```
General
                               7 Graphs
    Algorithms
                               8 2D Geometry
    Structures
                               9 3D Geometry
    Strings
                               10 Optimization
    Greedy
                               11 Python
    Math
                               12 Additional
      General
run.sh
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
for i in *.in; do
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
done
Header
// use better compiler options
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
// include everything
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
using namespace std;
using namespace __gnu_cxx; // rope
using namespace __gnu_pbds; // tree/trie
// common defines
#define fastio

→ ios_base::sync_with_stdio(0);cin.tie(0);
#define nostacklim rlimit RZ; getrlimit(3,&RZ)
 \rightarrow );RZ.rlim cur=-1;setrlimit(3,\&RZ);
#define DEBUG(v) cerr<<__LINE__<<": "<<#v<<" =
\Rightarrow "<<v<<'\n'; #define TIMER
→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
#define i128 _ int128
#define u128 unsigned i128
#define ld long double
// global variables
mt19937 rng((uint32_t)chrono::steady

    _clock::now().time_since_epoch().count());
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c - '0';
void read(int& n) {
   char c; n = 0; int s = 1;
   if ((c=getchar_unlocked())=='-') s = -1;
 else n = c - \sqrt{0}:
 while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c -
void read(ld& n) {
 char c; n = 0;

.ld m = 0, o = 1; bool d = false; int s = 1;

.if ((c=getchar_unlocked())=='-') s = -1;
 else if (c == '.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n')
  if (c == '.') d = true;
```

```
else if (d) { m=m*10+c-'0'; o*=0.1; } else n = n * 10 + c - '0':
 n = s * (n + m * o):
void read(double& n) {
 ld m; read(m); n = m;
void read(float& n) {
  ld m; read(m); n = m;
 void read(string& s) {
 char c; s = ""
 while((c=getchar_unlocked())!=' '&&c!='\n')
bool readline(string& s) {
 char c; s = ""
 while(c=getchar_unlocked()) {
  if (c == '\n') return true;
if (c == EOF) return false;
  s += c:
 return false:
void print(unsigned int n) {
 if (n / 10) print(n / 10);
 putchar unlocked(n % 10 + '0'):
void print(int n) {
 if (n < 0) { putchar_unlocked('-'); n*=-1; }
 print((unsigned int)n);
Common Structs
   n-dimension vectors
// Vec<2, int> v(n, m) = arr[n][m]

// Vec<2, int> v(n, m, -1) default init -1

template<int D, typename T>
struct Vec : public vector < Vec < D-1. T>> {
 template<typename... Args>
 Vec(int n=0, Args... args) : vector<Vec<D-1,
 \rightarrow T>>(n, Vec<D-1, T>(args...)) {}
};
template<typename T>
struct Vec<1, T> : public vector<T> {
   Vec(int n=0. T val=T()) : vector<T>(n. val) {}
   Algorithms
Binary Search
 // search for k in [p,n)
template<typename T>
int binsearch(T x[], int k, int n, int p = 0) {
 for (int i = n; i >= 1; i /= 2)

while (p+i < n && x[p+i] <= k) p += i;

return p; // bool: x[p] == k;
Min/Max Subarray
   max - compare = a < b, reset = a < 0
 // min - compare = a > b, reset = a > 0
// returns {sum, {start, end}}
pair<int, pair<int, int>>
     ContiguousSubarray(int* a, int size,
     bool(*compare)(int, int),
 bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
 \rightarrow 0, s = 0;
 for (int i = 0: i < size: i++) {
  cur += a[i];
  if ((*compare)(best, cur)) { best = cur;
    start = s; end = i; }
  if ((*reset)(cur)) { cur = 0; s = i + 1; }
 return {best, {start, end}};
Quickselect
 #define QSNE -999999
```

```
int partition(int arr[], int 1, int r)
 int x = arr[r], i = 1;
for (int j = 1; j <= r - 1; j++)
if (arr[j] <= x)
swap(arr[i++], arr[j]);
 swap(arr[i], arr[r]);
 return i:
// find k'th smallest element in unsorted array,
\rightarrow only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - 1 + 1)) return QSNE;
 swap(arr[1 + rng() % (r-1+1)], arr[r]);
 int pos = partition(arr, 1, r);
 if (pos-l==k-1) return arr[pos];
 if (pos-l>k-1) return qselect(arr,l,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
|}
|// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
→ on both axis
pair<int. int> saddleback search(int** arr. int
\hookrightarrow x, int y, int v) {
int i = x-1, j = 0;
while (i >= 0 && j < y) {
   if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i--: j++;
 return {-1, -1};
Ternary Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {
 while (b-a > 4) {
  int m = (a+b)/2;
if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
 . a = i;
return a:
#define TERNPREC 0.000001
double ternsearch (double a, double b, double
     (*f)(double)) {
 while (b-a > TERNPREC * 4) {
    double m = (a+b)/2:
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  > = m;
else b = m + TERNPREC;
 for (double i = a + TERNPREC; i <= b; i +=
    TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
double goldsection(double a, double b, double
 double g(double)) {
  double r = (sqrt(5)-1)/2, eps = 1e-7;
  double x1 = b - r*(b-a), x2 = a + r*(b-a);
  double f1 = f(x1), f2 = f(x2);
 while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
   a = x1; x1 = x2; f1 = f2;
x2 = a + r*(b-a); f2 = f(x2);
 return a:
```

```
3 Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums
\hookrightarrow O(\log n) updates, O(\log n) gets
struct Fenwick { int n: ll* tree:
 void update(int i, int val) {
 .++i;
while (i <= n) {
  tree[i] += val;</pre>
   i += i & (-i);
 Fenwick(int size) {
  n = size;
  tree = new ll[n+1];
for (int i = 1; i <= n; i++)
   .tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
 11 operator[](int i) {
  if (i < 0 || i > n) return 0;
  while (i>0)
  sum += tree[i];
i -= i & (-i);
  return sum;
 ll getRange(int a, int b) { return
    operator[](b) - operator[](a-1); }
Hashtable
// similar to unordered map, but faster
struct chash {
    const uint64_t C = (11)(2e18 * M_PI) + 71;
    ll operator()(11 x) const { return
    builtin bswap64(x*C): }
int main() {
  gp_hash_table<11,int,chash>
\rightarrow hashtable({},{},{},{},{1<<16});
for (int i = 0; i < 100; i++)

hashtable[i] = 200+i;

if (hashtable.find(10) != hashtable.end())

cout << hashtable[10];
Ordered Set
template <typename T>
using oset = tree<T,null type,less<T>,rb tree
    tag, tree order statistics node update>;
template <typename T, typename D> using omap = tree<T,D,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
int main()
 oset<int> o_set;
o_set.insert(5); o_set.insert(1);

    o_set.insert(3);
// get second smallest element
 cout << *(o_set.find_by_order(1));</pre>
 // number of elements less than k=4
cout << ' ' << o_set.order_of_key(4) << '\n';</pre>
 // equivalent with ordered map
 omap<int,int> o_map;
o_map[5]=1;o_map[1]=2;o_map[3]=3;
cout << (*(o_map.find_by_order(1))).first;
 cout << ' ' << o_map.order_of_key(4) << '\n';
```

```
Rope
                                                         // print things with prefix "1"
                                                         auto range = trie.prefix_range("1");
// O(\log n) insert, delete, concatenate
                                                         for (auto it = range.first; it !=
int main() {
 // generate rope
                                                         → range.second: it++)
 rope<int> v;
                                                          cout << *it << '
 for (int i = 0; i < 100; i++)
.v.push_back(i);
                                                        Wavelet Tree
 // move range to front
                                                        using iter = vector<int>::iterator;
 rope<int> copy = v.substr(10, 10);
v.erase(10, 10);
                                                        struct WaveletTree {
   Vec<2, int> C; int s;
 v.insert(copy.mutable_begin(), copy);
                                                          // sigma = highest value + 1
                                                         WaveletTree(vector<int>& a. int sigma) :
 // print elements of rope
for (auto it : v) cout << it << "";
                                                            s(sigma), C(sigma*2, 0) {
                                                          build(a.begin(), a.end(), 0, s-1, 1);
                                                         void build(iter b. iter e. int L. int U. int
Segment Tree
                                                          u) {
if (L == U) return
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
                                                          int M = (L+U)/2;
 typedef int T;
                                                           C[u].reserve(e-b+1); C[u].push back(0);
 static constexpr T UNIT = INT_MIN;
                                                          for (auto it = b; it != e; ++it)
C[u].push_back(C[u].back() + (*it<=M));
 T f(T a, T b) {
 if (a == UNIT) return b;
if (b == UNIT) return a;
                                                          auto p = stable_partition(b, e, [=](int
                                                           i){return i<=M;});
  return max(a,b);
                                                          build(b, p, L, M, u*2);
 int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
                                                          build(p, e, M+1, U, u*2+1);
                                                         // number of occurrences of x in [0,i)
\rightarrow n(n) {}
                                                         int rank(int x, int i) {
   int L = 0, U = s-1, u = 1, M, r;
   while (L != U) {
 SegmentTree(vector<T> arr)

    SegmentTree(arr.size()) {

 for (int i=0:i<arr.size():i++)
                                                           M = (L+U)/2;
r = C[u][i]; u*=2;

    update(i,arr[i]);

                                                           if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
                                                          return i:
                                                          ^{\prime\prime} number of occurences of x in [l,r)
 T query(int b, int e) { // query [b, e)
                                                         int count(int x, int 1, int r) {
  return rank(x, r) - rank(x, 1);
  Tra = UNIT, rb = UNIT;
  for (b+=n, e+=n; b<= b/=2, e/=2) {
    if (b % 2) ra = f(ra, s[b++]);
    if (e % 2) rb = f(s[--e], rb);
                                                         // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  return f(ra, rb):
                                                          while (L != U) {
   M = (L+U)/2;
 T get(int p) { return query(p, p+1); }
                                                           ri = C[u][1]; rj = C[u][r]; u*=2;
                                                           if (k \le rj-ri)^{n}l = ri, r = rj, U = M;
Sparse Table
                                                           else k -= řj-rí, l -= ŕi, r -= ŕj,
template < class T> struct SparseTable {
                                                           L = M+1, ++u;
 vector<vector<T>> m;
                                                          return U:
 SparseTable(vector<T> arr) {
  m.push back(arr);
  for (int k = 1: (1<<(k)) <= size(arr): k++)
                                                         // # elements between [x,y] in [l, r)
                                                         mutable int L, U;
  m.push_back(vector<T>(size(arr)-(1<(k)+1));
                                                         int range(int x, int y, int 1, int r) const {
