo) % mo;
return true:
     while (pos.first < R) {
          double cur = pos.first;
while (i < (int)in.size() &&</pre>
                                                            Count Digit Occurences
    in[i].first.first <= cur)</pre>
                                                            /stcount(n.d) counts the number of occurences of
                                                            \rightarrow a digit d in the range [0,n]*/
     max(pos,{in[i].first.second.in[i].second})
                                                            11 digit count(11 n, 11 d) {
    i++;
                                                            ll result = 0;
while (n != 0)
          if (pos.first == cur) return {};
          a.push_back(pos.second);
                                                             result += ((n\%10) == d?1:0):
                                                             n /= 10:
     return a:
                                                            return result:
6 Math
                                                            11 count(11 n, 11 d) {
    if (n < 10) return (d > 0 && n >= d);
Catalan Numbers
                                                            if ((n \% 10) != 9) return digit count(n, d) +
ull* catalan = new ull[1000000];
void genCatalan(int n, int mod) {
  catalan[0] = catalan[1] = 1;
  for (int i = 2; i <= n; i++) {</pre>
                                                            \rightarrow count(n-1, d):
                                                            return 10*count(n/10, d) + (n/10) + (d > 0):
 ior (int i = 2; i \lefta in, i = 1; catalan[i] = 0;
catalan[i] = i = 1; j >= 0; j--) {
catalan[i] += (catalan[j] * catalan[i-j-1])
                                                            Discrete Logarithm
                                                            int discretelog(int a. int b. int m) {
                                                            11 \text{ n} = \text{sgrt}(\text{m}) + 1, an = 1:
    % mod:
                                                            for (ll i = 0; i < n; ++i)
an = (an * a) % m; unordered_map<ll, il> vals;
 if (catalan[i] >= mod)
catalan[i] -= mod;
                                                            for (11 q = 0, cur = b; q <= n; q++) {
   vals[cur] = q;
// TODO: consider binomial coefficient method
                                                              cur = (cur * a) \% m:
Combinatorics (nCr, nPr)
                                                             for (ll p = 1, cur = 1; p \le n; p++) {
                                                              cur = (cur * an) % m;
// can optimize by precomputing factorials, and
                                                              if (vals.count(cur))
\hookrightarrow fact[n]/fact[n-r]
                                                               int ans = n * p - vals[cur];
ull nPr(ull n, ull r) {
                                                              return ans:
 for (ull i = n-r+1; i <= n; i++)
 .v *= i;
return v:
                                                            return -1:
úll nPr(ull n. ull r. ull m) {
                                                            Euler Phi / Totient
                                                            int phi(int n) {
 for (ull i = n-r+1: i <= n: i++)
  v = (v * i) \% m;
                                                            for (int i = 2; i * i <= n; i++) {
 return v;
                                                             if (n % i == 0) r -= r / i;
while (n % i == 0) n /= i;
ull nCr(ull n, ull r) {
```

```
if (n > 1) r = r / n:
                                                          // make sure you call genPrimes first
 return r:
                                                         // return vector of prime factor powers as
#define n 100000
                                                            vector v of size pi(n)
ll phi[n+1];
                                                         // so that v[i] = power of primes[i] dividing
void computeTotient() {
 for (int i=1; i<=n; i++) phi[i] = i;
                                                         \stackrel{\longrightarrow}{//} \stackrel{n!}{0(pi(n)} * log(n)) where pi(n) is prime
 for (int p=2; p<=n; p++) {
                                                         if (phi[p] == p) {
                                                         // so basically O(n) since pi(n) = O(n/log(n))
    phi[p] = p-1;
                                                          vector<int> factorize factorial(int n) {
   for (int i = 2*p; i<=n; i += p) phi[i] =
                                                          vector<int> factorization(primes.size(), 0);
     (phi[i]/p) * (p-1);
                                                          for(int i=0;i<primes.size() && primes[i] <=
                                                          \rightarrow n;i++) {
                                                          factorization[i] = legendre(n, primes[i]);
                                                          return factorization;
Factorials
// digits in factorial
                                                          // same thing but for C(n,k)
#define kamenetsky(n) (floor((n * log10(n /
                                                          vector<int> factorize_binom(int n, int k) {
 \hookrightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
                                                          vector<int> factorization(primes.size(), 0);
for(int i=0;i<primes.size() && primes[i] <=</pre>
// approximation of factorial
#define stirling(n) ((n == 1) ? 1 : sqrt(2 *

    n;i++) {

 \hookrightarrow M PI * n) * pow(n / M E, n))
                                                           factorization[i] = legendre(n, primes[i]) -
// natural log of factorial #define lfactorial(n) (lgamma(n+1))
                                                             legendre(k, primes[i]) - legendre(n-k,
                                                             primes[i]):
                                                         } return factorization;
Prime Factorization
// do not call directly
ll pollard_rho(ll n, ll s) {
                                                         Farey Fractions
 .11 x, y;
  x = y = rand() \% (n - 1) + 1;
                                                             generate 0 <= a/b <= 1 ordered, b <= n farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
 int head = 1, tail = 2;
while (true) {
    x = mult(x, x, n);
    x = (x + s) % n;
                                                         // Jureg(x) - 0/1 1/4 1/3 1/2 2/3 3/4 1/1
// length is sum of phi(i) for i = 1 to n
vector<pair<int, int>> farey(int n) {
                                                          int h = 0, k = 1, x = 1, y = 0, r;
vector<pair<int, int>> v;
  if (x == y) return n;
  11 d = _gcd(max(x - y, y - x), n);
if (1 < d && d < n) return d;
                                                           v.push_back({h, k});
  if (++head == tail) v = x, tail <<= 1:
                                                           r = (n-y)/k;
                                                           y += r*k; x += r*h;
swap(x,h); swap(y,k);
// call for prime factors
                                                          x = -x; y = -y;

