# IS5102 Database Management Systems

Lecture 10: Intermediate SQL

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(with thanks to Susmit Sarkar)

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# This week: More SQL

- ► Modifying Data
- Orderings and Aggregates
- Set Operations
- ► Integrity Constraints
- Views

- ► Set operations UNION, INTERSECT, and EXCEPT
  - ▶ Each of the above operations automatically eliminates duplicates
- ▶ To retain all duplicates use the corresponding multi-set versions

```
UNION ALL,
INTERSECT ALL
EXCEPT ALL
```

ightharpoonup Suppose a tuple occurs m times in  ${f r}$  and n times in  ${f s}$ , then, it occurs:

- ightharpoonup m+n times in r UNION ALL s
- $ightharpoonup \min(m,n)$  times in r INTERSECT ALL s
- $ightharpoonup \max(0,m-n)$  times in r EXCEPT ALL s

Remark: SQLite only supports UNION ALL, but not the other two multi-set versions

▶ Find courses that ran in Semester 2 of 2019 or in Semester 1 of 2020

```
SELECT course_id

FROM course_runs

WHERE semester = 2

AND year = 2019

UNION

SELECT course_id

FROM course_runs

WHERE semester = 1

AND year = 2020;
```

▶ Find courses that ran in Semester 2 of 2019 and in Semester 2 of 2020

```
SELECT course_id

FROM course_runs

WHERE semester = 2

AND year = 2019

INTERSECT

SELECT course_id

FROM course_runs

WHERE semester = 2

AND year = 2020;
```

▶ Find courses that ran in Semester 1 of 2019 but not in Semester 1 of 2020

```
SELECT course_id
FROM course_runs
WHERE semester = 1
AND year = 2019
EXCEPT
SELECT course_id
FROM course_runs
WHERE semester = 1
AND year = 2020;
```

We can allow a default value to be specified

#### Example:

```
CREATE TABLE student (
stud_id CHAR(5),
name VARCHAR(20) NOT NULL,
dept_id VARCHAR(20),
tot_cred NUMERIC(3,0) DEFAULT 0,
PRIMARY KEY (stud_id),
FOREIGN KEY (dept_id) REFERENCES department);
```

- ► The default value of tot\_cred is set to 0
- ▶ When a tuple is inserted if no value is provided its value is set to 0

# Referential Integrity

Ensuring that a value that appears in one relation for a given set of attributes also appears for a certain set of attributes in another relation.

Example: If MATH is a department ID appearing in one of the tuples in the instructor relation, then MATH also appears in some tuple in the department relation.

#### Formal Definition:

Let A be a set of attributes. Let R and S be two relations that contain attributes A and where A is the primary key of S. A is said to be a foreign key of R if for any values of A appearing in R these values also appear in S

# Cascading Actions in Referential Integrity

```
CREATE TABLE course (
 course_id VARCHAR(8),
 title VARCHAR(50),
 dept_id VARCHAR(20).
 credits NUMERIC(2,0),
 PRIMARY KEY (course_id),
 FOREIGN KEY (dept_id) references department
 );
```

# Cascading Actions in Referential Integrity

```
CREATE TABLE course (
   course_id VARCHAR(8),
  title VARCHAR(50),
  dept_id VARCHAR(20).
  credits NUMERIC(2,0),
  PRIMARY KEY (course id).
  FOREIGN KEY (dept_id) references department
   ON DELETE CASCADE
   ON UPDATE CASCADE
 ):
Alternative actions to cascade: SET NULL, SET DEFAULT
Demo: see cascading.sql
```

#### View Rationale

- ▶ In some cases, it is not desirable for all users to see the entire logical model (that is, all the actual relations stored in the database.)
- Consider a person who needs to know an instructors name and department, but not the salary. This person should see a relation described, in SQL, by

```
SELECT instr_id, instr_name, dept_id
FROM instructor;
```

- A view provides a mechanism to hide certain data from the view of certain users.
- Any relation that is not of the conceptual model but is made visible to a user as a "virtual relation" is called a view.
- ► A view is defined using the CREATE VIEW statement which has the form CREATE VIEW v AS < query expression > where <query expression> is any legal SQL expression. The view name is represented by v.

#### View Definition

- Once a view is defined, the view name can be used to refer to the virtual relation that the view generates.
- ▶ View definition is not the same as creating a new relation by evaluating the query expression.
- ▶ Rather, a view definition causes the saving of an expression; the expression is substituted into queries using the view.

```
A view of instructors without their salary
```

```
CREATE VIEW faculty AS
SELECT instr_id, instr_name, dept_id
FROM instructor;
```

Create a view of department salary totals

```
CREATE VIEW departments_total_salary(dept_code, total_salary) AS SELECT dept_id, SUM (salary) FROM instructor GROUP BY dept_id;
```

## Views Defined Using Other Views

```
CREATE VIEW acad_year_2020 AS

SELECT semester, course_id, dept_id, title
FROM course_runs NATURAL JOIN course
WHERE year = 2020;

CREATE VIEW cs_acad_year_2020 AS

SELECT semester, course_id, title
FROM acad_year_2020
WHERE dept_id= 'CS';
```

Add a new tuple to faculty view which we defined earlier

```
INSERT INTO faculty VALUES ('30765', 'James', 'CHEM');
```

This insertion must be represented by the insertion of the tuple

```
('30765', 'James', 'CHEM', NULL)
```

into the instructor relation.

**WARNING**: this feature is not supported in SQLite:

https://www.sqlite.org/omitted.html

Some updates cannot be translated uniquely:

```
CREATE VIEW instructor_info AS

SELECT instr_id, instr_name, building

FROM instructor, department

WHERE instructor.dept_id = department.dept_id;

INSERT INTO instructor_info

VALUES ('69987', 'Raul', 'Bute');

Which department, if there are multiple departments in Bute?
```

What if no department is in Bute?

## View Update restrictions

Most SQL implementations allow updates only on simple views

- ► The FROM clause has only one database relation.
- ► The SELECT clause contains only attribute names of the relation, and does not have any expressions, aggregates, or distinct specification.
- ▶ Any attribute not listed in the SELECT clause can be set to NULL
- ► The query does not have a GROUP BY or HAVING clause.

### Materialized Views

- ▶ When defining a view, simply create a physical table representing the view at the time of creation.
- ► Can be a cheaper option.
- ▶ How are updates handled to the "base" relations on which the view was defined?

Advantages	Disadvantages
Data independence	Update restriction
Improved security	Performance
Reduced complexity	
Convenience	
Customisation	
Data Integrity	

## Reading and practice

- ▶ [DBSC] Chapters 4-5, Database System Concepts, Silberschatz, Korth and Sudarshan
- ▶ [DBS] Chapter 7, Database Systems, Connolly and Begg
- ▶ [DBD] Chapters 15-16, Database Design, Watt and Eng
- Useful sites
  - http://www.w3schools.com/sql/
  - ► http://sqlzoo.net/wiki/Main\_Page