  for (int i = 0; i < size(arr)-(1<<k)+1; i
                                                          if (y < x \text{ or } r <= 1) return 0;
                                                          L = x; U = y;
 [k][i] = min(m[k-1][i],
                                                          return range(1, r, 0, s-1, 1);
\rightarrow m[k-1][i+(1<<(k-1))]:
}
// min of range [l,r]
                                                         int range(int 1, int r, int x, int y, int u)
                                                         → const {
                                                          if (y < L or U < x) return 0;
if (L <= x and y <= U) return r-l;
T query(int 1, int r) {
 int k = _-lg(r-l+1);
                                                          int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return \min(m[k][1], m[k][r-(1<< k)+1]);
                                                          return range(ri, rj, x, M, u*2) + range(1-ri
}
};
                                                            r-rj, M+1, y, u*2+1);
                                                          ^{\prime}// # elements <= x in [l, r]
                                                         int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);
typedef trie<string, null_type,

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>

→ trie_type;

int main() {
                                                             Strings
 // generate trie
 trie_type trie;
                                                        Aho Corasick
 for (int i = 0; i < 20; i++)
                                                           range of alphabet for automata to consider
 trie.insert(to string(i)); // true if new,
                                                           MAXC = 26, OFFC = 'a' if only lowercase
\hookrightarrow false if old
```

```
|const int MAXC = 256:
const int OFFC = 0:
struct aho_corasick {
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
  int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
  for (int i = 0; i < size; i++) {
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
      s[cur].go[c-OFFC] = s.size();
      s.push back(state());
     cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue int> sq;
for (int c = 0; c < MAXC; c++) {
   if (s[0].so[c] != 0) {
      ...s[s[0].so[c]].fail = 0;</pre>
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
   for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
      int failure = s[e].fail;
while (s[failure].go[c] == -1)
        failure = s[failure].fail;
      failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
      for (auto length : s[failure].out)
s[s[e].go[c]].out.insert(length);
      sq.push(s[e].go[c]);
 // list of {start pos, pattern id}
  vector<pair<int, int>> search(string text)
  vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {
  while (s[cur].go[text[i]-OFFC] == -1)
    cur = s[cur].fail;
cur = s[cur].go[text[i]-OFFC];
    if (s[cur].out.size())
    for (auto end : s[cur].out)
. toret.push_back({i - end.first + 1,
     end.second):
  return toret:
Boyer Moore
struct defint { int i = -1; };
vector<int> boyermoore(string txt, string pat)
 vector<int> toret; unordered_map<char, defint>string lcp(string* arr, int n, bool sorted =
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0; i < m; i++) badchar[pat[i]].i
 \rightarrow = i; int s = 0:
 while (s \leq n - m) {
  int j = m - 1:
```

while $(j \ge 0 \&\& pat[j] == txt[s + j]) j--;$

.if (j < 0) {

```
..toret.push_back(s);
   s += (s + m < n) ? m - badchar[txt[s +
   m]].i : 1:
 .} else
   s += \max(1, i - badchar[txt[s + i]].i):
 return toret:
English Conversion
const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven", "eight", "nine"};
const string teens[] ={"ten", "eleven",
    "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen";
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",

    "eighty", "ninety"};
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion", "quintillion", "sextillion",
    "septillion"};
string convert(int num, int carry) {
 if (num < 0) return "negative " +

    convert(-num, 0):

     (num < 10) return ones[num];
(num < 20) return teens[num % 10]
 if (num < 100) return tens[(num / 10) - 2] +
     (num\%10==0?"":"") + ones[num\%10]:
 if (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" + (num%100==0?"":" ") + convert(num % 100,
    0);
 return convert(num / 1000, carry + 1) + " " +
    mags[carry] + " " + convert(num % 1000.
    0):
string convert(int num) {
 return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = 0; i < n + 1; i++)
  next[i] = 0;
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
   j = next[j];
  if (j > 0 | pat[j] == pat[i])
   next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
  if (txt[i] == pat[j]) {
  if (++j == n)
    toret.push_back(i - j + 1);
  } else if (j > 0) {
...j = next[j];
 return toret;
Longest Common Prefix (array)
 // longest common prefix of strings in array
 → false) {
idise; l
if (n == 0) return "";
if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==

    arr[n-1][v])
    r += arr[0][v++];

 return r:
```

```
Longest Common Subsequence
                                                       unsigned int mwb(string s, set<string> dict) { | for (int j = 1; j <= n; j++)</pre>
                                                        int l = s.size();
string lcs(string a, string b) {
                                                        vector<unsigned int> arr(l+1, -1);
 int m = a.length(), n = b.length();
                                                        arr[0] = 0;
for (int i = 0; i < 1; i++) {
 int L[m+1][n+1];
 for (int i = 0; i <= m; i++) {
                                                         if (arr[i] != -1) {
 for (int j = 0; j <= n; j++) {
...if (i == 0 || j == 0) L[i][j] = 0;
...else if (a[i-1] == b[j-1]) L[i][j] =
                                                         for (auto e : díct) {
                                                           int L = e.size();
if (1 >= i + L) {
                                                             bool isGood = true;
\hookrightarrow L[i-1][j-1]+1;
                                                         .....for (int j = 0; isGood && j < L; j++)
.....if (s[i+j] != e[j])
   else L[i][j] = \max(L[i-1][j], L[i][j-1]);
                                                            isGood = false;
if (isGood)
 // return L[m][n]; // length of lcs
                                                            arr[i+L] = min(arr[i]+1, arr[i+L]);
 string out = "";
 int i = m - 1, j = n - 1;
 while (i >= 0 && j >= 0) {
 if (a[i] == b[j]) {
                                                        return arr[1];
   out = a[i--] + out;
                                                       Hashing
  else if (L[i][j+1] > L[i+1][j]) i--;
                                                       #define HASHER 27
  .else j--;
                                                       ull basicHash(string s) {
                                                        ull v = 0:
 return out;
                                                        for (auto c : s) v = (c - 'a' + 1) + v *
                                                        → HASHER;
                                                        return v;
Longest Common Substring
// l is array of palindrome length at that
                                                       const int MAXN = 1000001:
int manacher(string s, int* 1) {
                                                       ull base[MAXN] = {1};
void genBase(int n) {
 int n = s.length() * 2;
                                                        for (int i = 1; i \le n; i++)

base[i] = base[i-1] * HASHER;
 for (int i = 0, j = 0, k; i < n; i += k, j =
 \rightarrow max(j-k, 0)) {
                                                       struct advHash {
  ull v, l; vector<ull> wip;
  while (i \ge j \&\& i + j + 1 < n \&\& s[(i-j)/2]
 \Rightarrow = s[(i+j+1)/2]) j++;
                                                        advHash(string& s): v(0) {
  .1[i] = j;
                                                         wip = vector<ull>(s.length()+1);\
 for (k = 1; i >= k && j >= k && l[i-k] !=
                                                         \sin c = c \cos w
   i-k: k++)
                                                         for (int i = 0; i < s.length(); i++)</pre>
  l[i+k] = min(l[i-k], j-k);
                                                          wip[i+1] = (s[i] - 'a' + 1) + wip[i] *
                                                           HASHER:
 return *max_element(1, 1 + n);
                                                         1 = s.length(): v = wip[1]:
                                                        ull del(int pos, int len) {
Cyclic Rotation (Lyndon)
                                                         return v - wip[pos+len]*base[l-pos-len] +
// simple strings = smaller than its nontrivial
                                                           wip[pos]*base[1-pos-len];
\rightarrow suffixes
                                                        ull substr(int pos, int len) {
// lyndon factorization = simple strings
                                                         return del(pos+len, (l-pos-len)) -
\hookrightarrow factorized
 // "abaaba" -> "ab", "aab", "a"
                                                           wip[pos]*base[len];
vector<string> duval(string s) {
 int n = s.length();
                                                        ull replace(int pos, char c) {
                                                         return v - wip[pos+1]*base[l-pos-1] + ((c -
 vector<string> lyndon;
 for (int i = 0; i < n;) {
   int j = i+1, k = i;
   for (; j < n && s[k] <= s[j]; j++)
   if (s[k] < s[j]) k = i;
                                                            'a' + 1) + wip[pos] *
                                                           HASHER) *base[1-pos-1];
                                                        ull replace(int pos, string s) {
                                                         // can't increase total string size
  for (; i <= k; i += j - k)
                                                         ull r = v -
   lyndon.push back(s.substr(i,j-k));
                                                            wip[pos+s.size()]*base[l-pos-s.size()], c
                                                          wip[pos];
 return lyndon;
                                                         for (int i = 0; i < s.size(); i++)
c = (s[i]-'a'+1) + c * HASHER;
}
// lexicographically smallest rotation
                                                          return r + c * base[l-pos-s.size()];
int minRotation(string s) {
 int n = s.length(); s += s;
 auto d = duval(s); int i = 0, a = 0;
while (a + d[i].length() < n) a +=</pre>
                                                       Subsequence Count

    d[i++].length();

                                                        // "banana", "ban" >> 3 (ban, ba..n, b..an)
 while (i && d[i] == d[i-1]) a -=
                                                       ull subsequences(string body, string subs) {

    d[i--].length();

                                                        int m = subs.length(), n = body.length();
                                                        if (m > n) return 0;
 return a;