} while (k > 1);

v.push_back({1, 1});
void factorize(ll n, vector<ll> &divisor) {
 if (n == 1) return:
 if (isPrime(n)) divisor.push back(n);
                                                          return v;
 while (d >= n) d = pollard_rho(n, rand() % (n Fast Fourier Transform
 → - 1) + 1);
factorize(n / d, divisor);
                                                          const double PI = acos(-1):
                                                         void fft(vector<cd>& a, bool invert) {
  factorize(d, divisor);
                                                          int n = a.size():
                                                          for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1;
                                                           for (; j & bit; bit >>= 1) j ^= bit;
j ^= bit;
Factorize Factorials
   NOTE: count distinct divisors of n bu
                                                           if (i < j) swap(a[i], a[j]);
 // computing (q1+1)*(q2+1)*...*(qk+1)
// where gi are powers of primes pi dividing n
                                                          for (int len = 2: len <= n: len <<= 1) {
// use that and this code to solve
                                                           double ang = 2 * PI / len * (invert ? -1 :
    https://open.kattis.com/problems/divisors
                                                             1):
 // max power of a prime p dividing n!
                                                           cd wlen(cos(ang), sin(ang));
 // O(log(n))
                                                           for (int i = 0; i < n; i += len) {
int legendre(int n, int p) {
 int mx = 0;
                                                            for (int j = 0; j < len / 2; j++) {
  while(n>0) n/=p, mx+=n;
                                                           cd u = a[i+j], v = a[i+j+len/2] * w;
 return mx;
                                                            a[i+j] = u + v;
                                                             a[i+j+len/2] = u - v;
bitset<10000> sieve:
vector<int> primes;
                                                            ..w *= wlen:
// get all primes O(n log n)
// if dealing with small numbers
void genPrimes(int n) {
                                                          if (invert)
 sieve[0] = sieve[1] = 1;
                                                           for (auto& x : a)
 primes.push_back(2);
                                                           x /= n:
 for (int i = 3; i <= n; i+=2)
  if (i%2 != 0 && !sieve[i]) {
    primes.push_back(i);</pre>
                                                         vector<int> fftmult(vector<int> const& a,

  vector<int> const& b) {
  vector<cd> fa(a.begin(), a.end()),
    for (int j = i * 3; j <= n; j += i*2)
     sieve[j] = 1;

    fb(b.begin(), b.end());
```

```
int n = 1 << (32 - __builtin_clz(a.size() +</pre>
                                                    |\#define\ mult(a,b,m)\ ((ull)a*b\%m)

    b.size() - 1));
fa.resize(n); fb.resize(n);
                                                     #define msub(a,b,m) (a-b+((a < b)?m:0))
                                                    ll mpow(ll b, ll e, ll m) {
                                                     11 x = 1;
fft(fa, false); fft(fb, false);
                                                     while (e > 0) {
if (e % 2) x = (x * b) % m;
for (int i = 0; i < n; i++) fa[i] *= fb[i];
fft(fa, true);
                                                      b = (b * b) \% m;
 vector<int> toret(n):
                                                      e /= 2;
for (int i = 0; i < n; i++) toret[i] =

→ round(fa[i].real());
                                                     return x % m;
return toret;
                                                    ull mfac(ull n, ull m) {
  ull f = 1;
  for (int i = n; i > 1; i--)
Greatest Common Denominator
ll egcd(ll a, ll b, ll& x, ll& y) {
                                                      f = (f * i) \% m;
if (b == 0) { x = 1; y = 0; return a; }
                                                     return f;
11 gcd = egcd(b, a \% b, x, y);
x = a / b * v;
                                                     // if m is not guaranteed to be prime
swap(x, y);
                                                    ll minv(ll b, ll m) {
return gcd;
                                                     \overline{11} x = 0, y = 0;
                                                     if (egcd(b, m, x, y) != 1) return -1;
                                                     return (x % m + m) % m:
Kth Root (floor)
                                                    11 mdiv_compmod(int a, int b, int m) {
struct KthRoot {
                                                     if (__gcd(b, m) != 1) return -1;
vector<ull> pow[65]; // pow[k][x] =
                                                     return mult(a, minv(b, m), m);
\rightarrow pow(x+2,k) (k >= 4)
KthRoot() {
                                                     // if m is prime (like 10^9+7)
 for (ull t = 2; t < (1<<16); t++) {
  ull s = t*t; s = s*s;
for (int k = 4; ; k++) {
                                                    11 mdiv_primemod (int a, int b, int m) {
                                                     return mult(a, mpow(b, m-2, m), m);
 pow[k].push_back(s);
    if (__builtin_umulll_overflow(s,t.&s))
                                                     // tonelli shanks = sqrt(n) % m, m is prime

→ break:

                                                    ll legendre(ll a, ll m){
                                                     if (a % m==0) return 0;
if (m == 2) return 1;
                                                     return mpow(a, (m-1)/2, m);
 ull sqrt(ull n) const {
 if (n == -1ull) return (unsigned int)(-1);
                                                    11 msqrt(11 n, 11 m) {
 ull x = std::sqrt((double)n);
                                                     ll s = builtin ctzll(m-1), q = (m-111)>>s,
 return x*x > n? x-1 : x;
                                                        z = rand()\%(m-1)+1:
                                                     if (m == 2) return 1;
if (s == 1) return mpow(n,(m+1)/411,m);
ull cbrt(ull n) const {
ull x = 0, y = 0;
                                                     while (legendre(z,m)!=m-1) z = rand()\%(m-1)+1;
 for (int s = 63; s >= 0; s -= 3) {
                                                     11 c = mpow(z,q,m), r = mpow(n,(q+1)/2,m), t
  y = 3*x*(x+1)+1:
                                                     \rightarrow = mpow(n,q,m), M = s;
  if (y \le (n>>s)) n = y << s, x++;
                                                     while (t != 1){
                                                      "11 i=1, ts = (t * t) % m;
                                                      while (ts != 1) i++, ts = (ts * ts) % m; ll b = c;
 return x:
 // returns floor(n^(1/k)), k \ge 1
                                                      for (int j = 0; j < M-i-1; j++) b = (b * b) % Matrix Subarray Sums
 ull operator()(ull n, int k) {
                                                      r = r * b \% m; c = b * b \% m; t = t * c \% m;
 if (k == 1 \mid | n == 0) return n;
 if (k == 2) return sqrt(n);
                                                     \rightarrow M = i;
 if (k == 3) return cbrt(n);
                                                     return r:
 auto ub = upper_bound(pow[k].begin().
\rightarrow pow[k].end(), n);
                                                     Modulo Tetration
 return (ub-pow[k].begin())+1;
                                                    ll tetraloop(ll a, ll b, ll m) {
                                                     if(b == 0 | | a == 1) return 1;
                                                     ll w = tetraloop(a,b-1,phi(m)), r = 1;
Josephus Problem
                                                     for (;w;w/=2) {
                                                      -if (w&1) {
    r *= a; if (r >= m) r -= (r/m-1)*m;
    }
// 0-indexed, arbitrary k
int josephus(int n, int k) {
if (n == 1) return 0;
if (k == 1) return n-1;
                                                      a *= a; if (a >= m) a -= (a/m-1)*m;
if (k > n) return (joséphus(n-1,k)+k)%n;
                                                     return r;
int res = josephus(n-n/k,k)-n\%k;
                                                    int tetration(int a, int b, int m) {
  if (a == 0 || m == 1) return ((b+1)&1)%m;
return res + ((res<0)?n:res/(k-1));
\frac{1}{2} fast case if k=2, traditional josephus
                                                     return tetraloop(a,b,m) % m;
int josephus(int n) {
return 2*(n-(1<<(32-__builtin_clz(n)-1)));
                                                    Matrix
                                                    template<typename T>
                                                    struct Mat : public Vec<2, T> {
Least Common Multiple
                                                     int w, h;
#define lcm(a,b) ((a*b)/__gcd(a,b))
                                                     Mat(int x, int y) : Vec<2, T>(x, y), w(x),
                                                     \rightarrow h(y) {}
Modulo Operations
                                                     static Mat<T> identity(int n) { Mat<T> m(n,n);
#define MOD 100000007
                                                        for (int i=0;i<n;i++) m[i][i] = 1; return
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0))
                                                     \stackrel{\hookrightarrow}{\rightarrow} m; }
```