                                                        ull** arr = new ull*[m+1];
                                                        for (int i = 0; i \le m; i++) arr[i] = new
Minimum Word Boundary
                                                         \rightarrow ull[n+1];
// minimum word boundary
                                                        for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[0][i] = 1;
for (int i = 1; i <= m; i++)
// compose string s using words from dict
// NOTE: can reuse words from dict
```

```
arr[i][j] = arr[i][j-1] + ((body[j-1] ==
   subs[i-1])? arr[i-1][j-1] : 0);
 return arr[m][n];
Suffix Array + LCP
struct SuffixArray {
 vector<int> sa, 1cp;
 SuffixArray(string& s, int lim=256) {
  int n = s.length() + 1, k = 0, a, b;
  vector<int> x(begin(s), end(s)+1), y(n),
 \rightarrow ws(max(n, lim)), rank(n);
  sa = lcp = y;
  iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j)
 \rightarrow 2), lim = p) {
  p = j; iota(begin(y), end(y), n - j);
   for (int i = 0; i < (n); i++)
if (sa[i] >= i)
     .y[p++] = sa[i] - j;
   fill(begin(ws), end(ws), 0);
for (int i = 0; i < (n); i++) ws[x[i]]++;
   for (int i = 1; i < (lim); i++) ws[i] +=
    ws[i - 1];
   for (int i = n; i--;) sa[--ws[x[y[i]]]] =
    v[i];
   swap(x, y); p = 1; x[sa[0]] = 0;
for (int i = 1; i < (n); i++) {
   a = sa[i - 1]; b = sa[i];</pre>
    x[b] = (y[a] = y[b] \&\& y[a + j] = y[b +
    j]) ? p - 1 : p++;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
  for (int i = 0, j; i < n - 1; lcp[rank[i++]]
    = k
  for (k \&\& k--, j = sa[rank[i] - 1];
     s[i + k] == s[j + k]; k++);
 // smallest cyclic shift
 int cyclic() { return sa[0]; }
 // longest repeated substring
 pair<int,int> lrs() {
  int length = -1, index = -1;
  for (int i = 0; i < lcp.size(); i++) {
  if (lcp[i] > length) {
   length = lcp[i];
    index = sa[i]:
  return {index,length};
 }
// count distinct substrings, excluding empty
 int distincts() {
   int n = sa.size() - 1, r = n - sa[0];
   for (int i = 1; i < lcp.size(); i++)</pre>
  r += (n - sa[i]) - lcp[i - 1];
  return r:
 }// count repeated substrings, excluding empty
 int repeateds() {
  .int r' = 0;
  for (int i = 1; i < lcp.size(); i++)
  r += \max(lcp[i] - lcp[i-1], 0);
  return r;
Suffix Tree (Ukkonen's)
struct SuffixTree {
 // n = 2*len+10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
 void ukkadd(int i, int c) { suff:
```

```
if (r[v]<=q) {
  if (q==-1 || c==toi(a[q])) q++; else {
...l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
   p[m] = p[v]; t[m][c] = m+1; t[m][toi(a[q])] = v;
   l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
v=s[p[m]]; q=l[m];
   while (q < r[m]) { v = t[v][toi(a[q])];
   q+=r[v]-l[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
   q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
  fill(r,r+N,(int)(a).size());
 memset(s, 0, sizeof s);
memset(t, -1, sizeof t);
fill(t[1],t[1]+ALPHA,0);
s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
  for(int i=0; i < a. size(); i++)
    ukkadd(i,toi(a[i]));
 // Longest Common Substring between 2 strings
 // returns {length, offset from first string}
 pair<int. int> best:
 int lcs(int_node, int i1, int i2, int olen) {
 if (1[node] <= i1 && i1 < r[node]) return 1;
if (1[node] <= i2 && i2 < r[node]) return 2;
  int mask=0
   len=node?olen+(r[node]-l[node]):0;
  for(int c=0; c<ALPHA; c++) if
   (t[node][c]!=-1)
mask |= lcs(t[node][c], i1, i2, len);
  if (mask==3)
  best=max(best, {len,r[node]-len});
  return mask:
 static pair<int, int> LCS(string s, string t)
 \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2));
 st.lcs(0, s.size(), s.size()+t.size()+1, 0); return st.best;
String Utilities
void lowercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
   ::tolower);
void uppercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
\hookrightarrow ::toupper);
void trim(string &s) {
 s.erase(s.begin(),find_if_not(s.begin(),s
     .end(),[](int c){return
    isspace(c);}));
 s.erase(find_if_not(s.rbegin(),s.rend(),[](int

    c){return isspace(c);}).base(),s.end());

vector<string> split(string& s, char token) {
    vector<string> v; stringstream ss(s);
    for (string e;getline(ss,e,token);)
        v.push_back(e);
    return v;
    Greedy
```

```
Interval Cover
int (L,R) = interval [L,R], in = \{\{l,r\}, index\}\} string a; (L,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\} vector (l,R) = interval [L,R], in = \{\{l,r\}, index\}

    vector<pair<pair<double,double>,int>> in) {
```

```
int i = 0; pair<double,int> pos = {L,-1};
                                                         .re %= mo;
    vector<int> a;
                                                         re = (re + mo) \% mo;
     sort(begin(in), end(in));
                                                         return true:
     while (pos.first < R) {
          double cur = pos.first;
while (i < (int)in.size() &&</pre>
                                                        Count Digit Occurences
    in[i].first.first <= cur)</pre>
                                                         /*count(n,d) counts the number of occurences of
                                                         \rightarrow a digit d in the range [0,n]*/
     max(pos, {in[i].first.second,in[i].second})
                                                         ll digit_count(ll n, ll d) {
                                                         ll result = 0:
          if (pos.first == cur) return {}:
                                                         while (n != 0) {
   result += ((n%10) == d ? 1 : 0):
          a.push_back(pos.second);
                                                          n /= 10;
     return a;
                                                         return result;
6 Math
                                                        11 count(ll n, ll d) {
                                                         if (n < 10) return (d > 0 && n >= d);
if ((n % 10) != 9) return digit_count(n, d) +
 Catalan Numbers
ull* catalan = new ull[1000000];
count(n-1, d):
                                                         return 10*count(n/10, d) + (n/10) + (d > 0):
                                                         Discrete Logarithm
                                                        int discretelog(int a, int b, int m) {
                                                         11 n = sqrt(m) + 1, an = 1;

→ % mod:

                                                         for (ll i = 0; i < n; ++i)
an = (an * a) % m;
 if (catalan[i] >= mod)
catalan[i] -= mod;
                                                         unordered_map<11, 11> vals;
                                                         for (ll q = 0, cur = b; q <= n; q++) {
  vals[cur] = q;
 // TODO: consider binomial coefficient method
                                                          cur = (cur * a) \% m;
Combinatorics (nCr, nPr)
                                                         for (ll p = 1, cur = 1; p \le n; p++) {
                                                          cur = (cur * an) % m;
 // can optimize by precomputing factorials, and
                                                           if (vals.count(cur)) {
 \hookrightarrow fact[n]/fact[n-r]
                                                           int ans = n * p - vals[cur];
 ull nPr(ull n, ull r) {
                                                           return ans:
 for (ull i = n-r+1; i <= n; i++)
 .v *= i;
return v:
                                                         return -1;
ull nPr(ull n, ull r, ull m) {
                                                        Euler Phi / Totient
 ull v = 1;
for (ull i = n-r+1: i <= n: i++)
                                                         int phi(int n) {
  v = (v * i) \% m;
                                                         int^r = n;
                                                         for (int i = 2; i * i <= n; i++) {
    if (n % i == 0) r -= r / i;
    while (n % i == 0) n /= i;
 return v;
ull nCr(ull n, ull r) {
 long double \dot{v} = 1;
 for (ull i = 1; i <= r; i++)
v = v * (n-r+i) /i;
                                                         if (n > 1) r = r / n;
                                                         return r;
 return (ull)(v + 0.001);
                                                         #define n 100000
                                                        ll phi[n+1];
 // requires modulo math
                                                        void computeTotient() {
 // can optimize by precomputing mfac and
                                                         for (int i=1; i<=n; i++) phi[i] = i;

→ minv-mfac

                                                         for (int p=2; p<=n; p++) {
ull nCr(ull n, ull r, ull m) {
                                                          if (phi[p] == p) {
 return mfac(n, m) * minv(mfac(k, m), m) % m
 \rightarrow minv(mfac(n-k, m), m) % m:
                                                           phi[p] = p-1;
                                                           for (int i = 2*p; i<=n; i += p) phi[i] =
                                                             (phi[i]/p) * (p-1);
Multinomials
ll multinomial(vector<int>& v) {
    ll c = 1, m = v.empty() ? 1 : v[0];
    for(int i = 1; i < v.size(); i++)
        for (int j = 0; j < v[i]; j++)
        c = c * ++m / (j+1);
    }
</pre>
                                                        Factorials
                                                         // digits in factorial
                                                        #define kamenetsky(n) (floor((n * log10(n /
 return c;
                                                         \rightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
                                                        // approximation of factorial #define stirling(n) ((n == 1) ? 1 : sqrt(2 *
 Chinese Remainder Theorem
bool ecrt(l1* r, l1* m, int n, l1& re, l1& mo)
                                                         \rightarrow M_PI * n) * pow(n / M_E, n))
// natural log of factorial
                                                        #define lfactorial(n) (lgamma(n+1))
 d = egcd(mo, m[i], x, y);

if ((r[i] - re) % d != 0) return false;

x = (r[i] - re) / d * x % (m[i] / d);

re += x * mo;
                                                        Prime Factorization
                                                        // do not call directly
                                                        ll pollard_rho(ll n, ll s) {
  mo = mo / d * m[i];
                                                        11 x, y;
```

```
x = y = rand() \% (n - 1) + 1;
 int head = 1, tail = 2;
 while (true) {
 x = mult(x, x, n);

x = (x + s) \% n;
 if (x == y) return n;
 11 d = gcd(max(x - y, y - x), n);
 if (1 < d && d < n) return d;
if (++head == tail) y = x, tail <<= 1;
// call for prime factors
void factorize(ll n, vector<ll> &divisor) {
 if (n == 1) return;
 if (isPrime(n)) divisor.push_back(n);
  while (d' \ge n) d = pollard_rho(n, rand() % (n)
 \rightarrow -1) +1);
 factorize(n / d, divisor);
 factorize(d, divisor);
Factorize Factorials
   NOTE: count distinct divisors of n by
// computing (q1+1)*(q2+1)*...*(qk+1)
// where qi are powers of primes pi dividing n
// use that and this code to solve
→ https://open.kattis.com/problems/divisors
// max power of a prime p dividing n!