```
Mat<T>& operator-=(const Mat<T>& m) {
   for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
(*this)[i][j] -= m[i][j];
   return *this;
  Mat<T> operator*(const Mat<T>& m) {
   Mat < T > z(w.m.h):
   for (int i = 0; i < w; i++)

for (int j = 0; j < h; j++)

for (int k = 0; k < m.h; k++)

z[i][k] += (*this)[i][j] * m[j][k];
      return z:
  Mat<T> operator+(const Mat<T>& m) { Mat<T>
  \rightarrow a=*this; return a+=m; }
  Mat<T> operator-(const Mat<T>& m) { Mat<T>
   → a=*this; return a-=m; }
  Mat<T>& operator*=(const Mat<T>& m) { return
  *this = (*this)*m; }
  Mat<T> power(int n) {
   Mat<T> a = Mat<T>::identity(w), m=*this;
   for (:n:n/=2.m*=m) if (n\&1) a *=m:
   return a;
 Matrix Exponentiation
// F(n) = c[0]*F(n-1) + c[1]*F(n-2) + ...
// b is the base cases of same length c
ll matrix_exponentiation(ll n, vector<ll> c,
  \rightarrow vector<11> b) {
  if (nth < b.size()) return b[nth-1];
  Mat<11> a(c.size(), c.size()); 11 s = 0;
for (int i = 0; i < c.size(); i++) a[i][0] =
  → c[i];
  for (int i = 0; i < c.size() - 1; i++)
 \hookrightarrow a[i][i+1] = 1;
 a = a.power(nth - c.size());
for (int i = 0; i < c.size(); i++)
    s += a[i][0] * b[i];</pre>
  return s;
 template<class T> struct MatrixSum {
   Vec<2. T> p:
  MatrixSum(Vec<2, T>& v) {
   p = Vec<2,T>(v.size()+1, v[0].size()+1);
   for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[0].size(); j++)
p[i+1][j+1] = v[i][j] + p[i][j+1] +
      p[i+1][i] - p[i][i];
  f sum(int u, int 1, int d, int r) {
    return p[d][r] - p[d][1] - p[u][r] + p[u][1];
Mobius Function
const int MAXN = 10000000;
// mu[n] = 0 iff n has no square factors
 // 1 = even number prime factors, -1 = odd
 short mu[MAXN] = \{0,1\};
void mobius(){
  for (int i = 1; i < MAXN; i++)
    if (mu[i])</pre>
    for(int'j = i + i; j < MAXN; j += i)
      mu[j] -= mu[i];
 Minimum Excluded
int mex(set<int>& s) {
  auto i = s.begin(); int v = 0;
  while (i != s.end() && *i == val) i++, v++;
return v;
```

Mat<T>& operator+=(const Mat<T>& m) {

for (int i = 0; i < w; i++)

return *this;

for (int j = 0; j < h; j++)
 (*this)[i][j] += m[i][j];

```
#define nimAdd(a,b) ((a)^(b))
ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
    M[0]^{(M[0] << 16)}, M[1]^{(M[1] << 8)},
    M[2]^{(M[2]<<4)}, M[3]^{(M[3]<<2)},
if (i-- == 0) return a&b;
  int k=1<<i;
  ull s=nimMúl(a,b,i), m=M[5-i],
    t=nimMul(((a^(a>>k))&m)|(s&~m),
    ((b^(b>>k))&m)|(m&(\sim m>>1))<< k, i);
  return ((s^t)\&m)<\langle k|((s^(t)>k))\&m);
Permutation
// c = array size. n = nth perm. return index
vector<int> gen_permutation(int c, int n) {
 vector<int> idx(c), per(c), fac(c); int i;
 for (i = 0; i < c; i++) idx[i] = i;
for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i;
 for (i = c - 1; i >= 0; i--)

per[c-i-1] = idx[fac[i]],
  idx.erase(idx.begin() + fac[i]);
 return per:
// get what nth permutation of vector
int get_permutation(vector<int>& v) {
 int use = 0, i = 1, r = 0;
for (int e: v) {
   r = r * i++ + __builtin_popcount(use &

    -(1<<e));</pre>
 use |= 1 << e;
 return r;
\mathbf{Permutation} (string/multiset)
string freq2str(vector<int>& v) {
 string s;
for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[i]; j++)
LS += (char)(i + 'A');
return s;
// nth perm of multiset, n is 0-indexed
string gen_permutation(string s, 11 n) {
 vector<int> freq(26, 0);
 for (auto e : s) freq[e - 'A']++;
 for (int i = 0; i < \frac{1}{26}; i++) if (freq[i] > 0)
  freq[i]--; ll v = multinomial(freq);
 if (n < v) return (char)(i+'A') +
 → gen_permutation(freq2str(freq), n);
 freq[i]++; n -= v;
return "":
Miller-Rabin Primality Test
// Miller-Rabin primality test - O(10 log^3 n)
bool isPrime(ull n) {
  if (n < 2) return false;
if (n < 2) return false;
if (n == 2) return true;
if (n % 2 == 0) return false;
ull s = n - 1;
while (s % 2 == 0) s /= 2;
for (int i = 0; i < 10; i++) {</pre>
  ull temp = s;
  ull a = rand() \% (n - 1) + 1;
  ull mod = mpow(a, temp, n);
  while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
   mod = mult(mod, mod, n);
   temp *= 2;
  if (mod!=n-1&&temp%2==0) return false;
.}
.return true:
```

Nimber Arithmetic

```
Sieve of Eratosthenes
bitset<100000001> sieve:
// generate sieve - O(n log n)
void genSieve(int n) {
sieve[0] = sieve[1] = 1;
for (ull i = 3; i * i < n; i += 2)
if (!sieve[i])
 ...for (ull j = i * 3; j <= n; j += i * 2)
...sieve[j] = 1;</pre>
// query sieve after it's generated - O(1)
bool querySieve(int n) {
return n == 2 | | (n \% 2 != 0 \&\& !sieve[n]):
Compile-time Prime Sieve
const int MAXN = 100000:
template<int N>
struct Sieve
 bool sieve[N]:
 constexpr Sieve() : sieve() {
 sieve[0] = sieve[1] = 1;
for (int i = 2; i * i < N; i++)
   if (!sieve[i])</pre>
  for (int j = i * 2; j < N; j += i)
...sieve[j] = 1;
bool isPrime(int n) {
static constexpr Sieve<MAXN> s;
return !s.sieve[n]:
Simpson's / Approximate Integrals
// integrate f from a to b, k iterations // error <= (b-a)/18.0 * M * ((b-a)/2k)^4
// where M = max(abs(f^{**}(x))) for x in [a,b] // "f" is a function "double functionbulle function"
double Simpsons (double a, double b, int k,

    double (*f)(double)) {

 double dx = (b-a)/(2.0*k), t = 0;
for (int i = 0; i < k; i++)

t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
\leftrightarrow (*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector < double > solve Eq (double a, double b,

    double c) {
    vector<double> r:
}
 double z = b * b - 4 * a * c;
 if (z == 0)
 r.\widetilde{push}_{back(-b/(2*a))};
 else if (z > 0) {
 r.push back((sgrt(z)-b)/(2*a));
  r.push_back((sqrt(z)+b)/(2*a));
 return r;
\frac{1}{1/2}ax^3 + bx^2 + cx + d = 0, find x
vector < double > solve Eq (double a, double b,

    double c, double d) {
    vector < double > res;

 long double a1 = b/a, a2 = c/a, a3 = d/a;
long double q = (a1*a1 - 3*a2)/9.0, sq =
\rightarrow -2*sqrt(q);
long double r = (2*a1*a1*a1 - 9*a1*a2 +
\rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z <= 0) {
 theta = acos(r/sqrt(q*q*q));
  res.push_back(sq*cos(theta/3.0) - a1/3.0);
 res.push_back(sq*cos((theta+2.0*PI)/3.0) -
res.push_back(sq*cos((theta+4.0*PI)/3.0) -
\rightarrow a1/3.0);
 else {
```

```
res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
  res[0] = (res[0] + q / res[0]) *
   ((r<0)?1:-1) - a1 / 3.0;
 return res:
// linear diophantine equation ax + by = c,
\rightarrow find x and y
// infinite solutions of form x+k*b/g, y-k*a/g bool solveEq(11 a, 11 b, 11 c, 11 &x, 11 &y, 11
 g = egcd(abs(a), abs(b), x, y);
 if (c % g) return false;
x *= c / g * ((a < 0) ? -1 : 1);
y *= c / g * ((b < 0) ? -1 : 1);
return true;
 '/ m = \# equations, n = \# variables, a[m][n+1]
\hookrightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
\rightarrow a[i][n+1]
// find a solution of some kind to linear
\rightarrow equation
const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a >
vector<double> solveEq(double **a, int m, int
if (!zero(a[j][i])) {
    if (j != cur) swap(a[j], a[cur]);
    for (int sat = 0; sat < m; sat++) {
   if (sat == cur) continue;
      double num = a[sat][i] / a[cur][i];
     for (int sot = 0; sot <= n; sot++)
.a[sat][sot] -= a[cur][sot] * num;
     }
    break:
 for (int j = cur; j < m; j++)
  if (!zero(a[j][n])) return vector<double>();
 vector<double> ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
    if (sat < m && !zero(a[sat][i]))
    ans[i] = a[sat][n] / a[sat++][i];
    return ans;
   solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 is no
→ solutions
// `alls` is whether to find all solutions, or
\hookrightarrow any
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&
 \rightarrow b, Vec<1, double>& x, bool alls=false) {
 int n = A.size(), m = x.size(), rank = 0, br,
 vector<int> col(m); iota(begin(col), end(col)
if (bv <= eps) {
   for(int j = i; j < n; j++)
if (fabs(b[j]) > eps)
     return -1:
   break;
  swap(A[i], A[br]);
swap(b[i], b[br]);
  swap(col[i], col[bc]);
  for(int j = 0; j < n; j++)
```

```
swap(A[j][i], A[j][bc]);
  by = 1.0 / A[i][i];
for(int j = (alls)?0:i+1; j < n; j++) {
   .if (j != i) {
    double fac = A[j][i] * bv;
    b[j] = fac * b[i];
    for(int k = i+1; k < m; k++)
A[j][k] -= fac*A[i][k];
  rank++:
 if (alls) for (int i = 0; i < m; i++) x[i] =
 \rightarrow -DBL_MAX;
 for (int i = rank; i--;) {
  bool isGood = true;
  if (alls)
  for (int j = rank; isGood && j < m; j++)
if (fabs(A[i][j]) > eps)
  isGood = false;
b[i] /= A[i][i];
  if (isGood) x[col[i]] = b[i];
 if (!alls)

for(int j = 0; j < i; j++)

b[j] -= A[j][i] * b[i];
 return rank:
Graycode Conversions
ull graycode2ull(ull n) {
 for (; n; n = n >> 1) i ^= n;
 return i;
ull ull2graycode(ull n) {
return n ^ (n >> 1);
Date Utilities
// handles -4799-01-01 to 1465001-12-31
int date2int(int y, int m, int d){
return 1461*(y+4800+(m-14)/12)/4+367*(m-2-(m-14)/12)
   -14)/12*12)/12-3*((y+4900+(m-14)/12)/100)
pair<int,pair<int,int>> int2date(int x){
int n,i,j;
 x+=68569
 n=4*x/146097:
x-=(146097*n+3)/4;
i=(4000*(x+1))/1461001;
 x=1461*i/4-31;
 j=80*x/2447;
 return \{100*(n-49)+i+j/11, \{j+2-12*(j/11), \}
\rightarrow x-2447*i/80}}:
int dayOfWeek(int y, int m, int d){ //0=sunday
static int cal[]={0,3,2,5,0,3,5,1,4,6,2,4};
return (y+y/4-y/100+y/400+cal[m-1]+d)\%7;
Unix/Epoch Time
// O-indexed month/time, 1-indexed day
// minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
hour, int minute, int second) {
struct tm t; time_t epoch;
t.tm_year = year - 1900; t.tm_mon = month;
 t.tm_mday = day; t.tm_hour = hour;
t.tm_min = minute; t.tm_sec = second;
t.tm_isdst = 0; // 1 = daylights savings
 epoch = mktime(&t);
 return (ull)epoch;
vector<int> toDate(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return {t.tm_year+1900,t.tm_mon,t.tm_mday,t
```