// D(log(n))
int legendre(int n, int p) {
    int mx = 0;
while(n>0) n/=p, mx+=n;
    return mx;
bitset<10000> sieve:
vector<int> primes;
// get all primes O(n log n)
// if dealing with small numbers
void genPrimes(int n) {
    sieve[0] = sieve[1] = 1;
    primes.push_back(2);
    for (int i = 3; i <= n; i+=2
         if (i%2 != 0 && !sieve[i]) {
             primes.push_back(i);
             for (int j = i * 3; j <= n; j +=
→ i*2)
                 sieve[j] = 1;
  ' make sure you call genPrimes first
// return vector of prime factor powers as
\rightarrow vector v of size pi(n)
// so that v[i] = power of primes[i] dividing
\nearrow n! O(pi(n) * log(n)) where pi(n) is prime
   counting fn
// so basically O(n) since pi(n) = O(n/log(n))
vector<int> factorize_factorial(int n) {
    vector<int> factorization(primes.size(),
    for(int i=0;i<primes.size() && primes[i] <=</pre>
        factorization[i] = legendre(n,
    primes[i]);
    return factorization:
// same thing but for C(n,k)
vector<int> factorize_binom(int n, int k) {
    vector<int> factorization(primes.size(),
    for(int i=0;i<primes.size() && primes[i] <=</pre>
    n;i++) {
        factorization[i] = legendre(n,
    primes[i]) - legendre(k, primes[i]) -
    legendre(n-k, primes[i]);
```

```
return factorization:
Farev Fractions
   generate 0 \le a/b \le 1 ordered. b \le n
   farey(4) = 0/1 \ 1/4 \ 1/3 \ 1/2 \ 2/3 \ 3/4 \ 1/1
// length is sum of phi(i) for i = 1 to n
vector<pair<int, int>> farev(int n) {
 int h = 0, k = 1, x = 1, y = 0, r;
 vector<pair<int, int>> v;
 v.push_back({h, k});
 r = (n-y)/k;
  v += r*k: x += r*h:
swap(x,h); swap(y,k);
x = -x; y = -y;
} while (k > 1);
v.push_back({1, 1});
return v;
Fast Fourier Transform
#define cd complex<double>
const double PI = acos(-1):
void fft(vector<cd>& a, bool invert) {
 int n = a.size();
 for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1:
  for (; j & bit; bit >>= 1) j ^= bit;
  if (i < j) swap(a[i], a[j]);
 for (int len = 2; len <= n; len <<= 1) {
  double ang = 2 * PI / len * (invert ? -1 :
  cd wlen(cos(ang), sin(ang));
  for (int i = 0; i < n; i += len) {
   .cd w(1);
   for (int j = 0; j < len / 2; j++) {
  cd u = a[i+j], v = a[i+j+len/2] * w;
   a[i+j] = u + v;
a[i+j+len/2] = u - v;
   w *= wlen:
 if (invert)
 for (auto& x : a)
  x /= n:
vector<int> fftmult(vector<int> const& a.

  vector<int> const& b) {
  vector<cd> fa(a.begin(), a.end()),

    fb(b.begin(), b.end());
int n = 1 << (32 - __builtin_clz(a.size() +</pre>
 \rightarrow b.size() - 1));
fa.resize(n); fb.resize(n);
 fft(fa, false); fft(fb, false)
 for (int i = 0; i < n; i++) fa[i] *= fb[i];
 fft(fa, true);
 vector<int> toret(n):
 for (int i = 0; i < n; i++) toret[i] =

    round(fa[i].real());
    return toret:
Greatest Common Denominator
| ll egcd(ll a, ll b, ll& x, ll& y) {
if (b == 0) { x = 1; y = 0; return a; }
 ll gcd = egcd(b, a % b, x, y);
x = a / b * y;
swap(x, y);
return gcd;
```

```
Josephus Problem
                                                                                                         for (;w;w/=2) {
  if (w&1) {
 // O-indexed, arbitrary k
int josephus(int n, int k) {
 if (n == 1) return 0;
if (k == 1) return n-1;
if (k > n) return (josephus(n-1,k)+k)%n;
  int res = josephus(n-n/k,k)-n\%k;
                                                                                                          return r;
  return res + ((res<0)?n:res/(k-1));
 // fast case if k=2, traditional josephus
 int josephus(int n) {
 return 2*(n-(1<<(32-_builtin_clz(n)-1)));
                                                                                                        Matrix
Least Common Multiple
 #define lcm(a,b) ((a*b)/__gcd(a,b))
                                                                                                          int w, h;
Modulo Operations
                                                                                                          \rightarrow h(y) \{\}
 #define MOD 1000000007
 #define madd(a,b,m) (a+b-((a+b-m>=0)?m:0))
 #define mult(a,b,m) ((ull)a*b%m)
                                                                                                               m; }
 #define msub(a,b,m) (a-b+((a<b)?m:0))
11 mpow(11 b, 11 e, 11 m) {
 Il mpow(11 b, 11 s, -1 s
                                                                                                           return *this:
    e /= 2;
 return x % m;
 ull mfac(ull n, ull m) {
                                                                                                            return *this:
  for (int i = n; i > 1; i--)
   f = (f * i) \% m;
  return f;
 // if m is not guaranteed to be prime
ll_minv(ll b, ll m) {
  11 x = 0, y = 0;
   if (egcd(b, m, x, y) != 1) return -1;
                                                                                                                return z:
  return (x % m + m) % m;
 il mdiv compmod(int a, int b, int m) {
  if (__gcd(b, m) != 1) return -1;
  return mult(a, minv(b, m), m);
 // if m is prime (like 10^9+7)
11 mdiv_primemod (int a, int b, int m) {
  return mult(a, mpow(b, m-2, m), m);
}
// tonelli shanks = sqrt(n) % m, m is prime
                                                                                                            return á;
ll legendre(ll a, ll m){
  if (a % m==0) return 0;
  if (m == 2) return 1:
  return mpow(a, (m-1)/2, m);
 11 msqrt(11 n, 11 m) {
 ll s = __builtin_ctzll(m-1), q = (m-111)>>s,
 \rightarrow z = rand()%(m-1)+1;
  if (m == 2) return 1;
if (s == 1) return mpow(n,(m+1)/411,m);
  while (legendre(z,m)!=m-1) z = rand()\%(m-1)+1;
 11 c = mpow(z,q,m), r = mpow(n,(q+1)/2,m), t
                                                                                                          → c[i];
 \rightarrow = mpow(n,q,m), M = s;
   while (t != 1){
                                                                                                         \rightarrow a[i][i+1] = 1;
   11 i=1. ts = (t * t) % m:
    while (ts != 1) i++, ts = (ts * ts) % m;
  for (int'j = 0; j < M-i-1; j++) b = (b * b) %
                                                                                                         return s:
 r = r * b \% m; c = b * b \% m; t = t * c \% m;
 \hookrightarrow M = i;
  return r;
Modulo Tetration
11 tetraloop(ll a, ll b, ll m) {
 if(b == 0 | a == 1) return 1;
```

```
ll w = tetraloop(a,b-1,phi(m)), r = 1;
  r *= a; if (r >= m) r -= (r/m-1)*m;
  a *= a; if (a >= m) a -= (a/m-1)*m;
int tetration(int a, int b, int m) {
  if (a == 0 || m == 1) return ((b+1)&1)%m;
 return tetraloop(a,b,m) % m;
template<typename T>
struct Mat : public Vec<2. T> {
 Mat(int x, int y) : Vec<2, T>(x, y), w(x),
 static Mat<T> identity(int n) { Mat<T> m(n,n); Nimber Arithmetic
     for (int i=0;i<n;i++) m[i][i] = 1; return
 Mat<\hat{T}>\& operator+=(const Mat<T>\& m) {
  for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
   (*this)[i][j] += m[i][j];
 Mat<T>\& operator=(const Mat<T>\& m) {
  for (int i = 0; i < w; i++)
   for (int j = 0; j < h; j++)
(*this)[i][j] -= m[i][j];
 Mat<T> operator*(const Mat<T>& m) {
  Mat < T > z(w,m.h);
   for (int i = 0; i < w; i++)
   for (int j = 0; j < h; j++)
    for (int k = 0; k < m.h; k++)
z[i][k] += (*this)[i][j] * m[j][k];
 Mat<T> operator+(const Mat<T>& m) { Mat<T>
    a=*this; return a+=m; }
 Mat<T> operator-(const Mat<T>& m) { Mat<T>
    a=*this; return a-=m; }
 Mat<T>& operator*=(const Mat<T>& m) { return
  → *this = (*this)*m; }
 Mat<T> power(int n) {
  Mat<T> a = Mat<T>::identity(w),m=*this;
  for (;n;n/=2,m*=m) if (n\&1) a *= m;
Matrix Exponentiation
// F(n) = c[0]*F(n-1) + c[1]*F(n-2) + ...

// b is the base cases of same length c

ll matrix_exponentiation(ll n, vector<ll> c,
 vector<11> b) {
   if (nth < b.size()) return b[nth-1];
   Mat<11> a(c.size(), c.size()); ll s = 0;
   for (int i = 0; i < c.size(); i++) a[i][0] =</pre>
 for (int i = 0; i < c.size() - 1; i++)
 a = a.power(nth - c.size());
 for (int i = 0; i < c.size(); i++)

s += a[i][0] * b[i];
Matrix Subarray Sums
 template<class T> struct MatrixSum {
 Vec<2, T> p;
MatrixSum(Vec<2, T>& v) {
  p = Vec<2, T>(v.size()+1, v[0].size()+1);
  for (int i = 0; i < v.size(); i++)
   for (int j = 0; j < v[0].size(); j++)
```

```
...p[i+1][j+1] = v[i][j] + p[i][j+1] +
   p[i+1][j] - p[i][j];
 T sum(int u, int 1, int d, int r) {
 return p[d][r] - p[d][l] - p[u][r] + p[u][l];
Mobius Function
const int MAXN = 10000000;
// mu[n] = 0 iff n has no square factors
// 1 = even number prime factors, -1 = odd
short mu[MAXN] = \{0, 1\};
void mobius(){
 for (int i = 1; i < MAXN; i++)
. if (mu[i])
  for (int^{\prime}j = i + i; j < MAXN; j += i)
    mu[i] -= mu[i];
#define nimAdd(a,b) ((a)^(b))
ull nimMul(ull a, ull b, int i=6) {
static const ull M[]={INT_MIN>>32,
    M[0]^{(M[0] << 16)}, M[1]^{(M[1] << 8)},
    M[2]^{M}[2]^{(4)}, M[3]^{M}[3]^{(2)},
   M[4]^{(M[4]<<1)};
  if (i--==0) return a\&b;
  int k=1 << i;
  ull s=nimMul(a,b,i), m=M[5-i],
    t=nimMul(((a^(a>>k))&m)|(s&~m),
    ((b^(b>>k))&m)|(m&(~m>>1))<< k, i);
  return ((s^t)\&m)<< k | ((s^(t>>k))\&m);
Permutation
// c = array size, n = nth perm, return index
vector<int> gen_permutation(int c, int n) {
 vector<int> idx(c), per(c), fac(c); int i;
 for (i = 0; i < c; i++) idx[i] = i;
for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i;
for (i = c - 1; i >= 0; i--)
...per[c-i-1] = idx[fac[i]],
  idx.erase(idx.begin() + fac[i]);
 return per:
 // get what nth permutation of vector
int get permutation(vector<int>& v) {
 int use = 0, i = 1, r = 0;
 for (int e : v) {
 r = r * i++ + __builtin_popcount(use &
 \rightarrow -(1<<e)):
 use |= 1 << e;
 return r;
Permutation (string/multiset)
string freq2str(vector<int>& v) {
 string s;
 for (int i = 0; i < v.size(); i++)
  for (int j = 0; j < v[i]; j++)
   s += (char)(i + 'A');
// nth perm of multiset, n is 0-indexed
string gen_permutation(string s, ll n) {
 vector<int> freq(26, 0);
 for (auto e : s) freq[e - 'A']++;
 for (int i = 0; i < 26; i++) if (freq[i] > 0)
  freq[i]--; ll v = multinomial(freq);
  if (n < v) return (char)(i+'A') +
    gen_permutation(freq2str(freq), n);
  freq[i]++; n -= v;
 return "":
Miller-Rabin Primality Test
```

```
|// Miller-Rabin primality test - O(10 log^3 n)
bool isPrime(ull n) {
bool isPrime(ull n) {
    if (n < 2) return false;
    if (n == 2) return true;
    if (n % 2 == 0) return false;
    ull s = n - 1;
    while (s % 2 == 0) s /= 2;
    for (in f == 0, in < 10, in)</pre>
 for (int i = 0; i < 10; i++) {
  ull temp = s;
  ull a = rand() \% (n - 1) + 1;
  ull mod = mpow(a, temp, n);
while (temp!=n-1&&mod!=1&&mod!=n-1) {
   mod = mult(mod, mod, n);
   temp *= 2;
  if (mod!=n-1\&\&temp\%2==0) return false;
 .}
.return true:
Sieve of Eratosthenes
bitset<100000001> sieve;
// generate sieve - O(n log n)
void genSieve(int n) {
 sieve[0] = sieve[1] = 1:
 for (ull i = 3; i * i < n; i += 2)
. if (!sieve[i])
  for (ull_j = i * 3; j \le n; j += i * 2)
    sieve[j] = 1;
// query sieve after it's generated - 0(1)
bool querySieve(int n) {
 return n == 2 || (n % 2 != 0 && !sieve[n]);
Compile-time Prime Sieve
const int MAXN = 100000;
template<int N>
struct Sieve
 bool sieve[N];
 constexpr Sieve() : sieve() {
  sieve[0] = sieve[1] = 1;
  for (int i = 2; i * i < N; i++)
if (!sieve[i])
    for (int_j = i * 2; j < N; j += i)
      sieve[j] = 1;
bool isPrime(int n) {
   static constexpr Sieve<MAXN> s;
 return !s.sieve[n];
Simpson's / Approximate Integrals
// integrate f from a to b, k iterations 

// error <= (b-a)/18.0 * M * ((b-a)/2k)^4 

// where M = max(abs(f^{```}(x))) for x in [a,b] 

// "f" is a function "double func(double x)"
double Simpsons (double a, double b, int k,
 \rightarrow double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;
 for (int i = 0; i < k; i++)
  t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
 \rightarrow (*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
 // ax^2 + bx + c = 0, find x
vector<double> solveEq(double a, double b,
 → double c) {
vector<double> r;
 double z = b * b - 4 * a * c;
if (z == 0)
 r.push_back(-b/(2*a));
 else if (z > 0) {
r.push back((sqrt(z)-b)/(2*a));
  r.push_back((sqrt(z)+b)/(2*a));
 return r;
\frac{1}{2} / ax^3 + bx^2 + cx + d = 0, find x
vector < double > solveEq (double a, double b,

→ double c. double d) {
```

```
vector<double> res:
 long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
\rightarrow -2*sqrt(q);
long double r = (2*a1*a1*a1 - 9*a1*a2 +
\rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z \le 0) {
 theta = acos(r/sqrt(q*q*q));
res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push_back(sq*cos((theta+2.0*PI)/3.0) -
\rightarrow a1/3.0):
 res.push_back(sq*cos((theta+4.0*PI)/3.0) -
\rightarrow a1/3.0);
 else {
 res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
 res[0] = (res[0] + q / res[0]) *
\rightarrow ((r<0)?1:-1) - a1 / 3.0;
 return res;
\frac{1}{2}// linear diophantine equation ax + by = c,
\hookrightarrow find x and y
// infinite solutions of form x+k*b/g, y-k*a/g bool solveEq(ll a, ll b, ll c, ll &x, ll &y, ll
 g = egcd(abs(a), abs(b), x, y);
 if (c % g) return false;
x *= c / g * ((a < 0) ? -1 : 1);

y *= c / g * ((b < 0) ? -1 : 1);
 return true;
}// m = \# equations, n = \# variables, a[m][n+1]
\rightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
\rightarrow a[i][n+1]
// find a solution of some kind to linear
const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a >
→ -eps): }
vector<double> solveEq(double **a, int m, int
 int cur = 0;
for (int i = 0; i < n; i++) {
    for (int j = cur; j < m; j++) {
        if (!zero(a[j][i])) {
    if (j != cur) swap(a[j], a[cur]);
for (int sat = 0; sat < m; sat++) {</pre>
   ...if (sat == cur) continue;
  double num = a[sat][i] / a[cur][i];
for (int sot = 0; sot <= n; sot++)
a[sat][sot] -= a[cur][sot] * num;
    .}
.cur++;
    break
 for (int j = cur; j < m; j++)
 if (!zero(a[i][n])) return vector <double > ():
 vector<double > ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
    if (sat < m && !zero(a[sat][i]))
        ans[i] = a[sat][n] / a[sat++][i];
    return ans;
}
// solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions. -1 is no
→ solutions
// `alls` is whether to find all solutions, or
const double eps = 1e-12:
int solveEq(Vec<2, double>& A, Vec<1, double>&
\rightarrow b, Vec<1, double>& x, bool alls=false)
                                                               return (y+y/4-y/100+y/400+cal[m-1]+d)%7;
int n = A.size(), m = x.size(), rank = 0, br,
vector<int> col(m); iota(begin(col), end(col), Unix/Epoch Time
```

```
for(int i = 0; i < n; i++) {
   double v, bv = 0;
   for(int r = i; r < n; r++)</pre>
  for(int j = i; j < n; j++)
if (fabs(b[j]) > eps)
      return -1:
   swap(A[i], A[br]);
   swap(b[i], b[br]);
   swap(col[i], col[bc]);
  swap(cofi); Cofice();
for(int j = 0; j < n; j++)
  swap(A[j][i], A[j][bc]);
bv = 1.0 / A[i][i];
for(int j = (alls)?0:i+1; j < n; j++) {
    if (j != i) {</pre>
     double fac = A[j][i] * bv;
     b[i] = fac * b[i];
    for(int k = i+1; k < m; k++)
A[j][k] -= fac*A[i][k];
  rank++;
 if (alls) for (int i = 0; i < m; i++) x[i] =
 → -DBL_MAX;
 for (int i = rank; i--;) {
  bool isGood = true;
   for (int j = rank; isGood && j < m; j++)
    if (fabs(A[i][j]) > eps)
   isGood = false;
b[i] /= A[i][i];
  if (isGood) x[col[i]] = b[i];
if (!alls)
    for(int j = 0; j < i; j++)
    b[j] -= A[j][i] * b[i];</pre>
 return rank;
Gravcode Conversions
ull graycode2ull(ull n) {
 ull i = 0:
 for (; n; n = n >> 1) i ^= n;
ull ull2graycode(ull n) {
 return n ^ (n >> 1);
Date Utilities
// handles -4799-01-01 to 1465001-12-31
int date2int(int y, int m, int d){
 return 1461*(y+4800+(m-14)/12)/4+367*(m-2-(m
     -14)/12*12)/12-3*((y+4900+(m-14)/12)/100)
    /4+d-32075;
pair<int,pair<int,int>> int2date(int x){
 int n,i,j;
 x+=68569
 n=4*x/146097:
 x = (146097*n+3)/4;
 i=(4000*(x+1))/1461001:
 x=1461*i/4-31;
 j=80*x/2447;
 return \{100*(n-49)+i+j/11, \{j+2-12*(j/11), \}
    x-2447*i/80};
int dayOfWeek(int y, int m, int d){      //0=sunday
    static int cal[]={0,3,2,5,0,3,5,1,4,6,2,4};
```

```
|// 0-indexed month/time, 1-indexed day
// minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
 → hour, int minute, int second) {
 struct tm t; time_t epoch;
t.tm_year = year - 1900; t.tm_mon = month;
t.tm_mday = day; t.tm_hour = hour;
 t.tm_min = minute; t.tm_sec = second;
t.tm_isdst = 0: // 1 = daylights savings
 epoch = mktime(&t);
 return (ull)epoch;
vector<int> toDate(ull epoch) {
 time t e=epoch; struct tm t=*localtime(&e);
 return {t.tm_year+1900,t.tm_mon,t.tm_mday,t|
    .tm hour,t.tm min,t.tm sec};
int getWeekday(ull epoch) {
 time t e=epoch; struct tm t=*localtime(&e);
 return t.tm_wday; // 0-6, 0 = sunday
int getDavofYear(ull epoch) {
 time t e=epoch: struct tm t=*localtime(&e):
 return t.tm_yday; // 0-365
const int months[] =
 \rightarrow {31,28,31,30,31,30,31,30,31,30,31};
bool validDate(int year, int month, int day) {
     bool leap = !(year%(year%25?4:16));
     if (month >= 12) return false:
     return day <= months[month] + (leap &&
 \rightarrow month == 1):
```

Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))1+!(n-2), $!n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}$, $!n = [\frac{n!}{e}]$

In a partially ordered set, a chain is a subset of elements that are all comparable to each other An antichain is a subset where no two are comparable.

Dilworth's theorem states the size of a max-Dijkstra's imal antichain equals the size of a minimal chain cover of a partially ordered set S. The // use add_edge(..., true) for digraphs width of S is the maximum size of an antichain void add_edge(Vec<2, edge> &graph, int u, int in S, which is equal to the minimum number of chains needed to cover S, or the minimum of chains needed to cover S, or the minimum if (!directed) graph[v].push_back({v,u,w}); number of chains such that all elements are in at least one chain.

|Rosser's Theorem states the nth prime| \hookrightarrow number is greater than n * ln(n) for n > 1.