```
| int getWeekday(ull epoch) {
| time_t e=epoch; struct tm t=*localtime(&e);
| return t.tm_wday; // 0-6, 0 = sunday
| int getDayofYear(ull epoch) {
| time_t e=epoch; struct tm t=*localtime(&e);
| return t.tm_yday; // 0-365
| const int months[] =
| \int \{31,28,31,30,31,30,31,30,31,30,31\};
| bool validDate(int year, int month, int day) {
| bool leap = !(year%(year%25?4:16));
| if (month >= 12) return false;
| return day <= months[month] + (leap &&
| \int month == 1);
| }
| </pre>
```

Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n + 1)(!(n - 1) + !(n - 2)), $!n = n! \sum_{i=0}^{n} \frac{(-1)^{i}}{i!}$, $!n = [\frac{n!}{e}]$ In a partially ordered set, a chain is a subset of

In a partially ordered set, a chain is a subset of elements that are all comparable to eachother. An antichain is a subset where no two are comparable.

Dilworth's theorem states the size of a maximal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the *nth* prime number is greater than n * ln(n) for n > 1.

Nicomachi's Theorem states $1^3 + 2^3 + ... + n^3 = (1 + 2 + ... + n)^2$ and is equivalent to $(n\frac{n+1}{2})^2$.

Lagrange's Four Square Theorem states every natural number is the sum of the squares of four non-negative integers. This is a special case of the Fermat Polygonal Number Theorem where every positive integer is a sum of at most n s-gonal numbers. The nth s-gonal number $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$

7 Graphs

```
bitset<20001> marked; // change size as needed 2D Grid Shortcut
 void make_set(int _p) { p=_p, rank=0, sz=1; }
                                                      if(edges.size()==0) return false;
                                                       stack<int> st;
                                                                                                          marked[0] = 1;
                                                                                                                                                              for (edge ep : graph[0]) pq.push(ep); while(MST.size()!=graph.size()-1 &&
                                                       int a[] = \{-1, -1\}
                                                                                                                                                              \#define\ fordir(x,y,n,m)\ for(auto[dx,dy]:dir)if
BFS
                                                       for(int v=0; v<n; v++) {
  if(indeg[v]!=graph[v].size()) {</pre>
                                                                                                                                                                 (inbound(x+dx,n)&Binbound(y+dy,m))
// adjacency list named 'graph'
int visited[MAX];
                                                                                                             pq.size()!=0) {
                                                                                                                                                              const pair<int,int> dir[] =
                                                         bool b = indeg[v] > graph[v].size();
                                                                                                           edge e = pq.top(); pq.pop();
                                                         if (abs(((int)indeg[v])-((int)graph[v]
                                                                                                                                                                 \{\{1,0\},\{0,1\},\{-1,0\},\{0,-1\}\};
int parent[MAX];
                                                                                                           int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
int vc = 0:
                                                                                                                                                                  2D Geometry
                                                         .size())) > 1) return
vector<int> bfs(int start, int end) {
                                                                                                           else if(marked[u]) swap(u, v);
                                                         false
                                                                                                                                                              #define point complex<double>
                                                         if (a[b] != -1) return false;
 visited[start] = vc;
                                                                                                           for(edge ep : graph[u]) pq.push(ep);
                                                                                                                                                              #define EPS 0.0000001
                                                        _a[b] = v;
                                                                                                           marked[u] = 1
 parent[start] = -1;
                                                                                                                                                              #define sq(a) ((a)*(a))
                                                                                                           MST.push_back(e);
 queue<int> q;
                                                                                                                                                              #define c\bar{b}(a) ((a)*(a)*(a))
                                                                                                           total += e.w;
 q.push(start);
                                                                                                                                                              double dot(point a, point b) { return
                                                      int s = (a[0]!=-1 \&\& a[1]!=-1 ? a[0] :
 while (!q.empty()) {

    real(conj(a)*b); }

                                                         (a[0]=-1 \&\& a[1]=-1 ? edges[0].u : -1));
                                                                                                          return MST:
  int v = q.front(); q.pop();
                                                                                                                                                              double cross(point a, point b) { return
                                                       if(s==-1) return false;
  for (auto e : graph[v]) {
                                                       while(!st.empty() || !graph[s].empty()) {
                                                                                                                                                              \rightarrow imag(conj(a)*b); }
  if (visited[e] != vc) {
   visited[e] = vc;
                                                                                                                                                             struct line { point a, b; };
struct circle { point c; double r; };
struct segment { point a, b; };
                                                       if (graph[s].empty()) {
                                                                                                          Union Find
                                                         circuit.push_back(s); s = st.top();
                                                                                                         int uf find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
    q.push(e);
                                                        st.pop(); }
    parent[e] = v;
                                                                                                          return s[i].p;
                                                                                                                                                              struct triangle { point a, b, c; };
    if (e == end) goto DONE;
                                                        else {
                                                                                                                                                              struct rectangle { point tl, br; };
                                                         int w = edges[*graph[s].begin()].v;
                                                                                                         void uf_union(subset* s, int x, int y) {
                                                         graph[s].erase(graph[s].begin());
                                                                                                                                                              struct convex_polygon {
                                                                                                          int xp = uf_find(s, x), yp = uf_find(s, y);
                                                                                                                                                               vector<point points;
                                                        st.push(s); s = w;
                                                                                                          if (s[xp].rank > s[yp].rank) s[yp].p = xp,
                                                                                                                                                               convex_polygon(vector<point> points) :
 // path reconstruction
if (visited[end] != vc) return {};
                                                                                                          \Rightarrow s[xp].sz += s[yp].sz;
                                                                                                                                                                 points(points) {}
                                                       circuit.push_back(s);
                                                                                                          else if (s[xp].rank < s[yp].rank) s[xp].p =
                                                                                                                                                               convex_polygon(triangle a) {
 vector<int> path;
                                                       return circuit.size()-1==edges.size();
                                                                                                          \rightarrow yp, s[yp].sz += s[xp].sz;
                                                                                                                                                               points.push_back(a.a); points.push_back(a.b);
 for (int v = end; v != -1; v = parent[v])
                                                                                                          else s[yp].p = xp, s[xp].rank++, s[xp].sz +=
  path.push_back(v);
                                                                                                                                                                 points.push_back(a.c);
                                                                                                             s[vp].sz;
 return path;
                                                                                                                                                               convex_polygon(rectangle a) {
                                                     Flovd Warshall
                                                                                                         void uf_size(subset *s, int i) {
  return s[uf_find(s, i)].sz;
                                                                                                                                                               points.push_back(a.tl);
                                                     const ll inf = 1LL << 62;</pre>
Dijkstra's
                                                     #define \overline{FOR(i,n)} for (int \ i = 0; \ i < n; \ i++)
                                                                                                                                                                  points.push_back({real(a.tl),
const int inf = 20000001; // change as needed
                                                     void floydWarshall(Vec<2, 11>& m) {
                                                                                                                                                                 imag(a.br)});
// use add_edge(..., true) for digraphs
                                                     int n = m.size();
FOR(i,n) m[i][i] = min(m[i][i], OLL);
FOR(k,n) FOR(i,n) FOR(j,n) if (m[i][k] != inf
                                                                                                                                                               points.push_back(a.br);
                                                                                                         Bipartite Graph
void add_edge(Vec<2, edge> &graph, int u, int
                                                                                                                                                                  points.push_back({real(a.br),
→ v, int w, bool directed=true) {
                                                                                                                                                                  imag(a.tl)});
                                                                                                         A bipartite graph has "left" and "right" set of
 graph[u].push_back({u,v,w});
                                                        && m[k][j] != inf)
 if(!directed) graph[v].push_back({v,u,w});
                                                                                                         → nodes
Every edge has an endpoint in each set (L/R)
                                                      auto newDist = max(m[i][k] + m[k][j], -inf);
                                                      m[i][j] = min(m[i][j], newDist);
                                                                                                                                                              struct polygon {
                                                                                                         A matching is a subset of all edges
vector<int> dijkstra(Vec<2, edge> &graph, int
                                                                                                                                                              vector<point> points;
                                                                                                         Such that each vertex is an endpoint
 FOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n)
                                                                                                                                                               polygon(vector<point> points) :
                                                                                                         Of at most one edge in the subset
 vector<int> D(graph.size(), inf);
                                                      if (m[i][k] != inf && m[k][j] != inf)
                                                                                                         sqrt(V)*E time
                                                                                                                                                                 points(points) {}
 priority_queue<edge, vector<edge>,
                                                        m[i][j] = -inf;
                                                                                                                                                               polygon(triangle a) {
                                                                                                         tested on "piano lessons"