Nicomachi's Theorem states $1^3 + 2^3 + ... +$ $n^3 = (1 + 2 + ... + n)^2$ and is equivalent to $(n^{\frac{n+1}{2}})^2$.

Lagrange's Four Square Theorem states every natural number is the sum of the squares of four non-negative integers. This is a special case of the Fermat Polygonal Number **Theorem** where every positive integer is a sum of at most n s-gonal numbers. The nth| s-gonal number $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$ | }

```
Graphs
struct edge {
 int u,v,w;
 edge (int u.int v.int w) : u(u).v(v).w(w) {}
 edge (): u(0), v(0), w(0) {}
bool operator < (const edge &e1, const edge
bool operator > (const edge &e1, const edge
\rightarrow &e2) { return e1.w > e2.w; }
struct subset {
  int p, rank, sz;
 subset(int p) : p(p), rank(0), sz(1) {}
 subset() : p(0), rank(0), sz(0) {}
 void make_set(int _p) { p=_p, rank=0, sz=1; }
// adjacency list named 'graph'
int visited[MAX];
int parent [MAX];
int \bar{v}c = 0;
```

```
vector<int> bfs(int start, int end) {
    visited[start] = vc;
    parent[start] = -1;
    queue<int> q;
    q.push(start);
while (!q.empty()) {
         int v = q.front(); q.pop();
for (auto e : graph[v]) {
              if (visited[e] != vc) {
   visited[e] = vc;
                   q.push(e);
                   parent[e] = v;
                   if (e == end) goto DONE;
    DONE:
// path reconstruction
    if (visited[end] != vc) return {};
    vector<int> path;
    for (int v = end; v != -1; v = parent[v])
         path.push_back(v);
    return path;
```

```
vector<int> dijkstra(Vec<2, edge> &graph, int
   src) {
    vector<int> D(graph.size(), inf);
    priority_queue < edge, vector < edge >,
    greater < edge >> pq;
    pq.push({src,src,0});
D[src]=0;
    while(!pq.empty()) {
        edge e = pq.top(); pq.pop();
        for(int i=0;i<graph[v].size();i++) {</pre>
             int u = graph[v][i].v;
             if(D[v] + graph[v][i].w < D[u]) {
                D[u] = D[v] + graph[v][i].w;
                pq.push({src,u,D[u]});
    return D:
```

```
Eulerian Path
                                                                  for (edge e : edges)
#define edge_list vector<edge>
#define adj_sets vector<set<int>>
                                                                       if (D[e.u] != inf && D[e.u] + e.w < |
                                                                                                                      return false:
                                                             D[e.v])
                                                                                                                 |vector<pair<int, int>> bipartite_match(Vec<2,</pre>
struct EulerPathGraph {
                                                                            D[e.v] = -inf;
                                                              return D;
                                                                                                                      int> &G, int m) {
 adj_sets graph; // actually indexes incident
                                                                                                                      vector<int> L(G.size(), -1), R(m, -1);
                                                                                                                      V.reset();
bool running = true;
 edge_list edges; int n; vector<int> indeg;
                                                         Minimum Spanning Tree
 EulerPathGraph(int n): n(n) {
                                                                                                                      while (running) {
                                                            returns vector of edges in the mst
  indeg = *(new vector<int>(n,0));
                                                                                                                           running = false;
                                                            graph[i] = vector of edges incident to
  graph = *(new adj_sets(n, set<int>()));
                                                                                                                           V.reset():
                                                                                                                           for (int i=0;i<L.size();i++)
if (L[i] == -1)
                                                             places total weight of the mst in Stotal
 void add_edge(int u, int v) {
  graph[u].insert(edges.size());
indeg[v]++;
                                                          // if returned vector has size != n-1, there is
                                                                                                                                     running |= match(i, G, R, L);
                                                             no MST
                                                         vector<edge> mst(Vec<2, edge> graph, 11
                                                                                                                      vector<pair<int,int>> ret;
  edges.push_back(edge(u,v,0));
                                                             &total) {
tal = 0:
                                                                                                                      for (int i = 0; i < L.size(); ++i)
    if(L[i]!=-1) ret.push_back({i, L[i]});</pre>
 bool eulerian_path(vector<int> &circuit) {
  if(edges.size()==0) return false;
                                                          priority_queue<edge, vector<edge>,

→ greater<edge>> pq;

  stack<int> st;
int a[] = {-1, -1};
                                                          vector<edge> MST;
                                                          bitset<20001> marked; // change size as needed 2D Grid Shortcut
  for(int v=0; v<n; v++) {
    if(indeg[v]!=graph[v].size()) {
                                                          marked[0] = 1;
                                                                                                                  #define inbound(x,n) (0<=x&&x<n)
                                                          for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
                                                                                                                  #define fordir(x,y,n,m) for(auto[dx,dy]:dir)if
    bool b = indeg[v] > graph[v].size();
                                                                                                                  \hookrightarrow (inbound(x+dx,n)&\mathbb{U}inbound(y+dy,m))
   . if (abs(((int)indeg[v])-((int)graph[v])
                                                             pg.size()!=0) {
                                                                                                                  const pair<int,int> dir[] =
     .size())) > 1) return
                                                           edge e = pq.top(); pq.pop();
                                                                                                                  \hookrightarrow {{1,0},{0,1},{-1,0},{0,-1}};
    false;
if (a[b] != -1) return false;
                                                           int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
else if(marked[u]) swap(u, v);
  [a[b] = v;
                                                                                                                       2D Geometry
                                                                                                                 #define point complex<double>
#define EPS 0.0000001
                                                           for(edge ep : graph[u]) pq.push(ep);
                                                           marked[u] = 1;
MST.push_back(e);
  int s = (a[0]!=-1 \&\& a[1]!=-1 ? a[0] :
                                                                                                                 #define sq(a) ((a)*(a))
#define cb(a) ((a)*(a)*(a))
 \rightarrow (a[0]==-1 && a[1]==-1 ? edges[0].u : -1));
                                                           total += e.w;
  if(s==-1) return false;
                                                                                                                  double dot(point a, point b) { return
  while(!st.empty() || !graph[s].empty()) {
   if (graph[s].empty()) {
                                                          return MST;

    real(conj(a)*b); }

                                                                                                                  double cross(point a, point b) { return
    circuit.push back(s); s = st.top();
                                                         Union Find
                                                        Union Find
int uf find(subset* s, int i) {
   if (s[i].p!= i) s[i].p = uf_find(s, s[i].p);
   return s[i].p;
}

struct line { point a, p; r;
   struct circle { point c; double r; };
   struct segment { point a, b; };
   struct triangle { point a, b, c; };
   struct triangle { point a, b, c; };
}

    imag(conj(a)*b); }

    st.pop(); }
   else {
    int w = edges[*graph[s].begin()].v;
    graph[s].erase(graph[s].begin());
    st.push(s); s = w;
                                                                                                                  struct rectangle { point tl, br; };
                                                          int xp = uf_find(s, x), yp = uf_find(s, y);
if (s[xp].rank > s[yp].rank) s[yp].p = xp,
                                                                                                                  struct convex_polygon {
  circuit.push back(s):
                                                                                                                   vector<point> points;
                                                             s[xp].sz += s[vp].sz;
                                                                                                                   convex_polygon(vector<point> points) :
  return circuit.size()-1==edges.size():
                                                          else if (s[xp].rank < s[yp].rank) s[xp].p =
                                                                                                                   → points(points) {}