    greater < edge >> pq;

                                                                                                         sourced from
                                                                                                                                                               points.push_back(a.a); points.push_back(a.b);
 pq.push({src,src,0});
D[src]=0;
                                                                                                             https://codeforces.com/blog/entry/58048
                                                                                                                                                                 points.push_back(a.c);
                                                     Bellman Ford
 while(!pq.empty()) {
                                                     const int inf = 20000001:
                                                                                                         #define MAXNODES 1001
                                                                                                                                                               polygon(rectangle a) {
  edge e = pq.top(); pq.pop();
int v = e.v;
                                                     vector<11> bellman_ford(vector<edge> edges, int bitset<MAXNODES> V;
                                                                                                                                                               points.push_back(a.tl);
                                                                                                         bool match(int node, Vec<2,int> &G, vector<int>

ightarrow src, {	t int} V) {
  for(int i=0;i<graph[v].size();i++) {
                                                                                                                                                                  points.push_back({real(a.tl),
                                                     vector (11> D(V, inf);
                                                                                                          int u = graph[v][i].v;
                                                                                                                                                                 imag(a.br)});
                                                     D[src] = 0;
for (int i=1:i<=V-1:i++)
                                                                                                          if (V[node]) return false;
   if(D[v] + graph[v][i].w < D[u]) {
  D[u] = D[v] + graph[v][i].w;
                                                                                                                                                               points.push_back(a.br);
                                                                                                          V[node] = 1;
                                                      for (edge e : edges)
                                                                                                          for(auto vec : G[node]) {
  if (R[vec] == -1 || match(R[vec], G, R, L))
                                                                                                                                                                  points.push_back({real(a.br),
    pq.push({src,u,D[u]});
                                                       if (D[e.u] != inf \&\& D[e.u] + e.w < D[e.v])
                                                                                                                                                                  imag(a.tl)});
                                                      D[e.v] = D[e.u] + e.w;

// detect negative cycles: *typically* 2 is as
                                                                                                           L[node] = vec; R[vec] = node;
                                                                                                                                                               polygon(convex_polygon a) {
return D;
                                                        good as V-1 for this
                                                                                                            return true;
                                                                                                                                                               for (point v : a.points)
                                                     for (int i=1;i<=V-1;i++)
                                                                                                                                                                 points.push_back(v);
                                                      for (edge e : edges)
                                                                                                          return false:
                                                       if[D[e.u] = \inf \&\& D[e.u] + e.w < D[e.v])
Eulerian Path
                                                     D[è.v]
                                                                 = -inf:
#define edge_list vector<edge>
#define adj sets vector<set<int>>>
                                                                                                         vector<pair<<u>int</u>, <u>int</u>>> bipartite_match(Vec<2,
                                                                                                                                                              // triangle methods
                                                                                                                                                              double area_heron(double a, double b, double
                                                                                                           \rightarrow int> &G, int m) {
struct EulerPathGraph {
                                                                                                          vector<int> L(G.size(), -1), R(m, -1);
                                                                                                                                                                 c) {
  (a < b) swap(a, b);</pre>
 adj sets graph; // actually indexes incident
                                                                                                          V.reset();
bool running = true;
                                                     Minimum Spanning Tree
                                                                                                                                                               if (a < c) swap(a, c);
                                                      / returns vector of edges in the mst
 edge_list edges; int n; vector<int> indeg;
                                                                                                          while (running) {
                                                                                                                                                              if (b < c) swap(b, c);
                                                     // graph[i] = vector of edges incident to
 EulerPathGraph(int n): n(n) {
                                                                                                           running = false;
                                                                                                                                                              if (a > b + c) return -1:
                                                                                                           V.reset();
                                                                                                                                                              return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
  indeg = *(new vector<int>(n,0));
                                                        places total weight of the mst in Stotal
                                                                                                           for (int i=0;i<L.size();i++)
if (L[i] == -1)
  graph = *(new adj_sets(n, set<int>()));
                                                                                                                                                                 /16.0);
                                                     // if returned vector has size != n-1. there is
                                                                                                             running |= match(i, G, R, L);
 void add_edge(int u, int v) {
                                                                                                                                                              // seament methods
                                                     vector<edge> mst(Vec<2, edge> graph, 11
  graph[u].insert(edges.size());
                                                                                                                                                             double lengthsq(segment a) { return
                                                       &total) {
                                                                                                           vector<pair<int,int>> ret;
  indeg[v]++;
                                                                                                          for (int i = 0; i < L.size(); ++i)
if(L[i]!=-1) ret.push_back({i, L[i]});</pre>
                                                                                                                                                                  sq(real(a.a) - real(a.b)) + sq(imag(a.a)
  edges.push_back(edge(u,v,0));
                                                     priority_queue<edge, vector<edge>,
                                                                                                                                                                 imag(a.b)); }
                                                        greater<edge>> pq;
                                                                                                                                                              double length(segment a) { return
                                                      vector<edge> MST;
 bool eulerian path(vector<int> &circuit) {
                                                                                                                                                                 sqrt(lengthsq(a)); }
```

```
// circle methods double circumference(circle a) { return 2 * a.r } // -1 outside, 0 inside, 1 tangent, 2

    intersection
int intersection(circle a, circle b,
double area(circle a) { return sq(a.r) * M_PI;
\stackrel{
ightharpoonup}{
ightharpoonup} } // rectangle methods
                                                       vector<point>& inter) {
                                                    double d2 = norm(b.c - a.c), rS = a.r + b.r,
                                                    \rightarrow rD = a.r - b.r;
double width(rectangle a) { return
                                                    if (d2 > sq(rS)) return -1;

→ abs(real(a.br) - real(a.tl)); }

                                                    if (d2 < sq(rD)) return 0;
double height(rectangle a) { return
                                                    double ca = 0.5 * (1 + rS * rD / d2);
point z = point(ca, sqrt(sq(a.r) / d2 -

→ abs(imag(a.br) - real(a.tl)); }

double diagonal(rectangle a) { return
                                                    \rightarrow sq(ca)):

    sqrt(sq(width(a)) + sq(height(a))); }

                                                    inter.push_back(a.c + (b.c - a.c) * z);
double area (rectangle a) { return width(a) *
                                                    if (abs(imag(z)) > EPS) inter.push_back(a.c +
                                                    \rightarrow (b.c - a.c) * coni(z)):
double perimeter(rectangle a) { return 2 *
                                                    return inter.size();
   (width(a) + height(a)); }
// check if `a` fit's inside `b
                                                    // points of intersection
// swap equalities to exclude tight fits
                                                    vector<point> intersection(line a, circle c) {
bool doesFitInside(rectangle a, rectangle b) {
                                                    vector<point> inter;
c.c -= a.a;
a.b -= a.a;
int x = width(a), w = width(b), y = height(a),

→ h = height(b):
                                                    point m = a.b * real(c.c / a.b);
if (x > y) swap(x, y);
if (w > h) swap(w, h);
                                                    double d2 = norm(m - c.c);
                                                    if (d2 > sq(c.r)) return 0;
if (w < x) return false;
                                                     double l = sqrt((sq(c.r) - d2) / norm(a.b));
 if (y <= h) return true;
                                                    inter.push back(a.a + m + 1 * a.b);
double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
                                                    if (abs(1) > EPS) inter.push back(a.a + m
return sq(a) \le sq(b) + sq(c);
                                                    \rightarrow * a.b);
                                                    return inter:
// polygon methods
// negative area = CCW, positive = CW
                                                    // area of intersection
double area(polygon a) {
                                                    double intersection(rectangle a, rectangle b) {
double area = 0.0; int n = a.points.size();
                                                    double x1 = max(real(a.tl), real(b.tl)), y1 =
for (int i = 0, j = 1; i < n; i++, j = (j +
                                                     → max(imag(a.tl), imag(b.tl));
double x2 = min(real(a.br), real(b.br)), y2 =
 area += (real(a.points[j]-a.points[i]))*
                                                     → min(imag(a.br), imag(b.br));
   (imag(a.points[j]+a.points[i]));
                                                    return (x2 \le x1 \mid y2 \le y1)? 0:
return area / 2.0;
                                                    \rightarrow (x2-x1)*(y2-y1);
// get both unsigned area and centroid
pair < double, point > area_centroid(polygon a) { Convex Hull
int n = a.points.size():
                                                    bool cmp(point a, point b) {
double area = 0;
                                                    if (abs(real(a) - real(b)) > EPS) return
 point c(0, 0);
                                                       real(a) < real(b);
 for (int i = n - 1, j = 0; j < n; i = j++) {
                                                    if (abs(imag(a) - imag(b)) > EPS) return
 double v = cross(a.points[i], a.points[j]) /
                                                    \rightarrow imag(a) \stackrel{\checkmark}{\circ} imag(b);
                                                    return false:
 areá += v:
 c += (a.points[i] + a.points[j]) * (v / 3);
                                                    convex_polygon convexhull(polygon a) {
                                                    sort(a.points.begin(), a.points.end(), cmp);
c /= area:
                                                    vector<point> lower, upper;
return {area, c};
                                                     for (int i = 0; i < a.points.size(); i++) {
                                                     while (lower.size() >= 2 &&
                                                        cross(lower.back() - lower[lower.size()
Intersection
                                                        2], a.points[i] - lower.back()) < EPS)
// -1 coincide, 0 parallel, 1 intersection
int intersection(line a, line b, point& p) {
  if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
                                                       lower.pop back();
                                                      while (upper size() >= 2 &&
p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
                                                        cross(upper.back() - upper[upper.size()
- - a.a, b.b - b.a) * (b - a) + a;
                                                       2], a.points[i] - upper.back()) > -EPS)
                                                       upper.pop_back();
                                                      lower.push_back(a.points[i]);
 if (abs(cross(a.b - a.a, a.b - b.a)) > EPS)
                                                      upper.push back(a.points[i]);
→ return 0;
return -1:
                                                    lower.insert(lower.end(), upper.rbegin() + 1, \rightarrow b*v + c*z + d = 0
}
// area of intersection
                                                       upper.rend());
double intersection(circle a, circle b) {
                                                    return convex polygon(lower);
double d = abs(a.c - b.c);
if (d <= b.r - a.r) return area(a);
if (d <= a.r - b.r) return area(b);</pre>
                                                    Maximum Colinear Points
if (d \ge a.r + b.r) return 0;
                                                   const 11 range = 10000;
struct Slope { // a rational number with
double alpha = acos((sq(a.r) + sq(d) -
\rightarrow sq(b.r)) / (2 * a.r * d));
                                                     \rightarrow unsigned infinity (1,0)
                                                    .11 p, q;
double beta = acos((sq(b.r) + sq(d) - sq(a.r))
Slope(11 pP=0, 11 qP=0) {
                                                     if (qP==0) {
                                                      p = 1, q = 0;
return;
    alpha) + sq(b.r) * (beta - 0.5 * sin(2 *
⇒ beta));
```