                                                             yp, s[yp].sz += s[xp].sz;
                                                                                                                   convex_polygon(triangle a) {
                                                          else s[yp].p = xp, s[xp].rank++, s[xp].sz +=
                                                                                                                    points.push_back(a.a); points.push_back(a.b);
Flovd Warshall
                                                             s[yp].sz;
                                                                                                                     points.push_back(a.c);
const ll inf = 1LL << 62;
#define FOR(i,n) for (int i = 0; i < n; i++)
                                                         void uf_size(subset *s, int i) {
  return s[uf find(s, i)].sz;
                                                                                                                   convex_polygon(rectangle a) {
void floydWarshall(Vec<2, 11>& m) {
                                                                                                                    points.push_back(a.tl);
 int n = m.size();
                                                                                                                      points.push_back({real(a.tl),
 FOR(i,n) m[i][i]' = min(m[i][i], OLL);
FOR(k,n) FOR(i,n) FOR(j,n) if (m[i][k] != inf Bipartite Graph
                                                                                                                      imag(a.br)}):
                                                                                                                    points.push_back(a.br);
 auto newDist = max(m[i][k] + m[k][j], -inf);
                                                         A bipartite graph has "left" and "right" set of
                                                                                                                      points.push back({real(a.br).
                                                         \rightarrow nodes
Every edge has an endpoint in each set (L/R)
  m[i][j] = min(m[i][j], newDist);
                                                                                                                      imag(a.tl)});
                                                         A matching is a subset of all edges
 FOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n)
                                                         Such that each vertex is an endpoint
 . if (m[i][k] != inf && m[k][j] != inf)
                                                                                                                  struct polygon {
                                                         Of at most one edge in the subset
                                                                                                                   vector point points;

→ m[i][i] = -inf:
                                                         sgrt(V)*E time
                                                                                                                   polygon(vector point points) :
                                                         tested on "piano lessons"

    points(points) {}

                                                         sourced from
Bellman Ford
                                                                                                                   polygon(triangle a) {
const int inf = 20000001; vector<11> bellman_ford(vector<edge> edges, int */ #define MAXNODES 1001
                                                             https://codeforces.com/blog/entry/58048
                                                                                                                    points.push_back(a.a); points.push_back(a.b);
                                                                                                                      points.push back(a.c);
\hookrightarrow src, int V) {
                                                         bitset<MAXNODES> V;
bool match(int node, Vec<2,int> &G, vector<int>
                                                                                                                  polygon(rectangle a) {
     vector<ll> D(V,inf);
    D[src] = 0;
for (int i=1:i<=V-1:i++)
                                                                                                                    points.push_back(a.tl);
                                                             &R, vector<int> &L) {
if (V[node]) return false;
                                                                                                                      points.push_back({real(a.tl),
          for (edge e : edges)
                                                              V[node] = 1;
                                                                                                                      imag(a.br)});
              if (D[e.u] != inf && D[e.u] + e.w <
                                                              for(auto vec : G[node]) {
   if (R[vec] == -1 || match(R[vec], G, R,
                                                                                                                    points.push_back(a.br);
 \hookrightarrow D[e.v])
                   D[e.v] = D[e.u] + e.w;
                                                                                                                      points.push_back({real(a.br),
                                                             L)) {
     // detect negative cycles: *typically* 2 is
                                                                                                                      imag(a.tl)});
                                                                       L[node] = vec; R[vec] = node;
    as good as V-1 for this
    for (int i=1;i<=V-1;i++)
                                                                                                                  polygon(convex polygon a) {
```

```
for (point v : a.points)
  points.push_back(v);
   triangle methods
double area_heron(double a, double b, double
 \rightarrow c) {
if (a < b) swap(a, b);
 if (a < c) swap(a, c);
 if (b < c) swap(b, c);
 if (a > b + c) return -1;
return \operatorname{sqrt}((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)

    √16.0):

// segment methods
double lengthsq(segment a) { return
    sq(real(a.a) - real(a.b)) + sq(imag(a.a) -
    imag(a.b)): }
double length(segment a) { return

    sqrt(lengthsq(a)): }

// circle methods
double circumference(circle a) { return 2 * a.r
→ * M_PI; }
double area(circle a) { return sq(a.r) * M_PI;
// rectangle methods
double width(rectangle a) { return
   abs(real(a.br) - real(a.tl)); }
double height (rectangle a) { return

    abs(imag(a.br) - real(a.tl)); }

double diagonal(rectangle a) { return

    sqrt(sq(width(a)) + sq(height(a))); }

double area (rectangle a) { return width(a) >
→ height(a): }
double perimeter(rectangle a) { return 2 *
   (width(a) + height(a)); }
   check if 'a' fit's inside 'b
// swap equalities to exclude tight fits
bool doesFitInside(rectangle a, rectangle b) {
 int x = width(a), w = width(b), y = height(a),
\rightarrow h = height(b):
if (x > y) swap(x, y);
 if (w > h) swap(w, h);
 if (w < x) return false:
 if (y <= h) return true;
 double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
 return sq(a) \le sq(b) + sq(c):
// polygon methods
// negative area = CCW, positive = CW
double area(polygon a) {
 double area = 0.0; int n = a.points.size();
for (int i = 0, j = 1; i < n; i++, j = (j +
\rightarrow 1) % n)
area += (real(a.points[j]-a.points[i]))*
 (imag(a.points[j]+a.points[i]));
return area / 2.0;
// get both unsigned area and centroid
pair<double, point> area centroid(polygon a) {
 int n = a.points.size();
 double area = 0:
 point c(0, 0);
 for (int i = n - 1, j = 0; j < n; i = j++) {
  double v = cross(a.points[i], a.points[j]) /
 \stackrel{\rightarrow}{\text{area}} += v:
 c += (a.points[i] + a.points[j]) * (v / 3);
 c /= area:
 return {area, c};
```

```
while (lower.size() >= 2 &&
Intersection
                                                       cross(lower.back() - lower[lower.size() -
// -1 coincide, 0 parallel, 1 intersection
int intersection(line a, line b, point& p) {
                                                       2], a.points[i] - lower.back()) < EPS)
                                                      lower.pop_back();
if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
                                                     while (upper.size() >= 2 &&
p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
- - a.a, b.b - b.a) * (b - a) + a;
...return 1;
                                                       cross(upper.back() - upper[upper.size() -
                                                       2], a.points[i] - upper.back()) > -EPS)
                                                      upper.pop back();
 if (abs(cross(a.b - a.a, a.b - b.a)) > EPS)
                                                     lower.push_back(a.points[i]);
→ return 0:
                                                     upper.push_back(a.points[i]);
return -1:
                                                    lower.insert(lower.end(), upper.rbegin() + 1,
// area of intersection
                                                       upper.rend());
double intersection(circle a, circle b) {
                                                    return convex_polygon(lower);
double d = abs(a.c - b.c):
if (d <= b.r - a.r) return area(a);
if (d <= a.r - b.r) return area(b);</pre>
if (d \ge a.r + b.r) return 0;
double alpha = acos((sq(a.r) + sq(d) - b.r))
                                                   Maximum Colinear Points
                                                   const ll range = 10000;
\rightarrow sq(b.r)) / (2 * a.r * d));
double beta = acos((sq(b.r) + sq(d) - sq(a.r)))/*
                                                   given N points with coordinates in [-range,
compute the size of the largest subset of the
    alpha)) + sq(b.r) * (beta - 0.5 * sin(2 *
   beta));
                                                   for which all points in the subset are
                                                   \mapsto colinear.