```
11 g = \_gcd(pP, qP);
  pP /= g, qP /= g;
if(qP < 0) pP *= -1, qP *= -1;
 p = pP, q = qP;
 bool operator== (const Slope &other) const {
 return other.p == p && other.q == q:
námespace std {
 template<>
 struct hash<Slope> { // tupical
 → rectangular/lattice hash
  size_t operator() (const Slope &r) const {
  return (2*range+1) * (r.p + range) + r.q +
// n points in [-range, range]
// compute the largest colinear subset
int max_colinear_points(vector<pair<11,11>>
    &points) {
 if(points.size() <= 2) return points.size();</pre>
 int best = 0:
 unordered_map<Slope, int> counter;
 for(int i=0;i<points.size();i++) {
  for(int j=i+1; j<points.size(); j++) {</pre>
   Slope slope(points[i].second-points[j]
    .second,points[i].first-points[j].first);
   best = max(best, ++counter[slope]+1):
  if(i != points.size()-1) counter.clear():
return best:
    3D Geometry
struct point3d {
 double x, y, z;
 point3d operator+(point3d a) const { return
 \rightarrow {x+a.x, y+a.y, z+a.z}; }
 point3d operator*(double a) const { return
 \rightarrow {x*a, v*a, z*a}: }
 point3d operator-() const { return {-x, -y,
 → -z}; }
 point3d operator-(point3d a) const { return
\rightarrow *this + -a: }
 point3d operator/(double a) const { return
 *this * (1/a); }
 double norm() { return x*x + y*y + z*z; }
double abs() { return sqrt(norm()); }
 point3d normalize() { return *this /

    this->abs(): }

double dot(point3d a, point3d b) { return
\rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
point3d cross(point3d a, point3d b) { return
   \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
  a.x*b.y - a.y*b.x; }
struct line3d { point3d a, b; };
struct plane { double a, b, c, d; } // a*x +
struct sphere { point3d c; double r; };
#define sq(a) ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
double surface(circle a) { return 4 * sq(a.r)
double volume(circle a) { return 4.0/3.0 *
\hookrightarrow cb(a.r) * M PI: }
```

10 Optimization

// SameNumberOfOneBits, next permutation int snoob(int a) {

```
int b = a \& -a, c = a + b;
 return c | ((a^ c) >> 2) / b:
 // example usage
// example usage
int main() {
    char l1[] = {'1', '2', '3', '4', '5'};
    char l2[] = {'a', 'b', 'c', 'd'};
    int d1 = 5, d2 = 4;
    // prints 12345abcd, 1234a5bcd, ...
  int min = (1 << d1) -1, max = min << d2
 for (int i = min; i <= max; i = snoob(i)) {
   int p1 = 0, p2 = 0, v = i;
   while (p1 < d1 || p2 < d2) {
      cout << ((v & 1) ? 11[p1++] : 12[p2++]);
    .v /= 2;
   cout << '\n':
bool isPowerOf2(11 a) {
  return a > 0 && !(a & a-1);
 bool isPowerOf3(11 a) {
return a>0&&!(12157665459056928801u11%a);
bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
  return abs(x-round(x)) < 0.00000000001;</pre>
Fast Modulo
// faster modulo with constant modulus
struct FastMod {
  ull b, m;
 FastMod(ull b) : b(b), m(-1ULL / b) {}
ull reduce(ull a) { // a % b + (0 or b)}
return a - (ull)((_uint128_t(m) * a) >> 64)

    * b;
}
11 Python
```

Recursion Limit Removal (Basic)

import sys sys.setrecursionlimit(10**6)

Recursion Limit Removal (Advanced)

```
# @bootstrap over recursive function
# replace 'return' with 'yield'
# for when sys method does not work
from types import GeneratorType def bootstrap(f, stack=[]):
 def wrappedfunc(*args, **kwargs):
  if stack:
return f(*args, **kwargs)
  else:
to = f(*args, **kwargs)
    while True:
   if type(to) is GeneratorType:
      stack.append(to)
       to = next(to)
      else:
       stack.pop()
       if not stack:
 to = stack[-1].send(to)
return to
return wrappedfunc
# EXAMPLE recursive fibonacci
@bootstrap
def f(n):
   if (n < 2):
     yield n</pre>
 yield (yield f(n-1)) + (yield f(n-2))
```

Python 3 Compatibility

import svs from __future__ import division, print function if sys.version info[0] < 3: from _builtin import xrange as range from future_builtins import ascii, filter, → hex. map. oct. zip

```
12 Additional
 Judge Speed
 // kattis: 0.50s
// codeforces: 0.421s
// atcoder: 0.455s
 #include <bits/stdc++.h>
using namespace std;
 int v = 1e9/2, p = 1;
 int main() {
  for (int i = 1; i <= v; i++) p *= i;</pre>
 Judge Pre-Contest Checks
       int128 and float128 support?
 -does extra or missing whitespace cause WA?
 -documentation up to date?
- printer usage available and functional?
// each case tests a different fail condition
// try them before contests to see error codes
struct g { int arr[1000000]; g(){}};
 vector<g> a;
// 0=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
int judge(int n) {
   if (n == 0) exit(0);
   if (n == 2) while(1);
   if (n == 3) while(1) putchar_unlocked('a');
   if (n == 3) while(1) putchar_unlocked('a');
   if (n == 3) while(1) putchar_unlocked('a');
  if (n == 4) assert(0);
if (n == 5) 0 / 0;
if (n == 6) *(int*)(0) = 0;
  return n + judge(n + 1);
 GCC Builtin Docs
// 128-bit integer
int128 a;
unsigned __int128 b;
// 128-bit float
// minor improvements over long double
// minor improvements over long double
_float128 c;
// log2 floor
_lg(n);
// number of 1 bits
// can add ll like popcountll for long longs
_builtin_popcount(n);
// number of trailing zeroes
_builtin_ctz(n);
// number of leading zeroes
_builtin_clz(n);
// 1-indexed least significant 1 bit
_builtin_fs(n):
__builtin_ffs(n);
// parity of number
__builtin_parity(n);
 Limits
                             \pm 2147483647|\ \pm 2^{31}-1|10^9
  int
                                                         2^{32} - 1 | 10^{9}
                                4294967295
  uint
            \pm 9223372036854775807 | \pm 2^{63} - 1 | 10^{18}
  11
                                                          \frac{1}{2}^{64} - \frac{1}{10}^{19}
            18446744073709551615
 i128 \pm 170141183460469231... \pm 2^{127} - 110^{38}
 Complexity classes input size (per second):
 O(n^n) or O(n!)
                                                                    n \leq 10
 O(2^n)
                                                                    n \leq 30
 O(n^3)
                                                                n < 1000
  O(n^2)
                                                              n < 30000
                                                                  n < 10^6
  O(n\sqrt{n})
  O(n \log n)
                                                                  n \le 10^7
 O(n)
                                                                  n < 10^9
```