Time: O(N^2 \log(range)) average
// -1 outside, 0 inside, 1 tangent, 2
int intersection(circle a, circle b,
                                                   Solves https://open.kattis.com/problems

    vector<point>& inter) {

                                                       /maxcolinear in 0.62s (not
double d2 = norm(b.c - a.c), rS = a.r + b.r.
                                                       amazina)
→ rD = a.r - b.r;
if (d2 > sq(rS)) return -1;
                                                   struct Slope { // a rational number with
if (d2 < sq(rD)) return 0;
                                                       unsigned infinity (1.0)
double ca = 0.5 * (1 + rS * rD / d2);
                                                       11 p, q;
Slope(11 pP=0, 11 qP=0) {
point z = point(ca, sqrt(sq(a.r) / d2 -
\rightarrow sq(ca));
                                                           if(qP==0) {
inter.push_back(a.c + (b.c - a.c) * z);
                                                                p = 1, q = 0;
if (abs(imag(z)) > EPS) inter.push_back(a.c +
                                                                return;
\rightarrow (b.c - a.c) * coni(z)):
                                                            11 g = \_gcd(pP, qP);
return inter.size():
                                                           pP /= g, qP /= g;
if(qP < 0) pP *= -1, qP *= -1;
// points of intersection
                                                           p = pP, q = qP;
vector<point> intersection(line a, circle c) {
vector<point> inter;
c.c -= a.a;
a.b -= a.a;
                                                       bool operator == (const Slope &other) const
point m = a.b * real(c.c / a.b);
                                                            return other.p == p && other.q == q;
double d2 = norm(m - c.c);
if (d2 > sq(c.r)) return 0;
                                                   namespace std {
 double l = sqrt((sq(c.r) - d2) / norm(a.b));
                                                       template<>
 inter.push back(a.a + m + 1 * a.b);
                                                       struct hash<Slope> { // typical
if (abs(1) > EPS) inter.push back(a.a + m - 1
                                                       rectangular/lattice hash
\rightarrow * a.b);
                                                           size_t operator() (const Slope &r)
return inter;
}
// area of intersection
                                                       const {
                                                                return (2*range+1) * (r.p + range)
double intersection(rectangle a, rectangle b) {
                                                       + r.q + range;
double x1 = max(real(a.tl), real(b.tl)), y1

→ max(imag(a.tl), imag(b.tl));
double x2 = min(real(a.br), real(b.br)), y2 =
                                                   int max_colinear_points(vector<pair<11,11>>

→ min(imag(a.br), imag(b.br));
                                                       &points) {
return (x2 \le x1 \mid y2 \le y1) ? 0:
                                                       if(points.size() <= 2) return
\hookrightarrow (x2-x1)*(y2-y1);
                                                       points.size();
                                                       int best = 0;
unordered_map<Slope, int> counter;
Convex Hull
                                                       for(int i=0:i<points.size():i++) {</pre>
bool cmp(point a, point b) {
                                                           for(int j=i+1; j<points.size(); j++) {</pre>
if (abs(real(a) - real(b)) > EPS) return
                                                                Slope

→ real(a) < real(b);
</pre>
                                                       slope(points[i].second-points[j] |
if (abs(imag(a) - imag(b)) > EPS) return
                                                       .second,points[i].first-points[j].first);

    imag(a) < imag(b);
</pre>
                                                                best = max(best,
return false:
                                                       ++counter[slope]+1);
convex_polygon convexhull(polygon a) {
                                                            if(i != points.size()-1)
sort(a.points.begin(), a.points.end(), cmp);
vector<point> lower, upper;
                                                       counter.clear();
for (int i = 0; i < a.points.size(); i++) {
```

```
return best:
     3D Geometry
struct point3d {
 double x, y, z;
 point3d operator+(point3d a) const { return
 \rightarrow {x+a.x, y+a.y, z+a.z}; }
 point3d operator*(double a) const { return
 \rightarrow {x*a, y*a, z*a}; }
 point3d operator-() const { return {-x. -v.
    -z}; }
 point3d operator-(point3d a) const { return
 \rightarrow *this + -a: }
 point3d operator/(double a) const { return
 double abs() { return sqrt(norm()); }
 point3d normalize() { return *this /
 \rightarrow this->abs(): }
double dot(point3d a, point3d b) { return
\rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
point3d cross(point3d a, point3d b) { return
    \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
a.x*b.y - a.y*b.x}; }
struct line3d { point3d a, b; };
struct plane { double a, b, c, d; } // a*x +
\Rightarrow b*y + c*z + d = 0
struct sphere { point3d c; double r; };
#define sq(a) ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
double surface(circle a) { return 4 * sq(a.r)
double volume(circle a) { return 4.0/3.0 *
\hookrightarrow cb(a.r) * M_PI; }
10 Optimization
Snoob
// SameNumberOfOneBits, next permutation
int snoob(int a) {
  int b = a & -a, c = a + b;
  return c | ((a ^ c) >> 2) / b;
 // example usage
// example usage int main() {
   char l1[] = {'1', '2', '3', '4', '
   char l2[] = {'a', 'b', 'c', 'd'};
   int d1 = 5, d2 = 4;
   // prints 12345abcd, 1234a5bcd, ...
 int min = (1 << d1) - 1, max = min << d2;
 for (int i = min; i <= max; i = snoob(i)) {
   int p1 = 0, p2 = 0, v = i;
   while (p1 < d1 || p2 < d2) {
   cout << ((v & 1) ? l1[p1++] : l2[p2++]);
  v /= 2;
  cout << '\n':
bool isPowerOf2(ll a) {
return a > 0 && !(a & a-1);
bool isPowerOf3(11 a) {
  return a>0&&!(12157665459056928801ull%a);
bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
 return abs(x-round(x)) < 0.00000000001:
       Python
Recursion Limit Removal (Basic)
import sys
                                                            return n + judge(n + 1);
sys.setrecursionlimit(10**6)
```

```
Recursion Limit Removal (Advanced)
# @bootstrap over recursive function
# replace 'return' with 'yield'
# for when sys method does not work
from types import GeneratorType
def bootstrap(f, stack=[]):
 def wrappedfunc(*args, **kwargs):
 if stack:
.return f(*args, **kwargs)
   to = f(*args, **kwargs)
   while True:
...if type(to) is GeneratorType:
     stack.append(to)
    to = next(to)
     stack.pop()
     if not stack:
to = stack[-1].send(to)
return to
return wrappedfunc
# EXAMPLE recursive fibonacci
def f(n):
   if (n < 2):
    yield n</pre>
yield (yield f(n-1)) + (yield f(n-2))
Python 3 Compatibility
from _future_ import division, print_function if sys.version_info[0] < 3:
from _builtin_ import xrange as range from future_builtins import ascii, filter,
\hookrightarrow hex, map, oct, zip
12 Additional
Judge Speed
// kattis: 0.50s
// codeforces: 0.421s
// atcoder: 0.455s
#include <bits/stdc++.h>
using namespace std;
\frac{1}{\text{int}} v = 1e9/2, p = 1;
int main() {
for (int i = 1; i <= v; i++) p *= i; cout << p;
Judge Pre-Contest Checks
     int128 and float128 support?
does extra or missing whitespace cause WA?
documentation up to date?
printer usage available and functional?
// each case tests a different fail condition
// try them before contests to see error codes
struct g { int arr[1000000]; g(){}};
vector<g> a;
// O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
→ 6=SIGSEGV 7=recursive MLE int judge(int n) {
if (n == 0) exit(0);
if (n == 1) while(1);
if (n == 2) while(1) a.push_back(g());
 if (n == 3) while(1) putchar_unlocked('a');
 if (n == 4) assert(0);
if (n == 5) 0 / 0;
if (n == 6) * (int*)(0) = 0:
```

```
GCC Builtin Docs
// 128-bit integer
_int128 a;
unsigned _int128 b;
// 128-bit float
// 128-bit float
// minor improvements over long double
float128 c;
// log2 floor
lg(n);
// number of 1 bits
// can add ll like popcountll for long longs
builtin_popcount(n);
// number of trailing zeroes
builtin_ctz(n);
// number of leading zeroes
builtin_ctz(n);
// 1-indexed least significant 1 bit
builtin_ffs(n);
__builtin_ffs(n);
// parity of number
__builtin_parity(n);
 Limits
                              int
  uint
            \pm 9223372036854775807 | \pm 2^{63} - 1|10^{18}
Complexity classes input size (per second):
 O(n^n) or O(n!)
                                                                       n \leq 10
 O(2^n)
                                                                      n \leq 30
 O(n^3)
                                                                  n < 1000
 O(n^2)
                                                                n \le 30000
                                                                    n \le 10^6
n \le 10^7
 O(n\sqrt{n})
 O(n \log n)
```

 $n < 10^9$

O(n)